

## Agenda

Modular Arithmetic Intro

Count pairs whose sum mod  $m$  is 0

GCD Intro

Properties of GCD

Delete one

$$x \cdot y \cdot 5 =$$

## Modulo ( $\div$ )

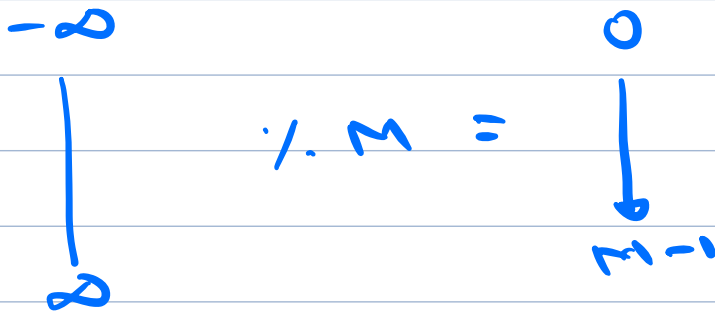
$A \div M = \text{Remainder}$  when  $A$  is divided by  $M$

Range of  $A \div M \rightarrow [0, m-1]$

Why do we need Mod?



Limit the range of the data



Int  $\rightarrow 2 \times 10^9$   
Long  $\rightarrow 4 \times 10^{18}$

ans  $\div 10^9 + 7$   
 $[0 \rightarrow 10^9 + 6]$   
(Int)

Properties of Mod (on arithmetic operators)

$$\textcircled{1} (a + b) \div m = (a \div m + b \div m) \div m$$

Eg.  $a = 9, b = 8, m = 5$

$$(9+8) \cdot 1.5 \quad \Bigg| \quad (9 \cdot 1.5 + 8 \cdot 1.5) \cdot 1.5$$

$$= 17 \cdot 1.5 = 2 \quad \Bigg| \quad (4 + 3) \cdot 1.5 = 7 \cdot 1.5$$

$$= 2$$

②  $(a \times b) \cdot m = (a \cdot m \times b \cdot m) \cdot m$

Eg.  $a=9, b=8 \quad m=5$

$$(9 \times 8) \cdot 1.5 \quad \Bigg| \quad (9 \cdot 1.5 \times 8 \cdot 1.5) \cdot 1.5$$

$$\Rightarrow 12 \cdot 1.5 \quad \Bigg| \quad \Rightarrow (4 + 3) \cdot 1.5$$

$$\Rightarrow 2 \quad \Bigg| \quad \Rightarrow 12 \cdot 1.5 = 2$$

③  $(a \cdot m) \cdot m = a \cdot m$

$$12 \cdot 1.5 = 2 \cdot 1.5 = 2 \cdot 1.5 = 2 \dots$$

④  $(a+m) \cdot m = (a \cdot m + m \cdot m) \cdot m$

$$= (a \cdot m) \cdot m$$

$$(a+m) \cdot m = a \cdot m$$

$$(18+5) \cdot 1.5$$

$$= 23 \cdot 1.5 = 3$$

$$18 \cdot 1.5 = 3$$

$$\underline{\text{|||||}} \quad \underline{\text{|||||}} \quad \underline{\text{|||||}} \quad \text{||||}$$

eg.  $(-4) \cdot 1.6 = (-4+6) \cdot 1.6 = 2 \cdot 1.6 = 2$

$$(-8) \cdot 1.6 = (-8+6) \cdot 1.6 = -2 \cdot 1.6 = (-2+6) \cdot 1.6$$

$$= 4 \cdot 6 = 4$$

$$(5) \quad (a - b) \% m = (a \% m - b \% m + m) \% m$$

For eg.  $a = 17$   $b = 8$   $m = 5$

$$\begin{aligned} \Rightarrow (17 - 8) \% 5 & \quad | \quad (17 \% 5 - 8 \% 5) \% 5 \\ \Rightarrow 9 \% 5 & = 4 \quad | \quad (2 - 3) \% 5 = (-1) \% 5 \\ & \quad \quad \quad = 4 \end{aligned}$$

