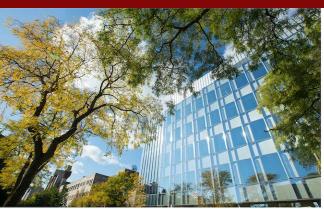
MPCS 52060 - Parallel Programming M6: Advanced Parallel Patterns and Techniques (Part 2)









Original slides from "The Art of Multiprocessor Programming" by Maurice Herlihy & Nir Shavit with modifications by Lamont Samuels



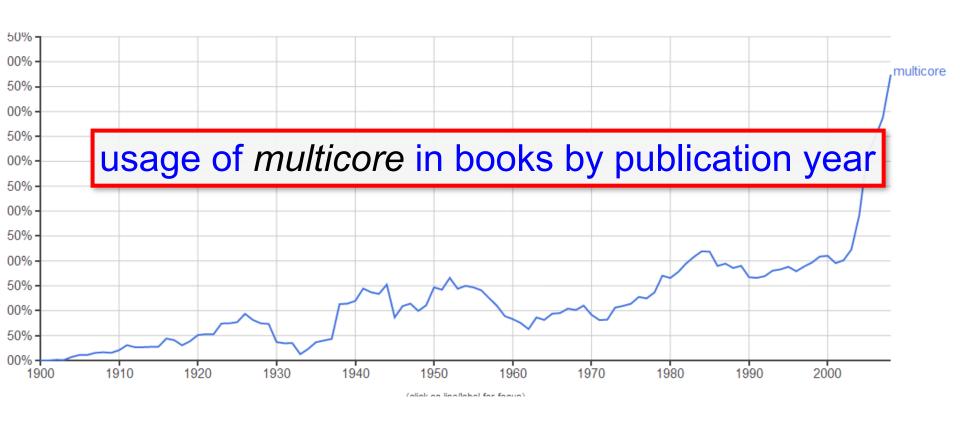
Data Parallelism

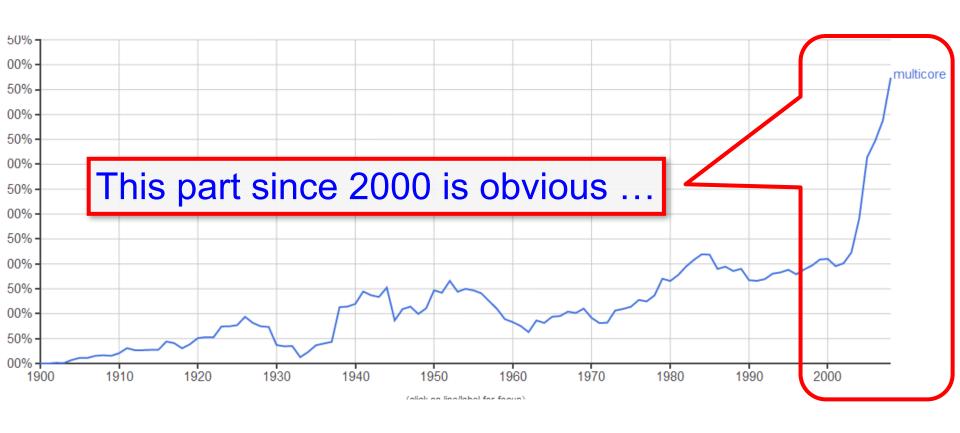
Ever Wonder ...

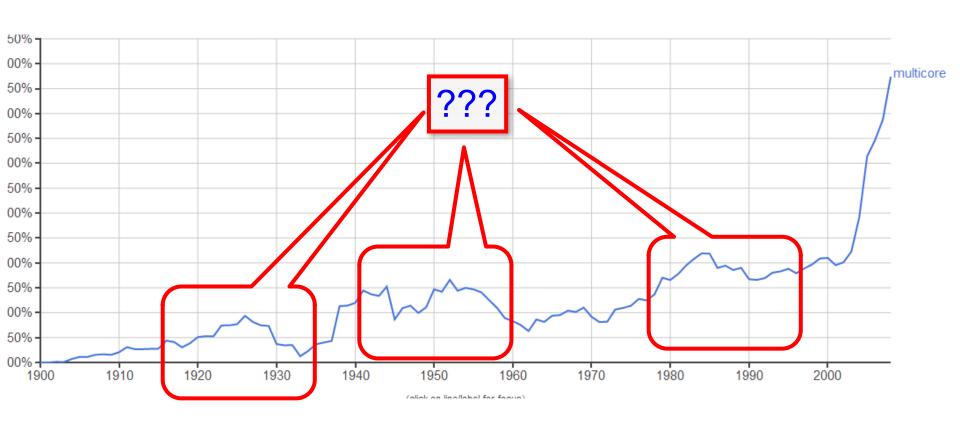
When did the term *multicore* become popular?

