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# **DNA Codec with Polar Codes, Huffman Compression and PGP Encryption**

**IT495 - DNA Storage**

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# Literature Review

# Literature Review

- **Channel polarization: A method for constructing capacity-achieving codes for symmetric binary-input memoryless channels Erdal Arıkan, Senior Member, IEEE**

This paper by Erdal Arıkan gives an idea about the construction of Error-correcting aka Capacity-achieving Polar Codes. We understood this paper's basic definition and implementation of Polar Encoder and Decoder. Based on the base formulas given in this, we are able to design efficient algorithms, that can be run in Python.

- **Introduction to Polar Codes, NPTEL - NOC IITM**

This YouTube playlist is by Prof. Andrew Thangaraj. Through this playlist, he explains What are Polar Codes, Polar Transforms, Polar-Encoder and Polar-Decoder the help of an algorithm. We took reference of this video to get familiar with the concepts of Polar Transform, Polar Encoder and Polar Decoder for carrying out our research in DNA Codec.

- **A Brief Introduction to Polar Codes, Henry D. PfisterY.**

This paper by Henry D. Pfister has basic details on kronecker delta products, designing efficient encoding and decoding algorithm and scheming polar code based on effective channel error rates. Also provides details on analysis on error rate of codes using monte carlo simulation of BSC, BEC channels.

# Literature Review

- **Kurniawan, A. Albone and H. Rahyuwibowo, "The design of mini PGP security," Proceedings of the 2011 International Conference on Electrical Engineering and Informatics.**

This paper gives an idea about creating a PGP application which has the conventional functions of encryption, decryption and digital signature. With help, we can plan to design the algorithm with some changes made in reference to considering the original system implementation.

- **“A Characterization of the DNA Data Storage Channel”, Reinhard Heckel<sup>1</sup>, Gediminas Mikutis<sup>2</sup> & Robert N.Grass**

This study describes the DNA data storage channel and how DNA can be utilised as an archival storage devices. It also underlines the limitations and imperfections in DNA synthesis, sequencing, and manipulation. We learned from others how vital it is to have a qualitative understanding of errors and molecule loss while designing a DNA storage system. In this paper, the error probabilities have also been described by analysing experimental data.

# Encrypted Error Correcting - DNA Codec System

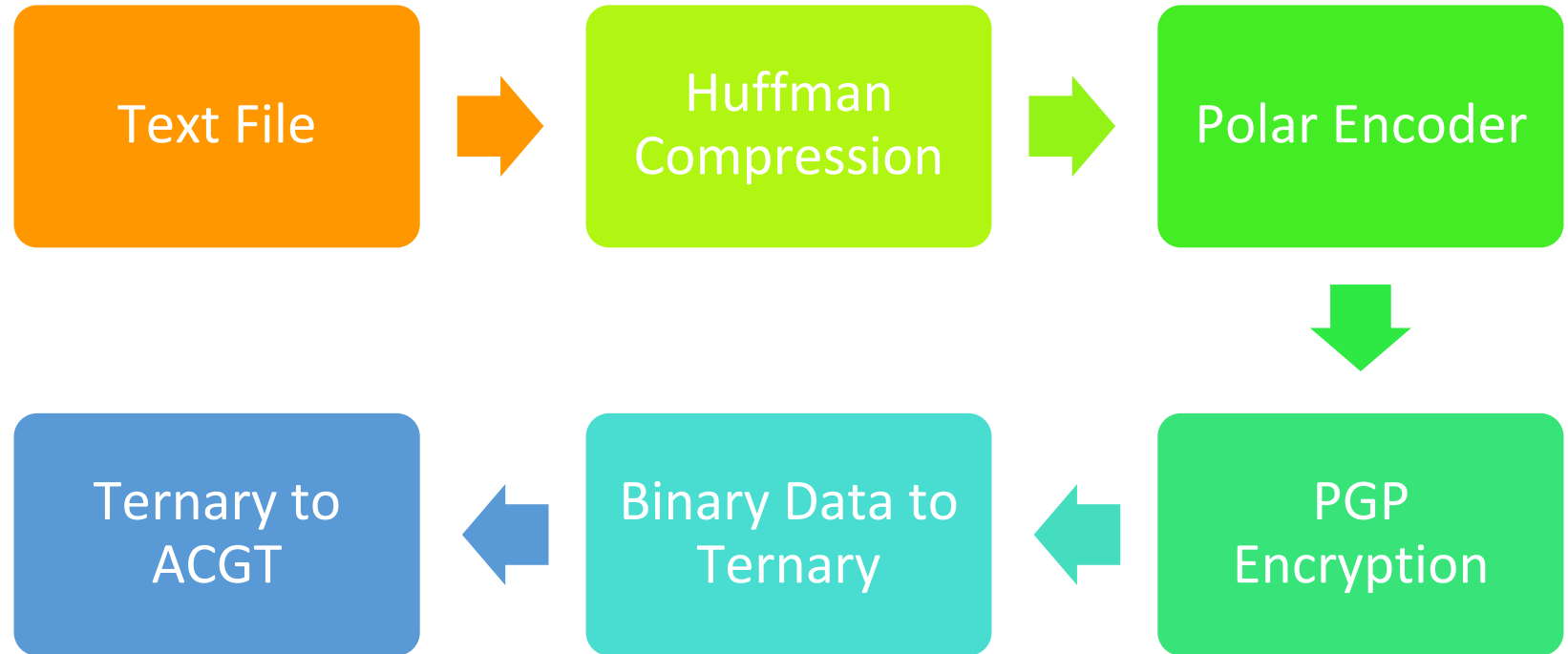
# DNA Codec

**Huffman Compression** - Compressing initial size of data is the best choice for Archival Storage like DNA!

**Polar Codes** - Linear Block Error Correcting Code, going to be used in 5G! Will provide the capability of error correctness for our DNA storage.

**PGP Encryption** – Laptops/PCs are now coming with encrypted Hard Disks, so why not have similar security for DNA storage? PGP Encryption on DNA leverages the capability of DNA of storing huge amounts of data only using a few grams. We are using the GnuPG library of python to do encryption-decryption. Kleopatra software is used for key – certificate generation.

# Encoder





**Text File:**

**Hello World!**

# Text File to Binary Data Conversion

ASCII to Binary: 01001000 01100101 01101100 01101100 01101111  
00100000 01010111 01101111 01110010 01101100 01100100 00100001

## Huffman Compression

Huffman Compression: 1100 1101 01 01 101 1110 1111 101 000 01 001  
100

# Polar Codes - Encoding

For Polar Codes Generator Matrix  $G_N$ ,

$$G_N = (I_{N/2} \otimes F) R_N (I_2 \otimes G_{N/2}), \text{ for } N \geq 2$$

Where,  $G_1 = I_1$  . Using (1) we can get  $G_2$  ,  $G_3$  , ... recursively.

We design an efficient encoding algorithm based on Generator Matrix.

