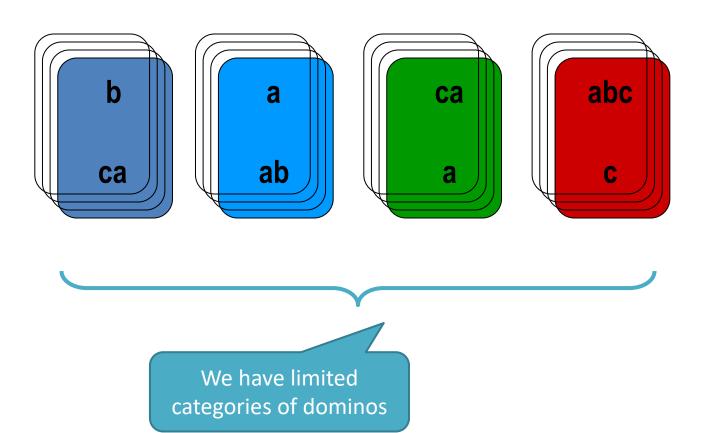
# CS 6041 Theory of Computation

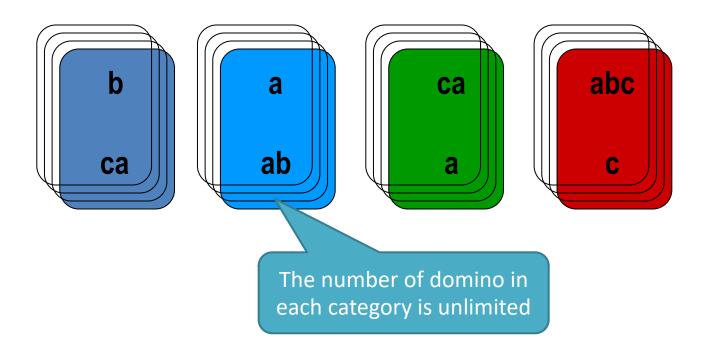
### Reducibility

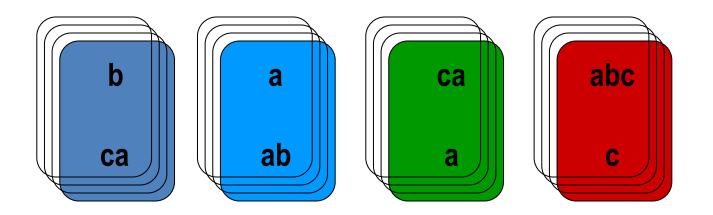
#### **Kun Suo**

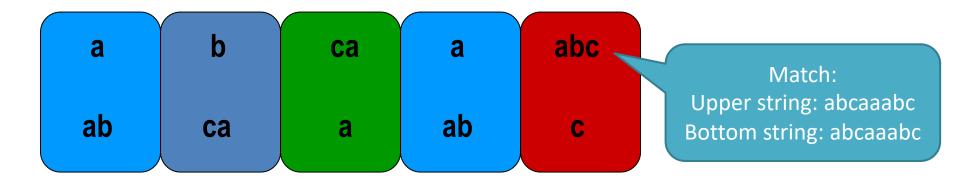
Computer Science, Kennesaw State University

https://kevinsuo.github.io/

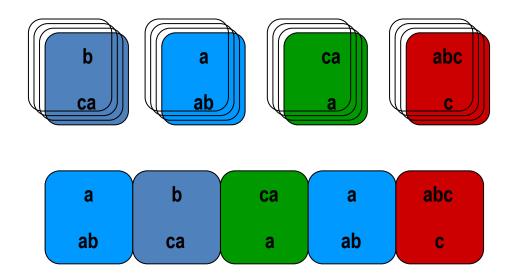








Whether a collection of dominos has a match



 PCP = {\langle P \rangle | P is an instance of the Post Correspondence Problem with a match}.

# **Description of PCP**

An individual domino

$$\left[\frac{a}{ah}\right]$$

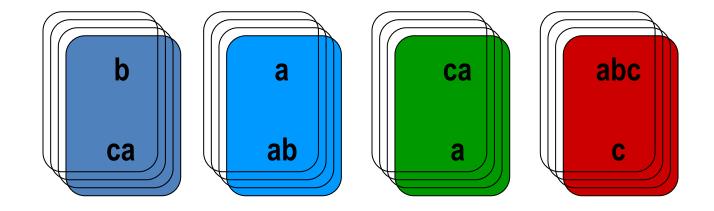
a

ab

### **Description of PCP**

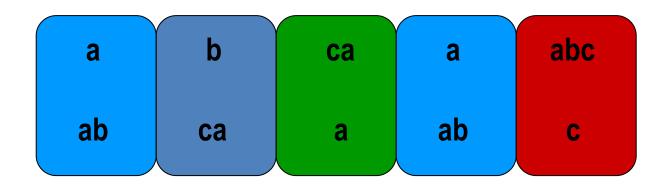
A collection of dominos

$$\left\{ \left[ \frac{b}{ca} \right], \left[ \frac{a}{ab} \right], \left[ \frac{ca}{a} \right], \left[ \frac{abc}{c} \right] \right\}$$



### **Description of PCP**

A match



### A collection without a match

For a given collection

$$\left\{ \left[ \frac{abc}{ab} \right], \left[ \frac{ca}{a} \right], \left[ \frac{acc}{ba} \right] \right\}$$

 It cannot contain a match because every top string is longer than the corresponding bottom string

### Theorem 5.15

PCP is undecidable

#### • Proof:

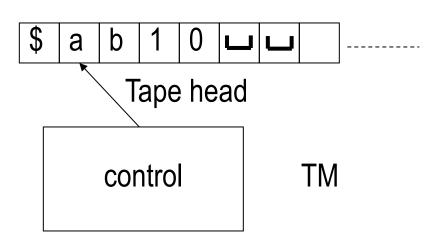
To simplify the problem, we create *Modified Post*Correspondence Problem (MPCP),

MPCP =  $\{\langle P \rangle \mid P \text{ is an instance of the PCP with a match that starts with the first domino}\}$ .

### Theorem 5.15

$$\begin{bmatrix} a \\ ab \end{bmatrix} \begin{bmatrix} b \\ ca \end{bmatrix} \begin{bmatrix} ca \\ a \end{bmatrix} \begin{bmatrix} a \\ ab \end{bmatrix} \begin{bmatrix} abc \\ c \end{bmatrix}$$

Input



• The PCP program is very similar with the  $A_{TM}$  program

#### • Proof:

Suppose PCP is decidable

We construct TM S to decide  $A_{TM}$  (Theorem 4.11:  $A_{TM}$  is undecidable)

$$M = (Q, \Sigma, \Gamma, \delta, q_0, q_{accept}, q_{reject})$$

S constructs an instance of the PCP P that has a match iff

M accepts w:



- (1) S first constructs an instance P' of the MPCP;
- (2) Transfer P' into P;

- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.1) Generating beginning configuration

Put 
$$\left[\frac{\#}{\#q_0w_1w_2...w_n\#}\right]$$
 into P' as  $\left[\frac{t_1}{b_1}\right]$ 

TM M = (Q,  $\Sigma$ ,  $\Gamma$ ,  $\delta$ ,  $q_0$ ,  $q_{accept}$ ,  $q_{reject}$ )

P'

#  $q_0$   $w_1$   $w_2$  ...  $w_n$  #

header

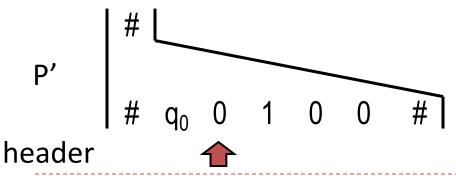
#### • Proof:

(1) S first constructs an instance P' of the MPCP;

#### (1.1) Generating beginning configuration

Suppose  $\Gamma = \{0,1,2, \_\}, w = 0100,$ 

Put the following into P': 
$$\left[\frac{\#}{\#q_00100\#}\right] = \left[\frac{t_1}{b_1}\right]$$



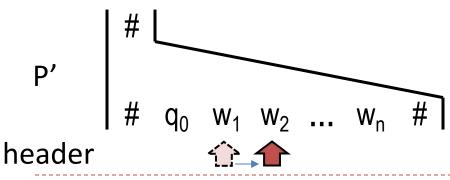
#### • Proof:

S first constructs an instance P' of the MPCP;

#### (1.2) The head move to the right

For each a,  $b \in \Gamma$  and q,  $r \in Q$ , where  $q \neq q_{reject}$ ,

if 
$$\delta(q, a) = (r, b, R)$$
, then put  $\left[\frac{qa}{br}\right]$  into P'



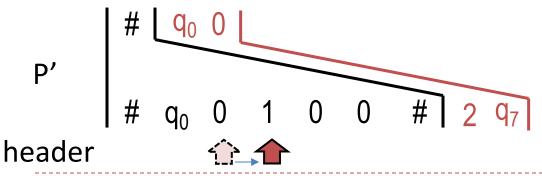
#### • Proof:

(1) S first constructs an instance P' of the MPCP;

#### (1.2) The head moves to the right

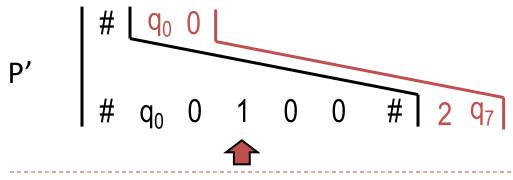
For each a,  $b \in \Gamma$  and q,  $r \in Q$ , where  $q \neq q_{reject}$ ,

if 
$$\delta(q_0, 0) = (q_7, 2, R)$$
, then put  $\left[\frac{q_0 0}{2q_7}\right]$  into P'



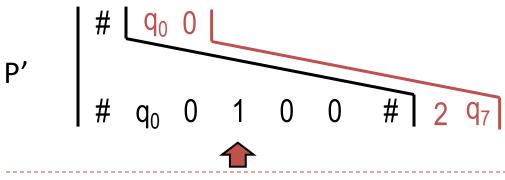
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

Discuss it later on.



