

Computer Science & IT

Data Structure & programming



Queue

Lecture No. 02

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Recap of Previous Lecture



Topic

Queue

Topic

Enqueue & dequeue

Topic

Topic

Topic

Topics to be Covered



Topic

Circular Queue

Topic

How to implement using stack

Topic

Topic

Topic



Topic : Question



#Q.

Compute the post fix equivalent of the following

expression $3 \times \log (x + 1) - a / 2$

$$= 3 * \log [x + 1] - a / 2$$

$$= 3 * [x + 1 \log] - a / 2$$

$$= 3 \times 1 + \log * - a / 2$$

$$= 3 \times 1 + \log * a / 2 -$$



Topic : Question



A function f defined on stacks of integers satisfies the following properties. $f(\phi) = 0$ and $f(\text{push}(S, i)) = \max(f(S), 0) + i$ for all stacks S and integers i . If a stack S contains the integers 2, -3, 2, -1, 2 in order from bottom to top, what is $f(S)$?

(A) 6

(B) 4

(C) 3

(D) 2

$$\begin{aligned} f(\underbrace{[2, -3, 2, -1, 2]}_1) &= f(\overset{\text{push}}{\overset{S}{[2, -3, 2, -1]}} \cdot 2) \\ &= \max(f(\overset{1}{S}), 0) + 2 \\ f([2, -3, 2, -1]) &= f(\overset{S}{\text{push}}([2, -3, 2], -1)) \\ &= \max(f(\overset{2}{[2, -3, 2]}, 0) - 1 \\ f([2, -3, 2]) &= f(\overset{S}{\text{push}}([2, -3], 2)) \\ &= \max(f(\overset{-1}{[2, -3]}, 0) + 2 \\ f([2, -3]) &= f(\overset{-1}{\text{push}}([2], -3)) = \max(f(\overset{2}{[2]}, 0) - 3 \end{aligned}$$

$$f([2]) = f(\text{push}(\square, 2))$$

$$= \max(f(\emptyset), 0) + 2$$

$$0 + 2 = \textcircled{2}$$



Topic : Circular Queue

```
void Enqueue(int data){  
    if((rear == Max-1 && front == 0) || (front == rear+1)){  
        printf("queue is full");  
        return;  
    }  
    rear = (rear+1) % Max;  
    a[rear] = data;  
    if(front == -1)  
        front = 0;  
}
```

$\text{front} = 2$
 $\text{rear} = 1$

$\text{front} := (\text{rear} + 1) \% \text{Max}$

$\text{front} = 0$ $\text{rear} = 4$

Single Condition



Topic : Circular Queue

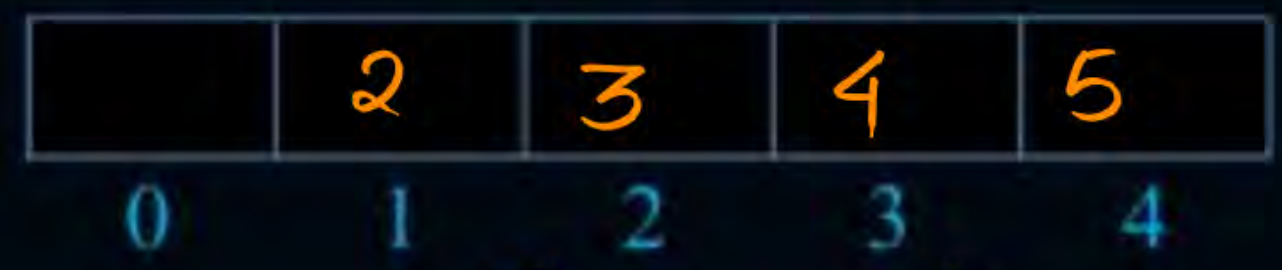
```
int dequeue() {  
    int data;  
    if (front == -1) { // if (IsEmpty())  
        printf("Queue is empty");  
        return -1;  
    }  
    data = a[front];  
    if (rear == front)  
        rear = front = -1;  
    else
```

```
        front = (front + 1) % Max;  
    return data;  
}
```

