

# CS 6041

## Theory of Computation

### Nondeterministic finite automata

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<https://kevinsuo.github.io/>

# Design a NFA for a language

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- Step 1: list all possible states
- Step 2: draw all the transitions between the states
- Step 3: add start and accept states



# Some practices

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- $L_1 = \{\text{Set of all strings that end with 0}\}, \Sigma = \{0,1\};$

*Can anyone draw the NFA?*



# Some practices

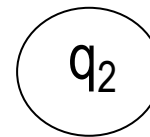
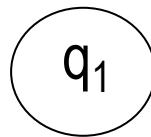
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- $L_1 = \{\text{Set of all strings that end with 0}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

q1: all the strings

q2: last letter is 0



NFA of  $L_1$



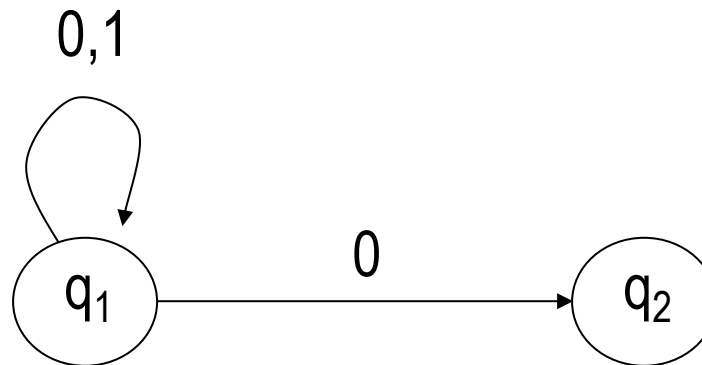
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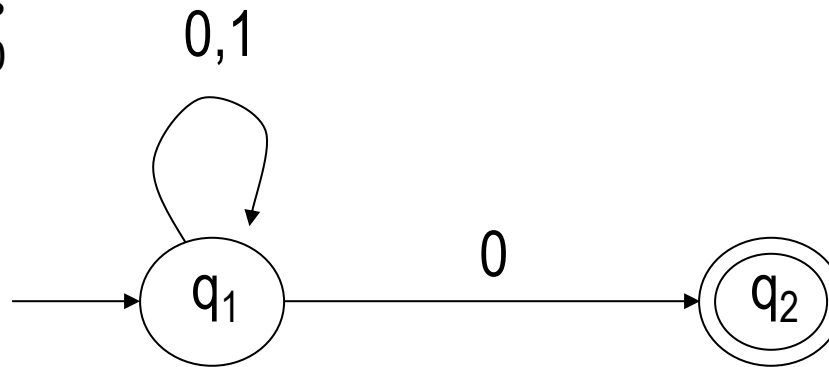
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*Can anyone draw the NFA?*

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NFA of  $L_1$

# Some practices

---

- $L_2 = \{\text{Set of all strings that start with 0}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*



# Some practices

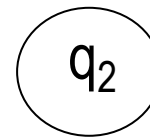
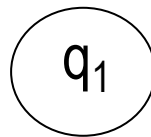
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*Can anyone draw the NFA?*

q1: empty string

q2: first letter is 0



NFA of  $L_2$





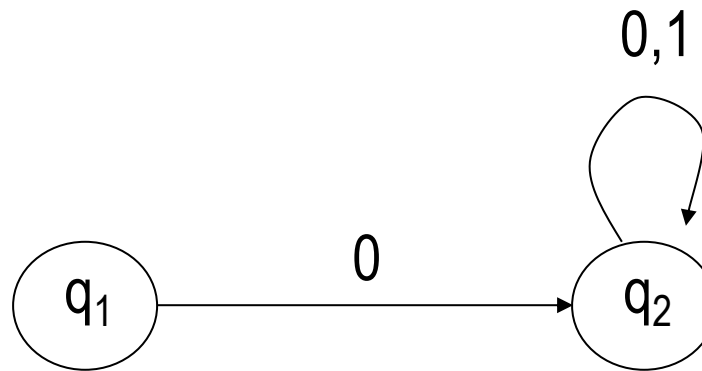
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NFA of  $L_2$

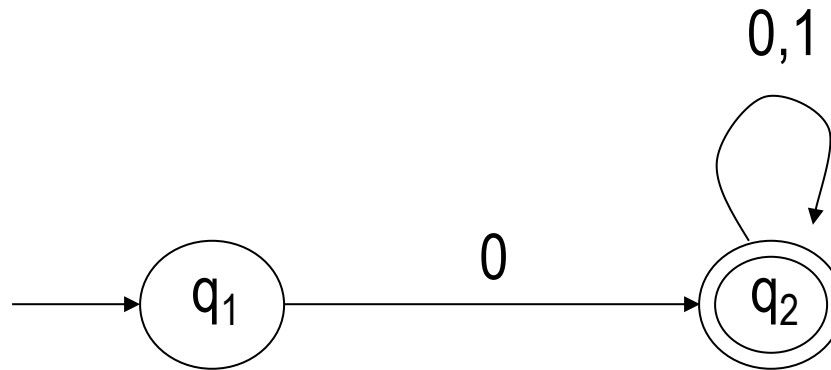
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- $L_2 = \{\text{Set of all strings that start with 0}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

q1: empty string  
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NFA of  $L_2$

# Some practices

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- $L_3 = \{\text{Set of all strings that length is 2}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

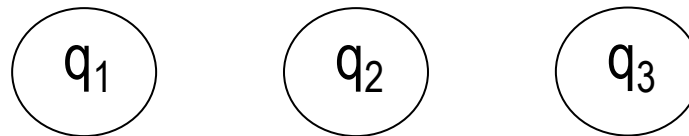


# Some practices

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- $L_3 = \{\text{Set of all strings that length is 2}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*



$q_1$ : strings that length is 0

$q_2$ : strings that length is 1

$q_3$ : strings that length is 2

NFA of  $L_3$

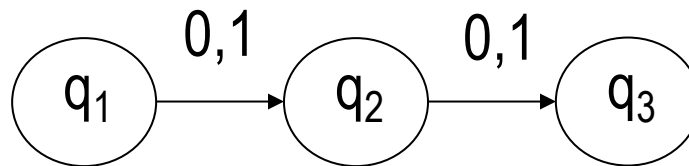


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*Can anyone draw the NFA?*



$q_1$ : strings that length is 0

$q_2$ : strings that length is 1

$q_3$ : strings that length is 2

NFA of  $L_3$

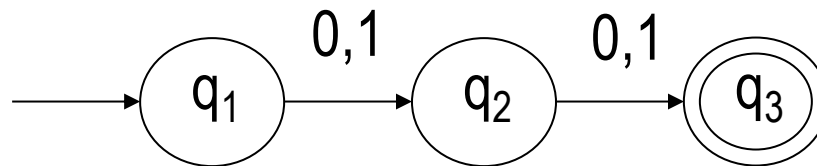


# Some practices

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- $L_3 = \{\text{Set of all strings that length is 2}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*



$q_1$ : strings that length is 0

$q_2$ : strings that length is 1

$q_3$ : strings that length is 2

NFA of  $L_3$



# Some practices

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- $L_4 = \{\text{Set of all strings that contain '0'}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*



# Some practices

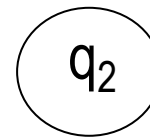
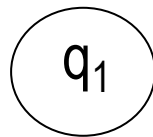
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- $L_4 = \{\text{Set of all strings that contain '0'}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

q1: all strings

q2: strings that contain 0



NFA of  $L_4$





# Some practices

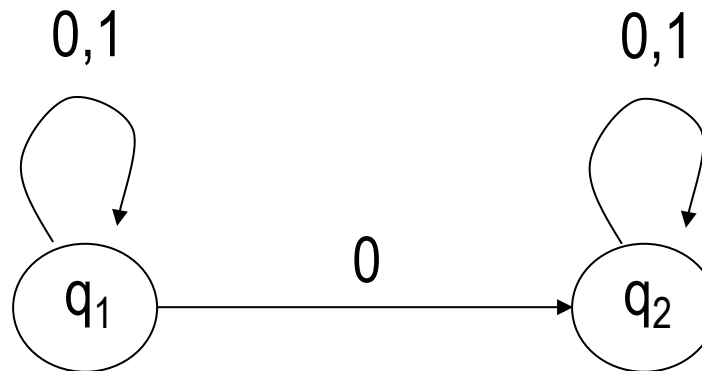
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NFA of  $L_4$

