

28<sup>th</sup>

Aug.

## 6 Recurrences

- ① # of binary strings that do not have 2 consecutive zeros (0s)

$$\hookrightarrow x^n S_n = x^n S_{n-1} + x^n S_{n-2} \quad n \geq 2 \quad S_0 = 1 \quad S_1 = 2$$

generating functions.

$$\hookrightarrow \sum_{n \geq 2} x^n S_n = \sum_{n \geq 2} x^{n-1} S_{n-1} + \sum_{n \geq 2} x^n S_{n-2} \quad n \geq 2$$

$$x^2(F(x) - S_0 - S_1 x)$$

$$\Rightarrow F(x) - 1 - 2x = x(F(x) - 1) + x^2 F(x)$$

$$\Rightarrow F(x) = \frac{1+x}{1+2x-x^2}$$

$$\Rightarrow \frac{1+2x}{(1-\alpha x)(1-\beta x)}$$

$$1+2x-x^2 = (1-\alpha x)(1-\beta x)$$

$$= 1 + \beta x^2 - (\alpha + \beta)x$$

$$\alpha + \beta = -1, \alpha\beta = -1, \frac{1}{\beta} + \beta = -1$$

$$\alpha, \beta = \frac{-1 \pm \sqrt{5}}{2}$$

$$\Rightarrow \frac{1+x}{(1-\alpha x)(1-\beta x)} = \frac{a}{1-\alpha x} + \frac{b}{1-\beta x}$$

$$\text{i.e. } 1+x = a - a\beta x + b - b\alpha x$$

$$\text{i.e. } a+b = 1$$

$$(a\beta + b\alpha) = -1$$

$$S_n = a_1 \alpha^n + b_1 \beta^n$$

for fibonacci

# Idea is to generate closed form of  $F(x)$

- ② # of ways of multiplying  $x_0 \dots x_n$

we got

$$\hookrightarrow C_n = C_0 C_{n-1} + C_1 C_{n-2} + \dots + C_{n-1} C_0$$

$$C_n = \sum_{i=0}^{n-1} C_i C_{n-i-1}$$

$\hookrightarrow$  multiply both sides by  $x^n$  (1)  
always first thing to do

$$\Rightarrow \sum_{n \geq 1} C_n x^n = \sum_{n \geq 1} \sum_{i=0}^{n-1} C_i C_{n-i-1} x^n$$

$$F(x) = C_0$$

first term is  $C_0 = 1$  $\hookrightarrow$  now if we have  $F(x) \cdot F(x)$ 

$$\Rightarrow (C_0 + C_1 x + C_2 x^2 + \dots)(C_0 + C_1 x + C_2 x^2 + \dots)$$

what is the coefficient of  $x^n$ 

$$\hookrightarrow C_0 C_0 + C_1 C_{n-1} + C_2 C_{n-2} + \dots + C_n C_0$$

 $\hookrightarrow$  write  $x^n$  as  $x^1 \cdot x^{n-1}$  in RHS

i.e.

$$\Rightarrow x \sum_{n \geq 1} \sum_{i=0}^{n-1} C_i C_{n-i-1} x^{n-1}$$

this is of the form  $F(x)^2$ .

$$\Rightarrow F(x) - C_0 = x F(x)^2$$

$$\Rightarrow F(x) - 1 = x F(x)^2$$

$$F(x) = \frac{(1 \pm \sqrt{1-4x})}{2x}$$

we will choose  $\frac{1 - \sqrt{1-4x}}{2x}$  why?

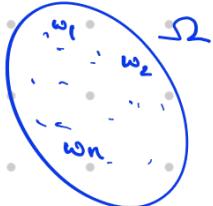
$$\hookrightarrow F(x) = \frac{1 - \sqrt{1-4x}}{2x}$$

$$(1+\alpha x)^{Y_2} = \sum_{n=0}^{\infty} (\alpha x)^n (Y_2)_n \quad \boxed{0}$$

# ★ PROBABILITY AND RANDOM VARIABLE

"Sample space"

where we are observing something and we get a set of outcomes  $\Omega$



for each outcome  $w \in \Omega$   
we assign it probability  $p(w)$ .

↳ Probability is a Total ordered function

$$P: \Omega \rightarrow \mathbb{R}$$

↳ Conditions satisfied

$$(1) P(w) \geq 0 \quad \forall w \in \Omega$$

$$(2) \sum_w P(w) = 1$$

↳ for quicksort: outcomes are a tree of path / partition made by algo.

# Infiniite Set of outcomes  $\rightarrow$  selecting a number  
↳ Prob of getting 0.3 from b/w (0,1)  $\approx \frac{1}{3}$

\* Event ( $E$ )

↳ Subset of  $\Omega$

↳  $E \subseteq \Omega$ .

$$P(E) = \sum_{w \in E} P(w).$$

\* Independent Event.

