



Great Ideas in Computer Architecture

Combinational Logic

Instructor: Jenny Song



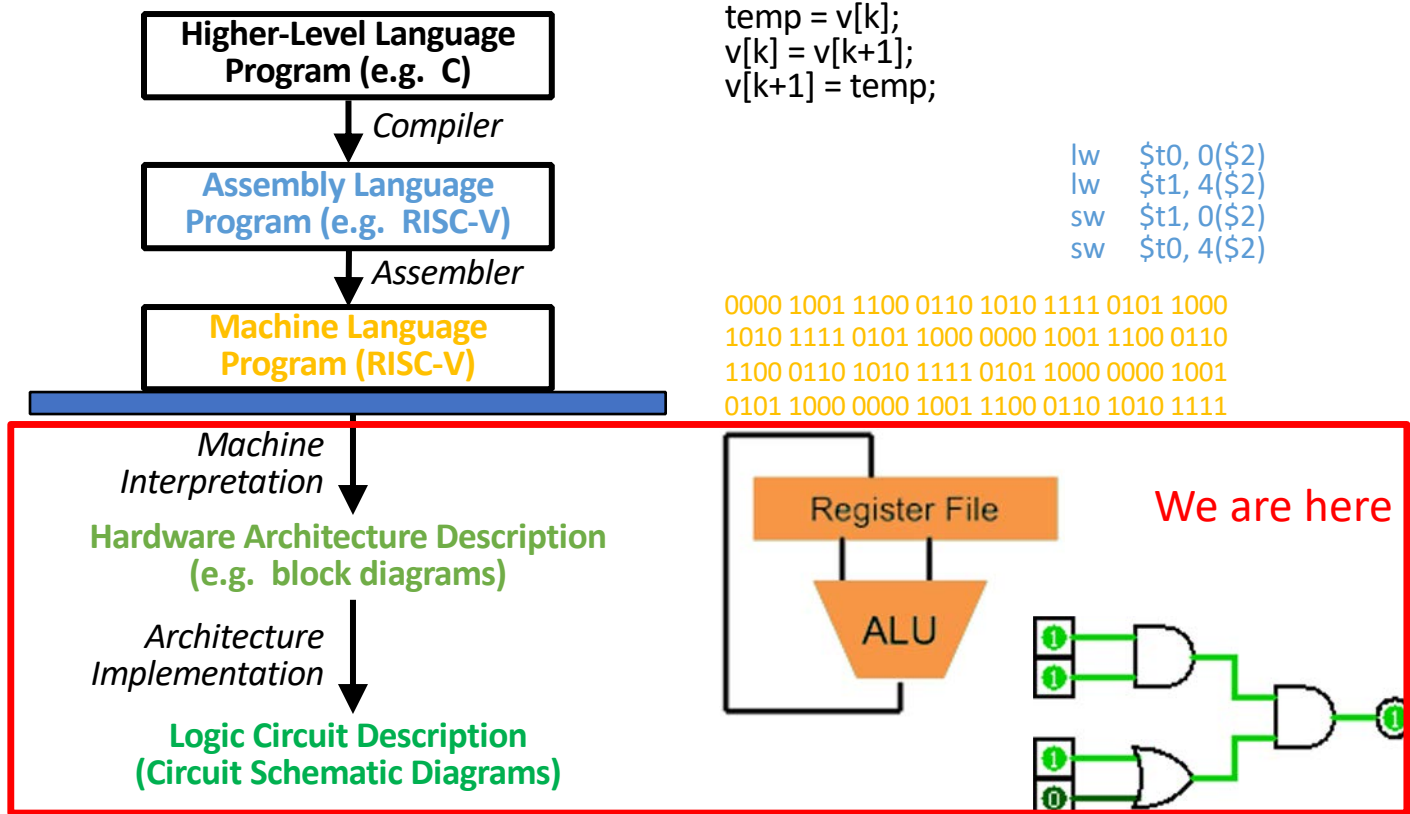
Review

- **Compiler** converts a single HLL file into a single assembly file $.c \rightarrow .s$
- **Assembler** removes pseudo-instructions, converts what it can to machine language, and creates a checklist for linker (relocation table) $.s \rightarrow .o$
 - Resolves addresses by making 2 passes (for internal forward references)
- **Linker** combines several object files and resolves absolute addresses $.o \rightarrow .out$
 - Enable separate compilation and use of libraries
- **Loader** loads the executable into memory and begins execution $.out \rightarrow [running]$

Agenda

- CALL Review
- **Hardware Design Overview**
- Switches and Transistors
- CMOS Networks
- Combinational Logic
 - Combinational Logic Gates
 - Truth Tables
 - Boolean Algebra
 - Circuit Simplification

Overview



Why study hardware design?