$$(-1) \% 5 = (-1 + 5) \% 5 = 4 \% 5 = 4$$

$$(6) \quad (a^b) \% m = (a \times a \times a \dots \text{b times}) \% m$$

$$= (a \% m) \times (a \% m) \times (a \% m) \dots$$

$\text{b times } \% m$

$$= ((a \% m)^b) \% m$$

$$\text{Quiz : } (37^{103} - 1) \% 12$$

$$\Rightarrow ((37^{103}) \% 12 - 1 \% 12 + 12) \% 12$$

$$\Rightarrow (1 - 1 + 12) \% 12 = 0$$

$$\textcircled{1} (a - b) \% m = (a \% m - b \% m + m) \% m$$

$$\textcircled{2} a^b \% m = ((a \% m)^b) \% m$$

$$\begin{aligned} 37^{103} \% 12 &= ((37 \% 12)^{103}) \% 12 \\ &= (1^{103}) \% 12 = 1 \end{aligned}$$

$$\text{If } A \% m = 0$$

A is completely divisible by m  
m is a factor of A

Prob 1: Given N array elements, find count of pairs  $(i, j)$  such that  $(a[i] + a[j]) \% m = 0$   
 $i \neq j$  and pair  $(i, j)$  is same as pair  $(j, i)$

2 idx  $(i, j) \rightarrow$  pair sum divisible by  $m$

$A = \langle 4, 3, 6, 3, 8, 12 \rangle$        $ans = 3$   
 $m = 6$

$$(6 + 12) \% 6 = 0$$

$$(3 + 3) \% 6 = 0$$

$$(4 + 8) \% 6 = 0$$

BF: Go to all unique pairs, calculate their sum. If their sum  $\% m == 0$   
cnt++

int ans = 0

```
for (i = 0 ; i < n ; i++)  
    for (j = i+1 ; j < n ; j++)  
        if (a[i] + a[j] % m == 0)  
            ans++  
return ans
```

TC:  $O(n^2)$   
SC:  $O(1)$

# Optimised Approach

$$(a+b) \% m = (a \% m + b \% m) \% m$$

$\downarrow$   $\downarrow$   
 $0$   $0$

Look for a pair whose sum is divisible by  $m$

sum of remainders should be divisible by  $m$

$$(x_1 + x_2) \% m = 0$$

$$x_1, x_2 \rightarrow [0 \quad m-1]$$

$$1 \rightarrow m-1$$

$$2 \rightarrow m-2$$

$$3 \rightarrow m-3$$

...

