

# Computer Science & Information Technology

## Theory of Computation

DPP: 2

### Regular Expression and Closure Properties

- Q1** Let  
 $L_1 = a^* b^*$   
 $L_2 = b^* a^*$   
 $L_3 = (a + b)^*$   
 $L_4 = a^* b^* a$   
 $L = (L_1 \cap L_2) - (L_3 \cup L_4)$   
 Number of strings in above language L will be \_\_\_\_\_.
- Q2** Consider a regular expression (R):  
 $R = (a + b)^* (a + b)^2 (a + b)^*$   
 How many equivalences classes are existing for above regular expression R?  
 (A) 2 (B) 3  
 (C) 4 (D) None
- Q3** Let L be any formal language. If  $L^*$  is regular language then what is L?  
 (A) L is regular.  
 (B) L is non-regular.  
 (C) L is CFL.  
 (D) None of these.
- Q4** Consider the following two statements:  
**[I]:** There exist a regular language  $L_1$ , such that for all language  $L_2$ ,  $L_1 \cup L_2$  is always regular.  
**[II]:** If all states of deterministic finite automata (DFA) except start state are final states then language accepted by DFA is  $\Sigma^+$ .  
 Which of the following is correct?  
 (A)  $S_1$  only.  
 (B)  $S_2$  only.  
 (C) Both  $S_1$  and  $S_2$  are true.
- (D) None of these.
- Q5** Consider the language L given by the regular expression  $(a + b)^* ab(a + b)^*$  over the alphabet  $\{a, b\}$ . What is the correct regular expression of  $\overline{L}$ ?  
 (A)  $(a + b)^* (ab + ba + bb + aa) + \epsilon$   
 (B)  $(a^* b^*)^* (ba + bb + aa) (a^* b^*)^* + a + b$   
 (C)  $(a + b)^* ba (a + b)^* + a + b$   
 (D)  $b^* a^*$
- Q6** For language  $L = \{\text{Every odd bit is } a\}$  On alphabet  $\Sigma = \{a, b\}$ . Which of the following is/are correct regular expression?  
 (A)  $(aa + ab)^* (\epsilon + a)$   
 (B)  $(aa + ab + ba + b)^* a$   
 (C)  $(aa + ba)^* (\epsilon + a + b)$   
 (D)  $(a(a + b))^* + (a(a + b))^* a$
- Q7** Let us consider the following regular expression  $R = a^* b^* + b^* a^*$ .  
 How many equivalence classes of expression that represent language are equivalent to regular expression R?
- Q8** Consider the following languages:  
 $L_1 = \{a^m b^n c^p \mid m, n, p \geq 0\}$ .  
 $L_2 = \{a^m b^m c^p \mid m, p \geq 0\}$ .  
 $L_3 = \{a^{2m} b^{2m} c^p \mid m, p \geq 0\}$ .  
 Which of the following is/are correct?  
 (A)  $L_1 \subseteq L_2$  and  $L_2 \subseteq L_1$ .  
 (B)  $L_2 \subseteq L_1$  and  $L_3 \subseteq L_1$ .  
 (C)  $L_3 \subseteq L_2$  and  $L_2 \subseteq L_1$ .  
 (D)  $L_2 \subseteq L_3$  and  $L_3 \subseteq L_1$



## Answer Key

Q1 0

Q2 (B)

Q3 (D)

Q4 (A)

Q5 (D)

Q6 (A, D)

Q7 6

Q8 (B, C)



## Hints & Solutions

**Q1 Text Solution:**

$$L_1 \cap L_2 = a^* + b^*$$

$$L_3 \cup L_4 = (a + b)^*$$

$$L = (a^* + b^*) - (a + b)^*$$

$$= \phi$$

Number of strings = 0

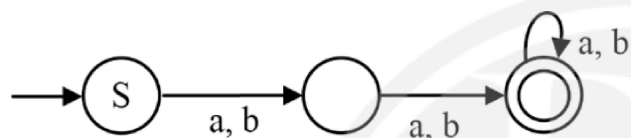
**Q2 Text Solution:**

$$R = (a + b)^* (a + b)^2 (a + b)^*$$

Number of equivalence classes in my hill Nerode

= Number of states in minimal DFA

DFA for R:



Number of states = 3

Number of equivalence classes = 3

**Q3 Text Solution:**

If  $L^*$  is regular,  $L$  may or may not be a regular.

