Contents		4 Geometry	6	6.5 Modular Linear Equation	18
		4.1 Geometry Template	6	6.6 Miller-Rabin and Pollard's Rho	18
1 Miscellaneous	1	4.2 Convex Hull		6.7 Berlekamp-Massey	18
1.1 VS Code Config	1	4.3 Closest Pair of Points			19
1.2 Day of Date	1	4.4 Smallest Enclosing Circle	11	6.9 Number Theoretic Transform	20
1.3 Number of Days since 1-1-1	1	4.5 Sutherland-Hodgman Algorithm		6.10 Gauss-Jordan	20
1.4 Enumerate Subsets of a Bitmask	2	4.6 Centroid of Polygon		6.11 Fibonacci Check	21
1.5 Fast IO	2	4.7 Pick Theorem		6.12 Derangement	21
1.6 Int to Roman	2				21
1.7 Josephus Problem	2	5 Graphs	13	6.14 Forbenius Number	21
1.8 Random Primes	2	5.1 Articulation Point and Bridge	13	6.15 Stars and Bars with Upper Bound	21
1.9 RNG	2	5.2 SCC and Strong Orientation			
2 Data Standard		5.3 Centroid Decomposition		7 Strings	21
2 Data Structures	2	5.4 Dinic's Maximum Flow		7.1 Aho-Corasick	21
2.1 2D Segment Tree	2	5.5 Minimum Cost Maximum Flow	14	7.2 Eertree	22
2.2 Fenwick RU-RQ	ა ე	5.6 Flows with Demands	14	7.3 Manacher's Algorithm	22
2.3 Heavy-Light Decomposition	3 4	5.7 Hungarian	15	7.4 Suffix Array	22
2.5 Persistent Segment Tree	4	5.8 Edmonds' Blossom			23
2.6 STL PBDS	5	5.9 Eulerian Path or Cycle	16		
2.7 Treap	5	5.10 Hierholzer's Algorithm	16	8 OEIS	24
2.8 Unordered Map Custom Hash	5	5.11 2-SAT		8.1 A000108 (Catalan)	24
2.9 Mo's on Tree	6				24
2.0 NO 5 OH TICC	Ü	6 Math	17	8.3 A000668 (Mersene Primes)	24
3 Dynamic Programming	6	6.1 Extended Euclidean GCD	17	8.4 A001434	24
3.1 DP Convex Hull	6	6.2 Generalized CRT	17	8.5 A018819	24
3.2 DP DNC	6	6.3 Generalized Lucas Theorem	17	8.6 A092098	24
3.3 DP Knuth-Yao	6	6.4 Linear Diophantine	17	8.7 A277402	25
		•			

1 Miscellaneous

1.1 VS Code Config

```
"command": "clear; ./runner ${file} ${fileBasenameNoExtension} 0"
"command": "clear; ./runner ${file} ${fileBasenameNoExtension} 1"
"command": "clear; rm -r .cache"
runner:
#!/bin/bash
cp=$1; uf=$3; tmp='.cache'; bp="$tmp/$2"; ch="$tmp/${2}-hash"
ti="${li}\nTime: %es"
rc() {
   mkdir -p "$tmp"
sch() {
    rc; echo -n "$(gh)" > "$ch"
gch() {
   if [[ -f "$ch" ]]; then cat $ch
   else echo -n 'NULL'; fi
gh() {
   sha256sum $cp
   oh=$(gch); nh=$(gh)
   if [[ "$oh" == "$nh" ]]; then return 1
   else return 0; fi
cc() {
   g++ $cp -02 -std=gnu++17 -Wall -Wextra -Wshadow -DLOCAL -o $bp
```

```
ma() {
    echo $li
    if nr; then
        if [[ -f "$bp" ]]; then rm "$bp"; fi
        sch; echo 'Compiling...'; echo $li; cc; echo $li
    if [[ -f "$bp" ]]; then
        echo 'Running...'; echo $li
        if (( $uf )); then command time -f "$ti" ./$bp < IN
        else command time -f "$ti" ./$bp; fi
        echo $li
    fi
}
ma</pre>
```

1.2 Day of Date

```
// 0-based
const vector<int> T = {
  0, 3, 2, 5, 0, 3,
  5, 1, 4, 6, 2, 4
}
int day(int d, int m, int y) {
  y -= (m < 3);
  return (y + y / 4 - y / 100 + y / 400 + T[m - 1] + d) % 7;
}</pre>
```

1.3 Number of Days since 1-1-1

```
int rdn(int d, int m, int y) {
```

```
if(m < 3)
  --y, m += 12;
return 365 * y + y / 4 - y / 100 + y / 400
       + (153 * m - 457) / 5 + d - 306;
```

1.4 Enumerate Subsets of a Bitmask

```
int x = 0;
do {
 // do stuff with the bitmask here
x = (x + 1 + \sim m) \& m;
} while(x != 0);
```

1.5 Fast IO

int read() {

```
char c;
 do {
   c = getchar_unlocked();
 } while(c < 33);</pre>
 int res = 0;
 int mul = 1;
 if(c == '-') {
   mul = -1;
   c = getchar_unlocked();
 while('0' <= c && c <= '9') {
   res = res * 10 + c - '0';
   c = getchar_unlocked();
 return res * mul;
void write(int x) {
 static char wbuf[10];
 if(x < 0) {
   putchar_unlocked('-');
   x = -x;
 int idx = 0;
 while(x) {
   wbuf[idx++] = x % 10;
   x /= 10;
 if(idx == 0)
   putchar_unlocked('0');
 for(int i = idx - 1; i >= 0; --i)
   putchar_unlocked(wbuf[i] + '0');
void write(const char* s) {
 while(*s) {
   putchar_unlocked(*s);
   ++s;
```

1.6 Int to Roman

```
const string R[] = {
      "M", "CM", "D", "CD", "C", "XC", "L", "XL", "X", "IX", "V", "IV", "I"
Page
\sim
    const int N[] = {
      1000, 900, 500, 400, 100, 90,
      50, 40, 10, 9, 5, 4, 1
```

```
};
string to_roman(int x) {
 if(x == 0) {
    return "0"; // Not decimal 0!
  string res = "";
  for(int i = 0; i < 13; ++i)
   while(x >= N[i])
     x -= N[i], res += R[i];
  return res;
```

1.7 Josephus Problem

```
ll josephus(ll n, ll k) { // O(k log n)
 if(n == 1)
   return 0;
 if(k == 1)
   return n - 1;
 if(k > n)
   return (josephus(n - 1, k) + k) % n;
 ll\ cnt = n / k;
 ll res = josephus(n - cnt, k);
 res -= n % k;
 if(res < 0)
   res += n;
 else
   res += res / (k - 1);
 return res;
int josephus(int n, int k) { // O(n)
 int res = 0;
 for(int i = 1; i <= n; ++i)
   res = (res + k) % i;
 return res + 1;
```

1.8 Random Primes

36671 74101 724729 825827 924997 1500005681 2010408371 2010405347

1.9 RNG

```
// RNG - rand_int(min, max), inclusive
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
template<class T>
T rand_int(T mn, T mx) {
 return uniform_int_distribution<T>(mn, mx)(rng);
```

2 Data Structures

2.1 2D Segment Tree

```
struct Segtree2D {
 struct Segtree {
   struct node {
     int l, r, val;
     node* lc, *rc;
     node(int _l, int _r, int _val = INF) : l(_l), r(_r), val(_val),
       lc(NULL), rc(NULL) {}
```

```
typedef node* pnode;
  pnode root;
  Segtree(int l, int r) {
    root = new node(l, r);
  void update(pnode& nw, int x, int val) {
    int l = nw - > l, r = nw - > r, mid = (l + r) / 2;
    if(l == r)
      nw->val = val;
    else {
      assert(l <= x && x <= r);
      pnode& child = x <= mid ? nw->lc : nw->rc;
        child = new node(x, x, val);
      else if(child->l <= x && x <= child->r)
        update(child, x, val);
      else {
        do {
          if(x \le mid)
            r = mid;
          else
            l = mid + 1;
          mid = (l + r) / 2;
        } while((x <= mid) == (child->l <= mid));</pre>
        pnode nxt = new node(l, r);
        if(child->l <= mid)</pre>
          nxt->lc = child;
        else
          nxt->rc = child;
        child = nxt;
        update(nxt, x, val);
      nw->val = min(nw->lc ? nw->lc->val : INF,
                    nw->rc ? nw->rc->val : INF);
  int query(pnode& nw, int x1, int x2) {
    if(!nw)
      return INF;
    int& l = nw->l, &r = nw->r;
    if(r < x1 || x2 < l)
      return INF;
    if(x1 <= l && r <= x2)
      return nw->val;
    int ret = min(query(nw->lc, x1, x2),
                  query(nw->rc, x1, x2));
    return ret;
  void update(int x, int val) {
    assert(root->l <= x && x <= root->r);
    update(root, x, val);
  int query(int l, int r) {
    return query(root, l, r);
};
struct node {
  int l, r;
  Segtree y;
  node* lc, *rc;
  node(int _l, int _r) : l(_l), r(_r), y(0, MAX),
    lc(NULL), rc(NULL) {}
typedef node* pnode;
```

```
pnode root;
  Segtree2D(int l, int r) {
   root = new node(l, r);
  void update(pnode& nw, int x, int y, int val) {
    int& l = nw->l, &r = nw->r, mid = (l + r) / 2;
   if(l == r)
     nw->y.update(y, val);
    else {
     if(x <= mid) {
        if(!nw->lc)
          nw->lc = new node(l, mid);
       update(nw->lc, x, y, val);
     } else {
       if(!nw->rc)
         nw->rc = new node(mid + 1, r);
       update(nw->rc, x, y, val);
     val = min(nw -> lc ? nw -> lc -> y.query(y, y) : INF,
                nw->rc ? nw->rc->y.query(y, y) : INF);
      nw->y.update(y, val);
  int query(pnode& nw, int x1, int x2, int y1, int y2) {
   if(!nw)
     return INF;
    int& l = nw->l, &r = nw->r;
   if(r < x1 || x2 < l)
     return INF:
   if(x1 <= l && r <= x2)
     return nw->y.query(y1, y2);
   int ret = min(query(nw->lc, x1, x2, y1, y2),
                  query(nw->rc, x1, x2, y1, y2));
   return ret;
  void update(int x, int y, int val) {
   assert(root->l <= x && x <= root->r);
   update(root, x, y, val);
  int query(int x1, int x2, int y1, int y2) {
    return query(root, x1, x2, y1, y2);
};
```

2.2 Fenwick RU-RQ

```
void updtRL(int l, int r, ll val) {
 updt(BIT1, l, val), updt(BIT1, r + 1, -val);
 updt(BIT2, l, val * (l - 1)), updt(BIT2, r + 1, -val * r);
ll query(int k) {
 return que(BIT1, k) * k - que(BIT2, k);
```

2.3 Heavy-Light Decomposition

```
struct HLD {
  vector<int> id, size, idx, up, root, st;
  vector<vector<int>> adj, chain;
  SegTree seg;
  HLD(const vector<vector<int>>& edges) :
```

```
decompose(0, -1);
    int cnt = 0;
    st.resize(chain.size());
    for(int i = 0; i < (int) chain.size(); ++i) {</pre>
      st[i] = cnt;
      cnt += chain[i].size();
  void precompute(int pos, int dad) {
    size[pos] = 1;
    up[pos] = dad;
    for(auto& i : adj[pos]) {
      if(i != dad) {
        precompute(i, pos);
        size[pos] += size[i];
   }
  void decompose(int pos, int dad) {
    if(id[pos] == -1) {
      id[pos] = chain.size();
      root.push_back(pos);
      chain.emplace_back();
    idx[pos] = chain[id[pos]].size();
    chain[id[pos]].push_back(pos);
    int mx = 0, heavy = -1;
    for(auto& i : adj[pos]) {
      if(i != dad && size[i] > mx) {
        mx = size[i];
        heavy = i;
    if(heavy != -1)
      id[heavy] = id[pos];
    for(auto& i : adj[pos]) {
      if(i != dad)
        decompose(i, pos);
 }
  void update(int ch, int l, int r, int val) {
    seg.update(st[ch] + l, st[ch] + r, val);
  int query(int ch, int l, int r, int val) {
    return seg.query(st[ch] + l, st[ch] + r, val);
};
// how to move from u to v
while(1) {
 if(hld.id[u] == hld.id[v]) {
    if(hld.idx[u] > hld.idx[v])
      swap(u, v);
    hld.update(hld.id[u], hld.idx[u], hld.idx[v], w);
    // or hld.query(hld.id[u], hld.idx[u], hld.idx[v]);
    break;
  if(hld.id[u] < hld.id[v])</pre>
    swap(u, v);
  hld.update(hld.id[u], 0, hld.idx[u], w);
 // or hld.query(hld.id[u], 0, hld.idx[u]);
 u = hld.up[hld.root[hld.id[u]]];
```

n(edges.size()), id(n, -1), size(n, -1), idx(n, -1),

up(n, -1), adj(edges), seg(n) {

precompute(0, -1);

