

## 2.9 PBLAS routines for matrix-matrix operations (Level 3)

### Scalable Universal Matrix Multiply Algorithm (SUMMA)

An alternative for Cannon algorithm represents the SUMMA algorithm:

- SUMMA = Scalable Universal Matrix Multiply Algorithm
- Slightly less efficient, but simpler and easier to generalize
- Presentation from *van de Geijn and Watts [1997]*
  - [www.netlib.org/lapack/lawns/lawn96.ps](http://www.netlib.org/lapack/lawns/lawn96.ps)
  - Similar ideas appeared many times
- Used in practice in PBLAS = Parallel BLAS
  - [www.netlib.org/lapack/lawns/lawn100.ps](http://www.netlib.org/lapack/lawns/lawn100.ps)

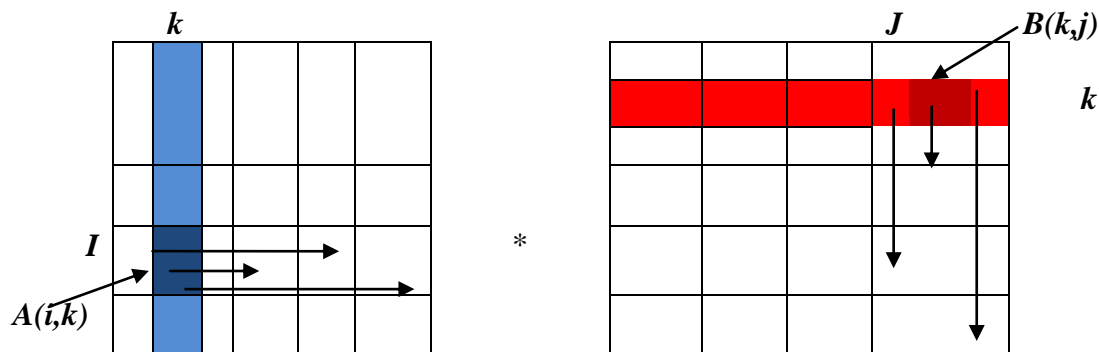
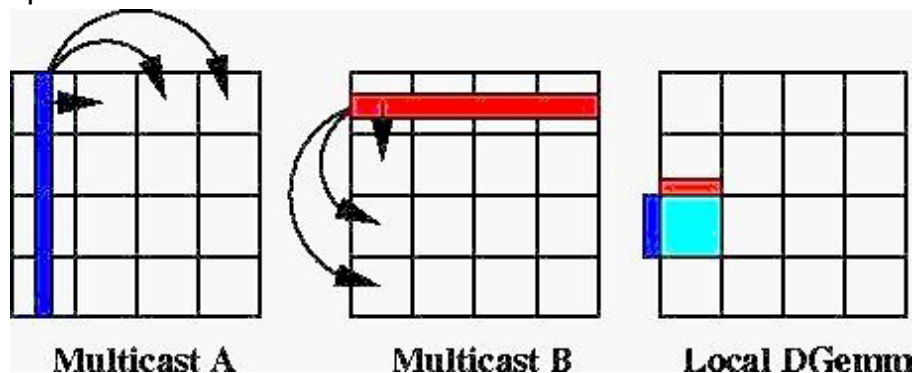
Naive matrix multiply

- For  $i = 0$  to  $n$
- For  $j = 0$  to  $n$
- For  $k = 0$  to  $n$
- $C[i,j] += A[i,k]*B[k,j]$

Calculates  $n^2$  dot products (inner products)  $C[i,j] = A[i,:]*B[:,j]$

The main steps of the algorithm:

- For each  $k$  (between 0 and  $n-1$ ),
  - Owner of partial row  $k$  broadcasts that row along its process column;
  - Owner of partial column  $k$  broadcasts that column along its process row



			C(I,J)

- $I, J$  represent all rows, columns owned by a processor
- $k$  is a single row or column (or a block of  $b$  rows or columns)
- $C(I,J) = C(I,J) + \sum_k A(I,k) * B(k,J)$

Assume a  $p_r$  by  $p_c$  processor grid ( $p_r = p_c = 4$  above)

Complete pseudo code of the algorithm.

// On each process  $P(i,j)$ :

- **For**  $k=0$  to  $n-1$  //... or  $n/b-1$  where  $b$  is the block size  
//...= # cols in  $A(I,k)$  and # rows in  $B(k,J)$ 
  - for all  $I = 1$  to  $p_r$  //... in parallel
  - owner of  $A(I,k)$  broadcasts it to whole processor row  
[adica, procesul  $p_r$  transmite  $A(I,k)$  tuturor proceselor de pe linia  $r$ ]
  - for all  $J = 1$  to  $p_c$  //... in parallel
  - owner of  $B(k,J)$  broadcasts it to whole processor column  
[adica, procesul  $p_c$  transmite  $B(k,J)$  tuturor proceselor de pe coloana  $c$ ]
  - Receive  $A(I,k)$  into Acol
  - Receive  $B(k,J)$  into Brow
  - $C(\text{myproc}, \text{myproc}) = C(\text{myproc}, \text{myproc}) + \text{Acol} * \text{Brow}$
- **Endfor**

We consider the following two matrices

A <sub>11</sub>	A <sub>12</sub>	A <sub>13</sub>	A <sub>14</sub>	A <sub>15</sub>	A <sub>16</sub>	A <sub>17</sub>	A <sub>18</sub>	A <sub>19</sub>
A <sub>21</sub>	A <sub>22</sub>	A <sub>23</sub>	A <sub>24</sub>	A <sub>25</sub>	A <sub>26</sub>	A <sub>27</sub>	A <sub>28</sub>	A <sub>29</sub>
A <sub>31</sub>	A <sub>32</sub>	A <sub>33</sub>	A <sub>34</sub>	A <sub>35</sub>	A <sub>36</sub>	A <sub>37</sub>	A <sub>38</sub>	A <sub>39</sub>
A <sub>41</sub>	A <sub>42</sub>	A <sub>43</sub>	A <sub>44</sub>	A <sub>45</sub>	A <sub>46</sub>	A <sub>47</sub>	A <sub>48</sub>	A <sub>49</sub>
A <sub>51</sub>	A <sub>52</sub>	A <sub>53</sub>	A <sub>54</sub>	A <sub>55</sub>	A <sub>56</sub>	A <sub>57</sub>	A <sub>58</sub>	A <sub>59</sub>
A <sub>61</sub>	A <sub>62</sub>	A <sub>63</sub>	A <sub>64</sub>	A <sub>65</sub>	A <sub>66</sub>	A <sub>67</sub>	A <sub>68</sub>	A <sub>69</sub>
A <sub>71</sub>	A <sub>72</sub>	A <sub>73</sub>	A <sub>74</sub>	A <sub>75</sub>	A <sub>76</sub>	A <sub>77</sub>	A <sub>78</sub>	A <sub>79</sub>
A <sub>81</sub>	A <sub>82</sub>	A <sub>83</sub>	A <sub>84</sub>	A <sub>85</sub>	A <sub>86</sub>	A <sub>87</sub>	A <sub>88</sub>	A <sub>89</sub>
A <sub>91</sub>	A <sub>92</sub>	A <sub>93</sub>	A <sub>94</sub>	A <sub>95</sub>	A <sub>96</sub>	A <sub>97</sub>	A <sub>98</sub>	A <sub>99</sub>

B <sub>11</sub>	B <sub>12</sub>	B <sub>13</sub>	B <sub>14</sub>	B <sub>15</sub>	B <sub>16</sub>	B <sub>17</sub>	B <sub>18</sub>	B <sub>19</sub>
B <sub>21</sub>	B <sub>22</sub>	B <sub>23</sub>	B <sub>24</sub>	B <sub>25</sub>	B <sub>26</sub>	B <sub>27</sub>	B <sub>28</sub>	B <sub>29</sub>
B <sub>31</sub>	B <sub>32</sub>	B <sub>33</sub>	B <sub>34</sub>	B <sub>35</sub>	B <sub>36</sub>	B <sub>37</sub>	B <sub>38</sub>	B <sub>39</sub>
B <sub>41</sub>	B <sub>42</sub>	B <sub>43</sub>	B <sub>44</sub>	B <sub>45</sub>	B <sub>46</sub>	B <sub>47</sub>	B <sub>48</sub>	B <sub>49</sub>

B <sub>51</sub>	B <sub>52</sub>	B <sub>53</sub>	B <sub>54</sub>	B <sub>55</sub>	B <sub>56</sub>	B <sub>57</sub>	B <sub>58</sub>	B <sub>59</sub>
B <sub>61</sub>	B <sub>62</sub>	B <sub>63</sub>	B <sub>64</sub>	B <sub>65</sub>	B <sub>66</sub>	B <sub>67</sub>	B <sub>68</sub>	B <sub>69</sub>
B <sub>71</sub>	B <sub>72</sub>	B <sub>73</sub>	B <sub>74</sub>	B <sub>75</sub>	B <sub>76</sub>	B <sub>77</sub>	B <sub>78</sub>	B <sub>79</sub>
B <sub>81</sub>	B <sub>82</sub>	B <sub>83</sub>	B <sub>84</sub>	B <sub>85</sub>	B <sub>86</sub>	B <sub>87</sub>	B <sub>88</sub>	B <sub>89</sub>
B <sub>91</sub>	B <sub>92</sub>	B <sub>93</sub>	B <sub>94</sub>	B <sub>95</sub>	B <sub>96</sub>	B <sub>97</sub>	B <sub>98</sub>	B <sub>99</sub>

As a result of distribution of these matrices using “2-D ciclic” algorithm on the grid of 2x3 processes (2 lines and 3 columns) with block dimensions 2x2, we obtain:

Matrix <b>A</b>	Matrix <b>B</b>
$\mathbf{A}_{(0,0)} = \begin{pmatrix} A_{11} & A_{12} & A_{17} & A_{18} \\ A_{21} & A_{22} & A_{27} & A_{28} \\ A_{51} & A_{52} & A_{57} & A_{58} \\ A_{61} & A_{62} & A_{67} & A_{68} \\ A_{91} & A_{92} & A_{97} & A_{98} \end{pmatrix}$	$\mathbf{B}_{(0,0)} = \begin{pmatrix} B_{11} & B_{12} & B_{17} & B_{18} \\ B_{21} & B_{22} & B_{27} & B_{28} \\ B_{51} & B_{52} & B_{57} & B_{58} \\ B_{61} & B_{62} & B_{67} & B_{68} \\ B_{91} & B_{92} & B_{97} & B_{98} \end{pmatrix}$
$\mathbf{A}_{(0,1)} = \begin{pmatrix} A_{13} & A_{14} & A_{19} \\ A_{23} & A_{24} & A_{29} \\ A_{53} & A_{54} & A_{59} \\ A_{63} & A_{64} & A_{69} \\ A_{93} & A_{94} & A_{99} \end{pmatrix}$	$\mathbf{B}_{(0,1)} = \begin{pmatrix} B_{13} & B_{14} & B_{19} \\ B_{23} & B_{24} & B_{29} \\ B_{53} & B_{54} & B_{59} \\ B_{63} & B_{64} & B_{69} \\ B_{93} & B_{94} & B_{99} \end{pmatrix}$
$\mathbf{A}_{(0,2)} = \begin{pmatrix} A_{15} & A_{16} \\ A_{25} & A_{26} \\ A_{55} & A_{56} \\ A_{65} & A_{66} \\ A_{95} & A_{96} \end{pmatrix}$	$\mathbf{B}_{(0,2)} = \begin{pmatrix} B_{15} & B_{16} \\ B_{25} & B_{26} \\ B_{55} & B_{56} \\ B_{65} & B_{66} \\ B_{95} & B_{96} \end{pmatrix}$
$\mathbf{A}_{(1,0)} = \begin{pmatrix} A_{31} & A_{32} & A_{37} & A_{38} \\ A_{41} & A_{42} & A_{47} & A_{48} \\ A_{71} & A_{72} & A_{77} & A_{78} \\ A_{81} & A_{82} & A_{87} & A_{88} \end{pmatrix}$	$\mathbf{B}_{(1,0)} = \begin{pmatrix} B_{31} & B_{32} & B_{37} & B_{38} \\ B_{41} & B_{42} & B_{47} & B_{48} \\ B_{71} & B_{72} & B_{77} & B_{78} \\ B_{81} & B_{82} & B_{87} & B_{88} \end{pmatrix}$