$$x \rightarrow m-x$$

$$(x_1 + x_2) \% 6 = 0$$

$$x_1, x_2 \rightarrow [0 \quad 5]$$

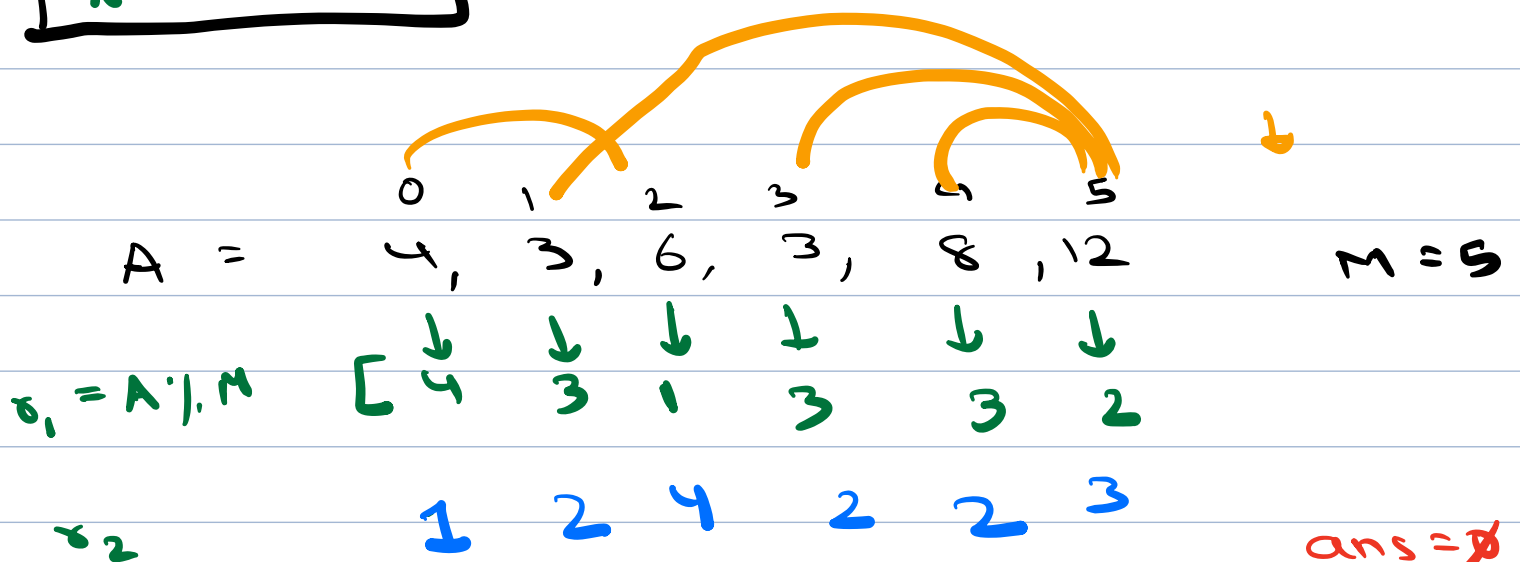
$$1 \rightarrow 5$$

$$2 \rightarrow 4$$

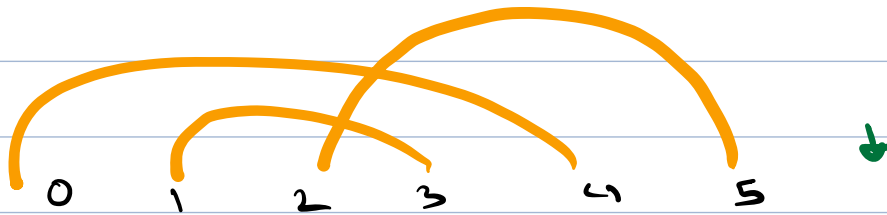
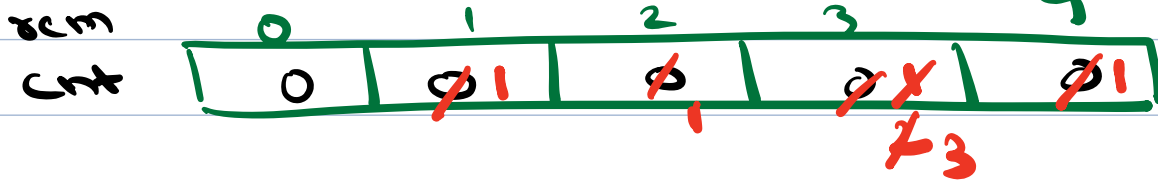
$$3 \rightarrow 3$$

$$0 \rightarrow 0$$

Observation :  $x_1 + x_2 = m$



ans = ~~0~~  
~~1~~  
4

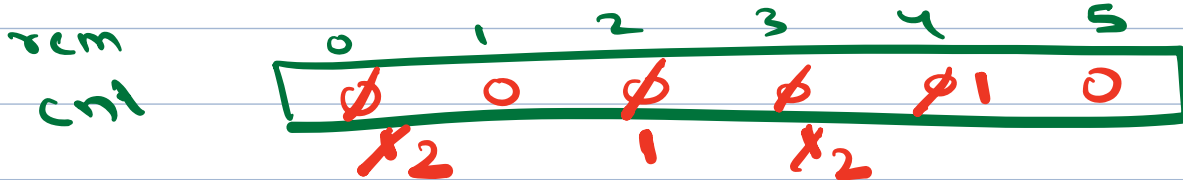


A = 1, 3, 6, 9, 8, 12, 9 M = 6

A % M =

4	3	0	3	2	0	3
↓	↓	↓	↓	↓	↓	
2	3	0	3	4	0	3

ans = 0  
1  
2  
3



ans = 3  
↓ + 2  
5

idx → rem      freq[rem]

int countPairs (int A[], int N, int M) {

int freq[M] = {0}    // idx → 0 to M-1

int ans = 0

for (int i = 0 ; i < N ; i++) {

int x1 = A[i] % M

int x2 = M - x1

if (x1 == 0)

x2 = 0

ans += freq[x2]

