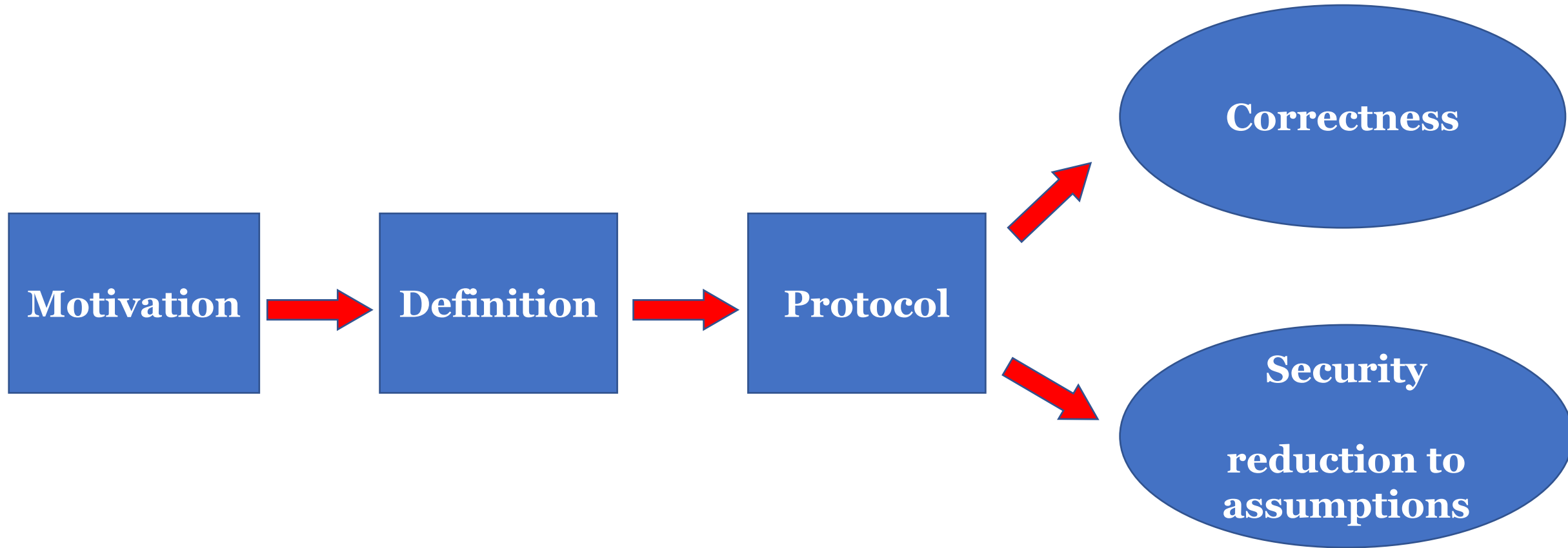


Secure Multi-Party Computation

Cryptography



I Reading antikontalk
Computer games???
to shopping more.



What do you do after
Computer games, ...
work?



Ideal scenario

Let's run a cryptographic
protöcöl.



We both
like sports!

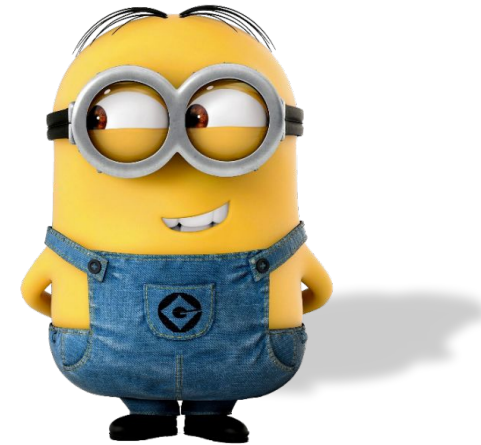
What do you do after
work?



