

# Theoretical Computer Science

## Tutorial Week 2

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# Agenda

- The naive set theory
  - **Basic definition**
  - Operations
  - Properties
  - Power Set and Cardinality
- Formal Languages
  - Alphabets and Strings
  - Languages
  - Operations

# The naive set theory

## Definition

$$A = \{x \in \mathbf{U} \mid P(x)\}$$

$$A = \{a_1, a_2, \dots, a_n\}$$

## Example

$$\{x \in \mathbb{Z} \mid x < 0\}$$

# The naive set theory

The naive set theory	Logic		
The empty set $\emptyset$	False <table><tr><td><math>F</math></td><td>0</td></tr></table>	$F$	0
$F$	0		
The universe <b>U</b>	True <table><tr><td><math>T</math></td><td>1</td></tr></table>	$T$	1
$T$	1		

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# The naive set theory

## Complement

$$A^c = \overline{A} = \{x \in \mathbf{U} \mid x \notin A\}$$

## Example

If  $\mathbf{U} = \{1, 2, 3, 4\}$  and  $A = \{1, 3\}$ , then

$$\overline{A} = \{2, 4\}$$

# The naive set theory

## Union

$$A \cup B = \{x \in \mathbf{U} \mid x \in A \vee x \in B\}$$

## Example

If  $A = \{1, 2, 3\}$  and  $B = \{2, 3, 4\}$ , then

$$A \cup B = \{1, 2, 3, 4\}$$

# The naive set theory

## Intersection

$$A \cap B = \{x \in \mathbf{U} \mid x \in A \& x \in B\}$$

## Example

If  $A = \{1, 2, 3\}$  and  $B = \{2, 3, 4\}$ , then

$$A \cap B = \{2, 3\}$$



# The naive set theory

## Difference

$$A \setminus B = \{x \in \mathbf{U} \mid x \in A \& x \notin B\}$$

## Example

If  $A = \{1, 2, 3\}$  and  $B = \{2, 3, 4\}$ , then

$$A \setminus B = \{1\}$$

# The naive set theory

## Definition

$$X \times Y = \{(x, y) \mid x \in X \& y \in Y\}$$

## Example

If  $A = \{1, 2, 3\}$  and  $B = \{a, b\}$ , then

$$A \times B = \{(1, a), (1, b), (2, a), (2, b), (3, a), (3, b)\}$$

# The naive set theory

## Definition

$$X_1 \times \cdots \times X_n = \{(x_1, \dots, x_n) \mid x_1 \in X_1 \& \dots \& x_n \in X_n\}$$

## Example

$$\underbrace{X \times \cdots \times X}_{n \text{ times}} = X^n$$

$$\mathbb{R}^3$$

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# The naive set theory

$$A \cap \emptyset = \emptyset \quad A \cup \emptyset = A$$

$$A \cap \mathbf{U} = A \quad A \cup \mathbf{U} = \mathbf{U}$$

$$A \cap \overline{A} = \emptyset \quad A \cup \overline{A} = \mathbf{U}$$

$$\overline{\overline{A}} = A$$

# The naive set theory

$A \cap B = B \cap A$ $A \cup B = B \cup A$	Commutativity
$A \cap (B \cap C) = (A \cap B) \cap C$ $A \cup (B \cup C) = (A \cup B) \cup C$	Associativity
$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	Distributivity
$\overline{A \cap B} = \overline{A} \cup \overline{B}$ $\overline{A \cup B} = \overline{A} \cap \overline{B}$	Morgan's laws

# The naive set theory

## Definition

$A \subseteq B$  if, for any  $x$ ,

$$x \in A \rightarrow x \in B$$

## Examples/Properties

For any  $A$ ,

$$\emptyset \subseteq A$$

$$A \subseteq \mathbf{U}$$

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# The naive set theory

## Definition

For a set  $A$ , the power of  $A$  is the set

$$2^A = \mathcal{P}(A) = \{B \mid B \subseteq A\}$$

## Examples

- 1) If  $A = \{a\}$  then  $\mathcal{P}(A) = \{\emptyset, \{a\}\}$
- 2) If  $A = \{a, b\}$  then  $\mathcal{P}(A) = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$

# The naive set theory

## Definition

Intuitively, the cardinality of a set  $A$ , denoted by  $|A|$ , is the number of elements of  $A$ .

