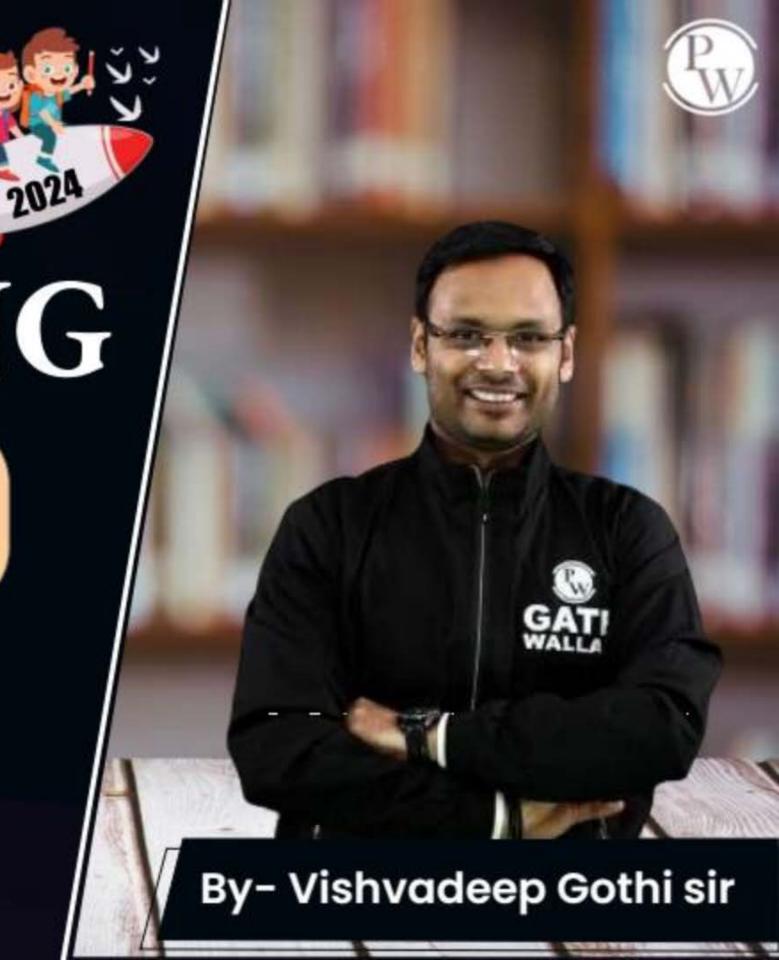
# CS & IT ENGING

**Operating System** 

Virtual Memory



Lecture - 05

# **Recap of Previous Lecture**









Topic Frame Allocation

Topic Thrashing

Topic Multilevel Paging

# **Topics to be Covered**











Topic

**Multilevel Paging** 

Topic

**Access Time in Multilevel Paging** 



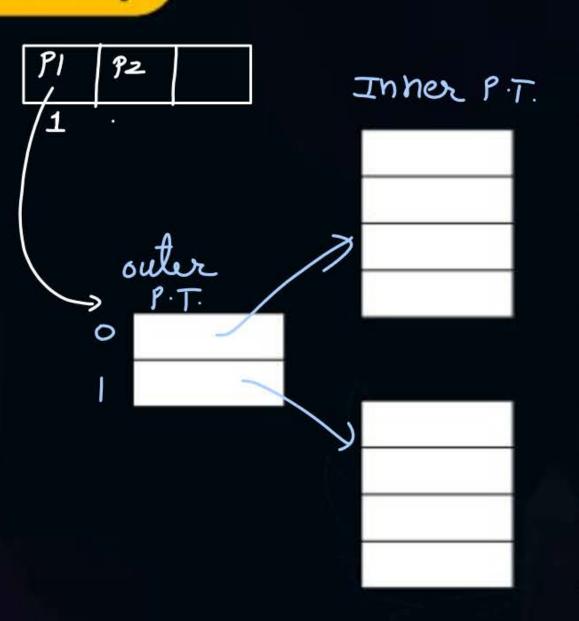
Page 111

# **Topic: Page Table in Memory**



Process	Page Table	
Page 000	000	
Page 001	001	
Page 010	010	
Page 011	011	
Page 100	100	
Page 101	101	
Page 110	110	

