2.9 PBLAS routines for matrix-matrix operations (Level 3) Scalable Universal Matrix Multiply Algorithm (SUMMA)

An alternative for Cannon algorithm represents the SUMMA algorithm:

- SUMMA = Scalable Universal Matrix Multiply Algorithm
- Slightly less efficient, but simpler and easier to generalize
- Presentation from van de Geijn and Watts [1997]
 - www.netlib.org/lapack/lawns/lawn96.ps
 - Similar ideas appeared many times
- Used in practice in PBLAS = Parallel BLAS
 - www.netlib.org/lapack/lawns/lawn100.ps

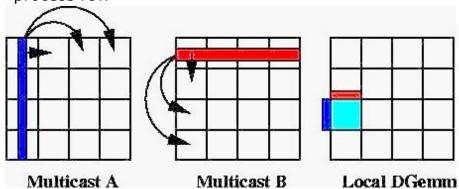
Naive matrix multiply

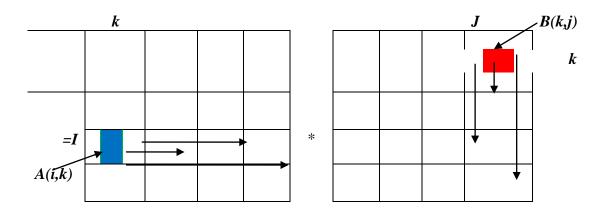
- For i = 0 to n
- For i = 0 to n
- For k = 0 to n
- C[i,j] += A[i,k]*B[k,j]

Calculates n^2 dot products (inner products) C[i,j] = A[i,:]*B[:,j]

The main steps of the algorithm:

- For each *k* (between 0 and n-1),
 - Owner of partial row k broadcasts that row along its process column;
 - Owner of partial column k broadcasts that column along its process row





	C(i,j)

- **I**, **J** represent all rows, columns owned by a processor
- **k** is a single row or column (or a block of **b** rows or columns)
- $C(I,J) = C(I,J) + \sum_{k} A(I,k) * B(k,J)$

Assume a p_r by p_c processor grid ($p_r = p_c = 4$ above) Complete pseudo code of the algorithm.

// On each process P(i,j):

- For k=0 to n-1 ... or n/b-1 where b is the block size
 - ...= # cols in A(I,k) and # rows in B(k,J)
 - for all I = 1 to p_r ... in parallel
 - owner of A(I,k) broadcasts it to whole processor row
 - for all J = 1 to p_c ... in parallel
 - owner of B(k,J) broadcasts it to whole processor column
 - Receive A(I,k) into Acol
 - Receive B(k, J) into Brow
 - C(myproc, myproc) = C(myproc, myproc) +Acol*Brow

• Endfor

We consider the following two matrices

A_{11}	A_{12}	A ₁₃	A_{14}	A_{15}	A_{16}	A ₁₇	A_{18}	A ₁₉
A_{21}	A ₂₂	A ₂₃	A ₂₄	A_{25}	A_{26}	A ₂₇	A_{28}	A ₂₉
A ₃₁	A ₃₂	A_{33}	A ₃₄	A ₃₅	A ₃₆	A ₃₇	A_{38}	A ₃₉
A ₄₁	A_{42}	A_{43}	A_{44}	A_{45}	A ₄₆	A ₄₇	A_{48}	A_{49}
A_{51}	A_{52}	A ₅₃	A ₅₄	A ₅₅	A ₅₆	A ₅₇	A_{58}	A ₅₉
A_{61}	A_{62}	A ₆₃	A ₆₄	A ₆₅	A ₆₆	A ₆₇	A ₆₈	A ₆₉
A ₇₁	A_{72}	A_{73}	A_{74}	A ₇₅	A_{76}	A ₇₇	A_{78}	A_{79}
A ₈₁	A_{82}	A_{83}	A_{84}	A_{85}	A_{86}	A_{87}	A_{88}	A_{89}
	A_{92}							

B_{11}	B ₁₂	B ₁₃	B ₁₄	B ₁₅	B ₁₆	B ₁₇	B_{18}	B ₁₉
B_{21}	B_{22}	B_{23}	B_{24}	B_{25}	B_{26}	B_{27}	B_{28}	B_{29}
B ₃₁	B_{32}	B ₃₃	B ₃₄	\mathbf{B}_{35}	\mathbf{B}_{36}	B ₃₇	\mathbf{B}_{38}	B ₃₉
B ₄₁	B_{42}	B ₄₃	\mathbf{B}_{44}	B_{45}	B ₄₆	B ₄₇	B_{48}	B ₄₉
B ₅₁	B ₅₂	B ₅₃	B ₅₄	B ₅₅	B ₅₆	B ₅₇	B ₅₈	B ₅₉
B_{61}	B ₆₂	B ₆₃	B ₆₄	B ₆₅	B ₆₆	B ₆₇	B ₆₈	B ₆₉
B ₇₁	\mathbf{B}_{72}	B ₇₃	B ₇₄	B ₇₅	B ₇₆	B ₇₇	\mathbf{B}_{78}	B_{79}
B_{81}	B_{82}	B ₈₃	B ₈₄	\mathbf{B}_{85}	\mathbf{B}_{86}	\mathbf{B}_{87}	\mathbf{B}_{88}	B ₈₉

As a result of distribution of these matrices using "2Dciclic" algorithm on the grid of 2x3 processes (2 lines and 3 columns) with block dimensions 2x2, we obtain:

2x2, we obtain:	1
Matrix A	Matrix B
$\mathbf{A_{(0,0)}} = \begin{bmatrix} A_{11} & A_{12} & A_{17} & A_{18} \\ A_{21} & A_{22} & A_{27} & A_{28} \\ A_{51} & A_{52} & A_{57} & A_{58} \\ A_{61} & A_{62} & A_{67} & A_{68} \\ A_{91} & A_{92} & A_{97} & A_{98} \end{bmatrix}$	$\mathbf{B_{(0,0)}} = \begin{pmatrix} B_{11} & B_{12} & B_{17} & B_{18} \\ B_{21} & B_{22} & B_{27} & B_{28} \\ B_{51} & B_{52} & B_{57} & B_{58} \\ B_{61} & B_{62} & B_{67} & B_{68} \\ B_{91} & B_{92} & B_{97} & B_{98} \end{pmatrix}$
$\mathbf{A_{(0,1)}} = \begin{vmatrix} A_{13} & A_{14} & A_{19} \\ A_{23} & A_{24} & A_{29} \\ A_{53} & A_{54} & A_{59} \\ A_{63} & A_{64} & A_{69} \\ A_{93} & A_{94} & A_{99} \end{vmatrix}$	$\mathbf{B_{(0,1)}} = \begin{pmatrix} B_{13} & B_{14} & B_{19} \\ B_{23} & B_{24} & B_{29} \\ B_{53} & B_{54} & B_{59} \\ B_{63} & B_{64} & B_{69} \\ B_{93} & B_{94} & B_{99} \end{pmatrix}$
$\begin{vmatrix} A_{15} & A_{16} \\ A_{25} & A_{26} \\ A_{55} & A_{56} \\ A_{65} & A_{66} \\ A_{95} & A_{96} \end{vmatrix}$	$\mathbf{B_{(0,2)}} = \begin{vmatrix} B_{15} & B_{16} \\ B_{25} & B_{26} \\ B_{55} & B_{56} \\ B_{65} & B_{66} \\ B_{95} & B_{96} \end{vmatrix}$
$\mathbf{A_{(1,0)}} = \begin{pmatrix} A_{31} & A_{32} & A_{37} & A_{38} \\ A_{41} & A_{42} & A_{47} & A_{48} \\ A_{71} & A_{72} & A_{77} & A_{78} \\ A_{81} & A_{82} & A_{87} & A_{88} \end{pmatrix}$	$\mathbf{B_{(1,0)}} = \begin{pmatrix} B_{31} & B_{32} & B_{37} & B_{38} \\ B_{41} & B_{42} & B_{47} & B_{48} \\ B_{71} & B_{72} & B_{77} & B_{78} \\ B_{81} & B_{82} & B_{87} & B_{88} \end{pmatrix}$
$\mathbf{A_{(1,1)}} = \begin{pmatrix} A_{33} & A_{34} & A_{39} \\ A_{43} & A_{44} & A_{49} \\ A_{73} & A_{74} & A_{79} \\ A_{83} & A_{84} & A_{89} \end{pmatrix}$	$\mathbf{B_{(1,1)}} = \begin{pmatrix} B_{33} & B_{34} & B_{39} \\ B_{43} & B_{44} & B_{49} \\ B_{73} & B_{74} & B_{79} \\ B_{83} & B_{84} & B_{89} \end{pmatrix}$

A _(1,2) =		A_{36} A_{46} A_{76}	B _(1,2) =	$ \begin{pmatrix} B_{35} \\ B_{45} \\ B_{75} \end{pmatrix} $	
	A_{85}	A_{86}		B_{85}	B_{86}

