

# CS 6041

# Theory of Computation

## Reducibility

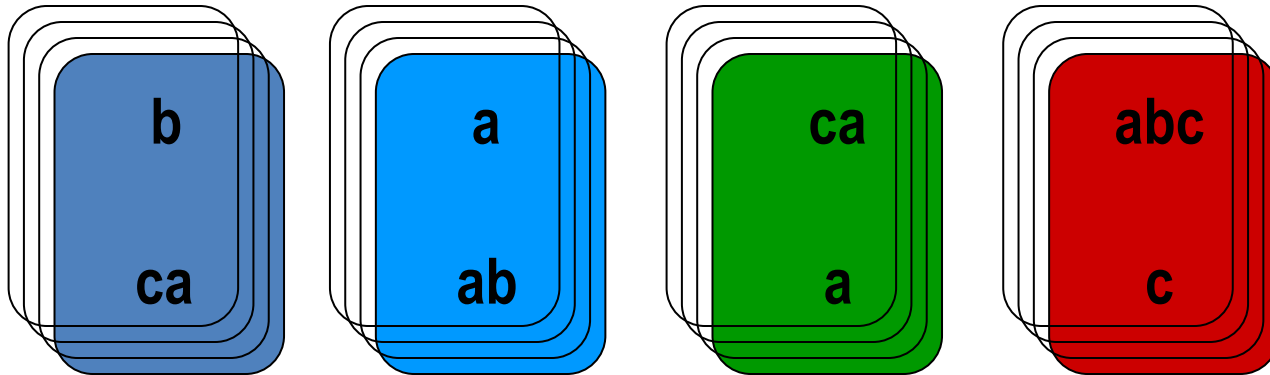
**Kun Suo**

Computer Science, Kennesaw State University

<https://kevinsuo.github.io/>

# Post Correspondence Problem (PCP)

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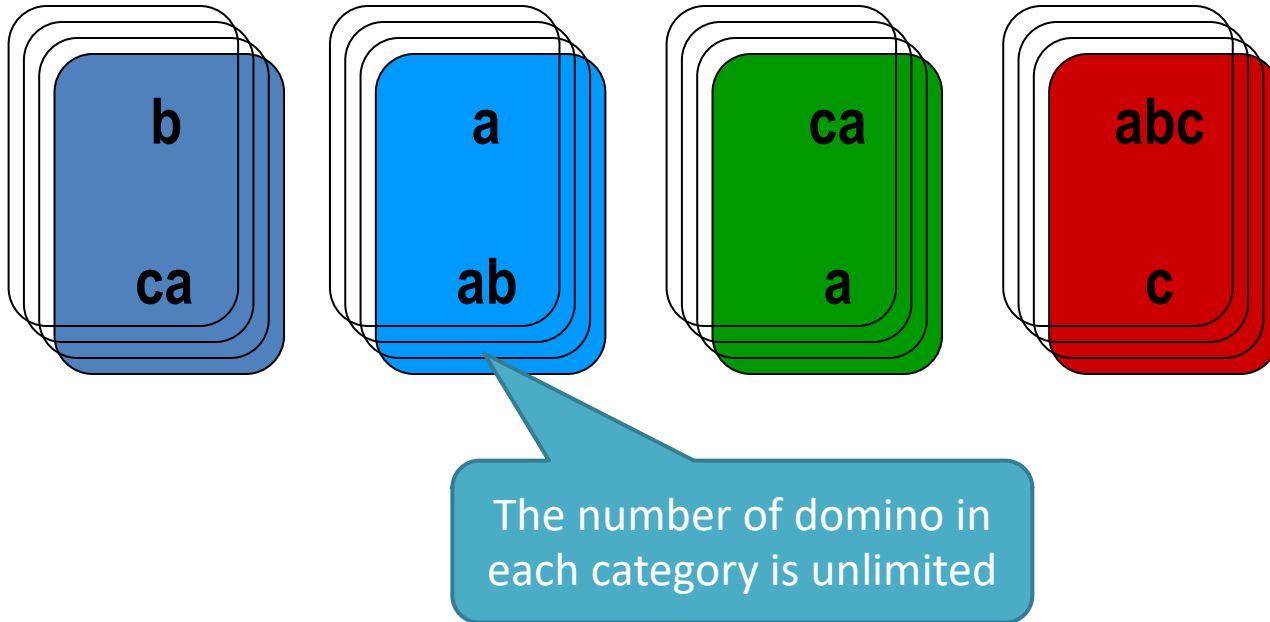


We have limited  
categories of dominos



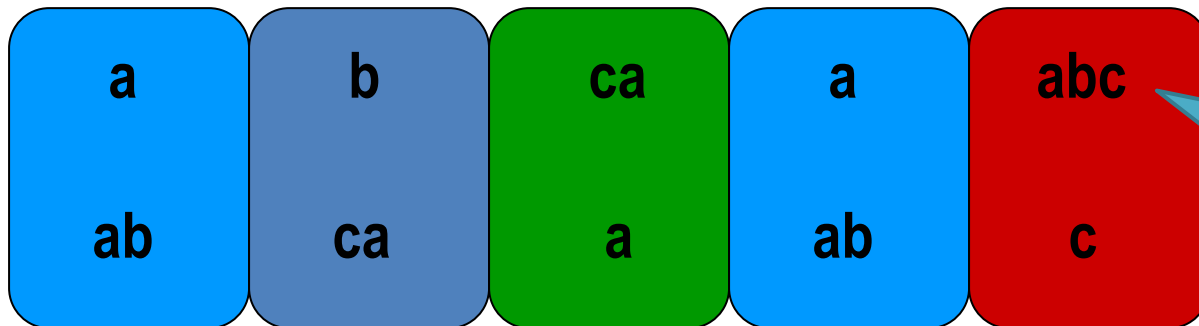
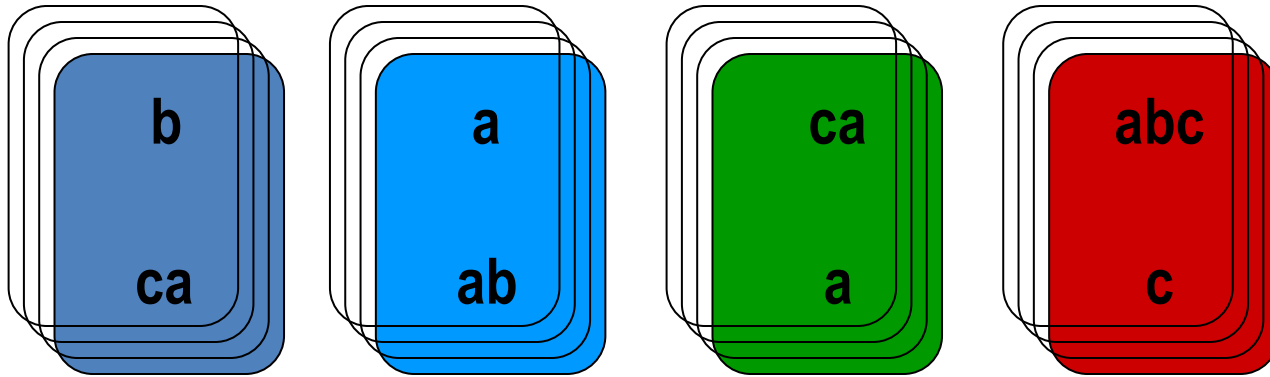
# Post Correspondence Problem (PCP)

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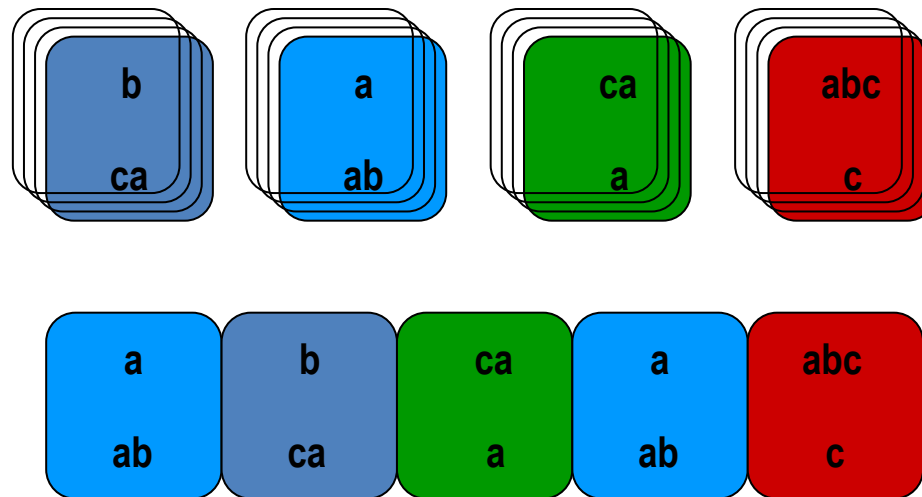