**Example 1:**  $L^* = (a + b)^*$  is regular,  $L = (a + b)$  is regular.

**Example 2:**  $L^* = \{a^P \mid P \text{ is prime}\}$  is regular but  $L = \{a^P \mid P \text{ is prime}\}$  is non-regular.

$\therefore$  Option (d) is correct.

**Q4 Text Solution:**

**S<sub>1</sub> True:**

$$L_1 = \sum^*$$

$$L_1 \cup L_2 = \sum^* \cup L_2 = \sum^* \text{ (Regular)}$$

**S<sub>2</sub> False:**

May or may not be  $\sum^+$

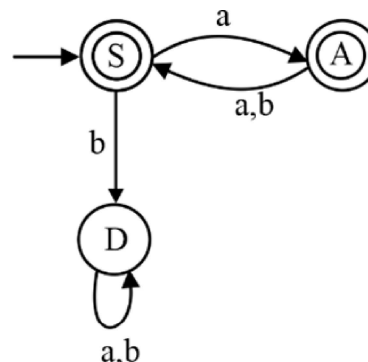
For example: DFA for language ending with "a" on alphabet {a, b}.

**Q5 Text Solution:**

$$L = (a + b)^* ab (a + b)^*$$

$$L = \{\text{containing 'ab' as a substring}\}$$

$$\bar{L} = \{b^* a^*\}$$

**Q6 Text Solution:**


F A:

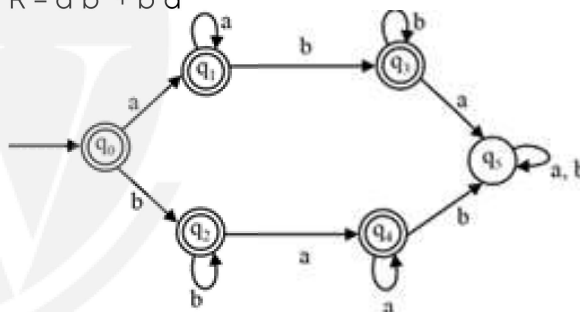
Regular expression =  $(a(a + b))^* a + (a(a + b))^* a$

$$= (a(a + b))^* (\epsilon + a) : (aa + ab)^* (\epsilon + a)$$

Hence, (a, d) are correct.

**Q7 Text Solution:**

$$R = a^* b^* + b^* a^*$$



$$R = a^* b^* + b^* a^*$$

$$= (\epsilon + aa^*) = (\epsilon + bb^*) + (\epsilon + bb^*) (\epsilon + aa^*)$$

$$= \epsilon + aa^* + bb^* + aa^* bb^* + bb^* aa^*$$

$$[\because a^* = (\epsilon + aa^*)]$$

Number of equivalence classes are equivalent to minimum number of states in DFA.

Regular expression for each state represents each equivalence class.

So,

$$[q_0] = \epsilon$$

$$[q_1] = aa^*$$



$$[q_2] = bb^*$$

$$[q_3] = aa^* + bb^*$$

$$[q_4] = bb^*aa^*$$

$$[q_5] = (aa^*bb^*a + bb^*aa^*b)(a + b)^*$$

**Q8 Text Solution:**

- $L_3 \subseteq L_1$  True
  - $L_2 \subseteq L_1$  True
  - $L_3 \subseteq L_2$  True
- (a) False                      (b) True  
(c) True                        (d) False



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