# ZigZag Conversion

## Problem

The string "PAYPALISHIRING" is written in a zigzag pattern on a given number of rows like this:

1	2	3	4	5	6	7
Р		Α		Н		N
Α	Р	L	S	ı	ı	G
Υ		ı		R		

And then read line by line: "PAHNAPLSIIGYIR"

#### example 1:

```
Input: s = "PAYPALISHIRING", numRows = 3
```

Output: "PAHNAPLSIIGYIR"

# example 2:

```
Input: s = "PAYPALISHIRING", numRows = 4
```

Output: "PINALSIGYAHRPI"

# Solution

# Approach 1: Sort by Row

## Intuition

By iterating through the string from left to right, we can easily determine which row in the Zig-Zag pattern that a character belongs to.

# Algorithm

We can use min(numRows, len(s)) lists to represent the non-empty rows of the Zig-Zag Pattern.

Iterate through s from left to right, appending each character to the appropriate row. The appropriate row can be tracked using two variables: the current row and the current direction.

The current direction changes only when we moved up to the topmost row or moved down to the bottommost row.

#### Java

```
class Solution {
  public String convert(String s, int numRows) {
    if (numRows == 1) return s;
}
```

## Python

```
class Solution(object):
    def convert(self, s, numRows):
        :type s: str
        :type numRows: int
        :rtype: str
        if numRows == 1:
            return s
        rows = ["" for i in range(numRows)]
        curRow = 0
        goingDown = False
        for i in range(len(s)):
            rows[curRow] += s[i]
            if (curRow == 0 or curRow == numRows - 1):
                goingDown = ~goingDown
            curRow += 1 if goingDown else -1
        ret = ""
        for row in rows:
            ret += row
        return ret
```

# Complexity Analysis

- Time Complexity: O(n), where n == len(s)
- Space Complexity: O(n)

#### Approach 2: Visit by Row

#### Intuition

Visit the characters in the same order as reading the Zig-Zag pattern line by line.

# Algorithm

Visit all characters in row 0 first, then row 1, then row 2, and so on...

For all whole numbers k,

- Characters in row 0 are located at indexes k(2\*numRows-2)
- Characters in row numRows 1 are located at indexes k(2\*numRows 2) + numRows 1
- Characters in inner row i are located at indexes k(2\*numRows-2)+i and (k+1)(2\*numRows-2)-i

Java(Official)

```
class Solution {
   public String convert(String s, int numRows) {

      if (numRows == 1) return s;

      StringBuilder ret = new StringBuilder();
      int n = s.length();
      int cycleLen = 2 * numRows - 2;

      for (int i = 0; i < numRows; i++) {
            for (int j = 0; j + i < n; j += cycleLen) {
                ret.append(s.charAt(j + i));
            if (i != 0 && i != numRows - 1 && j + cycleLen - i < n)
                ret.append(s.charAt(j + cycleLen - i));
            }
        return ret.toString();
    }
}</pre>
```

Java(My Own)

```
class Solution {
    public String convert(String s, int numRows) {
        String news = "";
         if(s.length() \leftarrow 0 \mid | numRows == 1)
             return s;
        int index = 0;
        for(int i = 0;i<numRows;i++){</pre>
             index = i;
             Boolean flag = true;
             int fold = 1;
             if(i==numRows-1)
                 fold = 2;
             while(index<s.length()){</pre>
                 news += s.charAt(index);
                 index = (fold * (numRows - 1) - index) * 2 + index;
                 if(i==0 \mid | i==numRows-1)
                      fold += 2;
                 else
                      fold ++;
             }
        }
        return news;
    }
}
```

# Python

```
class Solution:
    def convert(self, s, numRows):
        """
        :type s: str
        :type numRows: int
        :rtype: str
```

```
if numRows == 1:
    return s
strlen = len(s)
cycleLen = 2 * numRows - 2
ret = ""
for i in range(numRows):
    for j in range(0, strlen - i, cycleLen):
        ret += s[j+i]
        if (i != 0 and i != numRows - 1 and j + cycleLen - i < strlen):
        ret += s[j + cycleLen - i]
return ret</pre>
```

# **Complexity Analysis**

• Time Complexity: O(n), where n == len(s)

• Space Complexity: O(n)