$\mathbf{A}_{(1,1)} = \begin{pmatrix} A_{33} & A_{34} & A_{39} \\ A_{43} & A_{44} & A_{49} \\ A_{73} & A_{74} & A_{79} \\ A_{83} & A_{84} & A_{89} \end{pmatrix}$	$\mathbf{B}_{(1,1)} = \begin{pmatrix} B_{33} & B_{34} & B_{39} \\ B_{43} & B_{44} & B_{49} \\ B_{73} & B_{74} & B_{79} \\ B_{83} & B_{84} & B_{89} \end{pmatrix}$
$\mathbf{A}_{(1,2)} = \begin{pmatrix} A_{35} & A_{36} \\ A_{45} & A_{46} \\ A_{75} & A_{76} \\ A_{85} & A_{86} \end{pmatrix}$	$\mathbf{B}_{(1,2)} = \begin{pmatrix} B_{35} & B_{36} \\ B_{45} & B_{46} \\ B_{75} & B_{76} \\ B_{85} & B_{86} \end{pmatrix}$

So in a compact form we have

	0	1	2
0	$\mathbf{A}_{(0,0)}$	$\mathbf{A}_{(0,1)}$	$\mathbf{A}_{(0,2)}$
1	$\mathbf{A}_{(1,0)}$	$\mathbf{A}_{(1,1)}$	$\mathbf{A}_{(1,2)}$

	0	1	2
0	$\mathbf{B}_{(0,0)}$	$\mathbf{B}_{(0,1)}$	$\mathbf{B}_{(0,2)}$
1	$\mathbf{B}_{(1,0)}$	$\mathbf{B}_{(1,1)}$	$\mathbf{B}_{(1,2)}$

To determine the product  $\mathbf{C}=\mathbf{AB}$  we use the algorithm SUMMA. Processes will determine the following submatrices of the matrix  $\mathbf{C}$ :

$$\mathbf{C}_{(0,0)} = \begin{pmatrix} C_{11} & C_{12} & C_{17} & C_{18} \\ C_{21} & C_{22} & C_{27} & C_{28} \\ C_{51} & C_{52} & C_{57} & C_{68} \\ C_{61} & C_{62} & C_{67} & C_{68} \\ C_{91} & C_{92} & C_{97} & C_{98} \end{pmatrix} \quad \mathbf{C}_{(0,1)} = \begin{pmatrix} C_{13} & C_{14} & C_{19} \\ C_{23} & C_{24} & C_{29} \\ C_{53} & C_{54} & C_{59} \\ C_{63} & C_{64} & C_{69} \\ C_{93} & C_{94} & C_{99} \end{pmatrix} \quad \mathbf{C}_{(0,2)} = \begin{pmatrix} C_{15} & C_{16} \\ C_{25} & C_{26} \\ C_{55} & C_{56} \\ C_{65} & C_{66} \\ C_{95} & C_{96} \end{pmatrix}$$

$$\mathbf{C}_{(1,0)} = \begin{pmatrix} C_{31} & C_{32} & C_{37} & C_{38} \\ C_{41} & C_{42} & C_{47} & C_{48} \\ C_{71} & C_{72} & C_{77} & C_{78} \\ C_{81} & C_{82} & C_{87} & C_{88} \end{pmatrix} \quad \mathbf{C}_{(1,1)} = \begin{pmatrix} C_{33} & C_{34} & C_{39} \\ C_{43} & C_{44} & C_{49} \\ C_{73} & C_{74} & C_{79} \\ C_{83} & C_{84} & C_{89} \end{pmatrix} \quad \mathbf{C}_{(1,2)} = \begin{pmatrix} C_{35} & C_{36} \\ C_{45} & C_{46} \\ C_{75} & C_{76} \\ C_{85} & C_{86} \end{pmatrix}$$

Here, for example, the process (0,0) can do:

$$\begin{aligned} C_{11} &= A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} + A_{17}B_{71} + A_{18}B_{81} + A_{19}B_{91} \\ C_{12} &= A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} + A_{17}B_{72} + A_{18}B_{82} + A_{19}B_{92} \\ C_{17} &= A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} + A_{17}B_{77} + A_{18}B_{87} + A_{19}B_{97} \end{aligned}$$

$$C_{18}=A_{11}B_{18} + A_{12}B_{28}+ A_{13}B_{38} +A_{14}B_{48} +A_{15}B_{58}+ A_{16}B_{68}+ A_{17}B_{78} +A_{18}B_{88} +A_{19}B_{98}$$

$$C_{21}=A_{21}B_{11} + A_{22}B_{21}+ A_{23}B_{31} +A_{24}B_{41} +A_{25}B_{51}+ A_{26}B_{61}+ A_{27}B_{71} +A_{28}B_{81} +A_{29}B_{91}$$

$$C_{22}=A_{21}B_{12} + A_{22}B_{22}+ A_{23}B_{32} +A_{24}B_{42} +A_{25}B_{52}+ A_{26}B_{62}+ A_{27}B_{72} +A_{28}B_{82} +A_{29}B_{92}$$

$$C_{27}=A_{21}B_{17} + A_{22}B_{27}+ A_{23}B_{37} +A_{24}B_{47} +A_{25}B_{57}+ A_{26}B_{67}+ A_{27}B_{77} +A_{28}B_{87} +A_{29}B_{97}$$

$$C_{28}=A_{21}B_{18}+ A_{22}B_{28}+ A_{23}B_{38}+A_{24}B_{48} +A_{25}B_{58}+ A_{26}B_{68}+ A_{27}B_{77} +A_{28}B_{88} +A_{29}B_{98}$$

$$C_{51}=A_{51}B_{11} + A_{52}B_{21}+ A_{53}B_{31} +A_{54}B_{41} +A_{55}B_{51}+ A_{56}B_{61}+ A_{57}B_{71} +A_{58}B_{81} +A_{59}B_{91}$$

$$C_{52}=A_{51}B_{12} + A_{52}B_{22}+ A_{53}B_{32} +A_{54}B_{42} +A_{55}B_{52}+ A_{56}B_{62}+ A_{57}B_{72} +A_{58}B_{82} +A_{59}B_{92}$$

$$C_{57}=A_{51}B_{17} + A_{52}B_{27}+ A_{53}B_{37} +A_{54}B_{47} +A_{55}B_{57}+ A_{56}B_{67}+ A_{57}B_{77} +A_{58}B_{87} +A_{59}B_{97}$$

$$C_{58}=A_{51}B_{18}+ A_{52}B_{28}+ A_{53}B_{38}+A_{54}B_{48} +A_{55}B_{58}+ A_{56}B_{68}+ A_{57}B_{78} +A_{58}B_{88} +A_{59}B_{98}$$

$$C_{61}=A_{61}B_{11} + A_{62}B_{21}+ A_{63}B_{31} +A_{64}B_{41} +A_{65}B_{51}+ A_{66}B_{61}+ A_{67}B_{71} +A_{68}B_{81} +A_{69}B_{91}$$

$$C_{62}=A_{61}B_{12} + A_{62}B_{22}+ A_{63}B_{32} +A_{64}B_{42} +A_{65}B_{52}+ A_{66}B_{62}+ A_{67}B_{72} +A_{68}B_{82} +A_{69}B_{92}$$

$$C_{67}=A_{61}B_{17} + A_{62}B_{27}+ A_{63}B_{37} +A_{64}B_{47} +A_{65}B_{57}+ A_{66}B_{67}+ A_{67}B_{77} +A_{68}B_{87} +A_{69}B_{97}$$

$$C_{68}=A_{61}B_{18}+ A_{62}B_{28}+ A_{63}B_{38}+A_{64}B_{48} +A_{65}B_{58}+ A_{66}B_{68}+ A_{67}B_{77} +A_{68}B_{88} +A_{69}B_{98}$$

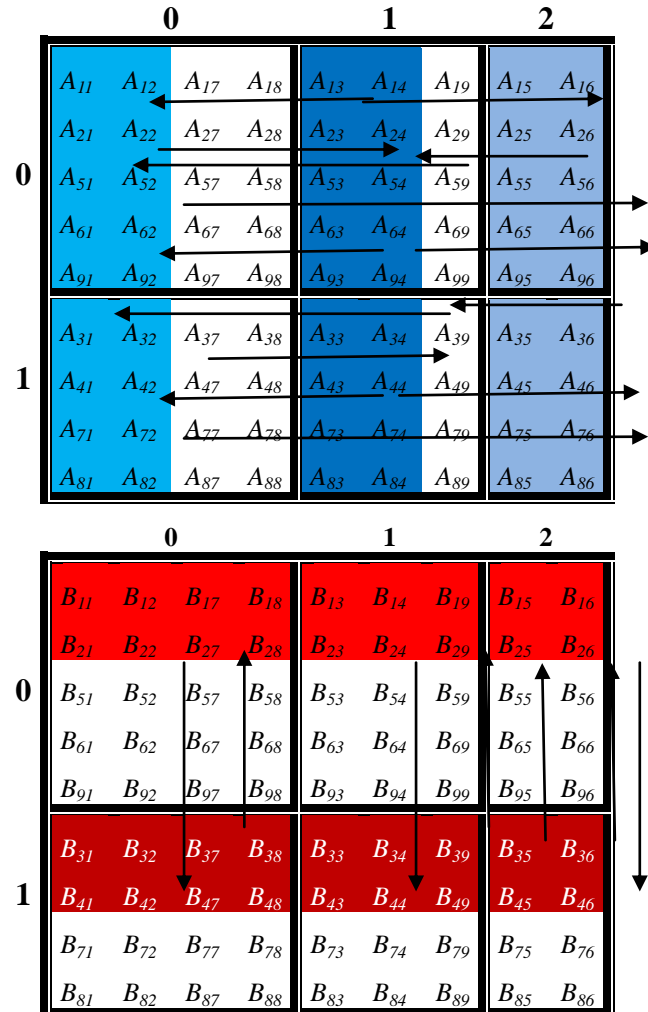
$$C_{91}=A_{91}B_{11} + A_{92}B_{21}+ A_{93}B_{31} +A_{94}B_{41} +A_{95}B_{51}+ A_{96}B_{61}+ A_{97}B_{71} +A_{98}B_{81} +A_{99}B_{91}$$

$$C_{92}=A_{91}B_{12} + A_{92}B_{22}+ A_{93}B_{32} +A_{94}B_{42} +A_{95}B_{52}+ A_{96}B_{62}+ A_{97}B_{72} +A_{98}B_{82} +A_{99}B_{92}$$

$$C_{97}=A_{91}B_{17} + A_{92}B_{27}+ A_{93}B_{37} +A_{94}B_{47} +A_{95}B_{57}+ A_{96}B_{67}+ A_{97}B_{77} +A_{98}B_{87} +A_{99}B_{97}$$

$$C_{98}=A_{91}B_{18}+ A_{92}B_{28}+ A_{93}B_{38}+A_{94}B_{48} +A_{95}B_{58}+ A_{96}B_{68}+ A_{97}B_{78} +A_{98}B_{88} +A_{99}B_{98}$$

We apply the algorithm SUMMA. The initial situation is:

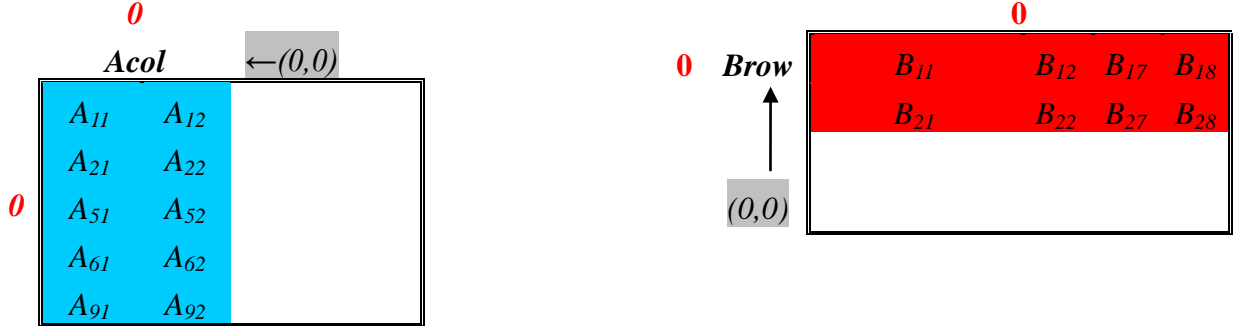


## The process (0,0)

On the base of algorithm SUMMA we obtain:

Let  $k=2$  (the block length, i.e. there are transmitted two rows and two columns).