2.4 Li Chao Tree

```
typedef long long int TD;
const TD INF = 100000000000000;
namespace LICHAO {
struct Node {
 TD m, c;
 Node* l, *r;
};
Node* newNode(Node* x = NULL) {
 Node* ret = (Node*)malloc(sizeof(Node));
 if(x)
   ret->m = x->m, ret->c = x->c;
 ret->l = ret->r = NULL;
 return ret;
void update(Node* k, TD l, TD r, TD m, TD c) {
 TD mid = l + r >> 1;
 bool le = m * l + c < k-> m * l + k->c;
 bool ri = m * mid + c < k->m * mid + k->c;
 if(ri)
   swap(k->m, m), swap(k->c, c);
 if(r - l <= 1)
   return;
  else if(le != ri)
   update((k->1) ? (k->1) : (k->1 = newNode(k)), l, mid, m, c);
   update((k->r) ? (k->r) : (k->r = newNode(k)), mid, r, m, c);
TD query(Node* k, TD l, TD r, TD p) {
 if(!k)
   return INF;
 if(r - l <= 1)
   return p * k->m + k->c;
 if(p < (l + r >> 1))
   return min(p * k->m + k->c, query(k->l, l, l + r >> 1, p));
   return min(p * k->m + k->c, query(k->r, l + r >> 1, r, p));
```

2.5 Persistent Segment Tree

```
class PersistentSegtree {
private:
 int n, ptr, sz;
 struct P {
   int val = 0, l, r;
 vector<P> node;
 vector<int> root;
  int newNode() {
   return ptr++;
 int copyNode(int idx) {
   node[ptr] = node[idx];
   return ptr++;
  int build(int l, int r) {
    int idx = newNode();
   if(l == r)
     return idx;
   node[idx].l = build(l, (l + r) / 2);
   node[idx].r = build((l + r) / 2 + 1, r);
   return idx;
  int update(int idx, int l, int r, int x, int val) {
    idx = copyNode(idx);
```

```
if(l == r) {
     node[idx].val += val;
     return idx;
   int mid = (l + r) / 2;
   if(x \le mid)
     node[idx].l = update(node[idx].l, l, mid, x, val);
     node[idx].r = update(node[idx].r, mid + 1, r, x, val);
    node[idx].val = node[node[idx].l].val + node[node[idx].r].val;
   return idx;
  int query(int idxl, int idxr, int l, int r, int x, int y) {
   if(y < l \mid | r < x)
     return 0;
   if(x <= l && r <= y)
     return node[idxr].val - node[idxl].val;
   int mid = (l + r) / 2;
   return query(node[idxl].l, node[idxr].l, l, mid, x, y)
          + query(node[idxl].r, node[idxr].r, mid + 1, r, x, y);
public:
 PersistentSegtree(int _n) : n(_n), ptr(0) {
   sz = 30 * n;
   node.resize(sz);
   root.push back(build(1, n));
 void update(int x, int val) {
   root.push_back(update(root.back(), 1, n, x, val));
 int query(int l, int r, int x, int y) {
   return query(root[l - 1], root[r], 1, n, x, y);
};
```

2.6 STL PBDS

2.7 Treap

```
// Complexity: O(log N) for split and merge
//
// empty treap: Treap* tr = nullptr;
// insert v at x: [l, r] = split(tr, x), m = Treap(v), merge lmr
// delete at x: [l, r] = split(tr, x), [m, r] = split(r, 1), merge lr
// lazy prop: propagate every time a node is accessed
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
using Key = int;
struct Treap {
   Key val;
   Treap* left;
   Treap* right;
```

```
int prio, sz;
 Treap() {}
 Treap(int _val);
int size(Treap* tr) {
 return tr ? tr->sz : 0;
void update(Treap* tr) {
 tr->sz = 1 + size(tr->left) + size(tr->right);
Treap::Treap(Key _val) :
 val(_val), left(nullptr), right(nullptr), prio(rng()) {
 update(this);
pair<Treap*, Treap*> split(Treap* tr, int sz) {
 if(!tr) return {nullptr, nullptr};
 int left sz = size(tr->left);
 if(sz <= left_sz) {</pre>
   auto [left, mid] = split(tr->left, sz);
   tr->left = mid;
   update(tr);
   return {left, tr};
 } else {
   auto [mid, right] = split(tr->right, sz - left_sz - 1);
   tr->right = mid;
   update(tr);
   return {tr, right};
Treap* merge(Treap* l, Treap* r) {
 if(!l)
   return r;
  if(!r)
   return l;
  if(l->prio < r->prio) {
   l->right = merge(l->right, r);
   update(l);
   return l;
 } else {
   r->left = merge(l, r->left);
   update(r);
   return r;
```

2.8 Unordered Map Custom Hash

2.9 Mo's on Tree

```
ST(u) \leq ST(v)
P = LCA(u, v)
If P = u, query [ST(u), ST(v)]
Else query [EN(u), ST(v)] + [ST(P), ST(P)]
```

3 Dynamic Programming

3.1 DP Convex Hull

```
/* dp[i] = min k < i \{dp[k] + x[i] * m[k]\}
   Make sure gradient (m[i]) is either non-increasing if min,
   or non-decreasing if max. x[i] must be non-decreasing. just sort */
// while this is true, pop back from dg. a=new line, b=last, c=2nd last
bool cekx(int a, int b, int c) {
 // if not enough, change to cross mul
  // if cross mul, beware of negative denominator, and overflow
  return (double)(y[b] - y[a]) / (m[a] - m[b]) \le (double)(y[c] - y[b]) /
         (m[b] - m[c]);
```

3.2 DP DNC

```
void f(int rem, int l, int r, int optl, int optr) {
 if(l > r)
  return;
 int mid = l + r >> 1;
 int opt = MOD, optid = mid;
 for(int i = optl; i <= mid && i <= optr; ++i) {</pre>
   if(dp[rem - 1][i] + c[i][mid] < opt) {</pre>
     opt = dp[rem - 1][i] + c[i][mid];
     optid = i;
 dp[rem][mid] = opt;
 f(rem, l, mid - 1, optl, optid);
 f(rem, mid + 1, r, optid, optr);
 return:
rep(i, 1, n)dp[1][i] = c[0][i];
rep(i, 2, k)f(i, i, n, i, n);
```

3.3 DP Knuth-Yao

```
// opt[i+1][j] <= opt[i][j] <= opt[i][j+1]
// dp[i][j] = min{k} dp[i][k]+dp[k][j]+cost[i][j]
for(int k = 0; k \le n; k++) {
 for(int i = 0; i + k <= n; i++) {
   if(k < 2)
      dp[i][i + k] = 0, opt[i][i + k] = i;
      int sta = opt[i][i + k - 1];
      int end = opt[i + 1][i + k];
      for(int j = sta; j <= end; j++) {</pre>
        if(dp[i][j] + dp[j][i + k] + cost[i][i + k] < dp[i][i + k]) {
          dp[i][i + k] = dp[i][j] + dp[j][i + k] + cost[i][i + k];
          opt[i][i + k] = j;
```

4 Geometry

4.1 Geometry Template

```
Don't Forgor the Proof Peko
Bina Nusantara University
TABLE OF CONTENT
0. Basic Rule
    0.1. Everything is in double
    0.2. Every comparison use EPS
    0.3. Every degree in rad
1. General Double Operation
    1.1. const double EPS=1E-9
    1.2. const double PI=acos(-1.0)
    1.3. const double INFD=1E9
    1.3. between d(double x,double l,double r)
        check whether x is between l and r inclusive with EPS
    1.4. same_d(double x,double y)
        check whether x=y with EPS
    1.5. dabs(double x)
        absolute value of x
2. Point
    2.1. struct point
        2.1.1. double x,y
            cartesian coordinate of the point
        2.1.2. point()
            default constructor
        2.1.3. point(double _x,double _y)
            constructor, set the point to (_x,_y)
        2.1.4. bool operator< (point other)
            regular pair <double, double > operator < with EPS
        2.1.5. bool operator == (point other)
            regular pair < double, double > operator == with EPS
    2.2. hypot(point P)
        length of hypotenuse of point P to (0,0)
    2.3. e_dist(point P1,point P2)
        euclidean distance from P1 to P2
    2.4. m_dist(point P1,point P2)
        manhattan distance from P1 to P2
    2.5. point rotate(point P,point O,double angle)
        rotate point P from the origin O by angle ccw
Vector
    3.1. struct vec
        3.1.1. double x,y
            x and y magnitude of the vector
        3.1.2. vec()
            default constructor
        3.1.3. vec(double _x,double _y)
            constructor, set the vector to (_x,_y)
        3.1.4. vec(point A, point B)
            constructor, set the vector to vector AB (A->B)
/*General Double Operation*/
const double PI = acos(-1.0);
const double INFD = 1E9;
double between_d(double x, double l, double r) {
 return (min(l, r) \le x + EPS \&\& x \le max(l, r) + EPS);
double same_d(double x, double y) {
  return between_d(x, y, y);
double dabs(double x) {
  if(x < EPS)
    return -x;
  return x;
/*Point*/
struct point {
  double x, y;
  point() {
```

```
return translate(point(P.x * cos(angle) - P.y * sin(angle),
                         P.x * sin(angle) + P.y * cos(angle)), v);
point mid(point P, point Q) {
 return point((P.x + Q.x) / 2, (P.y + Q.y) / 2);
double angle(point A, point O, point B) {
 vec OA(0, A), OB(0, B);
 return acos(dot(OA, OB) / sqrt(norm_sq(OA) * norm_sq(OB)));
int orientation(point P, point Q, point R) {
 vec PQ(P, Q), PR(P, R);
 double c = cross(PQ, PR);
 if(c < -EPS)
   return -1;
 if(c > EPS)
   return 1;
 return 0;
/*Line*/
struct line {
 double a, b, c;
 line() {
   a = b = c = 0.0;
 line(double _a, double _b, double _c) {
   a = _a;
   b = b;
   c = _c;
  line(point P1, point P2) {
   if(P1 < P2)
     swap(P1, P2);
   if(same_d(P1.x, P2.x))
     a = 1.0, b = 0.0, c = -P1.x;
   else
     a = -(P1.y - P2.y) / (P1.x - P2.x), b = 1.0, c = -(a * P1.x) - P1.y;
 line(point P, double slope) {
   if(same_d(slope, INFD))
     a = 1.0, b = 0.0, c = -P.x;
   else
     a = -slope, b = 1.0, c = -(a * P.x) - P.y;
 bool operator== (line other) {
   return same_d(a, other.a) && same_d(b, other.b) && same_d(c, other.c);
 double slope() {
   if(same_d(b, 0.0))
     return INFD;
   return -(a / b);
bool paralel(line L1, line L2) {
 return same_d(L1.a, L2.a) && same_d(L1.b, L2.b);
bool intersection(line L1, line L2, point& P) {
 if(paralel(L1, L2))
   return false;
  P.x = (L2.b * L1.c - L1.b * L2.c) / (L2.a * L1.b - L1.a * L2.b);
 if(same_d(L1.b, 0.0))
   P.y = -(L2.a * P.x + L2.c);
 else
   P.y = -(L1.a * P.x + L1.c);
 return true;
double pointToLine(point P, point A, point B, point& C) {
 vec AP(A, P), AB(A, B);
 double u = dot(AP, AB) / norm_sq(AB);
 C = translate(A, scale(AB, u));
```

