Quiz: question 1

What is the storage required to represent analysis solution to a matrix completion problem

minimize
$$\sum_{(i,j)\in E} (x_{ij} - b_{ij})^2$$
 s.t. $||X|| \le \alpha$

with variable $X \in \mathbf{R}^{m \times n}$ if all solutions have rank $\leq r$?

- $ightharpoonup O(mnr \log n)$
- ► *O*(*mnr*)
- $O(r(m+n)\log(m+n))$
- ightharpoonup O(r(m+n))
- $ightharpoonup O(r\log(m+n))$

Quiz: question 2

What is the main advantage of the Conditional Gradient method for designing an optimal storage algorithm?

- ▶ the rank of any iterate ≤ the rank of the solution
- the rank of the iterate grows by at most 1 at each iteration
- ► the iteration can be expressed in terms of a low dimensional "dual" variable
- the method converges linearly for underconstrained problems like matrix completion and phase retrieval

Sketchy Decisions: Convex Low-Rank Matrix Optimization with Optimal Storage

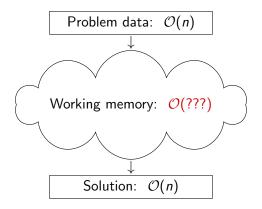
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Based on joint work with Alp Yurtsever (MIT), Volkan Cevher (EPFL), and Joel Tropp (Caltech)

Goal

Can we develop algorithms that provably solve a problem using **storage** bounded by the size of the **problem data** and the size of the **solution**?



Model problem: low rank matrix optimization

consider a convex problem with decision variable $X \in \mathbb{R}^{m \times n}$ compact matrix optimization problem:

$$\begin{array}{ll} \text{minimize} & f(\mathcal{A}X) \\ \text{subject to} & \|X\|_{S_1} \leq \alpha \end{array} \tag{CMOP}$$

- $A: \mathbb{R}^{m \times n} \to \mathbb{R}^d$
- $f: \mathbb{R}^d \to \mathbb{R}$ convex and smooth
- ▶ $||X||_{S_1}$ is Schatten-1 norm: sum of singular values

assume

- **compact specification**: problem data use $\mathcal{O}(n)$ storage
- **compact solution**: rank $X_{\star} = r$ constant

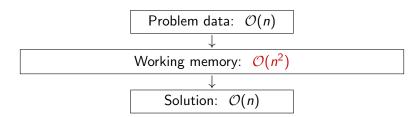
Note: Same ideas work for $X \succeq 0$

Are desiderata achievable?

minimize
$$f(AX)$$

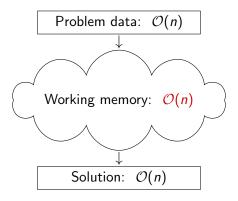
subject to $||X||_{S_1} \le \alpha$

CMOP, using any first order method:



Are desiderata achievable?

CMOP, using **SketchyCGM**:



Application: matrix completion

find X matching M on observed entries

minimize
$$\sum_{(i,j)\in\Omega} (X_{ij} - M_{ij})^2$$
 subject to $\|X\|_{S_1} \le \alpha$

- ightharpoonup m = rows, n = columns of matrix to complete
- $d = |\Omega|$ number of observations
- ▶ A selects observed entries X_{ij} , $(i,j) \in \Omega$
- $f(z) = \|z \mathcal{A}M\|^2$

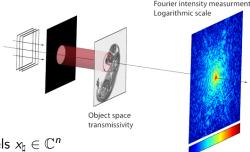
Matrix completion is a CMOP

find X matching M on observed entries

minimize
$$\sum_{(i,j)\in\Omega} (X_{ij} - M_{ij})^2$$
 subject to $||X||_{S_1} \le \alpha$

- ▶ compact specification if $d = \mathcal{O}(m+n)$ observations e.g., constant # observations / person
- compact solution if rank(X) constant i.e., constant # parameters / person
- ▶ in practice, usually find rank \ll 200 even with m and n in the millions. . .

Application: Phase retrieval



- ▶ image with n pixels $x_{\natural} \in \mathbb{C}^n$
- acquire noisy nonlinear measurements $b_i = |\langle a_i, x_1 \rangle|^2 + \omega_i$
- relax: if $X = x_{\natural}x_{\natural}^*$, then

$$|\langle a_i, x_{\natural} \rangle|^2 = x_{\natural} a_i^* a_i x_{\natural}^* = \operatorname{tr}(a_i^* a_i x_{\natural}^* x_{\natural}) = \operatorname{tr}(a_i^* a_i X)$$

recover image by solving

minimize
$$f(AX; b)$$
 subject to $\operatorname{tr} X = \alpha$
 $X \succ 0$.