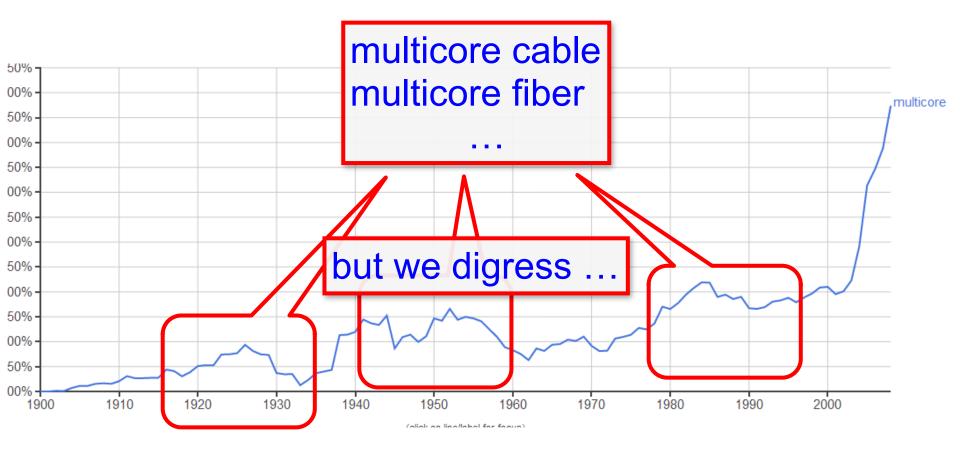
A multi-core processor is a single computing component with two or more independent actual central processing units, which are the units that read and execute program instructions.)

wikipedia

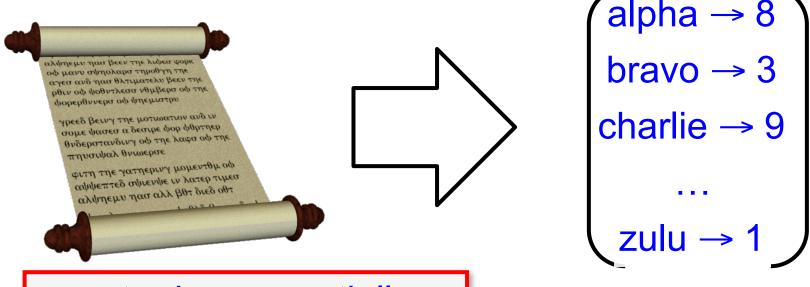








WordCount



easy to do sequentially ...

what about in parallel?

MapReduce

split text among mapping threads



chapter 1

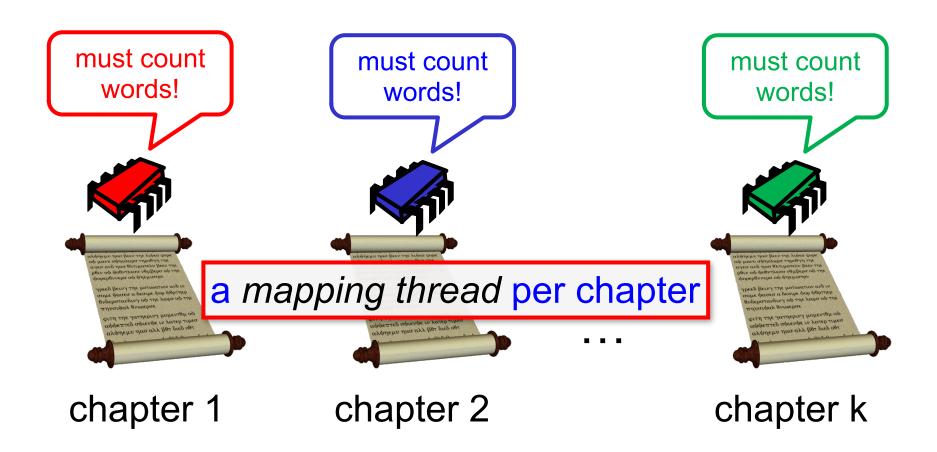


chapter 2

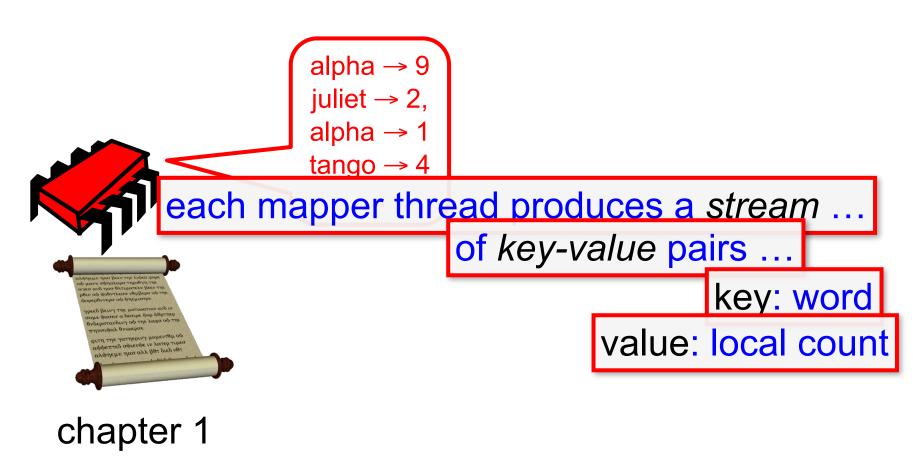


chapter k

Map Phase



Map Phase



```
abstract class Mapper<IN, K, V>
   extends RecursiveTask<Map<K, V>> {
   IN input;
   public void setInput(IN anInput) {
      input = anInput;
   }
}
```

```
abstract class Mapper<IN, K, V>
   extends RecursiveTask<Map<K, V>> {
   IN input;
   public void setInput(IN anInput) {
      input = anInput;
   }
}
```