P = translate(P, flip(v));

```
x = y = 0.0;
  point(double _x, double _y) {
   x = _x;
   y = _y;
  bool operator< (point other) {</pre>
    if(x < other.x + EPS)
      return true;
    if(x + EPS > other.x)
      return false;
    return y < other.y + EPS;</pre>
  bool operator== (point other) {
    return same_d(x, other.x) && same_d(y, other.y);
double e_dist(point P1, point P2) {
 return hypot(P1.x - P2.x, P1.y - P2.y);
double m dist(point P1, point P2) {
 return dabs(P1.x - P2.x) + dabs(P1.y - P2.y);
double pointBetween(point P, point L, point R) {
 return (e_dist(L, P) + e_dist(P, R) == e_dist(L, R));
bool collinear(point P, point L,
               point R) { //newly added(luis), cek 3 poin segaris
 return P.x * (L.y - R.y) + L.x * (R.y - P.y) + R.x * (P.y - L.y) ==
         0; // bole gnti "dabs(x)<"EPS
/*Vector*/
struct vec {
  double x, y;
 vec() {
   x = y = 0.0;
 vec(double _x, double _y) {
   x = _x;
   y = _y;
  vec(point A) {
   x = A.x;
   y = A.y;
 vec(point A, point B) {
   x = B.x - A.x;
    y = B.y - A.y;
vec scale(vec v, double s) {
 return vec(v.x * s, v.y * s);
vec flip(vec v) {
 return vec(-v.x, -v.y);
double dot(vec u, vec v) {
 return (u.x * v.x + u.y * v.y);
double cross(vec u, vec v) {
 return (u.x * v.y - u.y * v.x);
double norm_sq(vec v) {
 return (v.x * v.x + v.y * v.y);
point translate(point P, vec v) {
 return point(P.x + v.x, P.y + v.y);
point rotate(point P, point O, double angle) {
 vec v(0);
```

```
return e_dist(P, C);
double lineToLine(line L1, line L2) {
 if(!paralel(L1, L2))
   return 0.0;
 return dabs(L2.c - L1.c) / sqrt(L1.a * L1.a + L1.b * L1.b);
/*Line Segment*/
struct segment {
 point P, Q;
 line L;
 segment() {
   point T1;
   P = Q = T1;
   line T2;
   L = T2;
 segment(point _P, point _Q) {
   P = P;
   Q = Q;
   if(0 < P)
     swap(P, Q);
   line T(P, Q);
   L = T;
 bool operator== (segment other) {
   return P == other.P && 0 == other.0;
bool onSegment(point P, segment S) {
 if(orientation(S.P, S.Q, P) != 0)
   return false;
 return between_d(P.x, S.P.x, S.Q.x) && between_d(P.y, S.P.y, S.Q.y);
bool s_intersection(segment S1, segment S2) {
 double o1 = orientation(S1.P, S1.Q, S2.P);
 double o2 = orientation(S1.P, S1.Q, S2.Q);
 double o3 = orientation(S2.P, S2.Q, S1.P);
 double o4 = orientation(S2.P, S2.Q, S1.Q);
 if(o1 != o2 && o3 != o4)
   return true;
 if(o1 == 0 && onSegment(S2.P, S1))
   return true;
 if(o2 == 0 && onSegment(S2.Q, S1))
   return true;
 if(o3 == 0 && onSegment(S1.P, S2))
   return true;
 if(o4 == 0 && onSegment(S1.Q, S2))
   return true;
 return false;
double pointToSegment(point P, point A, point B, point& C) {
 vec AP(A, P), AB(A, B);
 double u = dot(AP, AB) / norm_sq(AB);
 if(u < EPS) {
   C = A;
   return e_dist(P, A);
 if(u + EPS > 1.0) {
   C = B;
   return e_dist(P, B);
 return pointToLine(P, A, B, C);
double segmentToSegment(segment S1, segment S2) {
 if(s_intersection(S1, S2))
   return 0.0;
 double ret = INFD;
 point dummy;
 ret = min(ret, pointToSegment(S1.P, S2.P, S2.Q, dummy));
 ret = min(ret, pointToSegment(S1.Q, S2.P, S2.Q, dummy));
```

```
ret = min(ret, pointToSegment(S2.P, S1.P, S1.Q, dummy));
  ret = min(ret, pointToSegment(S2.Q, S1.P, S1.Q, dummy));
 return ret:
/*Circle*/
struct circle {
 point P;
 double r;
  circle() +
   point P1;
   P = P1;
   r = 0.0;
  circle(point _P, double _r) {
   P = P;
   r = _r;
  circle(point P1, point P2) {
   P = mid(P1, P2);
   r = e_dist(P, P1);
  circle(point P1, point P2, point P3) {
   vector<point> T;
   T.clear();
   T.pb(P1);
   T.pb(P2);
   T.pb(P3);
    sort(T.begin(), T.end());
    P1 = T[0];
    P2 = T[1];
    P3 = T[2];
    point M1, M2;
    M1 = mid(P1, P2);
    M2 = mid(P2, P3);
    point Q2, Q3;
    Q2 = rotate(P2, P1, PI / 2);
    Q3 = rotate(P3, P2, PI / 2);
    vec P1Q2(P1, Q2), P2Q3(P2, Q3);
    point M3, M4;
    M3 = translate(M1, P1Q2);
    M4 = translate(M2, P2Q3);
    line L1(M1, M3), L2(M2, M4);
   intersection(L1, L2, P);
   r = e_{dist(P, P1)};
 bool operator==(circle other) {
    return (P == other.P && same_d(r, other.r));
bool insideCircle(point P, circle C) {
 return e_dist(P, C.P) <= C.r + EPS;</pre>
bool c_intersection(circle C1, circle C2, point& P1, point& P2) {
 double d = e_dist(C1.P, C2.P);
 if(d > C1.r + C2.r) {
   return false; //d+EPS kalo butuh
 if(d < dabs(C1.r - C2.r) + EPS)
   return false;
  double x1 = C1.P.x, y1 = C1.P.y, r1 = C1.r, x2 = C2.P.x, y2 = C2.P.y, r2 = C2.r;
 double a = (r1 * r1 - r2 * r2 + d * d) / (2 * d), h = sqrt(r1 * r1 - a * a);
  point T(x1 + a * (x2 - x1) / d, y1 + a * (y2 - y1) / d);
  P1 = point(T.x - h * (y2 - y1) / d, T.y + h * (x2 - x1) / d);
  P2 = point(T.x + h * (y2 - y1) / d, T.y - h * (x2 - x1) / d);
 return true:
bool lc_intersection(line L, circle 0, point& P1, point& P2) {
 double a = L.a, b = L.b, c = L.c, x = 0.P.x, y = 0.P.y, r = 0.r;
  double A = a * a + b * b, B = 2 * a * b * y - 2 * a * c - 2 * b * b * x,
        C = b * b * x * x + b * b * y * y - 2 * b * c * y + c * c - b * b * r * r;
  double D = B * B - 4 * A * C;
```

point T1, T2;

if(same_d(b, 0.0)) {

return false;

return true;

return true;

P1 = P2 = T1;

return true;

return false;

if(D < EPS)

D = sqrt(D);

P1 = T1;

P2 = T2;

if(!cek)

return true;

if(P1 == P2)

if(!b1)

if(!b2) P2 = P1;

return b1; if(b1 || b2) {

P1 = P2;

return true;

vector<point> T;
T.clear();

sort(T.begin(), T.end());

T.pb(A);

T.pb(B);

T.pb(C);

A = T[0];

B = T[1];

C = T[2];

vec BC(B, C);

return false;

/*Triangle*/

return false:

if(same_d(D, 0.0)) { T1.x = -B / (2 * A);

if(dabs(x - T1.x) + EPS > r)

P1 = P2 = point(T1.x, y);

P1 = point(T1.x, y - dy);

P2 = point(T1.x, y + dy);

T1.y = (c - a * T1.x) / b;

T1.x = (-B - D) / (2 * A);T1.y = (c - a * T1.x) / b;

T2.x = (-B + D) / (2 * A);

T2.y = (c - a * T2.x) / b;

 $if(same_d(T1.x - r - x, 0.0) | | same_d(T1.x + r - x, 0.0))$

double dx = dabs(T1.x - x), dy = sqrt(r * r - dx * dx);

bool sc_intersection(segment S, circle C, point& P1, point& P2) {

bool b1 = between_d(P1.x, x1, x2) && between_d(P1.y, y1, y2);

bool b2 = between_d(P2.x, x1, x2) && between_d(P2.y, y1, y2);

double ab = e_dist(A, B), bc = e_dist(B, C), ac = e_dist(C, A);

double r = t_area(A, B, C) / (t_perimeter(A, B, C) / 2);

double x1 = S.P.x, y1 = S.P.y, x2 = S.Q.x, y2 = S.Q.y;

bool cek = lc_intersection(S.L, C, P1, P2);

double t_perimeter(point A, point B, point C) {
 return e_dist(A, B) + e_dist(B, C) + e_dist(C, A);

circle t_inCircle(point A, point B, point C) {

double ratio = e_dist(A, B) / e_dist(A, C);

return sqrt(s * (s - ab) * (s - bc) * (s - ac));

double t_area(point A, point B, point C) {

double s = t_perimeter(A, B, C) / 2;

T1.x = c / a;

```
point P;
  P = translate(B, BC);
 line AP1(A, P);
  ratio = e_dist(B, A) / e_dist(B, C);
  vec AC(A, C);
  AC = scale(AC, ratio / (1 + ratio));
  P = translate(A, AC);
 line BP2(B, P);
 intersection(AP1, BP2, P);
 return circle(P, r);
circle t_outCircle(point A, point B, point C) {
 return circle(A, B, C);
/*Polygon*/
struct polygon {
 vector<point> P;
 polygon() {
   P.clear();
 polygon(vector<point>& _P) {
   P = P;
bool rayCast(point P, polygon& A) {
  point Q(P.x, 10000);
  line cast(P, Q);
  int cnt = 0;
  FOR(i, (int)(A.P.size()) - 1) {
   line temp(A.P[i], A.P[i + 1]);
    point I;
    bool B = intersection(cast, temp, I);
   if(!B)
     continue;
    else if(I == A.P[i] || I == A.P[i + 1])
   else if(pointBetween(I, A.P[i], A.P[i + 1]) && pointBetween(I, P, Q))
     cnt++;
 return cnt % 2 == 1;
// line segment p-q intersect with line A-B.
point lineIntersectSeg(point p, point q, point A, point B) {
 double a = B.y - A.y;
 double b = A.x - B.x;
 double c = B.x * A.y - A.x * B.y;
 double u = fabs(a * p.x + b * p.y + c);
 double v = fabs(a * q.x + b * q.y + c);
 return point((p.x * v + q.x * u) / (u + v), (p.y * v + q.y * u) / (u + v));
// cuts polygon Q along the line formed by point a -> point b
// (note: the last point must be the same as the first point)
vector<point> cutPolygon(point a, point b, const vector<point>& Q) {
 vector<point> P;
  for(int i = 0; i < (int)Q.size(); i++) {</pre>
    double left1 = cross(toVec(a, b), toVec(a, Q[i]));
    double left2 = 0;
   if(i != (int)Q.size() - 1)
     left2 = cross(toVec(a, b), toVec(a, Q[i + 1]));
   if(left1 > -EPS)
     P.push_back(Q[i]);
   if(left1 * left2 < -EPS)</pre>
     P.push_back(lineIntersectSeg(Q[i], Q[i + 1], a, b));
 if(!P.empty() && !(P.back() == P.front()))
    P.push_back(P.front());
  return P;
circle minCoverCircle(polygon& A) {
 vector<point> p = A.P;
```

BC = scale(BC, ratio / (1 + ratio));

Page

10

return a.first * b.second - a.second * b.first;

// (a-c) . (b-c)

TD dot(Pt a, Pt b, Pt c) {

TD v1 = a.first - c.first;

TD v2 = a.second - c.second;