freq[x1] ++



1  
return ans

TC:  $O(N)$   
SC:  $O(1)$

Break till 10:42

GCD  $\rightarrow$  Greatest Common Divisor

HCF  $\rightarrow$  Highest Common Factor

If  $x$  is a factor of  $A$

$$\Rightarrow A \div x = 0$$

$GCD(A, B)$  = Greatest factor that divides both  $A$  and  $B$

$$GCD(A, B) = x$$

①  $A \div x = 0$

②  $B \div x = 0$

③  $x$  is largest no. which divides both

$$GCD(15, 25) = 5$$

↓   ↓  
1   1  
3   5  
5   25  
15

$$GCD(12, 30) = 6$$

↓   ↓  
1   1  
2   2  
3   3  
4   5  
6   6  
12   10  
15  
30

$$GCD(a, b) = GCD(a, -b) = GCD(-a, b) = GCD(-a, -b)$$

$$GCD(0, 4) = 4$$

↓   ↓  
1   1  
2   2  
3   4  
...  
0

$\frac{0}{0} \rightarrow$  undefined

$$0 \div 1 = 0$$

$$0 \div 2 = 0 \dots$$

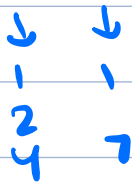
$$\text{GCD}(0, a) = a$$



$$\text{GCD}(0, 0) = \infty$$

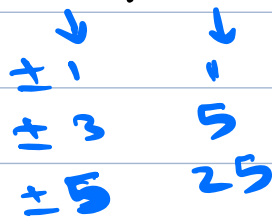


$$\text{GCD}(4, 7) = 1$$



(Co-primes : pair of nos. whose GCD is 1)

$$\text{GCD}(-15, 25) = 5$$



$$\text{GCD}(-15, 25) = \text{GCD}(15, 25)$$

Properties of GCD

$$\textcircled{1} \quad \text{GCD}(A, B) = \text{GCD}(B, A)$$

$$\textcircled{2} \quad \text{GCD}(0, A) = A$$

$$\textcircled{3} \quad \text{GCD}(A, B, C) = \text{GCD}(\text{GCD}(A, B), C)$$

$$= \text{GCD}(\text{GCD}(A, C), B)$$

$$= \text{GCD}(\text{GCD}(B, C), A)$$

$$\begin{aligned} \text{GCD}(2, 3, 4) &= \text{GCD}(\text{GCD}(2, 3), 4) = 1 \\ &= \text{GCD}(\text{GCD}(3, 4), 2) = 1 \end{aligned}$$

$$\textcircled{4} \quad \text{GCD}(1, A) = 1$$

$\downarrow$     $\downarrow$   
 $\vdots$   
 $A$

$$\textcircled{5} \quad \text{Given } A \geq B > 0$$

$$\boxed{\text{GCD}(A, B) = \text{GCD}(A - B, B)}$$

$$\begin{aligned} \text{GCD}(A, B) &= \text{GCD}(A - B, B) \\ &= \text{GCD}(A - B - B, B) \quad \xrightarrow{A-2B} \\ &= \text{GCD}(A - B - B - B, B) \quad \xrightarrow{A-3B} \\ &= \text{GCD}(A - 3B - B, B) \quad \xrightarrow{A-4B} \end{aligned}$$

$$\textcircled{6} \quad \boxed{\text{GCD}(A, B) = \text{GCD}(A \cdot 1 \cdot B, B)}$$

$$\begin{aligned} \text{eg. } \text{GCD}(17, 5) &= \text{GCD}(17 - 5, 5) \quad \text{12} \\ &= \text{GCD}(17 - 5 - 5, 5) \quad \text{7} \\ &= \text{GCD}(17 - 5 - 5 - 5, 5) \quad \text{2} \\ &= \text{GCD}(17 \cdot 1 \cdot 5, 5) \end{aligned}$$

