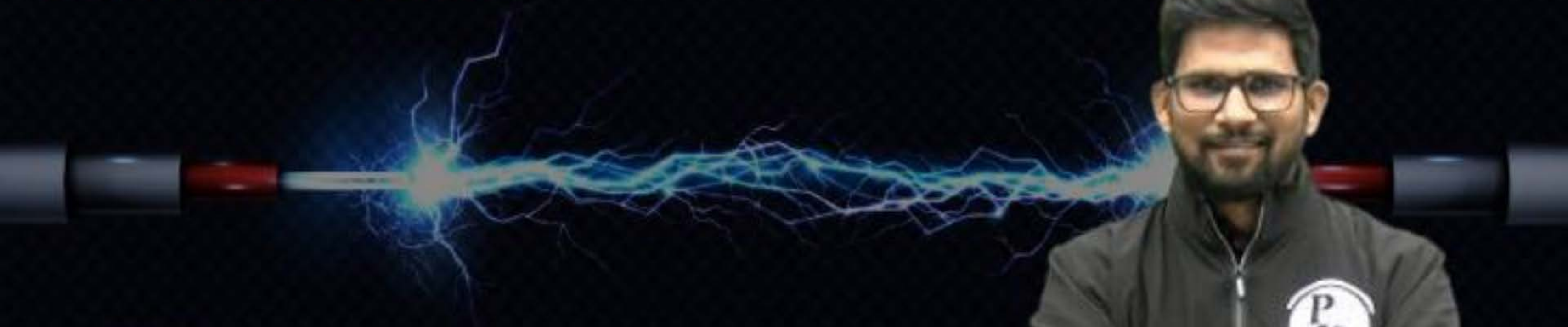


COMPUTER SCIENCE & IT

DIGITAL LOGIC




Lecture No. 10

Combinational Circuit



By- Chandan Gupta Sir



Recap of Previous Lecture

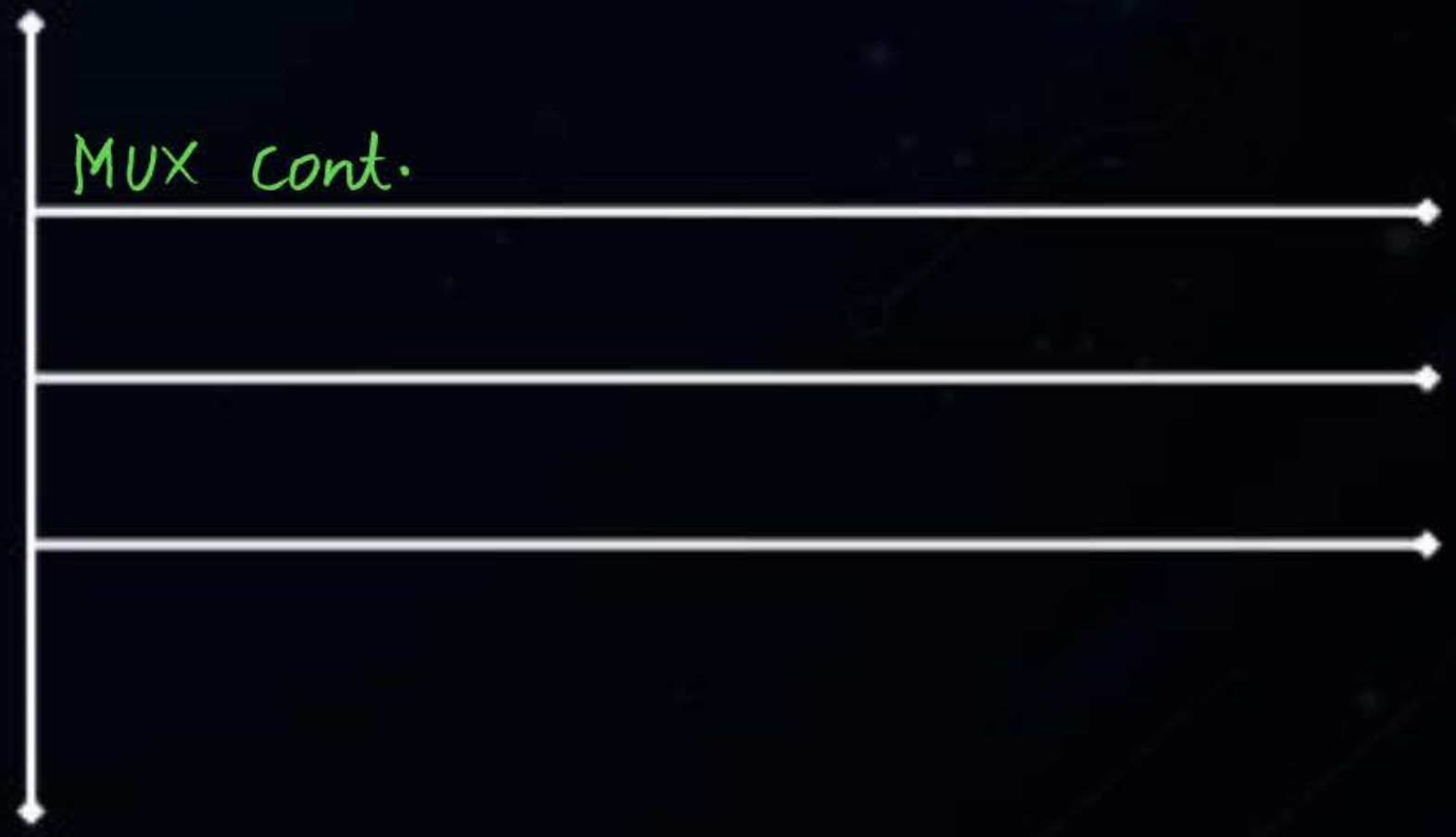
MUX

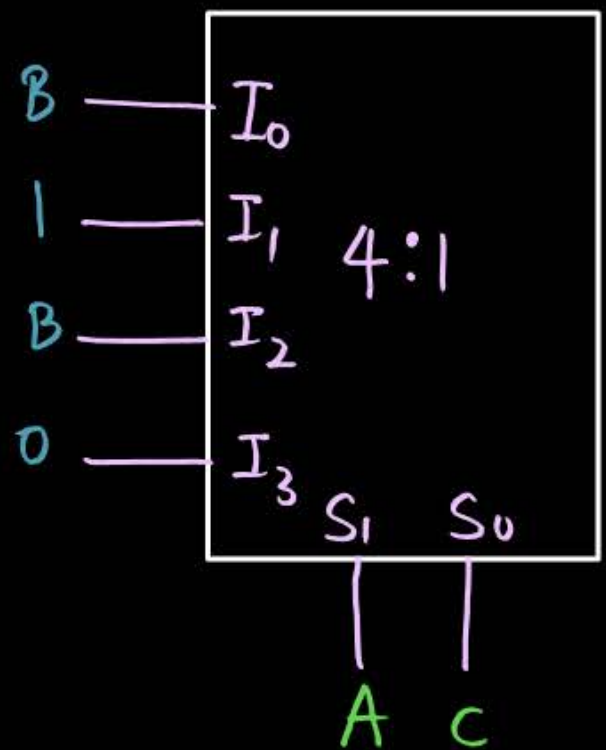




Topics to be Covered

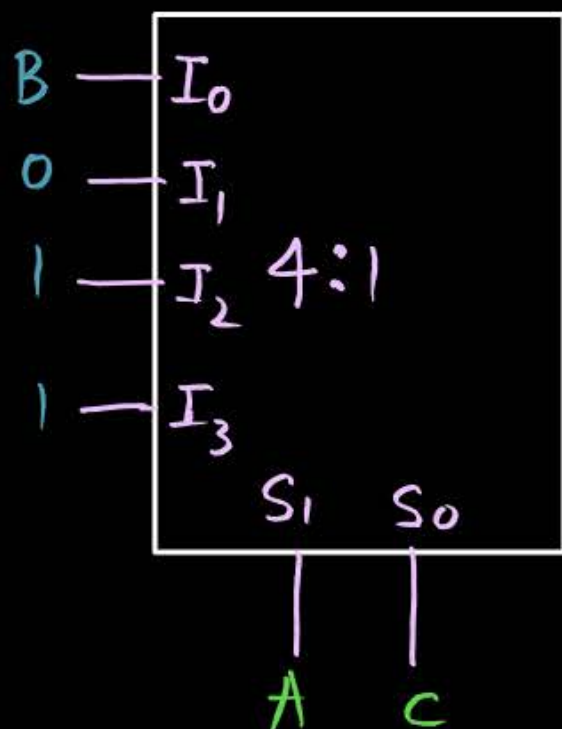
MUX cont.





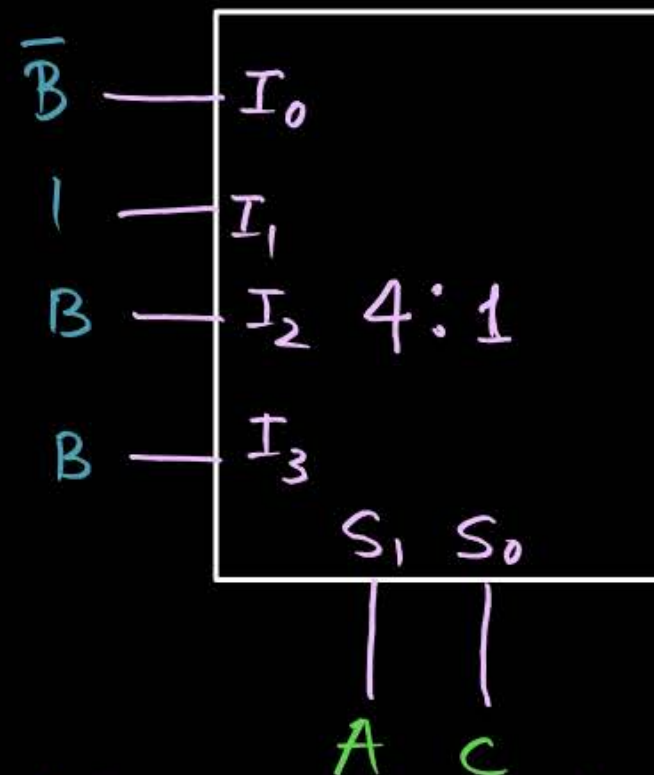
$$f_1(A, B, C) = \Sigma(1, 2, 3, 6)$$

	I_0	I_1	I_2	I_3
\overline{B}	0	1	4	5
B	2	3	6	7
	B	1	B	0



$$f_2(A, B, C) = \Sigma(2, 4, 5, 6, 7)$$

	I_0	I_1	I_2	I_3
\overline{B}	0	1	4	5
B	2	3	6	7
	B	0	1	1



$$f_3(A, B, C) = \Sigma(0, 1, 3, 6, 7)$$

	I_0	I_1	I_2	I_3
\overline{B}	0	1	4	5
B	2	3	6	7
	\overline{B}	1	B	B



- 4 : 1 MUX can be used to implement

$2^2 : 1$

All 2-variable function

Some of the
3-variable function

- 4 : 1 MUX + 1 NOT GATE can be used to implement

└→ All 3-variable function

- 8 : 1 MUX can be used to implement

$2^3:1$

All-3-Variable function

Some of the
4-Variable function

- 8 : 1 MUX + 1 NOT GATE can be used to implement

└→ All-4-variable function

- $2^n : 1$ MUX can be used to implement

All n -variable function.

