

CS & IT ENGINEERING



Operating System

Deadlock

Lecture -1

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Recap of Previous Lecture



Topic

Semaphore

Topic

Questions on Semaphore

Topic

Classical Problems on Synchronization

Topic

Bounded Buffer Problem

Topics to be Covered



Topic

Reader-Writer Problem

Topic

Dining Philosopher Problem

Topic

Deadlock

Reader-writer problem:-

A file is shared b/w multiple reader processes and multiple writer processes.



Topic : Reader-Writer Problem Solution

- If writer is accessing the file, then all other readers and writers will be blocked
- If any reader is reading, then other readers can read but writer will be blocked

currently	allowed	
	reader	writer
Reader	✓	✗
writer	✗	✗



Topic : Reader-Writer Problem Solution

- **Variables:**
 - mutex: Binary Semaphore to provide Mutual Exclusion
 - wrt: Binary Semaphore to restrict readers and writers if writing is going on
 - Readcount: Integer variable, denotes number of active readers
- **Initialization:**
 - mutex: 1
 - w r t : 1
 - Read count: 0



Topic : Writer() Process



wait(wrt)

// performs writing

signal(wrt)



Topic : Reader() Process

```
wait(mutex)
readcount ++;
if (readcount == 1)
{ wait(wrt);
}
signal(mutex)
// Performs Reading
wait(mutex)
readcount --;
if (readcount == 0)
{ signal(wrt); }
signal(mutex)
```

Case 1 :-

A writer^① comes & writes

A reader^① comes \Rightarrow stuck at wait(wrt)

A reader^② comes \Rightarrow stuck at wait(mutex)

A writer^② comes

\hookrightarrow stuck at wait(wrt)

wrt = 1 0

mutex = 1 0

readcount = ~~0~~ 1

Case 2:-

A reader^① comes & reads

A writer comes \Rightarrow stuck at wait(wrt)

