CS 3502 Operating Systems

Scheduling

Kun Suo

Computer Science, Kennesaw State University

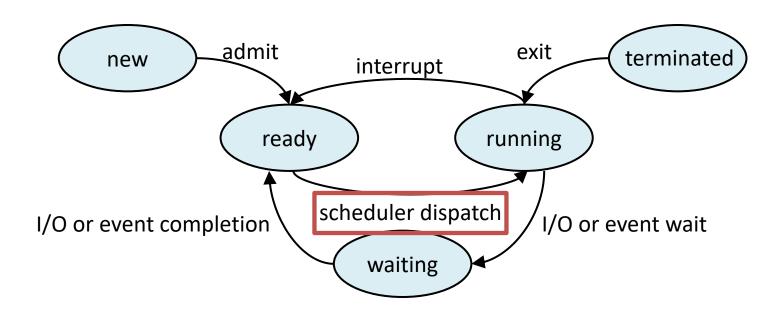
https://kevinsuo.github.io/

Outline

- Introduction to CPU scheduling
 - What is CPU scheduling
 - Why we need CPU scheduling
 - When scheduling happens
- Scheduling policies
 - FCFS, SJF, RR, Priority
 - Scheduling on multiple CPUs

What is CPU scheduling?

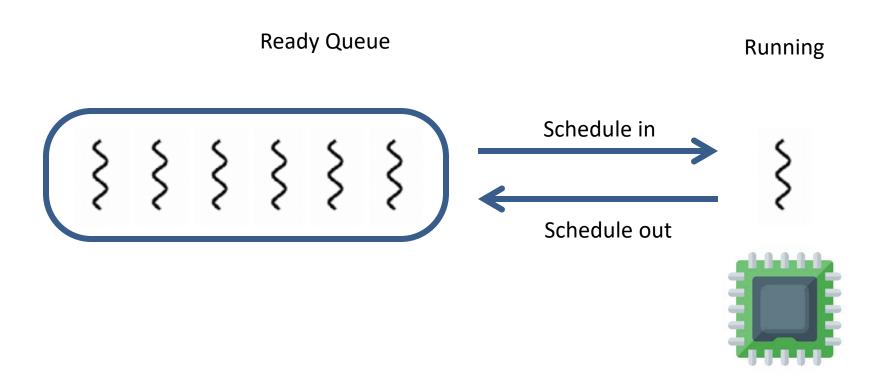
The five-state process model



CPU scheduling

Selects from among the processes/threads that are ready to execute, and allocates the CPU to it

What is CPU scheduling?

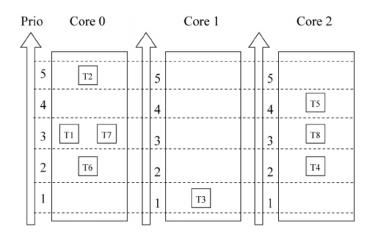


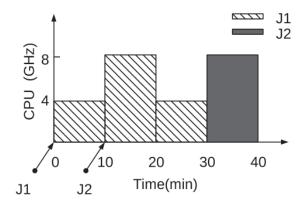
Why CPU scheduling?

- Terminal B Terminal C

 CPU
- In support of multiprogramming
 - uniprocessor systems
 - Time-sharing processor
 - multiprocessor systems
 - Efficiently distributing tasks
 - Real-time systems
 - Reliably guaranteeing deadlines





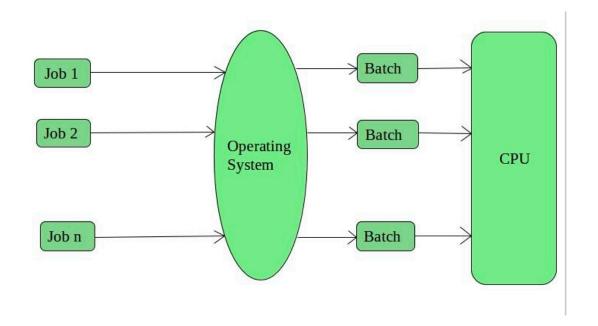


All systems

- Fairness giving each process a fair share of the CPU
- Policy enforcement seeing that stated policy is carried out
- Balance keeping all parts of the system busy

