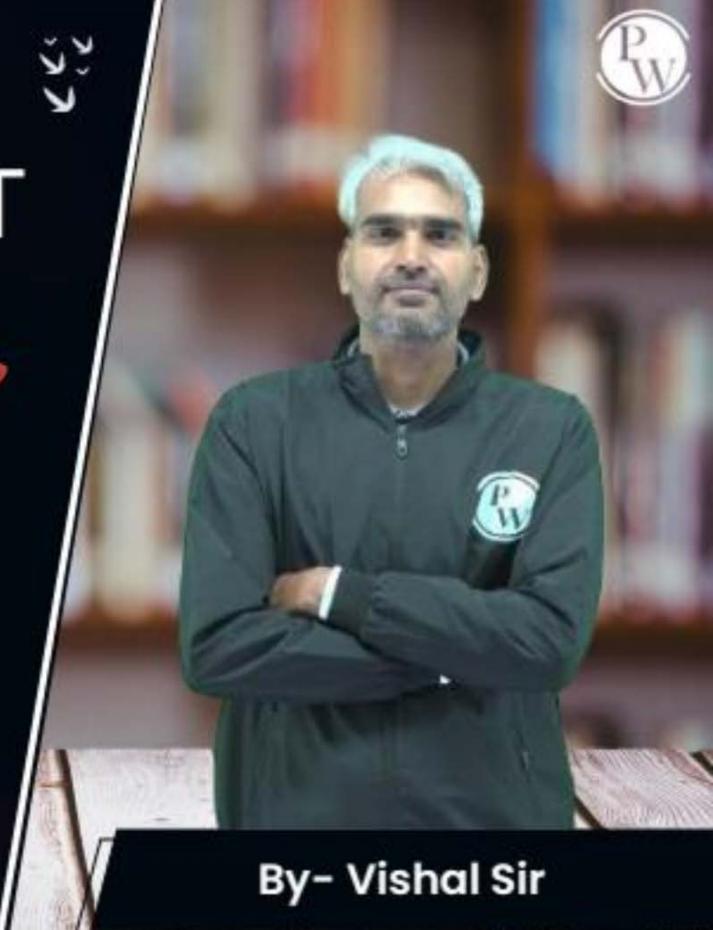
Computer Science & IT

Discrete Mathematics

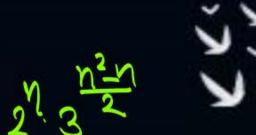
Set Theory & Algebra

Lecture No. 07













Symmetric, anti-symmetric and asymmetric relation



Topic

Topic

Transitive relation

if (aRb4 bRc) then aRc, 40,6,ceA

Topic

Complement of a relation

Topic

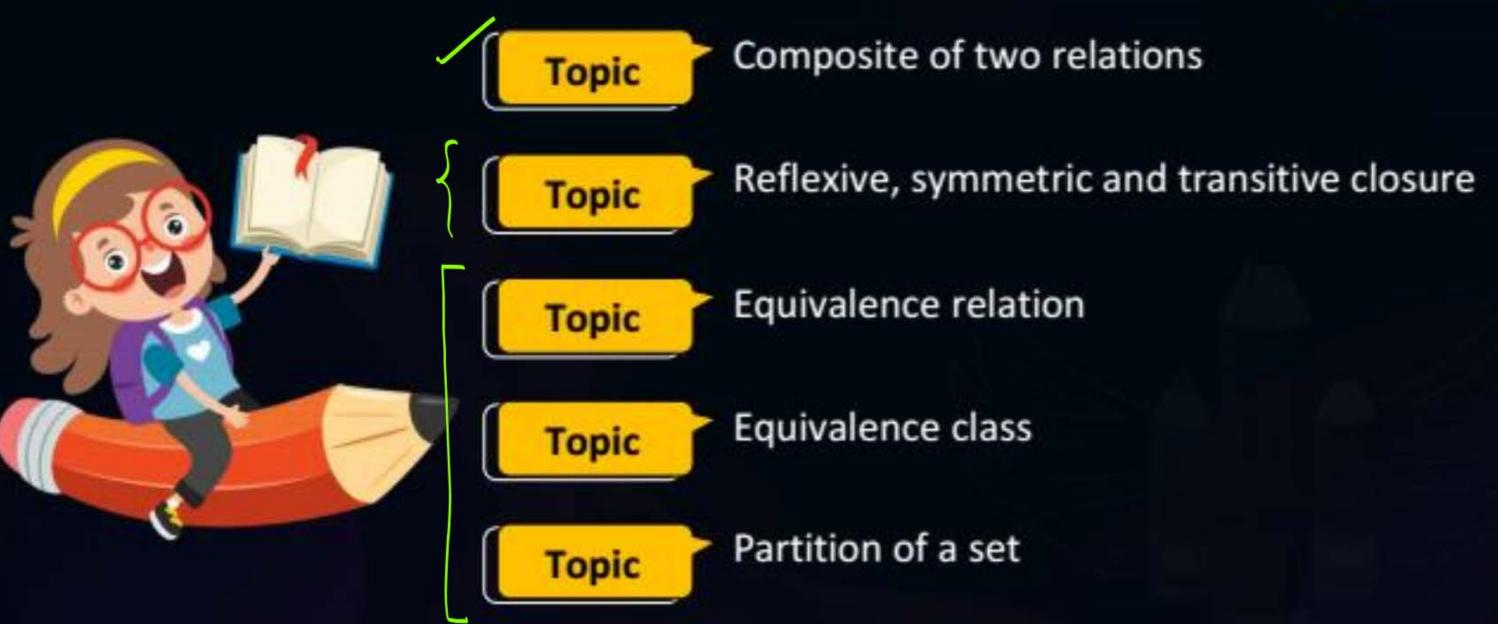
Inverse of a relation

Topics to be Covered







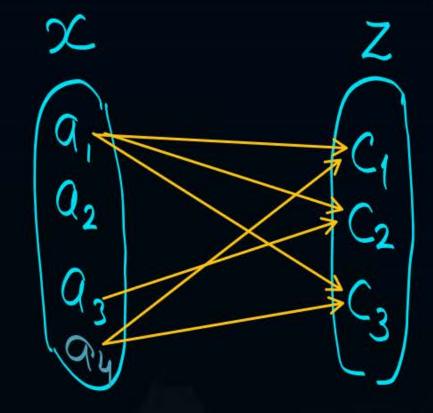




Topic: Composite of two relations







R is relation ton X to Y
i.e. R: X-y

S is a relation from y to Z

is S: Y > Z

Composite Relation
R;s: X→Z

 $R = \{(a_1, b_2), (a_1, b_4), (a_2, b_3), (a_3, b_5), (a_4, b_2)\}$ $\{S = \{(b_1, c_1), (b_1, c_3), (b_2, c_1), (b_2, c_3), (b_4, c_2), (b_5, c_2)\}$ $Identify \quad RjS = ?$