↳ for two Events  $E_1$  and  $E_2$

↳ These two events are independent if

$$P(E_1 \cap E_2) = P(E_1) \cdot P(E_2)$$

## EXAMPLES

(1) Throw 2 die

$E_1$  = sum of the value = 6

$E_2$  = the first die is odd

$$P(E_1) = \frac{5}{36} \quad P(E_2) = \frac{1}{3}$$

$$P(E_1 \cap E_2) = \frac{1}{18}$$

$$\therefore P(E_1) \cdot P(E_2) = \frac{5}{36} \times \frac{1}{3}$$

↳ hence they are not independent.

## \* Conditional prob.

↳ Put some condition on a event.

↳ suppose  $E$  are a set of outcomes that have happened.

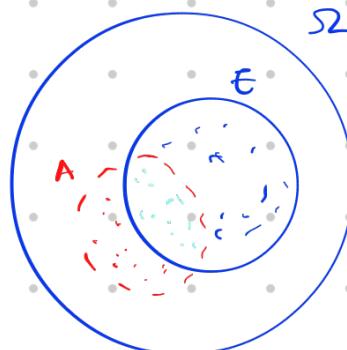
↳ now this set  $E$  becomes our  $\Omega$

↳  $\forall w \in E \quad P(w)$  changes as other outcomes from  $\Omega$  are not part of our system.

↳ we will have to scale them, that we will do by.

$$\forall w \in E \quad P'(w) = \frac{P(w)}{\sum_{w \in E} P(w)} \rightarrow P(E).$$

$$\Rightarrow P(w|E) = \begin{cases} \frac{P(w)}{P(E)} & \text{if } w \in E \\ 0 & \text{if } w \notin E \end{cases}$$



i.e. the prob of happening  $A \cap E$ ,  $\because A-E$  are all 0

↳  $P(A|E) \rightarrow$  prob of  $A$  given  $E$  has happened.

$$\therefore \sum_{w \in A \cap E} P(w) + \sum_{\substack{w \in A \\ w \notin E}} P(w)$$

$$\therefore \frac{P(A \cap E)}{P(E)} = P(A|E)$$

# If  $A$  and  $E$  are independent

$$P(A|E) = P(A)$$

## \* Basic facts

$$\textcircled{1} \quad P(A_1 \cup A_2) = P(A_1) + P(A_2) - P(A_1 \cap A_2)$$

Special Case for die example  
 $A_1 \rightarrow \{2, 4\}, A_2 = \{3, 5\}$

$$* A_1 \cap A_2 = \emptyset$$

$$\hookrightarrow P(A_1 \cup A_2) = P(A_1) + P(A_2)$$

\* Union bound  $\rightarrow$  really important on  
 $\hookrightarrow P(A_1 \cup A_2) \leq P(A_1) + P(A_2)$

## Examples

\textcircled{1} 4 envelope — name on it

4 letter — name on it.

\hookrightarrow we put any letter in any envelope.

\hookrightarrow Possible outcomes = 4!

$P(\text{Exactly } k \text{ letters are in right envelope})$

$$\hookrightarrow \frac{4C_2}{4!}$$

$\boxed{A} \boxed{B} \boxed{C} \boxed{D}$

$P(\text{at least } 3)$

$$\hookrightarrow P(\text{none correct}) + P(1 \text{ correct}) + P(2 \text{ correct}) + P(3 \text{ correct})$$

$$\text{or } 1 - P(4 \text{ correct}).$$

## \* Random Variable.

\hookrightarrow A random Variable  $X$  is a mapping from  $\Omega \rightarrow \mathbb{R}$

example

\textcircled{1} Throw 10 coins,  $X = \text{no. of heads}$ .

$$1 \leq 2^{\Omega} = 2^{10}, \quad X = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

\textcircled{2} we have  $n$  bins and  $n$  balls

$$\square \quad \square \quad \square \quad - \cdots - \quad \square \quad n \text{ bins}$$

$\underbrace{\quad}_{\text{n balls}}$

\hookrightarrow we can fill bins with any no. of balls

$X = \text{no. of bins that are empty}$

$$X = \{0, 1, 2, \dots, n-1\}$$

$$|\Omega| = n.$$

\textcircled{3}



you can move  $(x-1)$  or  $(x+1)$  why?

