

Agenda

- i) Design min Stack
- ii) Nearest Smaller
 ↳ on left
 ↳ on right
 } **
- iii) Largest Area histogram

Q.1 Design a Stack that supports push, pop, top and min functions in $O(1)$ time.

10 5 9 18 12 3 7 17

minStack st = new minStack();

st.push(10)

st.push(5)

st.push(9)

st.push(18)

st.push(12)

st.push(3)

st.push(7)

st.min() → 3

st.pop()

st.pop()

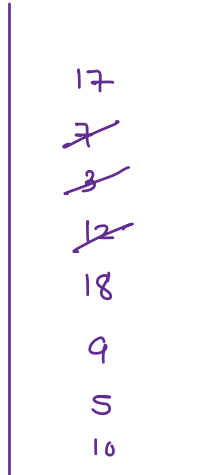
st.min() → 5

st.top() → 12

st.pop()

st.push(17)

st.top() → 17



class minStack {

void push(int x)

void pop()

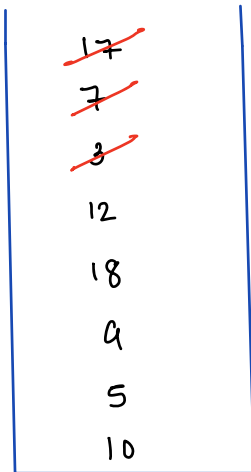
$O(1)$

int top()

int getmin()

}

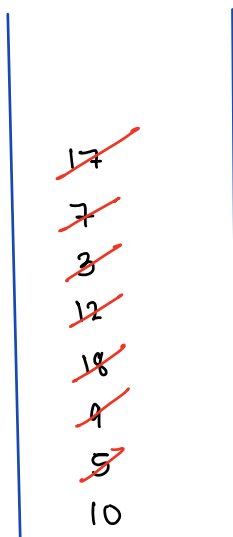
↓
10 5 4 18 12 3 7 17



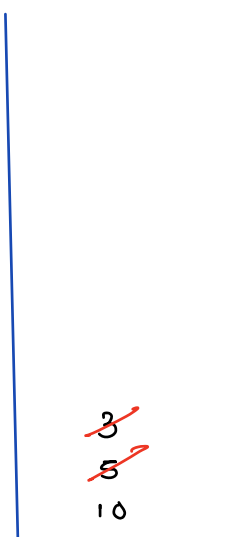
min = ~~10~~ ~~18~~ 3

idea: we need separate stack to manage min values.

↓
10 5 4 18 12 3 7 17



valst



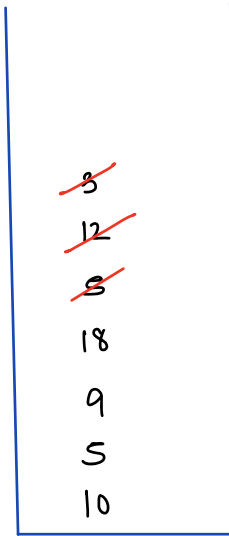
minst

$x \leq \text{minst}.\text{peek}()$

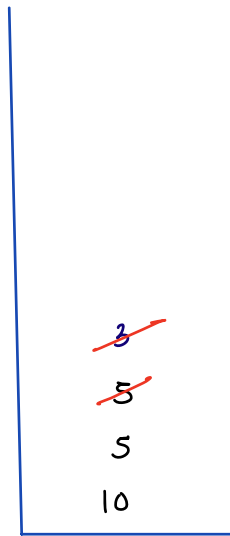
→ $\text{minst}.\text{push}(x);$

↓

10 5 9 18 5 12 3 7 3 17



valst



minst

```
class MinStack {  
    → two stack : valst, minst
```

```
void push(int x) {
```

```
    → manage size=0 condition
```

```
    → add x in valst
```

```
    → add x in minst ( x <= minst.peek() )
```

```
}
```

```
void pop() {
```

```
    → manage size=0 condition
```

```
    → temp = pop from valst
```

```
    → pop from minst ( temp == minst.peek() )
```

```
}
```

```
int top() {
```

```
    → manage size=0 condition
```

```
    → return valst.peek()
```

```
}
```

```
int getmin() {
```

```
    → manage size=0 condition
```

```
    → return minst.peek()
```

```
}
```

```
}
```

```

class Solution {
    Stack<Integer>st = new Stack<>();
    Stack<Integer>minSt = new Stack<>();

    public void push(int x) {
        if(st.size() == 0) {
            st.push(x);
            minSt.push(x);
        }
        else {
            st.push(x);

            if(x < minSt.peek()) {
                minSt.push(x);
            }
        }
    }

    public void pop() {
        if(st.size() == 0) {
            return;
        }
        else {
            int top = st.pop();

            if(top == minSt.peek()) {
                minSt.pop();
            }
        }
    }

    public int top() {
        if(st.size() == 0) {
            return -1;
        }
        else {
            return st.peek();
        }
    }

    public int getMin() {
        if(st.size() == 0) {
            return -1;
        }
        else {
            return minSt.peek();
        }
    }
}

```

Q.2 Nearest / next smaller on left.

for every element find out the value of nearest smaller on left.
↳ in terms of distance

A = [10 16 5 9 12 8 25 7 13]

ans -1 10 -1 5 9 5 8 5 7

A = [18 3 13 14 5 24 4]

ans -1 -1 3 13 3 5 3

Expected TC: $O(n)$

A = [10 16 5 9 12 8 25 7 13]

ans -1 10 -1 5 9 5 8 5 7

↓

What?

