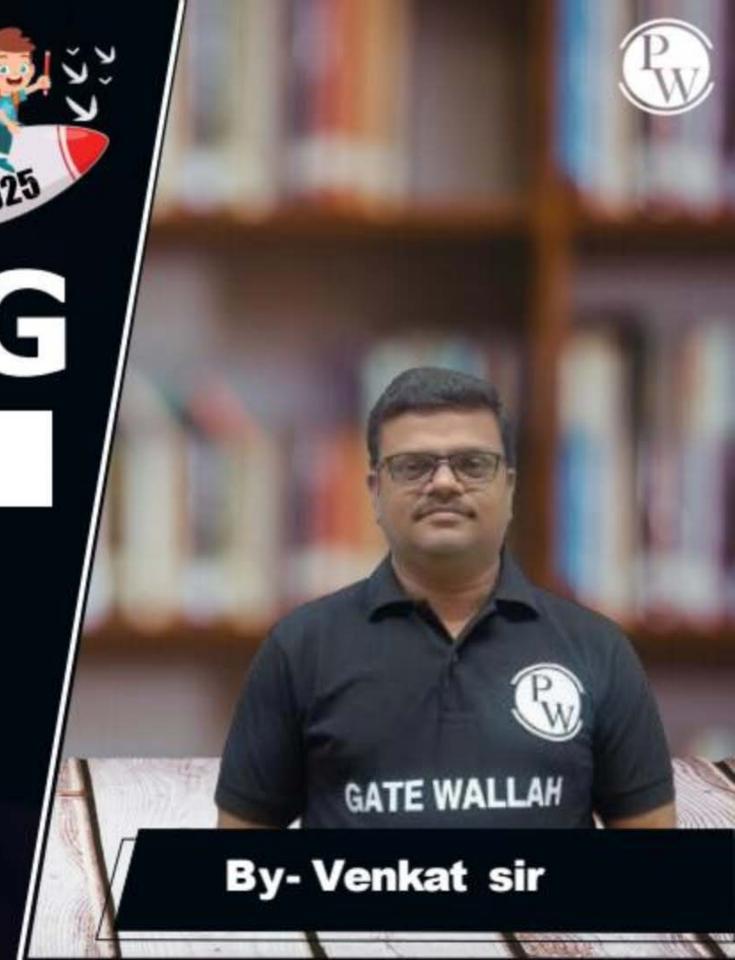
CS & IT ENGING

Theory of Computation

DFA

Lecture No.- 04



Recap of Previous Lecture





Topic DFA Design

Topic Possibility of DFA.

if Dependency exist. DFA possible no Dependency -> DFA possible if DFA possible they that language if Kyown as Regular Language.

if DFA not possible => Non Regular Language.

Topics to be Covered









Topic

Finite Automaton & Regular Languages.

Topic

Pushdown Automata & Context free Languages.

Topic

Turing Machine & Recursive Enumerable Languages.

Topic

Undecidability.

BOOKS:





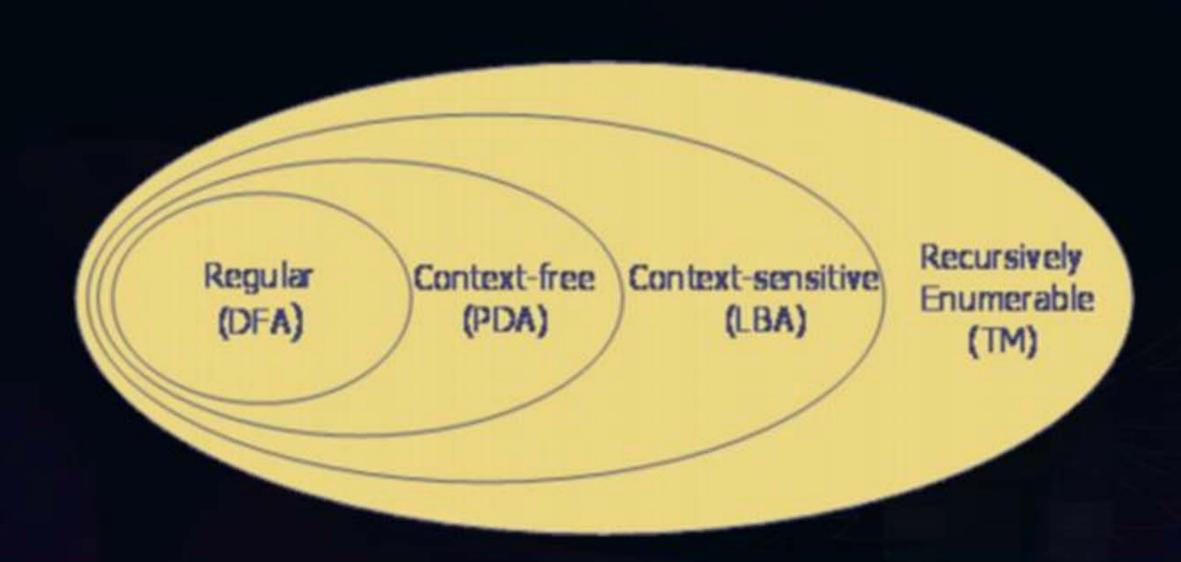






Topic: Theory of Computation







Topic: Finite Automata



It is a mathematical model which contains finite number of states and transitions.





Topic: Deterministic Finite Automata



DFA: It is a finite automata in which from every state on every input symbol exactly one transition should exits.