- Be able to answer to the question “How does a computer work?”
- We need some digital systems knowledge to build our own processor
- Understand how code is actually executed on a computer
- Understand capabilities and limitations of HW in general and processors in particular
- What processors can do fast and what they can't do fast (avoid slow things if you want your code to run fast!)
- Background for other more in-depth classes EECS151, CS152
- There is just so much you can do with standard processors: you may need to design own custom hardware for extra performance.

Synchronous Digital Systems (SDS)

Hardware of a processor, such as a RISC-V processor, is an example of a Synchronous Digital System

Synchronous:

- All operations coordinated by a central clock
 - “Heartbeat” of the system! (processor frequency)

Digital:

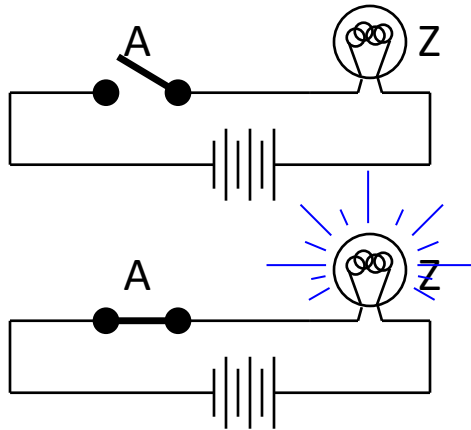
- Represent all values with two discrete values
- Electrical signals are treated as 1's and 0's
 - 1 and 0 are complements of each other
- High/Low voltage for True/False, 1/0

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Switches

- The basic element of physical implementations
- Convention: if input is a “1,” the switch is *asserted*



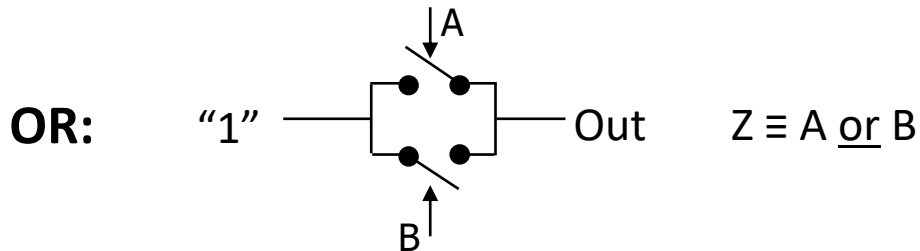
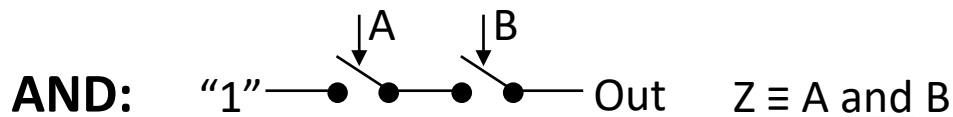
Open switch if A is “0” (unasserted)
and turn OFF light bulb (Z)

Close switch if A is “1” (asserted)
and turn ON light bulb (Z)

In this example, $Z \equiv A$.

Switch Logic

- Can compose switches into more complex ones (Boolean functions)
 - Arrows show action upon assertion (1 = close)



Historical Note

- Early computer designers built ad hoc circuits from switches
- Began to notice common patterns in their work: ANDs, ORs,
- Master's thesis (by Claude Shannon, 1940) made link between work and 19th Century Mathematician George Boole
- Called it “Boolean” in his honor
- Could apply math to give theory to hardware design, minimization...

Computers need switches

To create an electric computer:

- We need the ability to control switches with electricity
- Switches controlling groups of other switches
 - So electricity should flow through the switch

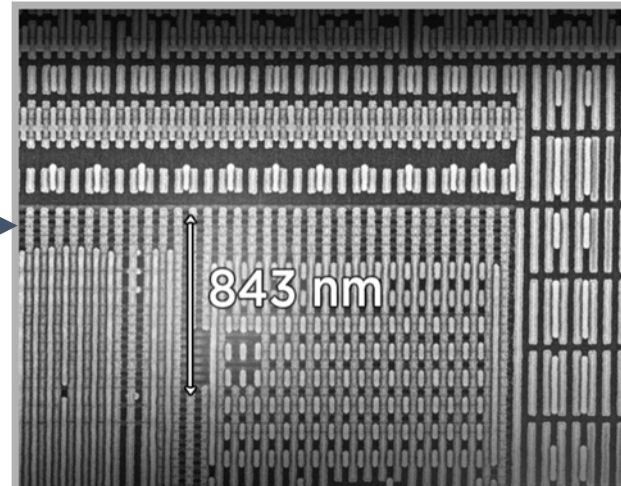
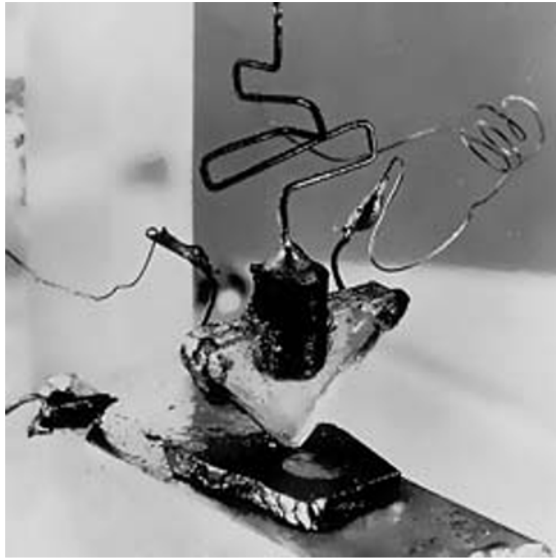
Transistors are hardware switches!

