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procedure CalculateKey(s)
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[01] return $[\min(g(s), rhs(s)) + h(s_{start}, s) + k_m; \min(g(s), rhs(s))];$

procedure Initialize(

- $\{02^{\circ}\}\ U = \emptyset;$
- $\{03'\}\ k_m = 0;$
- [04'] for all $s \in S \ rhs(s) = g(s) = \infty$;

D* Lite Pseudocode

filmal version)

- $\{05'\}\ rhs(s_{goal}) = 0;$
- [06'] U.Inserf(sgoal, CalculateKey(sgoal));

procedure UpdateVertex(u)

- $\{07^{\circ}\}\ \text{if } (u \neq s_{goal})\ rhs(u) = \min_{s' \in Succ(u)} (c(u,s') + g(s'));$
 - $\{08'\}$ if $(u \in U)$ U.Remove(u);
- [09'] if $(g(u) \neq rhs(u))$ U.Insert(u, CalculateKey(u));

procedure ComputeShortestPath(

- 10'} while (U.TopKey() \leq CalculateKey(s_{start}) OR $rhs(s_{start}) \neq g(s_{start})$)
 - $k_{old} = \text{U.TopKey}();$
 - u = U.Pop():
- if $(k_{old} \leq Calculate Key(u))$ 137
- U.Insert(u, CalculateKey(u));14.
- else if (g(u) > rhs(u))15:
 - g(u) = rhs(u);
- for all $s \in Pred(u)$ UpdateVertex(s);
- $g(u) = \infty$:
- for all $s \in Pred(u) \cup \{u\} \text{ UpdateVertex}(s);$

Orocedure Main()

- 21' $s_{last} = s_{start}$;
 - 22'} Initialize();
- 23'} ComputeShortestPath();
- 24'} while $(s_{start} \neq s_{goal})$ 25'} /* if $(g(s_{start}) = \infty)$ then there is no known path */
- $s_{start} = \arg \min_{s' \in Succ(s_{start})} (c(s_{start}, s') + g(s'))$
 - Move to sstart;
- Scan graph for changed edge costs; 287
 - if any edge costs changed 29.3
- for all directed edges (u, v) with changed edge costs $k_m = k_m + h(s_{last}, s_{start});$ Slast = Satart;
 - Update the edge cost c(u, v);
 - UpdateVertex(u);
- ComputeShortestPath();

If some edge costs have changed:

- Update km
- Update S_{Last}
- update the rhs-values and keys of the vertices potentially affected, as well as their memberships in the priority queue, if they have become locally consistent or inconsistent.
- Recalculate shortest path (line 35)