



CS & IT ENGINEERING

Theory of Computation

Regular Expression

DPP-02

Discussion Notes



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[NAT]

#Q. Let

$$L_1 = a^* b^*$$

$$L_2 = b^* a^*$$

$$L_3 = (a + b)^*$$

$$L_4 = a^* b^* a^*$$

$$L = (L_1 \cap L_2) - (L_3 \cup L_4)$$

Number of strings in above language L will be 0.

$$L = \underbrace{(a^* + b^*)}_{\downarrow} - \underbrace{(a + b)^*}_{\downarrow}$$

$$= 0$$

$$L = 0$$

$$L_1 = a^* (b^*)$$

$$L_2 = (b^*) a^*$$

$$L_1 \cap L_2 = a^* + b^*$$

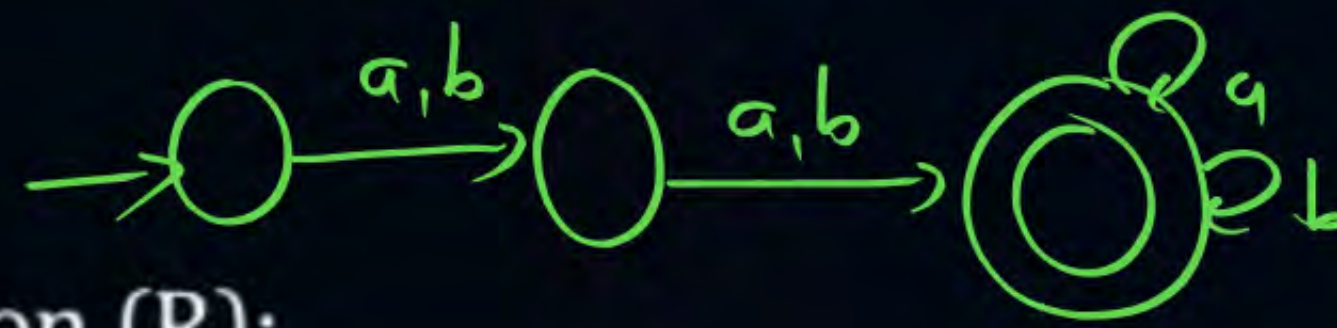
$$L_3 = (a + b)^*$$

$$L_4 = a^* b^* a^*$$

$$L_3 \cup L_4 = \underbrace{(a + b)^*}_{\downarrow} \cup \underline{a^* b^* a^*}$$

$$= \underline{(a + b)^*}$$

[MCQ]



#Q. Consider a regular expression (R):

$$R = (a + b)^* (a + b)^2 (a + b)^*.$$

How many equivalences classes are existing for above regular expression R? $(n+1)$

min DFA States

A 2

~~**B** 3~~

C 4

D None

No. of equivalence classes = No. of states of min DFA

[MCQ]

#Q. Let L be any formal language. If L^* is regular language then what is L ?

☐ A L is regular.

☐ B L is non-regular.

☐ C L is CFL.

☒ D None of these.

$$L^* = \text{regular} \quad L^* = (a+b)^* \checkmark$$

$$L = \text{regular} = (a+b) \rightarrow \text{reg}$$

L
 \downarrow
 may (or) may not be
 Regular

$$L = \{a^p \mid p \text{ is a prime number}\} \rightarrow \text{Non regular}$$

$$L^* = \{(a^p)^*\} \rightarrow \text{regular}$$

[MCQ]

