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# 1 Miscellaneous

# 1.1 VS Code Config

```
"command": "clear; ./runner ${file} ${fileBasenameNoExtension} 0"
"command": "clear; ./runner ${file} ${fileBasenameNoExtension} 1"
"command": "clear; rm -r .cache"
runner:
#!/bin/bash
cp=$1; uf=$3; tmp='.cache'; bp="$tmp/$2"; ch="$tmp/${2}-hash"
ti="${li}\nTime: %es"
rc() {
   mkdir -p "$tmp"
sch() {
    rc; echo -n "$(gh)" > "$ch"
gch() {
   if [[ -f "$ch" ]]; then cat $ch
   else echo -n 'NULL'; fi
gh() {
   sha256sum $cp
   oh=$(gch); nh=$(gh)
   if [[ "$oh" == "$nh" ]]; then return 1
   else return 0; fi
cc() {
   g++ $cp -02 -std=gnu++17 -Wall -Wextra -Wshadow -DLOCAL -o $bp
```

```
ma() {
    echo $li
    if nr; then
        if [[ -f "$bp" ]]; then rm "$bp"; fi
        sch; echo 'Compiling...'; echo $li; cc; echo $li
    if [[ -f "$bp" ]]; then
        echo 'Running...'; echo $li
        if (( $uf )); then command time -f "$ti" ./$bp < IN
        else command time -f "$ti" ./$bp; fi
        echo $li
    fi
}
ma</pre>
```

# 1.2 Day of Date

```
// 0-based
const vector<int> T = {
  0, 3, 2, 5, 0, 3,
  5, 1, 4, 6, 2, 4
}
int day(int d, int m, int y) {
  y -= (m < 3);
  return (y + y / 4 - y / 100 + y / 400 + T[m - 1] + d) % 7;
}</pre>
```

# 1.3 Number of Days since 1-1-1

```
int rdn(int d, int m, int y) {
```

```
if(m < 3)
  --y, m += 12;
return 365 * y + y / 4 - y / 100 + y / 400
       + (153 * m - 457) / 5 + d - 306;
```

## 1.4 Enumerate Subsets of a Bitmask

```
int x = 0;
do {
 // do stuff with the bitmask here
x = (x + 1 + \sim m) \& m;
} while(x != 0);
```

#### 1.5 Fast IO

int read() {

```
char c;
 do {
   c = getchar_unlocked();
 } while(c < 33);</pre>
 int res = 0;
 int mul = 1;
 if(c == '-') {
   mul = -1;
   c = getchar_unlocked();
 while('0' <= c && c <= '9') {
   res = res * 10 + c - '0';
   c = getchar_unlocked();
 return res * mul;
void write(int x) {
 static char wbuf[10];
 if(x < 0) {
   putchar_unlocked('-');
   x = -x;
 int idx = 0;
 while(x) {
   wbuf[idx++] = x % 10;
   x /= 10;
 if(idx == 0)
   putchar_unlocked('0');
 for(int i = idx - 1; i >= 0; --i)
   putchar_unlocked(wbuf[i] + '0');
void write(const char* s) {
 while(*s) {
   putchar_unlocked(*s);
   ++s;
```

# 1.6 Int to Roman

```
const string R[] = {
      "M", "CM", "D", "CD", "C", "XC", "L", "XL", "X", "IX", "V", "IV", "I"
Page
\sim
    const int N[] = {
      1000, 900, 500, 400, 100, 90,
      50, 40, 10, 9, 5, 4, 1
```

```
};
string to_roman(int x) {
 if(x == 0) {
    return "0"; // Not decimal 0!
  string res = "";
  for(int i = 0; i < 13; ++i)
   while(x >= N[i])
     x -= N[i], res += R[i];
  return res;
```

# 1.7 Josephus Problem

```
ll josephus(ll n, ll k) { // O(k log n)
 if(n == 1)
   return 0;
 if(k == 1)
   return n - 1;
 if(k > n)
   return (josephus(n - 1, k) + k) % n;
 ll\ cnt = n / k;
 ll res = josephus(n - cnt, k);
 res -= n % k;
 if(res < 0)
   res += n;
 else
   res += res / (k - 1);
 return res;
int josephus(int n, int k) { // O(n)
 int res = 0;
 for(int i = 1; i <= n; ++i)
   res = (res + k) % i;
 return res + 1;
```

#### 1.8 Random Primes

36671 74101 724729 825827 924997 1500005681 2010408371 2010405347

## 1.9 RNG

```
// RNG - rand_int(min, max), inclusive
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
template<class T>
T rand_int(T mn, T mx) {
 return uniform_int_distribution<T>(mn, mx)(rng);
```

### 2 Data Structures

### 2.1 2D Segment Tree

```
struct Segtree2D {
 struct Segtree {
   struct node {
     int l, r, val;
     node* lc, *rc;
     node(int _l, int _r, int _val = INF) : l(_l), r(_r), val(_val),
       lc(NULL), rc(NULL) {}
```

```
typedef node* pnode;
  pnode root;
  Segtree(int l, int r) {
    root = new node(l, r);
  void update(pnode& nw, int x, int val) {
    int l = nw - > l, r = nw - > r, mid = (l + r) / 2;
    if(l == r)
      nw->val = val;
    else {
      assert(l <= x && x <= r);
      pnode& child = x <= mid ? nw->lc : nw->rc;
        child = new node(x, x, val);
      else if(child->l <= x && x <= child->r)
        update(child, x, val);
      else {
        do {
          if(x \le mid)
            r = mid;
          else
            l = mid + 1;
          mid = (l + r) / 2;
        } while((x <= mid) == (child->l <= mid));</pre>
        pnode nxt = new node(l, r);
        if(child->l <= mid)</pre>
          nxt->lc = child;
        else
          nxt->rc = child;
        child = nxt;
        update(nxt, x, val);
      nw->val = min(nw->lc ? nw->lc->val : INF,
                    nw->rc ? nw->rc->val : INF);
  int query(pnode& nw, int x1, int x2) {
    if(!nw)
      return INF;
    int& l = nw->l, &r = nw->r;
    if(r < x1 || x2 < l)
      return INF;
    if(x1 <= l && r <= x2)
      return nw->val;
    int ret = min(query(nw->lc, x1, x2),
                  query(nw->rc, x1, x2));
    return ret;
  void update(int x, int val) {
    assert(root->l <= x && x <= root->r);
    update(root, x, val);
  int query(int l, int r) {
    return query(root, l, r);
};
struct node {
  int l, r;
  Segtree y;
  node* lc, *rc;
  node(int _l, int _r) : l(_l), r(_r), y(0, MAX),
    lc(NULL), rc(NULL) {}
typedef node* pnode;
```

```
pnode root;
  Segtree2D(int l, int r) {
   root = new node(l, r);
  void update(pnode& nw, int x, int y, int val) {
    int& l = nw->l, &r = nw->r, mid = (l + r) / 2;
   if(l == r)
     nw->y.update(y, val);
    else {
     if(x <= mid) {
        if(!nw->lc)
          nw->lc = new node(l, mid);
       update(nw->lc, x, y, val);
     } else {
       if(!nw->rc)
         nw->rc = new node(mid + 1, r);
       update(nw->rc, x, y, val);
     val = min(nw -> lc ? nw -> lc -> y.query(y, y) : INF,
                nw->rc ? nw->rc->y.query(y, y) : INF);
      nw->y.update(y, val);
  int query(pnode& nw, int x1, int x2, int y1, int y2) {
   if(!nw)
     return INF;
    int& l = nw->l, &r = nw->r;
   if(r < x1 || x2 < l)
     return INF:
   if(x1 <= l && r <= x2)
     return nw->y.query(y1, y2);
   int ret = min(query(nw->lc, x1, x2, y1, y2),
                  query(nw->rc, x1, x2, y1, y2));
   return ret;
  void update(int x, int y, int val) {
   assert(root->l <= x && x <= root->r);
   update(root, x, y, val);
  int query(int x1, int x2, int y1, int y2) {
    return query(root, x1, x2, y1, y2);
};
```

# 2.2 Fenwick RU-RQ

```
void updtRL(int l, int r, ll val) {
 updt(BIT1, l, val), updt(BIT1, r + 1, -val);
 updt(BIT2, l, val * (l - 1)), updt(BIT2, r + 1, -val * r);
ll query(int k) {
 return que(BIT1, k) * k - que(BIT2, k);
```

# 2.3 Heavy-Light Decomposition

```
struct HLD {
  vector<int> id, size, idx, up, root, st;
  vector<vector<int>> adj, chain;
  SegTree seg;
  HLD(const vector<vector<int>>& edges) :
```

```
precompute(0, -1);
    decompose(0, -1);
    int cnt = 0;
    st.resize(chain.size());
    for(int i = 0; i < (int) chain.size(); ++i) {</pre>
      st[i] = cnt;
      cnt += chain[i].size();
  void precompute(int pos, int dad) {
    size[pos] = 1;
    up[pos] = dad;
    for(auto& i : adj[pos]) {
      if(i != dad) {
        precompute(i, pos);
        size[pos] += size[i];
   }
  void decompose(int pos, int dad) {
    if(id[pos] == -1) {
      id[pos] = chain.size();
      root.push_back(pos);
      chain.emplace_back();
    idx[pos] = chain[id[pos]].size();
    chain[id[pos]].push_back(pos);
    int mx = 0, heavy = -1;
    for(auto& i : adj[pos]) {
      if(i != dad && size[i] > mx) {
        mx = size[i];
        heavy = i;
    if(heavy != -1)
      id[heavy] = id[pos];
    for(auto& i : adj[pos]) {
      if(i != dad)
        decompose(i, pos);
 }
  void update(int ch, int l, int r, int val) {
    seg.update(st[ch] + l, st[ch] + r, val);
  int query(int ch, int l, int r, int val) {
    return seg.query(st[ch] + l, st[ch] + r, val);
};
// how to move from u to v
while(1) {
 if(hld.id[u] == hld.id[v]) {
    if(hld.idx[u] > hld.idx[v])
      swap(u, v);
    hld.update(hld.id[u], hld.idx[u], hld.idx[v], w);
    // or hld.query(hld.id[u], hld.idx[u], hld.idx[v]);
    break;
  if(hld.id[u] < hld.id[v])</pre>
    swap(u, v);
  hld.update(hld.id[u], 0, hld.idx[u], w);
 // or hld.query(hld.id[u], 0, hld.idx[v]);
 u = hld.up[hld.root[hld.id[u]]];
```

n(edges.size()), id(n, -1), size(n, -1), idx(n, -1),

up(n, -1), adj(edges), seg(n) {

#### 2.4 Li Chao Tree

```
typedef long long int TD;
const TD INF = 100000000000000;
namespace LICHAO {
struct Node {
 TD m, c;
 Node* l, *r;
};
Node* newNode(Node* x = NULL) {
 Node* ret = (Node*)malloc(sizeof(Node));
 if(x)
   ret->m = x->m, ret->c = x->c;
 ret->l = ret->r = NULL;
 return ret;
void update(Node* k, TD l, TD r, TD m, TD c) {
 TD mid = l + r >> 1;
 bool le = m * l + c < k-> m * l + k->c;
 bool ri = m * mid + c < k->m * mid + k->c;
 if(ri)
   swap(k->m, m), swap(k->c, c);
 if(r - l <= 1)
   return;
  else if(le != ri)
   update((k->1) ? (k->1) : (k->1 = newNode(k)), l, mid, m, c);
   update((k->r) ? (k->r) : (k->r = newNode(k)), mid, r, m, c);
TD query(Node* k, TD l, TD r, TD p) {
 if(!k)
   return INF;
 if(r - l <= 1)
   return p * k->m + k->c;
 if(p < (l + r >> 1))
   return min(p * k->m + k->c, query(k->l, l, l + r >> 1, p));
   return min(p * k->m + k->c, query(k->r, l + r >> 1, r, p));
```

