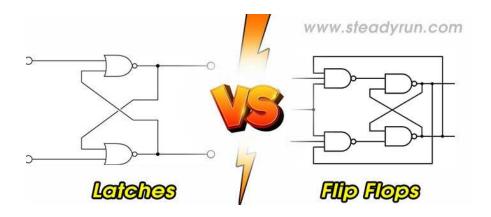
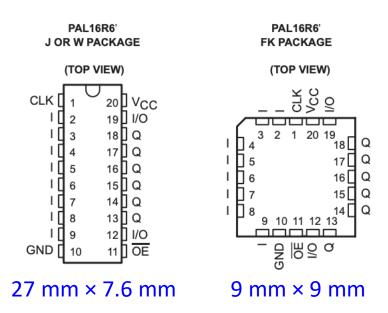
EE2000 Logic Circuit Design

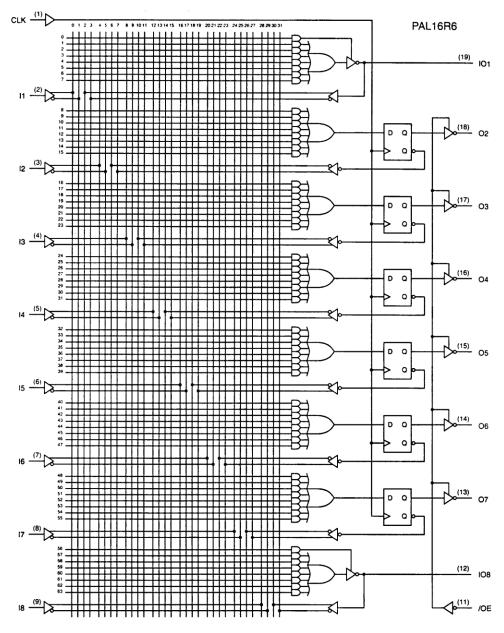
Lecture 7 – Latch and Flip-Flop Circuits



PAL16R6

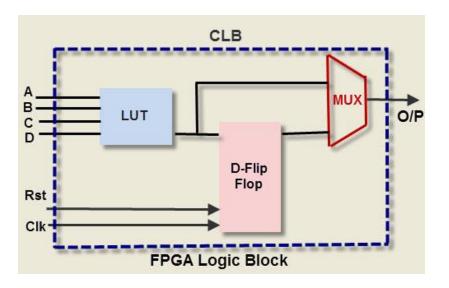


- A PAL device with 16 inputs, 6
 outputs with D-latch (1-bit memory
 device), and 2 un-buffered outputs
- 8 inputs are feedback from outputs
- This device can be used for designing a sequential circuit (Lecture 7)



Configurable Logic Block

- A configurable logic block (CLB) consists of a look-uptable (LUT)
- The LUT is for implementing the **truth table** of a logic function (e.g., 4-input Boolean function)
- Additional elements are for control or other functions:
 D-FF (store output), MUX (select output)



What will you learn?

- 7.1 Learn different latches Active High & Active Low
- 7.2 Learn different Flip-Flop (FF) SR-, D-, JK-, T-FFs
- 7.3 Learn the operation of a Master-Slave FF
- 7.4 Learn the operation of an Edge-Triggered FF
- 7.5 Learn the operation of Synchronous and Asynchronous Inputs

Two Classes of Logic Circuits

Combinational logic circuit

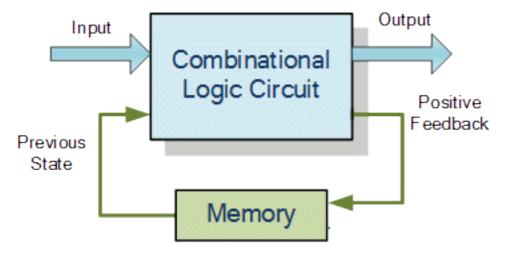
Output depends only on the inputs (As discussed in previous lectures)

Sequential logic circuit

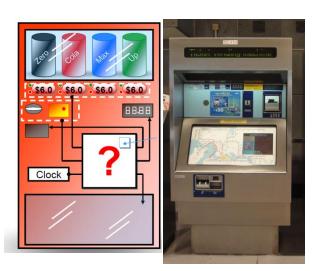
Output depends on present input + past history

Memory circuit (to store previous STATE information)

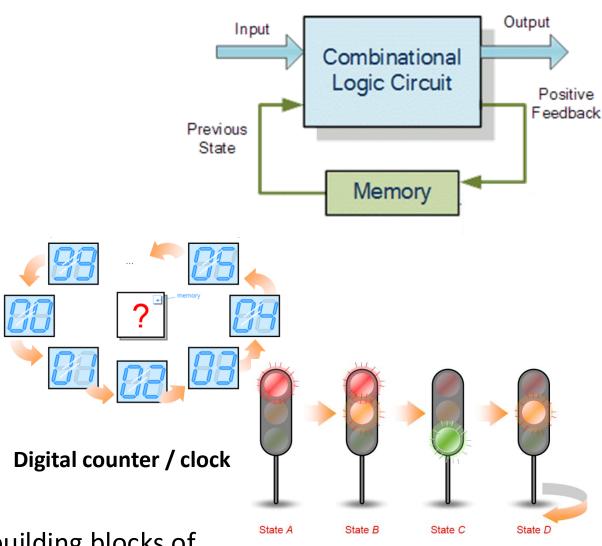
is required



Sequential Logic Circuit



Vending Machine Food/ drink / MTR tickets

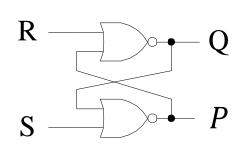


Latches and Flip-Flops are building blocks of sequential logic circuit

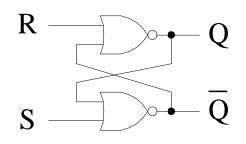
7.1 Latches

- Latch: A binary storage device, composed of logic gates with feedback
- The output state is controlled by its excitation input signals
- Two inputs: R stands for reset, S stands for set
- Two Outputs: Q stores the state of the latch (binary) value)
- Bistable circuit: Two states

Active High Latch



$$Q = (R + P)'$$
$$P = (S + Q)'$$



When S and R are 0,

$$Q = P'$$
 and $P = Q' = \bar{Q}$

The latch can store either

0
$$(Q = 0, \bar{Q} = 1)$$
 or

1
$$(Q = 1, \bar{Q} = 0)$$

• When S = 1, R = 0

$$Q=1$$
, $\overline{Q}=0$

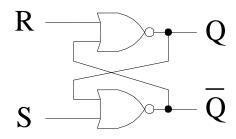
• When S = 0, R = 1

$$Q = 0, \bar{Q} = 1$$

When S and R are 1,

$$Q=0, \bar{Q}=0$$

Active High Latch

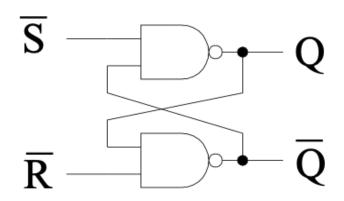




S	R	Q_{t+1} Q'_{t+1}		State
0	0	Q_t	Q_t'	Hold state
0	1	0	1	Reset state
1	0	1	0	Set state
1	1	0	0	Undefined state

t	1	2	3	4	5	6	7	8	9
S	1	0	0	0	0	1	1	1	0
R	0	0	1	0	0	0	1	0	0
Q	1	1	0	0	0	1	0	1	1
$ar{Q}$	0	0	1	1	1	0	0	0	0
State	S	Н	R	Н	Н	S	U	S	Н

Active Low Latch

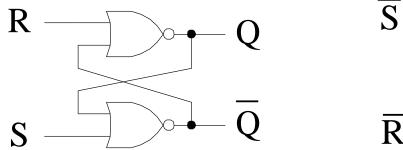


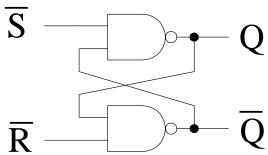
S'	R'	Q _{t+1}	Q' _{t+1}	State
1	1	Q_t	Q_t'	Hold state
1	0	0	1	Reset state
0	1	1	0	Set state
0	0	1	1	Undefined state

t	1	2	3	4	5	6	7	8	9
Ī	1	1	0	1	1	1	0	1	1
$ar{R}$	0	1	1	1	1	0	0	0	1
Q	0	0	1	1	1	0	1	0	0
$ar{Q}$	1	1	0	0	0	1	1	1	1
State	R	Н	S	Н	Н	R	U	R	Н