Match:  
Upper string: abcaaabc  
Bottom string: abcaaabc



# Post Correspondence Problem (PCP)

- Whether a collection of dominos has a match



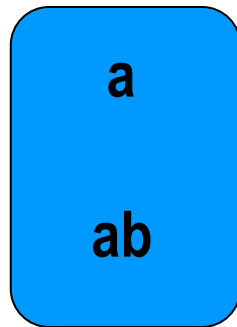
- $PCP = \{ \langle P \rangle \mid P \text{ is an instance of the Post Correspondence Problem with a match} \}.$

# Description of PCP

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An individual  
domino

$$\begin{bmatrix} a \\ \hline ab \end{bmatrix}$$

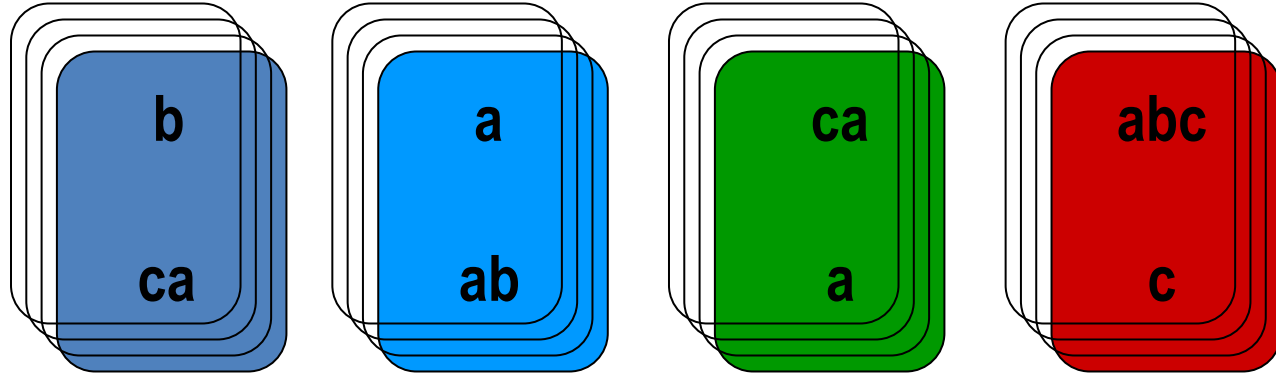


# Description of PCP

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A collection of  
dominos

$$\left\{ \begin{bmatrix} b \\ ca \end{bmatrix}, \begin{bmatrix} a \\ ab \end{bmatrix}, \begin{bmatrix} ca \\ a \end{bmatrix}, \begin{bmatrix} abc \\ c \end{bmatrix} \right\}$$

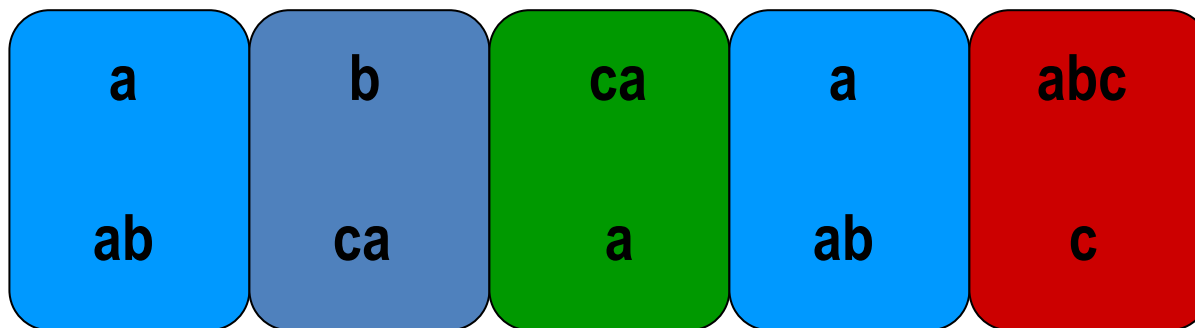
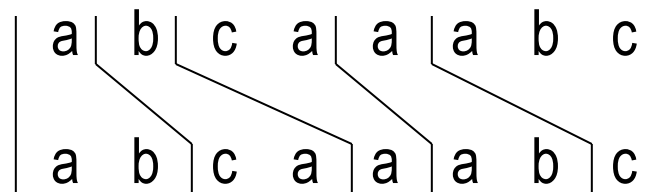


# Description of PCP

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A match

$$\begin{bmatrix} a \\ ab \end{bmatrix} \begin{bmatrix} b \\ ca \end{bmatrix} \begin{bmatrix} ca \\ a \end{bmatrix} \begin{bmatrix} a \\ ab \end{bmatrix} \begin{bmatrix} abc \\ c \end{bmatrix}$$





# A collection without a match

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- For a given collection

$$\left\{ \left[ \frac{abc}{ab} \right], \left[ \frac{ca}{a} \right], \left[ \frac{acc}{ba} \right] \right\}$$

- It cannot contain a match because every top string is longer than the corresponding bottom string



# Theorem 5.15

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- PCP is undecidable

- Proof:

To simplify the problem, we create *Modified Post Correspondence Problem (MPCP)* ,

$MPCP = \{ \langle P \rangle \mid P \text{ is an instance of the PCP with a match that starts with the first domino} \}$ .

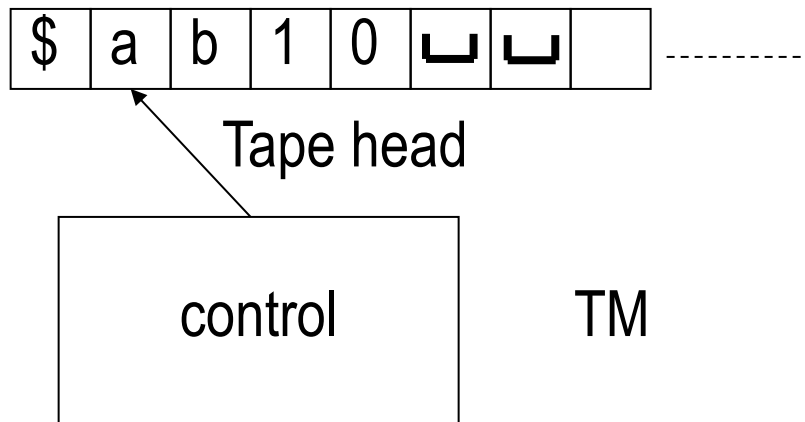


# Theorem 5.15

$$\left[ \frac{a}{ab} \right] \left[ \frac{b}{ca} \right] \left[ \frac{ca}{a} \right] \left[ \frac{a}{ab} \right] \left[ \frac{abc}{c} \right]$$

$$\begin{array}{|c|c|c|c|c|c|c|c|} \hline a & b & c & a & a & a & b & c \\ \hline a & b & c & a & a & a & b & c \\ \hline \end{array}$$

Input



- The PCP program is very similar with the  $A_{TM}$  program



# Theorem 5.15 proof

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- Proof:

Suppose PCP is decidable

We construct TM  $S$  to decide  $A_{TM}$  (Theorem 4.11:  $A_{TM}$  is undecidable)

$M = (Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$

$S$  constructs an instance of the PCP  $P$  that has a match iff

$M$  accepts  $w$ :  $\longleftarrow A_{TM}$

- ▶ (1)  $S$  first constructs an instance  $P'$  of the MPCP;
- ▶ (2) Transfer  $P'$  into  $P$ ;

PCP



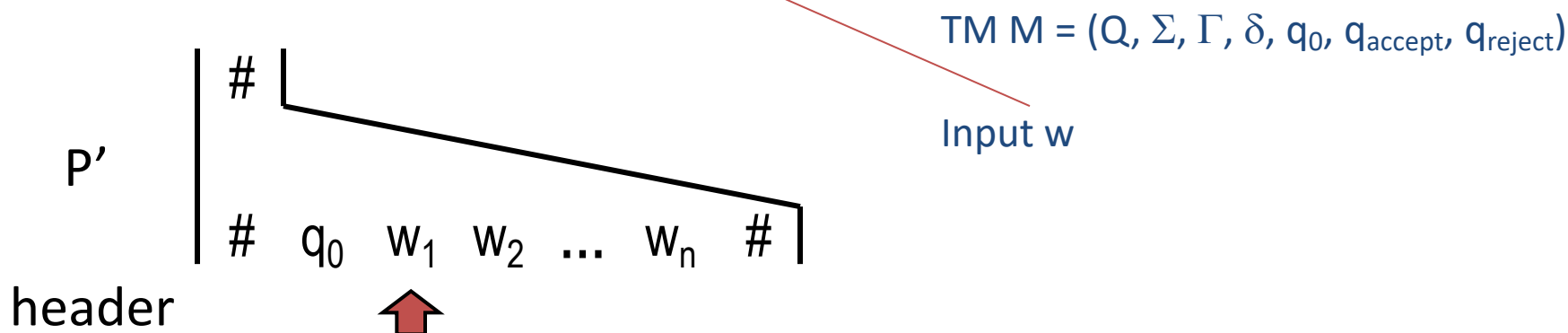
# Theorem 5.15 proof

- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.1) Generating beginning configuration**