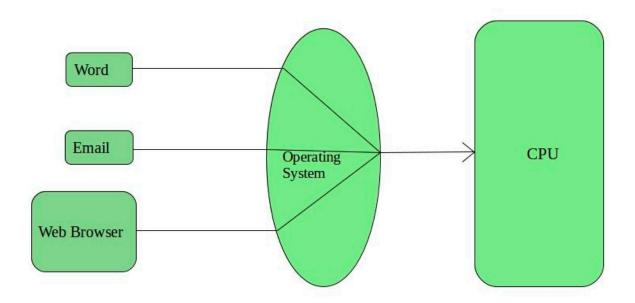
Batch systems

- Throughput maximize jobs per hour
- Turnaround time minimize time between submission and termination
- CPU utilization keep the CPU busy all the time



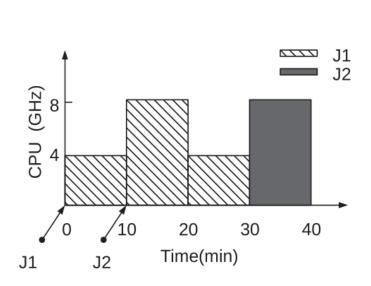
Interactive systems

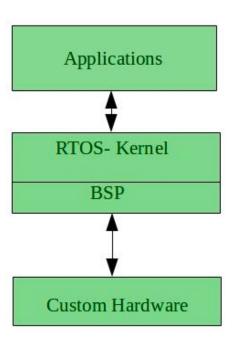
- Response time respond to requests quickly
- Proportionality meet users' expectations



Real-time systems

- Meeting deadlines avoid losing data
- Predictability avoid quality degradation in multimedia systems





All systems

- Fairness giving each process a fair share of the CPU
- Policy enforcement seeing that stated policy is carried out
- Balance keeping all parts of the system busy

Batch systems

- Throughput maximize jobs per hour
- Turnaround time minimize time between submission and termination
- CPU utilization keep the CPU busy all the time

Interactive systems

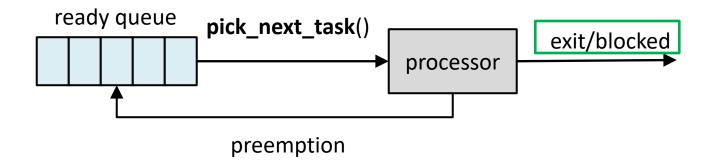
- Response time respond to requests quickly
- Proportionality meet users' expectations

Real-time systems

- Meeting deadlines avoid losing data
- Predictability avoid quality degradation in multimedia systems

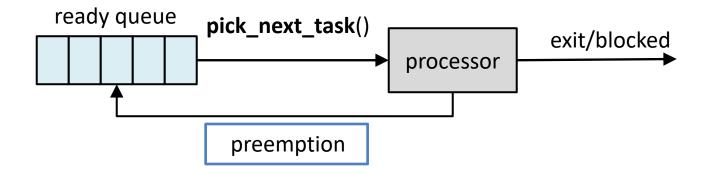
When scheduling happens?

- Non-preemptive
 - Scheduling only when current process terminates or gives up control



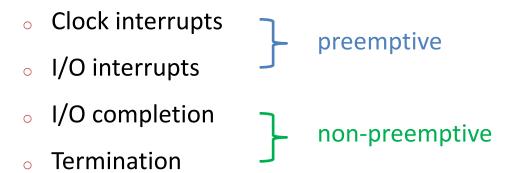
When scheduling happens?

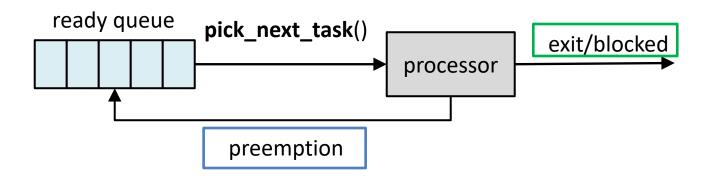
- Preemptive
 - Processes can be forced to give up control



When scheduling happens?