- Proof:
  - S first constructs an instance P' of the MPCP;
  - (1.4) For each  $a \in \Gamma$ ,

put 
$$\left[\frac{a}{a}\right]$$
 into P'

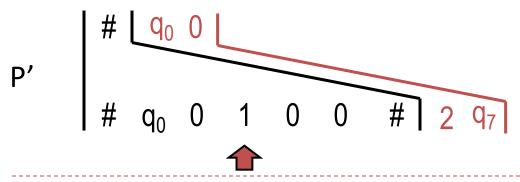


#### • Proof:

S first constructs an instance P' of the MPCP;

$$(1.4) \Gamma = \{0, 1, 2, \bot\},\$$

put 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} \frac{1}{1} \end{bmatrix}$ ,  $\begin{bmatrix} \frac{2}{2} \end{bmatrix}$  and  $\begin{bmatrix} \frac{1}{2} \end{bmatrix}$  into P'

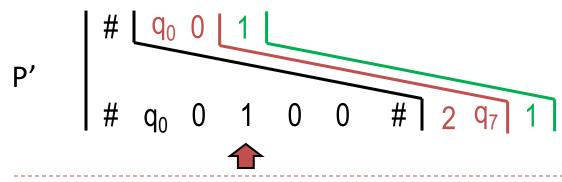


#### • Proof:

(1) S first constructs an instance P' of the MPCP;

$$(1.4) \Gamma = \{0, 1, 2, \_\},\$$

put 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} \frac{1}{1} \end{bmatrix}$ ,  $\begin{bmatrix} \frac{2}{2} \end{bmatrix}$  and  $\begin{bmatrix} \frac{1}{2} \end{bmatrix}$  into P'

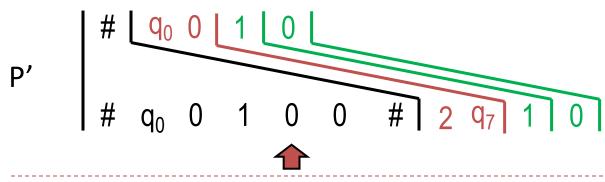


#### • Proof:

(1) S first constructs an instance P' of the MPCP;

$$(1.4) \Gamma = \{0, 1, 2, \bot\},\$$

put 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} \frac{1}{1} \end{bmatrix}$ ,  $\begin{bmatrix} \frac{2}{2} \end{bmatrix}$  and  $\begin{bmatrix} \frac{1}{2} \end{bmatrix}$  into P'

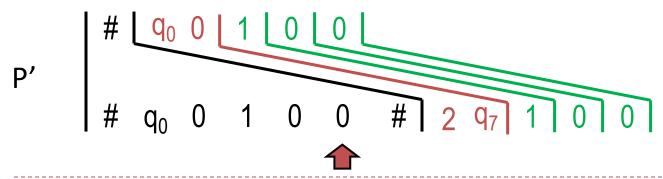


#### • Proof:

(1) S first constructs an instance P' of the MPCP;

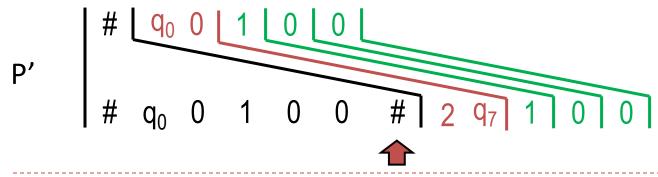
$$(1.4) \Gamma = \{0, 1, 2, \_\},\$$

put 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} \frac{1}{1} \end{bmatrix}$ ,  $\begin{bmatrix} \frac{2}{2} \end{bmatrix}$  and  $\begin{bmatrix} \frac{1}{2} \end{bmatrix}$  into P'



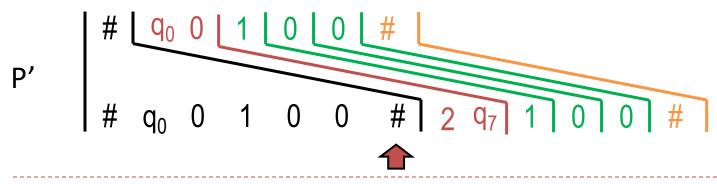
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.5) Copy # and Put \_ at the end of configuration

put 
$$\left[\frac{\#}{\#}\right]$$
 and  $\left[\frac{\#}{\#}\right]$  into P'



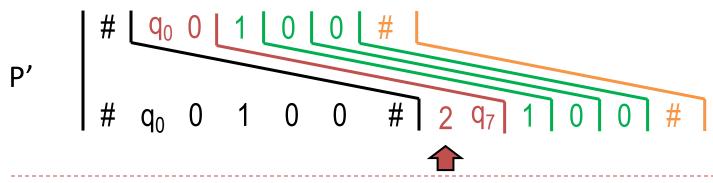
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.5) Copy # and Put \_ at the end of configuration

put 
$$\left[\frac{\#}{\#}\right]$$
 and  $\left[\frac{\#}{\#}\right]$  into P'



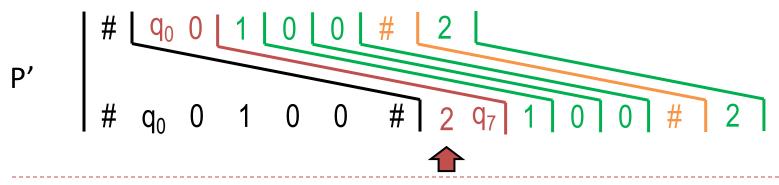
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.5) Copy # and Put \_ at the end of configuration

put 
$$\left[\frac{\#}{\#}\right]$$
 and  $\left[\frac{\#}{\#}\right]$  into P'



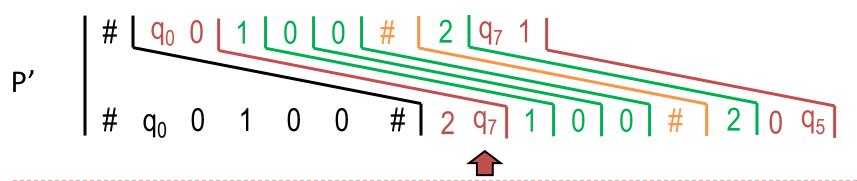
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.5) Copy # and Put \_ at the end of configuration

put 
$$\left[\frac{\#}{\#}\right]$$
 and  $\left[\frac{\#}{\#}\right]$  into P'



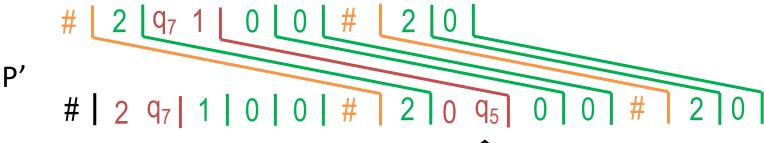
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.5) Copy # and Put \_ at the end of configuration

suppose 
$$\delta(q_7, 1) = (q_5, 0, R)$$
, then put  $\left[\frac{q_7 1}{0q_5}\right]$  into P'



- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.5) Copy # and Put \_ at the end of configuration

keep putting 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} -1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ -\# \end{bmatrix}$  into P'



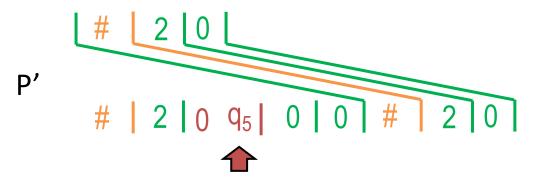
#### • Proof:

S first constructs an instance P' of the MPCP;

#### (1.3) The head moves to the left

For each a, b,  $c \in \Gamma$  and q,  $r \in Q$ , where  $q \neq q_{reject}$ ,

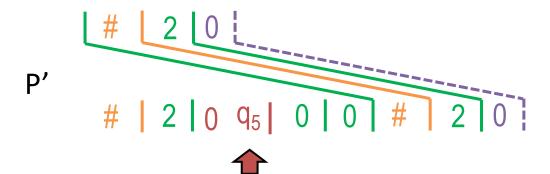
if  $\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{L})$ , then put  $\left[\frac{cqa}{rcb}\right]$  into P', c is the element on the left of q



- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

Suppose 
$$\delta(q_5, 0) = (q_9, 2, L)$$
, remove the old 0

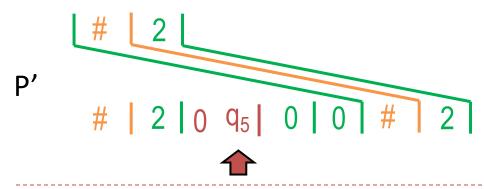
if 
$$\delta(q, a) = (r, b, L)$$
, then put  $\left[\frac{cqa}{rcb}\right]$  into P', c is the element on the left of q



- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

Suppose 
$$\delta(q_5, 0) = (q_9, 2, L)$$
, remove the old 0

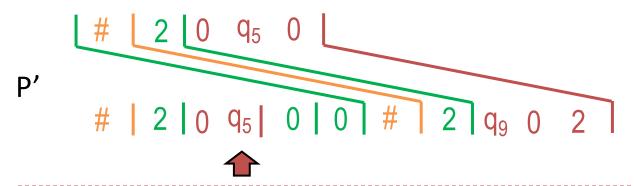
if 
$$\delta(q, a) = (r, b, L)$$
, then put  $\left[\frac{cqa}{rcb}\right]$  into P', c is the element on the left of q



- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

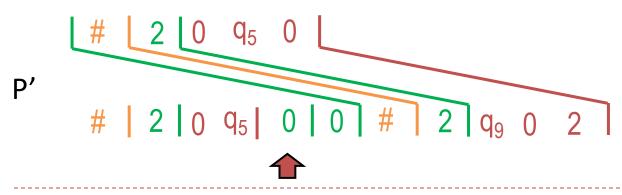
Suppose 
$$\delta(q_5, 0) = (q_9, 2, L)$$
, put  $\left[\frac{0q_50}{q_902}\right]$  into P'

if  $\delta(q, a) = (r, b, L)$ , then put  $\left[\frac{cqa}{rcb}\right]$  into P', c is the element on the left of q