Put  $\left[ \frac{\#}{\#q_0w_1w_2...w_n\#} \right]$  into  $P'$  as  $\left[ \frac{t_1}{b_1} \right]$



# Theorem 5.15 proof

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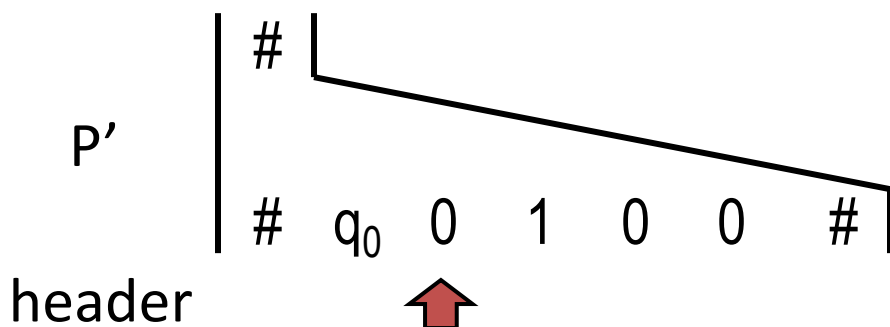
- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.1) Generating beginning configuration**

Suppose  $\Gamma = \{0, 1, 2, \_ \}$ ,  $w = 0100$ ,

Put the following into  $P'$  :  $\left[ \frac{\#}{\#q_0 0100\#} \right] = \left[ \frac{t_1}{b_1} \right]$



# Theorem 5.15 proof

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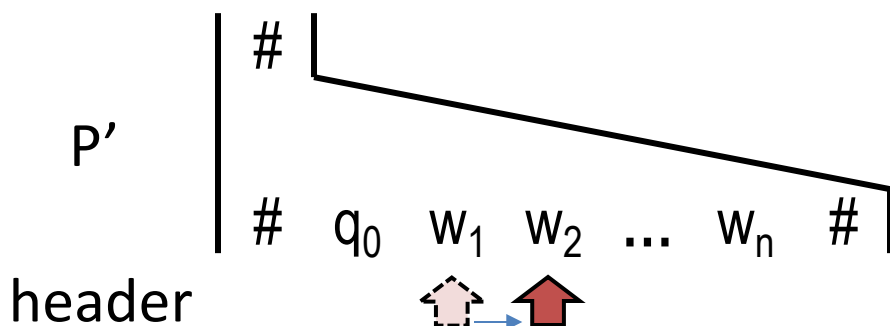
- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.2) The head move to the right**

For each  $a, b \in \Gamma$  and  $q, r \in Q$ , where  $q \neq q_{\text{reject}}$ ,

if  $\delta(q, a) = (r, b, R)$ , then put  $\left[ \frac{qa}{br} \right]$  into  $P'$



# Theorem 5.15 proof

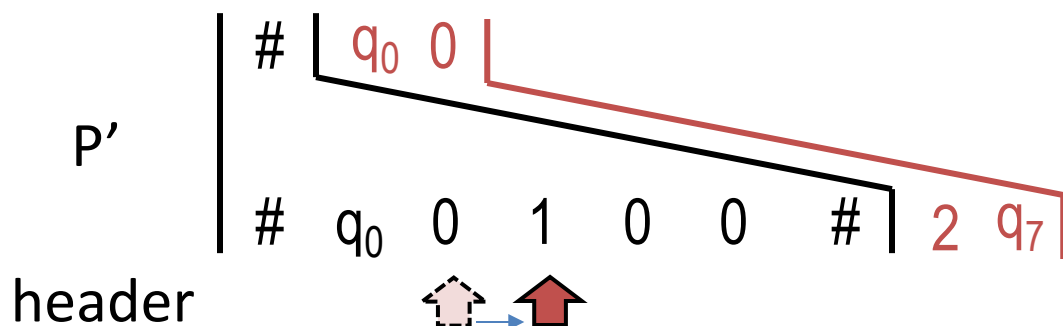
- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.2) The head moves to the right**

For each  $a, b \in \Gamma$  and  $q, r \in Q$ , where  $q \neq q_{\text{reject}}$ ,

if  $\delta(q_0, 0) = (q_7, 2, R)$ , then put  $\left[ \frac{q_0 0}{2 q_7} \right]$  into  $P'$





# Theorem 5.15 proof

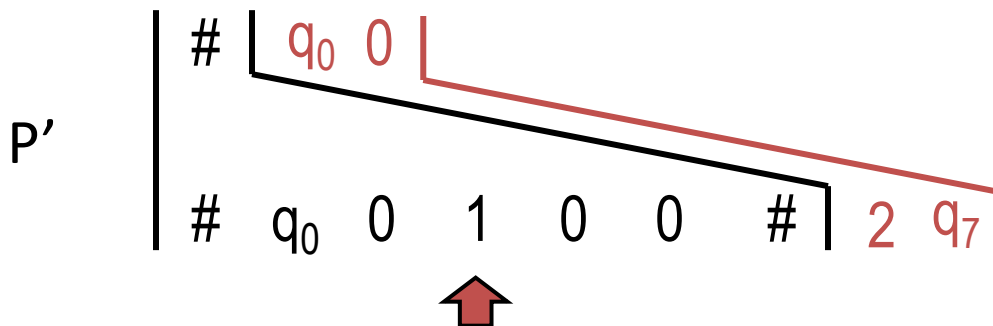
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.3) The head moves to the left**

Discuss it later on.



# Theorem 5.15 proof

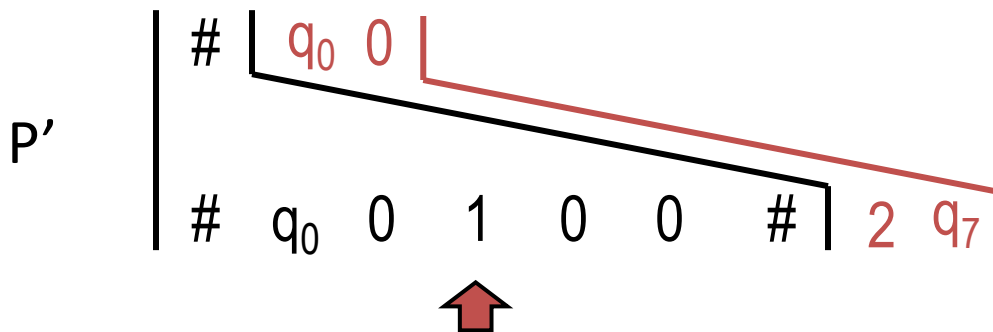
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.4) For each  $a \in \Gamma$ ,**

put  $\begin{bmatrix} a \\ a \end{bmatrix}$  into  $P'$



# Theorem 5.15 proof

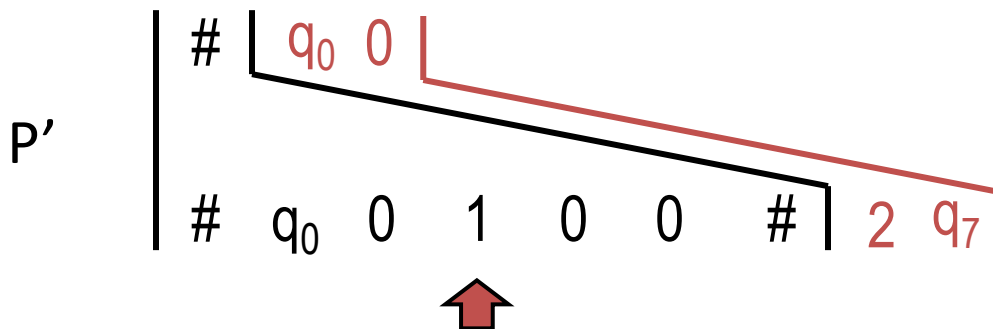
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

(1.4)  $\Gamma = \{0, 1, 2, \sqcup\}$ ,

put  $\begin{bmatrix} 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$  and  $\begin{bmatrix} \sqcup \\ \sqcup \end{bmatrix}$  into  $P'$