Topic: Deterministic Finite Automata



FORMAL DFA: DFA is defined as

DFA = $(Q, \sum, q_0, F, \delta)$

Q: Finite set of states

 Σ : Input alphabet

q₀: Initial state

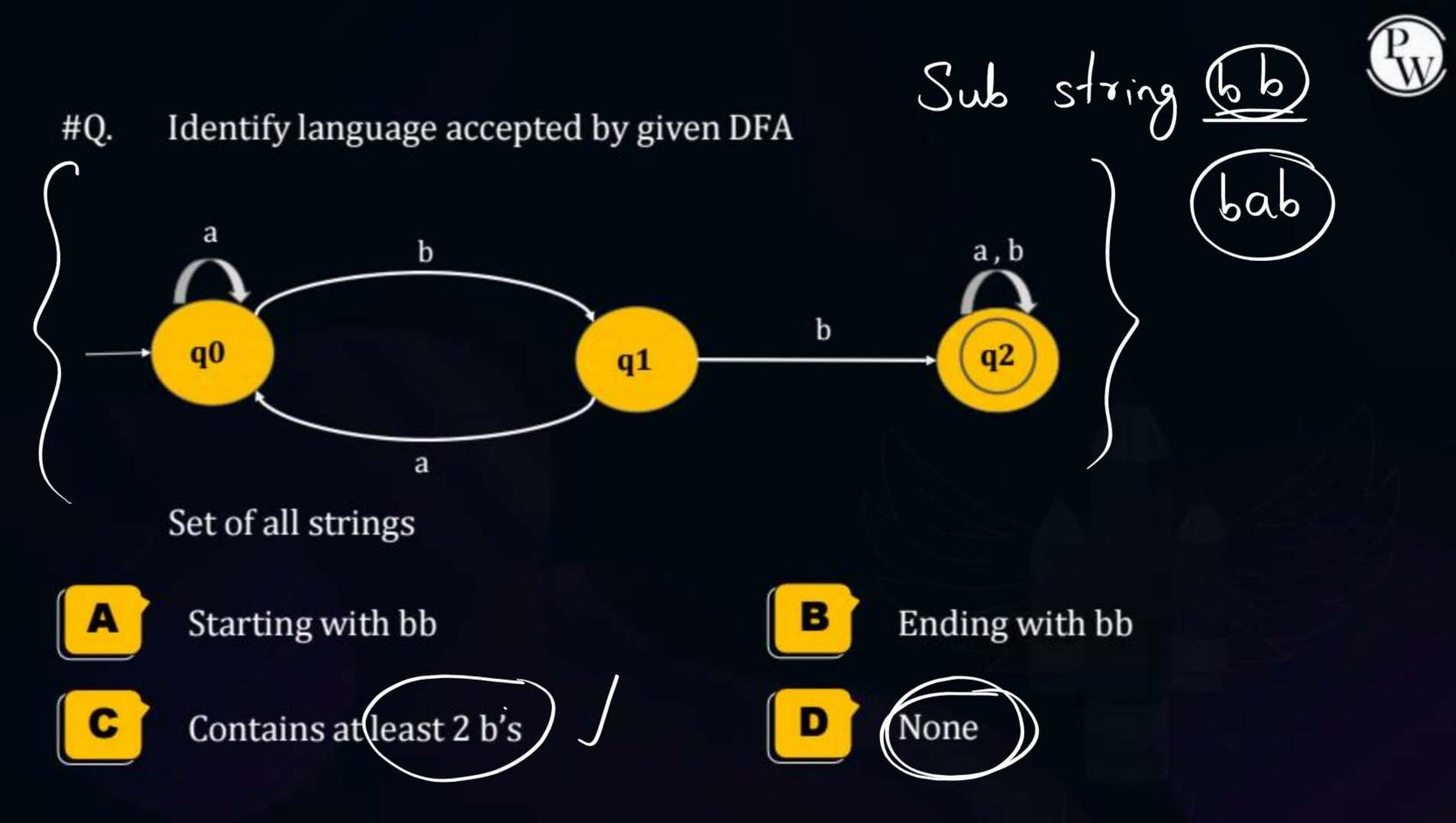
F: Set of final states

δ: Transition function Q * Σ → Q

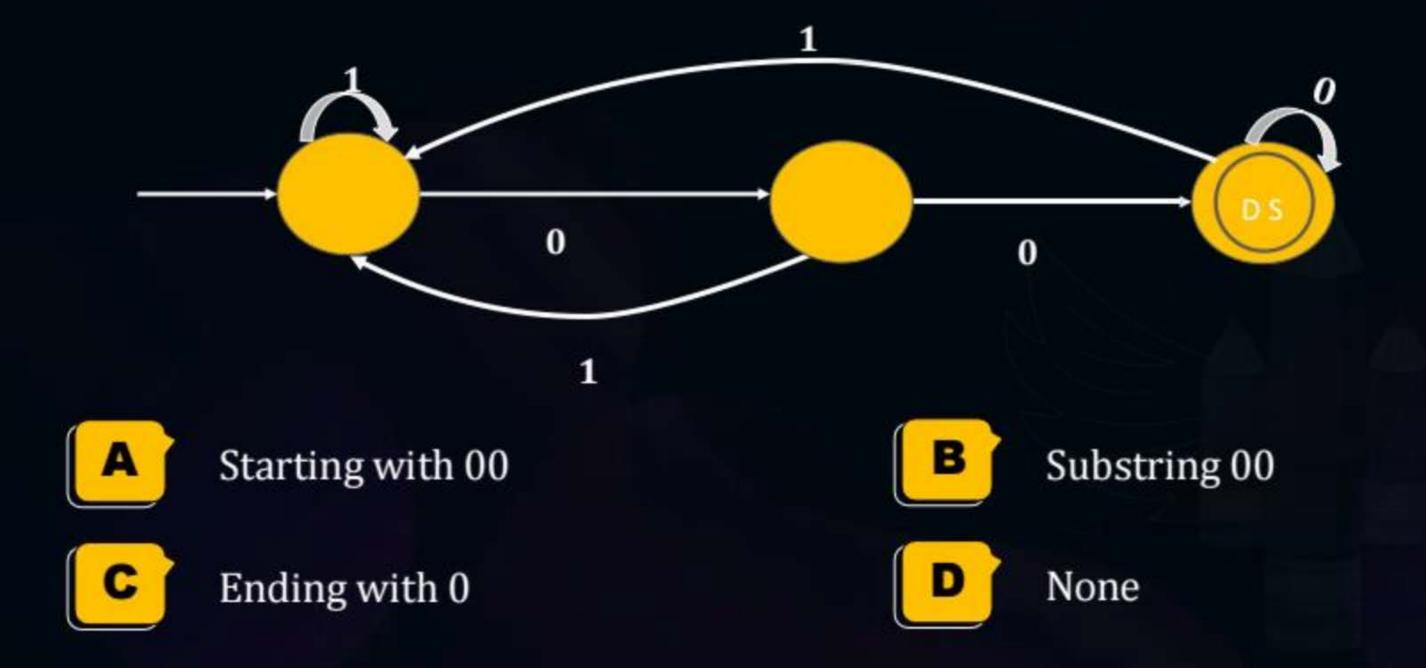


DFA acceptance method:

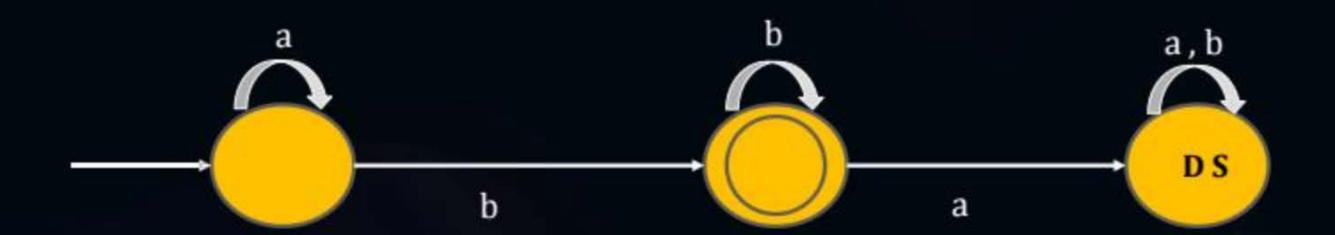
- Start at the Initial State: Begin at the initial state of the DFA.
- 2. Read Input Symbols: For each symbol in the input string, read it one by one.
- 3.Follow Transitions: Based on the current state and the input symbol being read, follow the transition defined by the transition function of the DFA. This transition function specifies the next state of the automaton for each combination of current state and input symbol.
- 4.Repeat Until End of Input: Continue this process of reading input symbols and following transitions until you reach the end of the input string.
- 5.Final State: Once you have processed all input symbols, check the current state of the DFA. If it is one of the accepting states then the input string is accepted. Otherwise, it is rejected.
- 6.Acceptance: If the DFA halts in an accepting state after reading the entire input string, then the Input is accepted.











