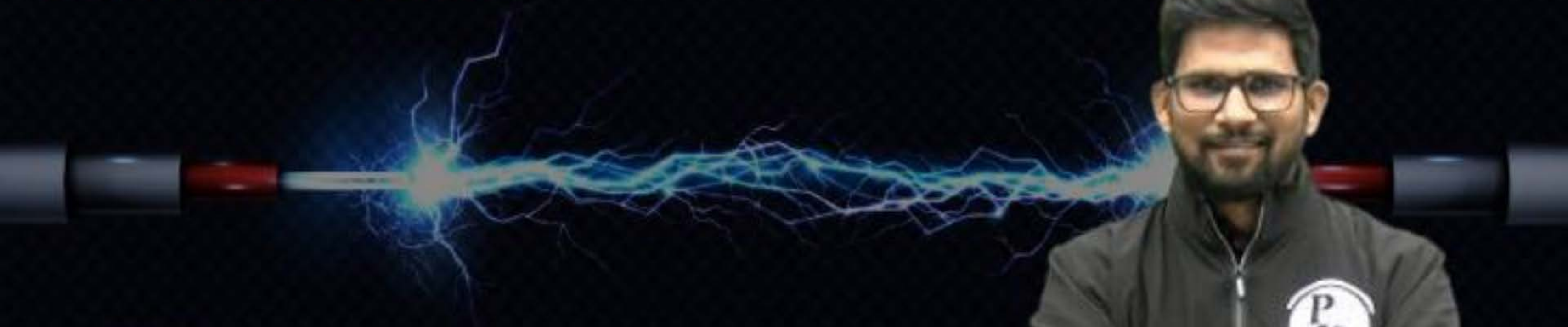


# COMPUTER SCIENCE & IT

## DIGITAL LOGIC



Lecture No. 03

**BOOLEAN THEOREMS AND  
GATES**

**By- Chandan Gupta Sir**







# Recap of Previous Lecture

Boolean Theorems & Questions Discussion



# Topics to be Covered

Logic gates





- $AB + A\bar{B}C = A[B + \bar{B}C] = A[(B + \bar{B}) \cdot (B + C)] = A \cdot (B + C) = AB + AC$

- $AB + B\bar{C} + \bar{A}C = B(A + \bar{C}) + \bar{A}C = \bar{A}C + (A + \bar{C})B$   
 $= P + \bar{P}B = (P + \bar{P}) \cdot (P + B) = (P + B)$   
 $= \bar{A}C + B$

- $(A + B + CD)(\bar{A} + B + \bar{C}D)(\bar{A} + B + C)$   
 $= B + (\bar{A} \cdot C) = (B + \bar{A}) \cdot (B + C) = (\bar{A} + B)(B + C)$

$$(A + B + CD)(\bar{A} + B + \bar{C} \cdot D \cdot C)$$

$$(A + B + CD)(\bar{A} + B)$$

$$= B + \bar{A}(A + CD)$$

$$= \underline{\underline{B + \bar{A}CD}} = (\bar{A} + B) \cdot (B + C) \cdot (B + D)$$