# Some practices

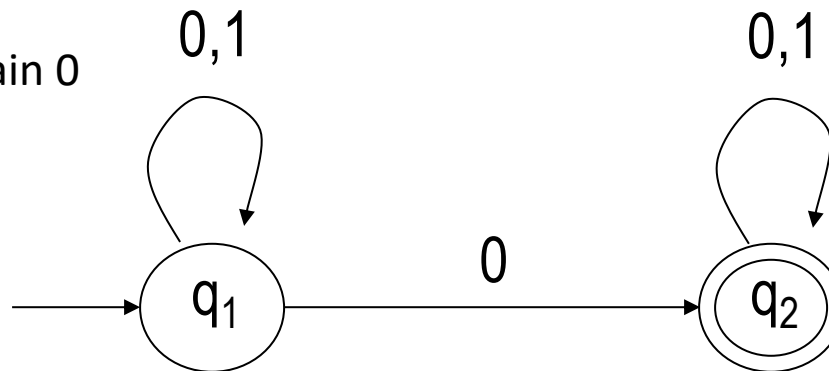
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*Can anyone draw the NFA?*

q1: all strings

q2: strings that contain 0



NFA of  $L_4$

# Some practices

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- $L_5 = \{\text{Set of all strings that starts with '10'}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

NFA of  $L_5$



# Some practices

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- $L_5 = \{\text{Set of all strings that starts with '10'}\}$ ,  $\Sigma = \{0,1\}$ ;

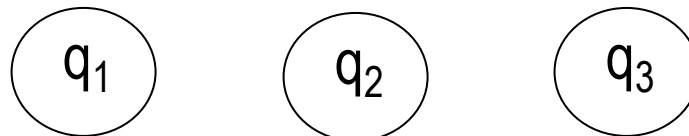
*Can anyone draw the NFA?*

q1: all strings

q2: strings that start with 1

q3: strings that start with 10

NFA of  $L_5$



# Some practices

---

- $L_5 = \{\text{Set of all strings that starts with '10'}\}$ ,  $\Sigma = \{0,1\}$ ;

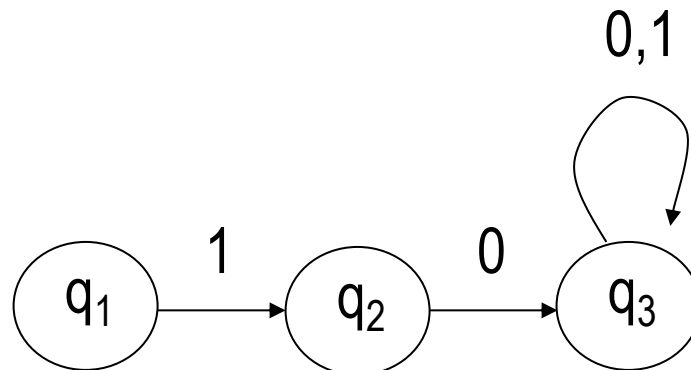
*Can anyone draw the NFA?*

q1: all strings

q2: strings that start with 1

q3: strings that start with 10

NFA of  $L_5$



# Some practices

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- $L_5 = \{\text{Set of all strings that starts with '10'}\}$ ,  $\Sigma = \{0,1\}$ ;

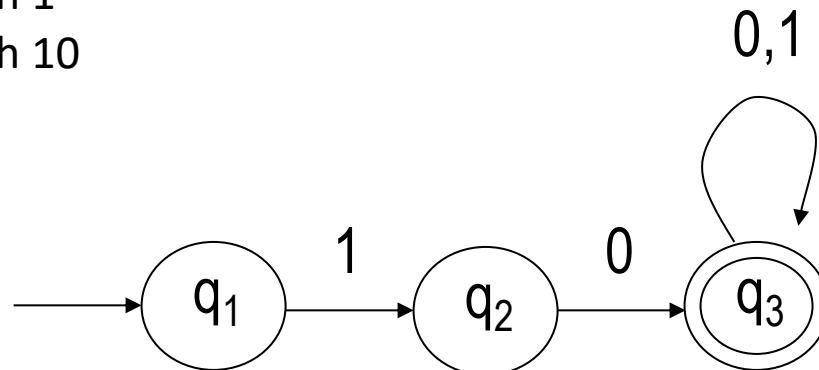
*Can anyone draw the NFA?*

q1: all strings

q2: strings that start with 1

q3: strings that start with 10

NFA of  $L_5$



# Some practices

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- $L_6 = \{\text{Set of all strings that contain '01'}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

NFA of  $L_6$



# Some practices

---

- $L_6 = \{\text{Set of all strings that contain '01'}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

q1: all strings

q2: strings that contain 0

q3: strings that contain 01

NFA of  $L_6$





# Some practices

---

- $L_6 = \{\text{Set of all strings that contain '01'}\}$ ,  $\Sigma = \{0,1\}$ ;

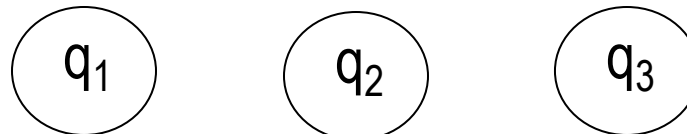
*Can anyone draw the NFA?*

q1: all strings

q2: strings that contain 0

q3: strings that contain 01

NFA of  $L_6$



# Some practices

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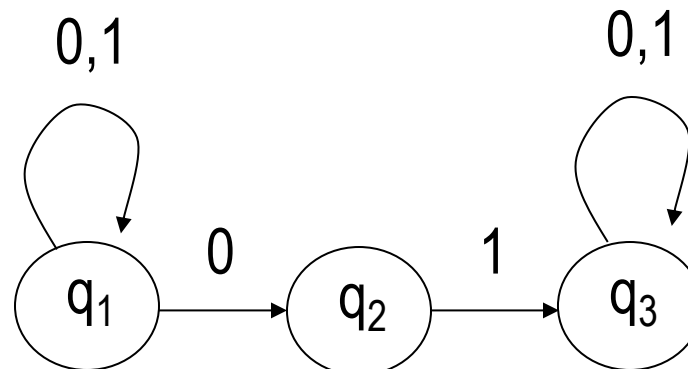
*Can anyone draw the NFA?*

q1: all strings

q2: strings that contain 0

q3: strings that contain 01

NFA of  $L_6$



# Some practices

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- $L_6 = \{\text{Set of all strings that contain '01'}\}$ ,  $\Sigma = \{0,1\}$ ;

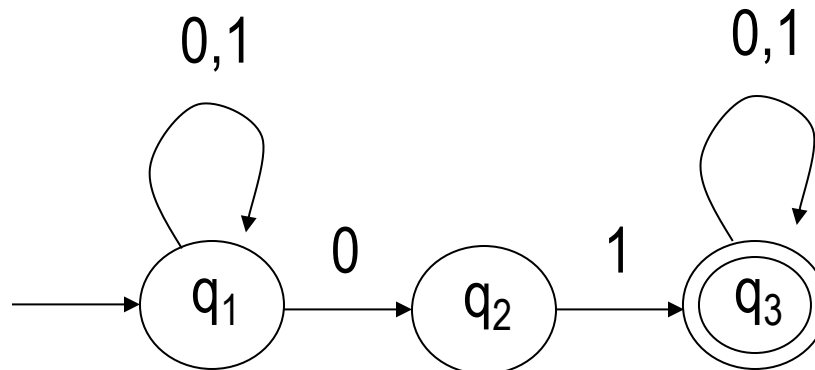
*Can anyone draw the NFA?*

q1: all strings

q2: strings that contain 0

q3: strings that contain 01

NFA of  $L_6$



# Some practices

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- $L_7 = \{\text{Set of all strings that end with '11'}\}$ ,  $\Sigma = \{0,1\}$ ;

