## Agenda

- 1) Bubble sort 3
  Algorithms
  2) Insertion sort
- 3) Merge 2 sorted arrays (oues)
- 4) Merge sort 3 Algorithm

What is sorting? Arranging the data

3 9 12 16 24 = inc. order

29 24 15 10 0 -2 => dec. order

# cg: 2 4 4 6 based on no. og

Jactors.

how to sort in Jawa

ind [] A = {3, 4, 1, 63;

Arrays. sort (A); { inc. order 3

11 A => 213463

Arrays-Soit (A) => T(: 0(nlogn)

## Bubble sort

3 8 6 2 4

itor: 3 6 2 4 8

it x 2 : 3 2 4 6 8

ito3: 3 2 4 6 8

itou; 2 3 4 6 8

3 8 6 2 4

2 4 6 8 8 8 i+(1: 3 8 6 2 4 6 1 2 3 4

2 8 6 itr2: 3 6 2 4 8 0 1 2 3 4

it 13: 2 3 4 6 8 9 4

its 4: 2 3 4 6 8

14 x : V= 1

ij (Arj) > Arj+17) {
swar;
3

code

```
void
     bubble-sout (int [] A) ?
     int n=A-length;
     Jox (int i=0; i< n-1; i++) }
                                                                  Tc: O(n2)
     Joselint j=0; j < n-i-1; j++) \bar{z}

ij(Arj) > Arj+13)

int temp = Arj;
Arj = Arj+13;
Arj = temp;
3
                                                                   Sc: 0(1)
3
       bubble-soft (int [] A) ?
void
     int n=A-length;
                                                   for (int i=0; i< n-1; i++) {
     iJ(Arij) > Arij+17)
int temp = Arij;
Arij = Arij+17;
Arij+17 = temp;
3
```

J

## Insertion sort

$$A = 3$$
 8 6 2 4 0 1 2 3 4

$$A = 3$$
 8 6 2 4 0 1 2 3 4

## ij(A137 > A13+13) 3

Swap 3

İ	
1	0 to 0
2	1 to 0
-	2 00 0
3	2 <del>t</del> o 0
4	3 to 0

```
roid insertion = sos+(int [DA)?

in+ n=A.long+n;
Jox(int i=1; i < n; i++)?
Jox(int j=i-1; j>=0; j--)?
SC: o(1)
int temp = A[j];
A[j+1] = temp;
Selse?
break;
Jox(int j=1; i < n; i++)?
A[j+1] = temp;
Jox(int j=1; i < n; i++)?
Jox(
```

3



roid insertion - sort (int [] A) ?
int n - A-length;

Jor(int i=1; i < n; i++)  $\bar{i}$ Jor(int j=i-1; j>=0; j--)  $\bar{i}$ ij(A[j] > A[j+1])  $\bar{i}$ int temp = A[j];

A[j] = A[j+1];

A[j+1] = temp;

3

else  $\bar{i}$ break;

θ= 7 4 2 5 0 1 2 3

1	o,				
1	0	4 72 0	λ γ	2 2	5 3
2	ەرا	2 1/ 0	24 7	7 / 2	5 3
3	عرا ر <sup>و</sup>	2	4	5 र्न 2	7 [   8   3

××

Q.1 Merge two sorted arrays.

Tc. O[n+m)

 $A \Rightarrow 2 5 9 12 15$ 

6 => 3 6 8 10 16 18

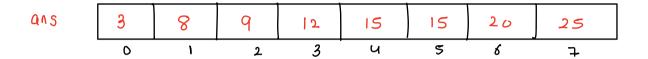
ans: В 

$$A = 2$$
 5 9 12 15 0 1 2 3 4

0ns: 2 3 5 6 8 9 10 12 15 16 18 0 1 2 3 4 5 6 7 8 9 10

j

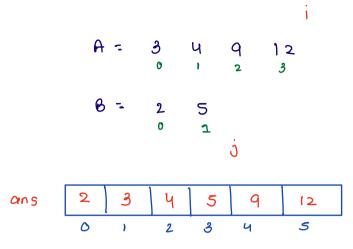
i



K

.

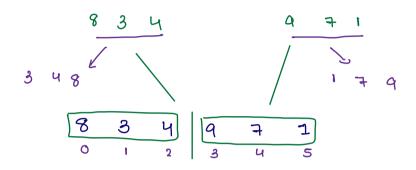
```
while(i < n \&\& j < m) {
    if(A[i] < B[j]) {</pre>
        //use A[i]
        ans[k] = A[i];
        i++;
        k++;
    else {
        //use B[j]
        ans[k] = B[j];
        j++;
        k++;
//if values are pending in A[]
while(i < n) {</pre>
    ans[k] = A[i];
    i++;
    k++;
//if values are pending in B[]
while(j < m) {
    ans[k] = B[j];
    j++;
    k++;
```



K

Merge sort TC: 0 (nlogn)

Li divide and ronguer



```
int mid = (dothi) |2;

int mid = (dothi) |2;

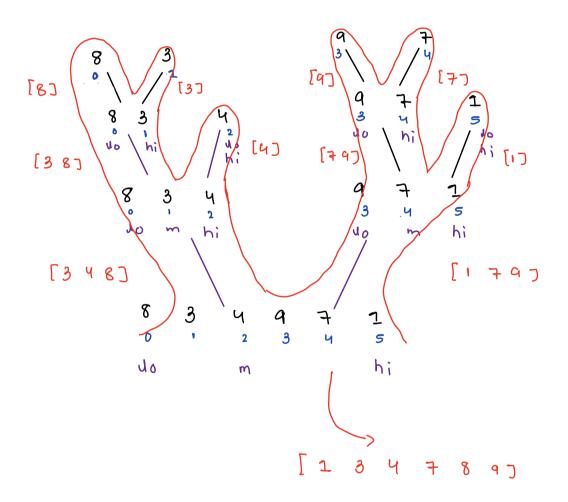
int() A = mergesort (arr, do, mid);

int() B = mergesort (arr, midt1, hi);

int() ans = merge (A, B);

return ans;
```

3



todo: understand complexities

T(: O(nlogn)

s(: 0 (n)

```
int mergesort (int[] arr, int do, int hi)

int mid = (dothi) | 2;

int [] sa = new int [];

int[] A = mergesort (arr, do, mid);

int[] B = mergesort (arr, midt1, hi);

int [] ans = merge(A,B);

return ans;
```

3

