

# COMPUTER SCIENCE & IT

## DIGITAL LOGIC




Lecture No. 09

Combinational Circuit



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# Recap of Previous Lecture

Comparator CKt







# Topics to be Covered

MUX



	$\bar{B}$	$B$
$\bar{A}$	1	1
$A$	1	1

$$f(A, B) = 1 = \Sigma(0, 1, 2, 3)$$

$$\begin{aligned}
 &= \bar{A}\bar{B} + \bar{A}B + A\bar{B} + AB \\
 &= \bar{A} + A \\
 &= 1
 \end{aligned}$$

Q.  $f_1(A, B, C) = \Sigma(2, 6, 7)$

$$f_2(A, B, C) = \Pi(4, 5, 6, 7) = \Sigma(0, 1, 2, 3)$$

$$f_3(A, B, C) = ?$$

$$\begin{array}{c} f_1 \\ f_2 \\ f_3 \end{array} \Rightarrow f = (f_1 + f_2 + f_3) = \Sigma(0, 1, 2, 3, 4, 5, 6, 7)$$

maximum

If  $f$  is logic '1', then number of possible maximum in function  $f_3$ .

$$f_4 = f_1 + f_2 = \Sigma(0, 1, 2, 3, 6, 7)$$

- a. 8
- ☒ b. 6
- c. 2
- d. 5

$$\begin{aligned}
 f_3 &= \Sigma(4, 5) + \Sigma(0, 1, 2, 3, 6, 7) \\
 (f_3) &= \Sigma(4, 5) \quad \downarrow \quad \downarrow \quad \downarrow \\
 &\quad \text{minimum} \quad \text{necessary} \quad \text{any of them} \\
 &\quad \text{SOP} \quad \quad \quad \text{is optional} \\
 f_3 &= \Pi(0, 1, 2, 3, 6, 7) \rightarrow \text{maxim} \\
 &\quad \quad \quad \text{max term}
 \end{aligned}$$

# [ MUX ]



- What is MUX ?

It is a combinational ckt having many i/p lines and one output line & on the basis of select line one of the i/p line is transferred to the o/p line.