111

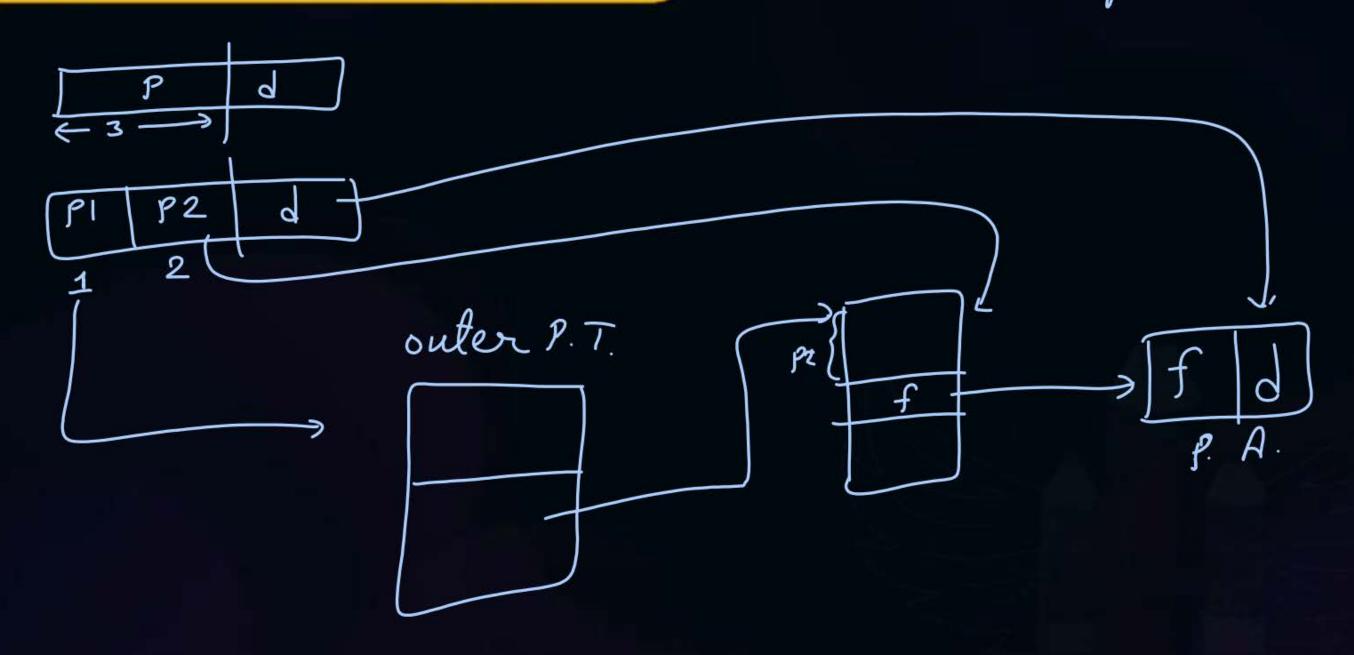




# **Topic: Multilevel Paging**

PTBR stores add of outer P.T.





P.T. Middle level P.T. outermost

outermost P.T. must be stored in single page.

If any level P.T. Cannot fit into a single Page, then level must be increased.





VA = 15 bits

Page size = 128 bytes =  $2^7 B \Rightarrow d = 7$  bits

Page table entry size = 2 bytes

## Number of levels in multilevel page 2.



VA

= 24 bits

Page size

= 1Kbytes =  $2^{10}$ B = 3 = 10 bits

Page table entry size = 4 bytes

### Number of levels in multilevel page 2. ?

no. of P.T. entries = 
$$\frac{2^{10} B}{2^{2} B} = \frac{8}{2}$$
  
stored in page



VA = 29 bits

Page size =  $2Kbytes = 2B \Rightarrow J = 11 bits$ 

Page table entry size = 4 bytes

### Number of levels in multilevel page 2. ?

no. of P.T. entries = 
$$\frac{2^{11}B}{4B} = 2^9$$
per page





VA = 40 bits

Page size =  $4Kbytes = 2^{12}B \Rightarrow d = 12 bits$ 

Page table entry size = 8 bytes

Number of levels in multilevel page 4 ?

no. of P.T.E. Per page = 
$$\frac{2^{12}}{2^3} = 2^9$$







#Q. Consider a virtual memory system with physical memory of 8GB, a page size of 8KB and 46-bit virtual address. Assume every page table exactly fits into a single page. If page table entry size is 4B then how many levels of page tables would be required.

no of P.T. entries Per page = 
$$\frac{2^{13}B}{2^2B} = 2^{11}$$



# GATE-2023



#Q. Consider a computer system with 57-bit virtual addressing using multi-level tree-structured page tables with L levels for virtual to physical address translation. The page size is 4KB(1KB=1024B) and a page table entry at any of the levels occupies 8 bytes.