## 2.5 Persistent Segment Tree

```
class PersistentSegtree {
private:
 int n, ptr, sz;
 struct P {
   int val = 0, l, r;
 vector<P> node;
 vector<int> root;
  int newNode() {
   return ptr++;
 int copyNode(int idx) {
   node[ptr] = node[idx];
   return ptr++;
  int build(int l, int r) {
    int idx = newNode();
   if(l == r)
     return idx;
   node[idx].l = build(l, (l + r) / 2);
   node[idx].r = build((l + r) / 2 + 1, r);
   return idx;
  int update(int idx, int l, int r, int x, int val) {
    idx = copyNode(idx);
```

```
if(l == r) {
     node[idx].val += val;
     return idx;
   int mid = (l + r) / 2;
   if(x \le mid)
     node[idx].l = update(node[idx].l, l, mid, x, val);
     node[idx].r = update(node[idx].r, mid + 1, r, x, val);
    node[idx].val = node[node[idx].l].val + node[node[idx].r].val;
   return idx;
  int query(int idxl, int idxr, int l, int r, int x, int y) {
   if(y < l \mid | r < x)
     return 0;
   if(x <= l && r <= y)
     return node[idxr].val - node[idxl].val;
   int mid = (l + r) / 2;
   return query(node[idxl].l, node[idxr].l, l, mid, x, y)
          + query(node[idxl].r, node[idxr].r, mid + 1, r, x, y);
public:
 PersistentSegtree(int _n) : n(_n), ptr(0) {
   sz = 30 * n;
   node.resize(sz);
   root.push back(build(1, n));
 void update(int x, int val) {
   root.push_back(update(root.back(), 1, n, x, val));
 int query(int l, int r, int x, int y) {
   return query(root[l - 1], root[r], 1, n, x, y);
};
```

## 2.6 STL PBDS

# 2.7 Treap

```
// Complexity: O(log N) for split and merge
//
// empty treap: Treap* tr = nullptr;
// insert v at x: [l, r] = split(tr, x), m = Treap(v), merge lmr
// delete at x: [l, r] = split(tr, x), [m, r] = split(r, 1), merge lr
// lazy prop: propagate every time a node is accessed
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
using Key = int;
struct Treap {
   Key val;
   Treap* left;
   Treap* right;
```

```
int prio, sz;
 Treap() {}
 Treap(int _val);
int size(Treap* tr) {
 return tr ? tr->sz : 0;
void update(Treap* tr) {
 tr->sz = 1 + size(tr->left) + size(tr->right);
Treap::Treap(Key _val) :
 val(_val), left(nullptr), right(nullptr), prio(rng()) {
 update(this);
pair<Treap*, Treap*> split(Treap* tr, int sz) {
 if(!tr) return {nullptr, nullptr};
 int left sz = size(tr->left);
 if(sz <= left_sz) {</pre>
   auto [left, mid] = split(tr->left, sz);
   tr->left = mid;
   update(tr);
   return {left, tr};
 } else {
   auto [mid, right] = split(tr->right, sz - left_sz - 1);
   tr->right = mid;
   update(tr);
   return {tr, right};
Treap* merge(Treap* l, Treap* r) {
 if(!l)
   return r;
  if(!r)
   return l;
  if(l->prio < r->prio) {
   l->right = merge(l->right, r);
   update(l);
   return l;
 } else {
   r->left = merge(l, r->left);
   update(r);
   return r;
```

# 2.8 Unordered Map Custom Hash

#### 2.9 Mo's on Tree

```
ST(u) \leq ST(v)
P = LCA(u, v)
If P = u, query [ST(u), ST(v)]
Else query [EN(u), ST(v)] + [ST(P), ST(P)]
```

# 3 Dynamic Programming

### 3.1 DP Convex Hull

```
/* dp[i] = min k < i \{dp[k] + x[i] * m[k]\}
   Make sure gradient (m[i]) is either non-increasing if min,
   or non-decreasing if max. x[i] must be non-decreasing. just sort */
// while this is true, pop back from dg. a=new line, b=last, c=2nd last
bool cekx(int a, int b, int c) {
 // if not enough, change to cross mul
  // if cross mul, beware of negative denominator, and overflow
  return (double)(y[b] - y[a]) / (m[a] - m[b]) \le (double)(y[c] - y[b]) /
         (m[b] - m[c]);
```

#### 3.2 DP DNC

```
void f(int rem, int l, int r, int optl, int optr) {
 if(l > r)
  return;
 int mid = l + r >> 1;
 int opt = MOD, optid = mid;
 for(int i = optl; i <= mid && i <= optr; ++i) {</pre>
   if(dp[rem - 1][i] + c[i][mid] < opt) {</pre>
     opt = dp[rem - 1][i] + c[i][mid];
     optid = i;
 dp[rem][mid] = opt;
 f(rem, l, mid - 1, optl, optid);
 f(rem, mid + 1, r, optid, optr);
 return:
rep(i, 1, n)dp[1][i] = c[0][i];
rep(i, 2, k)f(i, i, n, i, n);
```

#### 3.3 DP Knuth-Yao

```
// opt[i+1][j] <= opt[i][j] <= opt[i][j+1]
// dp[i][j] = min{k} dp[i][k]+dp[k][j]+cost[i][j]
for(int k = 0; k \le n; k++) {
 for(int i = 0; i + k <= n; i++) {
   if(k < 2)
      dp[i][i + k] = 0, opt[i][i + k] = i;
      int sta = opt[i][i + k - 1];
      int end = opt[i + 1][i + k];
      for(int j = sta; j <= end; j++) {</pre>
        if(dp[i][j] + dp[j][i + k] + cost[i][i + k] < dp[i][i + k]) {
          dp[i][i + k] = dp[i][j] + dp[j][i + k] + cost[i][i + k];
          opt[i][i + k] = j;
```

# 4 Geometry

## 4.1 Geometry Template

```
Don't Forgor the Proof Peko
Bina Nusantara University
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0. Basic Rule
    0.1. Everything is in double
    0.2. Every comparison use EPS
    0.3. Every degree in rad
1. General Double Operation
    1.1. const double EPS=1E-9
    1.2. const double PI=acos(-1.0)
    1.3. const double INFD=1E9
    1.3. between d(double x,double l,double r)
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    1.4. same_d(double x,double y)
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    1.5. dabs(double x)
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    2.1. struct point
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        2.1.2. point()
            default constructor
        2.1.3. point(double _x,double _y)
            constructor, set the point to (_x,_y)
        2.1.4. bool operator< (point other)
            regular pair <double, double > operator < with EPS
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    2.2. hypot(point P)
        length of hypotenuse of point P to (0,0)
    2.3. e_dist(point P1,point P2)
        euclidean distance from P1 to P2
    2.4. m_dist(point P1,point P2)
        manhattan distance from P1 to P2
    2.5. point rotate(point P,point O,double angle)
        rotate point P from the origin O by angle ccw
Vector
    3.1. struct vec
        3.1.1. double x,y
            x and y magnitude of the vector
        3.1.2. vec()
            default constructor
        3.1.3. vec(double _x,double _y)
            constructor, set the vector to (_x,_y)
        3.1.4. vec(point A, point B)
            constructor, set the vector to vector AB (A->B)
/*General Double Operation*/
const double PI = acos(-1.0);
const double INFD = 1E9;
double between_d(double x, double l, double r) {
 return (min(l, r) \le x + EPS \&\& x \le max(l, r) + EPS);
double same_d(double x, double y) {
  return between_d(x, y, y);
double dabs(double x) {
  if(x < EPS)
    return -x;
  return x;
/*Point*/
struct point {
  double x, y;
  point() {
```

```
return translate(point(P.x * cos(angle) - P.y * sin(angle),
                         P.x * sin(angle) + P.y * cos(angle)), v);
point mid(point P, point Q) {
 return point((P.x + Q.x) / 2, (P.y + Q.y) / 2);
double angle(point A, point O, point B) {
 vec OA(0, A), OB(0, B);
 return acos(dot(OA, OB) / sqrt(norm_sq(OA) * norm_sq(OB)));
int orientation(point P, point Q, point R) {
 vec PQ(P, Q), PR(P, R);
 double c = cross(PQ, PR);
 if(c < -EPS)
   return -1;
 if(c > EPS)
   return 1;
 return 0;
/*Line*/
struct line {
 double a, b, c;
 line() {
   a = b = c = 0.0;
 line(double _a, double _b, double _c) {
   a = _a;
   b = b;
   c = _c;
  line(point P1, point P2) {
   if(P1 < P2)
     swap(P1, P2);
   if(same_d(P1.x, P2.x))
     a = 1.0, b = 0.0, c = -P1.x;
   else
     a = -(P1.y - P2.y) / (P1.x - P2.x), b = 1.0, c = -(a * P1.x) - P1.y;
 line(point P, double slope) {
   if(same_d(slope, INFD))
     a = 1.0, b = 0.0, c = -P.x;
   else
     a = -slope, b = 1.0, c = -(a * P.x) - P.y;
 bool operator== (line other) {
   return same_d(a, other.a) && same_d(b, other.b) && same_d(c, other.c);
 double slope() {
   if(same_d(b, 0.0))
     return INFD;
   return -(a / b);
bool paralel(line L1, line L2) {
 return same_d(L1.a, L2.a) && same_d(L1.b, L2.b);
bool intersection(line L1, line L2, point& P) {
 if(paralel(L1, L2))
   return false;
  P.x = (L2.b * L1.c - L1.b * L2.c) / (L2.a * L1.b - L1.a * L2.b);
 if(same_d(L1.b, 0.0))
   P.y = -(L2.a * P.x + L2.c);
 else
   P.y = -(L1.a * P.x + L1.c);
 return true;
double pointToLine(point P, point A, point B, point& C) {
 vec AP(A, P), AB(A, B);
 double u = dot(AP, AB) / norm_sq(AB);
 C = translate(A, scale(AB, u));
```

