*	Data Structures and Algorithms (DSA) - Assignment Number - 1
GUARRIAN III II I	A DESCRIPTION OF THE PROPERTY
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Q-1.	What is history? Explain different methods of hist function calculation.
->	Hashing: Hashing: Stashing is a technique which use less key comparisons and search the dement in O(n) times in the wordst case and in an awarage case it will be done in O(1) time. This method generally used the hash function to map the key into a table, which is called a hash table.
	Hashing is a technique which use less key comparisons and
	ANYPARE CASE it will be done in O(1) time This method agreeably used
	the hash function to map the key into a table, which is called
0	s hush table.
	Types of hish function:
	Division method-
	Division method- In this the hash function is dependent upon the remainder of a division. For example: - if the record 52, 58, 99, 84, is to be placed in a hash table and let us take the table size is 10.
	in a hush table and let us take the table size is 10.
	Then: - h(key) = record% table size.
	1 2 = 52 % 10
	9 = 99% 10.

3 Mid Square method:
In this method firstly key is squared and then mid part of the result is taken as the findex. For example: lonsider that if we want to place a record of \$3101 and the size of table is \$1000. So 3101* 3101 = \$9616201 j. e. h (3101) = 162 (middle 3 digit) are combined to produce a hash key for example: consider a record of 124 65512 then it will be divided into parts i. £ 124,655 After dividing the part combine these parts by adding it. H(ky) = 124 + 655 + 120 = 791.Q-2 Construct Hash July of size 15 and resolve collision using open addressing technique linear problem and use hash function $h(n) = h \mod 15$. : Clace 35 on 5th position of hish table. 35 % 105 = 5 36% 105 = 6; Place 36 on 6th position of hush table.

25% 105 = 10; Place 25 on 10th position of hush table.

47% 15 = 2; Place 47 on 2nd position of hush table.

501% 15 = 11; Place 2501 on 11 on position of hush table. 2501%15 =11

			Page 1		
	65	% 15			
	A. ~ 11		Using Linear position check next free location linearwise.		
	198 7 Jn.,				
	29°	%15 =	14: Place 29 on 11 th novition of hush table		
	16% 15 = 1; Clace 16 on 1 st position lox hish table.				
	14	% 15 =(14)= collision		
•	As oth i	i hext	14; Clace 29 on 14th position of hush table. 1; Clace 16 on 1 st position lot hush table. 14; collision. 14 the collision check react free location linearcoise free location: I place 14 on 0th location of hush table.		
			for the state of t		
	99	%15 =	=(9)= collision		
	de 10 th	99%15 = 9 collision. : Using Linear probing check next free location linearwise. As 12th is next free location: Close & & on 12th location of his table.			
	NO 12 M	us mus	a fru () want auge to on 12 un want of his hour.		
		14			
	1	16	: The hash table is as follows:-		
	2	47			
9	3				
	5	35			
	6	36			
	7	65			
	8	100			
	9	129			
	11	2501			
	12	99			
	13				
	14)	29			