\hookrightarrow Set of outcomes: infinite (countable)

(4)

1st Sept

## Bayes Rule

↳  $P(A|B)$

If  $B$  has happened what is the probability of  $A$  to happen.

hypothesis → observation  
 $P(A|B) = \frac{P(A \cap B)}{P(B)}$

$$P(A|B) = \frac{P(B|A) \cdot P(A)}{P(B)}$$

↳ It is Expressing  $P(A|B)$  in terms of  $P(B|A)$

Example:

① There are 2 boxes with colored balls

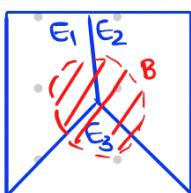


- ↳ choose a random bag and draw a ball
- ↳ The ball happens to be red. what is the probability that it chose bag 1?  
 ↳ A hypothesis
- ↳ B (observation)

$$\begin{aligned} \text{↳ } P(A|B) &= \frac{P(B|A) \cdot P(A)}{P(B)} \\ &= \frac{\frac{1}{2} \times \frac{1}{2}}{\frac{6}{20}} = \frac{5}{6}. \end{aligned}$$

② A set of Events  $E_1, E_2$  is said to be Partition the set of outcomes

$$\begin{aligned} E_1 \cap E_2 &= \emptyset \\ \cup E_i &= \Omega \end{aligned}$$



$$P_B[B] = P\{B \cap E_1\} + P\{B \cap E_2\} + \dots$$

or

$$P[B] = P\{B|E_1\}P(E_1) + P\{B|E_2\}P(E_2) + \dots$$

↳ In our example  $E_1 = \text{bag 1 is chosen}$   
 $E_2 = \text{bag 2 is chosen}$

$$P[B] = P\{B|E_1\}P(E_1) + P\{B|E_2\}P(E_2) + \dots$$

$$= \frac{1}{4} + \frac{1}{20} = \frac{3}{10}$$

## \* Random Variable. $[x : \Omega \rightarrow \mathbb{R}]$

Examples (contd.)

① Suppose we have a program which tosses  $x$  coins.

↳ and to get an output it takes some  $y$ , amount of time  $\forall x \in \Omega$ .

then  $x : \Omega \rightarrow \mathbb{R}$  } as per real life  
 $|\Omega| = 2^10$   
 $x : \# \text{ heads.}$  }  $x$  will be close to  $\frac{1}{2}$  (so)

Q. how do we capture the expected value of  $x$ .

↳ Perform experiment

$$\text{↳ } E[X] := \sum i \cdot P[x=i]$$

$$\text{# } P[x=i] \rightarrow \Pr(\text{outcomes } w \in \Omega \text{ where } x(w) = i)$$

↳ Suppose we conduct this experiment  $N$  times

Suppose  $X=x_j$  in the  $j^{\text{th}}$  experiment

$$\text{Output} = \frac{x_1 + x_2 + \dots + x_N}{N} \quad \text{let } P_i = \Pr[X=i].$$

③ Suppose in an "ideal world", we see  $i$  heads exactly  $P_i N$  times

↳ think with an ex:  
 like  $N=100, i=10$

$$\rightarrow \frac{x_1 + x_2 + \dots + x_N}{N} = \frac{\sum i \cdot P_i \cdot N}{N}$$

$$= \sum i \cdot P_i$$

↳ mathematical proof of ③

## Law of Large Numbers

Let  $X$  be a random variable.

Suppose we conduct this experiment  $N$  times, and  $x_i$  denotes the value of  $X$  in the  $i^{\text{th}}$  experiment.

Then,  $\frac{x_1 + x_2 + \dots + x_N}{N} \rightarrow \mathbb{E}X \text{ as } N \rightarrow \infty$

\* Examples

① Toss  $n$  coins  $X = \# \text{ heads}$

$$\mathbb{E}X = \sum_{i=0}^n i \cdot \Pr[X=i] = \sum_{i=0}^n i \cdot (\pi_i) (\gamma_i)^n \\ = \gamma_2$$

②  $n$  envelopes,  $n$  letters.

Randomly insert.

$X = \# \text{ letters in correct envelope}$ .

$$\mathbb{E}X = \sum_{i=0}^n i \left[ \frac{n!}{i!(n-i)!} \text{ (derangement)} \right]$$

no. of letters correctly placed.

no. of letters not " "

③ Derangements: how many permutations of  $\{1, 2, \dots, n\}$  are there where  $\sigma_i \neq i$  for each  $i$

↳ 2143

2341

⋮

$$\mathbb{E}X = \sum_{i=0}^n i \left[ \frac{n!}{i!(n-i)!} \text{ (derangement)} \right] = 1$$

③

3.1 : not biased coin.  $\Pr[H] = \Pr[T] = \frac{1}{2}$

keep tossing 1 coin till you see a head.

↳ H, TH, TTH, TTTH, ...  $T^{(x-1)}H$

$X$  are no. of coin tosses.

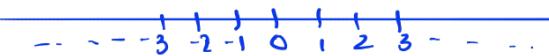
$$\mathbb{E}X = \sum_{i=1}^{\infty} i \cdot \frac{1}{2^i} = 2$$

3.2  $\Pr[H] = P(\text{Biased coin})$

$$\mathbb{E}X = \sum_{i=1}^{\infty} i \cdot (1-p)^{i-1} p = \frac{p}{1-p}$$

not getting heads

④ Random walk example.



equally likely to move left(L) or Right(R)  
starting from 0

$X = \text{distance from origin}$

$$\mathbb{E}X \approx \sqrt{n} \text{ DIY: But can?}$$

This means diff of L and R will be  $\sqrt{n}$

\* Linearity of Expectation.