Peko

point c; circle ret;

double cr = 0.0;

int i, j, k;
c = p[0];

```
for(i = 1; i < p.size(); i++) {</pre>
    if(e_dist(p[i], c) >= cr + EPS) {
      c = p[i], cr = 0;
      ret = circle(c, cr);
      for(j = 0; j < i; j++) {
        if(e_dist(p[j], c) >= cr + EPS) {
          c = mid(p[i], p[j]);
          cr = e_dist(p[i], c);
          ret = circle(c, cr);
          for(k = 0; k < j; k++) {
            if(e dist(p[k], c) >= cr + EPS) {
              ret = circle(p[i], p[j], p[k]);
              c = ret.P;
              cr = ret.r;
 return ret;
/*Geometry Algorithm*/
double DP[110][110];
double minCostPolygonTriangulation(polygon& A) {
 if(A.P.size() < 3)
   return 0;
  FOR(i, A.P.size()) {
    for(int j = 0, k = i; k < A.P.size(); j++, k++) {</pre>
      if(k < j + 2)
        DP[j][k] = 0.0;
      else {
        DP[j][k] = INFD;
        REP(l, j + 1, k - 1) {
          double cost = e_dist(A.P[j], A.P[k]) + e_dist(A.P[k], A.P[l]) + e_dist(A.P[l \leftrightarrow a.P[l])) + e_dist(A.P[l \leftrightarrow a.P[l]))
                         A.P[j]);
          DP[j][k] = min(DP[j][k], DP[j][l] + DP[l][k] + cost);
 return DP[0][A.P.size() - 1];
4.2 Convex Hull
typedef double TD;
                                    // for precision shits
namespace GEOM {
typedef pair<TD, TD> Pt;
                                  // vector and points
const TD EPS = 1e-9;
const TD maxD = 1e9;
TD cross(Pt a, Pt b, Pt c) {
                                  // right hand rule
 TD v1 = a.first - c.first;
                                  // (a-c) X (b-c)
 TD v2 = a.second - c.second;
 TD u1 = b.first - c.first;
 TD u2 = b.second - c.second;
 return v1 * u2 - v2 * u1;
TD cross(Pt a, Pt b) {
                                  // a X b
```

```
TD u1 = b.first - c.first;
  TD u2 = b.second - c.second;
 return v1 * u1 + v2 * u2;
TD dot(Pt a, Pt b) {
                                // a . b
 return a.first * b.first + a.second * b.second;
TD dist(Pt a, Pt b) {
 return sqrt((a.first - b.first) * (a.first - b.first) +
              (a.second - b.second) * (a.second - b.second));
TD shoelaceX2(vector<Pt>& convHull) {
 TD ret = 0;
  for(int i = 0, n = convHull.size(); i < n; i++)</pre>
    ret += cross(convHull[i], convHull[(i + 1) % n]);
 return ret;
vector<Pt> createConvexHull(vector<Pt>& points) {
  sort(points.begin(), points.end());
  vector<Pt> ret;
  for(int i = 0; i < points.size(); i++) {</pre>
    while(ret.size() > 1 &&
          cross(points[i], ret[ret.size() - 1], ret[ret.size() - 2]) < -EPS)</pre>
      ret.pop back();
    ret.push_back(points[i]);
  for(int i = points.size() - 2, sz = ret.size(); i >= 0; i--) {
   while(ret.size() > sz &&
          cross(points[i], ret[ret.size() - 1], ret[ret.size() - 2]) < -EPS)</pre>
     ret.pop_back();
   if(i == 0)
     break;
   ret.push_back(points[i]);
  return ret;
  bool isInside(Pt pv, vector<Pt>& x) { //using winding number
   int n = x.size(), wn = 0;
   x.push_back(x[0]);
    for(int i = 0; i < n; ++i) {
     if(((x[i + 1].first <= pv.first && x[i].first >= pv.first) ||
          (x[i + 1].first >= pv.first && x[i].first <= pv.first)) &&
          ((x[i + 1].second \leftarrow pv.second \& x[i].second \rightarrow pv.second) | |
           (x[i + 1].second >= pv.second && x[i].second <= pv.second))) {
        if(cross(x[i], x[i + 1], pv) == 0) {
          x.pop_back();
          return true;
    for(int i = 0; i < n; ++i) {
     if(x[i].second <= pv.second) {</pre>
       if(x[i + 1].second > pv.second && cross(x[i], x[i + 1], pv) > 0)
     else\ if(x[i+1].second <= pv.second && cross(x[i], x[i+1], pv) < 0)
    x.pop_back();
    return wn != 0;
bool isInside(Pt pv, vector<Pt>& x) { //using winding number
 int n = x.size(), wn = 0;
  x.push_back(x[0]);
  for(int i = 0; i < n; ++i) {</pre>
   if(((x[i + 1].first <= pv.first && x[i].first >= pv.first) ||
        (x[i + 1].first >= pv.first && x[i].first <= pv.first)) &&
        ((x[i + 1].second \le pv.second \& x[i].second >= pv.second) |
         (x[i + 1].second >= pv.second && x[i].second <= pv.second))) {
      if(cross(x[i], x[i + 1], pv) == 0) {
        x.pop_back();
        return true;
```

4.3 Closest Pair of Points

```
#define fi first
#define se second
typedef pair<int, int> pii;
struct Point {
 int x, y, id;
int compareX(const void* a, const void* b) {
 Point* p1 = (Point*)a, *p2 = (Point*)b;
 return (p1->x - p2->x);
int compareY(const void* a, const void* b) {
 Point* p1 = (Point*)a, *p2 = (Point*)b;
 return (p1->y - p2->y);
double dist(Point p1, Point p2) {
 return sqrt((double)(p1.x - p2.x) * (p1.x - p2.x) +
              (double)(p1.y - p2.y) * (p1.y - p2.y)
             );
pair<pii, double> bruteForce(Point P[], int n) {
 double min = 1e8;
 pii ret = pii(-1, -1);
  for(int i = 0; i < n; ++i)
   for(int j = i + 1; j < n; ++j)
      if(dist(P[i], P[j]) < min) {</pre>
        ret = pii(P[i].id, P[j].id);
        min = dist(P[i], P[j]);
 return pair<pii, double> (ret, min);
pair<pii, double> getmin(pair<pii, double> x, pair<pii, double> y) {
 if(x.fi.fi == -1 && x.fi.se == -1)
   return y;
 if(y.fi.fi == -1 && y.fi.se == -1)
    return x;
 return (x.se < y.se) ? x : y;
pair<pii, double> stripClosest(Point strip[], int size, double d) {
 double min = d;
 pii ret = pii(-1, -1);
  qsort(strip, size, sizeof(Point), compareY);
  for(int i = 0; i < size; ++i)</pre>
    for(int j = i + 1; j < size && (strip[j].y - strip[i].y) < min; ++j)</pre>
      if(dist(strip[i], strip[j]) < min) {</pre>
        ret = pii(strip[i].id, strip[j].id);
        min = dist(strip[i], strip[j]);
 return pair<pii, double>(ret, min);
pair<pii, double> closestUtil(Point P[], int n) {
 if(n \le 3)
    return bruteForce(P, n);
  int mid = n / 2;
```

```
Point midPoint = P[mid];
  pair<pii, double> dl = closestUtil(P, mid);
  pair<pii, double> dr = closestUtil(P + mid, n - mid);
  pair<pii, double> d = getmin(dl, dr);
  Point strip[n];
  int j = 0;
  for(int i = 0; i < n; i++)
   if(abs(P[i].x - midPoint.x) < d.second)</pre>
     strip[j] = P[i], j++;
 return getmin(d, stripClosest(strip, j, d.second));
pair<pii, double> closest(Point P[], int n) {
 qsort(P, n, sizeof(Point), compareX);
 return closestUtil(P, n);
Point P[50005];
int main() {
 int n;
  scanf("%d", &n);
  for(int a = 0; a < n; a++) {
   scanf("%d%d", &P[a].x, &P[a].y);
   P[a].id = a;
 pair<pii, double> hasil = closest(P, n);
 if(hasil.fi.fi > hasil.fi.se)
   swap(hasil.fi.fi, hasil.fi.se);
 printf("%d %d %.6lf\n", hasil.fi.fi, hasil.fi.se, hasil.se);
 return 0:
```

4.4 Smallest Enclosing Circle

```
// welzl's algo to find the 2d minimum enclosing circle of a set of points
// expected O(N)
// directions: remove duplicates and shuffle points, then call welzl(points)
struct Point {
 double x;
 double y;
struct Circle {
 double x, y, r;
 Circle() {}
 Circle(double _x, double _y, double _r): x(_x), y(_y), r(_r) {}
Circle trivial(const vector<Point>& r) {
 if(r.size() == 0)
    return Circle(0, 0, -1);
  else if(r.size() == 1)
   return Circle(r[0].x, r[0].y, 0);
  else if(r.size() == 2) {
   double cx = (r[0].x + r[1].x) / 2.0, cy = (r[0].y + r[1].y) / 2.0;
    double rad = hypot(r[0].x - r[1].x, r[0].y - r[1].y) / 2.0;
    return Circle(cx, cy, rad);
  } else {
    double x0 = r[0].x, x1 = r[1].x, x2 = r[2].x;
    double y0 = r[0].y, y1 = r[1].y, y2 = r[2].y;
   double d = (x0 - x2) * (y1 - y2) - (x1 - x2) * (y0 - y2);
   double cx = (((x0 - x2) * (x0 + x2) + (y0 - y2) * (y0 + y2)) / 2 *
                 (y1 - y2) - ((x1 - x2) * (x1 + x2) + (y1 - y2) * (y1 + y2)) / 2
                 * (y0 - y2)) / d;
    double cy = (((x1 - x2) * (x1 + x2) + (y1 - y2) * (y1 + y2)) / 2 *
                 (x0 - x2) - ((x0 - x2) * (x0 + x2) + (y0 - y2) * (y0 + y2)) / 2
                 * (x1 - x2)) / d;
    return Circle(cx, cy, hypot(x0 - cx, y0 - cy));
```

Peko

```
Circle welzl(const vector<Point>& p, int idx = 0, vector<Point> r = {}) {
    if(idx == (int) p.size() || r.size() == 3)
        return trivial(r);
    Circle d = welzl(p, idx + 1, r);
    if(hypot(p[idx].x - d.x, p[idx].y - d.y) > d.r) {
        r.push_back(p[idx]);
        d = welzl(p, idx + 1, r);
    }
    return d;
}
```

4.5 Sutherland-Hodgman Algorithm

```
// Complexity: linear time
// Ada 2 poligon, cari poligon intersectionnya
// poly_point = hasilnya, clipper = pemotongnya
#include<bits/stdc++.h>
using namespace std;
const double EPS = 1e-9;
struct point {
 double x, y;
 point(double _x, double _y): x(_x), y(_y) {}
struct vec {
 double x, y;
 vec(double _x, double _y): x(_x), y(_y) {}
point pivot(0, 0);
vec toVec(point a, point b) {
 return vec(b.x - a.x, b.y - a.y);
double dist(point a, point b) {
 return hypot(a.x - b.x, a.y - b.y);
double cross(vec a, vec b) {
 return a.x * b.y - a.y * b.x;
bool ccw(point p, point q, point r) {
 return cross(toVec(p, q), toVec(p, r)) > 0;
bool collinear(point p, point q, point r) {
 return fabs(cross(toVec(p, q), toVec(p, r))) < EPS;</pre>
bool lies(point a, point b, point c) {
 if((c.x >= min(a.x, b.x) \&\& c.x <= max(a.x, b.x)) \&\&
      (c.y >= min(a.y, b.y) \&\& c.y <= max(a.y, b.y)))
    return true;
  else
    return false;
bool anglecmp(point a, point b) {
 if(collinear(pivot, a, b))
    return dist(pivot, a) < dist(pivot, b);</pre>
  double d1x = a.x - pivot.x, d1y = a.y - pivot.y;
  double d2x = b.x - pivot.x, d2y = b.y - pivot.y;
 return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0;
point intersect(point s1, point e1, point s2, point e2) {
  double x1, x2, x3, x4, y1, y2, y3, y4;
 x1 = s1.x;
 y1 = s1.y;
  x2 = e1.x;
 y2 = e1.y;
  x3 = s2.x;
  y3 = s2.y;
  x4 = e2.x;
```

```
y4 = e2.y;
  double num1 = (x1 * y2 - y1 * x2) * (x3 - x4) - (x1 - x2) * (x3 * y4 - y3 * x4);
  double num2 = (x1 * y2 - y1 * x2) * (y3 - y4) - (y1 - y2) * (x3 * y4 - y3 * x4);
  double den = (x1 - x2) * (y3 - y4) - (y1 - y2) * (x3 - x4);
  double new x = num1 / den;
  double new_y = num2 / den;
 return point(new_x, new_y);
void clip(vector <point>& poly_points, point point1, point point2) {
 vector <point> new_points;
  new_points.clear();
  for(int i = 0; i < poly_points.size(); i++) {</pre>
    int k = (i + 1) % poly_points.size();
    double i_pos = ccw(point1, point2, poly_points[i]);
    double k pos = ccw(point1, point2, poly points[k]);
    //in in
    if(i_pos <= 0 && k_pos <= 0)
     new_points.push_back(poly_points[k]);
    else if(i pos > 0 && k pos <= 0) {
     new_points.push_back(intersect(point1, point2, poly_points[i],
                                     polv points[k]));
     new_points.push_back(poly_points[k]);
    else if(i pos <= 0 && k pos > 0) {
     new_points.push_back(intersect(point1, point2, poly_points[i],
                                     poly_points[k]));
    //out out
    else {
  poly_points.clear();
  for(int i = 0; i < new_points.size(); i++)</pre>
   poly_points.push_back(new_points[i]);
double area(const vector <point>& P) {
 double result = 0.0;
 double x1, y1, x2, y2;
  for(int i = 0; i < P.size() - 1; i++) {</pre>
   x1 = P[i].x;
   y1 = P[i].y;
   x2 = P[i + 1].x;
   y2 = P[i + 1].y;
   result += (x1 * y2 - x2 * y1);
 return fabs(result) / 2;
void suthHodgClip(vector <point>& poly_points, vector <point> clipper_points) {
 for(int i = 0; i < clipper_points.size(); i++) {</pre>
   int k = (i + 1) % clipper_points.size();
    clip(poly_points, clipper_points[i], clipper_points[k]);
vector<point> sortku(vector<point> P) {
 int P0 = 0;
  int i;
  for(i = 1; i < 3; i++) {
   if(P[i].y < P[P0].y \mid | (P[i].y == P[P0].y && P[i].x > P[P0].x))
     P0 = i;
  point temp = P[0];
  P[0] = P[P0];
  P[P0] = temp;
  pivot = P[0];
  sort(++P.begin(), P.end(), anglecmp);
  reverse(++P.begin(), P.end());
  return P;
```

```
int main {
  clipper_points = sortku(clipper_points);
  suthHodgClip(poly_points, clipper_points);
}
```