So in a compact form we have

	0	1	2
0	$A_{(0,0)}$	A _(0,1)	$A_{(0,2)}$
1	$A_{(1,0)}$	$A_{(1,1)}$	$A_{(1,2)}$

	0	1	2
0	$B_{(0,0)}$	$B_{(0,1)}$	$B_{(0,2)}$
1	$B_{(1,0)}$	$B_{(1,1)}$	$B_{(1,2)}$

To determine the product **C=AB** we use the algorithm SUMMA. Processes will determine the following submatrices of the matrix **C:**

$$\mathbf{C_{(0,0)}} = \begin{bmatrix} C_{11} & C_{12} & C_{17} & C_{18} \\ C_{21} & C_{22} & C_{27} & C_{28} \\ C_{51} & C_{52} & C_{57} & C_{68} \\ C_{61} & C_{62} & C_{67} & C_{68} \\ C_{91} & C_{92} & C_{97} & C_{98} \end{bmatrix} \qquad \mathbf{C_{(0,1)}} = \begin{bmatrix} C_{13} & C_{14} & C_{19} \\ C_{23} & C_{24} & C_{29} \\ C_{53} & C_{54} & C_{59} \\ C_{63} & C_{64} & C_{69} \\ C_{93} & C_{94} & C_{99} \end{bmatrix} \qquad \mathbf{C_{(0,2)}} = \begin{bmatrix} C_{15} & C_{16} \\ C_{25} & C_{26} \\ C_{55} & C_{56} \\ C_{65} & C_{66} \\ C_{95} & C_{96} \end{bmatrix}$$

$$\mathbf{C_{(1,0)}} = \begin{pmatrix} C_{31} & C_{32} & C_{37} & C_{38} \\ C_{41} & C_{42} & C_{47} & C_{48} \\ C_{71} & C_{72} & C_{77} & C_{78} \\ C_{81} & C_{82} & C_{87} & C_{88} \end{pmatrix} \quad \mathbf{C_{(1,1)}} = \begin{pmatrix} C_{33} & C_{34} & C_{39} \\ C_{43} & C_{44} & C_{49} \\ C_{73} & C_{74} & C_{79} \\ C_{83} & C_{84} & C_{89} \end{pmatrix} \quad \mathbf{C_{(1,2)}} = \begin{pmatrix} C_{35} & C_{36} \\ C_{45} & C_{46} \\ C_{75} & C_{76} \\ C_{85} & C_{86} \end{pmatrix}$$

Here for example

$$C_{11} = A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} + A_{17}B_{71} + A_{18}B_{81} + A_{19}B_{91} \\ C_{12} = A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} + A_{17}B_{72} + A_{18}B_{82} + A_{19}B_{92} \\ C_{17} = A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} + A_{17}B_{77} + A_{18}B_{87} + A_{19}B_{97} \\ C_{18} = A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} + A_{17}B_{78} + A_{18}B_{88} + A_{19}B_{98}$$

$$C_{21} = A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} + A_{27}B_{71} + A_{28}B_{81} + A_{29}B_{91}$$

$$C_{22} = A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} + A_{27}B_{72} + A_{28}B_{82} + A_{29}B_{92}$$

$$C_{27} = A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} + A_{27}B_{77} + A_{28}B_{87} + A_{29}B_{97}$$

$$C_{28} = A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{68} + A_{27}B_{77} + A_{28}B_{88} + A_{29}B_{98}$$

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\begin{array}{l} C_{51} = A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} + A_{57}B_{71} + A_{58}B_{81} + A_{59}B_{91} \\ C_{52} = A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} + A_{55}B_{52} + A_{56}B_{62} + A_{57}B_{72} + A_{58}B_{82} + A_{59}B_{92} \\ C_{57} = A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} + A_{57}B_{77} + A_{58}B_{87} + A_{59}B_{97} \\ C_{58} = A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48} + A_{55}B_{58} + A_{56}B_{68} + A_{57}B_{78} + A_{58}B_{88} + A_{59}B_{98} \\ C_{61} = A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{61} + A_{67}B_{71} + A_{68}B_{81} + A_{69}B_{91} \\ C_{62} = A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{42} + A_{65}B_{52} + A_{66}B_{62} + A_{67}B_{72} + A_{68}B_{82} + A_{69}B_{92} \\ C_{67} = A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} + A_{65}B_{57} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{87} + A_{69}B_{97} \\ C_{68} = A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48} + A_{65}B_{58} + A_{66}B_{68} + A_{67}B_{77} + A_{68}B_{88} + A_{69}B_{98} \\ C_{91} = A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} + A_{95}B_{51} + A_{96}B_{61} + A_{97}B_{71} + A_{98}B_{81} + A_{99}B_{91} \\ C_{92} = A_{91}B_{12} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} + A_{97}B_{77} + A_{98}B_{87} + A_{99}B_{97} \\ C_{98} = A_{91}B_{18} + A_{92}B_{27} + A_{93}B_{38} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} + A_{97}B_{77} + A_{98}B_{87} + A_{99}B_{98} \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} + A_{97}B_{78} + A_{98}B_{88} + A_{99}B_{98} \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} + A_{97}B_{78} + A_{98}B_{88} + A_{99}B_{98} \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} + A_{97}B_{78} + A_{98}B_{88} + A_{99}B_{98} \\ C_{98} = A_{91}B_{18} +
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We apply the algorithm SUMMA.

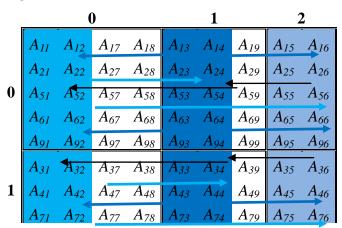
The main iterations are:

- I, J represent all rows, columns owned by a processor
- **k** is a single row or column (or a block of **b** rows or columns)
- $C(i,j) = C(i,j) + \sum_{k} A(i,k) \times B(k,j)$

Assume a p_r by p_c processor grid ($p_r = p_c = 4$ above). Need not be a square. Complete algorithm. On each process P(i,j):

- For k=0 to n-1 ... or n/b-1 where b is the block size = # cols in A(I,k) and # rows in B(k,J)
- for all I = 1 to p_r ... in parallel \checkmark owner of A(I,k) broadcasts it to whole processor row
- for all J = 1 to pc ... in parallel
 ✓ owner of B(k,J) broadcasts it to whole processor column
- Receive A(I,k) into Acol
- Receive B(k,J) into Brow
- C(myproc , myproc) = C(myproc , myproc) + Acol * Brow

The initial situation is:



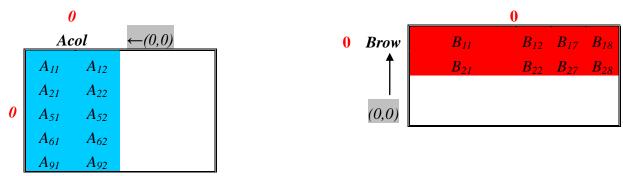
	A_{81}	A_{82}	A_{87}	A_{88}	A_{83}	A_{84}	A_{89}	A_{85}	A_{86}
		()			1		2	2
	B_{II}	B_{12}	B_{17}	B_{18}	B_{13}	B_{14}	B_{19}	B_{15}	B_{16}
	B_{21}	B_{22}							
0									
	B_{61}	B_{62}	B_{67}	B_{68}	B_{63}	B_{64}	B_{69}	B_{65}	B ₅₆ B ₆₆
	B_{91}	B_{92}	B_{97}	B_{98}	B_{93}	B_{94}	B_{99}	B_{95}	B_{96}
		B_{32}			B_{33}				
1	B_{41}	B_{42}			B_{43}				
	B_{71}	B_{72}			B_{73}				
	B_{81}	B_{82}	B_{87}	B_{88}	B_{83}	B_{84}	B_{89}	B_{85}	B_{86}

The process (0,0)

On the base of algorithm SUMMA we obtain:

Let k=2 (the block length, i.e. there are transmitted two rows and two columns).

• Iteration 0. There is no data transmission



Then it is possible to calculate the product $C_{(0,0)} = C_{(0,0)} + Acol*Brow$. So we obtain

$$C^{0}_{I,I} = A_{II}B_{II} + A_{I2}B_{2I} \quad C^{0}_{I,2} = A_{II}B_{12} + A_{12}B_{22} \quad C^{0}_{I,7} = A_{II}B_{17} + A_{12}B_{27} \quad C^{0}_{I,8} = A_{II}B_{18} + A_{12}B_{28}$$

$$C^{0}_{2,I} = A_{2I}B_{II} + A_{22}B_{2I} \quad C^{0}_{22} = A_{2I}B_{12} + A_{22}B_{22} \quad C^{0}_{2,7} = A_{2I}B_{17} + A_{22}B_{27} \quad C^{0}_{2,8} = A_{2I}B_{18} + A_{22}B_{28}$$

$$C^{0}_{5,I} = A_{5I}B_{II} + A_{52}B_{2I} \quad C^{0}_{52} = A_{5I}B_{12} + A_{52}B_{22} \quad C^{0}_{57} = A_{5I}B_{17} + A_{52}B_{27} \quad C^{0}_{58} = A_{5I}B_{18} + A_{52}B_{28}$$

$$C^{0}_{6I} = A_{6I}B_{II} + A_{62}B_{2I} \quad C^{0}_{62} = A_{6I}B_{12} + A_{62}B_{22} \quad C^{0}_{67} = A_{6I}B_{17} + A_{62}B_{27} \quad C^{0}_{68} = A_{6I}B_{18} + A_{62}B_{28}$$

$$C^{0}_{9I} = A_{9I}B_{II} + A_{92}B_{2I} \quad C^{0}_{92} = A_{9I}B_{I2} + A_{92}B_{22} \quad C^{0}_{97} = A_{9I}B_{17} + A_{92}B_{27} \quad C^{0}_{98} = A_{9I}B_{18} + A_{92}B_{28}$$

Finally, $C_{(0,0)} = C^{0}_{(0,0)}$

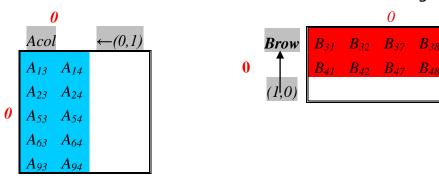
Therefore, at the end of this iteration the process (0,0) already has calculated $C_{68}=A_{61}B_{18}+A_{62}B_{28}$

 $\begin{array}{c} C_{11} = A_{11}B_{11} + A_{12}B_{21} \\ C_{12} = A_{11}B_{12} + A_{12}B_{22} \\ C_{17} = A_{11}B_{17} + A_{12}B_{27} \\ C_{18} = A_{11}B_{18} + A_{12}B_{28} \\ \end{array} \qquad \begin{array}{c} C_{91} = A_{91}B_{11} + A_{92}B_{211} \\ C_{92} = A_{91}B_{12} + A_{92}B_{222} \\ C_{97} = A_{91}B_{17} + A_{92}B_{277} \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} \\ \end{array}$

