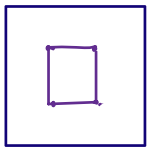
Agenda

1. Submatrix sum queries
2. Max submatrix sum
3. Sum of all submatrices

Submatrix : Part of matrix is a submatrix.



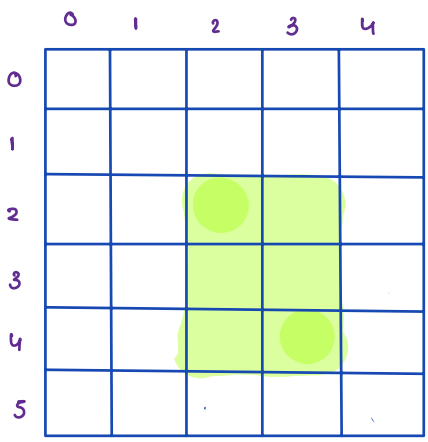
a) Single element is also a submatrix.

b) Complete matrix is also a submatrix.

c) We just need opposite corners to get submatrix.

Top left & bottom right OR

Top right & bottom left



TL	BR
2,2	4,3

Q.1 Given a N x M matrix and queries, for each query find submatrix sum.

queries

(TL)		(BR)		
x1	y1	x2	y2	ans
2	1	4	3	20
3	2	5	4	36
1	2	2	3	4

A ⇒

	0	1	2	3	4
0	7	1	-6	3	12
1	10	5	-1	0	9
2	6	4	-3	8	11
3	13	-8	-5	12	4
4	3	2	1	9	8
5	4	3	-2	6	3

i) Brute force : go on every query and then find sum that particular submatrix.

```
for (go on every query) {
        x1, y1      x2, y2
        int sum=0;
        for (i=x1; i<=x2; i++) {
                for (j=y1; j<=y2; j++) {
                        sum += A[i][j];
                }
        }
        sop ln (sum);
}
```

N * M

TC: O(N*M*Q)

ii) optimised → prefix Sum

1D array:

ps[i] ⇒ sum of elements from 0 to i.

2D array:

ps[i][j] ⇒ sum of submatrix from 0,0 (TL) to i,j (BR)

	0	1	2	3
0	3	2	4	1
1	-1	4	3	2
2	2	7	6	3

A

	0	1	2	3
0	3	5	9	10
1	2	8	15	18
2	4	17	30	36

PS

	0	1	2	3
0	3	2	4	1
1	-1	4	3	2
2	2	7	6	3

A

apply row-wise prefix sum →

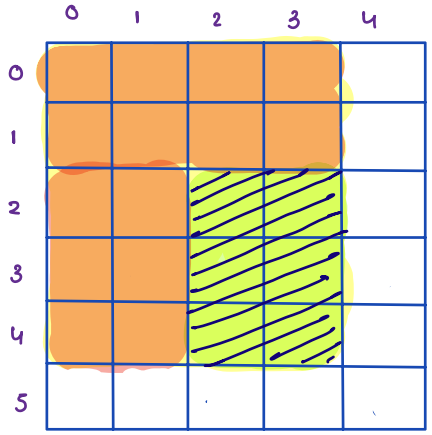
	0	1	2	3
0	3	5	9	10
1	-1	3	6	8
2	2	9	15	18

PS

apply col-wise prefix sum ↓

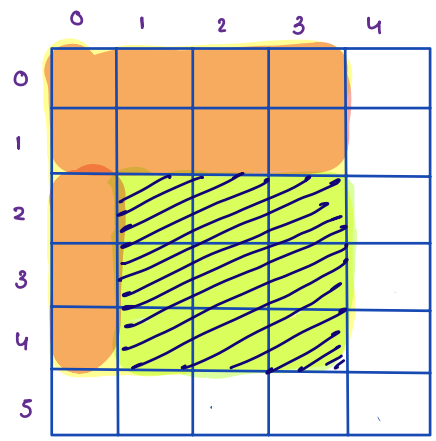
	0	1	2	3
0	3	5	9	10
1	2	8	15	18
2	4	17	30	36

```
int [][] prefixSum2D(int[][] A) {
    int n = A.length;
    int m = A[0].length;
    int[][] ps = new int[n][m];

    // apply prefix sum row by row
    for(int i=0; i<n; i++) {
        for(j=0; j<m; j++) {
            if(j==0) {
                ps[i][j] = A[i][j];
            }
            else {
                ps[i][j] = ps[i][j-1] + A[i][j];
            }
        }
    }

    // apply prefix sum col by col
    for(int j=0; j<m; j++) {
        for(int i=1; i<n; i++) {
            ps[i][j] = ps[i-1][j] + ps[i][j];
        }
    }

    return ps;
}
```



