

# **Specware 4.1 User Manual**

## **Specware 4.1 User Manual**

Copyright © 2001-2004 by Kestrel Development Corporation

Copyright © 2001-2004 by Kestrel Technology LLC

All rights reserved

The name Specware® is a registered trademark of Kestrel Development Corporation

# Table of Contents

<b>1. Installing Specware .....</b>	<b>1</b>
1.1. Contents of Distribution Package .....	1
1.2. System Requirements.....	1
1.2.1. Hardware .....	2
1.2.2. Operating system .....	2
1.3. Installation Instructions.....	2
1.4. Uninstalling.....	3
<b>2. Getting Started .....</b>	<b>5</b>
2.1. Starting Specware and the Specware Shell .....	5
2.2. Exiting Specware .....	5
<b>3. Usage Model .....</b>	<b>7</b>
3.1. Units .....	7
3.2. Interaction .....	7
<b>4. Defining Units .....</b>	<b>9</b>
4.1. Conceptual Model.....	9
4.1.1. Unit Identifiers .....	9
4.1.2. Unit Terms .....	9
4.2. Realization via the File System.....	9
4.2.1. The SWPATH Environment Variable .....	10
4.2.2. Single Unit in a File .....	10
4.2.3. Multiple Units in a File .....	11
4.3. Unit Definitions Are Managed Outside of Specware .....	12
<b>5. Processing Units .....</b>	<b>13</b>
5.1. Overview .....	13
5.2. Resolution of Unit Identifiers.....	13
5.2.1. SWPATH-Based Unit Identifier.....	14
5.2.2. Relative Unit Identifiers .....	15
5.3. Processing Commands .....	18
5.3.1. Processing a Unit .....	18
5.3.2. Processing a Multiple-Unit File .....	18
5.4. Unit-Specific Processing.....	19
5.4.1. Typechecking Specs.....	19
5.4.2. Proving Properties in Specs .....	20
5.4.2.1. The Proof Unit .....	20
5.4.2.2. Proof errors .....	21
5.4.2.3. Proof Log Files .....	21

5.4.2.4. Multiple Proofs .....	21
5.4.2.5. Generating Proof Units .....	21
5.4.2.6. Interrupting Snark .....	22
5.4.2.7. The Prover Base Library .....	22
5.4.2.8. The Experimental Nature of the Prover .....	22
5.4.3. Generating Lisp Code .....	23
5.4.4. Generating Java Code .....	23
5.4.4.1. The Option-Spec .....	24
5.4.4.2. Translation of Builtin Ops.....	25
5.4.4.3. MetaSlang/Java Interface .....	26
5.4.4.4. Type Conversion between Java and MetaSlang .....	29
5.4.5. Generating C Code.....	29
5.4.5.1. Generating, Compiling and Running the generated C code using "make" .....	30
5.4.5.2. Compiling and Running the generated C code without using "make" .....	31
5.4.5.3. Garbage Collector .....	32
5.4.5.4. Supplying a C "main" function .....	32
5.5. Auxiliary Commands .....	33
5.5.1. Displaying Unit Values .....	33
5.5.2. Inspecting and Clearing the Cache .....	33
5.5.3. Inspecting and Setting SWPATH.....	34
5.5.4. Generating Lisp Code .....	34
5.5.5. Miscellaneous Commands .....	35
<b>6. Lisp Code Generated from Specs .....</b>	<b>37</b>
6.1. Translation of Specware Names to Lisp Names .....	37
6.2. Arity and Currying Normalization .....	37
6.3. Representation of Other Types.....	39
<b>7. Debugging Generated Lisp Files.....</b>	<b>41</b>
7.1. Tracing .....	41
7.2. Breaking.....	41
7.3. Timing.....	42
7.4. Interrupting .....	43
<b>8. Emacs Usage in Specware .....</b>	<b>45</b>
8.1. Specware Mode Commands .....	45
8.2. Specware Interaction Commands .....	46
8.3. Other Useful Emacs Commands .....	46

# Chapter 1. Installing Specware

## 1.1. Contents of Distribution Package

The following programs and documents are included on the Specware 4.1 installation CD:

- `setup.exe` -- This program in the root directory of the CD is the Specware 4.1 installer for Windows. It should be launched automatically when the CD is inserted into the CD-ROM drive.
- `XEmacs` -- This folder contains the XEmacs `setup.exe` installer for Windows, as well as the packages necessary to install the program. XEmacs is the environment within which Specware 4.1 is designed to run.
- `SNARK` -- This folder contains the license agreement and modified source code for the theorem prover SNARK, which is built into Specware 4.1.

The following documentation is included with the distribution package:

- Specware 4.1 Quick Reference -- The Quick Reference gives an overview of processing commands and Metaslang language constructs.
- Specware 4.1 User Manual (this document) -- The User Manual serves as a quick guide to basic usage and concepts of Specware. After reading this, you should feel comfortable with the mechanics of running and using Specware.
- Specware 4.1 Tutorial -- The Tutorial will guide you through the process of specifying, refining and generating code in Specware. A comprehensive example provides step-by-step instructions for this development process.
- Specware 4.1 Language Manual -- The Language Manual discusses the Metaslang specification language and gives the grammar rules and meaning for each Metaslang language construct.

## **1.2. System Requirements**

### **1.2.1. Hardware**

Specware has relatively modest system requirements for simple projects. Of course, as with any development tool, as your projects being developed become more complex, you may wish to work on a more powerful machine. For average use, however, the following basic hardware configuration is recommended:

- CPU: 250 Mhz
- RAM: 128 MB total, at least 64 MB free for applications
- Disk space: 15 MB for base system, 10-50 MB for user projects

### **1.2.2. Operating system**

This version of Specware 4.1 has been tested to work with Windows XP, Windows 2000 and Windows NT 4.0.

## **1.3. Installation Instructions**

1. Insert the installation CD into your CD-ROM drive; the Specware 4.1 `setup.exe` installer wizard will be automatically launched.
2. After accepting the license agreement, the installer will try to find the path to the XEmacs `xemacs.exe` startup file under the `Program Files` directory. If the path is found, the installer continues. Otherwise, you will have the option to either install XEmacs from the distribution CD (click `Yes`) or to manually type in the full path to `xemacs.exe`, if it is installed elsewhere on your machine (click `No`). Clicking `Yes` will launch the XEmacs installer wizard. Select the "Install from Local Directory" option for quickest installation, specify `[CD-ROM drive]:\XEmacs\packages\` as the directory from which to install the packages, and click `Next` through the remaining steps in the wizard to accept the default configuration. After XEmacs has been installed, you will return to the Specware installer wizard.

3. Select the directory where you would like Specware to be installed (the default is C:\Program Files\Specware), and click Next to complete the installation. A shortcut to Specware will be placed on your Desktop as well as in the Program Files folder in the Start menu. Documentation, libraries and example code will be placed in the installation directory you selected.

## **1.4. Uninstalling**

To uninstall Specware 4.1, run the installer on the CD again and select the "Remove" option in the wizard, or use the Add/Remove Programs setting in the Control Panel.





# Chapter 2. Getting Started

Specware is a development environment that runs on top of Lisp. The following sections describe the Specware environment and the basic mechanisms for running Specware.

## 2.1. Starting Specware and the Specware Shell

To start Specware, first double-click the `Specware 4.1` shortcut on your Desktop, or select `Specware 4.1` from the `Start Menu -> Program Files -> Specware` folder. An XEmacs window comes up in which the Specware Shell is running. The Specware Shell starts with issuing a prompt (an asterisk “\*”), prompting the user to issue a Specware Shell command. All of the user interaction (see the next chapter) with Specware occurs at the Specware Shell prompt: the user issues a command; Specware processes it and shows any results or error messages, insofar as applicable; then prompts the user; and so on, until the end of the session.

Processing errors may cause the execution of the Specware Shell to be interrupted, throwing interaction into a Lisp shell. To return to the Specware Shell, select the appropriate restart action offered by Lisp. Alternatively, issue the Lisp command `:sw-shell`, which however means the existing context is lost.