- **A** $L = \{a^n b^m | n, m ≥ 1\}$
- C $L = \{a^n b^m | n, m \ge 0 \}$

- **B** $L = \{a^n b^m | n \ge 1, m \ge 0\}$
- D None

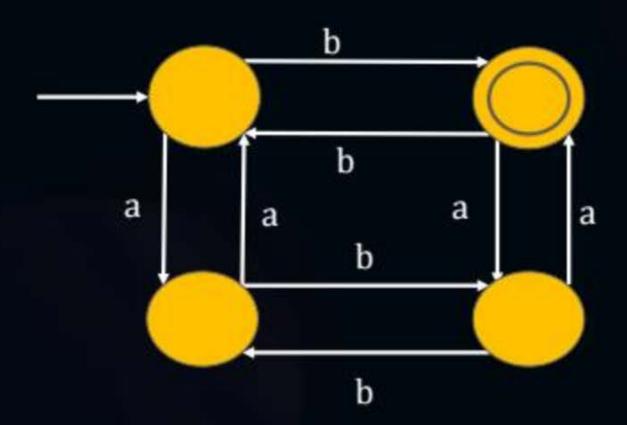




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- D None





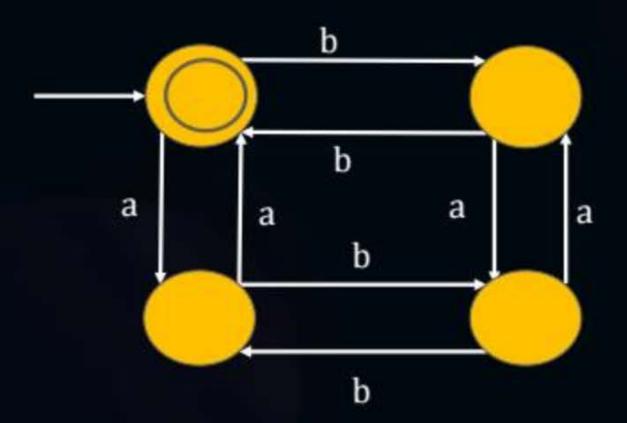
a's even and # b's even

a's odd and # b's even

a's odd and # b's odd

a's even and # b's odd





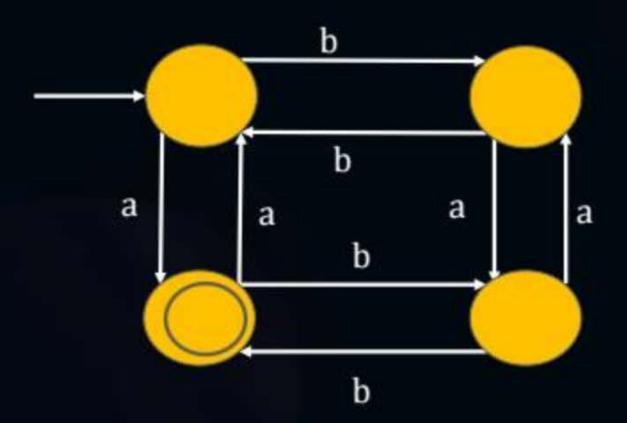
a's even and # b's even

a's odd and # b's even

a's odd and # b's odd

a's even and # b's odd





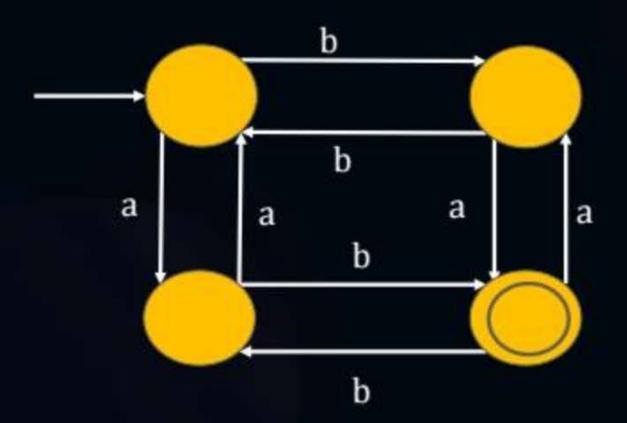
a's even and # b's even

a's odd and # b's odd

a's odd and # b's even

a's even and # b's odd





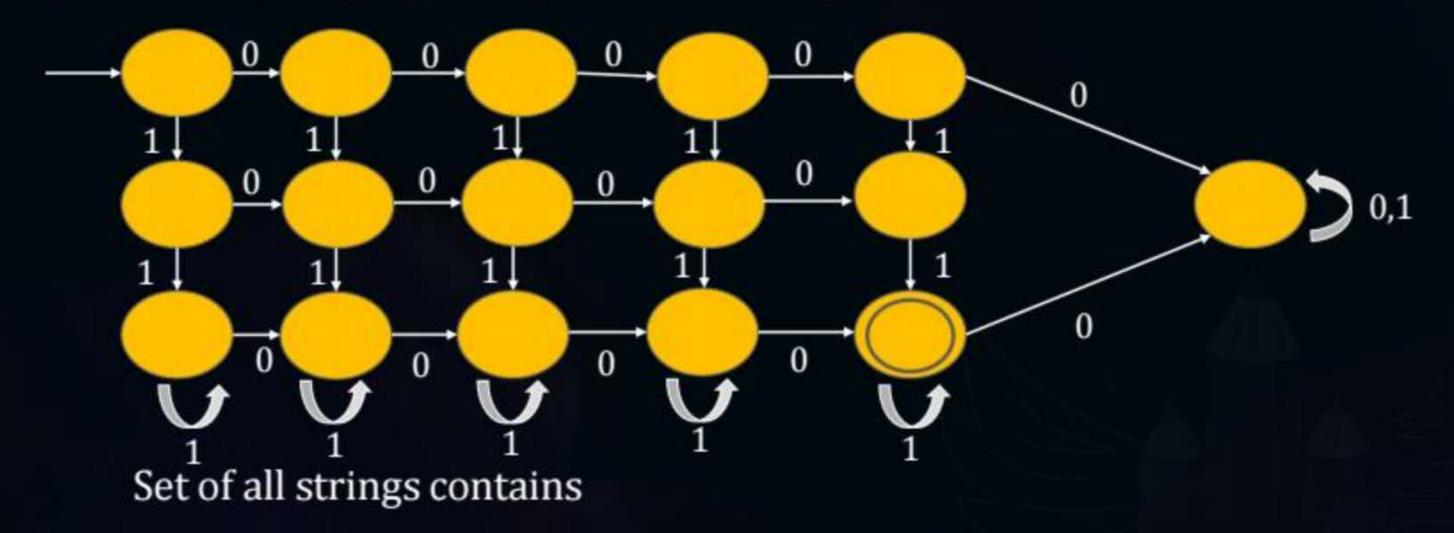
a's even and # b's even

a's odd and # b's odd

a's even and # b's odd

a's odd and # b's even





A Length of the string alteast 6

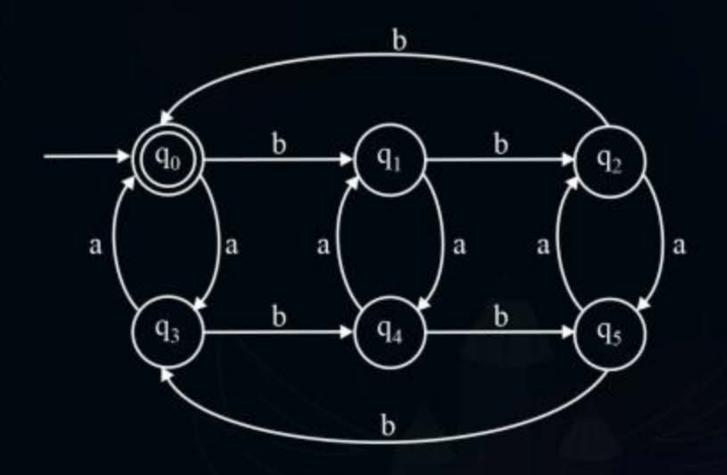
0's exactly4 and 1's atleast 3

