

# Module 1: AI for Channel Decoding

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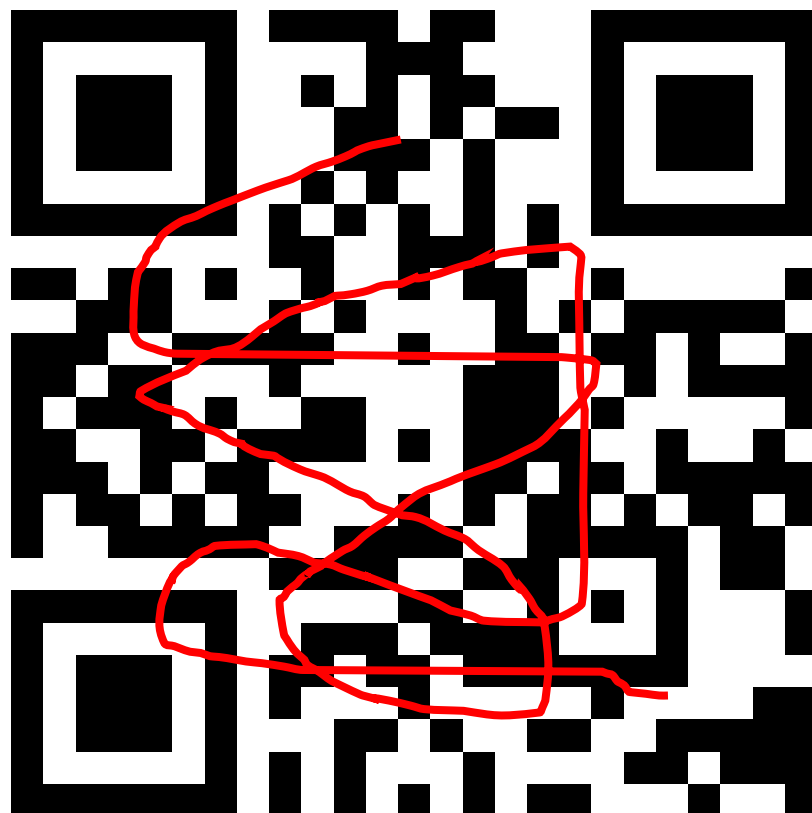
- You will learn:
  - How to use colab
  - Implement the whole communication system
  - Basics of channel coding
  - Support vector machine
  - Deep learning
- Grading:
  - Uncoded system 10%
  - Syndrome Decoding, 10% ML decoding 10%
  - Classification with support vector machine 20%
  - Deep Learning 20%
  - Mini project 20%
  - Report (no more than 10 pages) 10%

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# Error Correction Codes

# Terminology

- Message: Sequence of bits representing your data (link to the website)
- Codeword: Sequence of bits forming your QR code

- Example:

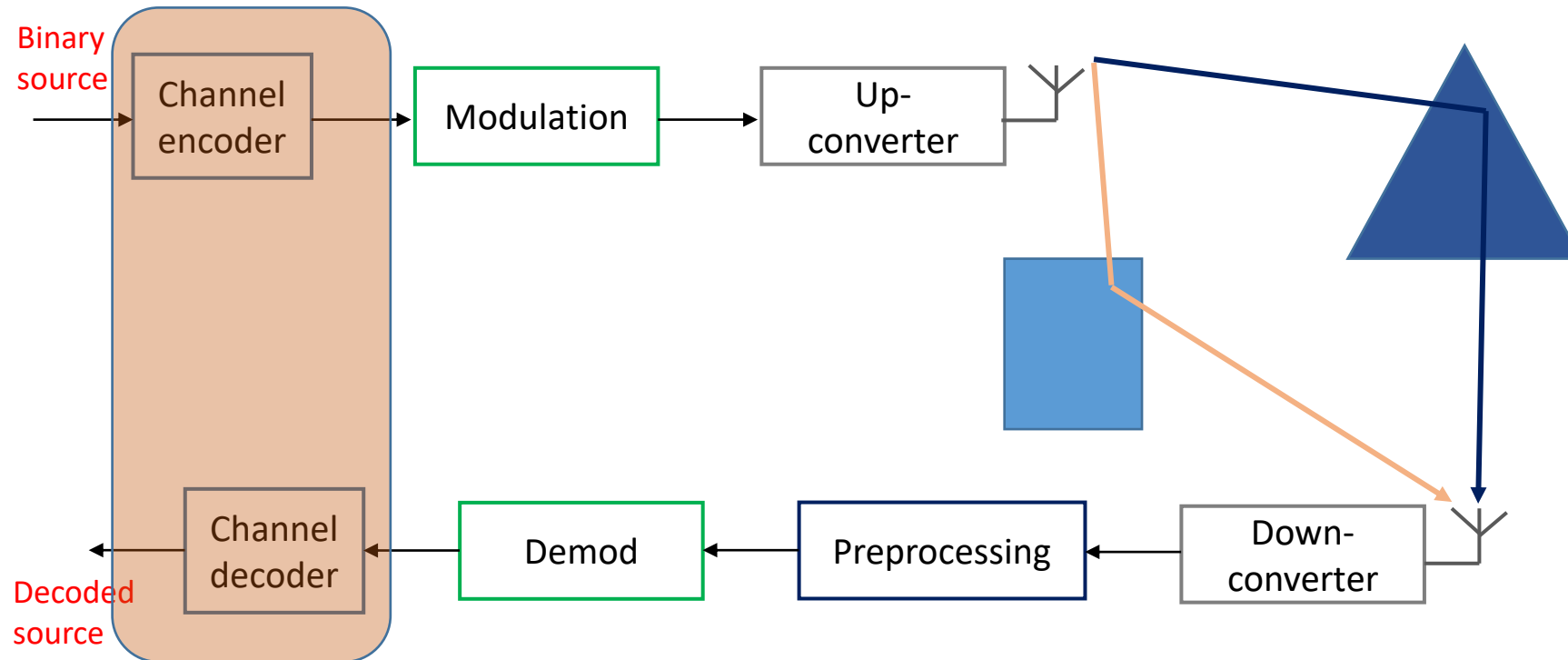
- Message = 1 1 0 1

- Codeword = 1 1 0 1 0 1 0

← These bits are added for error correction

It is everywhere!!! Even in string theory!!!

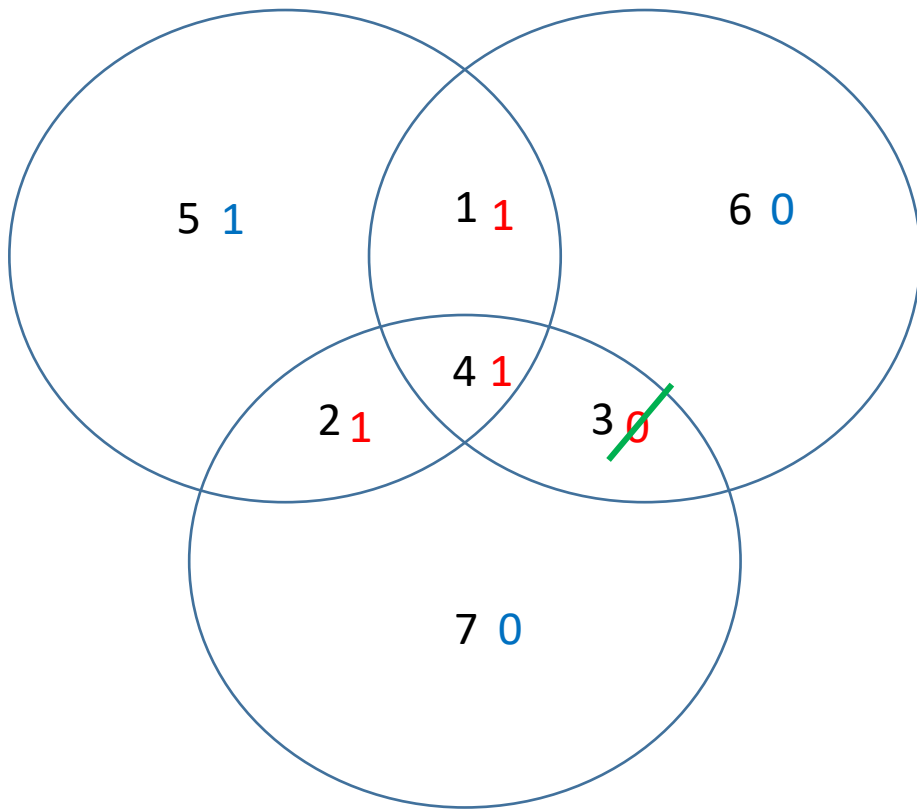
# A digital communication system over wireless channel