- Modern digital systems designed in CMOS
 - MOS: Metal-Oxide on Semiconductor
 - C for complementary: normally-open and normally-closed switches
- MOS transistors act as voltage-controlled switches

Transistors and CS61C

- The internals of transistors are important, but won't be covered in this class
 - Physical limitations relating to speed and power consumption
 - Can take EE16A/B, EE105, and EE140
- We will proceed with the abstraction of *Digital Logic* (0/1)
- It's also important to understand hardware **trends**

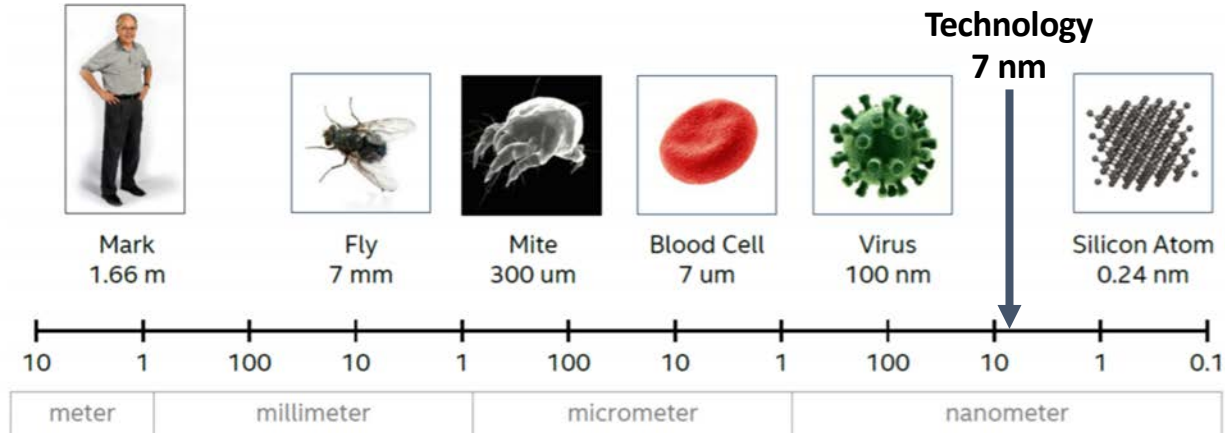
1947 → Today





2019 Processor Technology

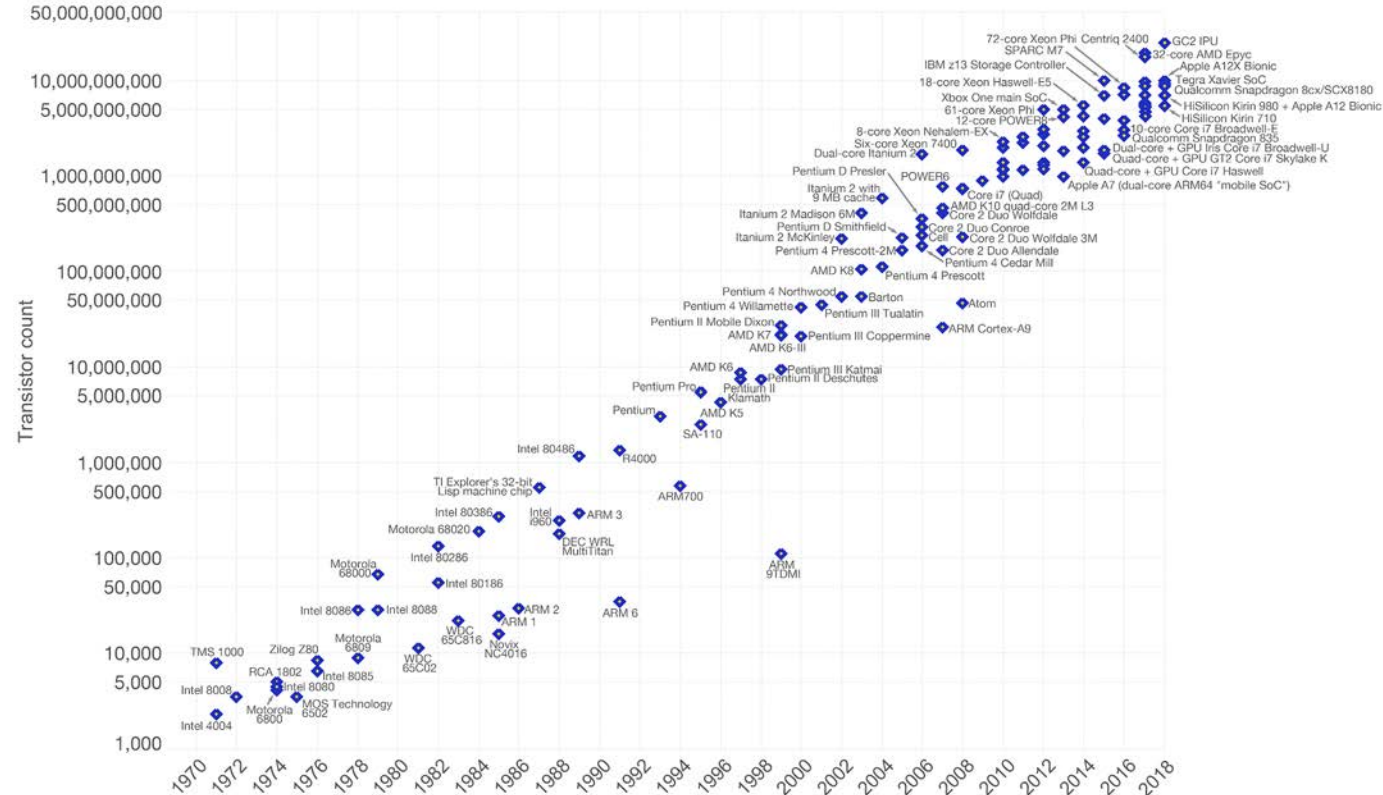
7 nm



Source: Mark Bohr, IDF14

Moore's Law – The number of transistors on integrated circuit chips (1971-2018)

