CS & IT ENGINEERING

THEORY OF COMPUTATION

Regular Languages



Lecture No.- 02

Recap of Previous Lecture







Topic

Regular Expression

Topic

Construction of Regular Expression

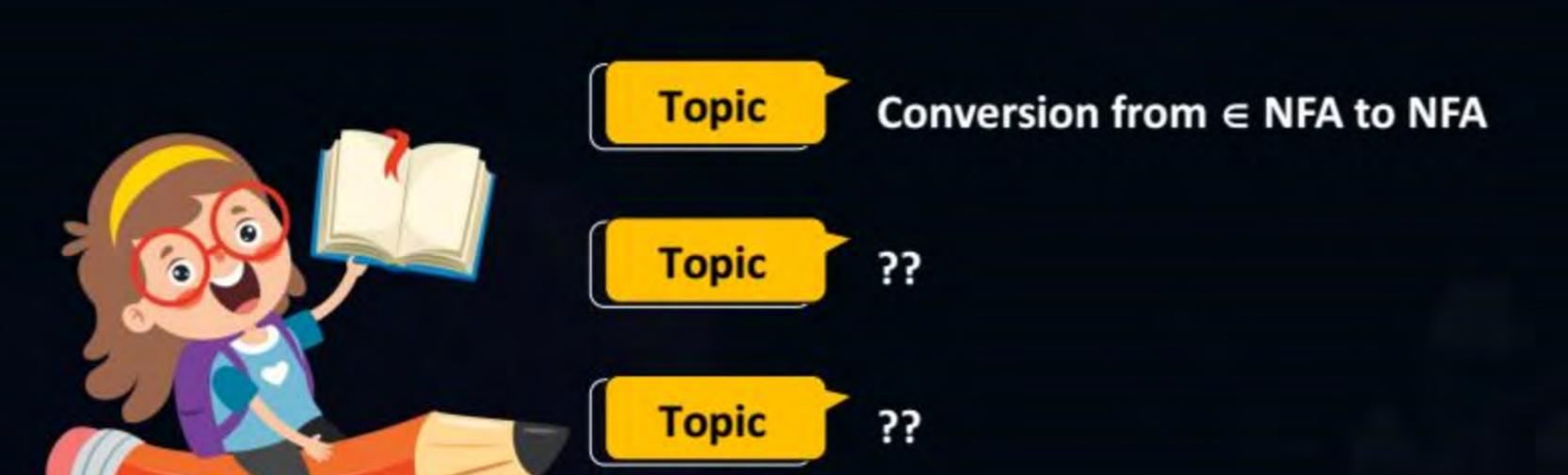
Topic

DFA States

Topics to be Covered





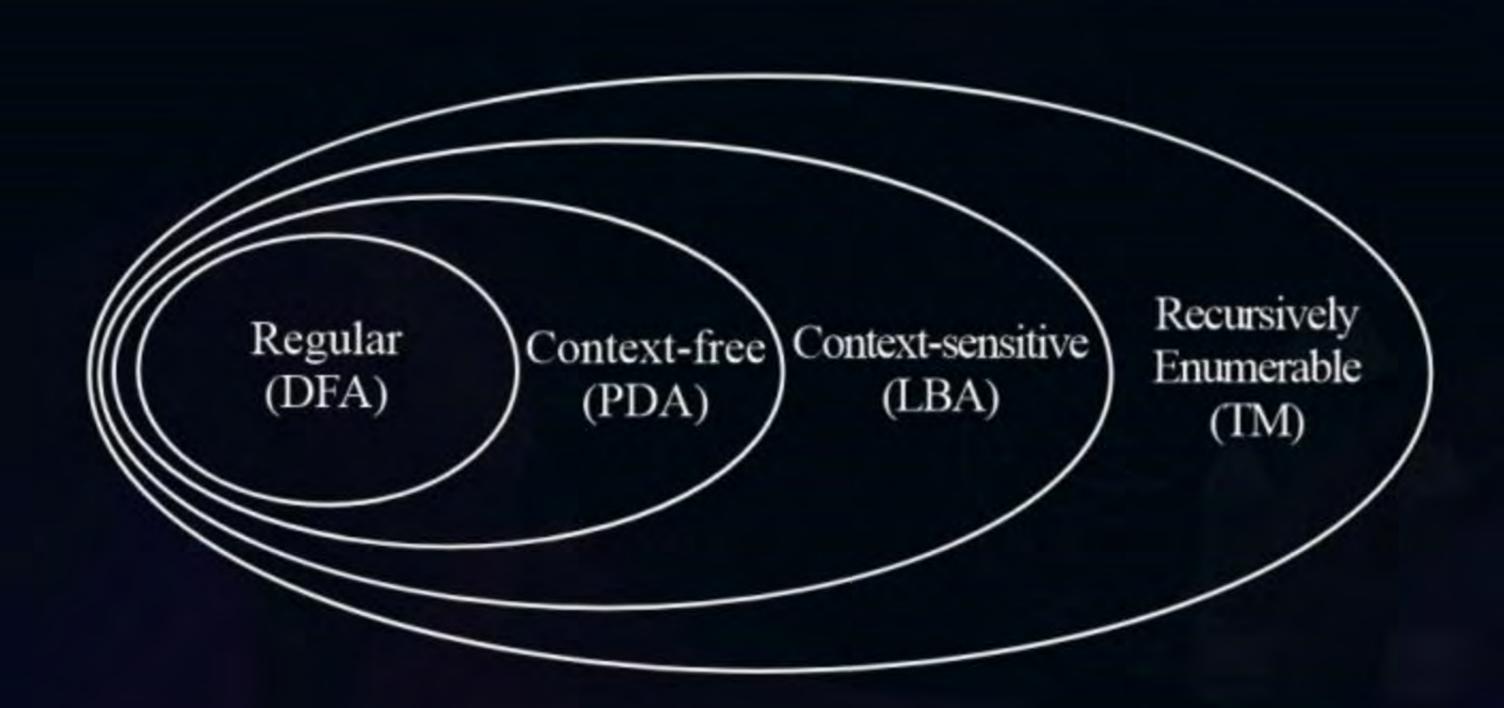


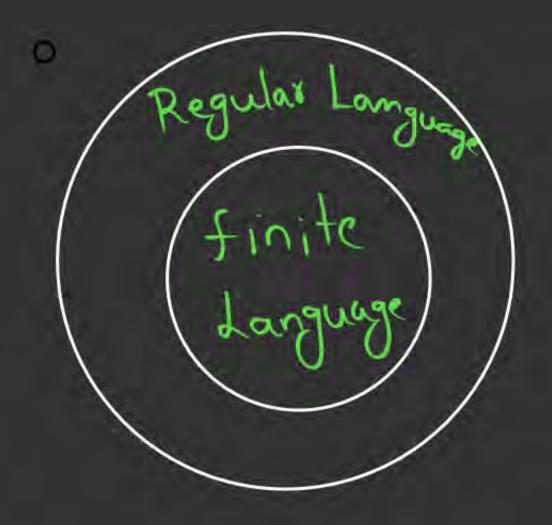
O(Finite danguage Every finite in Regular may be (Regular) Every Infinite danguages I may hot be Regular



Topic: Theory of Computation









Topic: Regular Language Detection



Which of these Languages are Regular

1.
$$L = \{a^nb^nc^n \mid 1 \le n \le 100 c\}$$

2.
$$L = \{a^nb^m \mid n + m = 10\}$$

3.
$$L = \{a^nb^m \mid n-m=5\}$$

4.
$$L = \{a^nb^m \mid n * m = 100\} \xrightarrow{finite} \{equal \}$$

5.
$$L = \{a^nb^m \mid n = 2m + 1\}$$

6.
$$L = \{a^nb^m \mid n > m\}$$

7.
$$L = \{a^n b^m \mid n > a_m a_m \}$$

Infinite Language De pendency exist Dependency not exist non Regular Regular



Topic: Regular Language Detection



8.
$$L = \{a^nb^m | n > m \text{ (or) } n < m\} = \{a^nb^m | n + m\} = Non Regular$$