¹image courtesy of Manuel Guizar-Sicairos

Phase retrieval is a CMOP

find X matching observations

minimize
$$f(AX; b)$$

subject to $\operatorname{tr} X = \alpha$
 $X \succeq 0$.

- ▶ compact specification if $d = \mathcal{O}(n)$ observations *e.g.*, constant # observations / pixels
- compact solution if rank(X) constant e.g., if correctly recover the rank-1 solution!

Why compact?

why a compact specification?

- data is expensive
- collect constant data per column (=user or sample)
- ▶ if solution is compact, compact specification should suffice

why a compact solution?

- the world is simple and structured
- ▶ given d observations, there is a solution with rank $\mathcal{O}(\sqrt{d})$ (Barvinok 1995, Pataki 1998)
- nice latent variable models are of log rank (Udell & Townsend 2019)

Optimal Storage

What kind of storage bounds can we hope for?

Assume black-box implementation of

$$A(uv^*)$$
 $u^*(A^*z)$ $(A^*z)v$

where $u \in \mathbb{R}^m$, $v \in \mathbb{R}^n$, and $z \in \mathbb{R}^d$

- ▶ Need $\Omega(m+n+d)$ storage to apply linear map
- ▶ Need $\Theta(r(m+n))$ storage for a rank-r approximate solution

Definition. An algorithm for the model problem has **optimal storage** if its working storage is

$$\Theta(d+r(m+n)).$$

Optimal Storage

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$$\Theta(d+r(m+n)).$$

If we write down X, we've already failed.

A brief biased history of matrix optimization (I)

- ▶ 1990s: Interior-point methods
 - ▶ Storage cost $\Theta((m+n)^4)$ for Hessian
- 2000s: Convex first-order methods (FOM)
 - ► (Accelerated) proximal gradient and others
 - Store matrix variable $\Theta(mn)$

```
(Interior-point: Nemirovski & Nesterov 1994; ...; First-order: Rockafellar 1976; Auslender & Teboulle 2006; ...)
```

A brief biased history of matrix optimization (I)

- ► 2008–Present: **Storage-efficient convex FOM**
 - Conditional gradient method (CGM) and extensions
 - ▶ Store matrix in low-rank form $\mathcal{O}(t(m+n))$ after t iterations
 - ▶ Requires storage $\Theta(mn)$ for $t \ge \min(m, n)$
 - Variants: prune factorization, or seek rank-reducing steps
- ► 2003–Present: **Nonconvex methods**
 - ▶ Burer–Monteiro factorization idea + various opt algorithms
 - Store low-rank matrix factors $\Theta(r(m+n))$
 - For guaranteed solution, need statistical assumptions or $\mathcal{O}(n^{3/2})$ storage

(CGM: Frank & Wolfe 1956; Levitin & Poljak 1967; Hazan 2008; Clarkson 2010; Jaggi 2013; ...; CGM+pruning: Rao Shah Wright 2015; Freund Grigas Mazumder 2017; ...; Nonconvex: Burer & Monteiro 2003; Keshavan et al. 2009; Jain et al. 2012; Bhojanapalli et al. 2015; Candès et al. 2014; Boumal et al. 2015; Bhojanapalli et al. 2018; Waldspurger & Waters 2018; ...)

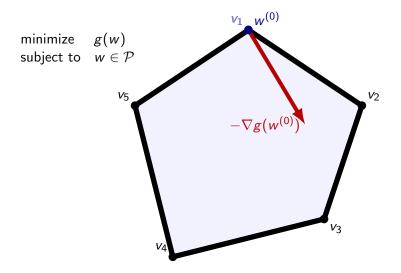
The dilemma

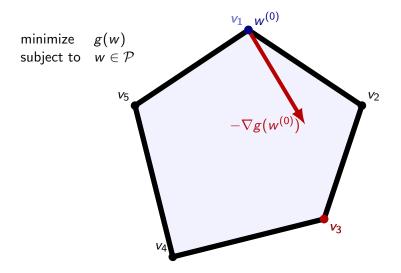
- convex methods: slow memory hogs with guarantees
- nonconvex methods: fast, lightweight, but brittle