CPU scheduling may take place at





Outline

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Scheduling Policies

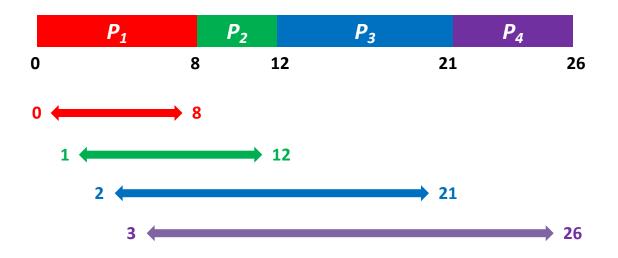
- Batch Systems
 - First-Come First-Serve (FCFS)
 - Shortest Job First
 - Shortest Remaining Time Next
- Interactive Systems
 - Round-Robin
 - Priority Scheduling
 - Multiple Queues
 - Shortest Process Next
 - Guaranteed Scheduling
 - Lottery Scheduling
- Real-time Systems
 - Rate Monotonic Scheduling
 - Earliest Deadline First Scheduling

Not exist best scheduling. It depends on your goals.



Turnaround time = End time - Arrival time

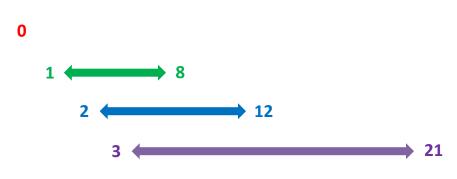
<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Response time = Start time – Arrival time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5





First-Come, First-Serve (FCFS)

CPU schedules the task that arrived earliest, non-preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average turnaround time = ((8-0)+(12-1)+(21-2)+(26-3)) / 4 = 15.25

First-Come, First-Serve (FCFS)

CPU schedules the task that arrived earliest, non-preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

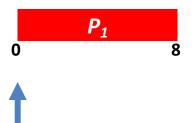


Average response time = (0+(8-1)+(12-2)+(21-3))/4 = 8.75

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

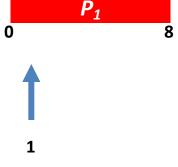


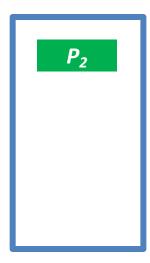


CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

pool

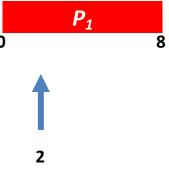




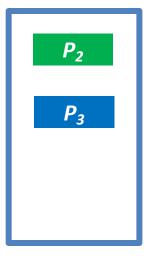
CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

non-preemptive



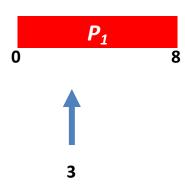
pool



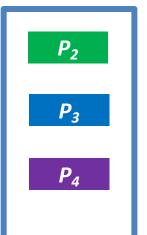
CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

non-preemptive

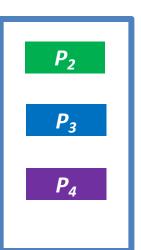


pool



CPU schedules the task with the shortest remaining time

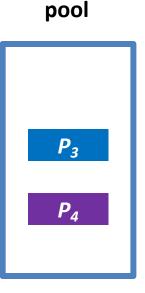
	<u>Process</u>	<u>Arrival Time</u>	Burst Time
	P_1	0	8
	P_2	1	4
	P_3	2	9
	P_4	3	5
non-preemptive	P_1		
	0	8	
			
		1	



pool

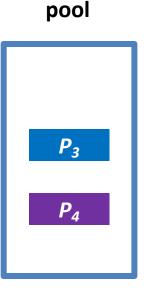
CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5
P ₁	8 12	



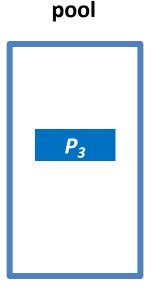
CPU schedules the task with the shortest remaining time

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P_1	0	8
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P_4	3	5
P ₁	8 12	



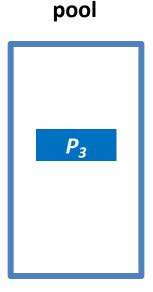
CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5
P_1	P ₂	P_4
0	8 12	17

