

# Lecture 38: Register Allocation

[Adapted from notes by R. Bodik and G. Necula]

## Topics:

- Memory Hierarchy Management
- Register Allocation:
  - Register interference graph
  - Graph coloring heuristics
  - Spilling
- Cache Management

# The Memory Hierarchy

Computers employ a variety of memory devices, trading off capacity, persistence, and speed (some years ago):

Device	Access time	Capacity
Registers	1 cycle	256-2000 bytes
Cache	1-100 cycles	256KB-16MB
Main memory	150-1000 cycles	32MB — >16GB
Disk	0.5-10M cycles	10GB — > 1TB
Disk farm	...	> 3PB

# Managing the Memory Hierarchy

- Programs are written as if there are only two kinds of memory: main memory and disk (variables and files).
- Programmer is responsible for moving data from disk to memory.
- Hardware is responsible for moving data between memory and caches
- Compiler is responsible for moving data between memory and registers (which the programmer usually doesn't see).
- Cache and register sizes are growing slowly: important to manage them well.
- Processor speed improves faster than memory speed and disk speed ???
- The cost of a cache miss is growing, and the widening gap is bridged with more caches.

# The Register Allocation Problem

- Our intermediate code style uses temporaries profligately, simplifying code generation and optimization, but complicating final translation to assembly
- Hence, the register allocation problem:
  - Rewrite the intermediate code to use fewer temporaries than there are machine registers
- So we must assign more temporaries to a register, without changing the program behavior

# An Example

Consider the program

```
a := c + d
e := a + b
f := e - 1
```

assuming that assumption that *a* and *e* die after use. Then,

- Can reuse *a* after *a + b*
- Same with temporary *e* after *e - 1*
- Can allocate *a*, *e*, and *f* all to one register (*r1*):

```
r1 := c + d
r1 := r1 + b
r1 := r1 - 1
```

# Basic Register Allocation Idea

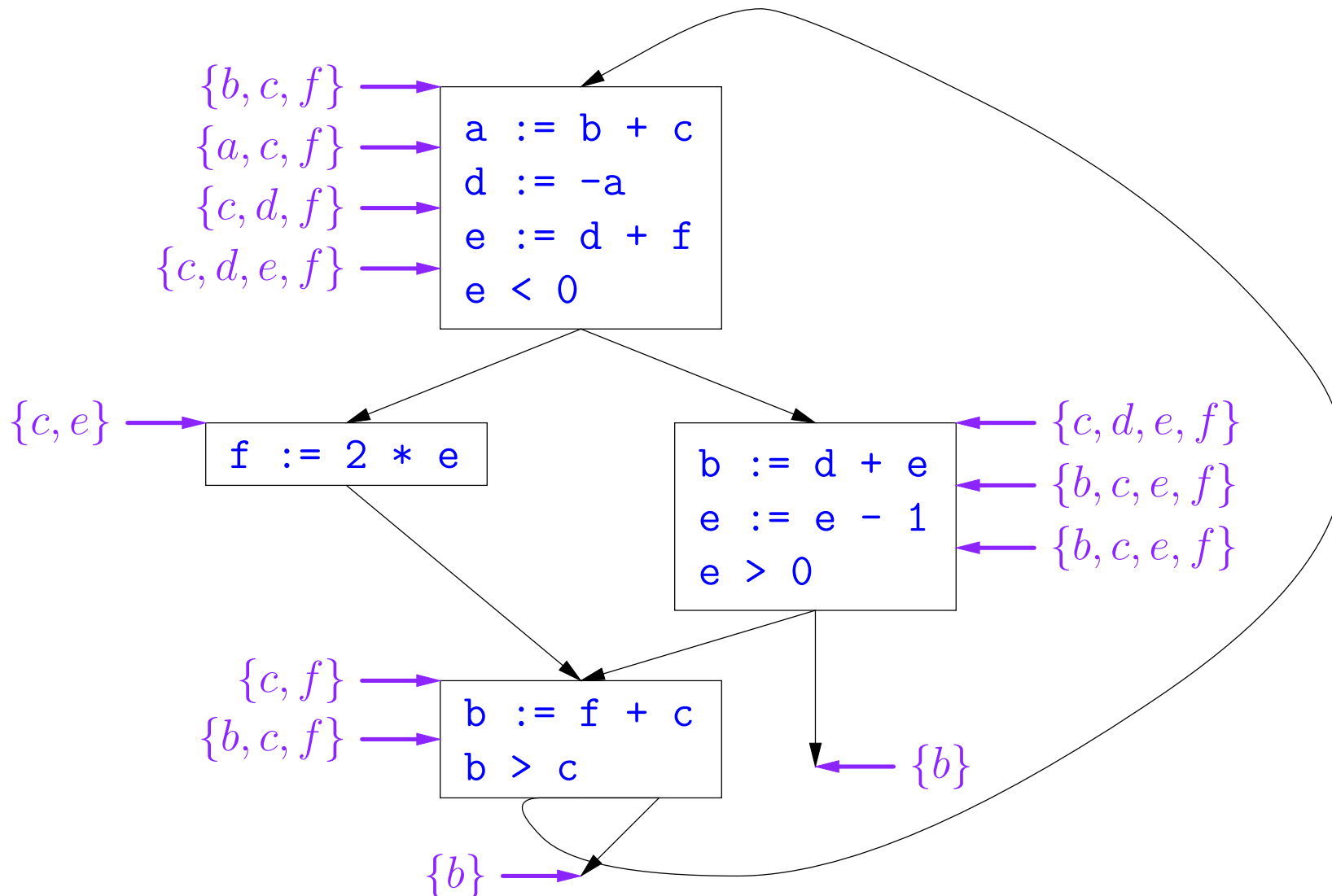
- So in general, since the value in a dead temporary is not needed for the rest of the computation,