Set A



Set B



Set $A \cap B$

Other elements remain secret to each other

Secure two-party computation

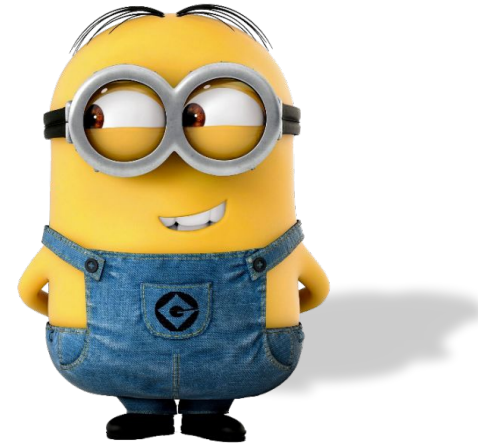
f

x



$f(x, y)$

y

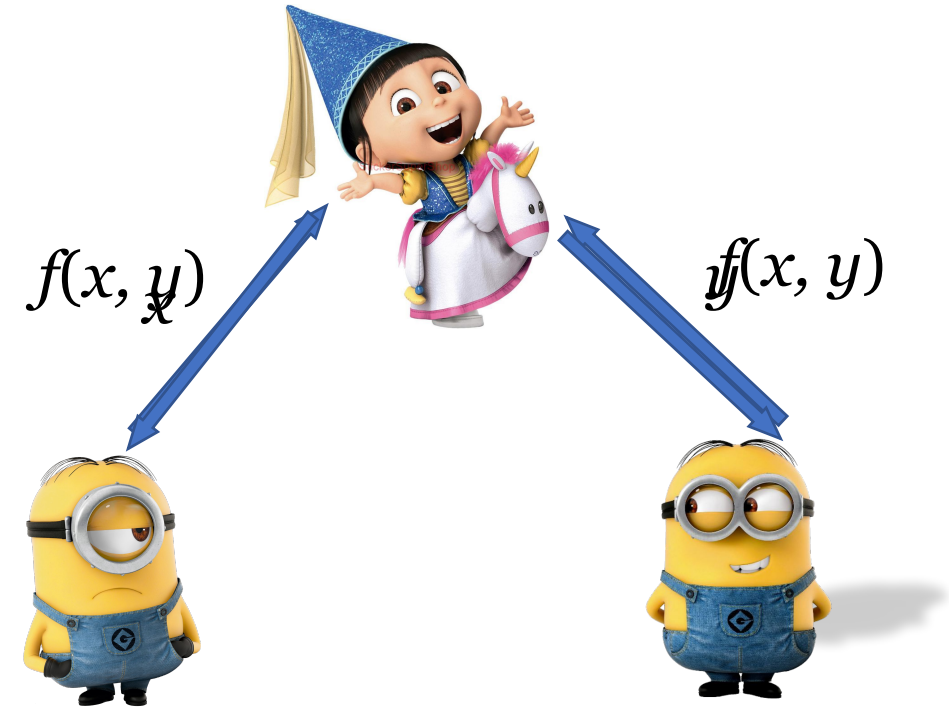
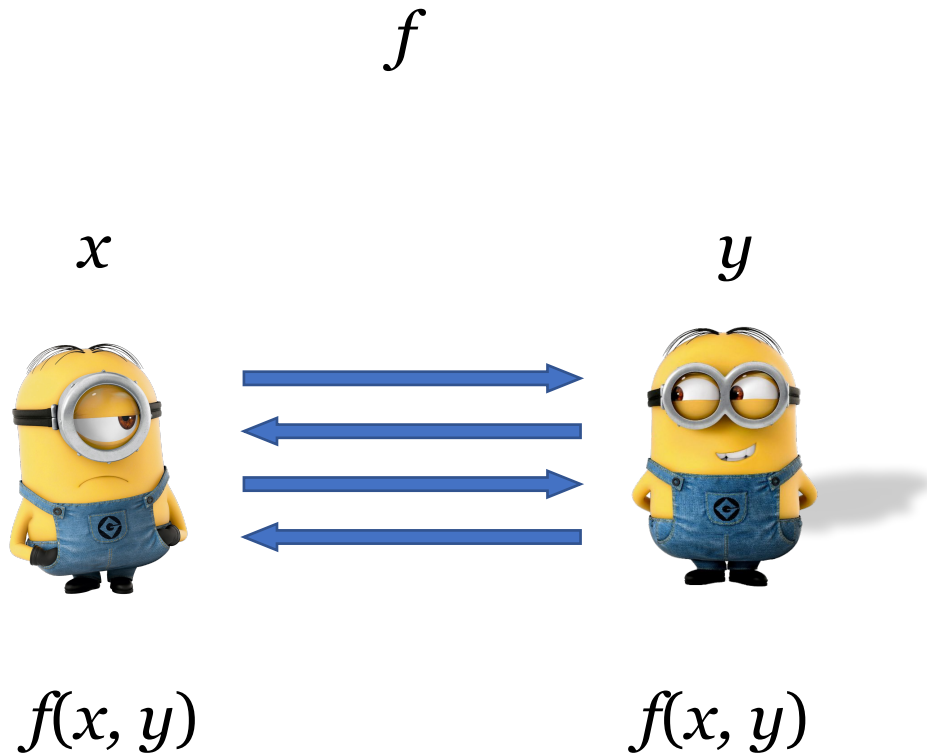


