SCALER

Interview Problems

Starting 9:05

TABLE OF CONTENTS

Revision

1. 'A' → 65; 'a' → 97, '0' - 48
2. Substring → a continuous part of a string.
3. "abc" → a, ab, abc, b, bc, c

   $N\left(\frac{N+1}{2}\right)$ → $3\left(\frac{3+1}{2}\right) = 6$
4. StringBuilder
5. "ana madam m"

< *Question* > : Given a binary array [ ]. We can almost replace a single 0 with 1. Find the maximum consecutive 1's we can get in the array[ ] after the replacement.

ex:- { 1, 1, 0, 1, 1, 0, 1, 1 }

5 → Output

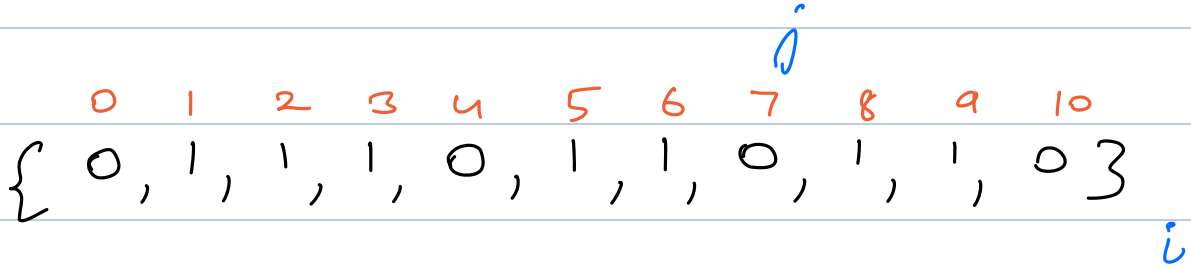
Q:- 1. { 1, 1, 0, 1, 1, 0, 1, 1, 1 }

6.

Q:- 2. { 0, 1, 1, 1, 0, 1, 1, 0, 1, 1, 0 }

6.

Approach:- Iterate over the array. Whenever a '0' is found, col. the consecutive ones on left (l) & right (r). Compare the global ans with (l+r+1). Update ans if needed.



ans = 0;

r = 0

l = 2

ans = max(ans, 0 + 2 + 1)

Code

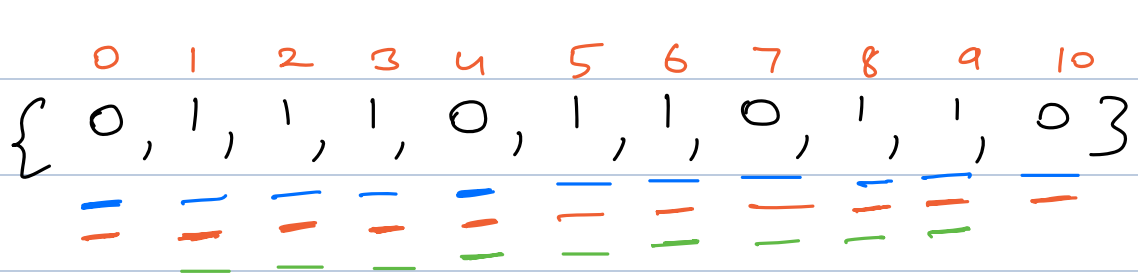
```
int solve (int arr[], int N) {
    int totalOnes = 0;
    for (i=0; i<N; i++) {
        if (arr[i] == 1) { totalOnes++; }
    }
    if (totalOnes == N) { return N; }

    int ans = 0;
    for (i=0; i<N; i++) {
        if (arr[i] == 0) {
            l=0; r=0;
            j = i+1;   # Consecutive 1's on right
            while (arr[j]==1 && j<N){
                j++; r++;
            }
            j = i-1;   # Consecutive 1's on left
            while (arr[j]==1 && j>=0){
                j--; l++;
            }
            ans = max(ans, l+r+1);
        }
    }

    return ans;
}
```

Q:3

T.C → O(N).



3 N

→ {0, 0, 1, 1, 0, 1, 1, 0}

2 2

Every element is getting accessed at max 3 times.

< Question > : Given a binary array [ ]. We can swap a single 0 with 1. Find the maximum consecutive 1's we can get in the array[ ] after almost 1 swap.

ex:- {1, 0, 1, 1, 0, 1}

Q:-4.