# Polar Encoding Calculation


$$G_1 = [1]$$

$$G_2 = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}$$

$$G_3 = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}$$

We get generator matrix using formula described. For coding part, we are using recursion and N can vary as we are not hardcoding the generator matrix.

# Polar Encoding Pseudocode

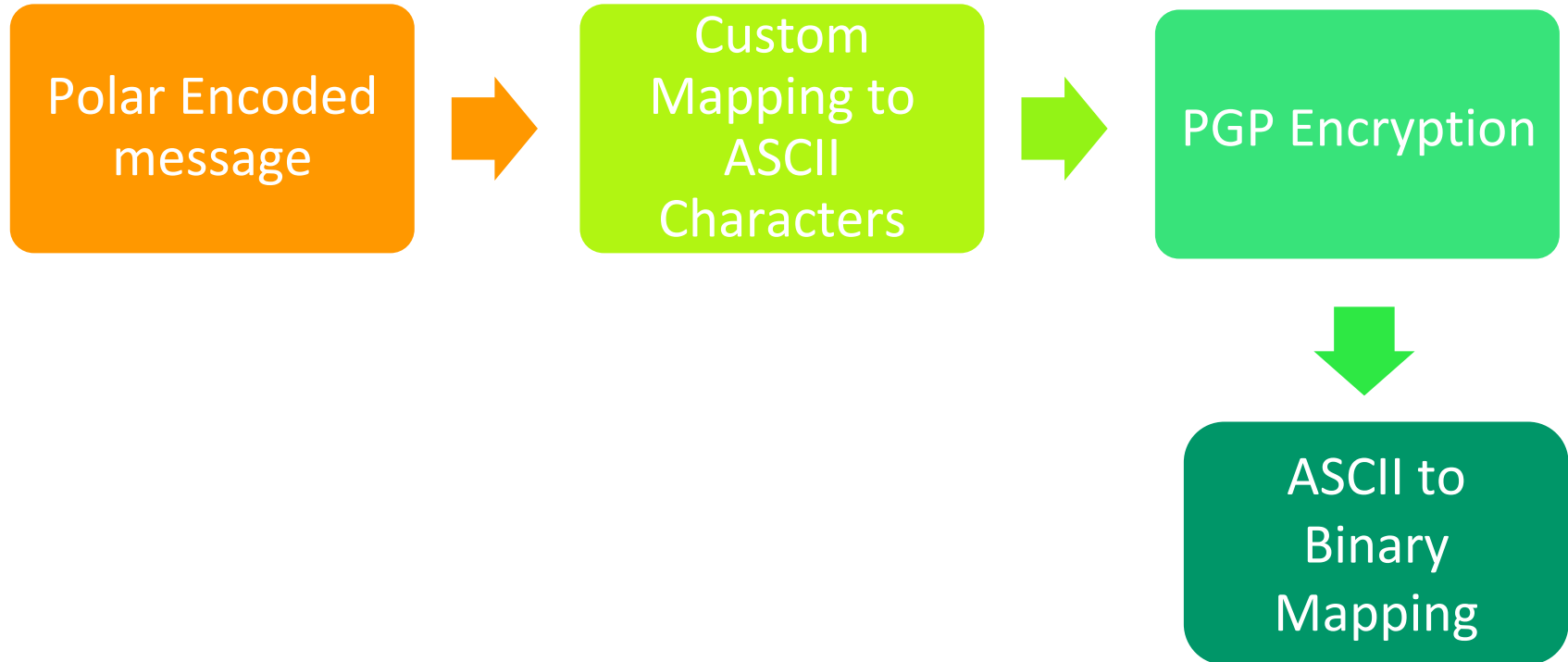
```
def polar_encoder(u):  
    if (len(u)==1):  
        x = u  
    else:  
         u1u2 = ((u[1::2]+u[2::2]) % 2)  
        u2 = u[2::2]  
        x = [polar_encoder(u1u2), polar_encoder(u2)]  
    return x
```

# Polar Encoding

Polar Encoder,  $N = 7$

```
1 1 0 1 1 0 0 0 1 1 1 0 0 1 0 0 1 0 0 0 0 0 1 0 0 1 1 1 0 0 1 0 0 1 0 0 1 1 1 0 1  
0 1 1 1 1 1 0 1 1 0 1 1 0 0 0 1 1 1 0 0 1 0 0 1 0 0 0 0 1 0 0 0 0 0 1 0 0 1 0 1 0  
1 1 1 0 0 0 1 1 1 0 0 0 1 0 1 0 0 0 0 1 0 0 1 1 0 1 1 1 1 0 0 1 1 1 0 1 0 0 1 1 1  
0 0 0 1 0
```

# PGP Encryption Block



# PGP Encryption

01101000010001100011010001000100011000100011000001000101011  
00110010000110101011001011010001100010101011101101110010100  
01010100110100000101010001011001000100000101110100001100110  
01101000011001001100010001100000101100000110000011110100101  
01110111011101010101011110000111011001110001001100010101100  
10110111101000001001101000100011001101101001100010010111101  
10010001110011011011110110011001001111011110010111001101010  
00001010001011000110100111001010001010110010110010000101111  
011000110100.....



# Binary to Ternary Conversion

$$y_1 = x_1$$

$$y_i = x_i + x_{i-1} \text{ Where } i = 2, 3, \dots, N$$

$x(i-1) \backslash x(i)$	0	1
0	0	1
1	1	2

# Binary to Ternary

01211100011001210012111001100110012100110012100001100111112  
101210110001211111210111211100121001111112211211221011110  
01111110121110000111111001112101100110000111221110001210121  
01211100012101101210011001210000111210000121000012221110111  
11221122112211111111112221000122112101221001101210011111210  
11121122211100001101211100110012101211211101210011011122211  
21011001221012112112221121012101101222112221011122101211111  
00001111001112100121110122101111001111121011121011000111222  
11210012111001210121110111211112112211221000111101211210001  
21110011001100012112100121122112111121110001101211211101221  
00012111121110110001112100011001210012101211210012111121011  
01210011110111101101221012100011011111211210001112210110110  
00121122112211211210011121110110121110121110012100110122101  
21122210001101221011000011122100111101221011000011012112211  
222111012112210012101210121011011110011110001211101222112..

