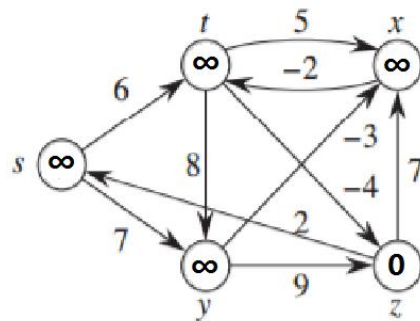
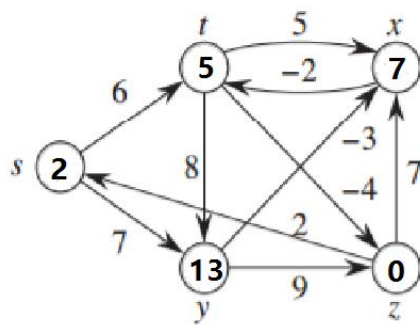


1.

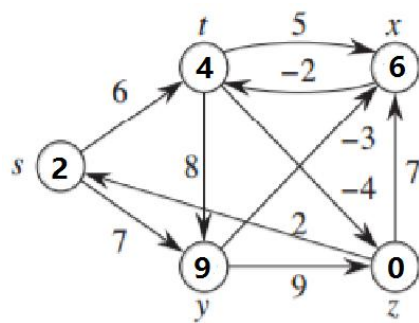
Initial:



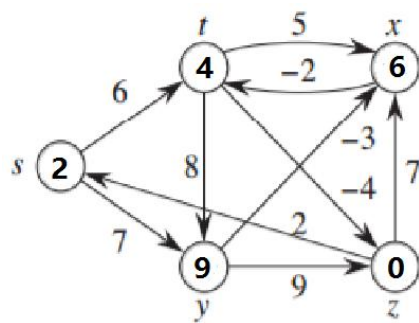
First Pass:



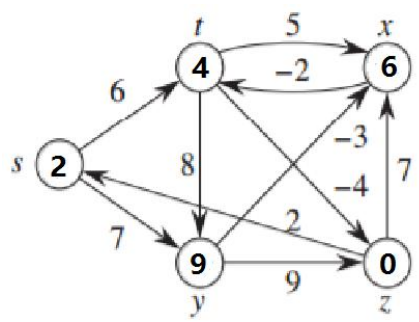
Second Pass:



Third Pass:



Fourth Pass:



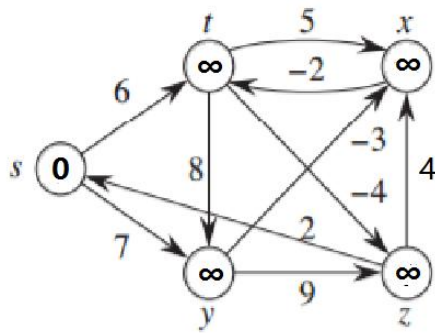
d	s	t	x	y	z
	∞	∞	∞	∞	0
1	2	5	7	13	0
2	2	4	6	9	0
3	2	4	6	9	0
4	2	4	6	9	0

π	s	t	x	y	z
	NIL	NIL	NIL	NIL	NIL
1	z	x	z	t	NIL
2	z	x	y	s	NIL
3	z	x	y	s	NIL
4	z	x	y	s	NIL

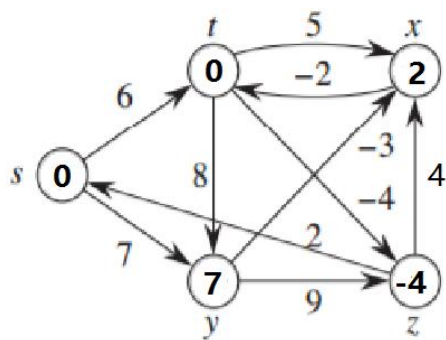
Return value is True.

Change weight of (z, x) to 4.

