Morday, March 16, 2020 1:16 PM

$$T(n) = aT(\frac{n}{b}) + f(n)$$

$$a = \# d \leq c \cdot problems$$

$$b = size \quad dench \quad subpoolshows$$

$$f(n) = \text{Time} \quad \text{To} \quad cirate \quad subpoolshows}$$

$$f(n) = \text{Time} \quad \text{To} \quad cirate \quad subpoolshows}$$

$$a \geq 1 \quad , \quad b \geq 1$$

$$T(n) = aT(\frac{n}{b}) + f(n)$$

$$if \quad f(n) \in \Theta(n^d)$$

$$so \quad we would \quad d \geq 0$$

$$Know \quad whor \quad d \quad is$$

Know who?
$$d$$
 is

$$\begin{cases}
\Theta(nd) & \text{if } a < b^d \\
\Theta(n^d | gn) & \text{if } a = b^d \\
\Theta(n^{log_ba}) & \text{if } a > b^d
\end{cases}$$

Example #1

$$T(n) = 9T(\frac{n}{3}) + h$$

epch of 5.20

$$T(n) = 71 (3) + 1$$

$$f(n) c \theta(n')$$

$$(n \epsilon \theta(n')) = 3 - 1$$

$$CA \neq 1 ?$$

$$a < b$$

$$q < 3$$

$$q < 3$$

$$N0$$

$$\frac{\text{Cose 2?}}{a = b^{3}}$$

$$9 = 3 \text{ NO}$$

$$T(n) \in \Theta(n^{\log_6 \alpha})$$

$$= \Theta(n^2)$$

$$= \Theta(n^2)$$

Master Method Example # 2

$$T(n) = T\left(\frac{2n}{3}\right) + 1$$

$$a = 1 b = \frac{3}{2}$$

$$f(n) \in \Theta(n^{\circ})$$

$$a < b^{\alpha}$$

$$1 < \left(\frac{3}{2}\right)^{\alpha}$$

$$a = b^{\alpha}$$

$$1 = \left(\frac{3}{2}\right)^{6}$$

we CAN USE COSE

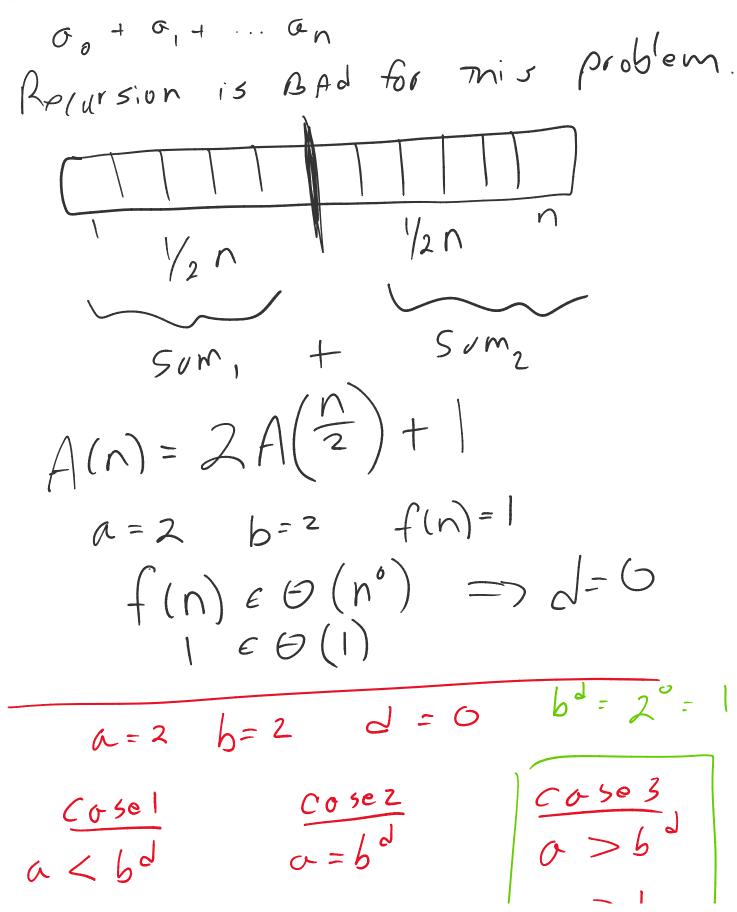
Case# 3?

Cose 2: By mm

$$=\Theta(1.|gn)$$
$$=\Theta(|gn)$$

Levitin's problem form the beginning of the divide-and-conquer chapter.

Monday, March 16, 2020 2:53 PM



$$a < b^{\alpha}$$

$$2 < 1$$

$$2 = 1$$

$$4 = 0$$

$$A(n) \in \Theta(n)$$

$$= \Theta(n)$$

$$= \Theta(n)$$

$$= \Theta(n)$$