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An effective hybrid EDA-based algorithm for solving multidimensional knapsack problem

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ARTICLE INFO

Keywords:

Multidimensional knapsack problem Estimation of distribution algorithm Probability model Hybrid algorithm Design of experiment

ABSTRACT

In this paper, an effective hybrid algorithm based on estimation of distribution algorithm (EDA) is proposed to solve the multidimensional knapsack problem (MKP). With the framework of EDA, the probability model is built with the superior population and the new individuals are generated based on probability model. In addition, an updating mechanism of the probability model is proposed and a mechanism for initializing the probability model based on the specific knowledge of the MKP is also proposed to improve the convergence speed. Meanwhile, an adaptive local search is proposed to enhance the exploitation ability. Furthermore, the influences of parameters are investigated based on Taguchi method of design of experiment and the importance of repair operator is also studied via simulation testing and comparisons. Finally, numerical simulation is carried out based on the benchmark instances, and the comparisons with some existing algorithms demonstrate the effectiveness of the proposed algorithm.

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1. Introduction

The multidimensional knapsack problem (MKP) is a wellknown NP-hard problem with wide engineering backgrounds. Many practical problems can be formulated as the MKP, such as cargo loading, cutting stock problems, resource allocation in computer systems, and economics problems (Martello & Toth, 1990). The study of the MKP has significant importance both academic and engineering fields. Early research on the MKP mainly focused on the exact solution methods, such as branch and bound method (Shih, 1979), dynamic programing algorithm (Toth, 1980), etc. However, the computational effort and memory requirement are tremendous when using the exact solution methods, so that they are insufficient in solving the large-scale problems. In recently years, many heuristic and meta-heuristic algorithms for solving the MKP have been proposed. Chu and Beasley (1998) proposed a heuristic based on genetic algorithm (GA), which required only a modest amount of computational effort. Vasquez and Vimont (2005) obtained some better results by combining linear programing (LP) with an efficient tabu search (TS). Kong, Tian, and Kao (2008) proposed a new ant colony optimization (ACO) approach to solve the MKP, called binary ant system (BAS). Very recently, Ke, Feng, Ren, and Wei (2010) developed an algorithm based on ACO called DMMAS to solve the MKP as an extension of max-min ant system that imposed lower and upper trail limits on pheromone values to avoid stagnation.

As a newly proposed population-based algorithm, Estimation of distribution algorithm (EDA) (Larranaga & Lozano, 2002) has gained an increasing study and wide applications during recent years. Population-based incremental learning (PBIL) (Baluja, 1994) is the earliest model of the EDA. According to the complexity of the model, the EDA can be classified as univariate model, bivariate model or multivariate model. The PBIL, univariate marginal distribution algorithm (UMDA) (MÄuhlenbein & Paass, 1996) and compact GA (CGA) (Harik, Lobo, & Goldberg, 1998) are univariate models, while mutual information maximization for input clustering (MIMIC) (De Bonet, Isbell, & Viola, 1997), combining optimizers with mutual information trees (COMIT) (Baluja & Davies, 1997) and bivariate marginal distribution algorithm (BMDA) (Pelikan & Mühlenbein, 1999) are bivariate models. The factorized distribution algorithms (FDA) (Mühlenbein & Mahnig, 1999), extended compact GA (ECGA) (Harik, 1999) and bayesian optimization algorithm (BOA) (Pelikan, Goldberg, & Cantú-Paz, 1999) are multivariate models. For more details about the EDA, please refer Larranaga and Lozano (2002).

So far the EDA has been applied to a variety of academic and engineering optimization problems, such as feature selection, cancer classification, quadratic assignment problem, machinery structure design, nurse rostering, and etc (Zhou & Sun, 2007). However, to the best of our knowledge, there is almost no research work about the EDA for solving the MKP except the relaxed complemental EDA for the MKP (Yang, Ouyang, & Quan, 2007). In this paper, we take the characteristic of the MKP to propose a hybrid EDA (HEDA) for solving the MKP. A probability model is built and an updating mechanism is proposed. In addition, a mechanism of

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initializing the probabilistic model is proposed based on the specific knowledge of the MKP in the framework of the EDA. Moreover, the influence of parameter setting and the importance of repair operator are both investigated. Finally, we use a set of benchmark instances to test the performances of the HEDA and to compare it with some existing methods to further demonstrate the effectiveness of the HEDA.