Moore's law describes the empirical regularity that the number of transistors on integrated circuits doubles approximately every two years. This advancement is important as other aspects of technological progress – such as processing speed or the price of electronic products – are linked to Moore's law.



Data source: Wikipedia (https://en.wikipedia.org/wiki/Transistor_count)
The data visualization is available at OurWorldinData.org. There you find more visualizations and research on this topic.

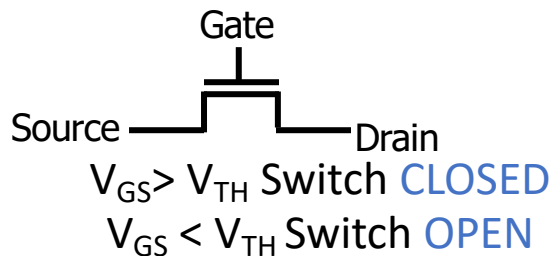
Licensed under CC-BY-SA by the author Max Roser.

Agenda

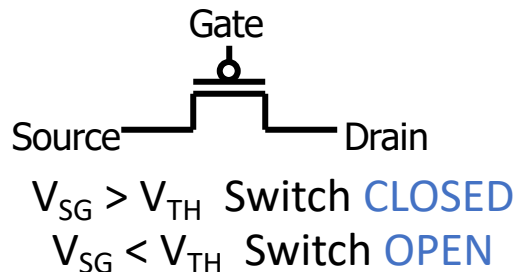
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Basics

