# CS & IT

## ENGINERING

Theory of Computation

DFA

Lecture No.- 01



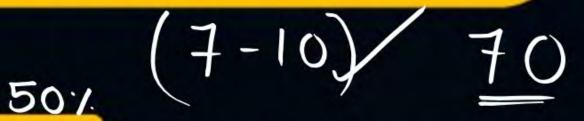
By- Venkat Rao sir

## THEORY OF COMPUTATION











Topic Finite Automaton & Regular Languages.

25%

Topic

Pushdown Automata & Context free Languages.



Turing Machine & Recursive Enumerable Languages.



Undecidability.

## **BOOKS:**





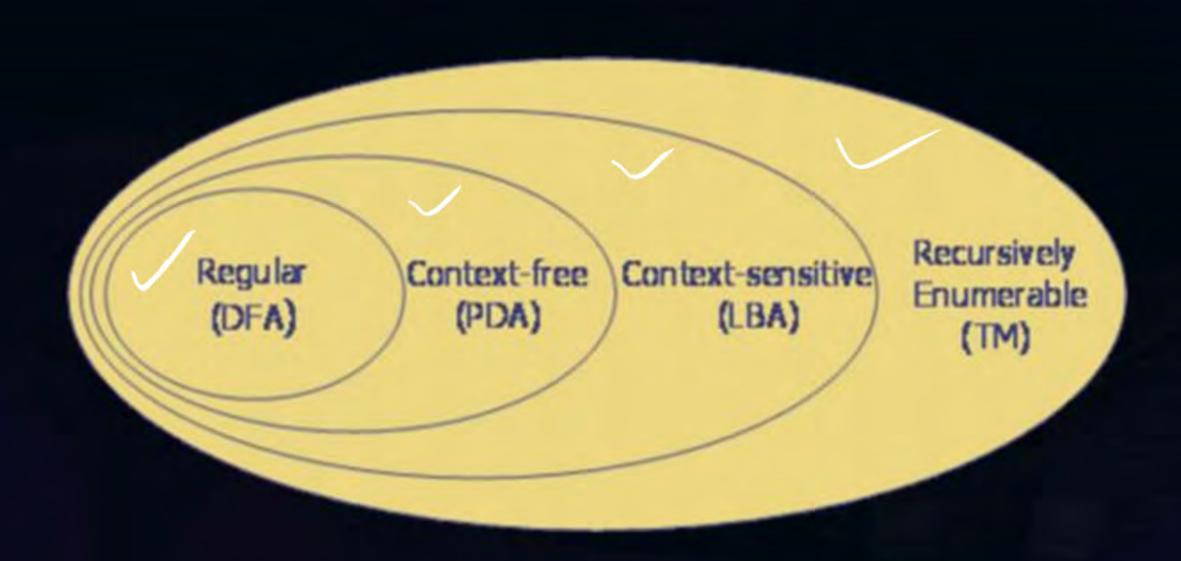






## **Topic: Theory of Computation**







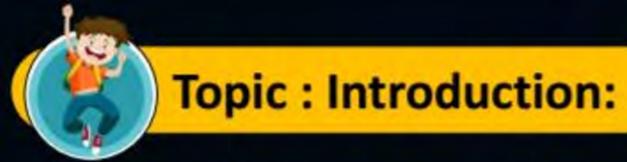
#### **Topic: Introduction:**



It is the mathematical study of computing machines and their capability

or

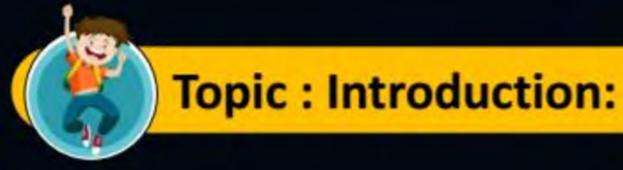
It is the study of automata theory and formal languages.





## **Applications of Theory of Computation:**

- \*Algorithm design and analysis
  - Compiler design
- Cryptography and network security
- Artificial intelligence and machine learning
- Database systems and query optimization
- Software verification and model checking





Decidable Problem Algo exist

Undecidable Problem no was

#### **Topic: Terminologies:**



Alphabet( $\Sigma$ ): Finite non-empty set of symbols

Ex:- 
$$\{a, b\}$$
 -  $\{a, 1, 2\}$  -  $\{\}$  -

#### **Topic: Terminologies:**



Alphabet( $\Sigma$ ): Finite non-empty set of symbols



#### **Topic: String:**



**String:** Finite sequence of symbols over the given alphabet  $\Sigma$ .

Ex:-



#### Topic: String:

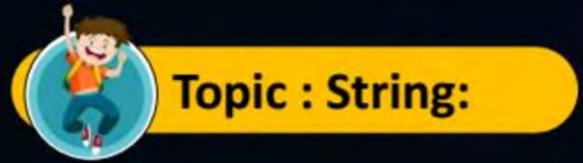




**Language**: Any set of strings over the given alphabet  $\Sigma = \{a, b\}$ .

$$L_6 = \{ \epsilon, (a, b) | aa, ab, ba, bb \} \longrightarrow infin$$

Complete Languege





Sub - String: Consecutive sequence of symbols over the given string.