$f(x, y)$

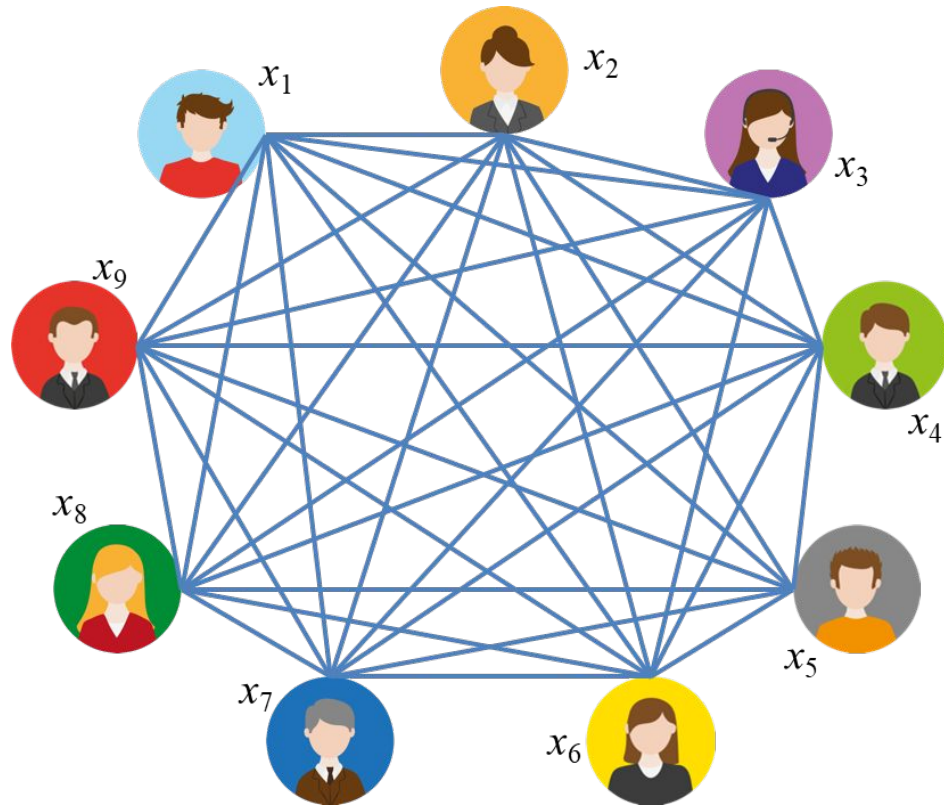


x and y remain secret

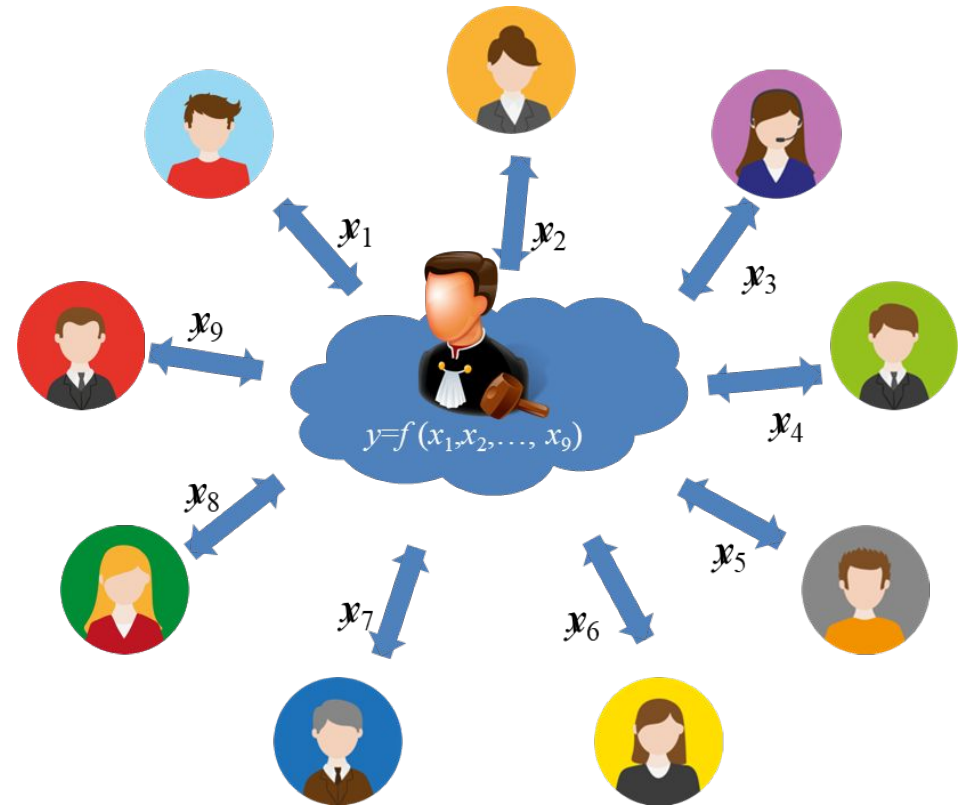
Secure two-party computation



Secure multi-party computation (MPC)



$$y = f(x_1, x_2, \dots, x_9)$$



Motivations: old days

- Millionaire problem
- Mental poker



Yao, Andrew Chi-Chih,
“Protocols for Secure Computations”,
Foundations of Computer Science (FOCS) 1982