# [ GATES ]

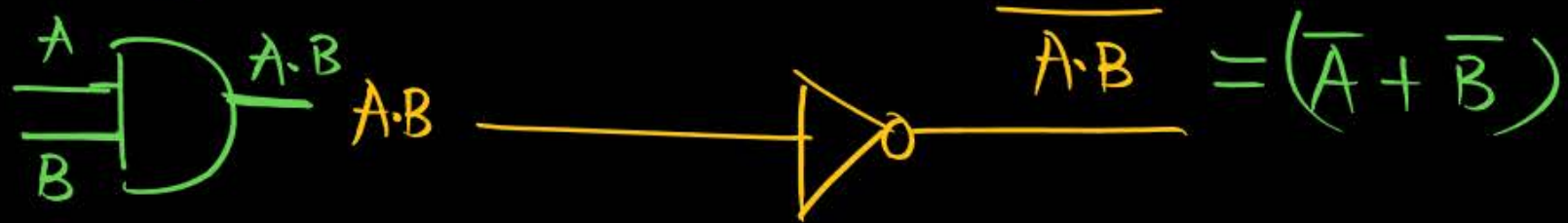
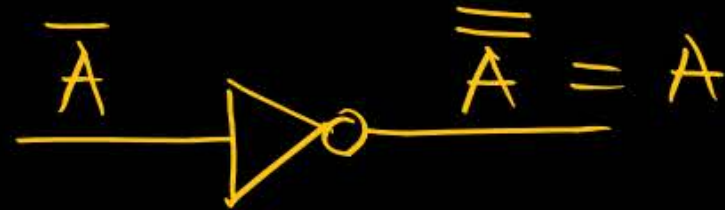
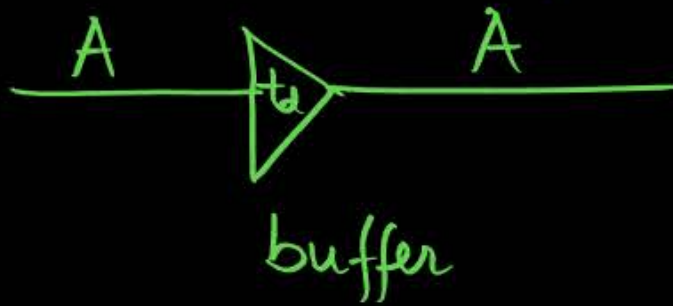


- Basic Gates : NOT, OR, AND
- Arithmetic Gates : XOR & XNOR
- Universal Gates : NAND & NOR

# [ NOT GATE ]



Representation of NOT GATE :



TOCCI

	Input <u>A</u>	Output Y
0	0	1
1	1	0

$$Y(A) = \sum(0) = \prod(1) \\ = \bar{A} = \bar{A}$$

# [ AND GATE ]



Representation :



	A	B	Output $y = A \cdot B$
0	0	0	0
1	0	1	0
2	1	0	0
3	1	1	1

$$\begin{aligned} \checkmark y(A,B) &= \sum 3 = \pi(0,1,2) = \frac{(A+B) \cdot (A+\bar{B})}{(\bar{A}+B)} \\ &= \underline{\underline{A \cdot B}} \\ &= [A + B \cdot \bar{B}] \\ &= A \cdot (\bar{A} + B) \\ &= \underline{\underline{A \cdot B}} \end{aligned}$$



- Commutative Law :



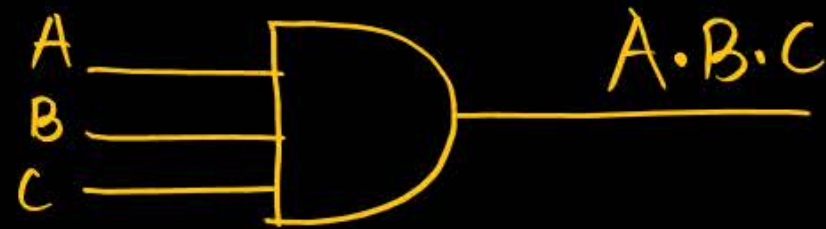
$$A \cdot B = B \cdot A$$

→ It holds commutative law.