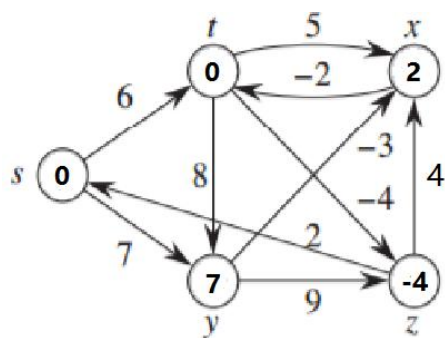
Initial:



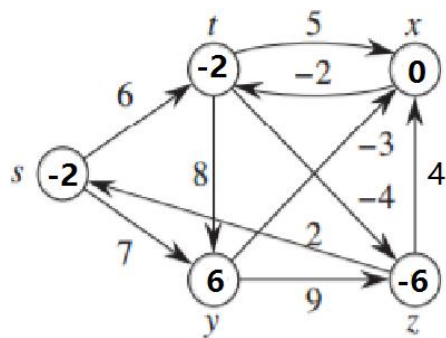
First Pass:



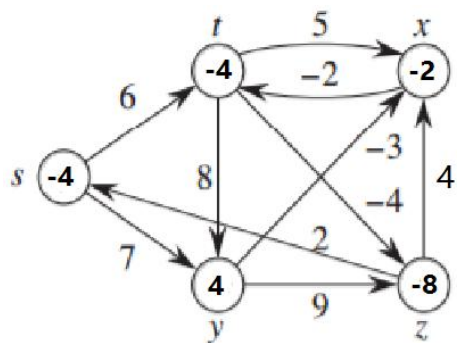
Second Pass:



Third Pass:



Fourth Pass:

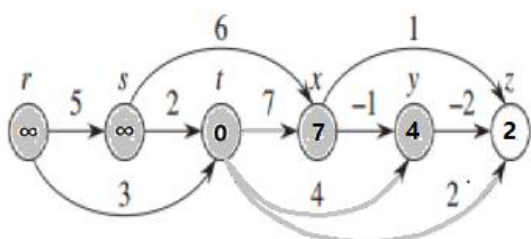
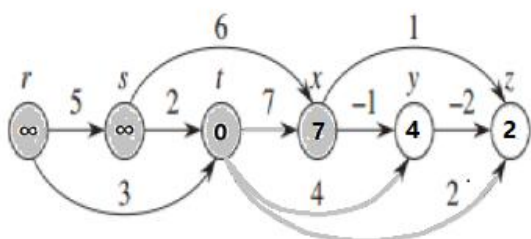
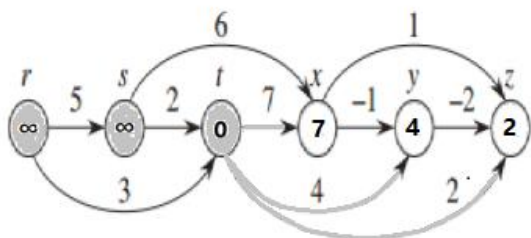
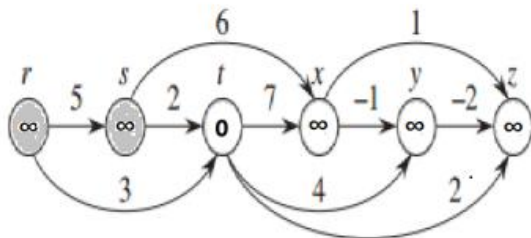
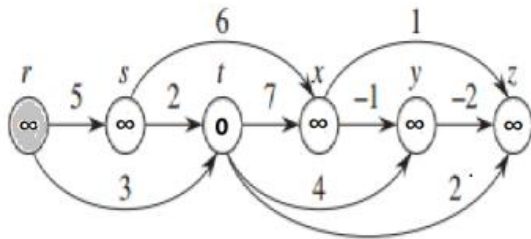
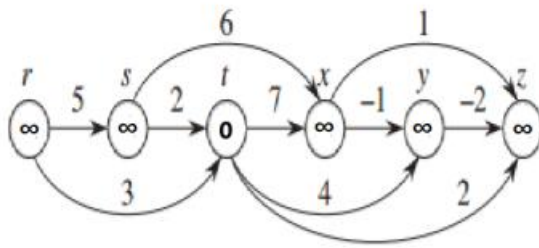


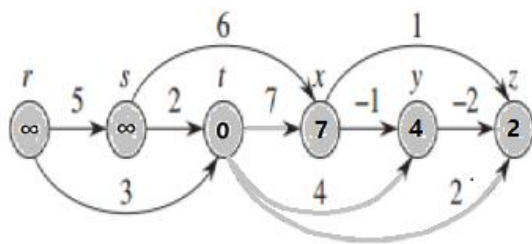
d	s	t	x	y	z
	0	∞	∞	∞	∞
1	0	2	4	7	-2
2	0	0	2	7	-4
3	-2	-2	0	6	-6
4	-4	-4	-2	4	-8