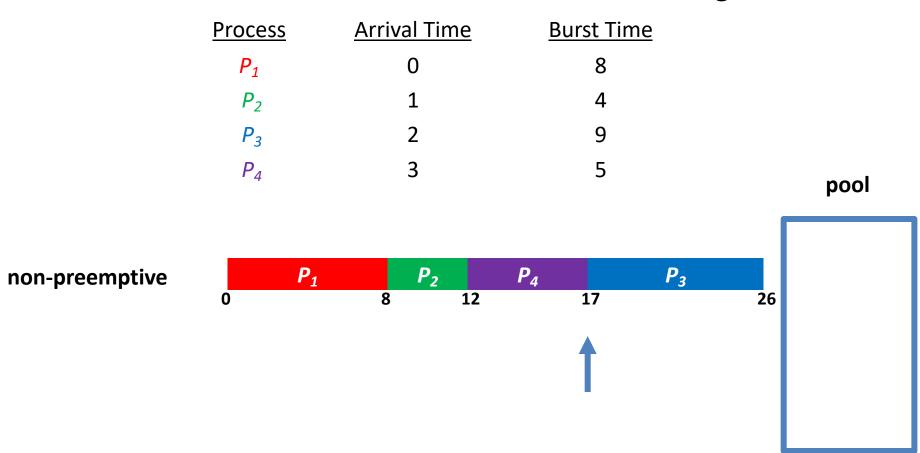


CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5
0	8 12	P ₄



CPU schedules the task with the shortest remaining time



Turnaround time = End time — Arrival time Response time = Start time — Arrival time

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

non-preemptive



Average turnaround time = ?

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5
· ·		

non-preemptive



Average turnaround time = ((8-0)+(12-1)+(26-2)+(17-3)) / 4 = 14.25

Turnaround time = End time — Arrival time Response time = Start time — Arrival time

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

non-preemptive



Average response time = ?

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
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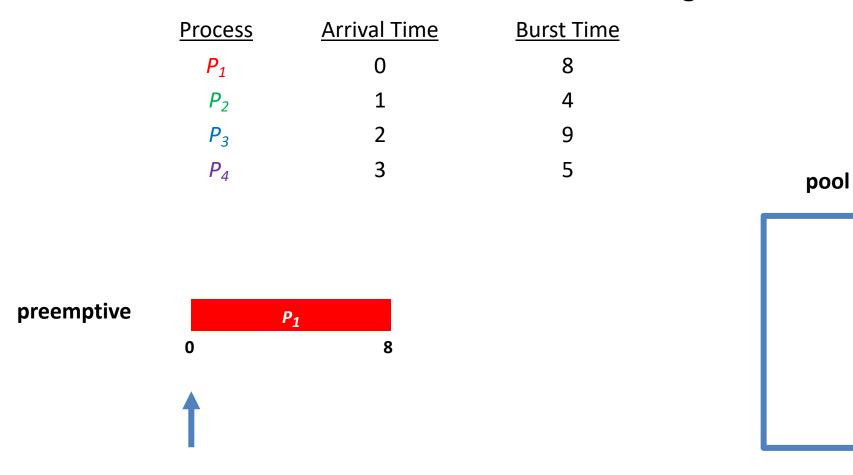
non-preemptive



Average response time = (0+(8-1)+(17-2)+(12-3))/4 = 7.75

Shortest Job First (SJF) with preemption

CPU schedules the task with the shortest remaining time



Shortest Job First (SJF) with preemption

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

preemptive P_1 0 1

Shortest Job First (SJF) with preemption

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

pool

preemptive

P₁ P₂

0 1 2

CPU schedules the task with the shortest remaining time

	<u>Process</u>	<u>Arrival Time</u>	Burst Time
	P_1	0	8
	P_2	1	4
	P_3	2	9
	P_4	3	5
preemptive	P ₁ P ₂ 0 1 2		
	1		
	P2	remain =3	



CPU schedules the task with the shortest remaining time

	<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
	P_1	0	8
	P_2	1	4
	P_3	2	9
	P_4	3	5
preemptive	P ₁ P ₂ 0 1	3	
		P2 remain =2	





$$P_4(5)$$

CPU schedules the task with the shortest remaining time

	<u>Process</u>	Arrival Time	Burst Time
	P_1	0	8
	P_2	1	4
	P_3	2	9
	P_4	3	5
preemptive	P ₁ P ₂ 0 1	5	
		†	







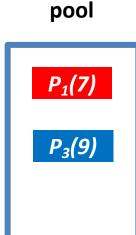
CPU schedules the task with the shortest remaining time

	<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
	P_1	0	8
	P_2	1	4
	P_3	2	9
	P_4	3	5
preemptive	P_1 P_2	P_4	
	0 1	5 10	



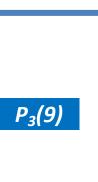
CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