Summary

Note that the input is <u>active high</u> for NOR gate implementation, whereas the input is <u>active low</u> for NAND gate implementation



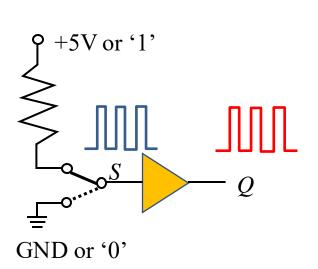


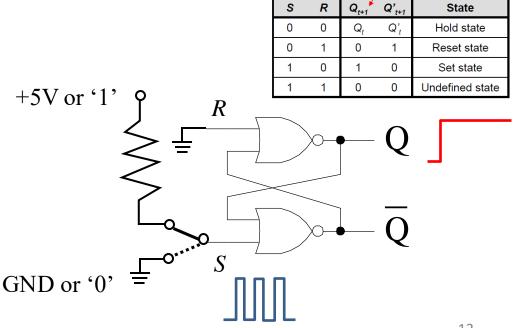
Require a signal "1" to change its state

Require a signal "0" to change its state

Application – Key Debouncing Circuit

When pressing a key or switch, its internal spring may cause the contact point to vibrate for a short period of time, causing a series of "ON" and "OFF" state before settling down



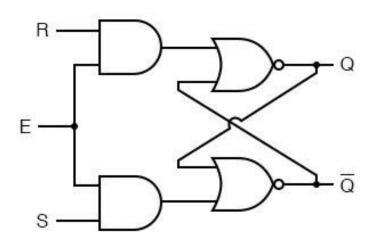


Next state of Q

Drawback - Noise

If inputs *S*, *R* are unstable (or noise), output produces an unstable short pulse

Solution: Gated Latch!

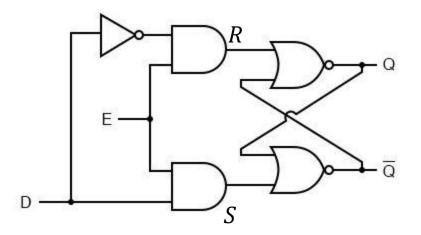


Ε	S	R	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Х	Х	Q_t	$\overline{Q_t}$	Hold
1	0	0	Q_t	$\overline{Q_t}$	Hold
1	0	1	0	1	Reset
1	1	0	1	0	Set
1	1	1	0	0	Undefined

Drawback – Undefined State

The latches will enter undefined state for some inputs (e.g. R = S = 1)

Solution: D Latch!



E	D	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Х	Q_t	$\overline{Q_t}$	Hold
1	1	1	0	Set
1	0	0	1	Reset

Remember in this way:

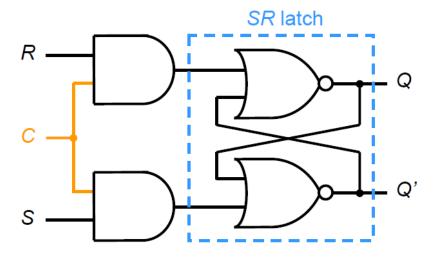
The value of *D* is transferred to the *Q* output

7.2 Flip-Flops

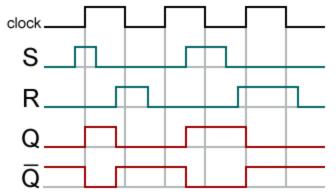
- Flip-Flop: A clocked binary storage device that the change of output state is controlled by the clock signal (synchronized)
- We will discuss the following four types of FF
 - Set-Reset FF (SR FF)
 - Data/Delay FF (D FF)
 - Jack Kilby FF (JK FF)
 - Toggle FF (T FF)

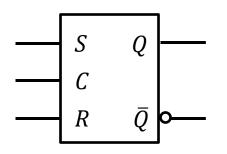
Set-Reset FF

A clocked signal is connected to the Enable input of a gated latch



С	S	R	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Х	Х	Q_t	$\overline{Q_t}$	Hold
1	0	0	Q_t	$\overline{Q_t}$	Hold
1	0	1	0	1	Reset
1	1	0	1	0	Set
1	1	1	0	0	Undefined

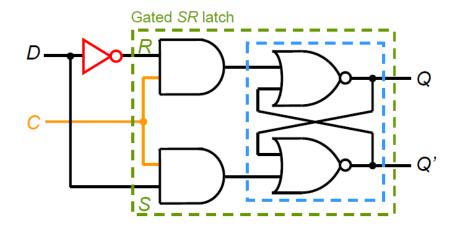




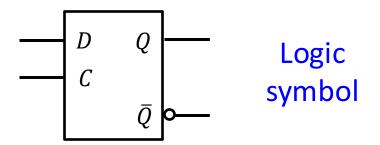
Logic symbol

Data/Delay FF

A clocked signal is connected to the Enable input of a D latch

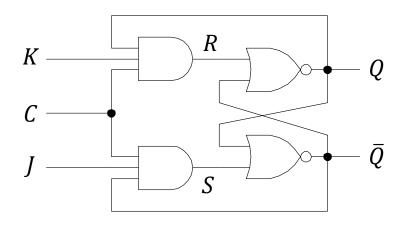


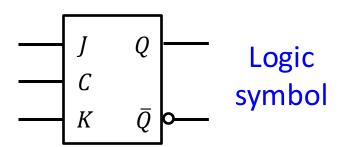
С	D	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Х	Q_t	$\overline{Q_t}$	Hold
1	1	1	0	Set
1	0	0	1	Reset



Jack-Kilby FF

Refined SR FF whereby the undetermined state of SR FF is defined; **J** is set and **K** is reset





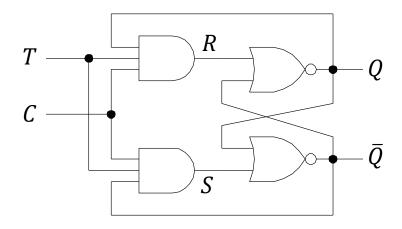
С	J	K	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Х	Х	Q_t	$\overline{Q_t}$	Hold
1	0	0	Q_t	$\overline{Q_t}$	Hold
1	0	1	0	1	Reset
1	1	0	1	0	Set
1	1	1	$\overline{Q_t}$	Q_t	Toggle

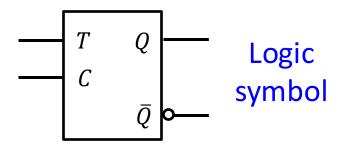


С	J	K	Q_t	$\overline{Q_t}$	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	1	0	1	1	0
1	1	1	1	0	0	1

Toggle FF

Output state changes for each clock pulse when **T** is in its active state





С	T	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Х	Q_t	$\overline{Q_t}$	Hold
1	0	Q_t	$\overline{Q_t}$	Hold
1	1	$\overline{Q_t}$	Q_t	Toggle



С	T	Q_t	$\overline{Q_t}$	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	0	1	1	0
1	1	1	0	0	1

Summary

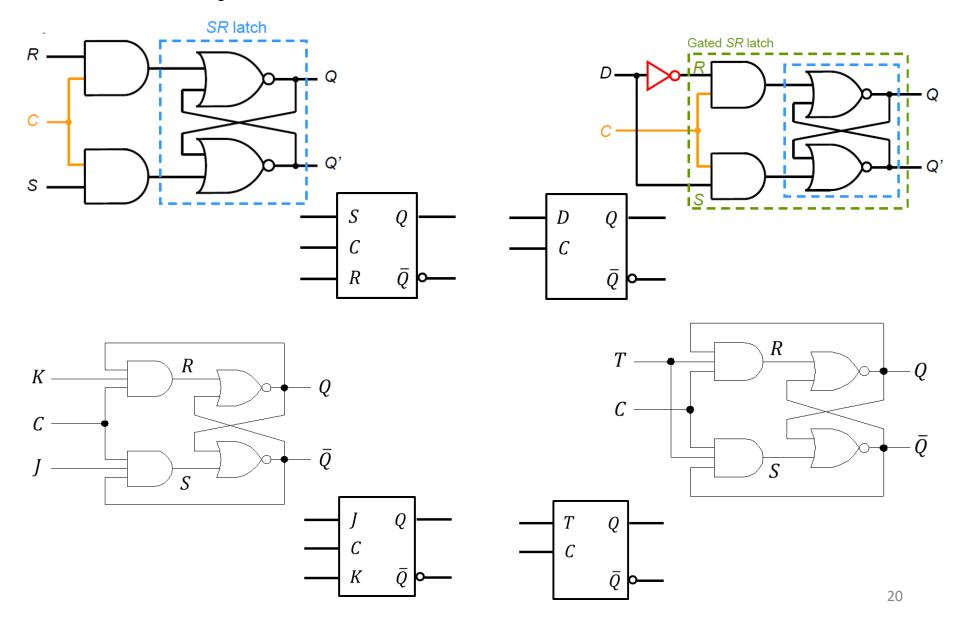


Table Given (Test and Exam)