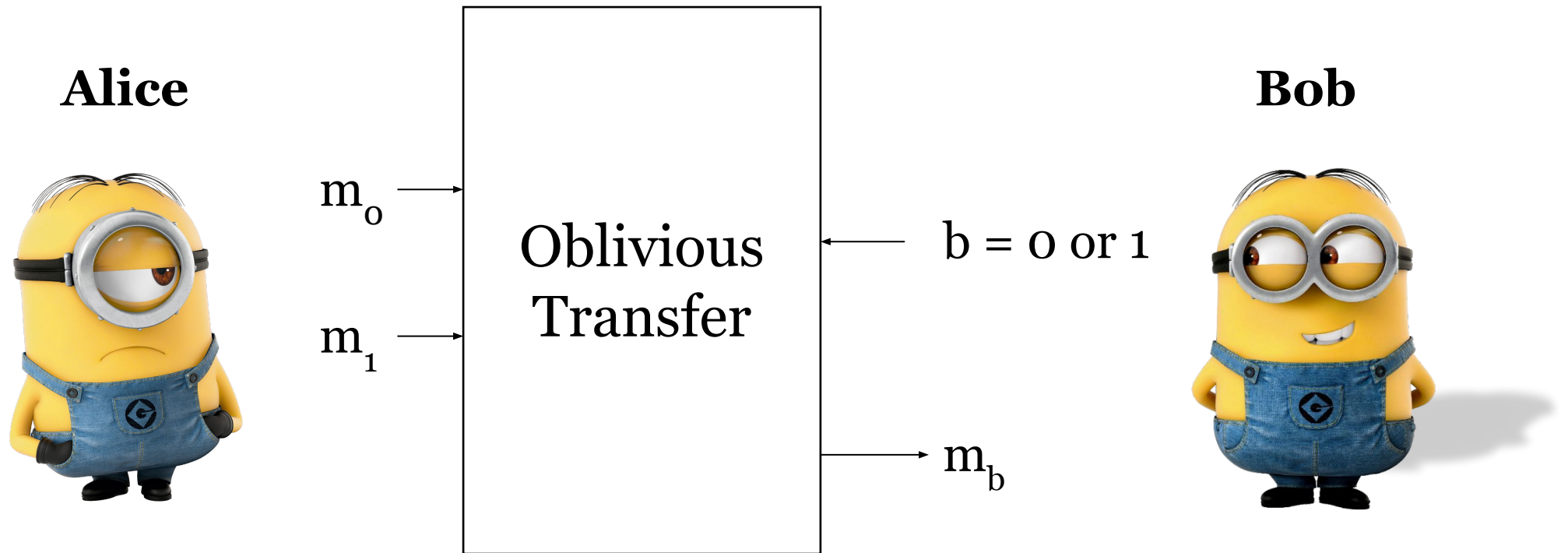
Motivations: now

- Auction
- Voting
- Privacy-preserving data mining/machine learning
- Distributed encryptions/signatures

