Inferencia de Tipos

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Inferencia de tipos

Motivación

Dada una expresión: ¿Tiene tipo? ¿Cuál es el tipo? ¿Es el más general? ¿Qué necesitamos saber del contexto?

Introducción

¿Tiene tipo? ¿Cuál es el tipo? ¿Es el más general? ¿Qué necesitamos saber del contexto?

- \bullet (λx . isZero(x)) true
- $\lambda x. succ(x)$
- $\lambda x. succ(y)$
- $\blacksquare \emptyset \rhd \lambda x : \mathsf{Nat}. x : \mathsf{Nat} \to \mathsf{Nat}$
- $\blacksquare \emptyset \rhd \lambda x : X_1.x : X_1 \rightarrow X_1$

Introducción

Generalidad

¿Qué significa ser el juicio *más general*? Que todos los juicios derivables para $\lambda x. x$ son instancias de $\emptyset > \lambda x: X_1. x: X_1 \to X_1$. Por ejemplo:

- $\blacksquare \emptyset > \lambda x : Nat. x : Nat \rightarrow Nat$
- $\emptyset \rhd \lambda x : Bool. x : Bool \rightarrow Bool$
-

Ejemplos a ojo

Inferir el juicio de tipado de las siguientes expresiones:

- 1 $\lambda x. y$
- 2 f true
- 3 iszero(x)

Algoritmo de Martelli Montanari

Determinar el resultado de aplicar el algoritmo MGU sobre las siguientes ecuaciones:

- $\blacksquare \ \mathsf{MGU}\{X_2 \to X_1 \to \mathsf{Bool} \stackrel{?}{=} X_2 \to X_3\}$
- $\textbf{2} \ \texttt{MGU}\{ \big(X_2 \to X_1 \big) \to \mathsf{Nat} \stackrel{?}{=} X_2 \to X_3 \}$

¿Qué tipo tienen las siguientes expresiones?

- 1 $\lambda f. \lambda x. f(f x)$
- $\mathbf{2} \times (\lambda x. \operatorname{succ}(x))$
- $\lambda x. x y x$

¿Qué tipo tienen las siguientes expresiones?

- $\mathbb{W}(\lambda x. x y x)$

Ejercicio

Dada la siguiente extensión al conjunto de términos para el cálculo λ con listas:

$$M ::= \ldots | map_{\sigma,\tau} | foldr_{\sigma,\tau}$$

La modificación al sistema de tipos es la introducción de dos axiomas de tipado para $map_{\sigma,\tau}$ y $foldr_{\sigma,\tau}$:

$$\mathbb{W}(\mathsf{map}) \stackrel{\mathrm{def}}{=} \emptyset \rhd \mathsf{map}_{X_1,X_2} : (X_1 \to X_2) \to [X_1] \to [X_2]$$

$$\mathbb{W}(\textit{foldr}) \stackrel{\text{def}}{=} \emptyset \rhd \textit{foldr}_{X_1,X_2} : (X_1 \to X_2 \to X_2) \to X_2 \to [X_1] \to X_2$$

siendo X_1 y X_2 variables de tipo frescas. Se asumen dadas las extensiones correspondientes para Exase y mgu. Usar el algoritmo $\mathbb{W}()$ con esta nueva extensión para tipar la siguiente expresión:

foldr map

$$\mathbb{W}(foldr\ map) = ??$$

$$\begin{split} & \mathbb{W}(\textit{foldr}) = \emptyset \rhd \textit{foldr}_{X_3, X_4} : (X_3 \to X_4 \to X_4) \to X_4 \to [X_3] \to X_4 \\ & \mathbb{W}(\textit{map}) = \emptyset \rhd \textit{map}_{X_1, X_2} : (X_1 \to X_2) \to [X_1] \to [X_2] \\ & S = \textit{MGU}\{(X_3 \to X_4 \to X_4) \to X_4 \to [X_3] \to X_4 \stackrel{?}{=} ((X_1 \to X_2) \to [X_1] \to [X_2]) \to X_5 \} \\ & \mapsto^1 \{X_3 \to X_4 \to X_4 \stackrel{?}{=} (X_1 \to X_2) \to [X_1] \to [X_2], \ X_4 \to [X_3] \to X_4 \stackrel{?}{=} X_5 \} \\ & \mapsto^1 \{X_3 \stackrel{?}{=} X_1 \to X_2, \ X_4 \to X_4 \stackrel{?}{=} [X_1] \to [X_2], \ X_4 \to [X_3] \to X_4 \stackrel{?}{=} X_5 \} \\ & \mapsto^4 \{X_4 \to X_4 \stackrel{?}{=} [X_1] \to [X_2], \ X_4 \to [X_1 \to X_2] \to X_4 \stackrel{?}{=} X_5 \} \mid \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{X_4 \stackrel{?}{=} [X_1], \ X_4 \stackrel{?}{=} [X_2], \ X_4 \to [X_1 \to X_2] \to X_4 \stackrel{?}{=} X_5 \} \mid \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{[X_1] \stackrel{?}{=} [X_2], \ [X_1] \to [X_1 \to X_2] \to [X_1] \stackrel{?}{=} X_5 \} \mid \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{[X_2] \to [X_2 \to X_2] \to [X_2] \stackrel{?}{=} X_5 \} \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{[X_2] \to [X_2 \to X_2] \to [X_2] \stackrel{?}{=} X_5 \} \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{[X_2] \to [X_2 \to X_2] \to [X_2] \stackrel{?}{=} X_5 \} \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1] \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_2] \to [X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1 \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{[X_1 \mid X_2 \mid X_2 \to X_2] \to [X_2] \mid \{X_2 \mid X_1 \} \circ \{[X_1 \mid X_4 \} \circ \{X_1 \to X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{X_1 \mapsto X_2 \mid X_2 \mapsto X_2 \mid X_2 \mid X_2 \mid X_3 \} \\ & \mapsto^4 \{\} \mid \{X_1 \mapsto X_2 \mid X_2 \mapsto X_2 \mid X_2$$

$$\mathbb{W}(\textit{foldr map}) = \emptyset \rhd \textit{foldr}_{X_2 \to X_2, [X_2]} \; \textit{map}_{X_2, X_2} \colon [X_2] \to [X_2 \to X_2] \to [X_2]$$

$$\mathbb{W}(\textit{foldr}) = \emptyset \rhd \textit{foldr}_{X_3,X_4} : (X_3 \to X_4 \to X_4) \to X_4 \to [X_3] \to X_4$$
$$\mathbb{W}(\textit{map}) = \emptyset \rhd \textit{map}_{X_1,X_2} : (X_1 \to X_2) \to [X_1] \to [X_2]$$

$$S = MGU\{(X_3 \to X_4 \to X_4) \to X_4 \to [X_3] \to X_4 \doteq ((X_1 \to X_2) \to [X_1] \to [X_2]) \to X_5\}$$

= $\{X_2 \to X_2 / X_3, [X_2] / X_4, X_2 / X_1, [X_2] \to [X_2 \to X_2] \to [X_2] / X_5\}$

Listas

$$\sigma ::= \dots \mid [\sigma]$$

$$M, N, O ::= \dots \mid [\]_{\sigma} \mid M :: N \mid Case \ M \ of \ [\] \leadsto N \ ; h :: t \leadsto O$$

$$\frac{\Gamma \rhd M : \sigma \qquad \Gamma \rhd N : [\sigma]}{\Gamma \rhd M :: N : [\sigma]}$$

$$\Gamma \rhd M : [\sigma] \qquad \Gamma \rhd N : \tau \qquad \Gamma \cup \{h : \sigma, t : [\sigma]\} \rhd O : \tau$$

 $\Gamma \rhd \textit{Case M of } [\] \leadsto \textit{N} \ ; \textit{h} :: \textit{t} \leadsto \textit{O} : \tau$

$$\mathbb{W}([\]) \stackrel{\mathrm{def}}{=} \emptyset \rhd [\]_X : [X] \qquad \text{con } X \text{ variable fresca}$$

$$\mathbb{W}(U :: V) \stackrel{\mathrm{def}}{=} S\Gamma_1 \cup S\Gamma_2 \rhd S(M :: N) : S\tau$$

$$\mathbb{W}(U) = \Gamma_1 \rhd M : \sigma$$

$$\mathbb{W}(V) = \Gamma_2 \rhd N : \tau$$

$$S = MGU\{\tau \stackrel{?}{=} [\sigma]\} \cup \{\sigma_1 \stackrel{?}{=} \sigma_2 \mid x : \sigma_1 \in \Gamma_1, x : \sigma_2 \in \Gamma_2\}$$

