

Assignment 4

Kevin G

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1. Stacking - push(8): Stack has [8].
push(2): Stack has [8, 2].
pop(): Top element removed [8].
push(pop()*2): pop 8 [2]. Multiply $8 * 2 = 16$. Add to stack [16].
push(10): Stack has [16, 10].
push(pop()/2): pop 10 [16]. Divide $10 / 2 = 5$. Add to stack [16, 5].

Final stack: [16, 5]

2. Queue - push(4): Queue has [4].
push(pop()+4): pop 4 [4]. Add $4 + 4 = 8$. Add to queue [8].
push(8): Queue has [8, 8].
push(pop()/2): pop 8 has [8]. Divide $8 / 2 = 4$. Add to queue [8, 4].
pop(): Front/first element 8 is removed. Queue [4].
pop(): remove last element. Queue [].

Final Queue: [].

3. Find in deque: Just like a DLL, we can modify the deque to be able to traverse from the end of the list. thus making making faster time for some operations. An algorithm could be:

```
Find Deque (Deque<Integer> q, int x) {  
    int left = 0;  
    int right = q.size() - 1;  
    Integer[] dArray = q.toArray(new Integer[0]);  
    while (left <= right) {  
        if (dArray[left] == x) { return left; }  
        if (dArray[right] == x) { return n - 1 - right; }  
        left++;  
        right--;  
    }  
    return -1;  
}
```


Assignment 4 Cont...

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7. Balanced Brackets - Time complexity is $O(n)$ where n is the length of the input string s . This is because we loop through each character once. This takes $O(n)$ time.

As for the rest of the stack operations, for each character we perform $O(1)$ operations like add, remove, and peek, since we do this for every character, the time is $O(n)$.

$$O(n) + O(n) = O(2n) = O(n)$$

Space complexity is $O(n)$ since a worse case would be when the stack is storing opening brackets which grows at size $O(n)$.

Decode String - Time complexity is $O(n)$ where n is the length of the output string. This function iterates over each character once so it's $O(n)$. The stack operations push and pop run in $O(1)$ complexity. Then there is appending string when reaching a closing bracket. While this gets repeated a k amount of times, all substrings are pushed and popped proportional to the length of output string. thus time complexity is $O(n)$.

Space complexity is $O(m+n)$ where m is the length of the input string and n is the length of the output string. the countStack and stringStack store intermediate counts and strings. the m is important to the depth of nested brackets.

Infix to Postfix: Time complexity is $O(n)$ where n is the length of infix array. the loop iterates through each character and does push and pop $O(1)$ operations at most once. each character gets appended which stays at $O(n)$.

Space Complexity is $O(n)$ where n is the infix array since we store the output postfix expression and have a stack that can hold n amounts of elements.