# CS & IT ENGING

Theory of Computation



Lecture No.- 06



# **Recap of Previous Lecture**









Complement **Topic** 

minimization of DFA Topic

# **Topics to be Covered**









Topic

Finite Automaton & Regular Languages.

Topic

Pushdown Automata & Context free Languages.

Topic

Turing Machine & Recursive Enumerable Languages.

Topic

Undecidability.



#### **Topic: Deterministic Finite Automata**



#### **FORMAL DFA:** DFA is defined as

DFA = 
$$(Q, \sum, q_0, F, \delta)$$

Q: Finite set of states

 $\Sigma$ : Input alphabet

q<sub>0</sub>: Initial state

F: Set of final states

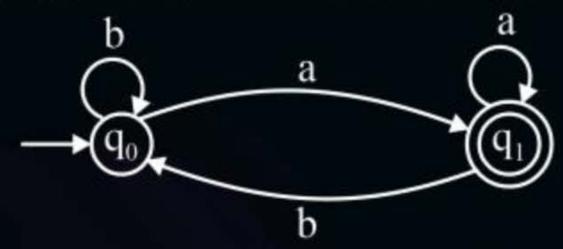
 $\delta$ : Transition function  $Q * \Sigma \rightarrow Q$ 



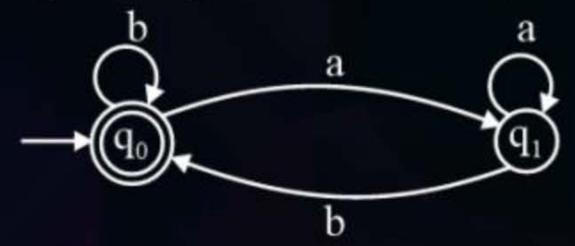
## **Topic: Complement of DFA**



By interchanging final and non final states we can convert into complement DFA.



Set of all strings ending with a after complement



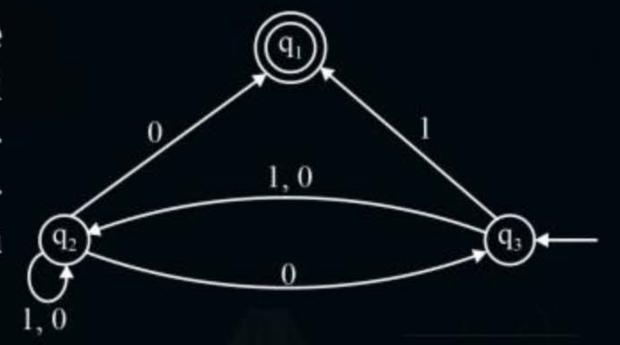
Set of all string .... Ending with

#### MCQ



#Q. Consider the NFA M shown below.

Let the language accepted by M be L. Let  $L_1$  be the language accepted by the NFA  $M_1$ , obtained by changing the accepting state of M to a non-accepting state and by changing the non-accepting state of M to accepting states. Which of the following statements is true?



**A** 
$$L_1 = \{0, 1\}^* - L$$

**B** 
$$L_1 = \{0, 1\}^*$$

$$L_1 = L$$



## **Topic: DFA Construction**



Construct minimal state DFA that accerpts all strings os 0's and 1's where each string ending with 00.

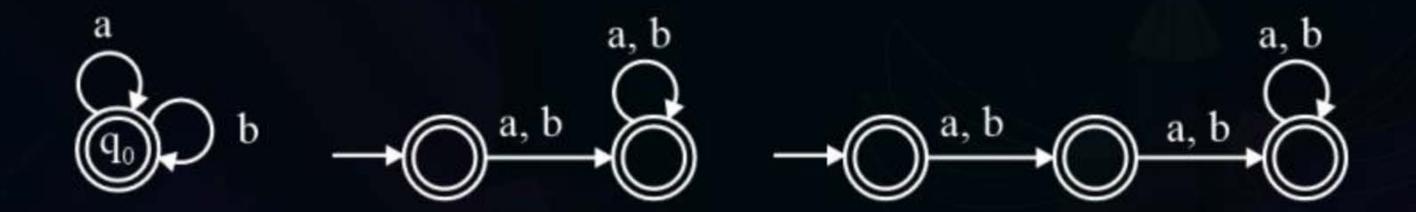


#### **Topic: Minimization of DFA**



→ For a given regular language even though many DFA exist but minimal state DFA is unique.