- Iteration 0. There is no data transmission



Then it is possible to calculate the product  $C_{(0,0)} = C_{(0,0)} + Acol * Brow$ . So we obtain

$C_{(0,0)}^0 =$	$C_{11}^0 = A_{11}B_{11} + A_{12}B_{21}$	$C_{12}^0 = A_{11}B_{12} + A_{12}B_{22}$	$C_{17}^0 = A_{11}B_{17} + A_{12}B_{27}$	$C_{18}^0 = A_{11}B_{18} + A_{12}B_{28}$
	$C_{21}^0 = A_{21}B_{11} + A_{22}B_{21}$	$C_{22}^0 = A_{21}B_{12} + A_{22}B_{22}$	$C_{27}^0 = A_{21}B_{17} + A_{22}B_{27}$	$C_{28}^0 = A_{21}B_{18} + A_{22}B_{28}$
	$C_{51}^0 = A_{51}B_{11} + A_{52}B_{21}$	$C_{52}^0 = A_{51}B_{12} + A_{52}B_{22}$	$C_{57}^0 = A_{51}B_{17} + A_{52}B_{27}$	$C_{58}^0 = A_{51}B_{18} + A_{52}B_{28}$
	$C_{61}^0 = A_{61}B_{11} + A_{62}B_{21}$	$C_{62}^0 = A_{61}B_{12} + A_{62}B_{22}$	$C_{67}^0 = A_{61}B_{17} + A_{62}B_{27}$	$C_{68}^0 = A_{61}B_{18} + A_{62}B_{28}$
	$C_{91}^0 = A_{91}B_{11} + A_{92}B_{21}$	$C_{92}^0 = A_{91}B_{12} + A_{92}B_{22}$	$C_{97}^0 = A_{91}B_{17} + A_{92}B_{27}$	$C_{98}^0 = A_{91}B_{18} + A_{92}B_{28}$

Finally,  $C_{(0,0)} = C_{(0,0)}^0$

Therefore, at the end of this iteration the process (0,0) already has calculated

$$\begin{aligned}C_{11} &= A_{11}B_{11} + A_{12}B_{21} \\C_{12} &= A_{11}B_{12} + A_{12}B_{22} \\C_{17} &= A_{11}B_{17} + A_{12}B_{27} \\C_{18} &= A_{11}B_{18} + A_{12}B_{28}\end{aligned}$$

$$\begin{aligned}C_{21} &= A_{21}B_{11} + A_{22}B_{21} \\C_{22} &= A_{21}B_{12} + A_{22}B_{22} \\C_{27} &= A_{21}B_{17} + A_{22}B_{27} \\C_{28} &= A_{21}B_{18} + A_{22}B_{28}\end{aligned}$$

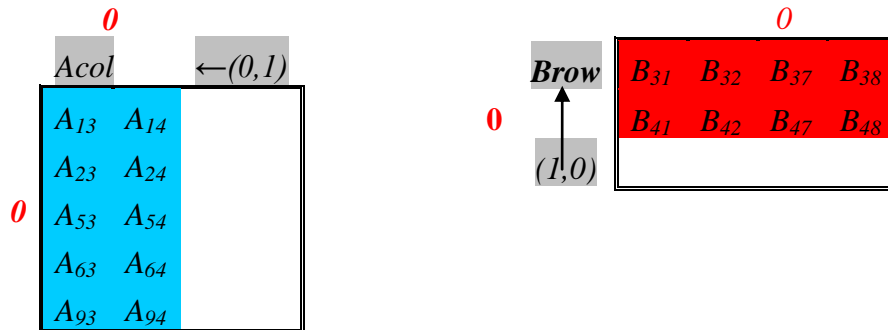
$$\begin{aligned}C_{51} &= A_{51}B_{11} + A_{52}B_{21} \\C_{52} &= A_{51}B_{12} + A_{52}B_{22}\end{aligned}$$

$$\begin{aligned}C_{57} &= A_{51}B_{17} + A_{52}B_{27} \\C_{58} &= A_{51}B_{18} + A_{52}B_{28}\end{aligned}$$

$$\begin{aligned}C_{61} &= A_{61}B_{11} + A_{62}B_{21} \\C_{62} &= A_{61}B_{12} + A_{62}B_{22} \\C_{67} &= A_{61}B_{17} + A_{62}B_{27} \\C_{68} &= A_{61}B_{18} + A_{62}B_{28}\end{aligned}$$

$$\begin{aligned}C_{91} &= A_{91}B_{11} + A_{92}B_{21} \\C_{92} &= A_{91}B_{12} + A_{92}B_{22} \\C_{97} &= A_{91}B_{17} + A_{92}B_{27} \\C_{98} &= A_{91}B_{18} + A_{92}B_{28}\end{aligned}$$

- Iteration 1. Here we have for *Acol* and *Brow* the following:



Then it is possible to calculate the product  $C_{(0,0)} = C_{(0,0)} + Acol * Brow$ . So we obtain

$C^1_{(0,0)}$	$C^1_{1,1} = A_{13}B_{31} + A_{14}B_{41}$	$C^1_{1,2} = A_{13}B_{32} + A_{14}B_{42}$	$C^1_{1,7} = A_{13}B_{37} + A_{14}B_{47}$	$C^1_{1,8} = A_{13}B_{38} + A_{14}B_{48}$
	$C^1_{2,1} = A_{23}B_{31} + A_{24}B_{41}$	$C^1_{2,2} = A_{23}B_{32} + A_{24}B_{42}$	$C^1_{2,7} = A_{23}B_{37} + A_{24}B_{47}$	$C^1_{2,8} = A_{23}B_{38} + A_{24}B_{48}$
	$C^1_{5,1} = A_{53}B_{31} + A_{54}B_{41}$	$C^1_{5,2} = A_{53}B_{32} + A_{54}B_{42}$	$C^1_{5,7} = A_{53}B_{37} + A_{54}B_{47}$	$C^1_{5,8} = A_{53}B_{38} + A_{54}B_{48}$
	$C^1_{6,1} = A_{63}B_{31} + A_{64}B_{41}$	$C^1_{6,2} = A_{63}B_{32} + A_{64}B_{42}$	$C^1_{6,7} = A_{63}B_{37} + A_{64}B_{47}$	$C^1_{6,8} = A_{63}B_{38} + A_{64}B_{48}$
	$C^1_{9,1} = A_{93}B_{31} + A_{94}B_{41}$	$C^1_{9,2} = A_{93}B_{32} + A_{94}B_{42}$	$C^1_{9,7} = A_{93}B_{37} + A_{94}B_{47}$	$C^1_{9,8} = A_{93}B_{38} + A_{94}B_{48}$

Therefore, at the end of this iteration the process (0,0) already has calculated

$$\begin{aligned}C_{11} &= A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} \\C_{12} &= A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} \\C_{17} &= A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} \\C_{18} &= A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48}\end{aligned}$$

$$\begin{aligned}C_{21} &= A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} \\C_{22} &= A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} \\C_{27} &= A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} \\C_{28} &= A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48}\end{aligned}$$

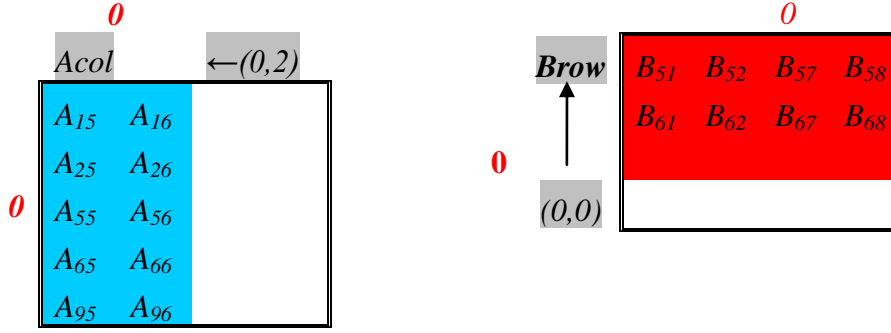
$$\begin{aligned}C_{51} &= A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} \\C_{52} &= A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} \\C_{57} &= A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47}\end{aligned}$$

$$C_{58} = A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48}$$

$$\begin{aligned}C_{61} &= A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} \\C_{62} &= A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{42} \\C_{67} &= A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} \\C_{68} &= A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48}\end{aligned}$$

$$\begin{aligned}C_{91} &= A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} \\C_{92} &= A_{91}B_{12} + A_{92}B_{22} + A_{93}B_{32} + A_{94}B_{42} \\C_{97} &= A_{91}B_{17} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} \\C_{98} &= A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48}\end{aligned}$$

- Iteration 2. Here we have for  $Acol$  and  $Bcol$  the following:



Then it is possible to calculate the product  $C_{(0,0)} = C_{(0,0)} + Acol * Brow$ . So we obtain

$C^2_{(0,0)}$	$C^2_{1,1} = A_{15}B_{51} + A_{16}B_{61}$	$C^2_{1,2} = A_{15}B_{52} + A_{16}B_{62}$	$C^2_{1,7} = A_{15}B_{57} + A_{16}B_{67}$	$C^2_{1,8} = A_{15}B_{58} + A_{16}B_{68}$
	$C^2_{2,1} = A_{25}B_{51} + A_{26}B_{61}$	$C^2_{2,2} = A_{25}B_{52} + A_{26}B_{62}$	$C^2_{2,7} = A_{25}B_{57} + A_{26}B_{67}$	$C^2_{2,8} = A_{25}B_{58} + A_{26}B_{68}$
	$C^2_{5,1} = A_{55}B_{51} + A_{56}B_{61}$	$C^2_{5,2} = A_{55}B_{52} + A_{56}B_{62}$	$C^2_{5,7} = A_{55}B_{57} + A_{56}B_{67}$	$C^2_{5,8} = A_{55}B_{58} + A_{56}B_{68}$
	$C^2_{6,1} = A_{65}B_{51} + A_{66}B_{61}$	$C^2_{6,2} = A_{65}B_{52} + A_{66}B_{62}$	$C^2_{6,7} = A_{65}B_{57} + A_{66}B_{67}$	$C^2_{6,8} = A_{65}B_{58} + A_{66}B_{68}$
	$C^2_{9,1} = A_{95}B_{51} + A_{96}B_{61}$	$C^2_{9,2} = A_{95}B_{52} + A_{96}B_{62}$	$C^2_{9,7} = A_{95}B_{57} + A_{96}B_{67}$	$C^2_{9,8} = A_{95}B_{58} + A_{96}B_{68}$

Therefore, at the end of this iteration the process  $(0,0)$  already has calculated

$$\begin{aligned} C_{11} &= A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} \\ C_{12} &= A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} \\ C_{17} &= A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} \\ C_{18} &= A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} \end{aligned}$$

$$\begin{aligned} C_{21} &= A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} \\ C_{22} &= A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} \\ C_{27} &= A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} \\ C_{28} &= A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{68} \end{aligned}$$

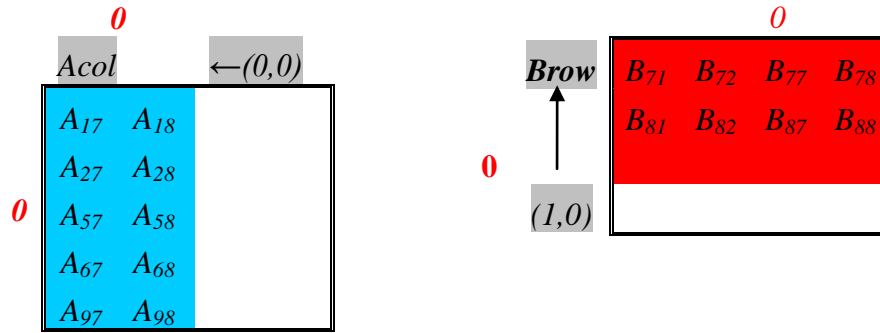
$$\begin{aligned} C_{51} &= A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} \\ C_{52} &= A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} + A_{55}B_{52} + A_{56}B_{62} \\ C_{57} &= A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} \\ C_{58} &= A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48} + A_{55}B_{58} + A_{56}B_{68} \end{aligned}$$

$$\begin{aligned} C_{61} &= A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{61} \\ C_{62} &= A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{42} + A_{65}B_{52} + A_{66}B_{62} \\ C_{67} &= A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} + A_{65}B_{57} + A_{66}B_{67} \\ C_{68} &= A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48} + A_{65}B_{58} + A_{66}B_{68} \end{aligned}$$

$$\begin{aligned} C_{91} &= A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} + A_{95}B_{51} + A_{96}B_{61} \\ C_{92} &= A_{91}B_{12} + A_{92}B_{22} + A_{93}B_{32} + A_{94}B_{42} + A_{95}B_{52} + A_{96}B_{62} \\ C_{97} &= A_{91}B_{17} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} \\ C_{98} &= A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} \end{aligned}$$

- Iteration 3. Here we have for  $Acol$  and  $Bcol$  the following:





Then it is possible to calculate the product  $C_{(0,0)} = C_{(0,0)} + A_{col} * B_{row}$ . So we obtain