## Examples

1.  $|\emptyset| = 0$
2. if  $A = \{2\}$  then  $|A| = 1$
2. if  $A = \{1, 2, 3\}$  then  $|A| = 3$
3.  $|\mathbb{N}| = |\mathbb{Z}| = |\mathbb{Q}| = \omega$
4.  $|\mathbb{R}| = 2^\omega$

# The naive set theory

$$|A \cup B| = |A| + |B| - |A \cap B|$$

$$|A \times B| = |A| \cdot |B|$$

# The naive set theory

## Properties

$$|A \times B| = |A| \cdot |B|$$

## Example

Let  $A = \{a_1, a_2, \dots, a_n\}$  and  $B = \{b_1, b_2, \dots, b_m\}$ .

	$a_1$	$a_2$	$\dots$	$a_n$
$b_1$	$(a_1, b_1)$	$(a_2, b_1)$	$\dots$	$(a_n, b_1)$
$b_2$	$(a_1, b_2)$	$(a_2, b_2)$	$\dots$	$(a_n, b_2)$
$\dots$	$\dots$	$\dots$	$\dots$	$\dots$
$b_m$	$(a_1, b_m)$	$(a_2, b_m)$	$\dots$	$(a_n, b_m)$

Obviously, the table contains  $n \times m$  elements.

# The naive set theory

## Properties

$$|A \times B| = |A| \cdot |B|$$

$$|A_1 \times A_2 \times \cdots \times A_n| = |A_1| \cdot |A_2| \cdot \cdots \cdot |A_n|$$

$$|A^n| = |A|^n$$

# The naive set theory

$$|2^A| = 2^{|A|}$$

$$A = \{a_1, a_2, \dots, a_{n-1}, a_n\}$$

$a_1$	$a_2$	$\dots$	$a_{n-1}$	$a_n$	<i>Subsets</i>
0	0	$\dots$	0	0	$\emptyset$
0	0	$\dots$	0	1	$\{a_n\}$
0	0	$\dots$	1	0	$\{a_{n-1}\}$
0	0	$\dots$	1	1	$\{a_{n-1}, a_n\}$
		$\dots$		$\dots$	
1	1	$\dots$	0	0	$\{a_1, a_2, \dots, a_{n-2}\}$
1	1	$\dots$	0	1	$\{a_1, a_2, \dots, a_{n-2}, a_n\}$
1	1	$\dots$	1	0	$\{a_1, a_2, \dots, a_{n-2}, a_{n-1}\}$
1	1	$\dots$	1	1	$A = \{a_1, a_2, \dots, a_{n-2}, a_{n-1}, a_n\}$

\*0 = " $a_i \notin A$ " & 1 = " $a_i \in A$ "

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## Definition

**Alphabet** is a finite set of symbols

## Examples

$$\{0, 1\}$$
$$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$$
$$\{a, b, c, d, \dots, x, y, z\}$$



## Definition

**String** over an alphabet  $\Sigma$  is a finite sequence of symbols in  $\Sigma$

## Examples

For  $\Sigma = \{0, 1\}$ ,

010011

11100011

## Definition

**String** over an alphabet  $\Sigma$  is a finite sequence of symbols in  $\Sigma$

## Examples

For  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$ ,

123456

666

2022

## Definition

**String** over an alphabet  $\Sigma$  is a finite sequence of symbols in  $\Sigma$

## Examples

For  $\Sigma = \{a, b, c, d, \dots, x, y, z\}$ ,

*peace*

*war*

*dfklgnkjrbgjrbg*

## Definition

**Length** of a string  $s$  is the number of symbols of  $s$   
and denotes as  $|s|$

## Examples

$$|peace| = 5$$

$$|war| = 3$$

$$|dfklgnkjrbgjrbbg| = 15$$

## Definition

$\epsilon$  is the **null** string (empty string) over any alphabet.