 $C_{21} = A_{21}B_{11} + A_{22}B_{21}$ $C_{22} = A_{21}B_{12} + A_{22}B_{22}$ $C_{27} = A_{21}B_{17} + A_{22}B_{27}$ $C_{28} = A_{21}B_{18} + A_{22}B_{28}$

 $C_{51} = A_{51}B_{11} + A_{52}B_{21}$ $C_{52} = A_{51}B_{12} + A_{52}B_{22}$ $C_{57} = A_{51}B_{17} + A_{52}B_{27}$ $C_{58} = A_{51}B_{18} + A_{52}B_{28}$

 $C_{61} = A_{61}B_{11} + A_{62}B_{21}$ $C_{62} = A_{61}B_{12} + A_{62}B_{22}$ $C_{67} = A_{61}B_{17} + A_{62}B_{27}$ • Iteration 1. Here we have for Acol and Bcol the following:



Then it is possible to calculate the product $C_{(0,0)} = C_{(0,0)} + Acol*Brow$. So we obtain

	'	'	- (0,0) - (0,0)	
	$C^{I}_{1,1} = A_{13}B_{31} + A_{14}B_{41}$	$C^{I}_{1,2} = A_{13}B_{32} + A_{14}B_{42}$	$C^{I}_{1,7} = A_{13}B_{37} + A_{14}B_{47}$	$C^{I}_{1,8} = A_{13}B_{38} + A_{14}B_{48}$
	$C^{I}_{2,I} = A_{23}B_{31} + A_{24}B_{41}$	$C^{I}_{22} = A_{23}B_{32} + A_{24}B_{42}$	$C^{I}_{2,7} = A_{23}B_{37} + A_{24}B_{47}$	$C^{I}_{2,8} = A_{23}B_{38} + A_{24}B_{48}$
$C^{1}_{(0,0)}=$	$C^{I}_{5,1}=A_{53}B_{31}+A_{54}B_{41}$	$C^{I}_{52} = A_{53}B_{32} + A_{54}B_{42}$	$C^{I}_{57} = A_{53}B_{37} + A_{54}B_{47}$	$C^{I}_{58} = A_{53}B_{38} + A_{54}B_{48}$
	$C^{I}_{61} = A_{63}B_{31} + A_{64}B_{41}$	$C^{I}_{62} = A_{63}B_{32} + A_{64}B_{42}$	$C^{I}_{67} = A_{63}B_{37} + A_{64}B_{47}$	$C^{I}_{68} = A_{63}B_{38} + A_{64}B_{48}$
	$C^{I}_{91} = A_{93}B_{31} + A_{94}B_{41}$	$C^{I}_{92} = A_{93}B_{32} + A_{94}B_{42}$	$C^{I}_{97} = A_{93}B_{37} + A_{94}B_{47}$	$C^{l}_{98} = A_{93}B_{38} + A_{94}B_{48}$

Therefore, at the end of this iteration the process (0,0) already has calculated

$$C_{11}=A_{11}B_{11}+A_{12}B_{21}+A_{13}B_{31}+A_{14}B_{41}$$

$$C_{12}=A_{11}B_{12}+A_{12}B_{22}+A_{13}B_{32}+A_{14}B_{42}$$

$$C_{17} = A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47}$$

$$C_{18} = A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48}$$

$$C_{21} = A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41}$$

$$C_{22} = A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42}$$

$$C_{27} {=} A_{21} B_{17} + A_{22} B_{27} {+} \ A_{23} B_{37} \ {+} A_{24} B_{47}$$

$$C_{28} = A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48}$$

$$C_{51} = A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41}$$

$$C_{52} = A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42}$$

$$C_{57} {=} A_{51} B_{17} + A_{52} B_{27} {+} A_{53} B_{37} + A_{54} B_{47}$$

$$C_{58} = A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48}$$

$$C_{61} = A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41}$$

$$C_{62} \! = \! A_{61}B_{12} + A_{62}B_{22} \! + A_{63}B_{32} + \! A_{64}B_{42}$$

$$C_{67} = A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47}$$

$$C_{68} = A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48}$$

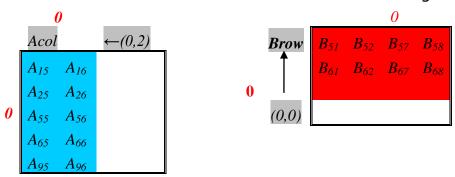
$$C_{91} = A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41}$$

$$C_{92} {=} A_{91} B_{12} + A_{92} B_{22} {+} \ A_{93} B_{32} \ {+} A_{94} B_{42}$$

$$C_{97} {=} A_{91} B_{17} + A_{92} B_{27} {+} \ A_{93} B_{37} \ {+} A_{94} B_{47}$$

$$C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48}$$

• Iteration 2. Here we have for Acol and Bcol the following:



Then it is possible to calculate the product $C_{(0,0)} = C_{(0,0)} + Acol*Brow$. So we obtain

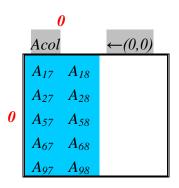
$$C^{2}_{I,I} = A_{15}B_{5I} + A_{16}B_{6I} \quad C^{2}_{I,2} = A_{15}B_{52} + A_{16}B_{62} \quad C^{2}_{I,7} = A_{15}B_{57} + A_{16}B_{67} \quad C^{2}_{I,8} = A_{15}B_{58} + A_{16}B_{68} \\ C^{2}_{2,I} = A_{25}B_{5I} + A_{26}B_{61} \quad C^{2}_{22} = A_{25}B_{52} + A_{26}B_{62} \quad C^{2}_{2,7} = A_{25}B_{57} + A_{26}B_{67} \quad C^{2}_{2,8} = A_{25}B_{58} + A_{26}B_{68} \\ C^{2}_{6I} = A_{65}B_{5I} + A_{66}B_{61} \quad C^{2}_{52} = A_{55}B_{52} + A_{66}B_{62} \quad C^{2}_{57} = A_{55}B_{57} + A_{56}B_{67} \quad C^{2}_{58} = A_{55}B_{58} + A_{56}B_{68} \\ C^{2}_{9I} = A_{95}B_{5I} + A_{96}B_{61} \quad C^{2}_{92} = A_{95}B_{52} + A_{96}B_{62} \quad C^{2}_{97} = A_{95}B_{57} + A_{96}B_{67} \quad C^{2}_{98} = A_{95}B_{58} + A_{96}B_{68}$$

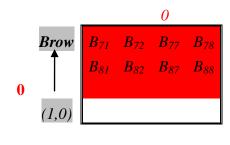
Therefore, at the end of this iteration the process (0,0) already has calculated

$$\begin{array}{l} C_{12} = A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} \\ C_{17} = A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} \\ C_{18} = A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} \\ \\ C_{21} = A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} \\ C_{22} = A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} \\ C_{27} = A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} \\ C_{28} = A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{68} \\ \\ C_{51} = A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} \\ C_{52} = A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} + A_{55}B_{52} + A_{56}B_{62} \\ C_{57} = A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} \\ C_{58} = A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{68} \\ \\ C_{61} = A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{61} \\ C_{62} = A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{41} + A_{65}B_{57} + A_{66}B_{62} \\ C_{67} = A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} + A_{65}B_{57} + A_{66}B_{67} \\ C_{68} = A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48} + A_{65}B_{58} + A_{66}B_{68} \\ \\ C_{91} = A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} + A_{95}B_{51} + A_{96}B_{61} \\ C_{92} = A_{91}B_{12} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} \\ \\ C_{97} = A_{91}B_{17} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{68} \\ \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} \\ \\ C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}$$

 $C_{11} = A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61}$

Iteration 3. Here we have for Acol and Bcol the following:





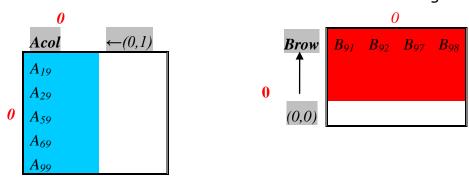
Then it is possible to calculate the product $C_{(0,0)} = C_{(0,0)} + Acol*Brow$. So we obtain

$$C^{3}_{1,1} = A_{17}B_{71} + A_{18}B_{81} \quad C^{3}_{1,2} = A_{17}B_{72} + A_{18}B_{82} \quad C^{3}_{1,7} = A_{17}B_{77} + A_{18}B_{87} \quad C^{3}_{1,8} = A_{17}B_{78} + A_{18}B_{88} \\ C^{3}_{2,1} = A_{27}B_{71} + A_{28}B_{81} \quad C^{3}_{22} = A_{27}B_{72} + A_{28}B_{82} \quad C^{3}_{2,7} = A_{27}B_{77} + A_{28}B_{87} \quad C^{3}_{2,8} = A_{27}B_{78} + A_{28}B_{88} \\ C^{3}_{61} = A_{67}B_{71} + A_{58}B_{81} \quad C^{3}_{52} = A_{57}B_{72} + A_{58}B_{82} \quad C^{3}_{57} = A_{57}B_{77} + A_{58}B_{87} \quad C^{3}_{58} = A_{57}B_{78} + A_{58}B_{88} \\ C^{3}_{61} = A_{67}B_{71} + A_{68}B_{81} \quad C^{3}_{62} = A_{67}B_{72} + A_{68}B_{82} \quad C^{3}_{67} = A_{67}B_{77} + A_{68}B_{87} \quad C^{3}_{68} = A_{67}B_{78} + A_{68}B_{88} \\ C^{3}_{91} = A_{97}B_{71} + A_{98}B_{81} \quad C^{3}_{92} = A_{97}B_{72} + A_{98}B_{82} \quad C^{3}_{97} = A_{97}B_{77} + A_{98}B_{87} \quad C^{3}_{98} = A_{97}B_{78} + A_{98}B_{88}$$