0's atleast 4 and # 1's exactly 2

D None



#Q. Consider the finite automaton m, m accepts all strings of a's and b's m which the number of a's and b's are respectively.



A Divisible by 3 and 2

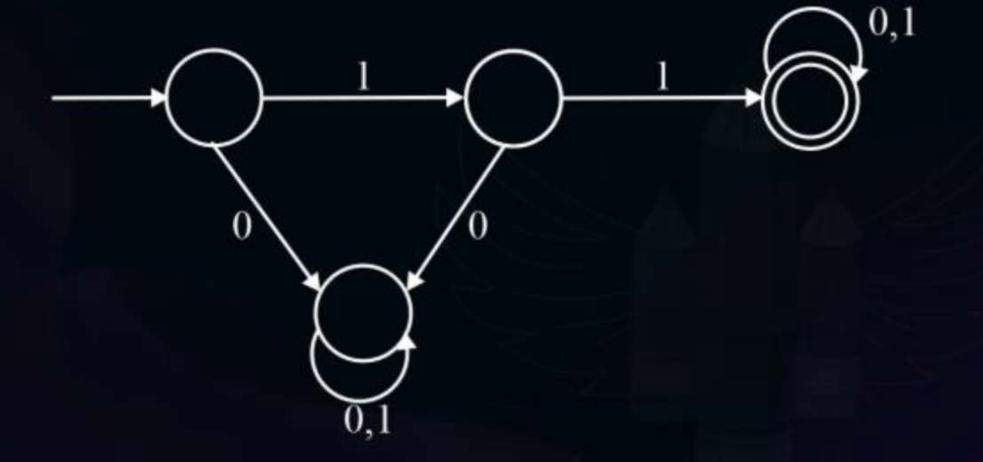
C Even and odd

B Odd and even

Divisible by 2 and 3

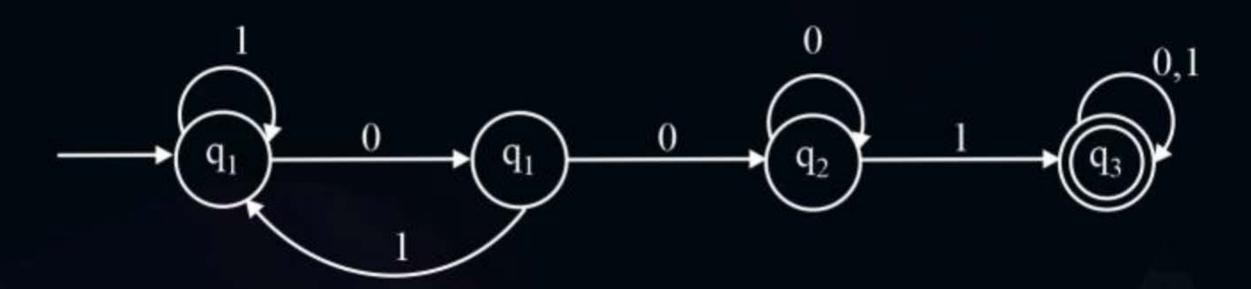


#Q. The number of binary strings of length n' accepted by the following finite automata is-





#Q. Consider the following deterministic finite automata m.



The language accepted by finite automata, which is obtained by interchanging final and non-final states in m, is-



A The set of all strings containing 001 as the subtrings

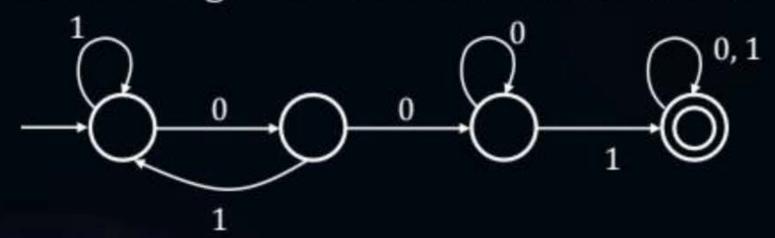
The set of all strings not containing 001 as the subtrings

The set of all strings ending with subtrings 001

The set of all strings ending with subtrings 001



#Q. Consider the following deterministic finite state automaton M.