P = translate(P, flip(v));

```
x = y = 0.0;
  point(double _x, double _y) {
   x = _x;
   y = _y;
  bool operator< (point other) {</pre>
    if(x < other.x + EPS)
      return true;
    if(x + EPS > other.x)
      return false;
    return y < other.y + EPS;</pre>
  bool operator== (point other) {
    return same_d(x, other.x) && same_d(y, other.y);
double e_dist(point P1, point P2) {
 return hypot(P1.x - P2.x, P1.y - P2.y);
double m dist(point P1, point P2) {
 return dabs(P1.x - P2.x) + dabs(P1.y - P2.y);
double pointBetween(point P, point L, point R) {
 return (e_dist(L, P) + e_dist(P, R) == e_dist(L, R));
bool collinear(point P, point L,
               point R) { //newly added(luis), cek 3 poin segaris
 return P.x * (L.y - R.y) + L.x * (R.y - P.y) + R.x * (P.y - L.y) ==
         0; // bole gnti "dabs(x)<"EPS
/*Vector*/
struct vec {
  double x, y;
 vec() {
   x = y = 0.0;
 vec(double _x, double _y) {
   x = _x;
   y = _y;
  vec(point A) {
   x = A.x;
   y = A.y;
 vec(point A, point B) {
   x = B.x - A.x;
    y = B.y - A.y;
vec scale(vec v, double s) {
 return vec(v.x * s, v.y * s);
vec flip(vec v) {
 return vec(-v.x, -v.y);
double dot(vec u, vec v) {
 return (u.x * v.x + u.y * v.y);
double cross(vec u, vec v) {
 return (u.x * v.y - u.y * v.x);
double norm_sq(vec v) {
 return (v.x * v.x + v.y * v.y);
point translate(point P, vec v) {
 return point(P.x + v.x, P.y + v.y);
point rotate(point P, point O, double angle) {
 vec v(0);
```

```
return e_dist(P, C);
double lineToLine(line L1, line L2) {
 if(!paralel(L1, L2))
   return 0.0;
 return dabs(L2.c - L1.c) / sqrt(L1.a * L1.a + L1.b * L1.b);
/*Line Segment*/
struct segment {
 point P, Q;
 line L;
 segment() {
   point T1;
   P = Q = T1;
   line T2;
   L = T2;
 segment(point _P, point _Q) {
   P = P;
   Q = Q;
   if(0 < P)
     swap(P, Q);
   line T(P, Q);
   L = T;
 bool operator== (segment other) {
   return P == other.P && 0 == other.0;
bool onSegment(point P, segment S) {
 if(orientation(S.P, S.Q, P) != 0)
   return false;
 return between_d(P.x, S.P.x, S.Q.x) && between_d(P.y, S.P.y, S.Q.y);
bool s_intersection(segment S1, segment S2) {
 double o1 = orientation(S1.P, S1.Q, S2.P);
 double o2 = orientation(S1.P, S1.Q, S2.Q);
 double o3 = orientation(S2.P, S2.Q, S1.P);
 double o4 = orientation(S2.P, S2.Q, S1.Q);
 if(o1 != o2 && o3 != o4)
   return true;
 if(o1 == 0 && onSegment(S2.P, S1))
   return true;
 if(o2 == 0 && onSegment(S2.Q, S1))
   return true;
 if(o3 == 0 && onSegment(S1.P, S2))
   return true;
 if(o4 == 0 && onSegment(S1.Q, S2))
   return true;
 return false;
double pointToSegment(point P, point A, point B, point& C) {
 vec AP(A, P), AB(A, B);
 double u = dot(AP, AB) / norm_sq(AB);
 if(u < EPS) {
   C = A;
   return e_dist(P, A);
 if(u + EPS > 1.0) {
   C = B;
   return e_dist(P, B);
 return pointToLine(P, A, B, C);
double segmentToSegment(segment S1, segment S2) {
 if(s_intersection(S1, S2))
   return 0.0;
 double ret = INFD;
 point dummy;
 ret = min(ret, pointToSegment(S1.P, S2.P, S2.Q, dummy));
 ret = min(ret, pointToSegment(S1.Q, S2.P, S2.Q, dummy));
```

```
ret = min(ret, pointToSegment(S2.P, S1.P, S1.Q, dummy));
  ret = min(ret, pointToSegment(S2.Q, S1.P, S1.Q, dummy));
 return ret:
/*Circle*/
struct circle {
 point P;
 double r;
  circle() +
   point P1;
   P = P1;
   r = 0.0;
  circle(point _P, double _r) {
   P = P;
   r = _r;
  circle(point P1, point P2) {
   P = mid(P1, P2);
   r = e_dist(P, P1);
  circle(point P1, point P2, point P3) {
   vector<point> T;
   T.clear();
   T.pb(P1);
   T.pb(P2);
   T.pb(P3);
    sort(T.begin(), T.end());
    P1 = T[0];
    P2 = T[1];
    P3 = T[2];
    point M1, M2;
    M1 = mid(P1, P2);
    M2 = mid(P2, P3);
    point Q2, Q3;
    Q2 = rotate(P2, P1, PI / 2);
    Q3 = rotate(P3, P2, PI / 2);
    vec P1Q2(P1, Q2), P2Q3(P2, Q3);
    point M3, M4;
    M3 = translate(M1, P1Q2);
    M4 = translate(M2, P2Q3);
    line L1(M1, M3), L2(M2, M4);
   intersection(L1, L2, P);
   r = e_dist(P, P1);
 bool operator==(circle other) {
    return (P == other.P && same_d(r, other.r));
bool insideCircle(point P, circle C) {
 return e_dist(P, C.P) <= C.r + EPS;</pre>
bool c_intersection(circle C1, circle C2, point& P1, point& P2) {
 double d = e_dist(C1.P, C2.P);
 if(d > C1.r + C2.r) {
   return false; //d+EPS kalo butuh
 if(d < dabs(C1.r - C2.r) + EPS)
   return false;
  double x1 = C1.P.x, y1 = C1.P.y, r1 = C1.r, x2 = C2.P.x, y2 = C2.P.y, r2 = C2.r;
 double a = (r1 * r1 - r2 * r2 + d * d) / (2 * d), h = sqrt(r1 * r1 - a * a);
  point T(x1 + a * (x2 - x1) / d, y1 + a * (y2 - y1) / d);
  P1 = point(T.x - h * (y2 - y1) / d, T.y + h * (x2 - x1) / d);
  P2 = point(T.x + h * (y2 - y1) / d, T.y - h * (x2 - x1) / d);
 return true:
bool lc_intersection(line L, circle 0, point& P1, point& P2) {
 double a = L.a, b = L.b, c = L.c, x = 0.P.x, y = 0.P.y, r = 0.r;
  double A = a * a + b * b, B = 2 * a * b * y - 2 * a * c - 2 * b * b * x,
        C = b * b * x * x + b * b * y * y - 2 * b * c * y + c * c - b * b * r * r;
  double D = B * B - 4 * A * C;
```

point T1, T2;

**if**(same\_d(b, 0.0)) {

return false;

return true;

return true;

P1 = P2 = T1;

return true;

return false;

if(D < EPS)

D = sqrt(D);

P1 = T1;

P2 = T2;

if(!cek)

return true;

if(P1 == P2)

if(!b1)

if(!b2) P2 = P1;

return b1; if(b1 || b2) {

P1 = P2;

return true;

vector<point> T;
T.clear();

sort(T.begin(), T.end());

T.pb(A);

T.pb(B);

T.pb(C);

A = T[0];

B = T[1];

C = T[2];

vec BC(B, C);

return false;

/\*Triangle\*/

return false:

if(same\_d(D, 0.0)) { T1.x = -B / (2 \* A);

if(dabs(x - T1.x) + EPS > r)

P1 = P2 = point(T1.x, y);

P1 = point(T1.x, y - dy);

P2 = point(T1.x, y + dy);

T1.y = (c - a \* T1.x) / b;

T1.x = (-B - D) / (2 \* A);T1.y = (c - a \* T1.x) / b;

T2.x = (-B + D) / (2 \* A);

T2.y = (c - a \* T2.x) / b;

 $if(same_d(T1.x - r - x, 0.0) | | same_d(T1.x + r - x, 0.0))$ 

double dx = dabs(T1.x - x), dy = sqrt(r \* r - dx \* dx);

bool sc\_intersection(segment S, circle C, point& P1, point& P2) {

bool b1 = between\_d(P1.x, x1, x2) && between\_d(P1.y, y1, y2);

bool b2 = between\_d(P2.x, x1, x2) && between\_d(P2.y, y1, y2);

double ab = e\_dist(A, B), bc = e\_dist(B, C), ac = e\_dist(C, A);

double r = t\_area(A, B, C) / (t\_perimeter(A, B, C) / 2);

double x1 = S.P.x, y1 = S.P.y, x2 = S.Q.x, y2 = S.Q.y;

bool cek = lc\_intersection(S.L, C, P1, P2);

double t\_perimeter(point A, point B, point C) {
 return e\_dist(A, B) + e\_dist(B, C) + e\_dist(C, A);

circle t\_inCircle(point A, point B, point C) {

double ratio = e\_dist(A, B) / e\_dist(A, C);

return sqrt(s \* (s - ab) \* (s - bc) \* (s - ac));

double t\_area(point A, point B, point C) {

double s = t\_perimeter(A, B, C) / 2;

T1.x = c / a;

```
point P;
  P = translate(B, BC);
 line AP1(A, P);
  ratio = e_dist(B, A) / e_dist(B, C);
  vec AC(A, C);
  AC = scale(AC, ratio / (1 + ratio));
  P = translate(A, AC);
 line BP2(B, P);
 intersection(AP1, BP2, P);
 return circle(P, r);
circle t_outCircle(point A, point B, point C) {
 return circle(A, B, C);
/*Polygon*/
struct polygon {
 vector<point> P;
 polygon() {
   P.clear();
 polygon(vector<point>& _P) {
   P = P;
bool rayCast(point P, polygon& A) {
  point Q(P.x, 10000);
  line cast(P, Q);
  int cnt = 0;
  FOR(i, (int)(A.P.size()) - 1) {
   line temp(A.P[i], A.P[i + 1]);
    point I;
    bool B = intersection(cast, temp, I);
   if(!B)
     continue;
    else if(I == A.P[i] || I == A.P[i + 1])
   else if(pointBetween(I, A.P[i], A.P[i + 1]) && pointBetween(I, P, Q))
     cnt++;
 return cnt % 2 == 1;
// line segment p-q intersect with line A-B.
point lineIntersectSeg(point p, point q, point A, point B) {
 double a = B.y - A.y;
 double b = A.x - B.x;
 double c = B.x * A.y - A.x * B.y;
 double u = fabs(a * p.x + b * p.y + c);
 double v = fabs(a * q.x + b * q.y + c);
 return point((p.x * v + q.x * u) / (u + v), (p.y * v + q.y * u) / (u + v));
// cuts polygon Q along the line formed by point a -> point b
// (note: the last point must be the same as the first point)
vector<point> cutPolygon(point a, point b, const vector<point>& Q) {
 vector<point> P;
  for(int i = 0; i < (int)Q.size(); i++) {</pre>
    double left1 = cross(toVec(a, b), toVec(a, Q[i]));
    double left2 = 0;
   if(i != (int)Q.size() - 1)
     left2 = cross(toVec(a, b), toVec(a, Q[i + 1]));
   if(left1 > -EPS)
     P.push_back(Q[i]);
   if(left1 * left2 < -EPS)</pre>
     P.push_back(lineIntersectSeg(Q[i], Q[i + 1], a, b));
 if(!P.empty() && !(P.back() == P.front()))
    P.push_back(P.front());
  return P;
circle minCoverCircle(polygon& A) {
 vector<point> p = A.P;
```

BC = scale(BC, ratio / (1 + ratio));

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return a.first \* b.second - a.second \* b.first;

// (a-c) . (b-c)

TD dot(Pt a, Pt b, Pt c) {

TD v1 = a.first - c.first;

TD v2 = a.second - c.second;

Peko

point c; circle ret;

double cr = 0.0;

int i, j, k;
c = p[0];

```
for(i = 1; i < p.size(); i++) {</pre>
    if(e_dist(p[i], c) >= cr + EPS) {
      c = p[i], cr = 0;
      ret = circle(c, cr);
      for(j = 0; j < i; j++) {
        if(e_dist(p[j], c) >= cr + EPS) {
          c = mid(p[i], p[j]);
          cr = e_dist(p[i], c);
          ret = circle(c, cr);
          for (k = 0; k < j; k++) {
            if(e dist(p[k], c) >= cr + EPS) {
              ret = circle(p[i], p[j], p[k]);
              c = ret.P;
              cr = ret.r;
 return ret;
/*Geometry Algorithm*/
double DP[110][110];
double minCostPolygonTriangulation(polygon& A) {
 if(A.P.size() < 3)
   return 0;
  FOR(i, A.P.size()) {
    for(int j = 0, k = i; k < A.P.size(); j++, k++) {</pre>
      if(k < j + 2)
        DP[j][k] = 0.0;
      else {
        DP[j][k] = INFD;
        REP(l, j + 1, k - 1) {
          double cost = e_dist(A.P[j], A.P[k]) + e_dist(A.P[k], A.P[l]) + e_dist(A.P[l \leftrightarrow a.P[l])) + e_dist(A.P[l \leftrightarrow a.P[l]))
                         A.P[j]);
          DP[j][k] = min(DP[j][k], DP[j][l] + DP[l][k] + cost);
 return DP[0][A.P.size() - 1];
4.2 Convex Hull
typedef double TD;
                                    // for precision shits
namespace GEOM {
typedef pair<TD, TD> Pt;
                                  // vector and points
const TD EPS = 1e-9;
const TD maxD = 1e9;
TD cross(Pt a, Pt b, Pt c) {
                                  // right hand rule
 TD v1 = a.first - c.first;
                                  // (a-c) X (b-c)
 TD v2 = a.second - c.second;
 TD u1 = b.first - c.first;
 TD u2 = b.second - c.second;
 return v1 * u2 - v2 * u1;
TD cross(Pt a, Pt b) {
                                  // a X b
```

```
TD u1 = b.first - c.first;
  TD u2 = b.second - c.second;
 return v1 * u1 + v2 * u2;
TD dot(Pt a, Pt b) {
                                // a . b
 return a.first * b.first + a.second * b.second;
TD dist(Pt a, Pt b) {
 return sqrt((a.first - b.first) * (a.first - b.first) +
              (a.second - b.second) * (a.second - b.second));
TD shoelaceX2(vector<Pt>& convHull) {
 TD ret = 0;
  for(int i = 0, n = convHull.size(); i < n; i++)</pre>
    ret += cross(convHull[i], convHull[(i + 1) % n]);
 return ret;
vector<Pt> createConvexHull(vector<Pt>& points) {
  sort(points.begin(), points.end());
  vector<Pt> ret;
  for(int i = 0; i < points.size(); i++) {</pre>
    while(ret.size() > 1 &&
          cross(points[i], ret[ret.size() - 1], ret[ret.size() - 2]) < -EPS)</pre>
      ret.pop back();
    ret.push_back(points[i]);
  for(int i = points.size() - 2, sz = ret.size(); i >= 0; i--) {
   while(ret.size() > sz &&
          cross(points[i], ret[ret.size() - 1], ret[ret.size() - 2]) < -EPS)</pre>
     ret.pop_back();
   if(i == 0)
     break;
   ret.push_back(points[i]);
  return ret;
  bool isInside(Pt pv, vector<Pt>& x) { //using winding number
   int n = x.size(), wn = 0;
   x.push_back(x[0]);
    for(int i = 0; i < n; ++i) {
     if(((x[i + 1].first <= pv.first && x[i].first >= pv.first) ||
          (x[i + 1].first >= pv.first && x[i].first <= pv.first)) &&
          ((x[i + 1].second \leftarrow pv.second \& x[i].second \rightarrow pv.second) | |
           (x[i + 1].second >= pv.second && x[i].second <= pv.second))) {
        if(cross(x[i], x[i + 1], pv) == 0) {
          x.pop_back();
          return true;
    for(int i = 0; i < n; ++i) {
     if(x[i].second <= pv.second) {</pre>
       if(x[i + 1].second > pv.second && cross(x[i], x[i + 1], pv) > 0)
     else\ if(x[i+1].second <= pv.second && cross(x[i], x[i+1], pv) < 0)
    x.pop_back();
    return wn != 0;
bool isInside(Pt pv, vector<Pt>& x) { //using winding number
 int n = x.size(), wn = 0;
  x.push_back(x[0]);
  for(int i = 0; i < n; ++i) {</pre>
   if(((x[i + 1].first <= pv.first && x[i].first >= pv.first) ||
        (x[i + 1].first >= pv.first && x[i].first <= pv.first)) &&
        ((x[i + 1].second \le pv.second \& x[i].second >= pv.second) |
         (x[i + 1].second >= pv.second && x[i].second <= pv.second))) {
      if(cross(x[i], x[i + 1], pv) == 0) {
        x.pop_back();
        return true;
```