*Can anyone draw the NFA?*

NFA of  $L_7$



# Some practices

---

- $L_7 = \{\text{Set of all strings that end with '11'}\}$ ,  $\Sigma = \{0,1\}$ ;

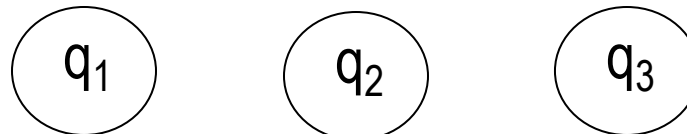
*Can anyone draw the NFA?*

q1: all strings

q2: strings that end with 1

q3: strings that end with 11

NFA of  $L_7$



# Some practices

---

- $L_7 = \{\text{Set of all strings that end with '11'}\}$ ,  $\Sigma = \{0,1\}$ ;

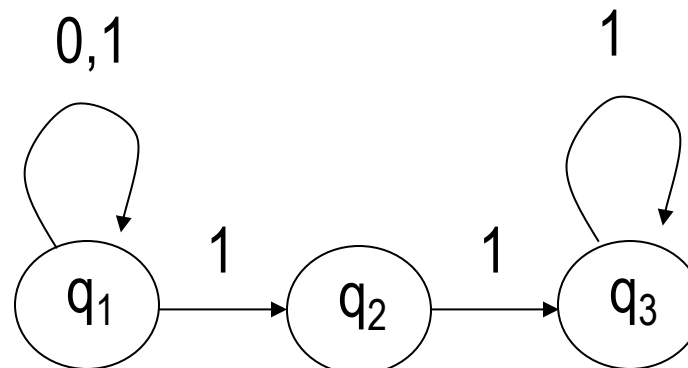
*Can anyone draw the NFA?*

q1: all strings

q2: strings that end with 1

q3: strings that end with 11

NFA of  $L_7$



# Some practices

---

- $L_7 = \{\text{Set of all strings that end with '11'}\}$ ,  $\Sigma = \{0,1\}$ ;

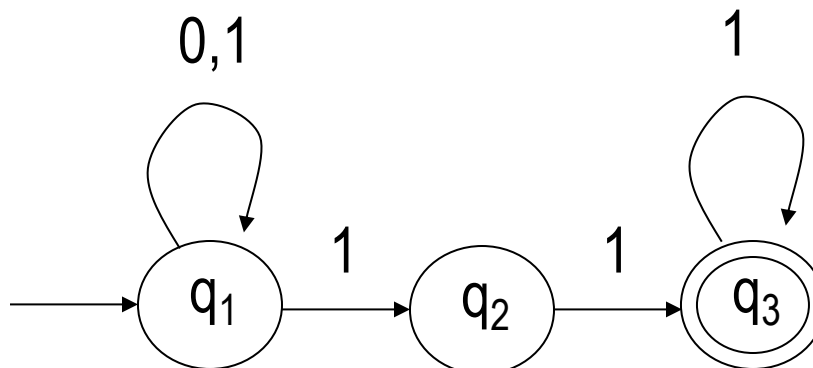
*Can anyone draw the NFA?*

q1: all strings

q2: strings that end with 1

q3: strings that end with 11

NFA of  $L_7$



# Equivalence of NFAs and DFAs

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- Definition of equivalence
  - two machines are *equivalent* if they recognize the same language





# Equivalence of NFAs and DFAs

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- Theorem: Every nondeterministic finite automaton has an equivalent deterministic finite automaton.

- Proof idea:

For given NFA, build an equivalent DFA

Use DFA to simulate NFA

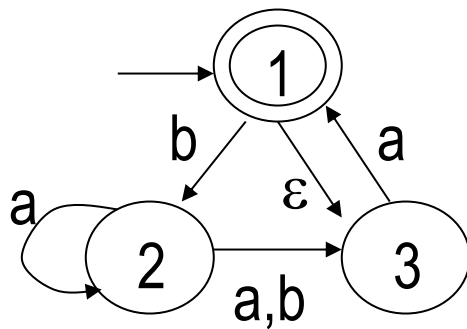
- DFA records all branches of NFA
- If NFA has  $k$  states, then DFA have totally  $2^k$  states (subset of  $k$  states in NFA)



# Equivalence of NFAs and DFAs

- Closure  $E(R)$  on  $\varepsilon$

Collection of states that can be reached from members of  $R$  by going only along  $\varepsilon$  arrows, including the members of  $R$  themselves.



$$E(\{1\}) = \{1, 3\}$$

# Equivalence of NFAs and DFAs

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- Step 1: Draw all the states in DFAs
- Step 2: Define the transitions in DFAs based on the NFAs
- Step 3: Define the start state and accept state in DFAs
- Step 4: Remove all inaccessible states

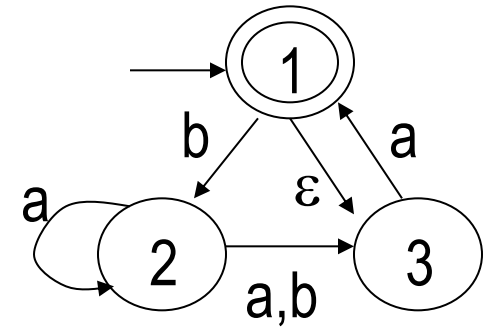


# Example

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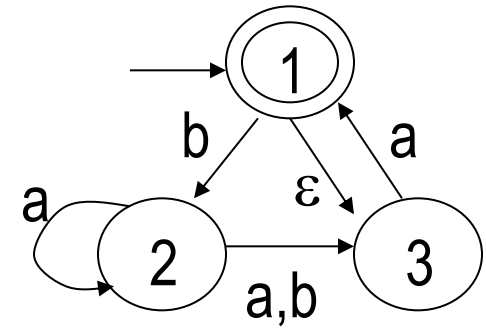
- NFA  $N_4 = (\{1,2,3\}, \{a,b\}, \delta, 1, \{1\})$

What is its equivalent DFA?



# Example

- List all subset states



$\emptyset$

$\{1\}$

$\{2\}$

$\{1,2\}$

$\{3\}$

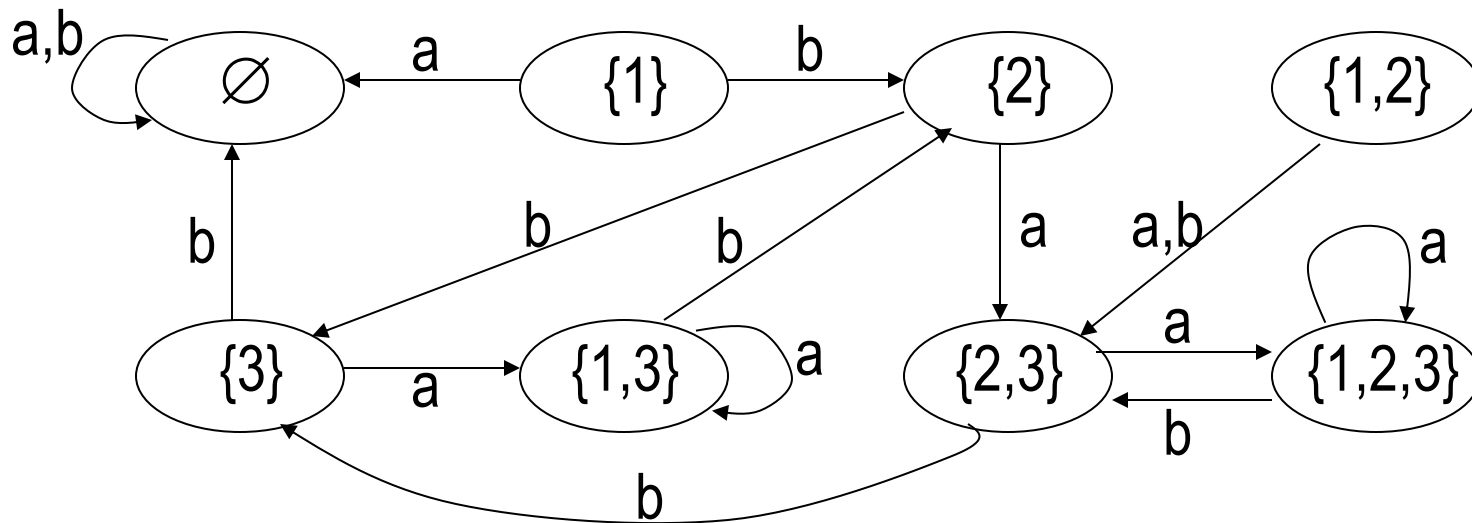
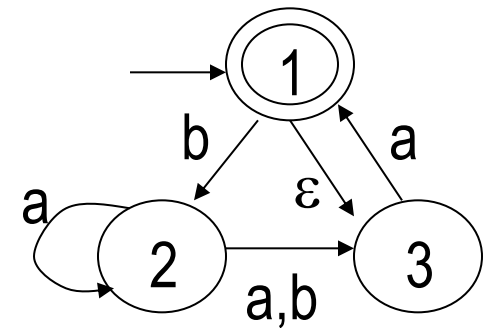
$\{1,3\}$

