

Agenda

- i) stack basics
- ii) Double character trouble
- iii) Expression evaluation **

Stack : It follows LIFO

↳ Last In first out

```
Stack < Integer > st = new Stack < > ();
```

↓
name of
variable

st.push(10);

st.push(20);

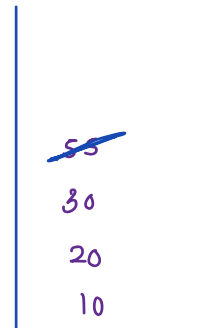
st.push(30);

st.push(55);

System.out.println(st.peak()); → 55

st.pop();

System.out.println(st.size()); → 3



all the above 4 functions are TC: O(1)

st. peek()	st. pop()
It gives us the topmost element of the stack.	It removes the topmost element of stack and also return us that removed value.

Stack < Integer > st = new Stack < > ();

st.push(10);

st.push(15);

st.push(35);

st.push(45);

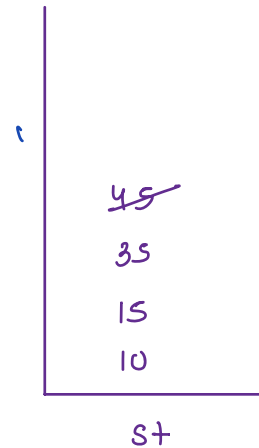
System.out.println(st.size()); → 4

System.out.println(st.peek()); → 45

System.out.println(st.size()); → 4

System.out.println(st.pop()); → 45

System.out.println(st.size()); → 3



Real life examples of stack

i) undo / redo

ii) ← (back) browsing

iii) Expression evaluate

Q. Double character trouble

⇒ Given a string, keep removing the same consecutive char until no more occurrence of same consecutive char remains.
Return the final answer string.

Note: Same consecutive chars are coming in doubled manner.

str = a b ~~bc~~ c

↓

a ~~b~~ b c

↓

a c

Expected Tc: $O(n)$

str = k m b ~~aab~~ m g j

↓

k m ~~b~~ b m g j

↓

k ~~m~~ m g j

↓

k g j

abccbc ↓

c
~~c~~
~~b~~
a

str = ca

ans = ac

kmbacbmjg ↓

j
g
~~a~~
~~b~~
~~m~~
k

str = jgk

ans = kgj

bcaacbzyxxgw ↓

w
~~x~~
~~y~~
z
~~a~~
~~c~~
~~b~~

str = wz

ans = zw

```
for ( ch of str ) {
```

```
    if ( st.size() == 0 || st.peek() != ch ) {
```

```
        st.push(ch);
```

```
    }  
    else if ( st.peek() == ch ) {
```

```
        st.pop();
```

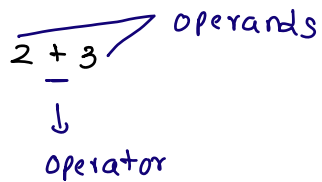
```
    }
```

```
}
```

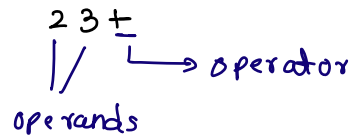
// travel stack and calculate final answer.

