Team notebook

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June 22, 2021

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2.1 Smallest Enclosing Circle

// Welzl's algorithm to find the smallest circle

```
// that encloses a group of poins in O(N * ITERS)
// returns {radius, x, y}
const int ITERS = 3e5;
const double INF = 1e12;
tuple<double, double, double> welzl(const vector<pair<int, int>>& points)
{
   double xt = 0, yt = 0;
   for(auto& [x, y] : points)
       xt += x;
       yt += y;
   }
   xt /= points.size();
   yt /= points.size();
   double p = 0.1;
   double mx_d;
   for(int i = 0; i < ITERS; ++i)</pre>
       mx_d = -INF;
       int mx_idx = -1;
       for(int j = 0; j < (int) points.size(); ++j)</pre>
           double cx = xt - points[j].first;
           double cy = yt - points[j].second;
           double cur = cx * cx + cy * cy;
           if(cur > mx d)
              mx_d = cur;
              mx_idx = j;
       }
       xt += (points[mx_idx].first - xt) * p;
       yt += (points[mx_idx].second - yt) * p;
       p *= 0.999;
   return {sqrt(mx_d), xt, yt};
```

3 Graphs

3.1 Articulation Point Bridge

```
// gr -> adj list
// vector vis, low -> initialize to -1
// int timer -> initialize to 0
void dfs(int pos, int dad = -1)
   vis[pos] = low[pos] = timer++;
   int kids = 0;
   for(auto& i : gr[pos])
       if(i == dad) continue;
       if(vis[i] >= 0)
          low[pos] = min(low[pos], vis[i]);
       else
       {
          dfs(i, pos);
          low[pos] = min(low[pos], low[i]);
          if(low[i] > vis[pos])
              is_bridge(pos, i)
          if(low[i] >= vis[pos] && dad >= 0)
              is_articulation_point(pos)
          ++kids;
   }
   if (dad == -1 && kids > 1)
       is_articulation_point(pos)
```

3.2 Dinic's Maximum Flow

```
// O(VE log(max_flow)) if scaling == 1
// O((V + E) sqrt(E)) if unit graph (turn scaling off)
// O((V + E) sqrt(V)) if bipartite matching (turn scaling off)
// indices are 0-based
const ll INF = 1e18;
struct Dinic
{
```

```
struct Edge
{
   int v;
   11 cap, flow;
   Edge(int _v, ll _cap) : v(_v), cap(_cap), flow(0) {}
};
int n;
11 lim;
vector<vector<int>> gr;
vector<Edge> e;
vector<int> idx, lv;
bool has_path(int s, int t)
   queue<int> q;
   q.push(s);
   lv.assign(n, -1);
   lv[s] = 0;
   while(!q.empty())
       int c = q.front();
       q.pop();
       if(c == t) break;
       for(auto& i : gr[c])
           ll cur_flow = e[i].cap - e[i].flow;
           if(lv[e[i].v] == -1 && cur_flow >= lim)
              lv[e[i].v] = lv[c] + 1;
              q.push(e[i].v);
          }
       }
   return lv[t] != -1;
}
ll get_flow(int s, int t, ll left)
   if(!left || s == t) return left;
   while(idx[s] < (int) gr[s].size())</pre>
       int i = gr[s][idx[s]];
       if(lv[e[i].v] == lv[s] + 1)
```

```
{
              11 add = get_flow(
                  e[i].v,
                  t,
                  min(left, e[i].cap - e[i].flow)
              if(add)
                  e[i].flow += add;
                  e[i ^ 1].flow -= add;
                  return add;
              }
          }
           ++idx[s];
       }
       return 0;
   }
   Dinic(int vertices, bool scaling = 1) // toggle scaling here
       : n(vertices), lim(scaling ? 1 << 30 : 1), gr(n) {}
   void add_edge(int from, int to, ll cap, bool directed = 1)
       gr[from].push_back(e.size());
       e.emplace_back(to, cap);
       gr[to].push_back(e.size());
       e.emplace_back(from, directed ? 0 : cap);
   }
   ll get_max_flow(int s, int t) // call this
       11 \text{ res} = 0;
       while(lim) // scaling
           while(has_path(s, t))
              idx.assign(n, 0);
              while(ll add = get_flow(s, t, INF)) res += add;
          lim >>= 1;
       }
       return res;
};
```