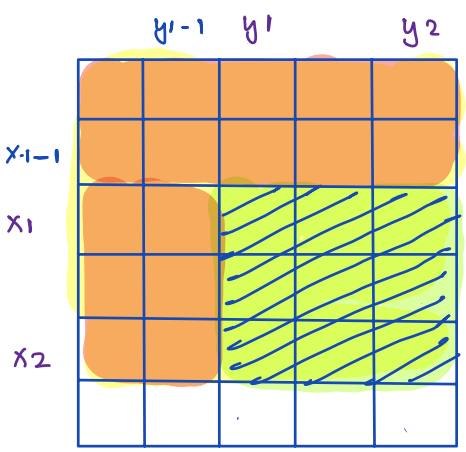
TL: 2,2 BR: 4,3

ps[4][3] - ps[1][3] - ps[4][1] + ps[1][1]



TL: 2,1 BR: 4,3

ps[4][3] - ps[1][3] - ps[4][1] + ps[0][1]

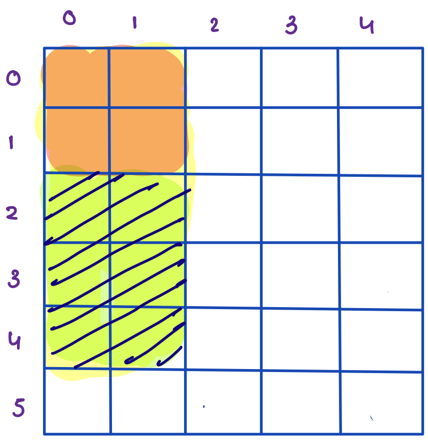


TL: x1, y1 BR: x2, y2

ps[x2][y2] - ps[x1-1][y2] - ps[x2][y1-1] + ps[x1-1][y1-1]

ps[x2][y2] - ps[x1-1][y2] - (ps[x2][y1-1] - ps[x1-1][y1-1])

ps[x2][y2] - ps[x1-1][y2] - ps[x2][y1-1] + ps[x1-1][y1-1]



TL	BR
(2,0)	(4,1)
x1, y1	x2, y2

ps[4][1] - ps[1][1] - ~~ps[4][-1]~~

\+ ~~ps[1][-1]~~

$ps[x2][y2] - ps[x1-1][y2] - ps[x2][y1-1] + ps[x1-1][y1-1]$

	0	1	2	3	4
0	3	2	4	1	6
1	-1	4	3	2	4
2	2	7	6	3	2
3	1	2	7	8	1

A

pfmat[][]

	0	1	2	3	4
0	3	5	9	10	16
1	2	8	15	18	28
2	4	17	30	36	48
3	5	20	40	54	67

ps

Query

x1	y1	x2	y2
1	2	2	4

$ps[2][4] - ps[0][4] - ps[2][1] + ps[0][1]$

$48 - 16 - 17 + 5 = 20$

```
void solve(int[][] A, int[][] Q) {
    int[][] PS = prefixSum2D(A);          } → 2*(N*M)

    for(int i=0; i< Q.length; i++) {      → Q
        int x1 = Q[i][0];
        int y1 = Q[i][1];
        int x2 = Q[i][2];
        int y2 = Q[i][3];
        // find sum of submatrix TL: x1,y1 & BR: x2,y2
        int sum = 0;
        sum += PS[x2][y2];

        if(x1 > 0)
        sum -= PS[x1-1][y2];

        if(y1 > 0)
        sum -= PS[x2][y1-1];

        if(x1 > 0 && y1 > 0)
        sum += PS[x1-1][y1-1];

        sop.println(sum);
    }
}
```

itr: $2(N*M)+Q$

TC: $O(N*M+Q)$

SC: $O(N*M)$

Q.2 Given row wise and column wise sorted matrix, find maximum submatrix sum.