input: document fragment

```
abstract class Mapper<IN, K, V>
   extends RecursiveTask<Map<K, V>> {
   IN input;
   public void setInput(IN anInput) {
      input = anInput;
   }
}
key: individual word
```

```
abstract class Mapper<IN, K, V>
   extends RecursiveTask<Map<K, V>> {
   IN input;
   public void setInput(IN anInput) {
      input = anInput;
   }
}
```

value: local count

```
abstract class Mapper<IN, K, V>
   extends RecursiveTask Map<K, V>> {
   IN input;
   public void setInput(IN anInput) {
     input = anInput;
   }
}
```

a task that runs in parallel with other tasks

```
abstract class Mapper<IN, K, V>
   extends RecursiveTask<Map<K, V>>
   IN input;
   public void setInput(IN anInput) {
     input = anInput;
   }
}
```

produces a map: word → count

```
abstract class Mapper<IN, K, V>
   extends RecursiveTask<Map<K, V>> {
   IN input;

   public void setInput(IN anInput) {
     input = anInput;
   }
}
```

initialize input: which document fragment?

```
class WordCountMapper extends
  mapreduce.Mapper<
      List<String>, String, Long
      > {
        ...
    }
```

```
class WordCountMapper extends
    mapreduce Mapper
List<String>, String, Long
> {
    document fragment is
list of words
```

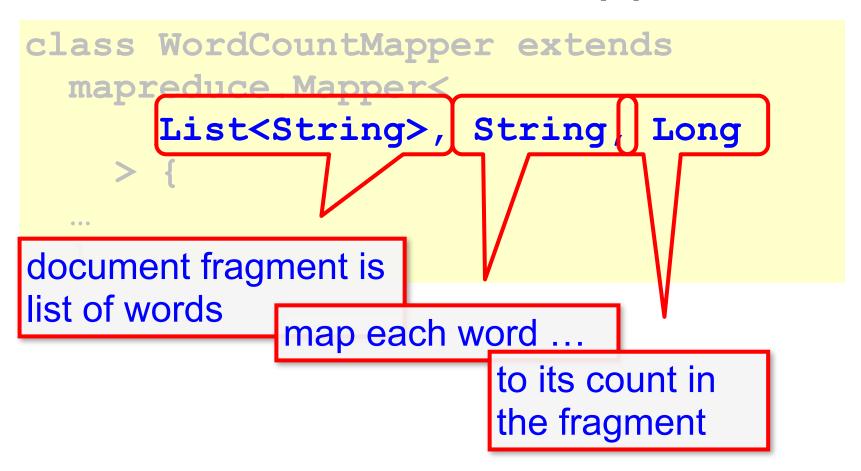
```
class WordCountMapper extends

mapreduce Mapper

List<String>, String, Long

document fragment is
list of words

map each word ...
```



```
Map<String,Long> compute() {
  Map<String,Long> map = new HashMap<>();
      for (String word : input) {
        map.merge(word,
                    1L,
                    (x, y) \rightarrow x + y);
      return map;
```

```
Map<String,Long> compute()
  Map<String,Long> map = new HashMap<>();
       for (String word : input) {
         map.merge (word,
                      (x, y) \rightarrow x + y);
    the compute () method constructs the
    local word count
```

```
Map<String, Long> compute
  Map<String,Long> map = new HashMap<>();
                           input)
        map.merge(word,
                    1L,
                    (x, y)
      return map;
       create a map to hold the output
```

```
Map<String,Long> compute() {
      for (String word : input)
         map.merge(word
                    1L,
      return map;
                       examine each word in the
                       document fragment
```

```
Map<String,Long> compute() {
  Map<String,Long> map = new HashMap<>();
       for (String word : input) {
          map.merge(word,
                       1L,
                       (x, y) \rightarrow x + y);
       return map;
          increment that word's
          count in the map
                                                27
                   ALL OF MURICIPIOCESSOL FTOGRAFILLING
```

```
Map<String,Long> compute() {
  Map<String,Long> map = new HashMap<>();
       for (String word : input) {
         map.merge(word,
                     1L,
                     (x, y) \rightarrow x + y);
                         when the local count is
      return map;
                         complete, return the map
```

Reduce Phase



a reducer thread merges mapper outputs

```
alpha → 2
juliet → 1
tango → 1
...
```

```
alpha → 1
foxtrot → 1
papa→ 1
tango → 1
```

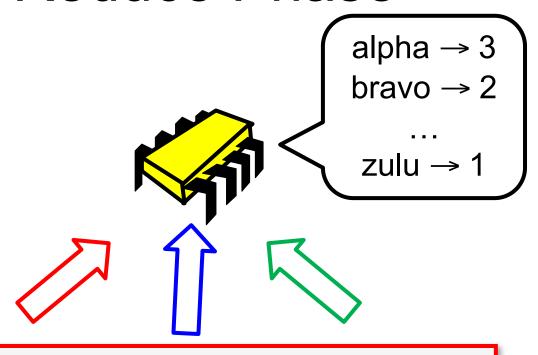
alpha → 1 oscar → 1, bravo → 2...







