Lucrare de Laborator Nr. 1 "Elaborarea unui program MPI pentru generarea jocurilor bimatriceale si determinarea situatiilor Nash de echilibru in strategii pure"

Consideram un joc in urmatoarea forma normala  $\Gamma = \langle I, J, A, B \rangle$ , unde  $I = \{1, 2, ..., n\}$  este multime de indici ale liniilor,  $J = \{1, 2, ..., m\}$  este multimea de indici ale coloanelor  $A = ||a_{ij}||_{i \in I}^{j \in J}$ ,  $B = ||b_{ij}||_{i \in I}^{j \in J}$  sunt matricele de utilitate ale jucatorului 1 si r 2, respectiv.

1) Generarea jocurilor bimatriceale.

Vom construi un sir de numere  $\mathbf{i} = i_1 i_2 ... i_j ... i_m$  si  $\mathbf{j} = j_1 j_2 ... j_i ... j_n$  unde elementele  $i_j = \mathbf{i}(j)$  (respectiv  $j_i = \mathbf{j}(i)$ ) semnifica urmatoarele: daca jucatorul 2 (respectiv jucatorul 1) alege coloana  $j \in J$  (alege linia  $i \in I$ ) atunci jucatorul 1 (respectiv jucatorul 2) va alege linia  $i_j \in I$  (va alege coloana  $j_i \in J$ ). Notam  $\mathbf{I} = \{\mathbf{i} = i_1 i_2 ... i_j ... i_m\}$  si  $\mathbf{J} = \{\mathbf{j} = j_1 j_2 ... j_i ... j_n\}$  multimea tuturor sirurilor de acest tip, si  $|\mathbf{I}| = n^m$  and  $|\mathbf{J}| = m^n$ 

Pentru orice pereche  $(\mathbf{i}, \mathbf{j})$  sa se genereze matricea  $AB(\mathbf{i}, \mathbf{j}) = \|(\mathfrak{a}_{ij}, \mathfrak{b}_{ij})\|_{i \in I}^{j \in J}$ , unde  $\mathfrak{a}_{ij} \equiv a_{i_j j_i}$ ,  $\mathfrak{b}_{ij} \equiv b_{i_j j_i}$ .

Example 0.1 Consideram jocul bimatriceal 
$$A = \begin{pmatrix} 3 & 5 & 4 \\ 6 & 7 & 2 \end{pmatrix}$$
,  $B = \begin{pmatrix} 0 & 5 & 1 \\ 4 & 3 & 2 \end{pmatrix}$ . Fie  $\mathbf{I} = \{\mathbf{i}^1, \mathbf{i}^2\}$  si  $\mathbf{J} = \{\mathbf{j}^1, \mathbf{j}^2\}$ , unde  $\mathbf{i}^1 = 1_1 1_2 2_3$ ,  $\mathbf{i}^2 = 1_1 2_2 1_3$  si  $\mathbf{j}^1 = 1_1 2_2$  and  $\mathbf{j}^2 = 2_1 1_2$ .

• Pentru  $\mathbf{i}^1 = 1_1 1_2 2_3, \mathbf{j}^1 = 1_1 2_2$ :

$$\begin{array}{l} \mathfrak{a}_{11} \equiv a_{i_{1}^{1}j_{1}^{1}} = a_{11} = 3; \ \mathfrak{a}_{12} \equiv a_{i_{2}^{1}j_{1}^{1}} = a_{11} = 3; \ \mathfrak{a}_{13} \equiv a_{i_{3}^{1}j_{1}^{1}} = a_{21} = 6 \\ \mathfrak{a}_{21} \equiv a_{i_{1}^{1}j_{2}^{1}} = a_{12} = 5; \ \mathfrak{a}_{22} \equiv a_{i_{2}^{1}j_{2}^{1}} = a_{12} = 5; \ \mathfrak{a}_{23} \equiv a_{i_{3}^{1}j_{2}^{1}} = a_{22} = 7 \\ \mathfrak{b}_{11} \equiv b_{i_{1}^{1}j_{1}^{1}} = b_{11} = 0; \ \mathfrak{b}_{12} \equiv b_{i_{2}^{1}j_{1}^{1}} = b_{11} = 0; \ \mathfrak{b}_{13} \equiv a_{i_{3}^{1}j_{1}^{1}} = b_{21} = 4 \\ \mathfrak{b}_{21} \equiv b_{i_{1}^{1}j_{2}^{1}} = b_{12} = 5; \ \mathfrak{b}_{22} \equiv b_{i_{2}^{1}j_{2}^{1}} = b_{12} = 5; \ \mathfrak{b}_{23} \equiv b_{i_{3}^{1}j_{2}^{1}} = b_{22} = 3 \end{array}$$