The remainder of the paper is organized as follows: In Section 2, the MKP is formulated. In Section 3, the basic EDA is introduced briefly. Then, the hybrid EDA for the MKP is proposed in Section 4. The influences of parameter setting and repair operator are investigated in Section 5, and computational results and comparisons are also provided. Finally we end the paper with some conclusions in Section 6.

2. Multidimensional knapsack problem

The multidimensional knapsack problem can be formulated as follows:

$$\text{Maximize } \sum_{i=1}^{n} p_{j} x_{j}$$
 (1)

Subject to
$$\sum_{i=1}^{n} r_{ij}x_{j} \leq b_{i}, \quad i = 1, 2, \dots, m,$$
 (2)

$$x_i \in \{0,1\}, \quad j = 1, 2, \dots, n$$
 (3)

where n is the number of items, and m is the number of knapsack constraints with capacities b_i (i = 1, 2, ..., m), associated weights r_{ij} and profits p_i .

The objective is to find a subset of items that yields a maximum profit without exceeding the resource capacities. Without loss of generality, it can be assumed that all the data r_{ij} , b_i and p_j are non-negative integers and $r_{ij} \leq b_i$, $\sum_{j=1}^n r_{ij} > b_i$ for all i = 1, 2, ..., m and j = 1, 2, ..., n.

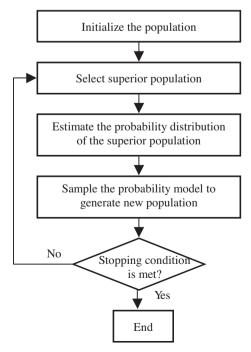


Fig. 1. The general framework of the EDA.

3. Estimation of distribution algorithm

Estimation of distribution algorithms (EDA) is a new paradigm in the field of Evolutionary Computation, which employs explicit probability distributions in optimization (Larranaga & Lozano, 2002). Compared with the genetic algorithms, the EDA reproduces new population implicitly instead of the crossover and mutation operators. In the EDA, a probability model of the most promising area is built by statistical information based on the searching experience, and then the probability model is used for sampling to generate the new individuals. Meanwhile, the probability model is updated in each generation with the potential individuals of the new population. In such an iterative way, the population evolves and finally satisfactory solutions can be obtained.

The general framework of the EDA is illustrated in Fig. 1.

The critical step of the above procedure is to estimate the probability distribution. The EDA makes use of the probability model to describe the distribution of the solution space and the updating process reflects the evolutionary trend of the population. Due to the difference of problem types, a proper probability model and the updating mechanism should be well developed to estimate the underlying probability distribution. Nevertheless, the EDA pays more attention to global exploration while its exploitation capability is limited. So, an effective EDA should balance the exploration and the exploitation abilities.

4. Hybrid EDA for MKP

In this section, we will propose a hybrid EDA (HEDA) for solving the MKP. First, we will introduce the encoding scheme, probability model, update mechanism, repair operator and local search strategy. Then, we will present the framework of proposed HEDA.

4.1. Encoding scheme

Every individual of the population is a solution of the MKP, which is represented by an n-bit binary string. The encoding scheme of the individual is illustrated in Fig. 2, where each bit $x_i = 1$ or 0.

4.2. Probability model and updating mechanism

The HEDA uses an n-dimensional vector $p(X) = [p(x_1), p(x_2), ..., p(x_n)]$ to represent the probability model of the solution space, where $p(x_j)$ is the probability of $x_j = 1$. Generally, the vector p(X) is initialized as p(X) = [0.5, 0.5, ..., 0.5], which ensures that the whole solution space can be sampled uniformly. To improve the converging speed of the HEDA, we propose an initialization mechanism based on the specific knowledge of the MKP.