## 2.2. Exiting Specware

To exit Specware, type `exit` or `quit` at the Specware prompt. This will terminate the Specware session.



# Chapter 3. Usage Model

## 3.1. Units

Simply put, the functionality provided by Specware consists in the capability to construct specs, morphisms, diagrams, code, proofs, and other entities. All these entities are collectively called *units*.

Some of the operations made available by Specware to construct units are fairly sophisticated. Examples are colimits, extraction of proof obligations, discharging of proof obligations by means of external theorem provers, and code generation.

The Metaslang language is the vehicle to construct units. The language has syntax to express all the unit-constructing operations that Specware provides. The user defines units in Metaslang, writing the definitions in `.sw` files. (This file extension comes from the first letters of the two syllables in “Specware”.)

Currently, the only way to construct units in Specware is by writing text in Metaslang. The `.sw` files that define units are edited outside of Specware, using XEmacs, Notepad, Vim, or any other text editor of choice. These files are processed by Specware by giving suitable commands to the Specware Shell. Future versions of Specware will include the ability to construct units by other means. For instance, instead of listing the nodes and edges of a diagram in text, it will be possible to draw the diagram on the screen.

## 3.2. Interaction

The interaction between the user and Specware takes place through the Specware Shell.

When `.sw` files are processed by Specware, progress and error messages are displayed in the XEmacs buffer containing the Specware Shell. In addition, the results of processing are saved into an internal cache that Specware maintains. Lastly, processing of certain kinds of units result in new files being created. For example, when Lisp code is generated from a spec, the code is deposited into a `.lisp` file.

From the Specware Shell it is possible to evaluate Metaslang expressions in the context of a given spec, either directly or through generated Lisp code.

Specware also features auxiliary commands to display information about units, inspect and clear the internal cache, and inspect and change the `SWPATH` environment variable, which determines how unit identifiers are resolved to `.sw` files.



# Chapter 4. Defining Units

## 4.1. Conceptual Model

A unit definition consists of a unit identifier and a unit term. The identifier identifies the unit and the term defines how the unit is constructed.

A project developed with Specware consists of a set of unit definitions, some of which may come from libraries. Units have unique identifiers within the project.

### 4.1.1. Unit Identifiers

A unit identifier is a finite, non-empty sequence of word symbols (word symbols are defined in the Metaslang grammar). The sequence of word symbols is essentially a “path” in a tree: the units comprising a project are organized in a tree.

This provides a convenient and simple way to organize the units comprising a project. Libraries are subtrees of the whole tree. Parallel development of different parts of a project can be carried out in different subtrees that can be later put together without risk of naming clashes.

### 4.1.2. Unit Terms

A unit term is text written in the Metaslang language. Metaslang features various ways to construct specs, morphisms, and all the other kinds of units. For instance, it is possible to construct a spec by explicitly listing its types, ops, and axioms. It is also possible to construct a spec by applying the colimit operation to a diagram of specs and morphisms.

A unit term may reference other units. For instance, a spec constructed by extending another one references the spec being extended.

References can be “SWPATH-based” or “relative”. In either case they are resolved to full unit identifiers of units in the tree, according to simple rules explained later.

## 4.2. Realization via the File System

The conceptual model just described is realized by means of the file system of the underlying operating system. The file system has essentially a tree structure. The tree of units comprising a Specware project is mapped to subtrees of the file system; the word symbols comprising a path are mapped to file and directory names.

Future versions of Specware will have a more sophisticated UI that will realize the conceptual model directly. Users will graphically see the units organized in a tree and they will be able to add, remove, move, and edit them. The mapping to the file system may even be made totally transparent to the user.

### 4.2.1. The `SWPATH` Environment Variable

The mapping of the conceptual unit tree to the file system is defined by the environment variable `SWPATH`. Similarly to the `PATH` environment variable in operating systems, `SWPATH` is a string consisting of a semicolon-separated list of absolute directory paths in the file system. See Section 5.5.3 for information on how to inspect and set `SWPATH`.

Roughly speaking, the unit tree consists of all the units defined in `.sw` files under the directories listed in `SWPATH`. The identifier of each unit is its path from the directory in `SWPATH` under which the file defining the unit is: if the unit is under a directory named `ub2`, its identifier is its absolute path in the file system “minus” the `ub2` prefix. This approximate statement is made precise and illustrated by examples below.

### 4.2.2. Single Unit in a File

The simplest way to define a unit is to write its term into a `.sw` file in the subtree of one of the directories in `SWPATH`. The identifier of the unit is the name of the file, without `.sw`, prefixed by the path from the directory in `SWPATH` to the file.

For example, suppose that `SWPATH` includes the directory `C:\users\me\specware`. The user creates a file named `A.sw` immediately under the directory `C:\users\me\specware\one\two`, containing the following text:

```
spec
  type X
endspec
```

The absolute path of the file in the file system is `C:\users\me\specware\one\two\A.sw`. The unit is a spec declaring just a type `x`.

The identifier of the unit is `/one/two/A`. Note that the path components are separated by “/” (forward slash), even though the underlying file system uses “\” (backward slash). Unit identifiers are sequences of word symbols separated by “/”, regardless of the underlying operating system.

### 4.2.3. Multiple Units in a File

It is also possible to put multiple units inside a `.sw` file. The file must be in the subtree of one of the directories in `SWPATH`. Instead of just containing a unit term, the file contains one or more unit definitions. A unit definition consists of a word symbol, “=” (equal), and a unit term.

This case works almost exactly as if the file were replaced by a directory with the same name (without `.sw`) containing one `.sw` file for every unit defined therein. Each such file has the word symbol of the unit as name (plus `.sw`) and the term of the unit as content.

The only difference between the case of multiple units per file and the almost equivalent case where the file is replaced by a directory containing single-unit files, is that in the former case the last separator is not “/” but “#” (sharp). (This is reminiscent of the URI syntax, where subparts of a document are referred to using “#”.)

Suppose, as in the previous example, that `SWPATH` includes the directory `C:\users\me\specware`. The user creates a file named `three.sw` immediately under the directory `C:\users\me\specware\one\two`, containing the following text:

```
B = spec
  type Y
endspec

three = spec
  import B
  type Z
endspec
```

This file defines two specs, one declaring just a type `Y`, the other, next to importing the first spec, declaring just a type `Z`. The identifier of the first spec is `/one/two/three#B`, the identifier of the second spec is `/one/two/three#three`.

As a particular instance of the case of multiple units per file, it is possible to have just one unit definition in the file. This is different from just having a unit term in a file. If the file contains a unit definition, then the word symbol at the left of “=” is part of the

unit's identifier, together with “#” and the file path (relative to the directory in `SWPATH`). If instead the file contains a unit term, then the identifier of the unit is the file path (relative to the directory in `SWPATH`), without any “#” and additional word symbol.

Despite the possibility of having one unit definition in a file, in this manual we use the term “multiple-unit file” to denote a file that contains one or more unit definitions. The term “single-unit file” is instead used to denote a file that only contains a unit term.

As a convenience, a unit in a multiple-unit file with the same name as the file (without the directory and extension) may be referred to with a URI for the file as a whole. For example, in the current case, the identifier `/one/two/three` refers to the same spec as `/one/two/three#three`. This feature supports a style of having one primary unit in a file with auxiliary units that are used to define the primary unit.

## 4.3. Unit Definitions Are Managed Outside of Specware

The `.sw` files are created, deleted, moved, and renamed by directly interacting with the file system of the underlying operating system.

The content of the `.sw` files can be edited with any desired text editor. A possibility is to use XEmacs, which is started when Specware is started and is used to interact with Specware. The XEmacs-Specware combo can be thought of as a (rather limited) Integrated Development Environment (IDE).

Note that unit definitions can be managed without running Specware at all. As described in the next chapter, Specware is used to process unit definitions. Future versions of Specware will provide true IDE functionality: unit definitions will be also managed within Specware, and the mapping to the file system could be even made transparent to the user.



# Chapter 5. Processing Units

## 5.1. Overview

Unit definitions are processed by Specware. The user instructs Specware to process units by supplying certain commands. Specware has access, via the Lisp runtime environment, to the underlying file system, so it can access the `.sw` files that define units. The environment variable `SWPATH` determines which `.sw` files are accessed by Specware to find unit definitions.

