# CS 6041 Theory of Computation

#### **Deterministic finite automata**

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https://kevinsuo.github.io/

### **Outline**

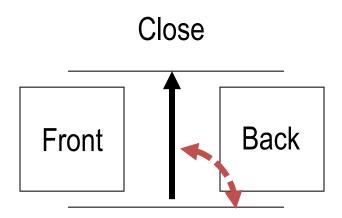
#### Finite Automata

- Definition
- Example
- Language of DFA
- Computation for DFAs

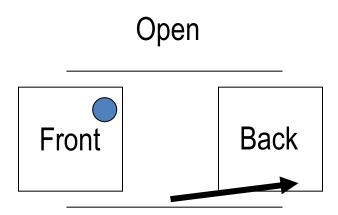
### Design DFAs

- Example
- Regular language
- Regular operation

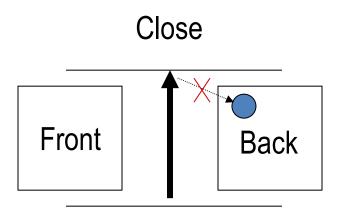
### **Automatic door**



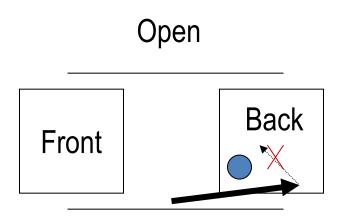
# **Automatic door (2)**



# **Automatic door (3)**



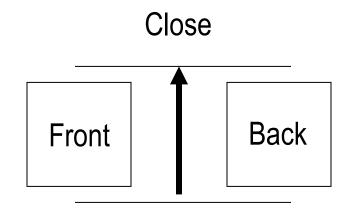
# **Automatic door (4)**



### Finite automata: automatic door

#### Two states

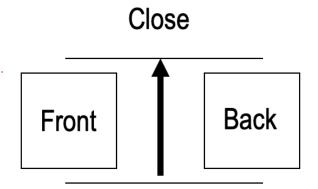
- Open
- Close



### Four inputs (4 possible cases)

- Front: a person is standing on the pad in front of the doorway
- Rear: a person is standing on the pad to the rear of the doorway
- Both: people are standing on both pads
- Neither: no one is standing on either pad

# State table and diagram

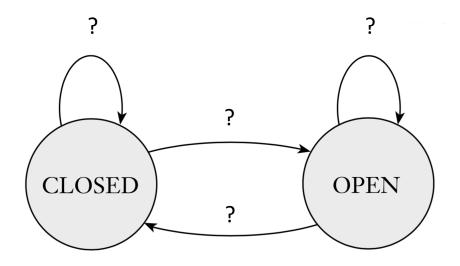


|        | Neither | Front | Rear   | Both   |
|--------|---------|-------|--------|--------|
| Closed | Closed  | Open  | Closed | Closed |
| Open   | Closed  | Open  | Open   | Open   |

# State table and diagram

input signal

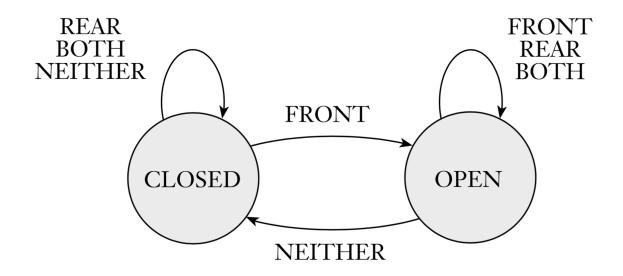
|        | NEITHER | FRONT | REAR   | ВОТН   |
|--------|---------|-------|--------|--------|
| CLOSED | CLOSED  | OPEN  | CLOSED | CLOSED |
| OPEN   | CLOSED  | OPEN  | OPEN   | OPEN   |