During the experiment, we find empirically that bias sampling towards the optimal solution can increase the probability to obtain good solutions. Thus, we choose a solution $S = (s_1, s_2, ..., s_n)$ with a similar construction to the optimal solution to guide the initialization of the probability vector as follows:

$$p(x_j) = \begin{cases} 0.5 + \beta, s_j = 1\\ 0.5 - \beta, s_j = 0 \end{cases}, \quad j = 1, 2, \dots, n$$
 (4)

where $\beta \in (0,0.5)$ represents the level of bias to the solution S that can be obtained by a simple procedure.

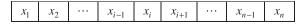


Fig. 2. Representation of a solution in HEDA.

Let $S' = (s'_1, s'_2, \dots, s'_n)$ be the solution of the linear programing (LP) relaxation problem (see Section 4.3) of the MKP. Then, S is constructed as follows:

$$s_j = \begin{cases} 1, s'_j = 1 \\ 0, else \end{cases}, \quad j = 1, 2, \dots, n$$
 (5)

In each generation of the HDEA, the new individuals are generated via sampling according to the probability vector p(X). For every x_j to be determined, a random number $r \in (0,1)$ is generated. Let $x_j = 1$ if $r < p(x_j)$, otherwise let $x_j = 0$. Once all the items are determined, an individual is constructed. In such a way, P individuals are constructed and their profits are calculated.

Then, the superior population that consists of the best N solutions is determined, and the probability vector p(X) is updated according to the following equation:

$$p_{l+1}(X) = (1 - \alpha)p_l(X) + \alpha \frac{1}{N} \sum_{k=1}^{N} X_l^k$$
 (6)

where $p_l(X)$ is the probability vector, and X_l^k is the kth individual of the superior population in the lth generation, and $\alpha \in (0,1)$ is the learning speed.

4.3. Repair operator

The above way to generate new solutions does not take into account the feasibility of the solutions that are sampled based on the probability model. The individuals could be illegal for violating some of the resource constraints. So, a repair procedure is needed if illegal individuals are constructed. In this sub-section, we provide two repair operators as follows:

Repair Operator 1 (R01): The general technique is a greedy-like heuristic based on the pseudo-utility ratios that are calculated by the surrogate duality approach (Pirkul, 1987). The general idea of this method is described briefly as follows:

The surrogate relaxation problem of the MKP can be defined as follows:

$$\text{Maximize } \sum_{j=1}^{n} p_j x_j \tag{7}$$

Subject to
$$\sum_{i=1}^{n} (\sum_{i=1}^{m} \omega_i r_{ij}) x_j \leq \sum_{i=1}^{m} \omega_i b_i,$$
 (8)

$$x_i \in \{0, 1\}, \quad j = 1, 2, \dots, n$$
 (9)

where $\omega = \{\omega_1, \omega_2, ..., \omega_m\}$ is a set of surrogate multipliers (or weights) of some positive real numbers. These weights can be obtained by solving the LP relaxation of the original MKP and the values of the dual variables are viewed as the weights. In other words, the weight $\omega_i(i=1,2,...,m)$ can be regarded as the shadow price of the i-th constraint in the LP relaxation of the MKP.

Having obtained the weights, the pseudo-utility ratio for each variable can be computed as follows:

$$u_j = \frac{p_j}{\sum_{i=1}^m \omega_i r_{ij}} \tag{10}$$

Before performing the repair procedure, the variables are ordered in the decreasing order according to their pseudo-utility ratios. The repair operator consists of two phases. The first phase examines each bit of the solution string in the increasing order of u_j and changes a bit from one to zero until the solution is legal. The second phase reverses the process by examining each bit in the decreasing order of u_j and changes a bit from zero to one as long as the feasibility is not violated. The pseudo-code of this repair operator is illustrated in Fig. 3.

```
Procedure Repair Operator 1

Let: I = \{1, 2, \dots, m\}.

1: Initialize R_i = \sum_{j=1}^n r_{ij}x_j, \forall i \in I

2: FOR j = n to 1 DO

3: IF (x_j = 1) and (R_i > b_i, \exists i \in I) THEN

4: Set x_j \leftarrow 0;

5: Set R_i \leftarrow R_i - r_{ij}, \forall i \in I;

6: END IF

7: END FOR

8: FOR j = 1 to n DO

9: IF (x_j = 0) and (R_i + r_{ij} \le b_i, \forall i \in I) THEN

10: Set x_j \leftarrow 1;

11: Set R_i \leftarrow R_i + r_{ij}, \forall i \in I.

12: END IF

13: END FOR
```

Fig. 3. Procedure of RO1.

