Topic 6: Mixed-Integer Nonlinear Optimization

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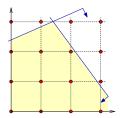


CENG60008 – Process Optimisation CENG70003 – Advanced Process Optimisation

Mixed-Integer Nonlinear Programming (MINLP)

f and g may not be linear or even convex functions

- That optimization problems with binary/ integer variables are more difficult to solve than continuous problems may appear counter-intuitive
- The feasible set is nonconnected and therefore inherently nonconvex



MINLP Solution Paradigm

Construct a sequence of related but simpler subproblems that converges (finitely) to the exact MINLP solution

- Use linearization / relaxation to construct the subproblems
 - ▶ Both the objective function and the feasible region
- The subproblems should yield valid bounds on the original MINLP
 - Lower bound for a minimize problem
 - Upper bound for a maximize problem
- The subproblems should be easier to solve than the original MINLP
 - Many subproblems may need to be solved!
- The number of subproblems solved should be much smaller than with the complete enumeration method
 - although pathological problems can always be constructed...

Contents

Solving MINLP models is challenging – We will learn some typical decomposition techniques for MINLPs with either convex or nonconvex participating functions

- Methods for MINLP with Convex Functions
 - Branch-and-Bound Algorithm
 - Outer-Approximation Algorithm
- Methods for MINLP with Nonconvex Functions



Recommended Readings:

- C.A. Floudas, Nonlinear and Mixed- Integer Optimization, Oxford University Press, 1995
- I.E. Grossmann, Review of nonlinear mixed-integer and disjunctive programming techniques, *Optimization & Engineering* 1:227-252, 2002

Main Solution Algorithms

$$\begin{aligned} & \underset{\mathbf{x}, \mathbf{y}}{\text{min}} \ f(\mathbf{x}, \mathbf{y}) \\ & \text{s.t.} \ \ \mathbf{g}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0} \\ & \quad \mathbf{x} \in [\mathbf{x}^{\mathrm{L}}, \mathbf{x}^{\mathrm{U}}] \subset \mathbb{R}^{n_{x}} \\ & \quad \mathbf{y} \in \{0, 1\}^{n_{y}} \end{aligned}$$

- f and g are nonlinear and convex functions
- Branch-and-Bound (BB)
 - Ravindran & Gupta (1985); Stubbs & Mehrotra (1999)
- Generalized Benders Decomposition (GBD)
 - Geoffrion (1972)
- Outer Approximation (OA)
 - ▶ Duran & Grossmann (1986); Fletcher & Leyffer (1994)
- Extended Cutting Plane (ECP)
 - Westerlund & Pettersson (1995)

Branch-and-Bound Algorithm

$$egin{aligned} \min_{\mathbf{x},\mathbf{y}} & f(\mathbf{x},\mathbf{y}) \ & ext{s.t.} & \mathbf{g}(\mathbf{x},\mathbf{y}) \leq \mathbf{0} \ & \mathbf{x} \in [\mathbf{x}^{\mathrm{L}},\mathbf{x}^{\mathrm{U}}] \subset \mathbb{R}^{n_{\mathrm{x}}} \ & \mathbf{y} \in \{0,1\}^{n_{\mathrm{y}}} \end{aligned}$$

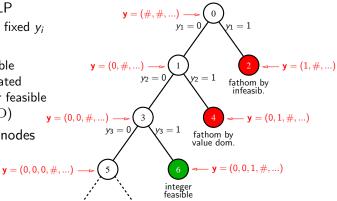
$$\begin{aligned} & \underset{\mathbf{x}, \mathbf{y}}{\text{min}} \ f(\mathbf{x}, \mathbf{y}) \\ & \text{s.t.} \ \ \mathbf{g}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0} \\ & \quad \mathbf{x} \in [\mathbf{x}^{\mathrm{L}}, \mathbf{x}^{\mathrm{U}}] \subset \mathbb{R}^{n_{x}} \\ & \quad \mathbf{y} \in [0, 1]^{n_{y}} \end{aligned}$$

Key Properties:

- If an NLP relaxation is infeasible, so is the MINLP it relaxes
 - ▶ But an infeasible MINLP could have a feasible NLP relaxation
- The optimal solution value of an NLP relaxation yields a lower bound on the MINLP it relaxes (minimize case)
- If the optimal solution to an NLP relaxation is integer feasible, it is also optimal for the MINLP it relaxes

Branch-and-Bound Algorithm

- Branch on fractional y_i
- Solve relaxed NLP
 - Mixed free / fixed y_i
- Fathom if:
 - Node infeasible
 - Node dominated
 - Node integer feasible (update UBD)
- Stop if no open nodes



Convexity is Critical!

Branch-and-Bound Example

Workshop. Consider the following MINLP model and derive the relaxed NLP model at the root node.

$$\begin{aligned} & \underset{\mathbf{x}, \mathbf{y}}{\min} \ x_1^2 + x_2^2 + y_1 + 1.5y_2 + 0.5y_3 \\ & \text{s.t.} \ (x_1 - 2)^2 - x_2 \leq 0 \\ & x_1 \geq 2y_1 \\ & x_1 - x_2 \leq 4(1 - y_2) \\ & x_1 \geq (1 - y_1) \\ & x_2 \geq y_2 \\ & x_1 + x_2 \geq 3y_3 \\ & y_1 + y_2 + y_3 \geq 1 \\ & 0 \leq x_1, x_2 \leq 4 \\ & y_1, y_2, y_3 \in \{0, 1\} \end{aligned}$$

NLP relax

$$\min_{\mathbf{x}, \mathbf{y}} x_1^2 + x_2^2 + y_1 + 1.5y_2 + 0.5y_3$$
s.t. $(x_1 - 2)^2 - x_2 \le 0$

$$x_1 \ge 2y_1$$

$$x_1 - x_2 \le 4(1 - y_2)$$

$$x_1 \ge (1 - y_1)$$

$$x_2 \ge y_2$$

$$x_1 + x_2 \ge 3y_3$$

$$y_1 + y_2 + y_3 \ge 1$$

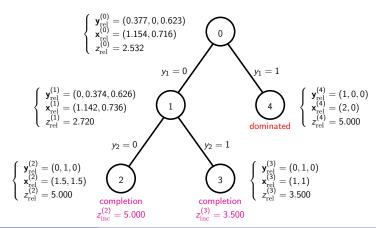
$$0 \le x_1, x_2 \le 4$$

$$0 \le y_1, y_2, y_3 \le 1$$

Optimal solution: $f^* = 3.5$, $x_1^* = x_2^* = 1$, $y_2^* = 1$, $y_1^* = y_2^* = 0$ Lower Bound: LBD⁰ \approx 2.532, $x_1^0 \approx 1.154, x_2^0 \approx 0.716,$ $y_1^0 \approx 0.377, y_2^0 = 0, y_3^0 \approx 0.623$

Branch-and-Bound Example

Workshop. Apply the branch-and-bound algorithm to the MINLP model, using GAMS to solve the relaxed NLP models – Branch on fractional variables and apply depth-first search.



Outer-Approximation Algorithm

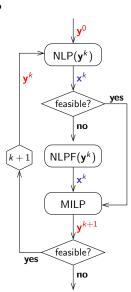
• Idea. Alternate solution between NLP and MILP

Primal problem $NLP(\mathbf{y}^k)$: set $\mathbf{y} = \mathbf{y}^k$ in f and \mathbf{g}

$$\mathbf{x}^k \in rg \min_{\mathbf{x}} \ f(\mathbf{x}, \mathbf{y}^k)$$
 s.t. $\mathbf{g}(\mathbf{x}, \mathbf{y}^k) \leq \mathbf{0}$ $\mathbf{x} \in [\mathbf{x}^L, \mathbf{x}^U] \subset \mathbb{R}^{n_{\mathbf{x}}}$

