

Homework 4

1. $S_0 \rightarrow 1S_01 \mid S\# \mid \#S$ CFG $G_1 = (\{S_0, S\}, \{1, \#\}, R, S_0)$
 $S \rightarrow 1S \mid 1$ w/ $R = \{S_0 \rightarrow 1S_01, S_0 \rightarrow S\#, S_0 \rightarrow \#S, S \rightarrow 1S, S \rightarrow 1\}$

2. $S \rightarrow 0S0 \mid 1S0 \mid 0X1 \mid 1X1$
 $X \rightarrow 0X0 \mid 1X0 \mid 0X1 \mid 1X1 \mid 0 \mid 1 \mid \epsilon$

CFG $G_2 = (\{S, X\}, \{0, 1\}, R, S)$

w/ $R = \{S \rightarrow 0S0, S \rightarrow 1S0, S \rightarrow 0X1, S \rightarrow 1X1, X \rightarrow 0X0, X \rightarrow 1X0, X \rightarrow 0X1, X \rightarrow 1X1, X \rightarrow 0, X \rightarrow 1, X \rightarrow \epsilon\}$

3. 2. $R = (R_1 \circ R_2)$ This is equivalent to the CFG that consists of all rules and variables from each of G_1 and G_2 with a new start variable S and one additional rule $S \rightarrow S_1S_2$.

Equivalence is because S derives the set of strings that have a prefix derived by S_1 and a suffix derived by S_2 . This is the language of $(R_1 \circ R_2)$.

3. $R = (R_i^*)$ This is equivalent to the CFG that consists of all rules and variables from G_i with a new start variable S and one additional rule $S \rightarrow SS_1 \mid \epsilon$.

Equivalence is because S derives any string that consists of a non-negative number of substrings which are derived by S_1 . This is the language of R_i^* .

4a. The grammar generates all strings that consist of positive number of substrings that are $a^n b^n$ with $n > 0$.

- b. The grammar is ambiguous because it can derive $ababab$ in two distinct leftmost derivations

1st $S \rightarrow SS \rightarrow SSS \rightarrow TSS \rightarrow abSS \rightarrow abTS \rightarrow ababS \rightarrow ababT \rightarrow ababab$

2nd $S \rightarrow SS \rightarrow TS \rightarrow abS \rightarrow abSS \rightarrow abTS \rightarrow ababS \rightarrow ababT \rightarrow ababab$

c. $S \rightarrow SS | T$
 $T \rightarrow aTb | ab$

step 1

$$S_0 \rightarrow S$$

$$S \rightarrow SS | T$$

$$T \rightarrow aTb | ab$$

step 2

$$S_0 \rightarrow SS$$

$$S \rightarrow SS | T$$

$$T \rightarrow aTb | ab$$

step 3

$$S_0 \rightarrow SS | aTb | ab$$

$$S \rightarrow SS | aTb | ab$$

$$T \rightarrow aTb | ab$$

step 4

$$S_0 \rightarrow SS | AZ | AB$$

$$S \rightarrow SS | AZ | AB$$

$$T \rightarrow AZ | AB$$

$$Z \rightarrow TB$$

$$A \rightarrow a$$

$$B \rightarrow b$$

CNF

$$S_0 \rightarrow SS | AZ | AB$$

$$S \rightarrow SS | AZ | AB$$

$$T \rightarrow AZ | AB$$

$$Z \rightarrow TB$$

$$A \rightarrow a$$

$$B \rightarrow b$$

Citations: Used class Notes & Textbook