4.6 Centroid of Polygon

```
C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i \ y_{i+1} - x_{i+1} \ y_i)
C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i \ y_{i+1} - x_{i+1} \ y_i)
```

4.7 Pick Theorem

A: Area of a simply closed lattice polygon B: Number of lattice points on the edges I: Number of points in the interior $A = I + \frac{B}{2} - 1$

5 Graphs

5.1 Articulation Point and Bridge

```
// gr -> adi list
// vector vis, low -> initialize to -1
// int timer -> initialize to 0
void dfs(int pos, int dad = -1) {
 vis[pos] = low[pos] = timer++;
 int kids = 0;
 for(auto& i : gr[pos]) {
   if(i == dad)
     continue:
   if(vis[i] >= 0)
     low[pos] = min(low[pos], vis[i]);
    else {
     dfs(i, pos);
     low[pos] = min(low[pos], low[i]);
     if(low[i] > vis[pos])
       is_bridge(pos, i)
       if(low[i] >= vis[pos] && dad >= 0)
         is_articulation_point(pos)
         ++kids;
 if(dad == -1 && kids > 1)
   is_articulation_point(pos)
```

5.2 SCC and Strong Orientation

```
#define N 10020
vector<int> adj[N];
bool vis[N], ins[N];
int disc[N], low[N], gr[N];
stack<int> st;
int id, grid;
void scc(int cur, int par) {
 disc[cur] = low[cur] = ++id;
 vis[cur] = ins[cur] = 1;
  st.push(cur);
  for(int to : adj[cur]) {
    //if (to==par) continue; // ini untuk SO(scc undirected)
    if(!vis[to])
      scc(to, cur);
    if(ins[to])
      low[cur] = min(low[cur], low[to]);
  if(low[cur] == disc[cur]) {
```

```
grid++; // group id
while(ins[cur]) {
    gr[st.tp] = grid;
    ins[st.tp] = 0;
    st.pop();
    }
}
```

5.3 Centroid Decomposition

```
int build_cen(int nw) {
  com_cen(nw, 0); //fungsi untuk itung size subtree
 int siz = sz[nw] / 2;
  bool found = false;
 while(!found) {
   found = true;
    for(int i : v[nw]) {
     if(!rem[i] && sz[i] < sz[nw]) {</pre>
       if(sz[i] > siz) {
          found = false;
          nw = i:
          break;
 big
  rem[nw] = true;
  for(int i : v[nw])if(!rem[i])
     par_cen[build_cen(i)] = nw;
 return nw;
```

5.4 Dinic's Maximum Flow

```
// O(VE log(max_flow)) if scaling == 1
// O((V + E) sqrt(E)) if unit graph (turn scaling off)
// O((V + E) sqrt(V)) if bipartite matching (turn scaling off)
// indices are 0-based
const ll INF = 1e18;
struct Dinic {
  struct Edge {
    int v;
    ll cap, flow;
    Edge(int _v, ll _cap): v(_v), cap(_cap), flow(0) {}
  };
  int n;
  ll lim;
  vector<vector<int>> gr;
  vector<Edge> e;
  vector<int> idx, lv;
  bool has_path(int s, int t) {
    queue<int> q;
    q.push(s);
    lv.assign(n, -1);
    lv[s] = 0;
    while(!q.empty()) {
      int c = q.front();
      q.pop();
      if(c == t)
        break;
      for(auto& i : gr[c]) {
        ll cur_flow = e[i].cap - e[i].flow;
        if(lv[e[i].v] == -1 && cur_flow >= lim) {
          lv[e[i].v] = lv[c] + 1;
```

```
q.push(e[i].v);
 return lv[t] != -1;
ll get_flow(int s, int t, ll left) {
  if(!left || s == t)
    return left;
  while(idx[s] < (int) gr[s].size()) {</pre>
    int i = gr[s][idx[s]];
    if(lv[e[i].v] == lv[s] + 1) {
      ll add = get_flow(e[i].v, t, min(left, e[i].cap - e[i].flow));
        e[i].flow += add;
        e[i ^ 1].flow -= add;
        return add;
    ++idx[s];
  return 0;
Dinic(int vertices, bool scaling = 1) : // toggle scaling here
  n(vertices), lim(scaling ? 1 << 30 : 1), <math>gr(n) {}
void add_edge(int from, int to, ll cap, bool directed = 1) {
  gr[from].push_back(e.size());
  e.emplace_back(to, cap);
  gr[to].push_back(e.size());
  e.emplace_back(from, directed ? 0 : cap);
ll get_max_flow(int s, int t) { // call this
  ll res = 0;
  while(lim) { // scaling
    while(has_path(s, t)) {
      idx.assign(n, 0);
      while(ll add = get_flow(s, t, INF))
        res += add;
    lim >>= 1;
 }
  return res;
```

5.5 Minimum Cost Maximum Flow

```
// 1-based index
template<class T>
using rpq = priority_queue<T, vector<T>, greater<T>>;
const ll INF = 1e18;
struct MCMF {
 struct Edge {
    int v;
    ll cap, cost;
    int rev;
    Edge(int _v, ll _cap, ll _cost, int _rev) :
      v(_v), cap(_cap), cost(_cost), rev(_rev) {}
  ll flow, cost;
 int st, ed, n;
 vector<ll> dist, H;
 vector<int> pv, pe;
```

```
vector<vector<Edge>> adj;
  bool dijkstra() {
   rpq<pair<ll, int>> pq;
   dist.assign(n + 1, INF);
   dist[st] = 0;
    pq.emplace(0, st);
    while(!pq.empty()) {
     auto [cst, pos] = pq.top();
     pq.pop();
     if(dist[pos] < cst)</pre>
       continue;
      for(int i = 0; i < (int) adj[pos].size(); ++i) {</pre>
       auto& e = adj[pos][i];
       int nxt = e.v;
       ll nxt cst = dist[pos] + e.cost + H[pos] - H[nxt];
        if(e.cap > 0 && nxt_cst < dist[nxt]) {</pre>
          dist[nxt] = nxt_cst;
          pe[nxt] = i;
          pv[nxt] = pos;
          pq.emplace(nxt_cst, nxt);
   return dist[ed] != INF;
  MCMF(int _n) : n(_n), pv(n + 1), pe(n + 1), adj(n + 1) {}
  void add_edge(int u, int v, ll cap, ll cst) {
   adj[u].emplace_back(v, cap, cst, adj[v].size());
   adj[v].emplace_back(u, 0, -cst, adj[u].size() - 1);
  pair<ll, ll> solve(int _st, int _ed) {
   st = _st, ed = _ed;
    flow = 0, cost = 0;
   H.assign(n + 1, 0);
    while(dijkstra()) {
     for(int i = 0; i <= n; ++i)
       H[i] += dist[i];
     ll f = INF;
      for(int i = ed; i != st; i = pv[i])
       f = min(f, adj[pv[i]][pe[i]].cap);
      flow += f;
     cost += f * H[ed];
      for(int i = ed; i != st; i = pv[i]) {
       auto& e = adj[pv[i]][pe[i]];
       e.cap -= f;
        adj[i][e.rev].cap += f;
   return {flow, cost};
};
```

5.6 Flows with Demands

```
let S0 be the source and T0 be the original sink
1. add 2 additional nodes, call them S1 and T1
2. connect S0 to nodes normally
3. connect nodes to TO normally
4. for each edge(U, V), cap = original cap - demand
5. for each node N:
  1. add an edge(S1, N), cap = sum of inward demand to N
   2. add an edge(N, T1), cap = sum of outward demand from N
6. add an edge(T0, S0), cap = INF
7. the above is not a typo!
8. run max flow normally
9. for each edge(S1, V) and (U, T1), check if flow == cap
```

if step #9 fails, then it is not possible to satisfy the given demand

Mathematically, let d(e) be the demand of edge e. Let V be the set of every vertex in the graph.

- $c'(S_1, v) = \sum_{u \in V} d(u, v)$ for each edge (s', v).
- $c'(v,T_1) = \sum_{v \in V} d(v,w)$ for each edge (v,t').
- c'(u,v) = c(u,v) d(u,v) for each edge (u,v) in the old network.
- $c'(T_0, S_0) = \infty$

5.7 Hungarian

```
template <typename TD> struct Hungarian {
 TD INF = 1e9; //max_inf
 vector<vector<TD> > adj; // cost[left][right]
 vector<TD> hl, hr, slk;
 vector<int> fl, fr, vl, vr, pre;
 deque<int> q;
 Hungarian(int _n) {
   n = n;
   adj = vector<vector<TD> >(n, vector<TD> (n, 0));
 int check(int i) {
   if(vl[i] = 1, fl[i] != -1)
     return q.push_back(fl[i]), vr[fl[i]] = 1;
     swap(i, fr[fl[i] = pre[i]]);
   return 0;
 void bfs(int s) {
   slk.assign(n, INF);
   vl.assign(n, 0);
   vr = vl;
   q.assign(vr[s] = 1, s);
   for(TD d;;) {
     for(; !q.empty(); q.pop_front()) {
       for(int i = 0, j = q.front(); i < n; i++) {</pre>
         if(d = hl[i] + hr[j] - adj[i][j], !vl[i] && d <= slk[i]) {</pre>
           if(pre[i] = j, d)
             slk[i] = d:
           else if(!check(i))
             return;
     d = INF;
     for(int i = 0; i < n; i++) if(!vl[i] && d > slk[i])
         d = slk[i];
     for(int i = 0; i < n; i++) {
       if(vl[i])
         hl[i] += d;
       else
         slk[i] -= d;
       if(vr[i])
         hr[i] -= d;
     for(int i = 0; i < n; i++) if(!vl[i] && !slk[i] && !check(i))</pre>
 TD solve() {
   fl.assign(n, -1);
   fr = fl;
   hl.assign(n, 0);
   hr = hl:
   pre.assign(n, 0);
   for(int i = 0; i < n; i++)
```

```
hl[i] = *max_element(adj[i].begin(), adj[i].begin() + n);
    for(int i = 0; i < n; i++)
     bfs(i);
    TD ret = 0;
    for(int i = 0; i < n; i++) if(adj[i][fl[i]])</pre>
        ret += adj[i][fl[i]];
    return ret;
}; //i will be matched with fl[i]
```

5.8 Edmonds' Blossom

```
// Maximum matching on general graphs in O(V^2 E)
// Indices are 1-based
// Stolen from ko osaga's cheatsheet
struct Blossom {
  vector<int> vis, dad, orig, match, aux;
  vector<vector<int>> conn;
  int t, N;
  queue<int> 0;
  void augment(int u, int v) {
    int pv = v;
    do {
      pv = dad[v];
      int nv = match[pv];
      match[v] = pv;
      match[pv] = v;
      v = nv;
    } while(u != pv);
  int lca(int v, int w) {
    while(true) {
      if(v) {
        if(aux[v] == t)
          return v;
        aux[v] = t;
        v = orig[dad[match[v]]];
      swap(v, w);
  void blossom(int v, int w, int a) {
    while(orig[v] != a) {
      dad[v] = w;
      w = match[v];
      if(vis[w] == 1) {
        Q.push(w);
        vis[w] = 0;
      orig[v] = orig[w] = a;
      v = dad[w];
  bool bfs(int u) {
    fill(vis.begin(), vis.end(), -1);
    iota(orig.begin(), orig.end(), 0);
    Q = queue<int>();
    Q.push(u);
    vis[u] = 0;
    while(!Q.empty()) {
      int v = Q.front();
      Q.pop();
      for(int x : conn[v]) {
        if(vis[x] == -1) {
          dad[x] = v;
```

vis[x] = 1;

if(!match[x]) {

return 1;

augment(u, x);

Q.push(match[x]);

blossom(x, v, a);blossom(v, x, a);

Blossom(int n) : // n = vertices

for(int i = 0; i <= n; ++i) {

conn[i].clear();

void add_edge(int u, int v) {

conn[u].push_back(v);

conn[v].push_back(u);

aux(n + 1), conn(n + 1), t(0), N(n) {

match[i] = aux[i] = dad[i] = 0;

int solve() { // call this for answer

return false;

vis[match[x]] = 0;

