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Point values are assigned for each question. Points earned: _____ / 100, = _____ %

"I pledge my honor that I have abided by the Stevens honor system"

- Find a tight upper bound for $f(n) = n^4 + 10n^2 + 5$. Write your answer here: $O(n^4)$ (4 points)

Prove your answer by giving values for the constants c and n_0 . Choose the smallest integer value possible for c . (4 points)

$$c=2$$

$$n_0=4$$

- Find an asymptotically tight bound for $f(n) = 3n^3 - 2n$. Write your answer here: $\Omega(n^3)$ (4 points)

Prove your answer by giving values for the constants c_1 , c_2 , and n_0 . Choose the tightest integer values possible for c_1 and c_2 . (6 points)

$$c_1=2$$

$$c_2=3$$

$$n_0=2$$

- Is $3n - 4 \in \Omega(n^2)$? Circle your answer: yes / **no**. (2 points)

If yes, prove your answer by giving values for the constants c and n_0 . Choose the smallest integer value possible for c . If no, derive a contradiction. (4 points)

$$0 \leq cn^2 \leq 3n - 4 \quad (\forall n \geq n_0)$$

$$3n - 4 \leq 3n - 4n \quad (\forall n \geq 1) = -1n$$

$$cn^2 \leq -1n \quad (\forall n \geq \max(n_0, 1))$$

$$cn \leq -1 \quad (\forall n \geq \max(n_0, 1))$$

$$n \leq -1/c \quad (\forall n \geq \max(n_0, 1))$$

n is able to grow to infinity meaning it is impossible to find a positive constant c such that n is bounded by the constant of $15/c$. This is impossible because n grows with no bound while $15/c$ remains a finite number. Therefore, the c we need doesn't exist and $f(n) = 3n - 4 \notin \Omega(n^2)$?

- Write the following asymptotic efficiency classes in **increasing** order of magnitude.

$O(n^2)$, $O(2n)$, $O(1)$, $O(n \lg n)$, $O(n)$, $O(n!)$, $O(n^3)$, $O(\lg n)$, $O(nn)$, $O(n^2 \lg n)$ (2 points each)

$O(1)$, $O(\lg n)$, $O(n)$, $O(n \lg n)$, $O(n^2)$, $O(n^2 \lg n)$, $O(n^3)$, $O(2^n)$, $O(n!)$,

- Determine the largest size n of a problem that can be solved in time t , assuming that the algorithm takes $f(n)$ milliseconds. Write your answer for n as an integer. (2 points each)

a. $f(n) = n$, $t = 1 \text{ second}$ 1000

b. $f(n) = n \lg n$, $t = 1 \text{ hour}$ 204094

c. $f(n) = n^2$, $t = 1 \text{ hour}$ 1897

d. $f(n) = n^3$, $t = 1 \text{ day}$ 441

e. $f(n) = n!$, $t = 1 \text{ minute}$ 8

- Suppose we are comparing two sorting algorithms and that for all inputs of size n the first algorithm runs in $4n^3$ seconds, while the second algorithm runs in $64n \lg n$ seconds. For which integer values of n does the first algorithm beat the second algorithm? $n \leq 6$ (4 points)

Explain in detail how you got your answer or paste code that solves the problem (2 point):

Set $4n^3 < 64n \lg n$ then simplified to $n^2 < 16 \lg n$. You then go through iterations of n plugging in each one till you reach a value that has a greater value on first algorithm than the second algorithm.

Below is a chart made to show each iteration:

n	n^2	$16 \lg(n)$
1	1	0
2	4	16
3	9	25

4	16	32
5	25	37
6	36	41
7	49	45

In this case that is at $n=7$ where the n^2 reaches 49 and right side is approximately 45 meaning that the second algorithm is now faster.

- Give the complexity of the following methods. Choose the most appropriate notation from among O , Θ , and Ω . (8 points each)

```
int function1(int n) {
    int count = 0;
    for (int i = n / 2; i <= n; i++) {
        for (int j = 1; j <= n; j *= 2) {
            count++;
        }
    }
    return count;
}
```

Answer: $\Theta(n \log n)$

```
int function2(int n) {
    int count = 0;
    for (int i = 1; i * i * i <= n; i++) {
        count++;
    }
    return count;
}
```

Answer: $\Theta(\sqrt[3]{n})$

```
int function3(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            for (int k = 1; k <= n; k++) {
                count++;
            }
        }
    }
}
```

```
    return count;
}
```

Answer: $\Theta(n^3)$

```
int function4(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            count++;
            break;
        }
    }
    return count;
}
```

Answer: $\Theta(n)$

```
int function5(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {
        count++;
    }
    for (int j = 1; j <= n; j++) {
        count++;
    }
    return count;
}
```

Answer: $\Theta(n)$