

Implementation of Databases

Chapter 1: Architectures of Database Systems

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Lecture

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Excercises

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1.1 Goals and Tasks of DBMS

Learning Goals

At the end of this section you will be able to

- ✓ explain, what is DBS and what is a DBMS
- ✓ name and describe the goals and tasks of a DBMS
- ✓ discuss the relationship between software applications and DBMS



What is a Database System?

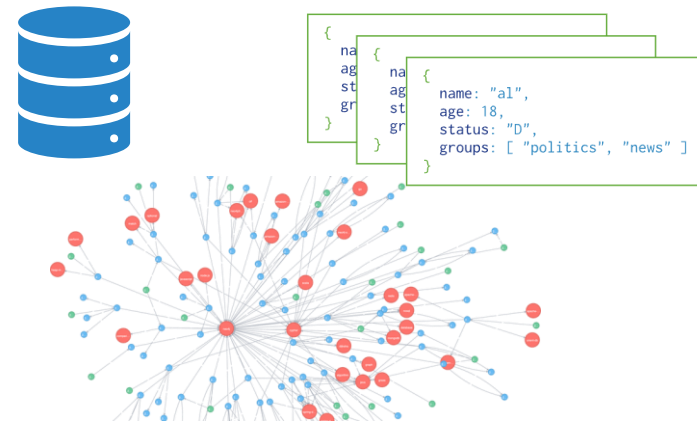
Database Management
System (DBMS)

*software to manage
databases & data*

+

Database

*self-describing collection
of related data*



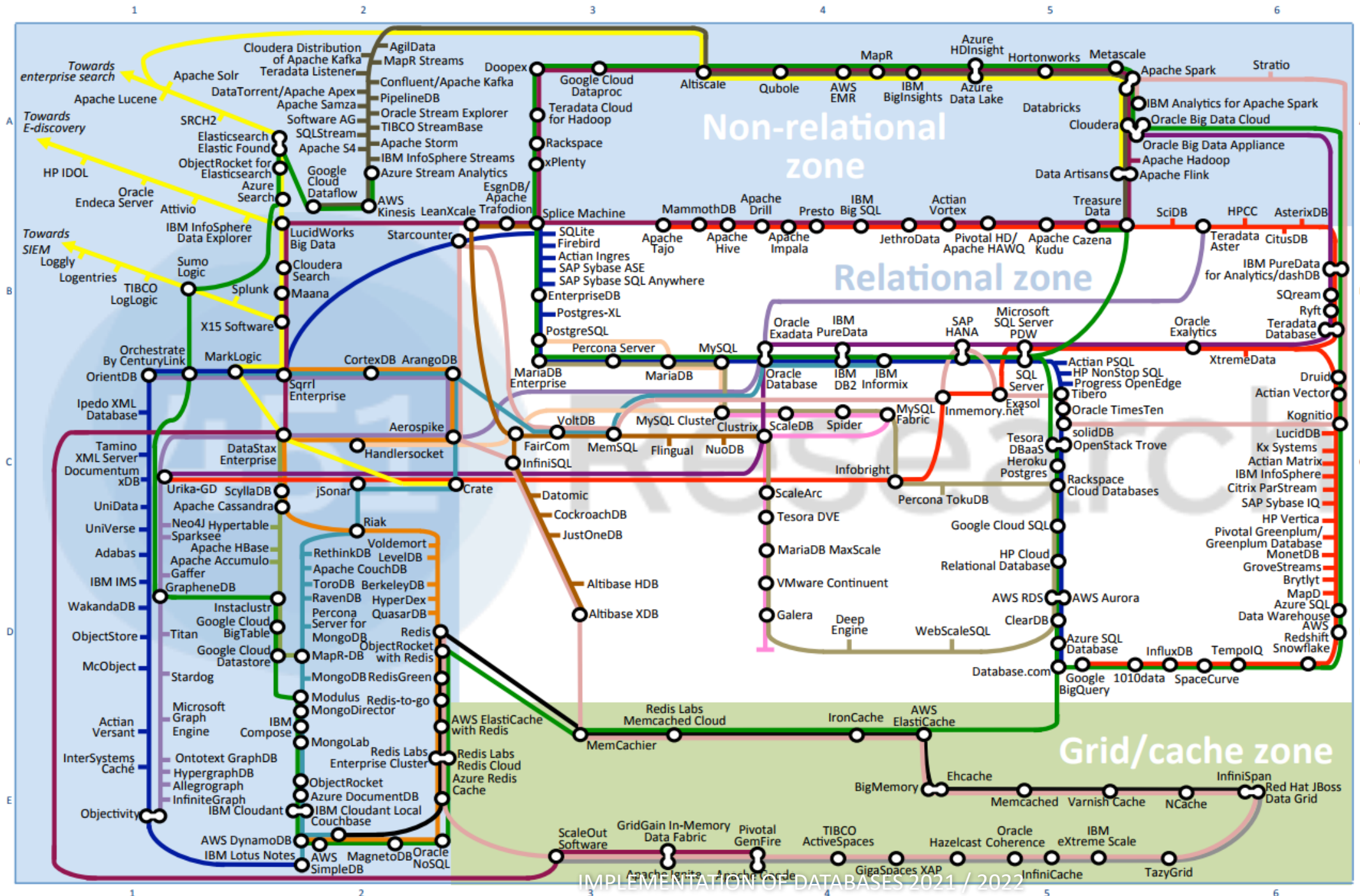
Data Platforms Map January 2016

Key:

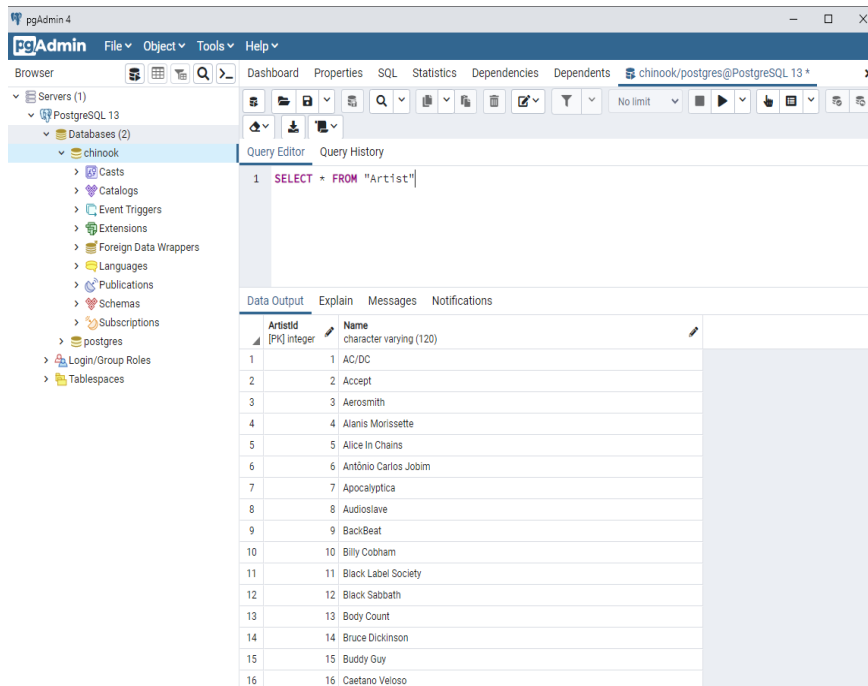
- General purpose
- Specialist analytic
- as-a-Service
- BigTables
- Graph
- Document
- Key value stores
- Key value direct access
- Hadoop
- MySQL ecosystem
- Advanced clustering/sharding
- New SQL databases
- Data caching
- Data grid
- Search
- Appliances
- In-memory
- Stream processing

<https://451research.com/state-of-the-database-landscape>

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Goals and Tasks (1)



The screenshot shows the pgAdmin 4 interface. On the left, the 'Servers' tree is expanded to show 'PostgreSQL 13' and 'Databases (2)', with 'chinook' selected. The 'Query Editor' is open, showing a query: `SELECT * FROM "Artist"`. The 'Data Output' tab is active, displaying a table with 16 rows. The table has two columns: 'ArtistId' (PK Integer) and 'Name' (character varying (120)).

ArtistId	Name
1	AC/DC
2	Accept
3	Aerosmith
4	Alanis Morissette
5	Alice In Chains
6	Antônio Carlos Jobim
7	Apocalyptica
8	Audioslave
9	BackBeat
10	Billy Cobham
11	Black Label Society
12	Black Sabbath
13	Body Count
14	Bruce Dickinson
15	Buddy Guy
16	Caetano Veloso

Data Independence

- Manage data independent of applications, make data available for different applications
- **Physical data independence:** logical schema is independent of physical structure
- **Logical data independence:** external schema (for users / applications) is independent of logical schema

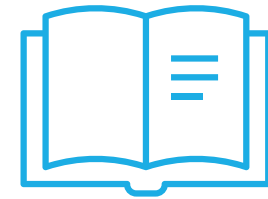
Data Manipulation & Retrieval

- DML: Data Manipulation Language
 - Retrieve / select
 - Insert / delete / update
- CRUD operations: create – read – update – delete

Goals and Tasks (2)

Structure definition and integrity assurance

- DDL: Data Definition Language
- Data dictionary / system catalog / metadata
- Integrity conditions / assertions / constraints



Protection of databases in multi-user mode

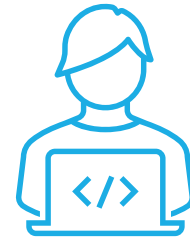
- Transaction management
 - **ACID**: Atomicity, Consistency, Isolation, Durability
- Recovery: Restart on error
- Data security and data protection



1.1 Goals and Tasks of DBMS (3)

Realization of User Interfaces

- Interactive end-user interface
- API: Application Programming Interface



Performance control

- Monitoring of the system load and runtime behavior
- Index
- Clustering / data aggregation



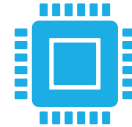
DBMS and Applications



In contrast to OS,
DBMS is an
application



Databases: Consistent
non-volatile memory



Application:
Presentation of data,
data processing



Application and
DBMS often run on
different computers



Communication:
Connections between
software systems
(partially a job of OS)



ISO-OSI Reference
Model for
communication

AP	7	Application
	6	Presentation
OS	5	Session
	4	Transport
	3	Network
	2	Data Link
	1	Physical



Quiz



1.2 Basic Architecture of a DBMS

Learning Goals

At the end of this section you will be able to

- ✓ Explain the basic architecture of a DBMS
- ✓ Name the most important layers and interfaces of the basic and the 5 layer models and describe their tasks
- ✓ Give an example, how a query is processed throughout the layers