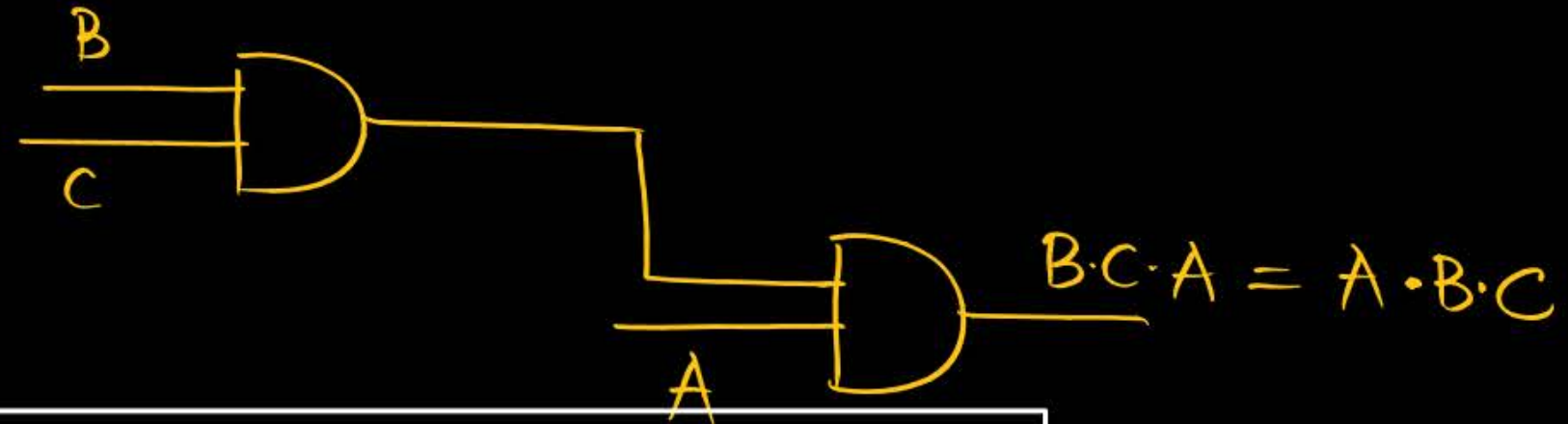
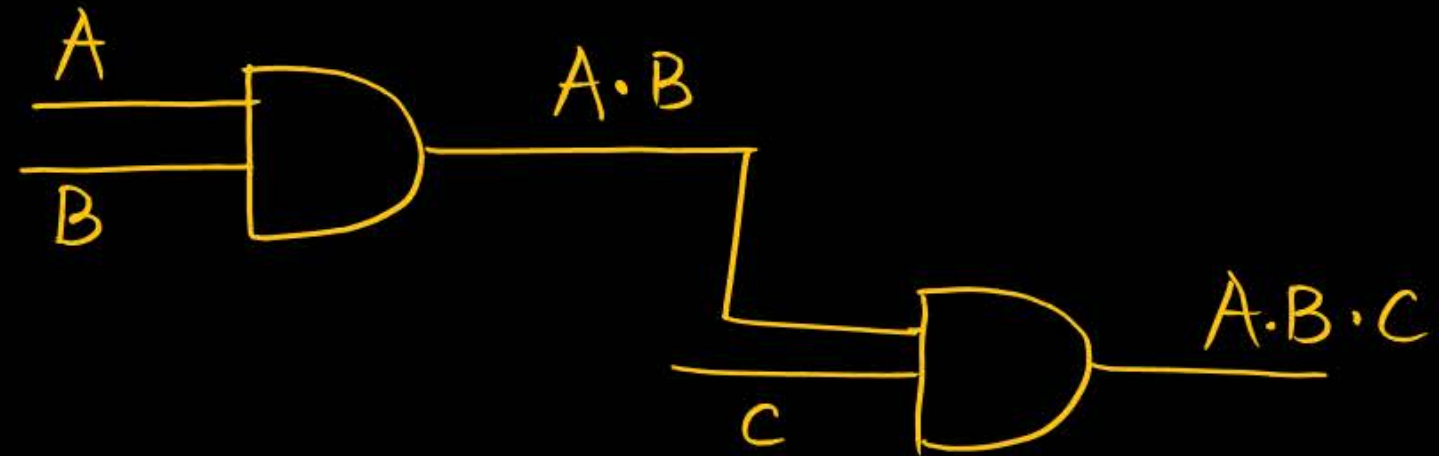
→ order of the variable is not important.