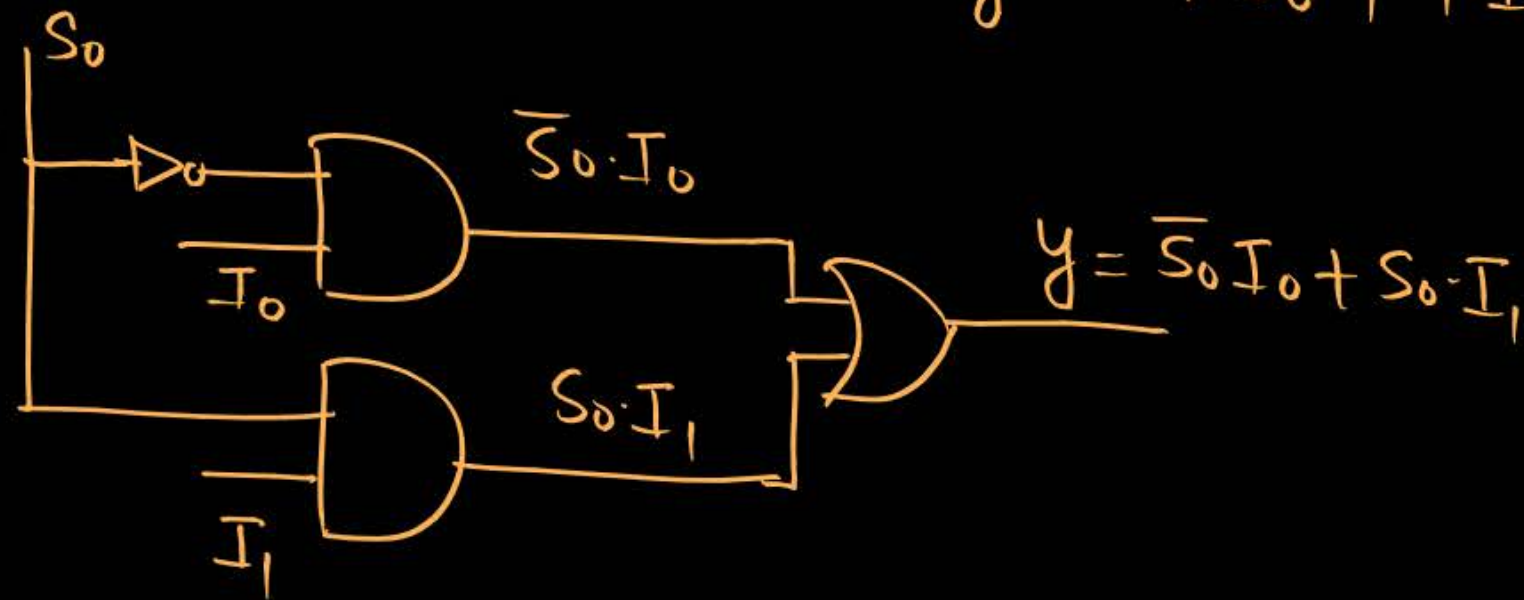
- 2:1 MUX



$$y = \bar{S}_0 \cdot I_0 + S_0 \cdot I_1$$

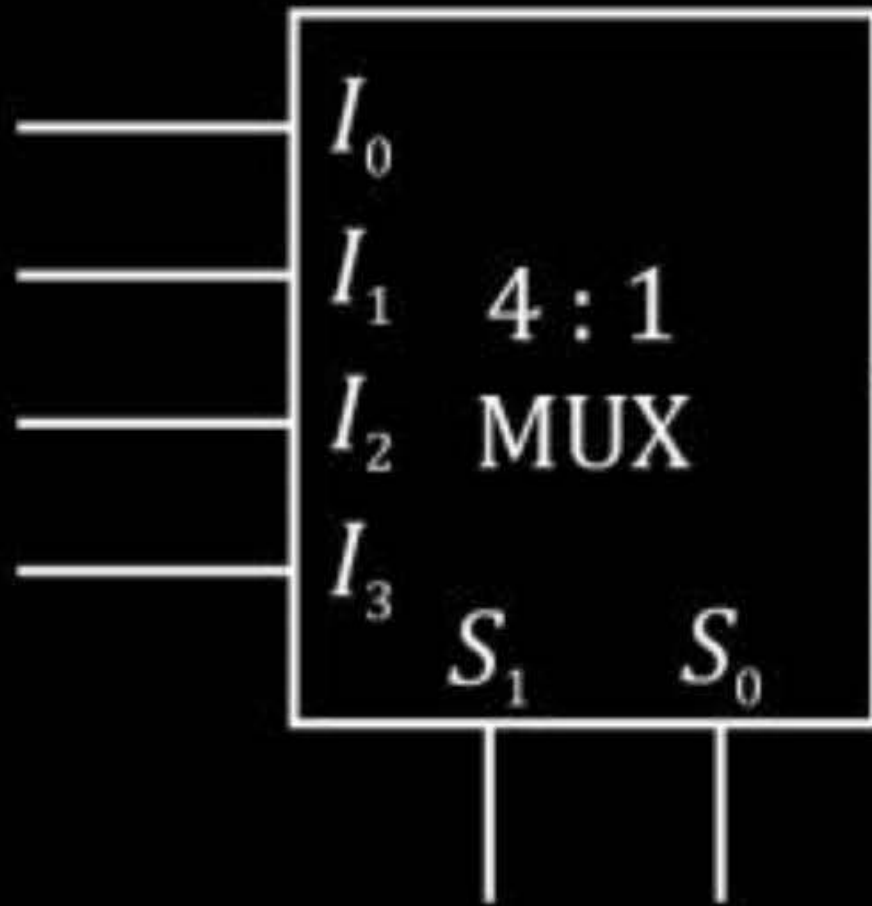
$$S_0 = 0 \Rightarrow y = 1 \cdot I_0 + 0 \cdot I_1 = I_0$$

$$S_0 = 1 \Rightarrow y = 0 \cdot I_0 + 1 \cdot I_1 = I_1$$



$$N \longrightarrow 2^N$$

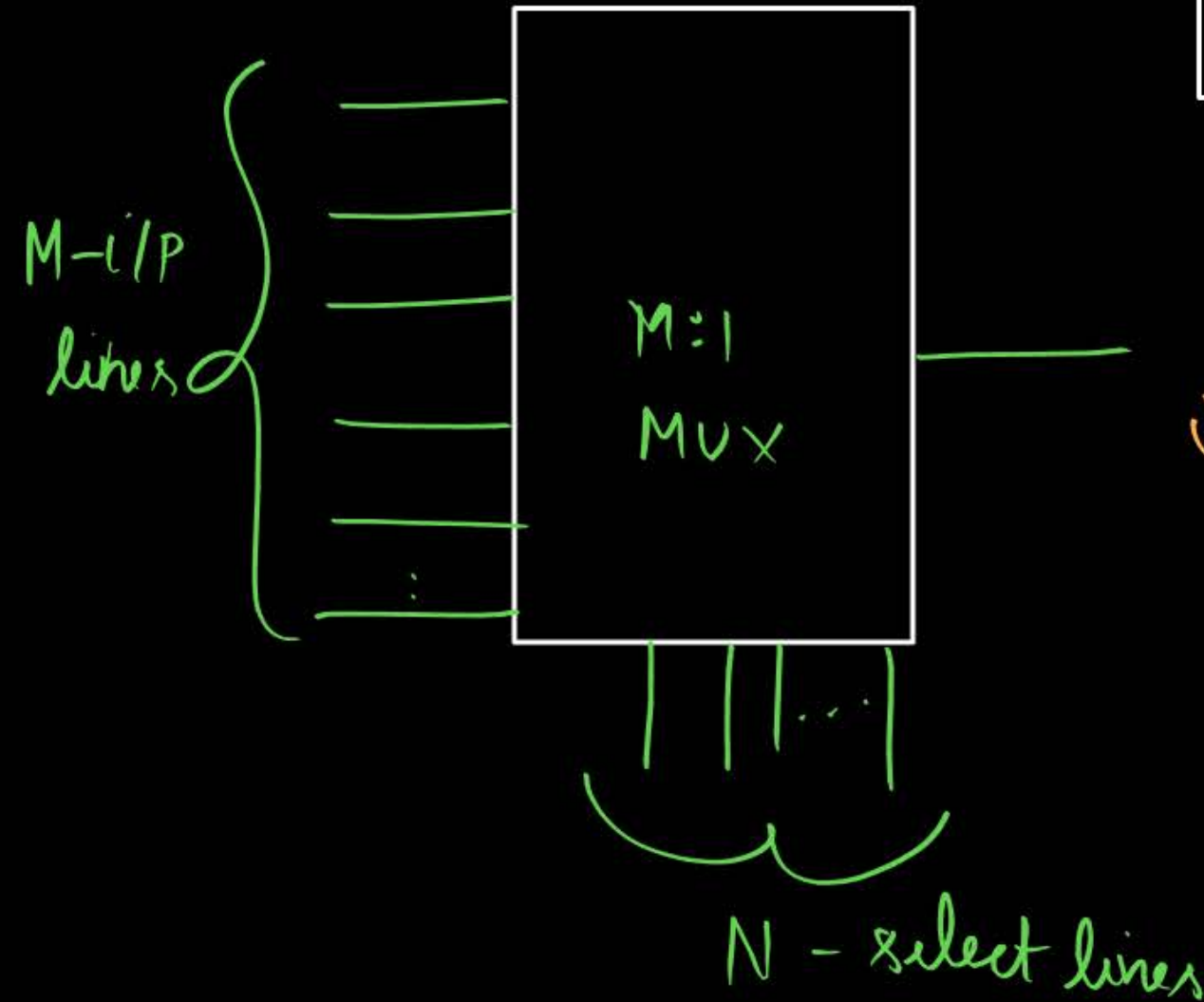
- 4:1 MUX



$$y = \overline{S_1} \overline{S_0} \cdot I_0 + \overline{S_1} S_0 \cdot I_1 + S_1 \overline{S_0} \cdot I_2 + S_1 S_0 \cdot I_3$$