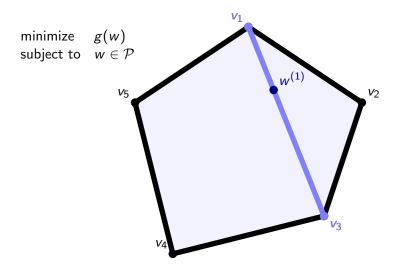
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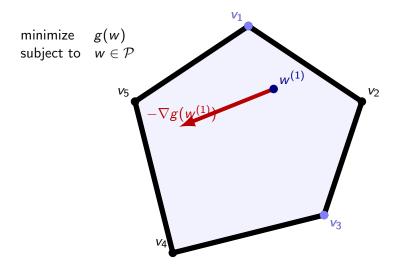
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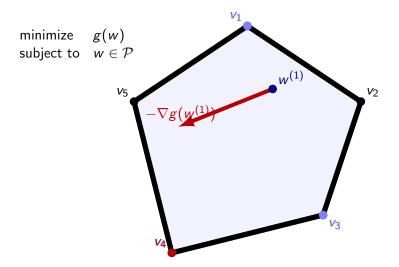
goal: low memory and guaranteed convergence

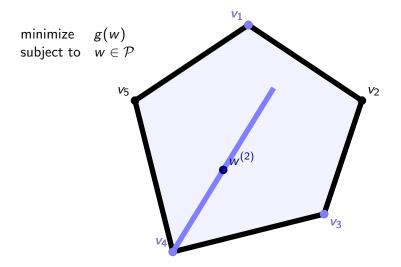












Conditional Gradient Method

minimize
$$f(AX)$$

subject to $||X||_{S_1} \le \alpha$

CGM. set
$$X^0 = 0$$
. for $t = 0, 1, ...$

- compute $G^t = \mathcal{A}^* \nabla f(\mathcal{A}X^t)$
- set search direction

$$H^t = \underset{\|X\|_{\mathcal{S}_1} \leq \alpha}{\operatorname{argmax}} \langle X, -G^t \rangle$$

- set stepsize $\eta^t = 2/(t+2)$
- update $X^{t+1} = (1 \eta^t)X^t + \eta^t H^t$

Conditional gradient method (CGM)

features:

relies on efficient linear optimization oracle to compute

$$H^t = \underset{\|X\|_{\mathcal{S}_1} \le \alpha}{\operatorname{argmax}} \langle X, -G^t \rangle$$

bound on suboptimality follows from subgradient inequality

$$f(\mathcal{A}X^{t}) - f(\mathcal{A}X^{\star}) \leq \langle X^{t} - X^{\star}, G^{t} \rangle$$

$$\leq \langle X^{t} - X^{\star}, \mathcal{A}^{*} \nabla f(\mathcal{A}X^{t}) \rangle$$

$$\leq \langle \mathcal{A}X^{t} - \mathcal{A}X^{\star}, \nabla f(\mathcal{A}X^{t}) \rangle$$

$$\leq \langle \mathcal{A}X^{t} - \mathcal{A}H^{t}, \nabla f(\mathcal{A}X^{t}) \rangle$$

to provide stopping condition

faster variants: linesearch, away steps, . . .

Linear optimization oracle for MOP

compute search direction

$$\mathop{\rm argmax}_{\|X\|_{S_1} \leq \alpha} \langle X, -G \rangle$$

▶ solution given by maximum singular vector of -G:

$$-G = \sum_{i=1}^{n} \sigma_{i} u_{i} v_{i}^{*} \implies X = \alpha u_{1} v_{1}^{*}$$

 \blacktriangleright use Lanczos method: only need to apply G and G^*

Algorithm 1 CGM for the model problem (CMOP)

```
Input: Problem data for (CMOP); suboptimality \varepsilon Output: Solution X_{\star}
```

```
function CGM
X \leftarrow 0
for t \leftarrow 0, 1, \dots do
(u, v) \leftarrow \texttt{MaxSingVec}(-\mathcal{A}^*(\nabla f(\mathcal{A}X)))
H \leftarrow -\alpha uv^*
if \langle \mathcal{A}X - \mathcal{A}H, \nabla f(\mathcal{A}X) \rangle \leq \varepsilon then break for
\eta \leftarrow 2/(t+2)
X \leftarrow (1-\eta)X + \eta H
return X
```

Two crucial ideas

To solve the problem using optimal storage:

▶ Use the low-dimensional "dual" variable

$$z_t = \mathcal{A}X_t \in \mathbb{R}^d$$

to drive the iteration.