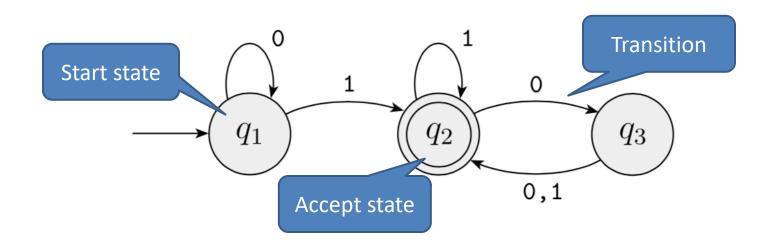
# State table and diagram

#### input signal

|        | NEITHER | FRONT | REAR   | BOTH   |
|--------|---------|-------|--------|--------|
| CLOSED | CLOSED  | OPEN  | CLOSED | CLOSED |
| OPEN   | CLOSED  | OPEN  | OPEN   | OPEN   |



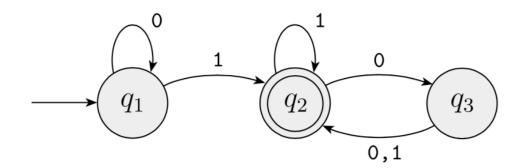
### State diagram of finite automata



- Start state: indicated by the arrow pointing at it from nowhere
- Accept state: is the one with a double circle
- Transitions: the arrows going from one state to another with input
- Output: is either accept or reject. If input is 1101, output is accept.

# Automata M<sub>1</sub>

State diagram

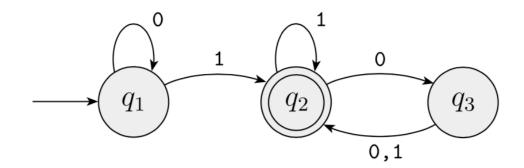


State transition table?

|       | 0 | 1 |
|-------|---|---|
| $q_1$ |   |   |
| $q_2$ |   |   |
| $q_3$ |   |   |

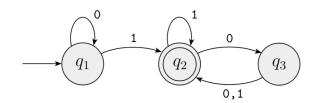
# Automata M<sub>1</sub>

State diagram

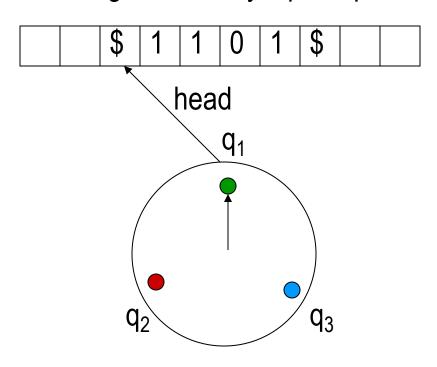


State transition table

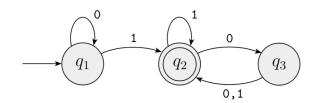
|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_3$ | $q_2$ |
| $q_3$ | $q_2$ | $q_2$ |