Therefore, at the end of this iteration the process (0,0) already has calculated

$$\begin{array}{l} C_{11} = A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} + A_{17}B_{71} + A_{18}B_{81} \\ C_{12} = A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} + A_{17}B_{72} + A_{18}B_{82} \\ C_{17} = A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} + A_{17}B_{77} + A_{18}B_{87} \\ C_{18} = A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} + A_{17}B_{78} + A_{18}B_{88} \\ C_{21} = A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} + A_{27}B_{71} + A_{28}B_{81} \\ C_{22} = A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} + A_{27}B_{77} + A_{28}B_{82} \\ C_{27} = A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} + A_{27}B_{77} + A_{28}B_{87} \\ C_{28} = A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{67} + A_{27}B_{77} + A_{28}B_{88} \\ C_{51} = A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} + A_{57}B_{71} + A_{58}B_{81} \\ C_{52} = A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{41} + A_{55}B_{57} + A_{56}B_{67} + A_{57}B_{77} + A_{28}B_{88} \\ C_{57} = A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} + A_{57}B_{77} + A_{58}B_{87} \\ C_{58} = A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48} + A_{55}B_{58} + A_{56}B_{67} + A_{57}B_{77} + A_{58}B_{88} \\ C_{61} = A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{81} \\ C_{62} = A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{33} + A_{64}B_{41} + A_{65}B_{57} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{87} \\ C_{68} = A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{41} + A_{65}B_{57} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{87} \\ C_{68} = A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{41} + A_{95}B_{57} + A_{66}B_{67} + A_{67$$

• Iteration 4. Here we have for *Acol* and *Bcol* the following:



Then it is possible to calculate the product C(0,0) = C(0,0) + Acol*Brow. So we obtain

$$C^{4}_{1,1} = A_{19}B_{91} \quad C^{4}_{1,2} = A_{19}B_{92} \quad C^{4}_{1,7} = A_{19}B_{97} \quad C^{4}_{1,8} = A_{19}B_{98}$$

$$C^{4}_{2,1} = A_{29}B_{91} \quad C^{4}_{22} = A_{29}B_{92} \quad C^{4}_{2,7} = A_{29}B_{97} \quad C^{4}_{2,8} = A_{29}B_{98}$$

$$C^{4}_{61} = A_{69}B_{91} \quad C^{4}_{62} = A_{69}B_{92} \quad C^{4}_{67} = A_{69}B_{97} \quad C^{4}_{68} = A_{69}B_{98}$$

$$C^{4}_{91} = A_{99}B_{91} \quad C^{4}_{92} = A_{99}B_{92} \quad C^{4}_{97} = A_{99}B_{97} \quad C^{4}_{98} = A_{99}B_{98}$$

Therefore, at the end of this iteration the process (0,0) already has calculated

```
C_{11} = A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} + A_{17}B_{71} + A_{18}B_{81} + A_{19}B_{91}
C_{12} = A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} + A_{17}B_{72} + A_{18}B_{82} + A_{19}B_{92}
C_{17} = A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} + A_{17}B_{77} + A_{18}B_{87} + A_{19}B_{97}
C_{18} = A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} + A_{17}B_{78} + A_{18}B_{88} + A_{19}B_{98}
C_{21} = A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} + A_{27}B_{71} + A_{28}B_{81} + A_{29}B_{91}
C_{22} = A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} + A_{27}B_{72} + A_{28}B_{82} + A_{29}B_{92}
C_{27} = A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} + A_{27}B_{77} + A_{28}B_{87} + A_{29}B_{97}
C_{28} = A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{68} + A_{27}B_{77} + A_{28}B_{88} + A_{29}B_{98}
C_{51} = A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} + A_{57}B_{71} + A_{58}B_{81} + A_{59}B_{91}
C_{52} = A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} + A_{55}B_{52} + A_{56}B_{62} + A_{57}B_{72} + A_{58}B_{82} + A_{59}B_{92}
C_{57} = A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} + A_{57}B_{77} + A_{58}B_{87} + A_{59}B_{97}
C_{58} = A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48} + A_{55}B_{58} + A_{56}B_{68} + A_{57}B_{78} + A_{58}B_{88} + A_{59}B_{98}
C_{61} = A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{61} + A_{67}B_{71} + A_{68}B_{81} + A_{69}B_{91}
C_{62} = A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{42} + A_{65}B_{52} + A_{66}B_{62} + A_{67}B_{72} + A_{68}B_{82} + A_{69}B_{92}
C_{67} = A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} + A_{65}B_{57} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{87} + A_{69}B_{97}
C_{68} = A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48} + A_{65}B_{58} + A_{66}B_{68} + A_{67}B_{77} + A_{68}B_{88} + A_{69}B_{98}
C_{91} = A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} + A_{95}B_{51} + A_{96}B_{61} + A_{97}B_{71} + A_{98}B_{81} + A_{99}B_{91}
C_{92} = A_{91}B_{12} + A_{92}B_{22} + A_{93}B_{32} + A_{94}B_{42} + A_{95}B_{52} + A_{96}B_{62} + A_{97}B_{72} + A_{98}B_{82} + A_{99}B_{92}
C_{97} = A_{91}B_{17} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} + A_{97}B_{77} + A_{98}B_{87} + A_{99}B_{97}
```

 $C_{98} = A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} + A_{97}B_{78} + A_{98}B_{88} + A_{99}B_{98}$

Purpose: PvGEMM performs one of the matrix-matrix operations

 $sub(C) := alpha \times op(sub(A)) \times op(sub(B)) + beta \times sub(C),$ where sub(C) denotes C(IC:IC+M-1,JC:JC+N-1),

```
op(sub(A)) = \begin{cases} A( \text{ IA} : \text{IA} + \text{M} - 1, \text{JA} : \text{JA} + \text{K} - 1) & \text{if TRANSA} = '\text{N'} \\ A( \text{ IA} : \text{IA} + \text{K} - 1, \text{JA} : \text{JA} + \text{M} - 1)' & \text{if TRANSA} = '\text{T} \\ A( \text{ IA} : \text{IA} + \text{K} - 1, \text{JA} : \text{JA} + \text{M} - 1)' & \text{if TRANSA} = '' \end{cases}
op(sub(B)) = \begin{cases} B( \text{ IB} : \text{IB} + \text{K} - 1, \text{JB} : \text{JB} + \text{N} - 1) & \text{if TRANSB} = '\text{N'} \\ B( \text{ IB} : \text{IB} + \text{N} - 1, \text{JB} : \text{JB} + \text{K} - 1)' & \text{if TRANSB} = '\text{T} \end{cases}
B( \text{ IB} : \text{IB} + \text{N} - 1, \text{JB} : \text{JB} + \text{K} - 1)' & \text{if TRANSB} = '\text{C'} \end{cases}
```

alpha and beta are scalars, and sub(A), sub(B) and sub(C) are distributed matrices, with op(sub(A)) an M-by-K distributed matrix, op(sub(B)) a K-by-N distributed matrix and sub(C) an M-by-N distributed matrix.

Arguments

```
TRANSA (global input) CHARACTER
```

The form of op(A) to be used in the matrix multiplication as follows:

TRANSA = 'N', op(A) = A,

 $TRANSA = 'T', op(A) = A^T$

 $TRANSA = 'C', op(A) = A^{T}.$

TRANSB (global input) CHARACTER

The form of op(B) to be used in the matrix multiplication as follows:

TRANSB = 'N', op(B) = B,

 $TRANSB = 'T', op(B) = B^T$

 $TRANSB = 'C', op(B') = B^{T}.$

M (global input) INTEGER

The number of rows of the distributed matrices op(sub(A)) and sub(C). $M \ge 0$.

N (global input) INTEGER

The number of columns of the distributed matrices op(sub(B)) and sub(C). N > 0.

K (global input) INTEGER

The number of columns of the distributed matrix op(sub(A)) and the number of rows of the distributed matrix op(B). $K \ge 0$.

ALPHA (global input) REAL/COMPLEX

On entry, ALPHA specifies the scalar alpha.

A (local input) array of dimension (LLD A, KLa)

where KLa is LOCq(JA+K-1) when TRANSA = 'N', and is LOCq(JA+M-1) otherwise. Before entry, this array must contain the local pieces of the distributed matrix SUC(A).

IA (global input) INTEGER

The global row index of the submatrix of the distributed matrix A to operate on.

JA (global input) INTEGER

The global column index of the submatrix of the distributed matrix A to operate on.

DESCA (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix A.

```
B (local input) array of dimension (LLD B, KLb)
        where KLb is LOC_a(JB+N-1) when TRANSB = 'N', and is LOC_a(JB+K-1) otherwise.
        Before entry this array must contain the local pieces of the distributed matrix
        IB (global input) INTEGER
        The global row index of the submatrix of the distributed matrix B to operate on.
JB (global input) INTEGER
        The global column index of the submatrix of the distributed matrix B to operate
DESCB (global and local input) INTEGER array of dimension 8
        The array descriptor of the distributed matrix B.
BETA (global input) REAL/COMPLEX
        On entry, BETA specifies the scalar beta. When BETA is supplied as zero then
        sub( Y ) need not be set on input.
C (local input/local output) array of dimension (LLD C, LOC_{\alpha}(JC+N-
        1))
        Before entry, this array must contain the local pieces of the distributed matrix
        sub( C ). On exit, the distributed matrix sub( C ) is overwritten by the M-by-N
        distributed matrix alpha*op(sub(A))*op(sub(B)) + beta*sub(C).
IC (global input) INTEGER
        The global row index of the submatrix of the distributed matrix C to operate on.
JC (global input) INTEGER
        The global column index of the submatrix of the distributed matrix C to operate
DESCC (global and local input) INTEGER array of dimension 8
        The array descriptor of the distributed matrix C.
```