System Architecture Goal: Modularization

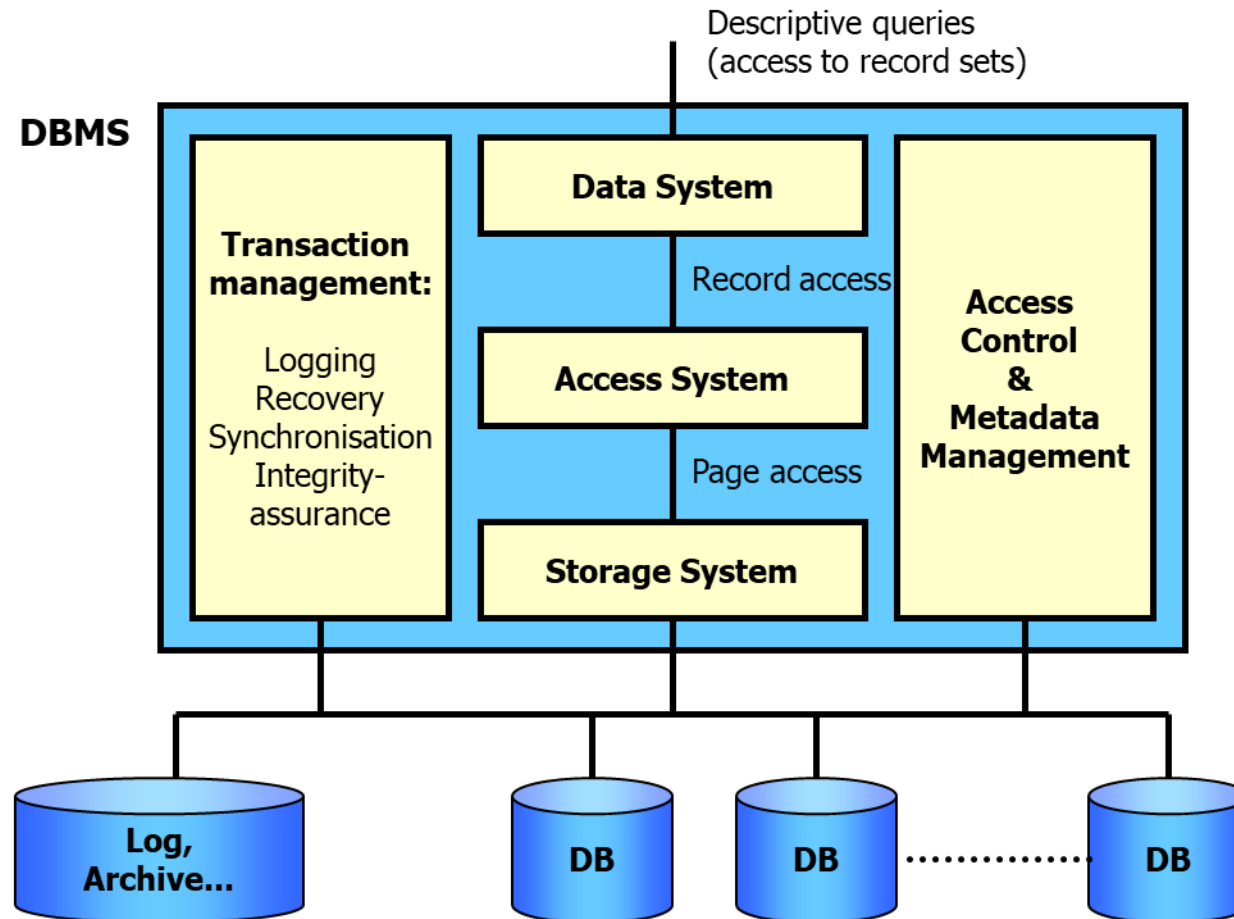


- Abstraction (concentrate on substantial points)
- Localization (of procedures & data)
- Information hiding principle / black box
- Completeness (on an abstract level)
- Verifiability

Concepts for modularization

- Functional abstraction
- Data abstraction
- Generic modules with objects and methods

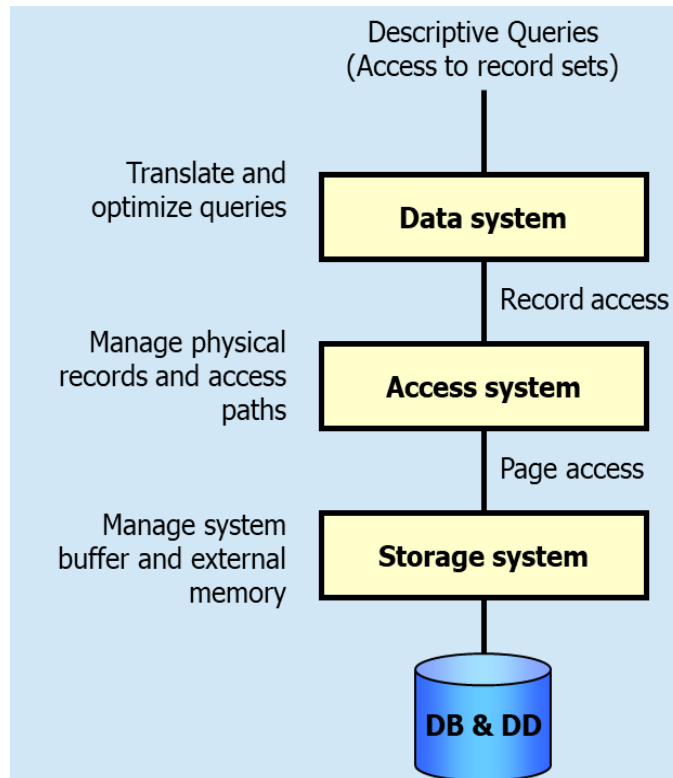
Simplified Architecture of a DBS



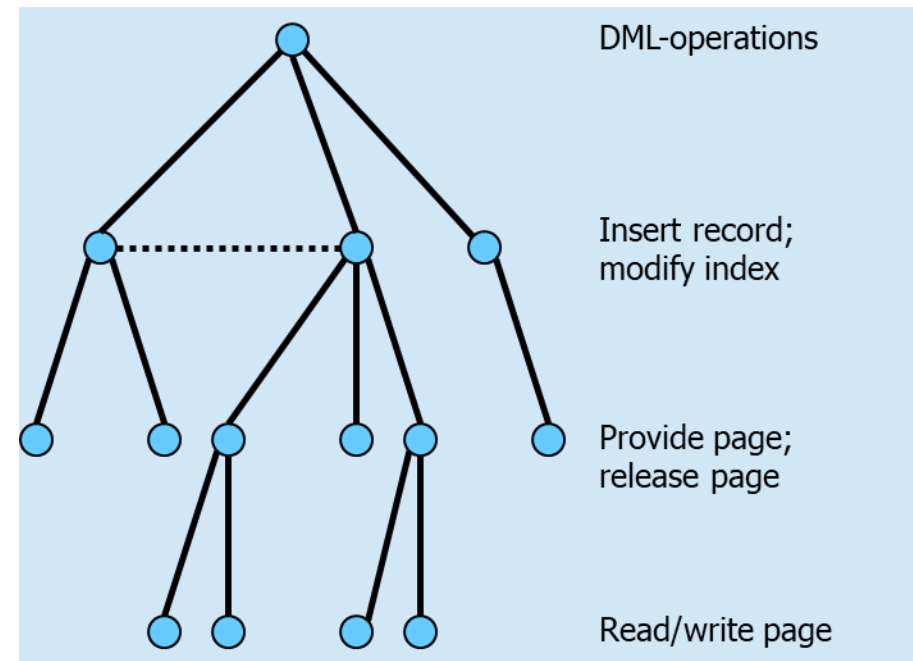
[Härder & Rahm, 2001]

Interaction of the Layers

Components



Control Flow



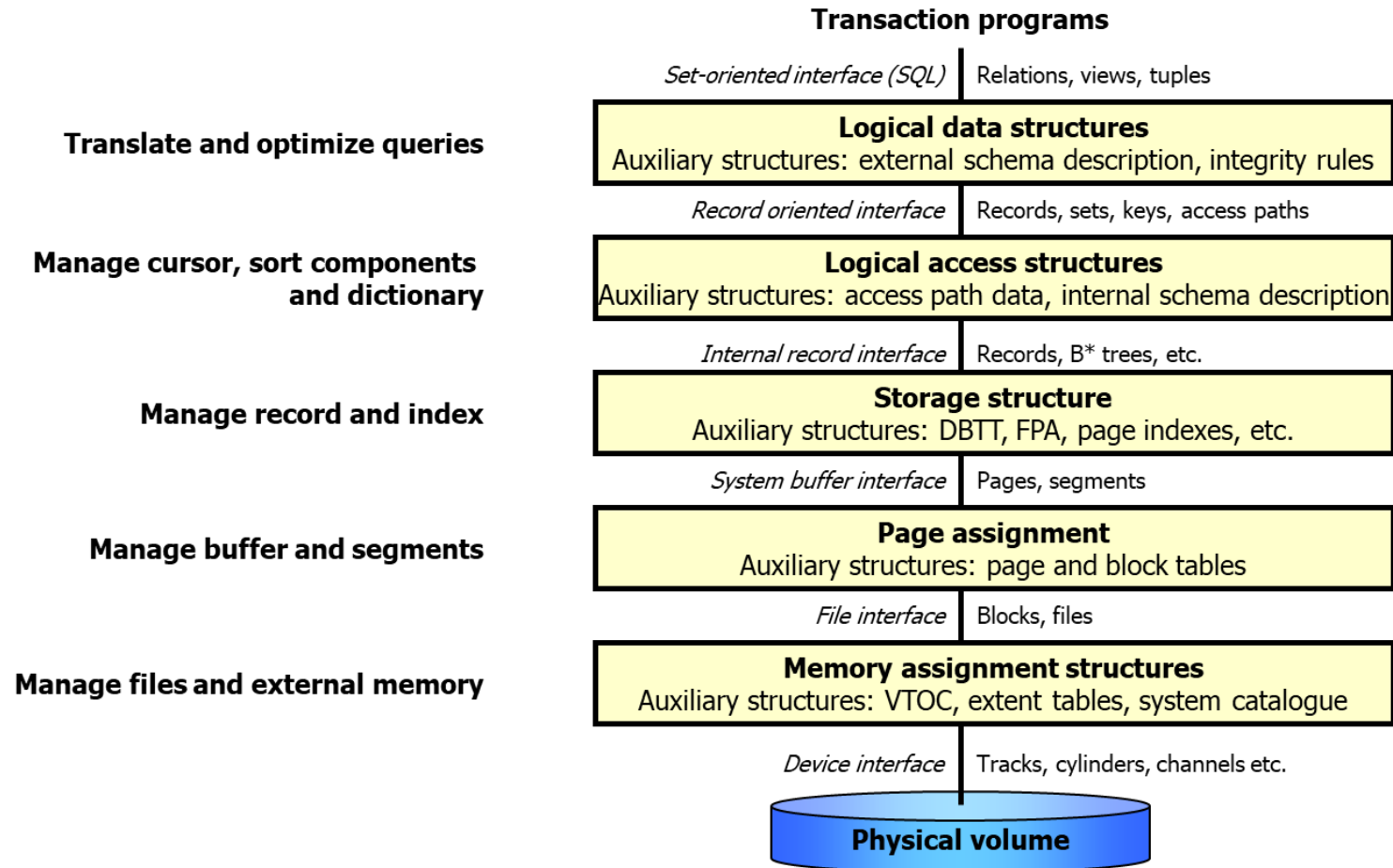
[Härder & Rahm, 2001]

Background Information on Hardware

	Capacity	Speed	Access time
L1-Cache	16-256 KB	300 GB/s	4 ns
L2-Cache	256 KB-4 MB	300 GB/s	10 ns
L3-Cache	few MBs	~50 GB/s	5-15 ns
RAM	GBs	Up to 10.000 MB/s	10-20 ns
Solid State Disc	GB – TB	200 MB/s	~0.1 ms
Hard disc	TBs	Up to 1.000 MB/s	7 ms
CD/DVD	640 MB – 20 GBs	10 MB/s	150 ms
Streamer	4 GB - >100 GB	2-10 MB/s	100 ms - >10 s
Network	-	1/10/100/1000 MB/s	ca. 1 ms