Total no of substring for the given string = n(n + 1)/2 + 1

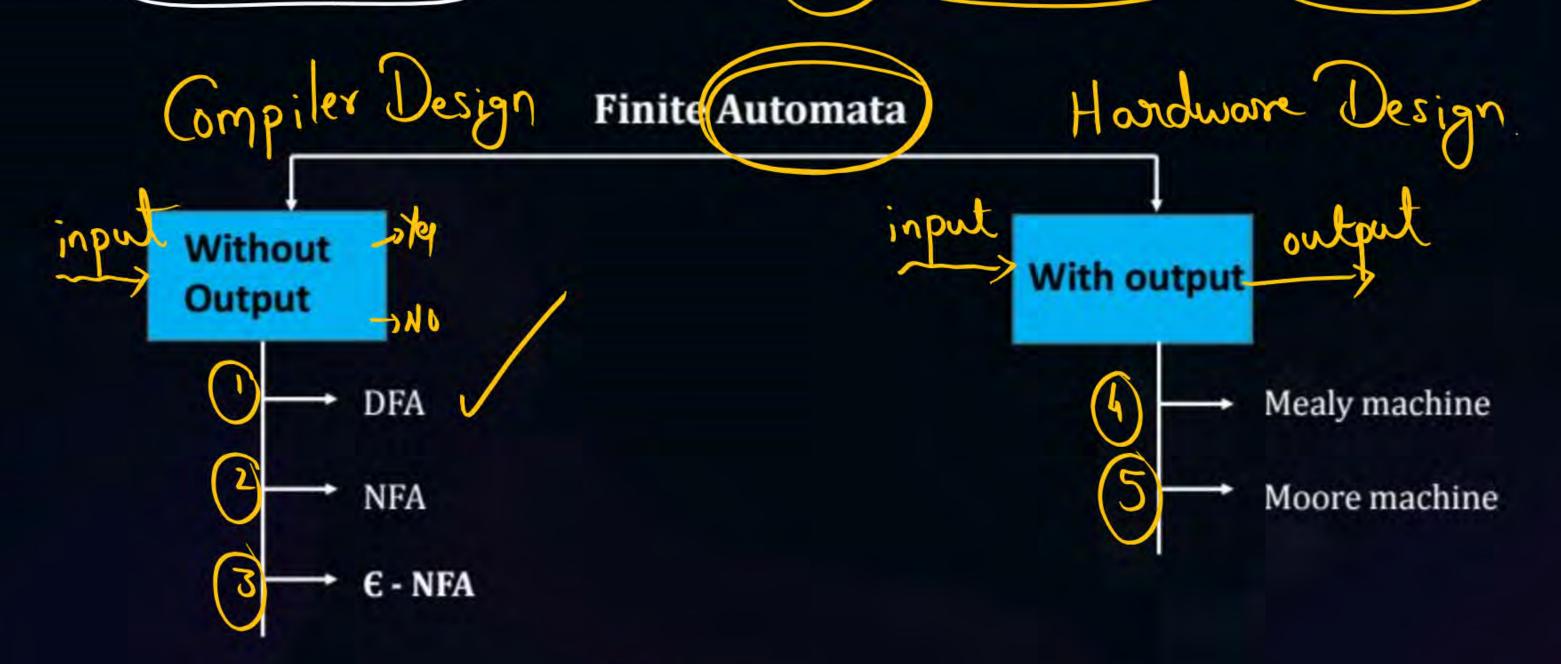




## Topic: Finite Automata



It is a mathematical model which contains finite number of states and transitions.



States FAN 0 transitions Lexical Analysis:



String Matching

Network Protocol Analysis:

Digital Circuit Design:

Regular Expression Engines:

**Natural Language Processing** 





DFA: It is a finite automata in which from every state on every input symbol exactly one transition should exits.





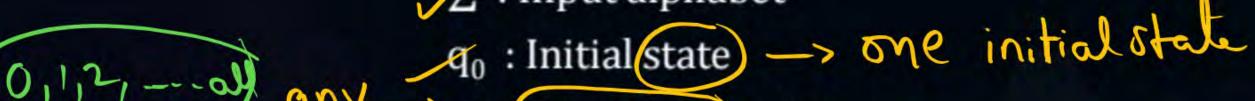
#### **FORMAL DFA:**

DFA is defined as

DFA = 
$$(Q, \Sigma, q_0, F, \delta)$$

Q: Finite set of states

芝: Input alphabet



F : Set of final states

$$\delta$$
: Transition function Q\*  $\Sigma$  → Q

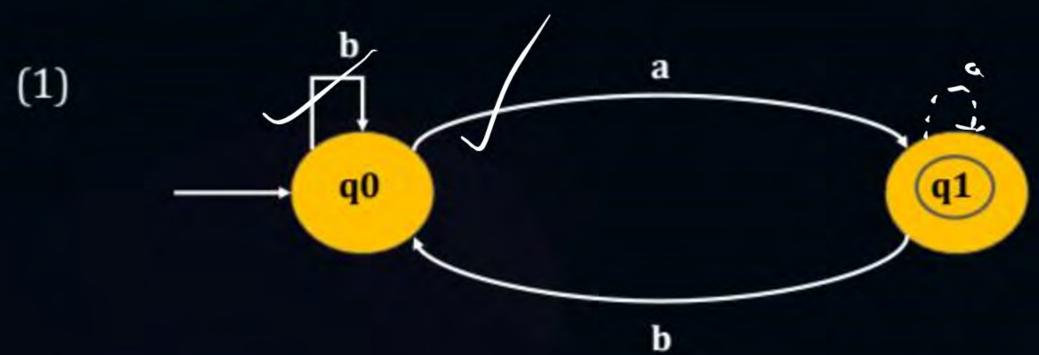




DFA?