Example 2.9.1. The using of function pdgemm_

```
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include "mpi.h"
#include <iostream>
#include <iomanip>
#include <fstream>
#include <sstream>
using namespace std;
#define AA(i,j) AA[(i)*m+(j)]
#define BB(i,j) BB[(i)*k+(j)]
\#define CC(i,j) CC[(i)*m+(j)]
extern "C"
{
void Cblacs pinfo( int* mypnum, int* nprocs);
void Cblacs get( int context, int request, int* value);
int Cblacs gridinit( int* context, char * order, int np row, int np col);
void Cblacs gridinfo( int context, int* np row, int* np col, int* my row, int*
                    my col);
void Cblacs gridexit( int context);
void Cblacs barrier(int, const char*);
void Cblacs exit( int error code);
void Cblacs pcoord(int, int, int*, int*);
int numroc (int *n, int *nb, int *iproc, int *isrcproc, int *nprocs);
int indx12g (int*, int*, int*, int*);
void descinit (int *desc, int *m, int *n, int *mb, int *nb, int *irsrc, int *icsrc,
              int *ictxt, int *lld, int *info);
void pdgemm (char *TRANSA, char *TRANSB, int *M, int *N, int *K, double *ALPHA, double
```

```
*A,int *IA,int *JA,int *DESCA, double * B, int * IB, int * JB, int
             *DESCB, double *BETA, double *C, int *IC, int *JC, int *DESCC);
} // extern "C"
int main(int argc, char **argv)
int i, j;
int iam, nprocs, nprow, npcol, myrow, mycol;
int descA[9],descB[9],descC[9];
int m,n,k,mb,nb,rsrc,csrc,ictxt,llda,lldb,lldc,info;
int lm, ln, lk;
int iloc, jloc;
int ZERO=0, ONE=1;
double alpha, beta;
m=5; n=5; k=5;
mb=2; nb=2;
double *AA = (double*)malloc(m*k*sizeof(double));// matricea "globala" pentru
            inmultirea AA*BB=CC
double *BB = (double*) malloc(k*n*sizeof(double));//matricea "globala" pentru
             inmultirea AA*BB=CC
double *CC = (double*) malloc(m*n*sizeof(double));//matricea "globala" pentru
             inmultirea AA*BB=CC
// initializarea marticelor globale
for(i=0;i<m;i++)
     for (j=0; j < k; j++)
            AA[i*m+j] = (i+j);
for(i=0;i<k;i++)
     for(j=0;j<n;j++)
            BB[i*k+j]=j+i;
nprow=2; npcol=2; // Astfel, programul se executa pe 4 procese
Cblacs pinfo(&iam, &nprocs);
Cblacs get(-1, 0, &ictxt);
Cblacs gridinit( &ictxt, "Row", nprow, npcol);
Cblacs gridinfo(ictxt, &nprow, &npcol, &myrow, &mycol);
rsrc=0; csrc=0;
  if (iam==0)
  printf("======= REZULT OF THE PROGRAM %s \n", argv[0]);
  cout << "Global matrix AA:\n";</pre>
              for (i = 0; i < m; ++i) {
                  for (j = 0; j < n; ++j) {
                      cout << setw(3) << *(AA + m*i + j) << " ";
                  cout << "\n";
              cout << endl;</pre>
  cout << "Global matrix BB:\n";</pre>
              for (i = 0; i < k; ++i) {
                  for (j = 0; j < n; ++j) {
                       cout << setw(3) << *(BB + k*i + j) << " ";
                  cout << "\n";
              }
              cout << endl;
 }
Cblacs_barrier(ictxt, "All");
int mA=numroc ( &m, &mb, &myrow, &rsrc, &nprow );
int kA = numroc ( &k, &nb, &mycol, &rsrc, &npcol );
int kB = numroc ( &k, &mb, &myrow, &rsrc, &nprow );
int nB = numroc ( &n, &nb, &mycol, &rsrc, &npcol );
int mC = numroc ( &m, &mb, &myrow, &rsrc, &nprow );
int nC = numroc_( &n, &nb, &mycol, &rsrc, &nprow );
```

```
&mb, &nb, &rsrc, &csrc, &ictxt, &mA, &info);
descinit (descA, &m, &k,
descinit_(descB, &k, &n,
                           &mb, &nb, &rsrc, &csrc, &ictxt, &kB, &info);
descinit (descC, &m, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &mC, &info);
double *A=(double*) malloc(mA*kA*sizeof(double)); //matricea locala
double *B = (double*) malloc(kB*nB*sizeof(double)); //matricea locala
double *C = (double*) malloc(mC*nC*sizeof(double)); //matricea locala
// se complecteaza cu valori matricele locale folosind algoritmul 2D-ciclic
for(iloc=0;iloc<mA;iloc++)</pre>
   for(jloc=0;jloc<kA;jloc++) {</pre>
   int fortidl = iloc + 1;
  int fortjdl = jloc + 1;
  i = indx12g (&fortid1, &mb, &myrow, &ZERO, &nprow)-1;
   j = indx12g (&fortjdl, &nb, &mycol, &ZERO, &npcol)-1;
  A[jloc*mA+iloc]=AA(i,j);
for(iloc=0;iloc<kB;iloc++)</pre>
   for(jloc=0;jloc<nB;jloc++) {</pre>
   int fortidl = iloc + 1;
   int fortjdl = jloc + 1;
   i = indx12g (&fortid1, &mb, &myrow, &ZERO, &nprow)-1;
   j = indxl2g (&fortjdl, &nb, &mycol, &ZERO, &npcol)-1;
   B[jloc*kB+iloc]=BB(i,j);
alpha = 1.0; beta = 0.0;
pdgemm ("No Transpose", "No Transpose", &m, &n, &k, &alpha, A, &ONE, &ONE, descA,
       B, &ONE, &ONE, descB, &beta, C, &ONE, &ONE, descC);
Cblacs barrier(ictxt, "All");
// Print out the matrix product C
for (int id = 0; id < nprocs; ++id)</pre>
if (id == iam) {
cout << "C loc on node " << iam << endl;</pre>
for (i = 0; i < mC; ++i)
 for (j = 0; j < nC; ++j)
 cout << setw(3) << *(C+mC*j+i) << " ";
 cout << endl;
Cblacs barrier(ictxt, "All");
}
// Se construieste matricea globala CC
for(iloc=0;iloc<mC;iloc++)</pre>
  for(jloc=0;jloc<nC;jloc++) {</pre>
  int fortidl = iloc + 1;
  int fortjdl = jloc + 1;
   i = indx12g (&fortid1, &mb, &myrow, &ZERO, &nprow)-1;
   j = indxl2g (&fortjdl, &nb, &mycol, &ZERO, &npcol)-1;
   CC(i,j) = C[jloc*mC+iloc];
Cblacs barrier(ictxt, "All");
//*** Tipar rezultate finale
for (int id = 0; id < nprocs; ++id)
 if (id == iam) {
 cout << "Global matrix CC on node "<< iam << endl;</pre>
              for (i = 0; i < m; ++i) {
                  for (j = 0; j < n; ++j)
                       cout << setw(3) << *(CC + m*i + j) << " ";
                  cout << "\n";
```

```
cout << endl;</pre>
Cblacs barrier(ictxt, "All");
/* Free memory */
   free(A);
   free(B);
   free(C);
   /* Release process grid */
   Cblacs gridexit(ictxt);
} /* main */
             The results are the following
[MI gr TPS1@hpc] $./mpiCC ScL -o Example2.9.1.exe Example2.9.1.cpp
[MI gr TPS1@hpc] $/opt/openmpi/bin/mpirun -n 4 -host compute-0-0,compute-0-1
               Example2.9.1.exe
[MI gr TPS1@hpc]$./mpiCC ScL -o pmm.exe pmm.cpp
[MI gr TPS1@hpc]$/opt/openmpi/bin/mpirun -n 4 -host compute-0-0,compute-0-1 pmm.exe
Global matrix AA:
 0 1 2 3
                4
 1 2 3 4 5
 2 3 4 5
                6
   4 5 6
 3
                7
 4 5 6 7
Global matrix BB:
 0 1 2 3
                4
 1 2 3 4 5
 2 3 4 5 6
 3 4 5 6
              7
 4 5 6 7
                8
C loc on node 0
 30 40 70
40 55 100
70 100 190
C loc on node 1
50 60
70 85
130 160
C loc on node 2
50 70 130
60 85 160
C loc on node 3
90 110
110 135
Global matrix CC on node 0
 30 40 0 0 70
 40 55
       0 0 100
 0 0 0 0 0
 0 0 0 0 0
```

70 100 0 0 190

0 70 85 0 0 0 0 0 0 0 0 0 0 0 130 160 0 Global matrix CC on node 2

0 0 50 60

0 0 0 0

50 70 0 0 130

0

0

0

Global matrix CC on node 1

0 0 0

0 0

0

```
60 85 0 0 160

0 0 0 0 0 0

Global matrix CC on node 3

0 0 0 0 0 0

0 0 0 0 0

0 0 90 110 0

0 0 110 135 0

0 0 0 0 0
```