Processing a unit causes the recursive processing of the units referenced in that unit's term. For instance, if a spec A extends a spec B which in turn extends a spec C, then when A is processed, B and C are also processed. There must be no circularities in the chain of unit dependencies.

Processing causes progress and/or error messages to be displayed in the XEmacs buffer containing the Progress messages inform the user that units were processed without error. Error messages provide information on the cause of errors, so that the user can fix them by editing the unit definitions. When an error occurs in the definition of some unit, Specware displays the `.sw` file containing the unit term in a separate XEmacs buffer, with the cursor positioned at the point of the erroneous text.

The processing of certain kinds of units also results in the creation of new files as an additional side effect. For instance, Lisp programs are a kind of unit, constructed by the `generate` operation of Metaslang. A side effect of processing one such unit is that the resulting code is written into a `.lisp` file.

When Specware processes a unit, it saves the processing results into an internal cache, associating the results with the unit's identifier. By using this cache, Specware avoids unnecessary re-computations: it only re-processes the units whose files have changed since the last time they were processed. From the point of view of the final result, this caching mechanism is completely transparent to the user. However, it improves the performance and response time of the system.

## 5.2. Resolution of Unit Identifiers

Unit terms may reference units in the form of unit identifiers. A unit identifier is resolved to the unit's term, which is contained in a `.sw` file. Unit identifiers are either

SWPATH-based or relative; these two kinds are syntactically distinguished from each other and are resolved in slightly different ways.

### 5.2.1. SWPATH-Based Unit Identifier

A SWPATH-based unit identifier starts with “/”, followed by a “/”-separated sequence of one or more path elements, where the last separator may be “#”. Examples are `/a/b/c`, `/d`, and `/e#f`.

Specware resolves a SWPATH-based unit identifier in the following steps:

1. If the unit identifier contains “#”, the “#” itself and the path element following it are removed, obtaining a “/”-separated sequence of one or more path elements, preceded by “/”. Otherwise, no removal takes place. Either way, the result of this first step is a “/”-separated sequence of path elements preceded by “/”.
2. The “/” signs of the “/”-separated sequence of path elements preceded by “/”, resulting from the previous step, are turned into “\”; in addition, `.sw` is appended at the end. The result of this second step is a (partial) path in the file system.
3. The path resulting from the previous step is appended after the first directory of SWPATH. If the resulting absolute path denotes an existing file, that is the result of this third step. Otherwise, the same attempt is made with the second directory of SWPATH (if any). Attempts continue until a directory is found in SWPATH such that the absolute path obtained by concatenating the directory with the result of the previous step denotes an existing file; such a file is the result of this step. If no such directory is found, the unit identifier cannot be resolved and an error is signaled by Specware.
4. There are two alternative steps here, depending on whether or not the original unit identifier contains “#”.
  - a. This is the case that the original unit identifier does not contain “#”. If the file resulting from the previous step is a single-unit file, i.e., it contains a unit term, that the final result of resolution. Otherwise, an error is signaled by Specware.
  - b. This is the case that the original unit identifier contains “#”. The file resulting from the previous step must be a multiple-unit file, i.e., it must contain a sequence of one or more unit definitions. If this is not the case, the unit identifier cannot be resolved and an error is signaled by Specware. If that is the case, a unit definition is searched in the file, whose path elements (to the left of “=”) is the same as the path element that follows “#” in the original unit identifier. If no such unit definition is found, the unit identifier cannot be

resolved and an error is signaled by Specware. If such a unit definition is found, its unit term (at the right of “=”) is the final result of resolution.

For example, consider a unit identifier `/a/b/c`. Since it does not contain “#”, the first step does not do anything. The result of the second step is `\a\b\c.sw`. Suppose that `SWPATH` is `C:\users\me\specware;C:\tmp`, that `C:\users\me\specware` does not contain any a subdirectory, and that `C:\tmp\a\b\c.sw` exists. The result of the third step is the file `C:\tmp\a\b\c.sw`. If such a file is a single-unit file, its content is the result of the fourth step.

As another example, consider a unit identifier `/e#f`. The result of the first step is `/e`. The result of the second step is `\e.sw`. Assuming that `SWPATH` is as before and that `C:\users\me\specware` contains a file `e.sw`, the file `C:\users\me\specware\e.sw` is the result of the third step. The file must be a multiple-unit file. Assuming that this is the case and that it contains a unit definition with path element `f`, its unit term is the result of the fourth step.

The directories in `SWPATH` are searched in order during the third step of resolution. So, in the last example, if the directory `C:\tmp` also contains a file `e.sw`, such a file is ignored. This features can be used, for example, to shadow selected library units that populate certain file system directories in `SWPATH`.

For example, suppose that the first directory in `SWPATH` is `C:\specware\libs` and that the directory `C:\specware\libs\data-structures` contains files `Sets.sw`, `Bags.sw`, `Lists.sw`, etc. defining specs of sets, bags, lists, etc. The unit identifier `/data-structures/Sets` resolves to the content of the file `C:\specware\libs\data-structures\Sets.sw`. If the user wanted to experiment with a slightly different version of the spec for sets, it is sufficient to prepend another directory to `SWPATH`, e.g. `C:\shadow-lib`, and to create that slightly different version of the spec for sets in `C:\shadow-lib\data-structures\Sets.sw`. The same unit identifier `/data-structures/Sets` will now resolve to the new version.

## 5.2.2. Relative Unit Identifiers

A relative unit identifier is a “/”-separated sequence of one or more path elements, where the last separator can be “#”. Examples are `a/b/c`, `d`, and `e#f`. So, `SWPATH`-based and relative unit identifiers can be distinguished by the presence or absence of “/” at the beginning.

The resolution of relative unit identifiers does not depend on `SWPATH`, but on the location of the file where the unit identifier occurs. There are two cases to consider: the

unit identifier occurring in a single-unit file and the unit identifier occurring in a multiple-unit file.

Suppose that the relative unit identifier occurs in a single-unit file. Then Specware attempts to resolve the unit identifier in the following steps:

1. If the unit identifier contains “#”, the “#” itself and the path element following it are removed, obtaining a “/”-separated sequence of one or more path elements. Otherwise, no removal takes place. Either way, the result of this first step is a “/”-separated sequence of path elements.
2. The “/” signs of the “/”-separated sequence of path elements, resulting from the previous step, are turned into “\”; in addition, `.sw` is appended at the end. The result of this second step is a (partial) path in the file system.
3. The path resulting from the previous step is appended after the absolute path of the directory of the file containing the relative unit identifier. If the resulting absolute path denotes an existing file, that is the result of this third step. Otherwise, the unit identifier cannot be resolved and an error is signaled by Specware.
4. There are two alternative steps here, depending on whether the original unit identifier contains “#” or not.
  - a. This is the case where the original unit identifier does not contain “#”. If the file resulting from the previous step is a single-unit file, i.e., it contains a unit term, that is the final result of resolution. Otherwise, an error is signaled by Specware.
  - b. This is the case that the original unit identifier contains “#”. The file resulting from the previous step must be a multiple-unit file, i.e., it must contain a sequence of one or more unit definitions. If this is not the case, the unit identifier cannot be resolved and an error is signaled by Specware. If that is the case, a unit definition is searched in the file, whose path element (at the left of “=”) is the same as the path element that follows “#” in the original unit identifier. If no such unit definition is found, the unit identifier cannot be resolved and an error is signaled by Specware. If instead such a unit definition is found, its unit term (to the right of “=”) is the final result of resolution.

So, resolution of a relative unit identifier occurring in a single-unit file is like resolution of a `SWPATH`-based unit identifier, except that the directory where the identifier occurs is used instead of `SWPATH`.