Let  $X$  be a random variable.

- Suppose there are other random variables

$X_1, X_2, X_3, \dots, X_K$  such that

$$X = X_1 + X_2 + \dots + X_K$$

$$\mathbb{E}X = \mathbb{E}X_1 + \mathbb{E}X_2 + \dots + \mathbb{E}X_K$$

Examples.

①  $n$  coin tosses

$X = \# \text{ heads}$

Define  $X_i = \begin{cases} 0 & \text{if } i^{\text{th}} \text{ coin toss is tails} \\ 1 & \text{if } i^{\text{th}} \text{ coin toss is heads} \end{cases}$

$$X = X_1 + X_2 + \dots + X_n$$

$$\mathbb{E}X = \mathbb{E}X_1 + \mathbb{E}X_2 + \dots + \mathbb{E}X_n$$

$$\mathbb{E}X_i = 1 \cdot \Pr[X_i=1] + 0 \cdot \Pr[X_i=0]$$

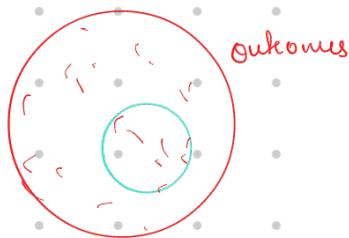
$$\mathbb{E}X_i = 1 \cdot \frac{1}{2} = \frac{1}{2}$$

$$\mathbb{E}X = \frac{n}{2}$$

4<sup>th</sup> Sept

## ★ Recap

$$\Pr[X=a] \equiv \Pr[\text{outcomes } \omega : X(\omega)=a]$$



$$E[X] = \sum_a a \Pr[X=a]$$

↳ To simplify these eq<sup>n</sup> we break them into simpler random variable.

i.e.  $X = X_1 + X_2 + \dots + X_n$

$E[X] = E[X_1 + X_2 + \dots + X_n] = E[X_1] + E[X_2] + \dots + E[X_n]$

↳ Linearity of Expectation.

## \* Proof of Linearity of Expectation

↳ prove for smaller expression

i.e. if  $X = Y+Z$

$$E[X] = E[Y+Z]$$

# we can use this to prove for multiple variables.

i.e.  $X = Y+Z+W$

$$X = Y+X'$$

$$= E[X] = E[Y+E[X']]$$

$$= E[X] = E[Y] + E[Z] + E[W].$$

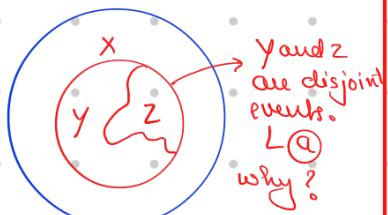
## Proof

$$\text{LHS} = E[X]$$

$$\sum_a a \Pr[X=a] \Rightarrow \sum_a [a \Pr\{Y+Z=a\}]$$

↳ for  $Y+Z$  to be equal to  $a$  only if  $y=b$ , then  $Z=a-b$

$$[Y+Z=a] = \bigcup_b [Y=b, Z=a-b].$$



by ④ we know that probability of union of disjoint events is sum of all events

i.e.  $\sum_a a \left( \sum_{b,c} \Pr\{Y=b, Z=c\} \right)$

$$= \sum_a a \left( \sum_{b,c} \Pr\{Y=b, Z=c\} \right) \because a \geq b+c$$

# our eq<sup>n</sup> is of form

$$\sum_a \sum_{b,c} f(b,c)$$

which is equivalent to

$\sum_{b,c} f(b,c) \because$  we are basically finding and summing all possible pairs of  $b, c$  which equal to some  $a$ .

$$= \sum_{b,c} [ \Pr\{Y=b, Z=c\} \cdot f(b,c) ]$$

$$= \sum_{b,c} b \underbrace{[ \Pr\{Y=b, Z=c\} ]}_{\textcircled{b}} + \sum_{b,c} c \underbrace{[ \Pr\{Y=b, Z=c\} ]}_{\textcircled{c}}$$

# we can say that

$$[Y=b] = \bigcup_c [Y=b, Z=c].$$

$$\Pr[Y=b] = \sum_c \Pr[Y=b, Z=c]. \rightarrow \textcircled{d}$$

↳ Saying as we toss a coin and ask what is the probability of first coin toss is heads.