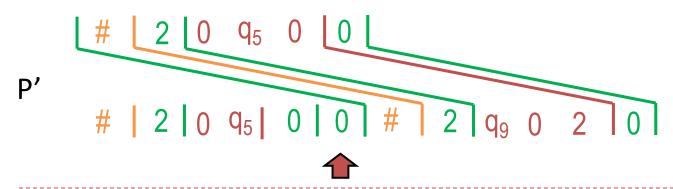
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

keep putting 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} -1 \\ -1 \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ -\# \end{bmatrix}$  into P'



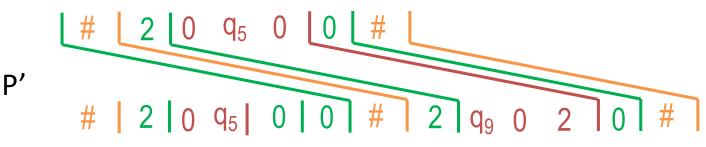
- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

keep putting 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} -1 \\ -1 \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ -\# \end{bmatrix}$  into P'



- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.3) The head moves to the left

keep putting 
$$\begin{bmatrix} 0 \\ 0 \end{bmatrix}$$
,  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 2 \end{bmatrix}$ ,  $\begin{bmatrix} -1 \\ 1 \end{bmatrix}$ ,  $\begin{bmatrix} \# \\ \# \end{bmatrix}$  and  $\begin{bmatrix} \# \\ -\# \end{bmatrix}$  into P'



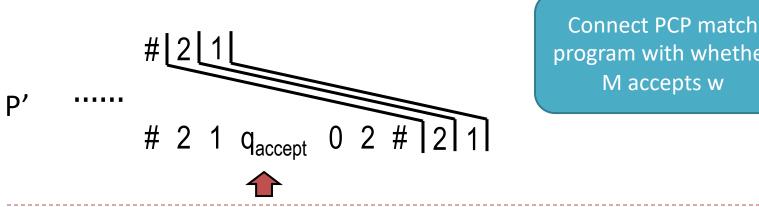


#### • Proof:

(1)S first constructs an instance P' of the MPCP;

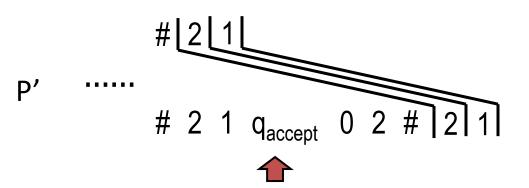
#### keep putting something into P' until M halts:

- if M rejects, S also rejects, means no match in PCP;
- if M accepts, add something to the top to match the bottom.



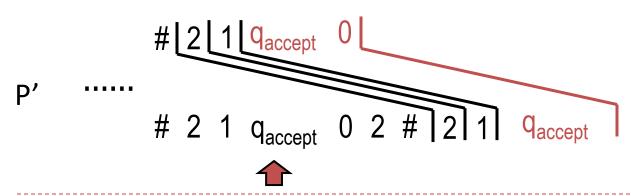
- (1) S first constructs an instance P' of the MPCP;
- (1.6) For  $a \in \Gamma$ ,

put 
$$\left[\frac{aq_{accept}}{q_{accept}}\right]$$
 and  $\left[\frac{q_{accept}a}{q_{accept}}\right]$  into P'



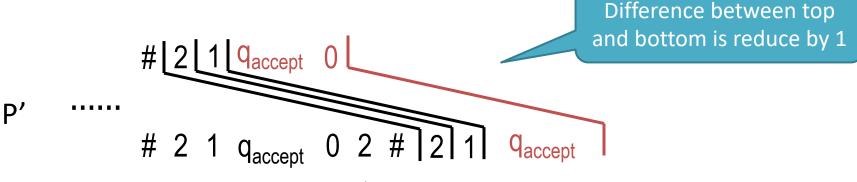
- S first constructs an instance P' of the MPCP;
- (1.6) For  $a \in \Gamma$ ,

put 
$$\left[\frac{aq_{accept}}{q_{accept}}\right]$$
 and  $\left[\frac{q_{accept}a}{q_{accept}}\right]$  into P'



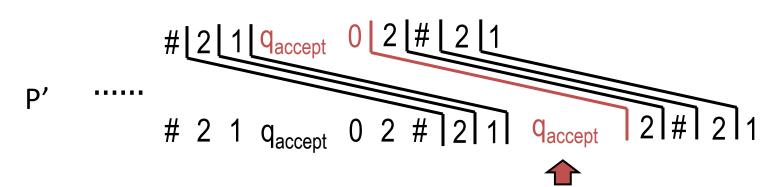
- (1) S first constructs an instance P' of the MPCP;
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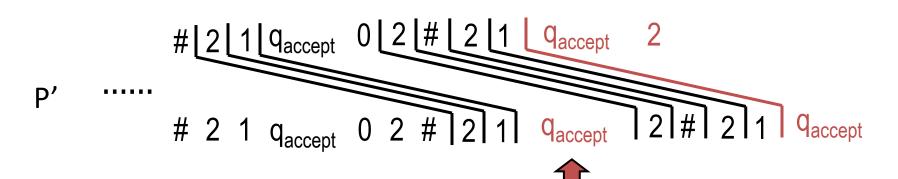
- S first constructs an instance P' of the MPCP;
- (1.6) For  $a \in \Gamma$ ,

put 
$$\left[\frac{aq_{accept}}{q_{accept}}\right]$$
 and  $\left[\frac{q_{accept}a}{q_{accept}}\right]$  into P'



- (1) S first constructs an instance P' of the MPCP;
- (1.6) For  $a \in \Gamma$ ,

put 
$$\left[\frac{aq_{accept}}{q_{accept}}\right]$$
 and  $\left[\frac{q_{accept}a}{q_{accept}}\right]$  into P'



- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.6) For  $a \in \Gamma$ ,

Keep doing until the top and bottom difference is 1.

- Proof:
  - (1) S first constructs an instance P' of the MPCP;
  - (1.7) Finish the match

put 
$$\left[\frac{q_{accept}##}{#}\right]$$
 into P'



#### • Proof:

(2) Transfer P' into P

P' is MPCP =  $\{\langle P \rangle | P \text{ is an instance of the PCP with a match that starts with the first domino}\}.$ 

Suppose  $u=u_1u_2...u_n$  to be any string of length n, define

- \*u = \*u1\*u2 \*...\*u3
- $u \star = u_1 * u_2 * ... * u_3 *$
- $\star u \star = *u_1 * u_2 * ... * u_3 *$

$$*u = *u1*u2 *...*u3$$

$$u \star = u_1 * u_2 * ... * u_3 *$$

$$\star u \star = *u_1 * u_2 * ... * u_3 *$$

#### • Proof:

(2) Transfer P' into P

If P' = 
$$\left\{ \begin{bmatrix} t_1 \\ b_1 \end{bmatrix}, \begin{bmatrix} t_2 \\ b_2 \end{bmatrix}, \dots, \begin{bmatrix} t_k \\ b_k \end{bmatrix} \right\}$$
, then let

$$\mathsf{P} = \left\{ \left[ \frac{\star t_1}{\star b_1 \star} \right], \left[ \frac{\star t_2}{b_2 \star} \right], \left[ \frac{\star t_3}{b_3 \star} \right], \dots, \left[ \frac{\star t_k}{b_k \star} \right], \left[ \frac{\star \Delta}{\Delta} \right] \right\}$$

First domino: The only element as beginning

The match in P must be in shape as  $\left[\frac{\star t_1}{\star b_1 \star}\right] \dots \left[\frac{\star \Delta}{\Delta}\right]$ 

• Proof:

Suppose PCP is decidable

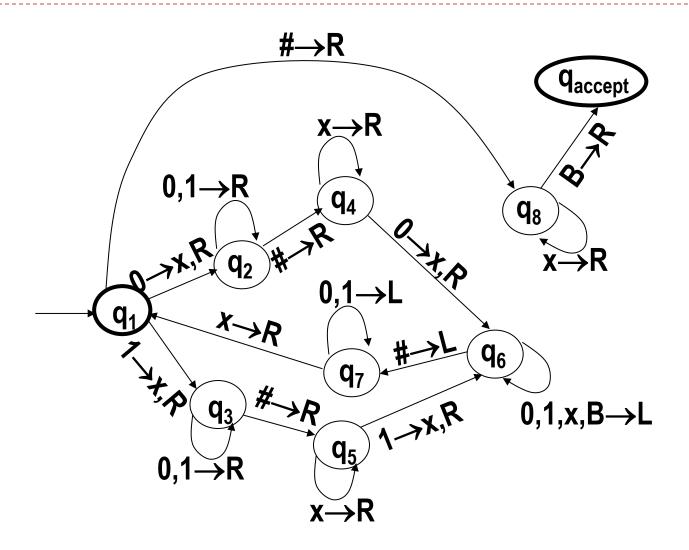
We construct TM S using PCP match or not to decide A<sub>TM</sub>

- (1) S first constructs an instance P' of the MPCP;
- ▶ (2) Transfer P' into P;

Theorem 4.11: A<sub>TM</sub> is undecidable. Contradiction!

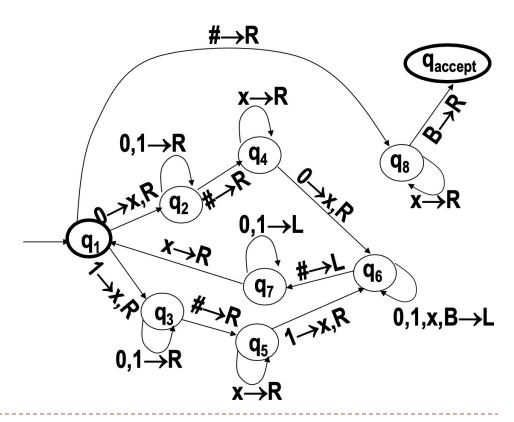
The suppose is wrong. Thus PCP is undecidable.