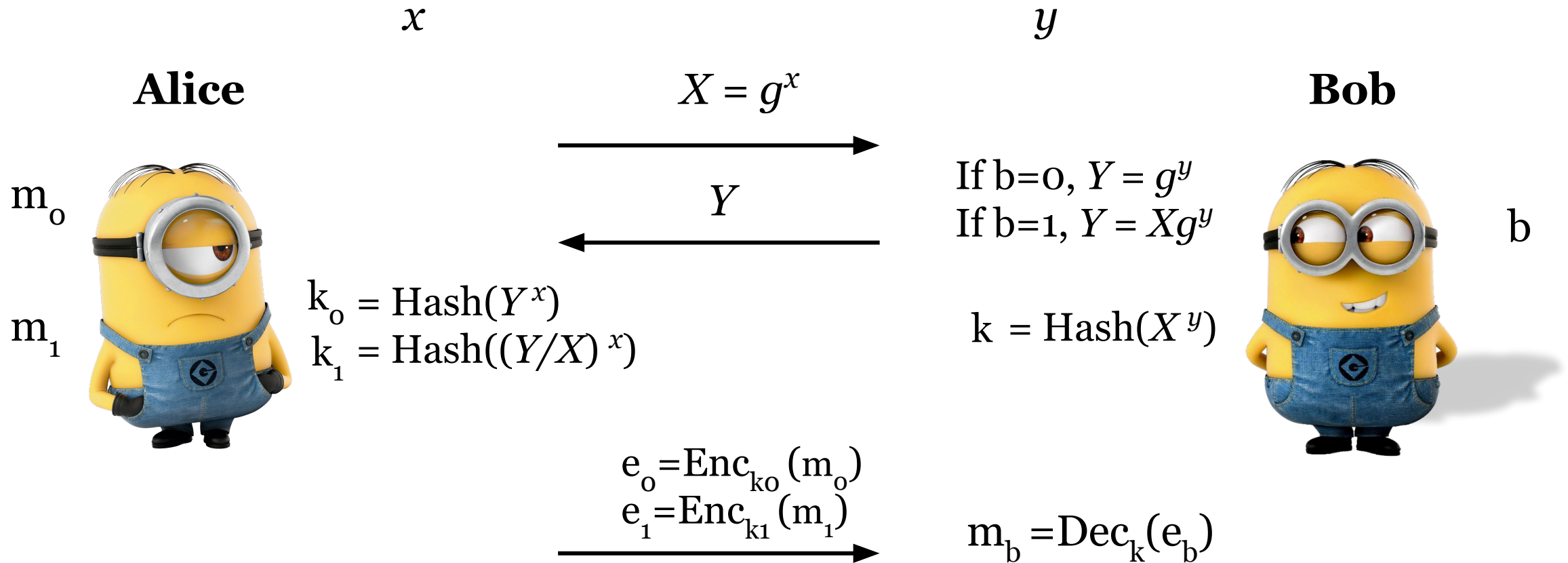
Existing solutions

- Yao's garbled circuit
- GMW protocol

Building block: oblivious transfer (OT)



Building block: oblivious transfer (OT)

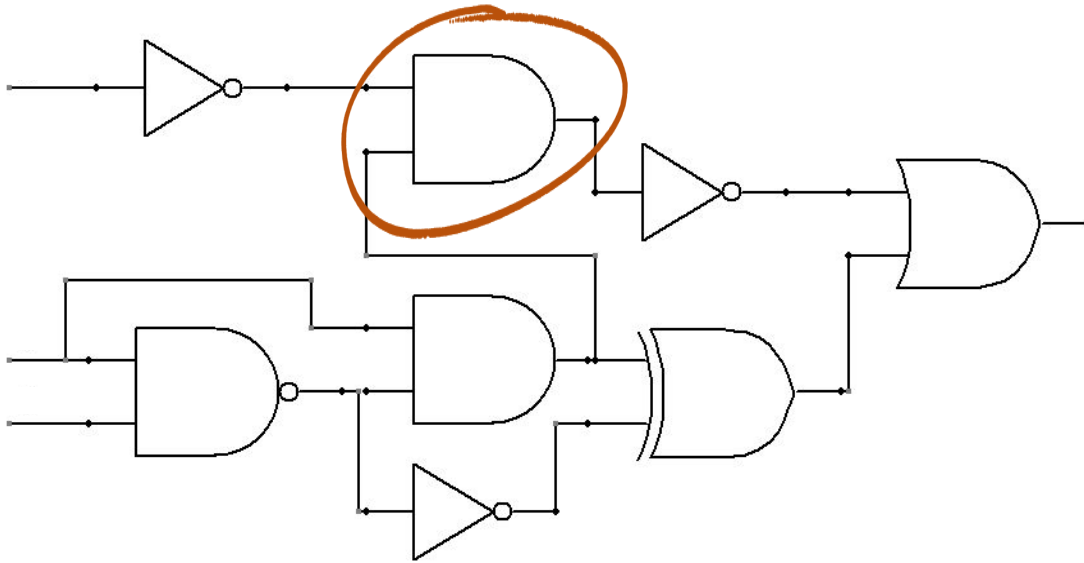


Yao's garbled circuit



Yao's garbled circuit

- What is a (Boolean) circuit?