• Pentru  $\mathbf{i}^2 = 1_1 2_2 1_3$ ,  $\mathbf{j}^1 = 1_1 2_2$ :

$$\begin{array}{l} \mathfrak{a}_{11} \equiv a_{i_{1}^{2}j_{1}^{1}} = a_{11} = 3; \ \mathfrak{a}_{12} \equiv a_{i_{2}^{2}j_{1}^{1}} = a_{21} = 6; \ a_{13} \equiv a_{i_{3}^{2}j_{1}^{1}} = a_{11} = 3 \\ \mathfrak{a}_{21} \equiv a_{i_{1}^{2}j_{2}^{1}} = a_{12} = 5; \ \mathfrak{a}_{22} \equiv a_{i_{2}^{2}j_{2}^{1}} = a_{22} = 7; \ \mathfrak{a}_{23} \equiv a_{i_{3}^{2}j_{2}^{1}} = a_{12} = 5 \\ \mathfrak{b}_{11} \equiv b_{i_{1}^{2}j_{1}^{1}} = b_{11} = 0; \ \mathfrak{b}_{12} \equiv b_{i_{2}^{2}j_{1}^{1}} = b_{21} = 4; \ \mathfrak{b}_{13} \equiv a_{i_{3}^{2}j_{1}^{1}} = b_{11} = 0 \\ \mathfrak{b}_{21} \equiv b_{i_{1}^{2}j_{2}^{1}} = b_{12} = 5; \ \mathfrak{b}_{22} \equiv b_{i_{2}^{2}j_{2}^{1}} = b_{22} = 3; \ \mathfrak{b}_{23} \equiv b_{i_{3}^{2}j_{2}^{1}} = b_{12} = 5 \end{array}$$

• Pentrur  $\mathbf{i}^1 = 1_1 1_2 2_3$ ,  $\mathbf{j}^2 = 2_1 1_2$ :

$$\begin{array}{l} \mathfrak{a}_{11} \equiv a_{i_{1}^{1}j_{1}^{2}} = a_{12} = 5; \ \mathfrak{a}_{12} \equiv a_{i_{2}^{1}j_{1}^{2}} = a_{12} = 5; \ \mathfrak{a}_{13} \equiv a_{i_{3}^{1}j_{1}^{2}} = a_{22} = 7 \\ \mathfrak{a}_{21} \equiv a_{i_{1}^{1}j_{2}^{2}} = a_{11} = 3; \ \mathfrak{a}_{22} \equiv a_{i_{2}^{1}j_{2}^{2}} = a_{11} = 3; \ \mathfrak{a}_{23} \equiv a_{i_{3}^{1}j_{2}^{2}} = a_{21} = 6 \\ \mathfrak{b}_{11} \equiv b_{i_{1}^{1}j_{1}^{2}} = b_{12} = 5; \ \mathfrak{b}_{12} \equiv b_{i_{2}^{1}j_{1}^{2}} = b_{12} = 5; \ \mathfrak{b}_{13} \equiv a_{i_{3}^{1}j_{1}^{2}} = b_{22} = 3 \\ \mathfrak{b}_{21} \equiv b_{i_{1}^{1}j_{2}^{2}} = b_{11} = 0; \ \mathfrak{b}_{22} \equiv b_{i_{2}^{1}j_{2}^{2}} = b_{11} = 0; \ \mathfrak{b}_{23} \equiv b_{i_{3}^{1}j_{2}^{2}} = b_{21} = 4 \end{array}$$

• For 
$$\mathbf{i}^2 = 1_1 2_2 1_3$$
,  $\mathbf{j}^2 = 2_1 1_2$ :  
 $\mathfrak{a}_{11} \equiv a_{i_1^2 j_1^2} = a_{12} = 5$ ;  $\mathfrak{a}_{12} \equiv a_{i_2^2 j_1^2} = a_{22} = 7$ ;  $\mathfrak{a}_{13} \equiv a_{i_3^2 j_1^2} = a_{12} = 5$   
 $\mathfrak{a}_{21} \equiv a_{i_1^2 j_2^2} = a_{11} = 3$ ;  $\mathfrak{a}_{22} \equiv a_{i_2^2 j_2^2} = a_{21} = 6$ ;  $\mathfrak{a}_{23} \equiv a_{i_3^2 j_2^2} = a_{11} = 3$   
 $\mathfrak{b}_{11} \equiv b_{i_1^2 j_1^2} = b_{12} = 5$ ;  $\mathfrak{b}_{12} \equiv b_{i_2^2 j_1^2} = b_{22} = 3$ ;  $\mathfrak{b}_{13} \equiv a_{i_3^2 j_1^2} = b_{12} = 5$   
 $\mathfrak{b}_{21} \equiv b_{i_1^2 j_2^2} = b_{11} = 0$ ;  $\mathfrak{b}_{22} \equiv b_{i_2^2 j_2^2} = b_{21} = 4$ ;  $\mathfrak{b}_{23} \equiv b_{i_3^2 j_2^2} = b_{11} = 0$ 