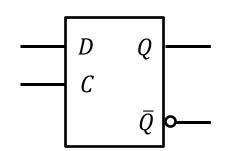
S	R	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	0	Q_t	$\overline{Q_t}$	Hold
0	1	0	1	Reset
1	0	1	0	Set
1	1	1	1	Undefined

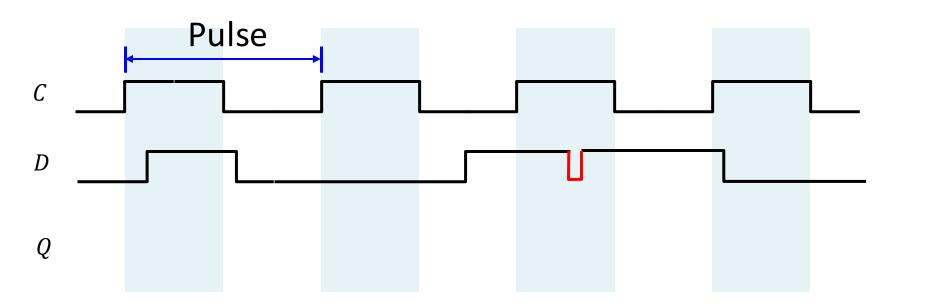
J	K	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	0	Q_t	$\overline{Q_t}$	Hold
0	1	0	1	Reset
1	0	1	0	Set
1	1	$\overline{Q_t}$	Q_t	Toggle

D	Q_{t+1}	$\overline{Q_{t+1}}$	State
1	1	0	Set
0	0	1	Reset

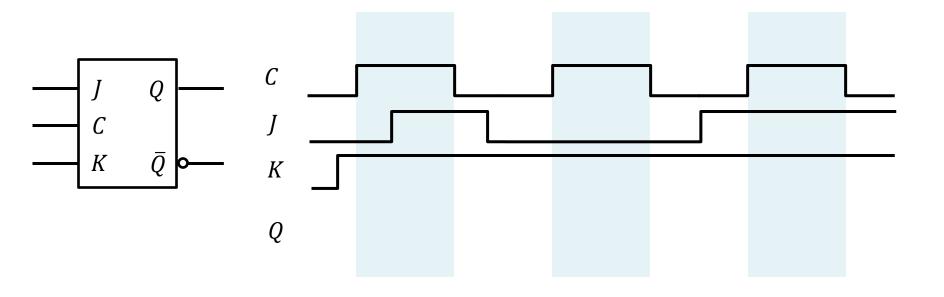
T	Q_{t+1}	$\overline{Q_{t+1}}$	State
0	Q_t	$\overline{Q_t}$	Hold
1	$\overline{Q_t}$	Q_t	Toggle

Exercise



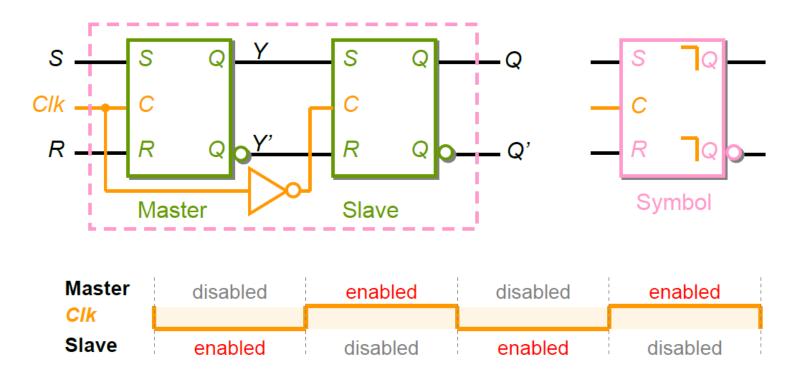


Exercise

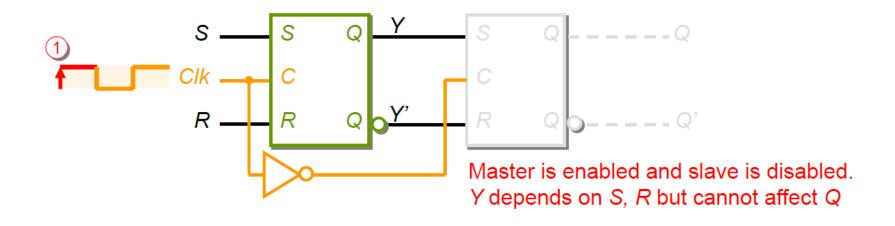


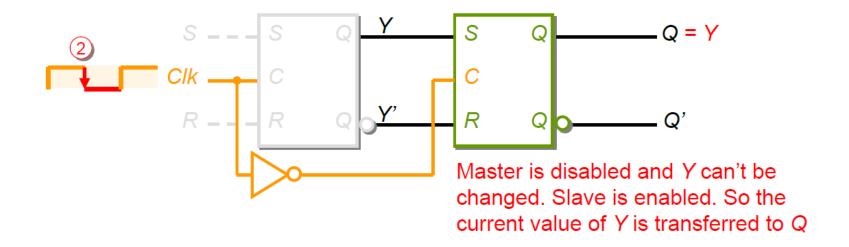
7.3 Master-Slave Flip-Flop

- A combination of two flip-flops together in a series configuration
- At any time, only one flip-flop is enabled



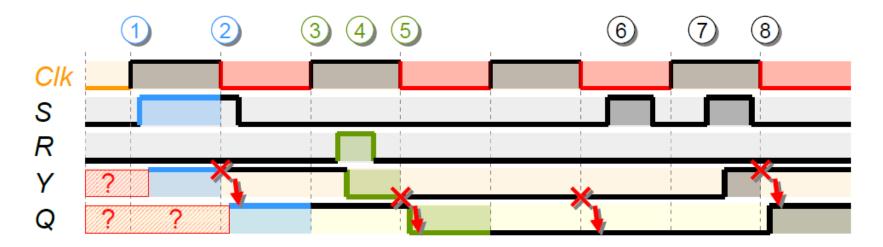
How Does It Work?





Timing Diagram

Initial inputs: *S* and *R* are LOW Initial state of *Y* and *Q* unknown



- 1 Clk switches to 1; master FF enabled; S = 1, R = 0, $\therefore Y = 1$
- 2 Clk switches to 0; slave FF enabled; Y = 1, $\therefore Q = 1$
- 3 Clk switches to 1; master FF enabled; S = 0, R = 0, HOLD
- (6) Clk is 0; master FF disabled; no change on Y.
- (7) Clk switches to 1; master FF enabled; S = 1, R = 0, $\therefore Y = 1$
- (8) Clk switches to 0; slave FF enabled; Y = 1, : Q = 1

