EC 504 Quiz 1 Solutions

- 1. Order the following functions in increasing order of asymptotic complexity, from smallest to largest.
 - 1. $\log(\log(n))$
 - $2. \log(n)$
 - 3. $\sum_{k=1}^{\log(n)} \frac{n^2}{2^k}$
 - 4. $n^{1/n}$

Solution: Here is the right order:

- $1. \ n^{1/n}, \, \text{because } \lim_{n \to \infty} n^{\frac{1}{n}} = \lim_{n \to \infty} e^{\frac{\log n}{n}} = \exp\left(\lim_{n \to \infty} \frac{\log n}{n}\right) = 1$
- 2. $\log(\log(n))$, because $\lim_{n\to\infty}\frac{\log\log n}{\log(n)}$ can be obtained using L'Hopital's rule as

$$\lim_{n \to \infty} \frac{\log \log n}{\log(n)} = \lim_{n \to \infty} \frac{\frac{d}{dn} \log \log n}{\frac{d}{dn} \log(n)} = \lim_{n \to \infty} \frac{\frac{1}{n \log(n)}}{\frac{1}{n}} = 0$$

- 3. $\log(n)$, as shown above.
- 4. $\sum_{k=1}^{\log(n)} \frac{n^2}{2^k}$, because $\sum_{k=1}^{\log(n)} \frac{n^2}{2^k} \in [\frac{n^2}{2}, n^2]$.
- 2. Order the following functions in increasing order of asymptotic complexity, from smallest to largest.
 - 1. $n^{1.8}$
 - 2. n^2
 - $3. \sum_{k=1}^{n} \frac{n^2}{k}$
 - 4. $3^{\ln(n)}$

Solution: Here is the right order:

- 1. $3^{\ln(n)}$
- 2. $n^{1.8}$
- 3. n^2
- 4. $\sum_{k=1}^{n} \frac{n^2}{k}$

because $3^{\ln(n)}=n^{\ln(3)},$ and $\ln(3)<1.8.$ Also, $\sum_{k=1}^n \frac{n^2}{k}\in\Theta(n^2\log(n)).$

- 3. Order the following functions in increasing order of asymptotic complexity, from smallest to largest.
 - 1. *n*!
 - $2. \ 2^n$
 - $3. n^n$
 - 4. $3^{\ln(n)}$

Solution: Here is the right order:

- 1. $3^{\ln(n)}$
- $2. \ 2^n$
- 3. n!
- 4. n^n

```
because n! \approx \sqrt{2\pi n} (n/e)^n.
```

4. Consider the following function:

```
public int A(int n) {
    if (n == 1) return 1;
    else return A(n-1) * A(n-1);
}
```

Let T(n) be the run time performance for an input n. Which of the following is true?

- a. $T(n) \in \Theta(2^n)$
- b. $T(n) \in \Theta(n^2)$
- c. $T(n) \in \Theta(n)$
- d. $T(n) \in \Theta(\log(n))$

Solution: The right answer is a. Let T(n) be the asymptotic time for the instance of size n. Then,

$$T(n) = 2 * T(n-1) + 1$$

with initial condition T(1) = 1.

By induction, T(2) = 3, T(3) = 7. Claim that $T(n) = 2^n - 1$. Then,

$$T(n+1) = 2T(n) + 1 = 2(2^{n} - 1) + 1 = 2^{n+1} + 1$$

which shows $T(n) \in \Theta(2^n)$.

5. Let T(n) denote the asymptotic run time of this function.

```
int B(int n) {
    if (n == 0) return 1;
    if (B(n/2) >= 10)
        return B(n/2) + 10;
    else
        return 10;
}
```

As n increases, which of the following is true?

- a. $T(n) \in \Theta(n)$
- b. $T(n) \in \Theta(n^2)$
- c. $T(n) \in \Theta(n \log(n))$
- d. $T(n) \in \Theta(2^n)$

Solution: The right answer is a. Note that B(0) = 1, and B(1) = 10. For n > 1, we always have B(n/2) >= 10, so we have to follow the top branch of the if statement.

Thus, we have the recursion:

$$T(n) = 2 * T(n/2) + c$$

for some constant c. The 2 happens because we have to evaluate B(n/2) once for the if condition, and another time for the return.

This fits the Master theorem model, we have that the solution is $T(n) \in \Theta(n)$.

6. Consider the function below. Let T(n) denote the asymptotic run time of this function depending on n.

```
public int D(int n) {
    if (n == 0) return 1;
    if (n== 1) return 3;
    return D(n-1) + D(n-2)*D(n-2);
}
```

As n increases, which of the following is true?

```
a. T(n) \in \Theta(n^2)
b. T(n) \in \Theta(n \log(n))
c. T(n) \in \Theta(n)
d. T(n) \in \Theta(2^n)
```

Solution: The right answer is d. Note that D(0) = 1, and D(1) = 3. For n > 1,

A recursion for the asymptotic run time is

$$T(n) = T(n-1) + 2T(n-2) + c$$

This is a constant coefficient difference equation, so we look for solutions of the form z^n for some roots z. Substituting into the equation, we find the roots must satisfy $z^2 - z - 2 = 0$, so the possible vales are z = 2, z = -1. The root with largest magnitude is 2, so $T(n) \in \Theta(2^n)$.

7. Consider the function below. Let T(n) denote the asymptotic run time of this function depending on n.

As n increases, which of the following is true?

```
a. T(n) \in \Theta(n)
b. T(n) \in \Theta(n^2)
c. T(n) \in \Theta(n^{2/3})
d. T(n) \in \Theta(n^{\log_3(2)})
```

Solution: The right answer is a. As before, let's derive a recursion: There are two calls to E(n/3), and there is a loop of size n. Hence, the recursion is

$$T(n) = 2T(n/3) + cn$$

Using the Master's theorem, the answer is $\Theta(n)$.