Reduce Phase



the reducer task produces a stream

of key-value pairs ...

key: word

value: word count

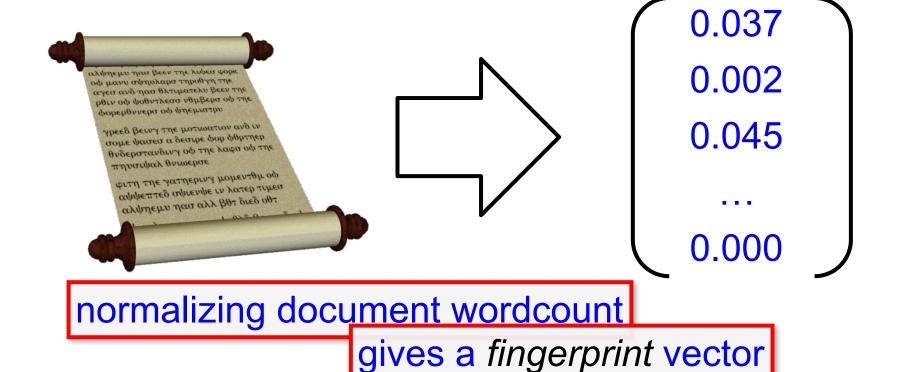
```
abstract class Reducer<K, V, OUT>
    extends RecursiveTask<OUT> {
 K key;
  List<V> valueList;
 public void setInput(
    K aKey,
    List<V> aList) {
    key = aKey;
    valueList = aList;
```

```
abstract class Reducer<K,
    extends RecursiveTask<OUT
 K key;
          valueList;
            setInput(
  public
    K aKey,
                     each reducer is given
    List<V> aList)
                     a single key (word)
    key = aKey;
    valueList = aList;
```

```
abstract class Reducer<K,
    extends RecursiveTask<OUT>
  K key;
  List<V> valueList;
           old set input (
    K aKey,
    List<V> aList
    key = aKey;
    valueLis
              and a list of associated values
              (word count per fragment)
```

```
abstract class Reducer<K,
    extends RecursiveTask<OUT>
  K key;
  List<V> valueList;
  public void setInput
    K aKey,
    List<V> aList)
    key = aKey;
    value It produces a single summary value
         (the total count for that word)
```

WordCount



Pseudocode Implementations

- See m6/concurrent/mapreduce/MapReduce.java
- See m6/concurrent/mapreduce/WordCount.java

Simple Video Game

- Prepare frame for display
 - By graphics coprocessor
- "soft real-time" application
 - Need at least 35 frames/second
 - OK to mess up rarely

Simple Video Game

```
while (true) {
  frame.prepare();
  frame.display();
}
```

Simple Video Game

```
while (true) {
  frame.prepare();
  frame.display();
}
```

- What about overlapping work?
 - 1st thread displays frame
 - 2nd prepares next frame

Two-Phase Rendering

```
while (true) {
  if (phase) {
    frame[0].display();
  } else {
    frame[1].display();
  }
  phase = !phase;
}
```

```
while (true) {
  if (phase) {
    frame[1].prepare();
  } else {
    frame[0].prepare();
  }
  phase = !phase;
}
```

Two-Phase Rendering

```
while (true) {
                         while (true) {
 if (phase) {
                          if (phase) {
  frame[0].display();
                            frame[1].prepare();
 frame[1].display();
                            frame[0].prepare();
 phase = !phase;
                                 =!phase;
```

even phases

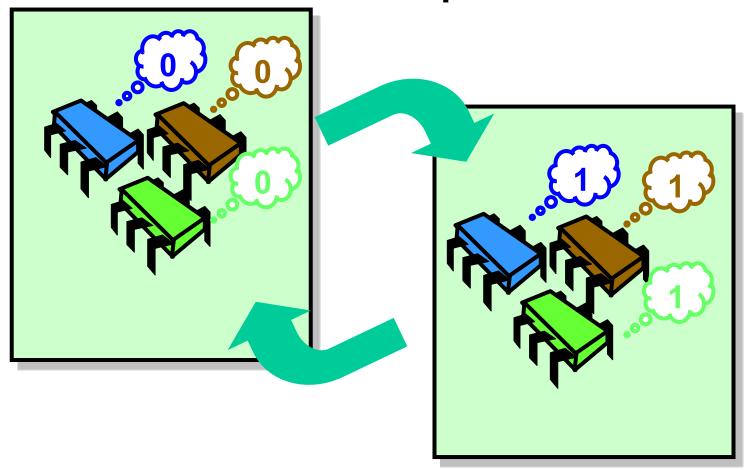
Two-Phase Rendering

```
while (true) {
                            while (true) {
 if (phase) {
                             if (phase) {
  frame[0].display();
                              frame[1].prepare();
 } else {
                             } else {
  frame[1].display();
                              frame[0].prepare();
 phase = !phase;
                                nase =
                        odd phases
```