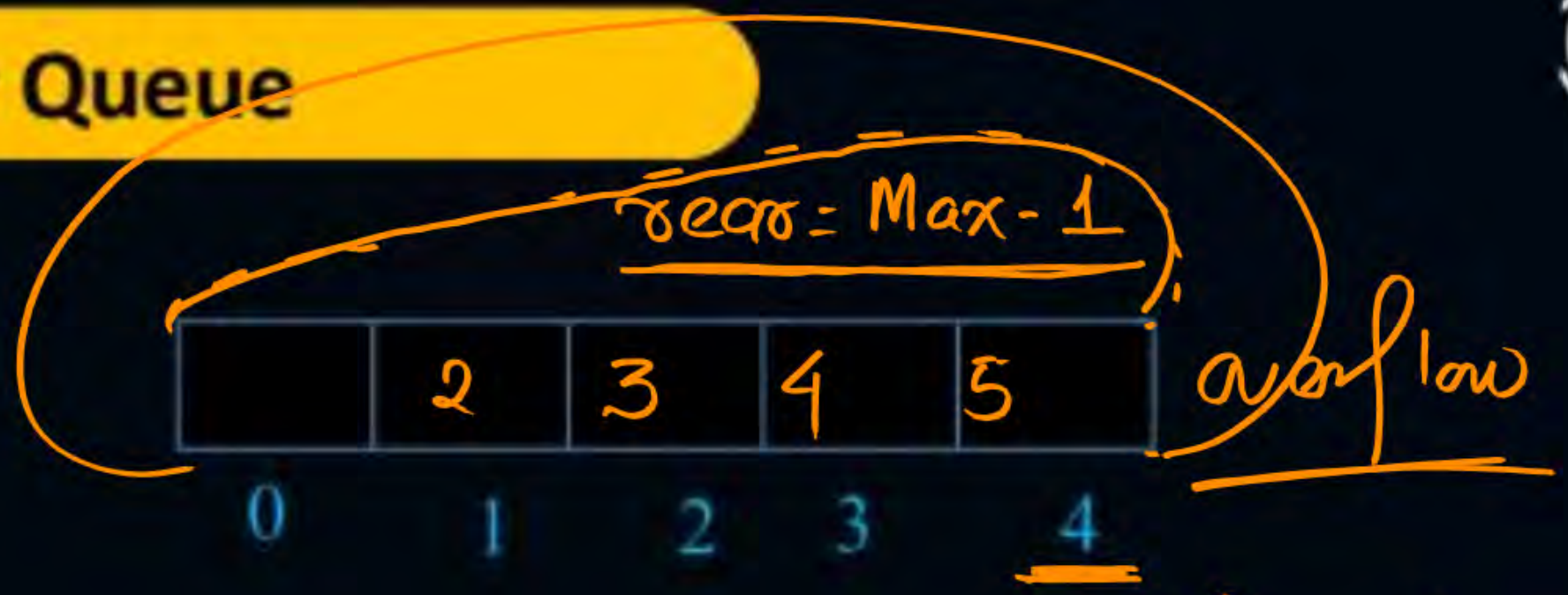



Topic: Circular Queue

front = 1, rear = 4



Enqueue(6)



Remainder



$$\begin{aligned} \text{rear} &= (\text{rear} + 1) \% \text{Max} \\ &= (4 + 1) \% 5 \\ &= 5 \% 5 = 0 \end{aligned}$$



Topic: Circular Queue

$\text{front} = 0, \text{rear} = 4$

1	2	3	4	5
0	1	2	3	4

$\text{front} = 1$
 $\text{rear} = 4$

Enqueue (6)

	2	3	4	5
0	1	2	3	4

Enqueue (6)

overflow $\text{rear} = \text{Max} - 1$
 $\text{front} = 0$ ✓

$\text{front} = 1$
 $\text{rear} = 0$

0	1	2	3	4

↑ first full condition

6	2	3	4	5
0	1	2	3	4

↑ Is queue is full?
① $\text{rear} < \text{front}$
② $\text{front} = \text{rear} + 1$ ✓



Topic: Circular Queue

rear = 0
front = 1

6	2	3	4	5
0	1	2	3	4

rear = 0
front = 2

6	2	3	4	5
0	1	2	3	4

rear = 0
front = 2
Dequeue()

