

Agenda

Prob 1 : Fast Power

Prob 2 : Print Array

Prob 3 : Indices Of an Array

Prob 4 : Check Palindrome

Revision

Quiz 1 : Recursion : A function that solves a problem by breaking it into smaller subproblems by calling itself

Quiz 2 : Data structure for fn call tracing : Stack



Quiz 3 : Base case for calculating factorial :

if ($N = 0$) return 1

Quiz 4 : TC and SC for factorial (recursion)

\downarrow
 N
 $N-1$
 \downarrow
 $N-2$
 \downarrow
 \dots
 0

TC: $O(N)$

SC: $O(N)$

Quiz 5: Fib(N) : Fib(N-1) + Fib(N-2)

1. Given two integers a and n , find a^n using recursion.

$$a = 2, n = 3$$

$$\text{ans} = 2^3 = 8$$

$$\downarrow \\ 2 \times 2 \times 2$$

$$a^n = a \times a \times a \dots \times a$$

$\underbrace{\quad\quad\quad}_{n \text{ times}}$

BF : Multiply ' a ' n times

TC: O(N)
SC: O(1)

Recursion

$$2^9 = \underbrace{2 \times 2 \times \dots \times 2}_{3^9 \text{ times}} \times 2$$

$$2^9 = 2^8 \times 2$$

$$a^n = a^{n-1} \times a$$

Base Case : If ($n == 0$)

return 1

// Given a and n , return a^n

int pow(int a, int n) {

If ($n == 0$)

return 1

return pow(a, n-1) \times a

→ 125

$a^n = a^0 \times a$

$\text{pow}(5, 3)$

$\downarrow \uparrow 2^5$

$\text{pow}(5, 2) \times 5$

$\downarrow \uparrow 2^5$

$\text{pow}(5, 1) \times 5$

$\downarrow \uparrow 5$

$\text{pow}(5, 0) \times 5$

1

// Assumption
 $f(n) <$

// Base case

// main logic

$\text{pow}(a, n)$

$a, n-1$

$a, n-2$

⋮

$a, 0$

$n+1$

TC: O(N)

SC: O(N)

Fn call stack

Optimized Approach

$$a^n = a^{n/2} \times a^{n/2}$$

$$7^6 = 7^3 \times 7^3$$

$$a^n \xrightarrow{n \text{ is even}} a^{n/2} \times a^{n/2}$$

$$7^6 = 7^3 \times 7^3$$

$$a^n \xrightarrow{n \text{ is odd}} a^{n/2} \times a^{n/2} \times a$$

$$2^{10} = 2^5 \times 2^5$$

$$2^{11} = 2^5 \times 2^5 \times 2$$

Assumption: Given a and n , return a^n

Main Logic: If n is even

$$\text{pow}(a, n) = \text{pow}(a, n/2) \times \text{pow}(a, n/2)$$

else

$$\text{pow}(a, n) = \text{pow}(a, n/2) \times \text{pow}(a, n/2) \times a$$

Base Case: If $(n == 0)$

return 1

$$2^{10} \\ \downarrow \quad \downarrow \\ 2^5 \times 2^5$$

$$2 \times 2^2 \times 2^2 \\ \swarrow \quad \swarrow$$

$$2^1 \times 2^1 \\ \swarrow \quad \swarrow \\ 2 \times 2^0 \times 2^0$$

$$2^{10} = 2^5 \times 2^5 \\ 32$$

$$2^{10} = 32 \times 32$$

// Given a and n, return a^n

```
int pow(int a, int n) {
```

If ($n == 0$)