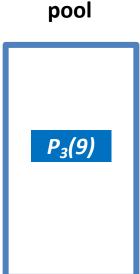
CPU schedules the task with the shortest remaining time

	<u>Process</u>		Arrival T	<u>ime</u>	<u>Bu</u>	rst Time	<u>,</u>
	P_1		0			8	
	P_2		1			4	
	P_3		2			9	
	P_4		3			5	
preemptive	P ₁ 0 1	P ₂ 5	P ₄	10	P ₁	17	



CPU schedules the task with the shortest remaining time

	<u>Process</u>	<u>Arrival Ti</u>	<u>me</u>	<u>Burst Time</u>
	P_1	0		8
	P_2	1		4
	P_3	2		9
	P_4	3		5
preemptive	P_1 P_2	P ₄	P ₁	ı
	0 1	5	10	17



CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
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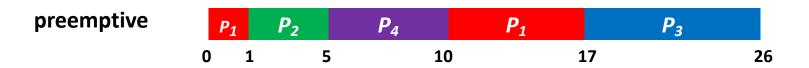


Average turnaround time = ?

Turnaround time = End time - Arrival time Response time = Start time - Arrival time

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average turnaround time = ((17-0)+(5-1)+(10-3)+(26-2))/4 = 13

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



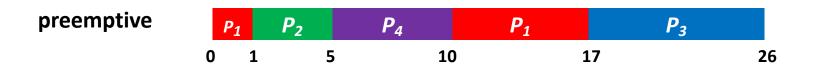


Average response time = ?

Turnaround time = End time - Arrival time Response time = Start time - Arrival time

CPU schedules the task with the shortest remaining time

<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average response time = (0+(1-1)+(5-3)+(17-2))/4 = 4.25

Process

 P_1

 P_3

 P_{4}

<u>Arrival Time</u>	Burst Time
0	8
1	4
2	9
3	5



Process

 P_1

 P_3

 P_{4}

Like FCFS, but with limited time slices, preemptive

<u>Arrival Time</u>	Burst Time
0	8
1	4
2	9
3	5



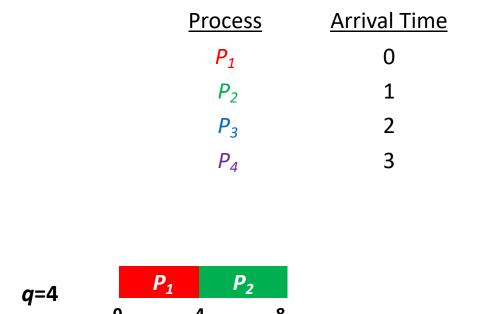
$$P_3(9)$$

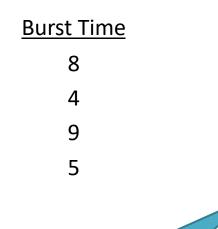
$$q=4$$
 P_1 $Q=4$

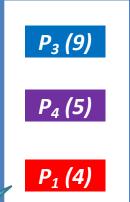


Like FCFS, but with limited time slices, preemptive

pool







P1 is put at the end of queue after scheduled out



Process

 P_1

 P_3

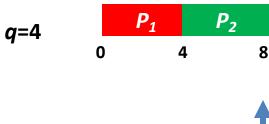
 P_{4}

Like FCFS, but with limited time slices, preemptive

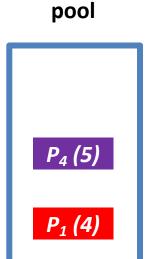
<u>Arrival Time</u>	Burst Time
0	8
1	4
2	9
3	5

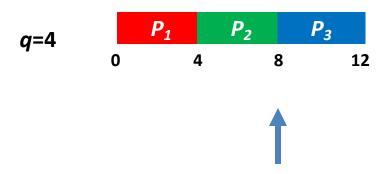




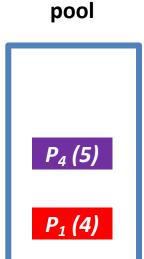


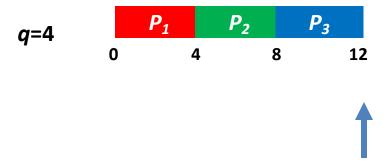
<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5