## Property

$$|\epsilon| = 0$$

## Definition

For two strings  $x$  and  $y$ , the concatenation  $x \cdot y$  is the operation of joining “end-to-end”.

## Examples

For  $x = 123$  and  $y = 987$ ,

$$x \cdot y = 123987$$

## Definition

For two strings  $x$  and  $y$ , the concatenation  $x \cdot y$  is the operation of joining “end-to-end”.

## Examples

For  $x = \textit{back}$  and  $y = \textit{end}$ ,

$$x \cdot y = \textit{backend}$$

## Property

$$(x \cdot y) \cdot z = x \cdot (y \cdot z)$$

## Examples

For  $x = ab$ ,  $y = cd$  and  $z = ef$ ,

$$(x \cdot y) \cdot z = (abcd) \cdot ef = abcdef$$

$$x \cdot (y \cdot z) = ab \cdot (cdef) = abcdef$$



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## Definition

The set of all strings over  $\Sigma$  is denoted by  $\Sigma^*$

## Examples

For  $\Sigma = \{0, 1\}$ ,

$$\Sigma^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, \dots\}$$

## Definition

A language  $L$  is a set of strings over an alphabet  $\Sigma$ .

## Equivalent definition

$$L \subseteq \Sigma^*$$

## Alphabet

For  $\Sigma = \{0, 1\}$ ,

$$\Sigma^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, \dots\}$$

## Languages

$$\begin{aligned} L_1 &= \{00000000, 00000001, \dots, 11111110, 11111111\} = \\ &= \{x \in \{0, 1\}^* \mid |x| = 8\} \end{aligned}$$

$$L_2 = \{0, 00, 01, 000, 001, 010, \dots\} = \{0 \cdot x \mid x \in \Sigma^*\}$$

## Alphabet

For  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$ ,

$$\Sigma^* = \mathbb{N} \cup \{\epsilon\}$$

## Languages

$$L_1 = \{0, 2, 4, 6, 8, 10, \dots\} = \{x \in \Sigma^* \mid x \text{ is even} \}$$

$$L_2 = \{2, 3, 5, 7, 13, \dots\} = \{x \in \Sigma^* \mid x \text{ is prime} \}$$

## Alphabet

For  $\Sigma = \{a, b, c, d, \dots, x, y, z\}$

## Languages

English, Italian, French,...

## Alphabet

For  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, +, -, =\}$

## Arithmetic

$\{0 + 0 = 0, 0 - 0 = 0, \dots, 12 + 32 = 44, \dots, 52 - 39 = 13, \dots\}$

## Alphabet

For  $\Sigma = \{A, B, C, \dots, ^c, \cup, \cap, \setminus, (, )\}$

## Formulas of the naive set theory

$$A^c, A \cup B, \dots, (A \cup B^c) \cap (A^c \cup B), \dots$$



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## Operations

$$\overline{L} = L^c$$

$$L_1 \cup L_2$$

$$L_1 \cap L_2$$

$$L_1 \setminus L_2$$

## Concatenation

$$L_1 \cdot L_2 = \{x \cdot y \mid x \in L_1 \text{ \& } y \in L_2\}$$

## Example

If  $L_1 = \{1, 2, 3\}$  and  $L_2 = \{a, b\}$ , then

$$L_1 \cdot L_2 = \{1a, 1b, 2a, 2b, 3a, 3b\}$$

## Concatenation

$$L_1 \cdot L_2 = \{x \cdot y \mid x \in L_1 \text{ \& } y \in L_2\}$$

## Example

If  $L_1 = \{1, 12\}$  and  $L_2 = \{\epsilon, 2\}$ , then

$$L_1 \cdot L_2 = \{1, 12, 122\}$$

Thank you for your attention!