return 1

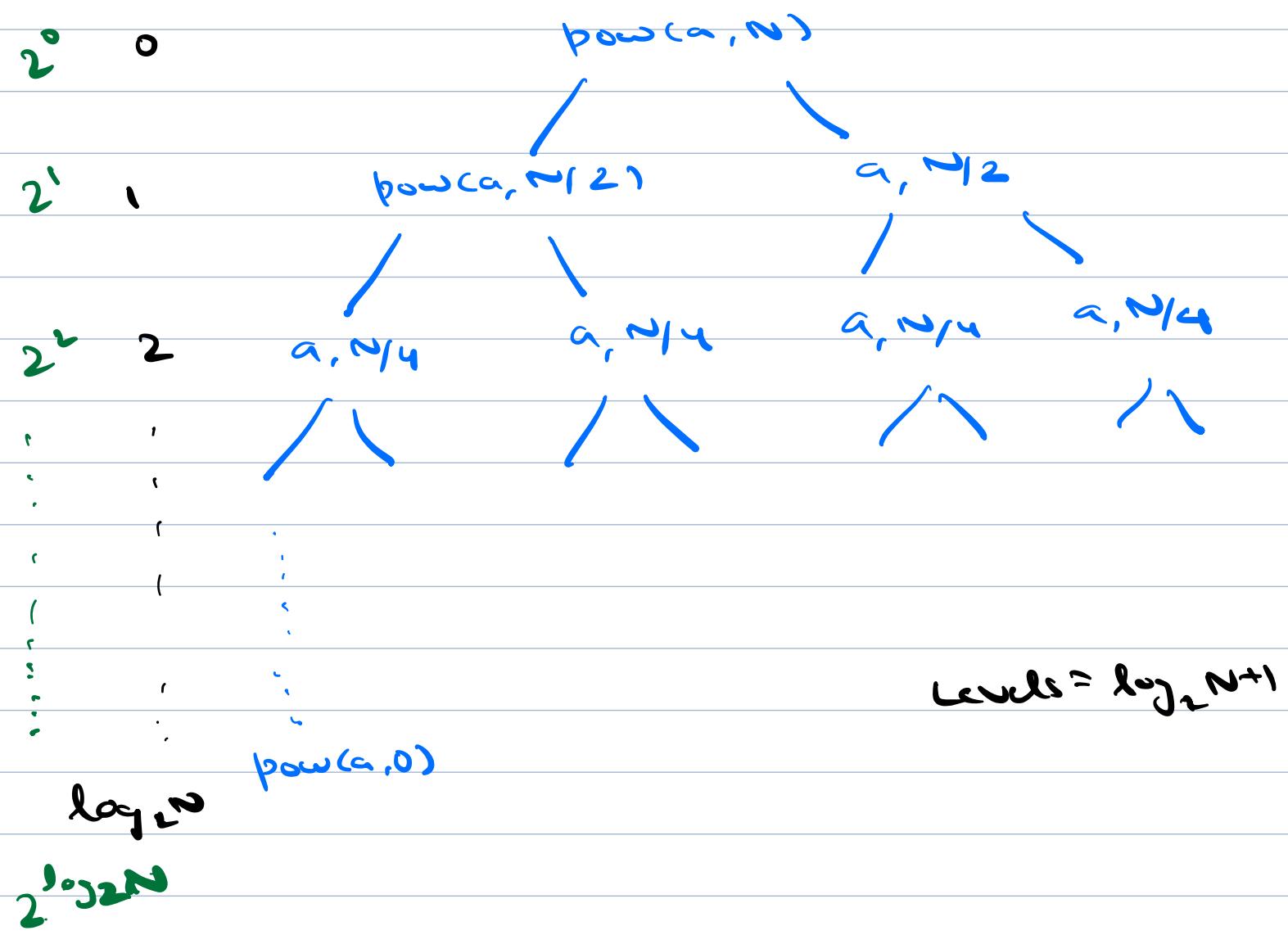
If ($n / 2 == 0$)

return pow(a, N/2) * pow(a, N/2)

else

return pow(a, N/2) * pow(a, N/2) * a

T



Total fn calls = $2^0 + 2^1 + \dots + 2^{\log_2 N}$

$$G_P = \frac{a(x^n - 1)}{x - 1}$$

$$= \frac{1(2^{\log_2 N+1} - 1)}{2 - 1}$$

$$= 2(2^{\log_2 N}) - 1$$

$$= 2N - 1$$

TC: $O(N)$

SC: $O(\log_2 N)$

$$\boxed{a^{\log_a b} = b}$$

Fast Power or Fast Exponentiation

// Given a and n, return a^n

```
int pow(int a, int n) {
```

If ($n == 0$)

return 1

↑ 8192

pow(2, 13)

int p = pow(a, n/2)