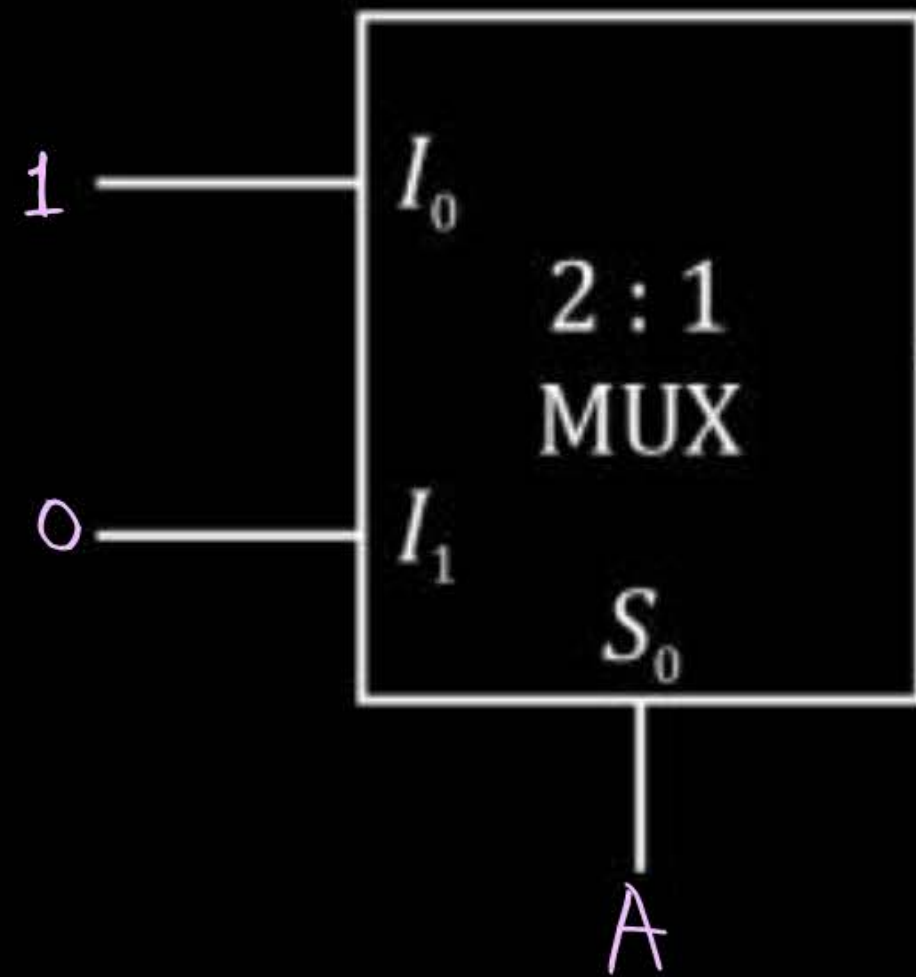
Some of the $(n+1)$ variable function

- $2^n : 1$ MUX + 1 NOT GATE can be used to implement

└→ All $(n+1)$ variable function.

[MUX as Universal Circuit]

- NOT GATE : $f(A) = \bar{A} = \sum(0)$

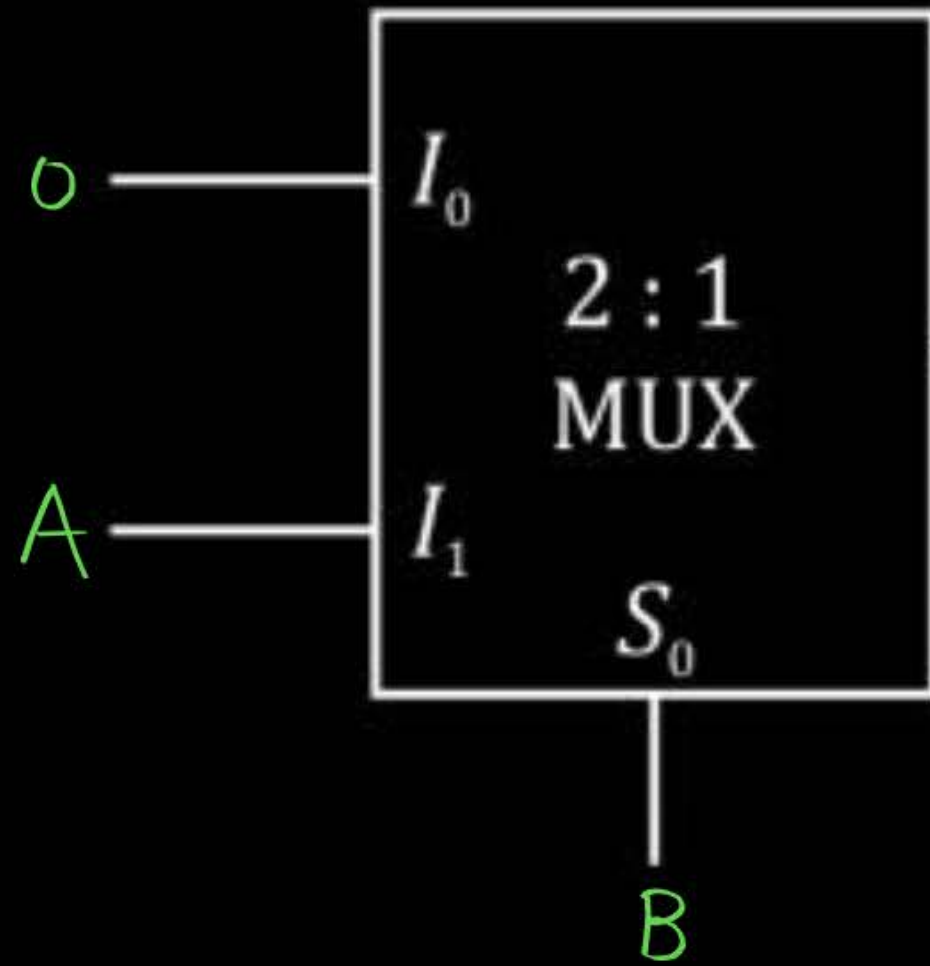


$$f(A=0) = 1, \quad I_0 = 1$$

$$f(A=1) = 0, \quad I_1 = 0$$

$$\bar{A} \cdot 1 + A \cdot 0 = \bar{A}$$

- AND GATE : $f(A, B) = A \cdot B = \sum(3) = \pi(0, 1, 2)$

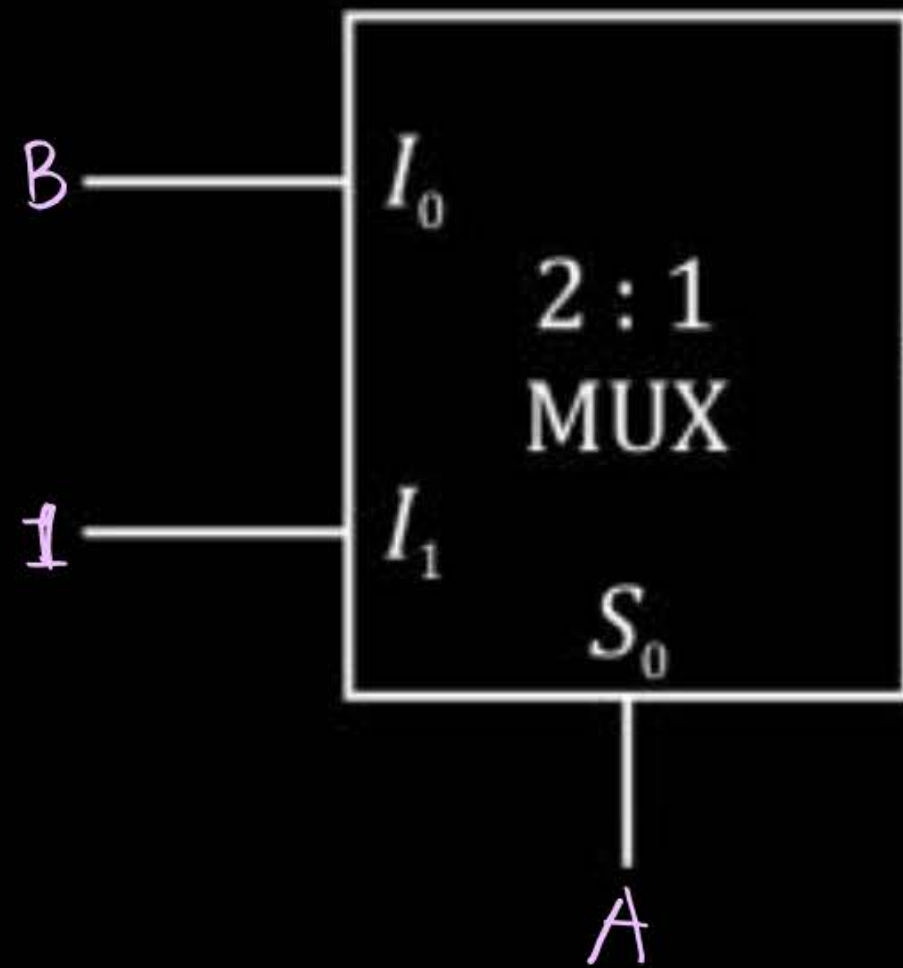


$$f(A, 0) = A \cdot 0 = 0, \quad I_0 = 0$$

$$f(A, B=1) = A \cdot 1 = A, \quad I_1 = A$$

$$\begin{aligned} &\overline{B} \cdot 0 + B \cdot A \\ &= A \cdot B \end{aligned}$$