Topic: Composite of two Relations

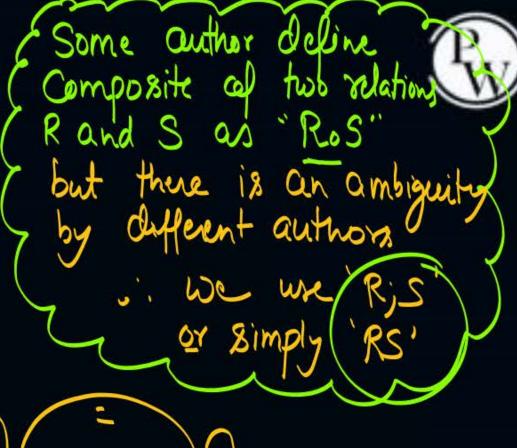
Let X,Y and Z be three sets, and

If R is a relation from X to Y, and S is a relation from Y to Z, then their composition R;S is the relation defined as,

R;S = {
$$(x,z) \mid (x,z) \in X \times Z$$
 and there exists $y \in Y$ such that $(x,y) \in R$ and $(y,z) \in S$ }

i.e., R;S is a relation from X to Z defined by the rule that $(x,z) \in R$;S if and only if there is an element

 $y \in Y$ such that $(x,y) \in R$ and $(y,z) \in S$





Topic: Reflexive Closure



Note: Replexive closure of R will be = R U DA

* If relation R is a reflexive relation on set A,
then reflexive closure of R will be relation R' itself