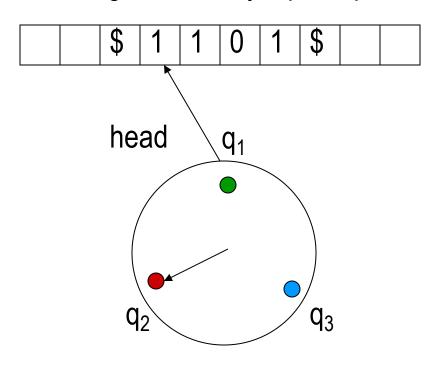
#### Single read-only input tape



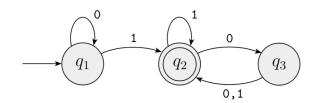
Finite state controller



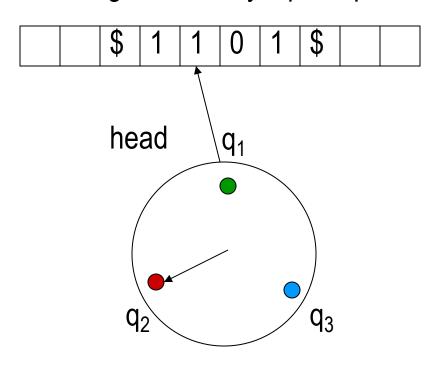
#### Single read-only input tape



Finite state controller

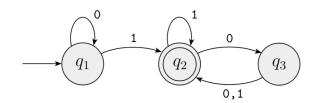


#### Single read-only input tape

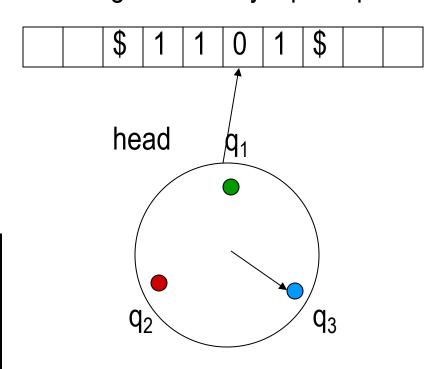


Finite state controller

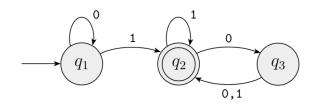
|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_3$ | $q_2$ |
| $q_3$ | $q_2$ | $q_2$ |



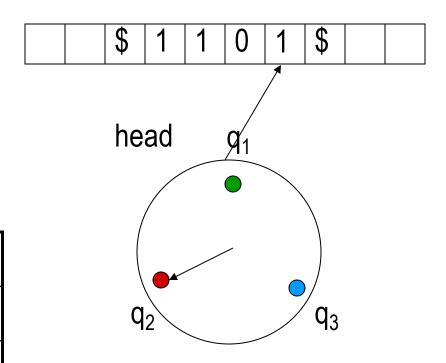
#### Single read-only input tape



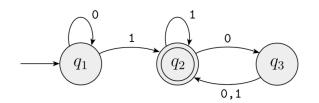
Finite state controller



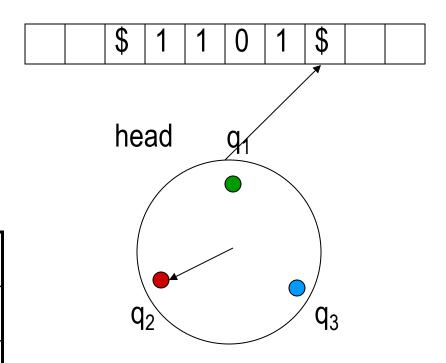
#### Single read-only input tape



Finite state controller



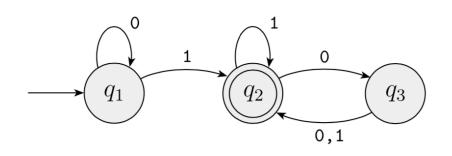
#### Single read-only input tape



Finite state controller

# What strings will M<sub>1</sub> accept

|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_3$ | $q_2$ |
| $q_3$ | $q_2$ | $q_2$ |



0\*1 0\*1 1\* 0\*1 1\*(01)\*

0\*1 1\*00 0\*1 1\*(00)\*

### **Definition of finite automaton**

- Finite automaton is a 5-tuple M=(Q, $\Sigma$ , $\delta$ ,q<sub>0</sub>,F)
  - Q: finite set called states
  - $\circ$   $\Sigma$ : finite set called the alphabet
  - $\delta$ : Q×Σ→Q, transition function
  - $\circ$  q<sub>0</sub>∈Q: start state
  - F⊆Q: accept states

• Language on M: L(M) =  $\{w \in \Sigma^* \mid \delta(q_0, w) \in F\}$ 

•  $M_2 = (\{q_1, q_2\}, \{0, 1\}, \delta, q_1, \{q_2\})$ 

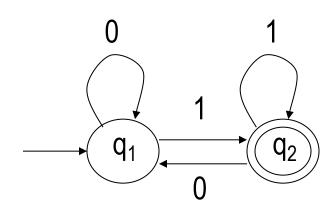
Can anyone draw the DFA?