- Associative Law :

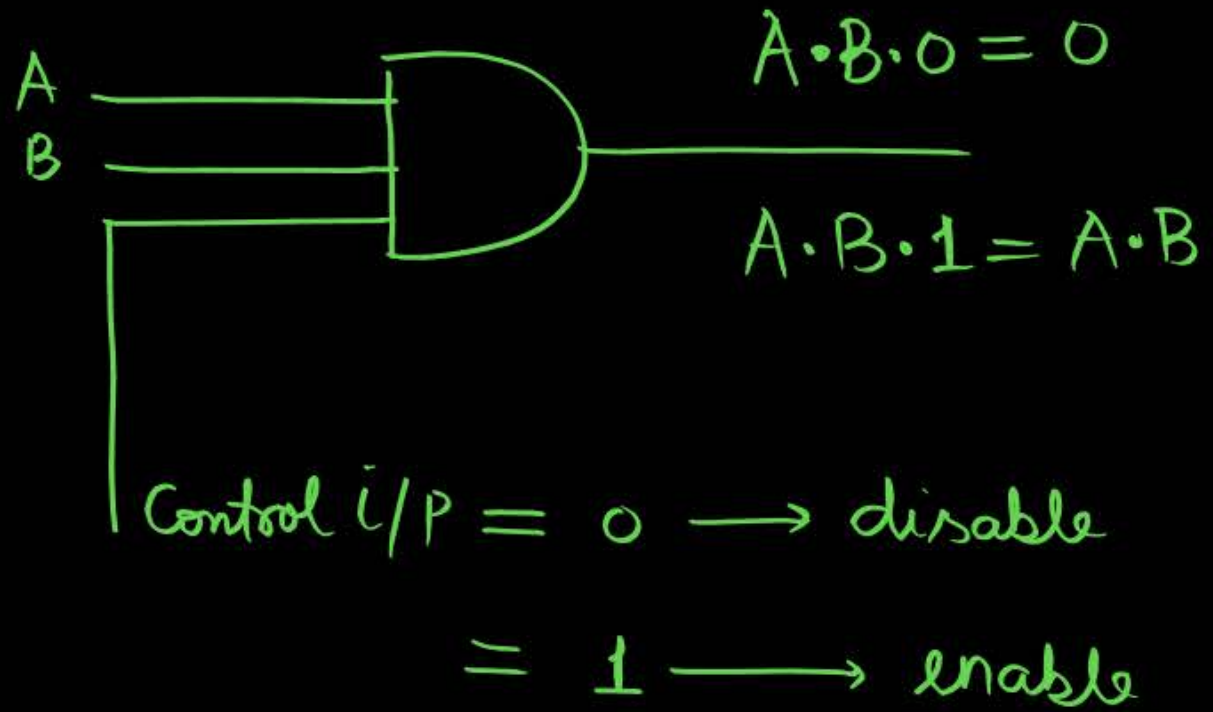


It holds  
associative law.



→ Multi i/p AND gate can be constructed using 2-i/p AND gates.

- Enable and Disable input for AND gate :



- IMP points :
- If any one of i/p line is at logic '0', then irrespective of other i/p lines O/P will be '0'.
- O/P will be '1' in only one case when all the i/p lines will be at logic '1'.



# [ OR GATE ]



Representation :