Suppose, instead, that the relative unit identifier occurs in a multiple-unit file. Then Specware attempts to resolve the unit identifier in the following steps:

1. If the relative unit identifier is a single path element, Specware attempts to find a unit definition with that path element inside the file where the unit identifier occurs. If such a unit definition is found, its unit term is the final result of resolution. Otherwise, the following steps are carried out:
2. If the unit identifier contains “#”, the “#” itself and the path element following it are removed, obtaining a “/”-separated sequence of one or more path elements. Otherwise, no removal takes place. Either way, the result of this first step is a “/”-separated sequence of path elements.
3. The “/” signs of the “/”-separated sequence of path elements, resulting from the previous step, are turned into “\”; in addition, .sw is appended at the end. The result of this second step is a (partial) path in the file system.
4. The path resulting from the previous step is appended after the absolute path of the directory of the file containing the relative unit identifier. If the resulting absolute path denotes an existing file, that is the result of this third step. Otherwise, the unit identifier cannot be resolved and an error is signaled by Specware.
5. There are two alternative steps here, depending on whether the original unit identifier contains “#” or not.
  - a. This is the case that the original unit identifier does not contain “#”. If the file resulting from the previous step is a single-unit file, i.e., it contains a unit term, that is the final result of resolution. Otherwise, an error is signaled by Specware.
  - b. This is the case that the original unit identifier contains “#”. The file resulting from the previous step must be a multiple-unit file, i.e., it must contain a sequence of one or more unit definitions. If this is not the case, the unit identifier cannot be resolved and an error is signaled by Specware. If that is the case, a unit definition is searched in the file, whose path element (at the left of “=”) is the same as the path element that follows “#” in the original unit identifier. If no such unit definition is found, the unit identifier cannot be resolved and an error is signaled by Specware. If instead such a unit definition is found, its unit term (to the right of “=”) is the final result of resolution.

So, resolution of a relative unit identifier occurring in a multiple-unit file is like resolution of a relative unit identifier occurring in a single-unit file, preceded by an

attempt to locate the unit in the file where the identifier occurs, only in case such a unit identifier is a path element.

## 5.3. Processing Commands

Each Specware Shell command consists of a keyword, the command name, followed by zero or more arguments. Several Specware Shell commands allow a variable number of arguments; e.g., zero or one. For many such commands, the zero-argument version means: use the last argument of the same kind last used for a Specware Shell command. In other cases, it means: use a default value for the argument. Which commands use which convention is detailed below.

A command entered by the user should be typed all on one line. The Return/Enter at the end of the line signals to the Specware Shell that the command must be executed.

### 5.3.1. Processing a Unit

The command to process a unit is:

```
proc [<unit-term>]
```

The Specware Shell command to process units is `proc`. The user supplies this command in the XEmacs buffer of the Lisp shell, followed by an argument. The argument is a unit reference, a piece of text defining an unnamed unit, or the name of a multiple-unit file. The whole command (including the argument) must fit in one line.

The unit term can be just a unit identifier can be `SWPATH`-based or relative. If it is `SWPATH`-based, Specware attempts to resolve it as explained in Section 5.2.1. If it is relative, Specware attempts to resolve it as explained in Section 5.2.2, as if the unit identifier occurred in a single-unit file of the current Lisp directory. Either way, if resolution fails an error is signaled by Specware.

If instead resolution succeeds, Specware parses and evaluates the unit term that results from resolution. Parsing and evaluation carry out the computations to construct the unit; they are Specware's "core" functionality. Parsing and evaluation implement the semantics of the Metaslang language.

If the unit term references other units, Specware recursively resolves the unit identifiers and parses and evaluates their unit terms.

### 5.3.2. Processing a Multiple-Unit File

The command to process a multiple-unit file is:

```
:sw <unit-identifier>
```

The unit identifier must not contain “#” and the file must not contain a unit of the same name as the file (without the directory and extension). Specware attempts to resolve the unit identifier. If it is a relative unit identifier, it is resolved as if it occurred inside a single-unit file in the Lisp current directory. However, the file obtained at the third step must be a multi-unit file, and not a single-unit file. If it is indeed a multi-unit file, Specware parses and evaluates all the unit definitions inside the file.

## 5.4. Unit-Specific Processing

This section describes what happens when specific kinds of units are processed.

### 5.4.1. Typechecking Specs

The user can construct specs by explicitly listing the types, ops, and axioms comprising the spec, possibly after importing one or more spec. When these spec definitions are processed, Specware typechecks all the expressions that appear in the spec.

Typechecking means checking that the expressions are type-correct, according to the rules of the Metaslang language.

In general, only some of the ops and variables that appear in an expression have explicit type information. Typechecking also involves reconstructing the types of those ops and variables that lack explicit type information.

Typechecking is an integrated process that checks the type correctness of expressions while reconstructing the missing type information. This is done by deriving and solving type constraints from the expression. For instance, if it is known that an op  $f$  has type  $A \rightarrow B$  then the type of the variable  $x$  in the expression  $f(x)$  must be  $A$ , and the type of the whole expression must be  $B$ .

If the missing type information cannot be uniquely reconstructed and/or if certain constraints are violated, Specware signals an error, indicating the textual position of the problematic expression.

Since the Metaslang type system features subtypes defined by arbitrary predicates, it is in general undecidable whether an expression involving subtypes is type-correct or not. When Specware processes a spec, it assumes that the type constraints arising from subtypes are satisfied, thus making typechecking decidable.

The proof obligations associated with a spec, which are extracted via the Metaslang `obligations` operation, include conjectures derived from the type constraints arising from subtypes. If all of these conjectures are discharged (using a theorem prover) then all the expressions in the spec are type-correct.

## 5.4.2. Proving Properties in Specs

Specware provides a mechanism for verifying the correctness of properties either specified in Metaslang specs or automatically generated as proof obligations arising from refinements or typechecking. Currently Specware comes packaged with the Snark first-order theorem prover. Interaction with Snark is through the proof unit described below.

### 5.4.2.1. The Proof Unit

The user invokes the Snark theorem prover by constructing a proof term. A typical proof term is of the form:

```
prove property in spec_term
  using hypothesis1, hypothesis2
  options "(use-paramodulation t)
          (use-resolution nil)
          (use-hyperresolution t)"
```

If this term is in file `proof.sw` then issuing the command `:sw proof` will result in translating `hypothesis1`, `hypothesis2` and `property` to Snark and then invoking the Snark prover to try to prove `property` from `hypothesis1` and `hypothesis2` using the options in the options list. Note that the properties `property`, `hypothesis1`, and `hypothesis2` must all be names of claims (i.e. axioms, conjectures, or theorems) that appear in the spec resulting from elaborating `spec_term`. Note also that `hypothesis1` and `hypothesis2` are required to appear earlier in that spec than `property`. Most users will omit the options list. Additionally, the `using` part can be omitted as well. In this case all the properties that appear in `spec_term` prior to `property` will be used to prove `property`.

After Snark completes Specware will report on the success or failure of the Snark proof.



### 5.4.2.2. Proof errors

Specware will report an error if `property` does not occur in `spec_term` or if one of the axioms do not occur in `spec_term` prior to `property`.

Snark will likely break into Lisp if the user inputs an incorrect option.

### 5.4.2.3. Proof Log Files

In the course of its execution Snark typically outputs a lot of information as well as a proof when it finds one. All this output can be overwhelming to the user, yet invaluable in understanding why his proofs succeeded or failed. To deal with all this output Specware redirects all the Snark output to log files. In our example above, which executed a proof in the file `proof.sw`, Specware will create a subdirectory called `snark` at the same level as `proof.sw`. In that directory a log file called `proof.log` will be created that contains all the Snark output.

### 5.4.2.4. Multiple Proofs

Just as there can be multiple units per file, there can be multiple proofs in single file. For example, in file `proof.sw` we could include more than one proof as follows:

```
p1 = prove prop1 using ax1, ax2
pla = prove prop1 using ax3
p2 = prove prop2
```

In this case Snark will be invoked three separate times, writing three different log files. In this case an additional subdirectory will be created under `snark`, called `proof`. The three log files will then be: `Snark/proof/p1.log`, `Snark/proof/pla.log`, and `Snark/proof/p2.log`.

### 5.4.2.5. Generating Proof Units

To facilitate the creation of proof units we provide the `:swpu` command. The `:swpu` commands takes a SpecCalculus term and generates proof units for all the conjectures and theorems in the spec denoted by the SpecCalculus term.