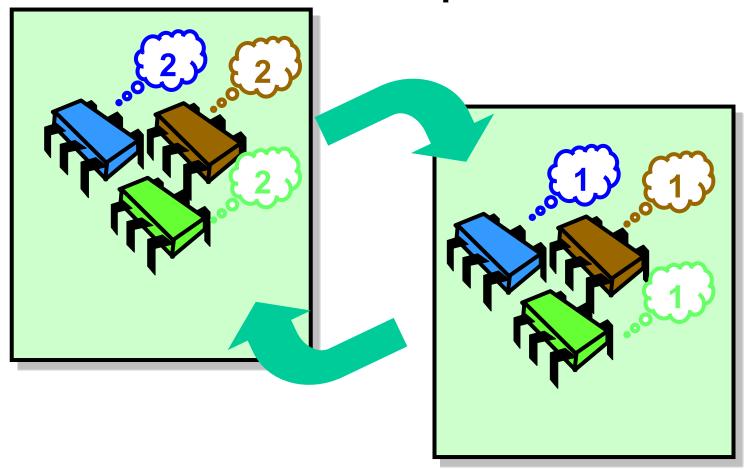
Synchronization Problems

- How do threads stay in phase?
- Too early?
 - "we render no frame before its time"
- Too late?
 - Recycle memory before frame is displayed