Repair Operator 2 (RO2): The above repair operator requires solving the LP relaxation of the MKP. So far, most mathematical tools are only able to solve the LP problems with small or medium scale, and it is difficult to implement that operator when the problem scale is large. Thus, we propose a new repair operator based on the specific knowledge of the MKP. With this technique, there is no need to solve the LP relaxation of the MKP in advance.

During the initialization phase, a simple matrix Q and an assistant matrix F are obtained. The element q_{ij} in Q is computed as follows:

$$q_{ij} = \frac{p_j}{r_{ii}} \tag{11}$$

Each row of the assistant matrix F is an arrangement of all the items, and f_{ij} represents the number of the jth item in the descending order of the elements in the ith row of matrix Q. To be specific, the pseudo-code is illustrated in Fig. 4.

Actually, q_{ij} is the jth item's pseudo-utility ratio of the ith constraint. For an illegal solution, this algorithm examines every constraint in the ascending order of q_{ij} and changes one bit from one to

```
Procedure Repair Operator 2
Let: I = \{1, 2, \dots, m\}.
 1: Initialize R_i = \sum_{j=1}^n r_{ij} x_j, \forall i \in I
 2: FOR i = 1 to m DO
 3:
        FOR j = n to 1 DO
 4:
            IF (x_{f_{ij}} = 1) and (R_i > b_i) THEN
               Set x_{f_{ij}} \leftarrow 0;
               Set R_i \leftarrow R_i - r_{i, f_{ii}}, \forall i \in I;
 7:
         END FOR
 9: END FOR
10: Find the constraint i that has the minimum of b_i - R_i.
        IF (x_{f_{ij}} = 0) and (R_i + r_{i, f_{ij}} \le b_i, \forall i \in I) THEN
            Set x_{f_{ij}} \leftarrow 1;
            Set R_i \leftarrow R_i + r_{i, f_{ii}}, \forall i \in I.
14:
15:
         END IF
16: END FOR
```

Fig. 4. Procedure of RO2.

zero until that constraint is satisfied. After that, it balances the remaining resources according to the minimum remaining resource. Clearly, this new repair operator is simple to implement, and especially it is not limited to the scale of problem.

4.4. Local search strategy

It is widely accepted that a local search procedure is efficient in improving the solutions generated by the EDA. In the HEDA, we use a simple method as the local search strategy.

In particular, five bits are randomly selected from the solution string, and they are changed form zero to one or from one to zero. When changing the last zero-valued bit to one, if the profit of the current solution is less than that of the best solution obtained so far, the new individual would not be constructed. Then, the new solution is repaired if it is illegal. Such a procedure is applied on the best individual of the current population *Y* times in every generation, where *Y* represents the intensity of local search.

Because the solutions obtained in the earlier generations are often of low quality, there is no need to use the local search too intensively. Thus, we design a special adaptive way to control *Y* to enhance the exploitation ability as follows:

$$Y = \begin{cases} 500, & generation < 50 \\ 800, & 50 \le generation < 100 \\ 1000, & generation \ge 100 \end{cases}$$
 (12)

4.5. Procedure of HEDA

With the design above, the procedure of the HEDA is illustrated in Fig. 5.

In summary, in the initial stage of evolution, the promising area of the solution space may be found by using the EDA based estimating and sampling. Then, local search strategy is performed in

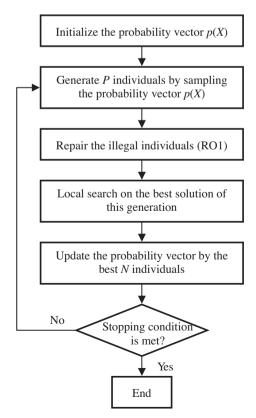


Fig. 5. The framework of the HEDA.

the "good" region to obtain better solutions. The benefits of the EDA and the local search are combined to balance global exploration and local exploitation. Moreover, the repair operator is applied on the illegal solutions if necessary.