$C^3_{(0,0)} =$	$C^3_{1,1} = A_{17}B_{71} + A_{18}B_{81}$	$C^3_{1,2} = A_{17}B_{72} + A_{18}B_{82}$	$C^3_{1,7} = A_{17}B_{77} + A_{18}B_{87}$	$C^3_{1,8} = A_{17}B_{78} + A_{18}B_{88}$
	$C^3_{2,1} = A_{27}B_{71} + A_{28}B_{81}$	$C^3_{2,2} = A_{27}B_{72} + A_{28}B_{82}$	$C^3_{2,7} = A_{27}B_{77} + A_{28}B_{87}$	$C^3_{2,8} = A_{27}B_{78} + A_{28}B_{88}$
	$C^3_{5,1} = A_{57}B_{71} + A_{58}B_{81}$	$C^3_{5,2} = A_{57}B_{72} + A_{58}B_{82}$	$C^3_{5,7} = A_{57}B_{77} + A_{58}B_{87}$	$C^3_{5,8} = A_{57}B_{78} + A_{58}B_{88}$
	$C^3_{6,1} = A_{67}B_{71} + A_{68}B_{81}$	$C^3_{6,2} = A_{67}B_{72} + A_{68}B_{82}$	$C^3_{6,7} = A_{67}B_{77} + A_{68}B_{87}$	$C^3_{6,8} = A_{67}B_{78} + A_{68}B_{88}$
	$C^3_{9,1} = A_{97}B_{71} + A_{98}B_{81}$	$C^3_{9,2} = A_{97}B_{72} + A_{98}B_{82}$	$C^3_{9,7} = A_{97}B_{77} + A_{98}B_{87}$	$C^3_{9,8} = A_{97}B_{78} + A_{98}B_{88}$

Therefore, at the end of this iteration the process (0,0) already has calculated

$$\begin{aligned}
 C_{11} &= A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} + A_{17}B_{71} + A_{18}B_{81} \\
 C_{12} &= A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} + A_{17}B_{72} + A_{18}B_{82} \\
 C_{17} &= A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} + A_{17}B_{77} + A_{18}B_{87} \\
 C_{18} &= A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} + A_{17}B_{78} + A_{18}B_{88}
 \end{aligned}$$

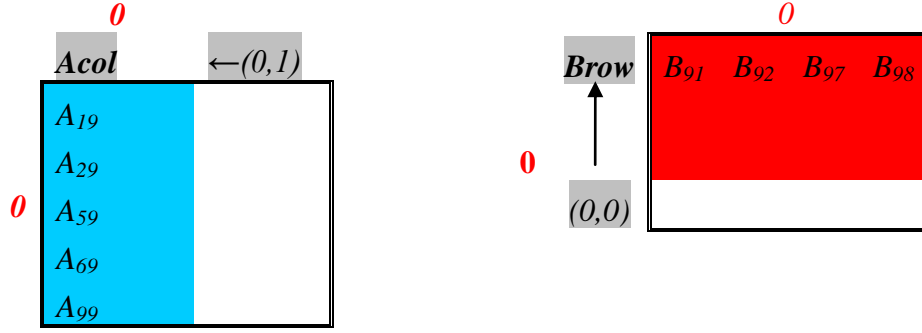
$$\begin{aligned}
 C_{21} &= A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} + A_{27}B_{71} + A_{28}B_{81} \\
 C_{22} &= A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} + A_{27}B_{72} + A_{28}B_{82} \\
 C_{27} &= A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} + A_{27}B_{77} + A_{28}B_{87} \\
 C_{28} &= A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{68} + A_{27}B_{77} + A_{28}B_{88}
 \end{aligned}$$

$$\begin{aligned}
 C_{51} &= A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} + A_{57}B_{71} + A_{58}B_{81} \\
 C_{52} &= A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} + A_{55}B_{52} + A_{56}B_{62} + A_{57}B_{72} + A_{58}B_{82} \\
 C_{57} &= A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} + A_{57}B_{77} + A_{58}B_{87} \\
 C_{58} &= A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48} + A_{55}B_{58} + A_{56}B_{68} + A_{57}B_{78} + A_{58}B_{88}
 \end{aligned}$$

$$\begin{aligned}
 C_{61} &= A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{61} + A_{67}B_{71} + A_{68}B_{81} \\
 C_{62} &= A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{42} + A_{65}B_{52} + A_{66}B_{62} + A_{67}B_{72} + A_{68}B_{82} \\
 C_{67} &= A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} + A_{65}B_{57} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{87} \\
 C_{68} &= A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48} + A_{65}B_{58} + A_{66}B_{68} + A_{67}B_{78} + A_{68}B_{88}
 \end{aligned}$$

$$\begin{aligned}
 C_{91} &= A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} + A_{95}B_{51} + A_{96}B_{61} + A_{97}B_{71} + A_{98}B_{81} \\
 C_{92} &= A_{91}B_{12} + A_{92}B_{22} + A_{93}B_{32} + A_{94}B_{42} + A_{95}B_{52} + A_{96}B_{62} + A_{97}B_{72} + A_{98}B_{82} \\
 C_{97} &= A_{91}B_{17} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} + A_{97}B_{77} + A_{98}B_{87} \\
 C_{98} &= A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} + A_{97}B_{78} + A_{98}B_{88}
 \end{aligned}$$

- Iteration 4. Here we have for  $Acol$  and  $Brow$  the following:



Then it is possible to calculate the product  $C(0,0) = C(0,0) + Acol * Brow$ . So we obtain

$C^4_{(0,0)}$	$C^4_{1,1} = A_{19}B_{91}$	$C^4_{1,2} = A_{19}B_{92}$	$C^4_{1,7} = A_{19}B_{97}$	$C^4_{1,8} = A_{19}B_{98}$
	$C^4_{2,1} = A_{29}B_{91}$	$C^4_{2,2} = A_{29}B_{92}$	$C^4_{2,7} = A_{29}B_{97}$	$C^4_{2,8} = A_{29}B_{98}$
	$C^4_{5,1} = A_{59}B_{91}$	$C^4_{5,2} = A_{59}B_{92}$	$C^4_{5,7} = A_{59}B_{97}$	$C^4_{5,8} = A_{59}B_{98}$
	$C^4_{6,1} = A_{69}B_{91}$	$C^4_{6,2} = A_{69}B_{92}$	$C^4_{6,7} = A_{69}B_{97}$	$C^4_{6,8} = A_{69}B_{98}$
	$C^4_{9,1} = A_{99}B_{91}$	$C^4_{9,2} = A_{99}B_{92}$	$C^4_{9,7} = A_{99}B_{97}$	$C^4_{9,8} = A_{99}B_{98}$

Therefore, at the end of this iteration the process  $(0,0)$  already has calculated

$$\begin{aligned}
 C_{11} &= A_{11}B_{11} + A_{12}B_{21} + A_{13}B_{31} + A_{14}B_{41} + A_{15}B_{51} + A_{16}B_{61} + A_{17}B_{71} + A_{18}B_{81} + A_{19}B_{91} \\
 C_{12} &= A_{11}B_{12} + A_{12}B_{22} + A_{13}B_{32} + A_{14}B_{42} + A_{15}B_{52} + A_{16}B_{62} + A_{17}B_{72} + A_{18}B_{82} + A_{19}B_{92} \\
 C_{17} &= A_{11}B_{17} + A_{12}B_{27} + A_{13}B_{37} + A_{14}B_{47} + A_{15}B_{57} + A_{16}B_{67} + A_{17}B_{77} + A_{18}B_{87} + A_{19}B_{97} \\
 C_{18} &= A_{11}B_{18} + A_{12}B_{28} + A_{13}B_{38} + A_{14}B_{48} + A_{15}B_{58} + A_{16}B_{68} + A_{17}B_{78} + A_{18}B_{88} + A_{19}B_{98}
 \end{aligned}$$

$$\begin{aligned}
 C_{21} &= A_{21}B_{11} + A_{22}B_{21} + A_{23}B_{31} + A_{24}B_{41} + A_{25}B_{51} + A_{26}B_{61} + A_{27}B_{71} + A_{28}B_{81} + A_{29}B_{91} \\
 C_{22} &= A_{21}B_{12} + A_{22}B_{22} + A_{23}B_{32} + A_{24}B_{42} + A_{25}B_{52} + A_{26}B_{62} + A_{27}B_{72} + A_{28}B_{82} + A_{29}B_{92} \\
 C_{27} &= A_{21}B_{17} + A_{22}B_{27} + A_{23}B_{37} + A_{24}B_{47} + A_{25}B_{57} + A_{26}B_{67} + A_{27}B_{77} + A_{28}B_{87} + A_{29}B_{97} \\
 C_{28} &= A_{21}B_{18} + A_{22}B_{28} + A_{23}B_{38} + A_{24}B_{48} + A_{25}B_{58} + A_{26}B_{68} + A_{27}B_{77} + A_{28}B_{88} + A_{29}B_{98}
 \end{aligned}$$

$$\begin{aligned}
 C_{51} &= A_{51}B_{11} + A_{52}B_{21} + A_{53}B_{31} + A_{54}B_{41} + A_{55}B_{51} + A_{56}B_{61} + A_{57}B_{71} + A_{58}B_{81} + A_{59}B_{91} \\
 C_{52} &= A_{51}B_{12} + A_{52}B_{22} + A_{53}B_{32} + A_{54}B_{42} + A_{55}B_{52} + A_{56}B_{62} + A_{57}B_{72} + A_{58}B_{82} + A_{59}B_{92} \\
 C_{57} &= A_{51}B_{17} + A_{52}B_{27} + A_{53}B_{37} + A_{54}B_{47} + A_{55}B_{57} + A_{56}B_{67} + A_{57}B_{77} + A_{58}B_{87} + A_{59}B_{97} \\
 C_{58} &= A_{51}B_{18} + A_{52}B_{28} + A_{53}B_{38} + A_{54}B_{48} + A_{55}B_{58} + A_{56}B_{68} + A_{57}B_{78} + A_{58}B_{88} + A_{59}B_{98}
 \end{aligned}$$

$$\begin{aligned}
 C_{61} &= A_{61}B_{11} + A_{62}B_{21} + A_{63}B_{31} + A_{64}B_{41} + A_{65}B_{51} + A_{66}B_{61} + A_{67}B_{71} + A_{68}B_{81} + A_{69}B_{91} \\
 C_{62} &= A_{61}B_{12} + A_{62}B_{22} + A_{63}B_{32} + A_{64}B_{42} + A_{65}B_{52} + A_{66}B_{62} + A_{67}B_{72} + A_{68}B_{82} + A_{69}B_{92} \\
 C_{67} &= A_{61}B_{17} + A_{62}B_{27} + A_{63}B_{37} + A_{64}B_{47} + A_{65}B_{57} + A_{66}B_{67} + A_{67}B_{77} + A_{68}B_{87} + A_{69}B_{97} \\
 C_{68} &= A_{61}B_{18} + A_{62}B_{28} + A_{63}B_{38} + A_{64}B_{48} + A_{65}B_{58} + A_{66}B_{68} + A_{67}B_{77} + A_{68}B_{88} + A_{69}B_{98}
 \end{aligned}$$

$$\begin{aligned}
 C_{91} &= A_{91}B_{11} + A_{92}B_{21} + A_{93}B_{31} + A_{94}B_{41} + A_{95}B_{51} + A_{96}B_{61} + A_{97}B_{71} + A_{98}B_{81} + A_{99}B_{91} \\
 C_{92} &= A_{91}B_{12} + A_{92}B_{22} + A_{93}B_{32} + A_{94}B_{42} + A_{95}B_{52} + A_{96}B_{62} + A_{97}B_{72} + A_{98}B_{82} + A_{99}B_{92} \\
 C_{97} &= A_{91}B_{17} + A_{92}B_{27} + A_{93}B_{37} + A_{94}B_{47} + A_{95}B_{57} + A_{96}B_{67} + A_{97}B_{77} + A_{98}B_{87} + A_{99}B_{97} \\
 C_{98} &= A_{91}B_{18} + A_{92}B_{28} + A_{93}B_{38} + A_{94}B_{48} + A_{95}B_{58} + A_{96}B_{68} + A_{97}B_{78} + A_{98}B_{88} + A_{99}B_{98}
 \end{aligned}$$

## Performance Analysis

On a 2-dimensional  $P \times Q$  processor grid, the communication time of SUMMA is doubled in order to broadcast  $T_B$  rows as well as  $T_A$  columns [aici indicele semnifica matricea]. Assume that the time for sending a column  $T_A$  and a row  $T_B$  to the next processor are  $t_{ca}$  and  $t_{cb}$ , respectively, and the time for multiplying  $T_A$  with  $T_B$  and adding the product to  $C$  is  $t_p$ . Using the Hockney model, transferring time is modeled by  $t_s = \alpha + \beta m$ , where  $\alpha$  is the latency for each message, and  $\beta$  is the transfer time per byte (or reciprocal of network bandwidth) we obtained that

$$t_{ca} = \alpha + \left( \frac{M}{P} k_b \right) \cdot \beta, \quad t_{cb} = \alpha + \left( \frac{N}{Q} k_b \right) \cdot \beta.$$

The  $t_p = 2 \left( \frac{M}{P} \times \frac{N}{Q} k_b \right) \gamma.$

So,

$$\begin{aligned} t_{summa}^{2D} &= K_g (2t_{ca} + 2t_{cb} + t_p) - t_{ca} + (Q - 2)t_{ca} - t_{cb} + (P - 2)t_{cb} = \\ &= K_g (2t_{ca} + 2t_{cb} + t_p) + (Q - 3)t_{ca} + (P - 3)t_{cb}. \end{aligned}$$

Here where  $\alpha$  is a communication start-up time,  $\beta$  is a data transfer time, and  $\gamma$  is a time for multiplication or addition,  $k_b$  is a block sizes,  $N$ -number of the rows in the matrix  $C$ ,  $M$ -number of the columns in the matrix  $C$ ,  $K_g = \lceil K/k_b \rceil$  -columns of blocks of  $A$  and  $K_g$  rows of blocks of  $B$ .