Topic: Symmetric Closure



Let R be a relation on set A, Symmetric closure of R is the 2 mallest symmetric Ruh
on set A that Contains R. eg: let A= {1,2,3} $R = \{(1,1), (1,2), (2,1), (3,1)\}$ Symmetric al R=S(1,1), (1,2), (2,1), (3,1), (1,3) Contain R Contain R

Slide

* Mote: - * Symmetric Closure of R = RUR

+ If relation R is a symmetric relation, then

Symmetric closure of R will be relation R itself.



Topic: Transitive Closure



Let R be a relation on set A, Transitive Closure af R is the smallest transitive relation On set A that Contains R.

Let A = {1,2,3} QH. because af (3,1) of (1,2) $R = \{(1,1), (1,2), (1,3), (2,2), (3,1)\}$ Find transitive closure of R. 1 iteration $R^* = \{ (1,1), (1,2), (1,3), (2,2), (3,1) \}$ all Paim of R* becourse $2^{\text{redifferentian}} = \{(1,1), (1,2), (1,3), (2,2), (3,1), (3,2), (3,3)\}$ (3,1) 4 (1,3)R*=R**, i. R* is transitive closure of R

If relation R is a transitive relation, then transitive Closure cel relation R Will be relation R itself 6 Reflexive transitive clusure: 7 Reflexive transitive clusure of relation R will be the relation which is Smallest relation on the set such that it is reflexive as well as transitive and contains R

let A: {1,2,3} Ø. $R = \{(1,1), (1,3), (2,2), (3,1), (3,2)\}$ find toansitive Closure of R $R^* = \{ (1,1), (1,3), (2,2), (3,1), (3,2), (1,2), (3,3) \}$ R**= {(1,1), (1,2), (1,3), (2,2), (3,1), (3,2), (3,3); \\Nothing new is added R*= R*X is R* is our toanitive Closure

A = {a,b,c,d} R=5(a.d), (b,a), (b,c), (C,a), (C,d), (d,c)}
Find transitive closure of R. $R'' = \{(a,d), (b,a), (b,c), (c,a), (c,d), (d,c), (q,c), (b,d), (c,c), (d,a), (d,d)\}$ Rearranges = {(a,c),(a,d),(b,a),(b,c),(b,d),(c,a),(c,c),(c,d),(d,a),(d,c),(d,d)} $R^{**} = \{(0,c), (a,d), (b,a), (b,c), (b,d), (c,a), (c,c), (c,d), (d,a), (d,c), (d,d), ((a,a), (c,c), (c,d), (d,d), (d,d), ((a,a), (c,c), (c,d), (d,d), (d$ Reamonged = {(a,a) (0,c) (a,d) (b,a) (b,c) (b,d) (c,a) (c,c) (c,d) (d,a) (d,c) (d,d) } R** = Sam as R**, i.e. R**= R** is toomitive closure of R Note: - Warshall's algorithm can be used to identify the transitive Closuse of a given relation



Topic: Equivalence relation



A relation R on set A is said to be an Equivalence relation if and only if relation is

P D Rellesive
P D symmetric
A 3 Transitive

 $\Delta_A = R_1 = \{(1,1), (2,2), (3,3)\}$ Reflexive \rightarrow in Equivalence Relative \rightarrow in Equivalence Relative \rightarrow In the smallest equivalence $AxA = R_2 = \{(1,1),(1,2),(1,3),(2,1),(2,2),(2,3),$ relation on set A (3,1),(3,2),(3,3)AXA is the Ax A is the Reflexive Symmetric Largest equivalence Symmetric Largestation on set A Transitive L R3 = {(1,1),(2,2),(3,3), (1,2),(2,1)} Replexive \ => 00 Equivalence Relh transitive 3

(3) $R_3 = \{(1,1),(2,2),(3,3),(1,2),(2,1)\}$ symmetric V = 0 or Equivalence Relative V = 0 or $V = \{(1,1),(2,2),(3,3),(3,3),(3,1)\}$ is Equivalence Relative V = 0 or V = 0