} else if(vis[x] == 0 && orig[v] != orig[x]) {

vis(n + 1), dad(n + 1), orig(n + 1), match(n + 1),

int a = lca(orig[v], orig[x]);

```
int ans = 0;
    vector<int> V(N - 1);
    iota(V.begin(), V.end(), 1);
    shuffle(V.begin(), V.end(), mt19937(0x94949));
    for(auto x : V) {
      if(!match[x]) {
       for(auto y : conn[x]) {
         if(!match[y]) {
           match[x] = y, match[y] = x;
            ++ans;
           break;
    for(int i = 1; i <= N; ++i) {
      if(!match[i] && bfs(i))
       ++ans;
   return ans;
};
5.9 Eulerian Path or Cycle
// else call on any vertex
// ans = path result
```

```
// finds a eulerian path / cycle
// visits each edge only once
// properties:
// - cycle: degrees are even
// - path: degrees are even OR degrees are even except for 2 vertices
// how to use: g = adjacency list g[n] = connected to n, undirected
// if there is a vertex u with an odd degree, call dfs(u)
vector<set<int>> g;
vector<int> ans;
```

```
void dfs(int u) {
  while(g[u].size()) {
    int v = *g[u].begin();
    g[u].erase(v);
    g[v].erase(u);
    dfs(v);
  ans.push_back(u);
```

5.10 Hierholzer's Algorithm

```
// Eulerian on Directed Graph
stack<int> path;
vector<int> euler;
inline void hierholzer() {
  path.push(0);
  int cur = 0;
  while(!path.empty()) {
   if(!adi[cur].empty()) {
      path.push(cur);
      int next = adj[cur].back();
      adj[cur].pob();
      cur = next;
   } else {
      euler.pb(cur);
      cur = path.top();
      path.pop();
 reverse(euler.begin(), euler.end());
```

5.11 2-SAT

```
struct TwoSAT {
 int n;
 vector<vector<int>> g, gr;
 vector<int> comp, topological_order, answer;
 vector<bool> vis;
 TwoSAT() {}
 TwoSAT(int _n) :
   n(n), g(2 * n), g(2 * n), comp(2 * n), answer(2 * n), vis(2 * n) {}
  void add_edge(int u, int v) {
   g[u].push_back(v);
   gr[v].push_back(u);
  // For the following three functions
  // int x, bool val: if 'val' is true, we take the variable to be x.
  // Otherwise we take it to be x's complement.
  // At least one of them is true
  void add_clause_or(int i, bool f, int j, bool p) {
   add_edge(i + (f ? n : 0), j + (p ? 0 : n));
   add_edge(j + (p ? n : 0), i + (f ? 0 : n));
  // Only one of them is true
  void add_clause_xor(int i, bool f, int j, bool p) {
   add_clause_or(i, f, j, p);
   add_clause_or(i, !f, j, !p);
  // Both of them have the same value
  void add_clause_and(int i, bool f, int j, bool p) {
   add_clause_xor(i, !f, j, p);
```

```
// Topological sort
void dfs(int u) {
  vis[u] = true;
  for(const auto& v : g[u])
    if(!vis[v])
      dfs(v);
  topological_order.push_back(u);
// Extracting strongly connected components
void scc(int u, int id) {
  vis[u] = true;
  comp[u] = id;
  for(const auto& v : gr[u])
    if(!vis[v])
      scc(v, id);
bool satisfiable() {
  fill(vis.begin(), vis.end(), false);
  for(int i = 0; i < 2 * n; i++)
    if(!vis[i])
      dfs(i);
  fill(vis.begin(), vis.end(), false);
  reverse(topological_order.begin(), topological_order.end());
  for(const auto& v : topological_order)
    if(!vis[v])
      scc(v, id++);
  // Constructing the answer
  for(int i = 0: i < n: i++) {
    if(comp[i] == comp[i + n])
      return false;
    answer[i] = (comp[i] > comp[i + n] ? 1 : 0);
  return true;
```

6 Math

6.1 Extended Euclidean GCD

```
// computes x and y such that ax + by = gcd(a, b) in O(log (min(a, b)))
// returns {gcd(a, b), x, y}
tuple<int, int, int> gcd(int a, int b) {
   if(b == 0) return {a, 1, 0};
   auto [d, x1, y1] = gcd(b, a % b);
   return {d, y1, x1 - y1* (a / b)};
}
```

6.2 Generalized CRT

```
template<typename T>
T extended_euclid(T a, T b, T& x, T& y) {
    if(b == 0) {
        x = 1;
        y = 0;
        return a;
}
T xx, yy, gcd;
gcd = extended_euclid(b, a % b, xx, yy);
x = yy;
y = xx - (yy * (a / b));
return gcd;
```

```
template<typename T>
T MOD(T a, T b) {
 return (a % b + b) % b;
// return x, lcm. x = a % n && x = b % m
template<typename T>
pair<T, T> CRT(T a, T n, T b, T m) {
 T _n, _m;
 T gcd = extended_euclid(n, m, _n, _m);
  if(n == m) {
    if(a == b)
      return pair<T, T>(a, n);
    else
      return pair<T, T>(-1, -1);
  } else if(abs(a - b) % gcd != 0)
    return pair<T, T>(-1, -1);
    T lcm = m * n / gcd;
    T \times = MOD(a + MOD(n \times MOD(n \times ((b - a) / gcd), m / gcd), lcm), lcm);
    return pair<T, T>(x, lcm);
```

6.3 Generalized Lucas Theorem

```
/*Special Lucas : (n,k) % p^x
 fctp[n] = Product of the integers less than or equal
 to n that are not divisible by p
 Precompute fctp*/
LL p
LL E(LL n, int m) {
 LL tot = 0;
 while(n != 0)
   tot += n / m, n /= m;
 return tot;
LL funct(LL n, LL base) {
 LL ans = fast(fctp[base], n / base, base) * fctp[n % base] % base;
 return ans;
LL F(LL n, LL base) {
 LL ans = 1:
 while(n != 0) {
   ans = (ans * funct(n, base)) % base;
   n /= p;
 return ans;
LL special_lucas(LL n, LL r, LL base) {
 p = fprime(base);
 LL pow = E(n, p) - E(n - r, p) - E(r, p);
 LL TOP = fast(p, pow, base) * F(n, base) % base;
 LL BOT = F(r, base) * F(n - r, base) % base;
 return (TOP * fast(BOT, totien(base) - 1, base)) % base;
//End of Special Lucas
```

6.4 Linear Diophantine

```
//FOR SOLVING MINIMUM ABS(X) + ABS(Y)
ll x, y, newX, newY, target = 0;
ll extGcd(ll a, ll b) {
   if(b == 0) {
      x = 1, y = 0;
      return a;
   }
ll ret = extGcd(b, a % b);
newX = y;
```

```
newY = x - y * (a / b);
 x = newX;
 y = newY;
 return ret;
ll fix(ll sol, ll rt) {
 ll ret = 0;
 //CASE SOLUTION(X/Y) < TARGET
 if(sol < target)</pre>
   ret = -floor(abs(sol + target) / (double)rt);
 //CASE SOLUTION(X/Y) > TARGET
 if(sol > target)
   ret = ceil(abs(sol - target) / (double)rt);
 return ret;
ll work(ll a, ll b, ll c) {
 ll gcd = extGcd(a, b);
 ll\ solX = x * (c / gcd);
 ll solY = y * (c / gcd);
 a /= gcd;
 b /= gcd;
 ll fi = abs(fix(solX, b));
 ll se = abs(fix(solY, a));
 ll lo = min(fi, se);
 ll hi = max(fi, se);
 ll ans = abs(solX) + abs(solY);
 for(ll i = lo; i <= hi; i++) {
   ans = min(ans, abs(solX + i * b) + abs(solY - i * a));
   ans = min(ans, abs(solX - i * b) + abs(solY + i * a));
 return ans;
```

6.5 Modular Linear Equation

```
// finds all solutions to ax = b (mod n)
vi modular_linear_equation_solver(int a, int b, int n) {
   int x, y;
   vi ret;
   int g = extended_euclid(a, n, x, y);
   if(!(b % g)) {
      x = mod(x * (b / g), n);
      for(int i = 0; i < g; i++)
         ret.push_back(mod(x + i * (n / g), n));
   }
   return ret;
}</pre>
```

6.6 Miller-Rabin and Pollard's Rho

```
namespace MillerRabin {
const vector<ll> primes = { // deterministic up to 2^64 - 1
 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37
ll gcd(ll a, ll b) {
 return b ? gcd(b, a % b) : a;
ll powa(ll x, ll y, ll p) { // (x ^ y) % p
 if(!y)
    return 1;
 if(y & 1)
    return ((\_int128) \times powa(x, y - 1, p)) \% p;
 ll temp = powa(x, y >> 1, p);
  return ((__int128) temp * temp) % p;
bool miller_rabin(ll n, ll a, ll d, int s) {
 ll x = powa(a, d, n);
 if(x == 1 || x == n - 1)
    return 0;
```

```
for(int i = 0; i < s; ++i) {
   x = ((\_int128) x * x) % n;
   if(x == n - 1)
      return 0;
 return 1;
bool is_prime(ll x) { // use this
 if(x < 2)
    return 0;
  int r = 0;
 ll d = x - 1;
  while((d & 1) == 0) {
   d >>= 1;
    ++r;
  for(auto& i : primes) {
   if(x == i)
     return 1;
   if(miller_rabin(x, i, d, r))
     return 0;
 return 1;
namespace PollardRho {
mt19937_64 generator(chrono::steady_clock::now()
                     .time_since_epoch().count());
uniform_int_distribution<ll> rand_ll(0, LLONG_MAX);
ll f(ll x, ll b, ll n) { // (x^2 + b) % n}
 return (((__int128) x * x) % n + b) % n;
ll rho(ll n) {
 if(n % 2 == 0)
   return 2;
  ll b = rand ll(generator);
  ll x = rand_ll(generator);
  ll y = x;
  while(1) {
   x = f(x, b, n);
   y = f(f(y, b, n), b, n);
    ll d = MillerRabin::gcd(abs(x - y), n);
    if(d != 1)
     return d;
void pollard_rho(ll n, vector<ll>& res) {
 if(n == 1)
    return;
  if(MillerRabin::is_prime(n)) {
    res.push_back(n);
    return;
 ll d = rho(n);
  pollard_rho(d, res);
  pollard_rho(n / d, res);
vector<ll> factorize(ll n, bool sorted = 1) { // use this
  vector<ll> res;
  pollard_rho(n, res);
  if(sorted)
    sort(res.begin(), res.end());
  return res;
```