<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



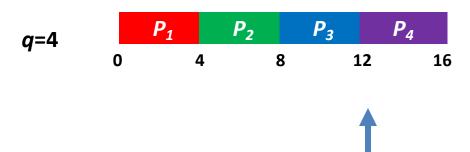


Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_{A}	3	5







Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_{A}	3	5







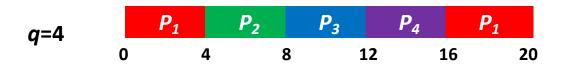


<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5





$$P_4$$
 (1)





<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



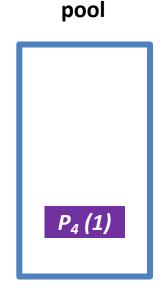


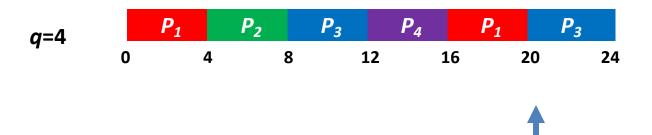
$$P_4$$
 (1)



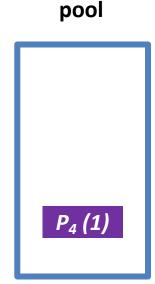


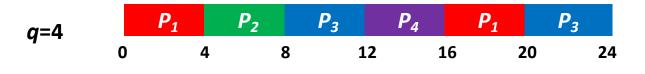
<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_{4}	3	5





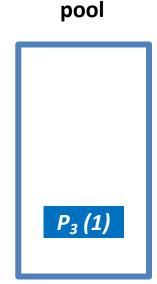
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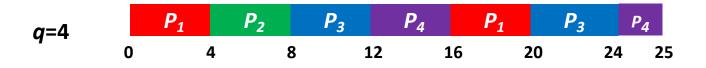






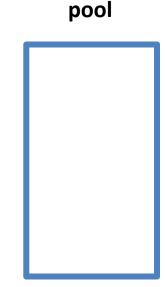
<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
P_1	0	8
P_2	1	4
P_3	2	9
P_{4}	3	5

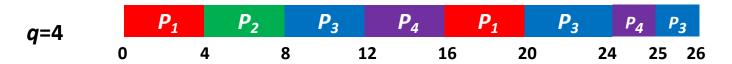






<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_{4}	3	5

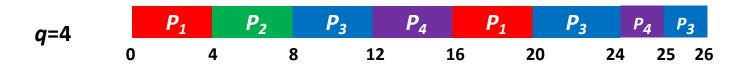






Like FCFS, but with limited time slices, preemptive

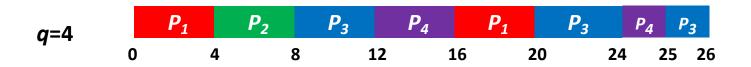
<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average turnaround time = ?

Like FCFS, but with limited time slices, preemptive

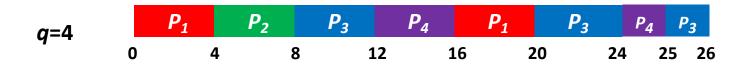
<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average turnaround time = ((20-0)+(8-1)+(26-2)+(25-3)) / 4 = 18.25

Like FCFS, but with limited time slices, preemptive

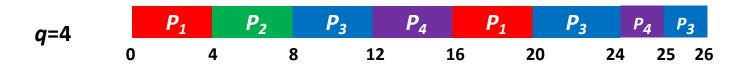
<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average response time = ?

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average response time = (0+(4-1)+(8-2)+(12-3))/4 = 4.5

Process

 P_1

 P_3

 P_{4}

<u>Arrival Time</u>	Burst Time
0	8
1	4
2	9
3	5



Select the task at the beginning

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

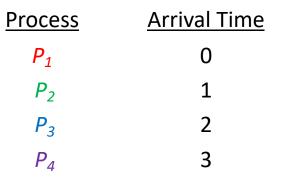


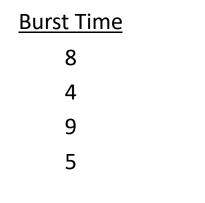


$$P_4(5)$$

Like FCFS, but with limited time slices, preemptive

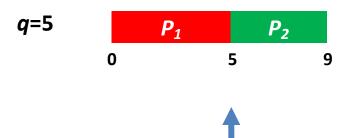
pool











P1 is put at the end of queue after scheduled out

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



$$P_4$$
 (5)

$$P_1(3)$$



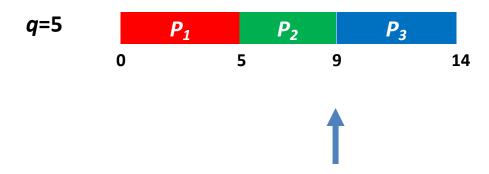


<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5





$$P_1(3)$$



<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5





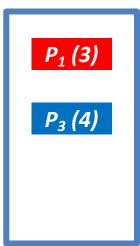
$$P_1$$
 (3)