 $\geq = \{a, b\}$ 

## Example of ::



not a DFA.







$$\leq = \{a, b\}$$

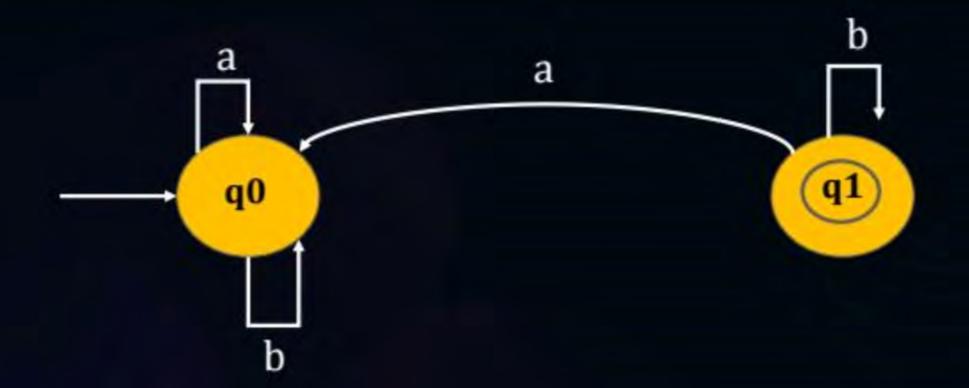
(2) 
$$\frac{a,b}{q_0}$$
 
$$\frac{\delta(q_0,a) = q_0(a)q_1}{b}$$





#### Example of DFA:



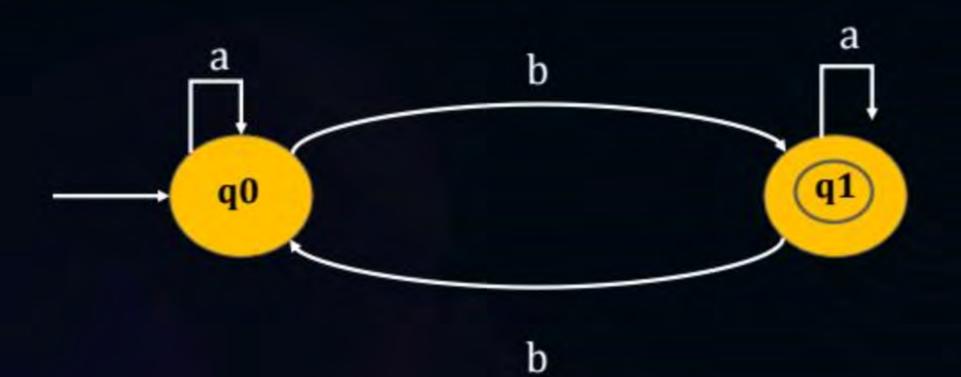






#### Example of DFA:

(4)



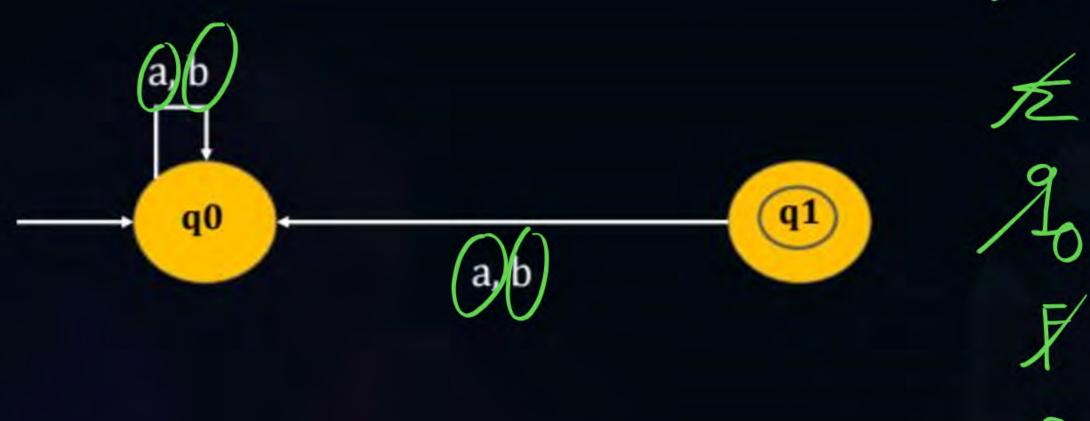




#### Example of DFA:



(5)