N-channel



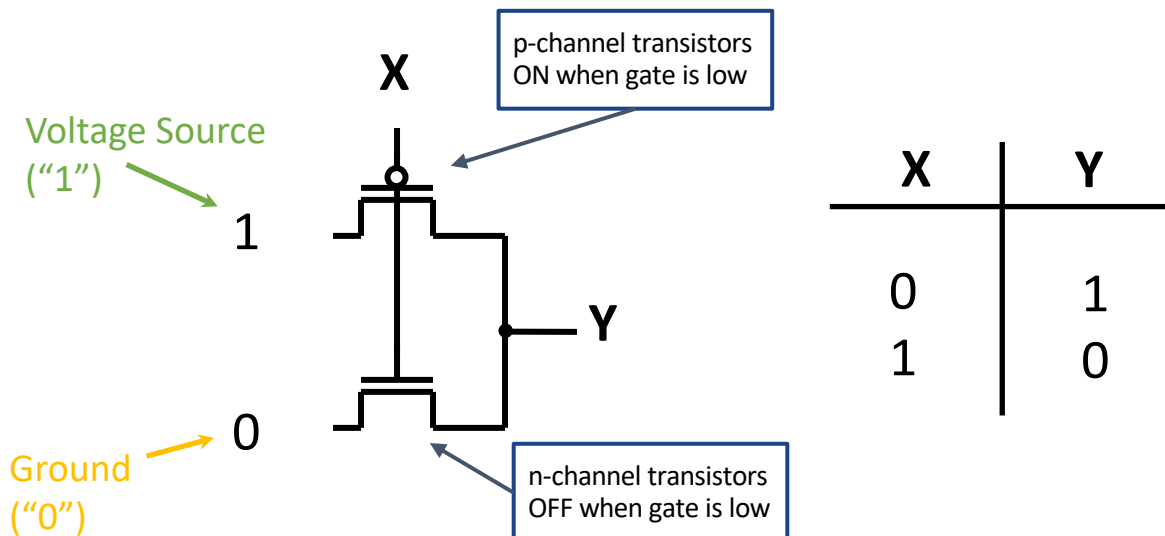
P-channel



- Three terminals: **Source**, **Gate**, and **Drain**
- Switch action: if voltage on gate terminal is (some amount) higher/lower than source terminal then conducting path established between drain and source terminals
- $V_{GS} = V_{\text{Gate}} - V_{\text{Source}}$; $V_{SG} = V_{\text{Source}} - V_{\text{Gate}}$; $V_{TH} = V_{\text{Threshold}}$

MOS Networks

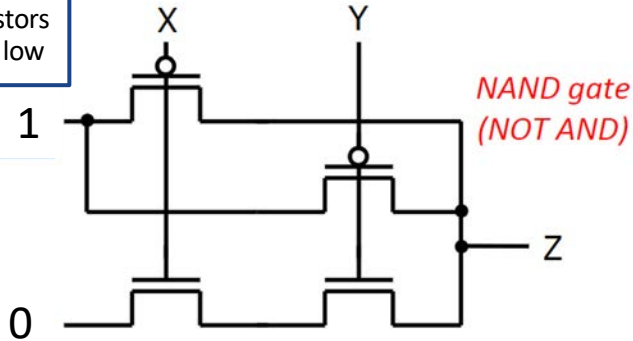
What is the relationship between X and Y?



Called an **inverter** or **NOT gate**

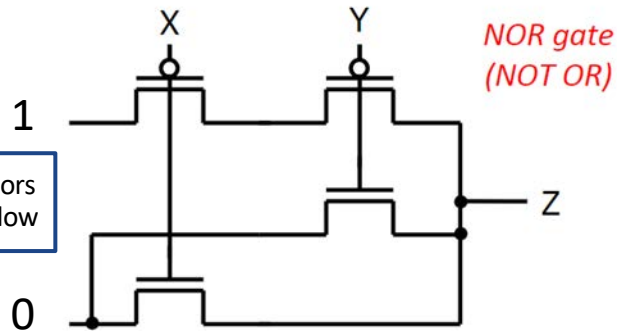
Two Input Networks

p-channel transistors
ON when gate is low



X	Y	Z
0	0	1
0	1	1
1	0	1
1	1	0

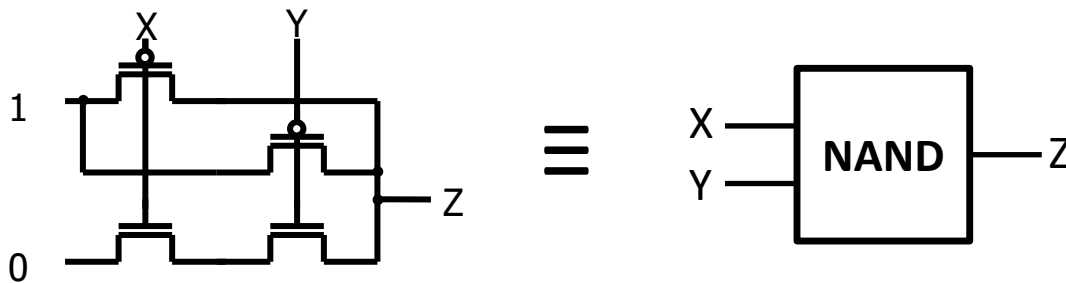
n-channel transistors
OFF when gate is low



X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	0

Abstraction: Block Diagrams

- In reality, chips composed of just transistors and wires
 - Small groups of transistors form useful building blocks, which we show as *blocks*



- Can combine to build higher-level blocks
 - You can build AND, OR, and NOT out of NAND!

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Type of Circuits

- *Digital Systems* consist of two basic types of circuits:

- Combinational Logic (CL)

- Output is a function of the inputs only, not the history of its execution
 - e.g. circuits to add A, B (ALUs)

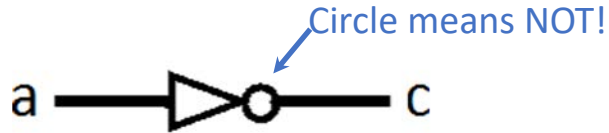
- Sequential Logic (SL)

- Circuits that “remember” or store information
 - a.k.a. “State Elements”
 - e.g. memory and registers (Registers)

Logic Gates (1/2)

- Special names and symbols:

NOT



a	c
0	1
1	0

AND



a	b	c
0	0	0
0	1	0
1	0	0
1	1	1

OR



a	b	c
0	0	0
0	1	1
1	0	1
1	1	1

Logic Gates (2/2)