6.7 Berlekamp-Massey

#include <bits/stdc++.h>

using namespace std;

```
#define pb push_back
typedef long long ll;
#define SZ 233333
const int MOD = 1e9 + 7; //or any prime
ll qp(ll a, ll b) {
 ll x = 1;
 a %= MOD;
  while(b) {
    if(b & 1)
     x = x * a % MOD;
    a = a * a % MOD;
    b >>= 1;
 return x;
namespace linear_seq {
vector<int> BM(vector<int> x) {
 //ls: (shortest) relation sequence (after filling zeroes) so far
  //cur: current relation sequence
  vector<int> ls, cur;
  //lf: the position of ls (t')
  //ld: delta of ls (v')
  int lf = -1, ld = -1;
  for(int i = 0; i < int(x.size()); ++i) {</pre>
    ll t = 0;
    //evaluate at position i
    for(int j = 0; j < int(cur.size()); ++j)</pre>
      t = (t + x[i - j - 1] * (ll)cur[j]) % MOD;
    if((t - x[i]) \% MOD == 0) {
      continue; //good so far
    //first non-zero position
    if(!cur.size()) {
      cur.resize(i + 1);
     lf = i;
      ld = (t - x[i]) % MOD;
      continue;
    //cur=cur-c/ld*(x[i]-t)
    ll k = -(x[i] - t) * qp(ld, MOD - 2) % MOD/*1/ld*/;
    vector<int> c(i - lf - 1); //add zeroes in front
    c.pb(k);
    for(int j = 0; j < int(ls.size()); ++j)</pre>
      c.pb(-ls[j]*k % MOD);
    if(c.size() < cur.size())</pre>
     c.resize(cur.size());
    for(int j = 0; j < int(cur.size()); ++j)</pre>
     c[j] = (c[j] + cur[j]) % MOD;
    //if cur is better than ls, change ls to cur
    if(i - lf + (int)ls.size() >= (int)cur.size())
      ls = cur, lf = i, ld = (t - x[i]) % MOD;
    cur = c;
  for(int i = 0; i < int(cur.size()); ++i)</pre>
   cur[i] = (cur[i] % MOD + MOD) % MOD;
  return cur;
int m; //length of recurrence
//a: first terms
//h: relation
ll a[SZ], h[SZ], t_[SZ], s[SZ], t[SZ];
//calculate p*q mod f
void mull(ll* p, ll* q) {
 for(int i = 0; i < m + m; ++i)
    t_{[i]} = 0;
  for(int i = 0; i < m; ++i) if(p[i])</pre>
      for(int j = 0; j < m; ++j)
        t_{i} = (t_{i} + j) + p[i] * q[j]) % MOD;
  for(int i = m + m - 1; i >= m; --i) if(t_[i])
      //miuns t_{i}x^{i-m}(x^m-\sum_{j=0}^{m-1} x^{m-j-1}h_{j})
```

```
for(int j = m - 1; ~j; --j)
        t_{i} - j - 1 = (t_{i} - j - 1) + t_{i} * h_{i} % MOD;
  for(int i = 0; i < m; ++i)
    p[i] = t_[i];
ll calc(ll K) {
  for(int i = m; ~i; --i)
   s[i] = t[i] = 0;
  //init
  s[0] = 1;
  if(m != 1)
   t[1] = 1;
  else
   t[0] = h[0];
  //binary-exponentiation
  while(K) {
   if(K & 1)
      mull(s, t);
    mull(t, t);
    K >>= 1;
  ll su = 0;
  for(int i = 0; i < m; ++i)</pre>
    su = (su + s[i] * a[i]) % MOD;
  return (su % MOD + MOD) % MOD;
int work(vector<int> x, ll n) {
 if(n < int(x.size()))</pre>
   return x[n];
  vector < int > v = BM(x);
  m = v.size();
  if(!m)
   return 0;
  for(int i = 0; i < m; ++i)
   h[i] = v[i], a[i] = x[i];
  return calc(n);
using linear_seq::work;
const vector<int> sequence = {
 0, 2, 2, 28, 60, 836, 2766
int main() {
 cout << work(sequence, 7) << '\n';</pre>
```

6.8 Fast Fourier Transform

```
using ld = double; // change to long double if reach 10^18
using cd = complex<ld>;
const ld PI = acos(-(ld)1);
void fft(vector<cd>& a, int sign = 1) {
 int n = a.size();
  ld theta = sign * 2 * PI / n;
  for(int i = 0, j = 1; j < n - 1; j++) {
    for(int k = n >> 1; k > (i ^= k); k >>= 1);
    if(j < i)
      swap(a[i], a[j]);
  for(int m, mh = 1; (m = mh << 1) <= n; mh = m) {
    int irev = 0;
    for(int i = 0; i < n; i += m) {
      cd w = exp(cd(0, theta * irev));
      for(int k = n >> 2; k > (irev ^= k); k >>= 1);
      for(int j = i; j < mh + i; j++) {
        int k = j + mh;
        cd x = a[j] - a[k];
        a[j] += a[k];
```

namespace FFT {

```
Page 20 of 25
```

```
a[k] = w * x;
 if(sign == -1) for(cd& i : a)
     i /= n;
vector<ll> multiply(vector<ll> const& a, vector<ll> const& b) {
 vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
 int n = 1;
 while(n < a.size() + b.size())</pre>
   n <<= 1;
 fa.resize(n);
 fb.resize(n);
 fft(fa);
 fft(fb);
 for(int i = 0; i < n; i++)
   fa[i] *= fb[i];
 fft(fa, -1);
 vector<ll> res(n);
 for(int i = 0; i < n; i++)
   res[i] = round(fa[i].real());
 return res;
```

6.9 Number Theoretic Transform

```
/* ---- Adjust the constants here ---- */
const int LN = 24; //23
const int N = 1 << LN:
typedef long long LL; // 2**23 * 119 + 1. 998244353
// `MOD` must be of the form 2**`LN` * k + 1, where k odd.
const LL MOD = 9223372036737335297; // 2**24 * 54975513881 + 1.
const LL PRIMITIVE_ROOT = 3; // Primitive root modulo `MOD`.
/* ---- End of constants ---- */
LL root[N];
inline LL power(LL x, LL y) {
 LL ret = 1;
 for(; y; y >>= 1) {
   if(y & 1)
     ret = (__int128) ret * x % MOD;
   x = (_int128) x * x % MOD;
 return ret;
inline void init_fft() {
 const LL UNITY = power(PRIMITIVE_ROOT, MOD - 1 >> LN);
 root[0] = 1;
 for(int i = 1; i < N; i++)
   root[i] = (__int128) UNITY * root[i - 1] % MOD;
 return;
// n = 2^{k} is the length of polynom
inline void fft(int n, vector<LL>& a, bool invert) {
 for(int i = 1, j = 0; i < n; ++i) {
   int bit = n >> 1;
   for(; j >= bit; bit >>= 1)
     j -= bit;
    i += bit;
    if(i < j)
      swap(a[i], a[j]);
  for(int len = 2; len <= n; len <<= 1) {
   LL wlen = (invert ? root[N - N / len] : root[N / len]);
    for(int i = 0; i < n; i += len) {</pre>
     LL w = 1:
      for(int j = 0; j<len >> 1; j++) {
       LL u = a[i + j];
```

```
LL v = (\_int128) a[i + j + len / 2] * w % MOD;
        a[i + j] = ((\_int128) u + v) % MOD;
       a[i + j + len / 2] = ((\_int128) u - v + MOD) % MOD;
       w = (__int128) w * wlen % MOD;
 if(invert) {
   LL inv = power(n, MOD - 2);
    for(int i = 0; i < n; i++)
     a[i] = (__int128) a[i] * inv % MOD;
 return;
inline vector<LL> multiply(vector<LL> a, vector<LL> b) {
 vector<LL> c;
 int len = 1 << 32 - __builtin_clz(a.size() + b.size() - 2);</pre>
 a.resize(len, 0);
 b.resize(len, 0);
  fft(len, a, false);
  fft(len, b, false);
  c.resize(len);
  for(int i = 0; i < len; ++i)
   c[i] = (__int128) a[i] * b[i] % MOD;
  fft(len, c, true);
 return c;
//FFT::init_fft(); wajib di panggil init di awal
```

6.10 Gauss-Jordan

```
// Gauss-Jordan elimination with full pivoting.
//
// Uses:
   (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT:
             a[][] = an nxn matrix
//
             b[][] = an nxm matrix
//
// OUTPUT:
                    = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT& a, VVT& b) {
 const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T det = 1;
  for(int i = 0; i < n; i++) {
    int pj = -1, pk = -1;
    for(int j = 0; j < n; j++) if(!ipiv[j])</pre>
        for(int k = 0; k < n; k++) if(!ipiv[k])</pre>
            if(pj == -1 \mid | fabs(a[j][k]) > fabs(a[pj][pk])) {
              pj = j;
              pk = k;
    if(fabs(a[pj][pk]) < EPS) {</pre>
      cerr << "Matrix is singular." << endl;</pre>
      exit(0);
```

ipiv[pk]++;

```
swap(a[pj], a[pk]);
   swap(b[pj], b[pk]);
   if(pj != pk)
     det *= −1;
   irow[i] = pj;
   icol[i] = pk;
   T c = 1.0 / a[pk][pk];
   det *= a[pk][pk];
   a[pk][pk] = 1.0;
   for(int p = 0; p < n; p++)
     a[pk][p] *= c;
   for(int p = 0; p < m; p++)
     b[pk][p] *= c;
   for(int p = 0; p < n; p++) if(p != pk) {
       c = a[p][pk];
       a[p][pk] = 0;
       for(int q = 0; q < n; q++)
         a[p][q] -= a[pk][q] * c;
       for(int q = 0; q < m; q++)
         b[p][q] -= b[pk][q] * c;
 for(int p = n - 1; p >= 0; p--) if(irow[p] != icol[p]) {
     for(int k = 0; k < n; k++)
       swap(a[k][irow[p]], a[k][icol[p]]);
 return det;
int main() {
 const int n = 4;
 const int m = 2;
 double A[n][n] = { {1, 2, 3, 4}, {1, 0, 1, 0}, {5, 3, 2, 4}, {6, 1, 4, 6} };
 double B[n][m] = { {1, 2}, {4, 3}, {5, 6}, {8, 7} };
 VVT a(n), b(n);
 for(int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
 double det = GaussJordan(a, b);
 // expected: 60
 cout << "Determinant: " << det << endl;</pre>
 // expected: -0.233333 0.166667 0.133333 0.0666667
              0.166667 0.166667 0.333333 -0.333333
              0.233333 0.833333 -0.133333 -0.0666667
 //
              0.05 -0.75 -0.1 0.2
 cout << "Inverse: " << endl;</pre>
 for(int i = 0; i < n; i++) {
   for(int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
   cout << endl;
 // expected: 1.63333 1.3
              -0.166667 0.5
              2.36667 1.7
              -1.85 -1.35
 cout << "Solution: " << endl;</pre>
 for(int i = 0; i < n; i++) {
   for(int j = 0; j < m; j++)
     cout << b[i][j] << ' ';
   cout << endl;</pre>
```

6.11 Fibonacci Check

```
bool is_fibonacci(int n) {
 return is_perfect_square(5 * n * n + 4)
         || is_perfect_square(5 * n * n - 4);
```

6.12 Derangement

```
der[0] = 1;
der[1] = 0;
for(int i = 2; i <= 10; ++i)
 der[i] = (ll)(i - 1) * (der[i - 1] + der[i - 2]);
```

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6.13 Bernoulli Number

$$\sum_{k=1}^{n} k^{p} = \frac{1}{p+1} \sum_{i=0}^{p} (-1)^{i} {p+1 \choose i} B_{i} n^{p+1-i} \qquad B_{m}^{+} = 1 - \sum_{k=0}^{m-1} {m \choose k} \frac{B_{k}^{+}}{m-k+1}$$

6.14 Forbenius Number

(X * Y) - (X + Y) and total count is (X - 1) * (Y - 1) / 2

6.15 Stars and Bars with Upper Bound

```
P = (1 - X^{r_1+1}) \dots (1 - X^{r_n+1}) = \sum_i c_i X^{e_i}
Ans = \sum_{i} c_i \binom{N - e_i + n - 1}{n - 1}
```

7 Strings

7.1 Aho-Corasick

```
const int K = 26;
struct Vertex {
 int next[K];
 bool leaf = 0;
 int p = -1, ans = 0;
  char pch;
  int link = -1, mlink = -1;
  //magic link, is the link to find the nearest leaf
 int go[K];
 Vertex(int p = -1, char ch = '$') : p(p), pch(ch) {
   fill(begin(next), end(next), -1);
    fill(begin(go), end(go), -1);
vector<Vertex> t;
int add_string(string const& s) {
 int v = 0;
  for(char ch : s) {
   int c = ch - 'a';
   if(t[v].next[c] == -1) {
     t[v].next[c] = t.size();
     t.emplace_back(v, ch);
   v = t[v].next[c];
 t[v].leaf = 1;
 return v;
int go(int v, char ch);
int get_link(int v) {
 if(t[v].link == -1) {
   if(v == 0 || t[v].p == 0)
     t[v].link = 0;
      t[v].link = go(get_link(t[v].p), t[v].pch);
  return t[v].link;
int get_mlink(int v) {
 if(t[v].mlink == -1) {
```

```
Don't Forgor the Proof Peko
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```

if(v == 0 || t[v].p == 0)

t[v].mlink = 0;

```
Page 22 of 25
```

```
else {
      t[v].mlink = go(get_link(t[v].p), t[v].pch);
      if(t[v].mlink && !t[t[v].mlink].leaf) {
        if(t[t[v].mlink].mlink == -1)
          get_mlink(t[v].mlink);
        t[v].mlink = t[t[v].mlink].mlink;
 return t[v].mlink;
int go(int v, char ch) {
 int c = ch - 'a';
 if(t[v].go[c] == -1) {
    if(t[v].next[c] != -1)
     t[v].go[c] = t[v].next[c];
    else
      t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
 return t[v].go[c];
//t.pb(Vertex());
7.2 Eertree
   Eertree - keep track of all palindromes and its occurences
  This code refers to problem Longest Palindromic Substring
https://www.spoj.com/problems/LPS/
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
struct node {
 int next[26];
 int sufflink;
 int len, cnt;
const int N = 1e5 + 69;
int n;
string s;
node tree[N];
int idx, suff;
int ans = 0;
void init_eertree() {
 idx = suff = 2;
 tree[1].len = -1, tree[1].sufflink = 1;
 tree[2].len = 0, tree[2].sufflink = 1;
bool add_letter(int x) {
 int cur = suff, curlen = 0;
 int nw = s[x] - 'a';
 while(1) {
    curlen = tree[cur].len;
   if(x - curlen - 1 >= 0 \&\& s[x - curlen - 1] == s[x])
    cur = tree[cur].sufflink;
  if(tree[cur].next[nw]) {
    suff = tree[cur].next[nw];
    return 0;
 tree[cur].next[nw] = suff = ++idx;
 tree[idx].len = tree[cur].len + 2;
```