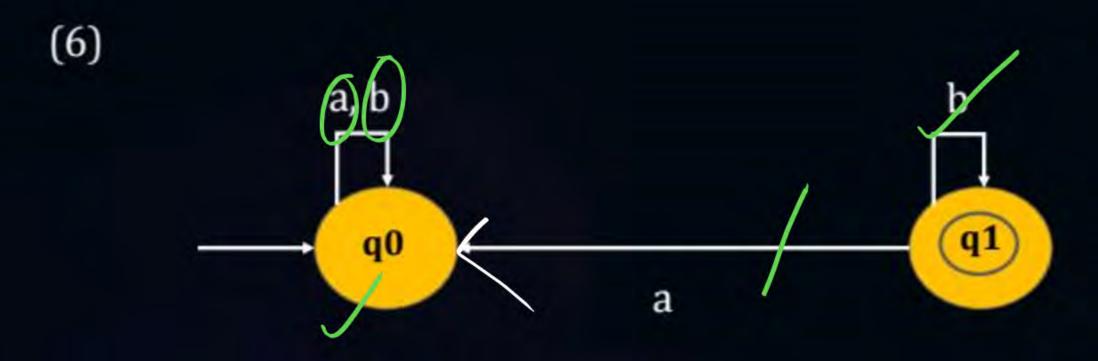


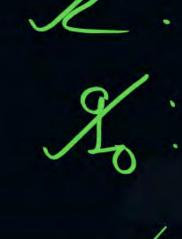


#### Example of DFA:









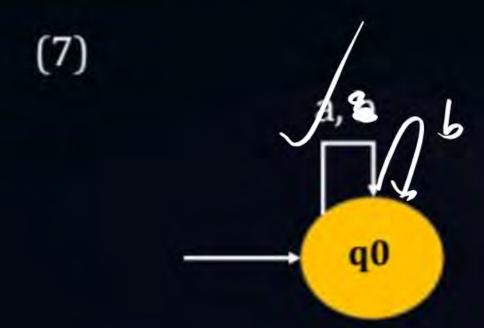




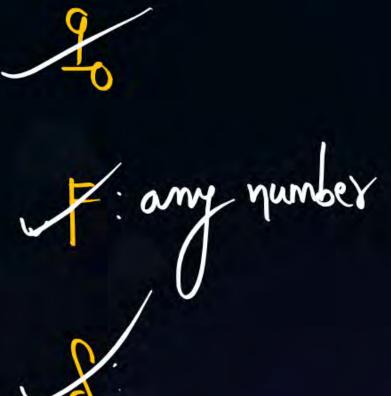


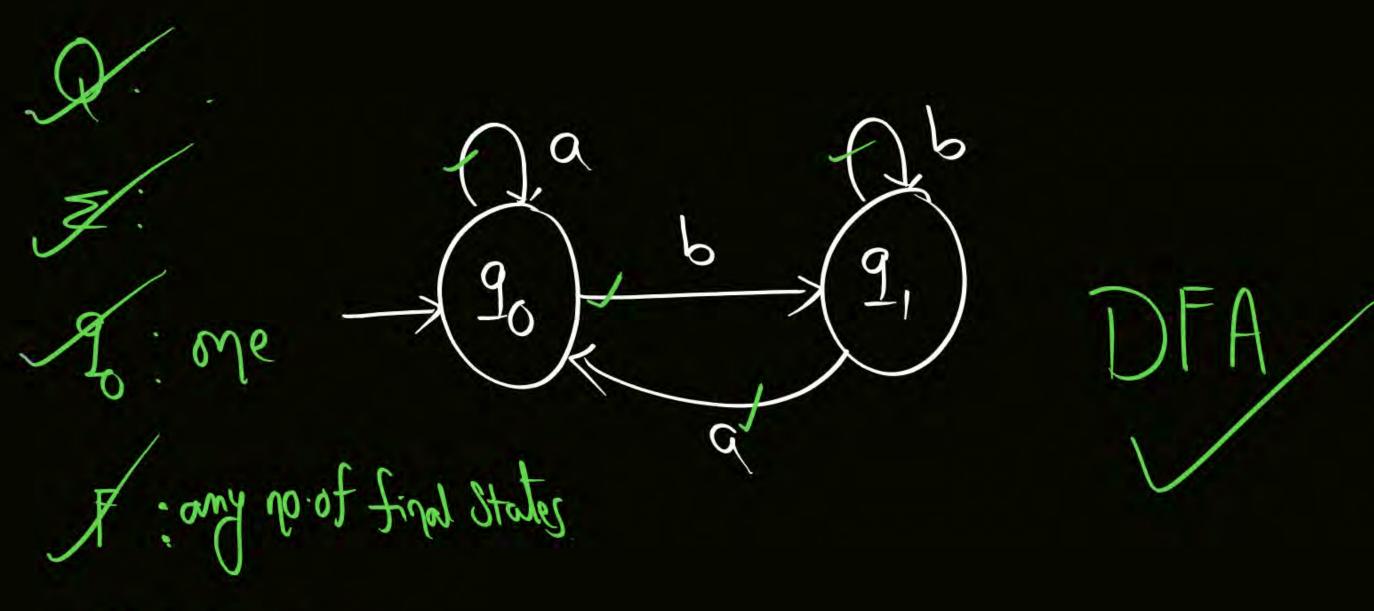












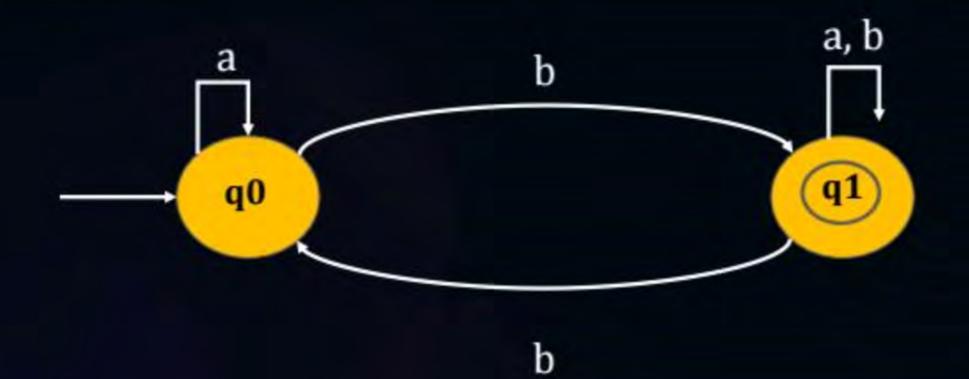
8:0x5>0





#### Example of DFA:

(8)