5. Computational results and comparisons

We choose three different sets of widely used benchmark instances to test the HEDA. Test set 1 consists of small size problems (available from ORLIB) with m=2 to 30 and n=15 to 105. Test set 2 consists of 11 large size problems (available from http://hces. bus.olemiss.edu/tools.html) with m=15 to 100 and n=100 to 2500. Test set 3 consists of 20 medium size problems (available from ORLIB) with m=5 or 10 and n=100. The algorithm is coded in C++ and run on a 3.2 GHz Intel Core i5 processor.

5.1. Parameters setting

The proposed HEDA contains several key parameters: the population size of each generation (P), the number of selected individual to update the probability model (N), the learning speed (α) and the bias level of initialization (β). To investigate the influence of those parameters on the performance of the algorithm, we implement the Taguchi method of design of experiment (DOE) (Montgomery, 2005) by using the problem 100.10.05. Combinations of different values of these parameters are shown in Table 1.

For each parameter combination, the HEDA is run 50 times independently with the maximum number 10,000 of the constructed solutions as the stopping criterion. The average response variable (ARV) value is the average of profit obtained by the HEDA in 50 times. According to the number of parameters and the number of factor levels, we choose the orthogonal array $L_{16}(4^4)$. That is, the total number of treatment is 16, the number of parameters is 4, and the number of factor levels is 4. The orthogonal array and the obtained ARV values are listed in Table 2.

Table 1Combinations of parameter values.

Parameters	Factor level			
	1	2	3	4
P	20	30	40	50
N (%P)	10	20	30	40
α	0.001	0.01	0.1	0.5
β	0	0.05	0.1	0.2

Table 2Orthogonal array and ARV values.

Experiment number	Facto	Factor			
	P	N	α	β	
1	1	1	1	1	22683.5
2	1	2	2	2	22687.4
3	1	3	3	3	22688.2
4	1	4	4	4	22686.7
5	2	1	2	3	22684.9
6	2	2	1	4	22678.4
7	2	3	4	1	22682.8
8	2	4	3	2	22681.5
9	3	1	3	4	22671.6
10	3	2	4	3	22679.8
11	3	3	1	2	22676.9
12	3	4	2	1	22676.9
13	4	1	4	2	22669.7
14	4	2	3	1	22668.4
15	4	3	2	4	22672.4
16	4	4	1	3	22672.0

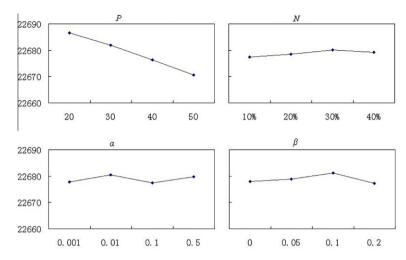


Fig. 6. Factor level trend of the HEDA.

Table 3 Response value.

Level	P	N	α	β
1	22686.5	22677.4	22677.7	22677.9
2	22681.9	22678.5	22680.4	22678.9
3	22676.3	22680.1	22677.4	22681.2
4	22670.6	22679.3	22679.7	22677.3
Delta	15.9	2.7	3.0	3.9
Rank	1	4	3	2

According to the orthogonal table, we illustrate the trend of each factor level in Fig. 6. Then, we figure out the response value of each parameter to analyze the significance rank of each parameter. The results are listed in Table 3.

From Table 3 it can be seen that P is the most significant parameter among the four parameters. In condition of a fixed maximum number of constructed solutions, small population size allows more generations so that it may provide deeper search. In addition, the significant rank of β is the second and the performance of the HEDA is better when β is more than zero, which demonstrates that the new initialization mechanism of the probability model is efficient. A better choose of the parameter combination is: P = 20, N = 6, $\alpha = 0.01$, and $\beta = 0.1$.