# Theorem 5.15 proof

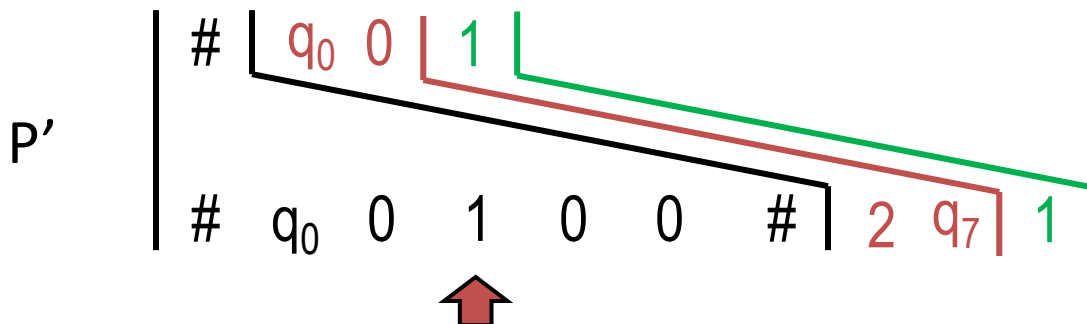
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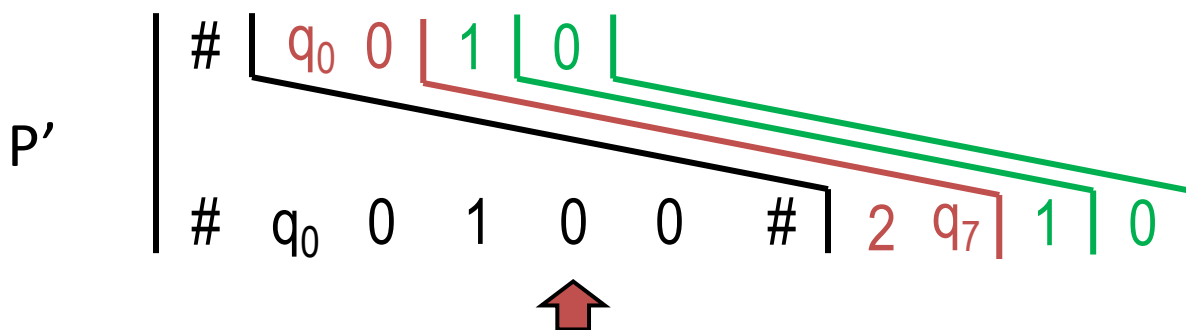
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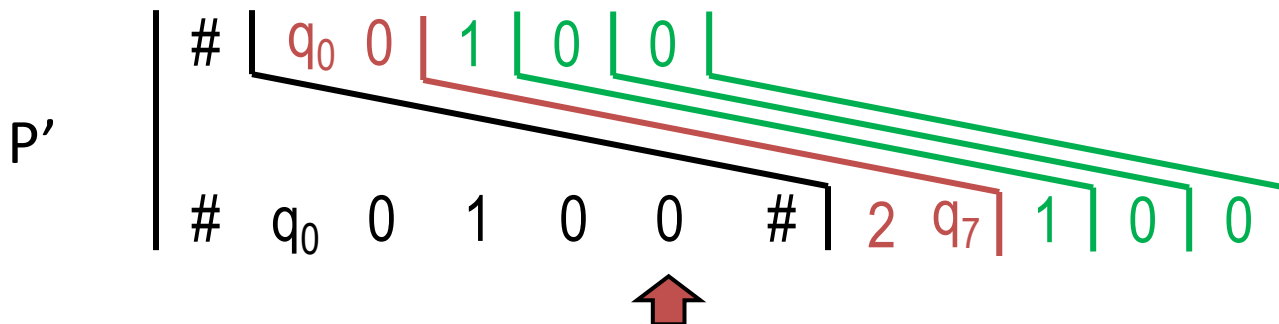
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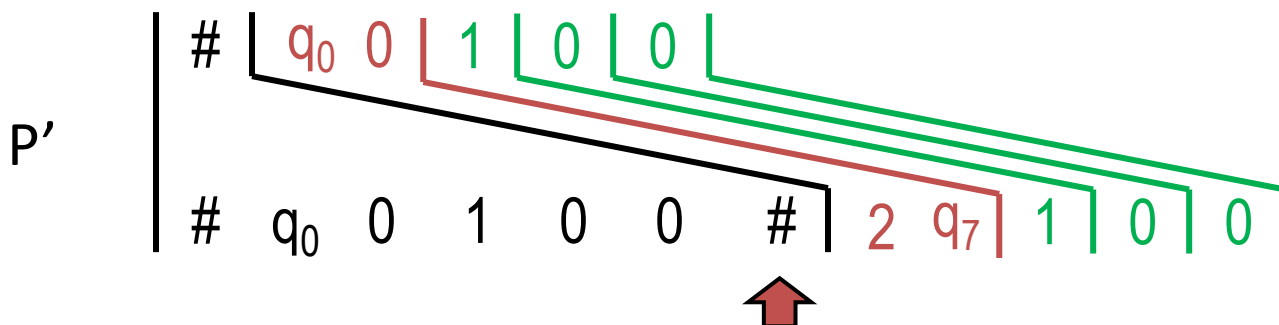
# Theorem 5.15 proof

- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.5) Copy # and Put  $\_$  at the end of configuration**

put  $\begin{bmatrix} \# \\ \_ \end{bmatrix}$  and  $\begin{bmatrix} \# \\ \_ \end{bmatrix}$  into  $P'$



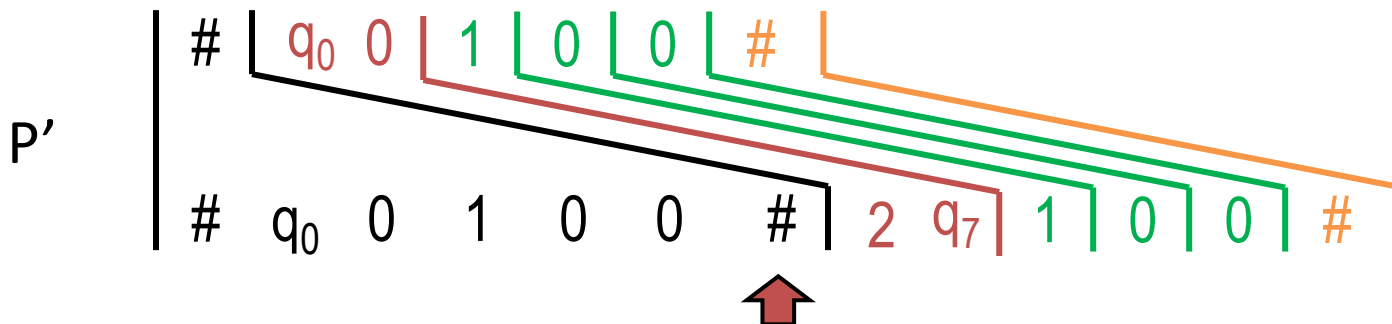
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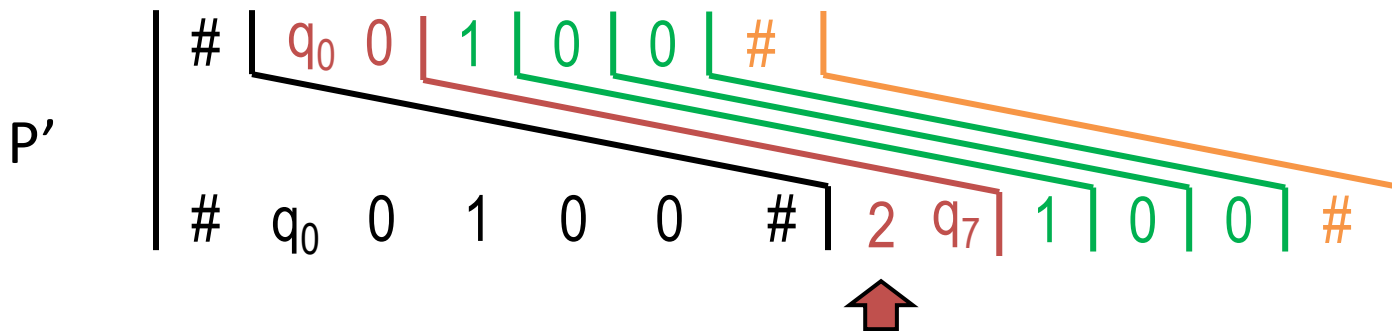
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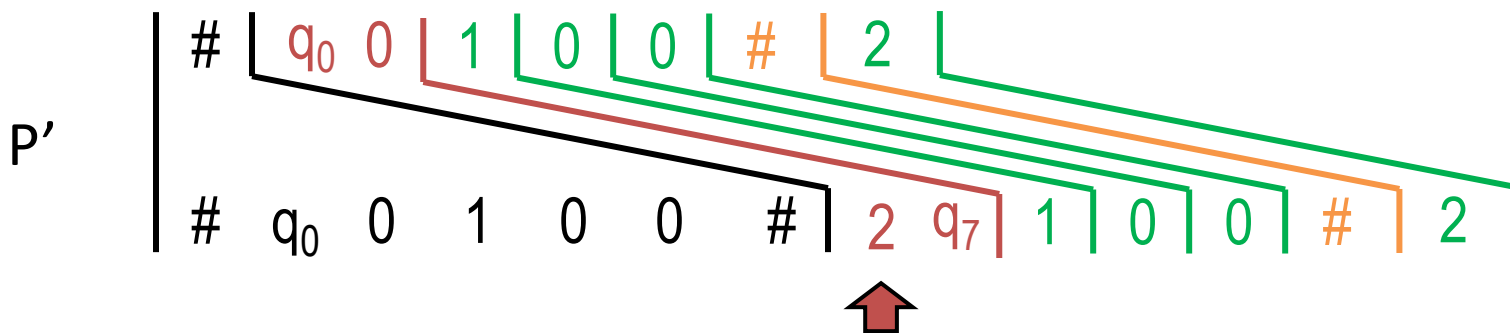
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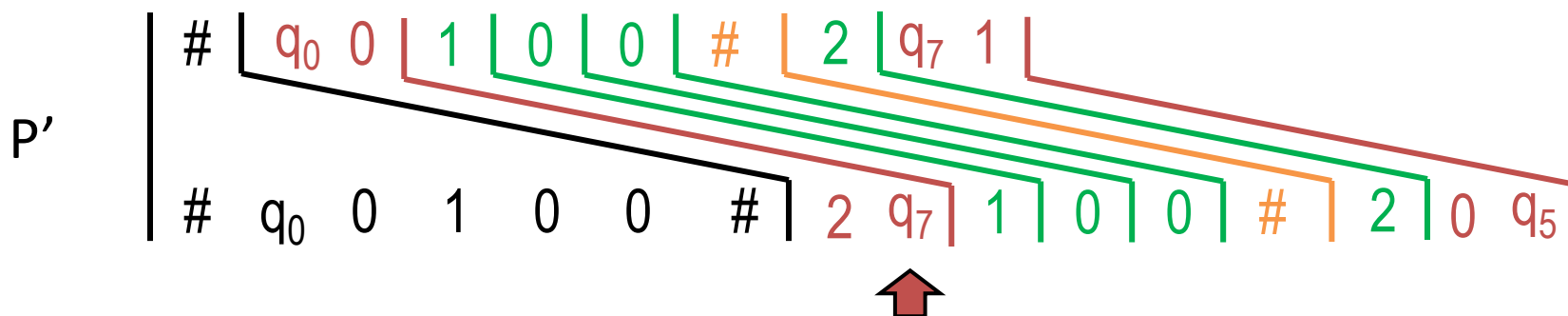
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.5) Copy # and Put  $\_$  at the end of configuration**

