

Refresher on
Algorithms

Keerat Kaur Guliani
ML Operations, TUSK

✕ Topics to cover today...

- + Divide and Conquer
- + Dynamic Programming
- + Greedy Approach

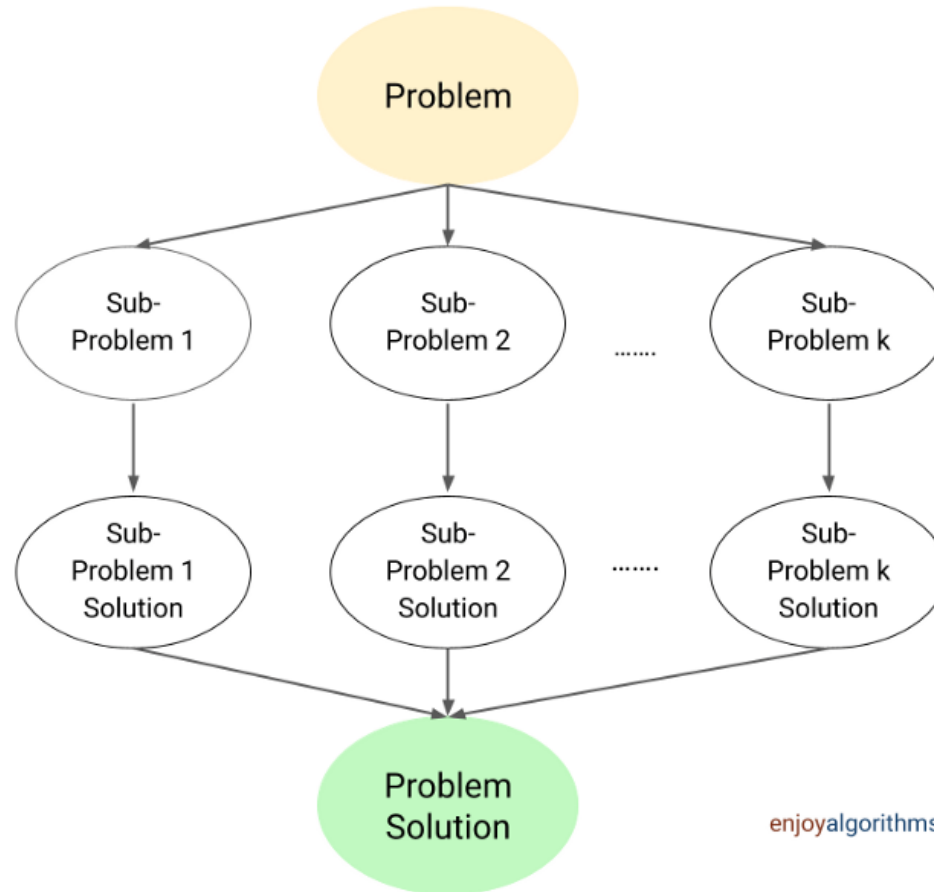


Divide and Conquer

Divide
Dividing the problem into
smaller sub-problems

Conquer
Solving each
sub-problems recursively

Combine
Combining sub-problem
solutions to build the original
problem solution



Divide and Conquer - Types

- + DAC using one subproblem – Decrease and Conquer

Eg: Binary Search

- + DAC using two subproblems

Eg: Merge Sort

Illustration: Maximum Subarray Sum

- + Given an integer array, find the subarray with the largest sum/return its sum.

```
Input: nums = [-2,1,-3,4,-1,2,1,-5,4]
```

```
Output: 6
```

```
Explanation: The subarray [4,-1,2,1] has the largest sum 6.
```

```
Input: nums = [1]
```

```
Output: 1
```

```
Explanation: The subarray [1] has the largest sum 1.
```

- + Not the same as maximum subsequence sum!