 $\mathbb{W}([\]) \stackrel{\text{def}}{=} \emptyset \rhd [\]_X : [X]$ con X variable fresca

$$\begin{split} \mathbb{W}(U :: V) & \stackrel{\mathrm{def}}{=} S\Gamma_1 \cup S\Gamma_2 \rhd S(M :: N) : S\tau \\ \mathbb{W}(U) = \Gamma_1 \rhd M : \sigma \\ \mathbb{W}(V) = \Gamma_2 \rhd N : \tau \\ S = MGU\{\tau \stackrel{?}{=} [\sigma]\} \cup \{\sigma_1 \stackrel{?}{=} \sigma_2 \mid x : \sigma_1 \in \Gamma_1, x : \sigma_2 \in \Gamma_2\} \\ \mathbb{W}(Case \ U \ of \ [\] \leadsto V \ ; h :: t \leadsto W) & \stackrel{\mathrm{def}}{=} \\ S\Gamma_1 \cup S\Gamma_2 \cup S\Gamma_{3'} \rhd S \ (Case \ M \ of \ [\] \leadsto N \ ; h :: t \leadsto O) : S\tau \\ \mathbb{W}(U) = \Gamma_1 \rhd M : \sigma \qquad \mathbb{W}(V) = \Gamma_2 \rhd N : \tau \qquad \mathbb{W}(W) = \Gamma_3 \rhd O : \rho \\ \tau_h = \left\{ \begin{array}{c} \alpha \ \text{si} \ h : \alpha \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right. \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right\} \right\} \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right\} \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right\} \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right\} \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no} \end{array} \right\} \tau_t = \left\{ \begin{array}{c} \beta \ \text{si} \ t : \beta \in \Gamma_3, \\ \text{var fresca si no}$$

Listas por comprensión

$$M ::= \ldots \mid [M \mid x \leftarrow M, M]$$

Consideremos el Cálculo Lambda extendido con las listas por comprensión vistas en la práctica 4.

La regla de tipado es la siguiente:

$$\frac{\Gamma \cup \{x : \sigma\} \, \rhd M \colon \tau \quad \Gamma \rhd N \colon [\sigma] \quad \Gamma \cup \{x : \sigma\} \, \rhd \, O \colon \mathsf{Bool}}{\Gamma \rhd [M \mid x \leftarrow N, O] \colon [\tau]}$$

Listas por Comprensión

$$\mathbb{W}([\ U\ |\ x\leftarrow V,W\])\stackrel{\mathrm{def}}{=} S\Gamma_{1'}\cup S\Gamma_2\cup S\Gamma_{3'}\ \rhd S\left([\ M\ |\ X\leftarrow N,O\]\right):S[\sigma_1]$$

$$\mathbb{W}(U)=\Gamma_1\rhd M:\sigma_1$$

$$\mathbb{W}(V)=\Gamma_2\rhd N:\sigma_2$$

$$\mathbb{W}(W)=\Gamma_3\rhd O:\sigma_3$$

$$\tau_{x1}=\left\{\begin{array}{ll}\alpha\ \mathrm{si}\ x:\alpha\in\Gamma_1,\\ \mathrm{var}\ \mathrm{fresca}\ \mathrm{si}\ \mathrm{no}\end{array}\right.$$

$$\tau_{x2}=\left\{\begin{array}{ll}\beta\ \mathrm{si}\ x:\beta\in\Gamma_3,\\ \mathrm{var}\ \mathrm{fresca}\ \mathrm{si}\ \mathrm{no}\end{array}\right.$$

$$\Gamma_{1'}=\Gamma_1\ominus\{x\}\qquad\Gamma_{3'}=\Gamma_3\ominus\{x\}$$

$$S=MGU(\{\tau_{x1}\stackrel{?}{=}\tau_{x2},\ \sigma_2\stackrel{?}{=}[\tau_{x1}],\ \sigma_3\stackrel{?}{=}\mathsf{Bool}\}$$

Dar el tipo de: [if x then $\underline{0}$ else $\underline{1} \mid x \leftarrow false :: iszero(x) :: [], true]$

 $\cup \{ \rho_1 \stackrel{?}{=} \rho_2 \mid v : \rho_1 \in \Gamma_i, \ v : \rho_2 \in \Gamma_i, \ i, j \in \{1', 2, 3'\} \} \}$

Fin

Preguntas?????