6		3	4	5
0	1	2	3	4

Enqueue(7)

rear = 1
front = 2

6	7	3	4	5
0	1	2	3	4

$rear = (rear + 1) \% Max$



Topic: Circular Queue

overflow

front = 2
rear = 1

front = 2, rear = 1

6	7	3	4	5
0	1	2	3	4

6	7	3	4	5
0	1	2	3	4

Enqueue(8)

←
front = rear + 1

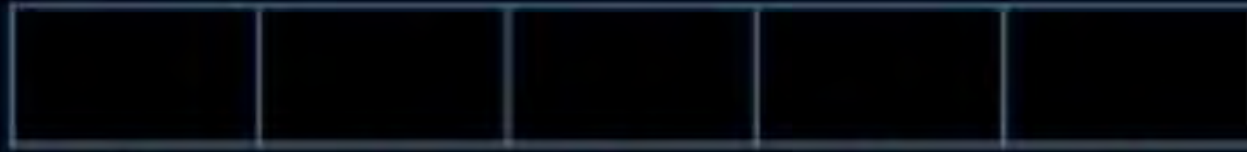
rear < front

0	1	2	3	4

0	1	2	3	4

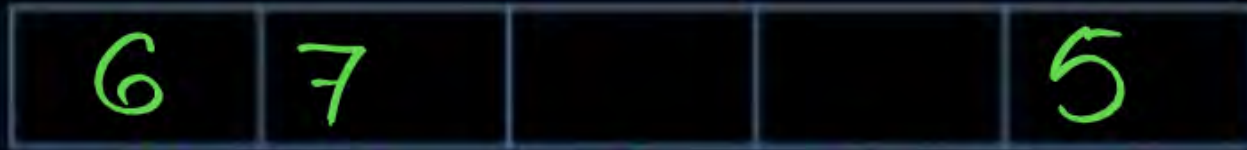


Topic: Circular Queue



0 1 2 3 4

front = 4, rear = 1



0 1 2 3 4

Enqueue(8)



0 1 2 3 4

front = 4, rear = 2



0 1 2 3 4

Condition = $\text{front} == (\text{rear} + 1) \% \text{Max}$
(4 == 2)



Topic: Circular Queue

front = 4, rear = 2

6	7	8		5
0	1	2	3	4

dequeue()

0	1	2	3	4

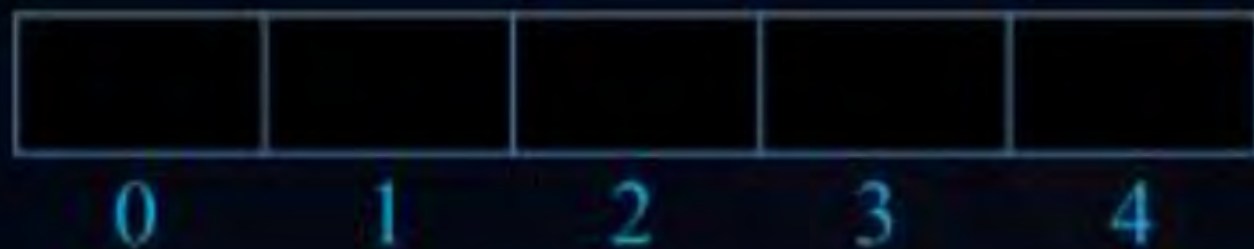
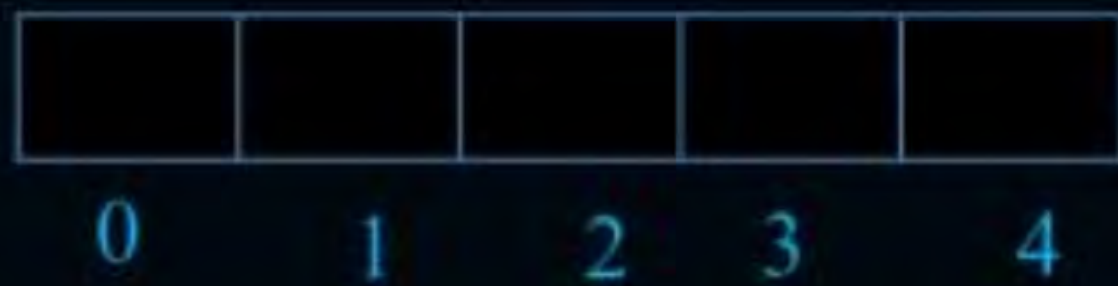
front = 0, rear = 2

6	7	8		5
0	1	2	3	4

0	1	2	3	4



Topic: Circular Queue





Topic : Circular Queue

if following operation are performed on queue of size 5.