 $R_6 = \left\{ (1,1), (2,2), (3,3), (1,2), (2,1), (1,3), (3,1) \right\}$ Rellesive V Symmetric but not transitive

but (23) € R6 i. Not transitive to make it transitive if we add (2,3) than for symmetricity we will have add (3,2) as well Ord it will become AXA i. No other equivalence rel' Possible on 8et A={1,2,3}



Topic: Equivalence Class



Let R' be an equivalence relation on set A, for any element $x \in A$ the equivalence class a element 'x' w.r.t. Equivalence relation R can be denoted by [x], and it is defined as. <u>ie</u> Equivalence class a

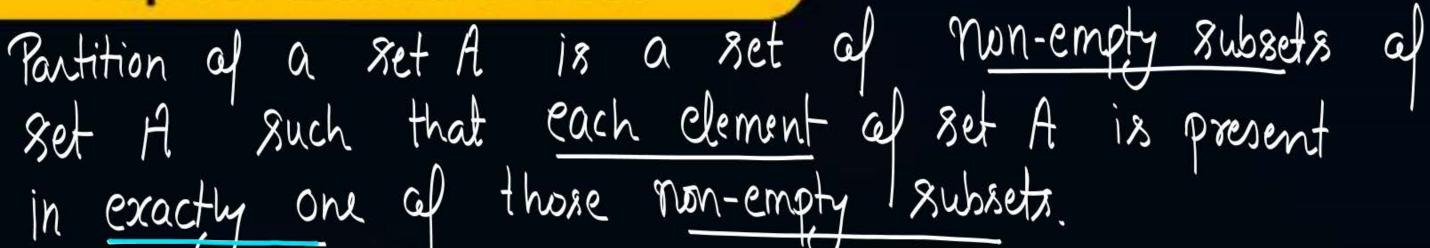
ie Equivalence class of clement'x' is a set of all elements which are related with x

1, 2, 3, 4, 57 let R is an Equivalence relation on set A. and $R = \{ (1,1), (2,2), (3,3), (4,4), (5,5), (1,3), (3,1), (4,5), (5,4) \}$ oflexive as well as transitive define Equivalence a every element at set A. wisit. equivalence Rel 'R' []=({1,3} Same There are three distinct equivalence Classes for elements of set A [2] = {2} je. {1,3} 4 {2} & {4,5} [3] = {3,1} [4] = ({4,5} Same

Equivalence class of element of ile [27] may be same as equivalence class af element y fix. [x] y even if $x \neq y$. $\begin{cases} \text{in the above eg. } \end{cases}$ pm [1] =[3] ={1,3} 2) The set of all distinct equivalence classes of elements a) set A w.r.t. on Equivalence relation R Creater a partition of set A. set al suitinct equivalence classes al elements above example is a partition set al set A={1,2,3,4,5}



Topic: Partition of a set



 (Ωl) let A be a non-empty set, and As, Az, Az, Az, -. - Ak non-empty subsets a set A, then $\{A_1, A_2, A_3, \dots, A_k\}$ is a partition of set Aif and only if $(A_i \cap A_j = \emptyset)$, $\forall i \neq j$ $\{A_1, A_2, A_3, \dots, A_k\}$ is a partition of set A $\{A_i \cap A_j = \emptyset\}$, $\{A_i \cap A_j = \emptyset\}$, $\{A_i \cap A_j = \emptyset\}$, $\{A_i \cap A_j = \emptyset\}$

Let $A = \{1, 2, 3, 4, 5\}$ which of the following is/one partitions of set A. 13,45 n ≤4,57 = 4 ≠ Ø (i. Not a partition) Q { {1,2}, {3,4}, {4,5}} 6) {{1,27, {35, {4,5}, {4},5}, {4}} are not allowed {i. Not a Partition} (E) \fightless \fightl

find the ho. Partitions of (1) When |A| = D(2) When |A| = 1 (3) When |A| = 2 (4) when |A|=3 (3) When |A| = 4 When 1A1 - 5 When AI =



2 mins Summary



Topic Composite of two relations

Topic Reflexive, symmetric and transitive closure

Topic Equivalence relation

Topic Equivalence class

Topic Partition of a set



THANK - YOU