Inverted versions are easier to implement in CMOS

NAND



NOR



XOR

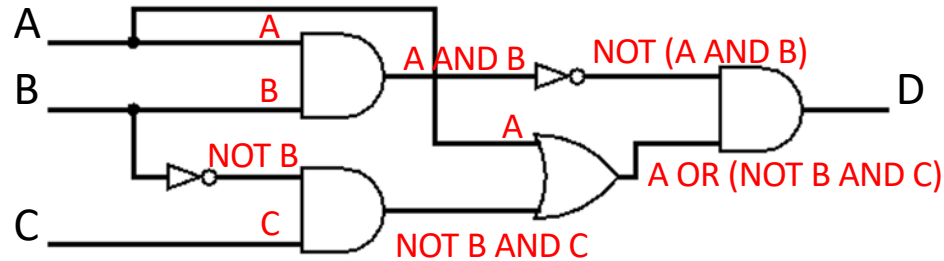


a	b	c
0	0	1
0	1	1
1	0	1
1	1	0

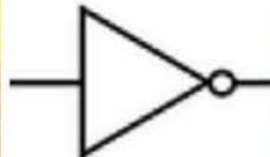
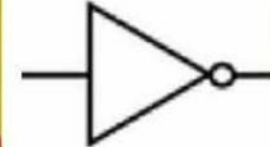
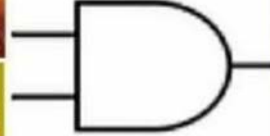
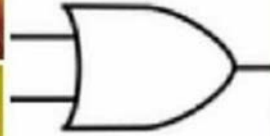
a	b	c
0	0	1
0	1	0
1	0	0
1	1	0

a	b	c
0	0	0
0	1	1
1	0	1
1	1	0

Combining Multiple Logic Gates



(NOT(A AND B)) AND (A OR (NOT B AND C))



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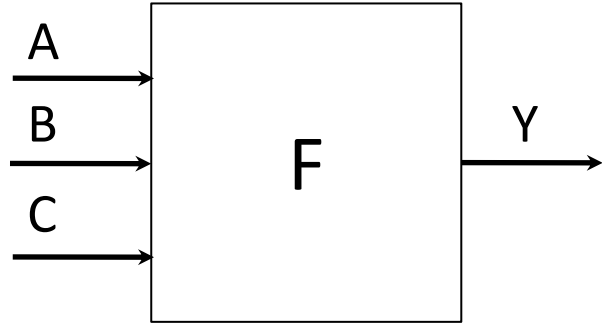
Representations of Combinational Logic

- ✓ Text Description
- ✓ Circuit Diagram
 - Transistors and wires
 - Logic Gates
- ✓ Truth Table
- ✓ Boolean Expression
- ✓ *All are equivalent*

Truth Tables

- Table that relates the inputs to a combinational logic circuit to its output
 - Output *only* depends on current inputs
 - Use abstraction of 0/1 instead of high/low V
 - Shows output for *every* possible combination of inputs
- How big?
 - 0 or 1 for each of N inputs, so 2^N rows

General Form



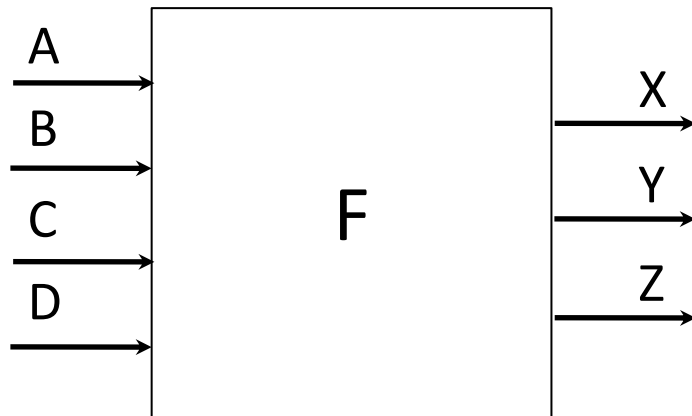
A	B	C	Y
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

8
Rows

If N inputs, how many distinct functions F do we have?

Function maps each row to 0 or 1,
so 2^R possible functions

Multiple Outputs



- For 3 outputs, just three indep. functions:
 $X(A,B,C,D)$, $Y(A,B,C,D)$, and $Z(A,B,C,D)$
 - Can show functions in separate columns without adding any rows!