suppose  $\delta(q_7, 1) = (q_5, 0, R)$ , then put  $\left[ \begin{smallmatrix} q_7 1 \\ 0 q_5 \end{smallmatrix} \right]$  into  $P'$



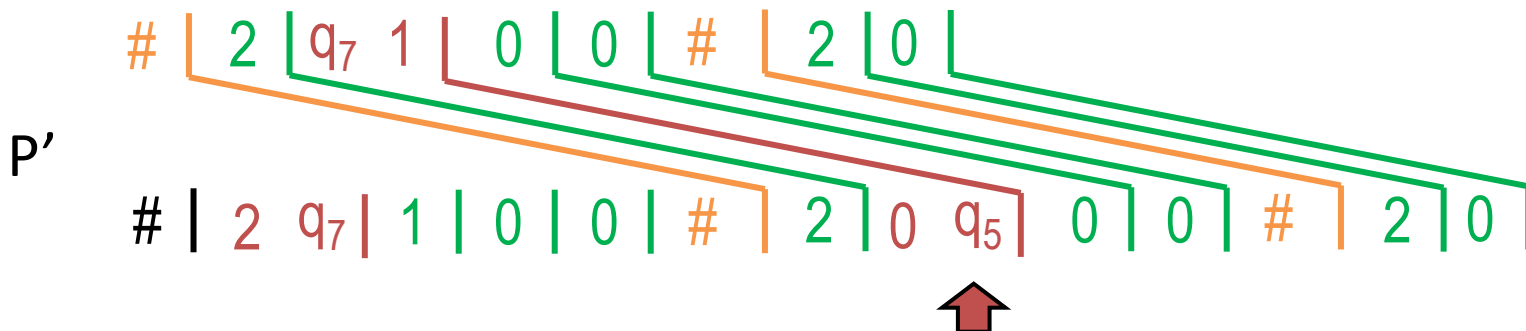
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.5) Copy # and Put  $\_$  at the end of configuration**

keep putting  $\begin{bmatrix} 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} \_ \\ \_ \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ \_ \end{bmatrix}$  into  $P'$



# Theorem 5.15 proof

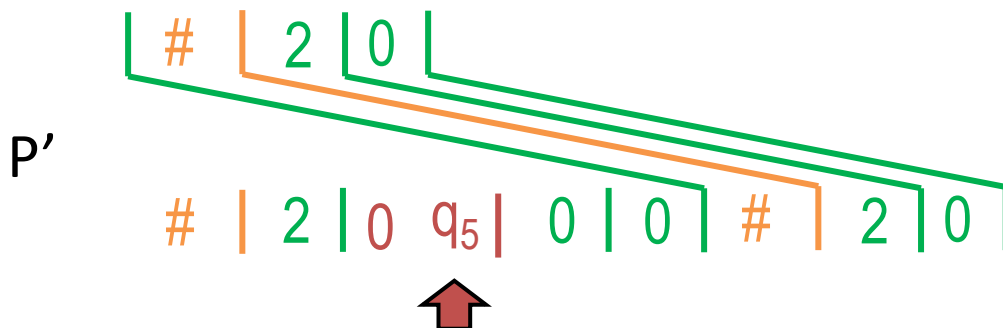
- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.3) The head moves to the left**

For each  $a, b, c \in \Gamma$  and  $q, r \in Q$ , where  $q \neq q_{\text{reject}}$ ,

if  $\delta(q, a) = (r, b, L)$ , then put  $\begin{bmatrix} cqa \\ rcb \end{bmatrix}$  into  $P'$ ,  $c$  is the element on the left of  $q$



# Theorem 5.15 proof

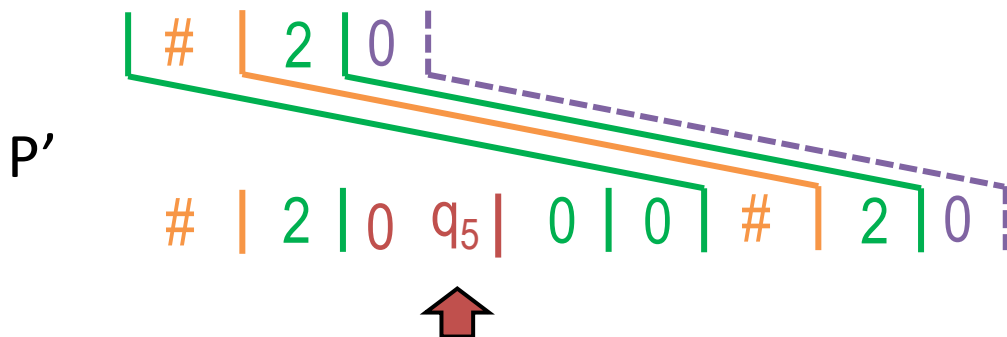
- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

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Suppose  $\delta(q_5, 0) = (q_9, 2, L)$ , remove the old 0

if  $\delta(q, a) = (r, b, L)$ , then put  $\left[ \frac{cqa}{rcb} \right]$  into  $P'$ ,  $c$  is the element on the left of  $q$