- OR GATE : $f(A, B) = (A + B) = \sum(1, 2, 3) = \pi(0)$



$$f(A=0, B) = 0 + B = B, \quad I_0 = B$$

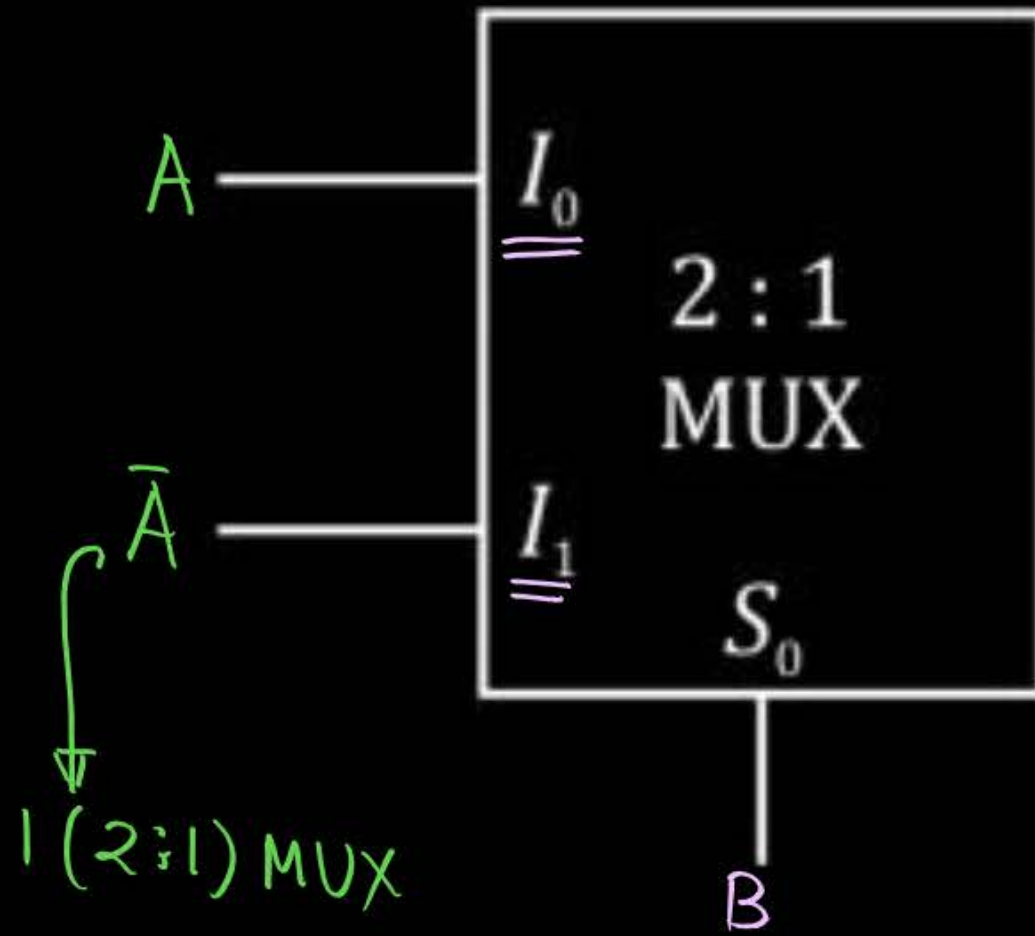
$$f(A=1, B) = 1 + B = 1, \quad I_1 = 1$$

$$\begin{aligned} & \bar{A} \cdot B + A \cdot 1 \\ &= A + (\bar{A} \cdot B) \\ &= (A + \bar{A}) \cdot (A + B) \\ &= (A + B) \end{aligned}$$

- XOR GATE : $f(A, B) = (A \oplus B) = \Sigma(1, 2) = \Pi(0, 3)$

$$f(A, B=0) = A \oplus 0 = A, \quad I_0 = A$$

$$f(A, B=1) = A \oplus 1 = \bar{A}, \quad I_1 = \bar{A}$$

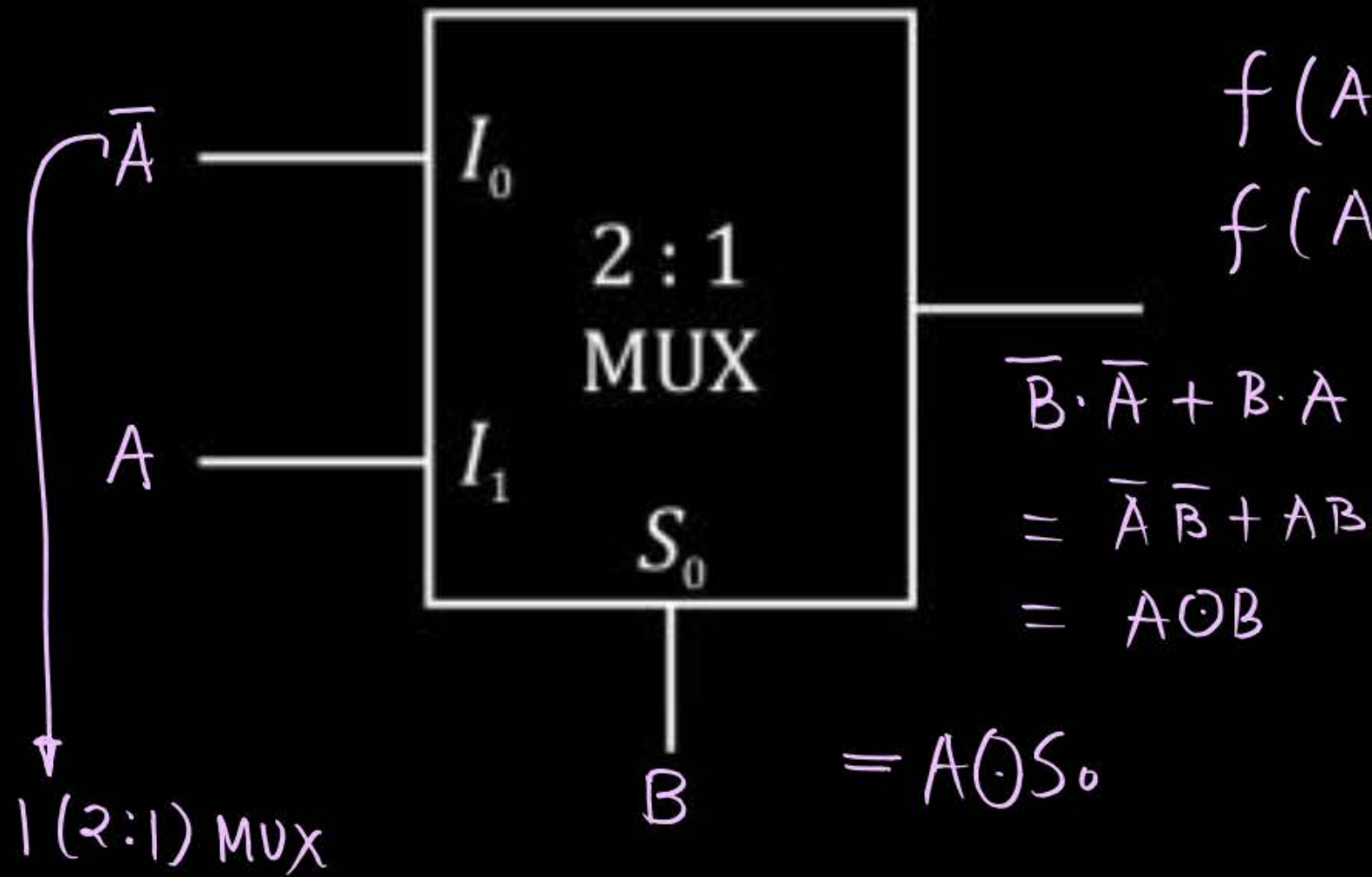


$$\begin{aligned} & \bar{B} \cdot A + B \cdot \bar{A} \\ &= \bar{A}B + A\bar{B} \\ &= A \oplus B \end{aligned}$$

$$A\bar{S}_0 + \bar{A}S_0 = A \oplus S_0$$

$$\Rightarrow A \oplus S_0$$

- XNOR GATE : $f(A, B) = (A \odot B) = \Sigma(0, 3) = \Pi(1, 2)$



$$f(A, B=0) = A \odot 0 = \bar{A}, \quad I_0 = \bar{A}$$

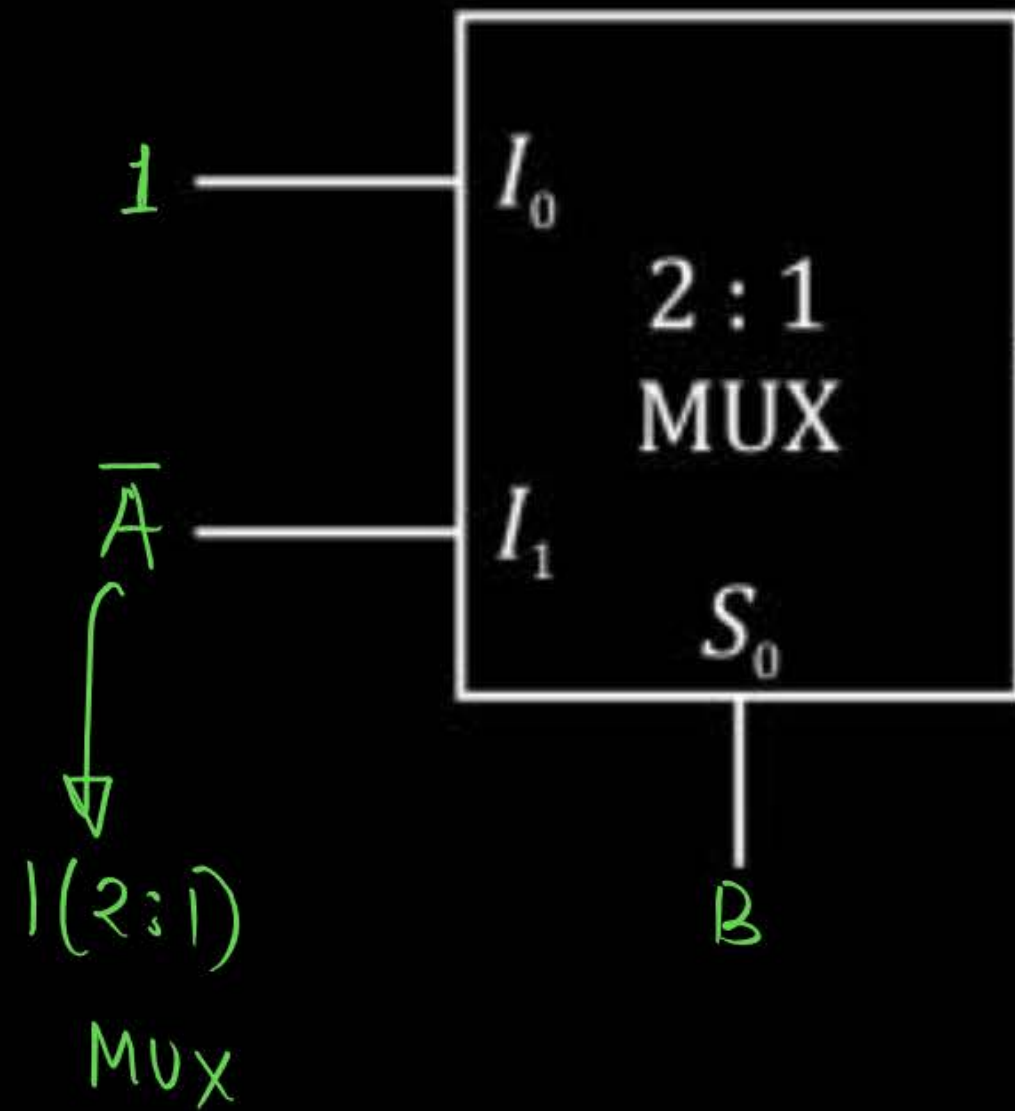
$$f(A, B=1) = A \odot 1 = A, \quad I_1 = A$$

$$\begin{aligned} & \bar{B} \cdot \bar{A} + B \cdot A \\ &= \bar{A} \bar{B} + AB \\ &= A \odot B \end{aligned}$$