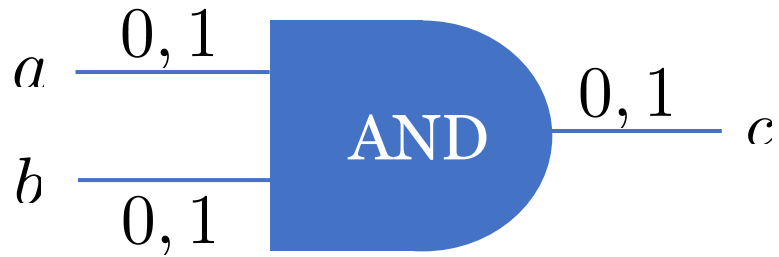
Wire <i>a</i>	Wire <i>b</i>	Wire <i>c</i>
0	0	0
0	1	0
1	0	0
1	1	1

Yao's garbled circuit

- Why circuit?
 - Captures all functions
 - Gate-by-gate paradigm

Yao's garbled circuit

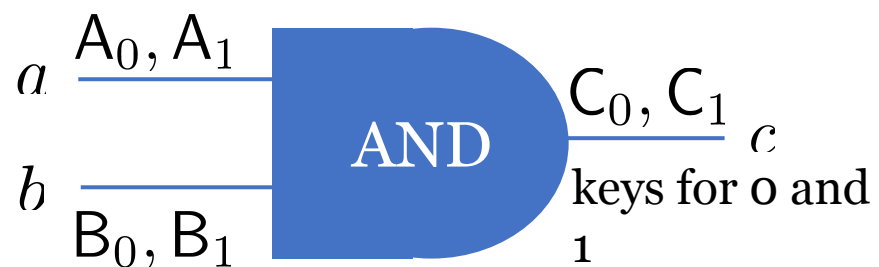
- How to garble a circuit?



Wire a	Wire b	Wire c
0	0	0
0	1	0
1	0	0
1	1	1

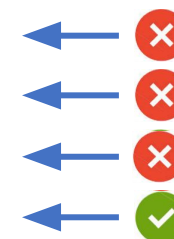
Yao's garbled circuit

- How to garble a circuit?



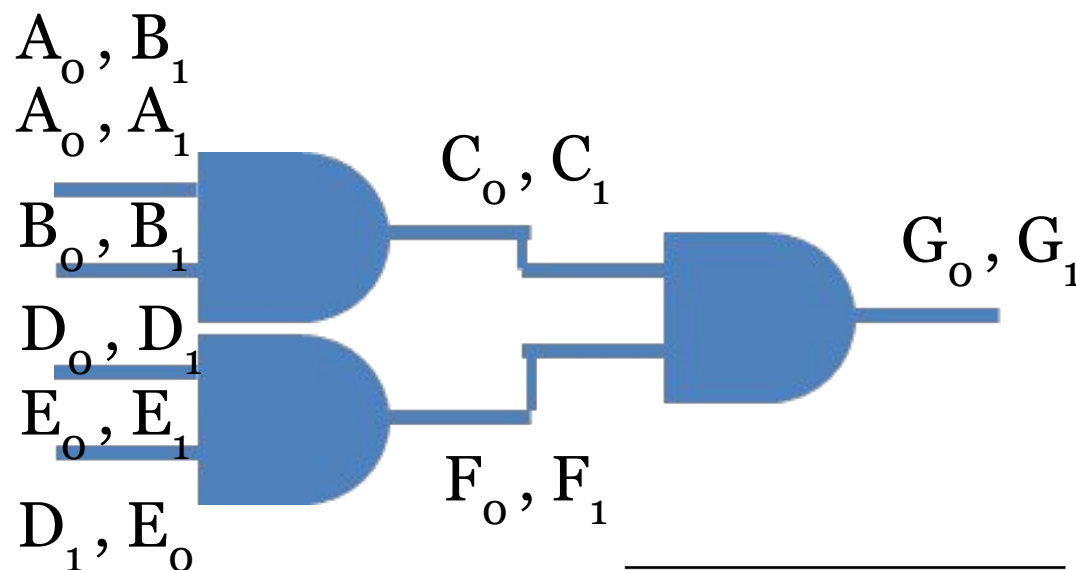
$A_1, B_0 \rightarrow C_0$
 $A_1, B_1 \rightarrow C_1$