Primal problem $NLPF(\mathbf{y}^k)$: if infeasible $NLP(\mathbf{y}^k)$

$$\begin{aligned} \mathbf{x}^k \in \arg\min_{\mathbf{x},\omega} \ \omega \\ \text{s.t. } \mathbf{g}(\mathbf{x},\mathbf{y}^k) &\leq \omega \\ \mathbf{x} \in [\mathbf{x}^{\mathrm{L}},\mathbf{x}^{\mathrm{U}}] \subset \mathbb{R}^{n_{\mathrm{x}}}, \ \omega \geq 0 \end{aligned}$$



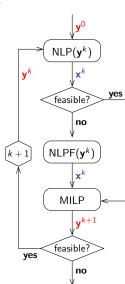
Outer-Approximation Algorithm

• Idea. Alternate solution between NLP and MILP

Master problem MILP: linearize f and $g @ (x^k, y^k)$

$$\begin{aligned} \mathbf{y}^{k+1} &\in \arg\min_{\mathbf{x},\mathbf{y},\eta} \, \eta \\ \text{s.t.} \ \forall j \leq k, \ \text{feasible NLP}(\mathbf{y}^j) \colon \\ f(\mathbf{x}^j,\mathbf{y}^j) + \boldsymbol{\nabla} f(\mathbf{x}^j,\mathbf{y}^j) \left[\begin{array}{c} \mathbf{x} - \mathbf{x}^j \\ \mathbf{y} - \mathbf{y}^j \end{array} \right] \leq \eta \\ \forall j \leq k \colon \\ \mathbf{g}(\mathbf{x}^j,\mathbf{y}^j) + \boldsymbol{\nabla} \mathbf{g}(\mathbf{x}^j,\mathbf{y}^j) \left[\begin{array}{c} \mathbf{x} - \mathbf{x}^j \\ \mathbf{y} - \mathbf{y}^j \end{array} \right] \leq \mathbf{0} \\ \mathbf{x} \in [\mathbf{x}^{\mathrm{L}},\mathbf{x}^{\mathrm{U}}] \subset \mathbb{R}^{n_{\mathrm{x}}}, \ \mathbf{y} \in \{0,1\}^{n_{\mathrm{y}}} \end{aligned}$$

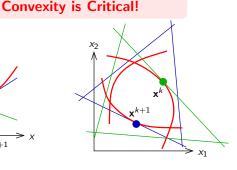
New cuts appended at each iteration: non-decreasing lower bound

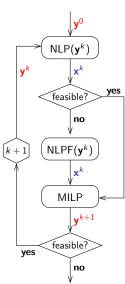


Outer-Approximation Algorithm

• Idea. Alternate solution between NLP and MILP

f(x) $x^{k} \qquad x^{k+1} \qquad x$





- OA Convergence and Performance
 - Finite termination
 - Fewer NLP subproblems than B&B is typical
 - ▶ Solving master MILP can be the bottleneck

Outer-Approximation Example

Workshop. Consider the following MINLP model and apply the OA algorithm, starting from $\mathbf{y}^0 = (1, 0, 0)$.

$$\min_{\mathbf{x}, \mathbf{y}} x_1^2 + x_2^2 + y_1 + 1.5y_2 + 0.5y_3$$
s.t. $(x_1 - 2)^2 - x_2 \le 0$

$$x_1 \ge 2y_1$$

$$x_1 - x_2 \le 4(1 - y_2)$$

$$x_1 \ge (1 - y_1)$$

$$x_2 \ge y_2$$

$$x_1 + x_2 \ge 3y_3$$

$$y_1 + y_2 + y_3 \ge 1$$

$$0 \le x_1, x_2 \le 4$$

$$y_1, y_2, y_3 \in \{0, 1\}$$

$$\min_{\mathbf{x},\mathbf{y}} x_1^2 + x_2^2 + y_1 + 1.5y_2 + 0.5y_3
s.t. (x_1 - 2)^2 - x_2 \le 0
x_1 \ge 2y_1
x_1 - x_2 \le 4(1 - y_2)
x_1 \ge (1 - y_1)
x_2 \ge y_2
x_1 + x_2 \ge 3y_3
y_1 + y_2 + y_3 \ge 1
0 \le x_1, x_2 \le 4
(y_1, y_2, y_3) = (1, 0, 0)$$