Ex: Complete Language:  $\Sigma^*$ 





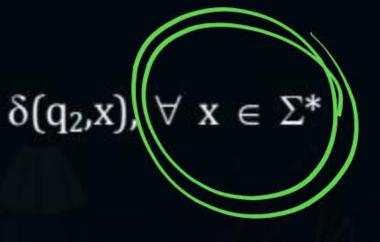
#### **Topic: Minimization Algorithm**



- State equivalence algorithm
- Table filling algorithm

#### **Equivalent States:**

Two states  $q_0$ ,  $q_1$  are said to be equivalent both  $\delta$  ( $q_0$ , x) and  $\delta(q_2,x)$ ,  $\forall$   $x \in \Sigma^*$  should result either final state or non final state.



$$\delta (q_1, x) = F$$
 $\delta (q_2, x) = F$ 
NF





Elimination inaccessible states.

#### inaccessible state:

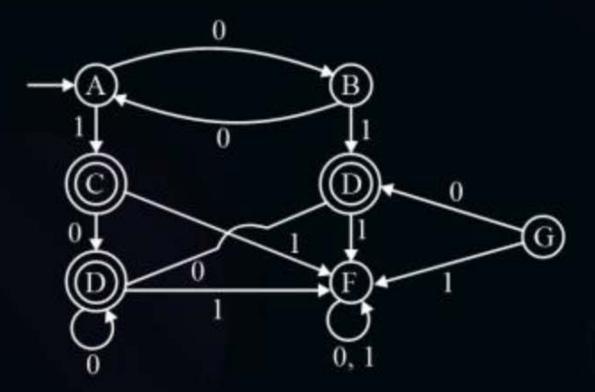
Any State which is not reachable from dead state is inaccessible state.

- Apply algorithm steps
- Merge single group into one state
- Construct new minimized DFA





#### Reduce states of following DFA



Step-1: Elimination inaccessible state.

Note: Dead state is different from inaccessible state.





Step:2

State	0	1
A	В	С
В	A	D
F	F	F
(C)	Е	F
<b>①</b>	E	F
(E)	Е	F

Algorithm:

1.  $\{A, B, F\} \{C, D, E\}$ 

2.

3.





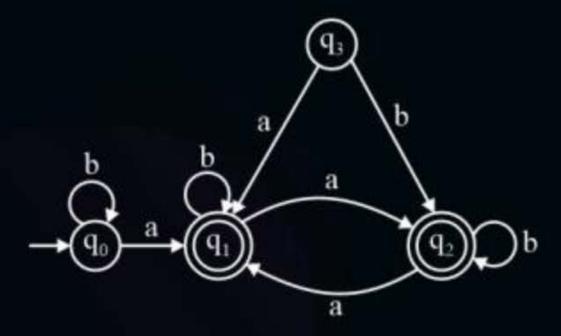
#### **Minimized DFA**







#### Consider the following Finite State Automation







Step 1: Eliminate q<sub>3</sub>

Step 2:

	a	b
$q_0$	$q_1$	$q_0$
$q_1$	q <sub>2</sub>	$q_1$
$(q_2)$	$q_1$	$q_2$

Algorithm step

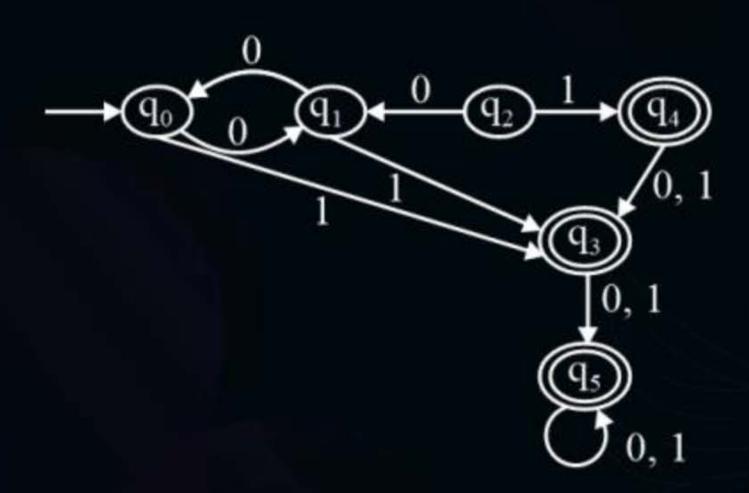
1. 
$$\{q_0\}\{q_1,q_2\}$$

2. 
$$\{q_0\}\{q_1,q_2\}$$

Minimum DFA







Minimize given DFA



## **Topic: Procedure**



Step 1: Eliminate

Step 2:

	a	b
$q_1$	$q_1$	$q_3$
$q_2$	$q_0$	$q_3$
$\overline{q_3}$	$q_5$	$q_5$
$(q_5)$	$q_5$	$q_5$