A reader^② comes \Rightarrow Reads

A reader^③ comes \Rightarrow Reads

Reader ② Exits ✓

Reader ① Exits ✓

Reader ③ Exits ✓

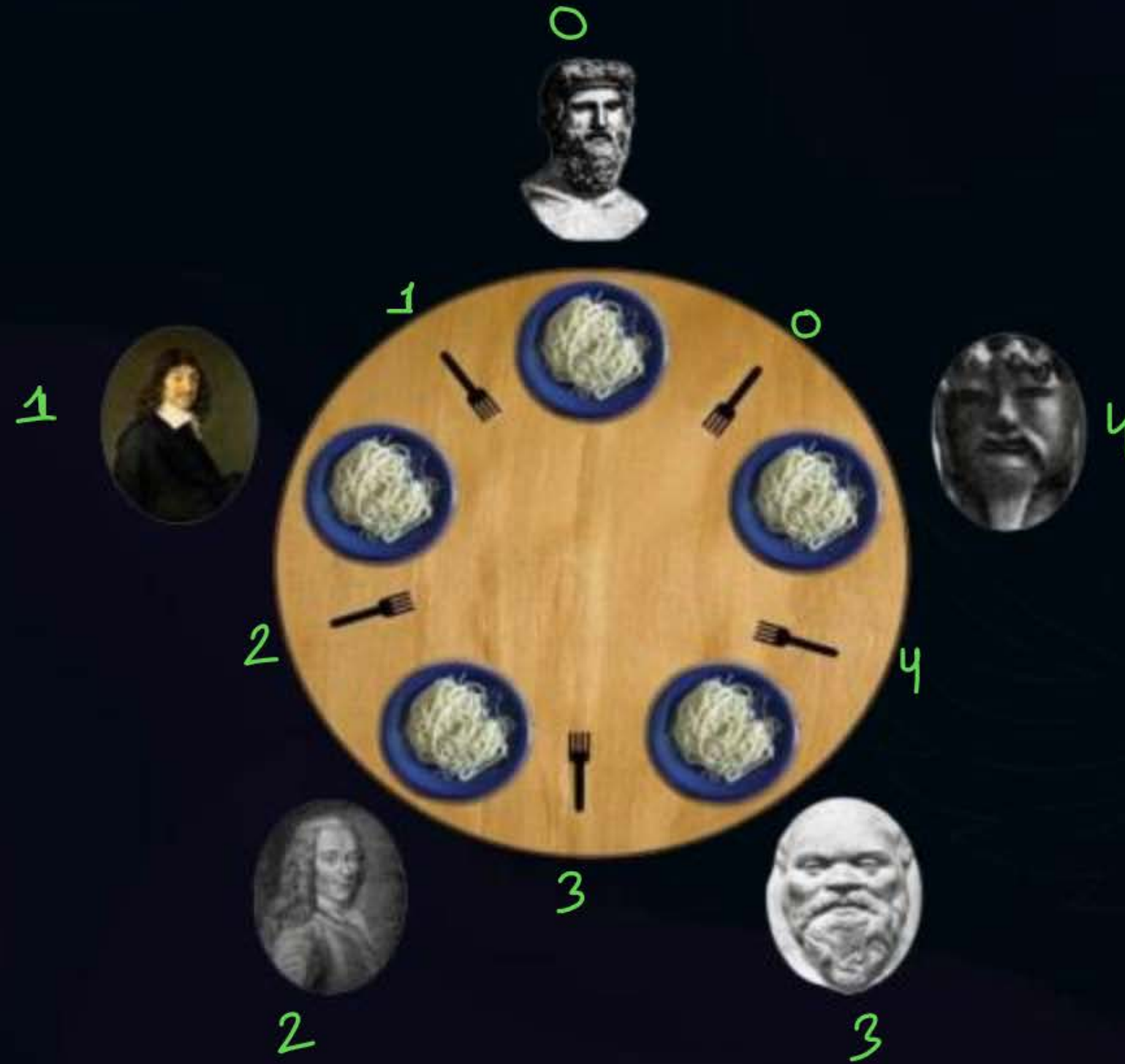
wrt = ~~1~~ 0 1

mutex = ~~1~~ 0 1 0 1 0 1 0 1 0

readcount = ~~0~~ 1 2 3 2 1 0



Topic : Dining Philosopher Problem





Topic : Dining Philosopher Problem

- K philosophers seated around a circular table
- There is one chopstick between each philosopher
- A philosopher may eat if he can pick up the two chopsticks adjacent to him
- One chopstick may be picked up by any one of its adjacent followers but not both



Topic : Dining Philosopher Problem Solution





Topic : Dining Philosopher Problem Solution

for k number of chopsticks, consider an array of binary semaphores of size k .

$$\text{chopstick}[k] = \{1, 1, 1, \dots, 1\};$$

each philosopher runs a process:-

$$p_0, p_1, \dots, p_{k-1}$$



Topic : Dining Philosopher Problem Solution

P_i

```
wait(chopstick[i])  
wait(chopstick[(i+1)%k])  
  
//eat  
signal(chopstick[i])  
signal(chopstick[(i+1)%k])
```

It can suffer from
deadlock.



if one-by-one all philosophers
pick one chopstick.



Topic : Dining Philosopher Problem Solution

Some of the ways to avoid deadlock are as follows –

1. There should be at most $(k-1)$ philosophers on the table



Topic : Dining Philosopher Problem Solution

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2. A philosopher should only be allowed to pick their chopstick if both are available at the same time



Topic : Dining Philosopher Problem Solution

Some of the ways to avoid deadlock are as follows –

1. There should be at most $(k-1)$ philosophers on the table
2. A philosopher should only be allowed to pick their chopstick if both are available at the same time
3. One philosopher should pick the left chopstick first and then right chopstick next; while all others will pick the right one first then left one