Recover solution from small (randomized) sketch.

Never write down X until it has converged to low rank.

Algorithm 2 CGM for the model problem (CMOP)

```
Input: Problem data for (CMOP); suboptimality \varepsilon Output: Solution X_{\star}
```

```
function CGM
X \leftarrow 0
for t \leftarrow 0, 1, \dots do
(u, v) \leftarrow \texttt{MaxSingVec}(-\mathcal{A}^*(\nabla f(\mathcal{A}X)))
H \leftarrow -\alpha uv^*
if \langle \mathcal{A}X - \mathcal{A}H, \nabla f(\mathcal{A}X) \rangle \leq \varepsilon \text{ then break for}
\eta \leftarrow 2/(t+2)
X \leftarrow (1-\eta)X + \eta H
return X
```

Introduce "dual variable" $z = AX \in \mathbb{R}^d$; eliminate X.

Algorithm 3 Dual CGM for the model problem (CMOP)

Input: Problem data for (CMOP); suboptimality ε **Output:** Solution X_{\star}

```
function DUALCGM z \leftarrow 0
for t \leftarrow 0, 1, \dots do (u, v) \leftarrow \texttt{MaxSingVec}(-\mathcal{A}^*(\nabla f(z)))
h \leftarrow \mathcal{A}(-\alpha u v^*)
if \langle z - h, \nabla f(z) \rangle \leq \varepsilon then break for \eta \leftarrow 2/(t+2)
z \leftarrow (1-\eta)z + \eta h
```

Introduce "dual variable" $z = AX \in \mathbb{R}^d$; eliminate X.

Algorithm 4 Dual CGM for the model problem (CMOP)

Input: Problem data for (CMOP); suboptimality ε **Output:** Solution X_{\star}

```
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if \langle z - h, \nabla f(z) \rangle \leq \varepsilon then break for \eta \leftarrow 2/(t+2)
z \leftarrow (1-\eta)z + \eta h
```

we've solved the problem...but where's the solution?

Two crucial ideas

1. Use the low-dimensional "dual" variable

$$z_t = \mathcal{A}X_t \in \mathbb{R}^d$$

to drive the iteration.

2. Recover solution from small (randomized) sketch.

How to catch a low rank matrix

if \hat{X} has the same rank as X^* , and \hat{X} acts like X^* (on its range and co-range), then \hat{X} is X^*

use single-pass randomized sketch (Tropp Yurtsever U Cevher 2017)

- see a series of additive updates
- remember how the matrix acts on random subspace
- reconstruct a low rank matrix that acts like X*
- ▶ storage cost for sketch and arithmetic cost of update are $\mathcal{O}(r(m+n))$; reconstruction is $\mathcal{O}(r^2(m+n))$

Single-pass randomized sketch

Draw and fix two independent standard normal matrices

$$\Omega \in \mathbb{R}^{n \times k}$$
 and $\Psi \in \mathbb{R}^{\ell \times m}$

with
$$k = 2r + 1$$
, $\ell = 4r + 2$.

Single-pass randomized sketch

Draw and fix two independent standard normal matrices

$$\Omega \in \mathbb{R}^{n \times k}$$
 and $\Psi \in \mathbb{R}^{\ell \times m}$

with
$$k = 2r + 1$$
, $\ell = 4r + 2$.

► The sketch consists of two matrices that capture the range and co-range of *X*:

$$Y = X\Omega \in \mathbb{R}^{n \times k}$$
 and $W = \Psi X \in \mathbb{R}^{\ell \times m}$

Single-pass randomized sketch

Draw and fix two independent standard normal matrices

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► The sketch consists of two matrices that capture the range and co-range of *X*:

$$Y = X\Omega \in \mathbb{R}^{n \times k}$$
 and $W = \Psi X \in \mathbb{R}^{\ell \times m}$

▶ Rank-1 updates to *X* can be performed on sketch:

Single-pass randomized sketch

Draw and fix two independent standard normal matrices

$$\Omega \in \mathbb{R}^{n \times k}$$
 and $\Psi \in \mathbb{R}^{\ell \times m}$

with k = 2r + 1, $\ell = 4r + 2$.