# Ternary to ACGT Conversion

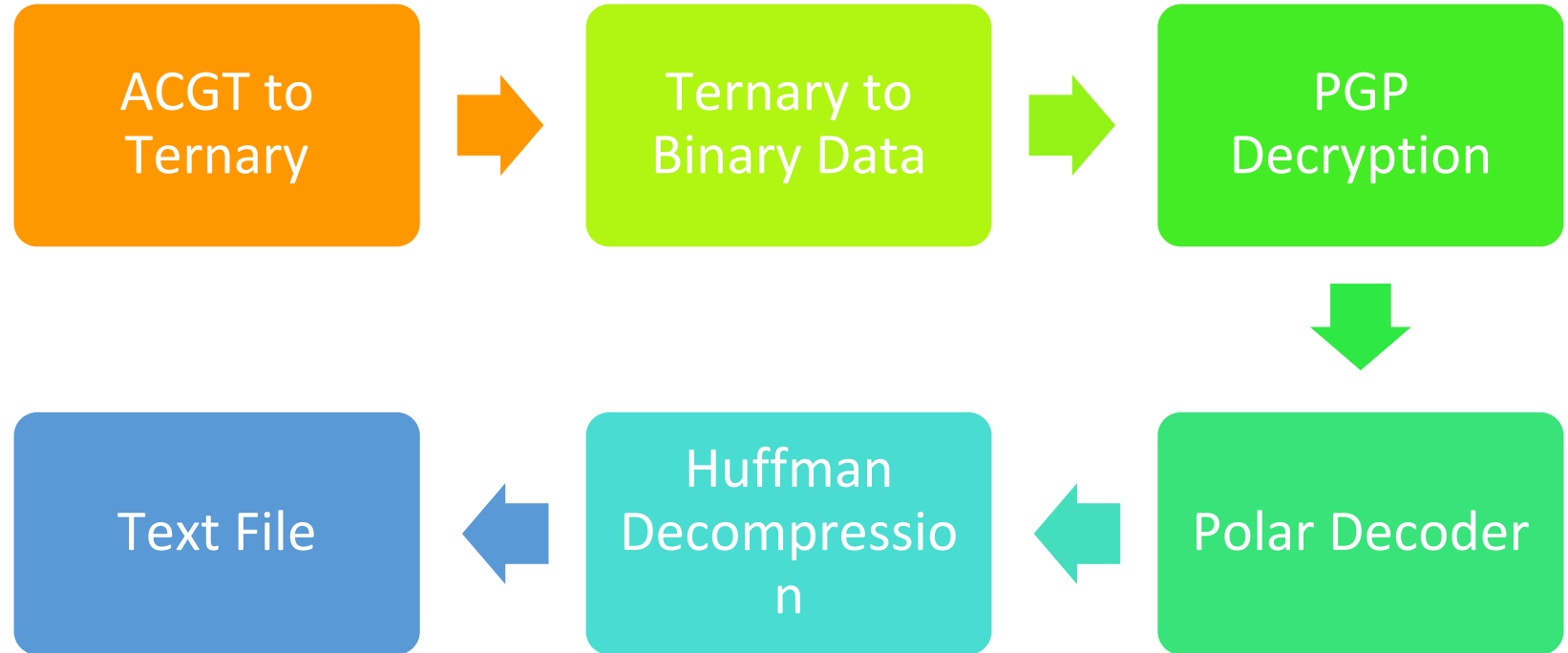
*Mapping of  $y_1 : 0 \rightarrow A, 1 \rightarrow C, 2 \rightarrow G$*

$y(i-1) \backslash x(i)$	0	1	2
A	C	G	T
C	G	T	A
G	T	A	C
T	A	C	G

# Ternary to ACGT

ACTGAGACGTCTACTGACGATCTCGTCTACTCGTCAGTAGACGATCGTACTCG  
TCTCTCAGTCAGTCTACGATCTCTCTGACTCTGAGACGATCGTCTCTCTCATCT  
GAGCAGTCTCTACTCTCTCGATCTCGTACTCTCTCGTCTCAGTCTACTCGTACT  
CTGCTCTACGATCGATCGATCTCGTAGCTAGACTGACGAGTAGCTACGTCTCA  
GTACGATCGTACTGCAGAGTCTCTCATCTGCTCATCTCTCTCTCTGCAGTACTG  
CTCAGTCATCGTCTAGCTACTCTCTGACTCTGAGCATCTCGTACTCGATCTCGT  
CTACTGACTGAGCTCTAGCTACTCGAGATGCTCAGTCTACTGCTAGCTCAGAT  
GCTCAGTCAGTCTAGCATCTGCAGTCTCATCGATCTCTCGTACTCTCGTCTCAG  
TAGCTCTAGCAGTCTCTACTCTCTGACTCTGACTCGTAGAGCATCTGACGATCT  
CGTCAGTCAGAGTCTCAGAGATCTGCTCATCGTAGAGACTGAGCTACGATCTC  
GTCTACTCGTAGCTCAGTAGCTCATCTGAGAGCTCTACGAGTCAGATCT..

# Decoder



# ACGT Data

ACTGAGACGTCTACTGACGATCTCGTCTACTCGTCAGTAGACGATCGTACTC  
GTCTCTCAGTCAGTCTACGATCTCTCTGACTCTGAGACGATCGTCTCTCTCAT  
CTGAGCAGTCTCTACTCTCTCGATCTCGTACTCTCTCGTCTCAGTCTACTCGT  
ACTCTGCTCTACGATCGATCGATCTCGTAGCTAGACTGACGAGTAGCTACGT  
CTCAGTACGATCGTACTGCAGAGTCTCTCATCTGCTCATCTCTCTCTCTGCAG  
TACTGCTCAGTCATCGTCTAGCTACTCTCTGACTCTGAGCATCTCGTACTCGA  
TCTCGTCTACTGACTGAGCTCTAGCTACTCGAGATGCTCAGTCTACTGCTAG  
CTCAGATGCTCAGTCAGTCTAGCATCTGCAGTCTCATCGATCTCTCGTACTCT  
CGTCTCAGTAGCTCTAGCAGTCTCTACTCTCTGACTCTGACTCGTAGAGCATC  
TGACGATCTCGTCAGTCAGAGTCTCAGAGATCTGCTCATCGTAGAGACTGAG  
CTACGATCTCGTCTACTCGTAGCTCAGTAGCTCATCTGAGAGCTCTACGAGT  
CAGATCT..

# ACGT to Ternary Conversion

*Mapping of  $y_1$  :  $0 \rightarrow A$ ,  $1 \rightarrow C$ ,  $2 \rightarrow G$*

$y(i) \backslash y(i-1)$	A	C	G	T
A	-	0	1	2
C	0	-	1	2
G	0	1	-	2
T	0	1	2	-

# ACGT to Ternary Conversion

0121110001100121001211100110011001210011001210000110011111210  
1210110001211111121011121110012100111111122112112210111100111  
1110121110000111111001112101100110000111221110001210121012111  
0001210110121001100121000011121000012100001222111011111221122  
1122111111111122210001221121012210011012100111112101112112221  
1100001101211100110012101211211101210011011122211210110012210  
1211211222112101210110122211222101112210121111100001111001112  
1001211101221011110011111210111210110001112221121001211100121  
0121110111211112112211221000111101211210001211100110011000121  
12100121122112111121110001101211211101221000121111211110110001  
1121000110012100121012112100121111210110121001111011110110122  
1012100011011111211210001112210110110001211221122112112100111  
2111011012111012111001210011012210121122210001101221011000011  
1221001111012210110000110121122112221110121122100121012101210  
11011110011110001211101222112..