Algorithm

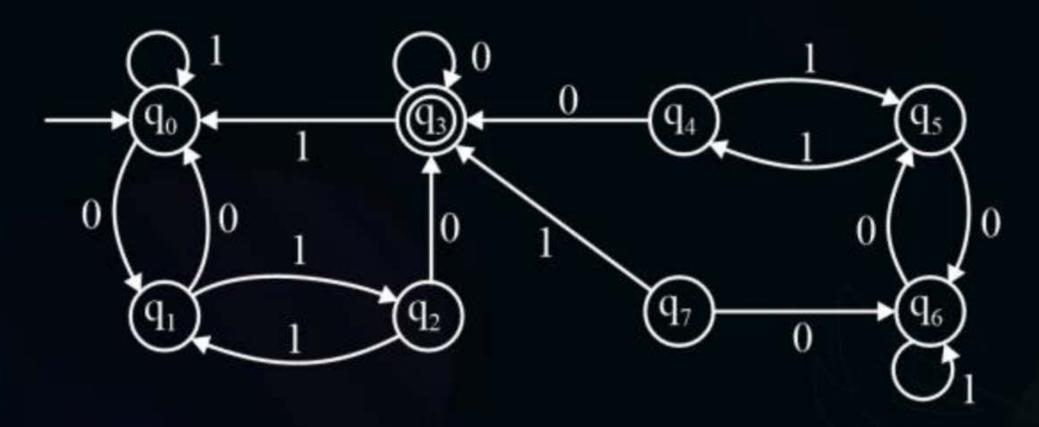
1.  $\{q_0, q_1\} \{q_3, q_5\}$ 

2.