5.2. Influence of repair operator

Next, we study the influence of the repair operator in the HEDA. Denote HEDA1 as the HEDA with RO1 and HEDA2 as the HEDA with RO2. We run both HEDA1 and HEDA2 20 times independently by using test set 2. Since test set 2 has medium or large scale, we set P = 200, N = 10, $\alpha = 0.01$, $\beta = 0.1$, and set the maximum number 100,000 of the constructed solutions as the stopping criterion.

The statistical results are summarized in Table 4, where Min.Dev and Ave.Dev represent the minimum and average percentage deviations from the best-known solutions, and Var. Dev denotes the variance of the deviations.

From Table 4, it can be seen that the performance of HEDA1 is better than that of HEDA2 on the first 8 problems. However, since RO1 depends on solving the LP relaxation problem of the MKP, it is very difficult to solve the problem with LP when the number of m+n is very large. So, HEDA2 can solve another 3 problems with larger scale while HEDA1 cannot solve those problems. The conclusion is that the repair operator is important to improve the performance of the algorithm. Comparing RO2, RO1 is superior to achieve better results but it can only applied to the problems with small or medium size.

5.3. Comparisons of HEDA with RCEDA

The relaxed complemental estimation of distribution algorithm (RCEDA) is the only existing EDA for solving the MKP (Yang et al., 2007). In this sub-section, we compare the HEDA with the RCEDA by using test set 1. For the HEDA, we set P=20, N=6, $\alpha=0.01$, and $\beta=0.1$. Both the algorithms are run 100 times independently. The maximum numbers of solutions constructed in the HEDA and the RCEDA are 5,000 and 50,000, respectively. That is, the HEDA uses less computational effort. The statistical results are listed in Table 5,

Table 4 Comparisons of HEDA1 with HEDA2.

Problem $n \times m$	Best known	HEDA1			HEDA2			
		Min. Dev	Ave. Dev	Var. Dev	Min. Dev	Ave. Dev	Var. Dev	
Mk_gk01	100 × 15	3766	0.2655	0.3173	0.0783	0.6107	0.9360	0.6555
Mk_gk02	100×25	3958	0.0758	0.1983	0.1194	0.7580	0.9790	0.4039
Mk_gk03	150×25	5650	0.1239	0.2018	0.0998	0.9381	1.1531	0.7828
Mk_gk04	150×50	5764	0.0867	0.3348	0.2916	1.0930	1.2673	0.7069
Mk_gk05	200×25	7557	0.1323	0.2772	0.2150	1.2439	1.4960	0.7483
Mk_gk06	200×50	7672	0.2998	0.4438	0.2339	1.1210	1.4436	1.1573
Mk_gk07	500×25	19215	0.5204	0.7203	0.8240	1.7018	1.8460	0.9873
Mk_gk08	500×50	18801	0.6117	0.7680	1.1028	1.6595	1.7260	0.2880
Mk_gk09	1500×25	58085				2.4585	2.5461	0.6813
Mk_gk10	1500×50	57292	-	-	-	2.0177	2.1170	0.7807
Mk_gk11	2500×100	95231	-	-	-	1.7043	1.7348	0.2184

Note: the bold values mean better results.

Table 5Comparisons of HEDA with RCEDA.

Problem	$n \times m$	Best known	RCEDA		HEDA	
			AVG SR		AVG	SR
Sento1	60 × 30	7772	7771.5	95	7772	100
Sento2	60×30	8722	8720.9	93	8772	100
Weing7	105×2	1095445	1095422.6	90	1095445	100
Weing8	105×2	624319	624316.8	80	624319	100
Knap15	15×10	4015	4012.7	95	4015	100
Knap20	20×10	6120	6117.9	97	6120	100
Knap39	39×5	10618	10615.3	95	10618	100
Knap28	28×10	12400	12389.4	98	12400	100

Note: the bold values mean better results.

Table 6Comparisons of HEDA with DMMAS.