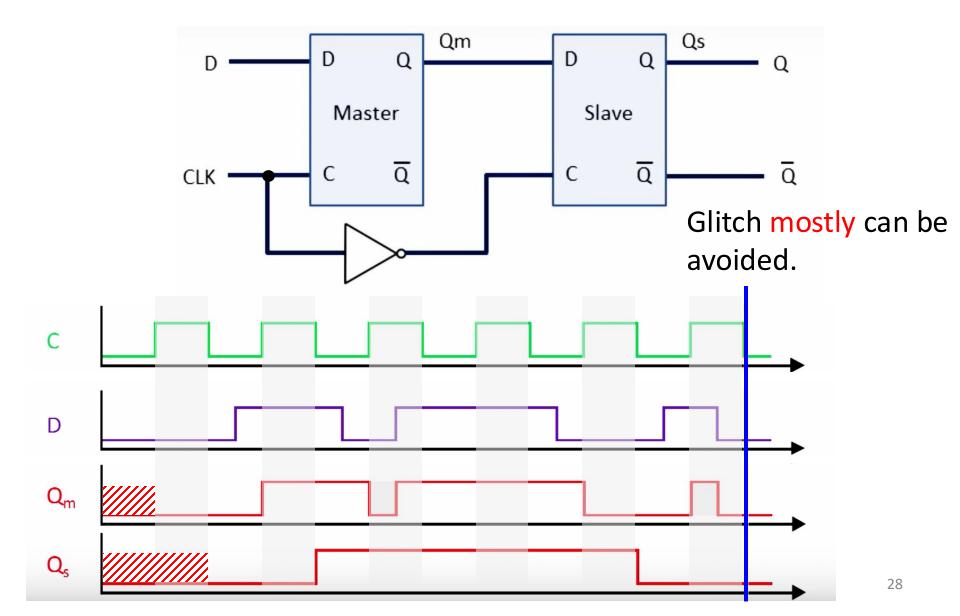
Characteristics

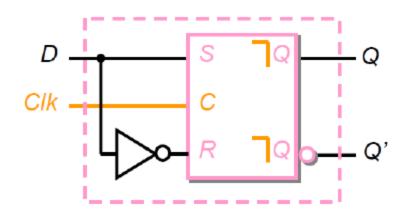


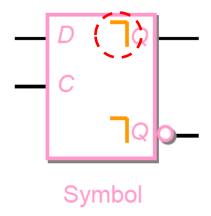
- Y is updated when Clk goes to and remains at 1
- Q is updated when Clk goes to 0 and remains unchanged until the next pulse

Therefore,

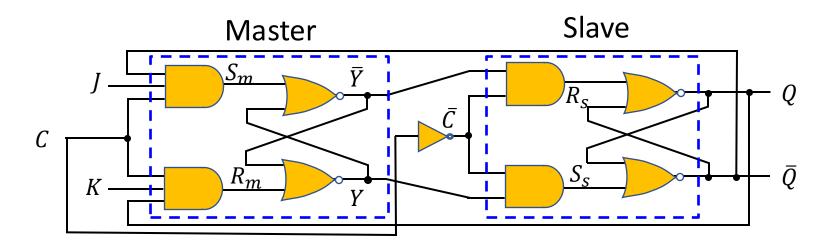
- > For each clock pulse, Q only changes state at most once
- > State change only occurs during the transition of Clk from 1 to 0



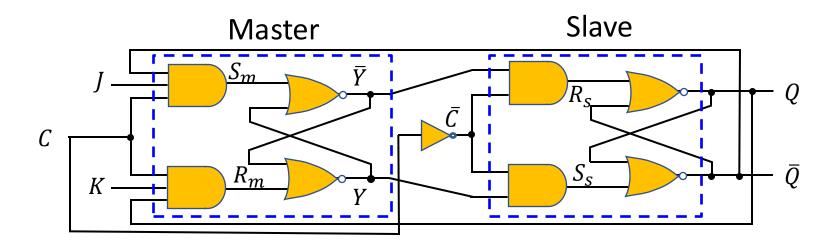


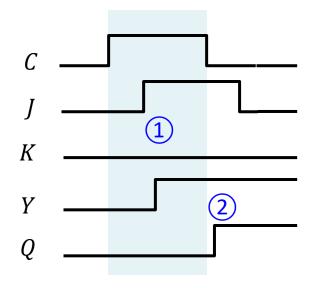


Inp	uts		Output	s
D	CIk	Next Q	Next Q'	Meaning
Х	0	Q	Q'	Hold
Х		D	D'	Set / Reset



- Two JK FFs connected in series
- Outputs are fedback to Master's inputs $Q \rightarrow K \& \bar{Q} \rightarrow J$
- *J* is SET and *K* is RESET
- Y is connected to SET of slave FF



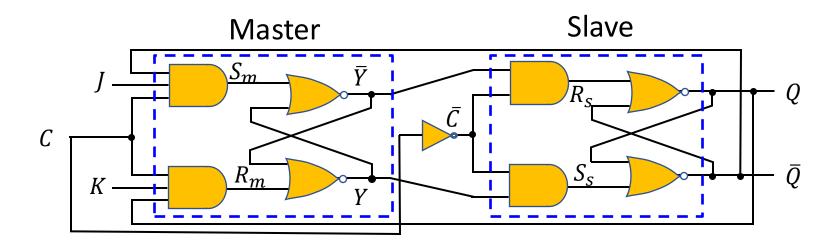


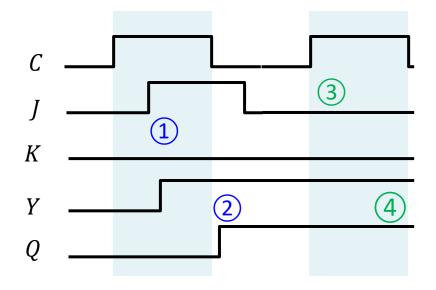
$$\mathbf{1}J = 1, K = 0,$$

$$S_m = \overline{Q} = 1, R_m = 0(SET)$$

$$\therefore Y = 1$$

②
$$Y(S_s) = 1, \overline{Y}(R_s) = 0$$
 (SET)
 $\therefore Q = 1$



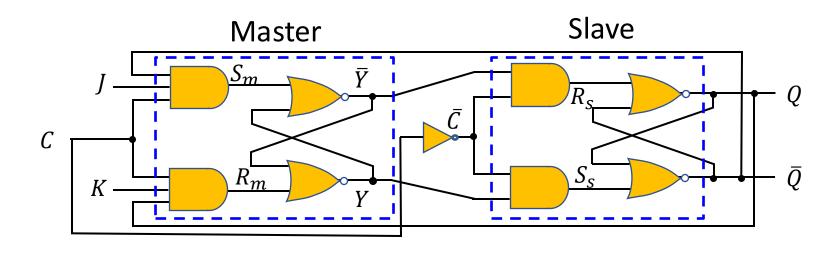


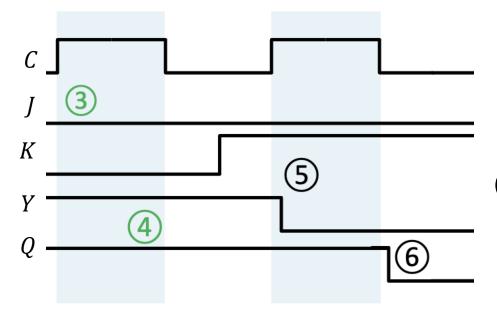
$$3J = 0, K = 0(HOLD)$$

$$\therefore Y = 1$$

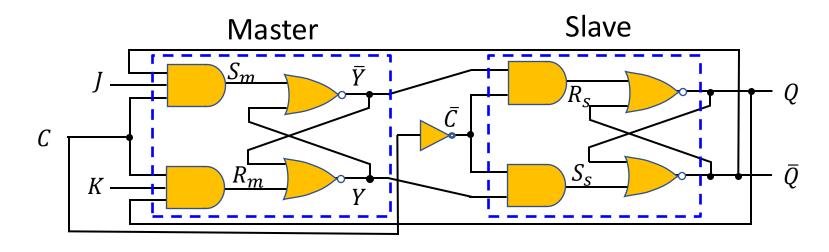
$$(4) Y(S_s) = 1, \overline{Y}(R_s) = 0 \text{ (SET)}$$