e.g.  $\text{gcd}(24, 16) = \text{gcd}(8, 16)$

$\downarrow$   
 $\text{gcd}(8, 16)$

$$\text{gcd}(A, B) = \text{gcd}(B, A \% B)$$

$$\text{gcd}(24, 16) = \text{gcd}(16, 8)$$

$$= \text{gcd}(8, 0) = 8$$

$$\text{gcd}(16, 24) = \text{gcd}(24, 16)$$

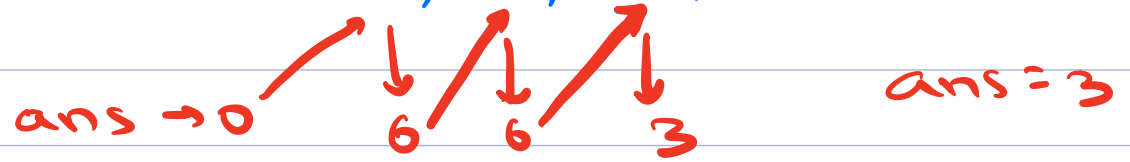
```
int gcd(int a, int b) {  
    if (b == 0) return a  
    return gcd(b, a % b)  
}
```

$$T.C: O(\log_2(\min(a, b)))$$

$$S.C: O(\log_2(\min(a, b)))$$

Prob: Calculate gcd of entire array

arr[3] = < 6, 12, 15 >



```
int ans = 0
```

```
for (i = 0; i < N; i++) {
```

```
    ans = gcd(ans, arr[i])
```

```
}  
return ans
```

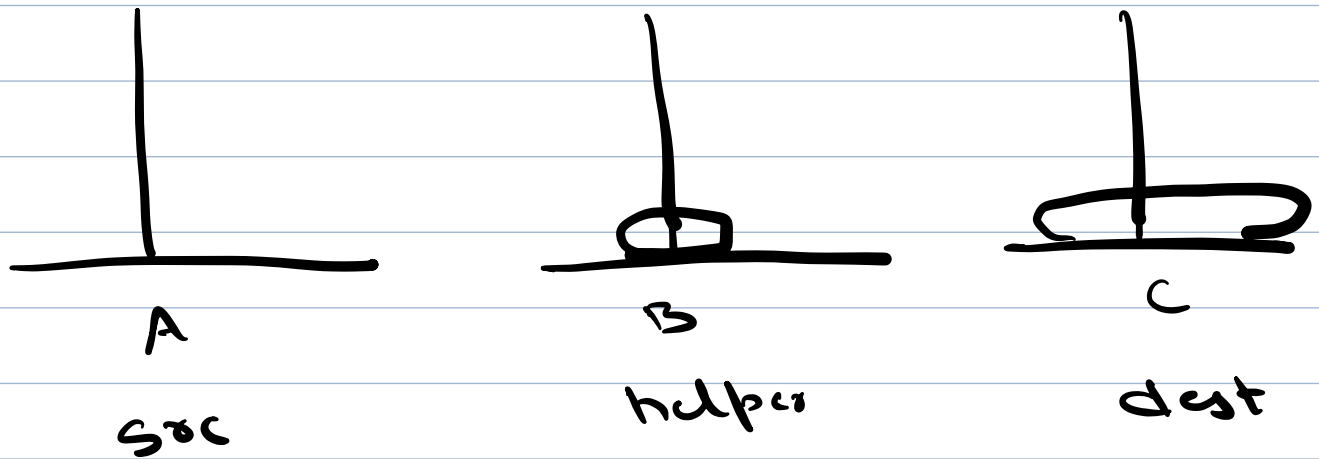
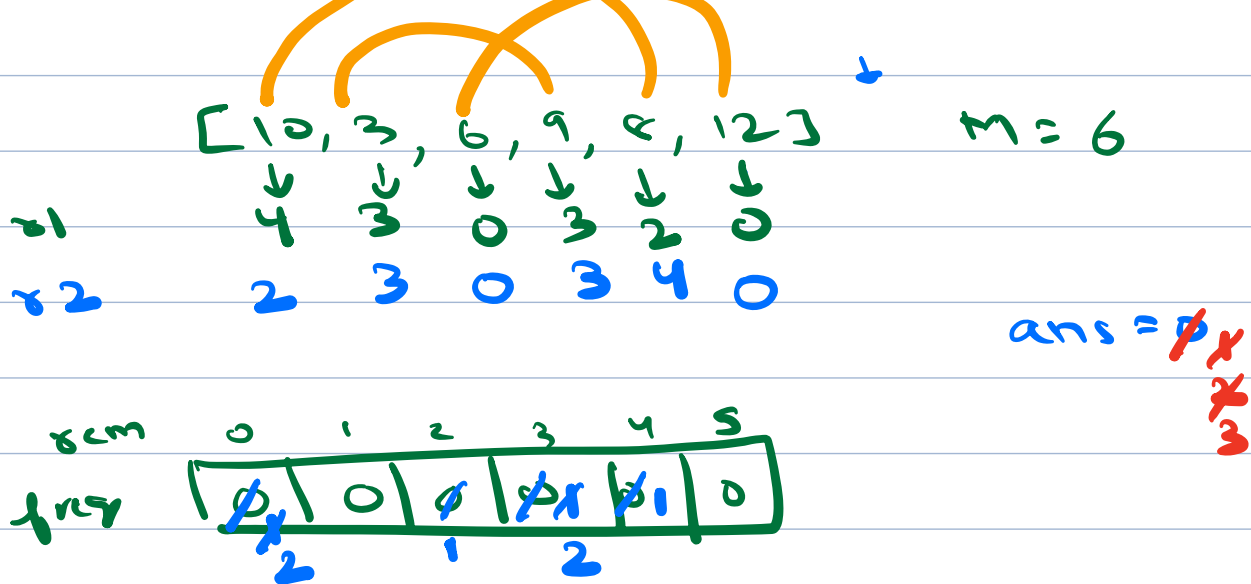
$$\text{gcd}(x, y) \xrightarrow{T_c} O(\log_2 \min(a, b))$$

