

[illegible]

- 1 sign bit
- biased exponent (single: 8 bit, double: 11 bit, extended: 15 bits)
- normalized mantissa (single: 23 bits, double: 52 bits, extended: 64 bits)
- You need to drop the 1 in the mantissa, becomes part of the exponent

Hamming Code:

- Required number of parity bits is  $\log_2 m + 1$

Test: does an AND operation sets CF to zero, SF to MSB and ZF, only if it is zero afterwards

AND = if both are 1 then 1

OR = Either one is One

XOR = they are different

WEEK 4:

CALL

- pushes the offset of the next instruction in the calling procedure onto the system stack.
- Copies the address of the called procedure into EIP
- Executes the called procedure until RET

RET

- Pops the top of stack into EIP
- Syntax RET n, n causes n to be added to the stack pointer after EIP is assigned a value (for variables passed to the stack)

PUSH

- Decrements the stack pointer by 4
- Actual decrements depends on operand

POP

- Copies value at ESP into a register or variable

## Week 5

Activation record:

- Area of the stack used for a procedure's return address, passed parameters, saved registers, and local variables
- Created by the following steps:
  - Calling program pushes arguments onto the stack and calls the procedure
  - The called procedure pushes EBP onto the stack, and sets EBP to ESP

Addressing Modes:

- Register Indirect: Access memory through address in a register
  - mov [edx+12], eax
- Indexed: array name, with "distance" to element in a register
  - mov list[edi], eax

- Base-indexed: starting address in one register, offset in another; add and access memory
  - mov eax, [edx + ecx]

Randomize procedure: must be called once at the beginning of the program.

RandomRange – Generates a random number in [0 .. N – 1]

- Pre: N > 0 in eax
- Post: random integer in [0 .. N-1] in eax
- Range = hi – lo + 1

Week 6

OFFSET

- returns the distance in bytes, of a label from the beginning of its enclosing segment.

PTR

- Overrides the default type of a label, provides the flexibility to access part of a variable
- EX:
  - myDouble DWORD 12345678h
  - mov ax, mydouble -> error
  - mov ax, WORD PTR myDouble -> 5678h
  - mov WORD PTR myDouble, 1357h -> saves 1357h

Little Endian order is used when storing data in memory: in memory 78h 56h 34h 12h  
mov al, BYTE PTR [myDouble + 1] = 56h

TYPE

- Returns the size, in bytes, of a single element of a data declaration
- var1 BYTE
- move eax, TYPE var1 ;1

LENGTHOF

- Counts the number of elements in a single data declaration
- List1 WORD 30 DUP (?) ;30
- Byte1 BYTE 10, 20, 30 ;3
- digitStr BYTE "1234567",0 ;8
- operator returns a value that is equivalent to multiplying LENGTHOF by TYPE

A data declaration spans multiple lines if each line ends with a comma.

mov edx, listD[esi \* TYPE listD]

Note: you can declare a pointer variable that contains the offset of another variable

Example:

List DWORD 100 DUP(?)

Ptr DWORD list

;Contains OFFSET list

## Two Dimensional Arrays:

Example:

Matrix DWORD 5 DUP (3 DUP (?) ); 15 elements

- An elements address is calculated as the base address plus an offset  
BaseAddress + elementSize \* [(row# \* elementsPerRow) + column#]

## String Primitives:

lodsb

- Moves byte at [esi] into the AL register
- Increments esi if direction flag is 0
- Decrements esi if the direction flag is 1

stosb

- Moves byte in the AL register to memory at [edi]
- Increments edi if direction flag is 0
- Decrements edi if direction flag is 1

cld

- Sets the direction flag to 0
- Causes esi and edi to be incremented by lodsb and stosb
- Use for moving "forward" through an array

std

- Sets direction flag to 1
- Causes esi and edi to be decremented by lodsb and stosb
- Used for moving "backward" through an array

ReadInt Algorithm:

Get str

X = 0

for k = 0 to (len(str) – 1)  
if 48 <= str[k] <= 57  
x = 10 \* x + (str[k] – 48)

else

break

Floating Point:

- Pushdown stack
- Operations are defined for the "top" one or two registers
- Registers referenced by name ST(x)
- ST = ST(0) = top of stack
- Instruction Format
  - OPCODE
  - OPCODE destination
  - OPCODE destination, source
- FINIT initialize FPU register stack
- FLD MemVar
  - Push ST(i) "down" to ST(i + 1) for i = 0 .. 6
  - Load ST(0) with Mem Var
- FST MemVar
  - Move top of stack to memory
  - Leave result in ST(0)
- FSTP MemVar
  - Pop top of stack to memory
  - Move ST(i) "up" to ST(i-1) for i=1..7
- FADD: Addition (pop top two, add, push result)
- FSUB: Subtraction
- FMUL: Multiplication
- FDIV: Division
- FDIVR: Division (reverses operands)
- FSIN: Sine (uses radians)
- FCOS: Cosine (uses radians)
- FSQRT: Square Root
- FABS: Absolute Value
- FYL2X: Y\*log2(X) X is in ST(0), Y is in ST(1)
- FYL2XP1: Y \* log2(X) + 1

## Week 7:

Procedure

- During assembly, procedure code is translated once
- During execution, control is transferred to the procedure at each call, may be called many times.

Macro

- Once defined, it can be invoked one or more times

- During assembly, entire macro code is substituted for each call

- A macro must be defined before it can be invoked

Macroname MACRO [param-1, param-2 ..]

Statement-list

ENDM

mWriteStr Macro buffer

push edx  
move dx, OFFSET buffer  
call WriteString  
pop edx

ENDM

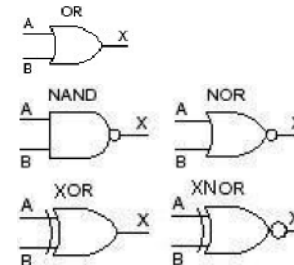
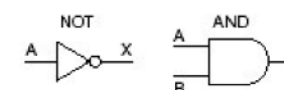
- Should specify that a label is LOCAL

Macro vs Procedure:

- Macros are very convenient
- Macros execute faster than procedure
- Macros are invoked by name
- If macro is called many times, the assembler produces "fat code"
- Use a macro for short code that is called "a few" times and uses only a few registers
- Use a procedure for more complex code or code that is called "many" times

## Boolean Expressions

Func	Log	Boolean
NOT(A)	~A	A/
AND(A,B)	A AND B	AB
OR(A,B)	A OR B	A+B
XOR(A,B)	A XOR B	A⊕B
NAND(A,B)	A NAND B	AB/
NOR(A,B)	A NOR B	A+B/
XNOR(A,B)	A XNOR B	A⊕B



Function of n binary variables has 2<sup>n</sup> possible combinations of values for the variables

Week 8:

- Internal Bus

- Control Unit, ALU, Registers, Addressing Unit communicate via a bus.
- Speed depends on bus width and bus length

Random access memory (RAM)

Read Only Memory (ROM)

- Clock Cycles

- Near light speed
- Clock cycle length determines CPU speed (mostly)

RISC – Reduced Instruction Set

Computer

- Clock Cycles
- Instruction executed directly by hardware
- Only LOAD and STORE instructions reference memory

**Multiprocessor parallelism** – Shared memory

**Multicomputer Parallelism** – distributed memory

- Multi-Processor
  - Difficult to build
  - Relatively Easy to Program
- Multi-Computer
  - Easy to build
  - Extremely difficult to program

Amdahl's Law

Speedup = n / (1 + (n – 1) f)

Total time = f\*T + (1-f)\*T / n

Max speed up = 1/f