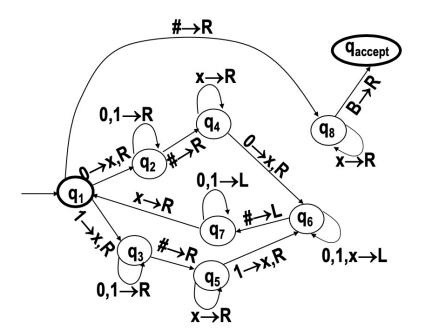
# **Example of TM:** B={ w#w | w∈{0,1}\* }

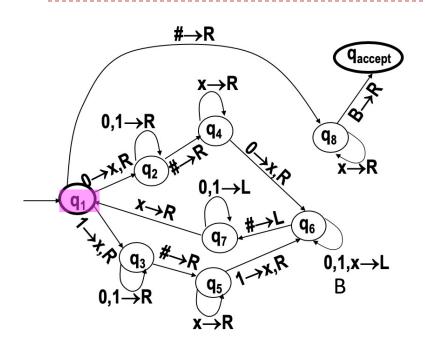


# Example of PCP to A<sub>TM</sub>

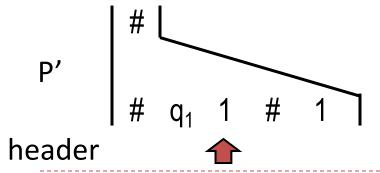
- $M = (Q, \Sigma, \Gamma, \delta, q_0, q_{accept}, q_{reject})$ 
  - Q= {q0, q1, q2}
  - $\Sigma = \{0, 1, x, \#\},$
  - $\Gamma = \{0, 1, x, \#, B\}$
  - o q0={q1}
  - $\circ$   $q_{accept} = q_{accept}$
  - q<sub>reject</sub> = all not q<sub>accept</sub>

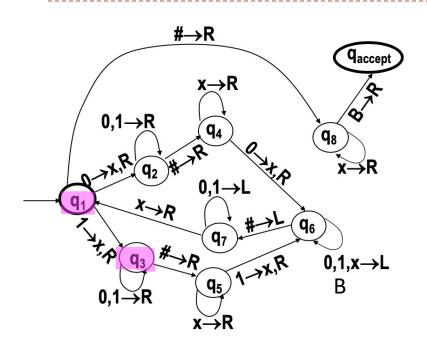




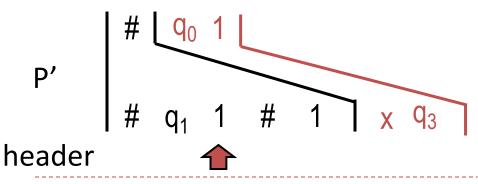


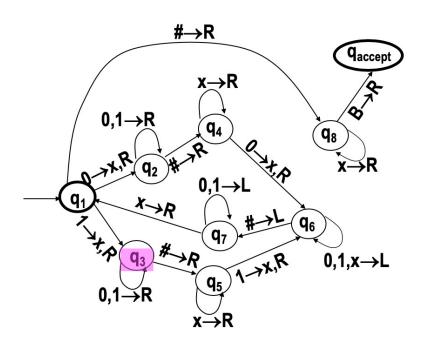
Put 
$$\left[\frac{\#}{\#q_0w_1w_2...w_n\#}\right]$$
 into P' as  $\left[\frac{t_1}{b_1}\right]$ 



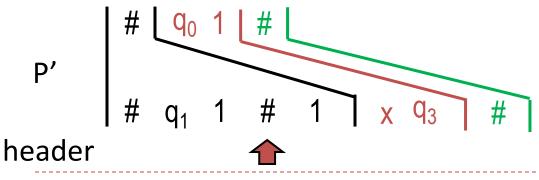


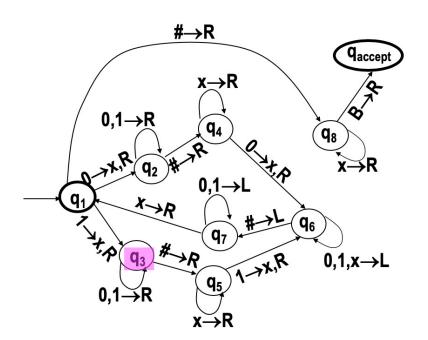
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



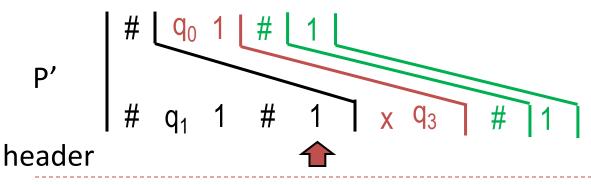


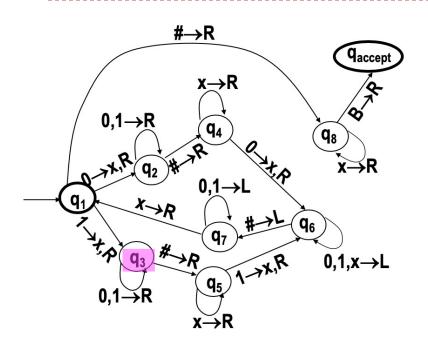
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



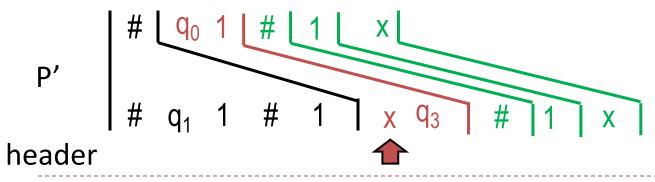


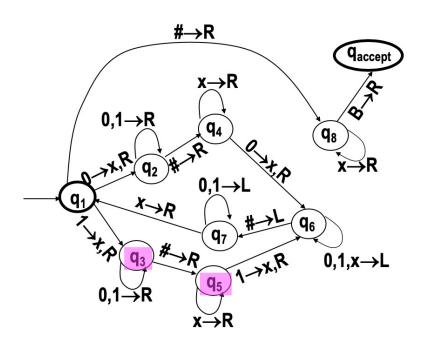
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



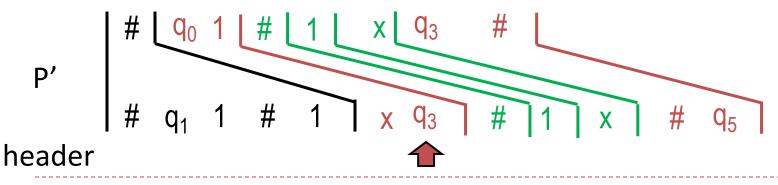


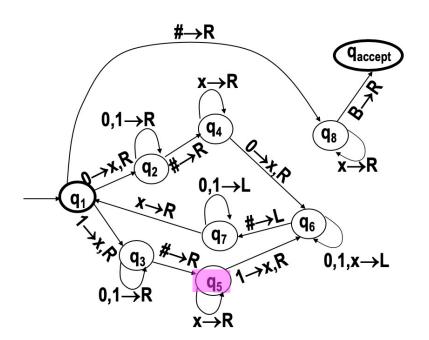
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



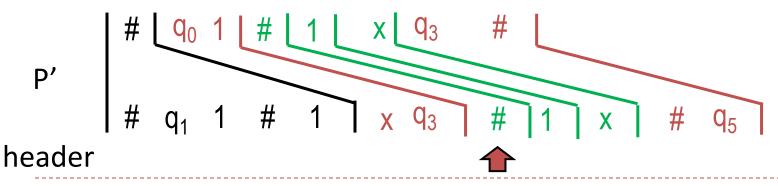


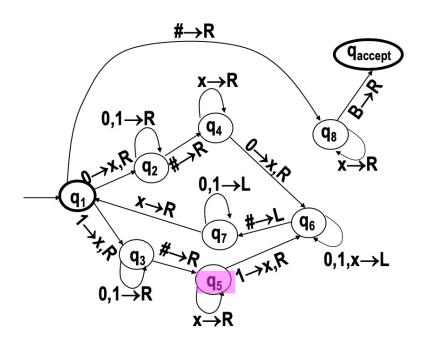
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



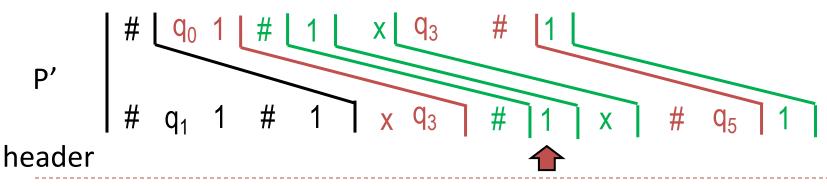


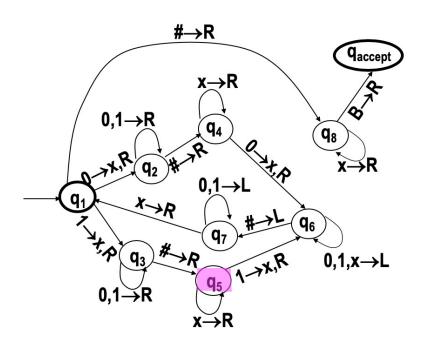
if 
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,  
then put  $\left[\frac{qa}{br}\right]$  into P'



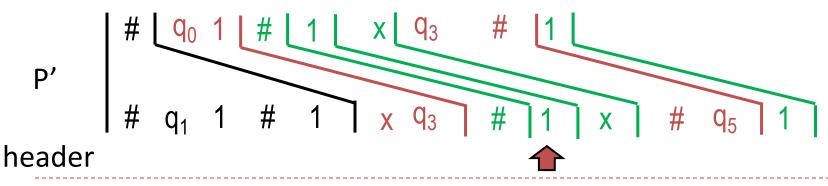


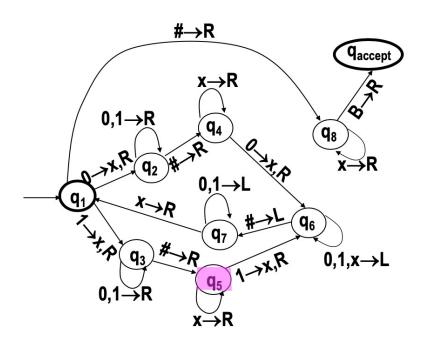
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