Upper Bound: UBD⁰ = 5,

$$x_1^0 = 2, x_2^0 = 0,$$

 $y_1^0 = 1, y_2^0 = 0, y_3^0 = 0$

 $x_1^* = x_2^* = 1$,

Optimal solution: $f^* = 3.5$,

 $y_1^* = 0, y_2^* = 1, y_2^* = 0$

Outer-Approximation Example

$$\begin{array}{l} \underset{x,y,\eta}{\min} \ \eta \\ \text{s.t.} \ 4x_1 + y_1 + 1.5y_2 + 0.5y_3 - 4 \leq \eta \\ & - x_2 \leq 0 \\ & x_1 \geq 2y_1 \\ & x_1 - x_2 \leq 4(1 - y_2) \\ & x_1 \geq (1 - y_1) \\ & x_2 \geq y_2 \\ & x_1 + x_2 \geq 3y_3 \\ & y_1 + y_2 + y_3 \geq 1 \\ & 0 \leq x_1, x_2 \leq 4 \\ & y_1, y_2, y_3 \in \{0, 1\} \\ & & \blacktriangleright \ \text{Lower Bound: LBD}^1 = 0.5, \\ & x_1^1 = 1, x_2^1 = 4, y_1^1 = 0, y_2^1 = 0, \\ & y_2^1 = 1 \end{array}$$

$$\min_{\mathbf{x}, \mathbf{y}} x_1^2 + x_2^2 + y_1 + 1.5y_2 + 0.5y_3$$
s.t. $(x_1 - 2)^2 - x_2 \le 0$

$$x_1 \ge 2y_1$$

$$x_1 - x_2 \le 4(1 - y_2)$$

$$x_1 \ge (1 - y_1)$$

$$x_2 \ge y_2$$

$$x_1 + x_2 \ge 3y_3$$

$$y_1 + y_2 + y_3 \ge 1$$

$$0 \le x_1, x_2 \le 4$$

$$(y_1, y_2, y_3) = (0, 0, 1)$$

Upper Bound: UBD¹ = 5, $x_1^1 = 1.5, x_2^1 = 1.5,$ $y_1^1 = 0, y_2^1 = 0, y_3^1 = 1$

Outer-Approximation Example

$$\min_{\mathbf{x}, \mathbf{y}, \eta} \eta$$
s.t. $4x_1 + y_1 + 1.5y_2 + 0.5y_3 - 4 \le \eta$

$$3x_1 + 3x_2 + y_1 + 1.5y_2 + 0.5y_3 - 4.5 \le \eta$$

$$- x_2 \le 0$$

$$- x_1 - x_2 + 1.75 \le 0$$

$$x_1 \ge 2y_1$$

$$x_1 - x_2 \le 4(1 - y_2)$$

$$x_1 \ge (1 - y_1)$$

$$x_2 \ge y_2$$

$$x_1 + x_2 \ge 3y_3$$

$$y_1 + y_2 + y_3 \ge 1$$

$$0 \le x_1, x_2 \le 4$$

$$y_1, y_2, y_3 \in \{0, 1\}$$

$$\blacktriangleright \text{ Lower Bound: LBD}^2 = 3,$$

$$x_1^2 = 1, x_2^2 = 1, y_1^2 = 0, y_2^2 = 1,$$

$$\min_{\mathbf{x},\mathbf{y}} x_1^2 + x_2^2 + y_1 + 1.5y_2 + 0.5y_3
\mathbf{s.t.} (x_1 - 2)^2 - x_2 \le 0
x_1 \ge 2y_1
x_1 - x_2 \le 4(1 - y_2)
x_1 \ge (1 - y_1)
x_2 \ge y_2
x_1 + x_2 \ge 3y_3
y_1 + y_2 + y_3 \ge 1
0 \le x_1, x_2 \le 4
(y_1, y_2, y_3) = (0, 1, 0)$$