Let S denote the set of seven-bit binary strings in which the first, the fourth, and the last bits are 1. The number of strings in S that are accepted by M is

A 1

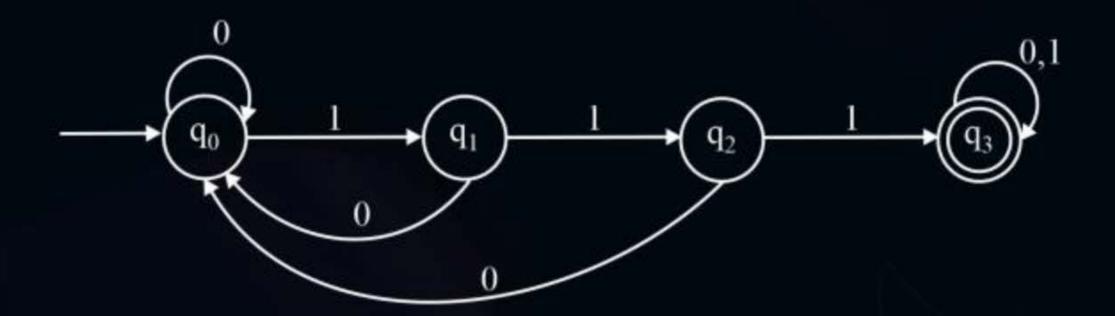
C 7

B 5

D 8



#Q. Consider the finite automata m.



The language accepted by m is, over the alphabet {0,1}



A The set of all strings containing three consecutive I's

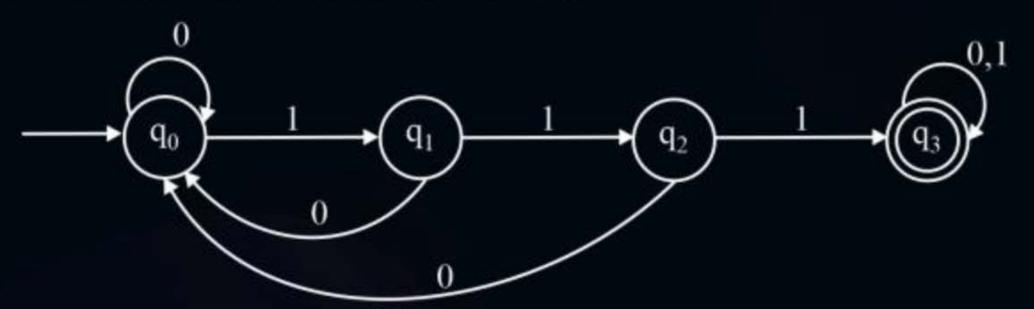
The set of all strings not containing three consecutive I's

The set of all strings beginnigs with three consecutive I's

The set of all strings ending with three consecutive I's



#Q. Consider the finite automata m.



Let S denotes the set of all six bit binary strings in which first and fourth bits are 1. The number of strings in S that are accepted by m is-

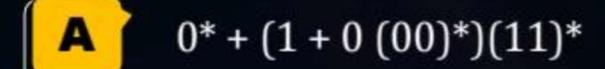
A 1

C 7

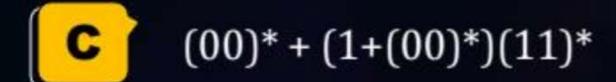
D 8



#Q. Let M be the 5-state NFA with ∈ - transitions shown in the diagram below. Which one of the following regular expression represents the language accepted by M?



B 0++1(11)*+0(11)*



D (00)* + 1(11)*







Construct DFA for the following Language.



#Q. Construct DFA for the language: $L = \{a^nb^m \mid n > m\}$



#Q. For which of the following languages DFA is possible

1.
$$L = \{a^n b^m \mid n, m \ge 1\}$$

2.
$$L = \{a^n b^n \mid n \ge 1\}$$

3.
$$L = \{a^n b^m \mid n < m\}$$

4.
$$L = \{a^n b^m \mid n \neq m\}$$

5.
$$L = \{a^n b^m c^{n+m} | n, m \ge 1\}$$

6.
$$L = \{a^n b^{2m} \mid n, m \ge 1\}$$



#Q. Which of the following are regular sets?

- 1. $\{a^n b^{2m} \mid n \ge 0, m \ge 0\}$
- 2. $\{a^n b^m \mid n = 2m\}$
- 3. $\{a^n b^m \mid n \neq m\}$
- 4. $\{x c y \mid x, y \in \{a, b\}^*\}$

A 1 and 4 only

C 1 only

B 1 and 3 only

D 4 only





Construct DFA for the following Language.

Instruct DFA for the following Language.

$$L = \{a^n b^m \mid n, m \ge 1\} \longrightarrow \text{Regular Language}$$

$$L = \{a^n b^n \mid n \ge 1\}$$

2.
$$L = \{a^n b^n \mid n \ge 1\}$$

3.
$$L = \{a^n b^m \mid n < m\}$$

4.
$$L = \{a^n b^m \mid n \neq m\} \longrightarrow \underline{\gamma \sigma \gamma} \quad \underline{\text{Regular}}$$

5.
$$L = \{a^n b^m c^{n+m} | n, m \ge 1\}$$

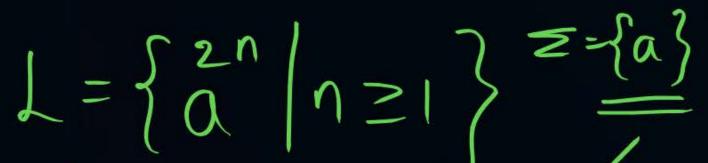
6.
$$L = \{a^n b^{2m} \mid n, m \ge 1\}$$





If comparision exist between symbols of language then DFA is not possible.