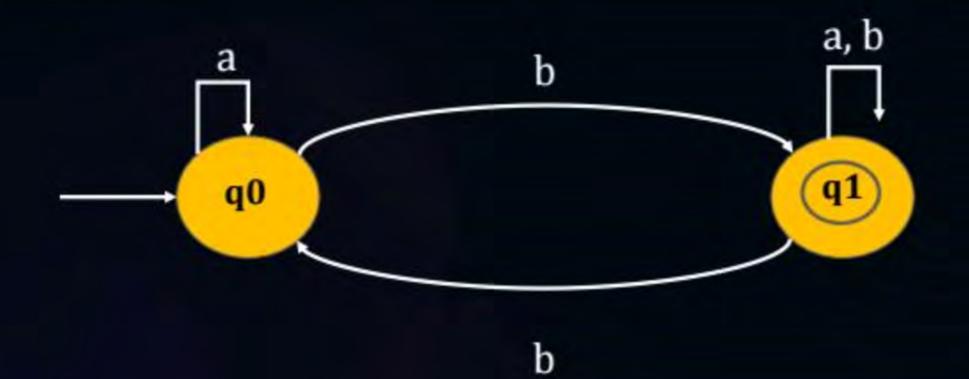






#### Example of DFA:

(8)



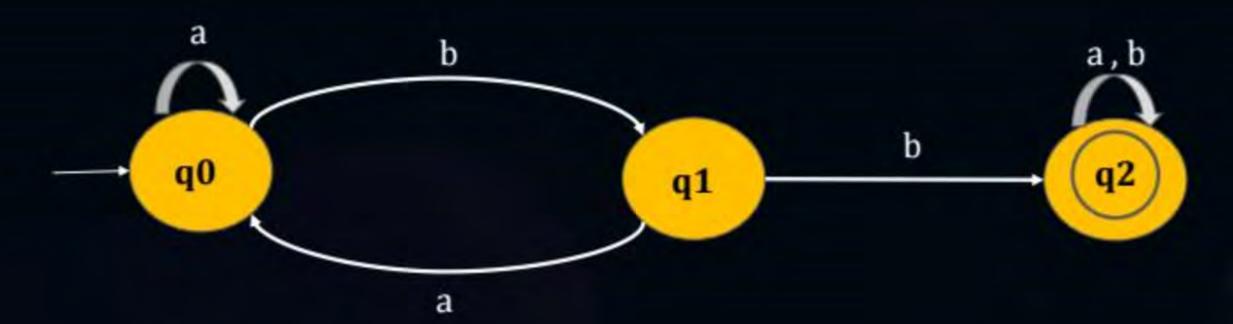


#### DFA acceptance method:

- 1.Start at the Initial State: Begin at the initial state of the DFA.
- 2.Read Input Symbols: For each symbol in the input string, read it one by one.
- 3.Follow Transitions: Based on the current state and the input symbol being read, follow the transition defined by the transition function of the DFA. This transition function specifies the next state of the automaton for each combination of current state and input symbol.
- 4.Repeat Until End of Input: Continue this process of reading input symbols and following transitions until you reach the end of the input string.
- 5.Final State: Once you have processed all input symbols, check the current state of the DFA. If it is one of the accepting states (states designated as final states), then the input string is accepted. Otherwise, it is rejected.
- 6.Acceptance: If the DFA halts in an accepting state after reading the entire input string, then the Input is accepted.