► The sketch consists of two matrices that capture the range and co-range of *X*:

$$Y = X\Omega \in \mathbb{R}^{n \times k}$$
 and $W = \Psi X \in \mathbb{R}^{\ell \times m}$

▶ Rank-1 updates to X can be performed on sketch:

▶ Both the storage cost for the sketch and the arithmetic cost of an update are $\mathcal{O}(r(m+n))$.

Recovery from sketch

To recover rank-r approximation \hat{X} from the sketch, compute

1.
$$Y = QR$$
 (tall-skinny QR)

2.
$$B = (\Psi Q)^{\dagger} W$$
 (small QR + backsub)

3.
$$\hat{X} = Q[B]_r$$
 (tall-skinny SVD)

Recovery from sketch

To recover rank-r approximation \hat{X} from the sketch, compute

1.
$$Y = QR$$
 (tall-skinny QR)

2.
$$B = (\Psi Q)^{\dagger} W$$
 (small QR + backsub)

3. $\hat{X} = Q[B]_r$ (tall-skinny SVD)

Theorem (Reconstruction (Tropp Yurtsever U Cevher, 2016))

Fix a target rank r. Let X be a matrix, and let (Y, W) be a sketch of X. The reconstruction procedure above yields a rank-r matrix \hat{X} with

$$\mathbb{E} \|X - \hat{X}\|_{F} \leq 2 \|X - [X]_{r}\|_{F}.$$

Similar bounds hold with high probability.

Previous work (Clarkson Woodruff 2009) algebraically but not numerically equivalent.

Recovery from sketch: intuition

let

$$Y = X\Omega \in \mathbb{R}^{n \times k}$$
 and $W = \Psi X \in \mathbb{R}^{\ell \times m}$

• if Q is an orthonormal basis for $\mathcal{R}(X)$, then

$$X = QQ^*X$$

- if $QR = X\Omega$, then Q is (approximately) a basis for $\mathcal{R}(X)$
- ightharpoonup and if $W = \Psi X$, we can estimate

$$W = \Psi X$$

$$\approx \Psi Q Q^* X$$

$$(\Psi Q)^{\dagger} W \approx Q^* X$$

▶ hence we may reconstruct X as

$$X \approx QQ^*X \approx Q(\Psi Q)^{\dagger}W$$

SketchyCGM

Algorithm 5 SketchyCGM for the model problem (CMOP)

```
Input: Problem data; suboptimality \varepsilon; target rank r Output: Rank-r approximate solution \hat{X} = U\Sigma V^*
```

```
function SketchyCGM
          SKETCH. INIT(m, n, r)
 2
          z \leftarrow 0
          for t \leftarrow 0, 1, \ldots do
                (u, v) \leftarrow \text{MaxSingVec}(-A^*(\nabla f(z)))
 5
               h \leftarrow \mathcal{A}(-\alpha uv^*)
               if \langle z - h, \nabla f(z) \rangle \leq \varepsilon then break for
               \eta \leftarrow 2/(t+2)
8
                z \leftarrow (1 - \eta)z + \eta h
                SKETCH.CGMUPDATE(-\alpha u, v, \eta)
10
          (U, \Sigma, V) \leftarrow \text{Sketch.Reconstruct}()
11
          return (U, \Sigma, V)
12
```

Guarantees

Suppose

- $\triangleright X_{\text{cgm}}^{(t)}$ is tth CGM iterate
- $\triangleright |X_{cgm}^{(t)}|_r$ is best rank r approximation to CGM solution
- $\hat{X}^{(t)}$ is SketchyCGM reconstruction after t iterations

Theorem (Convergence to CGM solution)

After t iterations, the SketchyCGM reconstruction satisfies

$$\mathbb{E} \, \| \hat{X}^{(t)} - X_{\operatorname{cgm}}^{(t)} \|_{\operatorname{F}} \leq 2 \, \| \lfloor X_{\operatorname{cgm}}^{(t)} \rfloor_r - X_{\operatorname{cgm}}^{(t)} \|_{\operatorname{F}} \, .$$

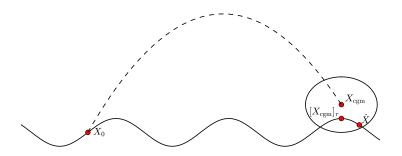
If in addition $X^* = \lim_{t \to \infty} X_{\text{cgm}}^{(t)}$ has rank r, then RHS $\to 0$!