#### 4.3 Closest Pair of Points

```
#define fi first
#define se second
typedef pair<int, int> pii;
struct Point {
 int x, y, id;
int compareX(const void* a, const void* b) {
 Point* p1 = (Point*)a, *p2 = (Point*)b;
 return (p1->x - p2->x);
int compareY(const void* a, const void* b) {
 Point* p1 = (Point*)a, *p2 = (Point*)b;
 return (p1->y - p2->y);
double dist(Point p1, Point p2) {
 return sqrt((double)(p1.x - p2.x) * (p1.x - p2.x) +
              (double)(p1.y - p2.y) * (p1.y - p2.y)
             );
pair<pii, double> bruteForce(Point P[], int n) {
 double min = 1e8;
 pii ret = pii(-1, -1);
  for(int i = 0; i < n; ++i)
   for(int j = i + 1; j < n; ++j)
      if(dist(P[i], P[j]) < min) {</pre>
        ret = pii(P[i].id, P[j].id);
        min = dist(P[i], P[j]);
 return pair<pii, double> (ret, min);
pair<pii, double> getmin(pair<pii, double> x, pair<pii, double> y) {
 if(x.fi.fi == -1 && x.fi.se == -1)
   return y;
 if(y.fi.fi == -1 && y.fi.se == -1)
    return x;
 return (x.se < y.se) ? x : y;
pair<pii, double> stripClosest(Point strip[], int size, double d) {
 double min = d;
 pii ret = pii(-1, -1);
  qsort(strip, size, sizeof(Point), compareY);
  for(int i = 0; i < size; ++i)</pre>
    for(int j = i + 1; j < size && (strip[j].y - strip[i].y) < min; ++j)</pre>
      if(dist(strip[i], strip[j]) < min) {</pre>
        ret = pii(strip[i].id, strip[j].id);
        min = dist(strip[i], strip[j]);
 return pair<pii, double>(ret, min);
pair<pii, double> closestUtil(Point P[], int n) {
 if(n \le 3)
    return bruteForce(P, n);
  int mid = n / 2;
```

```
Point midPoint = P[mid];
  pair<pii, double> dl = closestUtil(P, mid);
  pair<pii, double> dr = closestUtil(P + mid, n - mid);
  pair<pii, double> d = getmin(dl, dr);
  Point strip[n];
  int j = 0;
  for(int i = 0; i < n; i++)
   if(abs(P[i].x - midPoint.x) < d.second)</pre>
     strip[j] = P[i], j++;
 return getmin(d, stripClosest(strip, j, d.second));
pair<pii, double> closest(Point P[], int n) {
 qsort(P, n, sizeof(Point), compareX);
 return closestUtil(P, n);
Point P[50005];
int main() {
 int n;
  scanf("%d", &n);
  for(int a = 0; a < n; a++) {
   scanf("%d%d", &P[a].x, &P[a].y);
   P[a].id = a;
 pair<pii, double> hasil = closest(P, n);
 if(hasil.fi.fi > hasil.fi.se)
   swap(hasil.fi.fi, hasil.fi.se);
 printf("%d %d %.6lf\n", hasil.fi.fi, hasil.fi.se, hasil.se);
 return 0:
```

# 4.4 Smallest Enclosing Circle

```
// welzl's algo to find the 2d minimum enclosing circle of a set of points
// expected O(N)
// directions: remove duplicates and shuffle points, then call welzl(points)
struct Point {
 double x;
 double y;
struct Circle {
 double x, y, r;
 Circle() {}
 Circle(double _x, double _y, double _r): x(_x), y(_y), r(_r) {}
Circle trivial(const vector<Point>& r) {
 if(r.size() == 0)
    return Circle(0, 0, -1);
  else if(r.size() == 1)
   return Circle(r[0].x, r[0].y, 0);
  else if(r.size() == 2) {
   double cx = (r[0].x + r[1].x) / 2.0, cy = (r[0].y + r[1].y) / 2.0;
    double rad = hypot(r[0].x - r[1].x, r[0].y - r[1].y) / 2.0;
    return Circle(cx, cy, rad);
  } else {
    double x0 = r[0].x, x1 = r[1].x, x2 = r[2].x;
    double y0 = r[0].y, y1 = r[1].y, y2 = r[2].y;
   double d = (x0 - x2) * (y1 - y2) - (x1 - x2) * (y0 - y2);
   double cx = (((x0 - x2) * (x0 + x2) + (y0 - y2) * (y0 + y2)) / 2 *
                 (y1 - y2) - ((x1 - x2) * (x1 + x2) + (y1 - y2) * (y1 + y2)) / 2
                 * (y0 - y2)) / d;
    double cy = (((x1 - x2) * (x1 + x2) + (y1 - y2) * (y1 + y2)) / 2 *
                 (x0 - x2) - ((x0 - x2) * (x0 + x2) + (y0 - y2) * (y0 + y2)) / 2
                 * (x1 - x2)) / d;
    return Circle(cx, cy, hypot(x0 - cx, y0 - cy));
```

Peko

```
Circle welzl(const vector<Point>& p, int idx = 0, vector<Point> r = {}) {
    if(idx == (int) p.size() || r.size() == 3)
        return trivial(r);
    Circle d = welzl(p, idx + 1, r);
    if(hypot(p[idx].x - d.x, p[idx].y - d.y) > d.r) {
        r.push_back(p[idx]);
        d = welzl(p, idx + 1, r);
    }
    return d;
}
```

# 4.5 Sutherland-Hodgman Algorithm

```
// Complexity: linear time
// Ada 2 poligon, cari poligon intersectionnya
// poly_point = hasilnya, clipper = pemotongnya
#include<bits/stdc++.h>
using namespace std;
const double EPS = 1e-9;
struct point {
 double x, y;
 point(double _x, double _y): x(_x), y(_y) {}
struct vec {
 double x, y;
 vec(double _x, double _y): x(_x), y(_y) {}
point pivot(0, 0);
vec toVec(point a, point b) {
 return vec(b.x - a.x, b.y - a.y);
double dist(point a, point b) {
 return hypot(a.x - b.x, a.y - b.y);
double cross(vec a, vec b) {
 return a.x * b.y - a.y * b.x;
bool ccw(point p, point q, point r) {
 return cross(toVec(p, q), toVec(p, r)) > 0;
bool collinear(point p, point q, point r) {
 return fabs(cross(toVec(p, q), toVec(p, r))) < EPS;</pre>
bool lies(point a, point b, point c) {
 if((c.x >= min(a.x, b.x) \&\& c.x <= max(a.x, b.x)) \&\&
      (c.y >= min(a.y, b.y) \&\& c.y <= max(a.y, b.y)))
    return true;
  else
    return false;
bool anglecmp(point a, point b) {
 if(collinear(pivot, a, b))
    return dist(pivot, a) < dist(pivot, b);</pre>
  double d1x = a.x - pivot.x, d1y = a.y - pivot.y;
  double d2x = b.x - pivot.x, d2y = b.y - pivot.y;
 return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0;
point intersect(point s1, point e1, point s2, point e2) {
  double x1, x2, x3, x4, y1, y2, y3, y4;
 x1 = s1.x;
 y1 = s1.y;
  x2 = e1.x;
 y2 = e1.y;
  x3 = s2.x;
  y3 = s2.y;
  x4 = e2.x;
```

```
y4 = e2.y;
  double num1 = (x1 * y2 - y1 * x2) * (x3 - x4) - (x1 - x2) * (x3 * y4 - y3 * x4);
  double num2 = (x1 * y2 - y1 * x2) * (y3 - y4) - (y1 - y2) * (x3 * y4 - y3 * x4);
  double den = (x1 - x2) * (y3 - y4) - (y1 - y2) * (x3 - x4);
  double new x = num1 / den;
  double new_y = num2 / den;
 return point(new_x, new_y);
void clip(vector <point>& poly_points, point point1, point point2) {
 vector <point> new_points;
  new_points.clear();
  for(int i = 0; i < poly_points.size(); i++) {</pre>
    int k = (i + 1) % poly_points.size();
    double i_pos = ccw(point1, point2, poly_points[i]);
    double k pos = ccw(point1, point2, poly points[k]);
    //in in
    if(i_pos <= 0 && k_pos <= 0)
     new_points.push_back(poly_points[k]);
    else if(i pos > 0 && k pos <= 0) {
     new_points.push_back(intersect(point1, point2, poly_points[i],
                                     polv points[k]));
     new_points.push_back(poly_points[k]);
    else if(i pos <= 0 && k pos > 0) {
     new_points.push_back(intersect(point1, point2, poly_points[i],
                                     poly_points[k]));
    //out out
    else {
  poly_points.clear();
  for(int i = 0; i < new_points.size(); i++)</pre>
   poly_points.push_back(new_points[i]);
double area(const vector <point>& P) {
 double result = 0.0;
 double x1, y1, x2, y2;
  for(int i = 0; i < P.size() - 1; i++) {</pre>
   x1 = P[i].x;
   y1 = P[i].y;
   x2 = P[i + 1].x;
   y2 = P[i + 1].y;
   result += (x1 * y2 - x2 * y1);
 return fabs(result) / 2;
void suthHodgClip(vector <point>& poly_points, vector <point> clipper_points) {
 for(int i = 0; i < clipper_points.size(); i++) {</pre>
   int k = (i + 1) % clipper_points.size();
    clip(poly_points, clipper_points[i], clipper_points[k]);
vector<point> sortku(vector<point> P) {
 int P0 = 0;
  int i;
  for(i = 1; i < 3; i++) {
   if(P[i].y < P[P0].y \mid | (P[i].y == P[P0].y && P[i].x > P[P0].x))
     P0 = i;
  point temp = P[0];
  P[0] = P[P0];
  P[P0] = temp;
  pivot = P[0];
  sort(++P.begin(), P.end(), anglecmp);
  reverse(++P.begin(), P.end());
  return P;
```

```
int main {
  clipper_points = sortku(clipper_points);
  suthHodgClip(poly_points, clipper_points);
}
```

# 4.6 Centroid of Polygon

```
C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i \ y_{i+1} - x_{i+1} \ y_i)
C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i \ y_{i+1} - x_{i+1} \ y_i)
```

#### 4.7 Pick Theorem

A: Area of a simply closed lattice polygon B: Number of lattice points on the edges I: Number of points in the interior  $A = I + \frac{B}{2} - 1$ 