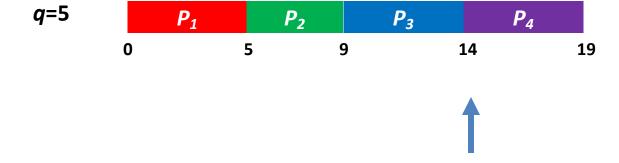




Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5

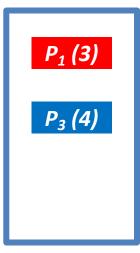




Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_{A}	3	5

pool

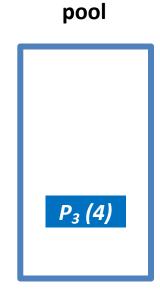


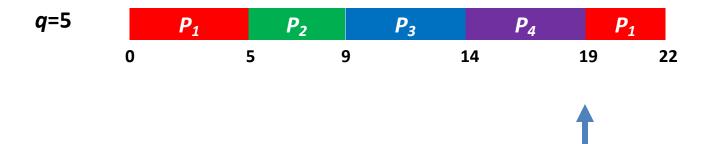




Like FCFS, but with limited time slices, preemptive

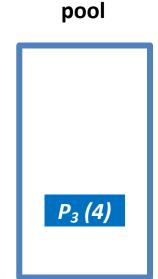
<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5





Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5







• Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



pool





Turnaround time = End time - Arrival time Response time = Start time - Arrival time

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average turnaround time = ?

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average turnaround time = ((22-0)+(9-1)+(26-2)+(19-2)) / 4 = 17.5

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average response time = ?

Like FCFS, but with limited time slices, preemptive

<u>Process</u>	<u>Arrival Time</u>	Burst Time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



Average response time = (0+(5-1)+(9-2)+(14-3))/4 = 5.5

Priority Scheduling

CPU schedules the highest priority (smaller value) first, FCFS within the same priority

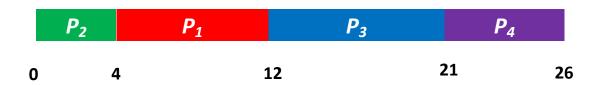
<u>Process</u>	<u>Priority</u>	Burst Time	
P_1	3	8	
P_2	1	4	Suppose arriving
P_3	4	9	time is same
P_4	2	5	



Priority Scheduling

CPU schedules the highest priority (smaller value) first, FCFS within the same priority

<u>Process</u>	<u>Priority</u>	Burst Time	
P_1	2	8	
P_2	1	4	Suppose arriving time is same
P_3	4	9	time is same
P_4	4	5	



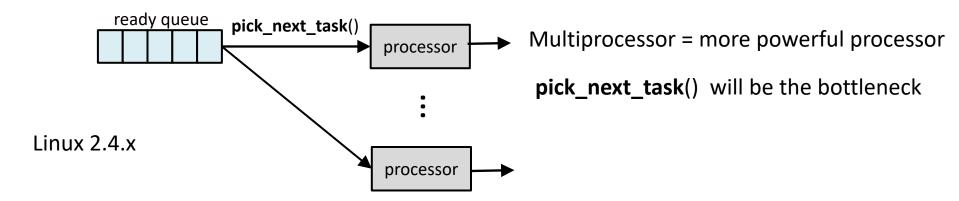
Comparison

	Turnaround time Response time	
FCFS	15.25	8.75
SJF-preemptive	13	4.25
RR (q=5)	17.5	5.5
Priority scheduling	N/A	N/A

	Throughput	Response time	Starvation
FCFS	TBD	TBD	No
SJF-preemptive	High	Good	Yes
RR	Can be low	Good	No
Priority scheduling	Can be high	Can be good	Can remove