To verify the results we use the sequential algorithm based on function dgemm. The program is presented below:

```
#include <string>
#include <iostream>
#include <stdio.h>
#include <math.h>
#include <iomanip>
#include <fstream>
#include <sstream>
using namespace std;
extern "C" {
void dgemm (const char &transa, const char &transb,
              const int &1, const int &n, const int &m, const double &alpha,
              const double *a, const int &lda, const double *b, const int &ldb,
              const double &beta, double *c, const int &ldc);
int main() {
int i, j, m=5, k=5, n=5;
  //Matrizen werden als eindimensionale Arrays gespeichert
  double *AA = new double [m*k];
  double *BB = new double[k*n];
  double *CC = new double [m*n];
  //Matrizen fuelln
for(i=0;i<m;i++)
     for (j=0; j < k; j++)
            AA[i*k+j]=i+j; //(10*i+j);
for(i=0;i<k;i++)
     for(j=0;j<n;j++)
            BB[i*n+j]=i+j;//AA[i*5+j]+5;
for( i=0 ; i<m*n ; i++)
    CC[i] = 0.0;
  cout << "Global matrix AA:\n";</pre>
              for (i = 0; i < m; ++i) {
                  for (j = 0; j < k; ++j) {
                       cout << setw(3) << *(AA + k*i + j) << " ";
                  cout << "\n";
              cout << endl;
  cout << "Global matrix BB:\n";</pre>
              for (i = 0; i < k; ++i) {
                  for (j = 0; j < n; ++j) {
                       cout << setw(3) << *(BB + n*i + j) << ";
                  cout << "\n";
              cout << endl;
  double alpha = 1.0;
  double beta = 0.0;
```

```
//dgemm berechnet C=alpha*A*B+beta*C
dgemm_('N','N',m,n,k,alpha,AA,m,BB,k,beta,CC,m);

cout << "Global matrix CC:\n " << endl;
for(i=0; i<m; i++) {
    for(j=0; j<n; j++)
cout << setw(3) << *(CC + n*i + j) << " ";
    cout << endl;
}
cout << endl;
delete[] AA;
delete[] BB;
delete[] CC;
return 0;
}</pre>
```

The results are the following:

```
[MI gr TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 1 -host compute-0-0,compute-0-4
dgemm1.exe
Global matrix AA:
 0 1 2 3
    2 3
 1
          4
              5
 2 3 4 5
              6
 3 4 5 6 7
 4 5 6 7
Global matrix BB:
 0 1 2 3
 1 2 3 4 5
 2 3 4 5 6
 3 4 5 6 7
 4 5 6 7 8
Global matrix CC:
30 40 50 60 70
40 55 70 85 100
50 70 90 110 130
60 85 110 135 160
70 100 130 160 190
```

Additional explanations for Example 2.9.1

In the program Example 2.9.1.cpp the parallelization on data level is "non-rational" in terms of usage of operative memory, because for all processes the operative memory is reserved for global matrices! Using the function pdgeadd we can develop a program, in which the operative memory is used already rationally. The program is presented below.

Example 2.9.2.

```
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include "mpi.h"
#include <iostream>
#include <iomanip>
#include <fstream>
#include <sstream>
using namespace std;
\#define AA(i,j) AA[(i)*n+(j)]
\#define BB(i,j) BB[(i)*k+(j)]
\#define CC(i,j) CC[(i)*n+(j)]
#define A(i,j) A[(i)*n+(j)]
#define B(i,j) B[(i)*n+(j)]
#define C(i,j) C[(i)*n+(j)]
static int MAX( int a, int b ){
        if (a>b) return(a); else return(b);
extern "C"
void Cblacs pinfo( int* mypnum, int* nprocs);
void Cblacs get( int context, int request, int* value);
int Cblacs gridinit( int* context, char * order, int np row, int np col);
void Cblacs gridinfo( int context, int* np row, int* np col, int* my row, int*
my col);
void Cblacs_gridexit( int context);
void Cblacs barrier(int, const char*);
void Cblacs exit( int error code);
void Cblacs pcoord(int, int, int*, int*);
int numroc (int *n, int *nb, int *iproc, int *isrcproc, int *nprocs);
int indx12g (int*, int*, int*, int*);
void descinit (int *desc, int *m, int *n, int *mb, int *nb, int *irsrc, int *icsrc,
int *ictxt, int *lld, int *info);
void pdgemm ( char *TRANSA, char *TRANSB, int *M, int *N, int *K, double *ALPHA, double
*A, int *IA, int *JA, int *DESCA,
             double * B, int * IB, int * JB, int * DESCB, double * BETA, double * C,
int * IC, int * JC, int * DESCC );
void pdgeadd (char *TRANS, int *M, int *N, double * ALPHA, double *A, int *IA, int *JA, int
*DESCA, double *BETA, double *C,
              int *IC, int *JC, int *DESCC);
} // extern "C"
int main(int argc, char **argv)
int i, j;
int iam, nprocs, nprow, npcol, myrow, mycol;
// descriptori pentri matrici locale
int descA[9],descB[9],descC[9];
// descriptori pentri matrici globale
int descAA[9],descBB[9],descCC[9];
```

```
int m,n,k,mb,nb,rsrc,csrc,ictxt,llda,lldb,lldc,info;
int lm, ln, lk;
int iloc, jloc;
int ZERO=0, ONE=1;
double zero=0.0E+0, one=1.0E+0;
int i one = 1, i zero = 0;
int 11d AA, 11d BB, 11d CC;
double alpha, beta;
double *AA, *BB, *CC, *A, *B, *C, *work, *tau;
m=5; n=5; k=5;
mb=2; nb=2;
nprow=2; npcol=2; // Astfel, programul se executa pe 4 procese
Cblacs pinfo(&iam, &nprocs);
Cblacs get(-1, 0, &ictxt);
Cblacs gridinit( &ictxt, "Row", nprow, npcol);
Cblacs gridinfo(ictxt, &nprow, &npcol, &myrow, &mycol);
rsrc=0; csrc=0;
if (iam==0)
AA = (double*) malloc(m*k*sizeof(double));// matricea "globala" (de dim. m*k) pentru
inmultirea AA*BB=CC
BB = (double*) malloc(k*n*sizeof(double));//matricea "globala" (de dim. k*n) pentru
inmultirea AA*BB=CC
CC = (double*) malloc(m*n*sizeof(double));//matricea "globala" (de dim. m*n) pentru
inmultirea AA*BB=CC
// initializarea marticelor globale
for(i=0;i<m;i++)
     for(j=0;j<k;j++)
            AA[i*k+j] = (10*i+j); //(i+j);
for(i=0;i<k;i++)
     for (j=0; j< n; j++)
            BB[i*n+j]=AA[i*5+j]+5; //(i+j);
else{
AA = NULL;
BB = NULL;
//other processes don't contain parts of A
  if (iam==0)
  printf("======= REZULT OF THE PROGRAM %s \n",arqv[0]);
  cout << "Global matrix AA:\n";</pre>
              for (i = 0; i < m; ++i) {
                  for (j = 0; j < k; ++j) {
                      cout << setw(3) << *(AA + k*i + j) << " ";
                  cout << "\n";
              }
              cout << endl;</pre>
  cout << "Global matrix BB:\n";</pre>
              for (i = 0; i < k; ++i) {
                  for (j = 0; j < n; ++j) {
                       cout << setw(3) << *(BB + n*i + j) << "";
                  cout << "\n";
              cout << endl;
  }
Cblacs barrier(ictxt, "All");
int mA = numroc_( &m, &mb, &myrow, &rsrc, &nprow );
```

```
int kA = numroc ( &k, &nb, &mycol, &rsrc, &npcol );
int kB = numroc ( &k, &mb, &myrow, &rsrc, &nprow );
int nB = numroc ( &n, &nb, &mycol, &rsrc, &npcol );
int mC = numroc ( &m, &mb, &myrow, &rsrc, &nprow );
int nC = numroc ( &n, &nb, &mycol, &rsrc, &nprow );
11d AA = MAX( numroc ( &m, &k, &myrow, &rsrc, &nprow ), 1 );
11d BB = MAX( numroc ( &k, &n, &myrow, &rsrc, &nprow ), 1 );
11d CC = MAX( numroc ( &m, &n, &myrow, &rsrc, &nprow ), 1 );
// Initialize discriptors (local matrix A is considered as distributed with blocking
parameters
// m, n, i.e. there is only one block - whole matrix A - which is located on process
descinit (descAA, &m, &k, &m, &k, &rsrc, &csrc, &ictxt, &lld AA, &info );
descinit (descBB, &k, &n, &k, &n, &rsrc, &csrc, &ictxt, &lld BB, &info );
descinit (descCC, &m, &n, &m, &n, &rsrc, &csrc, &ictxt, &lld CC, &info );
descinit_(descA, &m, &k, &mb, &nb, &rsrc, &csrc, &ictxt, &mA, &info);
descinit_(descB, &k, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &kB, &info);
descinit (descC, &m, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &mC, &info);
A = (double*) malloc(mA*kA*sizeof(double)); //matricea locala
B = (double*) malloc(kB*nB*sizeof(double)); //matricea locala
C = (double*) malloc(mC*nC*sizeof(double)); //matricea locala
// Call pdgeadd to distribute matrix (i.e. copy AA into A and BB into B)
pdgeadd ( "N", &m, &k, &one, AA, &i one, &i one, descAA, &zero, A, &i one, &i one,
descA);
pdgeadd ( "N", &k, &n, &one, BB, &i one, &i one, descBB, &zero, B, &i one, &i one,
descB);
// Tipar matricele locale A si B
for (int id = 0; id < nprocs; ++id)</pre>
Cblacs barrier(ictxt, "All");
if (id == iam) {
printf("Local A(%d*%d) on node %d (%d,%d) \n", kA, mA, iam, myrow, mycol);
//cout << "A on node " << iam << endl;
//for (i = 0; i < mA; ++i)
 for (j = 0; j < kA; ++j)
 //for (j = 0; j < kA; ++j)
 for (i = 0; i < mA; ++i)
 cout << setw(3) << *(A+mA*j+i) << " ";
 cout << endl;
printf("Local B(%d*%d) on node %d (%d,%d) \n", nB,kB,iam,myrow,mycol);
//cout << "B on node " << iam << endl;
//for (i = 0; i < kB; ++i)
for (j = 0; j < nB; ++j)
//for (j = 0; j < nB; ++j)
for (i = 0; i < kB; ++i)
cout << setw(3) << *(B+kB*j+i) << " ";
cout << endl;
}
cout << endl;
Cblacs barrier(ictxt, "All");
alpha = 1.0; beta = 0.0;
pdgemm ("No Transpose", "No
Transpose", &m, &n, &k, &alpha, A, &ONE, &ONE, descA, B, &ONE, &ONE, descB, &beta, C, &ONE, &ONE, desc
C);
Cblacs barrier(ictxt, "All");
// Print out the matrix product C
```

```
for (int id = 0; id < nprocs; ++id)</pre>
Cblacs barrier(ictxt, "All");
if (id == iam) {
printf("Local C(%d*%d) on node %d (%d,%d) \n", mC,nC,iam,myrow,mycol);
//cout << "C loc on node " << iam << endl;</pre>
for (i = 0; i < mC; ++i)
 for (j = 0; j < nC; ++j)
 cout << setw(3) << *(C+mC*j+i) << " ";
 cout << endl;
Cblacs barrier(ictxt, "All");
// Se construieste matricea CC
pdgeadd ( "N", &m, &n, &one, C, &i one, &i one, descC, &zero, CC, &i one, &i one,
descCC);
//*** Tipar rezultate finale
if (iam==0)
 {
   cout << "Global matrix CC =AA*BB:\n";</pre>
              //for (i = 0; i < m; ++i)
            for (j = 0; j < n; ++j) {
                  //for (j = 0; j < n; ++j)
                  for (i = 0; i < m; ++i) {
                      cout << setw(5) << *(CC + n*i + j) << " ";
                  cout << "\n";
              cout << endl;
/* Free memory */
    free(A);
    free(B);
    free(C);
if( myrow==0 && mycol==0 ) {
free( AA );
free (BB);
free ( CC);
    /* Release process grid */
    Cblacs gridexit(ictxt);
} /* main */
```