Problem	Best known	DMMAS			HEDA		
		Best	AVG	STD	Best	AVG	STD
5.100.00	24381	24381	24362	23.8	24381	24381	0.0
5.100.01	24274	24274	24273	6.2	24274	24274	0.0
5.100.02	23551	23551	23540	7.2	23551	22541	5.7
5.100.03	23534	23534	23482	14.9	23534	23524	11.0
5.100.04	23991	23991	23954	10.8	23991	23976	16.3
5.100.05	24613	24613	24608	6.3	24613	24612	5.7
5.100.06	25591	25591	25591	0.0	25591	25591	0.0
5.100.07	23410	23410	23404	13.3	23410	23375	7.1
5.100.08	24216	24216	24211	5.9	24216	24211	5.9
5.100.09	24411	24411	24406	13.8	24411	24411	0.0
10.100.00	23064	23064	23045	19.6	23064	23051	2.4
10.100.01	22801	22801	22742	40.8	22801	22739	25.2
10.100.02	22131	22131	22091	29.8	22131	22131	0.0
10.100.03	22772	22772	22710	37.6	22772	22772	0.0
10.100.04	22751	22751	22617	43.9	22751	22620	24.7
10.100.05	22777	22777	22663	40.3	22777	22683	30.0
10.100.06	21875	21875	21826	28.4	21875	21823	9.0
10.100.07	22635	22635	22557	31.0	22635	22552	27.7
10.100.08	22511	22438	22409	17.3	22511	22424	5.0
10.100.09	22702	22702	22696	32.7	22702	22702	0.0

Note: the bold values mean better results.

where SR denotes the successful ratio of the run that obtains the best-known value among the 100 independent runs.

From Table 5, it can be seen that the successful ratio and average profit obtained by the HEDA are better than those obtained by the RCEDA, even the maximum number of constructed solutions by the HEDA is only one tenth of that by the RCEDA. So, the conclusion is that our HEDA is superior to the existing EDA based algorithm for solving the MKP. Note that, it has been shown that the performance of the RCEDA is better than the hybrid GA (Yang et al., 2007). So, the performance of our HEDA is definitely better than that of the hybrid GA in the literature.

5.4. Comparisons of HEDA with DMMAS

Finally, we compare the HEDA with the DMMAS by using test set 3. For the HEDA, we set P = 20, N = 6, $\alpha = 0.01$, and $\beta = 0.1$. We set the maximum numbers of constructed solution 10,000 as the stopping criterion for the two algorithms. For every instance, the two algorithms are run 50 times independently. The statistical results are summarized in Table 6, including the best profit, average profit and the standard derivation.

Form Table 6, it can be seen that the HEDA can find all the best solutions of the twenty instances while the DMMAS can find the best solutions of nineteen instances. As for the average profit, the results of the HEDA are better than those of the DMMAS in 14 instances, the results of the HEDA are worse than those of the DMMAS in only 4 instances, and the results are the same in 2

instances. Especially, the HEDA can find the best-known solutions consistently in 7 instances. As for the STD, the results of the HEDA are better than that of the DMMAS in most instances. So, the conclusion is that our HEDA is superior to the DMMAS in solving the MKP.

6. Conclusions

In this paper, a hybrid estimation of distribution algorithm was proposed for solving the MKP. We designed a probability model with the superior population for the EDA to solve the MKP by generating the new individuals via sampling based on the probability model. We also designed an updating mechanism for the probability model and an initialization mechanism for probability vector to increase the convergence speed. To avoid the difficulty of the existing repair operator in solving large-scale problems, we presented a special repair operator. Moreover, an adaptive local search was designed to enhance the exploitation of the algorithm. The influence of parameter setting was investigated by using DOE based testing, and the two repair operators were also compared through simulation testing. We also compared our HEDA with the existing RCEDA and DMMAS, which showed that the HEDA was more effective in solving the MKP. The future work is to design more effective EDA for the MKP including the model and the repair operator as well as the local search and to study the EDA for solving other kinds of combinatorial optimization problems, such as scheduling problems.

Acknowledgments

This research is partially supported by National Science Foundation of China (61174189 and 60834004), Doctoral Program Foundation of Institutions of Higher Education of China (20100002110014), the National Key Basic Research and Development Program of China (No. 2009CB320602) and National Science and Technology Major Project of China (No. 2011ZX02504-008).

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