	A	B	$y = A + B$
0	0	0	0
1	0	1	1
2	1	0	1
3	1	1	1

$$y(A,B) = \sum(1,2,3) = \pi(0) = (A + B)$$

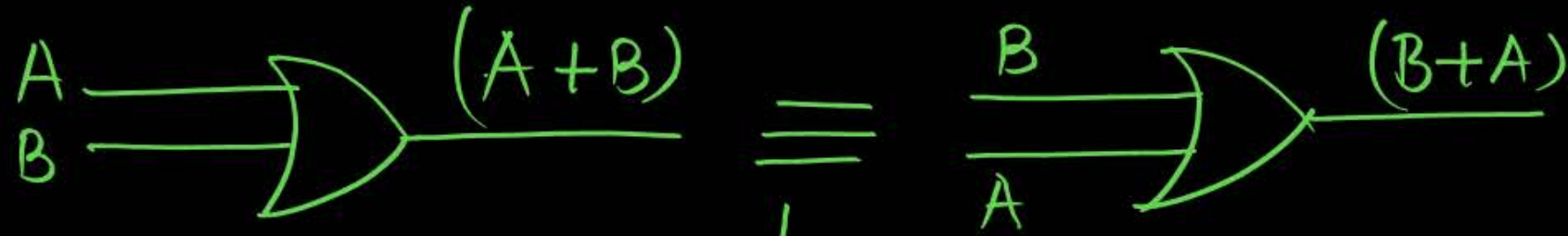
$$= \bar{A}B + A\bar{B} + AB$$

$$= \bar{A}B + A$$

$$= A + (\bar{A}B)$$