π	s	t	x	y	z
	NIL	NIL	NIL	NIL	NIL
1	NIL	x	y	s	t
2	NIL	x	x	s	t
3	z	x	x	s	t
4	z	x	x	s	t

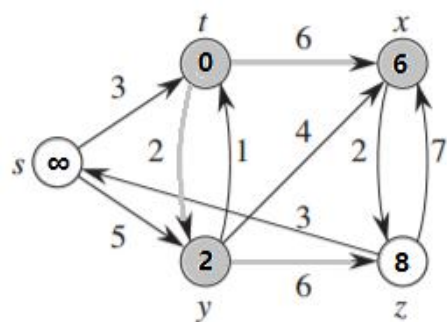
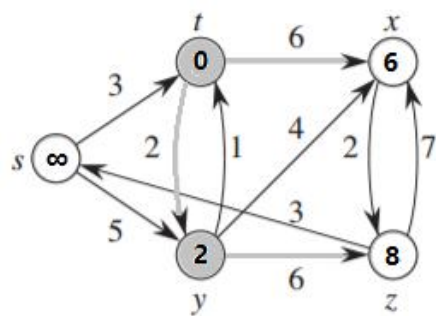
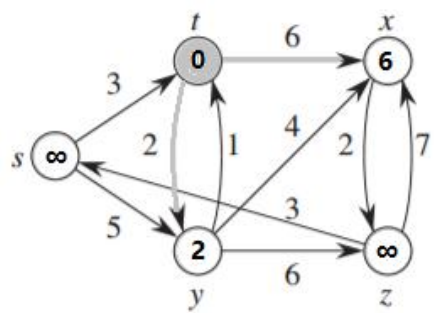
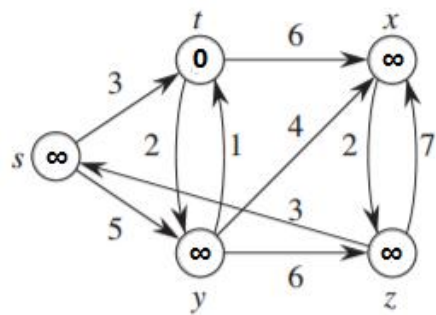
Return value is False.

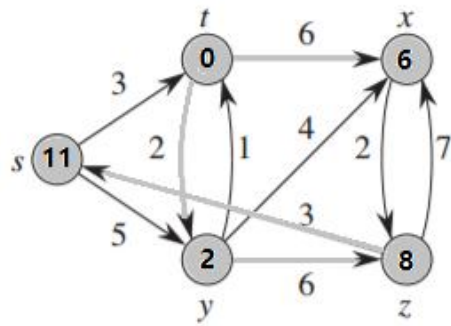
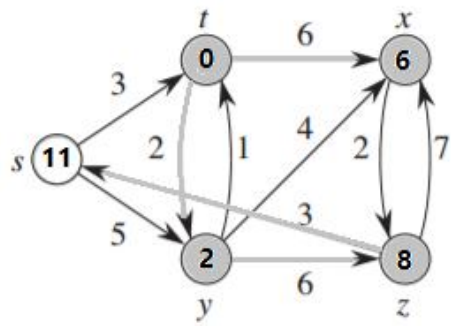
2.





3.





d	s	t	x	y	z
	∞	0	∞	∞	∞
1	∞	0	6	2	∞
2	∞	0	6	2	8
3	∞	0	6	2	8
4	11	0	6	2	8
5	11	0	6	2	8

π	s	t	x	y	z
	NIL	NIL	NIL	NIL	NIL
1	NIL	x	y	s	t
2	NIL	x	x	s	t
3	NIL	x	x	s	t
4	z	x	x	s	t
5	z	x	x	s	t

4.

RELIABILITY(G, r, x, y)

INITIALIZE-SINGLE-SOURCE(G, x) **Initialize to be negative infinite**

$S = \emptyset$

$Q = G.V$

while $Q \neq \emptyset$ **do**

$u = \text{EXTRACT-MAX}(Q)$

$S = S \cup \{u\}$

for each vertex $v \in G.Adj[u]$ **do**

if $v.d > u.d + r(u, v)$ **then**

$v.d = u.d + r(u, v)$

$v.\pi = u$

end if

end for

end while

while $y \neq x$ **do**

 Print y

$y = y.\pi$

end while

Print x

5.