# Theorem 5.15 proof

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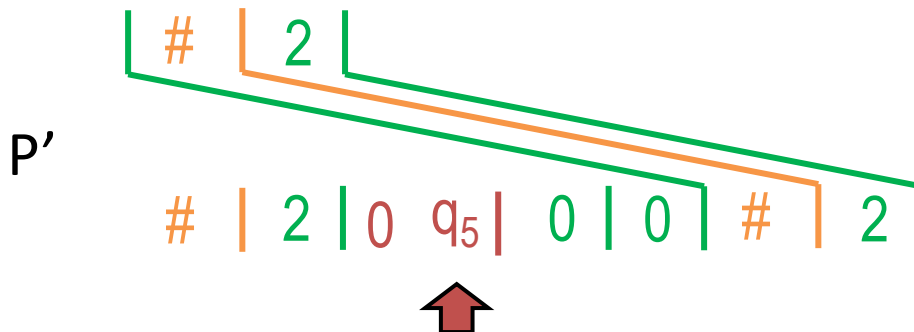
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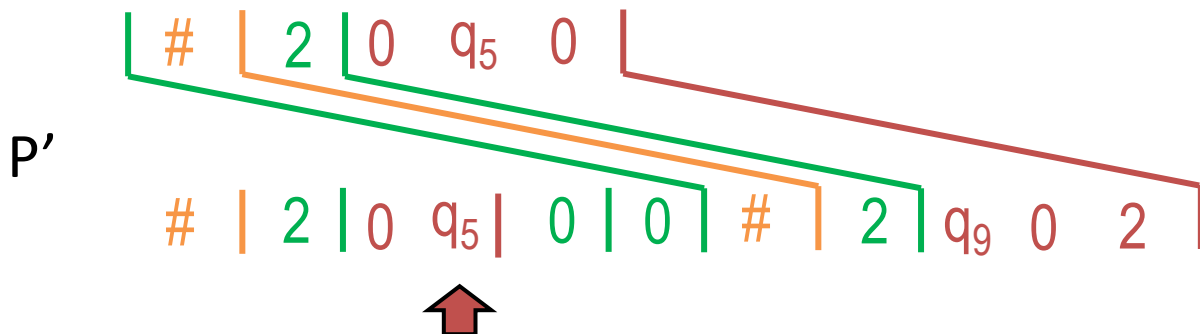
- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

**(1.3) The head moves to the left**

Suppose  $\delta(q_5, 0) = (q_9, 2, L)$ , put  $\begin{bmatrix} 0q_50 \\ q_902 \end{bmatrix}$  into  $P'$

if  $\delta(q, a) = (r, b, L)$ , then put  $\begin{bmatrix} cqa \\ rcb \end{bmatrix}$  into  $P'$ ,  $c$  is the element on the left of  $q$





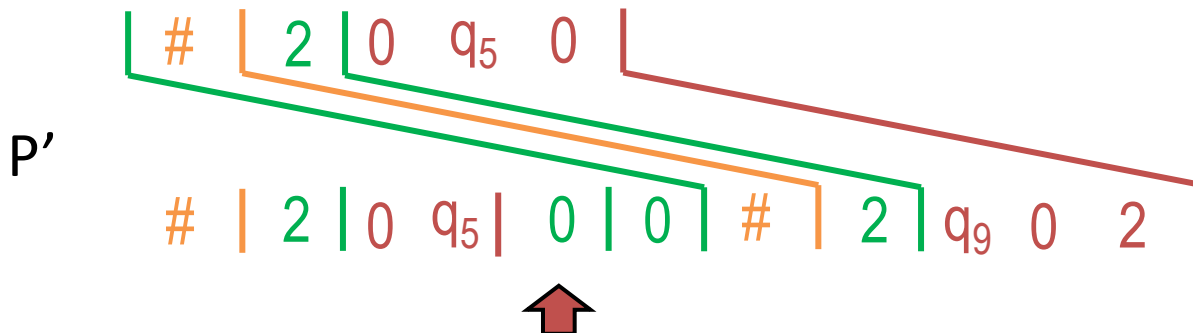
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keep putting  $\begin{bmatrix} 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} \_ \\ \_ \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ \_ \end{bmatrix}$  into  $P'$



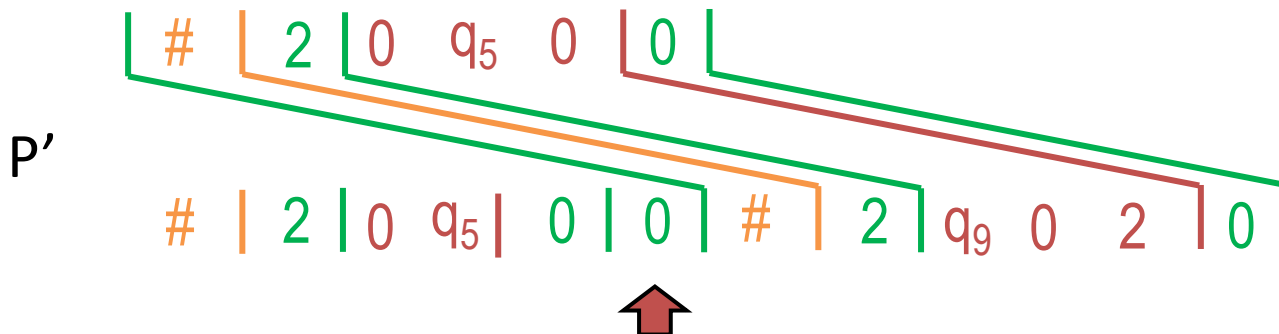
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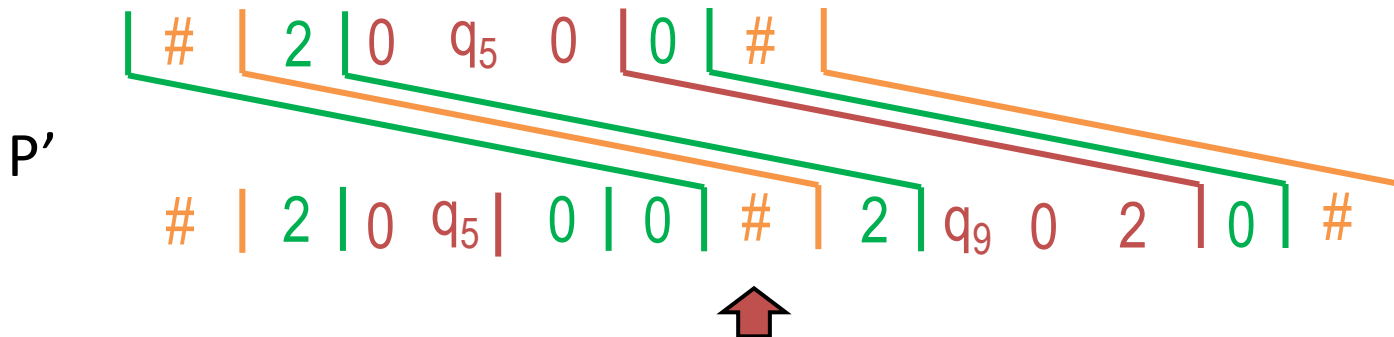
# Theorem 5.15 proof

- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

### (1.3) The head moves to the left

keep putting  $\begin{bmatrix} 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} \_ \\ \_ \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ \_ \end{bmatrix}$  into  $P'$



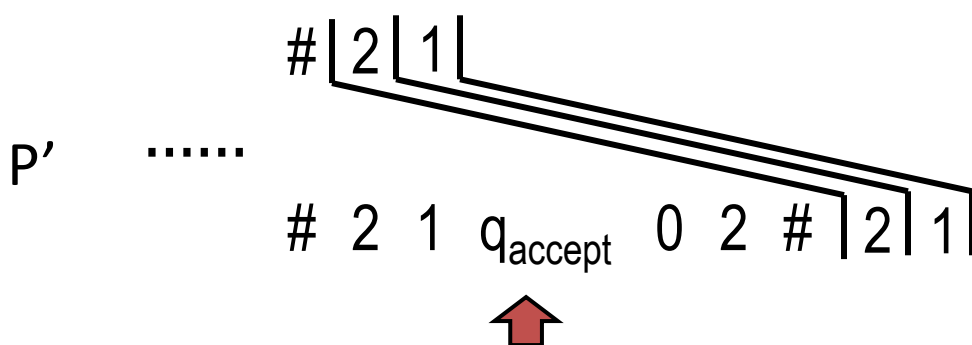
# Theorem 5.15 proof

- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

keep putting something into  $P'$  until M halts:

- ▶ if M rejects, S also rejects, means no match in PCP;
- ▶ if M accepts, add something to the top to match the bottom.