64 × 64 × 2 ↓ ↑ 8192

if ($n \% 2 == 0$)

return p × p

pow(2, 6)

else

return p × p × a

× × × ↓ ↑ 64

pow(2, 3)

2 × 2 × 2 ↓ ↑ 8

pow(2, 1)

1 × 1 × 2 ↓ ↑ 2

pow(2, 0)

pow (a, n)



pow ($a, n/2$)

TC : $O(\log_2 N)$

SC : $O(\log_2 N)$

↓
pow ($a, n/4$)



⋮

pow ($a, 0$)

2. Given an array of integers, write a recursive function to print all array elements.

A = [1, 2, 3, 4, 5]

O/P : 1 2 3 4 5

→ 1 2 3 → 5

parArray (arr, 0, N-1) → 0 N-2 N-1

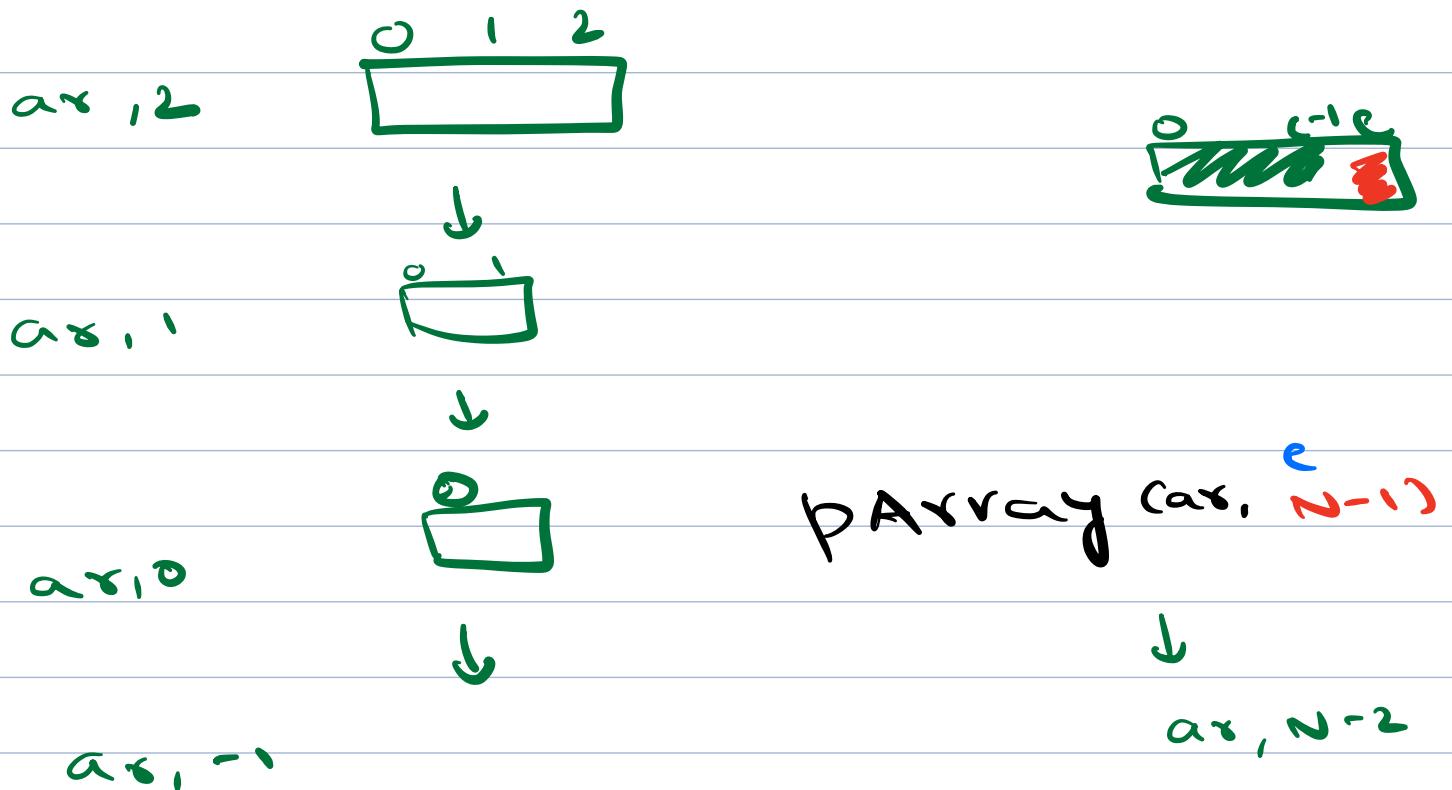
{
parArray (arr, 0, N-2)
cout << arr[N-1]}