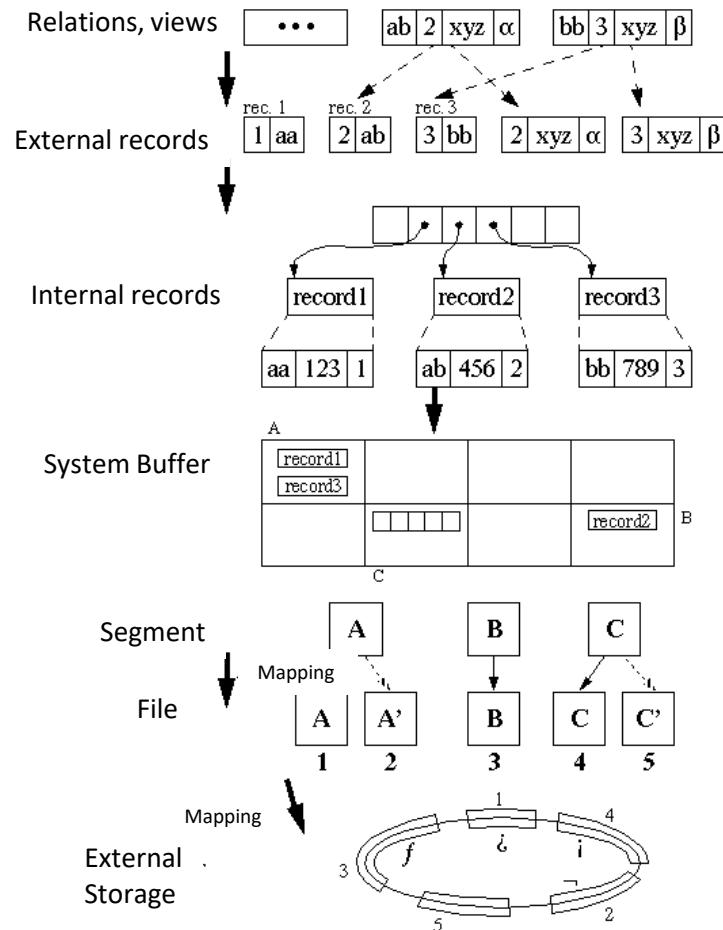
→ Prefer main memory, reduce disk accesses!

Five layer model of a DBS



[Härder & Rahm, 2001]

Example



```
SELECT p.Name FROM Professor p, Vorlesung v
WHERE
    p.Rang='C4' AND v.Titel='Logik'
    AND p.PersNr = v.gelesenVon
```

$$\{ p.name \mid p \in Professor \wedge \exists v \in Vorlesung \wedge p.Rang = 'C4' \wedge v.Titel = 'Logik' \wedge p.PersNr = v.gelesenVon \}$$

$$\pi_{Name}(\sigma_{Titel=Logik \wedge Rang=C4}(Professor \bowtie_{PersNr=gelesenVon} Vorlesung))$$

$$\pi_{Name}(\sigma_{Rang=C4}(Professor) \bowtie_{PersNr=gelesenVon} \sigma_{Titel=Logik}(Vorlesung))$$

```
OPEN CURSOR Vorlesung(Titel='Logik')
FIND NEXT record ...
OPEN CURSOR Professor(Rang='C4')
...
```

```
B+-TREE-SEARCH Vorlesung(Titel='Logik')
FETCH RECORD Vorlesung(...,gelesenVon)
...
B+-TREE-SEARCH Professor(PersNr=gelesenVon)
...
```

```
LOAD PAGE 123
WRITE PAGE 345
...
```

Data Independence: An Overview

Layer	What is hidden?
Logical data structures	Position indicator and explicit relations in the schema
Logical access paths	Number and kind of the physical access paths; internal representation of records
Storage structures	Management of buffers, logging
Page assignment structures	File mapping, indirect page assignment
Memory assignment structures	Technical features and technical details of external storage media

Problems:

Due to high specialization, functionality of operating system often not usable

- Segment-file mapping
- Paging
- Shadow memory
- Buffer management
- Dispatching



Quiz



1.3 Evolution of DBMS

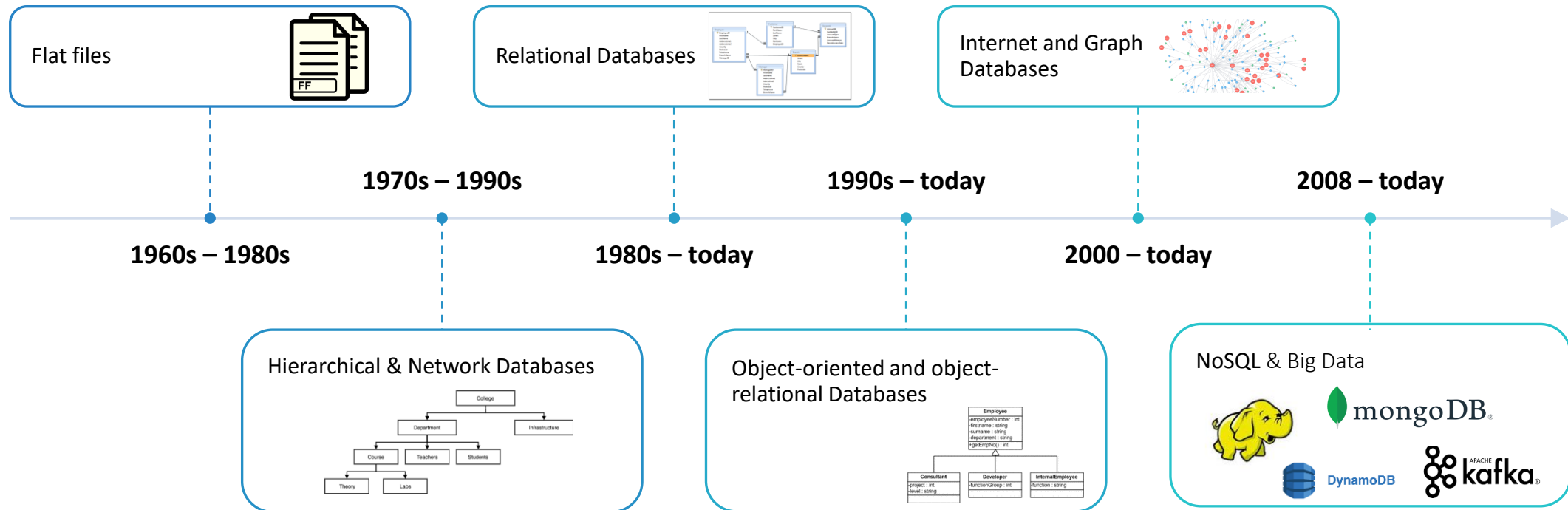
Learning Goals

At the end of this section you will be able to

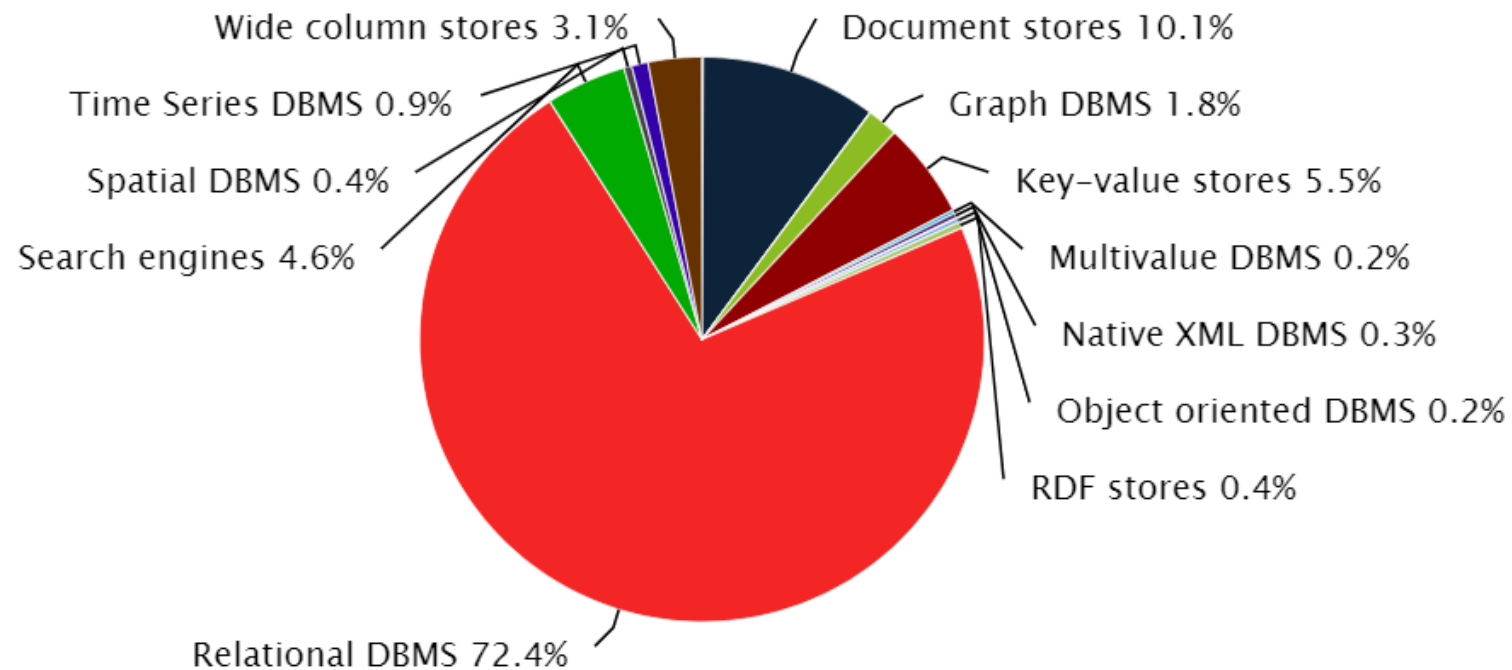
- ✓ name the most important stages in DBMS evolution
- ✓ classify a DBMS along different dimensions
- ✓ name different client-server architectures and discuss pros and cons
- ✓ explain tasks of client and server and name three different server types
- ✓ explain distributed database principles
- ✓ explain what is a big data architecture



Evolution Time Line



Data Model Popularity Today



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Classification Dimensions of DBMS [Elmasri & Navathe, 2017]

Data Model

Number of users

Number of sites

Costs & License

Types of
access path

Purpose

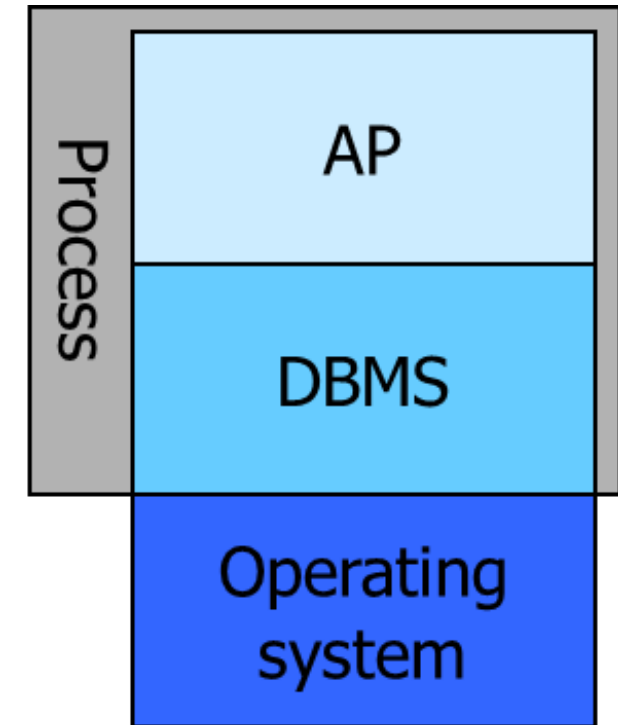
Embedded DBMS: 1-Layer-Architecture, Single User DBMS