Connect PCP match  
program with whether  
M accepts  $w$



# Theorem 5.15 proof

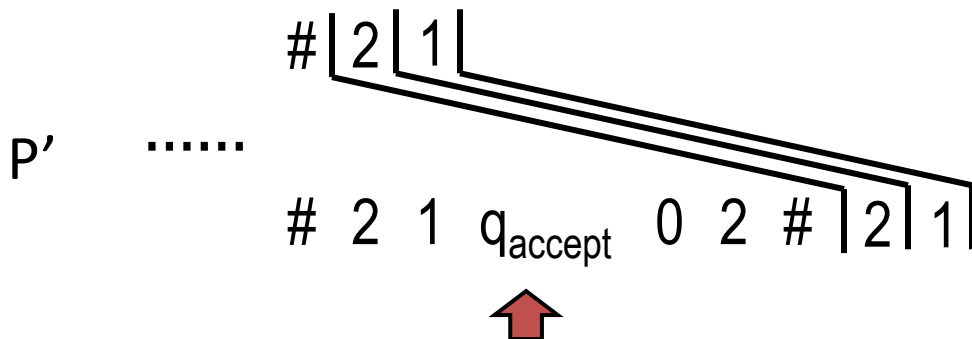
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

(1.6) For  $a \in \Gamma$ ,

put  $\left[ \frac{aq_{accept}}{q_{accept}} \right]$  and  $\left[ \frac{q_{accept}a}{q_{accept}} \right]$  into  $P'$



# Theorem 5.15 proof

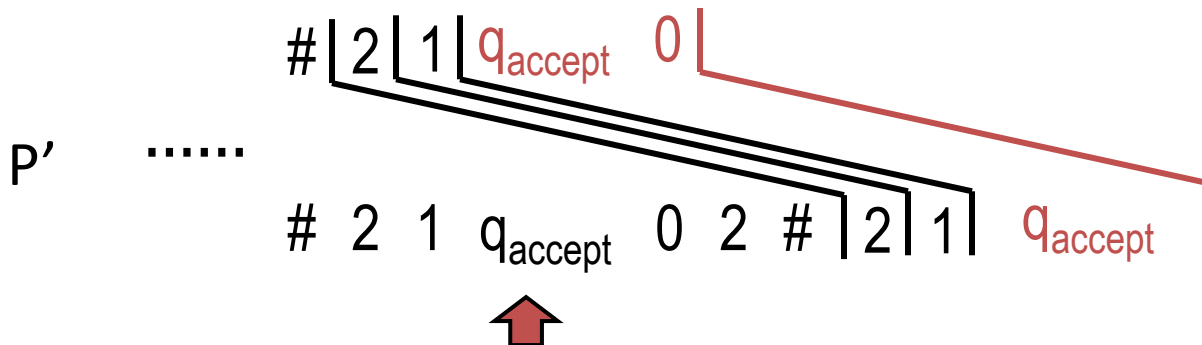
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(1) S first constructs an instance  $P'$  of the MPCP;

(1.6) For  $a \in \Gamma$ ,

put  $\left[ \frac{aq_{accept}}{q_{accept}} \right]$  and  $\left[ \frac{q_{accept}a}{q_{accept}} \right]$  into  $P'$



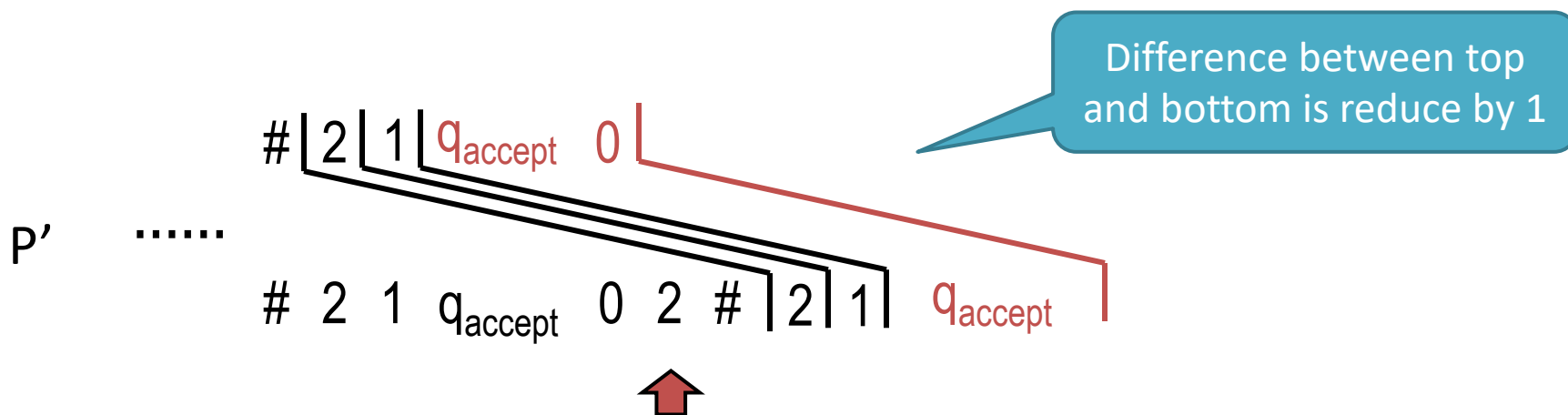
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# Theorem 5.15 proof

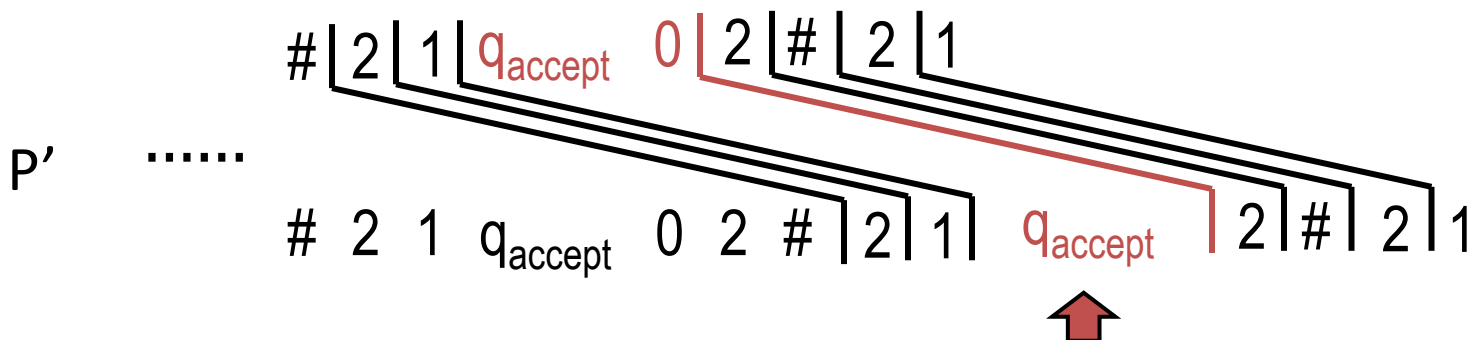
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# Theorem 5.15 proof

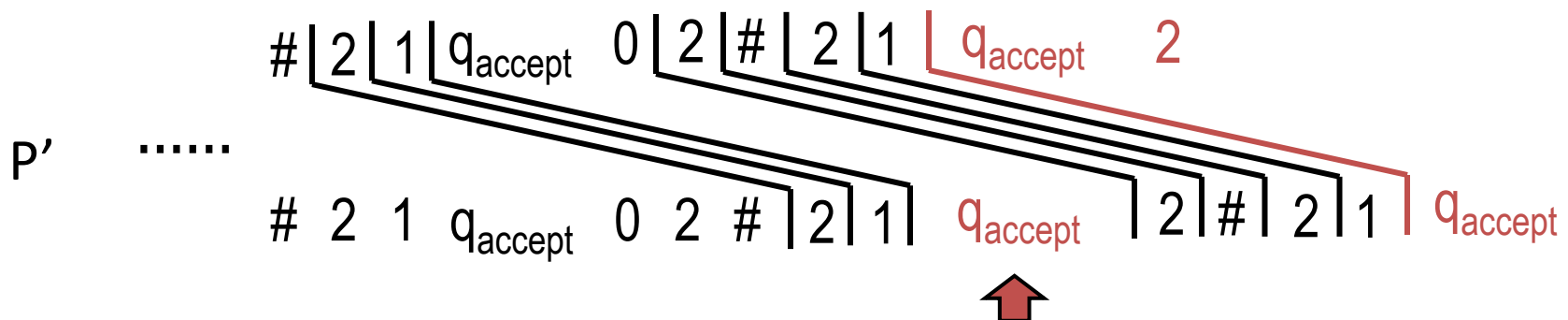
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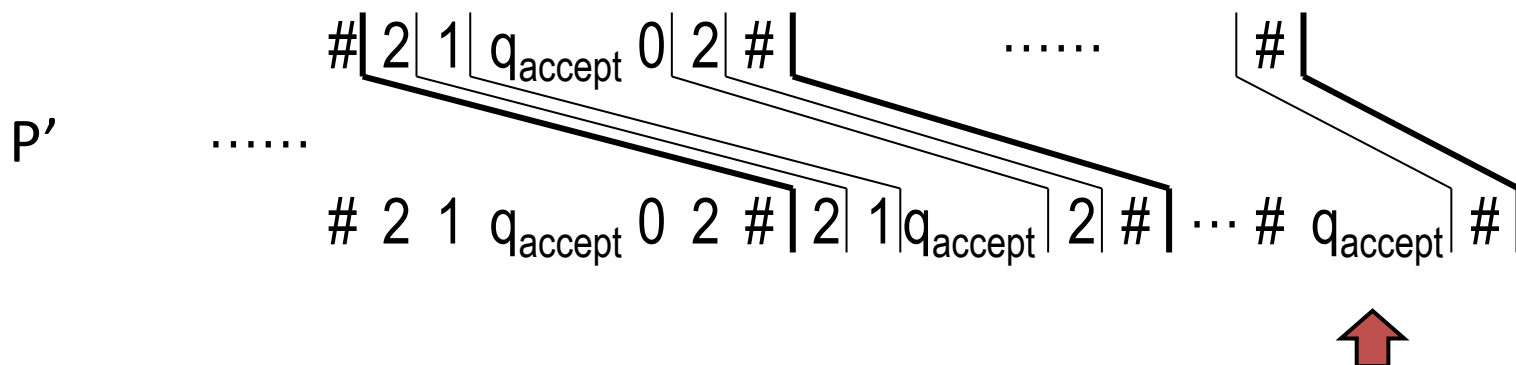
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- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

(1.6) For  $a \in \Gamma$ ,

Keep doing until the top and bottom difference is 1.