```
ans = max(ans, tree[idx].len);
  if(tree[idx].len == 1) +
   tree[idx].sufflink = 2;
   tree[idx].cnt = 1;
   return 1;
 while(1) {
   cur = tree[cur].sufflink;
    curlen = tree[cur].len;
   if(x - curlen - 1 >= 0 \&\& s[x - curlen - 1] == s[x]) {
     tree[idx].sufflink = tree[cur].next[nw];
 tree[idx].cnt = tree[tree[idx].sufflink].cnt + 1;
 return 1;
int main() {
 ios::sync_with_stdio(0);
 cin.tie(0);
 cin >> n >> s;
 init eertree();
  for(int i = 0; i < n; i++)</pre>
   add_letter(i);
 cout << ans << '\n';
 return 0;
```

7.3 Manacher's Algorithm

```
// Computes lps array. lps[i] means the longest palindromic substring centered at i (\leftarrow
    when i is even, it is between characters. when it is odd, it is on characters)lps↔
    [0] = 0; lps[1] = 1;
REP(i, 2, 2 * str.size()) {
 int l = i / 2 - lps[i] / 2;
 int r = (i - 1) / 2 + lps[i] / 2;
 while(1) { // widen
   if(l == 0 || r + 1 == str.size())
     break;
   if(str[l - 1] != str[r + 1])
     break;
   l--, r++;
 lps[i] = r - l + 1;
  // jump
 if(lps[i] > 2) {
   int j = i - 1, k = i + 1; // while lps[j] inside lps[i]
   while(lps[j] - j < lps[i] - i)</pre>
     lps[k++] = lps[j--];
    lps[k] = lps[i] - (i - j); // set lps[k] to edge of lps[i]
   i = k - 1; // jump to mirror, which is k
```

7.4 Suffix Array

```
// stores result in sa and lcp
// if lcp is needed, call SuffixArray(str, 1)
struct SuffixArray {
  int n;
  vector<int> sa, lcp, rnk, cnt;
  vector<pair<int, int>> p;
  SuffixArray(const string& s, bool calc_lcp = 0) :
    n(s.length()), sa(n), lcp(calc_lcp ? n : 0), rnk(n),
    cnt(max(n, 256)), p(n) {
    for(int i = 0; i < n; ++i)
        rnk[i] = s[i];
    iota(sa.begin(), sa.end(), 0);</pre>
```

```
for(int i = 1; i < n; i <<= 1)
      update_sa(i);
    if(!calc_lcp)
     return;
    vector<int> phi(n), plcp(n);
    phi[sa[0]] = -1;
    for(int i = 1; i < n; ++i)</pre>
      phi[sa[i]] = sa[i - 1];
    int l = 0;
    for(int i = 0; i < n; ++i) {</pre>
      if(phi[i] == -1)
        plcp[i] = 0;
      else {
        while((i + l < n) && (phi[i] + l < n)</pre>
              && (s[i + l] == s[phi[i] + l]))
        plcp[i] = l;
        l = max(l - 1, 0);
    for(int i = 0; i < n; ++i)
      lcp[i] = plcp[sa[i]];
  void update_sa(int len) {
    sort_sa(len);
    sort sa(0);
    for(int i = 0; i < n; ++i) p[i] = {rnk[i], rnk[(i + len) % n]};</pre>
    auto lst = p[sa[0]];
    rnk[sa[0]] = 0;
    int cur = 0;
    for(int i = 1; i < n; ++i) {
      if(lst != p[sa[i]]) {
        lst = p[sa[i]];
        ++cur;
      rnk[sa[i]] = cur;
  void sort_sa(int offset) {
    fill(cnt.begin(), cnt.end(), 0);
    for(int i = 0; i < n; ++i)
     ++cnt[rnk[(i + offset) % n]];
    int sum = 0;
    for(int i = 0; i < (int) cnt.size(); ++i) {</pre>
      int temp = cnt[i];
      cnt[i] = sum;
      sum += temp;
    vector<int> temp(n);
    for(int i = 0; i < n; ++i) {
      int cur = cnt[rnk[(sa[i] + offset) % n]]++;
      temp[cur] = sa[i];
       = move(temp);
};
```

7.5 Suffix Automaton

```
struct state {
  int len, link;
  map<char, int>next; //use array if TLE
};

const int MAXLEN = 100005;
  state st[MAXLEN * 2];
  int sz, last;

void sa_init() {
  sz = last = 0;
}
```

```
st[0].len = 0;
  st[0].link = -1;
  st[0].next.clear();
  ++sz;
void sa_extend(char c) {
 int cur = sz++;
  st[cur].len = st[last].len + 1;
  st[cur].next.clear();
  int p;
  for(p = last; p != -1 && !st[p].next.count(c); p = st[p].link)
    st[p].next[c] = cur;
  if(p == -1)
    st[cur].link = 0;
  else {
    int q = st[p].next[c];
    if(st[p].len + 1 == st[q].len)
      st[cur].link = q;
    else {
      int clone = sz++;
      st[clone].len = st[p].len + 1;
      st[clone].next = st[q].next;
      st[clone].link = st[q].link;
      for(; p != -1 \&\& st[p].next[c] == q; p = st[p].link)
        st[p].next[c] = clone;
      st[q].link = st[cur].link = clone;
  last = cur;
// forwarding
for(int i = 0; i < m; i++) {
 while(cur >= 0 && st[cur].next.count(pa[i]) == 0) {
    cur = st[cur].link;
   if(cur != -1)
     len = st[cur].len;
 if(st[cur].next.count(pa[i])) {
    len++:
    cur = st[cur].next[pa[i]];
 } else
    len = cur = 0;
// shortening abc -> bc
if(l == m) {
 1--:
 if(l <= st[st[cur].link].len)</pre>
    cur = st[cur].link;
// finding lowest and highest length
int lo = st[st[cur].link].len + 1;
int hi = st[cur].len;
//Finding number of distinct substrings
//answer = distsub(0)
LL d[MAXLEN * 2];
LL distsub(int ver) {
 LL tp = 1;
 if(d[ver])
    return d[ver];
  for(map<char, int>::iterator it = st[ver].next.begin();
     it != st[ver].next.end(); it++)
    tp += distsub(it->second);
  d[ver] = tp;
  return d[ver];
```

//Total Length of all distinct substrings

```
//call distsub first before call lesub
LL ans[MAXLEN * 2];
LL lesub(int ver) {
  LL tp = 0;
  if(ans[ver])
    return ans[ver];
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++)
    tp += lesub(it->second) + d[it->second];
  ans[ver] = tp;
  return ans[ver];
//find the k-th lexicographical substring
void kthsub(int ver, int K, string& ret) {
  for(map<char, int>::iterator it = st[ver].next.begin();
       it != st[ver].next.end(); it++) {
    int v = it->second;
    if(K <= d[v]) {
      if(K == 0) {
        ret.push_back(it->first);
      } else {
        ret.push_back(it->first);
        kthsub(v, K, ret);
        return;
    } else
      K -= d[v];
// Smallest Cyclic Shift to obtain lexicographical smallest of All possible
//in int main do this
int main() {
  string S;
  sa init();
  cin >> S; //input
  tp = 0:
  t = S.length();
  S += S:
  for(int a = 0; a < S.size(); a++)</pre>
    sa_extend(S[a]);
  minshift(0):
//the function
int tp, t;
void minshift(int ver) {
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++) {
    tp++;
    if(tp == t) {
      cout << st[ver].len - t + 1 << endl;</pre>
      break;
    minshift(it->second);
    break;
//end of function
// LONGEST COMMON SUBSTRING OF TWO STRINGS
string lcs(string s, string t) {
  sa_init();
  for(int i = 0; i < (int)s.length(); ++i)</pre>
    sa_extend(s[i]);
  int v = 0, l = 0,
      best = 0, bestpos = 0;
  for(int i = 0; i < (int)t.length(); ++i) {</pre>
```

```
while(v && ! st[v].next.count(t[i])) {
    v = st[v].link;
    l = st[v].length;
}
if(st[v].next.count(t[i])) {
    v = st[v].next[t[i]];
    ++l;
}
if(l > best)
    best = l, bestpos = i;
}
return t.substr(bestpos - best + 1, best);
}
```

8 OEIS

8.1 A000108 (Catalan)

Catalan numbers $f(n) = nCk(2n,n) / (n+1) = nCk(2n,n) - nCk(2n,n+1) = f(n-1) * 2*(2*n-1) / (n+1) \\ 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, 2674440, 9694845, 35357670, 129644790, 477638700, 1767263190, 6564120420, 24466267020, 91482563640, 343059613650, 1289904147324, 4861946401452, 18367353072152, 69533550916004, 263747951750360, 1002242216651368, 3814986502092304$

8.2 A000127

Maximal number of regions obtained by joining n points around a circle by straight lines $f(n) = (n^4 - 6*n^3 + 23*n^2 - 18*n + 24) / 24$ 1, 2, 4, 8, 16, 31, 57, 99, 163, 256, 386, 562, 794, 1093, 1471, 1941, 2517, 3214, 4048, 5036, 6196, 7547, 9109, 10903, 12951, 15276, 17902, 20854, 24158, 27841, 31931, 36457, 41449, 46938, 52956, 59536, 66712, 74519, 82993, 92171, 102091, 112792, 124314

8.3 A000668 (Mersene Primes)

Mersenne primes (of form 2^p - 1 where p is a prime)
3, 7, 31, 127, 8191, 131071, 524287, 2147483647, 2305843009213693951,
618970019642690137449562111, 162259276829213363391578010288127,
170141183460469231731687303715884105727

8.4 A001434

Number of graphs with n nodes and n edges.
0, 0, 1, 2, 6, 21, 65, 221, 771, 2769, 10250, 39243, 154658, 628635, 2632420, 11353457, 50411413, 230341716, 1082481189, 5228952960, 25945377057, 132140242356, 690238318754

8.5 A018819

Binary partition function: number of partitions of n into powers of 2 f(2m+1) = f(2m); f(2m) = f(2m-1) + f(m)1, 1, 2, 2, 4, 4, 6, 6, 10, 10, 14, 14, 20, 20, 26, 26, 36, 36, 46, 46, 60, 60, 74, 74, 94, 94, 114, 114, 140, 140, 166, 166, 202, 202, 238, 238, 284, 284, 330, 330, 390, 390, 450, 450, 524, 524, 598, 598, 692, 692, 786, 786, 900, 900, 1014, 1014, 1154, 1154, 1294, 1294

8.6 A092098

3-Portolan numbers: number of regions formed by n-secting the angles of an equilateral triangle. long long solve(long long n) {

Page 24 of 25

```
long long res = (n % 2 == 1 ? 3*n*n - 3*n + 1 : 3*n*n - 6*n + 6);
const int bats = n/2 - 1;
for (long long i=1; i<=bats; i++) for (long long j=1; j<=bats; j++) {
    long long num = i * (n-j) * n;
    long long denum = (n-i) * j + i * (n-j);
    res -= 6 * (num % denum == 0 && num / denum <= bats);
} return res;
}
1, 6, 19, 30, 61, 78, 127, 150, 217, 246, 331, 366, 469, 510, 625, 678, 817,
870, 1027, 1080, 1261, 1326, 1519, 1566, 1801, 1878, 2107, 2190, 2437, 2520,
2791, 2886, 3169, 3270, 3559, 3678, 3997, 4110, 4447, 4548, 4921, 5034, 5419,
5550, 5899, 6078, 6487</pre>
```

8.7 A277402

3-Portolan numbers: number of regions formed by n-secting the angles of an equilateral triangle. $a(n) = 6n + 6(n-1 - n\%2) + a(n-2); \quad f(n) = a(n) \text{ if } n \% \text{ 10 != 0 else } a(n) - 12 \\ 1, 6, 19, 30, 61, 78, 127, 150, 217, 234, 331, 366, 469, 510, 631, 678, 817, 870, 1027, 1074, 1261, 1326, 1519, 1590, 1801, 1878, 2107, 2190, 2437, 2514, 2791, 2886, 3169, 3270, 3571, 3678, 3997, 4110, 4447, 4554, 4921, 5046, 5419, 5550, 5941, 6078, 6487, 6630, 7057, 7194$