δ

|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_1$ | $q_2$ |

•  $M_2 = (\{q_1, q_2\}, \{0, 1\}, \delta, q_1, \{q_2\})$ 

#### Can anyone draw the DFA?



δ

|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_1$ | $q_2$ |

$$L(M_2) = \{ w \mid w \text{ ends with 1s } \}$$

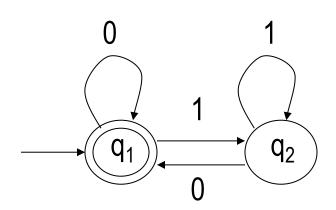
•  $M_3 = (\{q_1, q_2\}, \{0, 1\}, \delta, q_1, \{q_1\})$ 

Can anyone draw the DFA?

|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_1$ | $q_2$ |

•  $M_3 = (\{q_1, q_2\}, \{0, 1\}, \delta, q_1, \{q_1\})$ 

#### Can anyone draw the DFA?



|       | 0     | 1     |
|-------|-------|-------|
| $q_1$ | $q_1$ | $q_2$ |
| $q_2$ | $q_1$ | $q_2$ |

$$L(M_3) = \{ w \mid w = \varepsilon \text{ or } w \text{ ends with } 0s \}$$

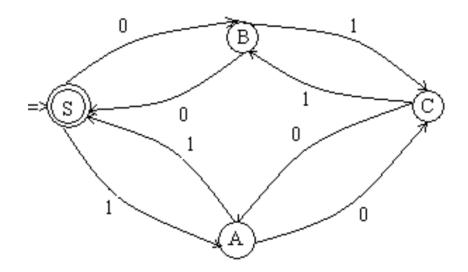
•  $M = (\{S,A,B,C\},\{0,1\},f,S,\{S\})$ 

- $\circ$  f(S,0)=B, f(B,0)=S
- $\circ$  f(S,1)=A, f(B,1)=C
- $\circ$  f(A,0)=C, f(C,0)=A
- f(A,1)=S, f(C,1)=B

Can anyone draw the DFA?

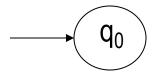
- $M = (\{S,A,B,C\},\{0,1\},f,S,\{S\})$ 
  - $\circ$  f(S,0)=B, f(B,0)=S
  - $\circ$  f(S,1)=A, f(B,1)=C
  - $\circ$  f(A,0)=C, f(C,0)=A
  - f(A,1)=S, f(C,1)=B

Can anyone draw the DFA?



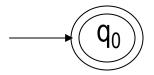
- Ø
- M =( $\{q_0\}, \{\}, f, q_0, \{\})$

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- M =( $\{q_0\},\{\},f,q_0,\{\})$

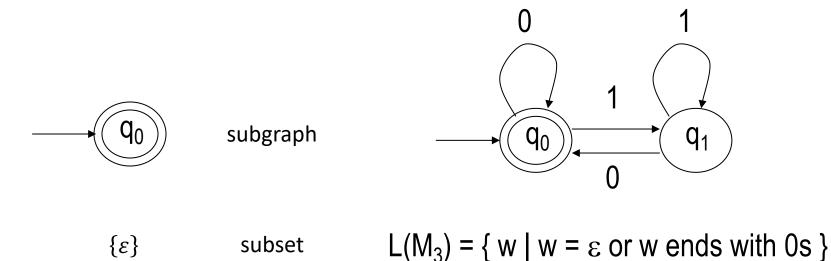


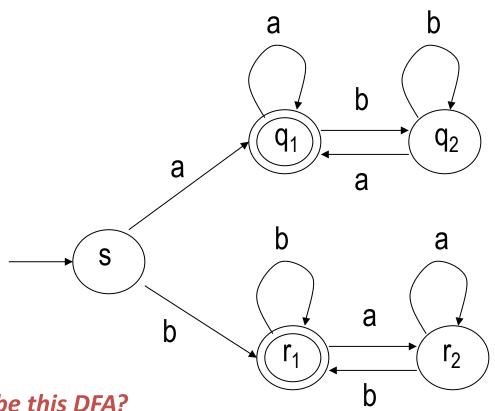
- {*E*}
- M =( $\{q_0\},\{\},f,q_0,\{q_0\}$ )