$$= A \odot S_0$$

$$\begin{aligned} & \bar{A} \bar{S}_0 + A S_0 \\ &= A \odot S_0 \end{aligned}$$

- NAND GATE : $f(A, B) = \overline{A \cdot B} = \sum(0, 1, 2) = \prod(3)$

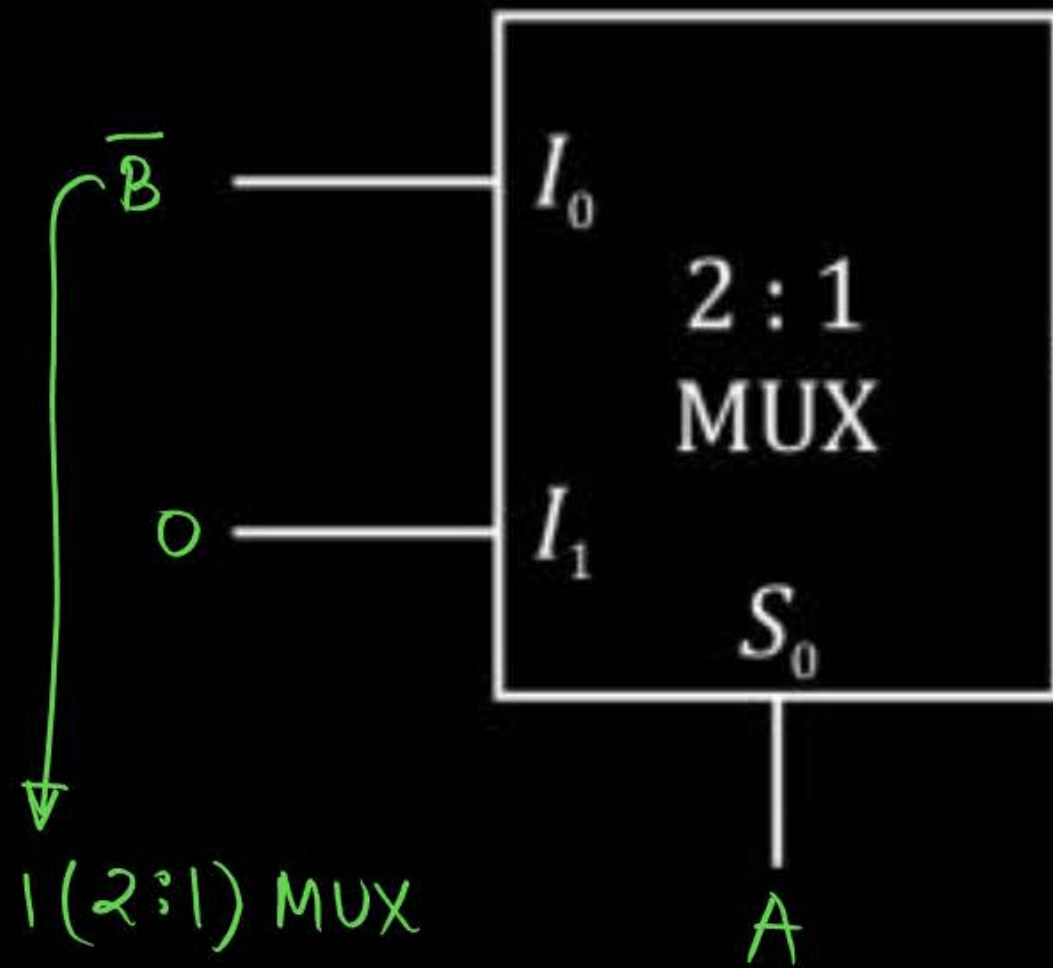


$$f(A, B=0) = \overline{A \cdot 0} = 1, I_0 = 1$$

$$f(A, B=1) = \overline{A \cdot 1} = \overline{A}, I_1 = \overline{A}$$

$$\begin{aligned} & \overline{B} \cdot 1 + (B \cdot \overline{A}) \\ & (\overline{B} + B) \cdot (\overline{B} + \overline{A}) \\ & = \overline{A} + \overline{B} \\ & = \overline{A \cdot B} \end{aligned}$$

- NOR GATE : $f(A, B) = \overline{A + B} = \Sigma(0) = \Pi(1, 2, 3)$



$$f(A=0, B) = \overline{0+B} = \overline{B}, \quad I_0 = \overline{B}$$

$$f(A=1, B) = \overline{1+B} = 0, \quad I_1 = 0$$

$$\begin{aligned} & \overline{A} \cdot \overline{B} + A \cdot 0 \\ &= \overline{A} \cdot \overline{B} \\ &= \overline{A+B} \end{aligned}$$

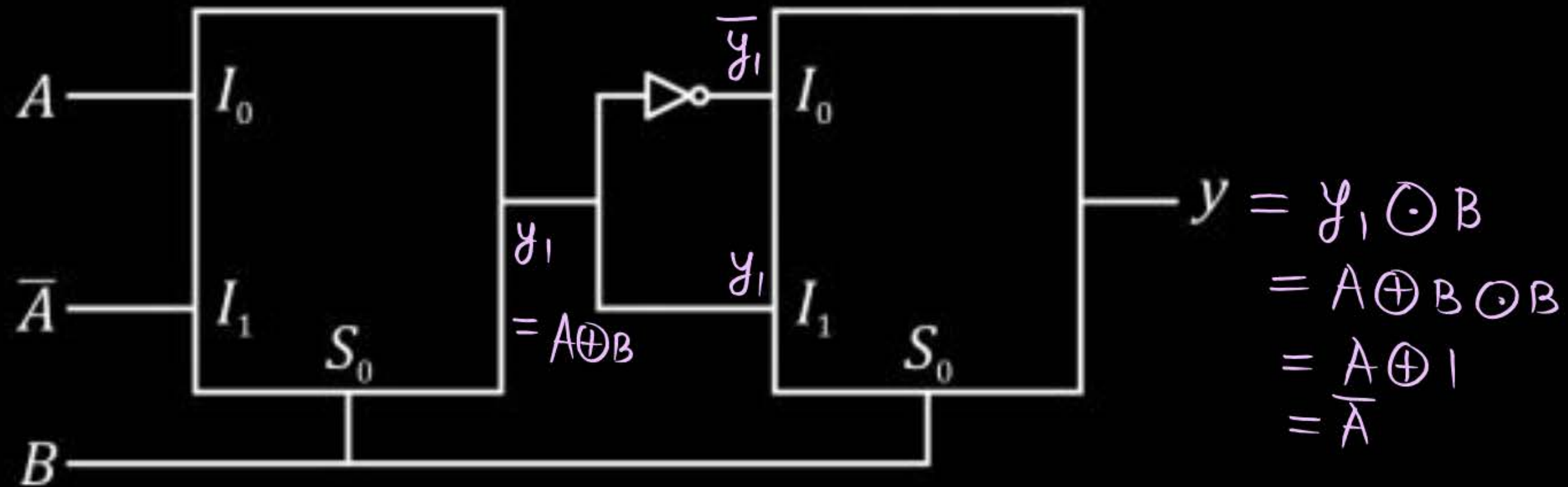
- IMP Notes :

- For implementing basic gates (NOT, AND, OR) we require only one (2:1) MUX.
- For implementing non-basic gates (XOR, XNOR, NAND, NOR), we require two (2:1) MUX.

[Questions On MUX]



A digital circuit is as given below :



Output y is

(a) $A \oplus B$

(b) $A \odot B$

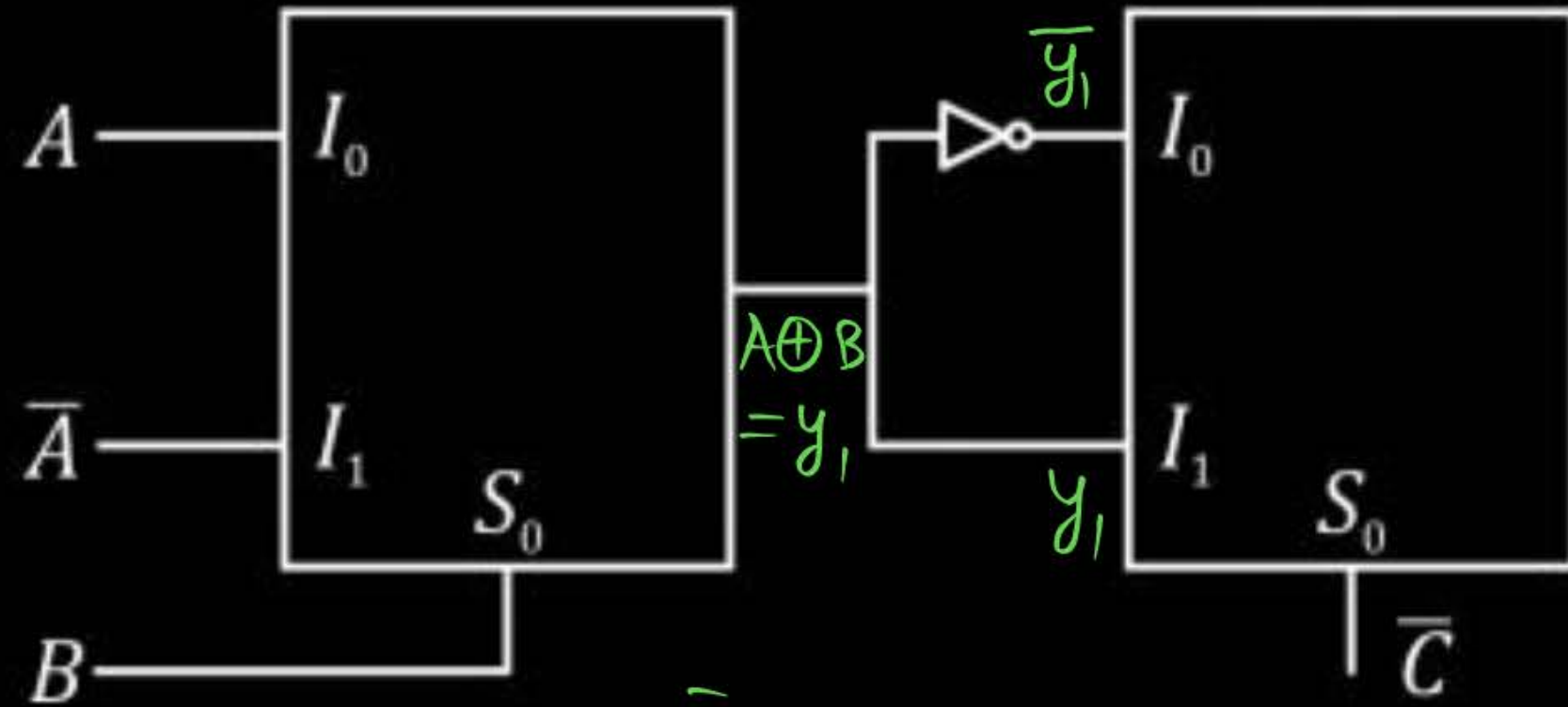
(c) $C.A$

~~(d) \bar{A}~~

[Question 2]



A digital circuit is as given below :