13
7
25
8
12
9
5
16
10
st

✓

A = [10 16 5 9 12 8 25 7 13]

ans = [-1 10 -1 5 9 5 8 5 7]

ans[n]

ans[0] = -1, st.push(A[0]);

for (i → 1 to A.length-1) {

while (st.size() > 0 & st.peek() >= A[i]) {
st.pop();

}

if (st.size() == 0) {

ans[i] = -1;

}

else {

ans[i] = st.peek();

}

st.push(A[i]);

}

13

7

~~25~~

~~8~~

~~12~~

~~9~~

5

~~16~~

~~10~~

st

TC: $O(n)$

SC: $O(n)$

Nearest / next smaller on left (Index based)

for every element find out the index of next nearest smaller on left.

↓

	0	1	2	3	4	5	6	7	8
A =	[10	16	5	9	12	8	25	7	13]
ans	-1	0	-1	2	3	2	5	2	7

for (i → 1 to A.length-1) {

while (st.size() > 0 && A[st.peek()] >= A[i]) {
 st.pop();

}

if (st.size() == 0) {

 ans[i] = -1;

}

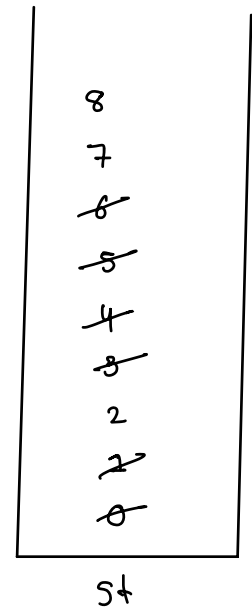
else {

 ans[i] = st.peek();

}

st.push(i);

}



Q.3 Nearest / next smaller on right.

For every element find the value of nearest / next small on right

A = [10 16 5 9 12 8 25 7 3]

ans 5 5 3 8 8 7 7 3 -1

same logic as that of next smaller left, the only change is direction of travelling should be right to left.

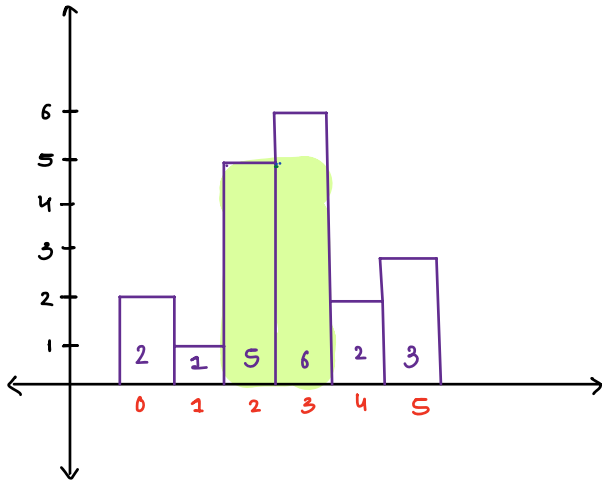


Nearest / next smaller on right (Index based)