# 5 Graphs

# 5.1 Articulation Point and Bridge

```
// gr -> adi list
// vector vis, low -> initialize to -1
// int timer -> initialize to 0
void dfs(int pos, int dad = -1) {
 vis[pos] = low[pos] = timer++;
 int kids = 0;
 for(auto& i : gr[pos]) {
   if(i == dad)
     continue:
   if(vis[i] >= 0)
     low[pos] = min(low[pos], vis[i]);
    else {
     dfs(i, pos);
     low[pos] = min(low[pos], low[i]);
     if(low[i] > vis[pos])
       is_bridge(pos, i)
       if(low[i] >= vis[pos] && dad >= 0)
         is_articulation_point(pos)
         ++kids;
 if(dad == -1 && kids > 1)
   is_articulation_point(pos)
```

# 5.2 SCC and Strong Orientation

```
#define N 10020
vector<int> adj[N];
bool vis[N], ins[N];
int disc[N], low[N], gr[N];
stack<int> st;
int id, grid;
void scc(int cur, int par) {
 disc[cur] = low[cur] = ++id;
 vis[cur] = ins[cur] = 1;
  st.push(cur);
  for(int to : adj[cur]) {
    //if (to==par) continue; // ini untuk SO(scc undirected)
    if(!vis[to])
      scc(to, cur);
    if(ins[to])
      low[cur] = min(low[cur], low[to]);
  if(low[cur] == disc[cur]) {
```

```
grid++; // group id
while(ins[cur]) {
    gr[st.tp] = grid;
    ins[st.tp] = 0;
    st.pop();
    }
}
```

# 5.3 Centroid Decomposition

```
int build_cen(int nw) {
  com_cen(nw, 0); //fungsi untuk itung size subtree
 int siz = sz[nw] / 2;
  bool found = false;
 while(!found) {
   found = true;
    for(int i : v[nw]) {
     if(!rem[i] && sz[i] < sz[nw]) {</pre>
       if(sz[i] > siz) {
          found = false;
          nw = i:
          break;
 big
  rem[nw] = true;
  for(int i : v[nw])if(!rem[i])
     par_cen[build_cen(i)] = nw;
 return nw;
```

#### 5.4 Dinic's Maximum Flow

```
// O(VE log(max_flow)) if scaling == 1
// O((V + E) sqrt(E)) if unit graph (turn scaling off)
// O((V + E) sqrt(V)) if bipartite matching (turn scaling off)
// indices are 0-based
const ll INF = 1e18;
struct Dinic {
  struct Edge {
    int v;
    ll cap, flow;
    Edge(int _v, ll _cap): v(_v), cap(_cap), flow(0) {}
  };
  int n;
  ll lim;
  vector<vector<int>> gr;
  vector<Edge> e;
  vector<int> idx, lv;
  bool has_path(int s, int t) {
    queue<int> q;
    q.push(s);
    lv.assign(n, -1);
    lv[s] = 0;
    while(!q.empty()) {
      int c = q.front();
      q.pop();
      if(c == t)
        break;
      for(auto& i : gr[c]) {
        ll cur_flow = e[i].cap - e[i].flow;
        if(lv[e[i].v] == -1 && cur_flow >= lim) {
          lv[e[i].v] = lv[c] + 1;
```

```
q.push(e[i].v);
 return lv[t] != -1;
ll get_flow(int s, int t, ll left) {
  if(!left || s == t)
    return left;
  while(idx[s] < (int) gr[s].size()) {</pre>
    int i = gr[s][idx[s]];
    if(lv[e[i].v] == lv[s] + 1) {
      ll add = get_flow(e[i].v, t, min(left, e[i].cap - e[i].flow));
        e[i].flow += add;
        e[i ^ 1].flow -= add;
        return add;
    ++idx[s];
  return 0;
Dinic(int vertices, bool scaling = 1) : // toggle scaling here
  n(vertices), lim(scaling ? 1 << 30 : 1), <math>gr(n) {}
void add_edge(int from, int to, ll cap, bool directed = 1) {
  gr[from].push_back(e.size());
  e.emplace_back(to, cap);
  gr[to].push_back(e.size());
  e.emplace_back(from, directed ? 0 : cap);
ll get_max_flow(int s, int t) { // call this
  ll res = 0;
  while(lim) { // scaling
    while(has_path(s, t)) {
      idx.assign(n, 0);
      while(ll add = get_flow(s, t, INF))
        res += add;
    lim >>= 1;
 }
  return res;
```

### 5.5 Minimum Cost Maximum Flow

```
// 1-based index
template<class T>
using rpq = priority_queue<T, vector<T>, greater<T>>;
const ll INF = 1e18;
struct MCMF {
 struct Edge {
    int v;
    ll cap, cost;
    int rev;
    Edge(int _v, ll _cap, ll _cost, int _rev) :
      v(_v), cap(_cap), cost(_cost), rev(_rev) {}
  ll flow, cost;
 int st, ed, n;
 vector<ll> dist, H;
 vector<int> pv, pe;
```

```
vector<vector<Edge>> adj;
  bool dijkstra() {
   rpq<pair<ll, int>> pq;
   dist.assign(n + 1, INF);
   dist[st] = 0;
    pq.emplace(0, st);
    while(!pq.empty()) {
     auto [cst, pos] = pq.top();
     pq.pop();
     if(dist[pos] < cst)</pre>
       continue;
      for(int i = 0; i < (int) adj[pos].size(); ++i) {</pre>
       auto& e = adj[pos][i];
       int nxt = e.v;
       ll nxt cst = dist[pos] + e.cost + H[pos] - H[nxt];
        if(e.cap > 0 && nxt_cst < dist[nxt]) {</pre>
          dist[nxt] = nxt_cst;
          pe[nxt] = i;
          pv[nxt] = pos;
          pq.emplace(nxt_cst, nxt);
   return dist[ed] != INF;
  MCMF(int _n) : n(_n), pv(n + 1), pe(n + 1), adj(n + 1) {}
  void add_edge(int u, int v, ll cap, ll cst) {
   adj[u].emplace_back(v, cap, cst, adj[v].size());
   adj[v].emplace_back(u, 0, -cst, adj[u].size() - 1);
  pair<ll, ll> solve(int _st, int _ed) {
   st = _st, ed = _ed;
    flow = 0, cost = 0;
   H.assign(n + 1, 0);
    while(dijkstra()) {
     for(int i = 0; i <= n; ++i)
       H[i] += dist[i];
     ll f = INF;
      for(int i = ed; i != st; i = pv[i])
       f = min(f, adj[pv[i]][pe[i]].cap);
      flow += f;
     cost += f * H[ed];
      for(int i = ed; i != st; i = pv[i]) {
       auto& e = adj[pv[i]][pe[i]];
       e.cap -= f;
        adj[i][e.rev].cap += f;
   return {flow, cost};
};
```

#### 5.6 Flows with Demands

```
let S0 be the source and T0 be the original sink
1. add 2 additional nodes, call them S1 and T1
2. connect S0 to nodes normally
3. connect nodes to TO normally
4. for each edge(U, V), cap = original cap - demand
5. for each node N:
  1. add an edge(S1, N), cap = sum of inward demand to N
   2. add an edge(N, T1), cap = sum of outward demand from N
6. add an edge(T0, S0), cap = INF
7. the above is not a typo!
8. run max flow normally
9. for each edge(S1, V) and (U, T1), check if flow == cap
```

if step #9 fails, then it is not possible to satisfy the given demand

Mathematically, let d(e) be the demand of edge e. Let V be the set of every vertex in the graph.

- $c'(S_1, v) = \sum_{u \in V} d(u, v)$  for each edge (s', v).
- $c'(v,T_1) = \sum_{v \in V} d(v,w)$  for each edge (v,t').
- c'(u,v) = c(u,v) d(u,v) for each edge (u,v) in the old network.
- $c'(T_0, S_0) = \infty$

# 5.7 Hungarian

```
template <typename TD> struct Hungarian {
 TD INF = 1e9; //max_inf
 vector<vector<TD> > adj; // cost[left][right]
 vector<TD> hl, hr, slk;
 vector<int> fl, fr, vl, vr, pre;
 deque<int> q;
 Hungarian(int _n) {
   n = n;
   adj = vector<vector<TD> >(n, vector<TD> (n, 0));
 int check(int i) {
   if(vl[i] = 1, fl[i] != -1)
     return q.push_back(fl[i]), vr[fl[i]] = 1;
     swap(i, fr[fl[i] = pre[i]]);
   return 0;
 void bfs(int s) {
   slk.assign(n, INF);
   vl.assign(n, 0);
   vr = vl;
   q.assign(vr[s] = 1, s);
   for(TD d;;) {
     for(; !q.empty(); q.pop_front()) {
       for(int i = 0, j = q.front(); i < n; i++) {</pre>
         if(d = hl[i] + hr[j] - adj[i][j], !vl[i] && d <= slk[i]) {</pre>
           if(pre[i] = j, d)
             slk[i] = d:
           else if(!check(i))
             return;
     d = INF;
     for(int i = 0; i < n; i++) if(!vl[i] && d > slk[i])
         d = slk[i];
     for(int i = 0; i < n; i++) {
       if(vl[i])
         hl[i] += d;
       else
         slk[i] -= d;
       if(vr[i])
         hr[i] -= d;
     for(int i = 0; i < n; i++) if(!vl[i] && !slk[i] && !check(i))</pre>
 TD solve() {
   fl.assign(n, -1);
   fr = fl;
   hl.assign(n, 0);
   hr = hl:
   pre.assign(n, 0);
   for(int i = 0; i < n; i++)
```

```
hl[i] = *max_element(adj[i].begin(), adj[i].begin() + n);
    for(int i = 0; i < n; i++)
     bfs(i);
    TD ret = 0;
    for(int i = 0; i < n; i++) if(adj[i][fl[i]])</pre>
        ret += adj[i][fl[i]];
    return ret;
}; //i will be matched with fl[i]
```

#### 5.8 Edmonds' Blossom

```
// Maximum matching on general graphs in O(V^2 E)
// Indices are 1-based
// Stolen from ko osaga's cheatsheet
struct Blossom {
  vector<int> vis, dad, orig, match, aux;
  vector<vector<int>> conn;
  int t, N;
  queue<int> 0;
  void augment(int u, int v) {
    int pv = v;
    do {
      pv = dad[v];
      int nv = match[pv];
      match[v] = pv;
      match[pv] = v;
      v = nv;
    } while(u != pv);
  int lca(int v, int w) {
    while(true) {
      if(v) {
        if(aux[v] == t)
          return v;
        aux[v] = t;
        v = orig[dad[match[v]]];
      swap(v, w);
  void blossom(int v, int w, int a) {
    while(orig[v] != a) {
      dad[v] = w;
      w = match[v];
      if(vis[w] == 1) {
        Q.push(w);
        vis[w] = 0;
      orig[v] = orig[w] = a;
      v = dad[w];
  bool bfs(int u) {
    fill(vis.begin(), vis.end(), -1);
    iota(orig.begin(), orig.end(), 0);
    Q = queue<int>();
    Q.push(u);
    vis[u] = 0;
    while(!Q.empty()) {
      int v = Q.front();
      Q.pop();
      for(int x : conn[v]) {
        if(vis[x] == -1) {
          dad[x] = v;
```

vis[x] = 1;

if(!match[x]) {

return 1;

augment(u, x);

Q.push(match[x]);

blossom(x, v, a);blossom(v, x, a);

Blossom(int n) : // n = vertices

for(int i = 0; i <= n; ++i) {

conn[i].clear();

void add\_edge(int u, int v) {

conn[u].push\_back(v);

conn[v].push\_back(u);

aux(n + 1), conn(n + 1), t(0), N(n) {

match[i] = aux[i] = dad[i] = 0;

int solve() { // call this for answer

return false;

vis[match[x]] = 0;

} else if(vis[x] == 0 && orig[v] != orig[x]) {

vis(n + 1), dad(n + 1), orig(n + 1), match(n + 1),

int a = lca(orig[v], orig[x]);

```
int ans = 0;
    vector<int> V(N - 1);
    iota(V.begin(), V.end(), 1);
    shuffle(V.begin(), V.end(), mt19937(0x94949));
    for(auto x : V) {
      if(!match[x]) {
       for(auto y : conn[x]) {
         if(!match[y]) {
           match[x] = y, match[y] = x;
            ++ans;
           break;
    for(int i = 1; i <= N; ++i) {
      if(!match[i] && bfs(i))
       ++ans;
   return ans;
};
5.9 Eulerian Path or Cycle
// else call on any vertex
// ans = path result
```

```
// finds a eulerian path / cycle
// visits each edge only once
// properties:
// - cycle: degrees are even
// - path: degrees are even OR degrees are even except for 2 vertices
// how to use: g = adjacency list g[n] = connected to n, undirected
// if there is a vertex u with an odd degree, call dfs(u)
vector<set<int>> g;
vector<int> ans;
```

```
void dfs(int u) {
  while(g[u].size()) {
    int v = *g[u].begin();
    g[u].erase(v);
    g[v].erase(u);
    dfs(v);
  ans.push_back(u);
```