$$L_1 = (a+b)^* \cup L_2$$

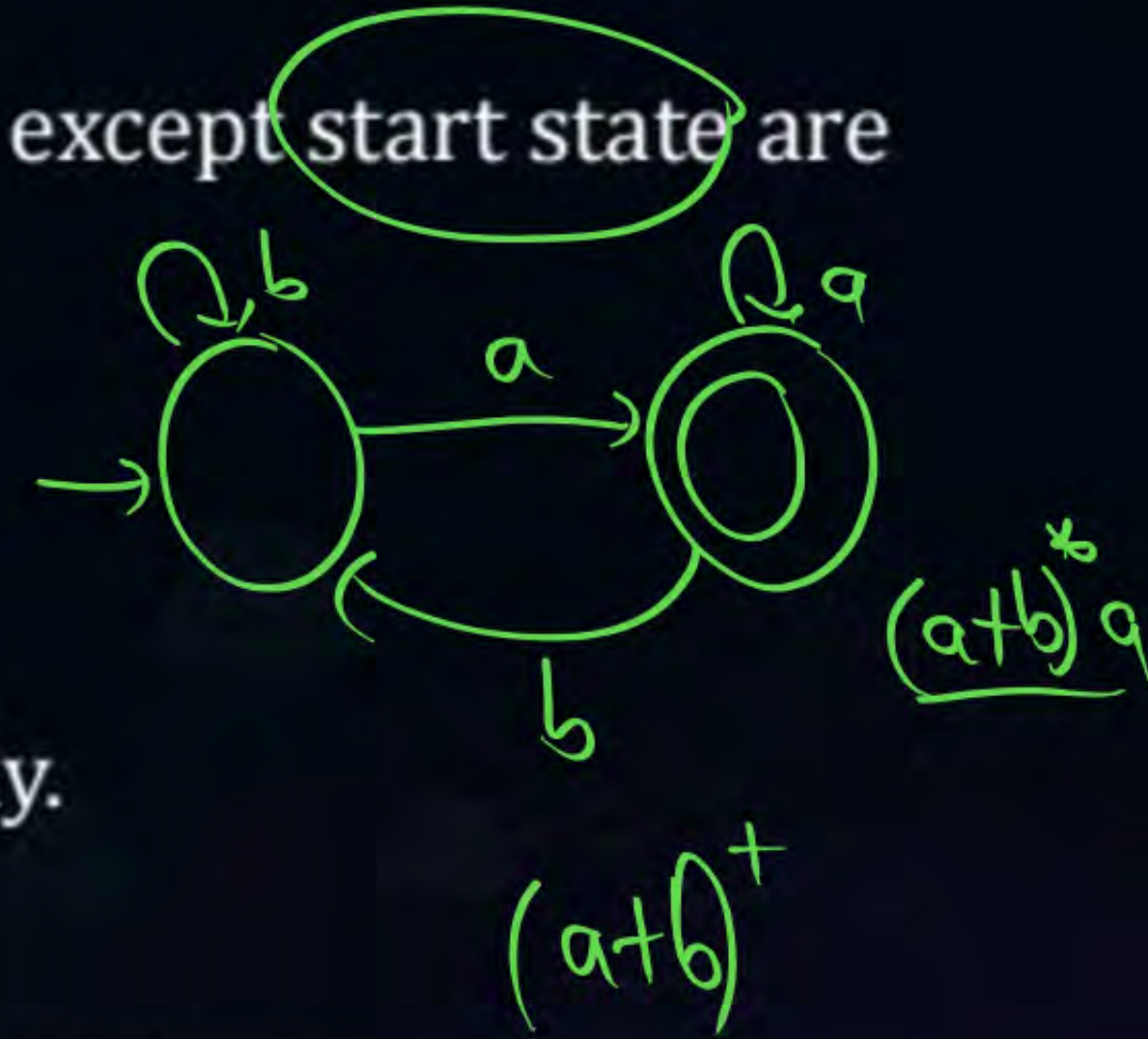


#Q. Consider the following two statements:

✓ (I). There exist a regular language L_1 , such that for all language L_2 , $L_1 \cup L_2$ is always regular.

✗ (II): If all states of deterministic finite automata (DFA) except start state are final states then language accepted by DFA is Σ^+ .

Which of the following is correct?



✓ A S_1 only.

B S_2 only.

C Both S_1 and S_2 are true.

D None of these.

[MCQ]

#Q. Consider the language L given by the regular expression $(a + b)^* ab(a + b)^*$ over the alphabet $\{a, b\}$. What is the correct regular expression of \overline{L} ?

☒ A $(a + b)^* (ab + ba + bb + aa) + \epsilon$

☒ B $(a^* b^*)^* (ba + bb + aa) (a^* b^*)^* + a + b$

☒ C $(a + b)^* ba(a + b)^* + a + b$

☒ D $b^* a^*$

$$L = (a+b)^* ab(a+b)^*$$

Sub string ab

$L^c =$ not having substring ab

[MSQ]

#Q. For language $L = \{\text{Every odd bit is } a\}$ On alphabet $\Sigma = \{a, b\}$. Which of the following is/are correct regular expression?

aa a a a a a a a b a b a b a

_{1 3 5 7}

☒ A $(\underline{aa} + \underline{ab})^* (\epsilon + a)$

☐ B $(aa + ab + \text{ba} + b)^* a$

☐ C $(aa + ba)^* (\epsilon + a + \text{b})$

☐ D $(a(a + b))^* + (a(a + b))^* a$

ba a

a b b b b b b

(A, D)

L,

[NAT]

No. of states of min DFA.