These proof units are written to a file that can then be elaborated to attempt proving all the conjectures in the spec. The proof file can also be edited to add options and usings to the proofs. By default the proofs are written in a file obtained from the unit id given to the `:swpu` command. For example, if the unit id is `/dir1/dir2/foo`, then the proof

units are written to `/dir1/dir2/foo_Proofs.sw`. Optionally the file for the proof units to be written to can be given as the second argument to the `:spwu` command.

By default proof units are generated not just for the conjectures explicitly in the spec, but also for all the proof obligations generated for the spec. The user can use the `:obligations` keyword to control the generation of proof obligations. For example, `:swpu spec :obligations nil`, will only generate proof units for the conjectures explicitly in spec.

#### 5.4.2.6. Interrupting Snark

As is the case with any first-order prover, Snark is likely to either loop forever or run for a longer time than the user can wait. The user can provide a time limit for Snark by using an appropriate option. However, there are likely to be times when the user wants to stop Snark in the middle of execution. He can do this by typing `Cntrl-C` `Cntrl-C` in the `*common-lisp*` buffer. This will then interrupt Snark and place the user in the Lisp debugger. He can exit the debugger by issuing the `:pop` command. A log file will still be written that he can look at if he needs to.

#### 5.4.2.7. The Prover Base Library

Specware has a base library that is implicitly imported by every spec. Unfortunately, the axioms in this library are not necessarily written to be useful by Snark. Instead of having Snark use these libraries we have created a separate base library for Snark. This library is located at `/Library/Base/ProverBase.sw`. The axioms in this spec are automatically sent to Snark as part of any proof.

#### 5.4.2.8. The Experimental Nature of the Prover

Our experience with the current prover interface is very new and as such we are still very much experimenting with it and don't expect it to be perfect at this point in time. Many simple theorems will be provable. Some that the user thinks should be might not, and the user will be required to add further hypothesis and lemmas that may seem unnecessary. We are currently working on making this interface as robust and predictable as possible, and welcome any feedback the user can offer.

One area where the user can directly experiment is with the axioms that make up the `ProverBase`. The axioms that make up an effective prover library are best determined by an experimental evolutionary process. The user is welcome to play with the axioms

in the `ProverBase`, by adding new ones or changing or deleting old ones. Keep in mind the goal is to have a single library that is useful for a wide range of proofs. Axioms that are specific to different proofs should be created in separate specs and imported where needed.

### 5.4.3. Generating Lisp Code

The user generates Lisp code from a spec by constructing a Lisp program unit via the Metaslang `generate` operation. The optional string given as argument must be an absolute file path in the file system. The generated Lisp code is deposited into that file. The `.lisp` file extension can be omitted.

The string argument can be omitted if the spec term given as argument is a unit identifier. In this case, Specware deposits the generated code into the file `A.lisp`, where `A` is the rightmost path element comprising the unit identifier. The `A.lisp` file is deposited in a `lisp` subdirectory immediately under the directory of the file containing the unit term of the spec identified by the unit identifier given as argument to `generate`.

For example, suppose that the first directory in `SWPATH` is `C:\tmp` and that a spec is defined in the single-unit file `C:\tmp\one\two\A.sw`. Suppose that the Lisp program unit is defined by:

```
generate lisp /one/two/A
```

Then the code is deposited into the file `C:\tmp\one\two\lisp\A.lisp`.

As another example, suppose that `SWPATH` is as before and that a spec is defined in the multiple-unit file `C:\tmp\one\two\f.sw`, and that `B` is the path element associated with the spec. Suppose that the Lisp program unit is defined by:

```
generate lisp /one/two/f#B
```

Then the code is deposited into the file `C:\tmp\one\two\lisp\B.lisp`.

### 5.4.4. Generating Java Code

As an experimental feature, Specware provides the possibility to generate Java code from executable MetaSlang specs. The Java code generator can either be called from the Specware Shell or using the `generate` construct inside a `.sw` file. In both cases, an

additional “option” spec can be supplied that is used to specify certain parameters. For the format of the option-spec see below.

Specware shell:

```
gen-java <UnitId> [<option-spec-id>]
```

Example:

```
gen-java MyExample MyExampleOptions
```

where MyExample.sw contains the spec that is to be translated into Java and MyExampleOptions.sw contains the option-spec (see below).

Inside a .sw file:

```
generate Java <spec-term> with <option-spec-id>
```

Example contents of a .sw file generating Java code:

```
let myspec = spec
...
endspec
in
let myoptions = spec
  def package = "com.mycompany.myapp"
  ...
endspec
in
generate Java myspec with myoptions
```

#### 5.4.4.1. The Option-Spec

The option spec is a regular MetaSlang spec that is used as an attribute store to be able to control certain parameters used to the Java code generator. The parameters are given by constant operator definitions inside the option-spec. The following list contains the operator names and types that are currently interpreted as parameters by the Java code generator:

Operator name & type	Used as	Default value
----------------------	---------	---------------

Operator name & type	Used as	Default value
package : String	Name of the Java package for all generated Java class files. The package name also determines the relative path of the generated .java files (see "basedir" parameter)	"specware.generated"
basedir : String	The base directory used for the generated Java class files. The full path is determine by this parameter and the relative path derived from the packge name. For instance, if basedir is "/a/b" and the package name is "c.d.e" then the generated Java class files would go into the directory "/a/b/c/d/e".	."
public : List String	The list of operators that are to be declared as "public" in the generated Java code. Only unqualified identifiers can be used in this list. The ops in this list determine the "entry-points" to the generated Java code, if it is embedded in another Java application.	[]
imports : List String	The list of imports that are to be used for all generated Java class files. Each element of this list has the usual format of the argument used by Java's import statement; e.,g., "java.util.*"	[]

Example option-spec:

```
spec
  def package = "test.sw.gen"
  def imports = ["java.util.*"]
  def public = ["test2"]
  def basedir = "~/myjavaapps"
endspec
```

#### 5.4.4.2. Translation of Builtin Ops

The following table shows the translation of some builtin MetaSlang operators into

Java:

MetaSlang	Java
String.writeLine(t) String.toScreen(t)	System.out.println(t)
String.concat(t1,t2), t1 ++ t2, t1 ^ t2	t1 + t2
String.newline	System.getProperty ("line.separator")
String.length	t.length()
String.substring(s,n1,n2)	s.substring(n1,n2)
Nat.toString(n) Nat.natToString(n) Nat.show(n) Integer.toString(n) Integer.intToString(n) Integer.show(n)	String.valueOf(n)
Nat.stringToNat(s) Integer.stringToInt(s)	Integer.parseInt(s)
t1 && t2	t1 && t2
t1    t2	t1    t2
t1 => t2	t1 ? t2 : true
t1 <=> t2	t1 ? t2 : !t2

#### 5.4.4.3. MetaSlang/Java Interface

In order to use Java methods and classes inside a MetaSlang spec, the following conventions are used by the Java code generator:

- Java Classes -- In order to use Java classes as sorts inside MetaSlang, you have to declare the sort without a definition term and add corresponding Java import statements using the option-spec (see above).

Example: use of the Java class "java.util.Vector"

In the spec:

```
...
sort Vector
...
```

```

op myvectorop: Vector -> Nat
def myvectorop(v) = ...
...

```

In the option-spec:

```

...
def imports = [ ..., "java.util.*", ... ]
...

```

The code generator interprets all sorts without a definition as base types, so that in this case the `op myvectorop` would become a static method in the generated `Primitive` class.

- Accessing Java Instance Methods -- Instance methods as well as static class methods can be called from MetaSlang using the following convention:

Assume, we want to use the instance method `foo(args)` defined in Java class `Bar`. First, the class must be known to MetaSlang by providing a sort declaration for the class. Then, an `op foo` must be declared with a signature that corresponds to the method's signature with an additional parameter preceeding the others. The sort of the parameter must be the class type:

```

sort Bar
op foo: Bar * T1 * ... * Tn -> T

```

where `T1 * ... * Tn -> T` is the original signature of `foo` without the additional parameter. The `T`'s are the translated MetaSlang sorts that correspond to the Java types occurring in `foo`'s signature; see the table below concerning the type conversion. In the MetaSlang code, a call to the instance method is created by the Java code generator whenever `foo` is called.

```

def mycode(...) =
  ...
  let b:Bar = ... in
  ...
  ... foo(b,arg1,...argn) ...