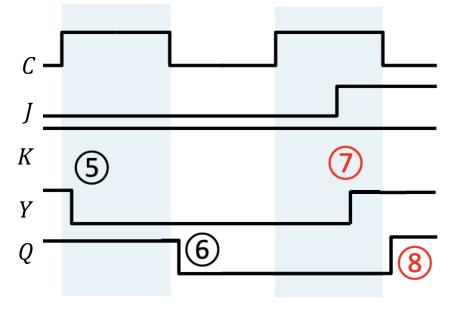
$$\therefore Q = 1$$





$$(5) J = 0, K = 1 (RESET)$$
∴ $Y = 0$

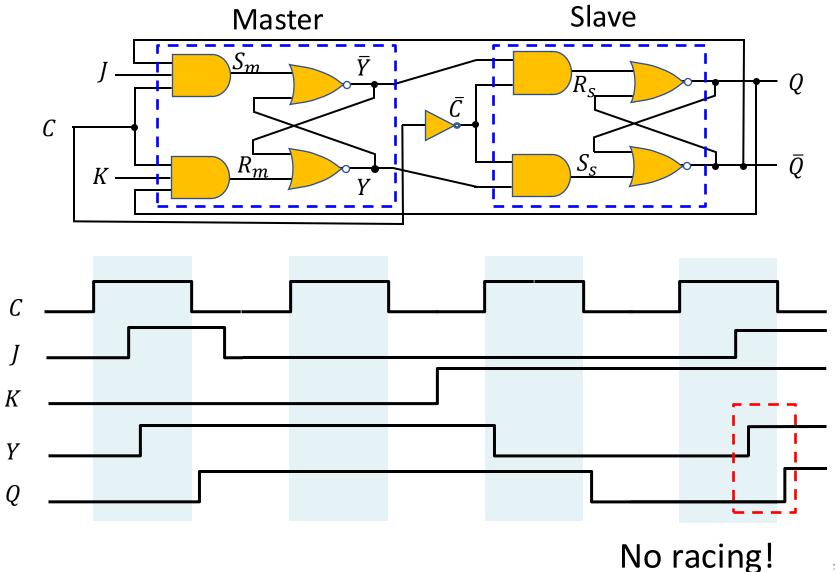




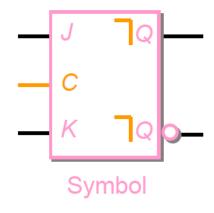
7
$$J = 1, K = 1$$

 $S_m = \overline{Q} = 1, R_m = 0 \text{(SET)}$
 $\therefore Y = 1 \text{ at the point C} \rightarrow 0$

$$\begin{array}{c} \text{(8)} \ Y(S_S) = 1, \overline{Y}(R_S) = 0 \\ \text{(SET)} \\ \therefore \ Q = 1 \end{array}$$



Inputs			Outputs		
J	K	CIk	Next Q	Next Q'	Meaning
Х	Х	0	Q	Q'	Hold
0	0	J	Q	Q'	Hold
0	1	J	0	1	Reset
1	0	J	1	0	Set
1	1		Q'	Q	Toggle



The same behaviors as master-slave *SR* flip-flop

Eliminated the undefined state

Problem

- For the presented MSFF, master is enabled during the period when the clock pulse is 1 (Level triggered
- Undesired behavior produced when inputs values of S and R (or J, K, D, T) change at that period, especially just before the clock pulse changes to 0
- Undesired output due to glitch could not be fully avoided
- Better solution: Edge-triggered FFs

7.4 Edge-triggered Flip-Flop

- Use one of the edges of the clock signal to read the input values
 - either positive or negative transition
 - triggering edge
- The response to the triggering edge at the output of the flip-flop is almost immediate
- The flip-flops remains unresponsive to the input change until the next triggering edge

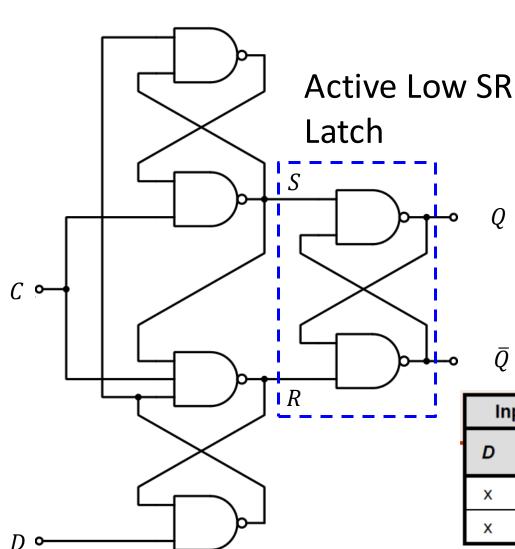
Triggering edge

- Triggered only during a signal transition from 0 to 1 (or from 1 to 0) on the clock
 - Positive (rising, leading) edge-triggering: 0-to-1 transition

Negative (falling, trailing) edge-triggering: 1-to-0 transition

Additional circuit is included to ensure it will only response to the input at the transition edge of the clock pulse

Positive-Edge-Triggered D-FF

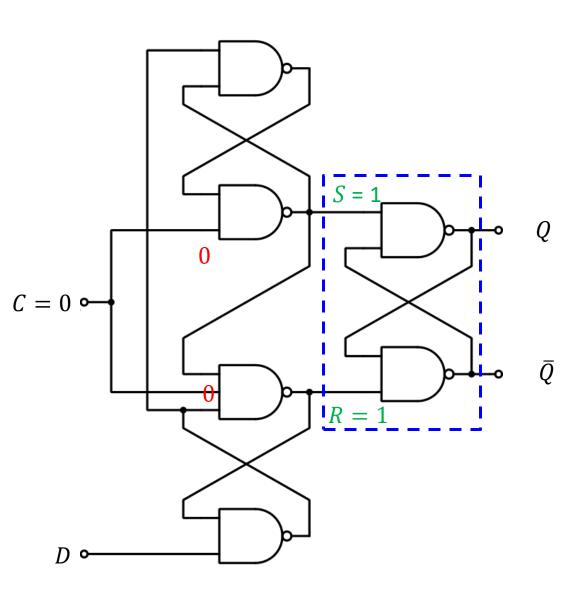


S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

Inputs		Outputs			
D	CIk	Next Q	Next Q'	State	
X	↓,0,1	Q	Q'	Hold	
X	1	D	D'	Set / Reset	

 \bar{Q}

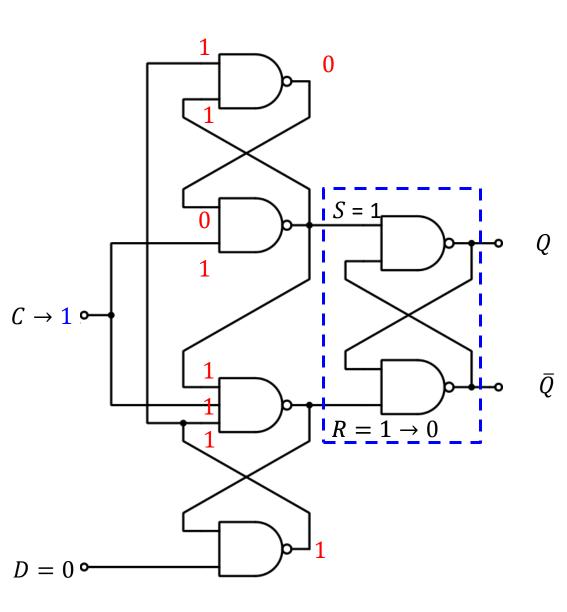
When C = 0



S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

HOLD!