$$= (A + \bar{A}) \cdot (A + B) = (A + B)$$

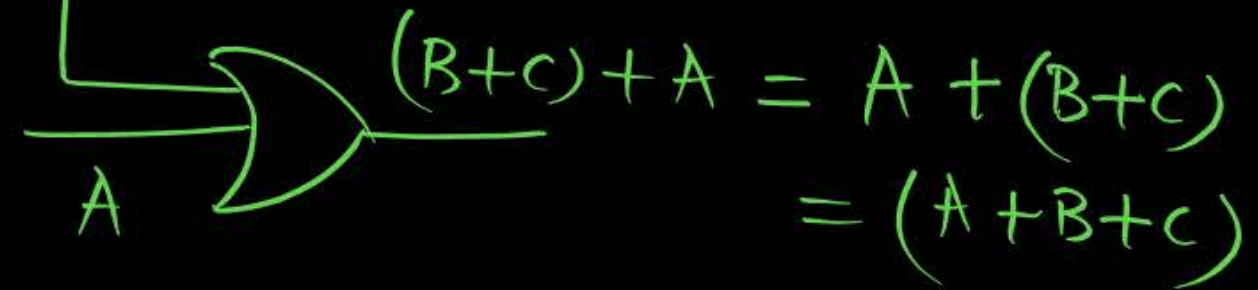
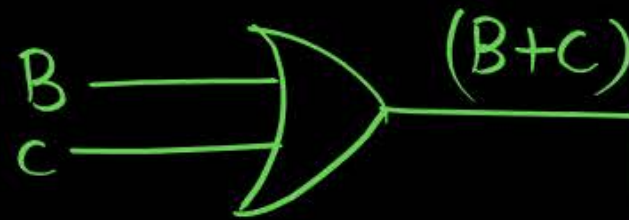
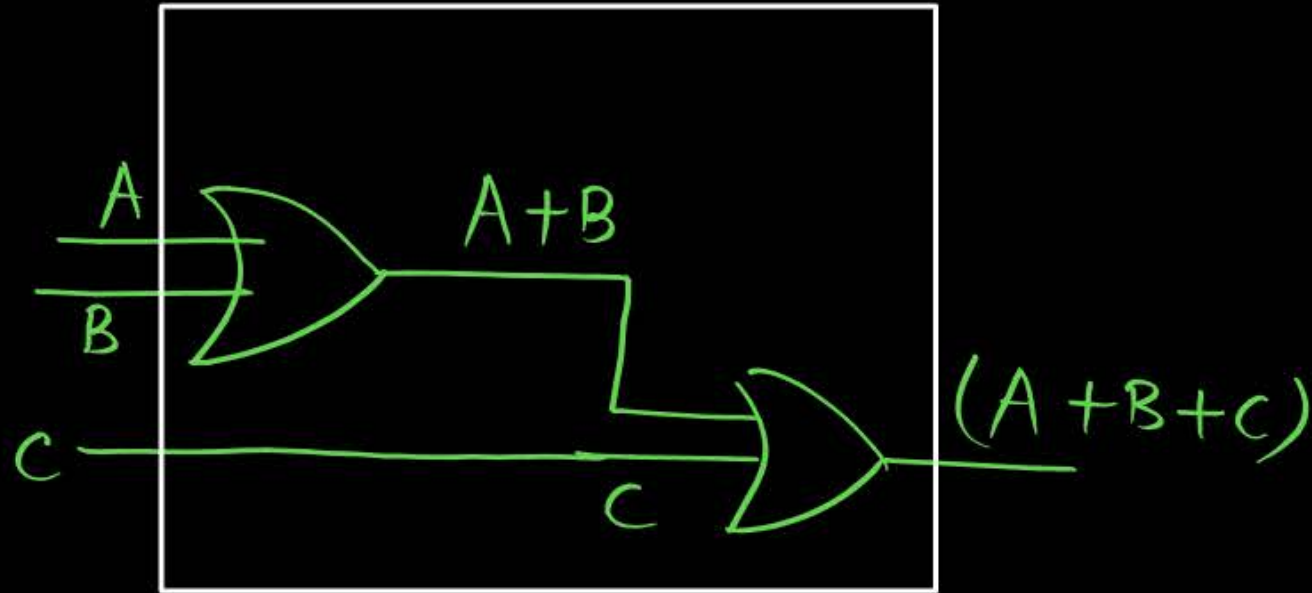
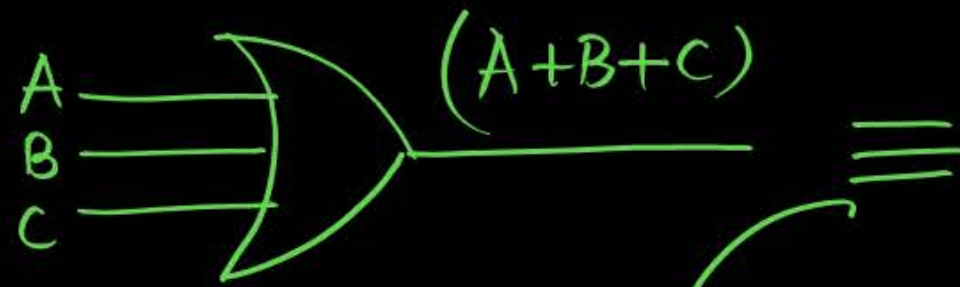
- Commutative Law :