$\{2,3\}$

$\{1,2,3\}$

# Example

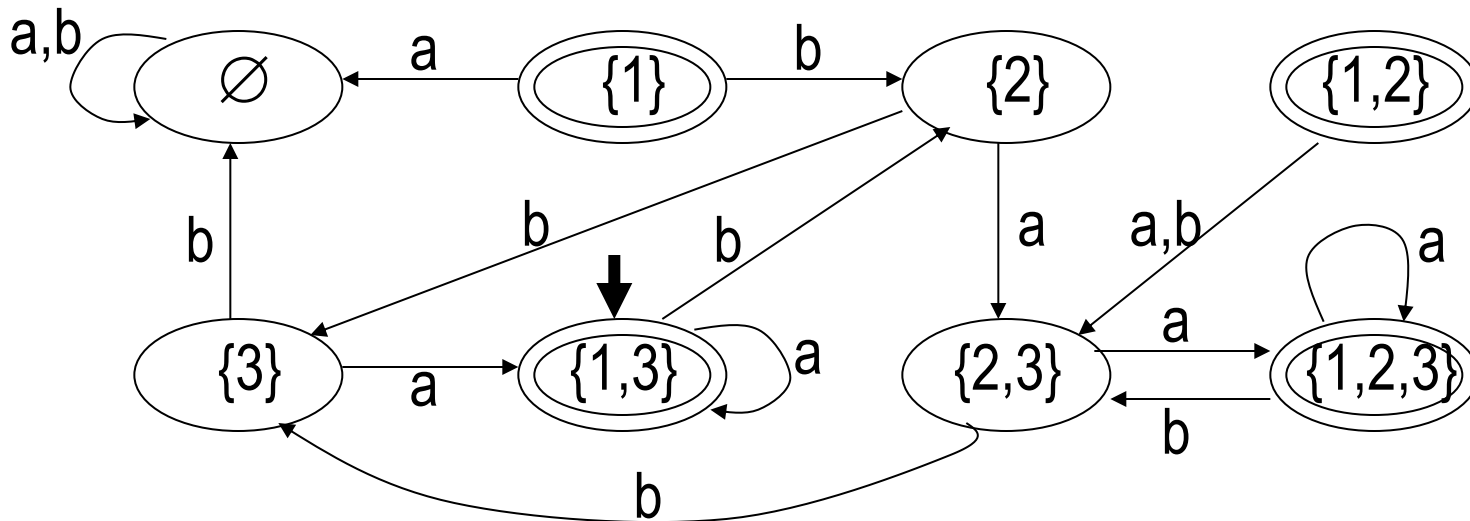
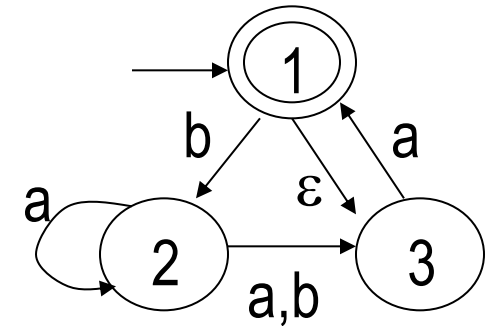
- Add transitions



# Example

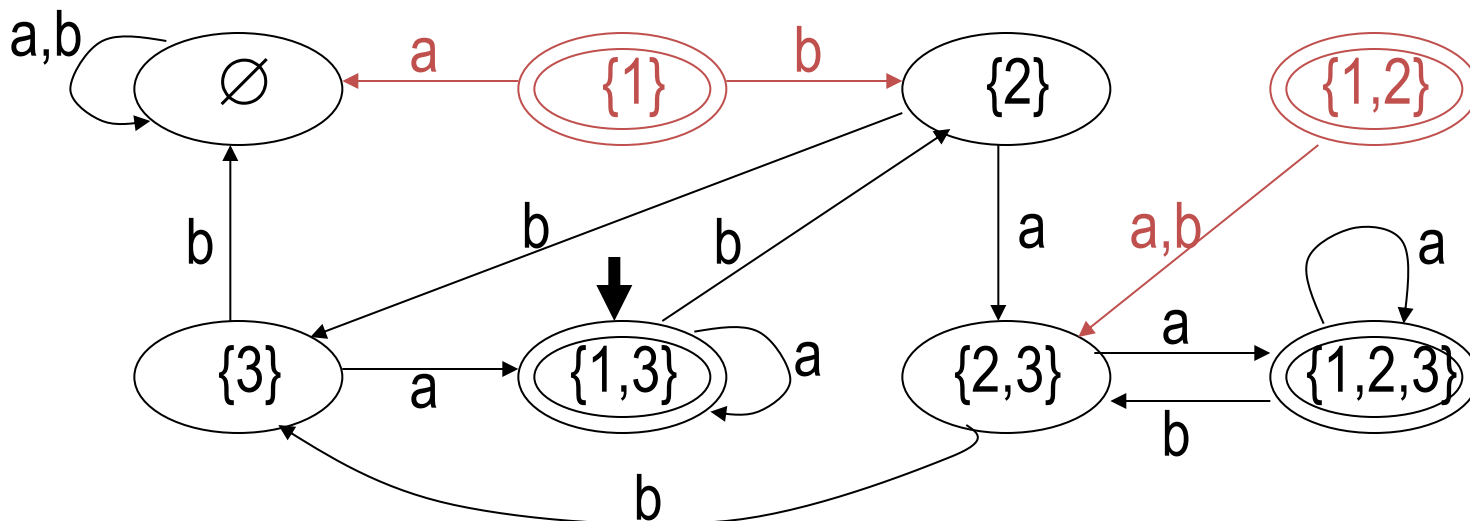
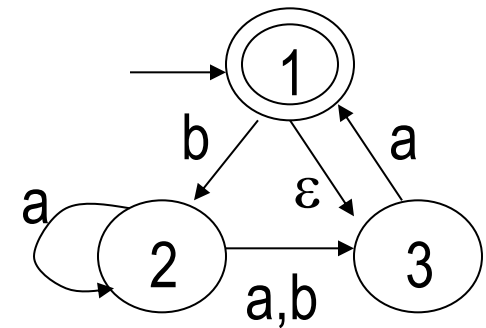
Collection of states that can be reached from members of  $R$  by going only along  $\varepsilon$  arrows, including the members of  $R$  themselves.