// Given an arr and int id, print \rightarrow from arr to id

```
void pArray (int arr[], int c) {
```

```
    if (c == -1) // or c < 0  
        return
```

```
pArray (arr, c-1)  
print (arr[c])
```



TC : O(N)

SC : O(N)

arr[-1]

$\text{pArray}(\text{arr}, 0, N-1) = \boxed{\text{print}(\text{arr}[0])}$
 $\text{pArray}(\text{arr}, 1, N-1)$

// Given an arr and start id, print arr from
 $s \rightarrow N-1$ id

void pArray (int arr[], int s) <

|
if (s == arr.size())
 return

print (arr[s])
pArray (arr, s+1)

}



$\text{pArray}(\text{arr}, 0)$

2. Given an array of integers, write a recursive function to print all array elements.

$A = [1, 2, 3, 4, 5]$

Fn to print mat of array

$\text{matArray}(\text{arr}, 0, N-1) = \text{mat}(\text{matArray}(\text{arr}, 0, N-2),$
 $\text{arr}[N-1])$

minⁿ⁻¹

// Given an arr and end id, find min from
base case id.

int minArray (int arr[], int c) <

if (c == -1) // or c < 0
return INT_MIN

return min(minArray (arr, c-1), arr[c])

min (0 → 0)

min Array (arr, 0) = min(minArray (arr, -1), arr[0])

Base Case : if c == 0

return arr[0]

Handle the case of array size = 0
properly

10:42

To find sum of array using recursion,
what recursive state is to be used?

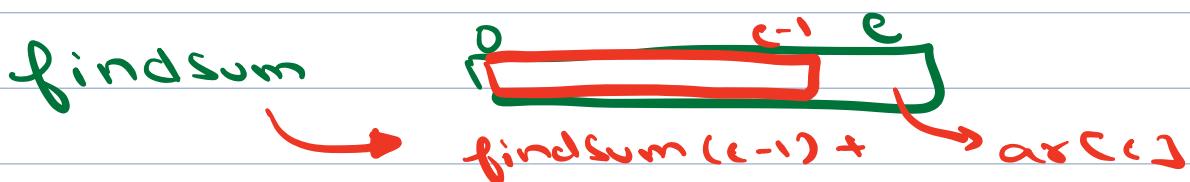
Hint : $\text{max}(\text{arr}[e])$, $\text{maxArray}(\text{arr}, e-1)$