- Channel enc/dec allows error correction by adding redundancy

How it works? How to create those bits?

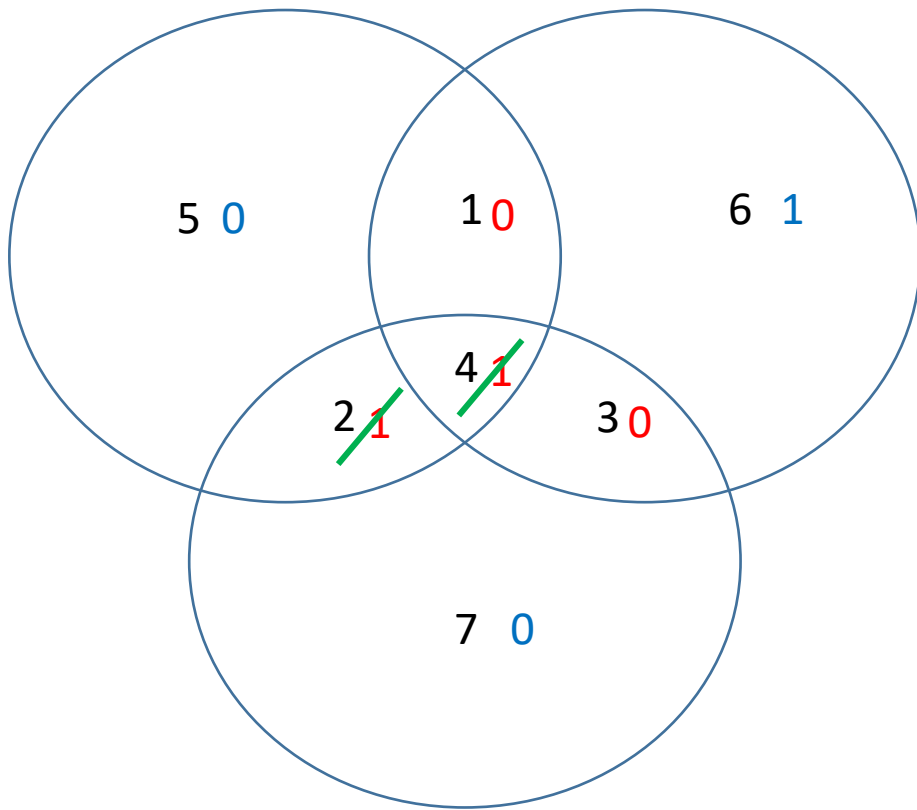
Index: 1 2 3 4 5 6 7  
Codeword: 1 1 0 1 1 0 0



The last 3 bits are added such that  
inside every circle the number of 1s are even