- Start state:  $E(\{1\}) = \{1,3\}$
- Accept state: all states with 1



# Example

- Remove inaccessible state

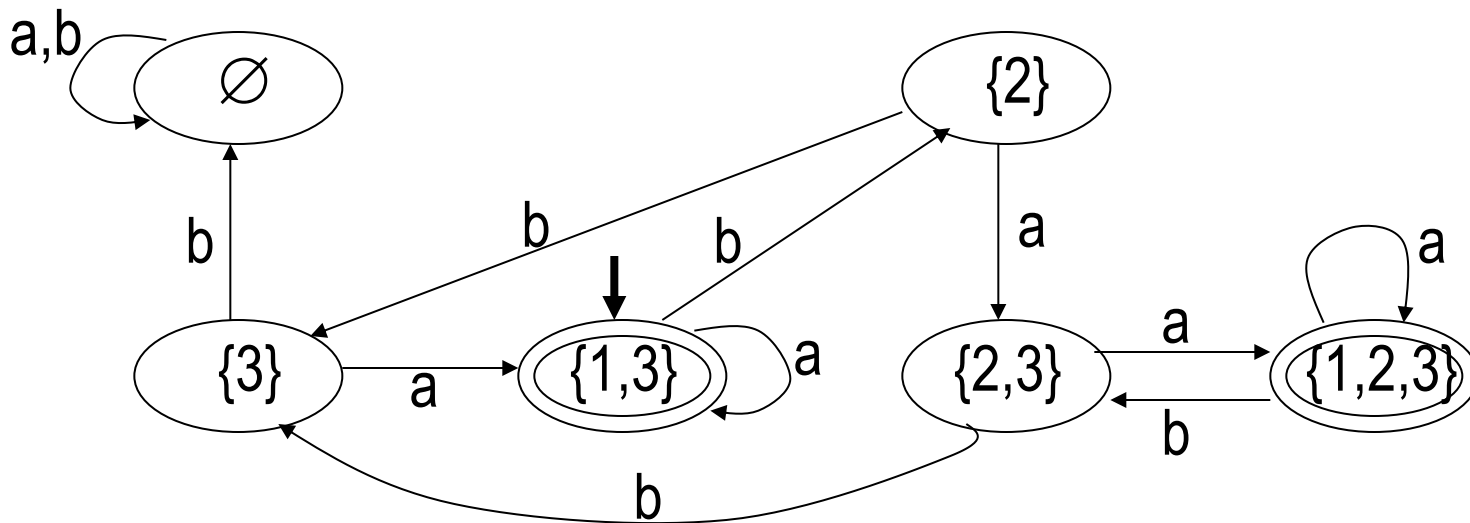
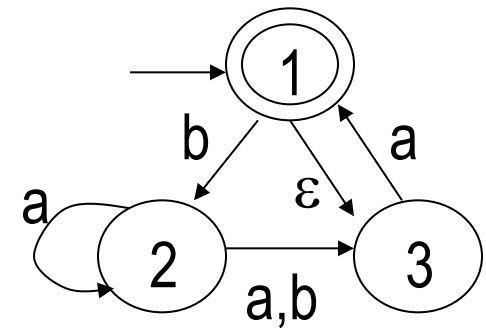




# Example

- Remove inaccessible state

- $\{1\}, \{1,2\}$



# Equivalence of NFAs and DFAs

- Theorem: Every NFA has an equivalent DFA.
- Proof:

let NFA  $N = (Q, \Sigma, \delta, q_0, F)$ ,

build DFA  $M = (Q', \Sigma, \delta', q_0', F')$ ,  $L(M) = L(N)$ .

$E(R) = \{q \mid \text{reach } q \text{ from } R \text{ by going along } \varepsilon\}$

let  $Q' = P(Q)$ . //  $2^k$  subsets of  $Q$ ,  $k = \text{size of } Q$ .

$q_0' = E(\{q_0\})$ . // Closure  $q_0$  on transition  $\varepsilon$

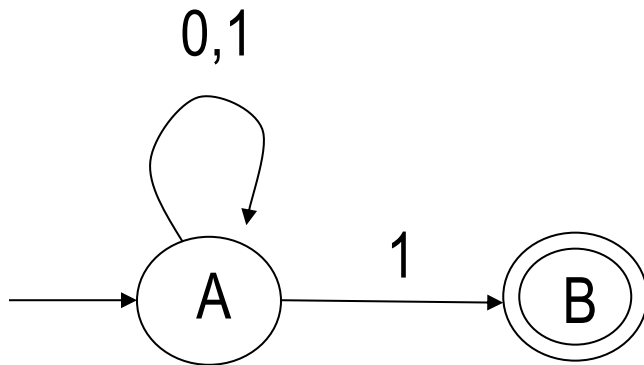
$F' = \{R \in Q' \mid R \cap F \neq \emptyset\}$ . //  $F'$  contains accept states in  $F$

$\delta'(R, a) = \bigcup_{r \in R} E(\delta(r, a))$ , for  $R \in Q'$  and  $a \in \Sigma$ . //  $\delta'$ : all states that can be reached by going along  $\varepsilon$  arrows after every step

# Practices: Conversion of NFA to DFA

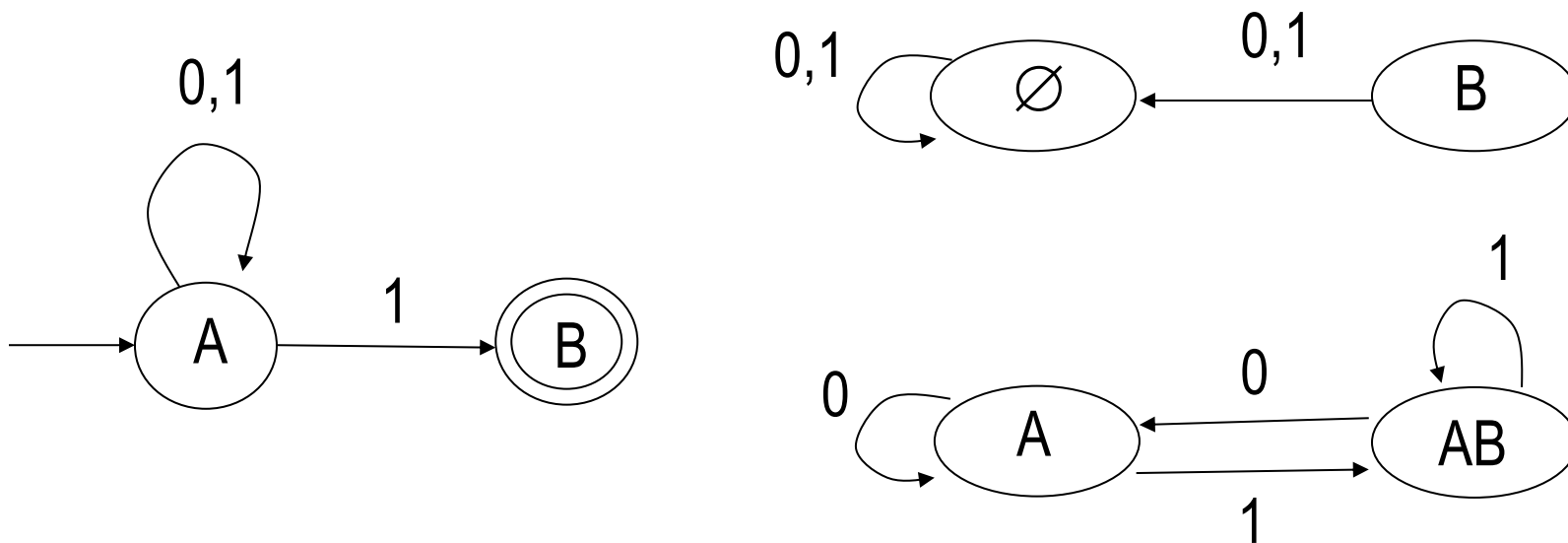
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- $L = \{\text{Set of all string that ends with '1'}\}$ ,  $\Sigma = \{0,1\}$ ;