Challenges on Emerging Hardware and Applications

Multi-processor → Single queue



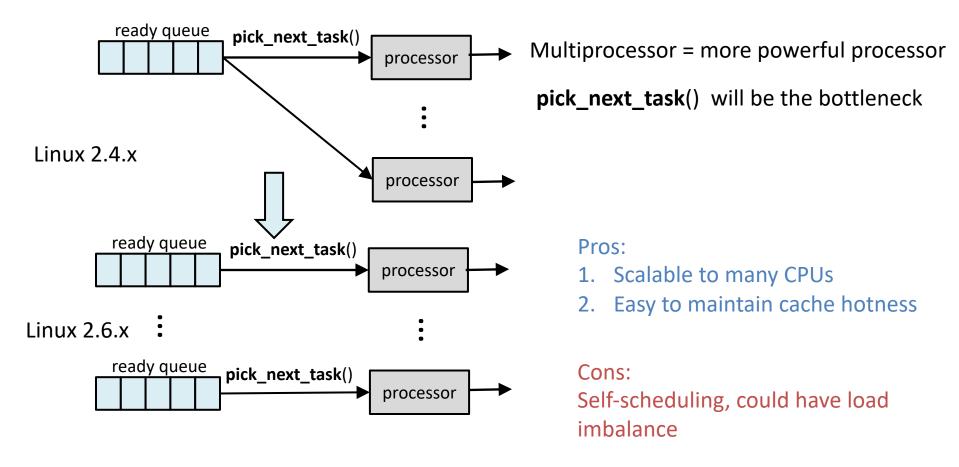
Pros:

- 1. Easy to implement
- 2. Perfect load balancing

Cons:

- 1. Scalability issues due to centralized synchronization
- 2. High overhead and low efficiency
- Hard to maintain cache hotness due to global scheduling

Challenges on Emerging Hardware and Applications

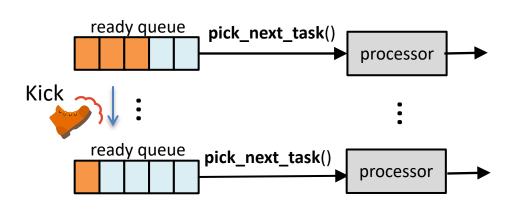


Overcome Load Imbalance

Push model

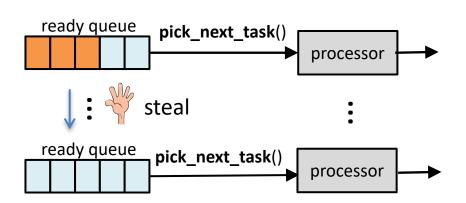
Every a while, a kernel thread checks load imbalance and move threads

Made by OS



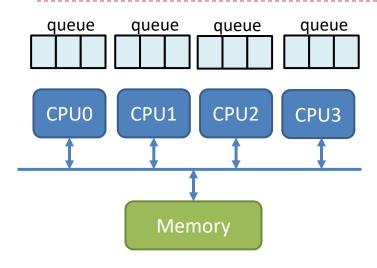
Pull model

Whenever a queue becomes empty, steal a thread from non-empty queues

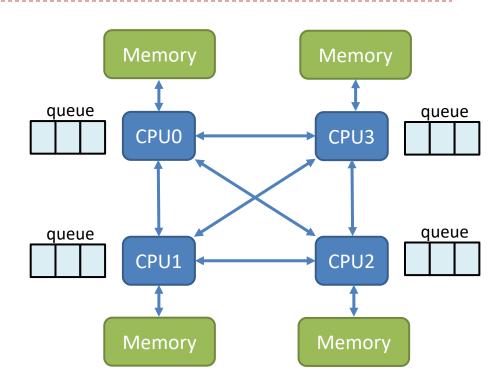


Made by local queue. Both are widely used

Load balance on SMP vs. NUMA



symmetric multiprocessing (SMP): The distance to memory is the same



Non-uniform memory access (NUMA): The distance to memory is different

Multi processor/core scheduling

The scheduling policy not only considers the fairness, throughput, etc., but also needs to consider the **hardware architecture** (e.g., locality)

Conclusion

- Introduction to CPU scheduling
 - What is CPU scheduling
 - Why we need CPU scheduling
 - When scheduling happens
- Scheduling policies
 - FCFS, SJF, RR, Priority
 - Scheduling on multiple CPUs