# 5.10 Hierholzer's Algorithm

```
// Eulerian on Directed Graph
stack<int> path;
vector<int> euler;
inline void hierholzer() {
  path.push(0);
  int cur = 0;
  while(!path.empty()) {
   if(!adi[cur].empty()) {
      path.push(cur);
      int next = adj[cur].back();
      adj[cur].pob();
      cur = next;
   } else {
      euler.pb(cur);
      cur = path.top();
      path.pop();
 reverse(euler.begin(), euler.end());
```

#### 5.11 2-SAT

```
struct TwoSAT {
 int n;
 vector<vector<int>> g, gr;
 vector<int> comp, topological_order, answer;
 vector<bool> vis;
 TwoSAT() {}
 TwoSAT(int _n) :
   n(n), g(2 * n), g(2 * n), comp(2 * n), answer(2 * n), vis(2 * n) {}
  void add_edge(int u, int v) {
   g[u].push_back(v);
   gr[v].push_back(u);
  // For the following three functions
  // int x, bool val: if 'val' is true, we take the variable to be x.
  // Otherwise we take it to be x's complement.
  // At least one of them is true
  void add_clause_or(int i, bool f, int j, bool p) {
   add_edge(i + (f ? n : 0), j + (p ? 0 : n));
   add_edge(j + (p ? n : 0), i + (f ? 0 : n));
  // Only one of them is true
  void add_clause_xor(int i, bool f, int j, bool p) {
   add_clause_or(i, f, j, p);
   add_clause_or(i, !f, j, !p);
  // Both of them have the same value
  void add_clause_and(int i, bool f, int j, bool p) {
   add_clause_xor(i, !f, j, p);
```

```
// Topological sort
void dfs(int u) {
  vis[u] = true;
  for(const auto& v : g[u])
    if(!vis[v])
      dfs(v);
  topological_order.push_back(u);
// Extracting strongly connected components
void scc(int u, int id) {
  vis[u] = true;
  comp[u] = id;
  for(const auto& v : gr[u])
    if(!vis[v])
      scc(v, id);
bool satisfiable() {
  fill(vis.begin(), vis.end(), false);
  for(int i = 0; i < 2 * n; i++)
    if(!vis[i])
      dfs(i);
  fill(vis.begin(), vis.end(), false);
  reverse(topological_order.begin(), topological_order.end());
  for(const auto& v : topological_order)
    if(!vis[v])
      scc(v, id++);
  // Constructing the answer
  for(int i = 0: i < n: i++) {
    if(comp[i] == comp[i + n])
      return false;
    answer[i] = (comp[i] > comp[i + n] ? 1 : 0);
  return true;
```

#### 6 Math

## 6.1 Extended Euclidean GCD

```
// computes x and y such that ax + by = gcd(a, b) in O(log (min(a, b)))
// returns {gcd(a, b), x, y}
tuple<int, int, int> gcd(int a, int b) {
   if(b == 0) return {a, 1, 0};
   auto [d, x1, y1] = gcd(b, a % b);
   return {d, y1, x1 - y1* (a / b)};
}
```

### 6.2 Generalized CRT

```
template<typename T>
T extended_euclid(T a, T b, T& x, T& y) {
    if(b == 0) {
        x = 1;
        y = 0;
        return a;
}
T xx, yy, gcd;
gcd = extended_euclid(b, a % b, xx, yy);
x = yy;
y = xx - (yy * (a / b));
return gcd;
```

```
template<typename T>
T MOD(T a, T b) {
 return (a % b + b) % b;
// return x, lcm. x = a % n && x = b % m
template<typename T>
pair<T, T> CRT(T a, T n, T b, T m) {
 T _n, _m;
 T gcd = extended_euclid(n, m, _n, _m);
  if(n == m) {
    if(a == b)
      return pair<T, T>(a, n);
    else
      return pair<T, T>(-1, -1);
  } else if(abs(a - b) % gcd != 0)
    return pair<T, T>(-1, -1);
    T lcm = m * n / gcd;
    T \times = MOD(a + MOD(n \times MOD(n \times ((b - a) / gcd), m / gcd), lcm), lcm);
    return pair<T, T>(x, lcm);
```

## 6.3 Generalized Lucas Theorem

```
/*Special Lucas : (n,k) % p^x
 fctp[n] = Product of the integers less than or equal
 to n that are not divisible by p
 Precompute fctp*/
LL p
LL E(LL n, int m) {
 LL tot = 0;
 while(n != 0)
   tot += n / m, n /= m;
 return tot;
LL funct(LL n, LL base) {
 LL ans = fast(fctp[base], n / base, base) * fctp[n % base] % base;
 return ans;
LL F(LL n, LL base) {
 LL ans = 1:
 while(n != 0) {
   ans = (ans * funct(n, base)) % base;
   n /= p;
 return ans;
LL special_lucas(LL n, LL r, LL base) {
 p = fprime(base);
 LL pow = E(n, p) - E(n - r, p) - E(r, p);
 LL TOP = fast(p, pow, base) * F(n, base) % base;
 LL BOT = F(r, base) * F(n - r, base) % base;
 return (TOP * fast(BOT, totien(base) - 1, base)) % base;
//End of Special Lucas
```

# 6.4 Linear Diophantine

```
//FOR SOLVING MINIMUM ABS(X) + ABS(Y)
ll x, y, newX, newY, target = 0;
ll extGcd(ll a, ll b) {
   if(b == 0) {
      x = 1, y = 0;
      return a;
   }
ll ret = extGcd(b, a % b);
newX = y;
```

```
newY = x - y * (a / b);
 x = newX;
 y = newY;
 return ret;
ll fix(ll sol, ll rt) {
 ll ret = 0;
 //CASE SOLUTION(X/Y) < TARGET
 if(sol < target)</pre>
   ret = -floor(abs(sol + target) / (double)rt);
 //CASE SOLUTION(X/Y) > TARGET
 if(sol > target)
   ret = ceil(abs(sol - target) / (double)rt);
 return ret;
ll work(ll a, ll b, ll c) {
 ll gcd = extGcd(a, b);
 ll\ solX = x * (c / gcd);
 ll solY = y * (c / gcd);
 a /= gcd;
 b /= gcd;
 ll fi = abs(fix(solX, b));
 ll se = abs(fix(solY, a));
 ll lo = min(fi, se);
 ll hi = max(fi, se);
 ll ans = abs(solX) + abs(solY);
 for(ll i = lo; i <= hi; i++) {
   ans = min(ans, abs(solX + i * b) + abs(solY - i * a));
   ans = min(ans, abs(solX - i * b) + abs(solY + i * a));
 return ans;
```

# 6.5 Modular Linear Equation

```
// finds all solutions to ax = b (mod n)
vi modular_linear_equation_solver(int a, int b, int n) {
   int x, y;
   vi ret;
   int g = extended_euclid(a, n, x, y);
   if(!(b % g)) {
      x = mod(x * (b / g), n);
      for(int i = 0; i < g; i++)
         ret.push_back(mod(x + i * (n / g), n));
   }
   return ret;
}</pre>
```

## 6.6 Miller-Rabin and Pollard's Rho

```
namespace MillerRabin {
const vector<ll> primes = { // deterministic up to 2^64 - 1
 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37
ll gcd(ll a, ll b) {
 return b ? gcd(b, a % b) : a;
ll powa(ll x, ll y, ll p) { // (x ^ y) % p
 if(!y)
    return 1;
 if(y & 1)
    return ((\_int128) \times powa(x, y - 1, p)) \% p;
 ll temp = powa(x, y >> 1, p);
  return ((__int128) temp * temp) % p;
bool miller_rabin(ll n, ll a, ll d, int s) {
 ll x = powa(a, d, n);
 if(x == 1 || x == n - 1)
    return 0;
```

```
for(int i = 0; i < s; ++i) {
   x = ((\_int128) x * x) % n;
   if(x == n - 1)
      return 0;
 return 1;
bool is_prime(ll x) { // use this
 if(x < 2)
    return 0;
  int r = 0;
 ll d = x - 1;
  while((d & 1) == 0) {
   d >>= 1;
    ++r;
  for(auto& i : primes) {
   if(x == i)
     return 1;
   if(miller_rabin(x, i, d, r))
     return 0;
 return 1;
namespace PollardRho {
mt19937_64 generator(chrono::steady_clock::now()
                     .time_since_epoch().count());
uniform_int_distribution<ll> rand_ll(0, LLONG_MAX);
ll f(ll x, ll b, ll n) { // (x^2 + b) % n}
 return (((__int128) x * x) % n + b) % n;
ll rho(ll n) {
 if(n % 2 == 0)
   return 2;
  ll b = rand ll(generator);
  ll x = rand_ll(generator);
  ll y = x;
  while(1) {
   x = f(x, b, n);
   y = f(f(y, b, n), b, n);
    ll d = MillerRabin::gcd(abs(x - y), n);
    if(d != 1)
     return d;
void pollard_rho(ll n, vector<ll>& res) {
 if(n == 1)
    return;
  if(MillerRabin::is_prime(n)) {
    res.push_back(n);
    return;
 ll d = rho(n);
  pollard_rho(d, res);
  pollard_rho(n / d, res);
vector<ll> factorize(ll n, bool sorted = 1) { // use this
  vector<ll> res;
  pollard_rho(n, res);
  if(sorted)
    sort(res.begin(), res.end());
  return res;
```

# 6.7 Berlekamp-Massey

#include <bits/stdc++.h>

using namespace std;

```
#define pb push_back
typedef long long ll;
#define SZ 233333
const int MOD = 1e9 + 7; //or any prime
ll qp(ll a, ll b) {
 ll x = 1;
 a %= MOD;
  while(b) {
    if(b & 1)
     x = x * a % MOD;
    a = a * a % MOD;
    b >>= 1;
 return x;
namespace linear_seq {
vector<int> BM(vector<int> x) {
 //ls: (shortest) relation sequence (after filling zeroes) so far
  //cur: current relation sequence
  vector<int> ls, cur;
  //lf: the position of ls (t')
  //ld: delta of ls (v')
  int lf = -1, ld = -1;
  for(int i = 0; i < int(x.size()); ++i) {</pre>
    ll t = 0;
    //evaluate at position i
    for(int j = 0; j < int(cur.size()); ++j)</pre>
      t = (t + x[i - j - 1] * (ll)cur[j]) % MOD;
    if((t - x[i]) \% MOD == 0) {
      continue; //good so far
    //first non-zero position
    if(!cur.size()) {
      cur.resize(i + 1);
     lf = i;
      ld = (t - x[i]) % MOD;
      continue;
    //cur=cur-c/ld*(x[i]-t)
    ll k = -(x[i] - t) * qp(ld, MOD - 2) % MOD/*1/ld*/;
    vector<int> c(i - lf - 1); //add zeroes in front
    c.pb(k);
    for(int j = 0; j < int(ls.size()); ++j)</pre>
      c.pb(-ls[j]*k % MOD);
    if(c.size() < cur.size())</pre>
     c.resize(cur.size());
    for(int j = 0; j < int(cur.size()); ++j)</pre>
     c[j] = (c[j] + cur[j]) % MOD;
    //if cur is better than ls, change ls to cur
    if(i - lf + (int)ls.size() >= (int)cur.size())
      ls = cur, lf = i, ld = (t - x[i]) % MOD;
    cur = c;
  for(int i = 0; i < int(cur.size()); ++i)</pre>
   cur[i] = (cur[i] % MOD + MOD) % MOD;
  return cur;
int m; //length of recurrence
//a: first terms
//h: relation
ll a[SZ], h[SZ], t_[SZ], s[SZ], t[SZ];
//calculate p*q mod f
void mull(ll* p, ll* q) {
 for(int i = 0; i < m + m; ++i)
    t_{[i]} = 0;
  for(int i = 0; i < m; ++i) if(p[i])</pre>
      for(int j = 0; j < m; ++j)
        t_{i} = (t_{i} + j) + p[i] * q[j]) % MOD;
  for(int i = m + m - 1; i >= m; --i) if(t_[i])
      //miuns t_{i}x^{i-m}(x^m-\sum_{j=0}^{m-1} x^{m-j-1}h_{j})
```

```
for(int j = m - 1; ~j; --j)
        t_{i} - j - 1 = (t_{i} - j - 1) + t_{i} * h_{i} % MOD;
  for(int i = 0; i < m; ++i)
    p[i] = t_[i];
ll calc(ll K) {
  for(int i = m; ~i; --i)
   s[i] = t[i] = 0;
  //init
  s[0] = 1;
  if(m != 1)
   t[1] = 1;
  else
   t[0] = h[0];
  //binary-exponentiation
  while(K) {
   if(K & 1)
      mull(s, t);
    mull(t, t);
    K >>= 1;
  ll su = 0;
  for(int i = 0; i < m; ++i)</pre>
    su = (su + s[i] * a[i]) % MOD;
  return (su % MOD + MOD) % MOD;
int work(vector<int> x, ll n) {
 if(n < int(x.size()))</pre>
   return x[n];
  vector < int > v = BM(x);
  m = v.size();
  if(!m)
   return 0;
  for(int i = 0; i < m; ++i)
   h[i] = v[i], a[i] = x[i];
  return calc(n);
using linear_seq::work;
const vector<int> sequence = {
 0, 2, 2, 28, 60, 836, 2766
int main() {
 cout << work(sequence, 7) << '\n';</pre>
```

