Group 14: Parallelisation of graph algorithms in Julia

Speaking Order:

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- 2. Aorthi Afroza
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Overview

Blain Cribb:

Julia

Aorthi Afroza:

- Considered algorithms
- Prim's algorithm

Kelsey Murray:

- A* Algorithm
- Julia benchmarking tools



- High-level language
- General-purpose
 - Data science
 - Machine learning
 - Fast numerical and scientific computing
 - Parallel computing
- Dynamically typed but with great performance
- Open-sourced with fantastic documentation
- "Looks like Python, feels like lisp, runs like Fortran"

Julia - Looks like Python

- Code is easy to read and write
- Designed to be easy to pick up for other language users
- Plenty of modern features that make coding easier:
 - Large base library
 - Garbage collection
 - Optional data typing
 - Interoperability

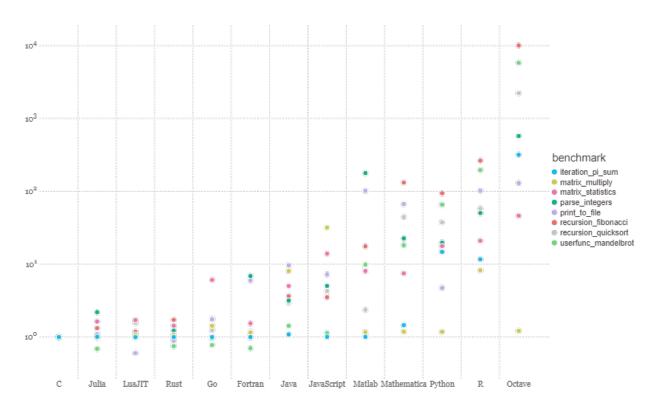
Julia - Feels like Lisp

- Support for metaprogramming
 - Code is stored as a data structure within the language
 - This code can be manipulated during runtime
 - Very useful for applications such as machine learning
- True Lisp-style macros unlike static macros
- Multiple dispatch for OOP like behaviour

Julia - Runs like Fortran

- Created to be fast from its inception
- Static analysis and just in time compilation allows for running of some code in a static manner
- Optional static typing means with smart programming we can achieve static-like speeds
- Multiple dispatch can help with the static analysis

Julia - Runs like Fortran



Julia - Parallelism

- Natively supported
- Multithreading (experimental)
- Loops are easily parallelisable through use of macros
- Supports atomic access for variables

Considered Algorithms

- Criteria for algorithms:
 - Different algorithm types; e.g. shortest path, traversal, minimum spanning tree
 - Algorithms with a considerable amount of literature out there
 - Successful existing implementations of their parallelised algorithms

Julia - LightGraphs.jl

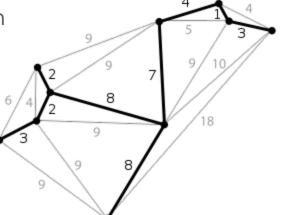
- A Julia Package that provides the framework for:
 - Building graphs
 - Traversing them
 - Building your own graphing algorithms
- Contains the module Parallel, which contains parallel implementations of:
 - Bellman Ford Shortest Paths
 - Dijkstra Shortest Paths
 - Floyd Warshall Shortest Paths
 - Johnson Shortest Paths

- o Bfs
- Greedy Color
- 5 Centrality measures
- 4 Distance measures

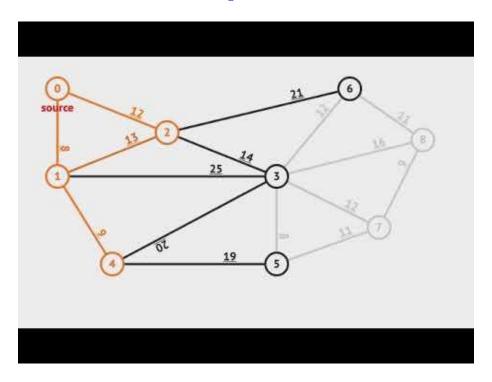
Prim's Algorithm

- Minimum spanning tree (MST)
 - A tree that connects all nodes in the graph
 - The least total cost among all trees that connect all the nodes

This may not be the shortest path

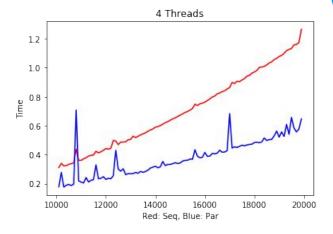


Prim's Algorithm: Example



Previous Attempts at Parallelising Prim's

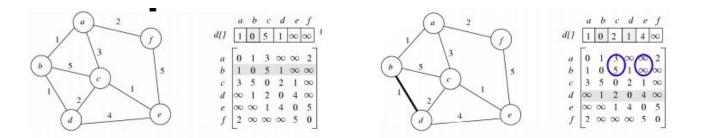
- PES University in India
- University of California, Irvine
 - Implemented in OpenMP



https://github.com/parthvshah/parallel-prims

- University of Singapore
 - Implemented in C++
 - This paper proposes an adaptation of Prim's Algorithm