- One process, one address space
- No concurrency control
- Simple crash recovery

AP: Application program with DB calls

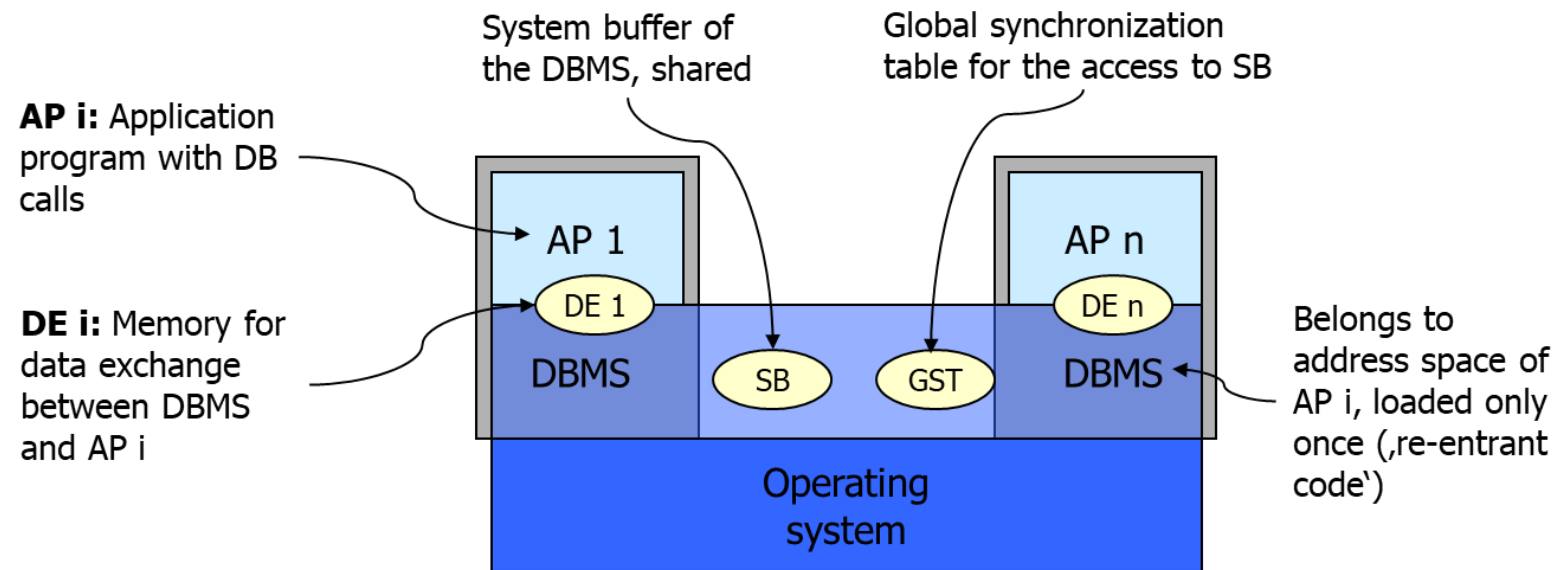
DBMS: Belongs to the address space of the AP

Example: PC database systems (MS Access)



Embedded DBMS: 1-Layer-Architecture – Multi-user DBMS

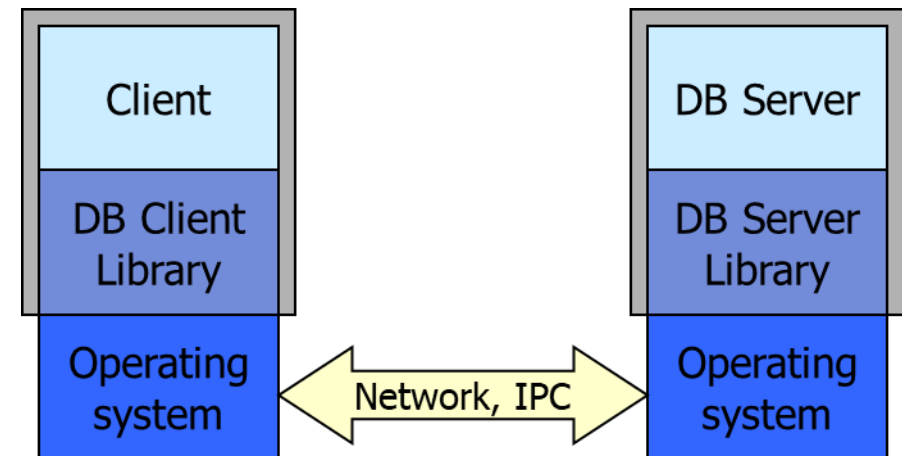
- Multiple processes, communication via shared address space
- Very efficient data exchange via shared memory, but
- No security concerning errors in the AP



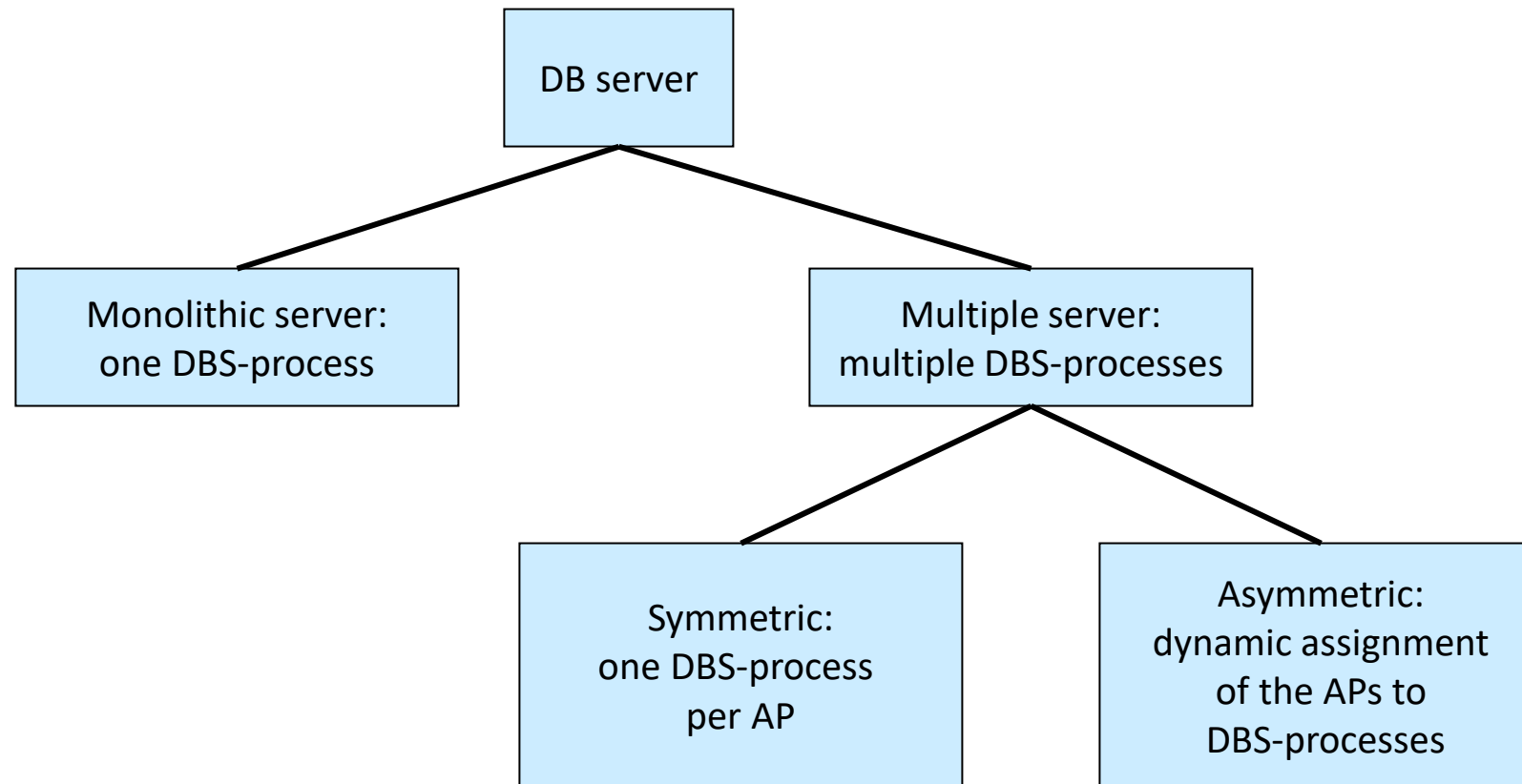
Example: IBM System/R* – research prototype

Embedded DBMS: 2-Layer-Architecture (“Client/Server”)

- Client and server totally separated \Rightarrow distributed access to DB
- Communication among clients and servers via a network or IPC.
- Specialized protocols used (JDBC, Net8, TCP/IP)
- Clear separation of client and server



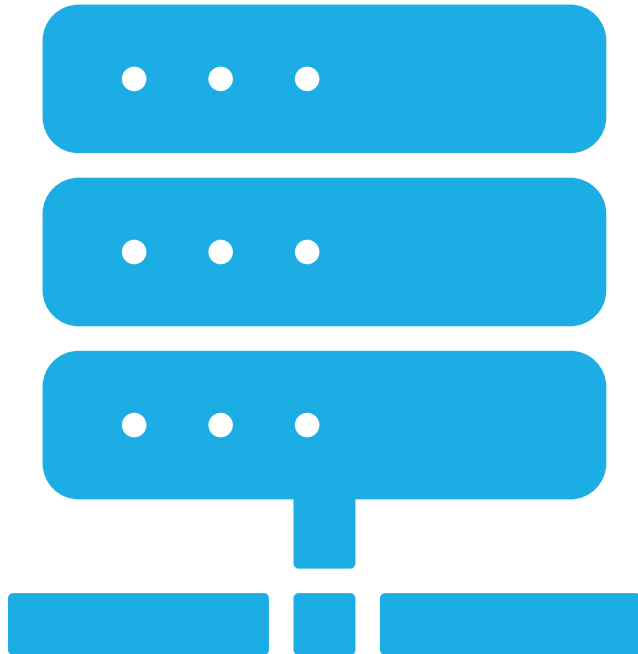
Classification in Client/Server Architecture





Multiple DB Server Processes

- Communication among the servers via “shared memory”
- Communication among clients and servers/dispatcher via OS-mechanisms (IPC) or network software
- **Symmetric assignment:**
Each client assigned to exact one server process.
Static assignment, fixed number n of servers stated in advance
 \Rightarrow maximal degree of parallelism is n
- **Asymmetric assignment:**
Each client assigned to a server process by a dispatcher.
Fixed number n of servers stated in advance, but degree of parallelism can be higher.