Any set of temporaries can share a single physical register if at most one is alive at any program point.

- This rule is easy to apply to basic blocks. General CFGs are considerably trickier.

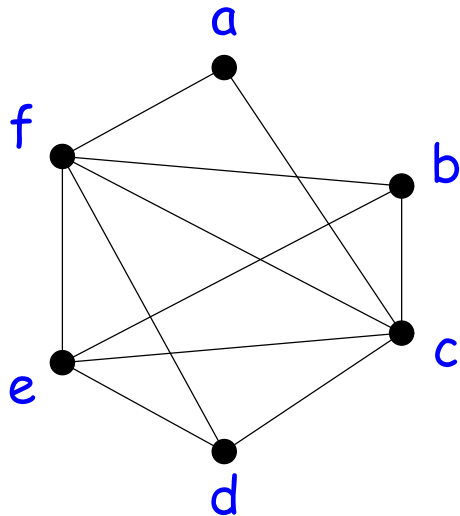
# Going Global: Allocation in CFGs (I)

First step is to compute live variables before each statement. In this example, assume that variable **b** is live at exit.



## Allocation in CFGs (II): Register Interference Graphs

- The sets in the previous slide indicate sets of virtual registers that are simultaneously alive at all points in the program, and therefore cannot share a physical register.
- Can summarize all these sets by constructing an undirected graph with a node for each virtual register, and an edge between any two virtual registers that appear together in the same set somewhere in the program.
- Call this the *register interference graph (RIG)*.



- The RIG extracts exactly the information needed to characterize legal register assignments
- Gives global (over the entire CFG) picture of the register requirements



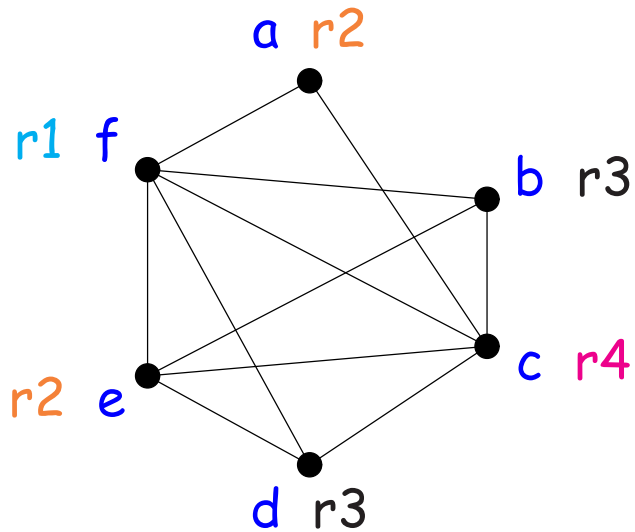
## Allocation in CFGs (III): Graph Coloring

- A *coloring* of a graph is an assignment of colors to nodes, such that nodes connected by an edge have different colors.
- A graph is *k-colorable* if it has a coloring with *k* colors.
- In our problem, *colors = registers*. That is,

If we have *k* available machine registers and our register interference graph is *k*-colorable, then the coloring gives us a register assignment.