✓ a $\text{arr}[e] + \text{findsum}(\text{arr}, e-1)$

b $\text{max}(\text{arr}[e]), \text{findsum}(\text{arr}, e-1)$

c $\text{arr}[e] + \text{max}(\text{arr}[e]), \text{findsum}(\text{arr}, e-1)$

d $\text{arr}[e] - \text{findsum}(\text{arr}, e-1)$



3. Given an array A of N integers and target B ,
find all indices at which B occurs in
array.

$$A = [4, 5, 3, 1, 5, 4, 5] \quad B = 5$$

ans: [1, 4, 6]

$$A = [1, 2, 3, 1, 1] \quad B = 1$$

ans: [0, 3, 4]

① ArrayList → GLOBAL var

② ArrayList → Input in function



getIndices (arr, 0, N-1, B) =

getIndices (arr, 0, N-2, B)

if (arr[N-1] == B)

ans.add(N-1)

// Given arr, B, and e idk, add all indices
from 0 → e where B is present in arr

void getIndices (int arr[], int e, int B, ^{list<int>} ans) {

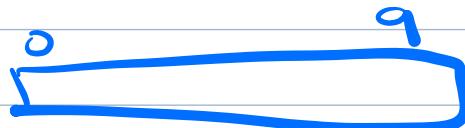
if (e == -1)
return

getIndices (arr, e-1, B, ans)

if (arr[e] == B) <
ans.add(e)

>

N=10



int main() {

list<int> ans = new ArrayList();

getIndices(ar, ar.size() - 1, B, ans)

TC: O(N)

SC: O(N)

↓
stack

A = [0 1 2 3 4,
1, 2, 3, 1, 1]

B = 1

N = 5

getIndices(ar, 4, 1, ans) C
↓ ↑ e B ans
[0, 3, 4]

getIndices(ar, 3, 1, ans) [0, 3]

↓ ↑

getIndices(ar, 2, 1, ans) [0]

↓ ↑

getIndices(ar, 1, 1, ans) C

↓ ↑

getIndices(ar, 0, 1, ans) ans
[0]

↓ ↑

getIndices (arr, -1, l, ans)

4. Given a string, write a recursive function to check if it is palindrome.

str = "radar" ans : true

str = "random" ans : false



0 → 6 str[0] == str[6]
↓
l → 5

// Given a str, check if its palindrome from s to e

bool isPalindrome (str, s, e)

```
if (s > e) {
    return true
}
if (str[s] != str[e])
    return false
```

return isPalindrome (str, s+1, e-1)

isPalindrome (str, 0, N-1)



$\boxed{0 \dots N-1}$

isPal (0, N-1)

$N/2$ for calls

TC : $O(N)$

SC : $O(N)$

↓
isPal (1, N-2)

↓

2, N-3

↓

⋮

$N/2, N/2$

Quiz

Output for $N = +3$

```
void solve (int N) {
    if (N == 0)
        return
    print (N)
    solve (N-1)
    print (N)
```

solve (2)
 3 $\boxed{2 \ 2 - 1 \ 1 \ 2 \ 3}$ 3
 3 2 1 1 2 3

Quiz

Output for $N = -3$

```
void solve (int N) {
    if (N == 0)
        return
    print (N)
    solve (N-1)
```

solve (-3)

↓

solve (-4)

↓

solve (-5)

↓
-6

O/P

$-3 \ -4 \ -5$
 stack overflow
 ↓
 MLE

- ① MCQs Recordings
- ② Visualizes Link
- ③ Optional Class

dec

$N \rightarrow 1$

solve(A) <

} printNos(A)
println()

// Given A, print all nos from 1->A

void printNos(A) <

} if (A == 0) return

printnos(A-1)

cout << A << " ";

$A = 4$

1 - 2 - 3 - 4

1 - 2 - 3 - 4 -

```

printNos (num, A) <
    if (num == A) <
        cout << A; return
    cout << num << " "
    printNos (num+1, A)

```



A - [] A -

3 - 2 - 1 - 1 - 2 - 3,

// Print A A-1 ... 1 1 2 ... ; A



values (A) <

if ($A \Rightarrow 0$)
return

```

print(A)
values (A-1)
print(A)

```

A A-1 | | . . . A-1 A

print (R)

print (A)