When $C = 0 \rightarrow 1 \& D = 0$

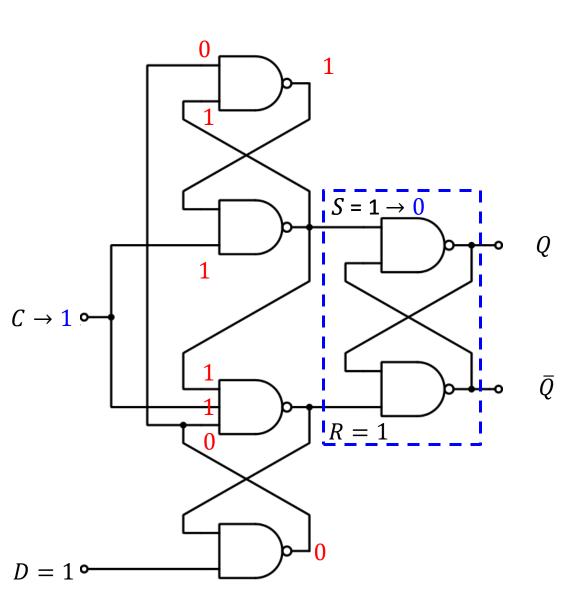


S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

HOLD → RESET

When $C = 0 \to 1 \& D$ = 0, $Q \to 0$

When $C = 0 \rightarrow 1 \& D = 1$

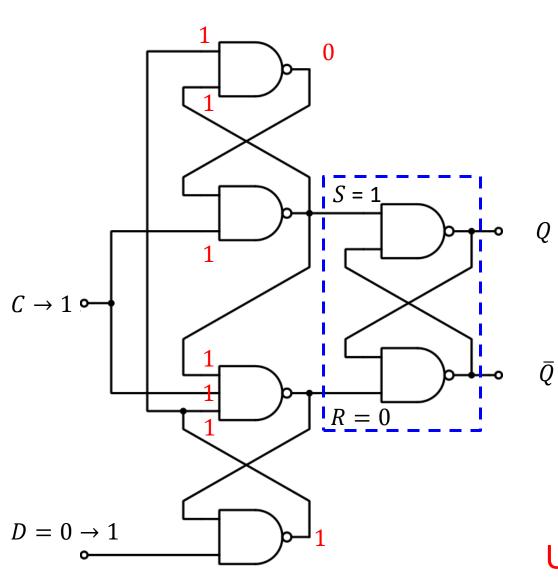


S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

 $HOLD \rightarrow SET$

When $C = 0 \to 1 \& D$ = 1, $Q \to 1$

When $C = 1 \& D = 0 \rightarrow 1$



S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

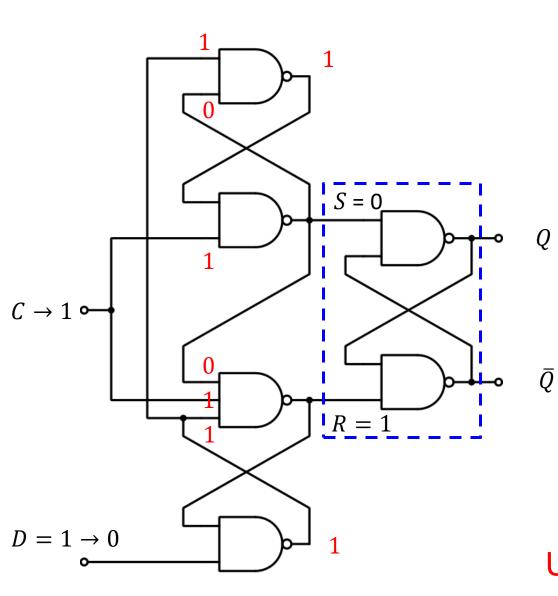
RESET → RESET

When
$$C = 1 \& D = 0$$

 $\rightarrow 1, Q = 0$

Unresponsive!

When $C = 1 \& D = 1 \rightarrow 0$



S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

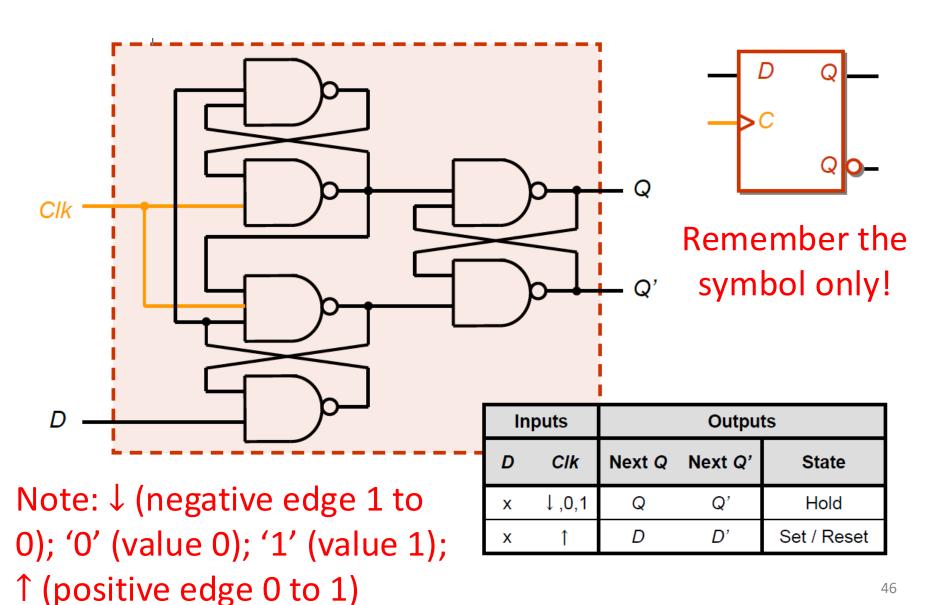
 $SET \rightarrow SET$

When
$$C = 1 \& D = 1$$

 $\rightarrow 0, Q = 1$

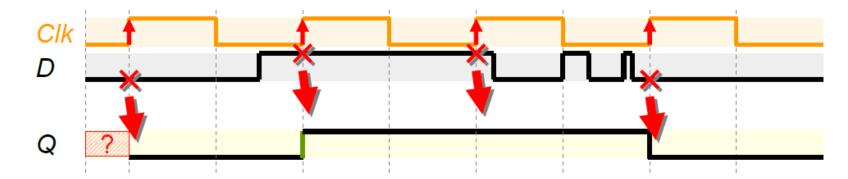
Unresponsive!

Positive-Edge-Triggered D-FF

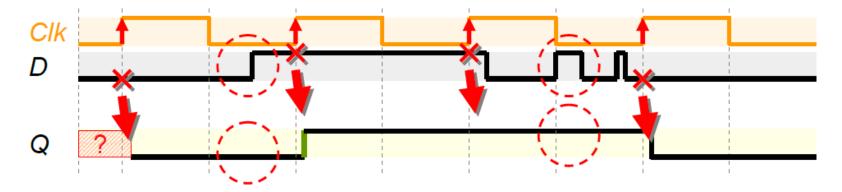


Timing Diagram

Timing Diagram of PET D FF (if no output delay)



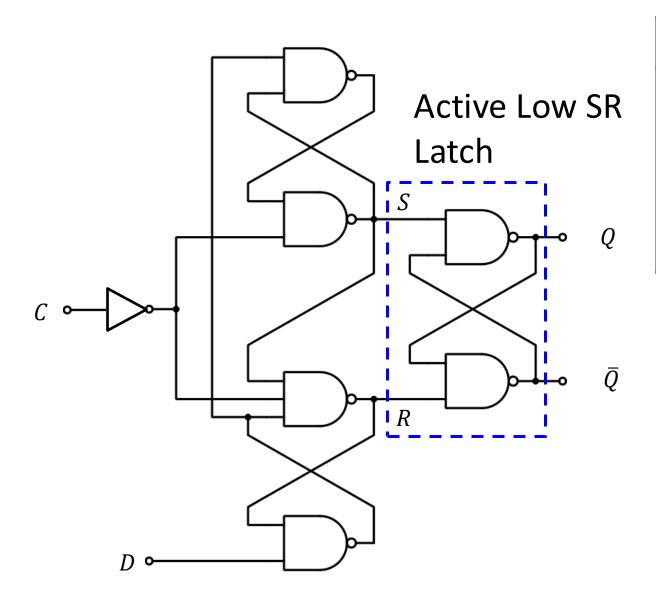
Timing Diagram of PET D FF



Thought

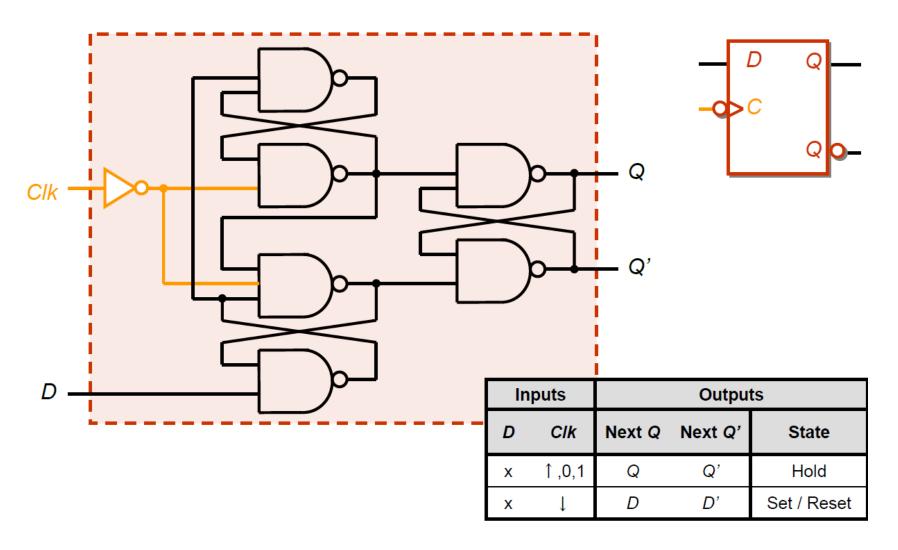
How to build a negative-edge triggered D-FF?