#### 6.8 Fast Fourier Transform

```
using ld = double; // change to long double if reach 10^18
using cd = complex<ld>;
const ld PI = acos(-(ld)1);
void fft(vector<cd>& a, int sign = 1) {
 int n = a.size();
  ld theta = sign * 2 * PI / n;
  for(int i = 0, j = 1; j < n - 1; j++) {
    for(int k = n >> 1; k > (i ^= k); k >>= 1);
    if(j < i)
      swap(a[i], a[j]);
  for(int m, mh = 1; (m = mh << 1) <= n; mh = m) {
    int irev = 0;
    for(int i = 0; i < n; i += m) {
      cd w = exp(cd(0, theta * irev));
      for(int k = n >> 2; k > (irev ^= k); k >>= 1);
      for(int j = i; j < mh + i; j++) {
        int k = j + mh;
        cd x = a[j] - a[k];
        a[j] += a[k];
```

namespace FFT {

```
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```

```
a[k] = w * x;
 if(sign == -1) for(cd& i : a)
     i /= n;
vector<ll> multiply(vector<ll> const& a, vector<ll> const& b) {
 vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
 int n = 1;
 while(n < a.size() + b.size())</pre>
   n <<= 1;
 fa.resize(n);
 fb.resize(n);
 fft(fa);
 fft(fb);
 for(int i = 0; i < n; i++)
   fa[i] *= fb[i];
 fft(fa, -1);
 vector<ll> res(n);
 for(int i = 0; i < n; i++)
   res[i] = round(fa[i].real());
 return res;
```

#### 6.9 Number Theoretic Transform

```
/* ---- Adjust the constants here ---- */
const int LN = 24; //23
const int N = 1 << LN:
typedef long long LL; // 2**23 * 119 + 1. 998244353
// `MOD` must be of the form 2**`LN` * k + 1, where k odd.
const LL MOD = 9223372036737335297; // 2**24 * 54975513881 + 1.
const LL PRIMITIVE_ROOT = 3; // Primitive root modulo `MOD`.
/* ---- End of constants ---- */
LL root[N];
inline LL power(LL x, LL y) {
 LL ret = 1;
 for(; y; y >>= 1) {
   if(y & 1)
     ret = (__int128) ret * x % MOD;
   x = (_int128) x * x % MOD;
 return ret;
inline void init_fft() {
 const LL UNITY = power(PRIMITIVE_ROOT, MOD - 1 >> LN);
 root[0] = 1;
 for(int i = 1; i < N; i++)
   root[i] = (__int128) UNITY * root[i - 1] % MOD;
 return;
// n = 2^{k} is the length of polynom
inline void fft(int n, vector<LL>& a, bool invert) {
 for(int i = 1, j = 0; i < n; ++i) {
   int bit = n >> 1;
   for(; j >= bit; bit >>= 1)
     j -= bit;
    i += bit;
    if(i < j)
      swap(a[i], a[j]);
  for(int len = 2; len <= n; len <<= 1) {
   LL wlen = (invert ? root[N - N / len] : root[N / len]);
    for(int i = 0; i < n; i += len) {</pre>
     LL w = 1:
      for(int j = 0; j<len >> 1; j++) {
       LL u = a[i + j];
```

```
LL v = (\_int128) a[i + j + len / 2] * w % MOD;
        a[i + j] = ((\_int128) u + v) % MOD;
       a[i + j + len / 2] = ((\_int128) u - v + MOD) % MOD;
       w = (__int128) w * wlen % MOD;
 if(invert) {
   LL inv = power(n, MOD - 2);
    for(int i = 0; i < n; i++)
     a[i] = (__int128) a[i] * inv % MOD;
 return;
inline vector<LL> multiply(vector<LL> a, vector<LL> b) {
 vector<LL> c;
 int len = 1 << 32 - __builtin_clz(a.size() + b.size() - 2);</pre>
 a.resize(len, 0);
 b.resize(len, 0);
  fft(len, a, false);
  fft(len, b, false);
  c.resize(len);
  for(int i = 0; i < len; ++i)
   c[i] = (__int128) a[i] * b[i] % MOD;
  fft(len, c, true);
 return c;
//FFT::init_fft(); wajib di panggil init di awal
```

#### 6.10 Gauss-Jordan

```
// Gauss-Jordan elimination with full pivoting.
//
// Uses:
   (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT:
             a[][] = an nxn matrix
//
             b[][] = an nxm matrix
//
// OUTPUT:
                    = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT& a, VVT& b) {
 const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T det = 1;
  for(int i = 0; i < n; i++) {
    int pj = -1, pk = -1;
    for(int j = 0; j < n; j++) if(!ipiv[j])</pre>
        for(int k = 0; k < n; k++) if(!ipiv[k])</pre>
            if(pj == -1 \mid | fabs(a[j][k]) > fabs(a[pj][pk])) {
              pj = j;
              pk = k;
    if(fabs(a[pj][pk]) < EPS) {</pre>
      cerr << "Matrix is singular." << endl;</pre>
      exit(0);
```

ipiv[pk]++;

```
swap(a[pj], a[pk]);
   swap(b[pj], b[pk]);
   if(pj != pk)
     det *= −1;
   irow[i] = pj;
   icol[i] = pk;
   T c = 1.0 / a[pk][pk];
   det *= a[pk][pk];
   a[pk][pk] = 1.0;
   for(int p = 0; p < n; p++)
     a[pk][p] *= c;
   for(int p = 0; p < m; p++)
     b[pk][p] *= c;
   for(int p = 0; p < n; p++) if(p != pk) {
       c = a[p][pk];
       a[p][pk] = 0;
       for(int q = 0; q < n; q++)
         a[p][q] -= a[pk][q] * c;
       for(int q = 0; q < m; q++)
         b[p][q] -= b[pk][q] * c;
 for(int p = n - 1; p >= 0; p--) if(irow[p] != icol[p]) {
     for(int k = 0; k < n; k++)
       swap(a[k][irow[p]], a[k][icol[p]]);
 return det;
int main() {
 const int n = 4;
 const int m = 2;
 double A[n][n] = { {1, 2, 3, 4}, {1, 0, 1, 0}, {5, 3, 2, 4}, {6, 1, 4, 6} };
 double B[n][m] = { {1, 2}, {4, 3}, {5, 6}, {8, 7} };
 VVT a(n), b(n);
 for(int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
 double det = GaussJordan(a, b);
 // expected: 60
 cout << "Determinant: " << det << endl;</pre>
 // expected: -0.233333 0.166667 0.133333 0.0666667
              0.166667 0.166667 0.333333 -0.333333
              0.233333 0.833333 -0.133333 -0.0666667
 //
              0.05 -0.75 -0.1 0.2
 cout << "Inverse: " << endl;</pre>
 for(int i = 0; i < n; i++) {
   for(int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
   cout << endl;
 // expected: 1.63333 1.3
              -0.166667 0.5
              2.36667 1.7
              -1.85 -1.35
 cout << "Solution: " << endl;</pre>
 for(int i = 0; i < n; i++) {
   for(int j = 0; j < m; j++)
     cout << b[i][j] << ' ';
   cout << endl;</pre>
```

# 6.11 Fibonacci Check

```
bool is_fibonacci(int n) {
 return is_perfect_square(5 * n * n + 4)
         || is_perfect_square(5 * n * n - 4);
```

# 6.12 Derangement

```
der[0] = 1;
der[1] = 0;
for(int i = 2; i <= 10; ++i)
 der[i] = (ll)(i - 1) * (der[i - 1] + der[i - 2]);
```

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#### 6.13 Bernoulli Number

$$\sum_{k=1}^{n} k^{p} = \frac{1}{p+1} \sum_{i=0}^{p} (-1)^{i} {p+1 \choose i} B_{i} n^{p+1-i} \qquad B_{m}^{+} = 1 - \sum_{k=0}^{m-1} {m \choose k} \frac{B_{k}^{+}}{m-k+1}$$

#### 6.14 Forbenius Number

(X \* Y) - (X + Y) and total count is (X - 1) \* (Y - 1) / 2

# 6.15 Stars and Bars with Upper Bound

```
P = (1 - X^{r_1+1}) \dots (1 - X^{r_n+1}) = \sum_i c_i X^{e_i}
Ans = \sum_{i} c_i \binom{N - e_i + n - 1}{n - 1}
```

# 7 Strings

# 7.1 Aho-Corasick

```
const int K = 26;
struct Vertex {
 int next[K];
 bool leaf = 0;
 int p = -1, ans = 0;
  char pch;
  int link = -1, mlink = -1;
  //magic link, is the link to find the nearest leaf
 int go[K];
 Vertex(int p = -1, char ch = '$') : p(p), pch(ch) {
   fill(begin(next), end(next), -1);
    fill(begin(go), end(go), -1);
vector<Vertex> t;
int add_string(string const& s) {
 int v = 0;
  for(char ch : s) {
   int c = ch - 'a';
   if(t[v].next[c] == -1) {
     t[v].next[c] = t.size();
     t.emplace_back(v, ch);
   v = t[v].next[c];
 t[v].leaf = 1;
 return v;
int go(int v, char ch);
int get_link(int v) {
 if(t[v].link == -1) {
   if(v == 0 || t[v].p == 0)
     t[v].link = 0;
      t[v].link = go(get_link(t[v].p), t[v].pch);
  return t[v].link;
int get_mlink(int v) {
 if(t[v].mlink == -1) {
```

```
Don't Forgor the Proof Peko
Bina Nusantara University
```

if(v == 0 || t[v].p == 0)

t[v].mlink = 0;

```
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```

```
else {
      t[v].mlink = go(get_link(t[v].p), t[v].pch);
      if(t[v].mlink && !t[t[v].mlink].leaf) {
        if(t[t[v].mlink].mlink == -1)
          get_mlink(t[v].mlink);
        t[v].mlink = t[t[v].mlink].mlink;
 return t[v].mlink;
int go(int v, char ch) {
 int c = ch - 'a';
 if(t[v].go[c] == -1) {
    if(t[v].next[c] != -1)
     t[v].go[c] = t[v].next[c];
    else
      t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
 return t[v].go[c];
//t.pb(Vertex());
7.2 Eertree
   Eertree - keep track of all palindromes and its occurences
  This code refers to problem Longest Palindromic Substring
https://www.spoj.com/problems/LPS/
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
struct node {
 int next[26];
 int sufflink;
 int len, cnt;
const int N = 1e5 + 69;
int n;
string s;
node tree[N];
int idx, suff;
int ans = 0;
void init_eertree() {
 idx = suff = 2;
 tree[1].len = -1, tree[1].sufflink = 1;
 tree[2].len = 0, tree[2].sufflink = 1;
bool add_letter(int x) {
 int cur = suff, curlen = 0;
 int nw = s[x] - 'a';
 while(1) {
    curlen = tree[cur].len;
   if(x - curlen - 1 >= 0 \&\& s[x - curlen - 1] == s[x])
    cur = tree[cur].sufflink;
  if(tree[cur].next[nw]) {
    suff = tree[cur].next[nw];
    return 0;
 tree[cur].next[nw] = suff = ++idx;
 tree[idx].len = tree[cur].len + 2;
```

```
ans = max(ans, tree[idx].len);
  if(tree[idx].len == 1) +
   tree[idx].sufflink = 2;
   tree[idx].cnt = 1;
   return 1;
 while(1) {
   cur = tree[cur].sufflink;
    curlen = tree[cur].len;
   if(x - curlen - 1 >= 0 \&\& s[x - curlen - 1] == s[x]) {
     tree[idx].sufflink = tree[cur].next[nw];
 tree[idx].cnt = tree[tree[idx].sufflink].cnt + 1;
 return 1;
int main() {
 ios::sync_with_stdio(0);
 cin.tie(0);
 cin >> n >> s;
 init eertree();
  for(int i = 0; i < n; i++)</pre>
   add_letter(i);
 cout << ans << '\n';
 return 0;
```