Expression Evaluation and conversion

Infix expression

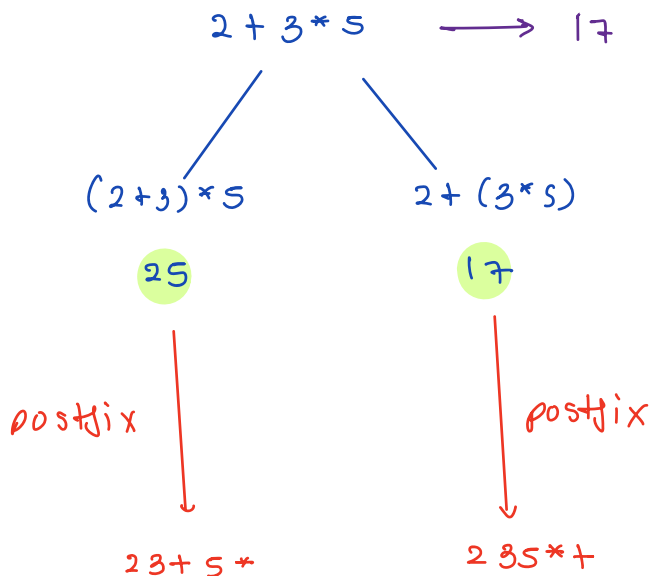
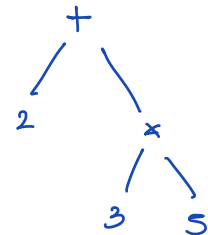


Postfix expression



Infix exp: $2 + 3 * 5 \Rightarrow 17$

Postfix exp: $2\ 3\ 5\ * +$



In postfix exp the order of operators is exactly same as the order in which the operators should get evaluated.

Advantages of postfix:

- no brackets
- no operator priority

Postfix evaluation

Given postfix expression, evaluate it and return final answer.

ops: *, /, -, +

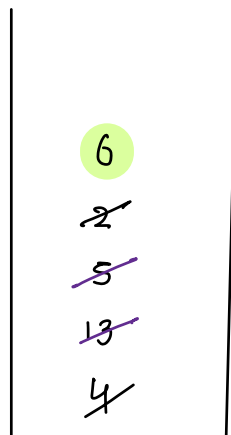
↓

eg1 ["2", "1", "+", "3", "*"] \Rightarrow 9

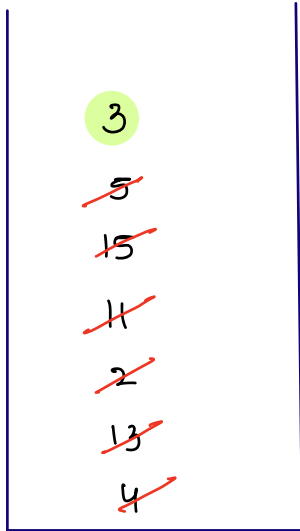


↓

eg2 ["4", "13", "5", "/", "+"] \Rightarrow 6



↓
["4", "13", "2", "-", "+", "5", "/"]



```
if (str is operator) {  
    // evaluate  
    int v2 = st.pop();  
    int v1 = st.pop();  
    calculate res of v1 str v2  
    st.push(res);  
}  
else {  
    int val = str to int  
    st.push(val);  
}
```


Infix to postfix conversion

lowercase letters \rightarrow operands

operators \rightarrow $^$, $+$, $-$, $*$, $/$

brackets

eg1 : "x^y / (a+b*c-d) - e + j*g-h" ↓

$^ \rightarrow$ power

priority

=

$^$

$/, *$

$+, -$

~~xy^abc*+d- / e-jg*+h-~~

~~h~~

~~xy^abc*+d- / e-jg*+~~

~~jg*~~

~~g~~

~~j~~

~~xy^abc*+d- / e-~~

~~e~~

~~xy^abc*+d- /~~

~~abc*+d-~~

~~d~~

~~abc*~~

~~bc*~~

~~c~~

~~b~~

~~a~~

~~xy^~~

~~y~~

~~x~~

operands

~~/~~
~~*~~
~~+~~
~~-~~
~~-~~
~~*~~
~~+~~
~~-~~
~~+~~
~~-~~

operator

if ch is operand then push it to operand stack

if ch is an opening bracket push it to operator stack

if ch is closing bracket

↳ evaluate till an opening brackets comes on
operator. peek()

if ch is an operator

↳ evaluate till higher or equal priority operators
are coming on operator. peek() but

→ stop if stack becomes empty

→ stop if operator. peek() becomes '('

add ch to operator stack

final evaluation