#Q. Identify language accepted by given DFA



Set of all strings

A Starting with bb

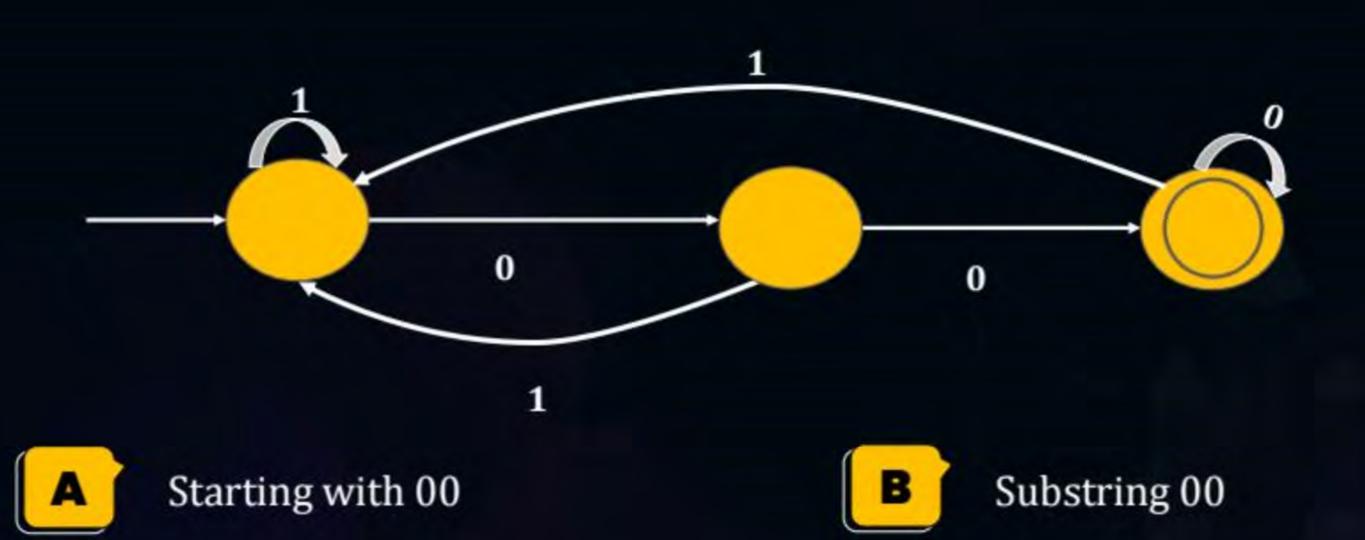
Contains at least 2 b's

B Ending with bb

D None

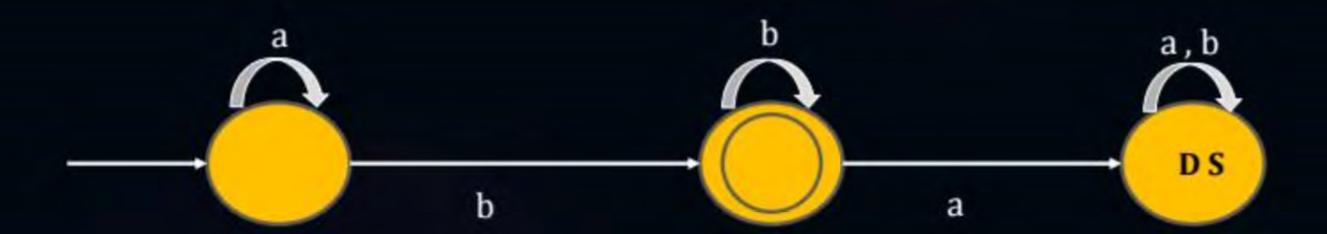


Ending with 0



None





- **A**  $L = \{a^n b^m | n, m ≥ 1\}$
- C  $L = \{a^n b^m | n, m \ge 0 \}$

- **B**  $L = \{a^n b^m | n \ge 1, m \ge 0\}$
- D None

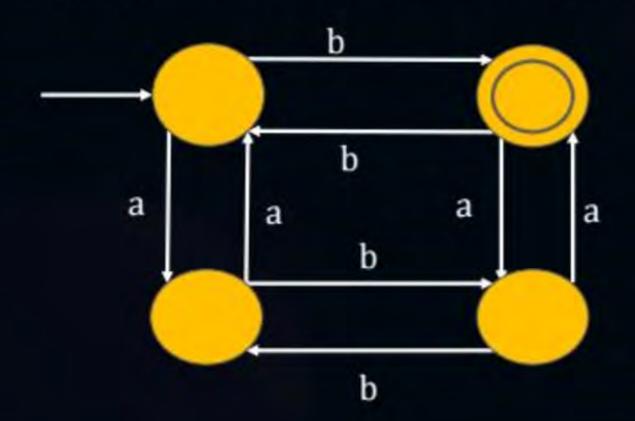




- **A**  $L = \{a^n b^m | n, m ≥ 1\}$
- C  $L = \{a^n b^m | n, m \ge 0 \}$

- **B**  $L = \{a^n b^m | n \ge 1, m \ge 0\}$
- D None





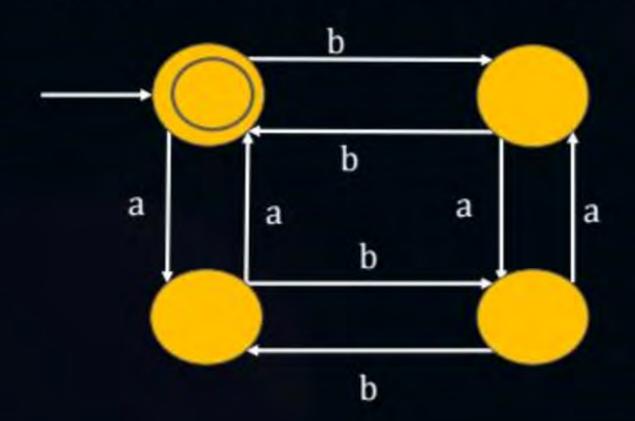
# a's even and # b's even

# a's odd and # b's odd

# a's odd and # b's even

# a's even and # b's odd





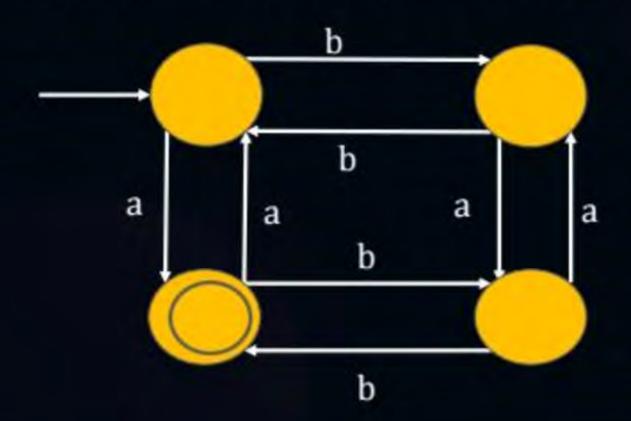