$$S_1=1, S_0=0, \quad y = 0 + 0 + I_2 + 0 = I_2$$

- Relation between number of input lines  $M$  and number of select lines  $N$  :



$$2^N \geq M$$

$$\log_2 2^N \geq \log_2 M$$

$$N \geq \log_2 M$$



# [ Higher order MUX using lower order MUX ]

- $4:1$   $\xrightarrow{\text{Using } 2:1 \text{ MUX}}$   $\frac{4}{2} + \frac{2}{2} = 3 = 2^2 - 1$   
 $2^2:1$
- $8:1 \text{ MUX}$   $\xrightarrow{\text{Using } 2:1 \text{ MUX}}$   $4 + 2 + 1 = 7 = 2^3 - 1$   
 $2^3:1$
- $16:1 \text{ MUX}$   $\xrightarrow{\text{Using } 2:1 \text{ MUX}}$   $8 + 4 + 2 + 1 = 15 = (2^4 - 1)$   
 $2^4:1$
- $2^n:1 \text{ MUX}$   $\xrightarrow{\text{Using } 2:1 \text{ MUX}}$   $(2^n - 1)$

Implement 8 : 1 MUX using 2 : 1 MUX and design the complete circuit.

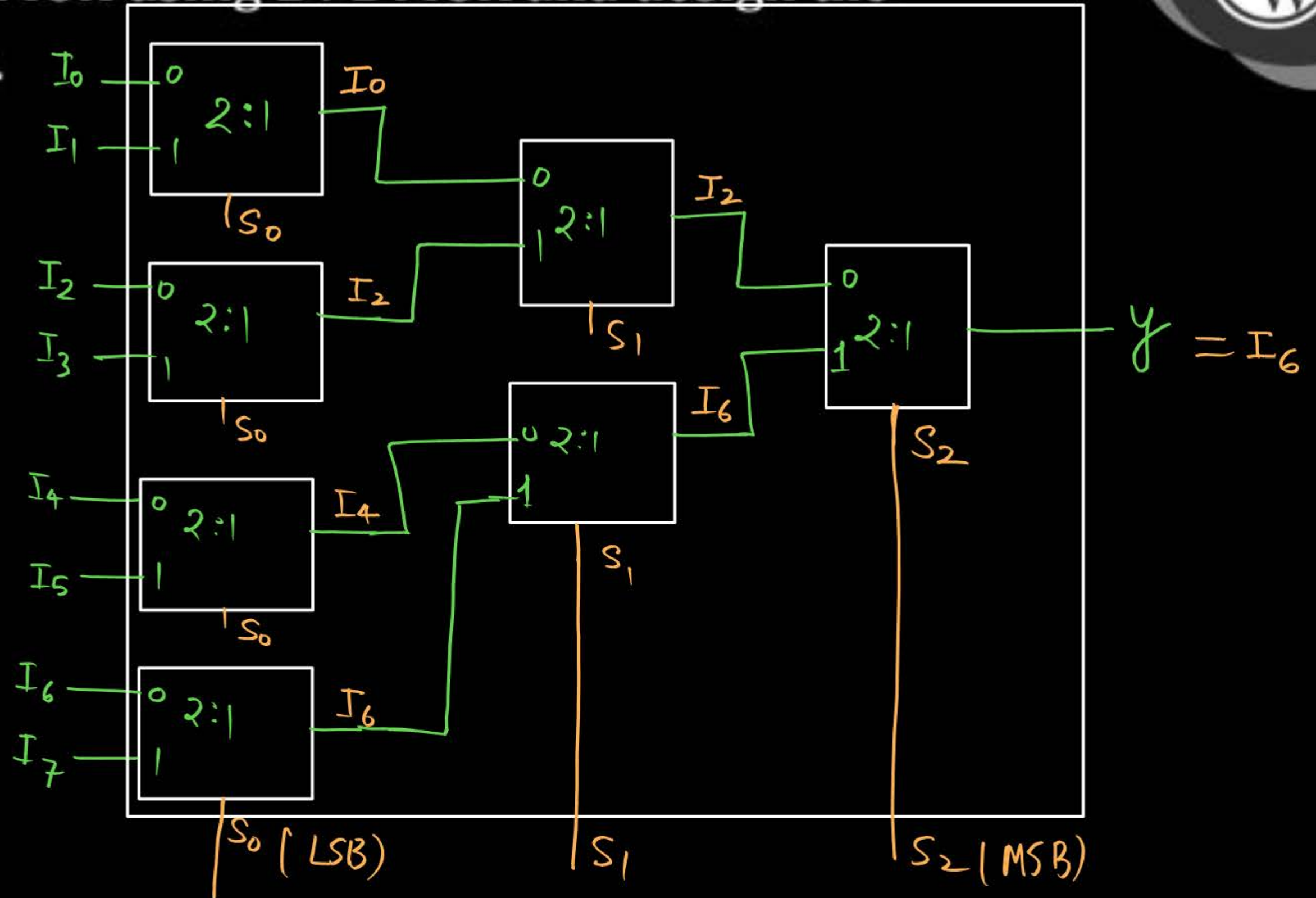
8:1

→ 8 i/p lines  $(I_0 - I_7)$

→  $S_2, S_1, S_0$

$8:1 \xrightarrow{2:1} 4+2+1$

$S_2 S_1 S_0 = 110$



- 16 : 1 MUX  $\xrightarrow{\text{Using 4 : 1 MUX}}$  4 + 1
- 64 : 1 MUX  $\xrightarrow{\text{Using 4 : 1 MUX}}$  16 + 4 + 1
- 64 : 1 MUX  $\xrightarrow{\text{Using 8 : 1 MUX}}$  8 + 1
- 256 : 1 MUX  $\xrightarrow{\text{Using 16 : 1 MUX}}$  16 + 1



