Sorting

Starting 9:05

TABLE OF CONTENTS

Revision Quizzes

1. Stacks / Heaps

2.

O/p - 50.

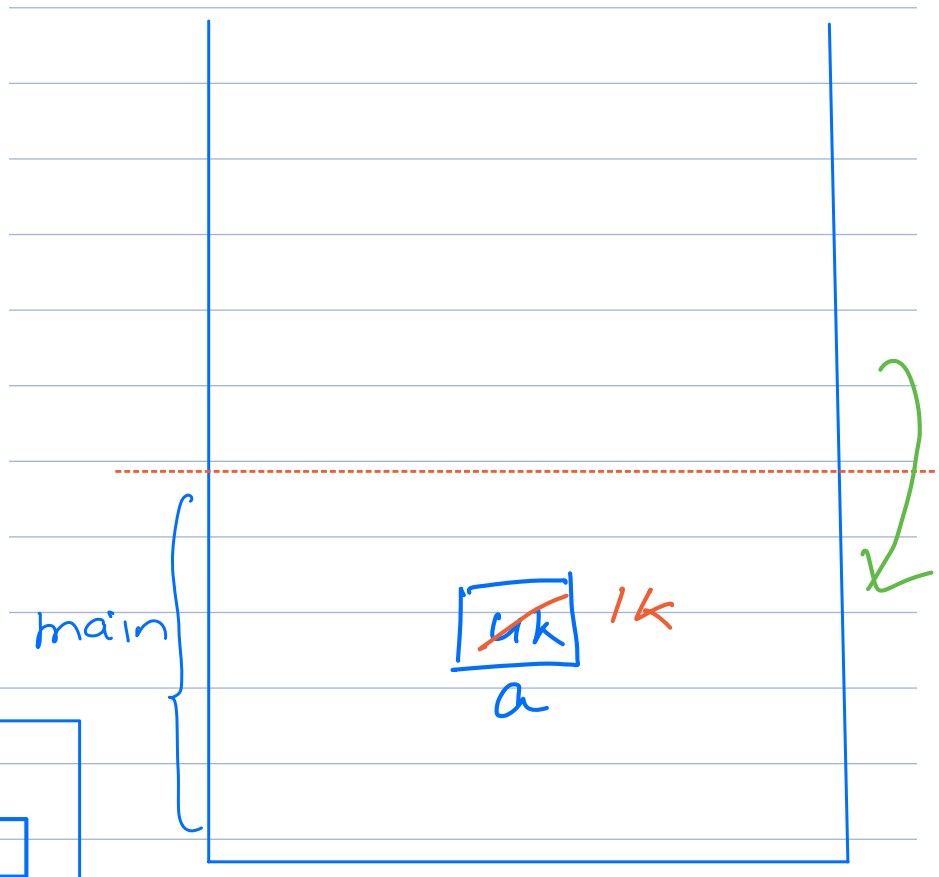
Predict the output :

Java

```java
static int[] fun(int[]a) {
    a = new int[2];
    a[0] = 50; a[1] = 60;
    return a;
}

public static void main(String args[]) {
    int[]a = {10,20,30};
    a = fun(a);
    System.out.println(a[0]);
}
```

Python

```python
def fun(a):
    a = [0, 0]
    a[0] = 50
    a[1] = 60
    return a

def main():
    a = [10, 20, 30]
    a = fun(a)
    print(a[0])

if __name__ == "__main__":
    main()
```



Sorting : An arrangement of data in a particular order based on some parameter.

Ex:- 1. $\{2, 3, 9, 12, 17, 19\}$

Parameter:- Value ↑

2. $\{19, 6, 5, 2, -1, -19\}$

Parameter:- Value ↓

Quiz 1

	0	1	2	3	4
arr :-	1	13	9	6	12
	1	2	3	4	6

Parameter:- Count of factors ↑

Note:- In today's class, we will be using in-built libraries to sort data.

Java :-

Arrays.sort(arr) # Static arrays

Collections.sort(arr) # ArrayList

C++ :- `sort(arr, arr + N);` # Static arrays

`sort(arr.begin(), arr.end());` # Vectors

Python :- `arr.sort();`

T.C. → O(N log N)

S.C → O(N)

Problem 1

Given N elements, at every step remove an array element.

Cost to remove an element = Sum of ~~array of~~ all the elements present in an array

Find minimum cost to remove all elements.

ex :- {2, 1, 4} (0, 1, 2)

Ways to removal

a) 1 → 1+2+4 = 7
2 → 2+4 = 6
4 → 4
17

b) 2 → 2+4+1 = 7
4 → 4+1 = 5
1 → 1
13

c) 4 → 1+2+4 = 7
2 → 1+2 = 3
1 → 1
11

↓ ans.