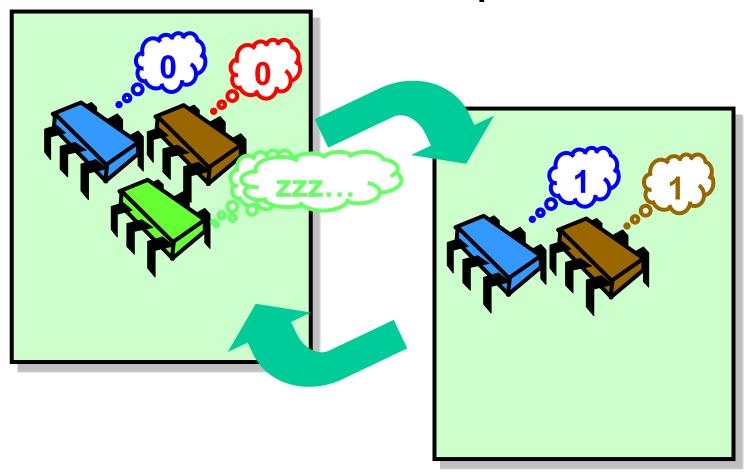
Ideal Parallel Computation



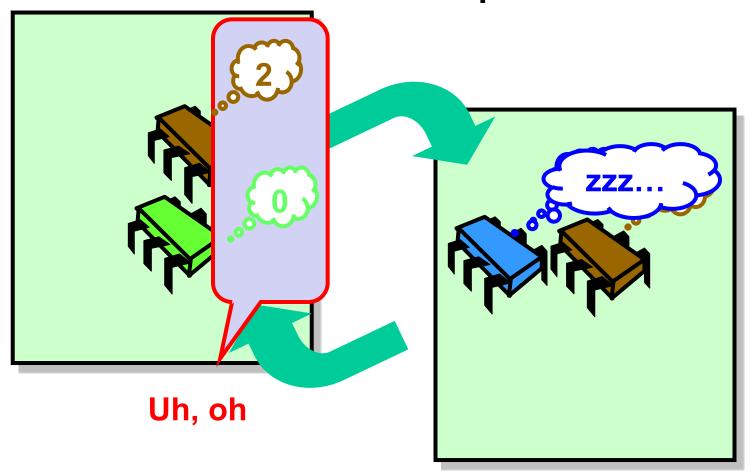
Ideal Parallel Computation



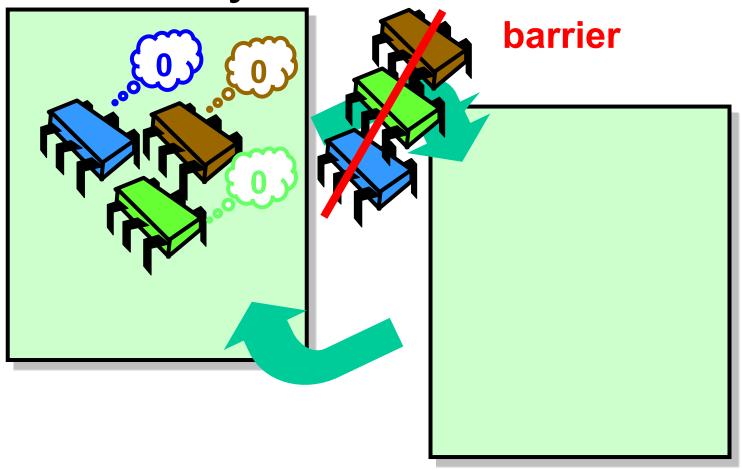
Real-Life Parallel Computation



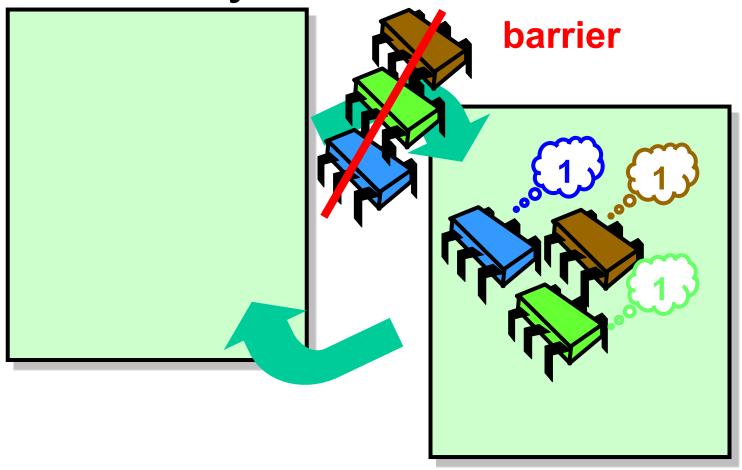
Real-Life Parallel Computation



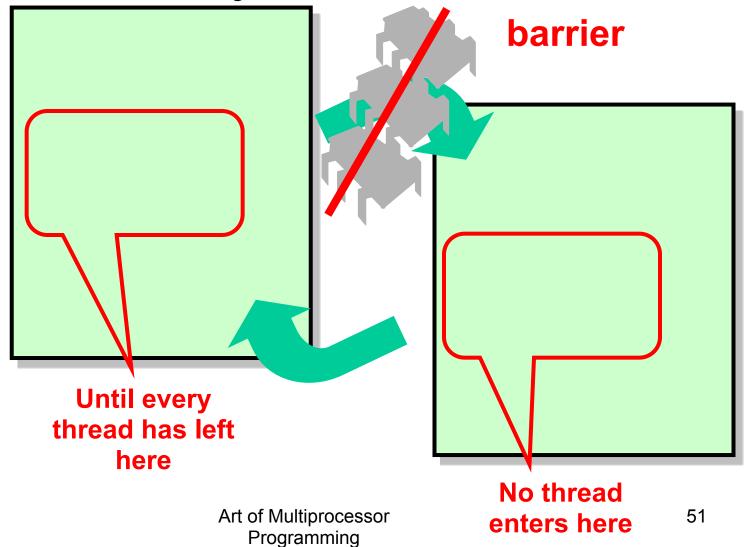
Barrier Synchronization



Barrier Synchronization



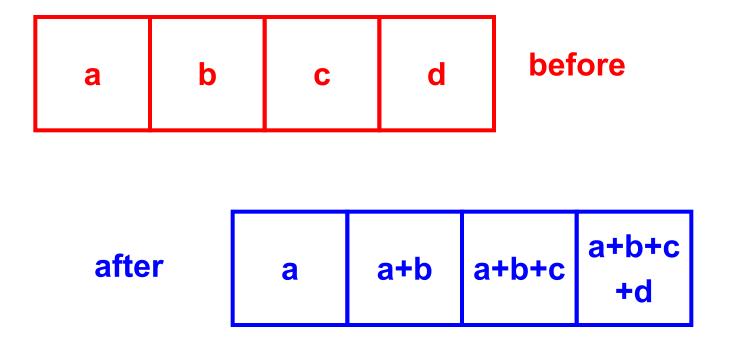
Barrier Synchronization



Why Do We Care?

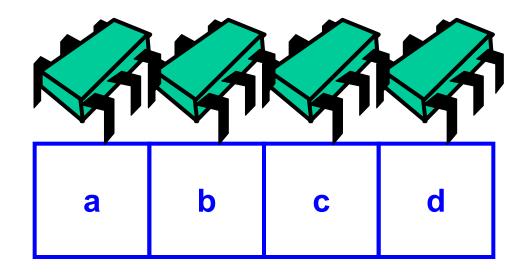
- Mostly of interest to
 - Scientific & numeric computation
- Elsewhere
 - Less common in systems programming
 - Still important topic