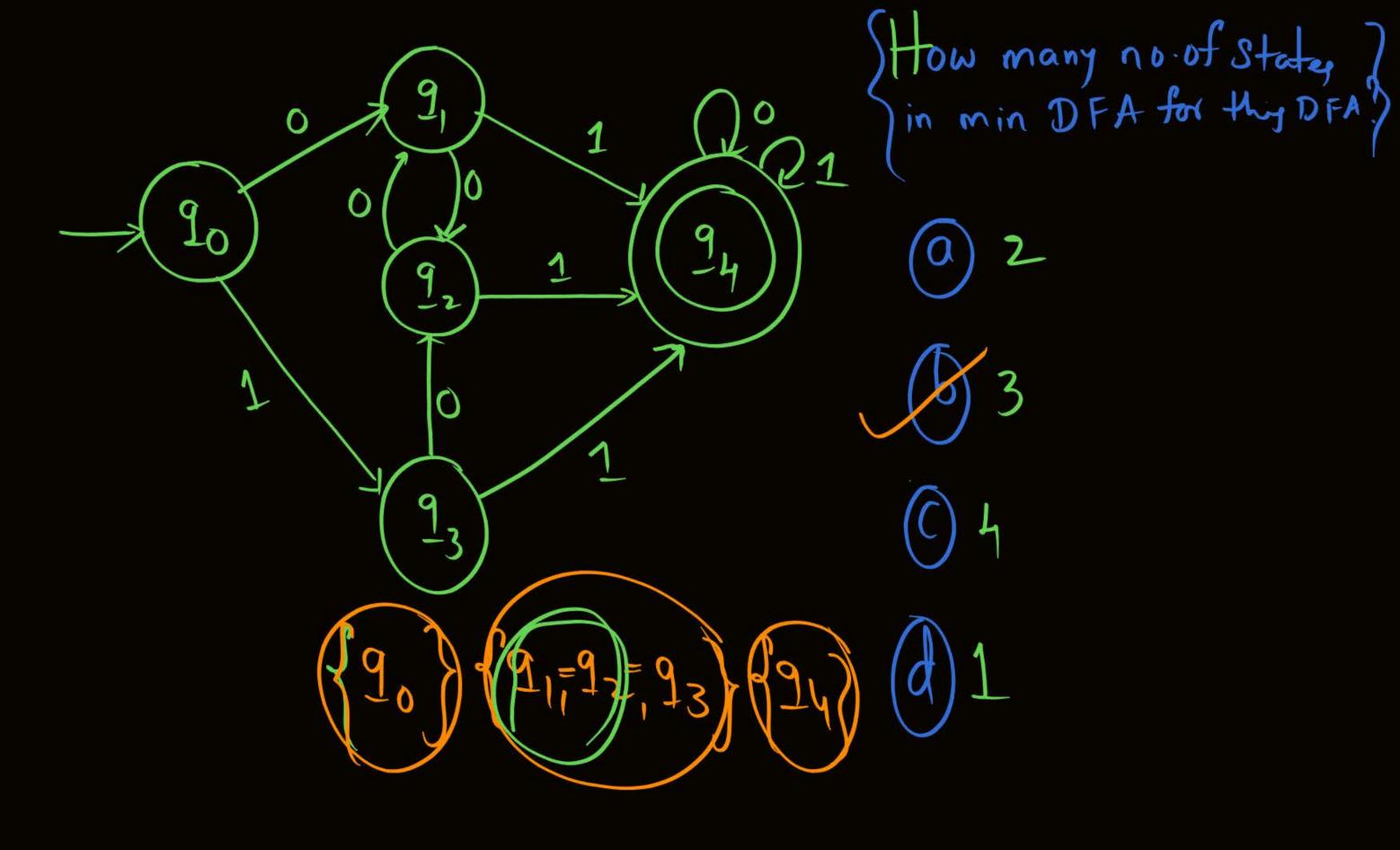
Minimum DFA







How many inaccessible states present in given DFA



Quhich of the following Statement is false! a) 243 are equivalent states (b) 1 & 2 are distinguishable " (C) 3 by are distinguishable " (d) Note.