# a's even and # b's even

# a's odd and # b's odd

# a's odd and # b's even

# a's even and # b's odd





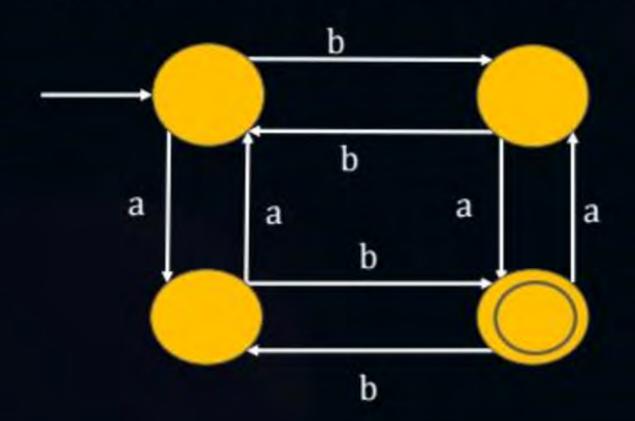
# a's even and # b's even

# a's even and # b's odd

# a's odd and # b's even

# a's odd and # b's odd





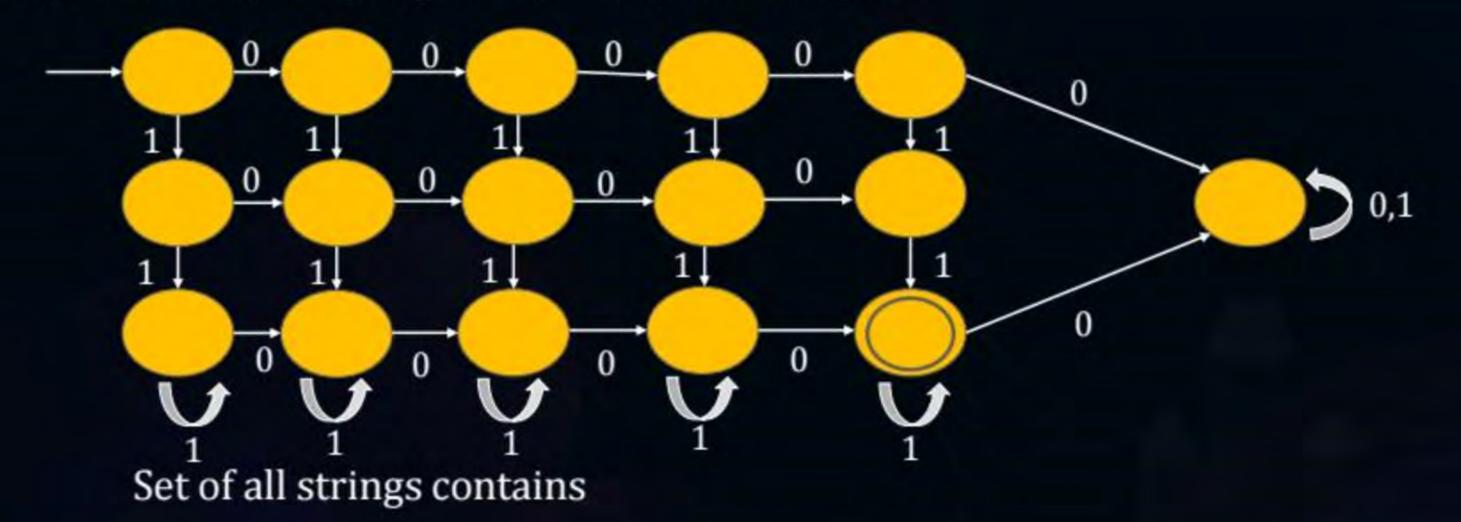
# a's even and # b's even

# a's odd and # b's odd

# a's odd and # b's even

# a's even and # b's odd





A Length of the string alteast 6

# 0's exactly4 and 1's atleast 3

# 0's atleast 4 and # 1's exactly 2

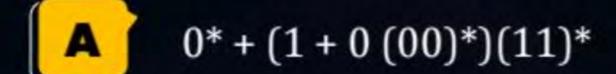
**D** None

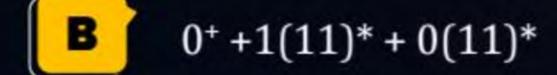


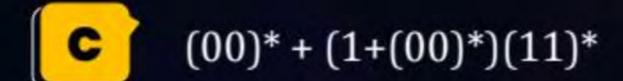
#Q. Let M be the 5-state NFA with ∈ - transitions shown in the diagram below. Which one of the following regular expression represents the language accepted by M?