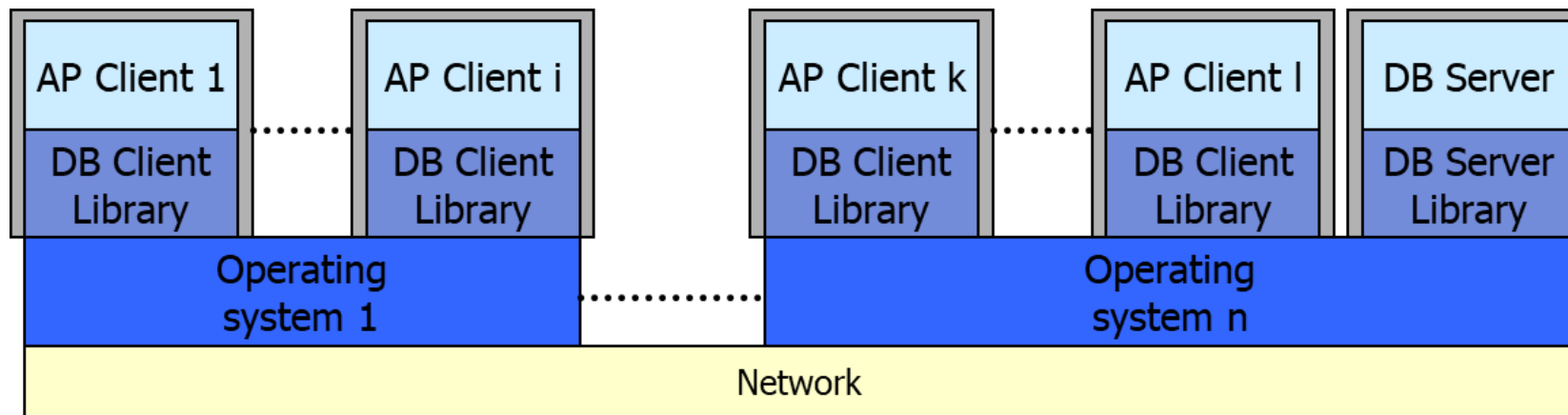


Single DB Server Process

- Synchronized access to system buffer and central system tables
- Server uses multi-threading (“re-entrant code”)
- Only one server process for many clients
- DB server process is preferred by OS

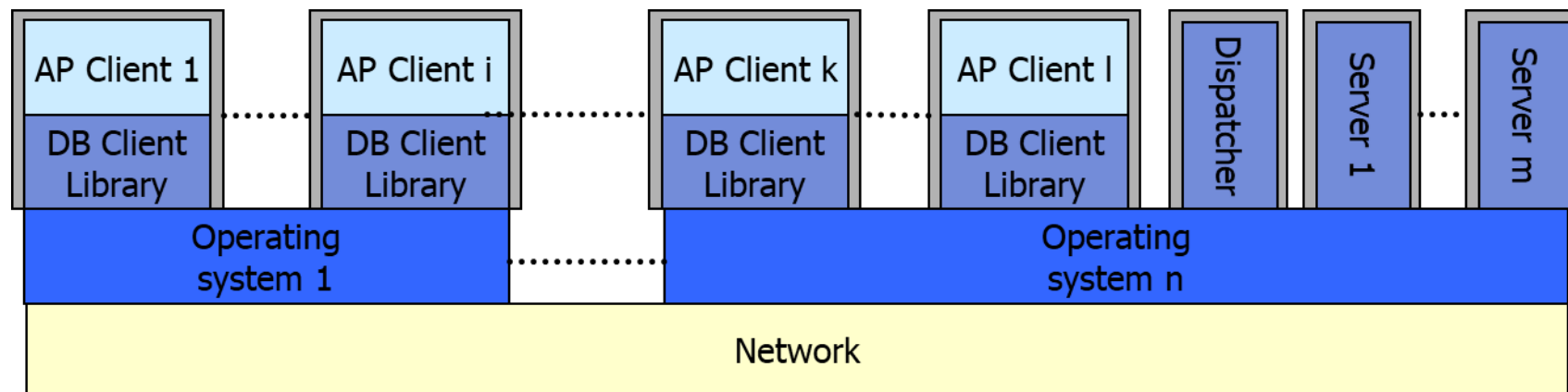
Monolithic Server

- Own resource management, duplicates OS functions
- Simple communication in the server via shared memory
- Example: PostgreSQL



Multiple Servers

- DBMS is a compound of different processes
- Communication via operating system or network
- Process scheduling by OS, advantageous in multi-processor computers, because OS manages processor allocation.
- Example: Oracle, IBM DB2, MySQL, MS SQL Server, Sybase



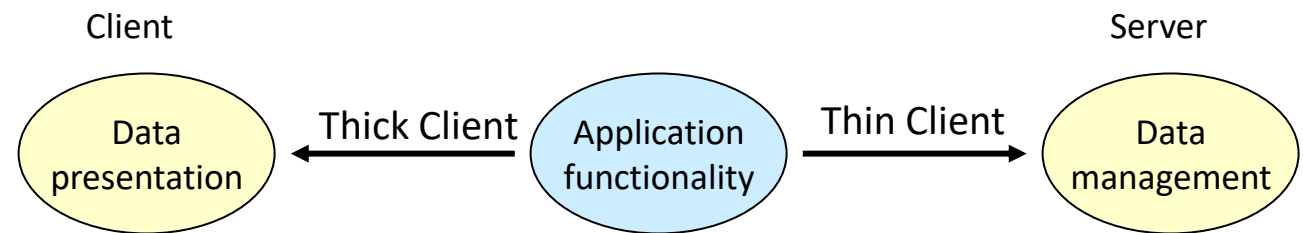
Tasks of Client and Server

Tasks of a server

- Data management
 - Relation Server
 - Object Server
 - Page Server
- Application functionality

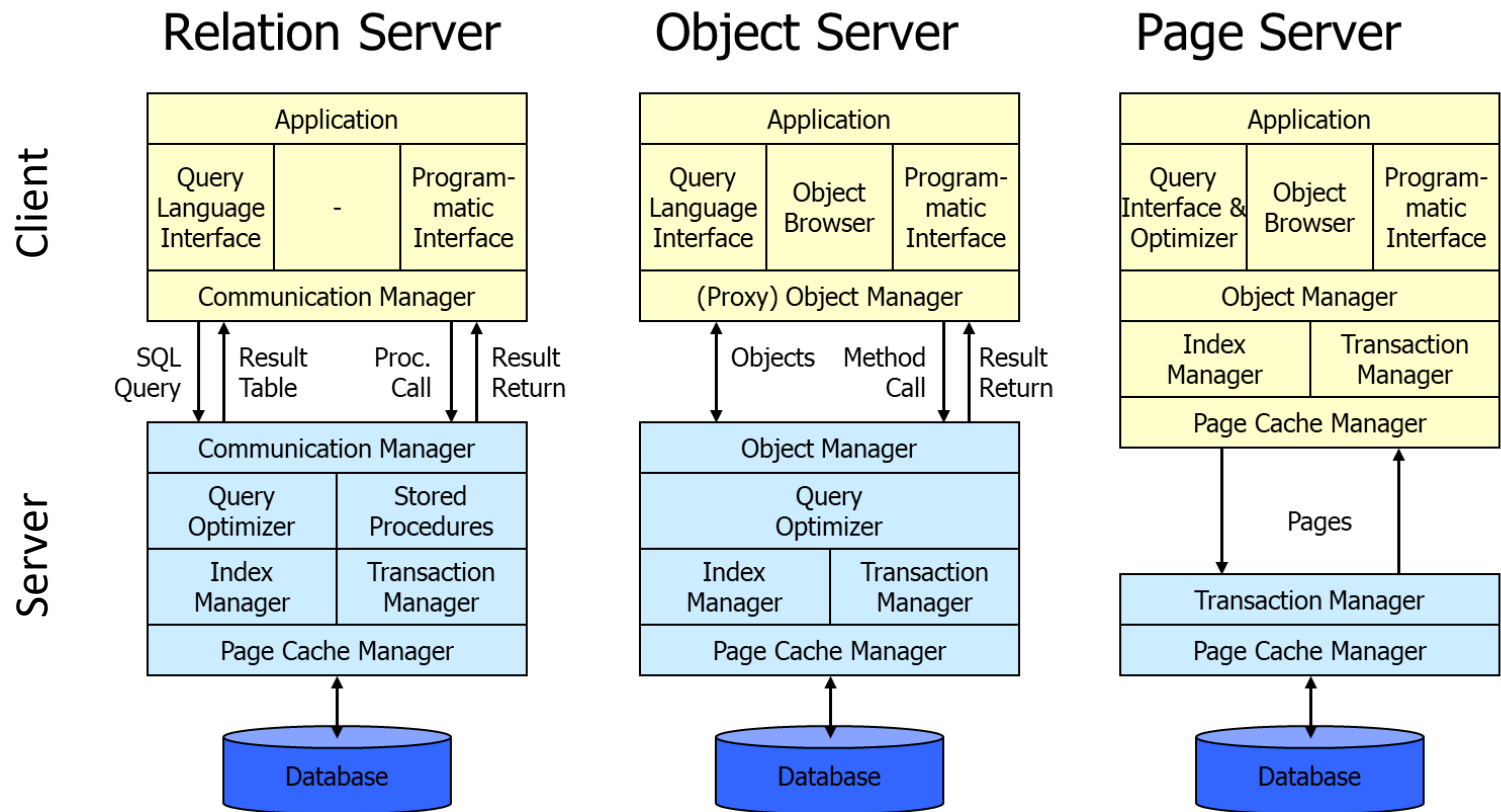
Tasks of a client

- Presentation of data
- Application functionality



Server Types

[Härder & Rahm, 2001]





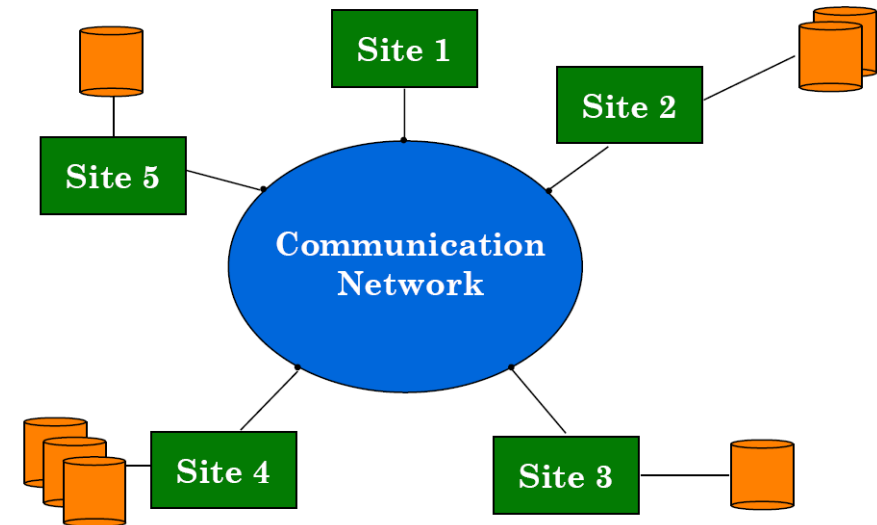
Quiz

Distributed Database Systems

- A distributed database (DDB) is a collection of multiple, *logically interrelated* databases distributed over a *computer network*
- A distributed database management system (D-DBMS) is the software system that permits the management of the distributed database and makes the distribution *transparent* to the users

Distributed database system (DDBS)

DDBS = DDB + D-DBMS



[Özsu & Valduriez, 2011]

Promises of D-DBMS

- **Transparent** management of distributed, fragmented, and replicated data
- Improved **reliability/availability** through distributed transactions
- Improved **performance**
- Easier and more economical system expansion
→ **Scalability**



Implicit Assumptions



Data stored at a number of sites

Each site *logically consists of a single processor.*



Processors at different sites are interconnected by a computer network

no multiprocessors → parallel database systems



Distributed database is a database, not a collection of files

Data is logically related as exhibited in the users' access patterns
→ relational data model