# Ternary to Binary Conversion

$$x_1 = y_1$$

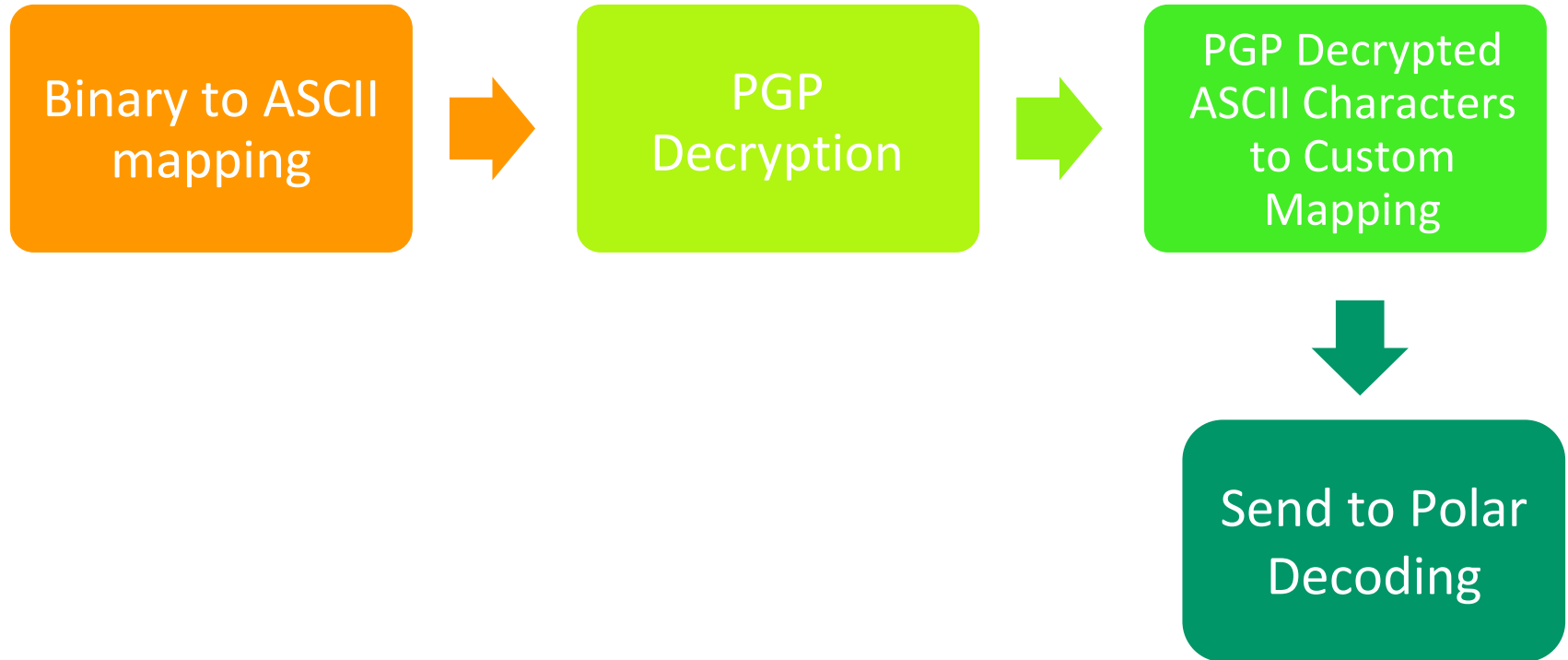
$$x_i = (y_i + x_{i-1}) \bmod 2, \quad \text{Where } i = 2, 3, \dots, N$$

$y(i) \backslash x(i-1)$	0	1	2
0	0	1	—
1	—	0	1

# Ternary to Binary Conversion

0110100001000110001101000100010001100010001100000100010101  
1001100100001101010110010110100011000101010111011011100101  
0001010100110100000101010001011001000100000101110100001100  
1100110100001100100110001000110000010110000011000001111010  
0101011101110111010101010111100001110110011100010011000101  
0110010110111101000001001101000100011001101101001100010010  
1111011001000111001101101111011001100100111101111001011100  
1101010000010100010110001101001110010100010101100101100100  
00101111011000110100.....

# PGP Decryption Block



# PGP Decryption

```
1 1 0 1 1 0 0 0 1 1 1 0 0 1 0 0 1 0 0 0 0 0 1 0 0 1 1 1 0 0 1 0 0 1 0 0 1 1 1 0  
1 0 1 1 1 1 1 0 1 1 0 1 1 0 0 0 1 1 1 0 0 1 0 0 1 0 0 0 0 1 0 0 0 0 0 1 0 0 1 0  
1 0 1 1 1 0 0 0 1 1 1 0 0 0 1 0 1 0 0 0 0 1 0 0 1 1 0 1 1 1 1 0 0 1 1 1 0 1 0 0  
1 1 1 0 0 0 1 0
```

# Polar Codes - Decoding

Worst Case time complexity for decoding block-length N code is  $O(N \log N)$ .

Decoder can be seen as N decision elements, one for each source element  $u_i$ . They are activated

in the order 1 to N according to the likelihood ratio,

$$L_N^{(i)}(y_1^N, \hat{u}_1^{i-1}) \triangleq \frac{W_N^{(i)}(y_1^N, \hat{u}_1^{i-1} | 0)}{W_N^{(i)}(y_1^N, \hat{u}_1^{i-1} | 1)}$$

# Polar Codes - Decoding

Decision at  $\hat{u}_i$  is generated as,

$$\hat{u}_i = \begin{cases} 0, & \text{if } L_N^{(i)}(y_1^N, \hat{u}_1^{i-1}) \geq 1 \\ 1, & \text{otherwise} \end{cases}$$

Efficient Decoding Algorithm is made based on equation (1) and (2).

# Polar Decoding

1100 1101 01 01 101 1110 1111 101 000 01 001 100

# Huffman Decompression

01001000 01100101 01101100 01101100 01101111 00100000  
01010111 01101111 01110010 01101100 01100100 00100001