The value of L is \_\_\_\_\_?

No. of PTE per page = 
$$\frac{4kB}{8B} = \frac{2^{12}}{2^3} = 2^9$$



$$VA = 26 \text{ bits}$$

Page table entry size = 4 bytes

2-level paging

Outer page table fits into a page exactly Page size = <u>1</u> Kbytes?

no of P.T.E. perpage = 
$$\frac{2^{x}B}{2^{2}B}$$

$$= 2^{x-2}$$





$$VA = 35 \text{ bits}$$

Page table entry size \equiv 8 bytes



3-level paging

Outer page table fits into a page exactly

Page size = 2 Kbytes?

no. of P.T.E. per page
$$= \frac{2^{x}B}{2^{3}B} = 2$$

$$3(x-3) + x = 35$$
  
 $4x - 9 = 75$   
 $x = 11$ 



# Topic: Question ATE - 2013



#Q. A computer uses 46-bit virtual address, 32-bit physical address, and a threelevel paged page table organization. The page table base register stores the base address of the first-level table (T1), which occupies exactly one page. Each entry of T1 stores the base address of a page of the second-level table (T2). Each entry of T2 stores the base address of a page of the third-level table (T3). Each entry of T3 stores a page table entry (PTE). The PTE is 32 bits in size. What is the size of a page in KB in this computer?

< L	A. 46 —	>	
PI PZ	Ps d	[	3(x-2)+x=466
x-2 x-2			3x - 6 + x = 46
			x = 13

P1 | P2 | d

Page 5:3e = 4 Bytes

no of pages in outer most p T = 1 no of -11 - Inner -11 - = 2 = 2 = 2

P1

3 pages needed to store p.T. across all levels

P.T. Size across all levels = 3 \* 4B = 12 Bytes





### Size of page tables across all levels?

		Page - 2 B
P1	P2	$ = 1 \times 3 $
1	8	$\int_{10}^{2} \int_{10}^{2} \int_{10}^{2$
) /		noutermost $P.T. = 1$ Inner $P.T. = 2^{p_1} = 2^1 = 2$
10	tal no. of	pages = 3





Size of page tables across all levels?

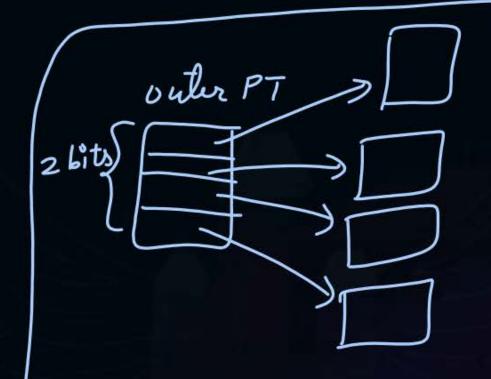
P1	P2	D
2	9	(12)