D-DBMS is a full-fledged DBMS

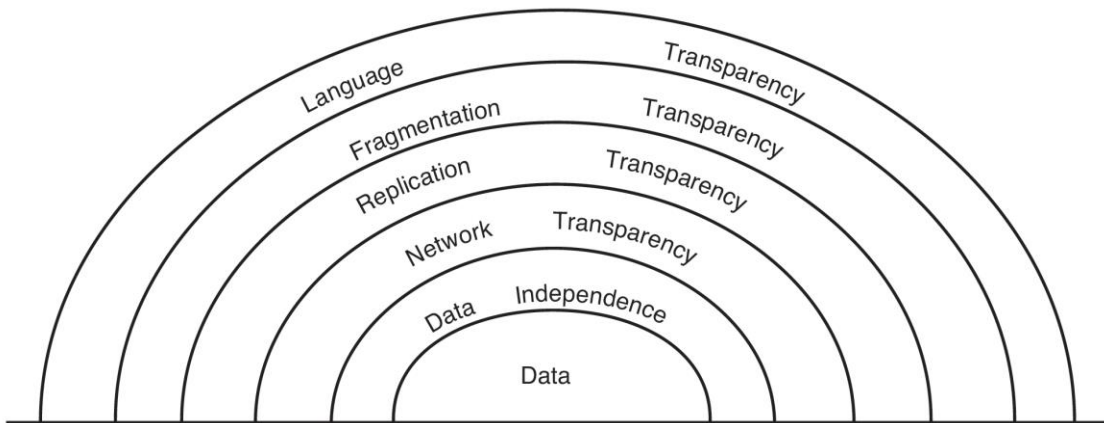
→ not remote file system, not a Transaction Processing system

Transparency

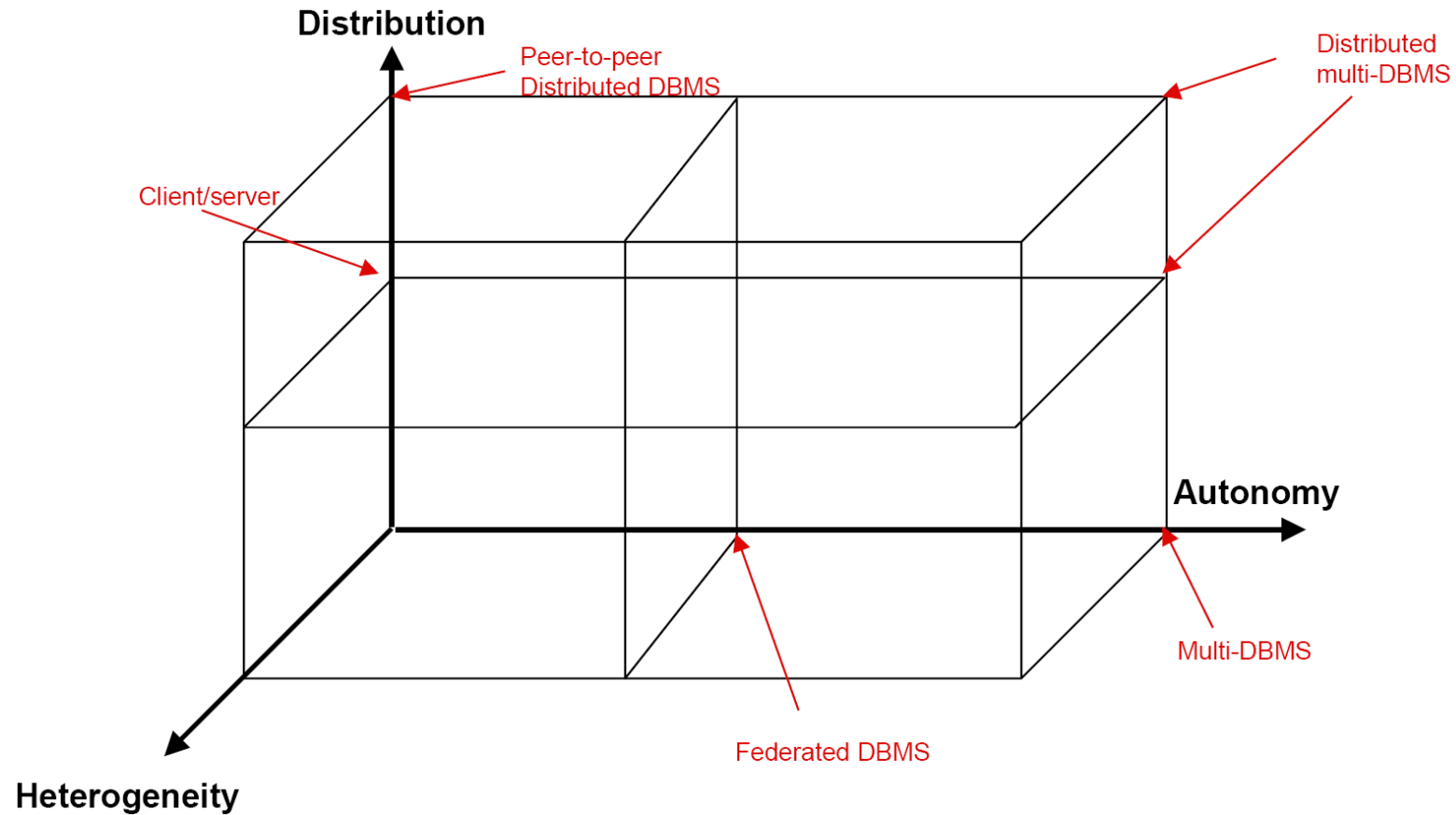
... is the separation of higher-level semantics of a system from lower-level implementation issues

Fundamental issue is to provide **data independence** in the distributed environment

- Network (distribution) transparency
- Replication transparency
- Fragmentation transparency
 - horizontal fragmentation: selection
 - vertical fragmentation: projection
 - hybrid



D-DBMS: Implementation Alternatives



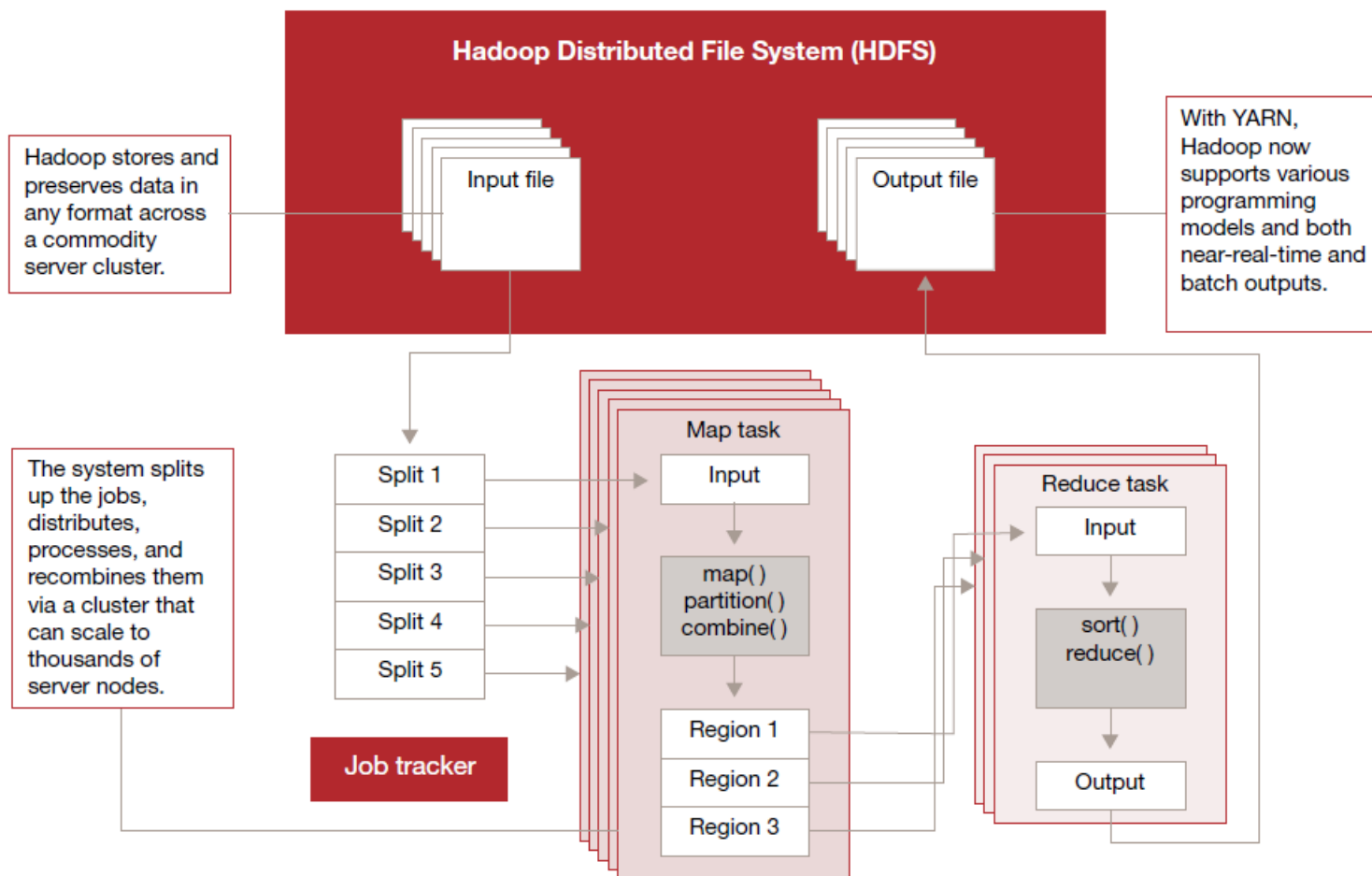
Big Data Architectures

- Big Data requires large, scalable, distributed architectures
- Four/Six Vs : Volume, Velocity, Variety, Veracity, [Hofstee & Nowka, 2013], (Value, Variability)
- Heterogeneity
 - Sources
 - Systems
 - Requirements
 - Client Applications

➔ Big Data Systems are complex eco-systems, independent components have to be integrated