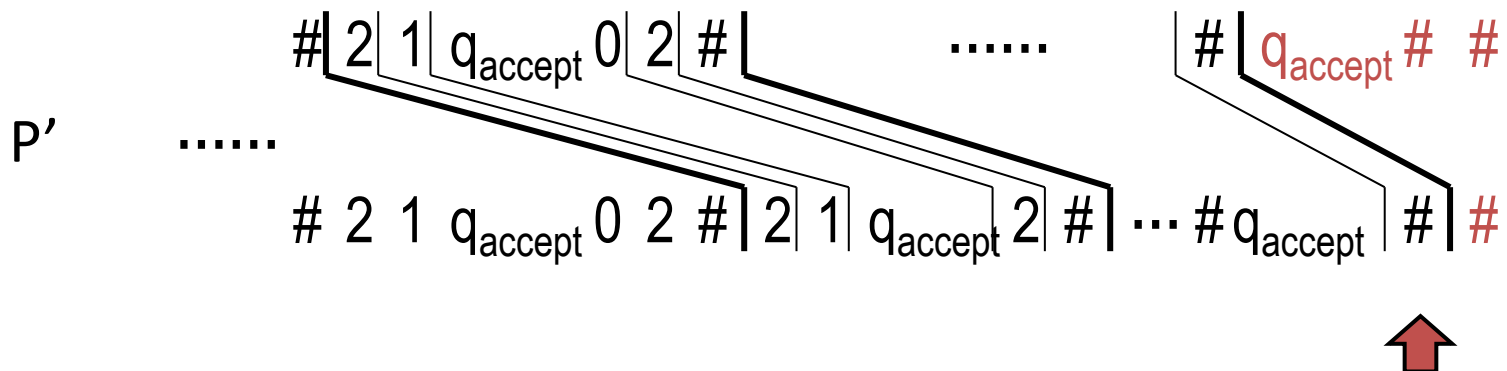
# Theorem 5.15 proof

- Proof:

(1) S first constructs an instance  $P'$  of the MPCP;

## (1.7) Finish the match

put  $\left[ \frac{q_{accept}^{##}}{\#} \right]$  into  $P'$



# Theorem 5.15 proof

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- Proof:

(2) Transfer  $P'$  into  $P$

$P'$  is MPCP =  $\{\langle P \rangle \mid P \text{ is an instance of the PCP with a match that starts with the first domino}\}$ .

Suppose  $u = u_1 u_2 \dots u_n$  to be any string of length  $n$ , define

- ▶  $\star u = \star u_1 \star u_2 \star \dots \star u_n$
- ▶  $u \star = u_1 \star u_2 \star \dots \star u_n \star$
- ▶  $\star u \star = \star u_1 \star u_2 \star \dots \star u_n \star$



# Theorem 5.15 proof

$$\star u = \star u_1 \star u_2 \star \dots \star u_3$$

$$u \star = u_1 \star u_2 \star \dots \star u_3 \star$$

$$\star u \star = \star u_1 \star u_2 \star \dots \star u_3 \star$$

- Proof:

(2) Transfer  $P'$  into  $P$

If  $P' = \left\{ \left[ \frac{t_1}{b_1} \right], \left[ \frac{t_2}{b_2} \right], \dots, \left[ \frac{t_k}{b_k} \right] \right\}$ , then let

$$P = \left\{ \left[ \frac{\star t_1}{\star b_1 \star} \right], \left[ \frac{\star t_2}{b_2 \star} \right], \left[ \frac{\star t_3}{b_3 \star} \right], \dots, \left[ \frac{\star t_k}{b_k \star} \right], \left[ \frac{\star \Delta}{\Delta} \right] \right\}$$

First domino: The only element as beginning

The match in  $P$  must be in shape as  $\left[ \frac{\star t_1}{\star b_1 \star} \right] \dots \left[ \frac{\star \Delta}{\Delta} \right]$



# Theorem 5.15 proof

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- Proof:

Suppose PCP is decidable

We construct TM  $S$  using PCP match or not to decide  $A_{TM}$

- ▶ (1)  $S$  first constructs an instance  $P'$  of the MPCP;
- ▶ (2) Transfer  $P'$  into  $P$ ;

Theorem 4.11:  $A_{TM}$  is undecidable. Contradiction!

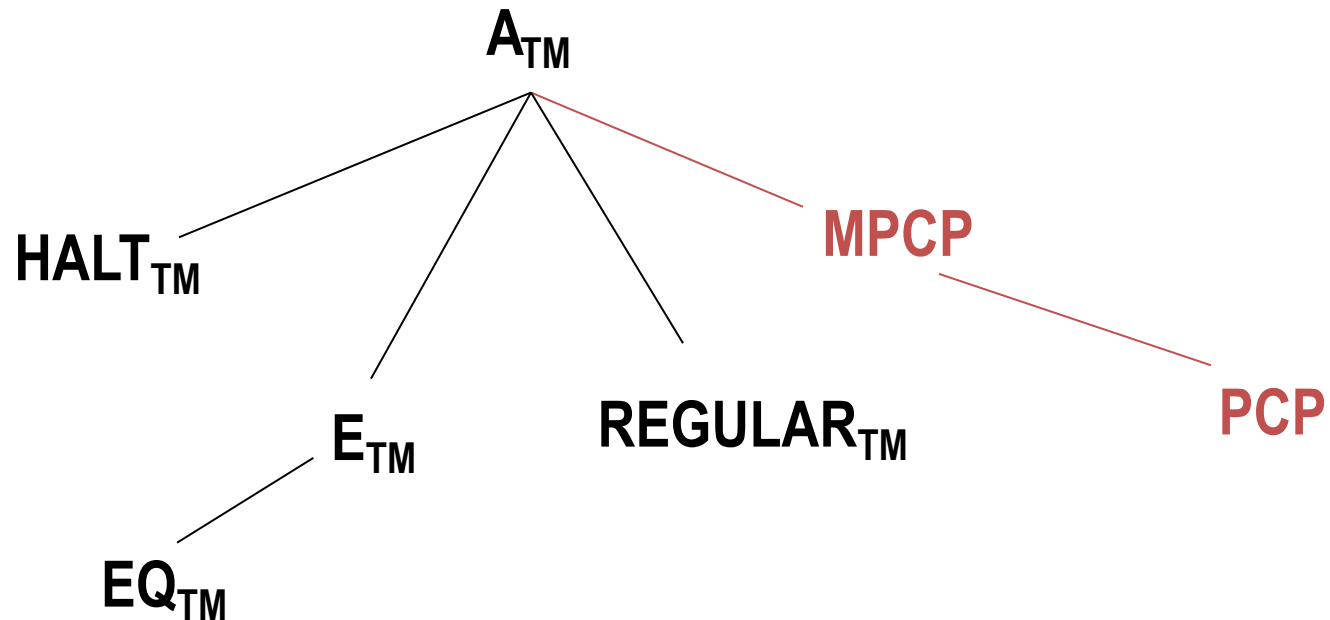
The suppose is wrong. Thus PCP is undecidable.



# Conclusion

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- Relationship of languages on reducibility



# Conclusion

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- Closure on operations

	<b>Complement</b> $\bar{A}$	<b>Intersection</b> $\cap$	<b>Union</b> $\cup$	<b>Star</b> $A^*$
<b>Regular/DFA/ NFA</b>	✓	✓	✓	✓
<b>CFL/ PDA</b>	×	×	✓	✓
<b>Turing- decidable TM</b>	✓	✓	✓	✓
<b>Turing- recognizable TM</b>	×	✓	✓	✓