(Tropp Yurtsever U Cevher, 2016)

Convergence when $rank(X_{cgm}) \le r$



Convergence when $rank(X_{cgm}) > r$



Guarantees (II)

Theorem (Convergence rate)

Fix $\kappa > 0$ and $\nu \geq 1$. Suppose the (unique) solution X_{\star} of (CMOP) has $\operatorname{rank}(X_{\star}) \leq r$ and

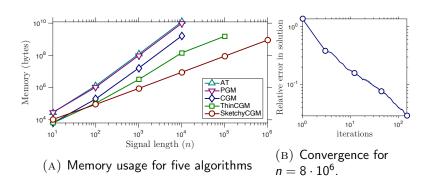
$$f(\mathcal{A}X) - f(\mathcal{A}X_{\star}) \ge \kappa \|X - X_{\star}\|_{\mathrm{F}}^{\nu} \quad \text{for all} \quad \|X\|_{S_{1}} \le \alpha. \quad (1)$$

Then we have the error bound

$$\mathbb{E} \left\| \hat{X}_t - X_\star
ight\|_{\mathrm{F}} \leq 6 \left(rac{2\kappa^{-1}C}{t+2}
ight)^{1/
u} \quad ext{for } t=0,1,2,\ldots$$

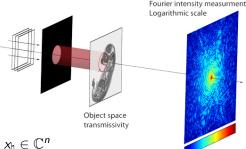
where C is the curvature constant (Eqn. (3), Jaggi 2013) of the problem (CMOP).

SketchyCGM is scalable



```
\begin{array}{rcl} \mathsf{PGM} &=& \mathsf{proximal\ gradient\ (via\ TFOCS\ (Becker\ Candès\ Grant,\ 2011))} \\ \mathsf{AT} &=& \mathsf{accelerated\ PGM\ (Auslander\ Teboulle,\ 2006)\ (via\ TFOCS),} \\ \mathsf{CGM} &=& \mathsf{conditional\ gradient\ method\ (Jaggi,\ 2013)} \\ \mathsf{ThinCGM} &=& \mathsf{CGM\ with\ thin\ SVD\ updates\ (Yurtsever\ Hsieh\ Cevher,\ 2015)} \\ \mathsf{SketchyCGM} &=& \mathsf{ours,\ using\ } r=1 \end{array}
```

Application: Phase retrieval



- ▶ image with n pixels $x_{\natural} \in \mathbb{C}^n$
- acquire noisy nonlinear measurements $b_i = |\langle a_i, x_{\dagger} \rangle|^2 + \omega_i$
- relax: if $X = x_{\natural}x_{\natural}^*$, then

$$|\langle a_i, x_{\natural} \rangle|^2 = x_{\natural} a_i^* a_i x_{\natural}^* = \operatorname{tr}(a_i^* a_i x_{\natural}^* x_{\natural}) = \operatorname{tr}(a_i^* a_i X)$$

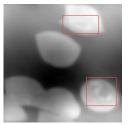
recover image by solving

minimize
$$f(AX; b)$$
 subject to $\operatorname{tr} X \leq \alpha$ $X \succeq 0$.

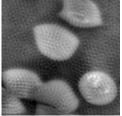
SketchyCGM is reliable

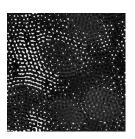
Fourier ptychography:

- imaging blood cells with A = subsampled FFT
- n = 25,600, d = 185,600
- ▶ rank(X_{\star}) ≈ 5 (empirically)



(A) SketchyCGM





(B) Burer–Monteiro (C) Wirtinger Flow

- brightness indicates phase of pixel (thickness of sample)
- red boxes mark malaria parasites in blood cells

Conclusion

SketchyCGM offers a proof-of-concept **convex method** with **optimal storage** for low rank matrix optimization using two new ideas:

- Drive the algorithm using a smaller (dual) variable.
- Sketch and recover the decision variable.

References:

- J. A. Tropp, A. Yurtsever, M. Udell, and V. Cevher. Randomized single-view algorithms for low-rank matrix reconstruction. SIMAX 2017.
- A. Yurtsever, M. Udell, J. A. Tropp, and V. Cevher. Sketchy Decisions: Convex Optimization with Optimal Storage. AISTATS 2017.