↳  $\Pr[\text{first heads}] = \sum \Pr[\text{first head, other heads}]$

$$\Pr[Y=b] = \Pr[Y=b, Z=H] + \Pr[Y=b, Z=T]$$

by using  $\textcircled{b}$

$$\sum_b b \underbrace{[ \Pr[Y=b, Z=c] ]}_{\textcircled{b}}$$

$$\sum_b b \cdot \Pr[Y=b] \rightarrow \text{by } \textcircled{d}$$

Similarly for  $\textcircled{c}$  we get  $\sum_c c \cdot \Pr[Z=c]$

hence

$$E[X] = \sum_b b \cdot \Pr[Y=b] + \sum_c c \cdot \Pr[Z=c]$$

$$E[X] = E[Y] + E[Z] \text{ hence proved.}$$

Prove:  $\mathbb{E}X = \mathbb{E}Y \cdot \mathbb{E}Z$

$$\hookrightarrow \sum_a \sum_{\substack{b,c \\ b+c=a}} \Pr[Y=b, Z=c] \quad \text{using the similar steps as before just we get } b+c=a$$

but now we cannot decompose this but suppose we could, then.

$$\Rightarrow \Pr[Y=b] \cdot \Pr[Z=c] \quad \text{(a)}$$

$$\Rightarrow \sum_{b,c} \Pr[Y=b] \Pr[Z=c] \quad (\text{b}, \text{c})$$

$$\Rightarrow \sum_{b,c} b \Pr[Y=b] \cdot c \Pr[Z=c] \quad \downarrow \text{how is it possible}$$

$$\Rightarrow \sum_b b \Pr[Y=b] \cdot \sum_c \Pr[Z=c] \cdot c \quad \begin{aligned} &\text{assume } \sum_{i=1}^n y_i \\ &= (x_1, y_1) + \\ &(x_2, y_2) + \\ &(x_3, y_3) + \\ &(x_4, y_4) \\ &\Rightarrow (x_1+x_2)(y_1+y_2) \end{aligned}$$

$$\Rightarrow \mathbb{E}Y \cdot \mathbb{E}Z.$$

for this to happen step (a) must be possible but in general it is not possible.

Example

$$X = \begin{cases} 1 & \text{if H} \\ 0 & \text{if T} \end{cases}$$

$Y = X, Z = X$  is  $X = YZ$ ?

↳ yes for this binary ex.

by  $\mathbb{E}X = \mathbb{E}Y$ .

$$\text{so, } \mathbb{E}X = \mathbb{E}Y = \mathbb{E}Z$$

$$= \mathbb{E}Y = \mathbb{E}Y \cdot \mathbb{E}Y$$

$\Rightarrow Y = Y^2$  which is not possible

hence this only works if  $Y$  and  $Z$  are independent i.e.

if  $Y$  and  $Z$  are independent.

$$\Pr[Y=b, Z=c] = \Pr[Y=b] \cdot \Pr[Z=c].$$

for sum, this is true regardless.

Example

(1) Toss 2 coins

$$X = \begin{cases} 1 & \text{if 1st coin is H} \\ 0 & \text{otherwise} \end{cases} \quad Y = \begin{cases} 1 & \text{if 2nd coin is H} \\ 0 & \text{otherwise} \end{cases}$$

$$\Pr[X=1, Y=1] = \frac{1}{4} = \Pr[X=H] \cdot \Pr[Y=H].$$

∴  $X$  and  $Y$  are independent.

(2) when throwing a dice

$$X = \begin{cases} 1 & \text{if number is even} \\ 0 & \text{otherwise} \end{cases}$$

$$Y = \begin{cases} 1 & \text{number is greater than 3} \\ 0 & \text{otherwise} \end{cases}$$

$X$  and  $Y$  are not independent.

i.e.

$$\Pr[X=1, Y=1] \neq \Pr[X=\text{even}] \cdot \Pr[Y=\text{gt 3}].$$

(3) Toss  $n$  coins

$X = \# \text{ heads}$

$$X = X_1 + X_2 + \dots + X_n$$

$$X_i = \begin{cases} 1 & \text{if } i^{\text{th}} \text{ coin toss is H} \\ 0 & \text{otherwise} \end{cases}$$

$$\mathbb{E}X = \mathbb{E}X_1 + \mathbb{E}X_2 + \dots + \mathbb{E}X_n$$

$$\mathbb{E}X = \mathbb{E}Y_1 + \mathbb{E}Y_2 + \dots + \mathbb{E}Y_n = n \cdot \mathbb{E}Y_1.$$