3

3.3 Edmonds' Blossom

```
// Maximum matching on general graphs in O(V^2 E)
// Indices are 1-based
// Stolen from ko_osaga's cheatsheet
struct Blossom
{
   vector<int> vis, dad, orig, match, aux;
   vector<vector<int>> conn;
   int t, N;
   queue<int> Q;
   void augment(int u, int v)
       int pv = v;
       do
           pv = dad[v];
           int nv = match[pv];
           match[v] = pv;
           match[pv] = v;
           v = nv;
       } while(u != pv);
   }
   int lca(int v, int w)
       ++t:
       while(true)
           if(v)
              if(aux[v] == t) return v;
              aux[v] = t;
              v = orig[dad[match[v]]];
           }
           swap(v, w);
       }
   }
```

```
void blossom(int v, int w, int a)
   while(orig[v] != a)
       dad[v] = w;
       w = match[v];
       if(vis[w] == 1)
          Q.push(w);
          vis[w] = 0;
       orig[v] = orig[w] = a;
       v = dad[w];
}
bool bfs(int u)
   fill(vis.begin(), vis.end(), -1);
   iota(orig.begin(), orig.end(), 0);
   Q = queue<int>();
   Q.push(u);
   vis[u] = 0;
   while(!Q.empty())
       int v = Q.front(); Q.pop();
       for(int x : conn[v])
          if(vis[x] == -1)
              dad[x] = v; vis[x] = 1;
              if(!match[x])
                  augment(u, x);
                  return 1;
              Q.push(match[x]);
              vis[match[x]] = 0;
          else if(vis[x] == 0 && orig[v] != orig[x])
              int a = lca(orig[v], orig[x]);
              blossom(x, v, a);
              blossom(v, x, a);
```

```
}
       }
   }
   return false;
Blossom(int n) : // n = vertices
   vis(n + 1), dad(n + 1), orig(n + 1), match(n + 1),
   aux(n + 1), conn(n + 1), t(0), N(n)
   for(int i = 0; i <= n; ++i)</pre>
       conn[i].clear();
       match[i] = aux[i] = dad[i] = 0;
   }
}
void add_edge(int u, int v)
   conn[u].push_back(v);
   conn[v].push_back(u);
}
int solve() // call this for answer
   int ans = 0;
   vector<int> V(N - 1);
   iota(V.begin(), V.end(), 1);
   shuffle(V.begin(), V.end(), mt19937(0x94949));
   for(auto x : V)
       if(!match[x])
           for(auto y : conn[x])
              if(!match[y])
                  match[x] = y, match[y] = x;
                  ++ans;
                  break;
           }
       }
   }
```

```
for(int i = 1; i <= N; ++i)
{
      if(!match[i] && bfs(i)) ++ans;
}
    return ans;
}
};</pre>
```

3.4 Eulerian Path or Cycle

```
// finds a eulerian path / cycle
// visits each edge only once
// properties:
// - cycle: degrees are even
// - path: degrees are even OR degrees are even except for 2 vertices
// how to use: g = adjacency list g[n] = connected to n, undirected
// if there is a vertex u with an odd degree, call dfs(u)
// else call on any vertex
// ans = path result
vector<set<int>> g;
vector<int> ans;
void dfs(int u)
   while(g[u].size())
       int v = *g[u].begin();
       g[u].erase(v);
       g[v].erase(u);
       dfs(v);
   ans.push_back(u);
```

