

Welcome 😊

Agenda: Constraints

Space Complexity

Arrays

2-3 questions

Importance of constraints

$$1 \leq N \leq 10^5$$

Complexity

Works ?

$$O(N^3)$$

✗

$$O(N^2 \log N)$$

✗

$$O(N^2)$$

✗

$$O(N \log N) \rightarrow$$

$$10^5 \log_2(10^5) \rightarrow 1.6 \times 10^6$$

✓

$$1 \leq N \leq 10^6$$

Complexity

Works ?

$$O(N^3)$$

✗

$$O(N^2 \log N)$$

✗

$$O(N^2)$$

✗

$$O(N \log N) \rightarrow$$

$$10^6 \log_2(10^6) \Rightarrow 2 \times 10^7$$

might
or
might not
work

$\leq 10^6$
✓

$10^7 - 10^8$
might
work
Not Sure

$> 10^8$
✗

Space Complexity

⇒ It is the amount of space additionally used by the algo., other than the input space to perform operations.

```
int func ( int arr[] , int n )  
{  
    . . .  
    . . .  
    . . .  
}
```

S.C

⇒

~~$O(N)$~~
 $O(1)$

⇒ I/p size is not considered in space complexity.

```
⑩  
func ( int N )  
{  
    int n = N      ⇒ 4B  
    int y = n + n   ⇒ 4B  
    long z = n + y  ⇒ 8B  
}
```

space = 16B
 $O(1)$

==

func (int N)

{ int n = N

$\Rightarrow 4B$

int y = n+n

$\Rightarrow 4B$

long z = n+y

$\Rightarrow 8B$

int arr[10]

$\Rightarrow 4 * 10 = 40$

int arr[N]

$\Rightarrow 4 * N$

}

\Rightarrow

$56B + 4 * N B$

$\Rightarrow O(N)$ space comp.

==

func (int N)

{ int n = N

$\Rightarrow 4B$

int y = n+n

$\Rightarrow 4B$

long z = n+y

$\Rightarrow 8B$

int arr[10]

$\Rightarrow 4 * 10 = 40$

int arr[N][N]

$\Rightarrow 4 * N * N$

}

\Rightarrow

$56B + 4 * N * N$

$\Rightarrow O(N^2)$ space comp.

Arrays

- ⇒ Collection of **same** types of data
 - ↳ primitive / complen.
- ⇒ **contiguous** memory block.

arr

0	1	2	3	4	5	6	7
3	4	7	2	1	8	6	5

⇒ Access index $i \Rightarrow arr[i]$

⇒ Declare array $\Rightarrow int\ arr[N]$ ↗ size of array,

first element $\rightarrow 0$

last element $\rightarrow \underline{\underline{N-1}}$

Q Print all the elements of the array

```
void printArray ( int arr[] , int N)
```

```
{
```

```
    for ( i=0 ; i<N ; i++)
```

```
    {
```

```
        print ( arr[i] )
```

```
    }
```

```
}
```

T.C $\Rightarrow O(N)$

S.C $\Rightarrow O(1)$

Q Given an array of size N. Reverse the entire array.

eg: A: { 1 2 3 4 5 6 }
%p { 2 5 4 3 2 1 }

A: { 1 2 3 4 5 6 }
 ↑ ↑
 { 5 2 3 4 1 6 }
 ↑ ↑
 { 5 4 3 2 1 6 }
 ↑↑

Code
void reverse (int arr[], int n)

i = 0 j = n-1

while (i < j)

{ // swap elements

int temp = arr[i]

arr[i] = arr[j]

arr[j] = temp

i++

j--

T.C $\Rightarrow O(N)$

S.C $\Rightarrow O(1)$

Q Given an array of size N , rotate array from last to first by K times.

eg: { 1 2 3 4 5 }

$K=1$ { 5 1 2 3 4 }

$K=2$ { 4 5 1 2 3 }

Brute force

Rotate the array one distance at a time.

T.C $O(N \times K)$

Optimized

{ 1 2 3 4 5 }

O/p $K=2$ { 4 5 1 2 3 }

1) Reverse the array

2) { 5 4 3 2 1 }

reverse

3) { 4 5 3 2 1 }

reverse

{ 4 5 1 2 3 }

	{	1	2	3	4	5	}
$k=1$	{	5	1	2	3	4	}
$k=2$	{	4	5	1	2	3	}
$k=3$	{	3	4	5	1	2	}
$k=4$	{	2	3	4	5	1	}
$k=5$	{	1	2	3	4	5	}
<u>$k=6$</u>	{	5	1	2	3	4	}

$$k = k \% N$$

H.W

$$T.C = N + N = O(N)$$

$$S.C = \underline{O(1)}$$

Dynamic Arrays

→ array with no fixed size.

→ automatically resize