# Graph Coloring: Example

Consider the sample RIG:



- There is *no* coloring with fewer than 4 colors
- There *are* 4-colorings of this graph

Before

```
.
A: a := b + c
d := -a
e := d + f
if e >= 0 jump C
B: f := 2 * e
jump D
C: b := d + e
e := e - 1
if e <= 0 jump E
D: b := f + c
if b <= c jump A
E:
```

After

```
A: r2 := r3 + r4
r3 := -r2
r2 := r3 + r1
if r1 >= 0 jump C
B: r1 := 2 * r2
jump D
C: r3 := r3 + r2
r2 := r2 - 1
if r2 <= 0 jump E
D: r3 := r1 + r4
if r3 <= r4 jump A
E:
```

## Allocation in CFGs (III): Computing Graph Colorings

- The remaining problem is to compute a coloring for the interference graph.
- Unfortunately, this problem is hard (NP-hard). No fast algorithms are known,
- And besides, a coloring might not exist for a given number of registers.
- For (1), we'll use heuristics.

# Graph Coloring Heuristic: Motivation

- Observation:
  - Pick a node  $t$  with  $< k$  neighbors in RIG.
  - Eliminate  $t$  and its edges from RIG.
  - If the resulting graph has a  $k$ -coloring then so does the original graph.
- Reason: whatever  $n \leq k - 1$  colors  $t$ 's neighbors have, we know we'll always be able to color  $t$  (since there are  $k$  colors). Therefore, eliminating  $t$  cannot affect the colorability of the other nodes.

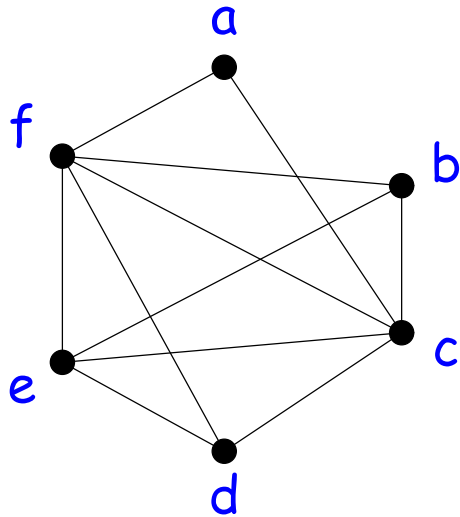
# Graph Coloring Heuristic

The following works well in practice:

- Pick a node  $t$  with  $< k$  neighbors.
- Push  $t$  on a stack and remove it from the RIG.
- Repeat until the graph has no nodes.
- Then start popping nodes from the stack and adding them back to the graph, assigning colors to each as we go (starting with the last node added).
- At each step, we know we can pick a color different from those assigned to already colored neighbors, by the observation on the last slide.

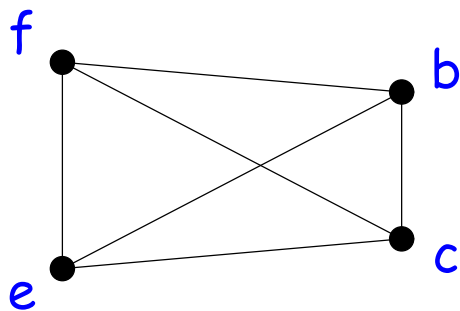
# Example of Using the Heuristic (I)

Start with our sample RIG and with  $k = 4$ :



Stack: []

Now remove  $a$  and then  $d$ , giving



Stack: [ $d, a$ ] (top on left)

Now all nodes have  $< 4$  neighbors; remove. Stack is [ $f, e, b, c, d, a$ ].

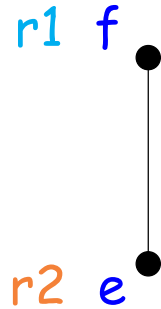
## Graph Coloring Example (2)

- Now we assign colors ...er, ...registers to: *f, e, b, c, d, a* in that order.
- At each step, guaranteed there's a free register.

*r1 f* ●

## Graph Coloring Example (2)

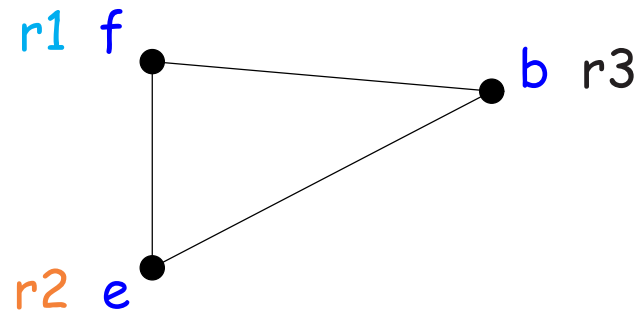
- Now we assign colors ...er, ...registers to: *f*, *e*, *b*, *c*, *d*, *a* in that order.
- At each step, guaranteed there's a free register.





