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Department of Artificial Intelligence & Data Science

Experiment No. 7
Kruskal's Algorithm
Date of Performance:
Date of Submission:

Experiment No. 7

Title: Kruskal's Algorithm.

Aim: To study and implement Kruskal's Minimum Cost Spanning Tree Algorithm.

Objective: To introduce Greedy based algorithms

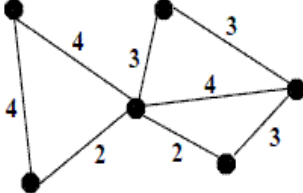
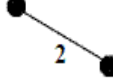
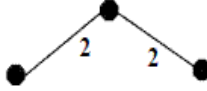
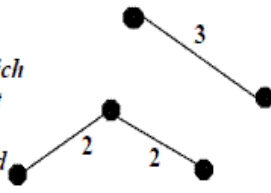
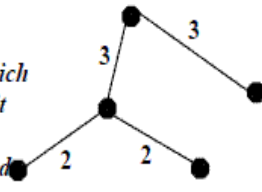
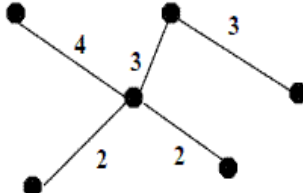
Theory:

Kruskal's algorithm finds a minimum spanning forest of an undirected edge-weighted graph. If the graph is connected, it finds a minimum spanning tree. (A minimum spanning tree of a connected graph is a subset of the edges that forms a tree that includes every vertex, where the sum of the weights of all the edges in the tree is minimized. For a disconnected graph, a minimum spanning forest is composed of a minimum spanning tree for each connected component.) It is a greedy algorithm in graph theory as in each step it adds the next lowest-weight edge that will not form a cycle to the minimum spanning forest.

Example:



Kruskal's Algorithm

<p>1 Given a network.....</p> 	<p>2 Choose the shortest edge (if there is more than one, choose any of the shortest).....</p> 	<p>3 Choose the next shortest edge and add it.....</p> 
<p>4 Choose the next shortest edge which wouldn't create a cycle and add it.</p> 	<p>5 Choose the next shortest edge which wouldn't create a cycle and add it.</p> 	<p>6 Repeat until you have a minimal spanning tree.</p> 

Algorithm and Complexity:



```
1  Algorithm Kruskal(E, cost, n, t)
2  // E is the set of edges in G. G has n vertices. cost[u, v] is the
3  // cost of edge (u, v). t is the set of edges in the minimum-cost
4  // spanning tree. The final cost is returned.
5  {
6      Construct a heap out of the edge costs using Heapify;
7      for i := 1 to n do parent[i] := -1;
8      // Each vertex is in a different set.
9      i := 0; mincost := 0.0;
10     while ((i < n - 1) and (heap not empty)) do
11     {
12         Delete a minimum cost edge (u, v) from the heap
13         and reheapify using Adjust;
14         j := Find(u); k := Find(v);
15         if (j ≠ k) then
16         {
17             i := i + 1;
18             t[i, 1] := u; t[i, 2] := v;
19             mincost := mincost + cost[u, v];
20             Union(j, k);
21         }
22     }
23     if (i ≠ n - 1) then write ("No spanning tree");
24     else return mincost;
25 }
```

Time Complexity is $O(n \log n)$, Where, n = number of Edges

Implementation:

```
#include <stdio.h>
#include <stdlib.h>
```

```
// Structure to represent a weighted edge in the graph
struct Edge {
    int src, dest, weight;
};
```

```
// Structure to represent a connected, undirected graph
struct Graph {
    int V, E;
    struct Edge* edge;
```



```
};
```

```
// Structure to represent a subset for union-find
```

```
struct Subset {
```

```
    int parent;
```

```
    int rank;
```

```
};
```

```
// Create a graph with V vertices and E edges
```

```
struct Graph* createGraph(int V, int E) {
```

```
    struct Graph* graph = (struct Graph*)malloc(sizeof(struct Graph));
```

```
    graph->V = V;
```

```
    graph->E = E;
```

```
    graph->edge = (struct Edge*)malloc(E * sizeof(struct Edge));
```

```
    return graph;
```

```
}
```

```
// Find set of an element i (uses path compression technique)
```

```
int find(struct Subset subsets[], int i) {
```

```
    if (subsets[i].parent != i)
```

```
        subsets[i].parent = find(subsets, subsets[i].parent);
```

```
    return subsets[i].parent;
```

```
}
```

```
// Union of two sets of x and y (uses union by rank)
```

```
void Union(struct Subset subsets[], int x, int y) {
```

```
    int xroot = find(subsets, x);
```

```
    int yroot = find(subsets, y);
```

```
    if (subsets[xroot].rank < subsets[yroot].rank)
```

```
        subsets[xroot].parent = yroot;
```

```
    else if (subsets[xroot].rank > subsets[yroot].rank)
```

```
        subsets[yroot].parent = xroot;
```

```
    else {
```

```
        subsets[yroot].parent = xroot;
```

```
        subsets[xroot].rank++;
```

```
}
```



}

// Compare function for qsort

```
int compare(const void* a, const void* b) {  
    struct Edge* aEdge = (struct Edge*)a;  
    struct Edge* bEdge = (struct Edge*)b;  
    return aEdge->weight - bEdge->weight;  
}
```

// Kruskal's algorithm function

```
void Kruskal(struct Graph* graph) {  
    int V = graph->V;  
    struct Edge result[V]; // Array to store the result MST  
    int e = 0; // Index variable for result array  
    int i = 0; // Index variable for sorted edges array
```

```
// Step 1: Sort all the edges in non-decreasing order of their weight  
qsort(graph->edge, graph->E, sizeof(graph->edge[0]), compare);
```

// Allocate memory for creating V subsets

```
struct Subset* subsets = (struct Subset*)malloc(V * sizeof(struct  
Subset));
```

// Create V subsets with single elements

```
for (int v = 0; v < V; v++) {  
    subsets[v].parent = v;  
    subsets[v].rank = 0;  
}
```

// Number of edges to be taken is equal to V-1

```
while (e < V - 1 && i < graph->E) {
```

// Step 2: Pick the smallest edge

```
struct Edge next_edge = graph->edge[i++];
```

```
int x = find(subsets, next_edge.src);
```

```
int y = find(subsets, next_edge.dest);
```



```
// If including this edge does not cause cycle, include it in result and  
increment the index of result for next edge
```

```
    if (x != y) {  
        result[e++] = next_edge;  
        Union(subsets, x, y);  
    }  
}
```

```
// Print the edges of MST  
printf("Edges of Minimum Spanning Tree:\n");  
for (i = 0; i < e; ++i)  
    printf("%d -- %d == %d\n", result[i].src, result[i].dest,  
result[i].weight);
```

```
    free(subsets);  
}
```

```
int main() {  
    int V, E;  
    printf("Enter number of vertices and edges: ");  
    scanf("%d %d", &V, &E);  
  
    struct Graph* graph = createGraph(V, E);  
  
    printf("Enter edge details (source destination weight):\n");  
    for (int i = 0; i < E; i++)  
        scanf("%d %d %d", &graph->edge[i].src, &graph->edge[i].dest,  
&graph->edge[i].weight);  
    Kruskal(graph);  
  
    return 0;  
}
```

Conclusion: Kruskal's algorithm has been successfully implemented