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- M =( $\{q_0\},\{\},f,q_0,\{q_0\}$ )

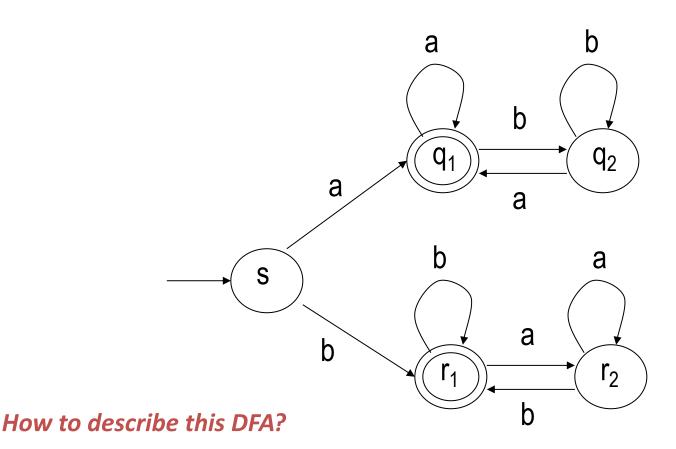


- {*E*}
- M =( $\{q_0\},\{\},f,q_0,\{q_0\}$ )

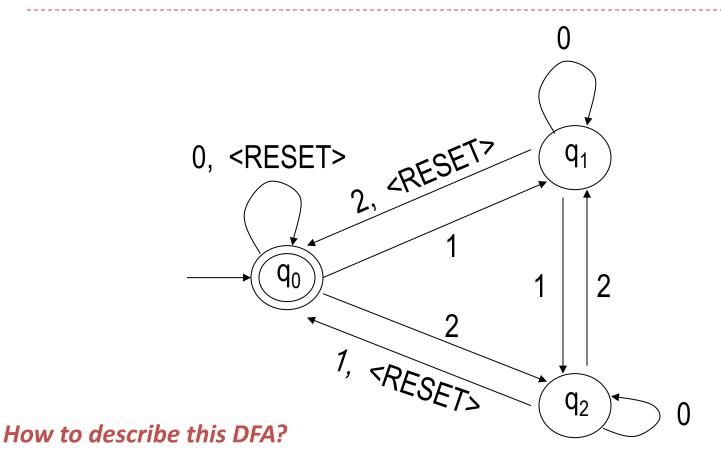




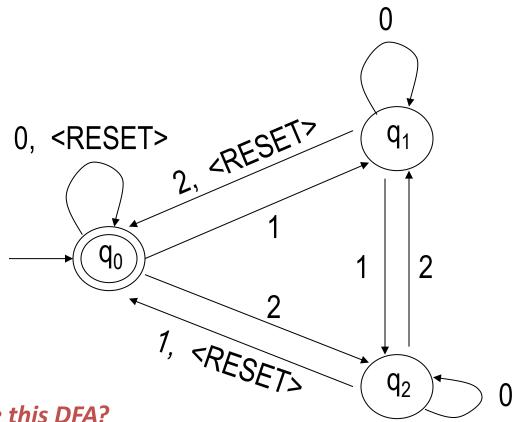
How to describe this DFA?



 $L(M_4) = \{ w \mid w \text{ starts and ends with the same letter} \}$ 



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How to describe this DFA?

 $L(M_5) = \{ w \mid after the last < RESET>, the sum of w is 0 modulo 3 \}$ 

### Formal definition of computation for DFAs

- Finite automaton  $M=(Q,\Sigma,\delta,q_0,F)$ ;
  - Input string w=w<sub>1</sub>w<sub>2</sub>...w<sub>n</sub>;
- Computation: a sequence of states r<sub>0</sub>,r<sub>1</sub>,..., r<sub>n</sub>;
  - $r_0 = q_0$ ; •  $\delta(r_i, w_{i+1}) = r_{i+1}$ ; (i=0,1,...,n-1)
- Accept computation: M accepts w
  - $\circ$   $r_n \in F$ ;
- M recognizes language A:
  - o L(M)={ w | M accepts w }