A =

	0	1	2	3
0	-20	-16	-4	8
1	-10	-8	2	14
2	-1	6	21	30
3	5	7	28	42

In ans submatrix, we need to include the max. element.

↓

the max element of A[][] is always present (n-1, m-1)

↓

BR is already fixed (n-1, m-1)

go on all valid TL and find out best answer possible.

A =

	0	1	2	3
0	-20	-16	-4	-1
1	-10	-8	-2	5
2	-4	2	4	8

A =

	0	1	2
0	-50	-40	-30
1	-25	-20	-15
2	-19	-14	-3

Top lefts

0,0
0,1
0,2
0,3
1,0
1,1
1,2
1,3
2,0
2,1
2,2
2,3

A =

	0	1	2	3
0	-20	-16	-4	-1
1	-10	-8	-2	5
2	-4	2	4	8

```
int solve(int[][] A) {
    int n = A.length;
    int m = A[0].length;
    int[][] PS = prefixSum2D(A);
    int x2 = n-1;
    int y2 = m-1;
    int ans = Integer.MIN_VALUE;
    // go on all possible top-left
    for(int i=0; i<n; i++) {
        for(int j=0; j<m; j++) {
            int x1 = i, y1 = j;
            // find sum of submatrix TL: x1,y1 & BR: x2,y2
            int sum = 0;
            sum += PS[x2][y2];
            if(x1>0)
            sum -= PS[x1-1][y2];
            if(y1>0)
            sum -= PS[x2][y1-1];
            if(x1>0 && y1>0)
            sum += PS[x1-1][y1-1];
            if(sum > ans) {
                ans = sum;
            }
        }
    }
    return ans;
}
```

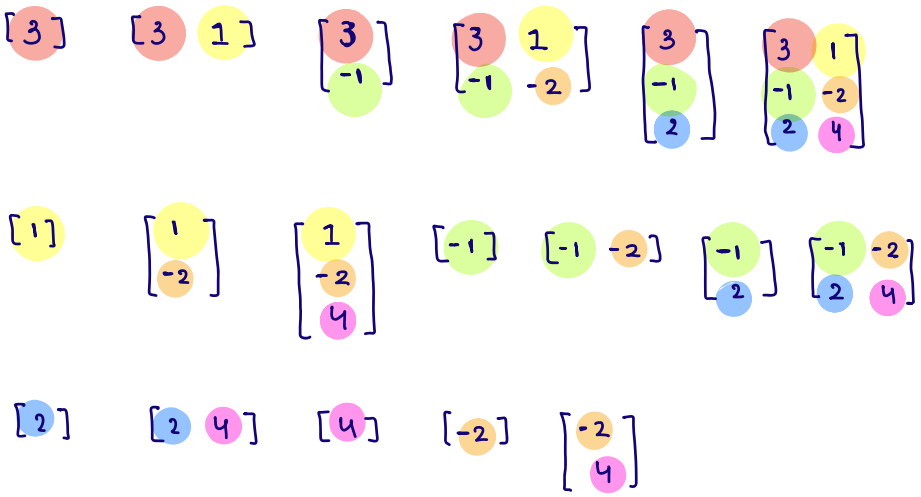
BR

→ all possible TL

Q-3 Given a N×M matrix, find **sum of all submatrices sum.**

A =

	0	1
0	3	1
1	-1	-2
2	2	4



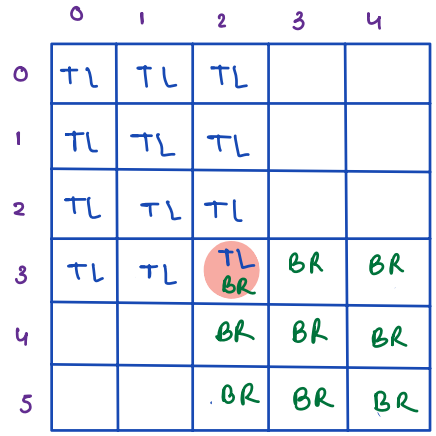
ans = 36

ans = 6*3 + 6*1 + 8*(-1) + 8*(-2) + 6*2 + 6*4

18 + 6 + -8 - 16 + 12 + 24 = 36

Conclusion: if A[i][j] is coming x submatrices so its contribution in ans will be

x * A[i][j]