Example: Parallel Prefix

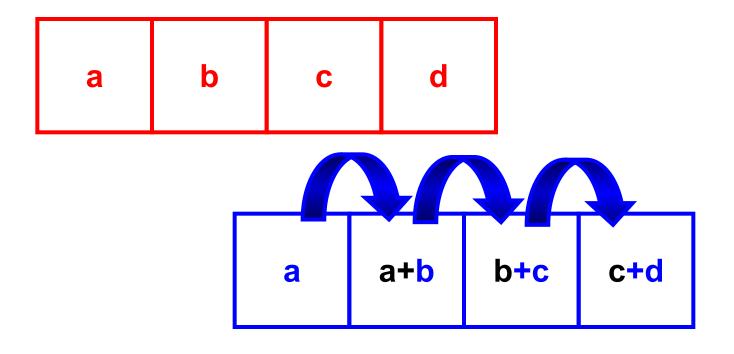


Parallel Prefix

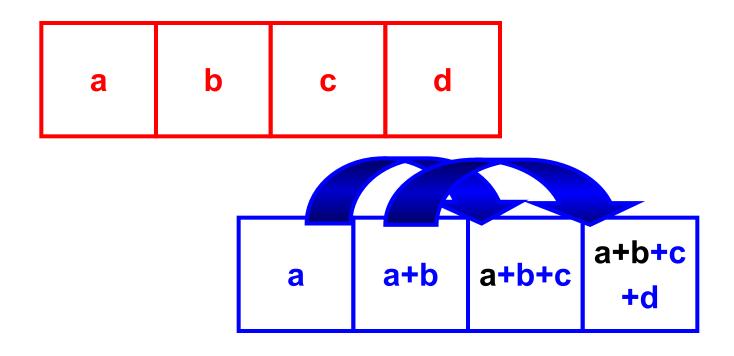
One thread Per entry



Parallel Prefix: Phase 1



Parallel Prefix: Phase 2



Parallel Prefix

- N threads can compute
 - Parallel prefix
 - Of N entries
 - In log₂ N rounds
- What if system is asynchronous?
 - Why we need barriers

```
class Prefix extends Thread {
 int[] a;
 int i;
 Barrier b;
void Prefix(int[] a,
               Barrier b, int i) {
  a = a;
  b = b;
  i = i;
```

```
class Prefix extends Thread {
int[] a;
int i;
 Barrier
void Prefi
              (int[] a,
                Barrier b, int i) {
  a = a;
                Array of input
  b = b;
                   values
  i = i;
```

```
class Prefix extends Thread {
 int[] a;
int i;
 Barrie
        efix(int[] a,
void Pr
                Barrier b, int i) {
  a = a;
 b = b;
              Thread index
 i = i;
```

```
class Prefix extends Thread {
 int[] a;
 int i;
Barrier b;
        afik(int[] a,
               Barrier b, int i) {
  a = a;
             Shared barrier
  b = b;
  i = i;
```

```
class Prefix extends Thread {
 int[] a;
 int i;
 Barrier b;
                          Initialize fields
 void Prefix(int[] a,
                     ier b, int i) {
```

Where Do the Barriers Go?

```
public void run() {
  int d = 1, sum = 0;
  while (d < N) {
   if (i >= d)
      sum = a[i-d];
   if (i >= d)
      a[i] += sum;
   d = d * 2;
}}
```

Where Do the Barriers Go?

```
public void run() {
 int d = 1, sum = 0;
 while (d < N) {
  if (i >= d)
    sum = a[i-d];
                   Make sure everyone reads
  b.await();
                      before anyone writes
  1f (1 >= d)
   a[i] += sum;
  d = d * 2;
}}}
```

Where Do the Barriers Go?

```
public void run() {
 int d = 1, sum = 0;
 while (d < N) {
  if (i >= d)
    sum = a[i-d];
                   Make sure everyone reads
  b.await();
                      before anyone writes
  1f (1 >= d)
   a[i] += sum;
                   Make sure everyone writes
  b.await()
                      before anyone reads
}}}
```

Barrier Implementations

- Cache coherence
 - Spin on locally-cached locations?
 - Spin on statically-defined locations?
- Latency
 - How many steps?
- Symmetry
 - Do all threads do the same thing?

```
public class Barrier {
 AtomicInteger count;
 int size;
 public Barrier(int n) {
  count = AtomicInteger(n);
  size = n;
 public void await() {
  if (count.getAndDecrement()==1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } } }
```

```
public class Barrier
AtomicInteger count;
 IIIL SIZE,
 public Barrier (1)
  count = AtomicInteg
  size = n;
                          Number of threads
                           not yet arrived
 public void await() {
  if (count.getAndDecrement() == 1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } }
```

```
public class Barrier {
 AtomicInteger count;
 int size;
                          Number of threads
                              participating
  count = AtomicInteger(n);
  size = n;
 public void await() {
  if (count.getAndDecrement() == 1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } }
```

```
public class Barrier {
                            Initialization
 AtomicInteger count;
 int size;
 public Barrier (int
  count = AtomicInteger(n);
  size = n;
 public void await() {
  if (count.getAndDecrement()==1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } }
```

```
public class Barrier {
 AtomicInteger count;
                            Principal method
 int size;
 public Barrier(int n) {
  count = AtomicInteger()
  size = n;
 public void await()
  <u>if (count.getAndDecre</u>ment() == 1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } }
```

```
public class Barrier {
 AtomicInteger count; If I'm last, reset fields
                           for next time
 int size;
 public Barrier(int n) {
  count = AtomicInteger(n);
  size = n;
 public void await()
  if (count.getAndDecrement() =
   count.set(size);
   while (count.get() != 0);
```

```
public class Barrier {
 AtomicInteger count;
 int size;
                          Otherwise, wait for
 public Barrier(int n) {
                           everyone else
  count = AtomicInteger(n)
  size = n;
 public void await()
  if (count.getAndDecrement() == 1) {
   count.set(size)
   alsa (
   while (count.get() != 0);
```

```
public class Barrier {
 AtomicInteger count;
 int size;
 public Barrier(int n) {
  count = AtomicInteger(n);
  size = n:
    What's wrong with this protocol?
 public void await()
  if (count.getAndDecrement()==1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } } }
```