(4)  $n$  letters in  $n$  envelopes

$X = \# \text{ number of letter that go in correct envelope}$

$$X = X_1 + X_2 + \dots + X_n$$

$$X_i = \begin{cases} 1 & \text{if } i^{\text{th}} \text{ letter is in correct envelope} \\ 0 & \text{otherwise.} \end{cases}$$

$$\Pr[X_i=1] = \frac{1}{n} \quad \begin{aligned} &\text{equally likely to go in any} \\ &\text{envelope} \end{aligned}$$

$$\text{i.e. } \mathbb{E}X_i = \frac{1}{n}.$$

$$\mathbb{E}X = \mathbb{E}X_1 + \mathbb{E}X_2 + \dots + \mathbb{E}X_n$$

$$= n \cdot \mathbb{E}X_1$$

$$\boxed{\mathbb{E}X = 1}$$

⑤ we are given an array of length  $n$  with distinct elements and we want to find the max.

i.e.  $\text{Max} = A[0]$   
for  $i \geq 1$  to  $n-1$   
if  $[A[i] > \text{max}]$   
 $\text{max} = A[i]$

⑥

6 In worst case we will execute ⑥  $n-1$  times.

6 Start with a random permutation of  $1, \dots, n$ .

$X = \# \text{ of times } ⑥ \text{ is executed.}$

what's  $\mathbb{E}X$ .

$X = X_1 + X_2 + \dots + X_{n-1}$  not independent because  
if max is at  $j$  the every  $i \leq j$  there is a selection.  
 $X_i = \begin{cases} 1 & \text{if } ⑥ \text{ is executed in } i^{\text{th}} \text{ iteration} \\ 0 & \text{otherwise.} \end{cases}$

$$\mathbb{E}X = \mathbb{E}X_1 + \mathbb{E}X_2 + \dots + \mathbb{E}X_{n-1}$$

$$\mathbb{E}X_i = \Pr[X_i = 1]. 1$$

↳ we will execute ⑥ at  $i^{\text{th}}$  iteration iff

$$A[i] > A[0] \text{ to } A[i-1]$$

$$\text{i.e. } \Pr[X_i = 1] = \frac{1}{i}$$

$$\mathbb{E}X = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n-1} \approx \log n.$$

$$\boxed{\mathbb{E}X \approx \log n}$$

## ⑥ Random Walk.

$X = \text{Position after } n \text{ steps.}$

$$X = Y_1 + Y_2 + \dots + Y_n. \quad \text{Independent.}$$

$Y_i = \begin{cases} 1 & \text{if move right on } i^{\text{th}} \text{ step.} \\ -1 & \text{if move left on } i^{\text{th}} \text{ step.} \end{cases}$

$$\mathbb{E}X_n = \mathbb{E}Y_1 + \mathbb{E}Y_2 + \dots + \mathbb{E}Y_n.$$

$$\mathbb{E}Y_i = 0 \quad \because 1 = Y_2 + (-1) \cdot Y_2 \quad \text{--- (a)}$$

$$\mathbb{E}X = 0.$$

now, if I want to calculate the distance  
we will need to ignore the sign  
i.e.  $\mathbb{E}(X_n^2)$

$$\text{Let } X_n = Y_1 + Y_2 + \dots + Y_n.$$

$$X_n^2 = (Y_1 + Y_2 + \dots + Y_n)^2$$

$$= X_n^2 = Y_1^2 + Y_2^2 + \dots + Y_n^2 + 2 \sum_{i \neq j} Y_i Y_j$$

$$\mathbb{E}(X_n^2) = \mathbb{E}(Y_1^2) + \mathbb{E}(Y_2^2) + \dots + 2 \sum_{i \neq j} \mathbb{E}(Y_i^2 Y_j)$$

$$\mathbb{E}(Y_i^2) = 1.$$

$\mathbb{E}(Y_i^2 Y_j) = \mathbb{E}Y_i^2 \cdot \mathbb{E}Y_j \quad \because \text{they are independent}$   
and  $\mathbb{E}Y_i^2, \mathbb{E}Y_j^2 = 0$   
by ⑥

i.e.

$$\mathbb{E}(X_n^2) = 1 + 1 + \dots + n \text{ times} + 0$$

$$\boxed{\mathbb{E}X_n^2 = n}$$

## ⑦ Randomized Quicksort.

① pick a random element

② divide the array in 3 parts. 

③ Sort both pieces.

# Every pair is going to be compared at most once  
i.e. running time is at most  $n^2$ .

↓ remaining.

8th Sept

## \* Linearity of Expectations

### ① R. Quicksort



Let  $X$  be the number of comparisons.

$\hookrightarrow$  all pairs are compared at most once.

Let us say that sorted order is

$$a_1 < a_2 < a_3 < \dots < a_n$$

$$X_{ij} = \begin{cases} 1 & \text{if } a_i \text{ and } a_j \text{ are compared} \\ 0 & \text{otherwise.} \end{cases}$$

$$\therefore E[X] = \sum_{\text{Pair}(i,j)} X_{ij}$$

$$\hookrightarrow E[X] = E \sum_{\text{Pair}(i,j)} X_{ij} \quad \# \text{ we can take expectation inside}$$

$$\hookrightarrow E[X] = \sum_{\text{Pair}(i,j)} E[X_{ij}] \rightarrow \Pr[X_{ij}=1] \rightarrow \Pr[\text{we compare } a_i \text{ and } a_j]$$

$\hookrightarrow$  we don't compare element for the bigger array when we split the array.