No. of pages to store outer P.T. = 1

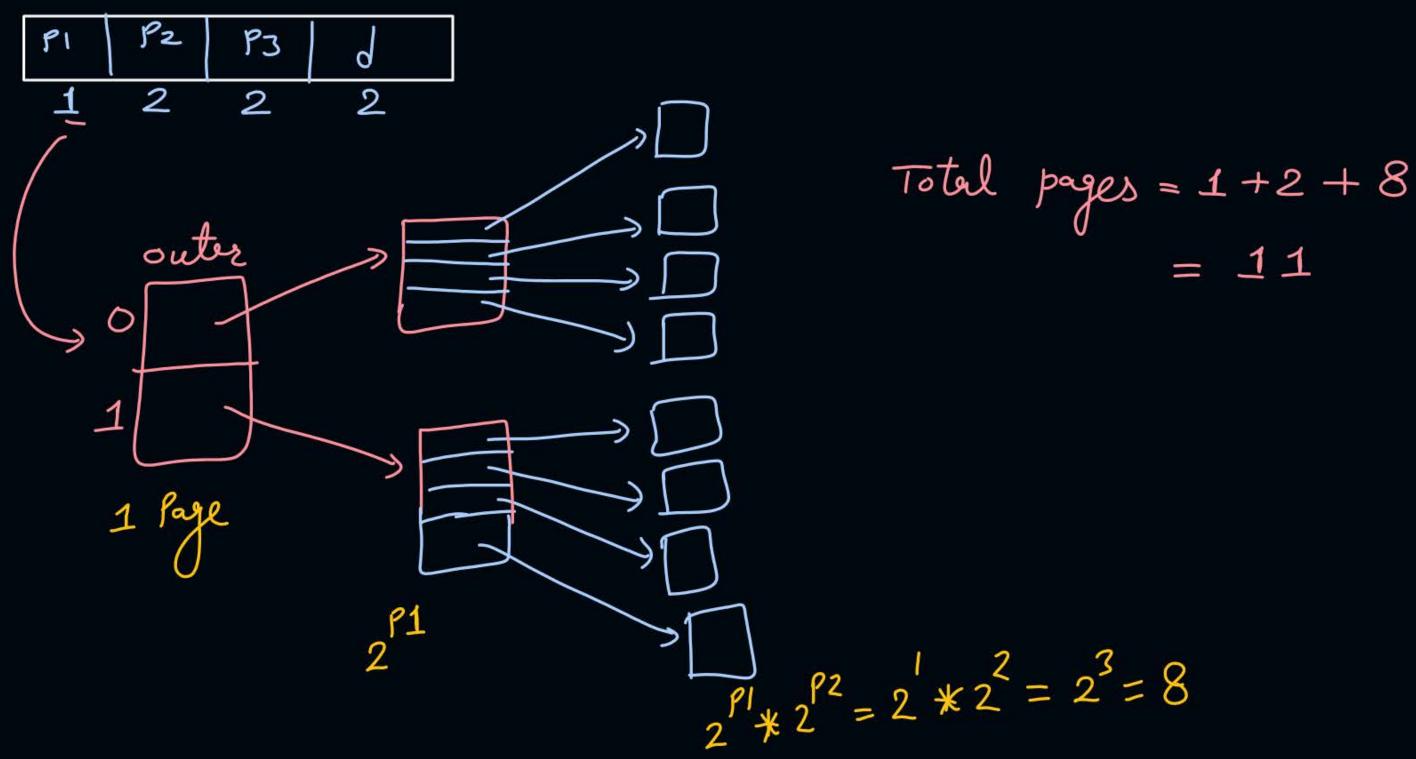
Inner P.T. = 2<sup>2</sup> = 4

Total = 5

P.T. size across all level = 5 \* 4KB = 20KB



<u>ex:</u>-







Size of page tables across all levels?

P1	P2	Р3	D
1	8	8	(10)

> 1 Page Size = 2 B = 1 kB

#Q.

	· · ·		1	Pag	e s	ge <i>=</i>	2	B =	2kB
PI	PZ	P3	1 9	ه و مو					
2	10	10	11						

no. of pages to store outermost PT = 1

11 \_\_\_\_\_ middle -11 - =  $2^2 = 4$ Total no. of pages = 4101

Total no. of pages = 4101

Total P.T. size across all levels = 4101 \* 2 kB = 8202 KB





#Q. Size of page Consider a three-level page table to translate a 39-bit virtual address to a physical address as shown below:

<b>~</b>	39-bits Virtual Address					
Level-1 Offset	Level-2 Offset	Level-3 Offset	Page Offset			
9 bits	9 bits	9 bits	12 bits			

The page size is  $4 \text{ KB} = (1 \text{KB} = 2^{10} \text{ bytes})$  and page table entry size at every level is 8 bytes. A process P is currently using 2 GB (1 GB =  $2^{30}$  bytes) virtual memory which OS mapped to 2 GB of physical memory. The minimum amount of memory required for the page table of P across all levels is \_\_\_\_\_ KB across all levels?







#Q. Consider a 32-bit system with 4 KB page size and page table entries of size 4 bytes each. Assume 1 KB =  $2^{10}$  bytes. The OS uses a 2-level page table for memory management, with the page table containing an outer page directory and an inner page table. The OS allocates a page for the outer page directory upon process creation. The OS uses demand paging when allocating memory for the inner page table, i.e. a page of the inner page table is allocated only if it contains at least one valid page table entry.

An active process in this system accesses 2000 unique pages during its execution, and none of the pages are swapped out to disk. After it completes the page accesses, let X denote the minimum and Y denote the maximum number of pages across the two levels of the page table of the process.

The value of X + Y is \_\_\_\_\_?



# 2 mins Summary



Topic

**Multilevel Paging** 

Topic

**Access Time in Multilevel Paging** 





# Happy Learning

THANK - YOU