{1, 1, 0, 1, 1, 1}

Approach :- Some as previous, just need to check for the case when the sequence you are trying to make have enough ones to be supplied.

```
int solve ( int arr[], int N) {

    int totalOnes = 0;
    for (i=0; i<N; i++) {
        if (arr[i] == 1) { totalOnes ++; }
    }
    if (totalOnes == N) { return N; }

    int ans = 0;
    for (i=0; i<N; i++) {
        if (arr[i] == 0) {
            l=0; r=0;
            j = i+1;   # Consecutive 1's on right.
            while (arr[j]==1 && j<N){
                j++; r++;
            }
            j = i-1;   # Consecutive 1's on left
            while (arr[j]==1 && j>=0){
                j--; l++;
            }
            if (l+r == totalOnes) {   totalOnes
                ans = max(ans, l+r);
            }
            else {
                ans = max(ans, l+r+1);
            }
        }
    }
    return ans;
}
```

T.C → O(N)

Majority Element

< *Question* > : Given array [N]. Find the majority element

↓

Elements which occurs more than N/2 times.

- *You can assume that majority element always exists.*

N = 3

ex :- { 2, 1, 4 } No majority element exists.

N = 8

ex :- { 3, 4, 3, 2, 4, 4, 4, 4 }

Atleast 5.

ans = 4.

N = 8

ex :- { 3, 3, 4, 2, 4, 4, 2, 4 }

No majority elem.

N = 11

Q:-5. {3, 4, 3, 6, 1, 3, 2, 5, 3, 3, 3}

Atleast 6.

ans = 3.

{1, 2, 3, 3, 3, 3, 3, 3, 4, 5, 6}.

Q:-7 I am always going to have a Style majority element.

Brute Force :- (1) Taking 2 loops. For each elem, count its frequency in another loop. T.C → O(N²).

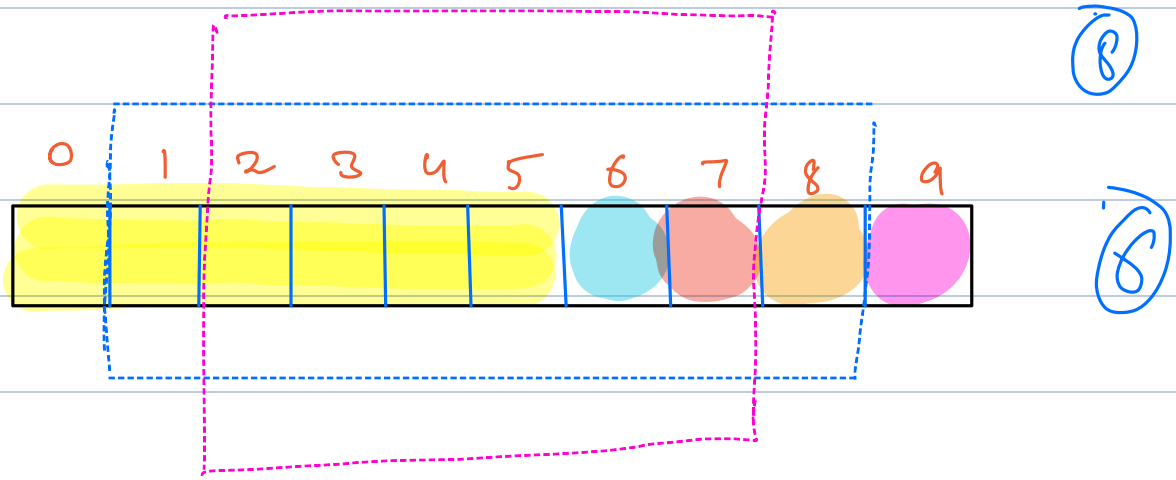
(2) Sort the array. T.C → O(N log N).

Iterate over the array & keep count of consecutive N/2 + 1 elements.

Observations

Moore's Voting Algorithm

1. There is going to be only 1 majority element.
2. Removing 2 distinct elements from my array has no impact on majority elem.



Another Example :-

OP : 🧍🧍🧍🧍🧍🧍🧍🧍🧍 = 9

YP : 🧍🧍🧍 = 3

RP : 🧍🧍 = 2

GP : 🧍🧍🧍 = 3

At least 9.

Removing	Count Of :- Orange	Yellow	Red	Green	Winner
	9	3	2	3	Orange
1 of 1R	8	3	1	3	Orange.
1 yr 1G	8	2	1	2	Orange

Algo:-

	0	1	2	3	4	5	6	7	8	9	10	i
ex:-	3	4	3	6	1	3	2	5	3	3	3	

majEleIndex = 0; ~~1 4 6 8~~

count = 1; ~~0 1 0 1 0 1 0 1 2~~ 3

Important Point:- The above algo gives us the potential majority element. We need to verify once whether the element pointed by majEleIndex has a frequency > N/2 or not.