Wire a	Wire b	Wire c
0	0	0
0	1	0
1	0	0
1	1	1



Yao's garbled circuit

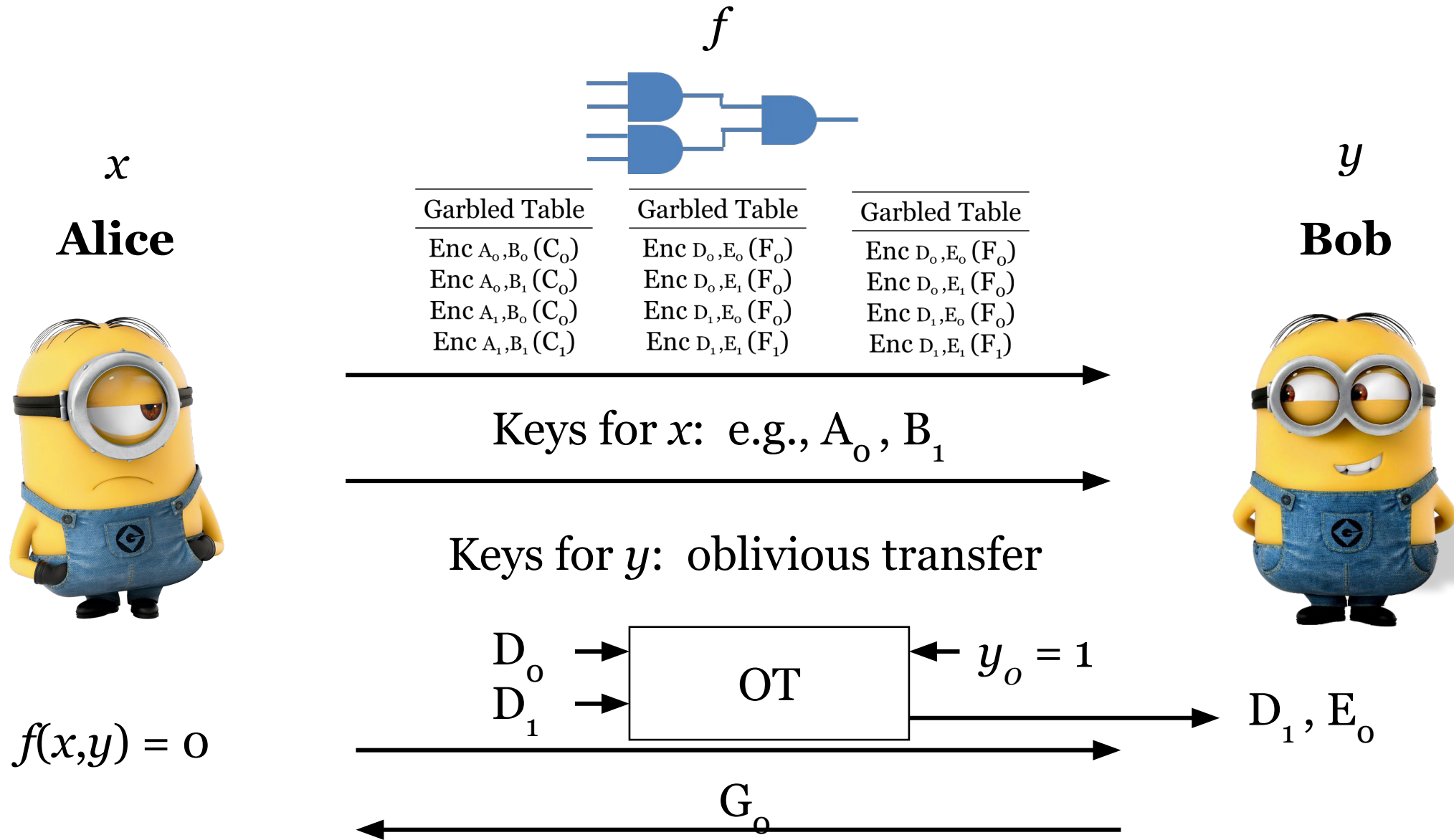
Garbled Table
$\text{Enc } A_0, B_0 (C_0)$
$\text{Enc } A_0, B_1 (C_0)$ ✓
$\text{Enc } A_1, B_0 (C_0)$
$\text{Enc } A_1, B_1 (C_1)$



Garbled Table
$\text{Enc } D_0, E_0 (F_0)$
$\text{Enc } D_0, E_1 (F_0)$
$\text{Enc } D_1, E_0 (F_0)$ ✓
$\text{Enc } D_1, E_1 (F_1)$

Garbled Table
$\text{Enc } C_0, F_0 (G_0)$ ✓
$\text{Enc } C_0, F_1 (G_0)$
$\text{Enc } C_1, F_0 (G_0)$
$\text{Enc } C_1, F_1 (G_1)$