```

Note, that a definition term must not be given for `foo`. Limitation: using `foo` as a function symbol in higher-order context will not yield the expected result.

- Accessing Java Class Methods -- The access to Java class methods is very similar to instance methods with the difference that instead of the type of the first argument, the qualifier of the `op` declaration is used to determine the class name. Therefore, in

general, it is not necessary to declare the class as a sort. Assume we want to access to class method `Math.abs()` from the Java library. We therefore declare the `abs` operator in MetaSlang as follows:

```
op Math.abs: Integer -> Nat
```

The code generator will then generate a call to the static method `Math.abs()` whenever it is used in the MetaSlang spec. The access to static methods has lower priority than the access to instance methods: if the first argument is a user type that is not refined in the spec, than the instance call is generated. In other words, a static method in class A with first argument of type A will not be accessible from MetaSlang. The latter situation is not very common, so it does not represent a limitation of the MetaSlang-Java interface in practice.

- Accessing Java Constructors -- The access of Java constructors follows the same principle as the access of class methods. The difference is that on the MetaSlang side, an op with a name having the prefix `new` and an appropriate result type must be declared. For instance, the Java class `Vector` declares a constructor with no arguments. If we want to use that in MetaSlang, we have to provide the following declarations:

```
sort Vector
op Vector.new: () -> Vector
```

Whenever `Vector.new()` is used as a term in the MetaSlang spec, a call to the corresponding Java constructor in the `Vector` class is generated. If the class has multiple constructors with different parameter lists, multiple `new` ops can be declared in the MetaSlang spec with different suffixes (e.g., `new_2`) The Java code ignores the suffixes, but they are essential for MetaSlang, which doesn't allow the redefinition of ops with different signatures.

In general, if multiple methods and constructors from a class in the Java library need to be accessed in the MetaSlang spec, it is a good idea to structure them using the qualifying] features of MetaSlang. For instance:

```
[[ |
Vector qualifying spec |
sort Vector |
op new: () -> Vector
|
op add: [a] Vector * a -> Vector |
op size: Vector -> Nat
|
endspec |
|
Math qualifying spec |
op max: Integer *
Integer -> Integer |
op min: Integer * Integer -> Integer |
.... |
endspec
```

and then importing the specs into the application spec that uses it. Future versions of the Specware system will provide a utility to convert a given Java class into a spec following the above conventions.



#### 5.4.4.4. Type Conversion between Java and MetaSlang

The following table shows the conversion of Java types to MetaSlang, which can be used when accessing Java methods from MetaSlang

Java	MetaSlang
int	Integer
boolean	Boolean
char	Char
void	( )
byte short float double	not implemented
Any Java class name	MetaSlang type with the same name (type must be declared in the spec)

#### 5.4.5. Generating C Code

As an experimental feature, Specware provides the possibility to generate C code from executable MetaSlang specs. The C generator has been tested to work under Linux as well as Windows the latter using the cygwin package. The C code generator can either be called from the Specware Shell or using the SpecCalc "generate" construct inside a .sw file. In both cases, an additional parameter can be supplied specifying the base name of the C source and header files into which the C code will be generated.

Specware shell

```
gen-c <UnitId> [<c-file-base-name>]
```

Example:

```
gen-c QuickSort QuickSort
```

where QuickSort.sw contains the spec that is to be translated into the C files QuickSort.h and QuickSort.c

Inside a .sw file:

```
generate c <spec-term> in "<c-file-base-name>"
```

Example contents of a .sw file generating c code:

```
generate c QuickSort in "QuickSort.c"
```

This would have the same result as the above command line variant. The suffix ".c" can be omitted or not.

#### 5.4.5.1. Generating, Compiling and Running the generated C code using "make"

The easiest and recommended way of generating C code and compiling it is by using the Specware Shell command

```
make <unit-id>
```

This command does the following things:

- it invokes the "gen-c" command on the given unit-id and uses the name of the unit-id as file name for the generated C code ("#"s are replaced by "\_"s). For example,

```
make foo#bar
```

would invoke "gen-c foo#bar foo\_bar.c"

- if the C code generation has been successful, a customized Makefile is generated into "swcmake.mk". This file will include references to the built-in Makerules and define the targets and dependencies in a way that it compiles and executable with the same name as the generated C-files with removed suffix, i.e. for the above example the name of the executable would be "foo\_bar".

By convention, if a file named <unit>\_main.c or <unit>\_test.c exists, it will be automatically included in the build process; <unit>\_test.c will only be used, if <unit>\_main.c does not exist. For the above example this would mean that, if a file named "foo\_bar\_main.c" exists, it would be included in the build.

In addition to the built-in Makerules file, the generated make-file "swcmake.mk" will also include a unit-specific make-file in the current directory called <executable>.mk, i.e. in the above example "foo\_bar.mk". This file can be used to set the make variables CFLAGS and USERFILES, which are used as follows:

CFLAGS:	the value of the CFLAGS variable is used in calls to the C compiler (gcc) and usually contains example-specific flags, e.g., optimizer flags. Example: <code>CFLAGS = -O3</code>
USERFILES:	the value of the USERFILES make-variable is used in calls for the final compilation and linking of the executable. It usually lists additional C-files (.o and/or .c files) that the example needs to be a fully stand-alone application.

Other make variables that are used in the generated/predefined rules are LDFLAGS (can be used to add additional libraries, etc.), CPPFLAGS (see below), and USEGC (see below).

- finally, the actual "make" shell-command is called with the generated make-file "swcmake.mk" top-level make-file. By default, the actual command called is "make" assuming a program with this name is available in the current shell settings. The system environment variable SPECWARE4\_MAKE can be used to overwrite this default behaviour: if SPECWARE4\_MAKE is set, the value of it is used as make command.

#### 5.4.5.2. Compiling and Running the generated C code without using "make"

The generated C code is designed to contain as minimal as possible references outside the generated code, but there are still some builtin routines that are referred to. For that reason, the C compiler needs a few extra arguments specifying system include paths and the location of the garbage collector library that is used in the generated code. The easiest way of compiling the generated code is by using a Makefile and including the supplied C generator system Makefile in it. The corresponding include statement in a user Makefile would then be as follows:

```
include $(SPECWARE4)/Languages/MetaSlang/CodeGen/C/Clib/Makerules
```

This file set the CPPFLAGS and CFLAGS make variables to include the paths and

library necessary for successfully compiling the generated code. If additions to this variables need to be done, one can either define the corresponding variable before the above include statement in the Makefile or use the "!=" assignment after the include statement to prevent "make" to recursively process the variable. Example

```
CPPFLAGS := -g -pg $(CPPFLAGS)
```

would be a valid statement for augmenting the CPPFLAGS after the include statement. See

```
$SPECWARE4/Languages/MetaSlang/CodeGen/C/Examples/Makefile
```

for an annotated example Makefile.

### 5.4.5.3. Garbage Collector

By default, the generated C code generates calls to the public-domain garbage collector written by H. Boehm (see [http://www.hpl.hp.com/personal/Hans\\_Boehm/gc/](http://www.hpl.hp.com/personal/Hans_Boehm/gc/)). The library needs to be built once on a fresh Specware4 tree and will then be used by the Specware-generated C code. The easiest way to build the gc-library is described in the example Makefile mentioned above: simply add the variable \$(GCLIB) to the list of dependencies in your main Makefile target. Alternatively, this can be done by hand by changing to the directory "\$SPECWARE4/Languages/MetaSlang/CodeGen/C/Clib/gc6.2" and then running "make". After successful completion of this command, a file named "gc.a" should be present in that directory.

To disable the garbage collector simply put the variable definition

```
USEGC = no
```

in front of the line including the above Makerules file. This will prevent the generated code from calling the allocation function of the garbage collector and the garbage collector library will not be bound to the executable.