Output y is

(a) $A \oplus B \oplus C$

(b) $A \oplus B \odot C$

(c) $A \oplus C$

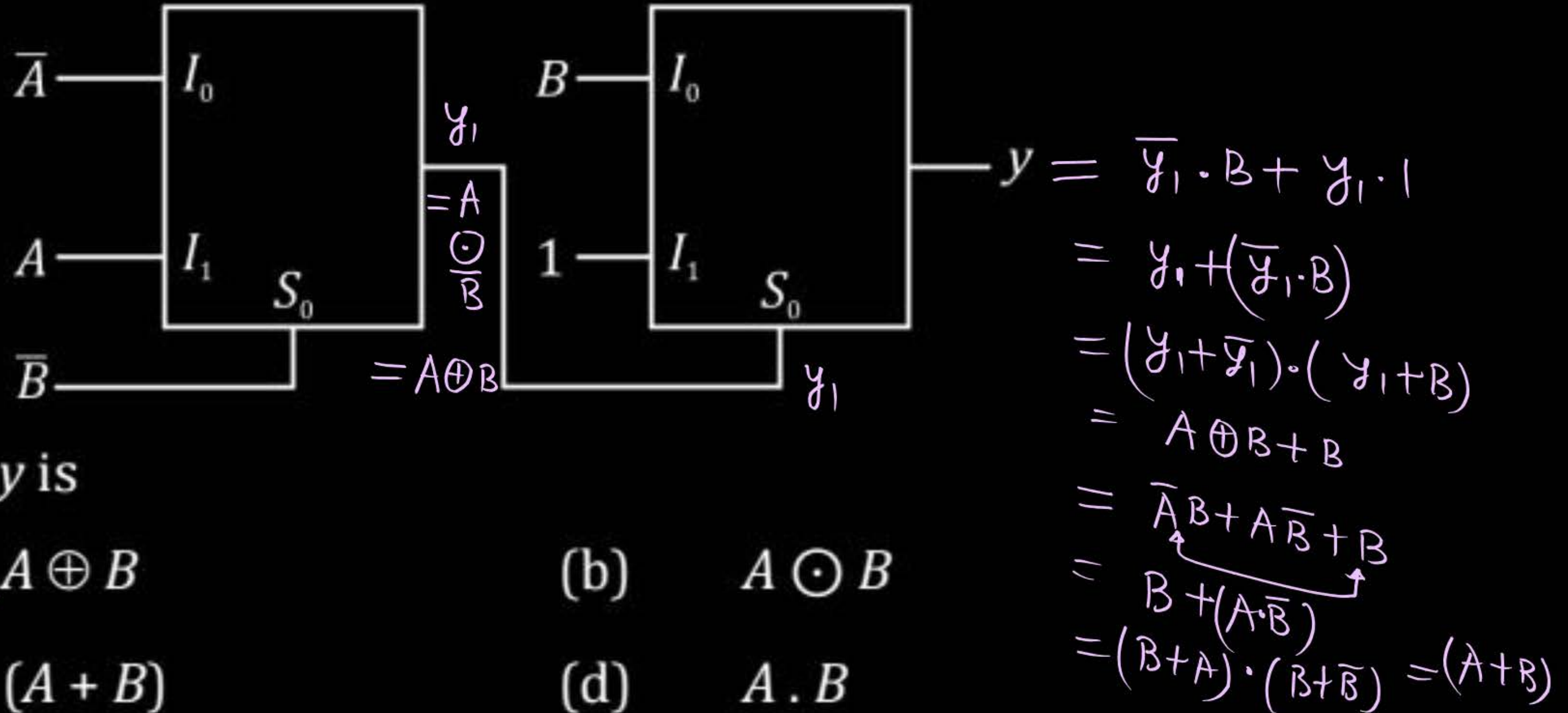
(d) $A \oplus B$

$$\begin{aligned}
 y &= y_1 \odot S_0 \\
 &= y_1 \odot \bar{C} \\
 &= A \oplus B \odot \bar{C} \\
 &= \overline{A \oplus B \odot C} \\
 &= A \odot B \odot C \\
 &= A \oplus B \oplus C
 \end{aligned}$$

Question 3



A digital circuit is as given below :



Output y is

(a) $A \oplus B$

(b) $A \odot B$

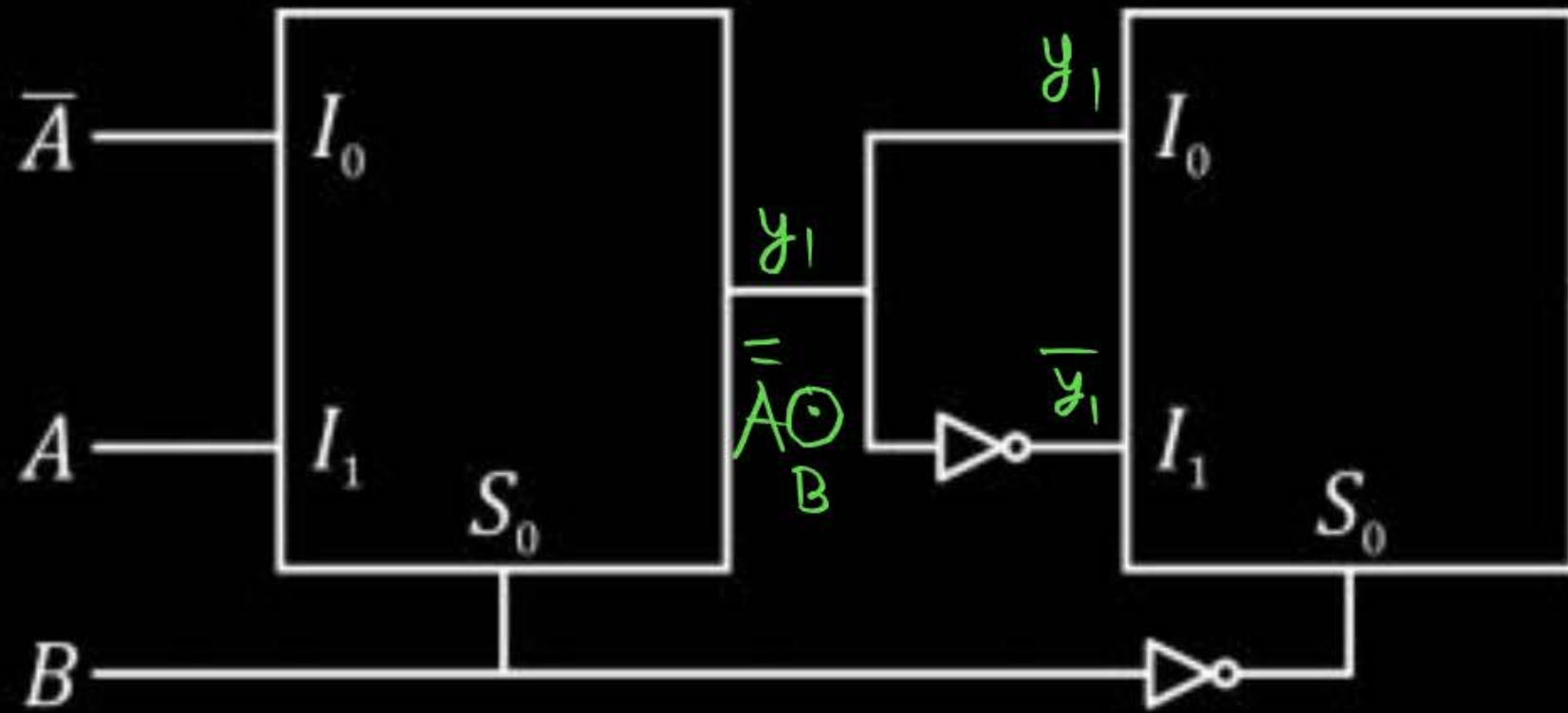
☒ (c) $(A + B)$

(d) $A \cdot B$

[Question 4]



A digital circuit is as given below :



$$\begin{aligned} y &= y_1 \oplus S_0 = y_1 \oplus \bar{B} \\ &= A \odot B \oplus \bar{B} \\ &= A \odot 1 \\ &= A \end{aligned}$$

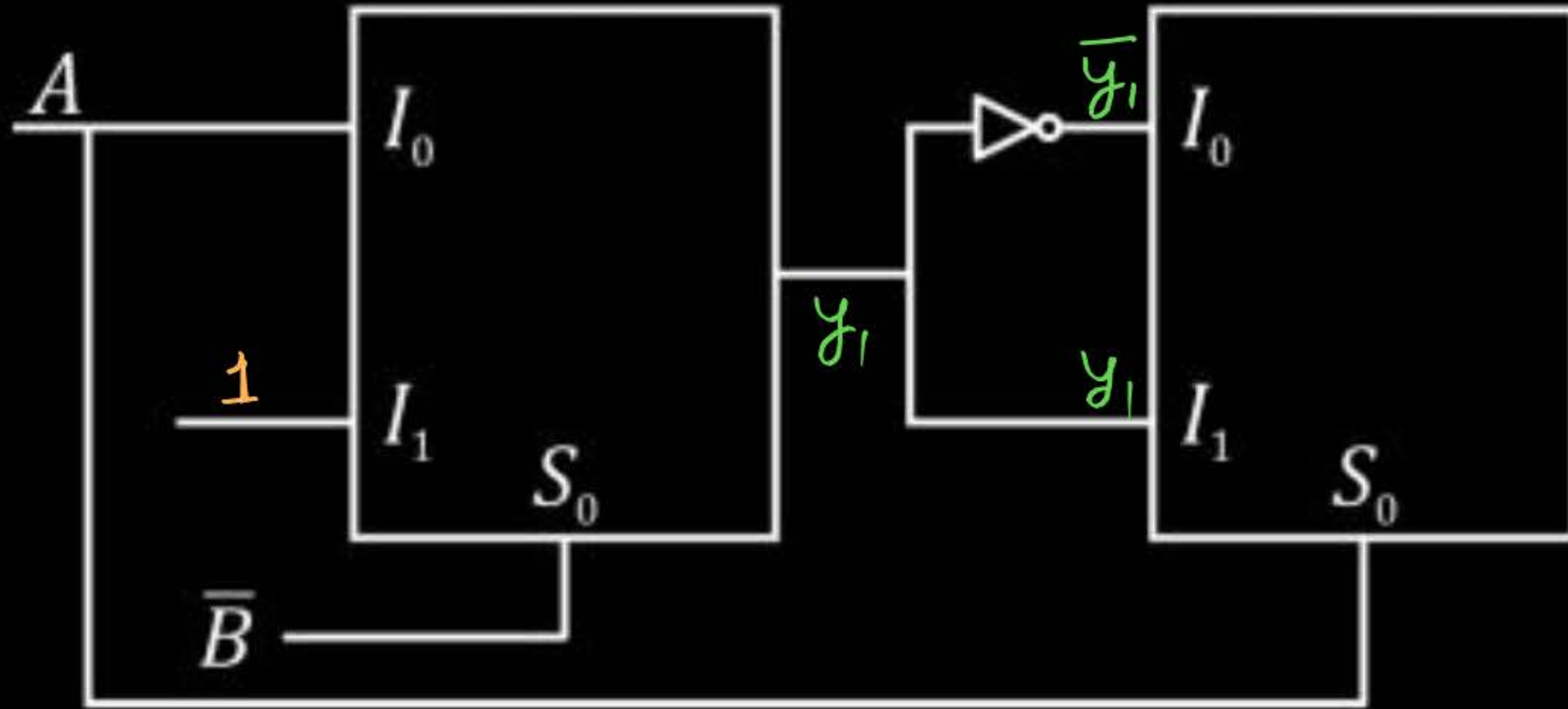
Output y is

- (a) B
- (b) A ✓
- (c) $(A + B)$
- (d) $A . B$

[Question 5]

$$\begin{aligned} y_1 &= \overline{S_0} A + S_0 \cdot 1 = B \cdot A + \overline{B} \\ &= \overline{B} + (B \cdot A) = (\overline{B} + A)(\overline{B} + B) \\ &= (A + \overline{B}) \end{aligned}$$