$$32:1 \xrightarrow{4:1} (8+2) 4:1 \text{ MUX} + 1(2:1) \text{ MUX} = 11(4:1) \text{ MUX}$$

but 1(4:1) MUX will be used as (2:1) MUX

$$128:1 \xrightarrow{8:1} (16+2) 8:1 \text{ MUX} + 1(2:1) \text{ MUX} \rightarrow 19(8:1) \text{ MUX}$$

but 1(8:1) MUX will be used as (2:1) MUX

$$256:1 \xrightarrow{8:1} \underbrace{(32+4)}_{36(8:1)} 8:1 \text{ MUX} + \underbrace{1(4:1) \text{ MUX}}_{\substack{\downarrow \\ 1(8:1) \text{ MUX is used as } (4:1) \text{ MUX}}}$$


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37(8:1) MUX

$$32:1 \xrightarrow{4:1}$$

└→ 32-i/p → lines  
└→  $S_4, S_3, S_2, S_1, S_0$

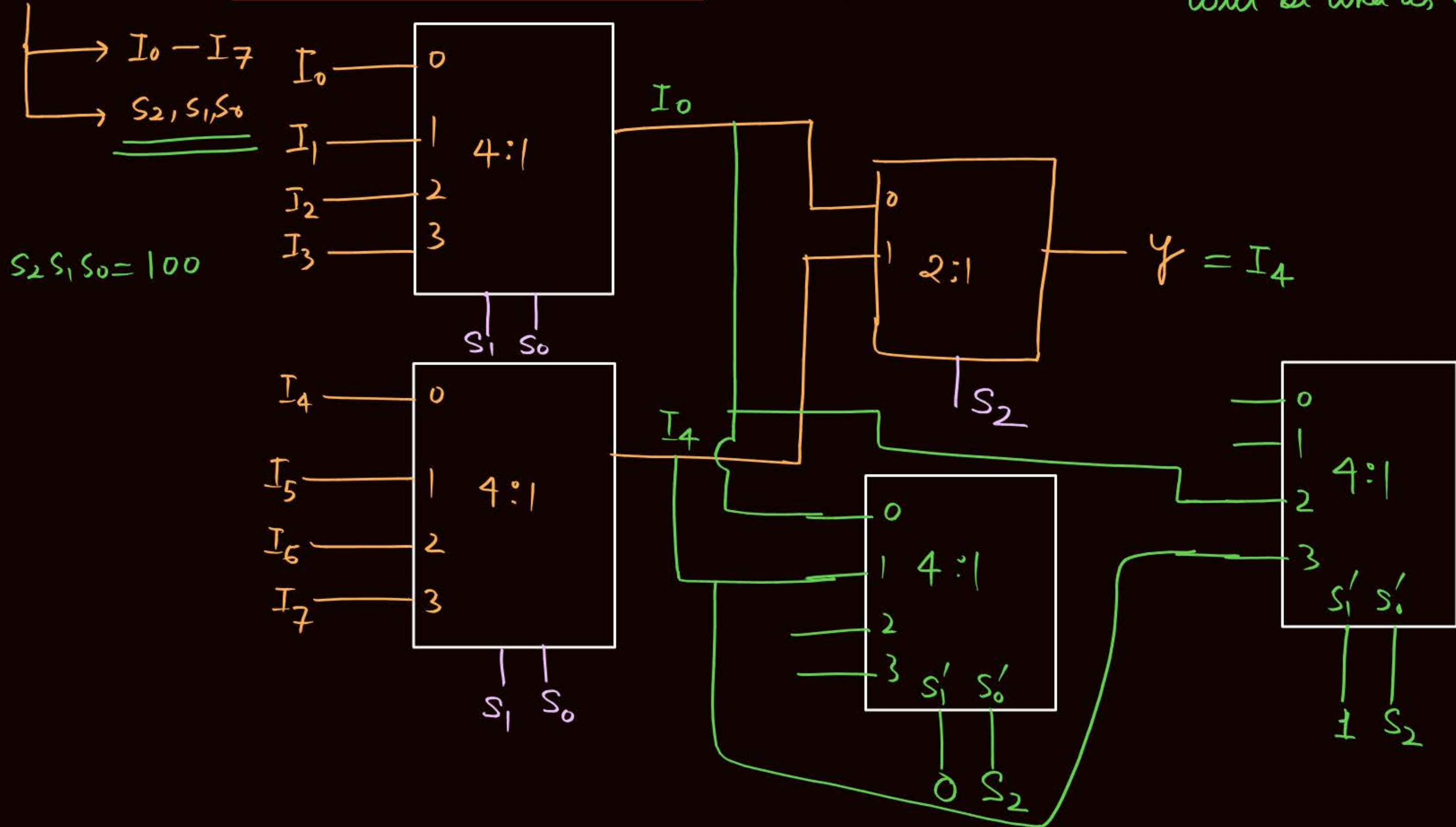
$$\boxed{8+2+1(2:1)}$$

4:1  
MUX

$$10(4:1) + \underline{1(4:1) \text{ MUX}}$$

will be used as  
2:1 MUX.