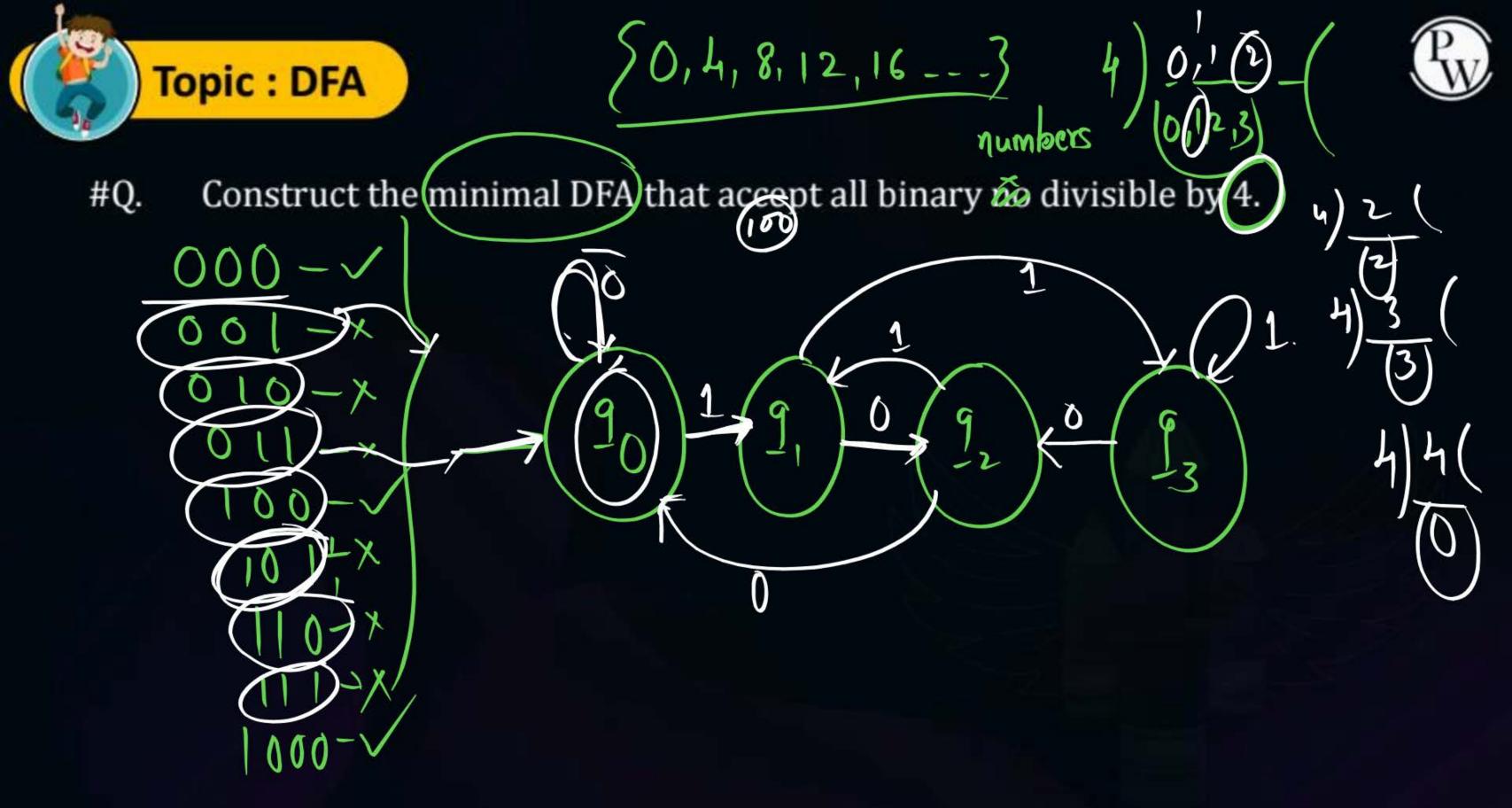
Consider the 5-state DFA M accepting the language  $L(M) \subset \text{subset } (0 + 1)$ shown below. For any string  $w \in (0 + 1)^*$  let  $n_0(w)$  be the number of 0's in w and n<sub>1</sub> (w) be the number of 1/s in w