Question: Convert the following statements into a Truth Table for: $(x \text{ XOR } y) \text{ OR } (\text{NOT } z)$

X	Y	Z	(A)	(B)	(C)	(D)
0	0	0	1	1	1	1
0	0	1	0	0	0	0
0	1	0	1	1	1	1
0	1	1	1	1	1	1
1	0	0	0	1	1	1
1	0	1	1	1	0	1
1	1	0	1	1	1	0
1	1	1	1	0	1	1

Question: Convert the following statements into a Truth Table for: $(x \text{ XOR } y) \text{ OR } (\text{NOT } z)$

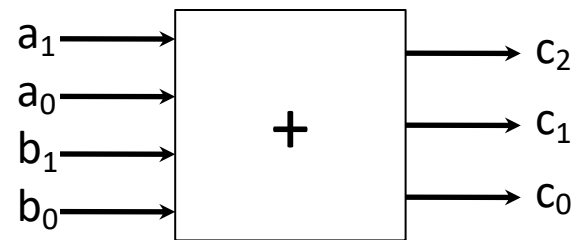
X	Y	Z	(A)	(B)	(C)	(D)
0	0	0	1	1	1	1
0	0	1	0	0	0	0
0	1	0	1	1	1	1
0	1	1	1	1	1	1
1	0	0	0	1	1	1
1	0	1	1	1	0	1
1	1	0	1	1	1	0
1	1	1	1	0	1	1

More Complicated Truth Tables

3-Input Majority

a	b	c	y
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

2-bit Adder



A		B		C		
a_1	a_0	b_1	b_0	c_2	c_1	c_0
			.			
			.			
			.			

How many rows?

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My Hand Hurts...

- Truth tables are huge
 - Write out EVERY combination of inputs and outputs (thorough, but inefficient)
 - Finding a particular combination of inputs involves scanning a large portion of the table
- There must be a shorter way to represent combinational logic
 - Boolean Algebra to the rescue!

a	b	c
00..0	00..0	00..0
00..0	00..1	00..1
.	.	.
.	.	.
.	.	.
.	.	.
.	.	.
11..1	11..1	11..1

Boolean Algebra

- Represent inputs and outputs as variables
 - Each variable can only take on the value 0 or 1
 - Overbar or \neg is NOT: “logical complement”
 - e.g. if A is 0, \bar{A} is 1. If A is 1, then $\neg A$ is 0
 - Plus (+) is 2-input OR: “logical sum”
 - Product (\cdot) is 2-input AND: “logical product”
 - All other gates and logical expressions can be built from combinations of these
- $\neg \mathbf{A} \mathbf{B} + \mathbf{A} \neg \mathbf{B} == (\text{NOT}(\mathbf{A} \text{ AND } \mathbf{B})) \text{ OR } (\mathbf{A} \text{ AND NOT } \mathbf{B})$

← For slides,
will use $\neg A$

Truth Table to Boolean Expression

- Read off of table
 - For 1, write variable name
 - For 0, write complement of variable

a	b	c
0	0	0
0	1	1
1	0	1
1	1	0

- *Sum of Products (SoP)*

- Take rows with 1's in output column, sum products of inputs
- $c = \neg a b + a \neg b$

- *Product of Sums (PoS)*

- Take rows with 0's in output column, product the sum of the *complements of the inputs*
- $c = (a + b) \cdot (\neg a + \neg b)$

We can show that these are equivalent!

$$= \neg(\neg a \neg b + a b)$$

Laws of Boolean Algebra

These laws allow us to perform simplification:

$x \cdot \bar{x} = 0$	$x + \bar{x} = 1$	complementarity
$x \cdot 0 = 0$	$x + 1 = 1$	laws of 0's and 1's
$x \cdot 1 = x$	$x + 0 = x$	identities
$x \cdot x = x$	$x + x = x$	idempotent law
$x \cdot y = y \cdot x$	$x + y = y + x$	commutativity
$(xy)z = x(yz)$	$(x + y) + z = x + (y + z)$	associativity
$x(y + z) = xy + xz$	$x + yz = (x + y)(x + z)$	distribution
$xy + x = x$	$(x + y)x = x$	uniting theorem
$\bar{x}y + x = x + y$	$(\bar{x} + y)x = xy$	uniting theorem v.2
$\overline{x \cdot y} = \bar{x} + \bar{y}$	$\overline{x + y} = \bar{x} \cdot \bar{y}$	DeMorgan's Law

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Manipulating Boolean Algebra

- SoP and PoS expressions can still be long
 - We wanted to have shorter representation than a truth table!
- Boolean algebra follows a set of rules that allow for simplification
 - Goal will be to arrive at the simplest equivalent expression
 - Allows us to build simpler (and faster) hardware

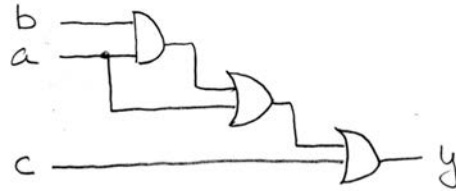
Faster Hardware?