### 5.4.5.4. Supplying a C "main" function

If you want to create a stand-alone C application using the Specware-generated code, you have to supply a main function. This can be done either by directly defining an unqualified MetaSlang operator "main" like this

```

op main: () -> ()
def main () ...

```

or by hand-coding a main C function in a separate C file, from where the Specware-generated code would be called. Note, that currently, passing command line arguments is not yet supported, if you define a MetaSlang main operator directly. See the Examples directory for examples of both a hand-written main c function that calls the generated code and a MetaSlang definition of the main operator.

## 5.5. Auxiliary Commands

### 5.5.1. Displaying Unit Values

When a unit definition is processed, a unit value is produced. For example, a spec is essentially a set of types, ops, and axioms. A spec can be constructed by means of various operations in the Metaslang language, but the final result is always a spec, i.e., a set of types, ops, and axioms.

The `:show` command is used to display the value of a unit. Like `proc`, it is also supplied in the XEmacs buffer of the Lisp shell and it is followed by a one-line argument. The argument is a unit identifier.

The command to display a unit is:

```
:show <unit-identifier>
```

This command first processes the unit, exactly as if it were `proc` `<unit-identifier>`. In addition, if processing does not yield errors, the value of the unit is displayed on screen. Of course, if the unit has been already processed via a `proc` command and its unit term has not been changed after that, the `:show` command will access the results saved in Specware's internal cache.

The `:show` command serves to inspect the value of units constructed via the Metaslang operations. This is especially useful for beginning users as an aid to clarify the semantics of such operations.

## 5.5.2. Inspecting and Clearing the Cache

As already mentioned, Specware maintains an internal cache that associates processing results with unit identifiers.

A list of the units currently present in the cache is displayed on screen via the following command:

```
:list
```

The cache is cleared (i.e., re-initialized) via the following command:

```
:sw-init
```

Normally, there is no need to use these commands.

## 5.5.3. Inspecting and Setting `SWPATH`

The value of the `SWPATH` environment variable is displayed on screen via the following command:

```
:swpath
```

The value of the `SWPATH` environment variable is changed via the following command:

```
:swpath <string>
```

The `<string>` must be a semicolon-separated list of absolute directory paths of the underlying operating system. For example, in order to set `SWPATH` to `C:\users\me` it is necessary to write `:swpath C:\users\me`.

Changes to `SWPATH` only apply to the currently running Specware session. If Specware is quit and then restarted, `SWPATH` loses the value assigned to it during the previous session, reverting to its default value.

## 5.5.4. Generating Lisp Code

Normally, Lisp code is generated by constructing a Lisp program unit via the Metaslang `generate` operation and by processing such unit via the `proc` command.

Specware provides an additional command to accomplish the same result without actually creating an explicit Lisp program unit. The command is:

```
:swl <unit-identifier> <string>
```

The unit identifier must resolve to a spec, which is processed by Specware. If the spec is successfully processed, Specware generates Lisp code from it (according to the semantics of `generate`) and deposits the resulting code into the file whose path is given by the string.

The second argument to `:swl`, i.e. the string, is optional. If it is not given, a file name is inferred as explained in Section 5.4.3 for `generate`.

```
:swll <unit-identifier>
```

This is like `:swl` except it generates code only for the local definitions of the spec given by the unit identifier and not any of the imports. It is intended for incremental development, so it compiles and loads the lisp definitions. Note that if you have not generated code for the imported specs and loaded them, trying to run the code generated by this command will give undefined function errors. Also, if the spec is unqualified but it is imported into a spec that is qualified, the package used will be `:SW-USER` instead of the package of the qualifier. To avoid this problem qualification can be added to the spec.

When developing a specware application you generate code for your application, compile, load and test it. Then if you make a modification to a single spec you can use the `:swll` to just load the generated code for the modification. This command is also useful when you just want to see the generated code for a particular spec.

```
:cl <lisp-file-name>
```

Compiles and loads a lisp file.

### 5.5.5. Miscellaneous Commands

```
:sw-help <command-name>
```

Prints brief documentation for the specware command `command-name`. If `command-name` is omitted it prints documentation for all the specware commands.

```
:dir <directory>
```

Gives directory listing of `directory`. If `directory` is omitted it uses the current directory.

```
:ls <directory>
```

## *Chapter 5. Processing Units*

Gives directory listing of directory. If directory is omitted it uses the current directory.



# Chapter 6. Lisp Code Generated from Specs

The translation of executable specs to Lisp code is straightforward for the most part as Lisp is a higher-order functional language. Functional expressions go to lambda expressions and most Specware types are implemented as Lisp lists and vectors apart from the strings, numbers, characters and booleans which are implemented by the corresponding Lisp datatypes. This guide is meant primarily to help the user in calling and debugging the functions generated from a spec, so we concentrate on the translation of op names to Lisp names and the implementation of types. The implementation details of procedural constructs such as pattern-matching are omitted. The interested user is free to examine the Lisp code itself, which is simple but verbose for pattern-matching constructs.

## 6.1. Translation of Specware Names to Lisp Names

Specware ops are implemented using Lisp `defuns` if they are functions, `defparameters` otherwise. Their names are upper-cased and put in the package with the same name as the qualifier, or `SW-USER` if unqualified. However, if the name is that of a built-in Lisp symbol, the name is prepended with the character "!" and not upper-cased. If the qualifier of the op is the same as a built-in Lisp package then `"-SPEC"` is appended to the spec name to get the package name. For example, the Lisp code for the spec:

```
Z qualifying spec
  def two: Nat = 2
  def add1(x:Nat): Nat = x + 1
endspec
```

is

```
(DEFPACKAGE "Z" )
(IN-PACKAGE "Z" )

(DEFPARAMETER TWO 2)
(DEFUN ADD1 (X) (INTEGER-SPEC::+-2 X 1))
```

## 6.2. Arity and Currying Normalization

All Specware functions are unary. Multiple argument functions are modeled using either functions with product domains, or curried functions. For efficiency we wish to exploit Common Lisp's support of n-ary functions. Arity normalization aims to minimize unpacking and packing of products when passing arguments to functions with product domains, and currying normalization aims to minimize closure creation when calling curried functions. The saving is particularly important for recursive functions where there is saving at each recursive call, and in addition, currying normalization may permit the Common Lisp compiler to do tail recursion optimization. The naming scheme does not require knowledge of the definition of a function when generating calls to the function.

For each function whose argument is a product, two entry points are created: a unary function whose name is derived from the op as described above, and an n-ary function whose name has "-n" appended. For example, for

```
op min : Integer * Integer -> Integer
```

there are two Lisp functions, #'MIN-2 and #'|!min|. A call with an explicit product is translated to the n-ary version, otherwise the unary version is used. For example, min(1,2) translates to (MIN-2 1 2), and foldr min inf 1 translates to (FOLDR-1-1-1 #'|!min| INF L). When generating Lisp for a definition, the form is examined to see whether the definition is naturally n-ary. If it is, then the primary definition is n-ary and the unary function is defined in terms of the n-ary function, otherwise the situation is reversed. For example, given the definition

```
def min(x,y) = if x <= y then x else y
```

we get the two Common Lisp definitions:

```
(DEFUN MIN-2 (X Y) (if (<= x y) x y))  
(DEFIN |!min| (X) (MIN-2 (CAR X) (CDR X)))
```

and given the definition

```
def multFG(x: Nat * Nat) = (F x) * (G x)
```

we get the two Common Lisp definitions:

```
(DEFUN MULTFG-2 (X Y) (MULTFG (CONS X Y)))  
(DEFUN MULTFG (X) (* (F X) (G X)))
```

For each curried function (i.e. for each function whose codomain is a function) there is an additional uncurried version of the function with "-1" added  $n$  times to the name where  $n$  is the number of curried arguments. For example, for

```
op foldr: [key,a,b] (a * b -> b) -> b -> map(key,a) -> b
```

there are two Lisp functions, #'FOLDR and #'FOLDR-1-1-1.

As with arity normalization, the definition of a curried function is examined to see whether it should be used to generate the curried or the uncurried version, with the other being defined in terms of this primary version.