5

4 Math

4.1 Euler's Totient

```
// Precompute up to n in O(n log log n)
vector<int> phi_1_to_n(int n)
    vector<int> phi(n + 1);
    phi[0] = 0;
    phi[1] = 1;
    for(int i = 2; i <= n; i++)</pre>
       phi[i] = i;
    for(int i = 2; i <= n; i++)</pre>
       if(phi[i] == i)
           for(int j = i; j <= n; j += i)</pre>
               phi[j] -= phi[j] / i;
    return phi;
}
// Calculate for a single n in O(sqrt(n))
ll totient(ll n)
    11 \text{ res} = 1:
    for(ll i = 2; i * i <= n; ++i)</pre>
        if(n \% i == 0)
           res *= i - 1;
           n /= i;
        while(n \% i == 0)
           res *= i:
           n /= i;
        }
    }
    if(n > 1) res *= n - 1;
    return res;
}
```

4.2 Extended Euclidean GCD

```
// computes x and y such that ax + by = gcd(a, b) in O(log (min(a, b)))
// returns {gcd(a, b), x, y}
tuple<int, int, int> gcd(int a, int b)
{
```

```
if(b == 0) return {a, 1, 0};
auto [d, x1, y1] = gcd(b, a % b);
return {d, y1, x1 - y1 * (a / b)};
}
```

4.3 Fibonacci Check

4.4 Matrix Multiplication

4.5 Miller-Rabin Pollard's Rho

```
namespace MillerRabin
{
   const vector<ll> primes = { // deterministic up to 2^64 - 1
      2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37
   };
   ll gcd(ll a, ll b)
```

}

```
return b ? gcd(b, a % b) : a;
   ll powa(ll x, ll y, ll p) // (x ^ y) % p
       if(!y) return 1;
       if(y & 1) return ((__int128) x * powa(x, y - 1, p)) % p;
       ll temp = powa(x, y >> 1, p);
       return ((__int128) temp * temp) % p;
   bool miller_rabin(ll n, ll a, ll d, int s)
      11 x = powa(a, d, n);
       if(x == 1 || x == n - 1) return 0;
       for(int i = 0; i < s; ++i)</pre>
       {
          x = ((_int128) x * x) % n;
          if(x == n - 1) return 0;
       }
       return 1;
   bool is_prime(ll x) // use this
       if(x < 2) return 0;
       int r = 0;
      11 d = x - 1;
       while((d \& 1) == 0)
          d >>= 1:
          ++r;
       for(auto& i : primes)
          if(x == i) return 1;
          if(miller_rabin(x, i, d, r)) return 0;
       }
       return 1;
   }
namespace PollardRho
   mt19937_64 generator(chrono::steady_clock::now()
                       .time_since_epoch().count());
```

```
uniform_int_distribution<1l> rand_ll(0, LLONG_MAX);
   11 f(11 x, 11 b, 11 n) // (x^2 + b) % n
       return (((__int128) x * x) % n + b) % n;
   ll rho(ll n)
   {
       if(n % 2 == 0) return 2;
       11 b = rand_ll(generator);
       11 x = rand_ll(generator);
       11 y = x;
       while(1)
          x = f(x, b, n);
          y = f(f(y, b, n), b, n);
          11 d = MillerRabin::gcd(abs(x - y), n);
          if(d != 1) return d;
   }
   void pollard_rho(ll n, vector<ll>& res)
       if(n == 1) return;
       if(MillerRabin::is_prime(n))
          res.push_back(n);
          return:
       }
       11 d = rho(n);
       pollard_rho(d, res);
       pollard_rho(n / d, res);
   vector<ll> factorize(ll n, bool sorted = 1) // use this
       vector<ll> res;
       pollard_rho(n, res);
       if(sorted) sort(res.begin(), res.end());
       return res;
   }
}
```

5 Miscellaneous

5.1 Dates

5.1.1 Day of Date

```
// O-based
const vector<int> T = {
    0, 3, 2, 5, 0, 3,
    5, 1, 4, 6, 2, 4
}
int day(int d, int m, int y)
{
    y -= (m < 3);
    return (y + y / 4 - y / 100 + y / 400 + T[m - 1] + d) % 7;
}</pre>
```

5.1.2 Number of Days since 1-1-1

5.2 Enumerate Subsets of a Bitmask

```
int x = 0;
do
{
    // do stuff with the bitmask here
    x = (x + 1 + ~m) & m;
} while(x != 0);
```