Maximum Subarray



Maximum sum
= $3+5+1+7+9$
= 25

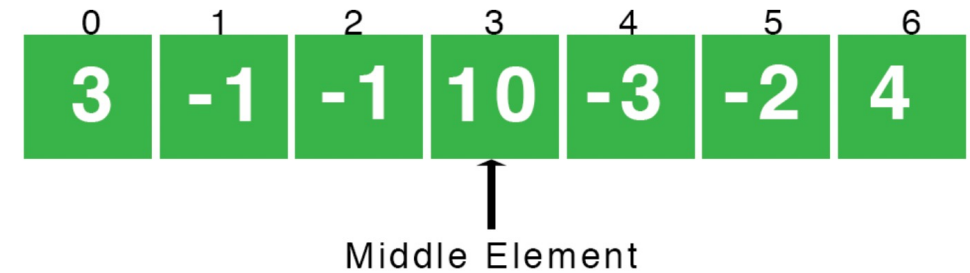
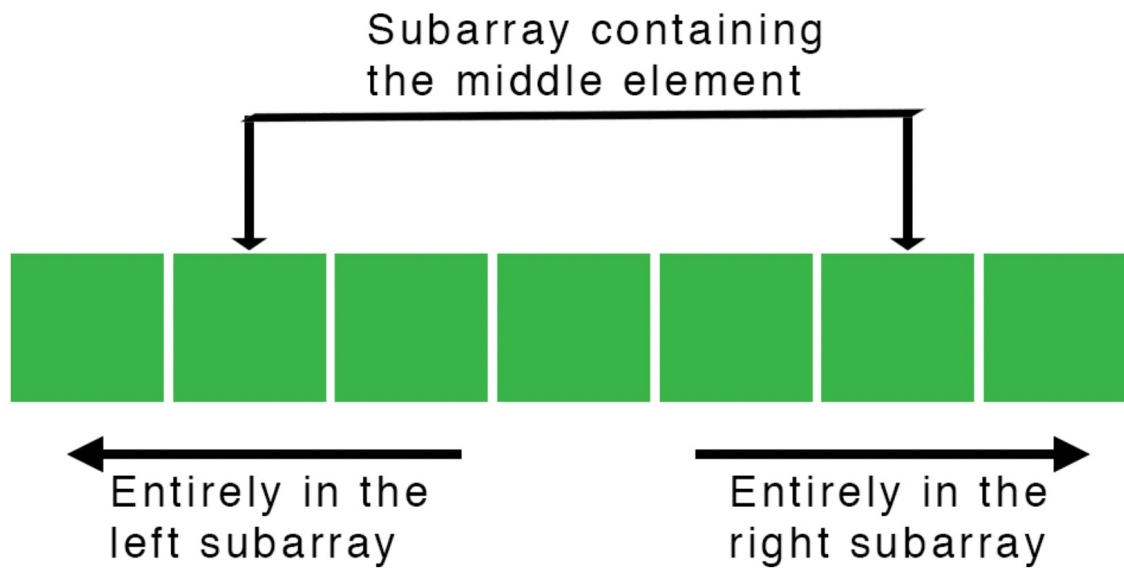


Maximum sum
= $3-1-1+10$
= 11

Naïve Method: Consider every possible subarray using 2 loops and return the overall maximum.

Time complexity: $O(n^2)$

Maximum Subarray with DAC



Left subarray

sum = 0

Start from index 3:

$10 > 0 \Rightarrow \text{sum} = 10$

$10 - 1 = 9 < \text{sum}(10)$

$9 - 1 = 8 < \text{sum}(10)$

$8 + 3 = 11 > \text{sum} \Rightarrow \text{sum} = 11$

Right subarray

sum = 0

Start from index 4:

$-3 < \text{sum}(0)$, sum is 0

$-3 - 2 = -5 < \text{sum}(0)$, sum is 0

$-5 + 4 = -1 < \text{sum}(0)$, sum is 0

Final crossing sum = $11 + 0 = 0$ from index 0 to 3



Code

$$T(n) = 2T(n/2) + \Theta(n)$$

Time: $O(n \log n)$

Space: $O(1)$

Dynamic Programming

What is DP?

- + Useful technique for optimization problems
- + Looks through all possible sub-problems and never recomputes the solution to any sub-problem
- + Guarantees correctness and efficiency

When can DP be applied?

- + Breaking down an optimization problem into simpler sub-problems (*optimal substructure*)
- + Storing the solution to each sub-problem so that each sub-problem is solved only once (*memoization*)

What exactly is a sub-problem?

- + Smaller version of the original problem
- + Sub-problems build on each other to obtain the solution to the original problem

Illustration 1: Longest Common Subsequence (LCS)

- + Given two strings s_1 and s_2 return the length of the longest common subsequence of characters between the two strings (need not be contiguous).

Input:

$s_1 = \text{"ADC"}$
 $s_2 = \text{"ABCD"}$

Output: 2

Explanation:

"ADC"
 "ABCD"

Both strings share the subsequence "A", "D".