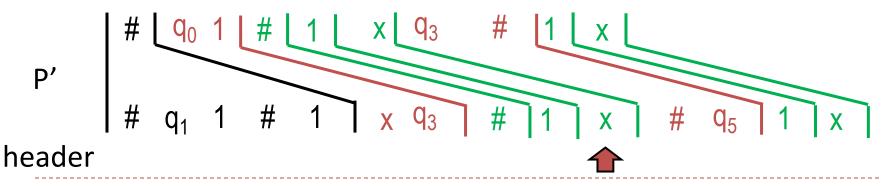


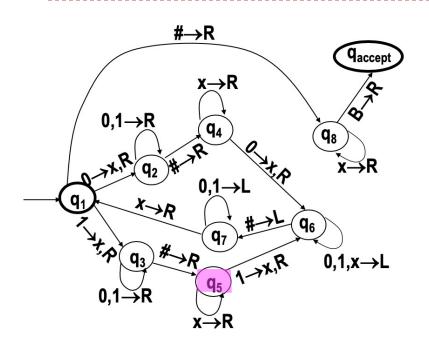
if 
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,  
then put  $\left[\frac{qa}{br}\right]$  into P'



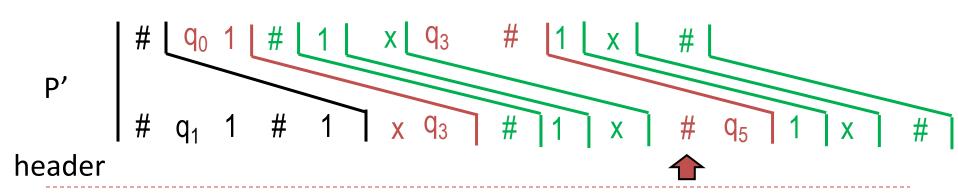


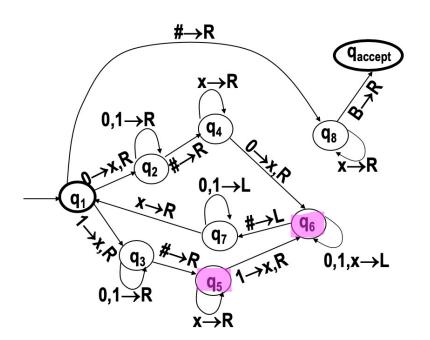
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



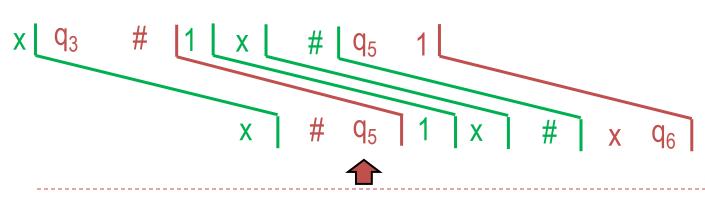


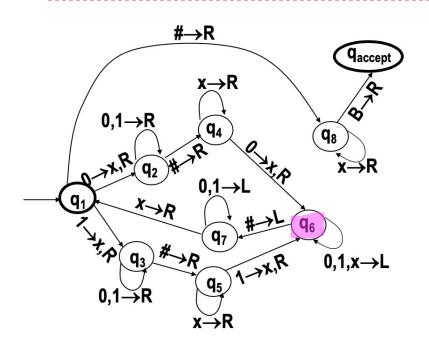
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



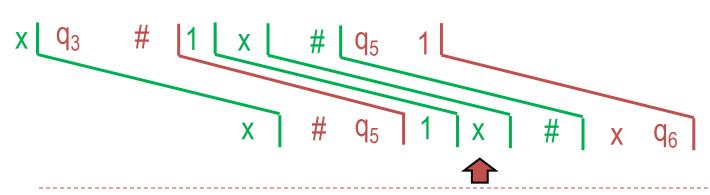


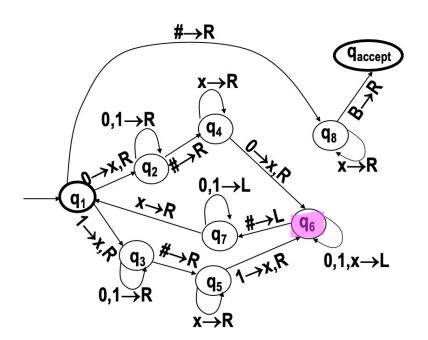
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



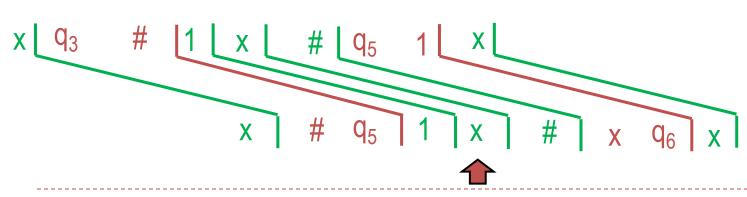


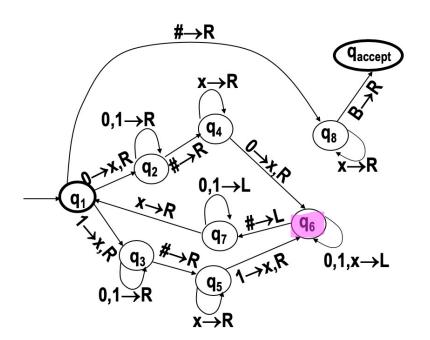
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



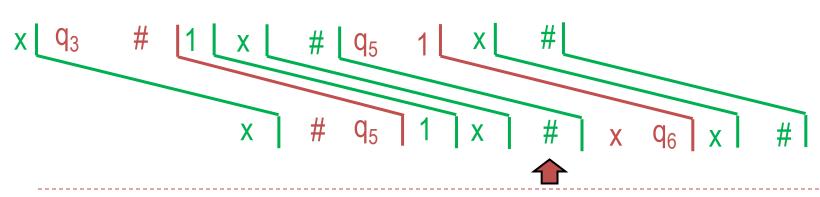


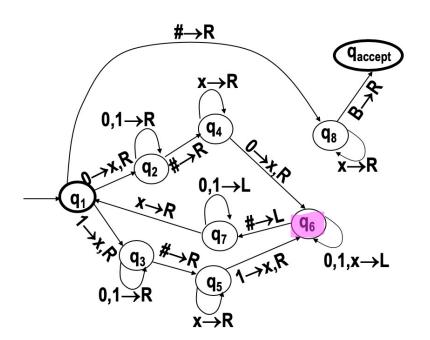
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



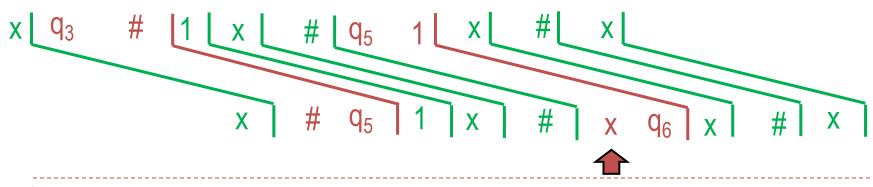


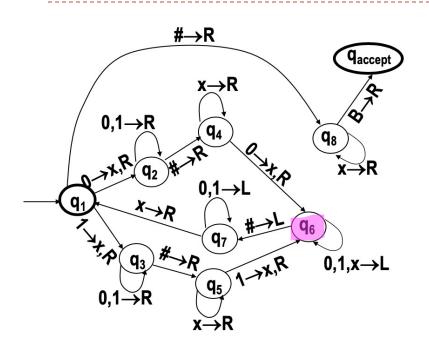
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

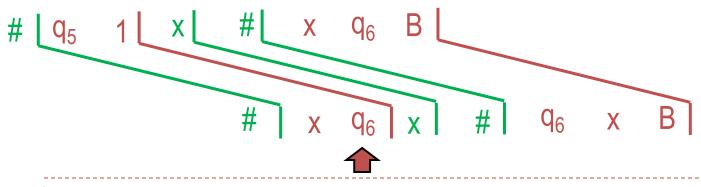


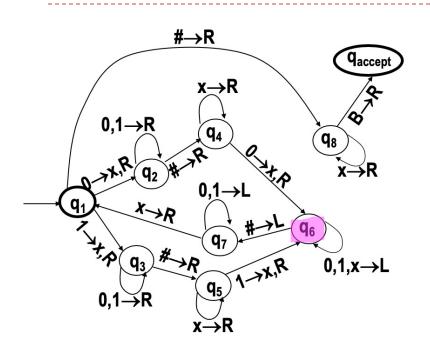


if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

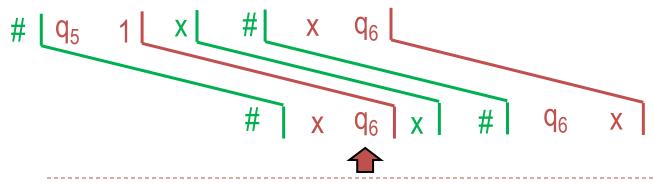


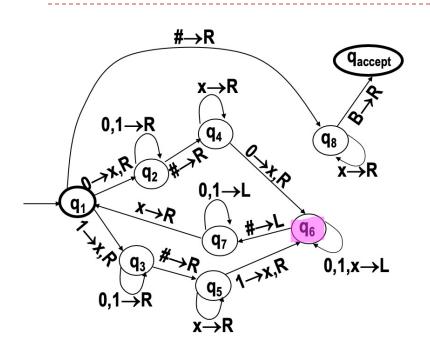


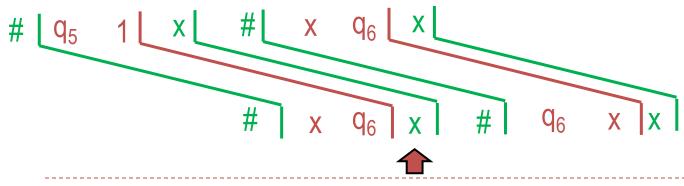


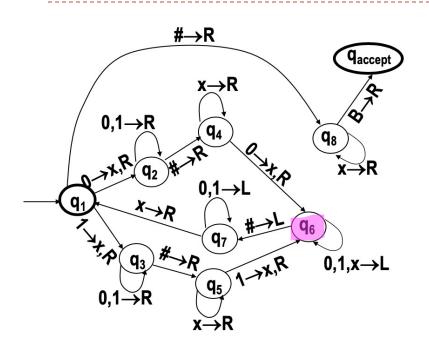


if  $\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{L})$ , then put  $\left[\frac{cqa}{rcb}\right]$  into P', c is the element on the left of q