- **Recall:** Everything we are dealing with is just an abstraction of transistors and wires
 - Inputs propagating to the outputs are voltage signals passing through transistor networks
 - There is always some *delay* before a CL's output updates to reflect the inputs
- Simpler Boolean expressions \leftrightarrow smaller transistor networks \leftrightarrow smaller circuit delays \leftrightarrow faster hardware

Boolean Algebraic Simplification Example

$$y = ab + a + c$$

Circuit Simplification



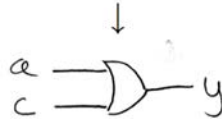
1) original circuit (Transistors and/or Gates)

$$\downarrow$$
$$y = ((ab) + a) + c$$

2) equation derived from original circuit

$$\downarrow$$
$$\begin{aligned} &= ab + a + c \\ &= a(b + 1) + c \\ &= a(1) + c \\ &= a + c \end{aligned}$$

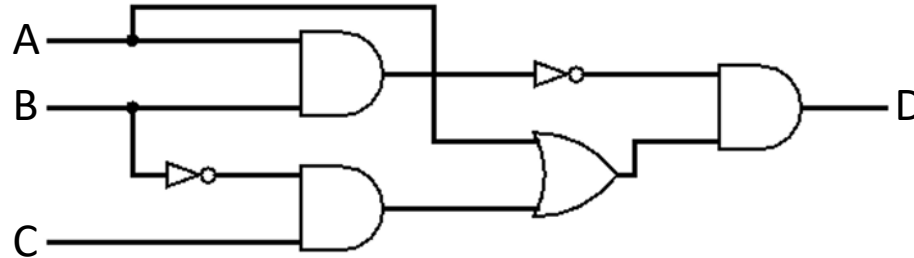
3) algebraic simplification



4) simplified circuit

Circuit Simplification Example (1/4)

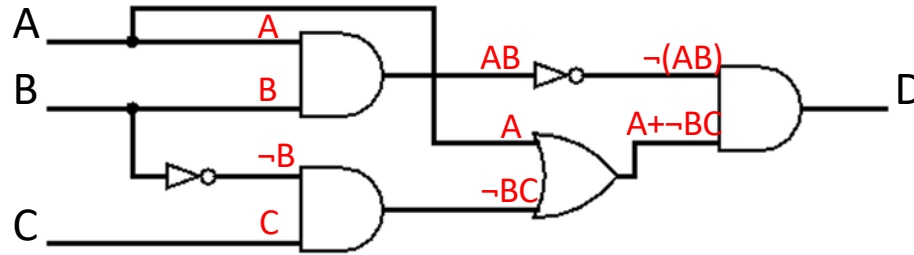
- Simplify the following circuit:



- Options:
 - 1) Test all combinations of the inputs and build the Truth Table, then use SoP or PoS
 - 2) Write out expressions for signals based on gates
 - Will show this method here

Circuit Simplification Example (2/4)

- Simplify the following circuit:



- Start from left, propagate signals to the right

Arrive at $D = \neg(AB)(A + \neg BC)$

Circuit Simplification Example (3/4)

- Simplify Expression:

$$D = \neg(AB)(A + \neg BC)$$

$$= (\neg A + \neg B)(A + \neg BC)$$

$$= \neg AA + \neg A\neg BC + \neg BA + \neg B\neg BC$$

$$= 0 + \neg A\neg BC + \neg BA + \neg B\neg BC$$

$$= \neg A\neg BC + \neg BA + \neg BC$$

$$= (\neg A + 1)\neg BC + \neg BA$$

$$\begin{aligned} &= \neg BC + \neg BA \\ &= \neg B(A + C) \end{aligned} \quad \left. \vphantom{\begin{aligned} &= \neg BC + \neg BA \\ &= \neg B(A + C) \end{aligned}} \right\} \text{Which of these is "simpler"?$$

DeMorgan's

Distribution

Complementarity

Idempotent Law

Distribution

Law of 1's

Distribution

Circuit Simplification Example (4/4)

- Draw out final circuit:

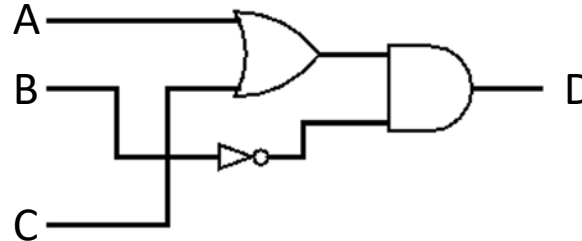
- $D = \neg BC + \neg BA = \neg B(A + C)$

5

3

How many gates
do we need for each?

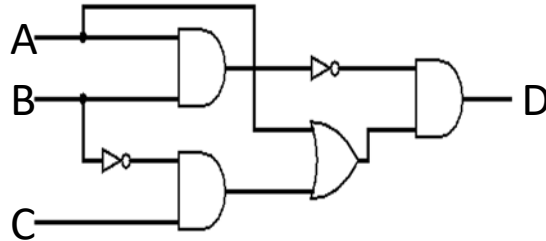
- Simplified Circuit:



- Reduction from 6 gates to 3!
 - Beore: $\neg(AB)(A + \neg BC)$

Summary

- Beginnings of hardware design and layered abstractions
- Transistors -> CMOS Networks -> Combinational Logic



- ✓ Text Description
- ✓ Circuit Diagram
 - Transistors and wires
 - Logic Gates
- ✓ Truth Table
- ✓ Boolean Expression

✓ *All are equivalent*

Converting Combinational Logic