→ It holds commutative law.

$$(A+B) = (B+A)$$

- Associative Law :

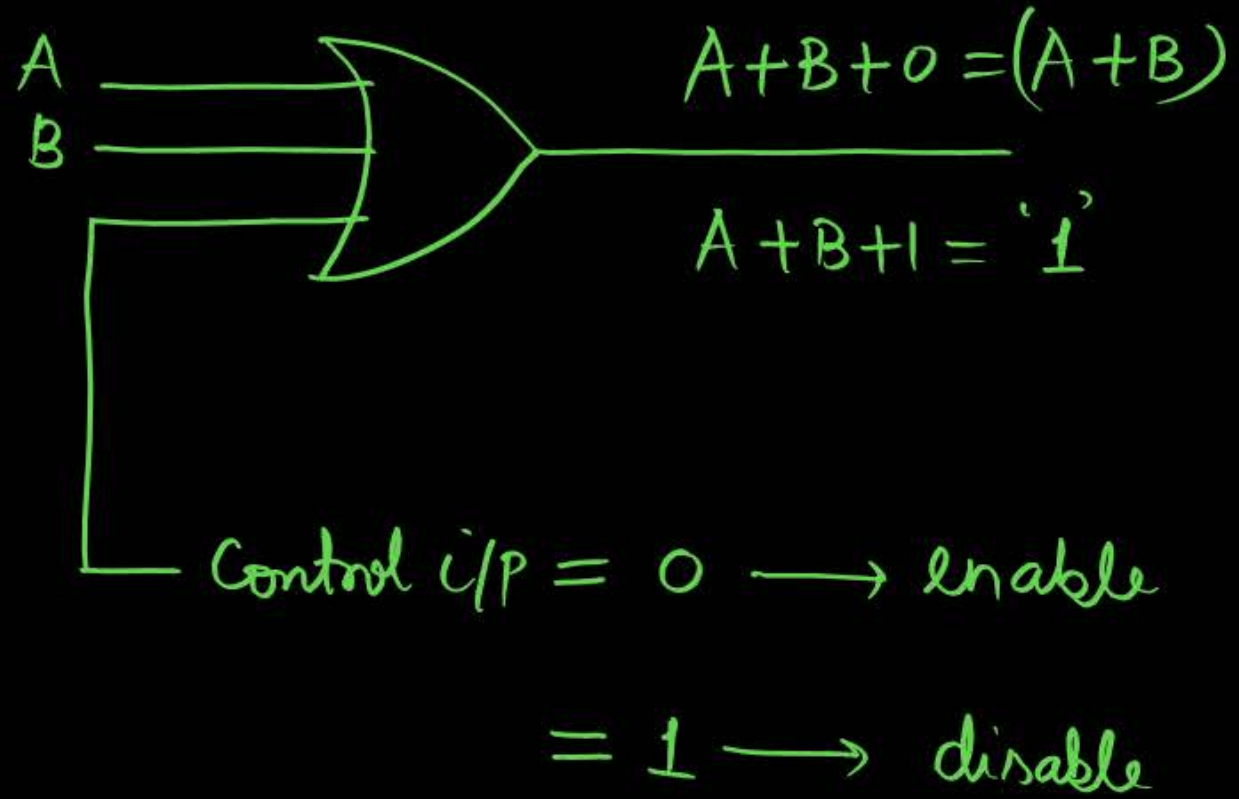


It holds associative law.

- Multi i/P OR gate can be constructed using 2-i/P OR gates.



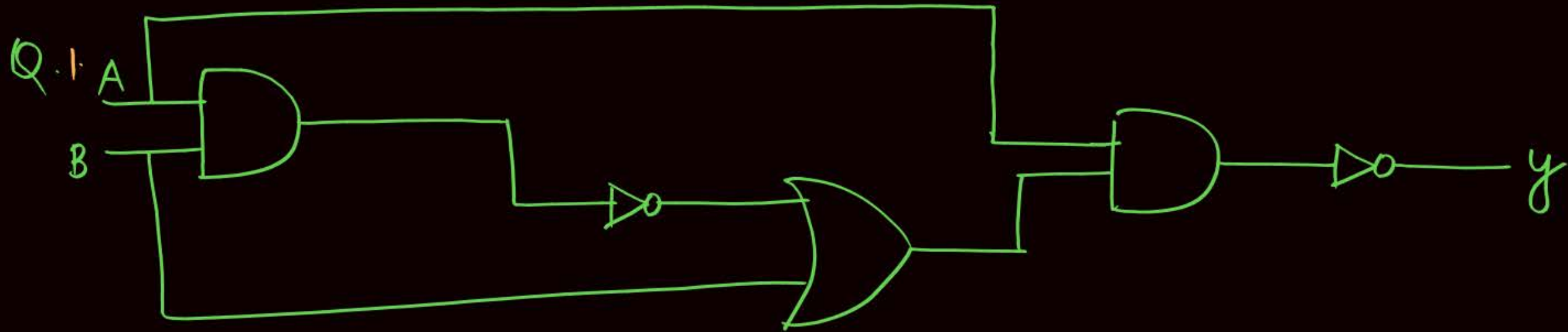
- Enable and disable input for 'OR gate' :