$8:1 \xrightarrow{4:1} 2(4:1)\text{MUX} + 1(2:1)\text{MUX} \Rightarrow (2+1) 4:1 \text{ MUX}$  but the last 4:1 MUX will be used as 2:1 MUX





# [ Implementation of Boolean function using MUX ]

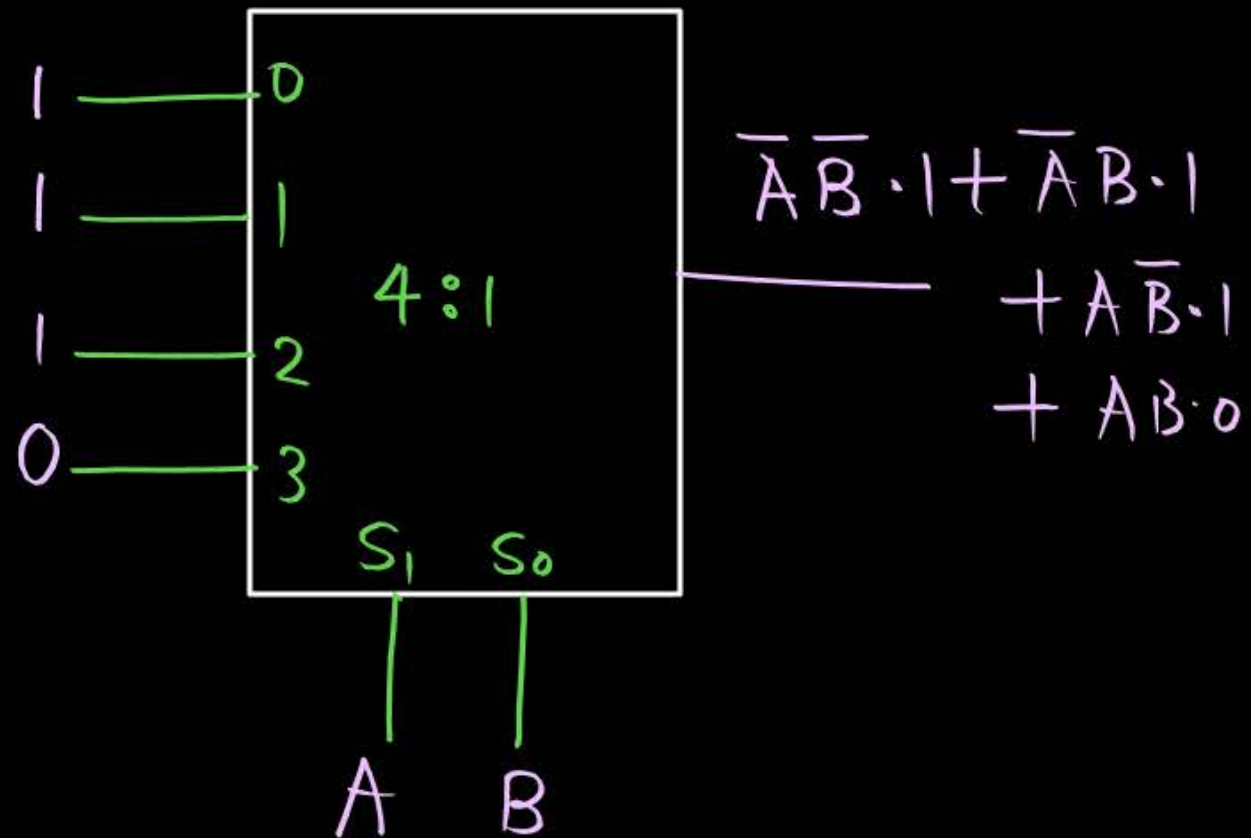


Q.  $f(A, B) = \Sigma(0, 1, 2) \rightarrow$  implement this logical function  $= \bar{A}\bar{B} + \bar{A}B + A\bar{B}$

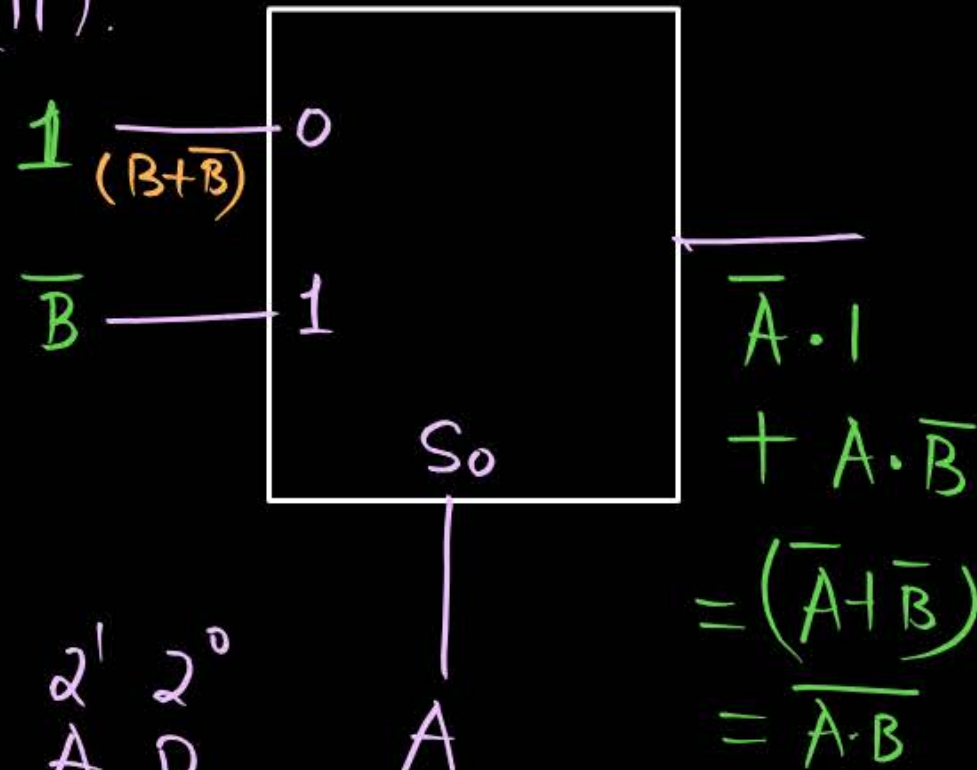
(i) using 4 : 1 MUX

(ii) using 2 : 1 MUX

(i).