Can we reduce the communication cost somehow? Obvious improvement [îmbunătățire evidentă]: instead of broadcasting single rows and columns, do block rows and columns.

---

```
PvGEMM (TRANSA, TRANSB, M, N, K, ALPHA, A, IA, JA, DESCA, B, IB, JB, DESCB, BET
A, C, IC, JC, DESCC)
void pdgemm_ (char *TRANSA, char *TRANSB, int *M, int *N, int *K, double
*ALPHA, double *A, int *IA, int *JA, int *DESCA, double *B, int *
IB, int *JB, int *DESCB, double *BETA, double *C, int *IC, int *JC,
int *DESCC );
```

**Purpose:** PvGEMM performs one of the matrix-matrix operations

$$sub(C) := \alpha \times op(sub(A)) \times op(sub(B)) + \beta \times sub(C),$$

where  $sub(C)$  denotes  $C(IC:IC+M-1, JC:JC+N-1)$ ,

$$\begin{aligned} op(sub(A)) &= \left\{ \begin{array}{ll} A( IA : IA + M - 1, JA : JA + K - 1 ) & \text{if } TRANSA = 'N' \\ A( IA : IA + K - 1, JA : JA + M - 1 )' & \text{if } TRANSA = 'T' \\ A( IA : IA + K - 1, JA : JA + M - 1 )' & \text{if } TRANSA = '' \end{array} \right\} \\ op(sub(B)) &= \left\{ \begin{array}{ll} B( IB : IB + K - 1, JB : JB + N - 1 ) & \text{if } TRANSB = 'N' \\ B( IB : IB + N - 1, JB : JB + K - 1 )' & \text{if } TRANSB = 'T' \\ B( IB : IB + N - 1, JB : JB + K - 1 )' & \text{if } TRANSB = 'C' \end{array} \right\} \end{aligned}$$

$\alpha$  and  $\beta$  are scalars, and  $sub(A)$ ,  $sub(B)$  and  $sub(C)$  are distributed matrices, with  $op(sub(A))$  an  $M$ -by- $K$  distributed matrix,  $op(sub(B))$  a  $K$ -by- $N$  distributed matrix and  $sub(C)$  an  $M$ -by- $N$  distributed matrix.

Arguments

TRANSA (global input) CHARACTER

The form of  $op(A)$  to be used in the matrix multiplication as follows:

TRANSA = 'N',  $op(A) = A$ ,

TRANSA = 'T',  $op(A) = A^T$ ,

$TRANSA = 'C', op(A) = A^T$ .  
 TRANSB (global input) CHARACTER  
*The form of  $op(B)$  to be used in the matrix multiplication as follows:*  
 $TRANSB = 'N', op(B) = B$ ,  
 $TRANSB = 'T', op(B) = B^T$ ,  
 $TRANSB = 'C', op(B) = B^T$ .  
 M (global input) INTEGER  
*The number of rows of the distributed matrices  $op(sub(A))$  and  $sub(C)$ .  $M \geq 0$ .*  
 N (global input) INTEGER  
*The number of columns of the distributed matrices  $op(sub(B))$  and  $sub(C)$ .  $N \geq 0$ .*  
 K (global input) INTEGER  
*The number of columns of the distributed matrix  $op(sub(A))$  and the number of rows of the distributed matrix  $op(B)$ .  $K \geq 0$ .*  
 ALPHA (global input) REAL/COMPLEX  
*On entry, ALPHA specifies the scalar alpha.*  
 A (local input) array of dimension (LLD\_A, KLa)  
*where KLa is  $LOC_q(JA+K-1)$  when  $TRANSA = 'N'$ , and is  $LOC_q(JA+M-1)$  otherwise. Before entry, this array must contain the local pieces of the distributed matrix  $sub(A)$ .*  
 IA (global input) INTEGER  
*The global row index of the submatrix of the distributed matrix A to operate on.*  
 JA (global input) INTEGER  
*The global column index of the submatrix of the distributed matrix A to operate on.*  
 DESCA (global and local input) INTEGER array of dimension 8  
*The array descriptor of the distributed matrix A.*  
 B (local input) array of dimension (LLD\_B, KLb)  
*where KLb is  $LOC_q(JB+N-1)$  when  $TRANSB = 'N'$ , and is  $LOC_q(JB+K-1)$  otherwise. Before entry this array must contain the local pieces of the distributed matrix  $sub(B)$ .*  
 IB (global input) INTEGER  
*The global row index of the submatrix of the distributed matrix B to operate on.*  
 JB (global input) INTEGER  
*The global column index of the submatrix of the distributed matrix B to operate on.*  
 DESCB (global and local input) INTEGER array of dimension 8  
*The array descriptor of the distributed matrix B.*  
 BETA (global input) REAL/COMPLEX  
*On entry, BETA specifies the scalar beta. When BETA is supplied as zero then  $sub(Y)$  need not be set on input.*  
 C (local input/local output) array of dimension (LLD\_C,  $LOC_q(JC+N-1)$ )  
*Before entry, this array must contain the local pieces of the distributed matrix  $sub(C)$ . On exit, the distributed matrix  $sub(C)$  is overwritten by the  $M$ -by- $N$  distributed matrix  $alpha*op(sub(A))*op(sub(B)) + beta*sub(C)$ .*  
 IC (global input) INTEGER  
*The global row index of the submatrix of the distributed matrix C to operate on.*  
 JC (global input) INTEGER  
*The global column index of the submatrix of the distributed matrix C to operate on.*  
 DESCC (global and local input) INTEGER array of dimension 8  
*The array descriptor of the distributed matrix C.*

**Example 2.9.1.** The using of function `pdgemm_`

```

#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include "mpi.h"
#include <iostream>
#include <iomanip>
#include <fstream>
#include <sstream>
using namespace std;
#define AA(i,j) AA[(i)*m+(j)]
#define BB(i,j) BB[(i)*k+(j)]
#define CC(i,j) CC[(i)*m+(j)]
extern "C"
{
void Cblacs_pinfo( int* mypnum, int* nprocs);
void Cblacs_get( int context, int request, int* value);
int Cblacs_gridinit( int* context, char * order, int np_row, int np_col);
void Cblacs_gridinfo( int context, int* np_row, int* np_col, int* my_row, int*
my_col);
void Cblacs_gridexit( int context);
void Cblacs_barrier(int, const char*);
void Cblacs_exit( int error_code);
void Cblacs_pcoord(int, int, int*, int*);
int numroc_( int *n, int *nb, int *iproc, int *isrcproc, int *nprocs);
int indx12g_(int*, int*, int*, int*, int*);
void descinit_(int *desc, int *m, int *n, int *mb, int *nb, int *irsrc, int *icsrc,
int *ictxt, int *lld, int *info);
void pdgemm_(char *TRANSA,char *TRANSB,int *M,int *N,int *K,double *ALPHA, double
*A,int *IA,int *JA,int *DESCA, double * B, int * IB, int * JB, int
*DESCB,double *BETA,double *C, int *IC, int *JC, int *DESCC );
} // extern "C"
int main(int argc, char **argv)
{
int i, j;
int iam,nprocs,nprow,npcol,myrow,mycol;
int descA[9],descB[9],descC[9];
int m,n,k,mb,nb,rsrc,csrc,ictxt,llda,lldb,lldc,info;
int lm,ln,lk;
int iloc,jloc;
int ZERO=0,ONE=1;
double alpha, beta;
m=5; n=5; k=5;
mb=2; nb=2;
double *AA = (double*)malloc(m*k*sizeof(double)); // matricea "globala" pentru
inmultirea AA*BB=CC
double *BB = (double*) malloc(k*n*sizeof(double)); //matricea "globala" pentru
inmultirea AA*BB=CC
double *CC = (double*) malloc(m*n*sizeof(double)); //matricea "globala" pentru
inmultirea AA*BB=CC
// initializarea matricelor globale
for(i=0;i<m;i++)
for(j=0;j<k;j++)
AA[i*m+j]=(i+j);
for(i=0;i<k;i++)
for(j=0;j<n;j++)
BB[i*k+j]=j+i;
nprow=2; npcil=2; // Astfel, programul se executa pe 4 procese
Cblacs_pinfo(&iam,&nprocs);
Cblacs_get(-1, 0, &ictxt);

```

```

Cblacs_gridinit( &ictxt, "Row", nprow, npcol);
Cblacs_gridinfo(ictxt, &nprow, &npcol, &myrow, &mycol);
rsrc=0; csrc=0;
if (iam==0)
{
printf("===== RESULT OF THE PROGRAM %s \n", argv[0]);
cout << "Global matrix AA:\n";
    for (i = 0; i < m; ++i) {
        for (j = 0; j < n; ++j) {
            cout << setw(3) << *(AA + m*i + j) << " ";
        }
        cout << "\n";
    }
    cout << endl;
cout << "Global matrix BB:\n";
    for (i = 0; i < k; ++i) {
        for (j = 0; j < n; ++j) {
            cout << setw(3) << *(BB + k*i + j) << " ";
        }
        cout << "\n";
    }
    cout << endl;
}
Cblacs_barrier(ictxt, "All");
int mA=numroc_( &m, &mb, &myrow, &rsrc, &nprow );
int kA = numroc_( &k, &nb, &mycol, &rsrc, &npcol );
int kB = numroc_( &k, &mb, &myrow, &rsrc, &nprow );
int nB = numroc_( &n, &nb, &mycol, &rsrc, &npcol );
int mC = numroc_( &m, &mb, &myrow, &rsrc, &nprow );
int nC = numroc_( &n, &nb, &mycol, &rsrc, &nprow );
descinit (descA, &m, &k, &mb, &nb, &rsrc, &csrc, &ictxt, &mA, &info);
descinit_(descB, &k, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &kB, &info);
descinit_(descC, &m, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &mC, &info);
double *A=(double*) malloc(mA*kA*sizeof(double)); //matricea locala
double *B = (double*) malloc(kB*nB*sizeof(double)); //matricea locala
double *C = (double*) malloc(mC*nC*sizeof(double)); //matricea locala
// se completeaza cu valori matricele locale folosind algoritmul 2D-ciclic
for(iloc=0;iloc<mA;iloc++){
    for(jloc=0;jloc<kA;jloc++){
        int fortidl = iloc + 1;
        int fortjdl = jloc + 1;
        i = indx12g_(&fortidl, &mb, &myrow, &ZERO, &nprow)-1;
        j = indx12g_(&fortjdl, &nb, &mycol, &ZERO, &npcol)-1;
        A[jloc*mA+iloc]=AA(i,j);
    }
}
for(iloc=0;iloc<kB;iloc++){
    for(jloc=0;jloc<nB;jloc++){
        int fortidl = iloc + 1;
        int fortjdl = jloc + 1;
        i = indx12g_(&fortidl, &mb, &myrow, &ZERO, &nprow)-1;
        j = indx12g_(&fortjdl, &nb, &mycol, &ZERO, &npcol)-1;
        B[jloc*kB+iloc]=BB(i,j);
    }
}
alpha = 1.0; beta = 0.0;
pdgemm_ ("No Transpose", "No Transpose", &m, &n, &k, &alpha, A, &ONE, &ONE, descA,
        B, &ONE, &ONE, descB, &beta, C, &ONE, &ONE, descC);
Cblacs_barrier(ictxt, "All");
// Print out the matrix product C
for (int id = 0; id < nprocs; ++id)
{
if (id == iam) {

```

```

cout << "C_loc on node " << iam << endl;
for (i = 0; i < mC; ++i)
{
    for (j = 0; j < nC; ++j)
        cout << setw(3) << *(C+mC*j+i) << " ";
    cout << endl;
}
Cblacs_barrier(ictxt, "All");
}
}
// Se construiesc matricea globala CC
for(iloc=0;iloc<mC;iloc++){
    for(jloc=0;jloc<nC;jloc++){
        int fortidl = iloc + 1;
        int fortjdl = jloc + 1;
        i = indx12g_(&fortidl, &mb, &myrow, &ZERO, &nprow)-1;
        j = indx12g_(&fortjdl, &nb, &mycol, &ZERO, &npcol)-1;
        CC(i,j)=C[jloc*mC+iloc];
    }
}
Cblacs_barrier(ictxt, "All");
/** Tipar rezultate finale
for (int id = 0; id < nprocs; ++id)
{
    if (id == iam) {
        cout << "Global matrix CC on node " << iam << endl;
        for (i = 0; i < m; ++i) {
            for (j = 0; j < n; ++j)
                cout << setw(3) << *(CC + m*i + j) << " ";
            cout << "\n";
        }
        cout << endl;
    }
}
Cblacs_barrier(ictxt, "All");
}
/* Free memory */
free(A);
free(B);
free(C);
/* Release process grid */
Cblacs_gridexit(ictxt);
} /* main */