$\hookrightarrow$  we compare when either  $a_i$  or  $a_j$  becomes Pivot before any other becomes the pivot.

$$\Pr[X=1] = \frac{2}{j-i+1} \quad \left. \begin{array}{l} \text{probability of selecting } i \\ \text{or } j \text{ as first pivot.} \end{array} \right.$$

$$\Rightarrow E[X] = \sum_{\text{Pair}(i,j)} \frac{2}{j-i+1} \quad \rightarrow \text{expected number of comp required in R. Quicksort.}$$

$$\hookrightarrow E[X] = 2 + \frac{2}{3} + \frac{2}{4} + \dots + \frac{2}{n}.$$

$$E[X] = 2(1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n})$$

$$E[X] = 2 \log n.$$

Q. Can we say that it is unlikely that random variable  $X$  is far from  $E[X]$ ?

### 6 MARKOV's INEQUALITY

$\hookrightarrow$  let  $X \geq 0$  be a Random variable

then  $\Pr[X > t] \leq \frac{E[X]}{t}$  where  $t > 0$

$\hookrightarrow$  lets say in a class of 100 average marks is 20 and we don't know what is most marks

then we can say.

not possible with negative marks.

$$\Pr[\text{Scoring more than 60}] \leq \frac{20}{60}$$

$$\Pr[X > 60] \leq \frac{1}{3}$$

### \* Proof.

$$E[X] = \sum_{i \geq 0} i \cdot \Pr[X=i].$$

$$= \sum_{\substack{i \geq 0 \\ i \geq t}} i \cdot \Pr[X=i] + \sum_{\substack{i \geq 0 \\ i \geq t}} i \cdot \Pr[X=i].$$

$$\geq \sum_{\substack{i \geq 0 \\ i \geq t}} i \cdot \Pr[X=i] \geq t \sum_{i \geq t} \Pr[X=i]$$

every  $i$  is at least  $t$  so we can make it constant.

$$E[X] \geq t \Pr[X \geq t]$$

$$\boxed{\Pr[X \geq t] \leq \frac{E[X]}{t}}$$

### Example

① Toss a coin 100 times.

$$\Pr[\text{getting more than 90 Heads}]$$

$$\Pr[X \geq 90] \leq ?$$

# Expectation doesn't pay the whole story.

Prob	A runs.	B
$\frac{1}{2}$	100	120

$$\frac{1}{2} = 0$$

Match 1.

$$E[X] = 20 \quad 20$$

Match 2.

$\hookrightarrow$  B is more Consistent while A has a lot more Variance

② We are given MCQ with +1 and -0 Marking Scheme

Strategy ①: Pick random from ABCD and then all answers are B.

Strategy ②: for each Question pick ABCD randomly

let  $X = \text{no. of correct answers}$

$$X_i = \begin{cases} 1 & \text{if } i^{\text{th}} \text{ question is right} \\ 0 & \text{otherwise.} \end{cases}$$

$$P(\text{i}^{\text{th}} \text{ question is correct}) = \frac{1}{4}$$

$$\mathbb{E}X = \frac{n}{4}$$

# but what if the exam setter selected all opt<sup>n</sup> to be same then

Strategy ① gets full marks with  $\frac{1}{4}$  Prob and 0 marks with  $\frac{3}{4}$  prob.

while Strategy ② gets the answer correctly consistently.

Q. how do we define Variance

↳ Intuitively it tells us how far away R or V goes away from the mean

Example

↳ Scores of a batsman.

$$5, 10, 15, 20, 15, 35, 40, 45, 0, 1, \dots$$

$$\mathbb{E}X = 20$$

$$[X_i - \mathbb{E}X]$$

$$\rightarrow -15, -10, -5, 0, \dots$$

↳ we cannot add as +ve and -ve cancel each other but hence we square and then add

i.e.

$$\text{Var}(X) = \mathbb{E}[X_i - \mathbb{E}X]^2 \text{ or } \sum P[X=i] \cdot (i - \mathbb{E}X)^2$$

or

$$\mathbb{E}(X^2 + u^2 - 2Xu)$$

$$= \mathbb{E}X^2 + u^2 - 2u\mathbb{E}X$$

$$\Rightarrow \mathbb{E}X^2 + u^2 - 2u^2 \quad \because \mathbb{E}X = u$$

$$\text{Var}(X) = \mathbb{E}X^2 - u^2$$

## Examples

① batsman

A	Prob	Score
$\frac{1}{5}$	100	
$\frac{4}{5}$	0	

$$\text{Mean} = 20$$

$$\text{Var} = \frac{1}{5}(100-20)^2 + \frac{4}{5}(0-20)^2$$

② Prob Score

1	20
0	0

$$\text{Var} = 1(20-20)^2 + 0(0-20)^2$$

$$\text{Var} = 0$$

② for question paper example from before examine sets A as solution for all questions.