MODIFIED-DIJKSTRA(G, w, s)

for each $v \in G.V$ **do** $v.d = VW + 1$ $v.\pi = NIL$

end for

$s.d = 0$

Initialize an array A of length $VW + 2$

$A[0].insert(s)$

Set $A[VW + 1]$ equal to a linked list containing every vertex except s

$k = 0$

for $i = 1$ to $|V|$ **do**

while $A[k] = NIL$ **do**

$k = k + 1$

end while

$u = A[k].head$

$A[k].delete(u)$

for each vertex $v \in G.Adj[u]$ **do**

if $v.d > u.d + w(u, v)$ **then**

$A[v.d].delete(v.list)$

$v.d = u.d + w(u, v)$

$v.\pi = u$

$A[v.d].insert(v)$

$v.list = A[v.d].head$

end if

end for

end for

6.

k	D^k
0	$\begin{pmatrix} 0 & \infty & \infty & \infty & -1 & \infty \\ 1 & 0 & \infty & 2 & \infty & \infty \\ \infty & 2 & 0 & \infty & \infty & -8 \\ -4 & \infty & \infty & 0 & 3 & \infty \\ \infty & 7 & \infty & \infty & 0 & \infty \\ \infty & 5 & 10 & \infty & \infty & 0 \end{pmatrix}$
1	$\begin{pmatrix} 0 & \infty & \infty & \infty & -1 & \infty \\ 1 & 0 & \infty & 2 & 0 & \infty \\ \infty & 2 & 0 & \infty & \infty & -8 \\ -4 & \infty & \infty & 0 & -5 & \infty \\ \infty & 7 & \infty & \infty & 0 & \infty \\ \infty & 5 & 10 & \infty & \infty & 0 \end{pmatrix}$

2	$\begin{pmatrix} 0 & \infty & \infty & \infty & -1 & \infty \\ 1 & 0 & \infty & 2 & 0 & \infty \\ 3 & 2 & 0 & 4 & 2 & -8 \\ -4 & \infty & \infty & 0 & -5 & \infty \\ 8 & 7 & \infty & 9 & 0 & \infty \\ 6 & 5 & 10 & 7 & 5 & 0 \end{pmatrix}$
3	$\begin{pmatrix} 0 & \infty & \infty & \infty & -1 & \infty \\ 1 & 0 & \infty & 2 & 0 & \infty \\ 3 & 2 & 0 & 4 & 2 & -8 \\ -4 & \infty & \infty & 0 & -5 & \infty \\ 8 & 7 & \infty & 9 & 0 & \infty \\ 6 & 5 & 10 & 7 & 5 & 0 \end{pmatrix}$
4	$\begin{pmatrix} 0 & \infty & \infty & \infty & -1 & \infty \\ -2 & 0 & \infty & 2 & -3 & \infty \\ 0 & 2 & 0 & 4 & -1 & -8 \\ -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & 7 & \infty & 9 & 0 & \infty \\ 3 & 5 & 10 & 7 & 2 & 0 \end{pmatrix}$
5	$\begin{pmatrix} 0 & 6 & \infty & 8 & -1 & \infty \\ -2 & 0 & \infty & 2 & -3 & \infty \\ 0 & 2 & 0 & 4 & -1 & -8 \\ -4 & 2 & \infty & 0 & -5 & \infty \\ 5 & 7 & \infty & 9 & 0 & \infty \\ 3 & 5 & 10 & 7 & 2 & 0 \end{pmatrix}$
6	$\begin{pmatrix} 0 & 6 & \infty & 8 & -1 & \infty \\ -2 & 0 & \infty & 2 & -3 & \infty \\ -5 & -3 & 0 & -1 & -6 & -8 \\ -4 & 2 & \infty & 0 & -5 & \infty \\ 5 & 7 & \infty & 9 & 0 & \infty \\ 3 & 5 & 10 & 7 & 2 & 0 \end{pmatrix}$

7.