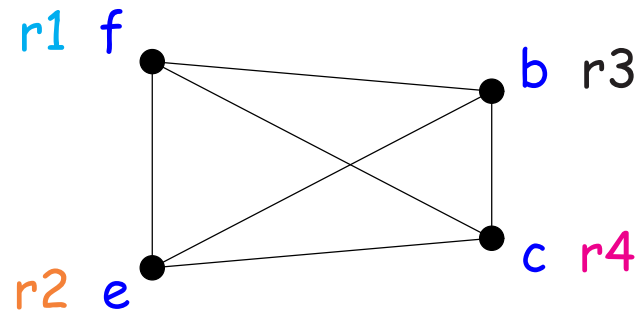
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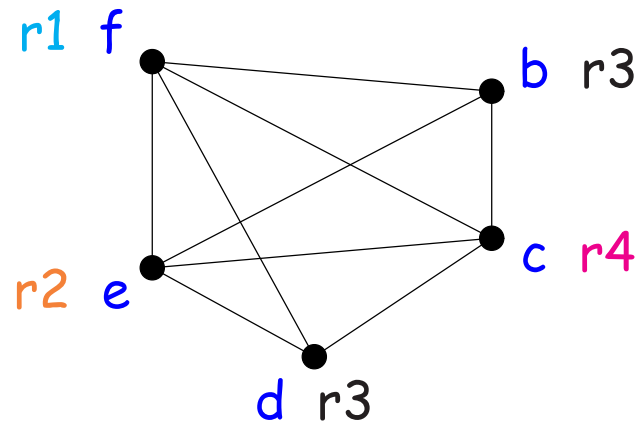
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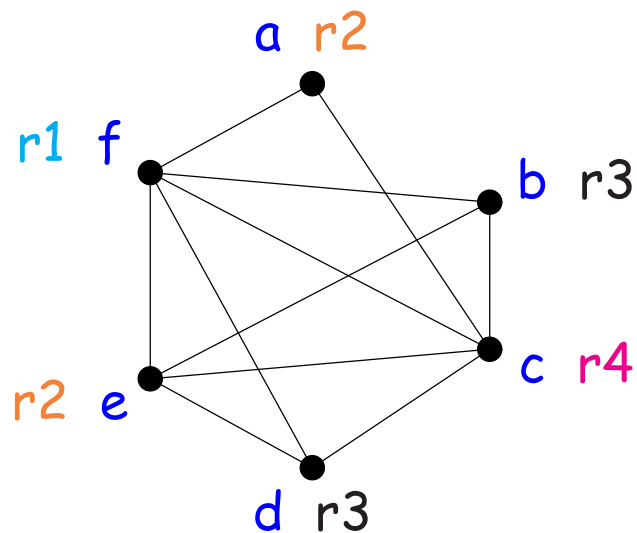
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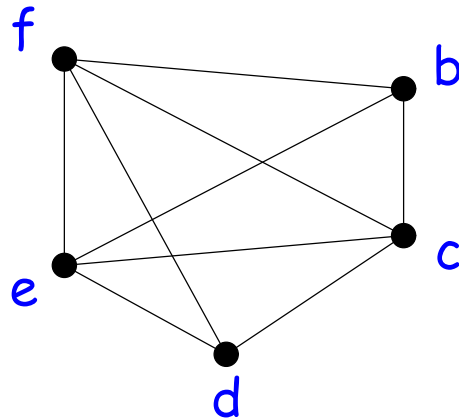
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# Spilling

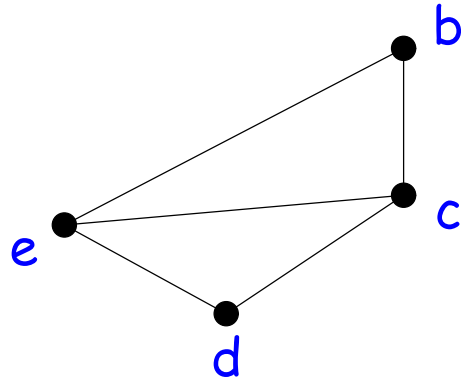
- What if during simplification we get to a state where all nodes have  $k$  or more neighbors?
- Example: try to find a 3-coloring of the RIG we've been using. After removing  $a$ , we get



- ...and now we are stuck, since all nodes have  $\geq 3$  neighbors.
- So, pick a node as a candidate for *spilling*, that is, to reside in memory.