Input:

$s_1 = \text{"DBC"}$
 $s_2 = \text{"CBD"}$

Output: 1

Explanation:

"DBC"
 "CBD"

or

"DBC"
 "CBD"

or

"DBC"
 "CBD"

Input:

$s_1 = \text{"ABCD"}$
 $s_2 = \text{"ABCD"}$

Output: 4

Explanation:

"ABCD"
 "ABCD"

Both strings share the subsequence "A", "B", "C", "D".

For example,

$$\begin{aligned} \text{lcs}(\text{'azb'}, \text{'aab'}) &= 1 + \text{lcs}(\text{'az'}, \text{'aa'}) = 1 + \max(\text{lcs}(\text{'az'}, \text{'a'}), \text{lcs}(\text{'a'}, \text{'aa'})) \\ &= 1 + \max(\max(\text{lcs}(\text{'az'}, \text{' '}), \text{lcs}(\text{'a'}, \text{'a'})), 1 + \text{lcs}(\text{' '}, \text{'a'})) \\ &= 1 + \max(1, 1) = 2 \end{aligned}$$

Subproblem decomposition using the DP Table

Find the LCS between AGGTAB and GXTXAYB.

		""	A	G	G2	T	A3	B
" "	0	0	0	0	0	0	0	0
G	0	0	1	1	1	1	1	
X	0	0	1	1	1	1	1	
T	0	0	1	1	2	2	2	
X	0	0	1	1	2	2	2	
A	0	1	1	1	2	3	3	
Y	0	1	1	1	2	3	3	
B	0	1	1	1	2	3	4	



Code

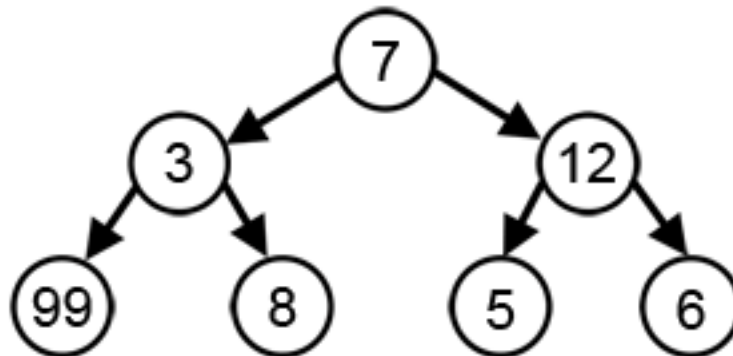
Time: $O(mn)$

Space: $O(mn)$

Where $m = \text{len}(s1)$, $n = \text{len}(s2)$

× Greedy Algorithm

- + An optimization algorithm that prioritizes the optimal choice at each step as it attempts to find the overall optimum to solve the entire problem.
- + An example where it works: Shortest path through a graph
- + An example where it does not work: To reach the largest sum in the graph



Greedy Algorithm

+ *When can a greedy algorithm be used?*

Greedy choice property: A global (overall) optimal solution can be reached by choosing the optimal choice at each step.

Optimal substructure: A problem has an optimal substructure if an optimal solution to the entire problem contains the optimal solutions to the sub-problems.

+ *Limitations of Greedy Algorithms:*

- May not be able to find the global optimum because only limited data is considered.
- Decision at each time step is made based on choices made so far, which may not include all possible future choices.

Illustration: Activity Selection Problem

- + Given N activities with their start time and end time. The task is to find the solution set having a maximum number of non-conflicting activities that can be executed within the given time, assuming only a single activity can be performed at a given time.

Input: start[] = [10, 12, 20]

end[] = [20, 25, 30]

Output: [0, 2]

Explanation: A maximum of two activities can be performed, i.e. Activity 0 and Activity 2[0-based indexing].

Input: start[] = [1, 3, 0, 5, 8, 5]

finish[] = [2, 4, 6, 7, 9, 9]

Output: [0, 1, 3, 4]

Explanation: A maximum of four activities can be performed, i.e. Activity 0, Activity 1, Activity 3, and Activity 4[0-based indexing].

The Greedy Choice

Choose activities with the earliest finish time. The earlier you finish, the more time you have for other activities!

+ Steps:

- Sort the activities according to their finishing time
- Select the first activity from the sorted array and print it
- Do the following for the remaining activities in the sorted array
 - + If the start time of this activity is greater than or equal to the finish time of the previously selected activity then select this activity and print it



Code

Time: $O(n \log n)$

Space: $O(1)$