FLOYD-WARSHALL(W)

$n = W.rows$

$D^{(0)} = W$

for $k = 1$ **to** n

let $D^{(k)} = (d_{ij}^{(k)})$, $\Phi^{(k)} = (\phi_{ij}^{(k)})$ be a new $n * n$ matrix

for $i = 1$ **to** n

for $j = 1$ **to** n

if $d_{ij}^{(k-1)} > d_{ik}^{(k-1)} + d_{kj}^{(k-1)}$

$d_{ij}^{(k)} = d_{ik}^{(k-1)} + d_{kj}^{(k-1)}$, $\phi_{ij}^{(k)} = k$

else

$d_{ij}^{(k)} = d_{ij}^{(k-1)}$, $\phi_{ij}^{(k)} = \phi_{ij}^{(k-1)}$

return $D^{(n)}$

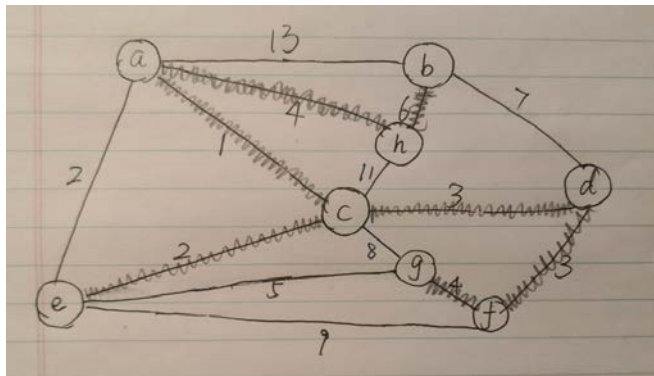
```

PRINT-ALL-PAIRS-SHORTEST-PATH( $\Phi$ , i, j)
  if i == j
    print i
  elseif  $\Phi_{ij} == \text{NIL}$ 
    print "no path between i and j"
  else
    PRINT-ALL-PAIRS-SHORTEST-PATH( $\Pi$ , i,  $\Phi_{ij}$ )
    print  $\Phi_{ij}$ 
    PRINT-ALL-PAIRS-SHORTEST-PATH( $\Pi$ ,  $\Phi_{ij}$ , j)

```

8.

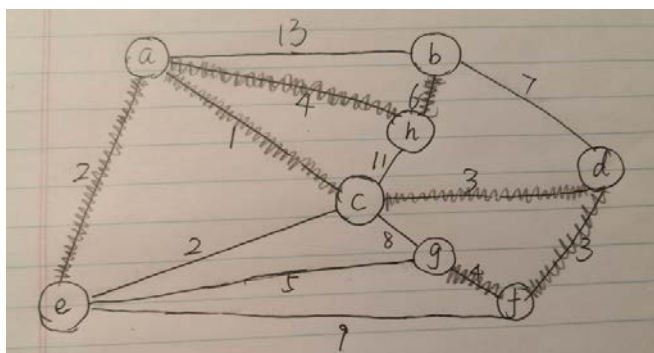
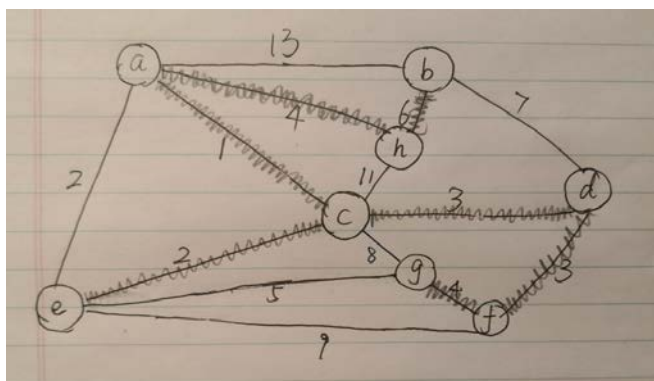
We can get one of the MST shown as follows:



So, the cost is $1 + 2 + 3 + 3 + 4 + 4 + 6 = 23$.

(b)

We can get 2 different MST:



(c)

For the second MST above, the order of the edges is (a, c) , (a, e) , (d, f) , (c, d) , (a, h) , (f, g) , (b, h) .

For (a, c) , We have $S = \{a\}$, $V - S = \{b, c, d, e, f, g, h\}$, so we add a safe edge (a, c) that $w(a, c) = 1$. And we will get $A = \{(a, c)\}$.