↳ asymmetry in picking up chopstick



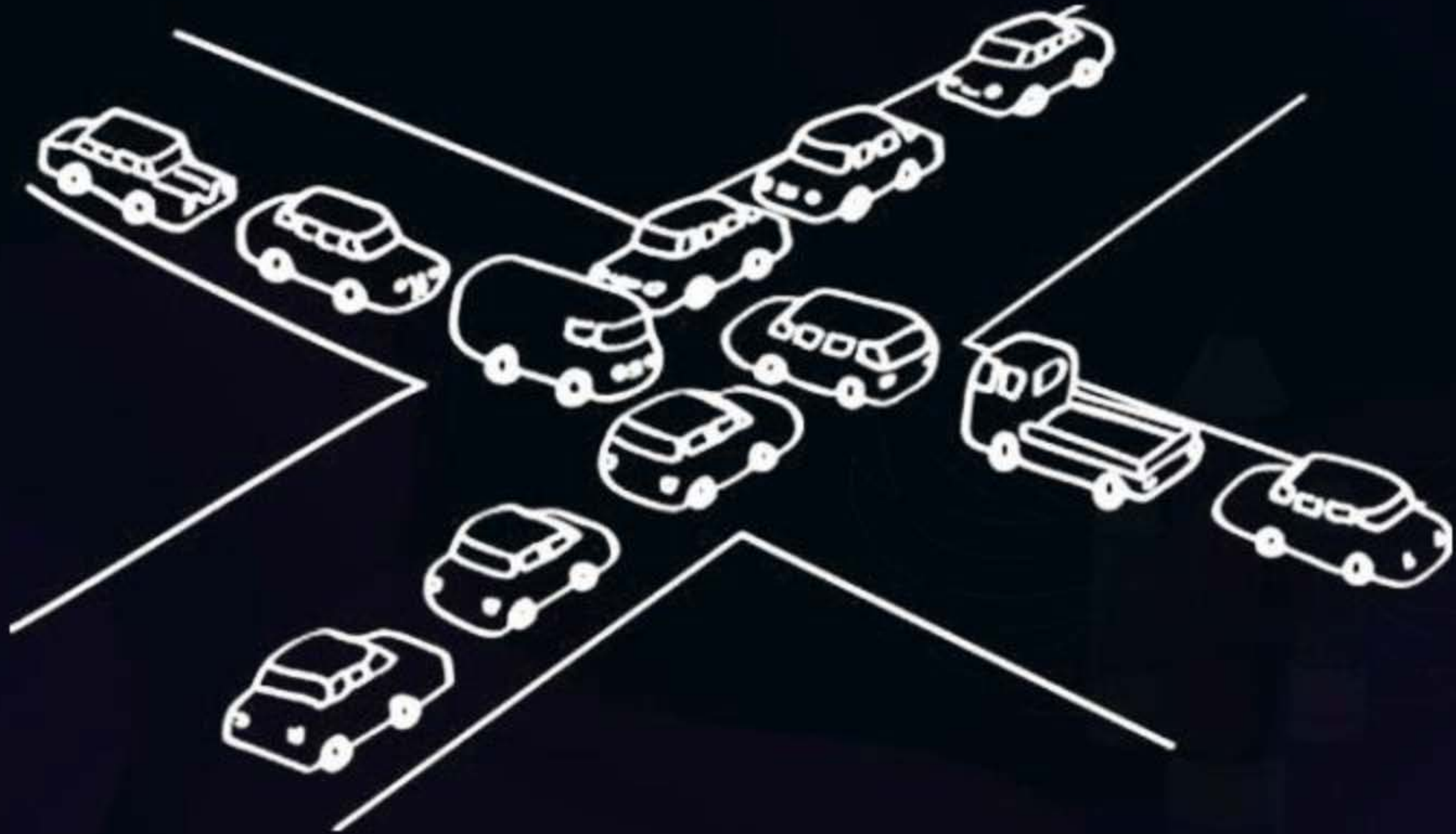
Topic : Deadlock



If two or more processes are waiting for such an event which is never going to occur



Topic : Deadlock





Topic : Operations on Resources

3 Operations on resources:

1. Request :- A process can request to OS to allocate a resource to it.
2. Use
3. Release
4. wait :- If resource is not available then process will wait.



2 mins Summary

Topic

Reader-Writer Problem

Topic

Dining Philosopher Problem

Topic

Deadlock



Happy Learning

THANK - YOU