5.3 Int to Roman

```
const string R[] = {
    "M", "CM", "D", "CD", "C", "XC", "L",
    "XL", "X", "IX", "V", "IV", "I"
};

const int N[] = {
    1000, 900, 500, 400, 100, 90,
    50, 40, 10, 9, 5, 4, 1
};

string to_roman(int x)
{
    if (x == 0) return "0"; // Not decimal 0!
    string res = "";
    for (int i = 0; i < 13; ++i)
        while (x >= N[i]) x -= N[i], res += R[i];
    return res;
}
```

5.4 Josephus Problem

```
11 josephus(11 n, 11 k) // O(k log n)
{
    if(n == 1) return 0;
    if(k == 1) return n - 1;
    if(k > n) return (josephus(n - 1, k) + k) % n;
    ll cnt = n / k;
    ll res = josephus(n - cnt, k);
    res -= n % k;
    if(res < 0) res += n;
    else res += res / (k - 1);
    return res;
}

int josephus(int n, int k) // O(n)
{
    int res = 0;
    for(int i = 1; i <= n; ++i)
        res = (res + k) % i;
    return res + 1;
}</pre>
```

6 Strings

6.1 Knuth-Morris-Pratt

```
// Constructs KMP failure function in O(n)
vector<int> kmp(const string& s)
{
    vector<int> res(s.length());
    int i = 1, j = 0;
    while(i < (int) s.length())
    {
        if(s[i] == s[j]) res[i++] = ++j;
        else if(j > 0) j = res[j - 1];
        else res[i++] = 0;
    }
    return res;
}
```

6.2 Suffix Array

```
// stores result in sa and lcp
// if lcp is needed, call SuffixArray(str, 1)
struct SuffixArray
{
   int n;
   vector<int> sa, lcp, rnk, cnt;
   vector<pair<int, int>> p;
   SuffixArray(const string& s, bool calc_lcp = 0) :
       n(s.length()), sa(n), lcp(calc_lcp ? n : 0), rnk(n),
       cnt(max(n, 256)), p(n)
       for(int i = 0; i < n; ++i) rnk[i] = s[i];</pre>
       iota(sa.begin(), sa.end(), 0);
       for(int i = 1; i < n; i <<= 1) update_sa(i);</pre>
       if(!calc_lcp) return;
       vector<int> phi(n), plcp(n);
       phi[sa[0]] = -1;
       for(int i = 1; i < n; ++i) phi[sa[i]] = sa[i - 1];</pre>
       int 1 = 0;
       for(int i = 0; i < n; ++i)
       {
```

```
if(phi[i] == -1) plcp[i] = 0;
       else
       Ł
           while ((i + 1 < n) \&\& (phi[i] + 1 < n)
                 && (s[i + 1] == s[phi[i] + 1])) ++1;
           plcp[i] = 1;
           1 = \max(1 - 1, 0);
   }
   for(int i = 0; i < n; ++i) lcp[i] = plcp[sa[i]];</pre>
void update_sa(int len)
{
   sort_sa(len); sort_sa(0);
   for(int i = 0; i < n; ++i) p[i] = {rnk[i], rnk[(i + len) % n]};</pre>
   auto lst = p[sa[0]];
   rnk[sa[0]] = 0;
   int cur = 0;
   for(int i = 1; i < n; ++i)</pre>
       if(lst != p[sa[i]])
           lst = p[sa[i]];
           ++cur;
       rnk[sa[i]] = cur;
}
void sort sa(int offset)
   fill(cnt.begin(), cnt.end(), 0);
   for(int i = 0; i < n; ++i) ++cnt[rnk[(i + offset) % n]];</pre>
   int sum = 0;
   for(int i = 0; i < (int) cnt.size(); ++i)</pre>
       int temp = cnt[i];
       cnt[i] = sum;
       sum += temp;
   vector<int> temp(n);
   for(int i = 0; i < n; ++i)
       int cur = cnt[rnk[(sa[i] + offset) % n]]++;
       temp[cur] = sa[i];
```

```
}
sa = move(temp);
}
```