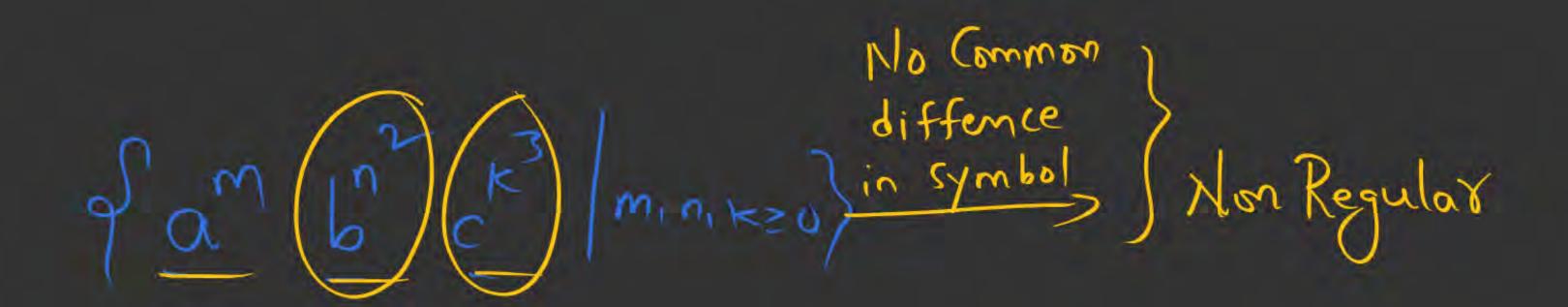
11.
$$L = \{a^n b^{2m} c^{3k}\}$$
 n, mm $k \ge 0\}$

12.
$$L = \{a^n b^{m^2} c^{k^3} | n, mk \ge 1\}$$
 -NO Dependency Non Regular

13.
$$L = \{a^n b^{m^2} c^{k^3} | n, mk \ge 1\}$$

14.
$$L = \left\{ \frac{a^{2^n}}{a^{3^n}} \middle| n \ge 0 \right\}$$

(Q) Which of the following in Regular? a) L= {(m2/2) m,n≥1} No Common det Non Regular (b) 1= {(n!) m! ck! | m, r, k > 1} -> Non Regular (C) [= {(an lm) (m+n) | m,n≥1} Dependent > Mon Regular (a) c= { and by ck lover. K lover. Finite > Regular



Common diff

(0) which of the following in Regular

(a)
$$L = \{a^n b^m | n \ge m \text{ and } n \le m\} = \{a^n b^n | n \ge i\}$$

(b)
$$L = \{a^{n}b^{m} | n > m (on) n < m\} = \{a^{n}b^{m}|n \neq m\}$$

$$Q = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p_m | u > m \text{ (and } u < w\} = \{a, p$$

(d) NWG

(a) Which of the following in Non Regular 2) L= {a^bm | (n+m) ig even} = (aa) (bb) + a(aa) b(bb) (b) L={a^bm/(n+m)ig odd}=(aa)*b(bb)*+a(aa)*(bb) (C) [= { a per | U= M_ } = { a per possible boundary Blow Board a) L= { 2m 13n | n, m > 10 dependency} Regular



Topic: Regular Language Detection





16.
$$L = \{a^{n^n} | n \ge 1\}$$
 $\longrightarrow \{a', a^2, a^3, \dots\}$ $\longrightarrow N$ on Regular

18.
$$L = \left\{ a^{100^{100^{100}}} \right\} \longrightarrow \left\{ 1 \right\} = finite \right\} \longrightarrow \text{Regular}$$

19.
$$L = \{(a^P)^* | p \text{ is prime number}\}$$

20.
$$L = \{a^p \mid p \text{ is prime number}\}$$

(a) Which of the following in Non Regular?

(b) Which of the following in Non Regular?

(c) Which of the following in Non Regular?

(d) Which of the following in Non Regular?

(e) L= fak | K in even number } Regular (b) L= { am Mary odd number} -> Regular (C) [= { (b) } | b in brime vamper) fatatatatatat-3-->
L= fat prime number) -> Non Regular

If a danguage if formed over 1 symbol.