Negative-Edge-Triggered D-FF



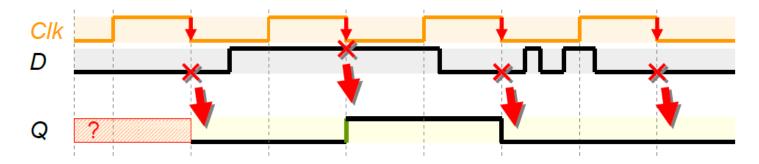
S	R	Q_{t+1}	$\overline{Q_{t+1}}$
1	1	Q_t	$\overline{Q_t}$
1	0	0	1
0	1	1	0
0	0	1	1

Negative-Edge-Triggered D-FF

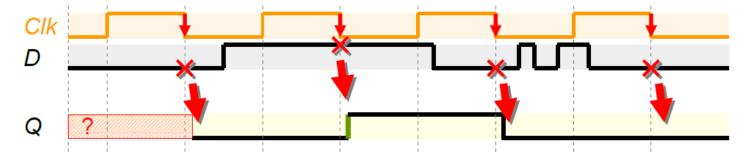


Timing Diagram

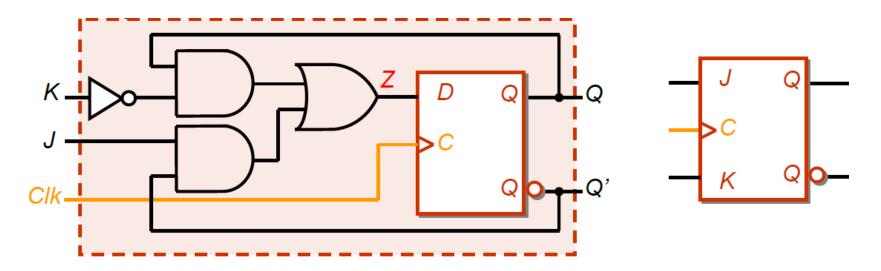
Timing Diagram of NET D FF (if no output delay)



Timing Diagram of NET D FF



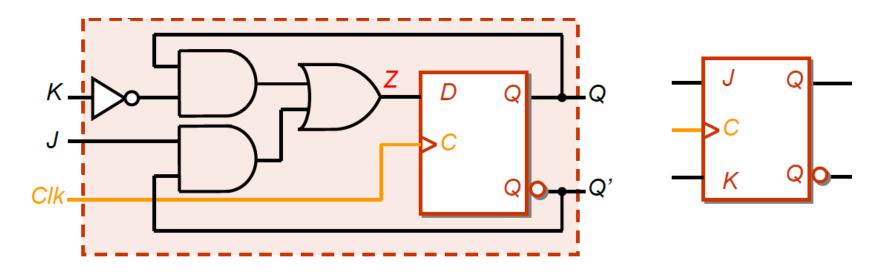
Edge-Triggered JK-FF



$$D = Q\overline{K} + J\overline{Q}$$

J	K	D	STATE
0	0	Q	HOLD
0	1	0	RESET
1	0	$Q + \bar{Q} = 1$	SET
1	1	$ar{Q}$	TOGGLE

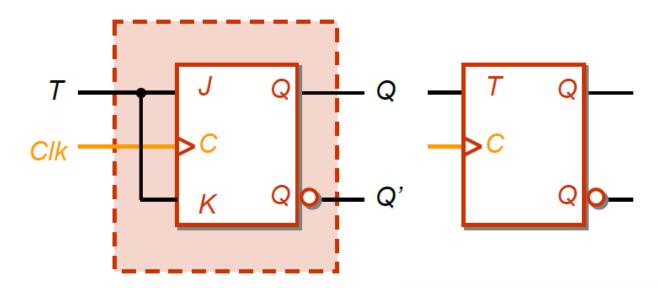
Edge-Triggered JK-FF



$$D = Q\overline{K} + J\overline{Q}$$

Inputs				Outputs	
J	K	CIk	Next Q	Next Q'	State
Х	Х	↓,0,1	Q	Q'	Hold
0	0	1	Q	Q'	Hold
0	1	1	0	1	Reset
1	0	1	1	0	Set
1	1	1	Q'	Q	Toggle

Edge-Triggered T-FF

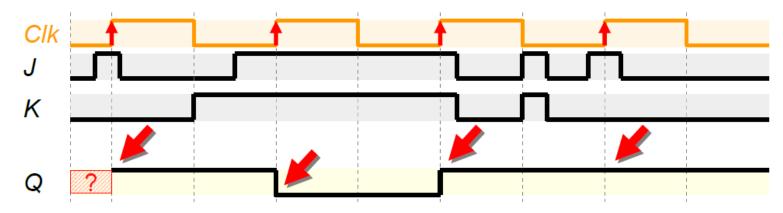


T	J	K	Q_{t+1}	STATE
0	0	0	Q_t	HOLD
1	1	1	$\overline{Q_t}$	TOGGLE

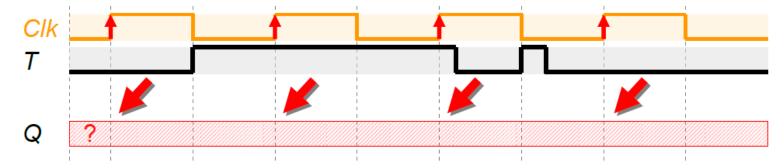
Inputs			Outputs	
T	CIk	Next Q Next Q'		State
Х	↓,0,1	Q	Q'	Hold
0	↑	Q	Q'	Hold
1	1	Q'	Q	Toggle

Timing Diagram

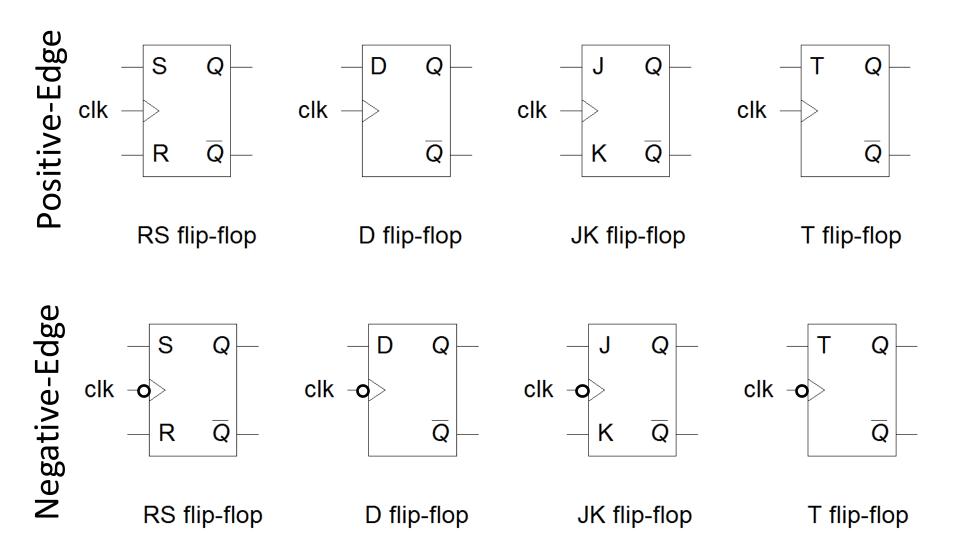
Timing Diagram of PET *JK* FF (if no output delay)



Timing Diagram of PET *T* FF (if no output delay)

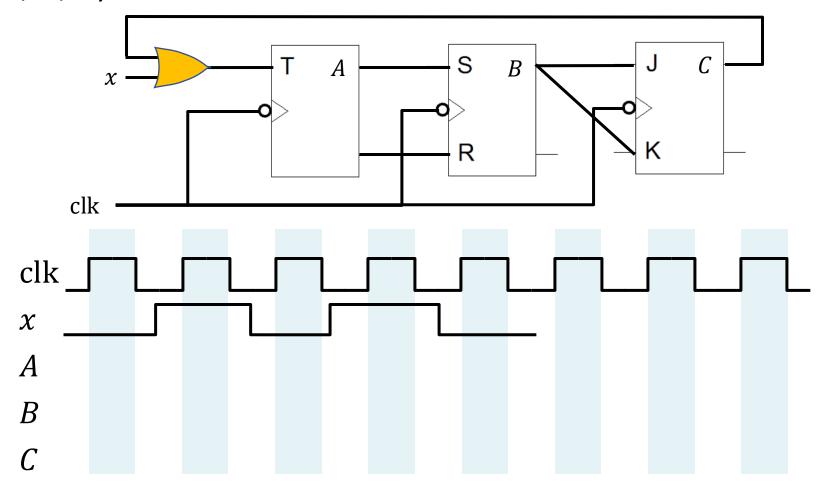


Logic Symbols



Exercise

Assume all FFs are initially 0 and no propagation delay, work out the timing diagram of the output of each FF (A, B, C).