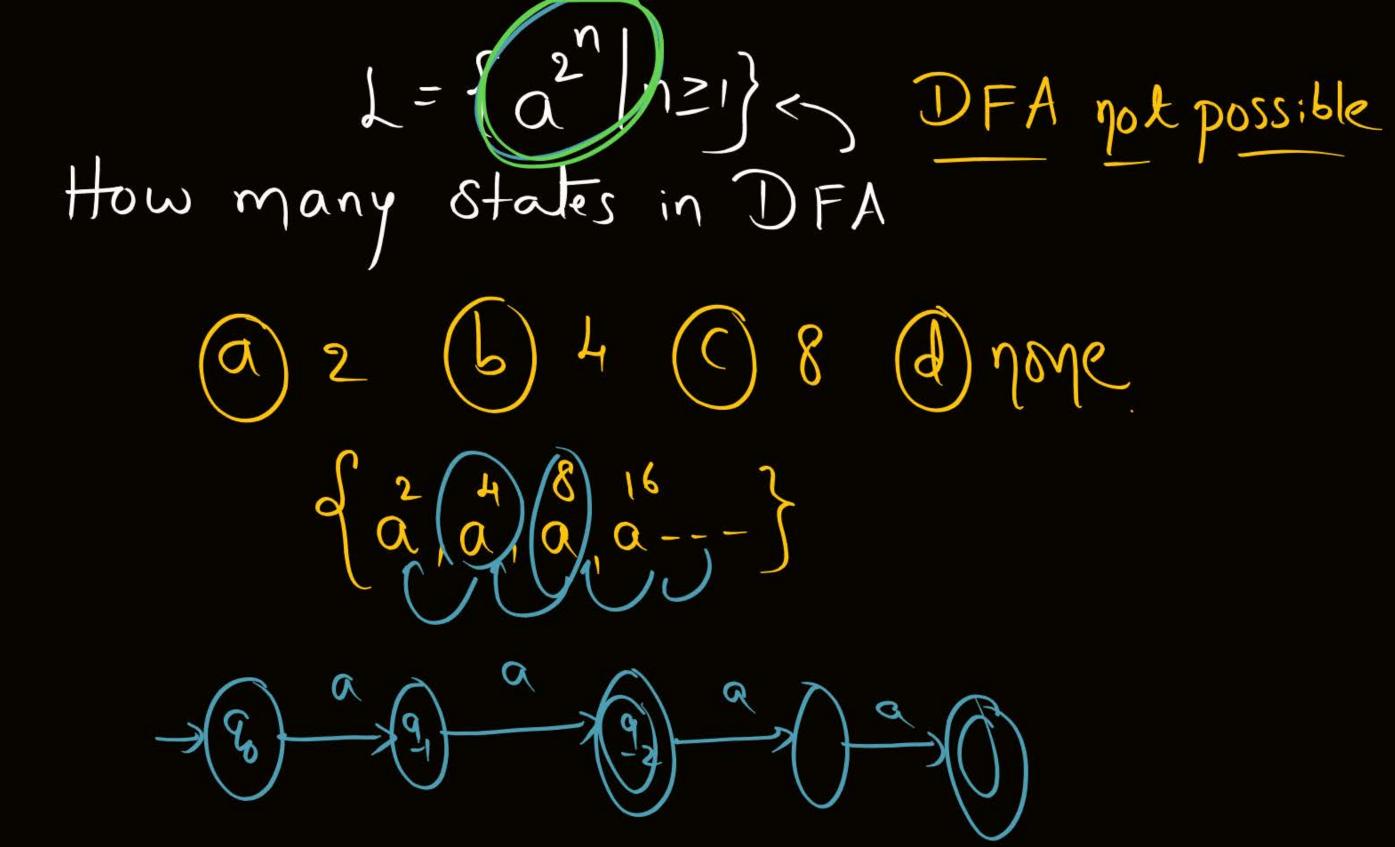
Construct DFA for the following Language.

$$\begin{array}{c|c}
 & 2 & 2 \\
\hline
 & 9 & 0 \\
\hline
 & 9 & 0
\end{array}$$





Construct DFA for the following Language. (a), a, a, a, ---)



2/3/6/8/10/12

which of the following in Regular?

not Reg

\[
\begin{align*}
\text{not} \\
\text{a} \\
\text{n} \\
\text{a} \\
\text{a} \\
\text{n} \\
\text{a} \\
\ $\frac{\text{Not Reg}}{\text{Col}} = \{ \alpha : | n \ge 1 \} = \{ \alpha, \alpha, \alpha, --- \}$ $\left(\frac{1}{2} \right) = \left\{ \begin{array}{c} 3 \\ 3 \end{array} \right\} = \left\{ \begin{array}{c} 3 \\ 3 \end{array} \right\} = \left\{ \begin{array}{c} 3 \\ 3 \end{array} \right\} = \left\{ \begin{array}{c} 3 \\ 3 \end{array} \right\}$ (d) Me

it a language y formed by 1 symbol Common différence exist -> DFA possible 70 11 11 DFA not
possible



















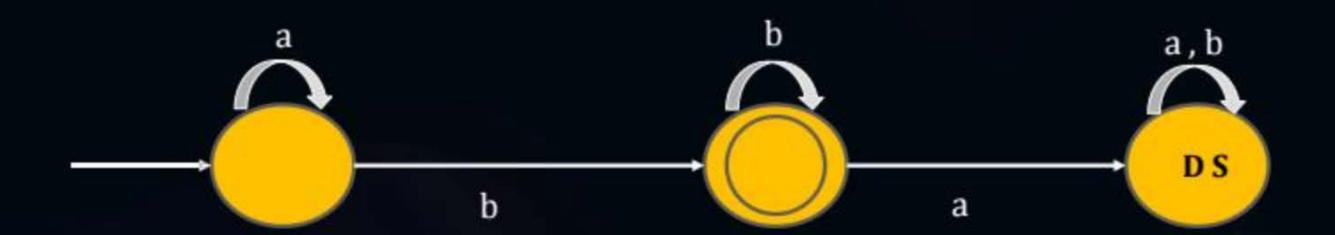








#Q. Identify language accepted by following DFA



- **A** $L = \{a^n b^m | n, m ≥ 1\}$
- C $L = \{a^n b^m | n, m \ge 0 \}$

- **B** $L = \{a^n b^m | n \ge 1, m \ge 0\}$
- D None



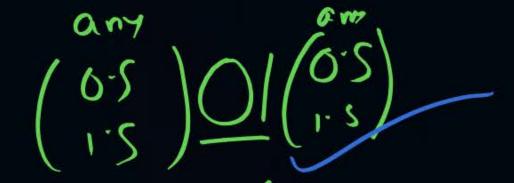
#Q. Identify language accepted by following DFA



- **A** $L = \{a^n b^m | n, m ≥ 1\}$
- C $L = \{a^n b^m | n, m \ge 0 \}$

- **B** $L = \{a^n b^m | n \ge 1, m \ge 0\}$
- D None

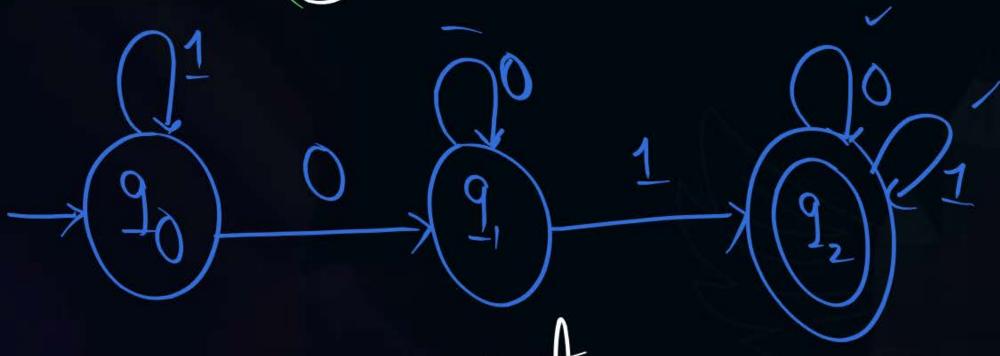






Construct minimal state DFA that accepts all strings os 0's and 1's where each

string contains substring 01.