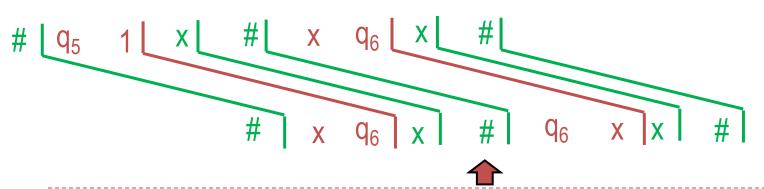


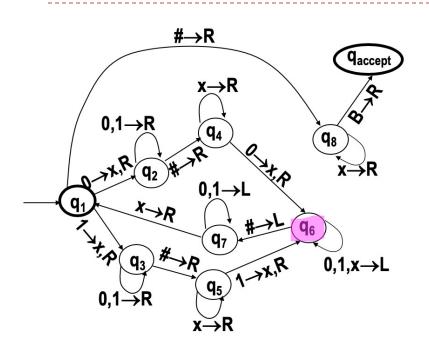


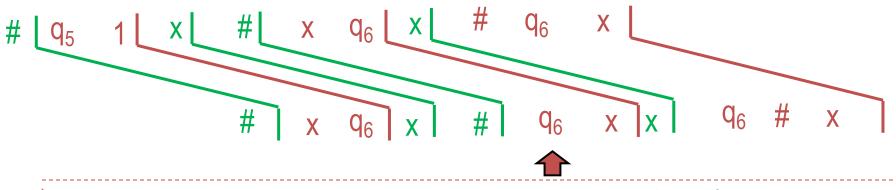


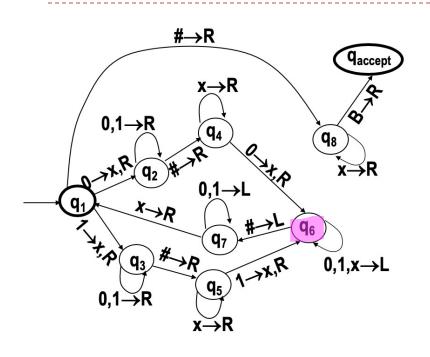


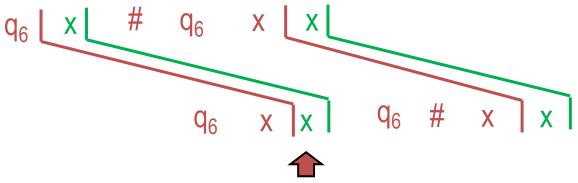
if  $\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{L})$ , then put  $\left[\frac{cqa}{rcb}\right]$  into P', c is the element on the left of q

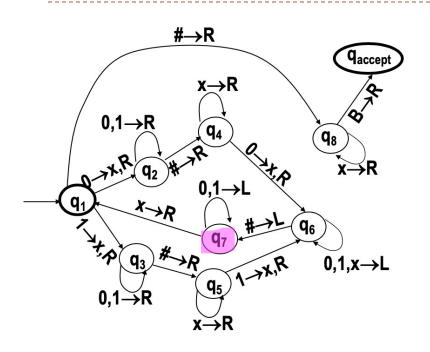


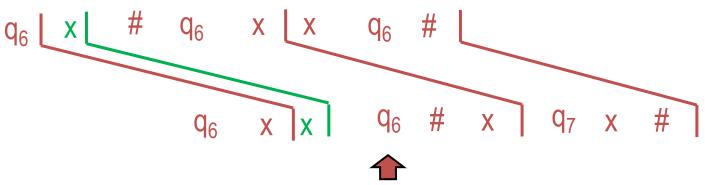


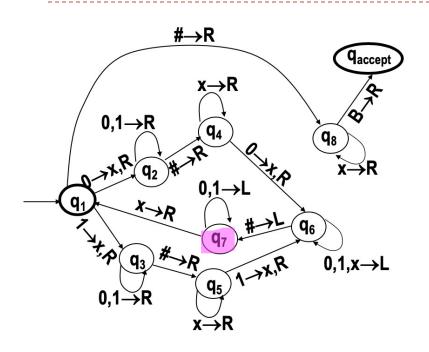




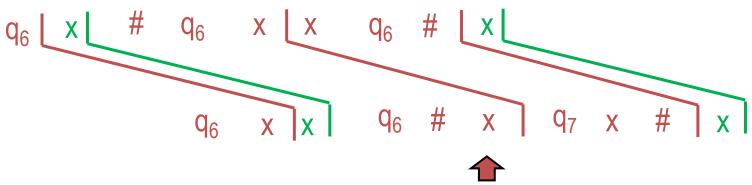


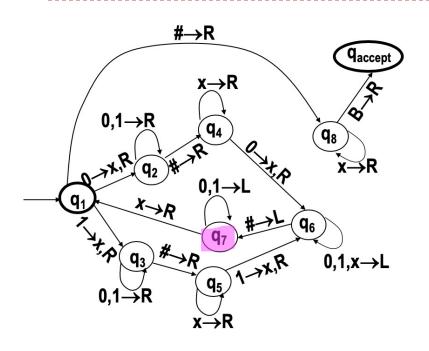




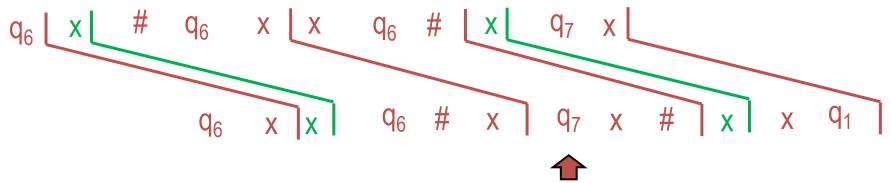


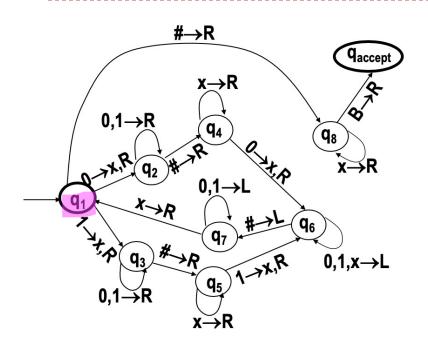
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



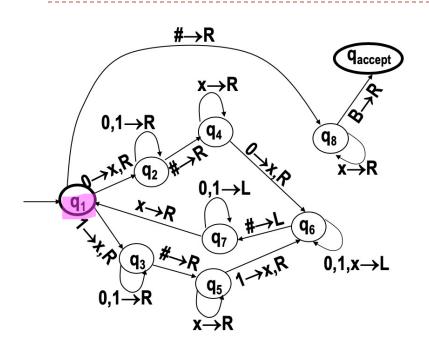


if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

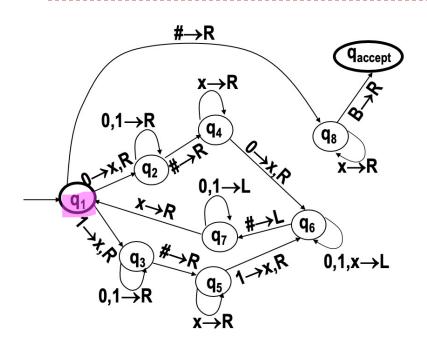




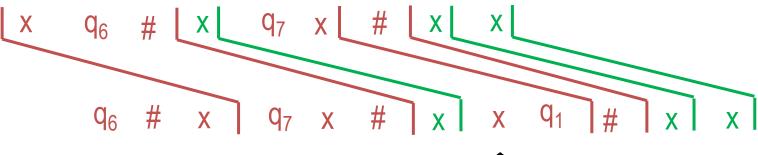
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

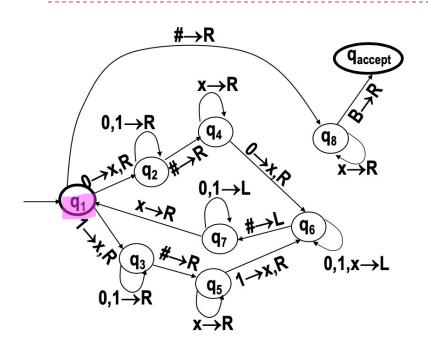


if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

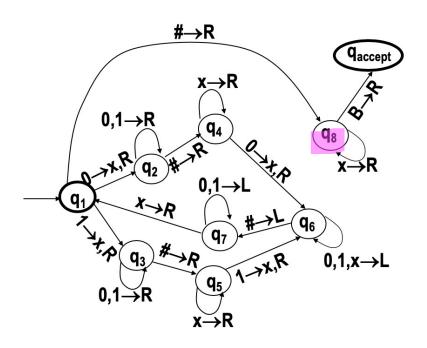


if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

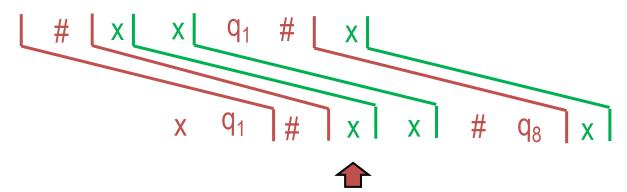


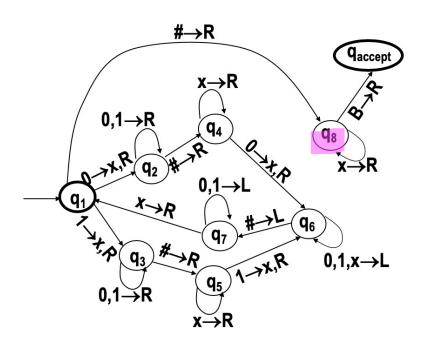


if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'

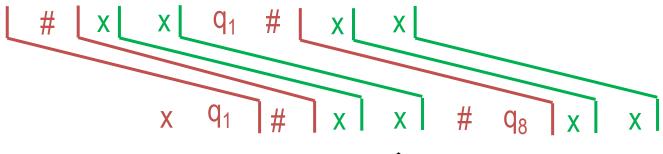


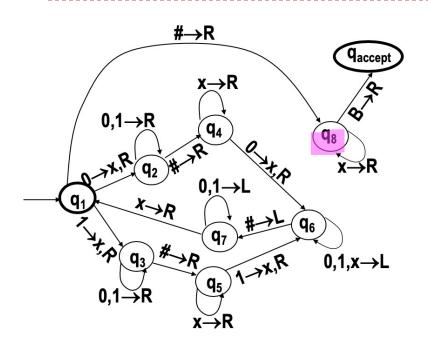
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