+ 3 more ;

1, 4, 2 → 15
2, 1, 4 → 16
4, 1, 2 → 12

Quiz 2:-

arr :- { 4 , 6 , 1 }
0 1 2

a) 1 — 1+4+6 = 11
4 — 4+6 = 10
6 — 6
27

b) 1 — 1+4+6 = 11
6 — 6+4 = 10
4 — 4
25

c) 4 — 1+4+6 = 11
1 — 1+6 = 7
6 — 6
24

a) 4 — 1+4+6 = 11
6 — 6+1 = 7
1 — 1
19

b) 6 — 1+4+6 = 11
1 — 1+4 = 5
4 — 4
20

c) 6 — 1+4+6 = 11
4 — 1+4 = 5
1 — 1
17

Generalising :-

ex:- $\{a, b, c, d\}$ (0, 1, 2, 3)

Let's remove in the order

	Cost
a	a + b + c + d
b	b + c + d
c	c + d
d	d

$a + 2b + 3c + 4d$

$a > b > c > d$ to minimise

Solution Approach

Sort the array in decreasing order.

Iterate through the array & accumulate the value of $arr[i] * (i+1)$ in the global sum.

Quiz 3

arr :- {3, 5, 1, -3}

↓ reverse - sort()

0 1 2 3

{5, 3, 1, -3}

ans += a[i] * (i+1)

i

ans = 5 + 6 + 3 + (-12) = (2)

$2^{10} = 1024$

10^9

4 a b .

Code

```
int findMinimumCost(int[] arr, int N) {

    NlogN { reverse_sort(arr); } N

    int ans = 0;

    N { for (i=0; i<N; i++) {
            ans += arr[i] * (i+1);
        }

    } return ans;
```

T.C → O(NlogN).
S.C → O(N).

Problem 2

(distinct elements)

Given N array elements, calculate number of noble integers.

An element element in arr [ ] is said to be noble if { count of smaller elements = element itself }

ex:- {1, -5, 3, 5, -10, 4} (0 1 2 3 4 5)

ans = 3.

Q-4. {-3, 0, 2, 5} (0 1 2 3) ans = 1.

Idea 1:- For every element, Iterate over the entire array to find out the count of elements smaller than the curr elem.

$$T.C \rightarrow O(N^2)$$

Idea 2:- Sort the array in increasing order.

Iterate through the array & increment ans if $(arr[i] == i)$

ex:- $\{1, -5, 3, 5, -10, 4\}$ (indices 0 1 2 3 4 5)

Sorted arr :- $\{-10, -5, 1, 3, 4, 5\}$ (indices 0 1 2 3 4 5).

Code

```
N lgN { Sort (arr);   int ans = 0;
N     { for (i = 0; i < N; i++) {
3       | if (arr[i] == i) {
1       |     ans++;
3       }
3     }
      return ans;
```

$T.C \rightarrow O(N \lg N)$

$S.C \rightarrow O(N)$.

Variation

What if the elements are not distinct ?

Q:- 5

arr :- {-10, 1, 1, 3, 100} (0 1 2 3 4)

ans = 3.

Q:- 6

arr :- {-10, 1, 1, 2, 4, 4, 4, 8, 10} (0 1 2 3 4 5 6 7 8)

ans = 5.

Q:- 7

arr :- {-3, 0, 2, 2, 5, 5, 5, 5, 8, 8, 10, 10, 10, 14} (0 1 2 3 4 5 6 7 8 9 10 11 12 13)

ans = 7.

Approach:- Sort the array & iterate on above array. While iterating

a) The curr. element is same as previous element :- The previous element is going to determine whether a[i] is noble or not.

b) The curr element is diffent to previous element :- The index i of that element will determine whether a[i] is noble or not

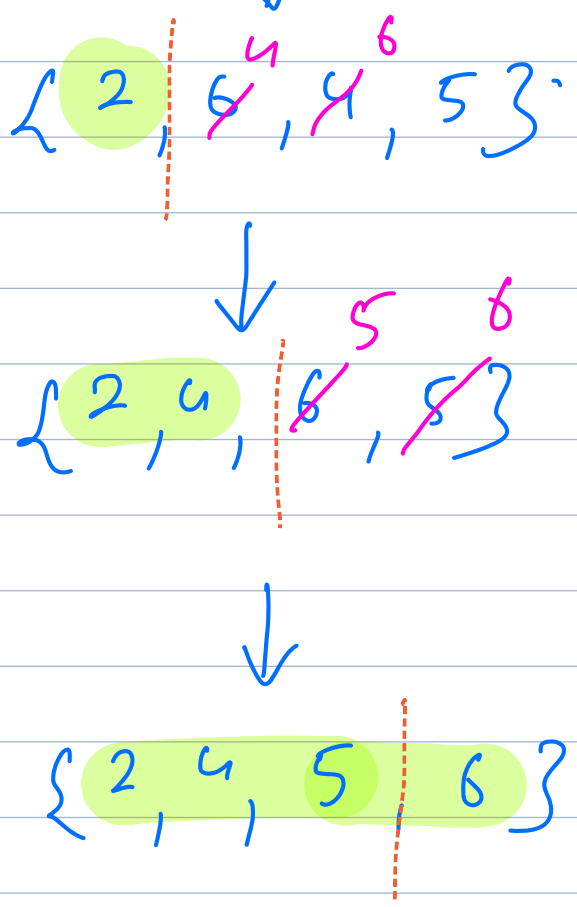
Code → H.W. 10:25

Selection Sort

3
2.
1

Arranging Students Height-wise to form an assembly line.