(ii).



$2^1$	$2^0$
A	B
0	0
0	1
1	0
1	1

$$f(A, B) = \bar{A}\bar{B} + \bar{A}B + A\bar{B}$$

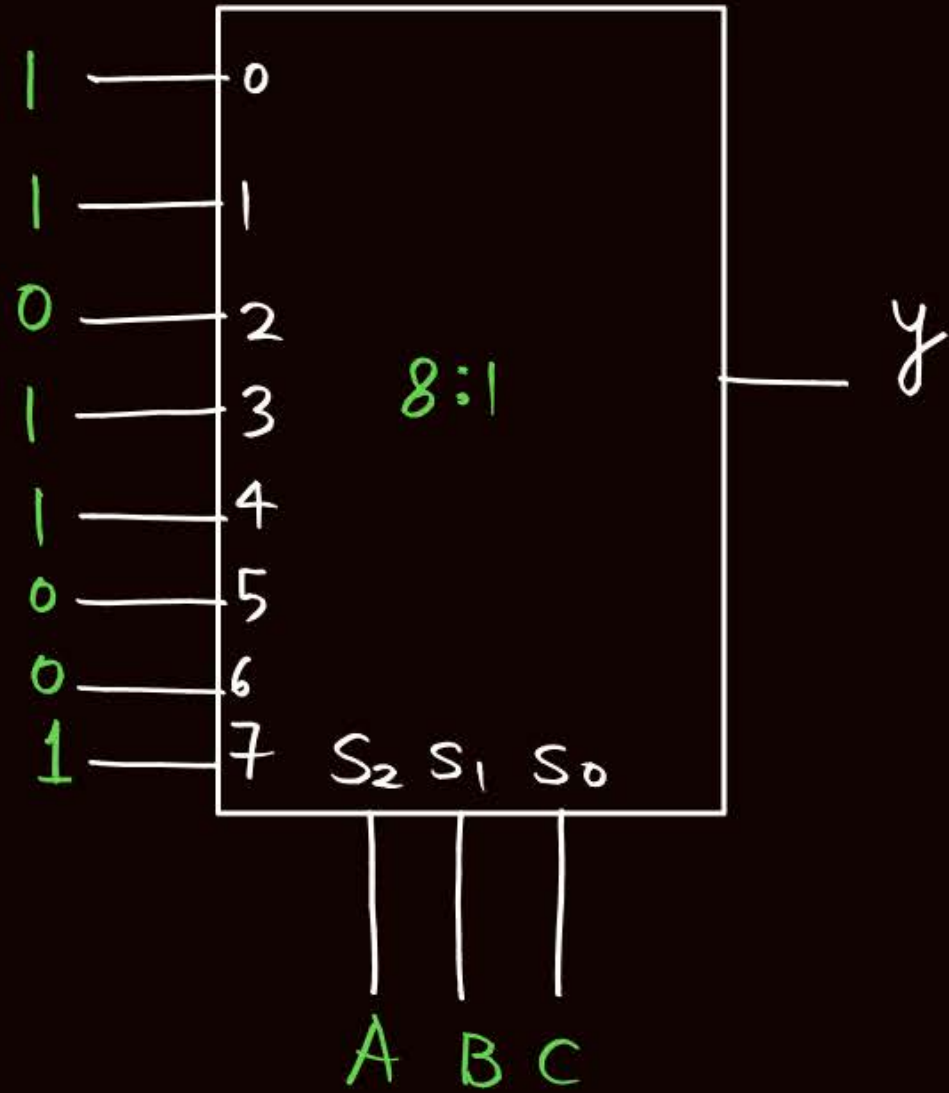
$$f(0, B) = \bar{B} + B + 0 = 1 \rightarrow I_0 = 1$$