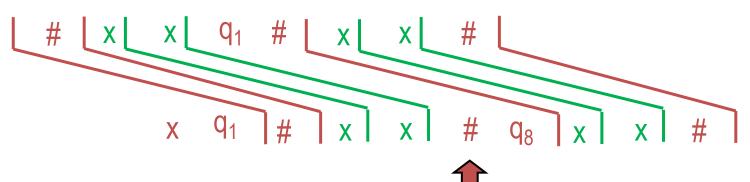


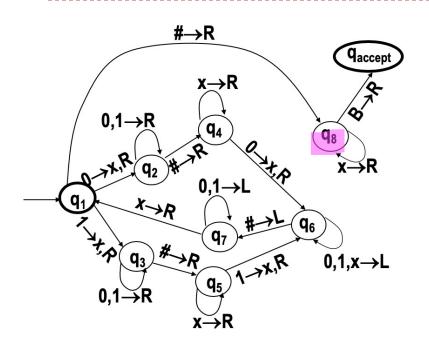
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



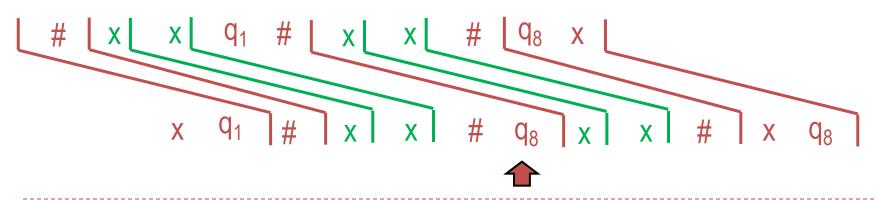


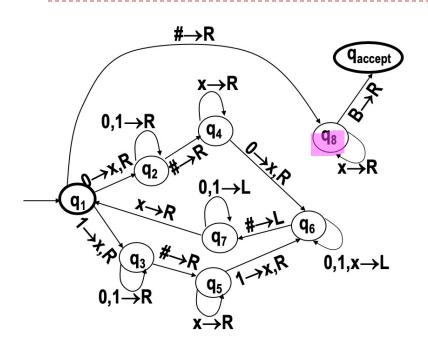
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



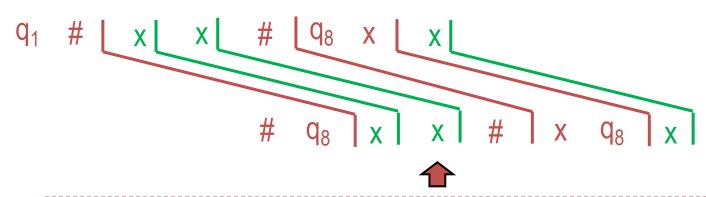


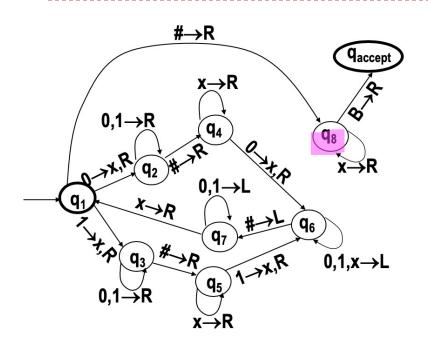
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



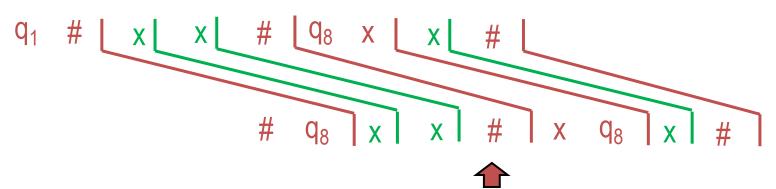


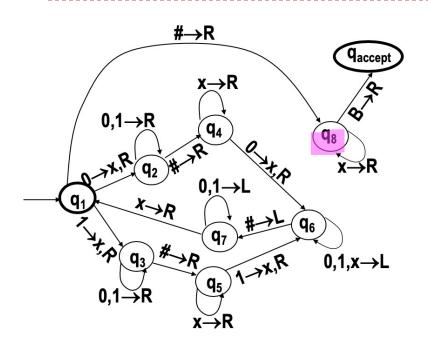
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



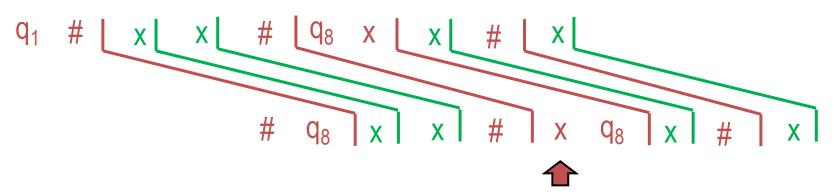


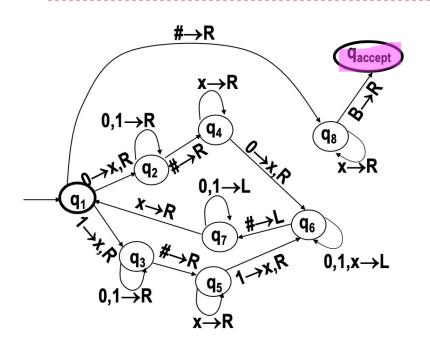
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



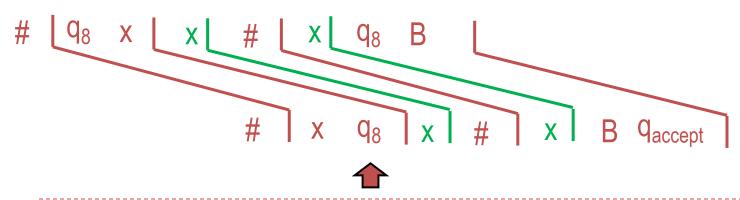


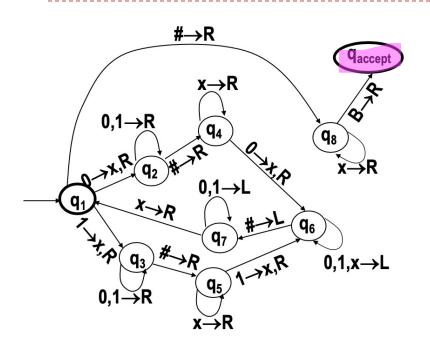
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



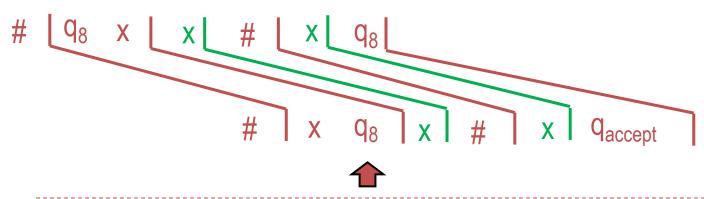


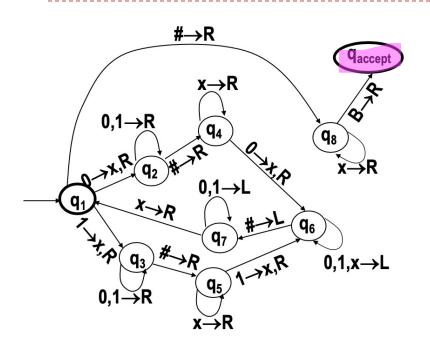
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



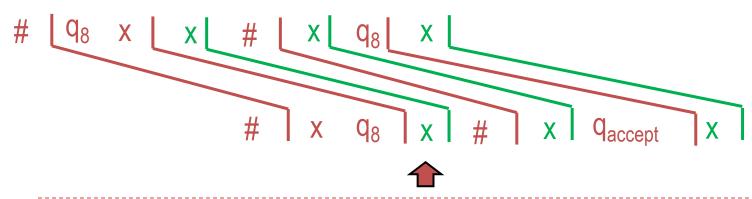


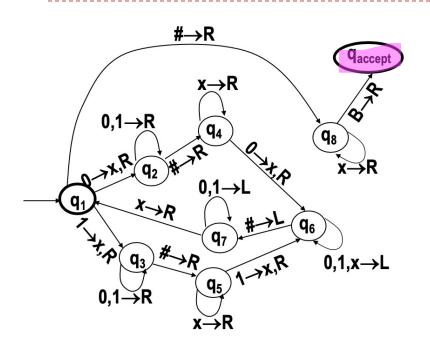
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



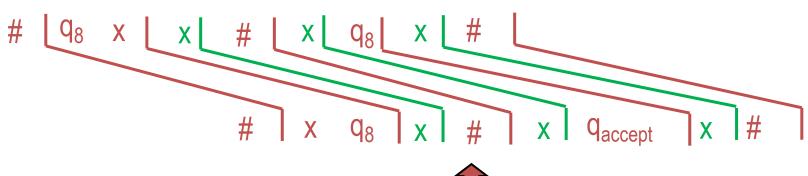


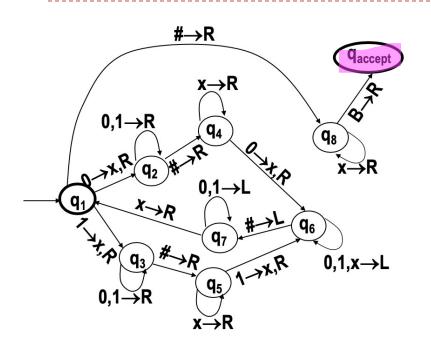
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