If Common différence exist (A.P)-> Régular

(0) Which of the following in Regular! (a) [= { a | n = 1} = { a | a | a | a | a | a | a | - - } - Non Regular (b) d = { an | nz|} -> { a', a², a³, a'--} -> Non Regular $\frac{1}{2} \left\{ \frac{1}{2} \left(\frac{1}{2} \left$

(a) which of the following y Regular?

(b) which of the following y Regular?

(c) which of the following y Regular?

(d) L= {1, 2, 4, 8, 16 - - - 2 - - } = {1 | n > 0} > Non Regular all these numbers written in Unary (b) [= {00,000,1000c au these numbers written in binary) (10) > Regular

(c) [= d 3, | v=0) -> Mw Sedinar.

a) Nous

Drood [1, 1111, 18 - - -] = No Common diff) Non Reg binary? 1, 2, 4, 8 - - - } 00000 10000



Topic: Regular Language Detection



21.
$$L = \{a^k | k \text{ is even number}\} - \frac{Regular}{}$$

(22)
$$L = \{ ww^R \mid w \not\models (a+b)^* \} \longrightarrow Non Regular$$

23.
$$L = \left\{ ww^R \middle| \begin{array}{l} w \in \{a, b\}^* \\ w \in \{a, b\}^* \end{array} \right\}$$

24.
$$L = \{wbw^R | w \in \{a\}^*\}$$

25.
$$L = \{x | x \in \{a, b\}^* \ n_a(x) \ \text{mod } 3 = n_b(x) \ \text{mod } 2\}$$

26.
$$L = \{x | x \in \{a, b\}^* n_a(x) \mod 2 > n_b(x) \mod 3\}$$

27.
$$L = \{x | x \in \{a, b\}^* \ n_a(x) \ \text{mod } 3 \neq n_b(x) \ \text{mod} 3\} \}$$

28.
$$L = \{x | x \in \{a, b, c\}^* n_a(x) \neq n_b(x)\}$$

Palindrome danguages L= { W W W W W R malayalam)
WR

 $L = \{WW^{R}\} W \in (0)^{k} \rightarrow \text{Regular}$ $W \in (a+b)^{k} \rightarrow \text{Non Regular}$ Dependency

(3)
$$L = \left\{ WWX \left| W, X \in (a+b)^{+} \right\} \right\}$$

(A)
$$L = \{W \times W^R | W_1 \times \{(a+b)^t\}$$

$$(5) L = \{ wwrwr|we(a+b)t \}$$



$$\begin{array}{ll}
U_1 \times e(a+b) \\
U_2 \times e(a+b)
\end{array}$$

$$\begin{array}{ll}
U_1 \times e(a+b) \\
U_2 \times e(a+b)
\end{array}$$

$$\begin{array}{ll}
U_1 \times e(a+b) \\
U_2 \times e(a+b)
\end{array}$$

$$\begin{array}{ll}
U_1 \times e(a+b) \\
U_2 \times e(a+b)
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U_2 \times e(a+b)
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U_1 \times e(a+b) \\
U_2 \times e(a+b)
\end{array}$$

$$\begin{array}{ll}
U_1 \times e(a+b) \\
U_2 \times e(a+b)
\end{array}$$

$$\begin{array}{ll}
U_2 \times e(a+b) \\
U_3 \times e(a+b)
\end{array}$$

$$\left\{ \begin{array}{l} W(X)W^{R} \\ W = (a+b)^{*} \\ W = (a+b)^{*} \\ W = 0 \end{array} \right\}$$

$$W = 0$$

$$W =$$

$$\begin{array}{cccc}
(a+b)^* & (a+b)^* & (a+b)^* \\
W = & (a+b)^* & (a+b)^* & (a+b)^* \\
W = & (a+b)^* & (a+b)^* & (a+b)^* \\
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W = & (a+b)^* & (a+b)^* & (a+b)^* & (a+b)^* & (a+b)^* & (a+b)^* \\
W = & (a+b)^* & (a+b)^*$$

(a+b) U a b = (a+b)

Super
$$a = a$$



2 mins Summary



Topic One

Topic Two

Topic Three

Topic Four

Topic Five



THANK - YOU