For every element find the idx of nearest / next small on right

0.4 Largest area histogram

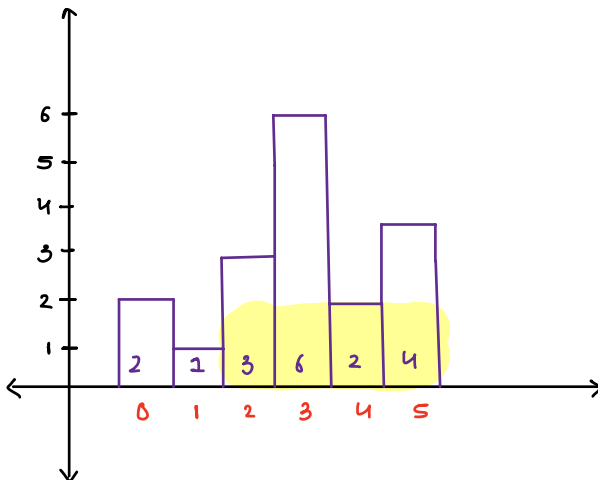
$A = [2 \ 1 \ 5 \ 6 \ 2 \ 3]$



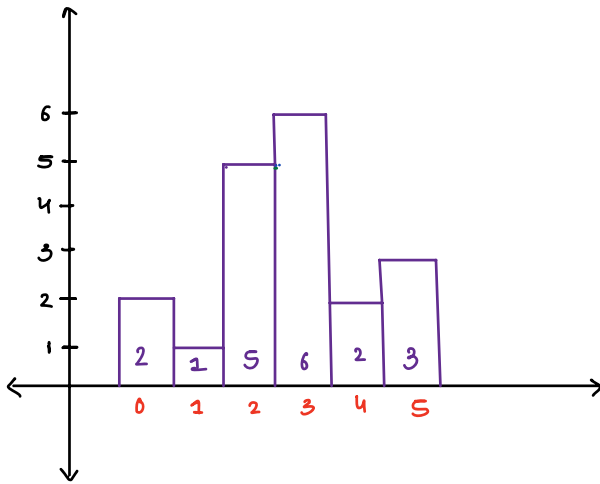
Return the area of
Largest rectangle
possible?

$$\text{ans} = 5 * 2 = 10$$

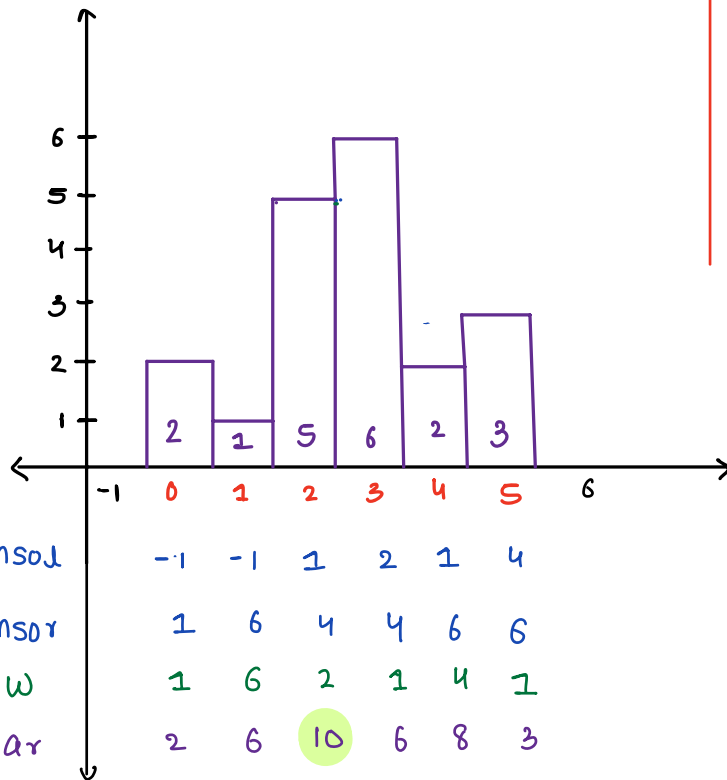
$A = [2 \ 1 \ 3 \ 6 \ 2 \ 4]$



$$\text{ans} = 2 * 4 = 8$$



i	h	w	ar
0	2	1	2
1	1	6	6
2	5	2	10
3	6	1	6
4	2	4	8
5	3	1	3



nsol	-1	-1	1	2	1	4
nsor	1	6	4	4	6	6
w	1	6	2	1	4	1
ar	2	6	10	6	8	3

$$h = A[i]$$

$$w = nsor[i] - nsol[i] - 1$$

$$ar = h * w$$

if (ar > max) {

max = ar;

}

Doubts

((a + b)) redundant +

→ 1 (yes redundant brace is present)

(a + (a + b))

→ 0 (no redundant brackets)

str: (a+(a+b))

for (i → 0 to str.length()-1)

if (ch == '(')

st.push(ch);

}

else if (ch == '+' || ch == '-' ||

ch == '*' || ch == '/')

st.push(ch);

}

else if (ch == ')')

// keep count of content (ops) till opening

int cnt = 0;

while (st.peek() != '(')

cnt++;

st.pop();

}

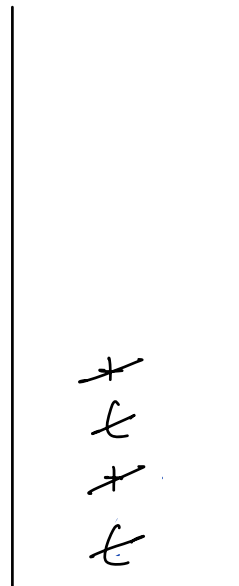
st.pop();

if (cnt == 0) return 1;

}

}

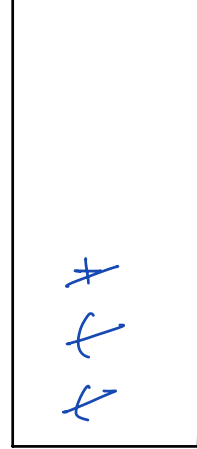
return 0;



st
(char)

cnt = 0
1

$(a+b)$ ↓



cnt = 0
↳ return 1