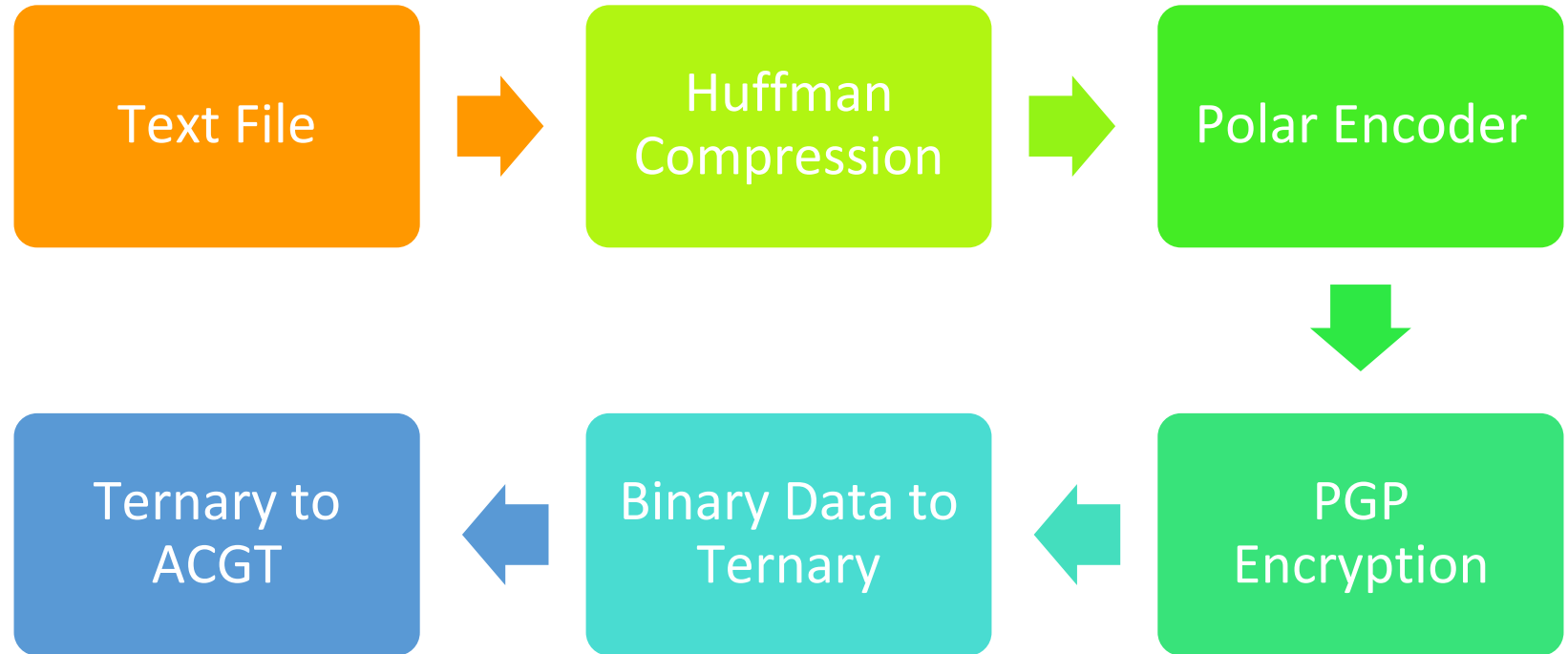
## Binary to ASCII

**Hello World!**

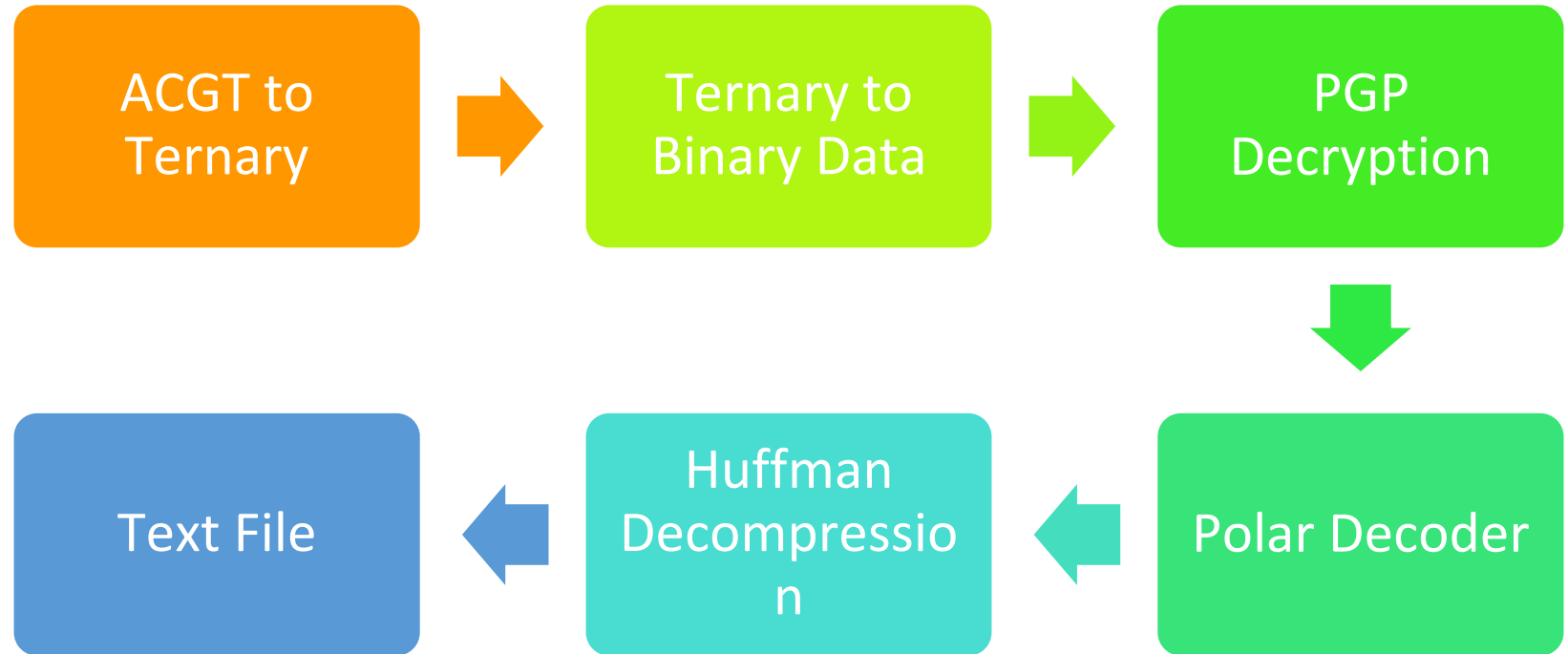


# User Application - UI Walkthrough

## Encoder



## Decoder



# Applications

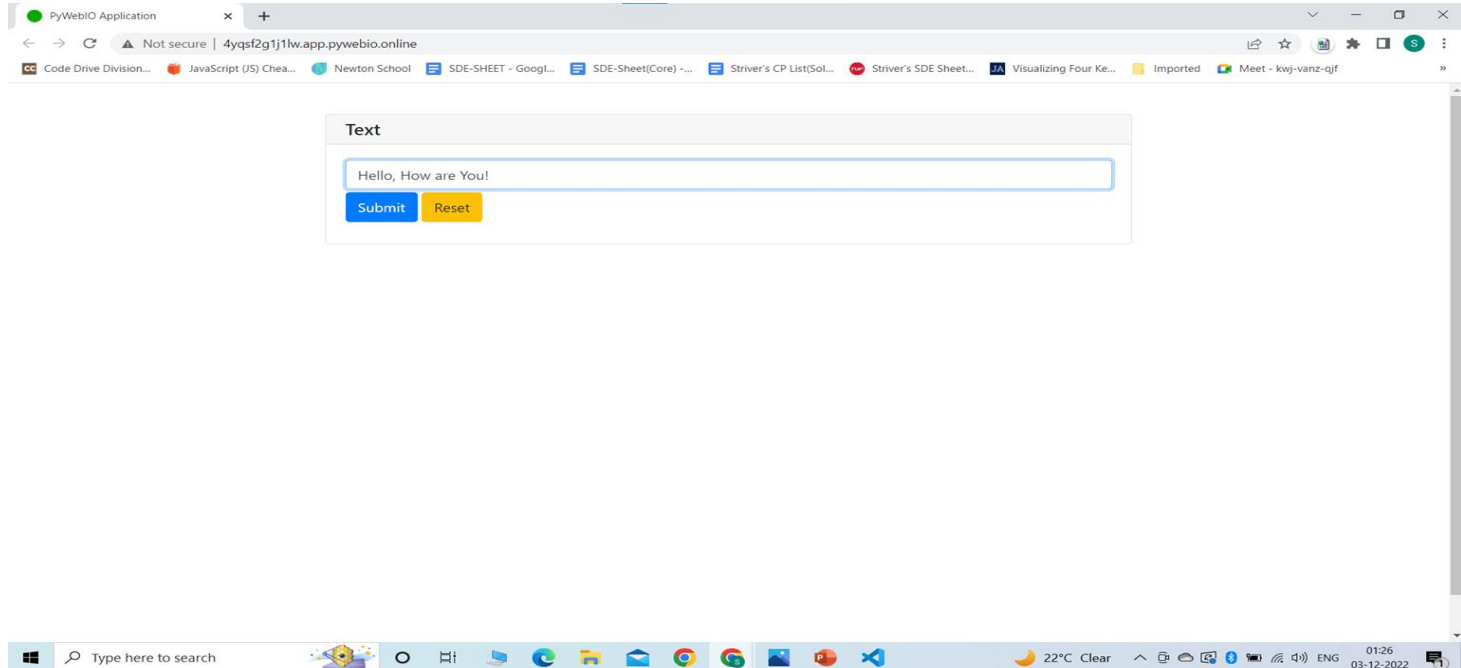
Encoder Web App: <http://hn8qenyu4brw.app.pywebio.online/>

Decoder Web App: <http://hc2zz2pry0cj.app.pywebio.online/>

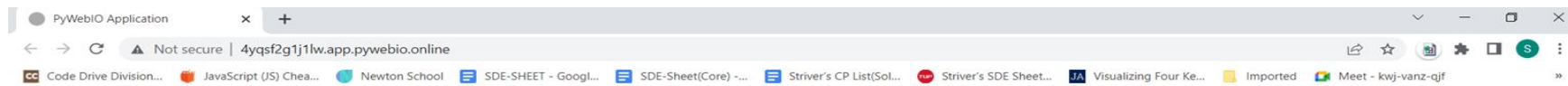
GitHub Repository: [https://github.com/KathanS/DNA\\_Codec](https://github.com/KathanS/DNA_Codec)

# **User Application - System Design Overview**

# Encoding App Web interface



# Blockwise Output [Encoder]



## Final DNA String

```