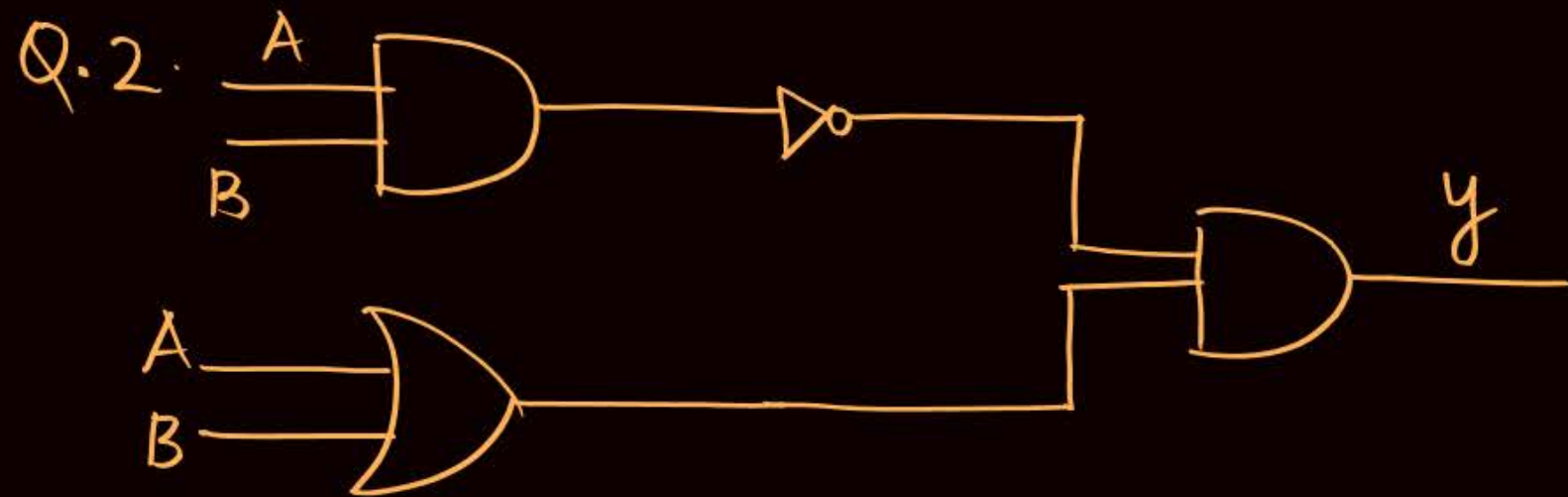
- IMP Points :

- If any one of the i/p line is at logic '1' then irrespective of other i/p lines o/p will be '1'.
- O/P will be '0' only in one case when all the i/p lines will be at logic '0'.

H.W.



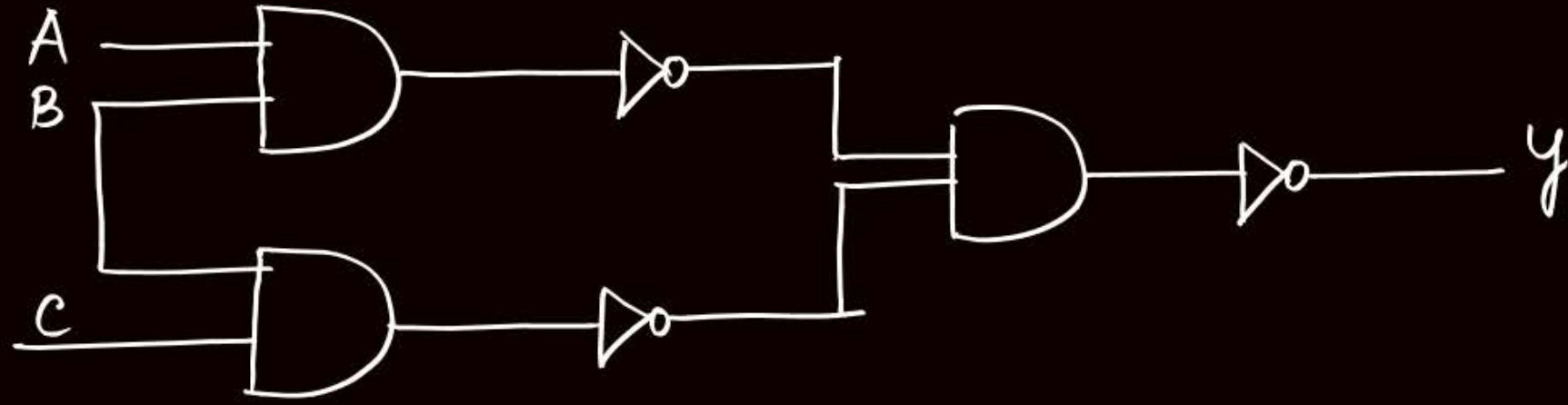
O/P y of the ckt ?



find o/p the y in SOP and POS format.



Q.



→ find out the O/P  $y$  and implement it using AND & OR only.



## 2 Minute Summary

→ Basic gates & its properties.

**Thank you**

**GW**  
*Soldiers !*