Suppose 1 bit is erased, we can fix it by  
checking the parity of each circle

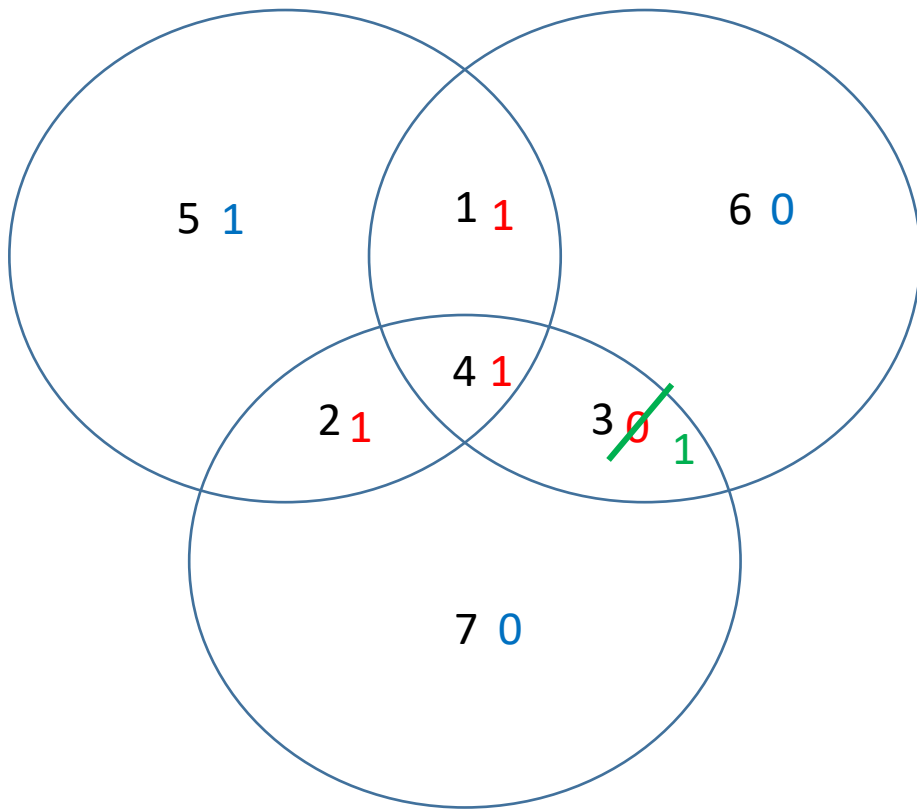
Index: 1 2 3 4 5 6 7  
Codeword: 0 1 0 1 ? ? ?



It's your turn

What if 2 bits are erased?

Index: 1 2 3 4 5 6 7  
Codeword: 1 1 0 1 1 0 0



It can also correct 1 bit flip

Suppose 1 bit is flipped, we can fix it by flipping 1 bit to meet all constraints



# $(n, k)$ -Linear Block Codes

- $k$ -bit message  $\mathbf{m}$ ,  $n$ -bit codeword  $\mathbf{c}$
- Relationship:  $\mathbf{c} = \mathbf{mG}$
- The code  $C$  contains all ( $2^k$  in total) such codewords
  - $C$  is the row space of  $\mathbf{G}$  ( $k \times n$ )
  - Call it a generator matrix
- There exists  $\mathbf{H}$  ( $n - k \times n$ ) such that  $\mathbf{cH}^T = \mathbf{Hc}^T = \mathbf{0}$ 
  - Rows of  $\mathbf{H}$  span the nullspace of  $\mathbf{G}$
  - Call it a parity check matrix