Upper Bound: UBD² = 3.5,

$$x_1^2 = 1, x_2^2 = 1,$$

 $y_1^2 = 0, y_2^1 = 2, y_3^2 = 0$

Outer-Approximation Example

min η

```
x,y,\eta
s.t. 4x_1 + y_1 + 1.5y_2 + 0.5y_3 - 4 \le \eta
     3x_1 + 3x_2 + v_1 + 1.5v_2 + 0.5v_3 - 4.5 \le n
     2x_1 + 2x_2 + v_1 + 1.5v_2 + 0.5v_3 - 2 \le n
      -x_2 < 0
      -x_1-x_2+1.75 < 0
      -2x_1-x_2+3 \le 0
     x_1 > 2v_1
     x_1 - x_2 \le 4(1 - y_2)
     x_1 > (1 - y_1)
     x_2 > v_2
     x_1 + x_2 > 3v_3
     y_1 + y_2 + y_3 > 1
     0 < x_1, x_2 < 4
     y_1, y_2, y_3 \in \{0, 1\}
```

Convergence within 3 iterations

Lower Bound: $LBD^3 = 3.5,$ $x_1^3 = 1, x_2^3 = 1, y_1^3 = 0,$ $y_2^3 = 1, y_2^3 = 0$

Contents

Solving MINLPs is challenging in general – We will learn some typical decomposition techniques for MINLPs with either convex or nonconvex participating functions

- Methods for MINLP with Convex Functions
 - Branch-and-Bound Algorithm
 - Outer-Approximation Algorithm
- Methods for MINLP with Nonconvex Functions

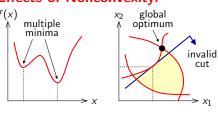


Recommended Readings:

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- I.E. Grossmann, Review of nonlinear mixed-integer and disjunctive programming techniques, *Optimization & Engineering* 1:227-252, 2002

Main Solution Algorithms

Effects of Nonconvexity:



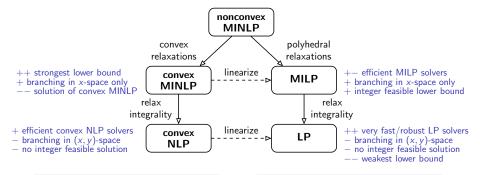
Heuristic Approaches:

- Use convex MINLP techniques, assuming valid lower bound
- Add slacks to the relaxations/ linearizations
- May remain tractable for large-scale nonconvex MINLP!

Rigorous Approaches: Heavily reliant on convex relaxation techniques!

- Gmin- α BB (Adjiman et al, 1997, 2000)
- BARON (Ryoo & Sahinidis, 1995; Tawarmalani & Sahinidis, 2002)
- Reformulation/relaxation (Smith & Pantelides, 1999; Zamora & Grossmann, 1999)
- OA for nonconvex MINLP (Kesavan et al, 2004)

Lower-Bounding Techniques



Convex MINLP Relaxation

$$\min_{\mathbf{x},\mathbf{y}} \, \mathbf{f}^{\mathrm{cv}}(\mathbf{x},\mathbf{y})$$

s.t.
$$\mathbf{g}^{cv}(\mathbf{x},\mathbf{y}) \leq \mathbf{0}$$

$$\boldsymbol{x} \in [\boldsymbol{x}^L, \boldsymbol{x}^U], \ \boldsymbol{y} \in \{0, 1\}$$

Convex NLP Relaxation

$$\min_{\mathbf{x},\mathbf{y}} \mathbf{f}^{cv}(\mathbf{x},\mathbf{y})$$

s.t.
$$\mathbf{g}^{cv}(\mathbf{x},\mathbf{y}) \leq \mathbf{0}$$

$$extbf{x} \in [extbf{x}^{\mathrm{L}}, extbf{x}^{\mathrm{U}}], \ extbf{y} \in [\![0, 1]\!]$$

Upper-Bounding Techniques

Local Solution of Nonconvex MINLP

Conceptually straightforward, but:

- Lacks robustness
- Comes with high computational burden
- 2 Local Solution of Nonconvex NLP

Requires a fixing of binary variables:

- Solution of convex MINLP/MINLP lower bound
- Rounding of convex NLP/LP lower bound