## Example of Spilling

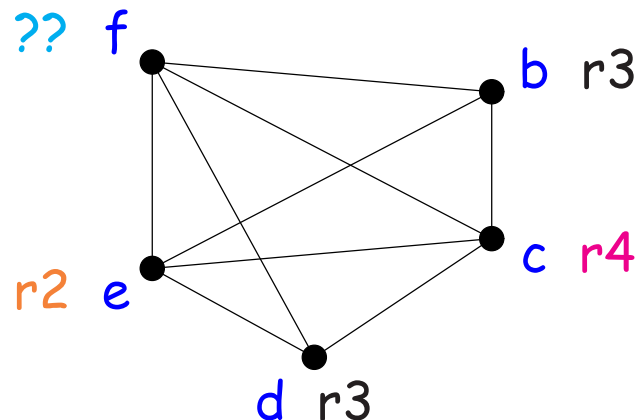
- Assume that  $f$  is picked as a candidate. When we remove it from the graph:



- Simplification now succeeds. We end up with the stack  $[e, c, b, d, f, a]$

## Example of Spilling (II)

- On the assignment phase we get to the point when we have to assign a color to  $f$
- Sometimes, it just happens that among the 4 neighbors of  $f$  we use  $< 3$  colors (*optimistic coloring*) ...



- ...but not this time.

## Example of Spilling (III)

- Since optimistic coloring failed we must spill register  $f$ : Allocate a memory location call it  $fa$  as the *home* of  $f$  (typically in the current stack frame).
- Before each operation that uses  $f$ , insert

$f := *fa$

- After each operation that defines (assigns to)  $f$ , insert

$*fa := f$

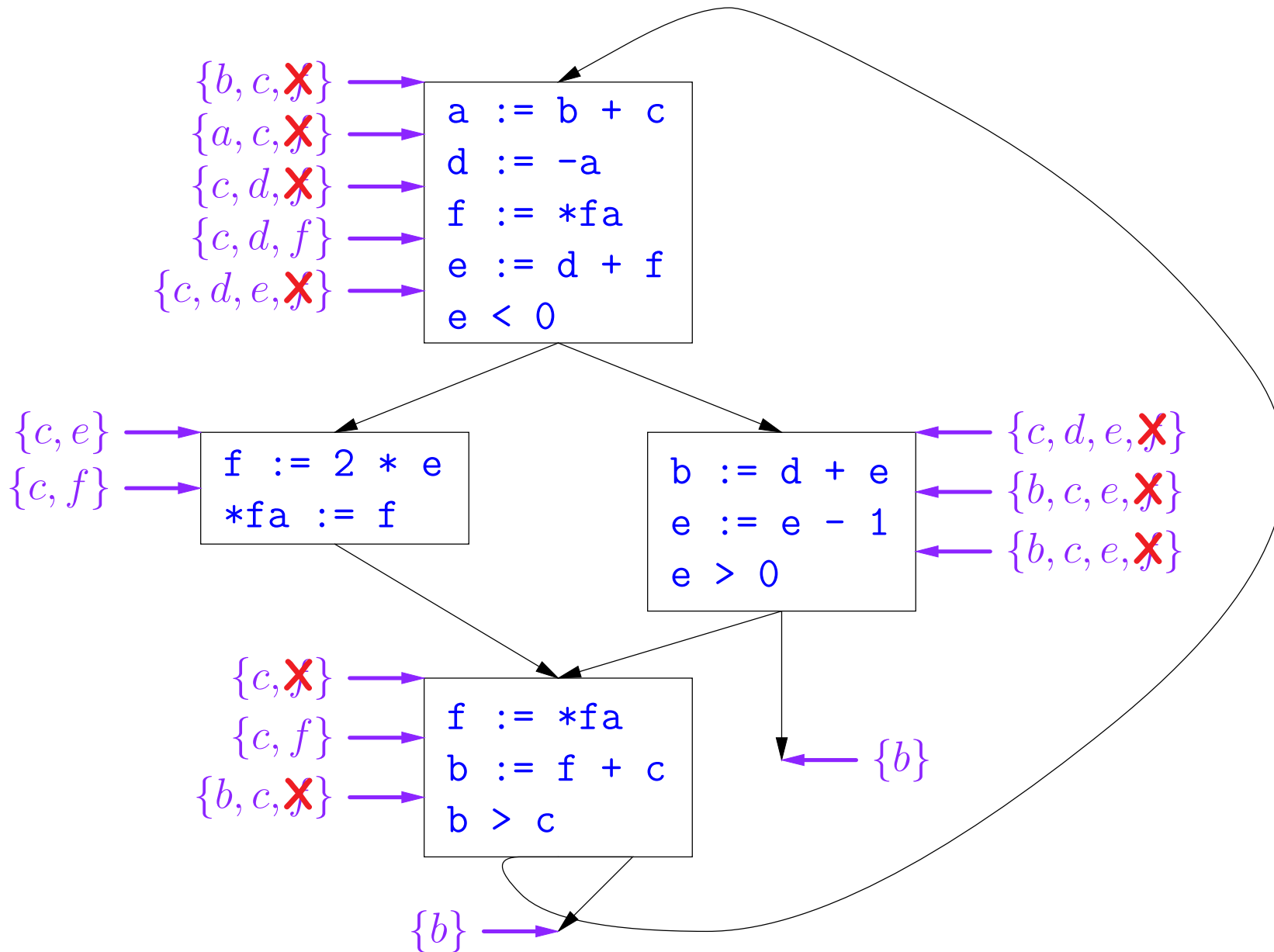
- This gives us:

```
A: a := b + c
d := -a
f := *fa
e := d + f
if e >= 0 jump C
B: f := 2 * e
*fa := f
jump D
```

```
C: b := d + e
e := e - 1
if e <= 0 jump E
f := *fa
D: b := f + c
if b <= c jump A
E:
```