The results are the following:

```
20 21 22 23 24
 30 31 32 33 34
 40 41 42 43 44
Global matrix BB:
 5 6 7 8 9
15 16 17 18 19
25 26 27 28 29
35 36 37 38
               39
45 46 47 48 49
Local A(3*3) on node 0 (0,0)
0 1 4
10 11 14
40 41 44
Local B(3*3) on node O(0,0)
 5 6 9
15 16 19
45 46 49
Local A(2*3) on node 1 (0,1)
20 21 24
30 31 34
Local B(2*3) on node 1 (0,1)
25 26 29
35 36 39
Local A(3*2) on node 2 (1,0)
 2 3
12 13
42 43
Local B(3*2) on node 2 (1,0)
 7 8
17 18
47 48
Local A(2*2) on node 3 (1,1)
22 23
32 33
Local B(2*2) on node 3 (1,1)
27 28
37 38
Local C(3*3) on node O(0,0)
800 1800 4800
835 1885 5035
940 2140 5740
Local C(3*2) on node 1 (0,1)
2800 3800
2935 3985
3340 4540
Local C(2*3) on node 2 (1,0)
870 1970 5270
905 2055 5505
Local C(2*2) on node 3 (1,1)
3070 4170
3205 4355
Global matrix CC =AA*BB:
 800 1800 2800 3800 4800
 835 1885 2935 3985 5035
 870 1970 3070 4170 5270
```

```
905 2055 3205 4355 5505
940 2140 3340 4540 5740
```

The accuracy of the results is confirmed by the sequential program:

```
[MI gr TPS1@hpc]$ ./mpiCC ScL -o dgemm1.exe dgemm1.cpp
[MI gr TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 1 -host compute-0-0,compute-0-4
dgemm1.exe
Global matrix AA:
 0 1 2 3 4
10 11 12 13 14
20 21 22 23 24
30 31 32 33 34
40 41 42 43 44
Global matrix BB:
 5 6 7 8 9
15 16 17 18 19
25 26 27 28 29
35 36 37 38 39
45 46 47 48 49
Global matrix CC:
800 1800 2800 3800 4800
835 1885 2935 3985 5035
870 1970 3070 4170 5270
905 2055 3205 4355 5505
940 2140 3340 4540 5740
```

Purpose: PvSYMM performs one of the distributed matrix-matrix operations $sub(C) := alpha \times sub(A) \times sub(B) + beta \times sub(C)$, or $sub(C) := alpha \times sub(B) \times sub(A) + beta \times sub(C)$, where sub(C) denotes C(IC:IC+M-1,JC:JC+N-1),

sub(B) denotes B(IB:IB+M-1,JB:JB+N-1).

Alpha and beta are scalars, sub(A) is a symmetric distributed matrix and sub(B) and sub(C) are M-by-N distributed matrices.

Arguments

```
SIDE (global input) CHARACTER

On entry, SIDE specifies whether the symmetric distributed matrix sub( A )

appears on the left or right in the operation as follows:

SIDE = 'L' sub( C ) := alpha*sub( A )*sub( B ) + beta*sub( C ),

SIDE = 'R' sub( C ) := alpha*sub( B )*sub( A ) + beta*sub( C ),

UPLO (global input) CHARACTER

On entry, UPLO specifies whether the upper or lower triangular part of the
```

symmetric distributed matrix sub(A) is to be referenced.

M (global input) INTEGER

The number of rows to be operated on i.e., the number of rows of the distributed submatrix sub(C). $M \ge 0$.

N (global input) INTEGER

The number of columns to be operated on i.e the number of columns of the distributed submatrix sub(C). $N \ge 0$.

ALPHA (global input) REAL/COMPLEX

On entry, ALPHA specifies the scalar alpha.

A (local input) array of dimension (LLD_A, $LOC_q(JA+NA-1)$)

Before entry this array contains the local pieces of the symmetric distributed matrix sub(A), such that when UPLO = 'U', the NA-by-NA upper triangular part of the distributed matrix sub(A) must contain the upper triangular part of the symmetric distributed matrix and the strictly lower triangular part of sub(A) is not referenced, and when UPLO = 'L', the NA-by-NA lower triangular part of the distributed matrix sub(A) must contain the lower triangular part of the symmetric distributed matrix and the strictly lower triangular part of sub(A) is not referenced.

IA (global input) INTEGER

The global row index of the submatrix of the distributed matrix A to operate on.

JA (global input) INTEGER

The global column index of the submatrix of the distributed matrix A to operate on.

DESCA (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix A.

B (local input) array of dimension (LLD B, $LOC_q(JB+N-1)$)

Before entry, this array contains the local pieces of the distributed matrix sub(B). IB (global input) INTEGER

The global row index of the submatrix of the distributed matrix B to operate on. JB (global input) INTEGER

The global column index of the submatrix of the distributed matrix B to operate

DESCB (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix B.

BETA (global input) REAL/COMPLEX

On entry, BETA specifies the scalar beta. When BETA is supplied as zero then sub(C) need not be set on input.

C (local input/local output) array of dimension (LLD_C, LOCq(JC+N1))

Before entry, this array must contain the local pieces of the distributed matrix $sub(\ C\)$. On exit, the distributed matrix $sub(\ C\)$ is overwritten by the M-by-N updated distributed matrix.

IC (global input) INTEGER

The global row index of the submatrix of the distributed matrix C to operate on.

JC (global input) INTEGER

The global column index of the submatrix of the distributed matrix C to operate on.

DESCC (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix C.

Purpose: PvTRAN transposes a distributed matrix $sub(C) = beta \times sub(C) +$

alpha x op(sub(A)) where sub(C) denotes C(IC:IC+M-1,JC:JC+N-1), sub(A) denotes A(IA:IA+N-1,JA:JA+M-1), op(A) denotes AT. Beta is a scalar, sub(C) is an M-by-N distributed matrix, sub(A) is an N-by-M distributed matrix.

Arguments

M (global input) INTEGER

The number of rows to be operated on i.e., the number of rows of the distributed submatrix sub(C). $M \ge 0$.

N (global input) INTEGER

The number of columns to be operated on i.e the number of columns of the distributed submatrix sub(C). $N \ge 0$.

ALPHA (global input) REAL

On entry, ALPHA specifies the scalar alpha.

A (local input) REAL array of dimension (LLD_A, LOCq(JA+M-1))

This array contains the local pieces of the distributed matrix sub(A).

IA (global input) INTEGER

The global row index of the submatrix of the distributed matrix A to operate on.

JA (global input) INTEGER

The global column index of the submatrix of the distributed matrix A to operate on.

DESCA (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix A.

BETA (global input) REAL

On entry, BETA specifies the scalar beta. When BETA is supplied as zero then sub(C) need not be set on input.

C (local input/local output) array of dimension (LLD_C,LOCq(JC+N1))

This array contains the local pieces of the distributed matrix sub(C). On exit, the distributed matrix sub(C) is over- written by the updated matrix.

IC (global input) INTEGER

The global row index of the submatrix of the distributed matrix C to operate on.

JC (global input) INTEGER

The global column index of the submatrix of the distributed matrix C to operate on.

DESCC (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix C.

```
PvTRSM(SIDE, UPLO, TRANSA, DIAG, M, N, ALPHA, A, IA, JA, DESCA, B, IB, JB, DESCB)

void pdtrsm_(F_CHAR_T SIDE, F_CHAR_T UPLO, F_CHAR_T TRANS, F_CHAR_T DIAG, int *M, int *N, double *ALPHA, double *A, int *IA, int *JA, int *DESCA, double *B, int *IB, int *JB, int *DESCB)
```

Purpose: PvTRSM solves one of the distributed matrix equations $op(sub(A)) \times X = alpha \times sub(B)$, or $X \times op(sub(A)) = alpha \times sub(B)$, where

sub(B) denotes B(IB:IB+M-1,JB:JB+N-1), alpha is a scalar, X and sub(B) are an M-by-N distributed matrix, sub(A) is a unit, or non-unit, upper or lower triangular distributed matrix and op(A) is one of op(A) = A or $op(A) = A^T$. The distributed matrix X is overwritten on sub(B).

Arguments

SIDE (global input) CHARACTER

On entry, SIDE specifies whether op(A) appears on the left or right of X as follows:

SIDE = 'L', op(sub(A))*X = alpha*sub(B), SIDE = 'R', X*op(sub(A)) = alpha*sub(B).

UPLO (global input) CHARACTER

On entry, UPLO specifies whether the distributed matrix sub(A) is an upper or lower triangular distributed matrix.