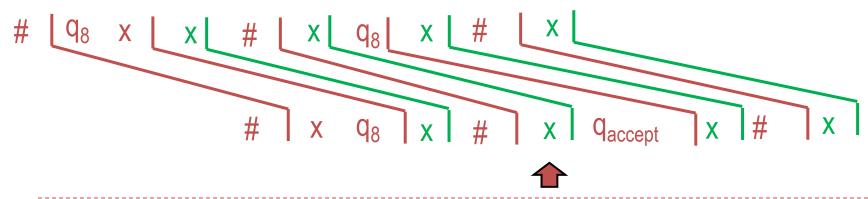


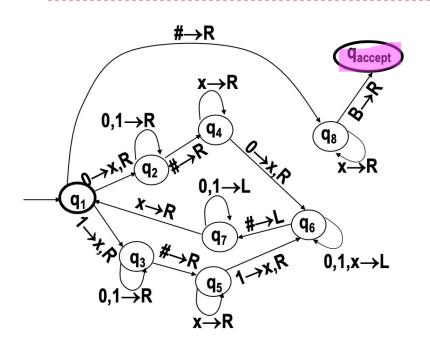
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



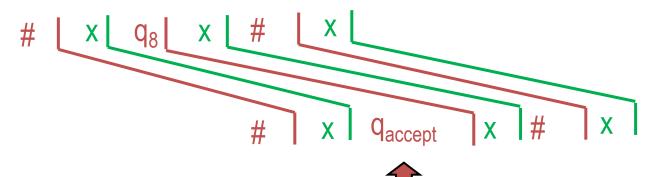


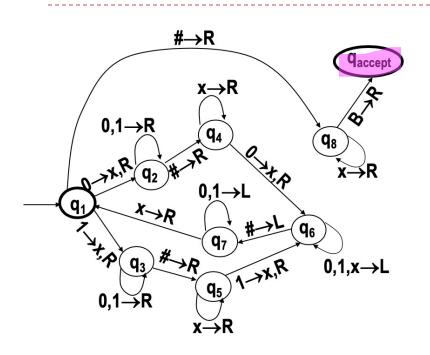
if 
$$\delta(q, a) = (r, b, R)$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



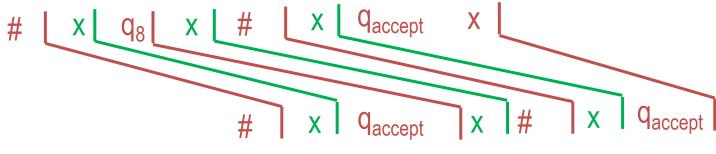


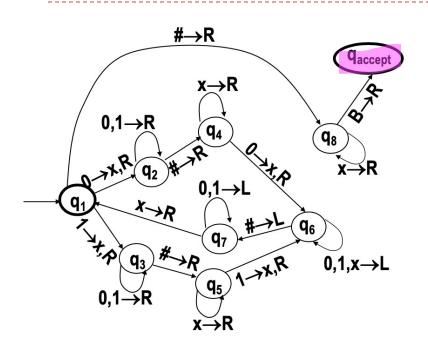
if 
$$\delta(\mathbf{q}, \mathbf{a}) = (\mathbf{r}, \mathbf{b}, \mathbf{R})$$
,  
then put  $\left[\frac{qa}{br}\right]$  into P'



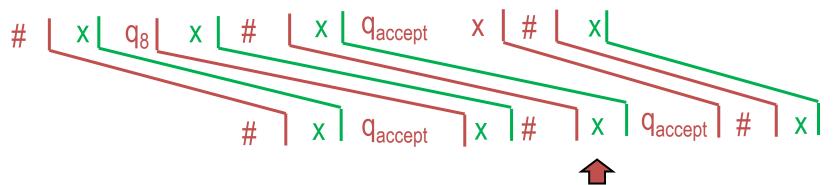


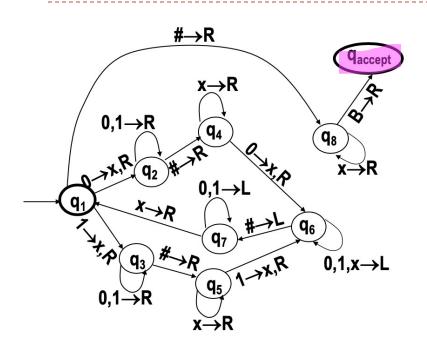
put 
$$\left[\frac{aq_{accept}}{q_{accept}}\right]$$
 and  $\left[\frac{q_{accept}a}{q_{accept}}\right]$  into P'



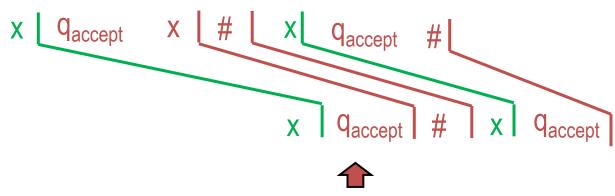


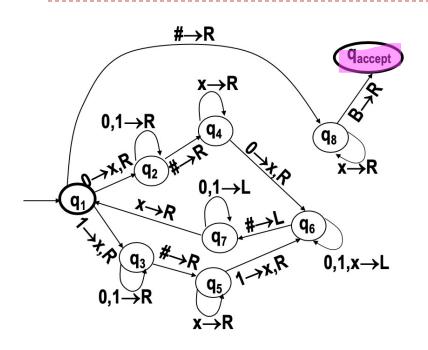




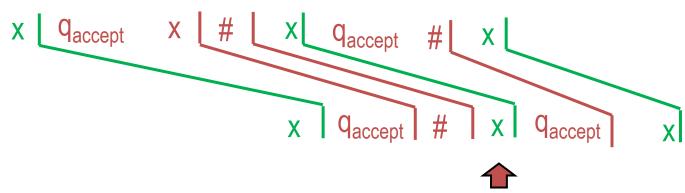


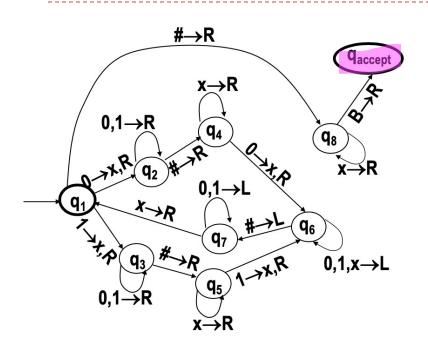




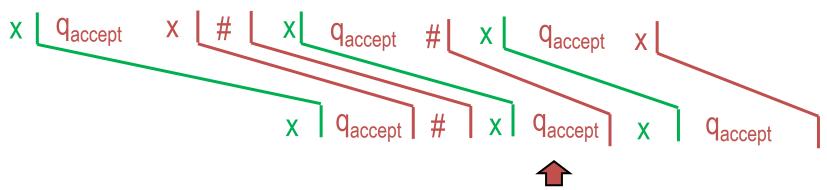


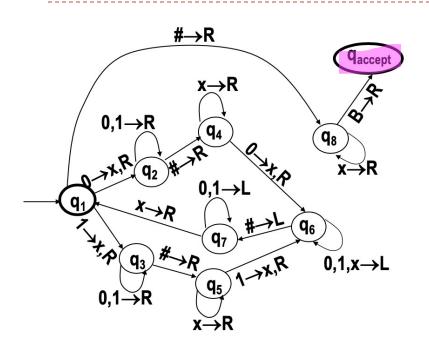




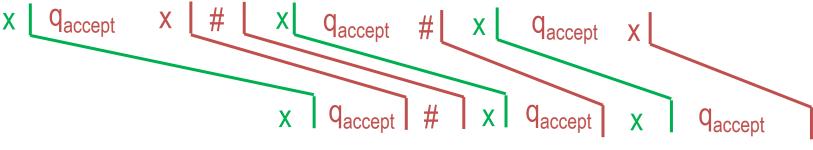


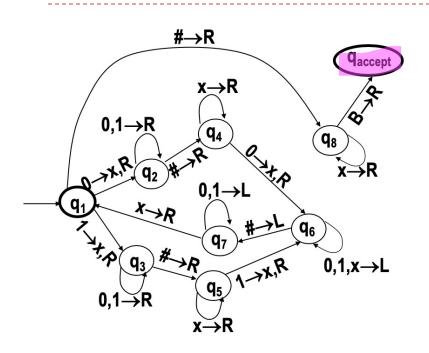




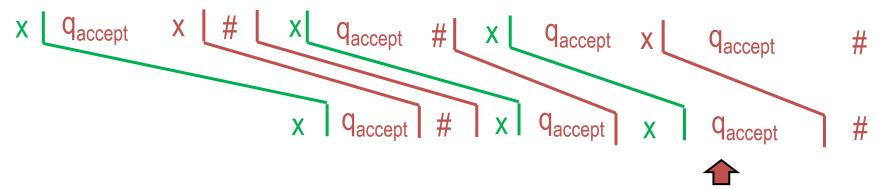


Finish the match, put  $\left[\frac{q_{accept}^{\#\#}}{\#}\right] \text{ into P'}$ 



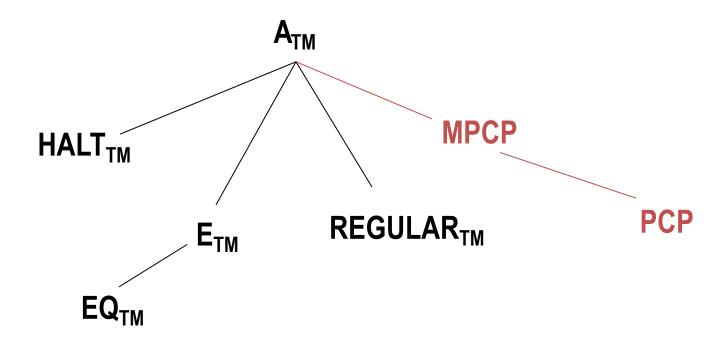


Finish the match, put  $\left[\frac{q_{accept}^{\#}}{\#}\right]$  into P'



#### Conclusion

Relationship of languages on reducibility



#### **Conclusion**

#### Closure on operations

	Complement $\overline{A}$	Intersection ∩	<b>Union</b> ∪	Star <i>A</i> *
Regular/DFA/ NFA	√	~	<b>√</b>	<b>√</b>
CFL/ PDA	×	×	<b>√</b>	~
Turing- decidable TM	√	<b>√</b>	√	<b>√</b>
Turing- recognizable TM	×	<b>√</b>	√	<b>√</b>