For (a, e) , We have $S = \{a, c\}$, $V - S = \{b, d, e, f, g, h\}$, so we add a safe edge (a, e) that $w(a, e) =$

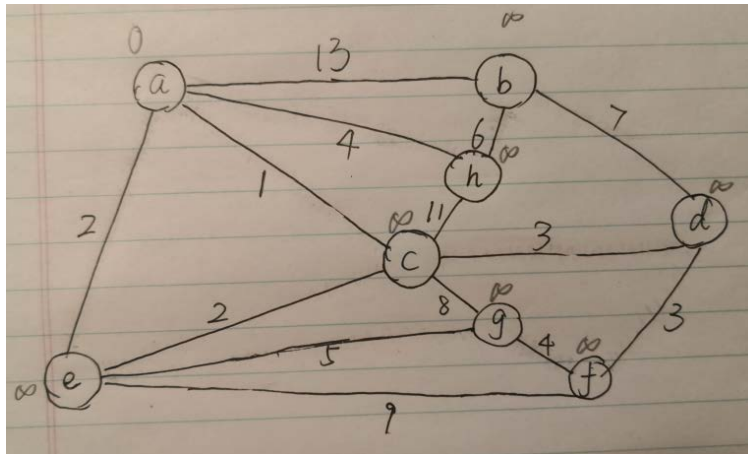
2. And we will get $A = \{(a, c), (a, e)\}$.

For (d, f) , We have $S = \{d\}$, $V - S = \{a, b, c, e, f, g, h\}$, so we add a safe edge (d, f) that $w(d, f) =$

3. And we will get $A = \{(a, c), (a, e), (d, f)\}$.

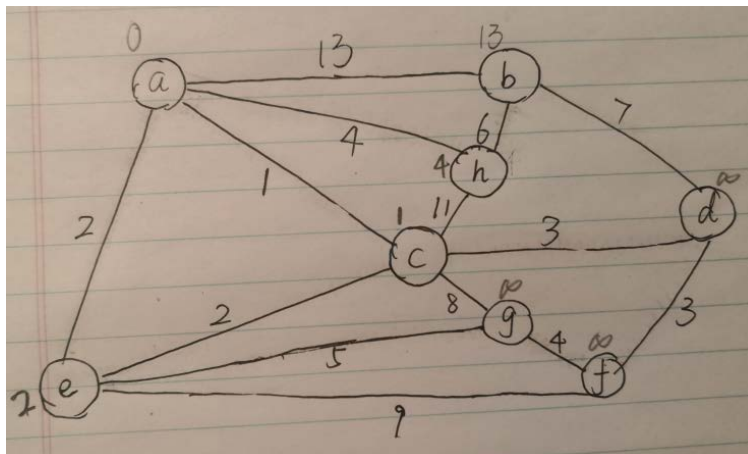
9.

At the beginning: $A = \{\}$, $Q = \{a, b, c, d, e, f, g, h\}$ and $a.key = 0$, other keys are ∞ .

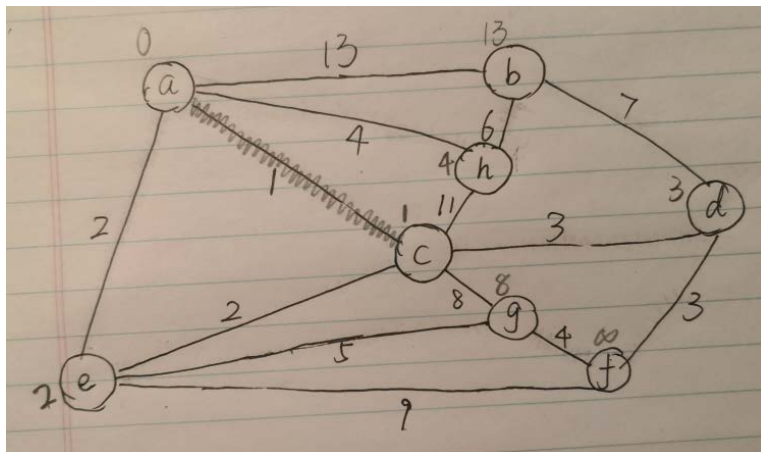


Because $a.key$ is the smallest, we will extract it and update keys for those vertices adjacent to a .

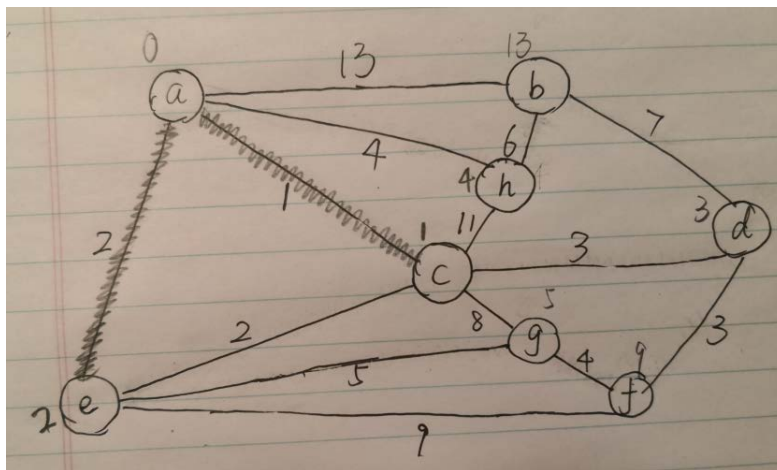
Now, $A = \{\}$, $Q = \{b, c, d, e, f, g, h\}$.