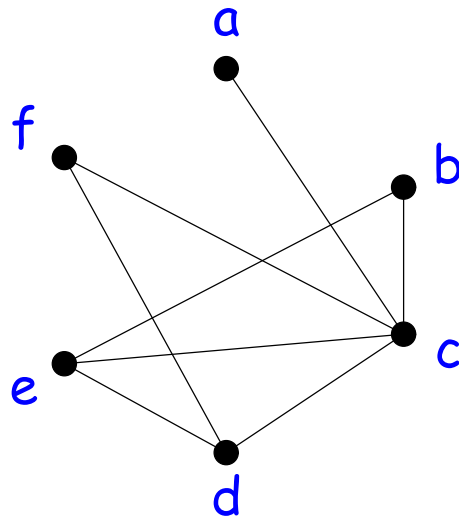


# Recomputing Liveness Information



# A New RIG

- The new liveness information is almost as before, except that that  $f$  is live only
  - Between an  $f := *fa$  and the next instruction, and
  - Between a  $\text{store } f, fa$  and the preceding instruction.
- That is, spilling reduces the live range of  $f$ , and thus the registers it interferes with, giving us this RIG:



- And this graph is 3-colorable (left to the reader).

# What to Spill?

- In general, additional spills might be required to allow a coloring.
- The tricky part is deciding what to spill. Possible heuristics:
  - Spill temporaries with most conflicts
  - Spill temporaries with few definitions and uses
  - Avoid spilling in inner loops

# Caches

- Compilers are very good at managing registers (much better than programmers: the *C register* declaration is really obsolete).
- Caches are another matter. The problem is still left to programmers, and it is still an open question whether compilers can do much in general to improve performance
- But they can (and a few do) perform some simple cache optimization

# Cache Optimization

- Consider the loop

```
for(j = 1; j < 10; j += 1)
    for(i = 1; i < 1000000; i += 1)
        a[i] *= b[i]
```

- Why does this have terrible cache performance?
- On the other hand,

```
for(i = 1; i < 1000000; i += 1)
    for(j = 1; j < 10; j += 1)
        a[i] *= b[i]
```

computes the same thing, but with much better (possibly 10x) performance [again why?].

- Compilers can do this: *loop interchange*.

## Cache Optimization (II)

- Other kinds of memory layout decisions possible, such as *padding* rows of a matrix with extra bytes to avoid cache conflicts when traversing a column (or row in FORTRAN) of a matrix. [Why might that help?]
- *Prefetching* instructions on some hardware can inform cache of anticipated future memory fetches so that they can proceed in parallel. Again, it is possible for compilers to supply these to a limited extent.

# Summary

- Both because it eases code generation, greatly improves performance, and because it is difficult for programmers to do it for themselves, register allocation is a “must have” optimization in production compilers for standard procedural languages.
- Graph coloring is a powerful register allocation scheme that compilers can apply automatically
- Good cache management could give even larger payoffs, but so far is difficult.