

Agenda

1. Longest Substring without Repeating Characters
2. First non repeating Element
3. Subarray with sum 0
4. Subarray with sum k

old → New ✓

M	18 Nov	M	→ Break
W	20 Nov	W	→ Class
	22 Nov	F	→ class

Next Fri → Break
(22 Nov)

2. Given a string s, find length of the longest substring without repeating characters.

str: " abc abc bb "

ans: 3

str: " b b b b b "

ans: 1

str: " p wwk w "

ans: 3

Brute Force: Go to all substrings, check if substr is valid (without duplicates), compare its length with ans and keep max in ans.

$\approx N^2$ substr

```
int ans = 0
for (s _____) <
    for (e _____) <
        // s e
        for (i = s; i <= e; i++)
            put all chars in hs
        if (hs.size() == e - s + 1)
            ans = max(ans, e - s + 1)
```

$$[a, b] = b - a + 1$$

TC: $O(N^3)$

SC: $O(N)$

Optimized :

ans = 5 bcade

0 1 2 3 4 5 6 7
a b c a d e c g

ans = ~~1~~ ~~2~~ ~~3~~ ~~4~~ 5

0 1 2 3 4 5 6 7
a b c a d e c g

s = 0

ans = ~~0~~ ~~1~~ ~~2~~ ~~3~~ 5

d	b
c	a

 d e c g

// str, int n

int ans = 0

HashSet <char> hs

int s = 0

for (e = 0 ; e < n ; e++) <

while (hs.contains(str[e]) == true) <

hs.remove(str[s])

s++

hs.add(str[e])

ans = max (ans, hs.size())

substr size

return ans

0 1 2 3 4
a b c d b

a	b
c	d

Tc: $O(N)$

Every char can be processed twice,
added once and removed once from
hashtxt

Sc: $O(\min(N, M))$

Str \rightarrow a to z

M \rightarrow size of character set
(ASCII \rightarrow 128 chars)

2. Find the first non-repeating element. ^{unique} from start

Ex 1 $arr[6] = \langle 1, 2, 3, 1, 2, 5 \rangle$ ans = 3

Ex 2 $arr[8] = \langle 4, 3, 3, 2, 5, 6, 4, 5 \rangle$ ans = 2

Ex 3 $arr[7] = \langle 2, 6, 8, 4, 7, 2, 9 \rangle$ ans = 6

Idea:

1. Insert all elements in hm
2. Iterate on hm and get 1st key with freq = 1

HM	
elem	freq
1	2
2	2
3	1
5	1

[Note - in hashmap, insertion order is not maintained; when we print hashmap we'll get any order]

Idea:

1. Insert all elements in hm
2. Iterate on array and get elem with freq = 1

```
// int arr[], int N
```

```
HashMap <int, int> hm
```

```
for (i=0; i<n; i++) <
```

```
    if (hm.containsKey(arr[i]) == true)
```

```
        hm[arr[i]]++
```

```
    else <
```

```
        | hm.insert(arr[i], 1)
```

```
>
```

```
for (i=0; i<N; i++) <
```

```
    int freq = hm[arr[i]]
```

```
    if (freq == 1)
```

```
        return arr[i]
```

```
>
```

```
return -1 // no element is  
unique
```

10:35

TC: $O(N)$

SC: $O(N)$

3. Given an array of N elements, check if there exists a subarray with sum equal to 0.

$N = 10$ ar: 2 2 1 -3 4 3 1 -2 -3 2

ans: true

$N = 10$ ar: 2 3 4 -2

ans: false

Brute Force: Go to all subarrays, iterate and calculate sum. If any subarray has 0 sum, return true otherwise at the end return false

1 subarray $\rightarrow O(N)$ (Iterate and get sum)