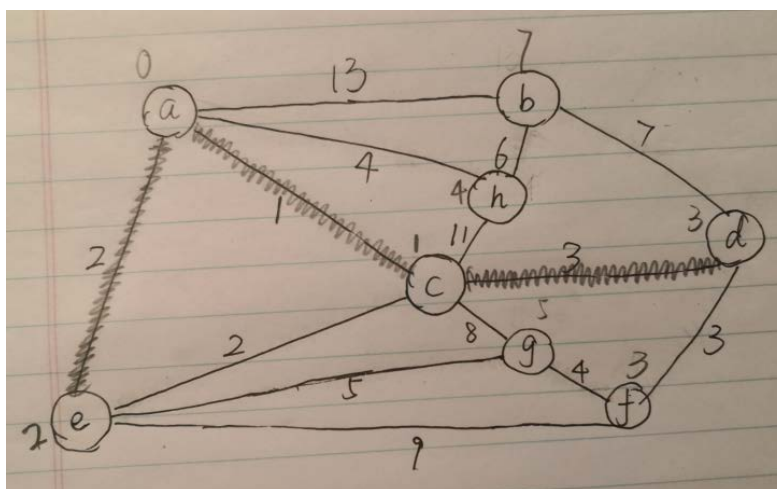
Because $c.key$ is the smallest, we will extract it and add the edge (a, c) to A . Also, updating keys for those vertices adjacent to c . Now, $A = \{(a, c)\}$, $Q = \{b, d, e, f, g, h\}$.



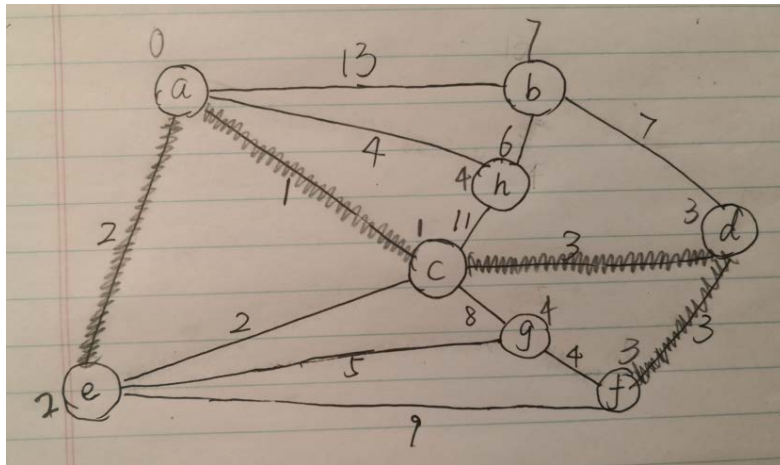
Because e.key is the smallest, we will extract it and add the edge (a, e) to A. Also, updating keys for those vertices adjacent to e. Now, $A = \{(a, c), (a, e)\}$, $Q = \{b, d, f, g, h\}$.



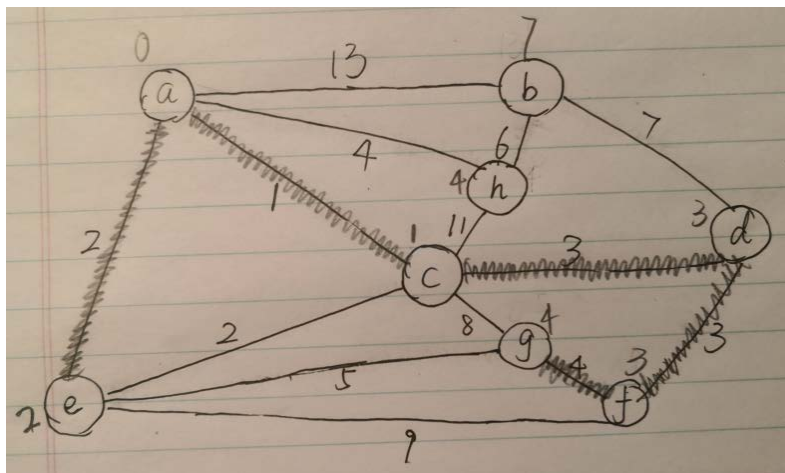
Because d.key is the smallest, we will extract it and add the edge (c, d) to A. Also, updating keys for those vertices adjacent to d. Now, $A = \{(a, c), (a, e), (c, d)\}$, $Q = \{b, f, g, h\}$.