Implementation Plan: Prim's



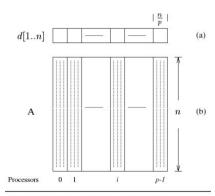
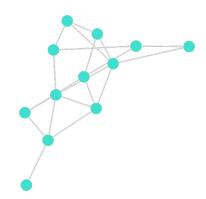
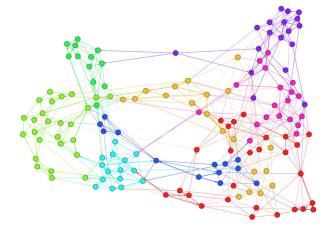


Figure 7.6 The partitioning of the distance array d and the adjacency matrix A among p processors. Copyright (r) 1994 Benjamin/Cummings Publishing Co.

Testing and Benchmarking: Prim's

- Datasets:
 - Densely connected, large, undirected graphs
- Metrics
 - Speed-up
 - Correctness





A* Algorithm

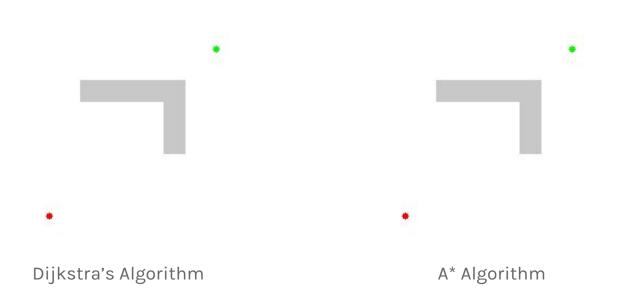
Pathfinding algorithm

- Popular in gaming where there's one source, one destination, and possibly obstacles
- o "Informed" Dijkstra's algorithm
- Cost of a considered path takes into account heuristic as well as edge weight

Best-first search, but not greedy

 Doesn't solely rely on what seems best at that time (the heuristic) but also includes the exact cost of the potential path (the edge weights)

A* Algorithm: Example

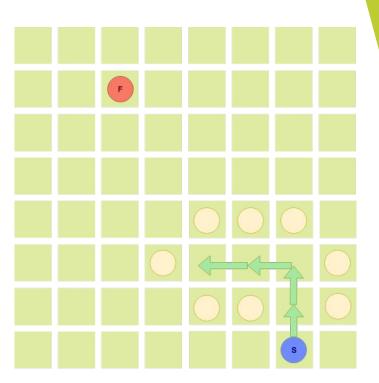


Parallel A* in Literature

First Approach:

- Looks at entire current best path
- Investigates all the nodes surrounding the path



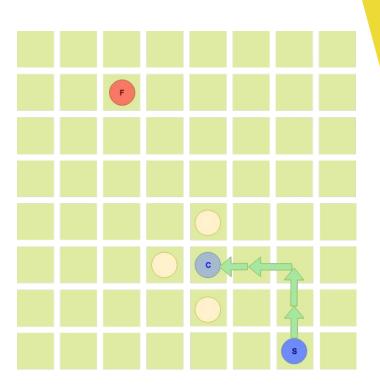


Parallel A* in Literature

Second Approach:

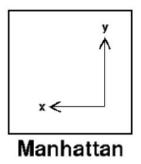
- Looks at the current end node of the current best path
- Investigates nodes
 directly adjacent to that
 end node
 start node

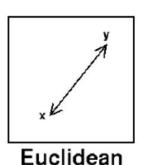


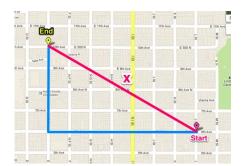


Implementation Plan: A*

- Can we move diagonally? Will this make a difference to the speedup?
 - If we can move diagonally: Euclidean distance as heuristic
 - If we can't: Manhattan distance as heuristic

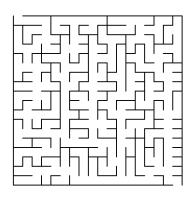


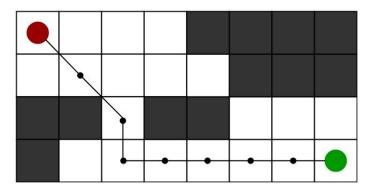




Testing and Benchmarking: A*

- Metrics:
 - Speedup
 - Correctness: Do we care if the exact path is the same?
 - Ideally parallel is at least as good as serial but we don't care if the path isn't exactly the same
- Datasets:
 - Large grid maps with obstacles





Julia Benchmarking Tools

@time macro

Output:

```
0.244729 seconds (294.16 k allocations: 14.614 MiB, 5.74% gc time)
```

@benchmark macro

Output:

```
BenchmarkTools.Trial:
   memory estimate: 2.13 KiB
   allocs estimate: 19
   -----
   minimum time: 1.770 μs (0.00% GC)
   median time: 2.170 μs (0.00% GC)
   mean time: 3.924 μs (37.21% GC)
   maximum time: 9.772 ms (99.92% GC)
   ------
   samples: 10000
   evals/sample: 10
```

Thank you!