N^2 " $\rightarrow O(N^2 \times N)$

TC: $O(N^3)$

SC: $O(1)$

bf

TC: $O(N + N^2)$
 $\approx O(N^2)$

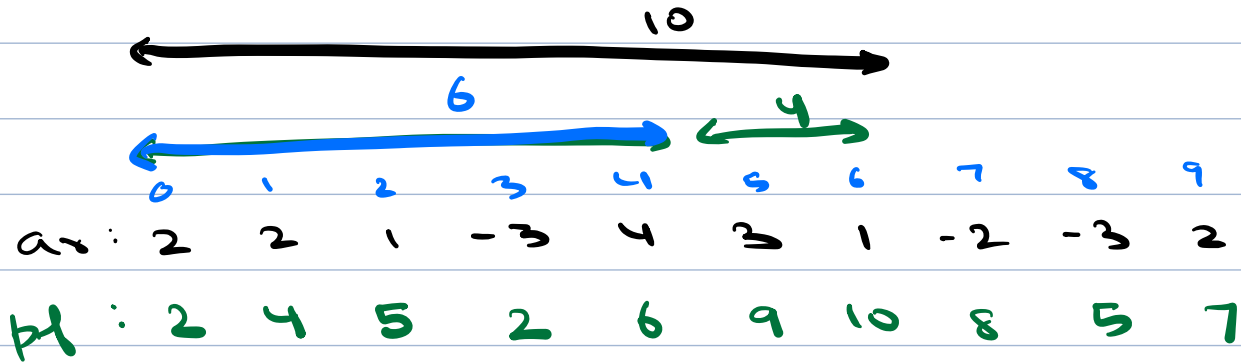
SC: $O(N)$

carry forward

TC: $O(N^2)$

SC: $O(1)$

Optimization: Sum of a subarray: Prefix Sum



* subarray

$$\text{sum}(l \rightarrow r) = 0$$

$$\text{sum}(0 \rightarrow r)$$

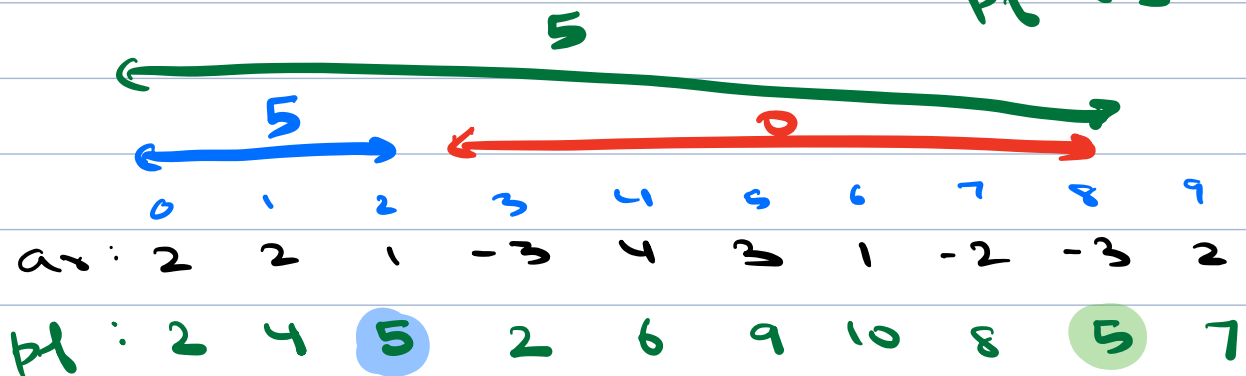
$$\text{pf}[r] = 0$$

$$\text{sum}(l \rightarrow r)$$

$$\text{pf}[r] - \text{pf}[l-1] = 0$$

$$\Rightarrow \text{pf}[r] = \text{pf}[l-1]$$

Duplicates in pf[]




```
// int arr[], int n
```

```
HashSet<int> hs; int sum=0
```

```
for (int i=0; i<n; i++) {
```

```
    sum = sum + arr[i]
```

```
    hs.insert(sum)
```

```
}
```

```
if (hs.size() < N || hs.contains(0) == true)
```

```
    return true
```

```
else
```

```
    return false
```

TC : $O(N)$

SC : $O(N)$

↓

hashset

4. Given an array of N integers, check if there exists a subarray with sum k .

0 1 2 3 4 5 6 7 8
 $a[] = 2 \quad 3 \quad 9 \quad -4 \quad 1 \quad 5 \quad 6 \quad 2 \quad 5$

$k = 11$ **true** $\langle 2, 3, 9, -4, 1 \rangle \langle 5, 6 \rangle$

$k = 15$ **true** $\langle -4, 1, 5, 6, 2, 5 \rangle$

$a: [5 \quad 10 \quad 20 \quad 100 \quad 105] \quad k = 110$ **false**

Optimization: Sum of a subarray: Prefix Sum

subarray $\text{sum}(l \rightarrow r) = k$

$\text{sum}(0 \rightarrow r)$

$pf[r] = k$

$\text{sum}(l \rightarrow r)$

$pf[r] - pf[l-1] = k$

$a - b = k$

A pair in $pf[]$
whose diff is k

$\Rightarrow b = a - k$

$a[] = 2 \quad 3 \quad 4 \quad -4 \quad 1 \quad 5 \quad 6 \quad 2 \quad 5$
 $K = 11$

$pf \rightarrow 2 \quad 5 \quad 14 \quad 10 \quad 11 \quad 16$

partner
 $\downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
 $-9 \quad -6 \quad 3 \quad -1 \quad 0 \quad 5$
 $\times \quad \times \quad \times \quad \times \quad \times \quad \checkmark$

Hs (all $pfsum[i]$ of past)

2	5
14	10
11	

$\star pf[i] = K = 11$

$sum(0-i) = K$

$// int a[], int N, int K$
 Hashset $<int>$ hs $// pfsum[i]$

$int sum = 0$

for (int $i = 0$; $i < n$; $i++$) {

$sum = sum + a[i]$

if (sum == K)

return true

if (hs.contains(sum - K) == true)

return true

hs.insert(sum)

}
 return false

$TC : O(N)$

$SC : O(N)$
 hashset

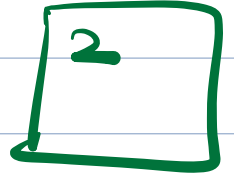
Doubts

$K = 1$

a-k

pl

↓
2 1 2 1
2 3
↓ ↓
1 2
✓



① Assignment

Approach → Pen & Paper

Notes / En

↓
Code