Prob	marks (marks) <sup>2</sup>
$\frac{1}{4}$	$n$ $n^2$
$\frac{3}{4}$	0      0

$$\mathbb{E}X = \frac{n}{4}$$

$$\text{Var} = \mathbb{E}X^2 - \mathbb{E}X^2$$

$$= \frac{n^2}{4} - \frac{n}{4} = \frac{n}{4}(n-1)$$

\* Observations.

$$① \mathbb{E}KX = K\mathbb{E}X$$

$$② \text{Var}(KX) = K^2 \text{Var}(X)$$

③ if  $X_1, X_2, \dots, X_n$  all are independent

$$\Rightarrow \text{Var}(X_1) + \text{Var}(X_2) + \dots + \text{Var}(X_n) = \mathbb{E}(X - \mu)^2 = \text{Var}(X)$$

↳ Proof

$$u_i = \mathbb{E}X_i$$

$$\mu = u_1 + u_2 + u_3 + \dots + u_n \quad (\text{Linearity})$$

$$\mathbb{E}(X - \mu)^2 = \mathbb{E}[(X_1 + X_2 + \dots + X_n) - (u_1 + u_2 + \dots + u_n)]^2 = \mathbb{E}[(X_1 - u_1)^2 + (X_2 - u_2)^2 + \dots + (X_n - u_n)^2]$$

$$\Rightarrow \mathbb{E}(X_1 - u_1)^2 + \mathbb{E}(X_2 - u_2)^2 + \dots + \mathbb{E}(X_n - u_n)^2 + 2 \sum_{i \neq j} \mathbb{E}(X_i - u_i)(X_j - u_j)$$

$$\mathbb{E}(x_1 - \mu_1)^2 + \mathbb{E}(x_2 - \mu_2)^2 + \dots + \mathbb{E}(x_n - \mu_n)^2$$

$$+ 2 \sum_{i \neq j} \mathbb{E}[(x_i - \mu_i)(x_j - \mu_j)]$$

$\underbrace{\quad}_{0}$

④ Let  $X$  be a random variable

$Y = \frac{x_1 + x_2 + \dots + x_n}{n}$  where  $x_i$  is an independent copy of  $x$ .

$$\text{Var}(Y) = \frac{1}{n^2} [\text{Var}(x_1 + \dots + x_n)]$$

$$= \frac{1}{n^2} [\text{Var}(x_1) + \text{Var}(x_2) + \dots + \text{Var}(x_n)]$$

$$\Rightarrow \frac{n \text{Var}(x)}{n^2} \Rightarrow \boxed{\frac{\text{Var}(x)}{n}}$$

↳ taking average of multiple experiments reduces the Variance.

### Examples

①  $n$  times tossing of coins

$X \rightarrow$  no. of heads.  $\rightarrow$  independent

$$\mathbb{E}x = \frac{n}{2} \quad X = x_1 + x_2 + \dots + x_n$$

$x_i = \begin{cases} 1 & \text{if } i\text{th is head} \\ 0 & \text{otherwise} \end{cases}$

$$\text{Var}(x) = \text{Var}(x_1) + \dots + \text{Var}(x_n)$$

$$\text{Var}(x_i) = \mathbb{E}x_i^2 - (\mathbb{E}x_i)^2$$

$$\text{Var}(x_i) = 1 \cdot \frac{1}{2} - \left(\frac{1}{2}\right)^2$$

$$= \frac{1}{4}$$

$$\boxed{\text{Var}(x) = n \cdot \frac{1}{4}}$$

# If a random variable has small variance then chances of deviating away from  $\mathbb{E}x$  are small.

↳ ①  $\Pr[X \geq t] \leq \frac{\mathbb{E}x}{t}$

②  $\Pr[|X - \mu| \geq t] \leq \frac{\text{Var}(x)}{t^2}$  [Chebychev]

↳  $\Pr[(X - \mu)^2 \geq t^2] \leq \frac{\mathbb{E}y = \text{Var}(x)}{t^2}$

↳  $\Pr[X \geq 3\frac{n}{4}]$  where  $X$  is no. of heads in  $n$  tosses

①  $\Pr(X \geq 3\frac{n}{4}) \leq \frac{\mathbb{E}x}{3\frac{n}{4}} = \frac{2}{3}$  ↳ Due to square we get sharper bound.

②  $\Pr(|X - \frac{n}{4}| \geq \frac{n}{4}) \leq \frac{\text{Var}(x)}{(\frac{n}{4})^2} = \frac{4}{n}$