```

The results are the following

```

[MI_gr_TPS1@hpc]$./mpiCC_ScL -o Example2.9.1.exe Example2.9.1.cpp
[MI_gr_TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 4 -host compute-0-0,compute-0-1
Example2.9.1.exe
[MI_gr_TPS1@hpc]$./mpiCC_ScL -o pmm.exe pmm.cpp
[MI_gr_TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 4 -host compute-0-0,compute-0-1 pmm.exe

```

Global matrix AA:

0	1	2	3	4
1	2	3	4	5
2	3	4	5	6
3	4	5	6	7
4	5	6	7	8

Global matrix BB:

0	1	2	3	4
1	2	3	4	5
2	3	4	5	6
3	4	5	6	7
4	5	6	7	8

```

C_loc on node 0
30 40 70
40 55 100
70 100 190
C_loc on node 1
50 60
70 85
130 160
C_loc on node 2
50 70 130
60 85 160
C_loc on node 3
90 110
110 135
Global matrix CC on node 0
30 40 0 0 70
40 55 0 0 100
0 0 0 0 0
0 0 0 0 0
70 100 0 0 190
Global matrix CC on node 1
0 0 50 60 0
0 0 70 85 0
0 0 0 0 0
0 0 0 0 0
0 0 130 160 0
Global matrix CC on node 2
0 0 0 0 0
0 0 0 0 0
50 70 0 0 130
60 85 0 0 160
0 0 0 0 0
Global matrix CC on node 3
0 0 0 0 0
0 0 0 0 0
0 0 90 110 0
0 0 110 135 0
0 0 0 0 0

```

To verify the results we use the sequential algorithm based on function dgemm. The program is presented below:

```

#include <string>
#include <iostream>
#include <stdio.h>
#include <math.h>
#include <iomanip>
#include <fstream>
#include <sstream>
using namespace std;

extern "C" {
    void dgemm_(const char &transa, const char &transb,
                const int &l, const int &n, const int &m, const double &alpha,
                const double *a, const int &lda, const double *b, const int &ldb,
                const double &beta, double *c, const int &ldc);
}

int main() {
    int i,j,m=5,k=5,n=5;
    //Matrizen werden als eindimensionale Arrays gespeichert
    double *AA = new double[m*k];

```



```

double *BB = new double[k*n];
double *CC = new double[m*n];
//Matrizen fuellen
for(i=0;i<m;i++)
    for(j=0;j<k;j++)
        AA[i*k+j]=i+j; //(10*i+j);
for(i=0;i<k;i++)
    for(j=0;j<n;j++)
        BB[i*n+j]=i+j;//AA[i*5+j]+5;

for( i=0 ; i<m*n ; i++)
    CC[i] = 0.0;
cout << "Global matrix AA:\n";
    for (i = 0; i < m; ++i) {
        for (j = 0; j < k; ++j) {
            cout << setw(3) << *(AA + k*i + j) << " ";
        }
        cout << "\n";
    }
    cout << endl;
cout << "Global matrix BB:\n";
    for (i = 0; i < k; ++i) {
        for (j = 0; j < n; ++j) {
            cout << setw(3) << *(BB + n*i + j) << " ";
        }
        cout << "\n";
    }
    cout << endl;
double alpha = 1.0;
double beta = 0.0;
//dgemm berechnet C=alpha*A*B+beta*C
dgemm_('N','N',m,n,k,alpha,AA,m,BB,k,beta,CC,m);

cout << "Global matrix CC:\n " << endl;
for(i=0 ; i<m ; i++) {
    for(j=0 ; j<n ; j++)
cout << setw(3) << *(CC + n*i + j) << " ";
    cout << endl;
}
cout << endl;

delete[] AA;
delete[] BB;
delete[] CC;

return 0;
}

```

The results are the following:

```

[MI_gr_TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 1 -host compute-0-0,compute-0-4
dgemm1.exe
Global matrix AA:
 0  1  2  3  4
 1  2  3  4  5
 2  3  4  5  6
 3  4  5  6  7
 4  5  6  7  8

Global matrix BB:

```

0	1	2	3	4
1	2	3	4	5
2	3	4	5	6
3	4	5	6	7
4	5	6	7	8

Global matrix CC:

30	40	50	60	70
40	55	70	85	100
50	70	90	110	130
60	85	110	135	160
70	100	130	160	190

In the program Example2.9.1.cpp the parallelization on data level is "non-rational" in terms of usage of operative memory, because for all processes the operative memory is reserved for global matrices! Using the function `pdgeadd` we can develop a program, in which the operative memory is used already rationally. The program is presented below.

### **Example 2.9.2.**

```
#include <stdio.h>
#include <string.h>
#include <stdlib.h>
#include "mpi.h"
#include <iostream>
#include <iomanip>
#include <fstream>
#include <sstream>
using namespace std;
#define AA(i,j) AA[(i)*n+(j)]
#define BB(i,j) BB[(i)*k+(j)]
#define CC(i,j) CC[(i)*n+(j)]
#define A(i,j) A[(i)*n+(j)]
#define B(i,j) B[(i)*n+(j)]
#define C(i,j) C[(i)*n+(j)]
static int MAX( int a, int b ){
    if (a>b) return(a); else return(b);
}

extern "C"
{

void Cblacs_pinfo( int* mypn, int* nprocs);
void Cblacs_get( int context, int request, int* value);
int Cblacs_gridinit( int* context, char * order, int np_row, int np_col);
void Cblacs_gridinfo( int context, int* np_row, int* np_col, int* my_row, int*
my_col);
void Cblacs_gridexit( int context);
void Cblacs_barrier(int, const char*);
void Cblacs_exit( int error_code);
void Cblacs_pcoord(int, int, int*, int*);
int numroc_( int *n, int *nb, int *iproc, int *isrcproc, int *nprocs);
int indxl2g_(int*, int*, int*, int*, int*);
void descinit_(int *desc, int *m, int *n, int *mb, int *nb, int *irsrc, int *icsrc,
int *ictxt, int *lld, int *info);
void pdgemm_( char *TRANSA, char *TRANSB, int *M, int *N, int *K, double *ALPHA, double
*A, int *IA, int *JA, int *DESCA,
```

```

        double * B, int * IB, int * JB, int * DESCB, double * BETA, double * C,
int * IC, int * JC, int * DESCC );
void pdgeadd_(char *TRANS, int *M, int *N, double * ALPHA, double *A, int *IA, int *JA, int
*DESCA, double *BETA, double *C,
        int *IC, int *JC, int *DESCC);

} // extern "C"

int main(int argc, char **argv)
{
int i, j;
int iam, nprocs, nprow, npcol, myrow, mycol;
// descriptori pentru matrici locale
int descA[9], descB[9], descC[9];
// descriptori pentru matrici globale
int descAA[9], descBB[9], descCC[9];
int m, n, k, mb, nb, rsrc, csrc, ictxt, llda, lladb, llcdc, info;
int lm, ln, lk;
int iloc, jloc;
int ZERO=0, ONE=1;
double zero=0.0E+0, one=1.0E+0;
int i_one = 1, i_zero = 0;
int lld_AA, lld_BB, lld_CC;
double alpha, beta;
double *AA, *BB, *CC, *A, *B, *C, *work, *tau;
m=5; n=5; k=5;
mb=2; nb=2;
nrow=2; npcol=2; // Astfel, programul se executa pe 4 procese
Cblacs_pinfo(&iam, &nprocs);
Cblacs_get(-1, 0, &ictxt);
Cblacs_gridinit( &ictxt, "Row", nrow, npcol);
Cblacs_gridinfo(ictxt, &nrow, &npcol, &myrow, &mycol);
rsrc=0; csrc=0;
if ( iam==0 )
{
AA = (double*) malloc(m*k*sizeof(double)); // matricea "globala" (de dim. m*k) pentru
inmultirea AA*BB=CC
BB = (double*) malloc(k*n*sizeof(double)); //matricea "globala" (de dim. k*n) pentru
inmultirea AA*BB=CC
CC = (double*) malloc(m*n*sizeof(double)); //matricea "globala" (de dim. m*n) pentru
inmultirea AA*BB=CC
// initializarea matricelor globale
for(i=0; i<m; i++)
    for(j=0; j<k; j++)
        AA[i*k+j]=(10*i+j); // (i+j);
for(i=0; i<k; i++)
    for(j=0; j<n; j++)
        BB[i*n+j]=AA[i*5+j]+5; // (i+j);
}
else{
AA = NULL;
BB = NULL;
//other processes don't contain parts of A
}
if (iam==0)
{
printf("===== REZULT OF THE PROGRAM %s \n", argv[0]);
cout << "Global matrix AA:\n";
    for (i = 0; i < m; ++i) {
        for (j = 0; j < k; ++j) {
            cout << setw(3) << *(AA + k*i + j) << " ";

```

```

        }
        cout << "\n";
    }
    cout << endl;
    cout << "Global matrix BB:\n";
    for (i = 0; i < k; ++i) {
        for (j = 0; j < n; ++j) {
            cout << setw(3) << *(BB + n*i + j) << " ";
        }
        cout << "\n";
    }
    cout << endl;
}

Cblacs_barrier(ictxt, "All");
int mA = numroc_( &m, &mb, &myrow, &rsrc, &nprow );
int kA = numroc_( &k, &nb, &mycol, &rsrc, &npcol );
int kB = numroc_( &k, &mb, &myrow, &rsrc, &nprow );
int nB = numroc_( &n, &nb, &mycol, &rsrc, &npcol );
int mC = numroc_( &m, &mb, &myrow, &rsrc, &nprow );
int nC = numroc_( &n, &nb, &mycol, &rsrc, &nprow );
lld_AA = MAX( numroc_( &m, &k, &myrow, &rsrc, &nprow ), 1 );
lld_BB = MAX( numroc_( &k, &n, &myrow, &rsrc, &nprow ), 1 );
lld_CC = MAX( numroc_( &m, &n, &myrow, &rsrc, &nprow ), 1 );
// Initialize descriptors (local matrix A is considered as distributed with blocking
parameters
// m, n, i.e. there is only one block - whole matrix A - which is located on process
(0,0) )
descinit_(descAA, &m, &k, &m, &k, &rsrc, &csrc, &ictxt, &lld_AA, &info );
descinit_(descBB, &k, &n, &k, &n, &rsrc, &csrc, &ictxt, &lld_BB, &info );
descinit_(descCC, &m, &n, &m, &n, &rsrc, &csrc, &ictxt, &lld_CC, &info );
descinit_(descA, &m, &k, &mb, &nb, &rsrc, &csrc, &ictxt, &mA, &info);
descinit_(descB, &k, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &kB, &info);
descinit_(descC, &m, &n, &mb, &nb, &rsrc, &csrc, &ictxt, &mC, &info);
A = (double*) malloc(mA*kA*sizeof(double)); //matricea locala
B = (double*) malloc(kB*nB*sizeof(double)); //matricea locala
C = (double*) malloc(mC*nC*sizeof(double)); //matricea locala
// Call pdgeadd_ to distribute matrix (i.e. copy AA into A and BB into B )
pdgeadd_( "N", &m, &k, &one, AA, &i_one, &i_one, descAA, &zero, A, &i_one, &i_one,
descA);
pdgeadd_( "N", &k, &n, &one, BB, &i_one, &i_one, descBB, &zero, B, &i_one, &i_one,
descB);
// Tipar matricele locale A si B
for (int id = 0; id < nprocs; ++id)
{
    Cblacs_barrier(ictxt, "All");
    if (id == iam) {
        printf("Local A(%d*%d) on node %d (%d,%d) \n", kA, mA, iam, myrow, mycol);
        //cout << "A on node " << iam << endl;
        //for (i = 0; i < mA; ++i)
        for (j = 0; j < kA; ++j)
        {
            //for (j = 0; j < kA; ++j)
            for (i = 0; i < mA; ++i)
                cout << setw(3) << *(A+mA*j+i) << " ";
            cout << endl;
        }
        printf("Local B(%d*%d) on node %d (%d,%d) \n", nB, kB, iam, myrow, mycol);
        //cout << "B on node " << iam << endl;
        //for (i = 0; i < kB; ++i)
        for (j = 0; j < nB; ++j)