➔ Hadoop is not a Big Data system, it is just a component

A basic Hadoop architecture for scalable data lake infrastructure

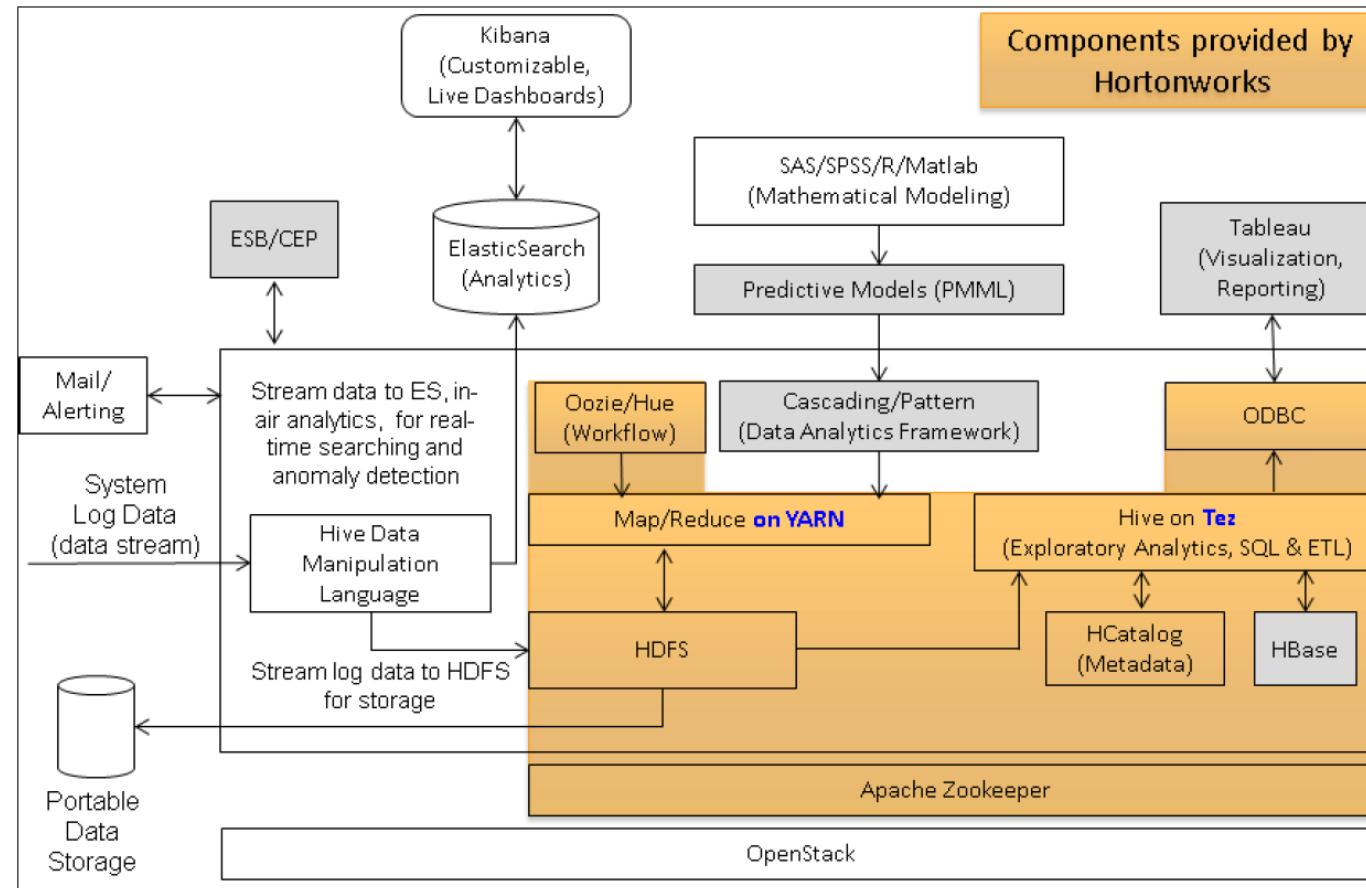


Source: Electronic Design, 2012, and Hortonworks, 2014

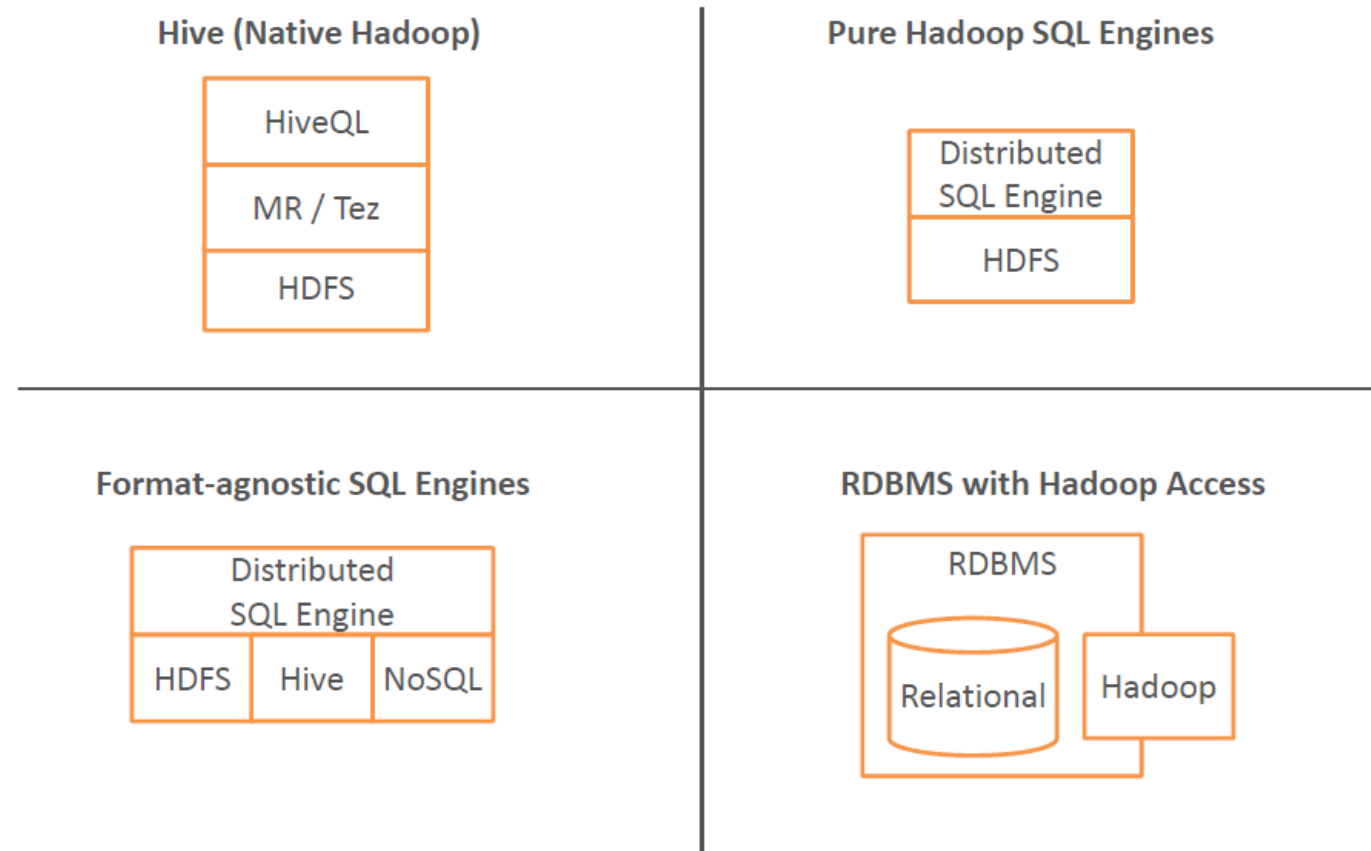
Source: <https://www.pwc.com/us/en/technology-forecast/2014/cloud-computing/assets/pdf/pwc-technology-forecast-data-lakes.pdf>

Hadoop as a Basic Big Data Platform

Example Big Data Architecture [Boci & Thistlethwaite, 2015]



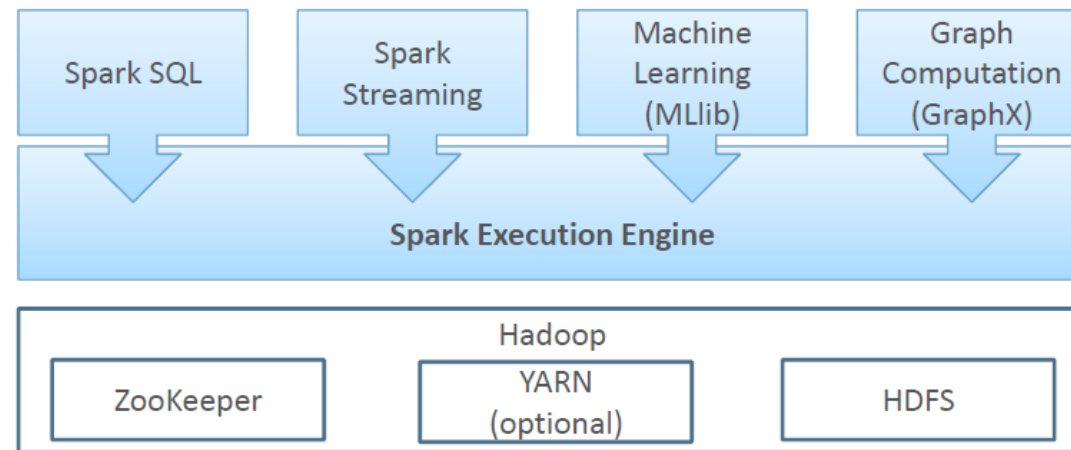
Architectures for Hadoop with SQL



[Albrecht, 2015]

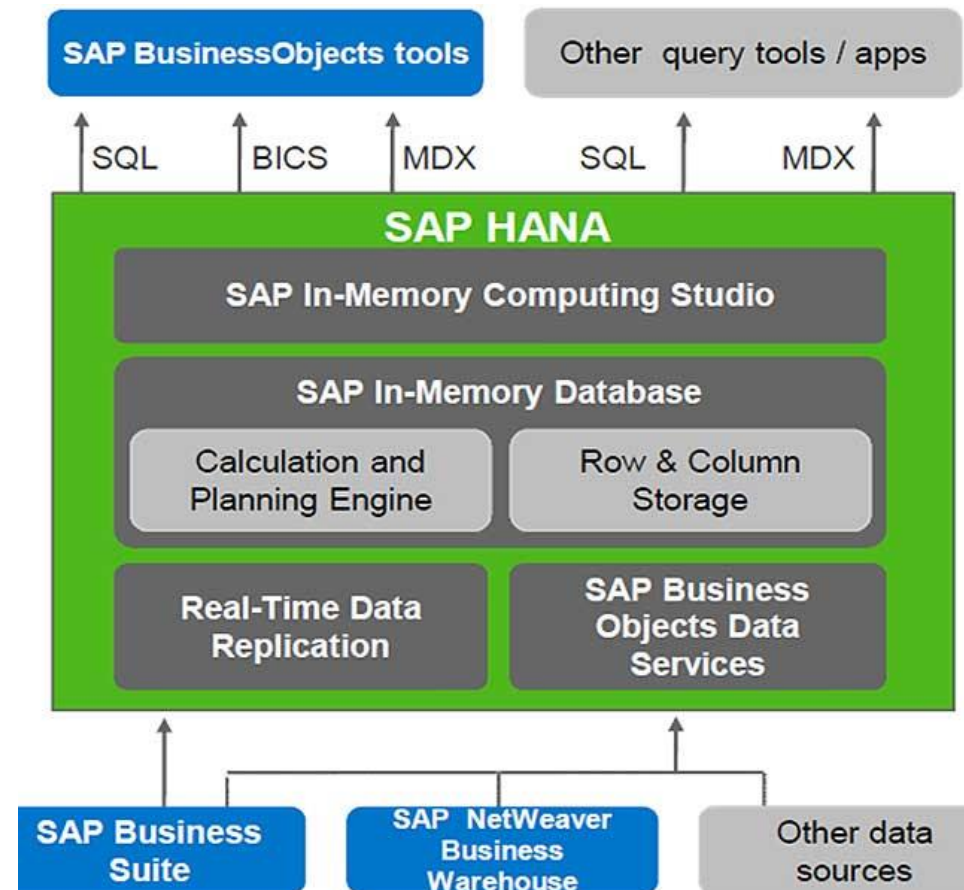
In-Memory Database Systems

- Distributed data processing, but do as much as possible in-memory and avoid I/O operations (to disk or distributed file system)
- Example: Apache Spark
 - Distributed In-Memory Computing Framework
 - General framework for all kinds of SQL and non-SQL analytics