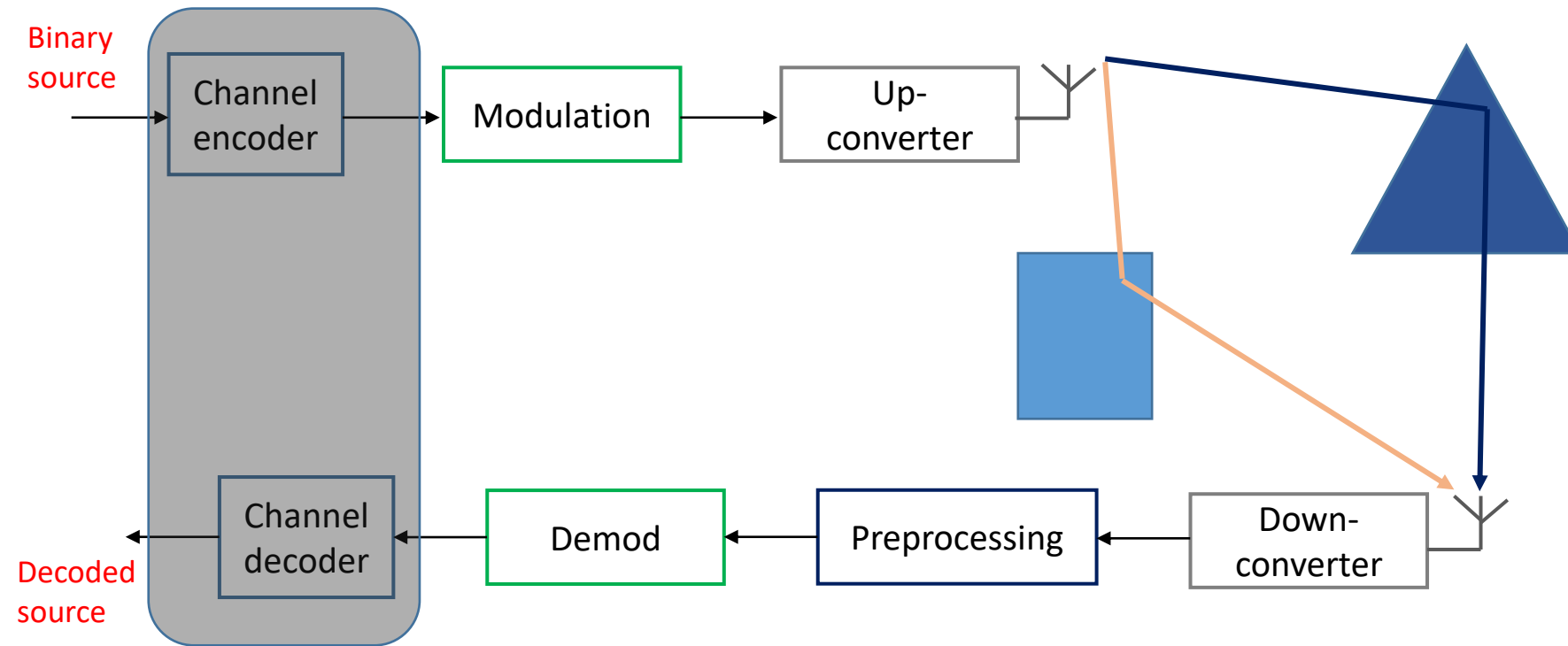
# $(7,4)$ -Hamming Codes

- 4-bit message  $\mathbf{m}$ , 7-bit codeword  $\mathbf{c}$
- Relationship:  $\mathbf{c} = \mathbf{mG}$

- $$\mathbf{G} := \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{pmatrix}$$

- $$\mathbf{H} := \begin{pmatrix} 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \end{pmatrix}$$

# Lab for Today – Uncoded system



# System Model

- When transmit, we map coded bits to baseband signal
- Binary phase shift keying (BPSK) for sending the bit  $c_i$

$$x_i = \sqrt{P}(2c_i - 1), \quad \mathbf{x} = [x_1, \dots, x_n]$$

- Additive white Gaussian noise (AWGN) channel

$$y_i = x_i + w_i, \quad w_i \sim N(0, N_0/2), \quad \mathbf{y} = [y_1, \dots, y_n]$$

# BER in Uncoded System

- Detection of  $x_i$  from  $y_i$

$$\hat{x}_i = \text{sign}(y_i) \quad \text{and} \quad \hat{r}_i = (\hat{x}_i + 1)/2$$

- Demodulate  $\hat{x}_i$  to  $\hat{m}_i$

- Error if  $\hat{m}_i \neq m_i$

- Bit error rate (BER):  $p_e = \sum 1(\{\hat{m}_i \neq m_i\})/n$

- Plot BER as a function of  $E_b/N_0$  where  $E_b$  is energy per bit