3 States

Substring (DFA) 2 length Substring -> 3 11 - Ant States



Topic: DFA Construction (0) 3 (6) 4 (7) 5 (9) 6









Construct minimal state DFA that accerpts all strings os 0's and 1's where each string contains substring 001 7001,0001,1001,0011--3



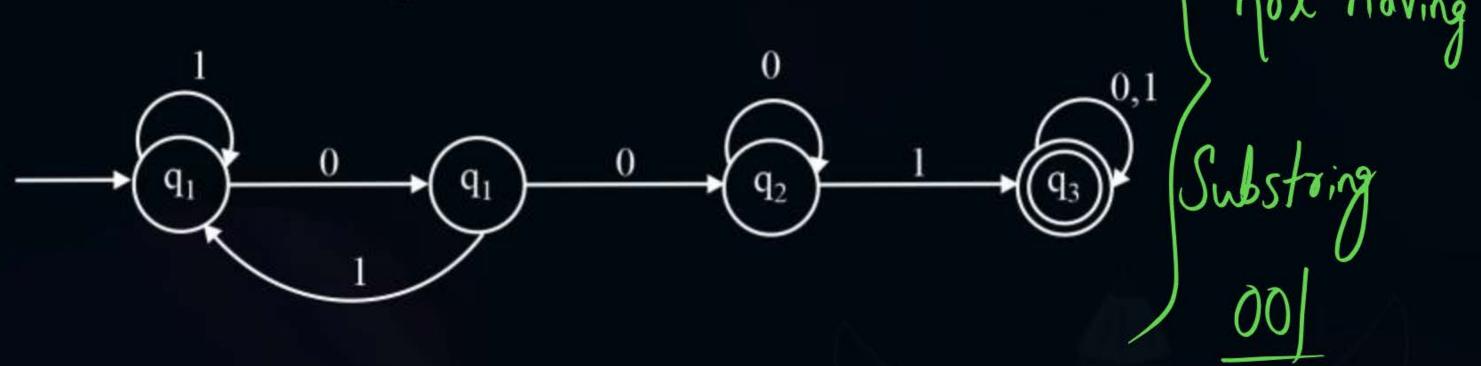


Construct minimal DFA that accerpts all strings os 0's and 1's where each string contains substring 0101

5 Stales

MCQ





The language accepted by finite automata, which is obtained by interchanging final and non-final states in m, is-



A The set of all strings containing 001 as the subtrings

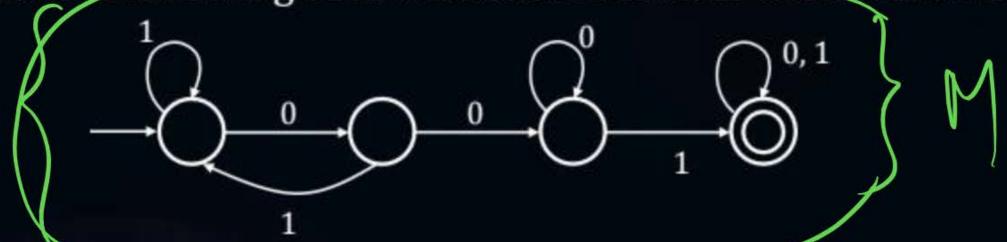
The set of all strings not containing 001 as the subtrings

The set of all strings ending with subtrings 001

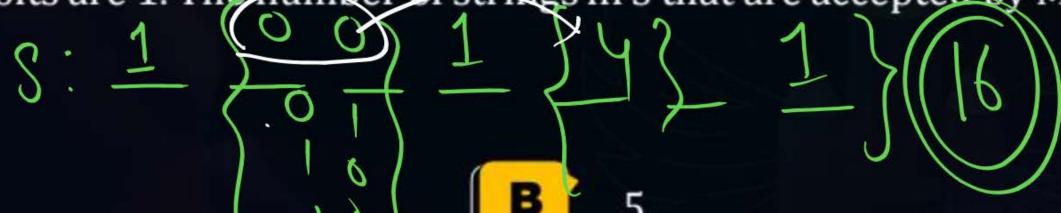
The set of all strings ending with subtrings 001



Consider the following deterministic finite state automaton M. #Q.



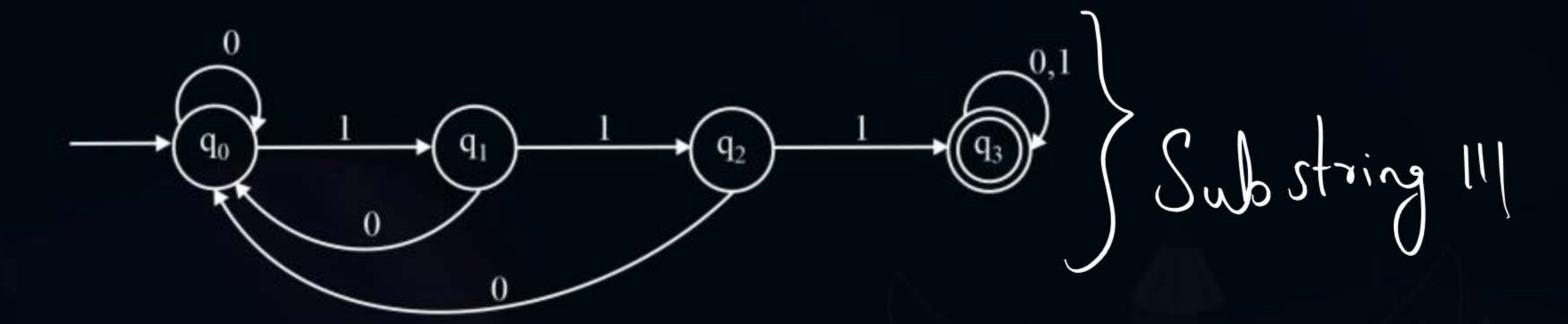
Let S denote the set of seven-bit binary strings in which the first the fourth and the last bits are 1. The number of strings in S that are accepted by M is



MCQ



#Q. Consider the finite automata m.