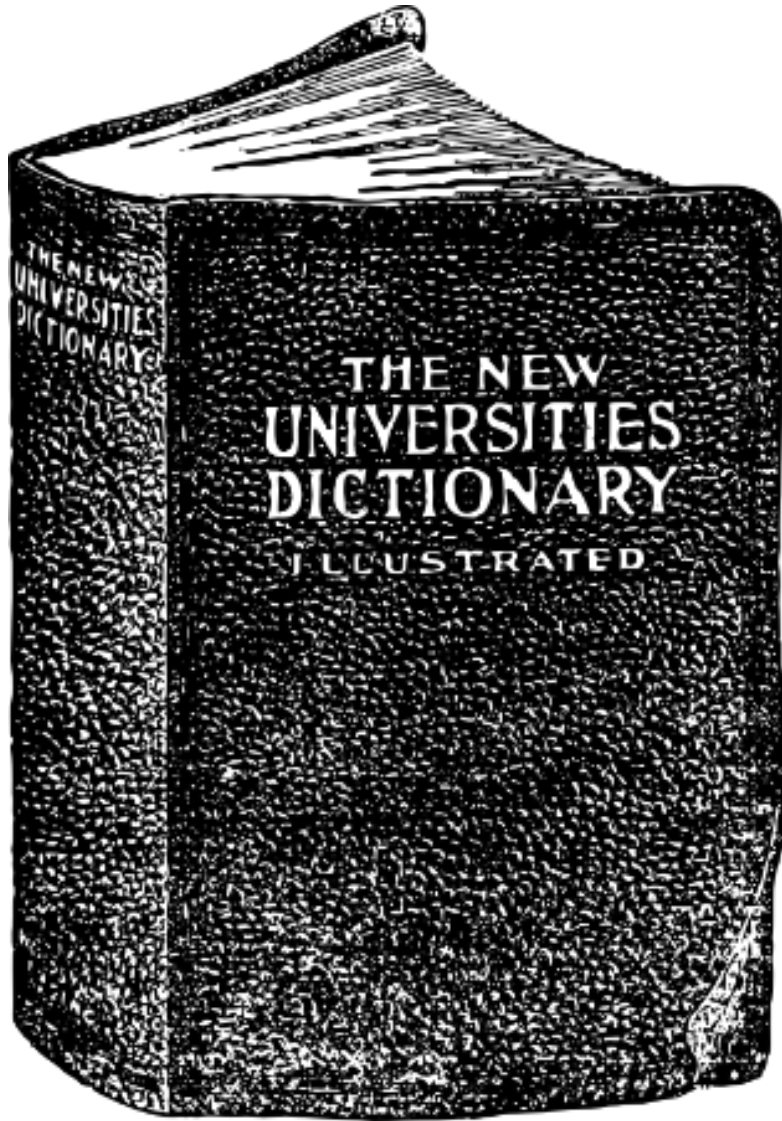


[Albrecht, 2015]

Another Example for an In-Memory Database System



<https://www.stechies.com/overview-sap-hana-database-architecture/>



DBMS with Data Dictionary

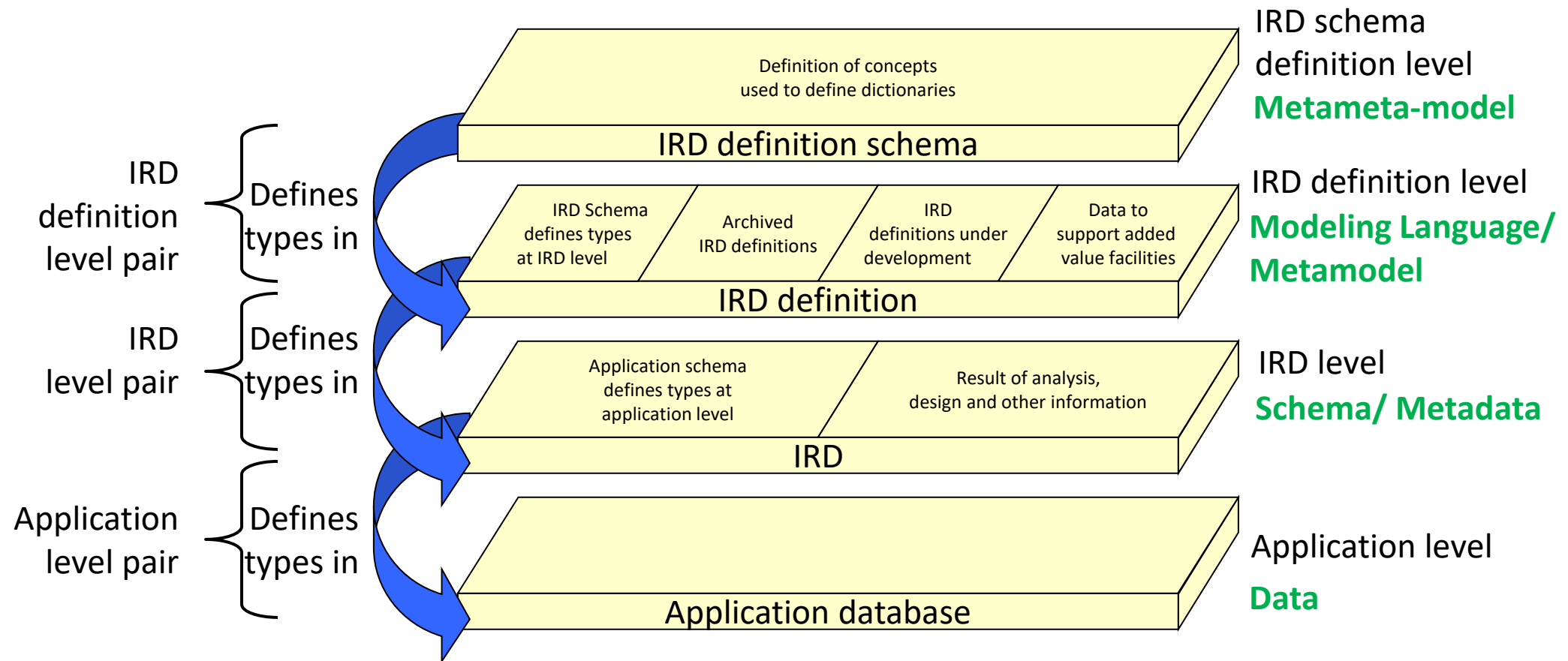
Problem: Each mapping step means loss of semantics

- Realisation of operations requires information about
 - Schemas
 - Integrity constraints
 - Index structures
 - Access authorization
 - ...

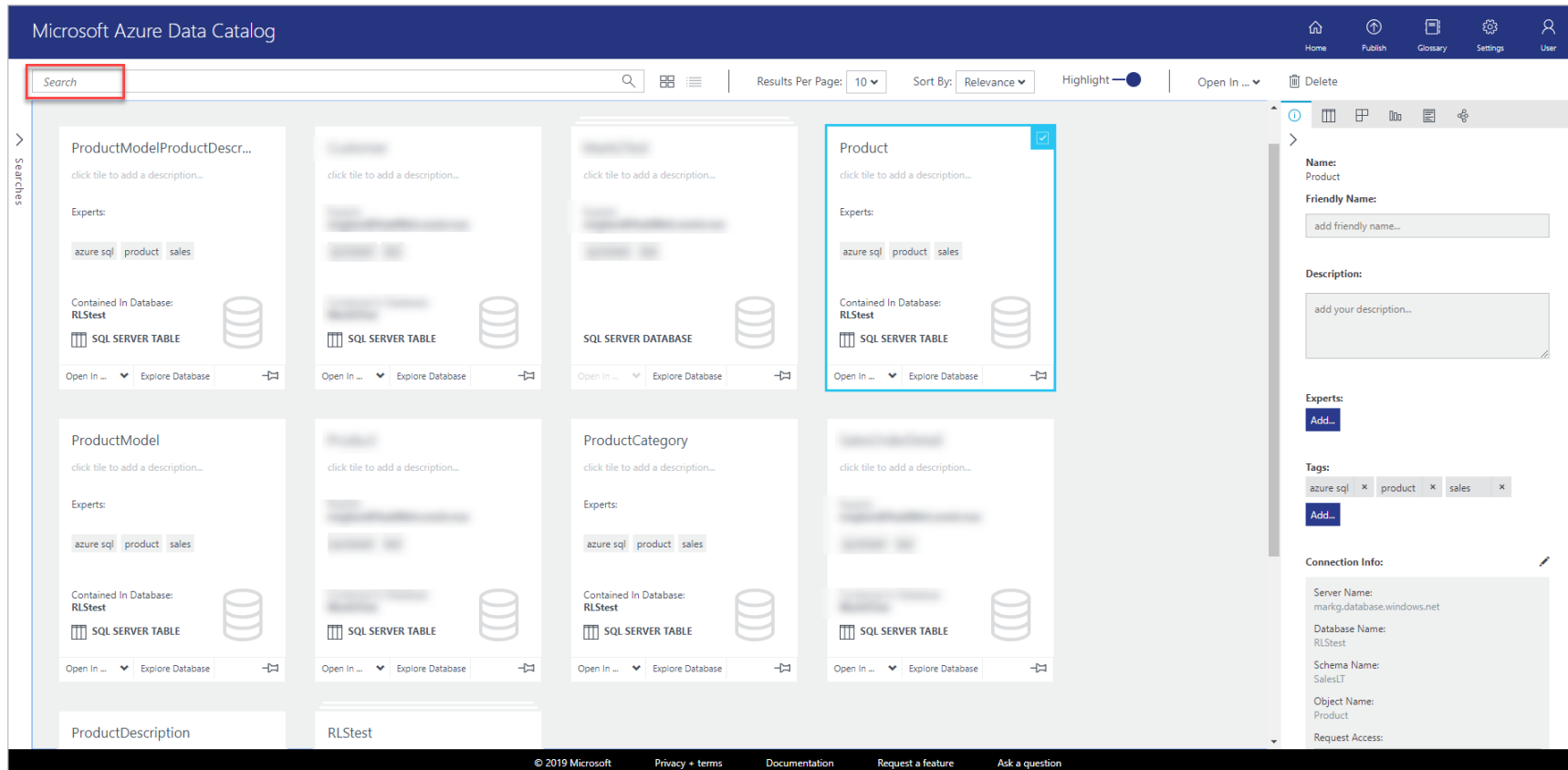
Approach: Comprehensive management by data dictionary

- Internal in DB (uniform model)
- Stand-alone module (fast and specialized service)

IRDS Framework Standard (ISO 10027:1990)



Example: Azure Data Catalog



<https://docs.microsoft.com/de-de/azure/data-catalog/register-data-assets-tutorial>



Quiz

Chapter 1 - Review Questions

- Explain the concept of data independence!
- What are the layers of the five layered DBMS architecture?
- How is a query processed in a DBMS?
- Why is transaction management not assigned to a single layer in the DBMS architecture?
- What are the four levels in the IRDS architecture?
- What is distribution transparency in distributed database systems?
- What are basic characteristics of Big Data?

References & Further Reading

Parts of the slides are based on course material by

- Prof. Dr. Matthias Jarke (Information Systems and Databases, RWTH Aachen University)
- Prof. Dr. Christoph Quix (Wirtschaftsinformatik und Data Science, Hochschule Niederrhein)

Further Reading

[Albrecht, 2015] J. Albrecht: Processing Big Data with SQL on Hadoop. TDWI, 2015.

[Boci & Thistlethwaite, 2015] Boci, E. & Thistlethwaite, S.: A novel big data architecture in support of ADS-B data analytic *Proc. Integrated Communication, Navigation, and Surveillance Conference (ICNS)*, **2015**, C1-1-C1-8

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