ACTCTGAGACTCTGAGACTCTGAGACTCTGAGACTCTGAGAGACGTCTAGACGAGAGAGTAGCAGAGTCTAGAGACTGCTACT
CGTACGAGAGTACGAGTAGCAGAGAGTACGTCTACGTAGACTGAGAGACGAGAGAGAGTTCAGAGAGTACGAGAGTACGAGAGTA
GCAGAGTAGAGACTCTGAGACTCTGAGACTCTGAGACTCTGAGACTCTGAGACGTCAGAGTACTCTCGTACTGAGACGCTCTTA
GCTCTACGAGTAGCTACTGAGACGAGTAGACGATGCTACGATGCTAGAGACGTAGAGCTACTGAGAGCTAGAGCTCTCTAGAG
AGAGAGACTCGAGAGTCAGTAGACTCGTAGACGTAGAGAGACTGAGACGTAGAGAGACGAGATCGAGTAGACGTAGAGAGA
CGAGTCTCATGAGCTCATCGAGTCATGAGCTCATGAGACGTAGAGACTGAGAGCAGATGAGACTGCTACTGACGAGATCTGCTA
CTGCTAGAGCATCGAGAGAGTTCAGAGTAGCTAGAGCTCTACTGAGACGTCATCGAGAGATCGAGTCAGAGTAGCATCGTAGCA
GATGAGACTCGAGAGAGATCGAGTACTCGTCAGATCGTCTCTCAGAGAGAGAGATGAGCAGATGCTCTAGCAGAGAGATCTCT
```