As well as producing more efficient code, the currying normalization makes code easier to debug using the Common Lisp trace facility. For example if a function has a call of the form `foldr x y z`, this call is implemented as `(FOLDR-1-1-1 x y z)`, so you can trace `FOLDR-1-1-1` to find out how it is being called and what it is returning.

## 6.3. Representation of Other Types

`Character` and `String` types are represented as Lisp characters and strings, `Nat` and `Integer` as Lisp integers, lists are represented using Lisp lists, and `Boolean` `true` and `false` by the symbols `T` and `NIL`.

Sums are represented as the cons of the constructor name in keyword package and the fields of the constructor.

Binary products are implemented as cons cells (except for function arguments which are described in the previous section): `CONS` to construct and `CAR` and `CDR` to access the first and second fields. Non-binary products are implemented as vectors: constructed using `VECTOR` and the  $i$ th element accessed by `(SVREF x i-1)`.

Records are implemented the same as products with the order of the fields being alphabetic in the field names.

Restrictions and comprehensions are implemented using their supersort.

A quotient is represented as a vector of three elements: the quotient tag (which is the value of the Lisp variable `SLANG-BUILT-IN:QUOTIENT-TAG`), the representation of the quotient relation, and the actual value in the underlying type.



# Chapter 7. Debugging Generated Lisp Files

## 7.1. Tracing

If you need to debug your application, there a number of useful Lisp facilities you should be aware of. The simplest trick is to trace some functions you care about to see what they are doing.

```
(trace foo)
  This will display the arguments to foo each time it is
  called, and will display the results each time it returns.
```

```
(untrace foo)
  This will turn off any tracing on foo.
```

## 7.2. Breaking

If you need a more detail view of runtime behavior, you might want to BREAK some functions you care about.

```
(trace (foo :break-all t))
```

This will invoke the debugger each time foo is called, and upon each exit from foo.

Once you arrive in the debugger, the following commands are most useful:

:down <n>	Move to a deeper frame	<n> is optional
:up <n>	Move to a higher frame	<n> is optional
:zoom <n>	Display n frames	<n> is optional
:pri	Enter a dialog that lets you set printer control variables. For example, setting depth to 5 and length to 10 will let you see the top level structure of expressions, while suppressing deep expressions and the tails of long expressions. Note that you'll need to	

specify values for many contexts, but just hitting return leaves a value unchanged. You will likely want to modify the trace, debugger, and current values.

```
(pprint *)    Pretty print the expression for the current
              frame. Note that this only works immediately
              after arriving at a frame, e.g. via :down 0
              if necessary.

:cont        Continue as if nothing happened.
:restart      Resume execution at this frame.
:reset       Return to Lisp top level. (E.g., bail out
              to try again.)
:exit        Exit from Lisp to operating system -- ends
              session.

:help        Online documentation.

(misc ...)    The debugger is is a read/eval/print loop, so
              arbitrary Lisp forms will be evaluated (in the
              current dynamic context).

(untrace foo) Stop entering debugger when foo is called.
              Note that you may still enter the debugger
              for each exit from calls to foo already
              recursively in progress.
```

## 7.3. Timing

If you are curious about the overall performance of your application, the `time` macro will provide some quick information:

```
(time (foo nil))
  This will report the time and space used by foo, e.g.:
```

```
USER(1): (time (list 1 2 3))
; cpu time (non-gc) 0 msec user, 0 msec system
; cpu time (gc)      0 msec user, 0 msec system
; cpu time (total)  0 msec user, 0 msec system
; real time   231 msec
; space allocation:
```

```
; 6 cons cells, 0 symbols, 0 other bytes, 0 static bytes  
(1 2 3)
```

Note that `time` is transparent, i.e., it returns whatever its argument would return, including multiple values, etc., so it is safe to intersperse it nearly anywhere.

Common Lisp has more facilities for rolling your own timers: see the generic Common Lisp documentation, or contact Kestrel Technologies.

## 7.4. Interrupting

Finally, note that a useful trick in Lisp is to start your application, e.g. `(foo nil)`, then at an appropriate time hit Cntrl-C. This will interrupt your application and put you into the debugger. From there you can enter the command `:zoom` to see the top of the stack. That can often be quite revealing.





# Chapter 8. Emacs Usage in Specware

There are a number of emacs commands for Specware usage. We also list some useful general-purpose commands. There is a Specware mode for editing Specware files.

## 8.1. Specware Mode Commands

These are the commands available when editing `.sw` files. Some of them are available in the Specware menu in the menubar.

Key	Name	Description
tab	sw:indent-line	Indents line based on context.
linefeed	newline-and-indent	Same as return followed by tab.
M-tab	sw:back-to-outer-indent	Indents line for outer expression.
M-C-q	sw:indent-sexp	Indents parenthesized expression following point.
M-C-\	sw:indent-region	Indents region.
M-*	sw:switch-to-lisp	Switch to <code>*common-lisp*</code> buffer.
M-.	sw:meta-point	Prompts for op or type name (with default based on symbol at point), and goes to its definition. Searches only the units that have been loaded. First searches for definitions visible from the current unit.
M-,	sw:continue-meta-point	If previous meta-point command returned more than one definition, go to the next definition. Can be repeated.
M-	sw:electric-pipe	Adds the skeleton of a new case to a case statement, properly indented.
C-c ;	sw:comment-region	Comment out the region. With a negative argument uncomments the region.

Key	Name	Description
C-c p	sw:process-current-file	Does :sw on current file. Also available in <code>dired</code> where it applies to the file on the current line. (This command will only work properly if the top-level directory for the file is in <code>SWPATH</code> ).
C-c C-p	sw:process-unit	Prompts for unitId to process. Defaults to unitId for the current file.
C-c g	sw:generate-lisp	Does :swl on current file. With an argument compiles and loads the generated lisp.
C-c C-l	sw:generate-lisp	Does :swll on current file (generates lisp for local definitions of current file and compiles and loads them).
C-c !	cd-current-directory	Does a :cd in the <code>*common-lisp*</code> buffer to the directory of the current file. Also available in <code>dired</code> .

## 8.2. Specware Interaction Commands

These are the commands available in the `*common-lisp*` buffer.

Key	Name	Description
M-*	sw:switch-to-lisp	When already in the <code>*common-lisp*</code> buffer, goes to the previous source file buffer.
M-.	sw:meta-point	Prompts for op or type name (with default based on symbol at point), and goes to its definition. Searches only the units that have been loaded.
M-p	fi:pop-input	Gets previous input.
M-r	fi:re-search-backward-expression	Gets previous input matching regular expression (prompted for). Typically you can get a previous input by typing in a small substring. Repeating the command without changing the expression will find earlier matches.

## 8.3. Other Useful Emacs Commands

These are a few general Emacs commands which are useful when using Specware. Commands with no key sequences are executed using `M-x name`. XEmacs can be customized using the Options menu. For example, to make `delete` delete the selected region (as in most Word Processing Programs), mouse Options, Editing, Active Regions. To make the changes permanent, mouse Options, Save Options to Init File.

Key	Name	Description
C-h	help	Help options. C-h m gives help about the current mode. Also see Help menu in menubar.
M-/	dabbrev-expand	Does symbol completion based on nearby words in buffer. Repeated key presses find additional completions.
C-sh-middle	mode-motion-copy	Copies the (highlighted) identifier or expression under the mouse to point.
	igrep	Grep for string in files. Brings up buffer with matching lines. Mouse middle on a line to go to it. (In Windows requires installation of Cygwin.)
	igrep-find	Like igrep but searches in all subdirectories.
	fgrep	Like igrep except uses fgrep to search.
	fgrep-find	Like igrep-find except uses fgrep to search. all subdirectories.
C-x C-f	find-file	Prompts for file to edit. tab does filename completion.
C-x d	dired	Prompts for directory to edit. Note that commands <code>sw:process-current-file</code> and <code>cd-current-directory</code> described in the previous section are available in dired mode.
	viper-mode	Vi emulation mode for people who like to edit using vi commands. Documentation is available under C-h i.