# 7.3 Manacher's Algorithm

```
// Computes lps array. lps[i] means the longest palindromic substring centered at i (\leftarrow
    when i is even, it is between characters. when it is odd, it is on characters)lps↔
    [0] = 0; lps[1] = 1;
REP(i, 2, 2 * str.size()) {
 int l = i / 2 - lps[i] / 2;
 int r = (i - 1) / 2 + lps[i] / 2;
 while(1) { // widen
   if(l == 0 || r + 1 == str.size())
     break;
   if(str[l - 1] != str[r + 1])
     break;
   l--, r++;
 lps[i] = r - l + 1;
  // jump
 if(lps[i] > 2) {
   int j = i - 1, k = i + 1; // while lps[j] inside lps[i]
   while(lps[j] - j < lps[i] - i)</pre>
     lps[k++] = lps[j--];
    lps[k] = lps[i] - (i - j); // set lps[k] to edge of lps[i]
   i = k - 1; // jump to mirror, which is k
```

# 7.4 Suffix Array

```
// stores result in sa and lcp
// if lcp is needed, call SuffixArray(str, 1)
struct SuffixArray {
  int n;
  vector<int> sa, lcp, rnk, cnt;
  vector<pair<int, int>> p;
  SuffixArray(const string& s, bool calc_lcp = 0) :
    n(s.length()), sa(n), lcp(calc_lcp ? n : 0), rnk(n),
    cnt(max(n, 256)), p(n) {
    for(int i = 0; i < n; ++i)
        rnk[i] = s[i];
    iota(sa.begin(), sa.end(), 0);</pre>
```

```
for(int i = 1; i < n; i <<= 1)
      update_sa(i);
    if(!calc_lcp)
     return;
    vector<int> phi(n), plcp(n);
    phi[sa[0]] = -1;
    for(int i = 1; i < n; ++i)</pre>
      phi[sa[i]] = sa[i - 1];
    int l = 0;
    for(int i = 0; i < n; ++i) {</pre>
      if(phi[i] == -1)
        plcp[i] = 0;
      else {
        while((i + l < n) && (phi[i] + l < n)</pre>
              && (s[i + l] == s[phi[i] + l]))
        plcp[i] = l;
        l = max(l - 1, 0);
    for(int i = 0; i < n; ++i)
      lcp[i] = plcp[sa[i]];
  void update_sa(int len) {
    sort_sa(len);
    sort sa(0);
    for(int i = 0; i < n; ++i) p[i] = {rnk[i], rnk[(i + len) % n]};</pre>
    auto lst = p[sa[0]];
    rnk[sa[0]] = 0;
    int cur = 0;
    for(int i = 1; i < n; ++i) {
      if(lst != p[sa[i]]) {
        lst = p[sa[i]];
        ++cur;
      rnk[sa[i]] = cur;
  void sort_sa(int offset) {
    fill(cnt.begin(), cnt.end(), 0);
    for(int i = 0; i < n; ++i)
     ++cnt[rnk[(i + offset) % n]];
    int sum = 0;
    for(int i = 0; i < (int) cnt.size(); ++i) {</pre>
      int temp = cnt[i];
      cnt[i] = sum;
      sum += temp;
    vector<int> temp(n);
    for(int i = 0; i < n; ++i) {
      int cur = cnt[rnk[(sa[i] + offset) % n]]++;
      temp[cur] = sa[i];
       = move(temp);
};
```

#### 7.5 Suffix Automaton

```
struct state {
  int len, link;
  map<char, int>next; //use array if TLE
};

const int MAXLEN = 100005;
  state st[MAXLEN * 2];
  int sz, last;

void sa_init() {
  sz = last = 0;
}
```

```
st[0].len = 0;
  st[0].link = -1;
  st[0].next.clear();
  ++sz;
void sa_extend(char c) {
 int cur = sz++;
  st[cur].len = st[last].len + 1;
  st[cur].next.clear();
  int p;
  for(p = last; p != -1 && !st[p].next.count(c); p = st[p].link)
    st[p].next[c] = cur;
  if(p == -1)
    st[cur].link = 0;
  else {
    int q = st[p].next[c];
    if(st[p].len + 1 == st[q].len)
      st[cur].link = q;
    else {
      int clone = sz++;
      st[clone].len = st[p].len + 1;
      st[clone].next = st[q].next;
      st[clone].link = st[q].link;
      for(; p != -1 \&\& st[p].next[c] == q; p = st[p].link)
        st[p].next[c] = clone;
      st[q].link = st[cur].link = clone;
  last = cur;
// forwarding
for(int i = 0; i < m; i++) {
 while(cur >= 0 && st[cur].next.count(pa[i]) == 0) {
    cur = st[cur].link;
   if(cur != -1)
     len = st[cur].len;
 if(st[cur].next.count(pa[i])) {
    len++:
    cur = st[cur].next[pa[i]];
 } else
    len = cur = 0;
// shortening abc -> bc
if(l == m) {
 1--:
 if(l <= st[st[cur].link].len)</pre>
    cur = st[cur].link;
// finding lowest and highest length
int lo = st[st[cur].link].len + 1;
int hi = st[cur].len;
//Finding number of distinct substrings
//answer = distsub(0)
LL d[MAXLEN * 2];
LL distsub(int ver) {
 LL tp = 1;
 if(d[ver])
    return d[ver];
  for(map<char, int>::iterator it = st[ver].next.begin();
     it != st[ver].next.end(); it++)
    tp += distsub(it->second);
  d[ver] = tp;
  return d[ver];
```

//Total Length of all distinct substrings

```
//call distsub first before call lesub
LL ans[MAXLEN * 2];
LL lesub(int ver) {
  LL tp = 0;
  if(ans[ver])
    return ans[ver];
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++)
    tp += lesub(it->second) + d[it->second];
  ans[ver] = tp;
  return ans[ver];
//find the k-th lexicographical substring
void kthsub(int ver, int K, string& ret) {
  for(map<char, int>::iterator it = st[ver].next.begin();
       it != st[ver].next.end(); it++) {
    int v = it->second;
    if(K <= d[v]) {
      if(K == 0) {
        ret.push_back(it->first);
      } else {
        ret.push_back(it->first);
        kthsub(v, K, ret);
        return;
    } else
      K -= d[v];
// Smallest Cyclic Shift to obtain lexicographical smallest of All possible
//in int main do this
int main() {
  string S;
  sa init();
  cin >> S; //input
  tp = 0:
  t = S.length();
  S += S:
  for(int a = 0; a < S.size(); a++)</pre>
    sa_extend(S[a]);
  minshift(0):
//the function
int tp, t;
void minshift(int ver) {
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++) {
    tp++;
    if(tp == t) {
      cout << st[ver].len - t + 1 << endl;</pre>
      break;
    minshift(it->second);
    break;
//end of function
// LONGEST COMMON SUBSTRING OF TWO STRINGS
string lcs(string s, string t) {
  sa_init();
  for(int i = 0; i < (int)s.length(); ++i)</pre>
    sa_extend(s[i]);
  int v = 0, l = 0,
      best = 0, bestpos = 0;
  for(int i = 0; i < (int)t.length(); ++i) {</pre>
```

```
while(v && ! st[v].next.count(t[i])) {
    v = st[v].link;
    l = st[v].length;
}
if(st[v].next.count(t[i])) {
    v = st[v].next[t[i]];
    ++l;
}
if(l > best)
    best = l, bestpos = i;
}
return t.substr(bestpos - best + 1, best);
}
```

### 8 OEIS

## 8.1 A000108 (Catalan)

Catalan numbers  $f(n) = nCk(2n,n) / (n+1) = nCk(2n,n) - nCk(2n,n+1) = f(n-1) * 2*(2*n-1) / (n+1) \\ 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, 2674440, 9694845, 35357670, 129644790, 477638700, 1767263190, 6564120420, 24466267020, 91482563640, 343059613650, 1289904147324, 4861946401452, 18367353072152, 69533550916004, 263747951750360, 1002242216651368, 3814986502092304$ 

#### 8.2 A000127

Maximal number of regions obtained by joining n points around a circle by straight lines  $f(n) = (n^4 - 6*n^3 + 23*n^2 - 18*n + 24) / 24$  1, 2, 4, 8, 16, 31, 57, 99, 163, 256, 386, 562, 794, 1093, 1471, 1941, 2517, 3214, 4048, 5036, 6196, 7547, 9109, 10903, 12951, 15276, 17902, 20854, 24158, 27841, 31931, 36457, 41449, 46938, 52956, 59536, 66712, 74519, 82993, 92171, 102091, 112792, 124314

## 8.3 A000668 (Mersene Primes)

Mersenne primes (of form 2^p - 1 where p is a prime)
3, 7, 31, 127, 8191, 131071, 524287, 2147483647, 2305843009213693951,
618970019642690137449562111, 162259276829213363391578010288127,
170141183460469231731687303715884105727

#### 8.4 A001434

Number of graphs with n nodes and n edges.
0, 0, 1, 2, 6, 21, 65, 221, 771, 2769, 10250, 39243, 154658, 628635, 2632420, 11353457, 50411413, 230341716, 1082481189, 5228952960, 25945377057, 132140242356, 690238318754

#### 8.5 A018819

Binary partition function: number of partitions of n into powers of 2 f(2m+1) = f(2m); f(2m) = f(2m-1) + f(m)1, 1, 2, 2, 4, 4, 6, 6, 10, 10, 14, 14, 20, 20, 26, 26, 36, 36, 46, 46, 60, 60, 74, 74, 94, 94, 114, 114, 140, 140, 166, 166, 202, 202, 238, 238, 284, 284, 330, 330, 390, 390, 450, 450, 524, 524, 598, 598, 692, 692, 786, 786, 900, 900, 1014, 1014, 1154, 1154, 1294, 1294

## 8.6 A092098

3-Portolan numbers: number of regions formed by n-secting the angles of an equilateral triangle. long long solve(long long n) {

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```
long long res = (n % 2 == 1 ? 3*n*n - 3*n + 1 : 3*n*n - 6*n + 6);
const int bats = n/2 - 1;
for (long long i=1; i<=bats; i++) for (long long j=1; j<=bats; j++) {
    long long num = i * (n-j) * n;
    long long denum = (n-i) * j + i * (n-j);
    res -= 6 * (num % denum == 0 && num / denum <= bats);
} return res;
}
1, 6, 19, 30, 61, 78, 127, 150, 217, 246, 331, 366, 469, 510, 625, 678, 817,
870, 1027, 1080, 1261, 1326, 1519, 1566, 1801, 1878, 2107, 2190, 2437, 2520,
2791, 2886, 3169, 3270, 3559, 3678, 3997, 4110, 4447, 4548, 4921, 5034, 5419,
5550, 5899, 6078, 6487</pre>
```

### 8.7 A277402

3-Portolan numbers: number of regions formed by n-secting the angles of an equilateral triangle.  $a(n) = 6n + 6(n-1 - n\%2) + a(n-2); \quad f(n) = a(n) \text{ if } n \% \text{ 10 != 0 else } a(n) - 12 \\ 1, 6, 19, 30, 61, 78, 127, 150, 217, 234, 331, 366, 469, 510, 631, 678, 817, 870, 1027, 1074, 1261, 1326, 1519, 1590, 1801, 1878, 2107, 2190, 2437, 2514, 2791, 2886, 3169, 3270, 3571, 3678, 3997, 4110, 4447, 4554, 4921, 5046, 5419, 5550, 5941, 6078, 6487, 6630, 7057, 7194$