```

```

{
//for (j = 0; j < nB; ++j)
for (i = 0; i < kB; ++i)
cout << setw(3) << *(B+kB*j+i) << " ";
cout << endl;
}
cout << endl;
}
Cblacs_barrier(ictxt, "All");
}
alpha = 1.0; beta = 0.0;
pdgemm ("No Transpose", "No
Transpose", &m, &n, &k, &alpha, A, &ONE, &ONE, descA, B, &ONE, &ONE, descB, &beta, C, &ONE, &ONE, desc
C);
Cblacs_barrier(ictxt, "All");
// Print out the matrix product C
for (int id = 0; id < nprocs; ++id)
{
Cblacs_barrier(ictxt, "All");
if (id == iam) {
printf("Local C(%d*%d) on node %d (%d,%d) \n", mC, nC, iam, myrow, mycol);
//cout << "C_loc on node " << iam << endl;
for (i = 0; i < mC; ++i)
{
for (j = 0; j < nC; ++j)
cout << setw(3) << *(C+mC*j+i) << " ";
cout << endl;
}
Cblacs_barrier(ictxt, "All");
}
}
// Se construiesc matricea CC

pdgeadd_ ("N", &m, &n, &one, C, &i_one, &i_one, descC, &zero, CC, &i_one, &i_one,
descCC);

/** Tipar rezultate finale

if (iam==0)
{
cout << "Global matrix CC =AA*BB:\n";
//for (i = 0; i < m; ++i)
for (j = 0; j < n; ++j) {
//for (j = 0; j < n; ++j)
for (i = 0; i < m; ++i) {
cout << setw(5) << *(CC + n*i + j) << " ";
}
cout << "\n";
}
cout << endl;
}

/* Free memory */
free(A);
free(B);
free(C);
if( myrow==0 && mycol==0 ){
free( AA );
free( BB);
free( CC);
}

```

```

    /* Release process grid */
    Cblacs_gridexit(ictxt);
} /* main */

```

The results are the following:

```

[MI_gr_TPS1@hpc]$ ./mpiCC_ScL -o Example2.9.2.exe Example2.9.2.cpp
[MI_gr_TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 4 -host compute-0-0,compute-0-4
Example2.9.2.exe
Global matrix AA:
  0   1   2   3   4
10 11 12 13 14
20 21 22 23 24
30 31 32 33 34
40 41 42 43 44

Global matrix BB:
  5   6   7   8   9
15 16 17 18 19
25 26 27 28 29
35 36 37 38 39
45 46 47 48 49

Local A(3*3) on node 0 (0,0)
  0   1   4
10 11 14
40 41 44
Local B(3*3) on node 0 (0,0)
  5   6   9
15 16 19
45 46 49

Local A(2*3) on node 1 (0,1)
20 21 24
30 31 34
Local B(2*3) on node 1 (0,1)
25 26 29
35 36 39

Local A(3*2) on node 2 (1,0)
  2   3
12 13
42 43
Local B(3*2) on node 2 (1,0)
  7   8
17 18
47 48

Local A(2*2) on node 3 (1,1)
22 23
32 33
Local B(2*2) on node 3 (1,1)
27 28
37 38

Local C(3*3) on node 0 (0,0)
800 1800 4800

```

```

835 1885 5035
940 2140 5740
Local C(3*2) on node 1 (0,1)
2800 3800
2935 3985
3340 4540
Local C(2*3) on node 2 (1,0)
870 1970 5270
905 2055 5505
Local C(2*2) on node 3 (1,1)
3070 4170
3205 4355
Global matrix CC =AA*BB:
  800  1800  2800  3800  4800
  835  1885  2935  3985  5035
  870  1970  3070  4170  5270
  905  2055  3205  4355  5505
  940  2140  3340  4540  5740

```

The accuracy of the results is confirmed by the sequential program:

```

[MI_gr_TPS1@hpc]$ ./mpiCC_ScL -o dgemml.exe dgemml.cpp
[MI_gr_TPS1@hpc]$ /opt/openmpi/bin/mpirun -n 1 -host compute-0-0,compute-0-4
dgemml.exe
Global matrix AA:
  0   1   2   3   4
 10  11  12  13  14
 20  21  22  23  24
 30  31  32  33  34
 40  41  42  43  44

Global matrix BB:
  5   6   7   8   9
 15  16  17  18  19
 25  26  27  28  29
 35  36  37  38  39
 45  46  47  48  49

Global matrix CC:

800 1800 2800 3800 4800
835 1885 2935 3985 5035
870 1970 3070 4170 5270
905 2055 3205 4355 5505
940 2140 3340 4540 5740

```

```

PvSYMM(SIDE,UPLO,M,N,ALPHA,A,IA,JA,DESCA,B,IB,JB,DESCB,BETA,
        C,IC,JC,DESCC)
void pdsymm_(F_CHAR_T SIDE,F_CHAR_T UPLO,int *M,int *N,
             double *ALPHA,double *A,int *IA,int *JA,int *DESCA,
             double *B,int *IB,int *JB,int *DESCB,double *BETA,
             double *C,int *IC,int *JC,int *DESCC )

```

**Purpose:** PvSYMM performs one of the distributed matrix-matrix operations  
 $sub( C ) := alpha \times sub( A ) \times sub( B ) + beta \times sub( C ),$  or  
 $sub( C ) := alpha \times sub( B ) \times sub( A ) + beta \times sub( C ),$   
 where  $sub( C )$  denotes  $C(IC:IC+M-1,JC:JC+N-1),$

$$\text{sub} ( A ) \text{ denotes } \begin{cases} A( IA : IA + M - 1, JA : JA + M - 1 ) & \text{if } SIDE = 'L', \\ A( IA : IA + N - 1, JA : JA + N - 1 ) & \text{if } SIDE = 'R' \end{cases}$$

$\text{sub}( B )$  denotes  $B( IB : IB + M - 1, JB : JB + N - 1 )$ .

Alpha and beta are scalars,  $\text{sub}( A )$  is a symmetric distributed matrix and  $\text{sub}( B )$  and  $\text{sub}( C )$  are M-by-N distributed matrices.

### Arguments

SIDE (global input) CHARACTER

*On entry, SIDE specifies whether the symmetric distributed matrix  $\text{sub}( A )$  appears on the left or right in the operation as follows:*

*SIDE = 'L'  $\text{sub}( C ) := \alpha * \text{sub}( A ) * \text{sub}( B ) + \beta * \text{sub}( C )$ ,*

*SIDE = 'R'  $\text{sub}( C ) := \alpha * \text{sub}( B ) * \text{sub}( A ) + \beta * \text{sub}( C )$ ,*

UPLO (global input) CHARACTER

*On entry, UPLO specifies whether the upper or lower triangular part of the symmetric distributed matrix  $\text{sub}( A )$  is to be referenced.*

M (global input) INTEGER

*The number of rows to be operated on i.e., the number of rows of the distributed submatrix  $\text{sub}( C )$ .  $M \geq 0$ .*

N (global input) INTEGER

*The number of columns to be operated on i.e the number of columns of the distributed submatrix  $\text{sub}( C )$ .  $N \geq 0$ .*

ALPHA (global input) REAL/COMPLEX

*On entry, ALPHA specifies the scalar alpha.*

A (local input) array of dimension (LLD\_A, LOC<sub>q</sub>(JA+NA-1))

*Before entry this array contains the local pieces of the symmetric distributed matrix  $\text{sub}( A )$ , such that when UPLO = 'U', the NA-by-NA upper triangular part of the distributed matrix  $\text{sub}( A )$  must contain the upper triangular part of the symmetric distributed matrix and the strictly lower triangular part of  $\text{sub}( A )$  is not referenced, and when UPLO = 'L', the NA-by-NA lower triangular part of the distributed matrix  $\text{sub}( A )$  must contain the lower triangular part of the symmetric distributed matrix and the strictly lower triangular part of  $\text{sub}( A )$  is not referenced.*

IA (global input) INTEGER

*The global row index of the submatrix of the distributed matrix A to operate on.*

JA (global input) INTEGER

*The global column index of the submatrix of the distributed matrix A to operate on.*

DESCA (global and local input) INTEGER array of dimension 8

*The array descriptor of the distributed matrix A.*

B (local input) array of dimension (LLD\_B, LOC<sub>q</sub>(JB+N-1))

*Before entry, this array contains the local pieces of the distributed matrix  $\text{sub}( B )$ .*

IB (global input) INTEGER

*The global row index of the submatrix of the distributed matrix B to operate on.*

JB (global input) INTEGER

*The global column index of the submatrix of the distributed matrix B to operate on.*

DESCB (global and local input) INTEGER array of dimension 8

*The array descriptor of the distributed matrix B.*

BETA (global input) REAL/COMPLEX

*On entry, BETA specifies the scalar beta. When BETA is supplied as zero then  $\text{sub}( C )$  need not be set on input.*

C (local input/local output) array of dimension (LLD\_C, LOC<sub>q</sub>(JC+N-1))

*Before entry, this array must contain the local pieces of the distributed matrix*



*sub( C ). On exit, the distributed matrix sub( C ) is overwritten by the M-by-N updated distributed matrix.*

IC (global input) INTEGER

*The global row index of the submatrix of the distributed matrix C to operate on.*

JC (global input) INTEGER

*The global column index of the submatrix of the distributed matrix C to operate on.*

DESCC (global and local input) INTEGER array of dimension 8

*The array descriptor of the distributed matrix C.*

```
PvTRAN(M,N,ALPHA,A,IA,JA,DESCA,BETA,C,IC,JC,DESCC)
```

```
void pdtran_(int *M, int *N, double *ALPHA, double *A, int *IA,  
             int *JA, int *DESCA, double *BETA, double *C, int  
             *IC, int *JC, int *DESCC)
```

Purpose: PvTRAN transposes a distributed matrix  $\text{sub}(C) = \text{beta} \times \text{sub}(C) + \text{alpha} \times \text{op}(\text{sub}(A))$  where  $\text{sub}(C)$  denotes  $C(\text{IC}:\text{IC}+\text{M}-1, \text{JC}:\text{JC}+\text{N}-1)$ ,  $\text{sub}(A)$  denotes  $A(\text{IA}:\text{IA}+\text{N}-1, \text{JA}:\text{JA}+\text{M}-1)$ ,  $\text{op}(A)$  denotes  $A^T$ . Beta is a scalar,  $\text{sub}(C)$  is an M-by-N distributed matrix,  $\text{sub}(A)$  is an N-by-M distributed matrix.

### Arguments

M (global input) INTEGER

*The number of rows to be operated on i.e., the number of rows of the distributed submatrix sub( C ).  $M \geq 0$ .*

N (global input) INTEGER

*The number of columns to be operated on i.e the number of columns of the distributed submatrix sub( C ).  $N \geq 0$ .*

ALPHA (global input) REAL

*On entry, ALPHA specifies the scalar alpha.*

A (local input) REAL array of dimension (LLD\_A, LOCq(JA+M-1))

*This array contains the local pieces of the distributed matrix sub( A ).*

IA (global input) INTEGER

*The global row index of the submatrix of the distributed matrix A to operate on.*

JA (global input) INTEGER

*The global column index of the submatrix of the distributed matrix A to operate on.*

DESCA (global and local input) INTEGER array of dimension 8

*The array descriptor of the distributed matrix A.*

BETA (global input) REAL

*On entry, BETA specifies the scalar beta. When BETA is supplied as zero then sub( C ) need not be set on input.*

C (local input/local output) array of dimension (LLD\_C, LOCq(JC+N-1))

*This array contains the local pieces of the distributed matrix sub( C ). On exit, the distributed matrix sub( C ) is over- written by the updated matrix.*

IC (global input) INTEGER

*The global row index of the submatrix of the distributed matrix C to operate on.*

JC (global input) INTEGER

*The global column index of the submatrix of the distributed matrix C to operate on.*

DESCC (global and local input) INTEGER array of dimension 8

*The array descriptor of the distributed matrix C.*

```
PvTRSM(SIDE,UPLO,TRANSA,DIAG,M,N,ALPHA,A,IA,JA,DESCA,  
        B,IB,JB,DESCB )
```

```
void pdtrsm_(F_CHAR_T SIDE, F_CHAR_T UPLO, F_CHAR_T TRANS,
```

```

F_CHAR_T DIAG, int *M, int *N, double *ALPHA, double
*A, int *IA, int *JA, int *DESCA, double *B, int
*IB, int *JB, int *DESCB)

```

**Purpose:** PvTRSM solves one of the distributed matrix equations

$op(sub(A)) \times X = \alpha \times sub(B)$ , or  $X \times op(sub(A)) = \alpha \times sub(B)$ , where

$$sub(A) \text{ denotes } \begin{cases} A(IA : IA + M - 1, JA : JA + M - 1) & \text{if SIDE} = 'L', \\ A(IA : IA + N - 1, JA : JA + N - 1) & \text{if SIDE} = 'R' \end{cases}$$

$sub(B)$  denotes  $B(IB:IB+M-1,JB:JB+N-1)$ ,  $\alpha$  is a scalar,  $X$  and  $sub(B)$  are an  $M$ -by- $N$  distributed matrix,  $sub(A)$  is a unit, or non-unit, upper or lower triangular distributed matrix and  $op(A)$  is one of  $op(A) = A$  or  $op(A) = A^T$ . The distributed matrix  $X$  is overwritten on  $sub(B)$ .