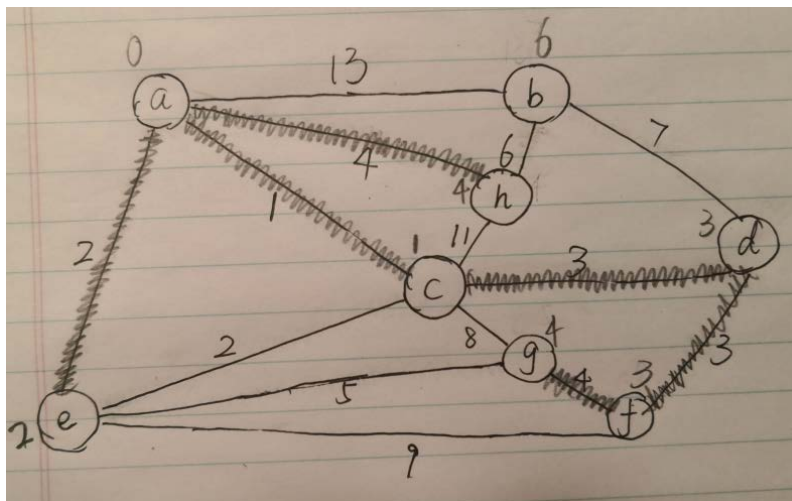
Because f.key is the smallest, we will extract it and add the edge (d, f) to A. Also, updating keys for those vertices adjacent to f. Now, $A = \{(a, c), (a, e), (c, d), (d, f)\}$, $Q = \{b, g, h\}$.



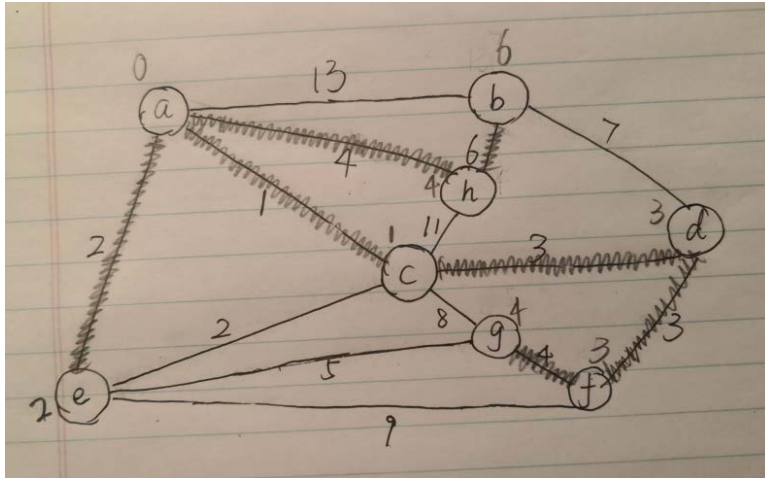
Because h.key and g.key are the smallest, we choose g according to the order of alphabet, we will extract it and add the edge (f, g) to A. Also, updating keys for those vertices adjacent to g. Now, $A = \{(a, c), (a, e), (c, d), (d, f), (f, g)\}$, $Q = \{b, h\}$.



Because h.key is the smallest, we will extract it and add the edge (a, h) to A. Also, updating keys for those vertices adjacent to h. Now, $A = \{(a, c), (a, e), (c, d), (d, f), (a, h)\}$, $Q = \{b\}$.



Because $b.key$ is the smallest, we will extract it and add the edge (h, b) to A . Also, updating keys for those vertices adjacent to b . Now, $A = \{(a, c), (a, e), (c, d), (d, f), (a, h), (b, h)\}$, $Q = \{\}$.



Because $Q = \emptyset$, we use Prim algorithm get the MST, which is shown as above. And the order the vertex are removed from Q is a, e, d, f, g, h, b .