TRANSA (global input) CHARACTER

The form of op(A) to be used in the matrix multiplication as follows:

TRANSA = 'N', op(A) = A, $TRANSA = 'T', op(A) = A^{T},$ $TRANSA = 'C', op(A) = A^{T}.$

DIAG (global input) CHARACTER

On entry, DIAG specifies whether or not sub(A) is unit triangular as follows:

DIAG = 'U', sub(A) is assumed to be unit triangular,

DIAG = 'N', sub(A) is not assumed to be unit triangular.

M (global input) INTEGER

The number of rows to be operated on i.e., the number of rows of the distributed submatrix sub(B). $M \ge 0$.

N (global input) INTEGER

The number of columns to be operated on i.e., the number of columns of the distributed submatrix sub(B). $N \ge 0$.

ALPHA (global input) REAL/COMPLEX

On entry, ALPHA specifies the scalar alpha.

A (local input) array of dimension (LLD A, LOCq(JA+NA-1)

Before entry with UPLO = 'U', the leading NA-by-NA upper triangular part of the distributed matrix sub(A) must contain the local pieces of the upper triangular distributed matrix and its strictly lower triangular part is not referenced. Before entry with UPLO = 'L', the leading NA-by-NA lower triangular part of the distributed matrix sub(A) must contain the lower triangular distributed matrix and its strictly upper triangular part is not referenced. Note that when DIAG = 'U', the diagonal elements of sub(A) are not referenced either, but are assumed to be unity.

IA (global input) INTEGER

The global row index of the submatrix of the distributed matrix A to operate on.

JA (global input) INTEGER

The global column index of the submatrix of the distributed matrix A to operate on.

DESCA (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix A.

B (local input) array of dimension (LLD B, $LOC_q(JB+N-1)$)

Before entry, this array contains the local pieces of the distributed matrix sub(B). On exit, sub(B) is overwritten by the solution distributed matrix.

IB (global input) INTEGER

The global row index of the submatrix of the distributed matrix B to operate on.

JB (global input) INTEGER

The global column index of the submatrix of the distributed matrix B to operate on.

DESCB (global and local input) INTEGER array of dimension 8

The array descriptor of the distributed matrix B.

PvSYMM(SIDE, UPLO, M, N, ALPHA, A, IA, JA, DESCA, B, IB, JB, DESCB, BETA, C, IC, JC, DESCC)

Purpose: Performs a scalar-matrix-matrix product (one matrix operand is symmetric) and adds the result to a scalar-matrix product for distribute matrices. The operation is defined as

```
sub(C):=ALPHA*sub(A)*sub(B)+BETA*sub(C) or sub(C):=ALPHA*sub(B)*sub(A)+BETA*sub(C),
```

where ALPHA and BETA are scalars, sub(A) is a symmetric distributed matrix, sub(A)=A(IA:IA+M-1,JA:JA+M-1), if SIDE='L', and sub(A)=A(IA:IA+N-1,JA:JA+N-1), if SIDE='R', sub(B) and sub(C) are M-by-N distributed matrices, sub(B)=B(IB:IB+M-1,JB:JB+N-1), sub(C)=C(IC:IC+M-1,JC:JC+N-1).

Input Parameters

SIDE (GLOBAL) CHARACTER*1.

Specifies whether the symmetric distributed matrix SUB(A) appears on the left or right in the operation:

 $if \ side = 'L' \ or 'L', \ then \ SUB(C) := ALPHA*SUB(A) *SUB(B) + BETA*SUB(C);$ $if \ SIDE = 'R' \ or 'r', \ then \ SUB(C) := ALPHA*SUB(B) *SUB(A) + BETA*SUB(C).$

UPLO (global) CHARACTER*1.

Specifies whether the upper or lower triangular part of the symmetric distributed matrix SUB(A) is used:

if UPLO = 'U' or 'U', then the upper triangular part is used;

if uplo = 'L' or 'l', then the lower triangular part is used.

M (global) INTEGER.

Specifies the number of rows of the distribute submatrix SUB(C), $M \ge 0$.

N(global) INTEGER.

Specifies the number of columns of the distribute submatrix SUB(C), $M \ge 0$.

ALPHA (global) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

Specifies the scalar ALPHA.

A (local) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM. Array, DIMENSION (LLD_A, $LOC_q(JA+NA-1)$).

Before entry this array must contain the local pieces of the symmetric distributed matrix SUB(A), such that when UPLO = 'U' or 'U', the NA-by-NA upper triangular part of the distributed matrix SUB(A) must contain the upper triangular part of the symmetric distributed matrix and the strictly lower triangular part of SUB(A) is not referenced, and when UPLO = 'L' or 'L', the NA-by-NA lower triangular part of the distributed matrix SUB(A) must contain the lower triangular part of the symmetric distributed matrix and the strictly upper triangular part of SUB(A) is not referenced.

IA, JA (global) INTEGER.

The row and column indices in the distributed matrix A indicating the first row and the first column of the submatrix SUB(A), respectively.

DESCA (global and local) INTEGER array of dimension 8.

The array descriptor of the distributed matrix A.

B (local) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

Array, DIMENSION (LLD_B, LOC_q(JB+N-1)). Before entry this array must contain the local pieces of the distributed matrix SUB(B).

IB, JB (global) INTEGER.

The row and column indices in the distributed matrix B indicating the first row and the first column of the submatrix SUB(B), respectively.

DESCB (global and local) INTEGER array of dimension 8.

The array descriptor of the distributed matrix B.

BETA (global) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

Specifies the scalar BETA. When BETA is set to zero, then SUB(C) need not be set on input.

C (local) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

Array, DIMENSION (LLD_C, LOCQ(JC+N-1)). Before entry this array must contain the local pieces of the distributed matrix SUB(C).

IC, JC (global) INTEGER.

The row and column indices in the distributed matrix c indicating the first row and the first column of the submatrix SUB(c), respectively.

DESCC (global and local) INTEGER array of dimension 8.

The array descriptor of the distributed matrix c

Output Parameters

C

Overwritten by the m-by-n updated matrix.

Purpose: The routine copies all or part of a distributed matrix A to another distributed matrix B. No communication is performed, PvLACPY performs a local copy sub(A)=sub(B), where sub(A) denotes A(IA:IA+M-1,JA:JA+N-1) and sub(B) denotes B(IB:IB+M-1,JB:JB+N-1).

Input Parameters

UPLO (global) CHARACTER.

Specifies the part of the distributed matrix sub (A) to be copied:

- = 'U': Upper triangular part; the strictly lower triangular part of sub(A) is not referenced;
- = 'L': Lower triangular part; the strictly upper triangular part of sub(A) is not referenced.

Otherwise: all of the matrix sub (A) is copied.

M (global) INTEGER.

The number of rows to be operated on, that is, the number of rows of the distributed submatrix sub(A) $(M \ge 0)$.

N (global) INTEGER.

The number of columns to be operated on, that is, the number of columns of the distributed submatrix sub(A) $(N \ge 0)$.

A (local) REAL for PSLACPY, DOUBLE PRECISION for PDLACPY, COMPLEX for PCLACPY, COMPLEX*16 for PZLACPY.

Pointer into the local memory to an array of DIMENSION(LLD_A , $LOC_c(JA+N-1)$). On entry, this array contains the local pieces of the distributed matrix sub(A).

IA, JA (global) INTEGER.

The row and column indices in the global array a indicating the first row and the first column of the submatrix sub(A), respectively.

DESCA (global and local) INTEGER array, DIMENSION (DLEN $_$).

The array descriptor for the distributed matrix A.

IB, JB (global) INTEGER.

The row and column indices in the global array ${\tt B}$ indicating the first row and the first column of ${\tt Sub}\,({\tt B})$ respectively.

DESCB (global and local) INTEGER array, DIMENSION (DLEN_).

The array descriptor for the distributed matrix A.

Output Parameters

B (local) REAL for PSLACPY, DOUBLE PRECISION for PDLACPY, COMPLEX for PCLACPY, COMPLEX*16 for PZLACPY.

Pointer into the local memory to an array of DIMENSION (LLD_B, LOCC(JB+N-1)). This array contains on exit the local pieces of the distributed matrix sub(B) set as follows:

```
if UPLO='U', B(IB+I-1,JB+J-1)=A(IA+I-1,JA+J-1), 1\le I\le J, 1\le J\le N; if UPLO='L', B(IB+I-1,JB+J-1)=A(IA+I-1,JA+J-1), J\le I\le M, 1\le J\le N; otherwise, B(IB+I-1,JB+J-1)=A(IA+I-1,JA+J-1), 1\le I\le M, 1\le J\le N.
```

The PBLAS Level 3 routines perform distributed matrix-matrix operations. Table "PBLAS Level 3 Routine Groups and Their Data Types" lists the PBLAS Level 3 routine groups and the data types associated with them.

PBLAS Level 3 Routine Groups and Their Data Types

Routine Group	Data Types	Description			
p?geadd	s, d, c, z	Distributed matrix-matrix sum of general matrices			
p?tradd	s, d, c, z	Distributed matrix-matrix sum of triangular matrices			
p?gemm	s, d, c, z	Distributed matrix-matrix product of general matrices			
p?hemm	c, z	Distributed matrix-matrix product, one matrix is Hermitian			
p?herk	c, z	Rank-k update of a distributed Hermitian matrix			
p?her2k	c, z	Rank-2k update of a distributed Hermitian matrix			
p?symm	s, d, c, z	Matrix-matrix product of distributed symmetric matrices			
p?syrk	s, d, c, z	Rank-k update of a distributed symmetric matrix			
p?syr2k	s, d, c, z	Rank-2k update of a distributed symmetric matrix			
p?tran	s, d	Transposition of a real distributed matrix			
p?tranc	c, z	Transposition of a complex distributed matrix (conjugated)			
p?tranu	c, z	Transposition of a complex distributed matrix			
p?trmm	s, d, c, z	Distributed matrix-matrix product, one matrix is triangular			
p?trsm	s, d, c, z	Solution of a distributed matrix equation, one matrix is triangular			