

Homework 9

1. Claim: $\text{REGULAR}_{\text{TM}}$ is not recognizable and not corecognizable.

Proof:

From Theorem 5.3, $\neg \text{ATM} \leq_m \text{REGULAR}_{\text{TM}}$ with reduction $f(\langle M, w \rangle) = \langle M, w \rangle$. Thus $\overline{\text{ATM}} \leq_m \overline{\text{REGULAR}_{\text{TM}}}$ and $\overline{\text{ATM}}$ is not recognizable since ATM is recognizable and ATM is not decidable. So $\text{REGULAR}_{\text{TM}}$ is not corecognizable since $\overline{\text{REGULAR}_{\text{TM}}}$ is not recognizable by inheritance of $\overline{\text{ATM}} \leq_m \overline{\text{REGULAR}_{\text{TM}}}$.

Also from $\text{ATM} \leq_m \text{REGULAR}_{\text{TM}} \Leftrightarrow \overline{\text{ATM}} \leq_m \overline{\text{REGULAR}_{\text{TM}}}$, $\overline{\text{ATM}} \leq_m \text{REGULAR}_{\text{TM}}$ and $\text{REGULAR}_{\text{TM}}$ inherits $\overline{\text{ATM}}$'s unrecognizability and thus $\text{REGULAR}_{\text{TM}}$ is also not recognizable.

2. \Rightarrow Assume A is recognizable by M . Then $A \leq_m \text{ATM}$ with reduction $f(w) = \langle M, w \rangle$ because the assumption that M recognizes A means $w \in A$ iff $w \in L(M)$.

\Leftarrow Assume $A \leq_m \text{ATM}$. Then we can run the universal Turing machine on reduction f , $f(x) = \langle M, w \rangle$ to recognize A . $x \in A$ iff an ATM recognizer accepts $f(x) = \langle M, w \rangle$. Thus A is recognizable.

3. Not recognizable. $f(\langle M, w \rangle) = \langle M, w \rangle$ is a mapping reduction for $\overline{\text{ATM}} \leq_m J$ because $f(x) \in J \Leftrightarrow x \in \overline{\text{ATM}}$, thus J inherits $\overline{\text{ATM}}$'s unrecognizability.

Not corecognizable. $f(\langle M, w \rangle) = \langle M, w \rangle$ is a mapping reduction for $\text{ATM} \leq_m \bar{J}$ because $f(x) \in \bar{J} \Leftrightarrow x \in \text{ATM}$ and \bar{J} then inherits ATM 's unrecognizability.

4. Claim: EQ_{LBA} is corecognizable and not recognizable, therefore not decidable

Proof:

Corecognizable

Build recognizer for EQ_{LBA} as follows

On input $\langle M_1, M_2 \rangle$:

If $\langle M_1, M_2 \rangle$ are not a pair of LBAs: accept

For $l = 0, 1, 2, \dots$

For each string w of length $|w| = l$

Check if M_1, M_2 accept w after $\leq |Q_1| l / |P|^{k_1} \leq |Q_2| l / |P|^{k_2}$

If one accepts and not the other: accept

Not Decidable

Suppose EQ_{LBA} is decidable. Since every CFG

can be represented by LBA, create two LBAs M_1 and M_2 .

Equivalent to two CFGs. Check if $\langle M_1, M_2 \rangle$

are in EQ_{LBA} . If yes: accept. Else reject

Contradiction: EQ_{CFG} is not decidable, thus EQ_{LBA} is not decidable

Not recognizable: Since EQ_{LBA} is not decidable and is corecognizable, EQ_{LBA} must not be recognizable.

5a. We cannot prove undecidability of EQ_{CFG} by reduction from EQ_{LBA} because EQ_{CFG} is decidable

b. The mapping reduction that shows $ALL_{CFG} \leq_m EQ_{CFG}$ is

Let M_1 be a TM that decides EQ_{CFG} and we construct M_2 to decide ALL_{CFG} as follows

M_2 = On input $\langle G \rangle$ where G is a CFG

1. Run M_1 on input $\langle G, G \rangle$, G_0 is CFG with $L(G_0) = \Sigma^*$
2. If M_1 accepts: accept. If M_1 rejects: reject.

c. If is a mapping reduction $\overline{A_{TM}} \leq A_{LLCF}$

Citation: Used Class Notes, textbook, online lecture slides from Stanford and University of Iowa