Putting everything together



Properties of Yao's garbled circuit

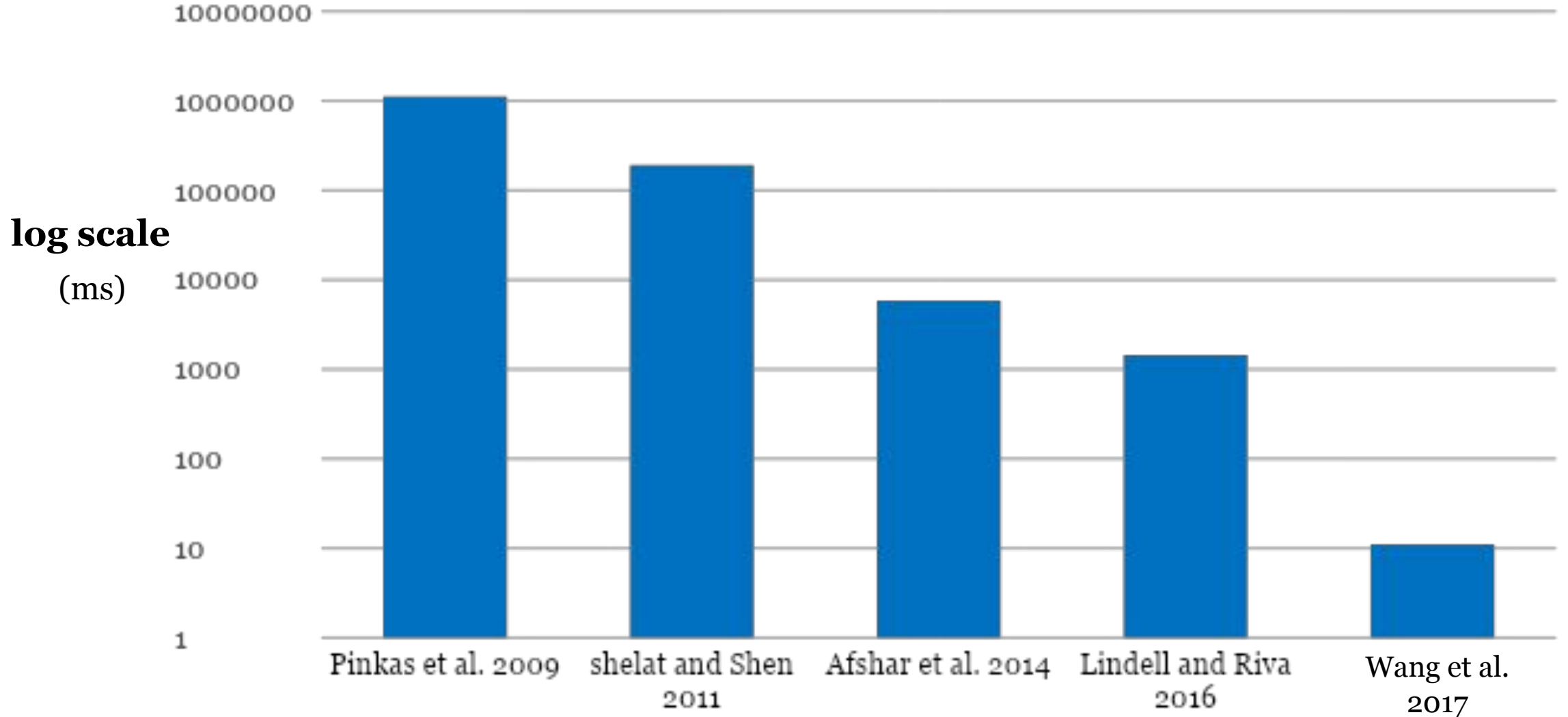
- Constant round

- Boolean circuits

Why not larger domain?

Optimizations

AES circuit



Point and permute

Garbled Table	
$\text{Enc}_{A_0, B_0}(C_0)$	✗
$\text{Enc}_{A_0, B_1}(C_0)$	✗
$\text{Enc}_{A_1, B_0}(C_0)$	✓
$\text{Enc}_{A_1, B_1}(C_1)$	✗

$A_1, B_0 \rightarrow C_0$

Solution: append the “select bit” to the keys

Row reduction

Garbled Table	
$\text{Hash}(A_0 B_0) \oplus (C_0)$	✗
$\text{Hash}(A_1 B_0) \oplus (C_0)$	✗
$\text{Hash}(A_0 B_1) \oplus (C_0)$	✓
$\text{Hash}(A_1 B_1) \oplus (C_1)$	✗

$A_1, B_0 \rightarrow C_0$

$$C_0 = \text{Hash}(A_0 || B_0)$$

Free XOR

XOR + AND is universal: construct NAND!!!

Garbled Table	A_0, A_1	$A, A \oplus R$
$\text{Enc}_{A_0, B_0}(C_0)$	B_0, B_1	$B, B \oplus R$
$\text{Enc}_{A_0, B_1}(C_0)$		
$\text{Enc}_{A_1, B_0}(C_0)$		
$\text{Enc}_{A_1, B_1}(C_1)$		

Half gates

- 1 input known to 1 party
 - Sender (garbler) knows 1 input
 - Receiver (evaluator) knows 1 input
- $c = a \wedge b = a \wedge (r \oplus r \oplus b) = a \wedge r \oplus a \wedge (r \oplus b)$