$$f(1, B) = \bar{B}, I_1 = \bar{B}$$

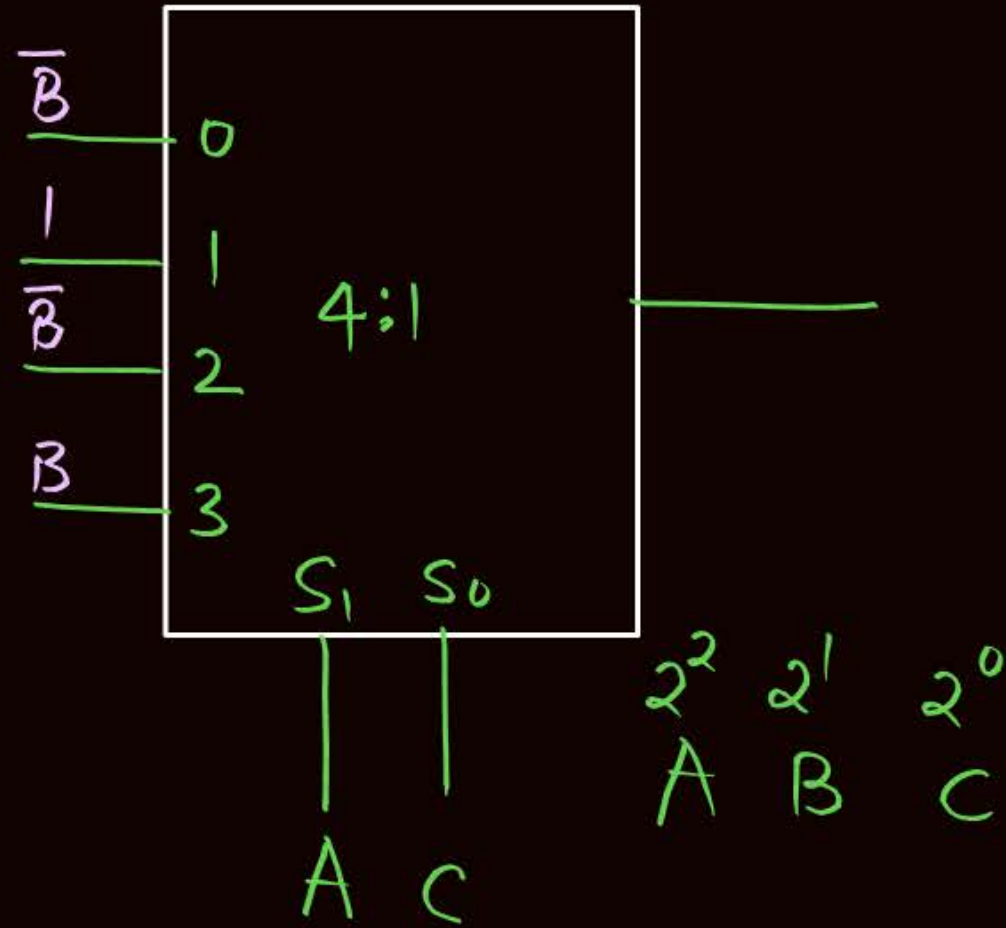
	$I_0$	$I_1$
$\bar{B}$	0	2
B	1	3
	1	$\bar{B}$