ex:- {5, 6, 4, 2, 3}
0 1 2 3
2 5



Code

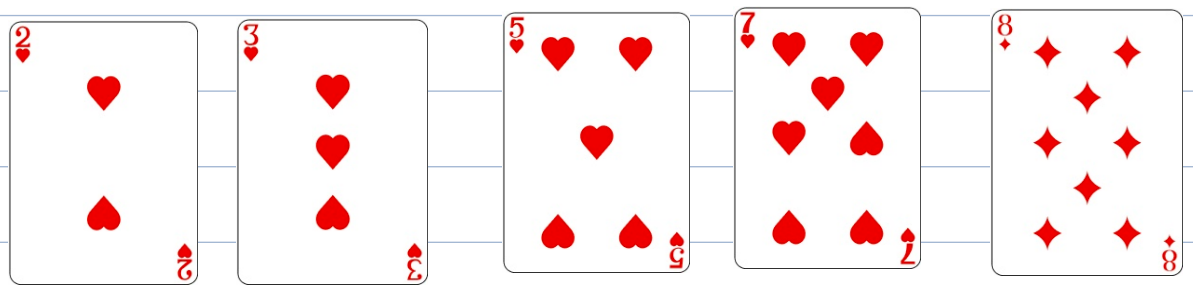
```
void selectionSort(int arr[], int N){
    for(i=0; i<N-1; i++){
        mnIndex = i;
        for(j=i+1; j<N; j++){
            if(arr[j] < arr[mnIndex]){
                mnIndex = j;
            }
        }
        swap(arr[i], arr[mnIndex]);
    }
}
```

arr : {5, 6, 4, 2, 3, 6}

T.C → $O(N^2)$

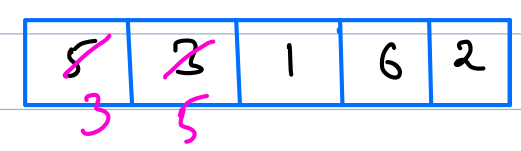
S.C → $O(1)$

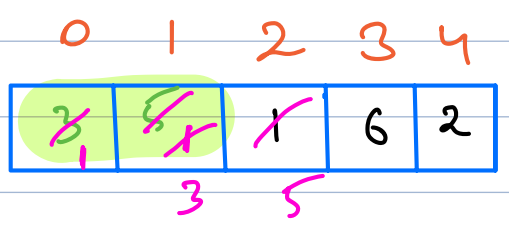
Insertion Sort

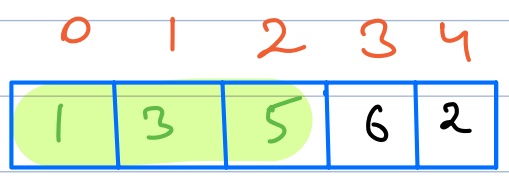


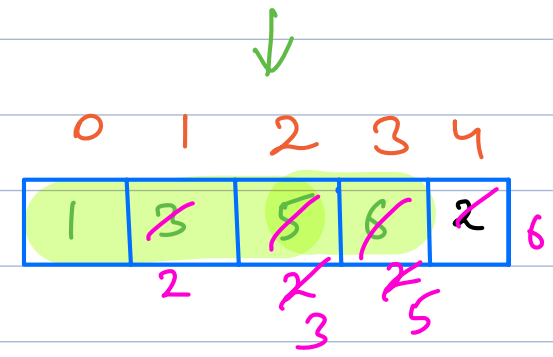
0 1 2 3 4

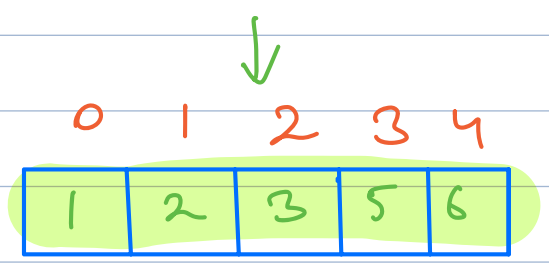
ex:-











Code

```
for (i=1 ; i< N; i++){
    for (j = i-1 ; j >= 0; j--){
        if (arr[j] > arr[j+1]){
            swap(arr[j], arr[j+1]);
        }
        else { break; }.
    }
```

3

ex:- { 1, 2 3, 5, 6, 2 3. }

0 1 2 3 4 — i

j

↓

{ 1, 2, 3, 5, 6 3. }

T.C → $O(N^2)$

ex:- { 5, 4, 3, 2, 1 }

i

ex:- { 1, 2, 3, 4, 5 }.

N-1

T.C → O(N) ⟶ Best Case.