Reuse

```
Barrier b = new Barrier(n);
while (mumble()) {
work();
b.await()
synchronize
repeat
```

```
public class Barrier {
 AtomicInteger count;
 int size;
 public Barrier(int n) {
  count = AtomicInteger(n);
  size = n;
 public void await() {
  if (count.getAndDecrement()==1) {
   count.set(size);
  } else {
   while (count.get() != 0);
 } } } }
```

```
public class Barrier
 AtomicInteger count
                         Waiting for
 int size;
                       Phase 1 to finish
 public Barrier (int n)
  count = AtomicInteger(n);
  size = n;
 public void await() {
  if (count.getAndDecrement
   count.set(size);
  } else {
   while (count
```

```
Phase 1
public c
                           Waiting for
 AtomicIr is so over
                        Phase 1 to finish
           rrier(int n) {
             .tomicInteger(n);
      unt.set(size):
    else {
   while (count.get()
                                               78
```

```
public c
              Prepare for
               phase 2
                              ZZZZZ.
 AtomicI
 int sizo
            arrier(int n) {
             AtomicInteger(n);
            oid await() {
       (count.getAndDecrement()
    <del>c</del>ount.<del>set(size)</del>
    while (count.get()
                                                   79
```

Uh-Oh

```
public class Barri
                        Waiting for
 AtomicInteger
                      Phase 2 to finish
 int size;
          arrior (int
 pub
            Atomic
                       Waiting for
                     Phase 1 to finish
 public vold await() {
  if (count\getAndDecrement
   count.set(size);
    else
   while (count.get
```

Basic Problem

- One thread "wraps around" to start phase 2
- While another thread is still waiting for phase 1
- One solution:
 - Always use two barriers

```
public class Barrier {
 AtomicInteger count;
 int size;
 volatile boolean sense = false;
 threadSense = new ThreadLocal<br/>boolean>...
 public void await {
  boolean mySense = threadSense.get();
  if (count.getAndDecrement()==1) {
   count.set(size); sense = mySense
  } else {
   while (sense != mySense) {}
 threadSense.set(!mySense)}}
```

```
Completed odd or
public class Barrier {
                                  even-numbered
 AtomicInteger count;
                                      phase?
 int size;
volatile boolean sense = false;
 public void await {
 boolean mySense = threadSense.get();
  if (count.getAndDecrement()==1) {
   count.set(size); sense = mySense
  } else {
   while (sense != mySense) {}
 threadSense.set(!mySense)}}
```

```
public class Barrier {
                                Store sense for
 AtomicInteger count;
                                  next phase
 int size;
 volatile boolean sense =
threadSense = new ThreadLocal<boolean>...
 public void await {
  boolean mySense = threadSense.get();
  if (count.getAndDecrement()==1) {
   count.set(size); sense = mySense
  } else {
   while (sense != mySense) {}
 threadSense.set(!mySense)}}
```

```
public class Barrier {
AtomicInteger count;
                     Get new sense determined
 int size;
volatile boolean sense = fby last phase
 threadSense = new ThreadLocal<poolean>...
 public void await {
 boolean mySense = threadSense.get();
   f (count getAndDegrement() -- 1)
   count.set(size); sense = mySense
  } else {
   while (sense != mySense) {}
 threadSense.set(!mySense)}}
```

```
public class Barrier {
 AtomicInteger count;
 int size;
                        If I'm last, reverse sense
 volatile boolean sense
                              for next time
 threadSense = new Thre
 public void await {
  boolean mySense = threadSense.vet();
  if (count.getAndDecrement()==1) {
   count.set(size); sense = mySense
    else d
   while (sense != mySense) {}
 threadSense.set(!mySense)}}
```

```
public class Barrier {
 AtomicInteger count;
                              Otherwise, wait for
 int size;
                                 sense to flip
 volatile boolean sense = f
 threadSense = new ThreadLocal \( \)boolean \( \)...
 public void await {
  boolean mySense = thread%ense.get();
  if (count.getAndDecrement()==1) {
   count.set(size); senge = my sense
    else {
   while (sense != mySense) {}
 threadSense.set(!mySense)}}
```

```
public class Barrier {
                              Prepare sense for
 AtomicInteger count;
                                 next phase
 int size;
 volatile boolean sense = fale
 threadSense = new ThreadLocal<br/>
<br/>
boolean>...
 public void await {
  boolean mySense = threadSense.get();
  if (count.getAndDecrepent()==1) {
   count.set(size); serise = mySense
  } else {
   while (sense != mySense)
 threadSense.set(!mySense)}}}
```