A digital circuit is as given below :



$$\begin{aligned} A \cdot (A + \overline{B}) \\ \overline{A} + A\overline{B} &= A(1 + \overline{B}) \\ &= A \end{aligned}$$

$$\begin{aligned} y &= y_1 \odot A \\ &= (A + \overline{B}) \odot A \\ &= (\overline{A + \overline{B}}) \cdot \overline{A} + (A + \overline{B})A \\ &= \overline{A} \cdot B \cdot \overline{A} + A \\ &= \overline{A}B + A \\ &= A + (\overline{A} \cdot B) \\ &= (A + \overline{A})(A + B) \\ &= (A + B) \end{aligned}$$

The output y is

(a) $A \odot B$

(b) $\overline{A}B$

(c) $(A + B)$

(d) $A\overline{B}$

[Question 6]

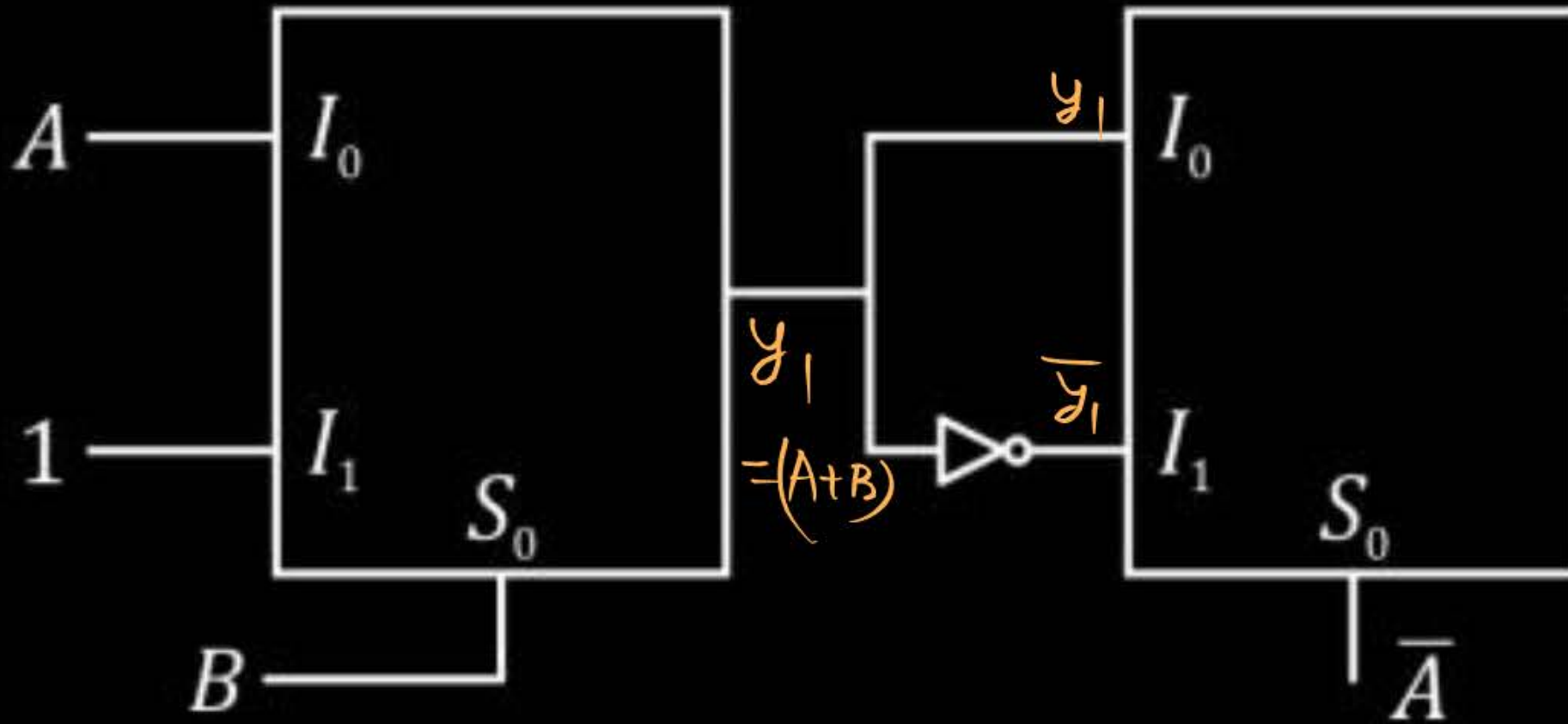
$$y_1 = \overline{S_0} \cdot A + S_0 \cdot 1 = (\overline{B} \cdot A) + B$$

$$= (B + \overline{B}) \cdot (B + A)$$

$$= (A + B)$$



A digital circuit is as given below :



Output y is

(a) 1

(b) $A \odot B$

(c) $\overline{A} \overline{B} = (A + \overline{B})$

(d) $\overline{A \cdot B}$

$$y = y_1 \oplus S_0$$

$$= y_1 \oplus \overline{A}$$

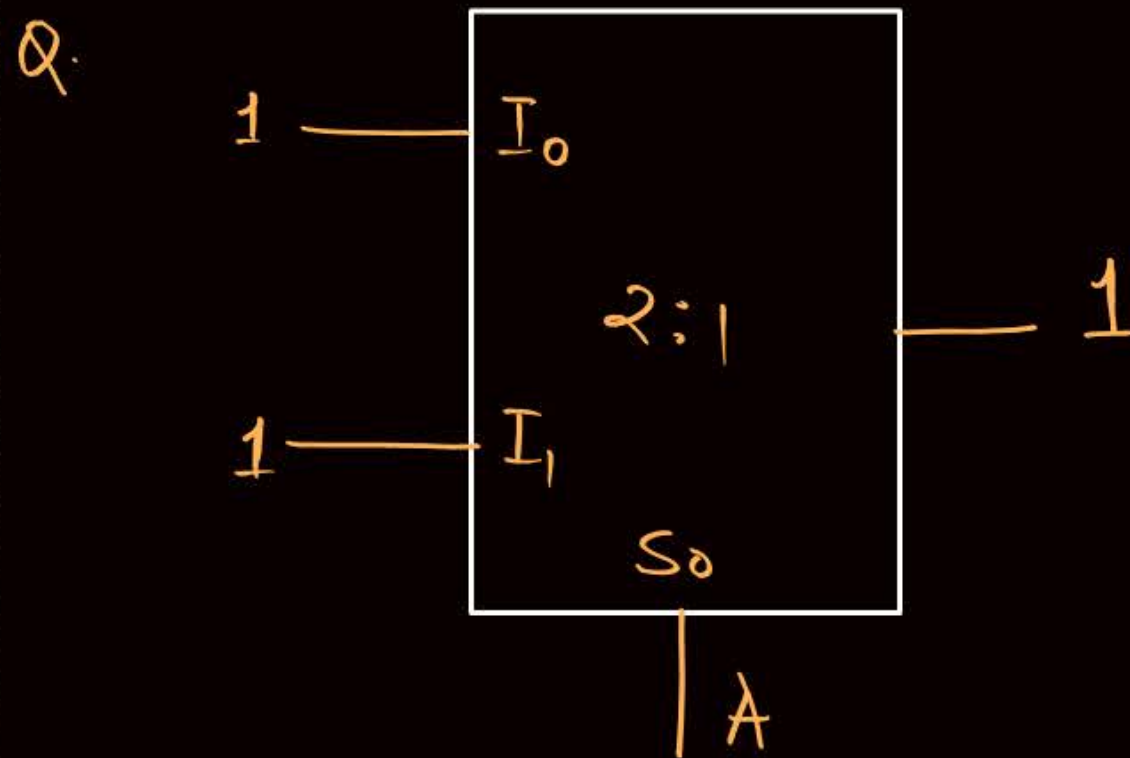
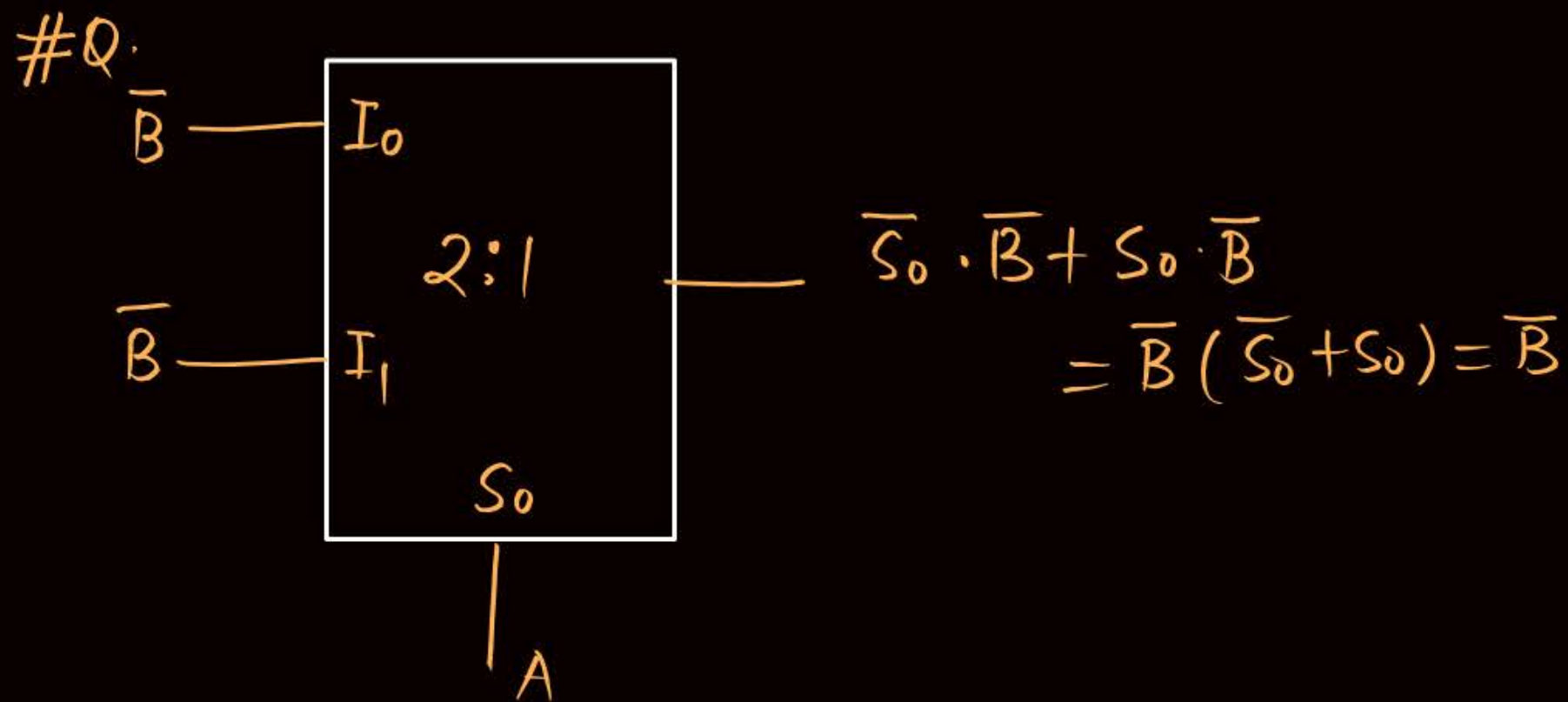
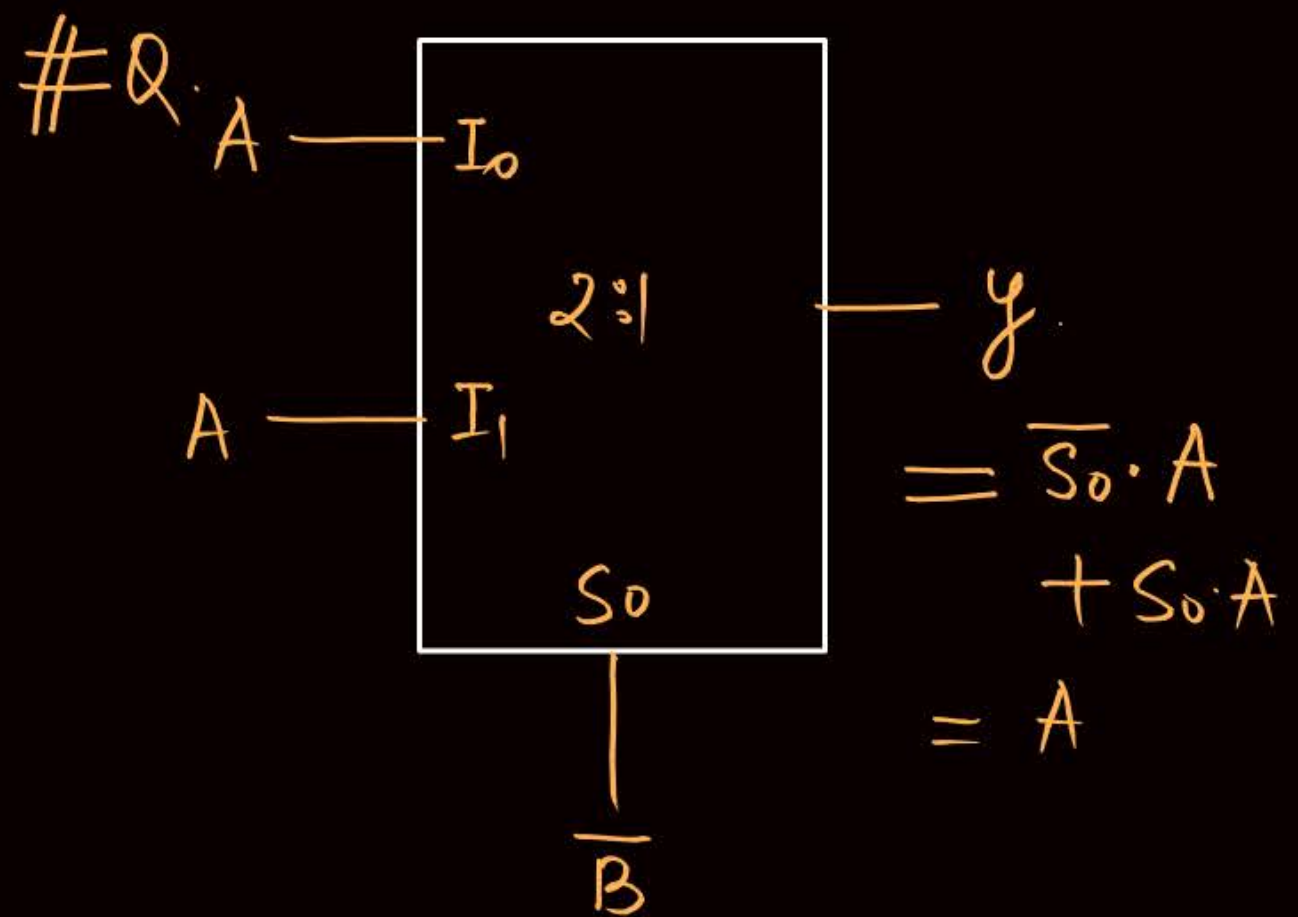
$$= (A + B) \oplus \overline{A}$$