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7.5 Synchronous & Asynchronous Inputs

- Inputs of flip-flops are categorized into two types
- Synchronous: Only affect the output in conjunction with a clock pulse (all inputs to FFs presented so far)
- Asynchronous: Inputs that can affect the output state independent of the clock pulse

Asynchronous Inputs

- The output will respond to the sync. inputs only if all async. inputs are inactive
- Particularly useful for bringing an FF into a desired initial state before normal clocked operation
- Preset (denoted by PRE) and Clear (denoted CLR) - Used to forcibly set and reset the state of flip-flop

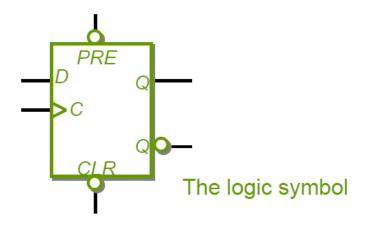
Example

The corresponding function table

	Inp	uts		Outputs			
PRE'	CLR'	CIk	D	Next Q	Next Q'	Meaning	
0	1	Х	X	1	0	Preset	
1	0	Х	Х	0	1	Clear	
0	0	Х	Х	1*	1*	Undefined	
1	1	1	X	D	D'	Set / Reset	
1	1	↓ ,0,1	Х	Q	Q'	Hold	

*Unpredictable behavior will result if both *PRE*' and *CLR*' set to 0 simultaneously

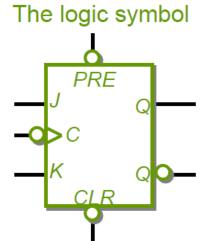
The same behaviors as *D* flip-flops



Example

The function table of an negative-edge-triggered *JK* flip-flop with asynchronous preset and clear inputs

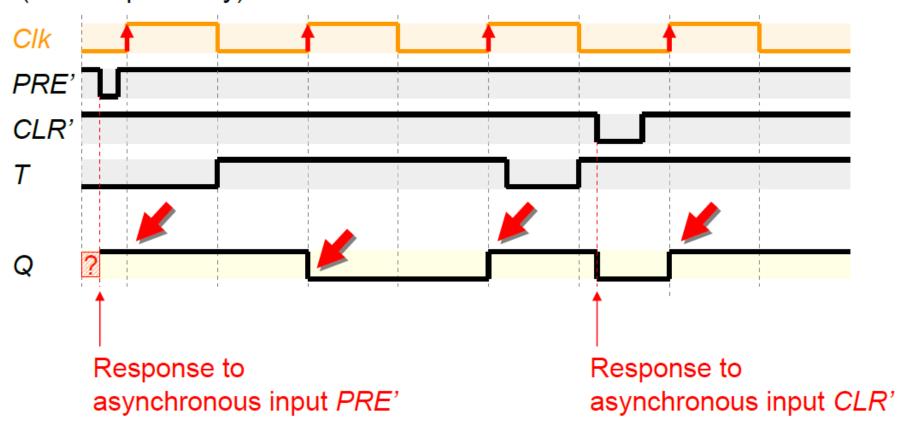
		Inputs		Outputs			
PRE'	CLR'	CIk	J	K	Next Q	Next Q'	Meaning
L	Н	Х	Х	X	Н	L	Preset
Н	L	Х	Х	Х	L	Н	Clear
L	L	Χ	Х	X	H*	H*	Undefined
Н	Н	\downarrow	L	L	Q	Q'	Hold
Н	Н	↓	Н	L	Н	L	Set
Н	Н	\downarrow	L	Н	L	Н	Reset
Н	Н	1	Н	Н	Q'	Q	Toggle
Н	Н	1,0,1	Х	Х	Q	Q'	Hold



The same behaviors as *JK* flip-flops

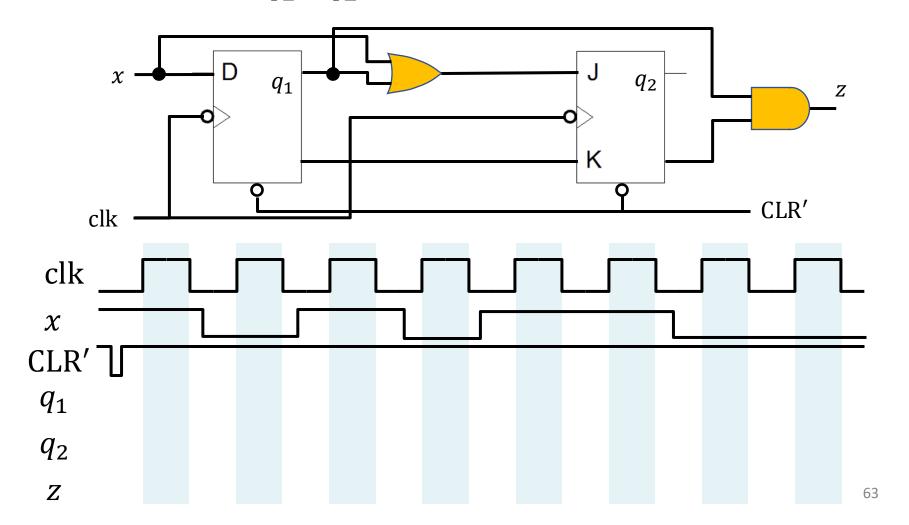
Timing Diagram

Timing Diagram of PET *T* FF with asynchronous preset and clear inputs (if no output delay)



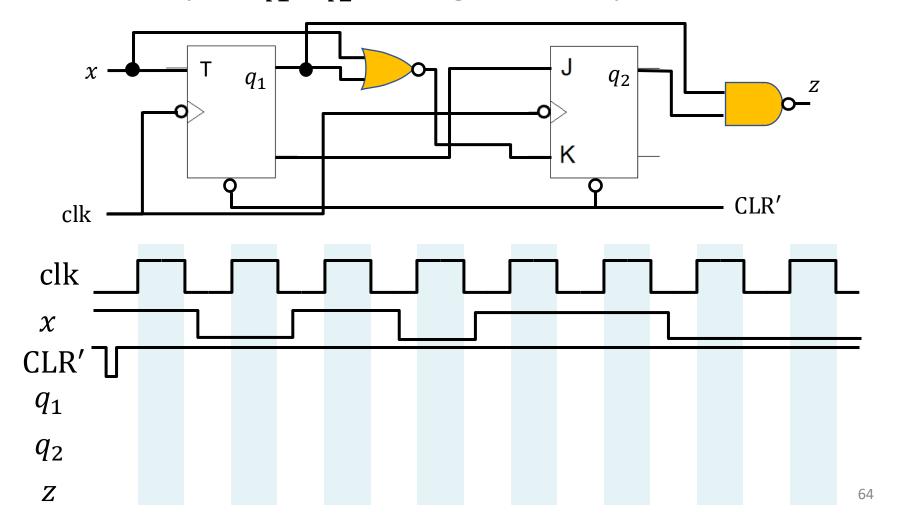
Exercise

Work out the timing diagram for the state of each flip-flop and the output $(q_1 \, , \, q_2 \, , \, z)$. (Ignore delay)



Exercise

Work out the timing diagram for the state of each flip-flop and the output $(q_1 \, , \, q_2 \, , \, z)$. (Ignore delay)



Summary

- ➤ Latches & Flip-Flops: Basic building block (Memory) of Sequential system
- > Flip-Flop (FF): One FF can store 1-bit data
 - Types: SR, D, JK, T, Master-Slave
 - Control inputs to change, set, reset or hold the value of Q
 - Triggering: Pulse vs Edge
 - Asynchronous inputs