Astfel am obtinut urmatoarele jocuri bimatriceale

$$AB\left(\mathbf{i}^{1},\mathbf{j}^{1}\right) = \begin{pmatrix} (3,0) & (3,0) & (6,4) \\ (5,5) & (5,5) & (7,3) \end{pmatrix},$$

$$AB\left(\mathbf{i}^{2},\mathbf{j}^{1}\right) = \begin{pmatrix} (3,0) & (6,4) & (3,0) \\ (5,5) & (7,3) & (5,5) \end{pmatrix},$$

$$AB\left(\mathbf{i}^{1},\mathbf{j}^{2}\right) = \begin{pmatrix} (5,5) & (5,5) & (7,3) \\ (3,0) & (3,0) & (6,4) \end{pmatrix},$$

$$AB\left(\mathbf{i}^{2},\mathbf{j}^{2}\right) = \begin{pmatrix} (5,5) & (7,3) & (5,5) \\ (3,0) & (6,4) & (3,0) \end{pmatrix}.$$

Sa se elaboreze un program MPI care va realiza urmatoarele:

- Procesul MPI cu rancul 0 va initializa matricele initiale A si B, si va transmite tuturor proceselor din comunicatorul MPI\_COM\_WORLD aceste matrici.
- $\bullet$  Fiecare proces MPI cu rancul k, folosind un algoritm combinatorial, va genera perechea de siruri

$$(\mathbf{i}^k = i_1^k i_2^k ... i_i^k ... i_m^k, \mathbf{j}^k = j_1^k j_2^k ... j_i^k ... j_n^k)$$
.

Generarea sirurilor **i** este echivalent cu urmatoarea problema: pentru o multime de numere naturale  $\{1, 2, ..., i, ...n\}$ , sa se genereze siruri de lungimea m din aceste numere. De exemplu daca  $I = \{1, 2\}$  si  $J = \{1, 2, 3\}$  atunci

$$\mathbf{I} = \{1_1 1_2 1_3, 2_1 2_2 2_3, 1_1 1_2 2_3, 1_1 2_2 1_3, 2_1 1_2 1_3, 1_1 2_2 2_3, 2_1 1_2 2_3, 2_1 2_2 1_3\},\$$

$$\mathbf{J} = \{1_1 1_2, 2_1 2_2, 3_1 3_2, 1_1 2_2, 2_1 1_2, 1_1 3_2, 3_1 1_2, 2_1 3_2, 3_1 2_2\}.$$

So, to construct the set I for n=2, m=3, we must: a) generate the strings (1,1,1) and (2,2,2); b) having the numbers  $\{1,2\}$  to generate all the sub-strings of length 3 with the elements in this set, that is the strings (1,1,2), (1,2,1), (2,1,1), (1,2,2), (2,1,2), (2,2,1). In mathematics, a multiset (or bag) is a generalization of the concept of a set that, unlike a set, allows multiple instances of the multiset's elements. For example,  $\{a,a,b\}$  and  $\{a,b\}$  are different multisets although they are the same set. However, order is important, so  $\{a,a,b\}$  and  $\{a,b,a\}$  are the different multiset. It can be easily noticed that any informational extended strategy is nothing more than a multiset, so their generation actually consists in generating multisets .

- Fiecare proces MPI cu rancul k construieste perechea de matrici  $(A^k, B^k)$ , unde  $A^k = \|\mathfrak{a}_{ij}\|_{i \in I}^{j \in J}$ ,  $B^k = \|\mathfrak{b}_{ij}\|_{i \in I}^{j \in J}$ , si  $\mathfrak{a}_{ij} \equiv a_{i_j j_i}$ ,  $\mathfrak{b}_{ij} \equiv b_{i_j j_i}$ .
- 2) Determinarea solutiilor de tip Nash pentru jocurile generate. Vezi fisierul "Modele de programare paralela pe clustere Programare MPI.pdf", pag. 113