$$= \overline{(A + B)} \cdot \overline{A} + (A + B) \cdot \overline{\overline{A}}$$

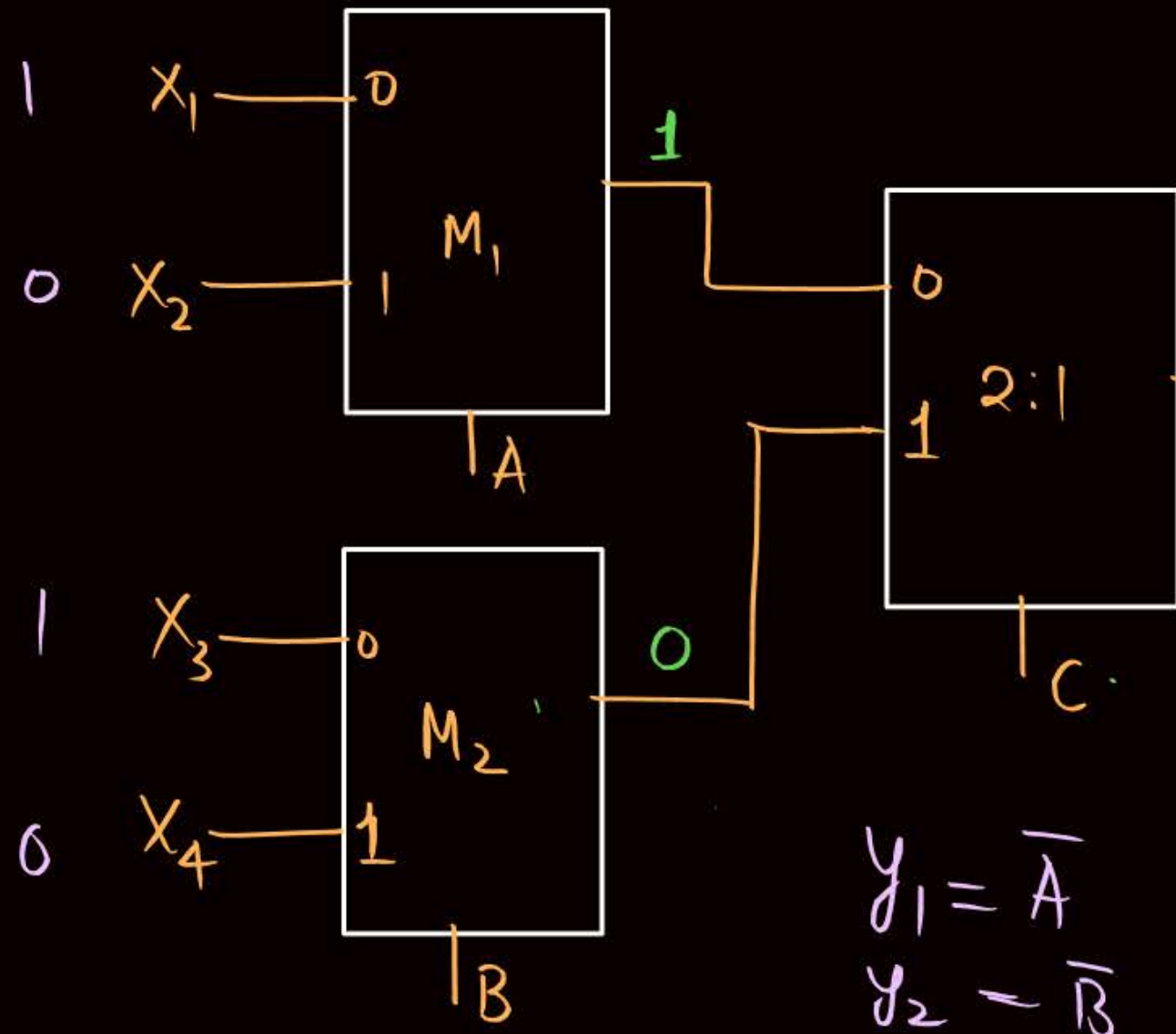
$$= \overline{A} \overline{B} \cdot \overline{A} + A$$

$$= \overline{A} \overline{B} + A$$

$$(A + \overline{B}) = A + (\overline{A} \overline{B}) = (A + \overline{A}) \cdot (A + \overline{B})$$



Q. Consider a digital circuit consisting of three 2:1 MUX M_1 , M_2 & M_3 as shown.
 X_1 & X_2 are input of M_1 , X_3 & X_4 are input of M_2 . A , B , C are select lines of M_1 , M_2 & M_3 respectively.



For an instance $X_1=1, X_2=1, X_3=0, X_4=0$
 the number of combinations ^{A, B, C} that gives O/P
 $Y=1$ _____.

