For NP, we defined it as the class of decision problems for each of which there exists a polynomial-time non-deterministic algorithm.

ProdTwoPrimes(n):

```
if prime_checking(n) then return false prod \leftarrow false Non-deterministically pick x and y \in \{2, ..., n-1\} If n = x * y and prime\_checking(x) and prime\_checking(y) then prod \leftarrow true  return prod
```

Let ProdTwoPrimes be the algorithm for the decision problem "given integer $n \ge 2$, is n the product of exactly two primes", since prime_checking is in **P**, the algorithm is a polynomial-time non-deterministic algorithm, hence ProdTwoPrimes is in **NP**.

Now, by tweaking this algorithm a little bit, invert its output,

${f Co_ProdTwoPrimes}(n)$:

```
if prime_checking(n) then return false prod \leftarrow false Non-deterministically pick x and y \in \{2, ..., n-1\} If n = x * y and prime_checking(x) and prime_checking(y) then prod \leftarrow true  return not prod
```

This makes the algorithm $Co_ProdTwoPrimes$ to output exactly the opposite output of ProdTwoPrimes. And since prime_checking is in **P**, this algorithm is a polynomial-time non-deterministic algorithm, hence it's also in **NP**. Since both the algorithms are in **NP**, the problem ProdTwoPrimes is in ProdTwoPrimes in ProdTwoPrimes is in ProdTwoPrimes in ProdTwoPrimes