Arguments

SIDE (global input) CHARACTER

On entry, SIDE specifies whether  $op(A)$  appears on the left or right of  $X$  as follows:

SIDE = 'L',  $op(sub(A)) * X = \alpha * sub(B)$ ,

SIDE = 'R',  $X * op(sub(A)) = \alpha * sub(B)$ .

UPLO (global input) CHARACTER

On entry, UPLO specifies whether the distributed matrix  $sub(A)$  is an upper or lower triangular distributed matrix.

TRANSA (global input) CHARACTER

The form of  $op(A)$  to be used in the matrix multiplication as follows:

TRANSA = 'N',  $op(A) = A$ ,

TRANSA = 'T',  $op(A) = A^T$ ,

TRANSA = 'C',  $op(A) = A^T$ .

DIAG (global input) CHARACTER

On entry, DIAG specifies whether or not  $sub(A)$  is unit triangular as follows:

DIAG = 'U',  $sub(A)$  is assumed to be unit triangular,

DIAG = 'N',  $sub(A)$  is not assumed to be unit triangular.

M (global input) INTEGER

The number of rows to be operated on i.e., the number of rows of the distributed submatrix  $sub(B)$ .  $M \geq 0$ .

N (global input) INTEGER

The number of columns to be operated on i.e., the number of columns of the distributed submatrix  $sub(B)$ .  $N \geq 0$ .

ALPHA (global input) REAL/COMPLEX

On entry, ALPHA specifies the scalar  $\alpha$ .

A (local input) array of dimension (LLD\_A, LOCq(JA+NA-1))

Before entry with UPLO = 'U', the leading NA-by-NA upper triangular part of the distributed matrix  $sub(A)$  must contain the local pieces of the upper triangular distributed matrix and its strictly lower triangular part is not referenced. Before entry with UPLO = 'L', the leading NA-by-NA lower triangular part of the distributed matrix  $sub(A)$  must contain the lower triangular distributed matrix and its strictly upper triangular part is not referenced. Note that when DIAG = 'U', the diagonal elements of  $sub(A)$  are not referenced either, but are assumed to be unity.

IA (global input) INTEGER

The global row index of the submatrix of the distributed matrix A to operate on.

JA (global input) INTEGER

The global column index of the submatrix of the distributed matrix A to operate on.

DESCA (global and local input) INTEGER array of dimension 8

*The array descriptor of the distributed matrix A.*

B (local input) array of dimension (LLD<sub>B</sub>, LOC<sub>q</sub>(JB+N-1))  
*Before entry, this array contains the local pieces of the distributed matrix sub( B ). On exit, sub( B ) is overwritten by the solution distributed matrix.*

IB (global input) INTEGER  
*The global row index of the submatrix of the distributed matrix B to operate on.*

JB (global input) INTEGER  
*The global column index of the submatrix of the distributed matrix B to operate on.*

DESCB (global and local input) INTEGER array of dimension 8  
*The array descriptor of the distributed matrix B.*

```
PvSYMM(SIDE,UPLO,M,N,ALPHA,A,IA,JA,DESCA,B,IB,JB,DESCB,BETA,C,IC,
,JC,DESCC)
void pdsymm_(F_CHAR_T SIDE, F_CHAR_T UPLO, int *M, int *N,
double *ALPHA, double *A, int *IA, int *JA, int *DESCA,
double *B, int *IB, int *JB, int *DESCB, double *BETA,
double *C, int *IC, int *JC, int *DESCC)
```

**Purpose:** Performs a scalar-matrix-matrix product (one matrix operand is symmetric) and adds the result to a scalar-matrix product for distribute matrices. The operation is defined as

$$\text{sub}(C) := \text{ALPHA} * \text{sub}(A) * \text{sub}(B) + \text{BETA} * \text{sub}(C) \text{ or}$$

$$\text{sub}(C) := \text{ALPHA} * \text{sub}(B) * \text{sub}(A) + \text{BETA} * \text{sub}(C),$$

where ALPHA and BETA are scalars, sub(A) is a symmetric distributed matrix, sub(A)=A(IA:IA+M-1,JA:JA+M-1), if SIDE='L', and sub(A)=A(IA:IA+N-1,JA:JA+N-1), if SIDE='R', sub(B) and sub(C) are M-by-N distributed matrices, sub(B)=B(IB:IB+M-1,JB:JB+N-1), sub(C)=C(IC:IC+M-1, JC:JC+N-1).

### Input Parameters

SIDE (GLOBAL) CHARACTER\*1.  
*Specifies whether the symmetric distributed matrix SUB(A) appears on the left or right in the operation:*  
*if side = 'L' or 'l', then SUB(C) := ALPHA \* SUB(A) \* SUB(B) + BETA \* SUB(C) ;*  
*if SIDE = 'R' or 'r', then SUB(C) := ALPHA \* SUB(B) \* SUB(A) + BETA \* SUB(C) .*

UPLO (global) CHARACTER\*1.  
*Specifies whether the upper or lower triangular part of the symmetric distributed matrix SUB(A) is used:*  
*if UPLO = 'U' or 'u', then the upper triangular part is used;*  
*if uplo = 'L' or 'l', then the lower triangular part is used.*

M (global) INTEGER.  
*Specifies the number of rows of the distribute submatrix SUB(C) ,  $M \geq 0$ .*

N(global) INTEGER.  
*Specifies the number of columns of the distribute submatrix SUB(C) ,  $M \geq 0$ .*

ALPHA (global) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.  
*Specifies the scalar ALPHA.*

A (local) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM. Array, DIMENSION (LLD<sub>A</sub>, LOC<sub>q</sub>(JA+NA-1)) .

*Before entry this array must contain the local pieces of the symmetric distributed matrix SUB(A) , such that when UPLO = 'U' or 'u', the NA-by-NA upper triangular part of the distributed matrix SUB(A) must contain the upper triangular part of the symmetric distributed matrix and the strictly lower triangular part of*

*SUB(A) is not referenced, and when UPLO = 'L' or 'L', the NA-by-NA lower triangular part of the distributed matrix SUB(A) must contain the lower triangular part of the symmetric distributed matrix and the strictly upper triangular part of SUB(A) is not referenced.*

IA, JA (global) INTEGER.

*The row and column indices in the distributed matrix A indicating the first row and the first column of the submatrix SUB(A), respectively.*

DESCA (global and local) INTEGER array of dimension 8.

*The array descriptor of the distributed matrix A.*

B (local) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

*Array, DIMENSION (LLD\_B, LOC<sub>q</sub>(JB+N-1)). Before entry this array must contain the local pieces of the distributed matrix SUB(B).*

IB, JB (global) INTEGER.

*The row and column indices in the distributed matrix B indicating the first row and the first column of the submatrix SUB(B), respectively.*

DESCB (global and local) INTEGER array of dimension 8.

*The array descriptor of the distributed matrix B.*

BETA (global) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

*Specifies the scalar BETA. When BETA is set to zero, then SUB(C) need not be set on input.*

C (local) REAL for PSSYMM, DOUBLE PRECISION for PDSYMM, COMPLEX for PCSYMM, DOUBLE COMPLEX for PZSYMM.

*Array, DIMENSION (LLD\_C, LOC<sub>q</sub>(JC+N-1)). Before entry this array must contain the local pieces of the distributed matrix SUB(C).*

IC, JC (global) INTEGER.

*The row and column indices in the distributed matrix C indicating the first row and the first column of the submatrix SUB(C), respectively.*

DESCC (global and local) INTEGER array of dimension 8.

*The array descriptor of the distributed matrix C*

## Output Parameters

C

Overwritten by the m-by-n updated matrix.

```
PvLACPY(UPLO,M,N,A,IA,JA,DESCA,B,IB,JB,DESCB)
void pdlapy_(char *uplo,int *m,int *n,double *a,int *ia,int
             *ja,int *desca,double *b,int *ib,int *jb,int
             *descb);
```

**Purpose:** The routine copies all or part of a distributed matrix A to another distributed matrix B. No communication is performed, PvLACPY performs a local copy sub(A)=sub(B), where sub(A) denotes A(IA:IA+M-1,JA:JA+N-1) and sub(B) denotes B(IB:IB+M-1,JB:JB+N-1).

## Input Parameters

UPLO (global) CHARACTER.

*Specifies the part of the distributed matrix sub(A) to be copied:*

= 'U': Upper triangular part; the strictly lower triangular part of sub(A) is not referenced;

= 'L': Lower triangular part; the strictly upper triangular part of sub(A) is not referenced.

*Otherwise: all of the matrix sub(A) is copied.*

M (global) INTEGER.

*The number of rows to be operated on, that is, the number of rows of the distributed submatrix  $sub(A)$  ( $M \geq 0$ ).*

N (global) INTEGER.

*The number of columns to be operated on, that is, the number of columns of the distributed submatrix  $sub(A)$  ( $N \geq 0$ ).*

A (local) REAL for PSLACPY, DOUBLE PRECISION for PDLACPY, COMPLEX for PCLACPY, COMPLEX\*16 for PZLACPY.

*Pointer into the local memory to an array of DIMENSION( $LLD\_A$ ,  $LOC_c(JA+N-1)$ ). On entry, this array contains the local pieces of the distributed matrix  $sub(A)$ .*

IA, JA (global) INTEGER.

*The row and column indices in the global array  $a$  indicating the first row and the first column of the submatrix  $sub(A)$ , respectively.*

DESCA (global and local) INTEGER array, DIMENSION ( $DLEN\_$ ).

*The array descriptor for the distributed matrix  $A$ .*

IB, JB (global) INTEGER.

*The row and column indices in the global array  $B$  indicating the first row and the first column of  $sub(B)$  respectively.*

DESCB (global and local) INTEGER array, DIMENSION ( $DLEN\_$ ).

*The array descriptor for the distributed matrix  $A$ .*

#### **Output Parameters**

B (local) REAL for PSLACPY, DOUBLE PRECISION for PDLACPY, COMPLEX for PCLACPY, COMPLEX\*16 for PZLACPY.

*Pointer into the local memory to an array of DIMENSION ( $LLD\_B$ ,  $LOCC(JB+N-1)$ ). This array contains on exit the local pieces of the distributed matrix  $sub(B)$  set as follows:*

*if UPLO='U',  $B(IB+I-1, JB+J-1) = A(IA+I-1, JA+J-1)$ ,  $1 \leq I \leq J$ ,  $1 \leq J \leq N$ ;*

*if UPLO='L',  $B(IB+I-1, JB+J-1) = A(IA+I-1, JA+J-1)$ ,  $J \leq I \leq M$ ,  $1 \leq J \leq n$ ;*

*otherwise,  $B(IB+I-1, JB+J-1) = A(IA+I-1, JA+J-1)$ ,  $1 \leq I \leq M$ ,  $1 \leq J \leq N$ .*

The PBLAS Level 3 routines perform distributed matrix-matrix operations. Table "PBLAS Level 3 Routine Groups and Their Data Types" lists the PBLAS Level 3 routine groups and the data types associated with them.

### **PBLAS Level 3 Routine Groups and Their Data Types**

<b>Routine Group</b>	<b>Data Types</b>	<b>Description</b>
p?geadd	s, d, c, z	Distributed matrix-matrix sum of general matrices
p?tradd	s, d, c, z	Distributed matrix-matrix sum of triangular matrices
p?gemm	s, d, c, z	Distributed matrix-matrix product of general matrices
p?hemm	c, z	Distributed matrix-matrix product, one matrix is Hermitian
p?herk	c, z	Rank-k update of a distributed Hermitian matrix
p?her2k	c, z	Rank-2k update of a distributed Hermitian matrix
p?symm	s, d, c, z	Matrix-matrix product of distributed symmetric matrices
p?syrk	s, d, c, z	Rank-k update of a distributed symmetric matrix

<b>Routine Group</b>	<b>Data Types</b>	<b>Description</b>
p?syr2k	s, d, c, z	Rank-2k update of a distributed symmetric matrix
p?tran	s, d	Transposition of a real distributed matrix
p?tranc	c, z	Transposition of a complex distributed matrix (conjugated)
p?tranu	c, z	Transposition of a complex distributed matrix
p?trmm	s, d, c, z	Distributed matrix-matrix product, one matrix is triangular
p?trsm	s, d, c, z	Solution of a distributed matrix equation, one matrix is triangular