Enqueue(11), Enqueue(12), Enqueue(13)

dequeue(), dequeue(), Enqueue(16), Enqueue(20)

Enqueue(29) Enqueue(70), $x = \text{dequeue}()$

if y is Index of front & z is Index of rear then

$x + y - z$ is 15

11), Enqueue(12), Enqueue(13)

), dequeue(), Enqueue(16), Enqueue(20)

(29) Enqueue(70), x = dequeue()

Index of front \geq is Index of rear then

15

rear = 2

11	12	13	
0	1	2	3

29	70	13	16
0	1	2	3

front = ~~2~~ 3

rear = 1

X

13-



Implementation of Queue Using Stack

- * Single stack can't implement using push, pop
- * Auxiliary storage, another stack is used



Implementation of Queue Using Stack



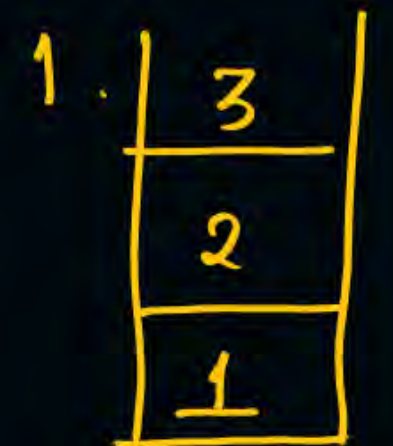
1. Suppose S_1, S_2 are two stacks

Algorithm — Enqueue = push(S_1)

dequeue() = 1 Empty S_1 by popping element from S_1
and push S_2

2 pop element from S_2 and return it

3. Restore Content of S_1 by popping element
from S_2 & pushing it in S_1 ✓



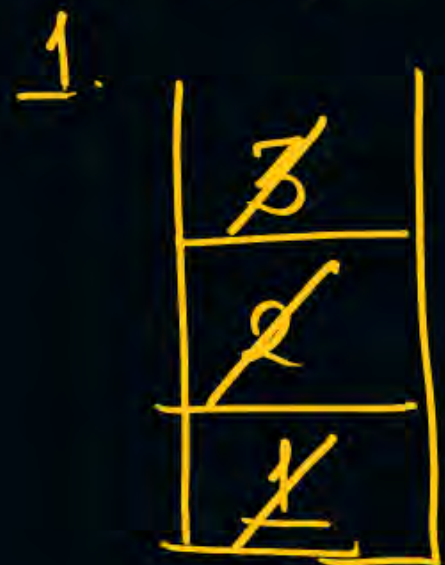
Enqueue(1) = push(1)

Enqueue(2) = push(2)

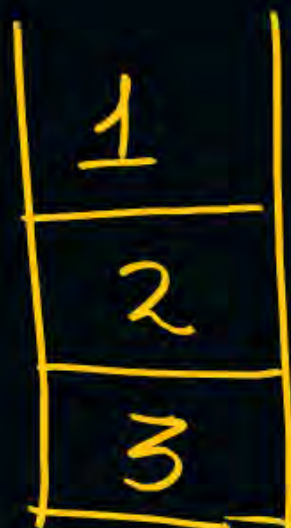
Enqueue(3) = push(3)

S₁

dequeue()



S₁

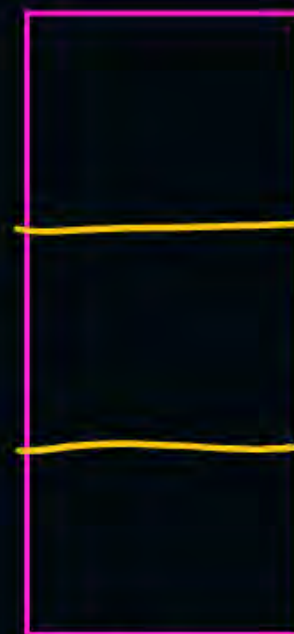


S₂

3.



S₁



S₂



S₁



S₂

return 1

THANK - YOU