Secret Code to be used for Decoding

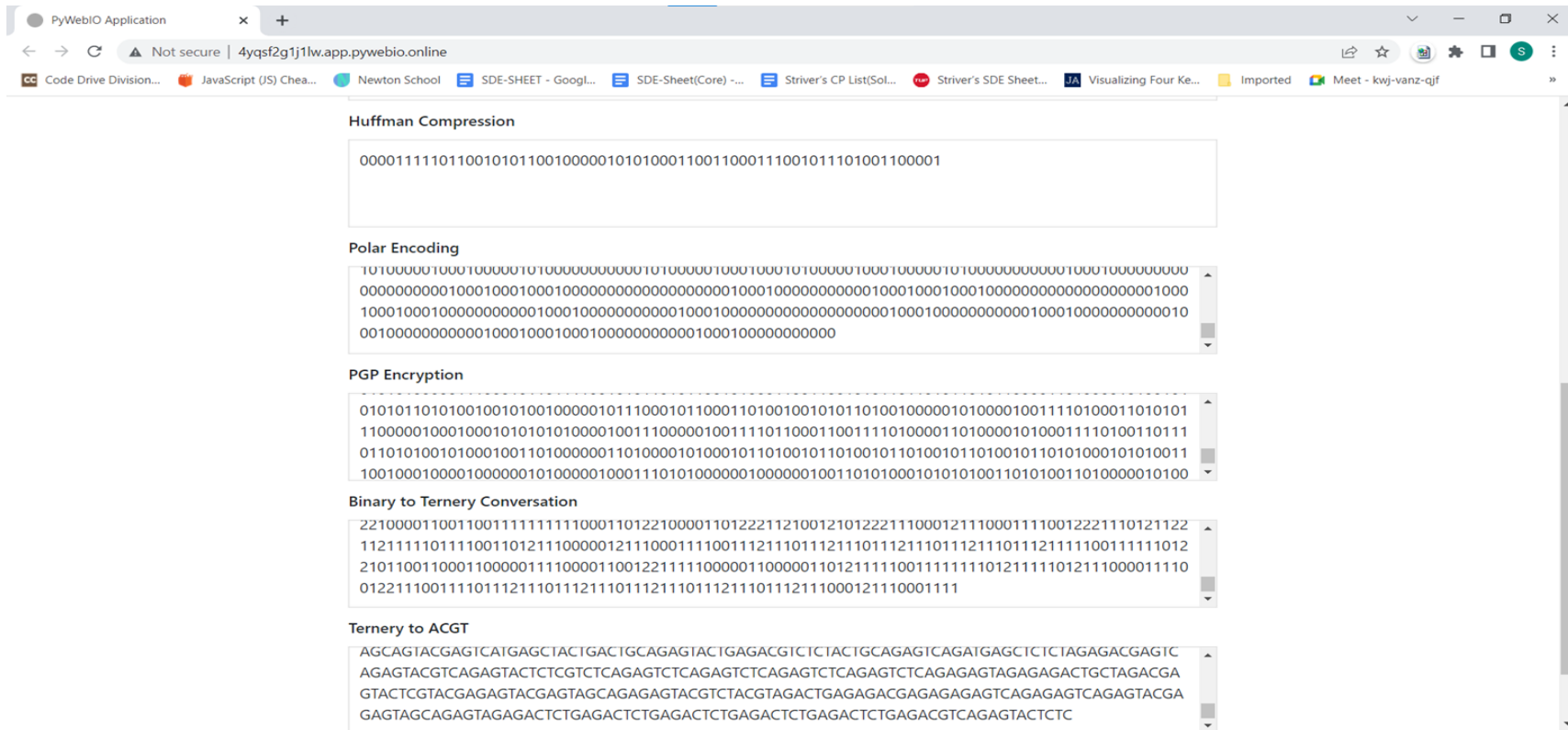
qwrmy8219

## Blockwise Ouput

Original Text

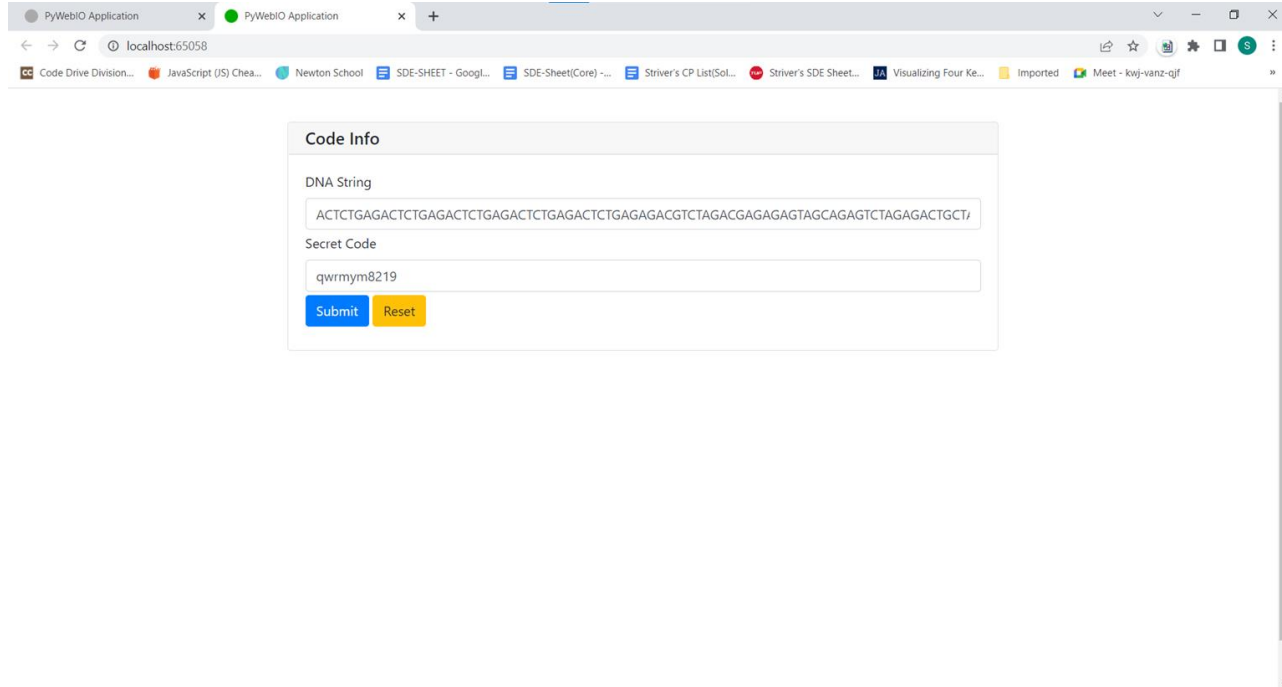
Hello, How are You!

## Blockwise Output [Encoder]





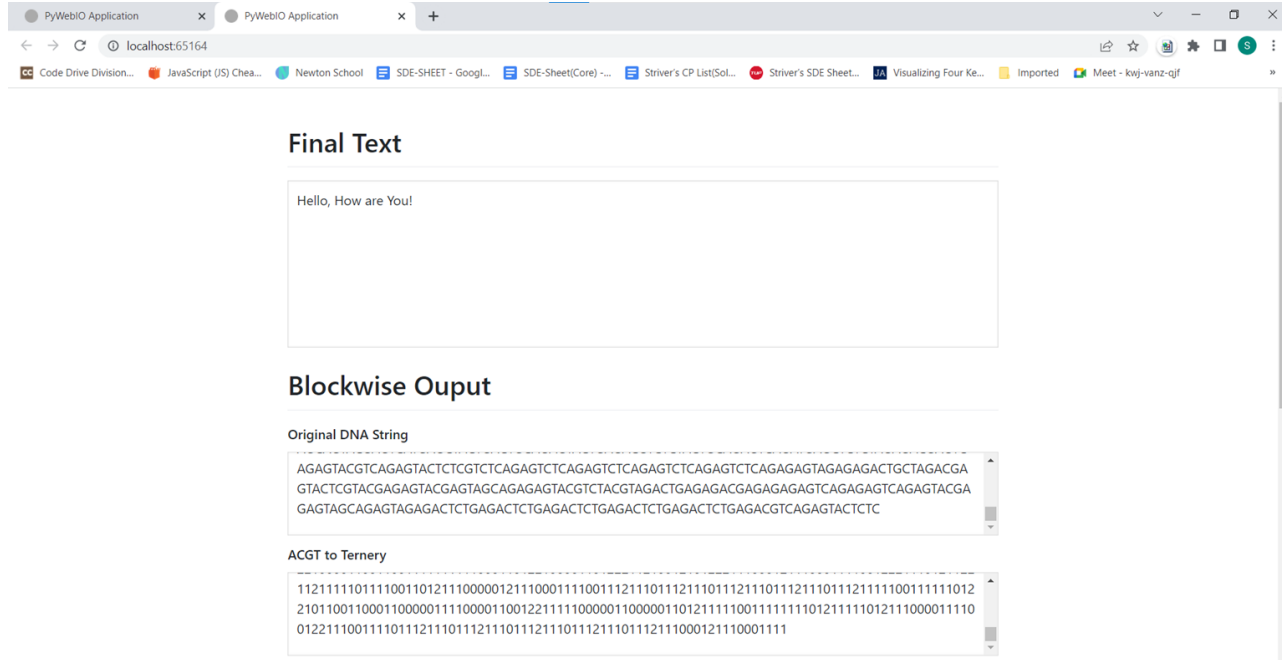
# Decoding App Web Interface



The screenshot shows a web browser window with two tabs, both titled 'PyWebIO Application'. The address bar displays 'localhost:65058'. The browser's tab bar shows several open tabs, including 'Code Drive Division...', 'JavaScript (JS) Chea...', 'Newton School', 'SDE-SHEET - Googl...', 'SDE-Sheet(Core) ~...', 'Striver's CP List(Sol...', 'Striver's SDE Sheet...', 'Visualizing Four Ke...', 'Imported', and 'Meet - kvj-vanz-qj'. The main content area features a 'Code Info' form with the following elements:

- Code Info** (Section Header)
- DNA String** (Label)
- 
- Secret Code** (Label)
- 
- 
-

# Blockwise Output [Decoder]

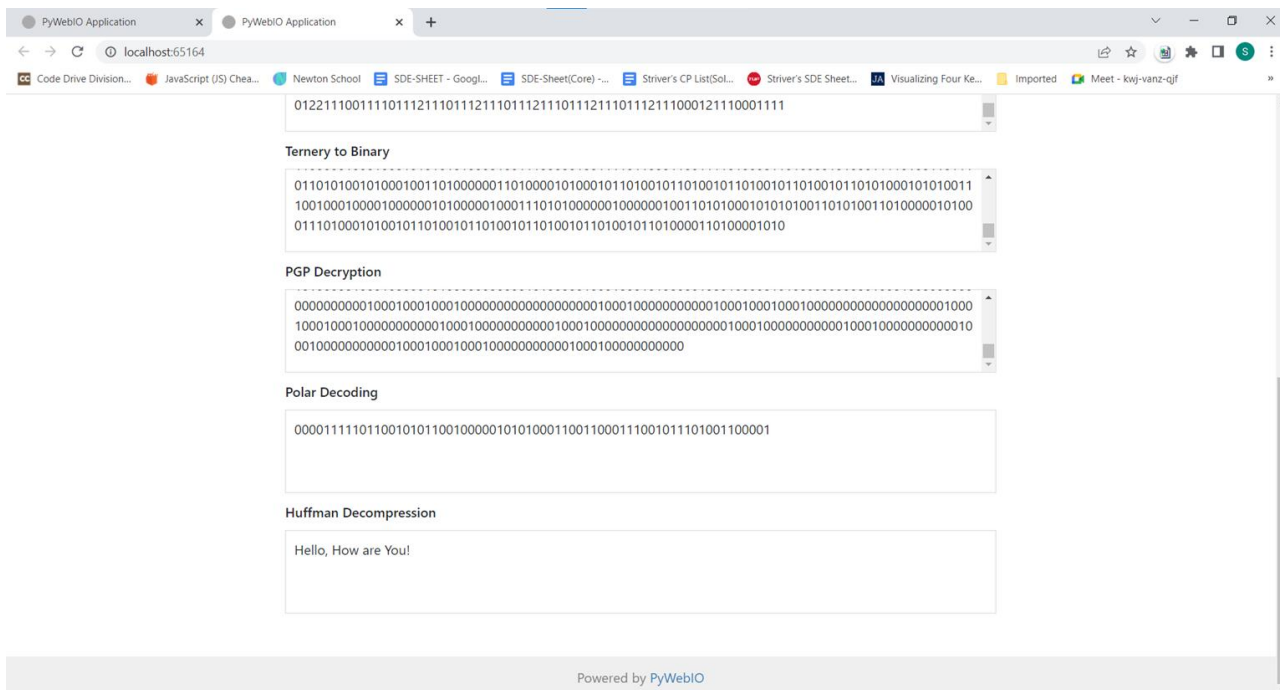


The screenshot shows a web browser window with two tabs, both titled 'PyWebIO Application'. The address bar shows 'localhost:65164'. The browser's taskbar at the bottom lists several open applications: Code Drive Division..., JavaScript (JS) Chea..., Newton School, SDE-SHEET - Googl..., SDE-Sheet(Core) ~..., Striver's CP List(Sol..., Striver's SDE Sheet..., Visualizing Four Ke..., Imported, and Meet - kvj-vanz-qif.

The main content area of the browser displays a web application interface. It features a section titled 'Final Text' with a text box containing the message 'Hello, How are You!'. Below this is a section titled 'Blockwise Ouput' (note the typo). Under 'Blockwise Ouput', there are two sub-sections, each with a scrollable text area:

- Original DNA String**:  
AGAGTACGTCAGAGTACTCTCGTCTCAGAGTCTCAGAGTCTCAGAGTCTCAGAGTCTCAGAGTAGAGAGACTGCTAGACGA  
GTACTCGTACGAGAGTACGAGTAGCAGAGAGTACGTCTACGTAGACTGAGAGACGAGAGAGTCTCAGAGAGTCAGAGTACGA  
GAGTAGCAGAGTAGAGACTCTGAGACTCTGAGACTCTGAGACTCTGAGACTCTGAGACGTCAGAGTACTCTC
- ACGT to Ternery** (note the typo):  
112111110111100110121110000012111000111100111211101112111011121110111211101112111100111111012  
21011001100011000001111000011001221111100000110000011012111110011111110121111012111000011110  
01221110011110111211101112111011121110111211101112111000121110001111

## Blockwise Output [Decoder]



# Microservices like Architecture

- Encoder Web App outputs DNA string and Secret Code, which will be required for Decoding.
- We are using Python Pickles to store outputs of intermediate blocks, this way system gets decoupled and one block can pass the message having secret code to next block after finishing its task.
- We are able to scale each block independently, this increases robustness and overall performance of the system.
- **Information Density**  
Without PGP: ~2.5  
With PGP: ~6

# References

1. Channel polarization: A method for constructing capacity-achieving codes for symmetric binary-input memoryless channels Erdal Arıkan, Senior Member, IEEE, <https://arxiv.org/pdf/0807.3917.pdf>
2. Introduction to Polar Codes, NPTEL - NOC IITM, [\(1225\) Introduction to Polar Codes: Polar Transform - YouTube](#)
3. A Brief Introduction to Polar Codes, Henry D. Pfister, [polar.pdf \(duke.edu\)](#)
4. A few concepts of DNA Codec - <https://youtube.com/playlist?list=PLZ3QKRH1yB54D8HioxHVe7mrEx85yArOR>
5. PGP Encryption/Decryption - [What is PGP Encryption and How Does It Work? \(varonis.com\)](#)
6. Huffman Compression - [Huffman Encoding & Python Implementation | by Yağmur Çiğdem Aktaş | Towards Data Science](#)
7. The design of mini PGP security, <https://ieeexplore.ieee.org/document/6021726?arnumber=6021726>