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b)  $a^n b^m c^m d^{2n}$  with  $n, m$  any natural numbers

$S \rightarrow aSdd \mid F \mid \lambda$

$F \rightarrow bFc \mid \lambda$

Explanation:  $aSdd$  ensures that there are twice as many  $d$  as  $a$  and  $bFc$  ensures that the number of  $b$  and  $c$  is the same.

c)  $a^n b^m$  with  $0 \leq n \leq m \leq 2n$

$S \rightarrow aSb \mid aSbb \mid \lambda$

Explanation:  $aSbb$  ensures that the maximum number of  $b$  is always twice the number of  $a$  but  $aSb$  allows for any number of  $b$  less than  $2n$ .

d)  $a^m b^n c^k$  where  $(m=n)$  or  $(m=k)$

$S \rightarrow T \mid U$

$T \rightarrow aTbC \mid \lambda$

$C \rightarrow cC \mid \lambda$

$U \rightarrow aUc \mid B \mid \lambda$

$B \rightarrow bB \mid \lambda$

Explanation: This grammar is the Union of  $a^m b^m$  and  $a^m c^m$ .

e)  $\{a^i b^j c^k d^k \mid i, k \geq 0\}$

$S \rightarrow TU \mid \lambda$

$T \rightarrow aTb$

$U \rightarrow cUb$

Explanation: This grammar is the concatenation of  $a^i b^j$  and  $c^k d^k$ .

f)  $\{a^i b^j c^k d^m \mid i, j, k, m \geq 0, \text{ and } (i = j \text{ or } k = m)\}$

$S \rightarrow F \mid Q$

$F \rightarrow TU$

$T \rightarrow aT \mid Tb \mid \lambda$

$U \rightarrow cUd \mid \lambda$

$Q \rightarrow NM$

$N \rightarrow aNb \mid \lambda$

$M \rightarrow cM \mid Md \mid \lambda$

Explanation: This grammar is the concatenation  $a^i b^j c^k d^m$  where  $i$  could equal  $j$  or  $k = m$ .

a)

$E \rightarrow E + T \mid T$

$T \rightarrow T * I \mid I$

$I \rightarrow a \mid b \mid c$

b)

$E \rightarrow E + T \mid E - T \mid T$

$T \rightarrow T * I \mid I$

$I \rightarrow a \mid b \mid c$

c)

$E \rightarrow E + T \mid E - T \mid T$

$T \rightarrow T * U \mid U$

$U \rightarrow I \uparrow U \mid I$

$I \rightarrow a \mid b \mid c$

d)

$C \rightarrow C = E \mid C < E \mid E$

$E \rightarrow E + T \mid E - T \mid T$

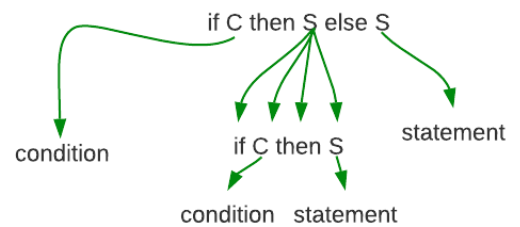
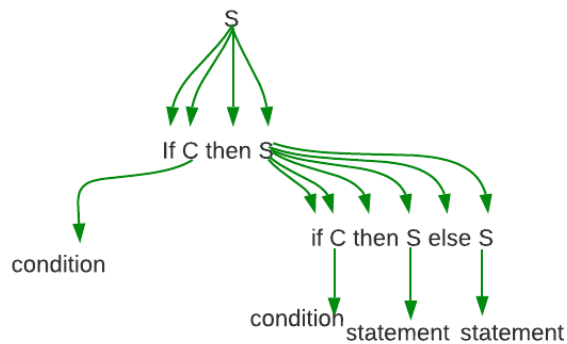
$T \rightarrow T * U \mid U$

$U \rightarrow I \uparrow U \mid I$

$I \rightarrow a \mid b \mid c$

203. PLAmb

1)



Both trees represent the string "if condition then if condition then statement else statement"

```
2) if (1 > x) {
    if (2 < x) {
        print("in second if");
    } else {
        print("In first if");
    }
}
```

}  
OR

```
If (1 > x){  
    If(2 < x){  
        print("in second if");  
    }  
} else{  
    print("in neither if");  
}
```

3)  
 $S \rightarrow \text{if } C \text{ then } S \mid \text{if } C \text{ then } T \text{ else } S \mid \text{statement}$   
 $T \rightarrow \text{if } C \text{ then } T \text{ else } T \mid \text{statement}$   
 $C \rightarrow \text{condition}$

The above grammar resolves the issue of dangling else.

209. ElimUseless

$S \rightarrow dS \mid aS \mid a$

217. GNFPPractice

$E \rightarrow 0^*T + E \mid 1^*T + E \mid 0 + E \mid 1 + E \mid$   
 $T \rightarrow 0^*T \mid 1^*T \mid 0 \mid 1$