$$\text{gcd}(\text{ans}, \text{arr}[i]) \rightarrow O(\log_2 \text{max\_ele})$$

$$T_c : O(N \times \log_2 \text{max\_ele})$$

---

$$\text{LCM}(a, b) = \frac{a \times b}{\text{gcd}(a, b)}$$



- ① Move S disc      src → helper
- ② Move L disc      src → dest
- ③ Move S disc      helper → dest



① Move smaller  $N-1$  discs

src  $\rightarrow$  helper

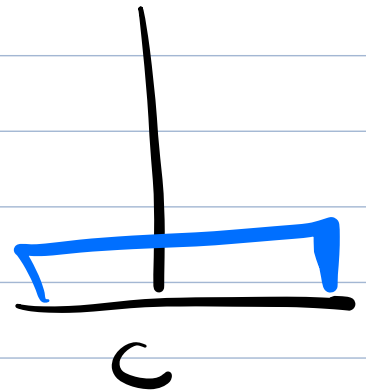
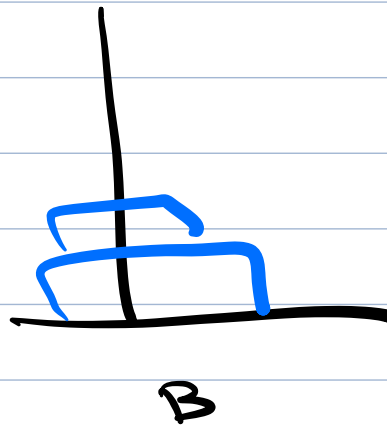
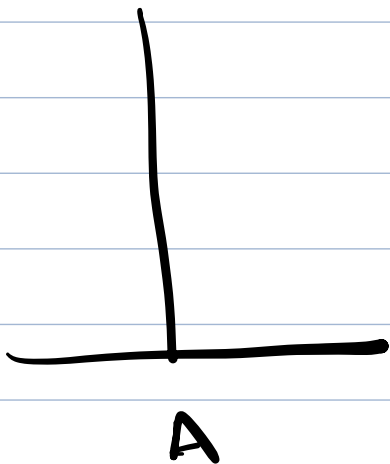
② Move  $L$  disc

src  $\rightarrow$  Dest

③ Move smaller  $N-1$  discs

helper  $\rightarrow$  disk

$N=3$



1  
3

src  
A

Helper  
B

Dest  
C



2

A

C

B

①

1

A  $\rightarrow$  C

②



3

<

2

B

A

C

// Given n discs, it will move src  $\rightarrow$  dest  
using helper

```
void toh (int N, char src, char help, char dest) {
```

```
    if (N == 0) return
```

```
    toh (N-1, src, dest, help)
```

```
    printf (src  $\rightarrow$  dest)
```

```
    toh (N-1, help, src, dest)
```

```
}
```

```
toh (N, A, B, C)
```