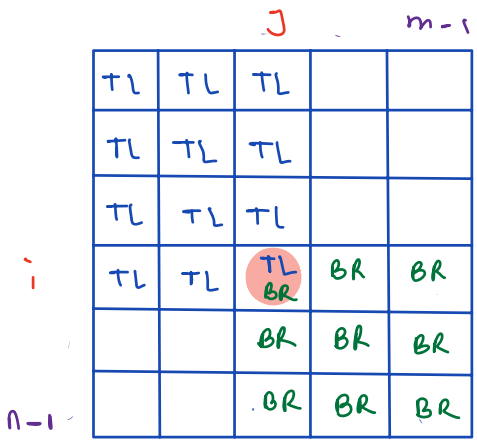


valid TL → 12

valid BR → 9

total submatrices = 12 * 9

= 108



valid TL : $(i+1)*(j+1)$

valid BR : $(n-i)*(m-j)$

total submatrices in which A[i][j] is present

=> $(i+1)*(j+1)*(n-i)*(m-j)$

```
int solve (int [][] A) {
    int n = A.length;
    int m = A[0].length;
    int ans = 0;
    for (int i=0; i<n; i++) {
        for (int j=0; j<m; j++) {
            int freq = (i+1)*(j+1)*(n-i)*(m-j);
            ans += freq*A[i][j];
        }
    }
    return ans;
}
```

Doubts

n = 3

m = 2

A =

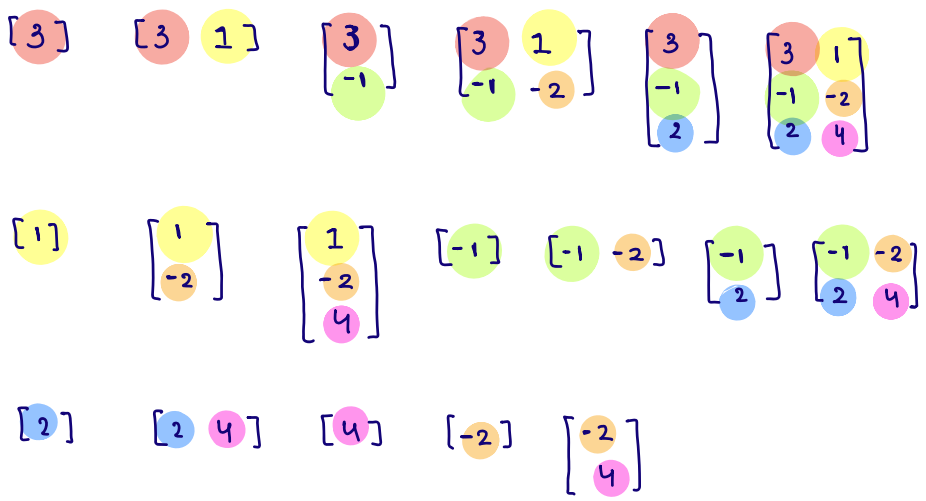
	0	1
0	3	1
1	-1	-2
2	2	4

```
int ans= 0;
for(int i=0; i<n; i++) {
    for(int j=0; j<m; j++) {
        int freq = (i+1)*(j+1)*(n-i)*(m-j);
        ans+= freq* A[i][j];
    }
}
```

3

3

i j	freq	impact
0 0	1*1*3*2=6	18
0 1	1*2*3*1=6	6
1,0	2*1*2*2=8	-8
1,1	2*2*2*1=8	-16
2,0	3*1*1*2=6	12
2,1	3*2*1*1=6	24
		36



6*3 + 6*1 + 8*(-1) + 8*(-2) + 6*2 + 6*4