#Q.  $f(A, B, C) = \sum(0, 1, 3, 4, 7)$

- (i) Using 8:1 MUX  
(ii) Using 4:1 MUX

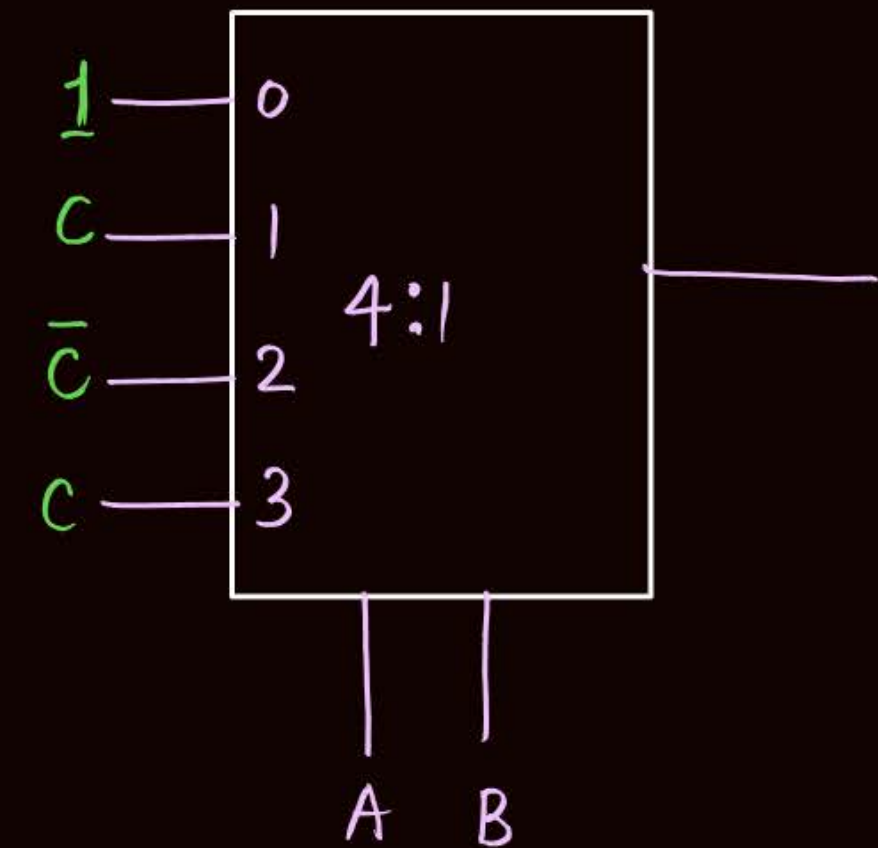


(ii) 4:1

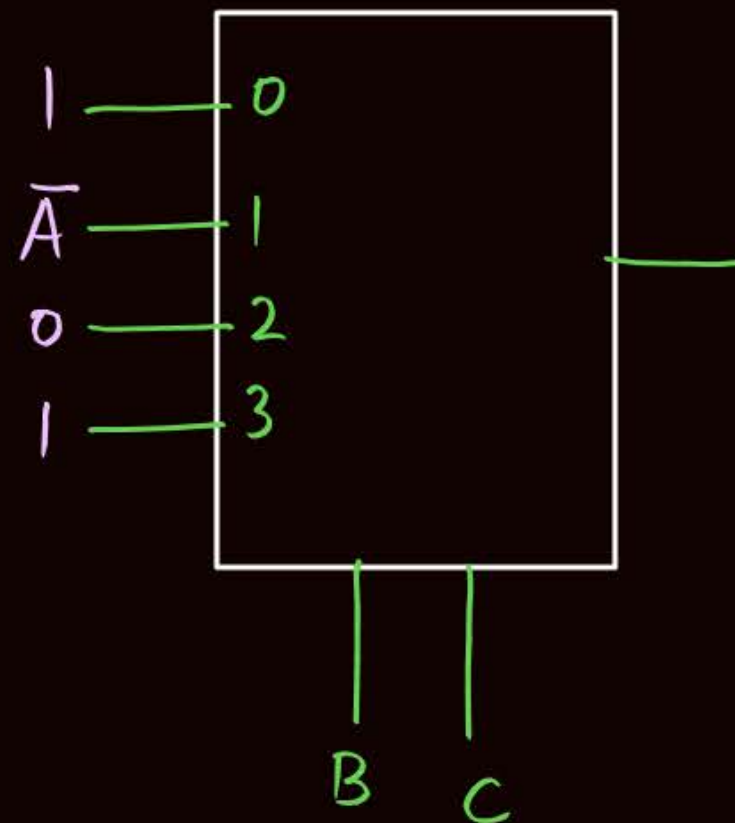


		$I_0$	$I_1$	$I_2$	$I_3$
0	$\overline{B}$	0	1	4	5
1	$B$	2	3	6	7
		$\overline{B}$	$B$	$\overline{B}$	$B$





	$I_0$	$I_1$	$I_2$	$I_3$
0 $\bar{C}$	0	2	4	6
1 C	1	3	5	7
	1	C	$\bar{C}$	C



	$I_0$	$I_1$	$I_2$	$I_3$
0 $\bar{A}$	0	1	2	3
1 A	4	5	6	7
	1	$\bar{A}$	0	1

$$f(A, B, C) = \sum(0, 1, 3, 4, 7)$$

$$f(A, 0, 0) = \sum(0, 4) = \bar{A}\bar{B}\bar{C} + A\bar{B}\bar{C}$$

$$f(A, 0, 1) = \sum 1 = \bar{A} + A = 1 \rightarrow I_0 = 1$$

$$= \bar{A}\bar{B}C$$

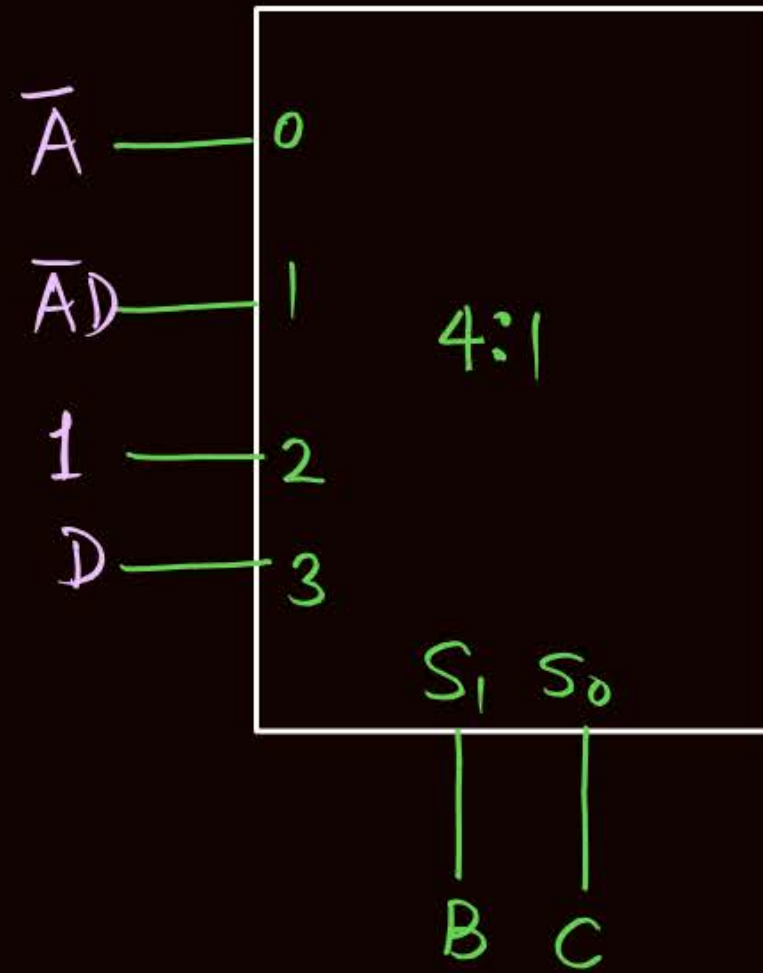
$$f(A, 1, 0) = 0, I_2 = 0, \bar{A} \cdot 1 \cdot 1 = \bar{A} \rightarrow I_1 = \bar{A}$$

$$f(A, 1, 1) = \sum(3, 7) = \bar{A} + A = 1 \rightarrow I_3 = 1$$



Q.  $f(A, B, C, D) = (\bar{A} + B)(\bar{C} + D) = \Sigma$

using 4:1 MUX.



		$I_0$	$I_1$	$I_2$	$I_3$
00	$\bar{A}\bar{D}$	0	2	4	6
01	$\bar{A}D$	1	3	5	7
10	$A\bar{D}$	8	10	12	14
11	$AD$	9	11	13	15
	$\bar{A}D + A\bar{D}$				

A	B	C	D
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0

$$\left. \begin{aligned} f(A, 0, 0, D) &= \bar{A} \rightarrow I_0 = \bar{A} \\ f(A, 0, 1, D) &= \bar{A}D \rightarrow I_1 = \bar{A}D \end{aligned} \right\} \begin{aligned} f(A, 1, 0, D) &= 1 \\ f(A, 1, 1, D) &= D \end{aligned}$$

$$I_2 = 1$$

$$I_3 = D$$

Q.  $f(A, B, C) = \Sigma (0, 1, 2, 3, 4, 7)$  H.W

(i) implement it using 8 : 1 MUX

(ii) implement it using 4 : 1 MUX

- AB as select line
- BC as select line
- AC as select line



## 2 Minute Summary

→ MUX



Thank you

**GW**  
*Soldiers !*