The language accepted by m is, over the alphabet {0,1}



Sub string

- The set of all strings containing three consecutive 's
- The set of all strings not containing three consecutive I's
- The set of all strings beginnigs with three consecutive l's
- The set of all strings ending with three consecutive I's

MCQ

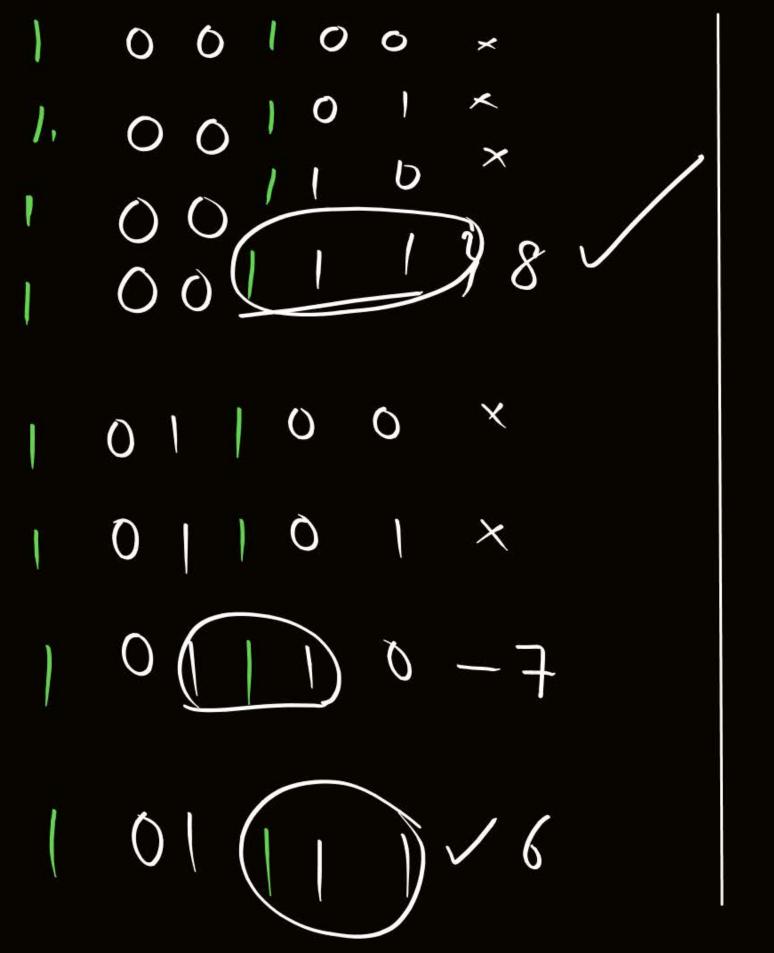


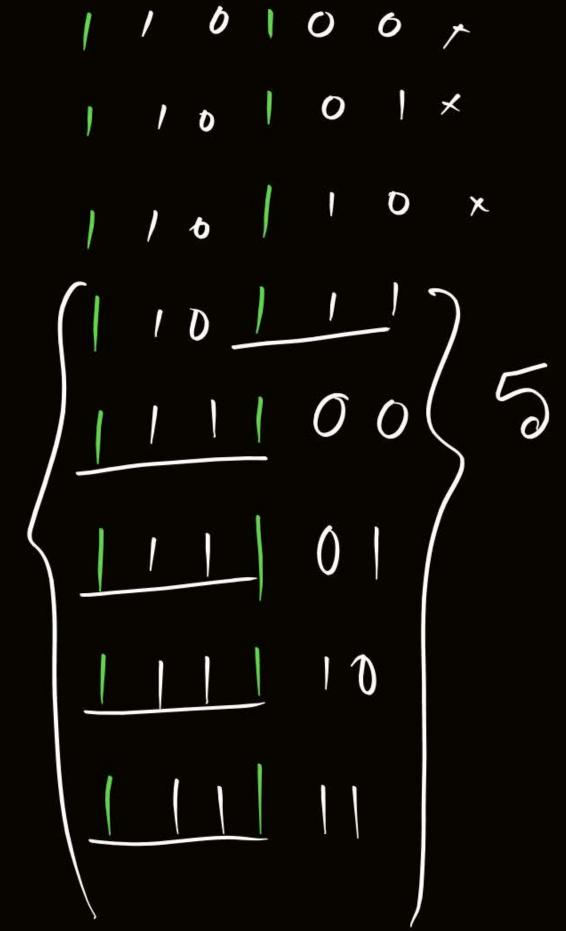
#Q. Consider the finite automata m.



Let S denotes the set of all six bit binary strings in which first and fourth bits are 1. The number of strings in S that are accepted by m is-











Construct DFA for

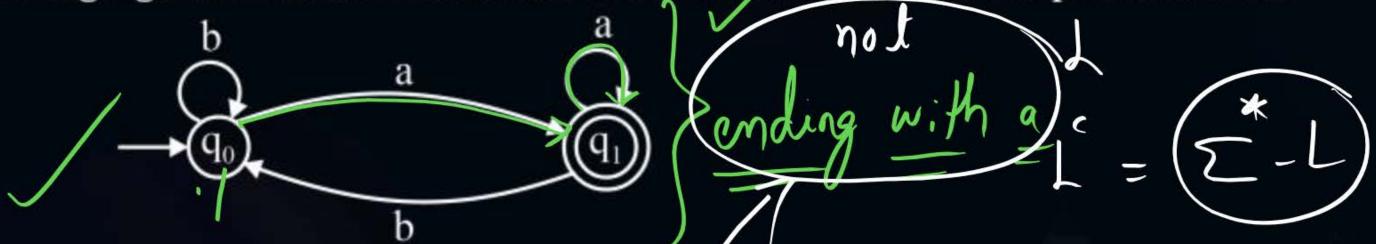


Topic: Complement of DFA

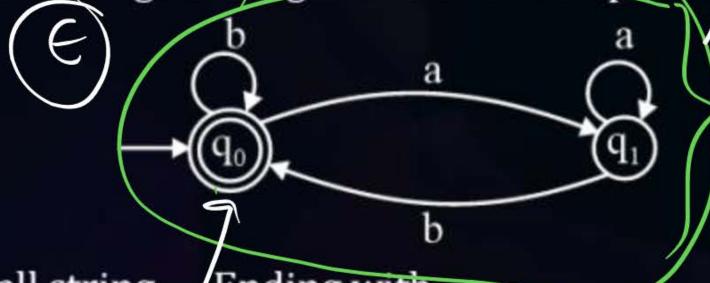




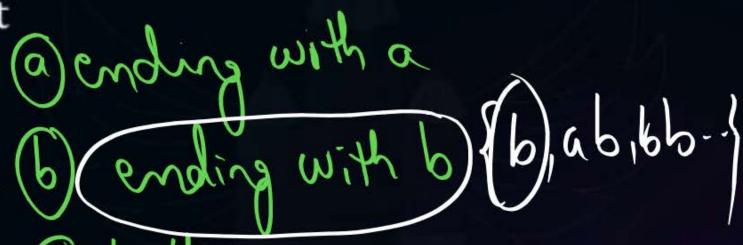
By interchanging final and non final states we can convert into complement DFA.



Set of all strings ending with a after complement



Set of all string ... Ending with



not having Sub string ool.









VENKAT SIR PW

members, 3 online

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Info

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Invite Link



Notifications









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ending with a



THANK - YOU