$$Y = \bar{C} \cdot 1 + C \cdot 0$$

$$= \bar{C} \cdot \bar{A}$$

$$+ C \bar{B}$$

$$= \bar{A} \bar{C} + \bar{B} C$$

$$= \sum (0, 1, 2, 5)$$

0	0	0
0	1	0
1	0	0
1	1	0

\bar{A}	\bar{C}	
0	0	0 = 0
0	1	0 = 2

\bar{B}	C	
0	0	1 → 1
1	0	1 → 5

$$Y_1 = \bar{A}$$

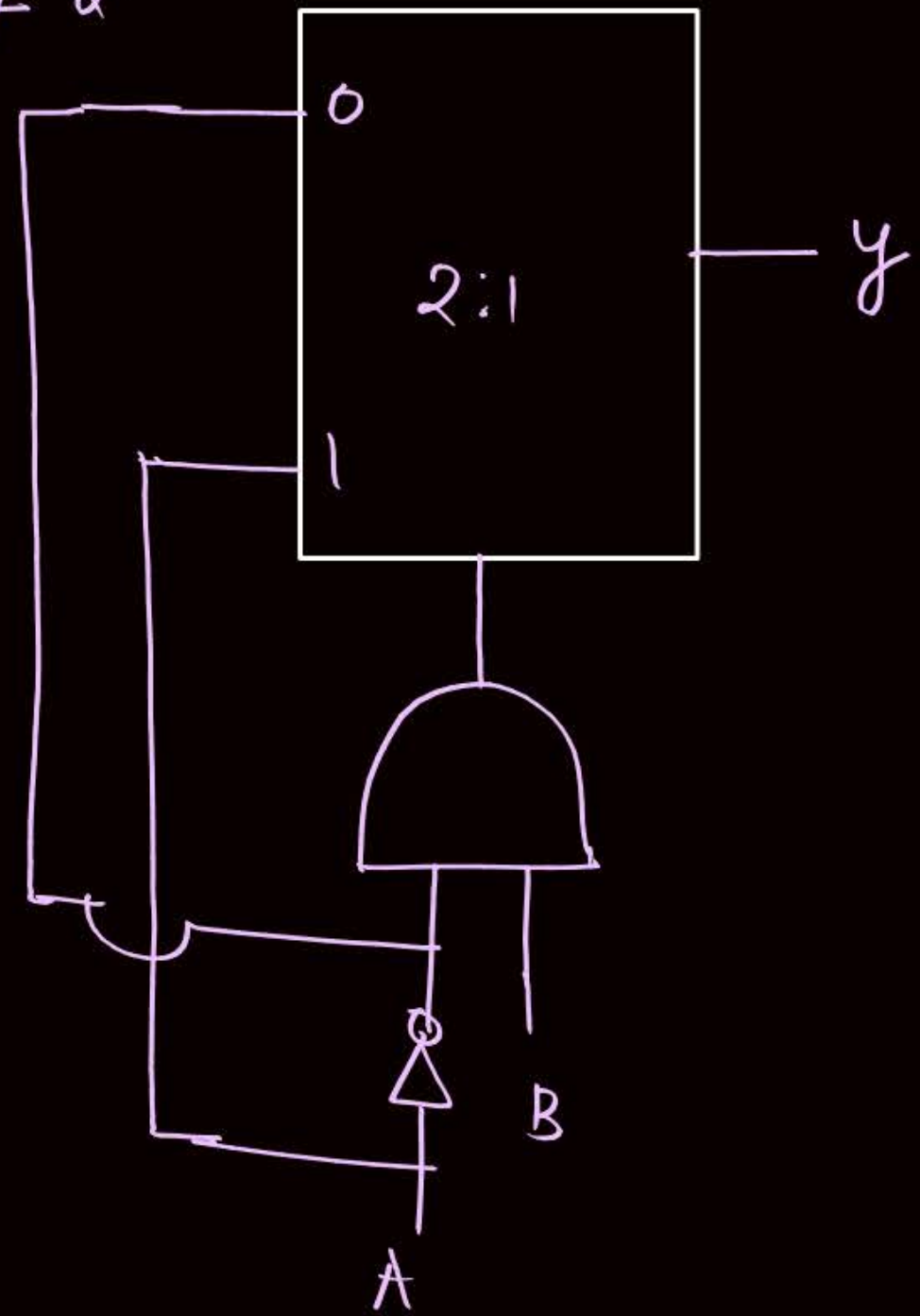
$$Y_2 = \bar{B}$$

$$y = \overline{c}$$

$$= \Sigma(0, 2, 4, 6)$$

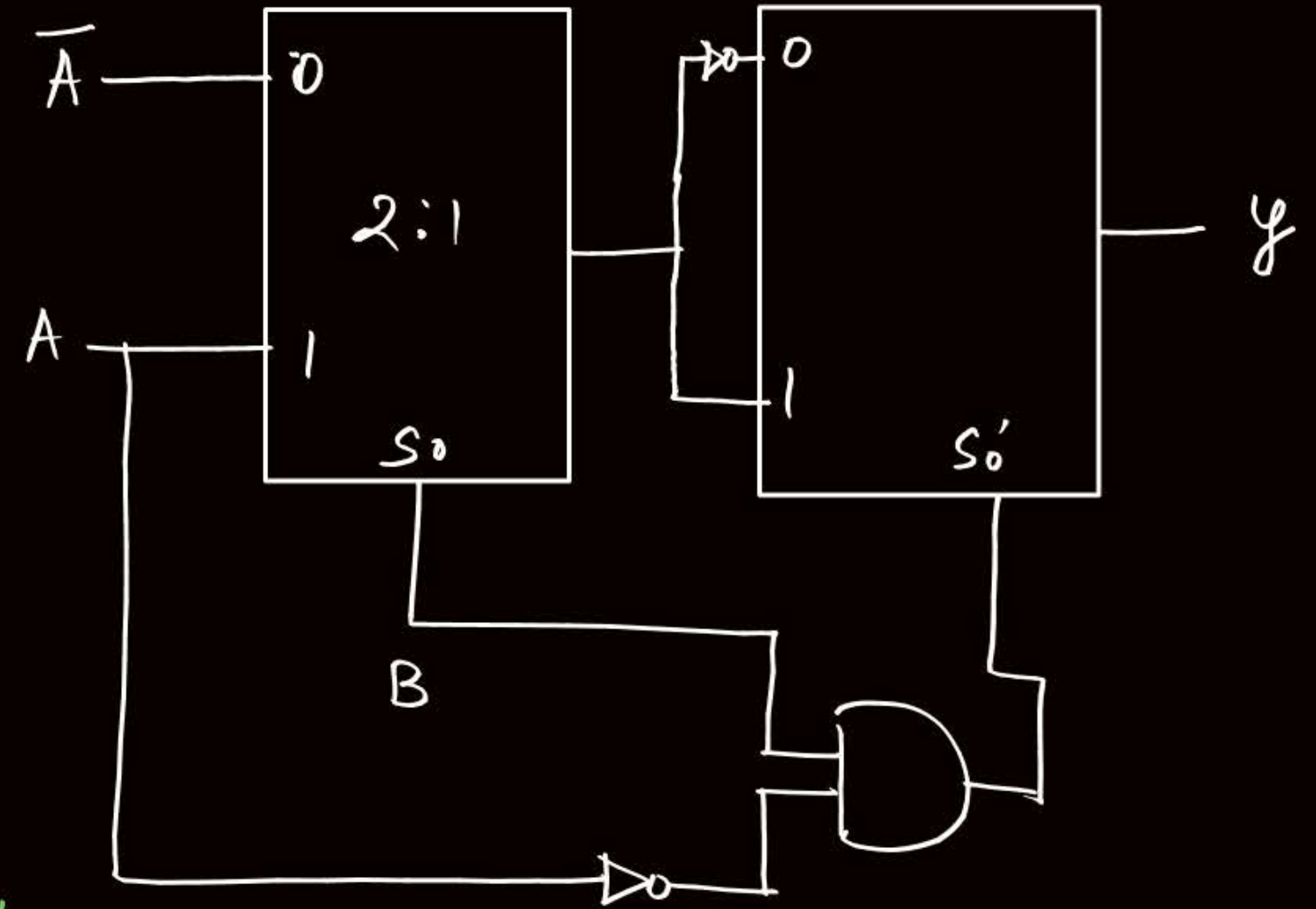
		\overline{c}
0	0	0
0	1	0
1	0	0
1	1	0

Q.



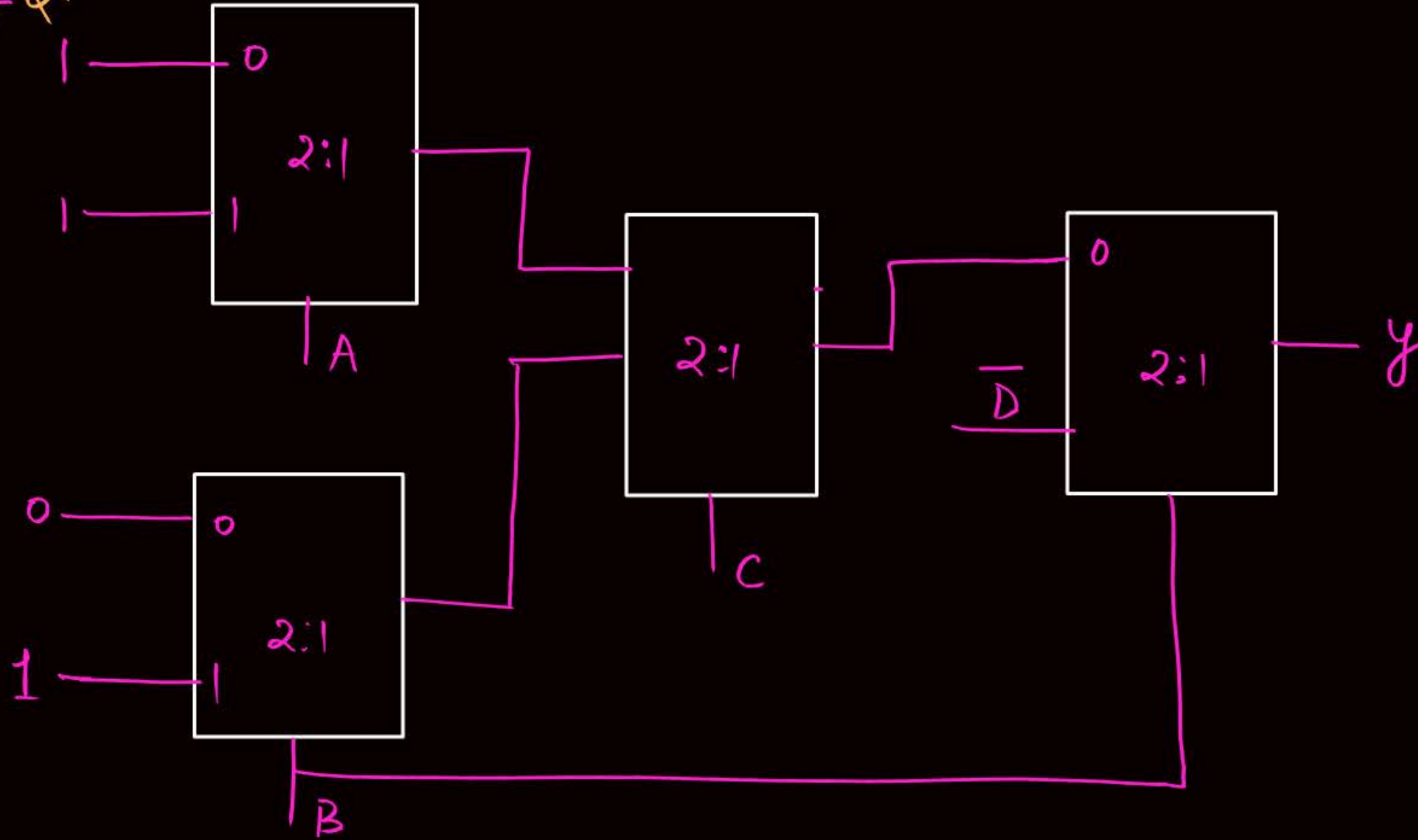
- O/P y is
- a. $\overline{A+B}$
 - b. $\overline{A \cdot B}$
 - c. $A \overline{B}$
 - d. $\overline{A} \cdot B$

Q.



- O/P y is
- a. \overline{AB}
 - b. $\overline{A+B}$
 - c. $A \overline{B}$
 - d. $\overline{A} B$

#Q.



No. of combinations of
A, B, C, D for which O/P
 $Y = 0$ _____.



2 Minute Summary

→ MUX & Question discussion.

Thank you

GW
Soldiers !

