

9) a) 1. while:

$$2^{i-1} < n$$

paso 1  $c=1$

$$i-1 < \log_2(n)$$

paso 2  $c=2.1$

$$i < \log_2(n) + 1$$

paso 3  $c=2^2.1$

paso i  $c=2^{i-1}.1$

$$T(n) = c_1 + \sum_{i=1}^{\log_2(n)} c_2 = c_1 + \log_2(n) \cdot c_2$$

b)

$T(n)$  es de  $O(\log_2(n))$

$$\left. \begin{array}{l} c_1 \leq \log_2(n) \\ c_1 \leq c_1 \cdot \log_2(n) \end{array} \right\} \begin{array}{l} c = c_1 \\ n_0 = 1 \end{array} \quad \left. \begin{array}{l} \log_2(n) \leq \log_2(n) \\ \log_2(n) \cdot c_2 \leq \log_2(n)^{c_2} \end{array} \right\} \begin{array}{l} c = c_2 \\ n_0 = 1 \end{array}$$

$$T(n) \leq \log_2(n) \cdot c_1 + \log_2(n) \cdot c_2$$

$$T(n) \leq \log_2(n) \cdot (c_1 + c_2)$$

$$T(n) \leq \log_2(n) \cdot c$$

$$T(n) \leq \log_2(n) \cdot c \text{ con } c = (c_1 + c_2), \forall n \geq n_0 \text{ con } n_0 = 1$$



a) 2.

paso 1  $c = n$

$$\frac{n}{2^{i-1}} > 1$$

paso 2  $c = \frac{n}{2}$

$$n > 2^{i-1}$$

paso 3  $c = \frac{n}{2^2}$

$$\log_2(n) > i-1$$

paso i  $c = \frac{n}{2^{i-1}}$

$$i-1 < \log_2(n)$$

$$i < \log_2(n) + 1$$

$$c_1 + \sum_{i=1}^{\log_2(n)} c_2 = c_1 + \log_2(n) \cdot c_2 \quad \text{b) tiene } O(\log_2(n))$$

$$c_1 \leq \log_2(n) \quad \left. \begin{array}{l} c = c_1 \\ n_0 = 1 \end{array} \right\}$$

$$c_1 \leq \log_2(n) \cdot c_1 \quad \left. \begin{array}{l} c = c_1 \\ n_0 = 1 \end{array} \right\}$$

$$T(n) \leq \log_2(n) \cdot (c_1 + c_2)$$

$$T(n) \leq \log_2(n) \cdot c$$

$$\log_2(n) \leq \log_2(n) \quad \left. \begin{array}{l} c = c_2 \\ n_0 = 1 \end{array} \right\}$$

$$\log_2(n) \cdot c_2 \leq \log_2(n) \cdot c_2 \quad \left. \begin{array}{l} c = c_2 \\ n_0 = 1 \end{array} \right\}$$

$$T(n) \leq O(n) \text{ con } c = (c_1 + c_2), \forall n > n_0, n_0 = 1$$