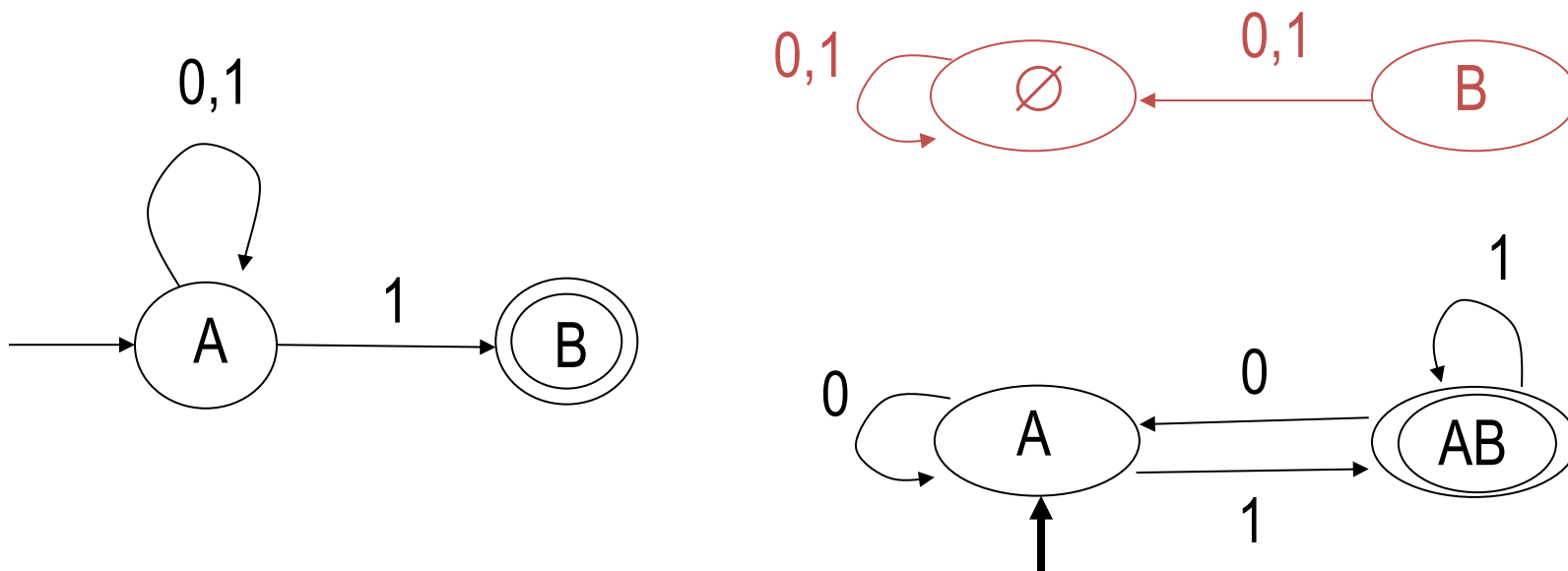
# Practices: Conversion of NFA to DFA

- $L = \{\text{Set of all string that ends with '1'}\}$ ,  $\Sigma = \{0,1\}$ ;



# Practices: Conversion of NFA to DFA

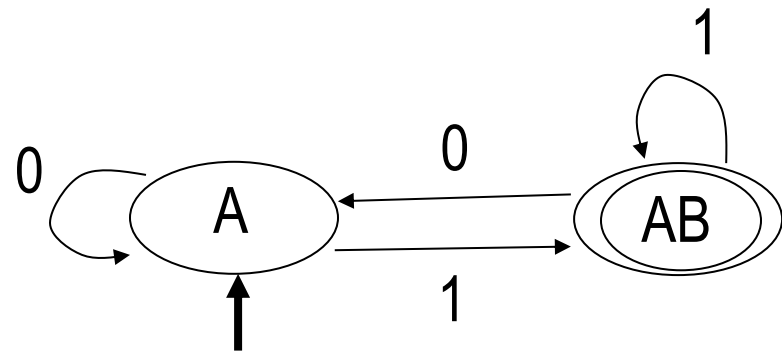
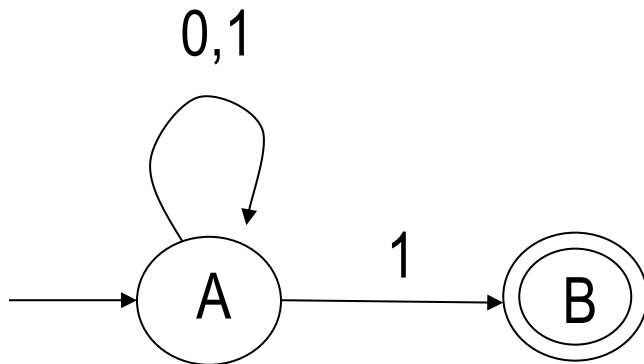
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# Practices: Conversion of NFA to DFA

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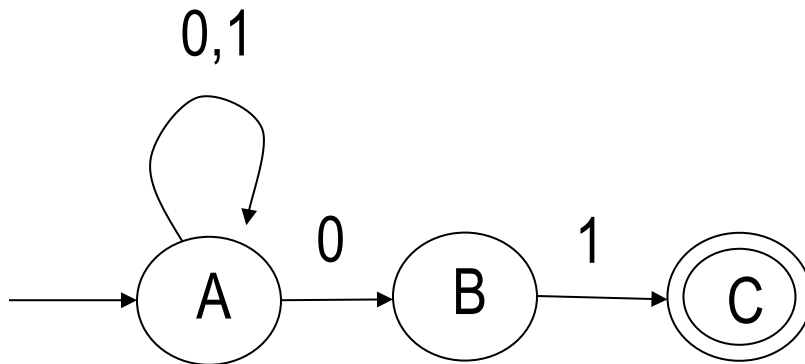
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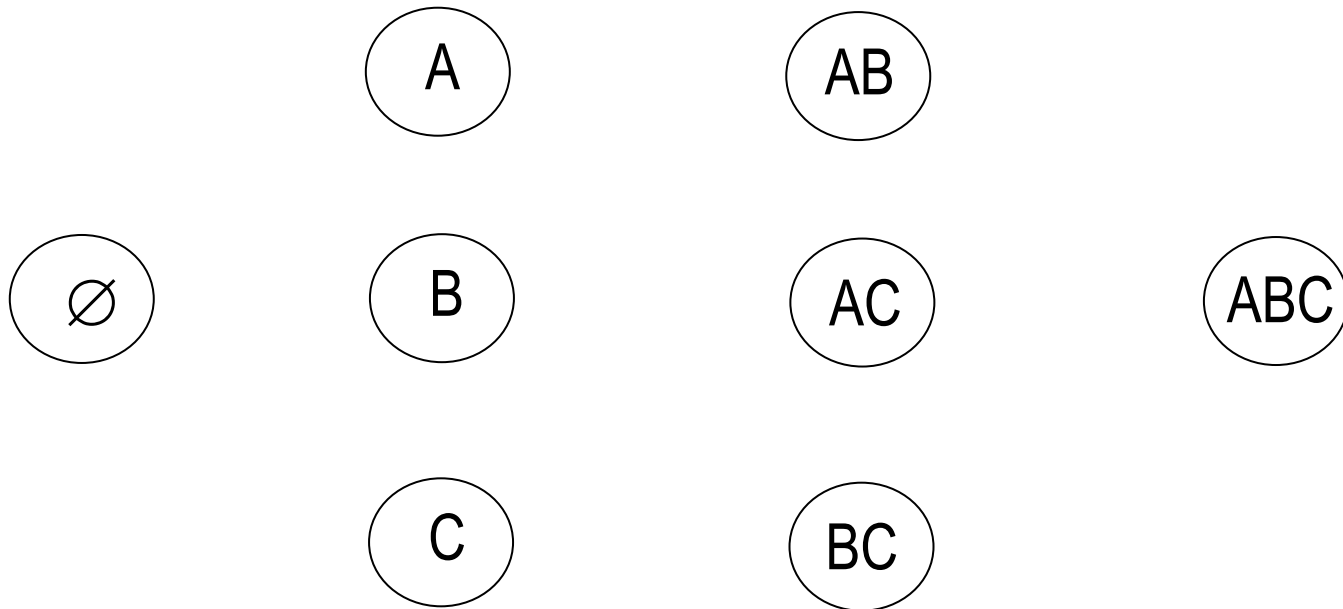
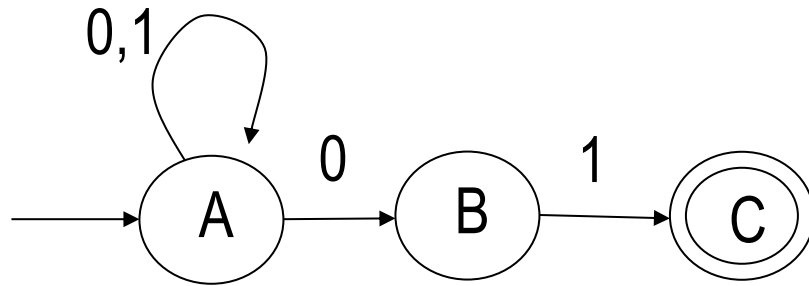
# Practices: Conversion of NFA to DFA