Code.

```
int majEleIndex = 0;
int count = 1;

for (i=1 ; i< N; i++) {
    if (count == 0){
        majEleIndex = i;
        count = 1;
    }
    else {
        if (arr[i] == arr[majEleIndex]){
            count++;
        } else {
            count--;
        }
    }
}
```

majEleIndex is pointing at a potential majority Element.

So, verify it.

```
int cout = 0;
for (i=0; i< N; i++) {
    if (arr[i] == arr[majEleIndex]){
        cout++;
    }
}
if (cout > N/2) {
    return arr[majEleIndex];
}
else {
    return -1; # No maj element exists.
}
```

{3, 3, 3, 4}.

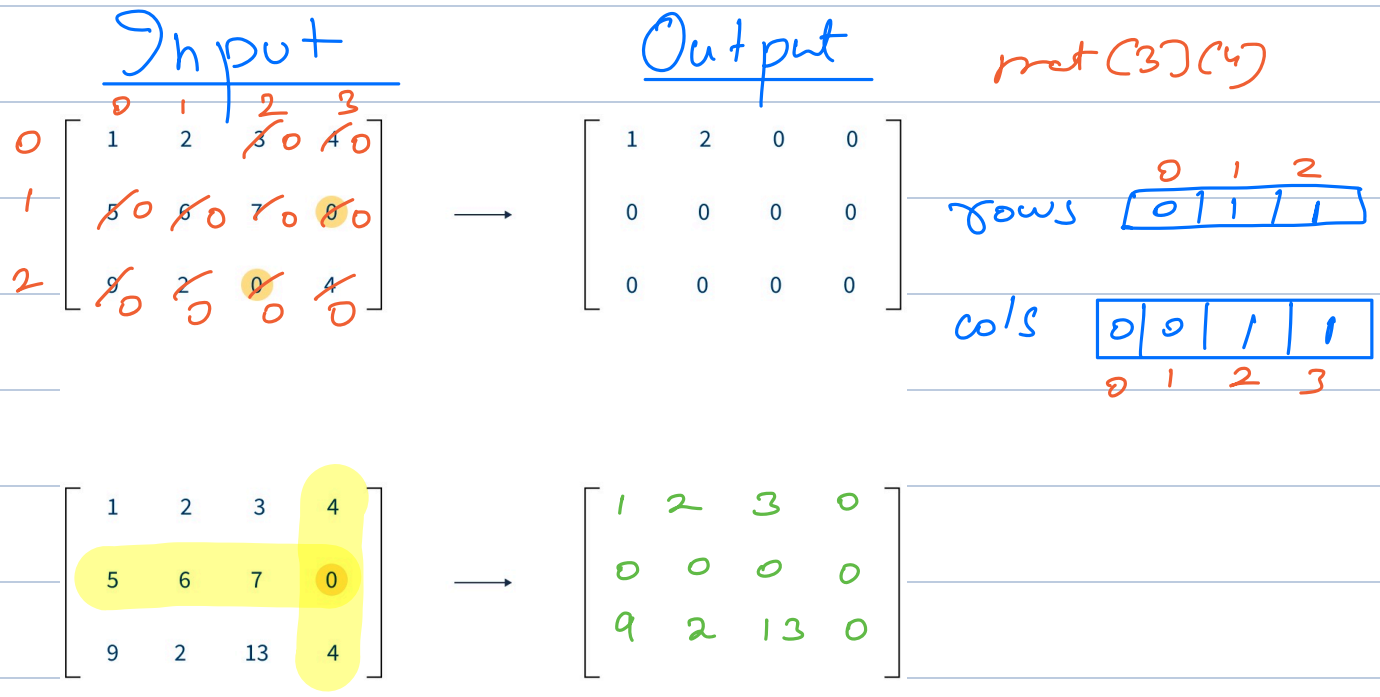
10:50

T.C → O(N)

S.C → O(1)

< *Question* > : Given array [N] [M].

Make all elements in a row and column zero if arr[i][j] = 0



Approach:-

Take two arrays, 1 of size rows & other of size columns.

Initial both these arrays with 0.

Iterate on every cell of the matrix.

if arr[i][j] == 0;

update row[i] = 1 & col[j] = 1;

Iterate again over the entire matrix. For every cell (i,j) if either row(i) is 1 or col(j) is 1, mark that cell to 0:-

Todo → Code.

T.C → O(N*M)
S.C → O(N+M)

90 + PSP

4 Q's:-

1 → Assignment.
2 Q → Easy / Medium.
1 Q → Medium / Hard.

Passing Criteria → 75+ Marks.