Which of the following statements is/are FALSE?



- States 2 and 4 are distinguishable in
- States 2 and 5 are distinguishable in  $M \rightarrow folse$
- Any string w with  $n_0(w) = n_1(w)$  is in  $L(M) \rightarrow m_0^2$
- States 3 and 4 are distinguishable in M





Table

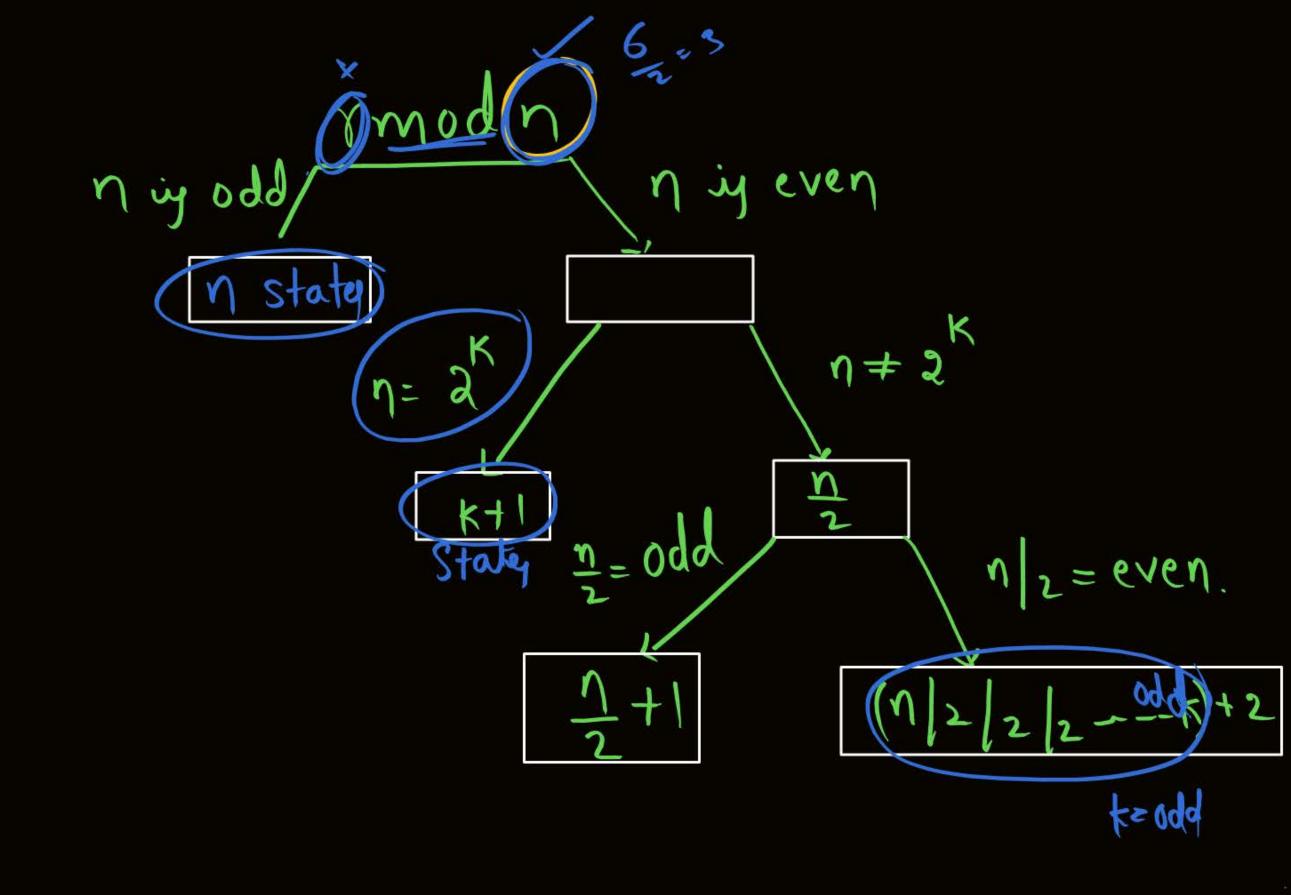




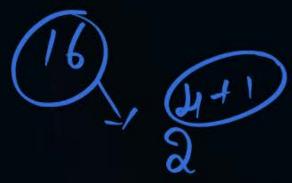
#Q. Construct the minimal DFA that accept all binary no divisible by . 5

		0	1	~
7(9	9	20	21	
9	1	92	93	
9	2	9 <sub>4</sub>	%	
9	3	9,	gr	
0	4	4/3	The	

minimization



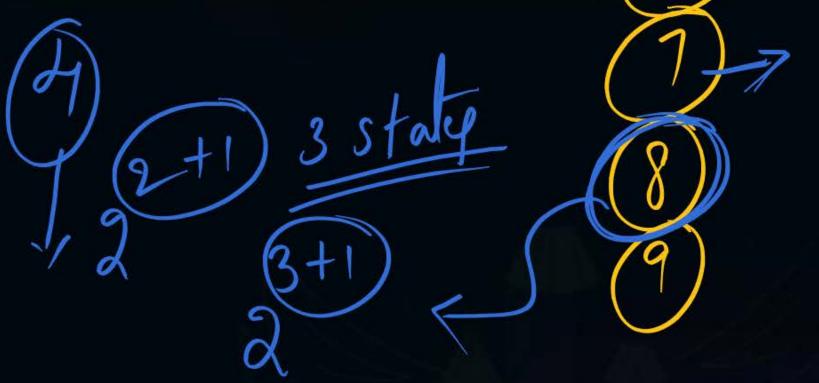








#Q. Construct the minimal DFA that accept all binary no divisible by . 6







#Q. Construct the minimal DFA that accept all binary no divisible by

$$3 \rightarrow 3$$
 $4 \rightarrow 3$ 
 $5 \rightarrow 5$ 
 $6 \rightarrow 3$ 
 $3 \rightarrow 3$ 
 $3 \rightarrow 5$ 
 $6 \rightarrow 3$ 
 $3 \rightarrow 5$ 
 $5 \rightarrow 10$ 
 $10 \rightarrow 10$ 
 $2 = (5+) = 6$ 





#Q. Construct the minimal DFA that accept all binary no divisible by

$$12 - \frac{12}{2} = \frac{6}{2} = \frac{3+2}{2} = \frac{5}{5}$$





#Q. Construct the minimal DFA that accept all binary no divisible by





Construct the minimal DFA that accept all strings of a's and b's where

- 1. Each string ending with b.
- Each string start with a and end with b.
- 3. Each string starting and ending with different symbol.
- 4. Each string starting and ending with same symbol.





Construct the minimal DFA that accept all string a's and b's where

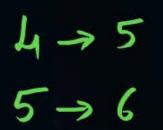
- 1. Length of string exactly 4.
- 2. Number of a's length of string atleast 4.
- 3. Length of string atmost 4.
- 4. Length of string divisible by 4.
- Number of a's exactly 5.  $\eta + 2 \eta + 2$ 
  - 6. Number of b's exactly 2.
  - 7. Number of a's divisible by 3.
- 8. Number of b's not divisible by 4.
- 9. Length of the string even.





#Q. Length of string exactly 4.











4-06



#Q. Length of string atmost 4.

atmosk 7-3 (1+2) states

Disposibi

Sunday

@ 10 am