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- $L = \{\text{Set of all string that ends with '01'}\}$ ,  $\Sigma = \{0,1\}$ ;

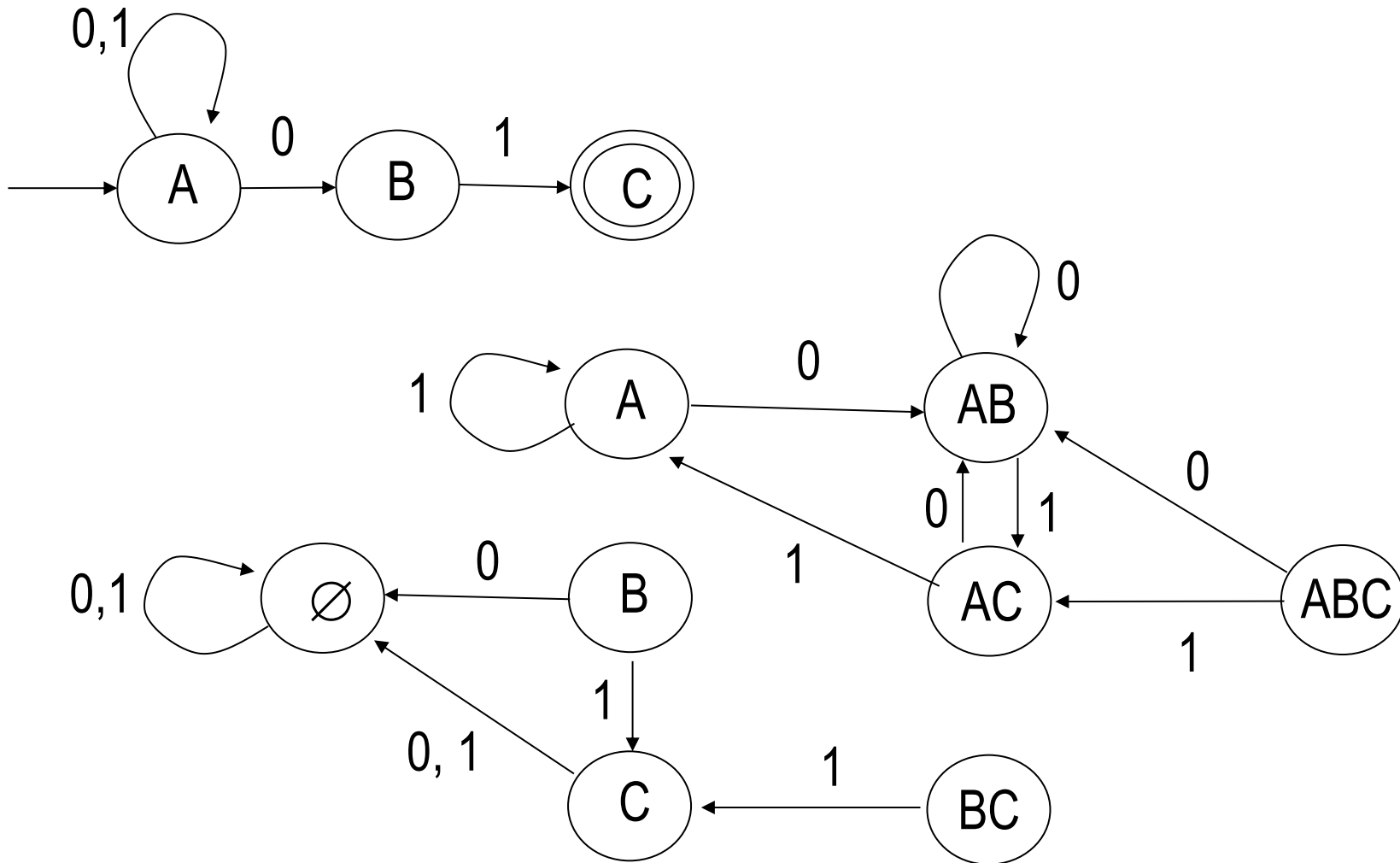


# Practices: Conversion of NFA to DFA

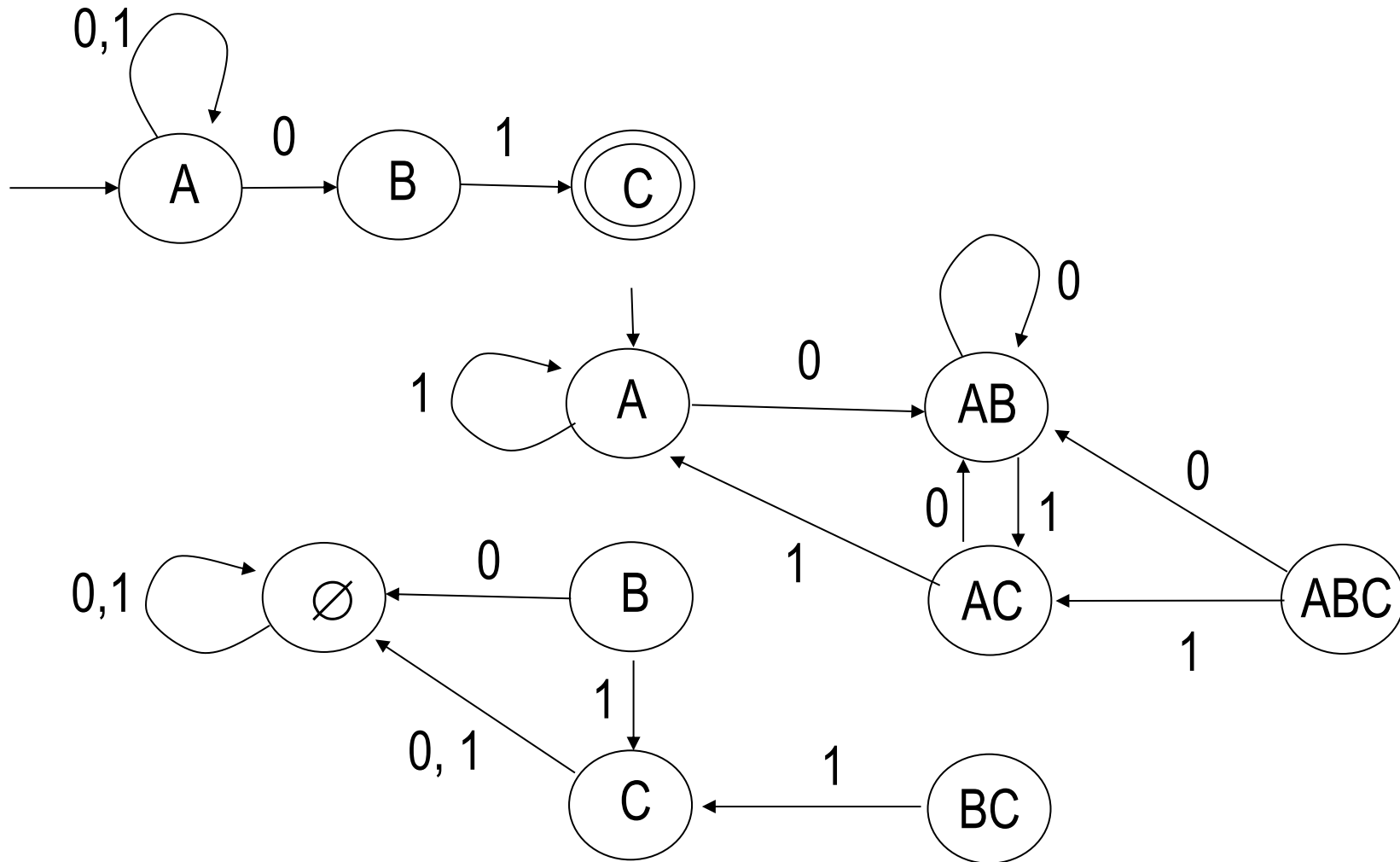




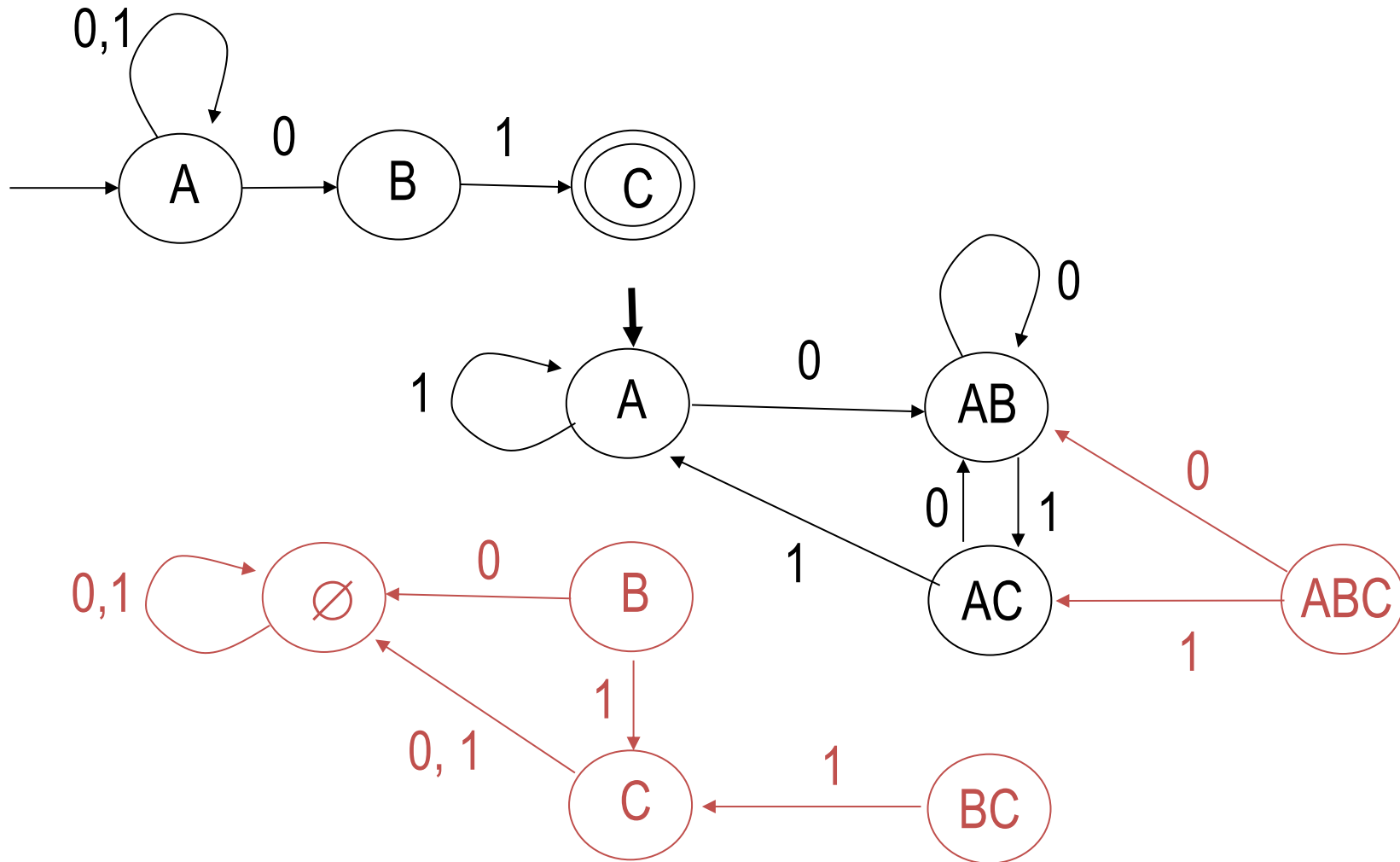
# Practices: Conversion of NFA to DFA



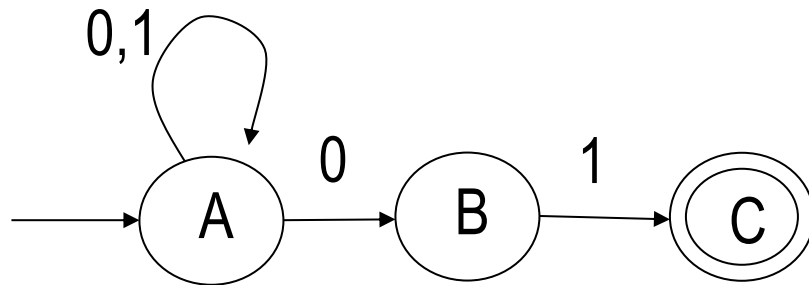
# Practices: Conversion of NFA to DFA



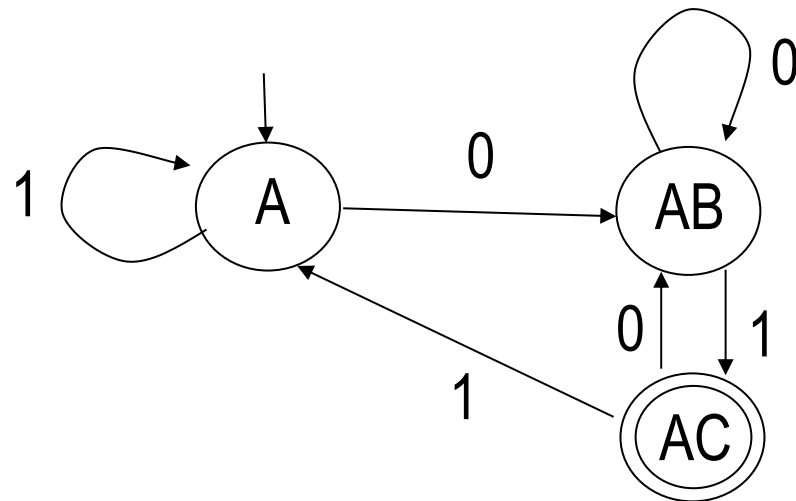
# Practices: Conversion of NFA to DFA



# Practices: Conversion of NFA to DFA



DFA



# Equivalence of NFAs and DFAs

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- Step 1: Draw all the states in DFAs
- Step 2: Define the transitions in DFAs based on the NFAs
- Step 3: Define the start state and accept state in DFAs
- Step 4: Remove all inaccessible states



# Draw NFA online

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- <http://madebyevan.com/fsm/>
  - Add a state: double-click on the canvas
  - Add an arrow: shift-drag on the canvas
  - Move something: drag it around
  - Delete something: click it and press the delete key (not the backspace key); On Laptop/Macbook, please press “Fn” + “Delete/backspace”.
  - Make accept state: double-click on an existing state
  - Add start state: shift-drag on the canvas to one state
  - Type numeric subscript: put an underscore before the number (like "S\_0")
  - Type greek letter: put a backslash before it (like "\beta")