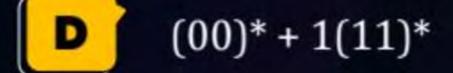
0

€



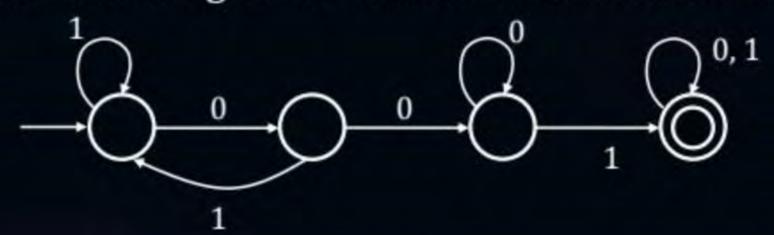








#Q. Consider the following deterministic finite state automaton M.



Let S denote the set of seven-bit binary strings in which the first, the fourth, and the last bits are 1. The number of strings in S that are accepted by M is

A 1

B

D



#### **Topic: DFA Construction**



#### Construct DFA for the following Language.

1. 
$$L = \{a^n b^m \mid n, m \ge 1\}$$

2. 
$$L = \{a^n b^n \mid n \ge 1\}$$

3. 
$$L = \{a^n b^m \mid n < m\}$$

4. 
$$L = \{a^n b^m \mid n \neq m\}$$

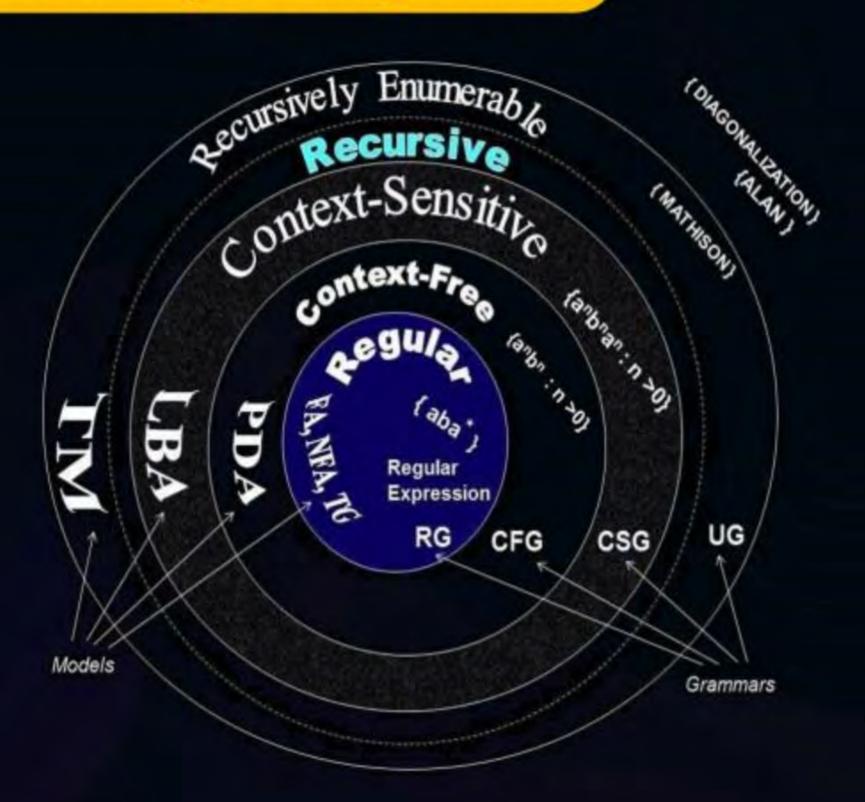
5. 
$$L = \{a^n b^m c^{n+m} | n, m \ge 1\}$$

6. 
$$L = \{a^n b^{2m} \mid n, m \ge 1\}$$



## **Topic: Theory of Computation**







#### **Topic: Expressive Power**



Number of languages accepted by particular automata is knowns as expressive power.

- Expressive power of NFA and DFA same. Hence every NFA is converted into DFA.
- 2. Expressive power of NPDA is more than DPDA. Hence conversion not possible
- Expressive power of DTM and NTM is same.



#Q. Let  $D_p$  are number of languages accepted by DFA and DPDA respectively. Let  $N_p$   $N_p$  are number of languages accepted NFA and NPDA respectively. Which of the following is true.

	$N_f = D_f$
A	$N_p = D_p$

$$\begin{array}{|c|c|} \hline & N_f \supset D_f \\ N_p \supset D_p \end{array}$$

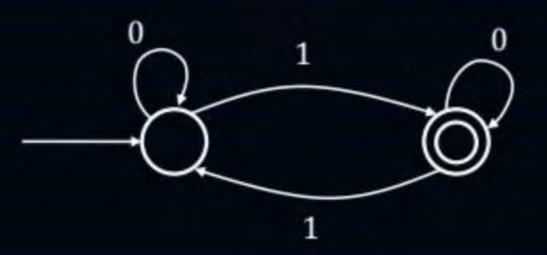


#Q. In which of the cases stated below the following statement is false? "Every nondeterministic machine M<sub>1</sub> there exists an equivalent deterministic machine M<sub>2</sub> recognizing the same language"

- A M<sub>1</sub> is non deterministic FA
- B M<sub>1</sub> is non deterministic turing machine
- M<sub>1</sub> Is non deterministic PDA
- D None



#Q. Which one of the following regular expressions is equivalent to the language accepted by the DFA given below?





## THANK - YOU