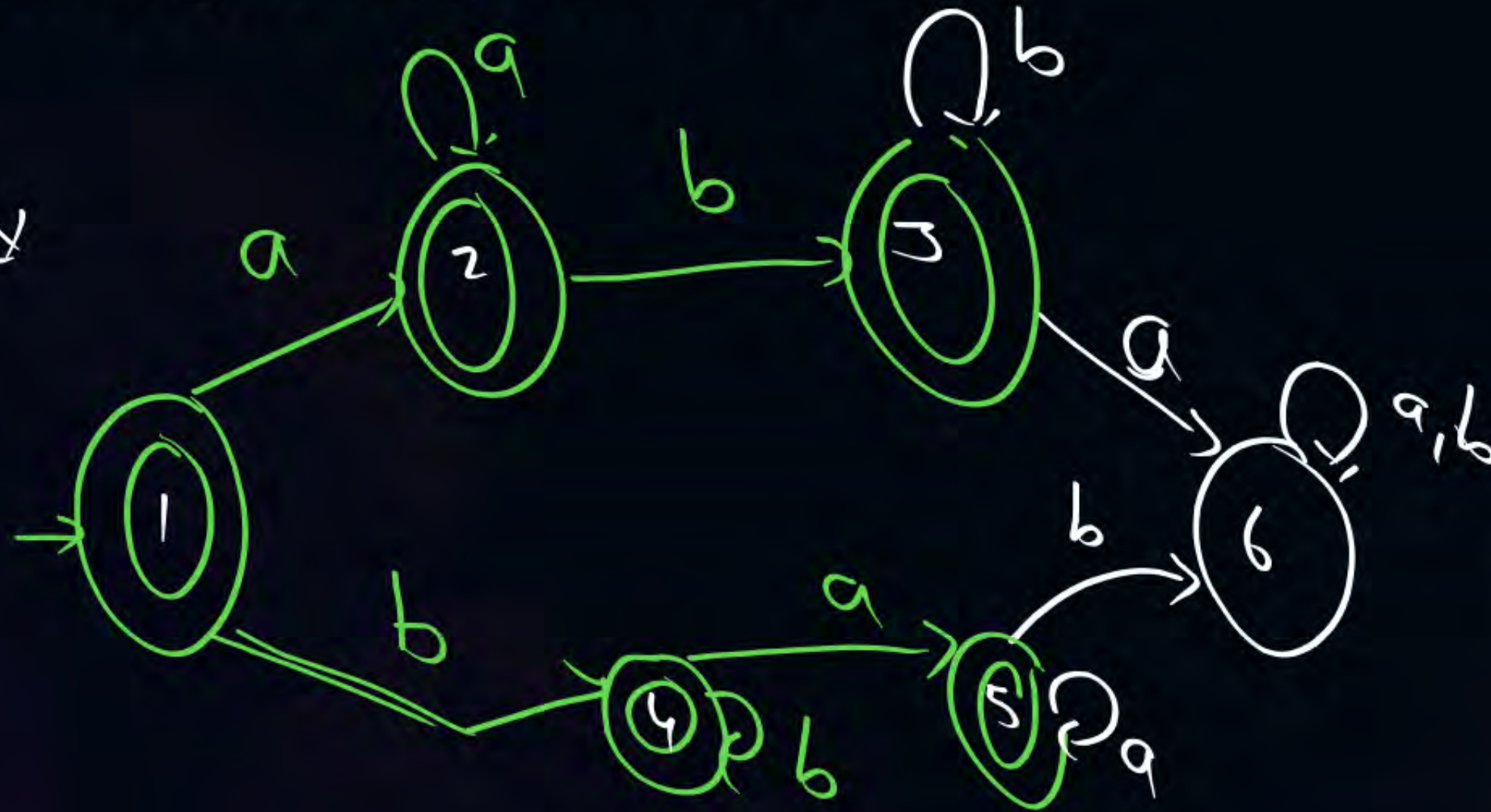
#Q.

Let us consider the following regular expression

$$R = a^*b^* + b^*a^*$$

Ans: 6 ✓

How many equivalence classes of expression that represent language are equivalent to regular expression R?



$$L_2 \subseteq L_1$$

$$L_3 \subseteq L_2$$

#Q. Consider the following languages:

$$L_1 = \{a^m b^n c^p \mid m, n, p \geq 0\}.$$

$$L_2 = \{a^m b^m c^p \mid m, p \geq 0\}.$$

$$L_3 = \{a^{2m} b^{2m} c^p \mid m, p \geq 0\}.$$

Which of the following is/are correct?

~~(B, C)~~

~~A~~ $L_1 \subseteq L_2$ and $L_2 \subseteq L_1$

\checkmark C $L_3 \subseteq L_2$ and $L_2 \subseteq L_1$

\checkmark B $L_2 \subseteq L_1$ and $L_3 \subseteq L_1$

~~D~~ $L_2 \subseteq L_3$ and $L_3 \subseteq L_1$



THANK - YOU