CS 6041 Theory of Computation

Context-free language

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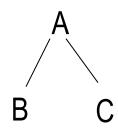
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https://kevinsuo.github.io/

Chomsky normal form (CNF)

- CNF: only allow CFG in the following forms
 - \circ S $\rightarrow \epsilon$
 - \circ A \rightarrow BC
 - A→a

Parse tree for CNF is only binary tree



- Here
 - A,B,C are variables
 - B,C are not start variables
 (start variable does not exist on the right side)
 - a is terminal

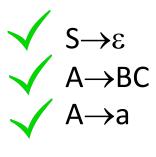
Chomsky normal form (CNF)

CNF: only allow CFG in the following forms

Only start variable S \circ S $\rightarrow \epsilon$ can generate ε $A \rightarrow BC$

Variables can only generate:

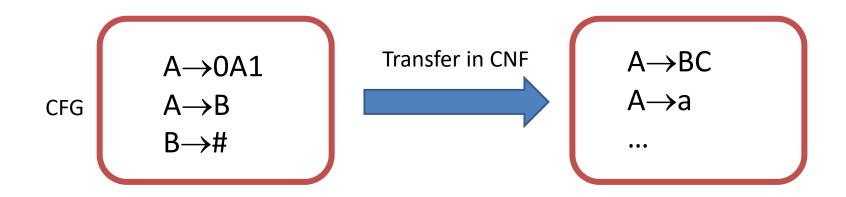
- 1, two variables
- 2, single terminal





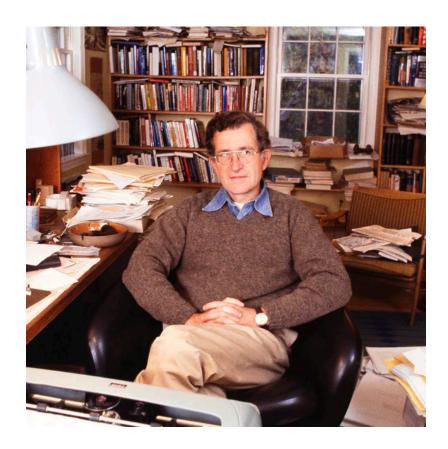
Chomsky normal form (CNF)

 Theorem: Any CFL is generated by a CFG in CNF



Noam Chomsky

http://linguistics.mit.edu/user/chomsky/



Chomsky normal form example

G₆:
$$S \rightarrow ASA \mid aB$$
,
 $A \rightarrow B \mid S$
 $B \rightarrow b \mid \varepsilon$

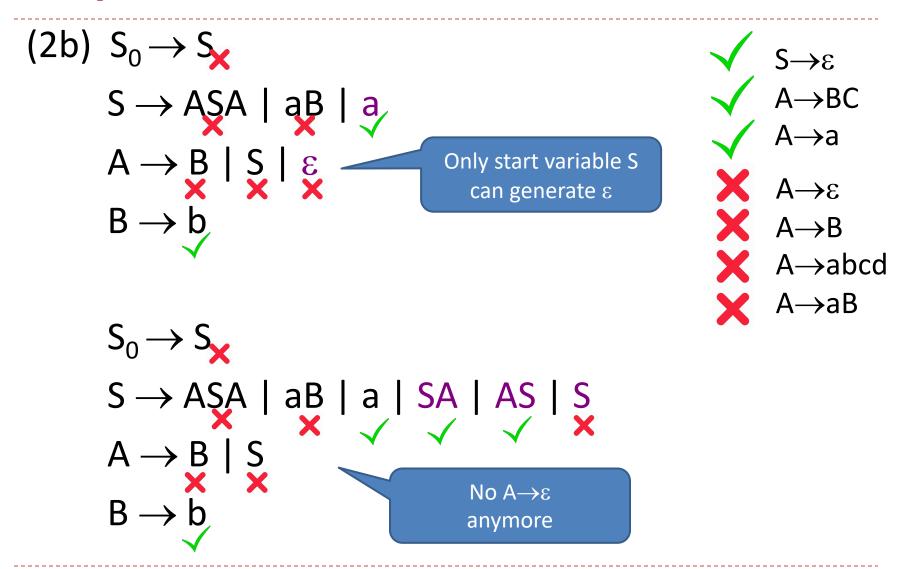
Get the CNF for G₆



G₆:
$$S \rightarrow ASA \mid aB$$
,
 $A \rightarrow B \mid S$
 $B \rightarrow b \mid \varepsilon$
 $A \rightarrow B$
 $A \rightarrow B$
 $A \rightarrow B$
 $A \rightarrow B$
 $A \rightarrow aB$
 $A \rightarrow aB$

(2a)
$$S_0 \rightarrow S_X$$

 $S \rightarrow ASA \mid aB$
 $A \rightarrow B \mid S$
 $B \rightarrow b \mid \varepsilon$
Conly start variable S
can generate ε
 $S \rightarrow ASA \mid aB \mid a$
 $A \rightarrow B \mid S \mid \varepsilon$
 $S \rightarrow ASA \mid aB \mid a$
 $A \rightarrow B \mid S \mid \varepsilon$
 $S \rightarrow b \mid S \mid \varepsilon$
 $S \rightarrow b \mid S \mid \varepsilon$



(3a)
$$S_0 \rightarrow S_{\bullet}$$

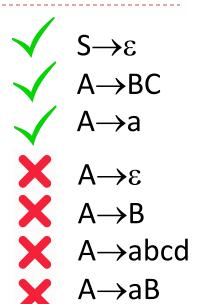
 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS \mid S_{\bullet}$
 $A \rightarrow B \mid S_{\bullet}$
 $B \rightarrow b$
 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow B \mid S_{\bullet}$
 $A \rightarrow B \mid S_{\bullet}$
 $A \rightarrow B \mid S_{\bullet}$

(3b)
$$S_0 \rightarrow S_*$$

 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow B \mid S_*$
 $B \rightarrow b$

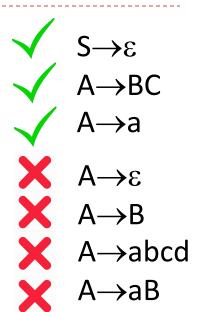
$$S_0 \rightarrow ASA \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow B \mid S$
 $B \rightarrow b$



(3c)
$$S_0 \rightarrow ASA \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow B \mid S$
 $B \rightarrow b$
 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$

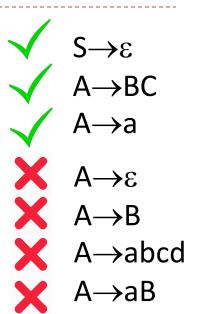


(3d)
$$S_0 \rightarrow ASA \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow b \mid S$
 $B \rightarrow b$

$$S_0 \rightarrow ASA \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow b \mid ASA \mid aB \mid a \mid SA \mid AS$
 $B \rightarrow b$

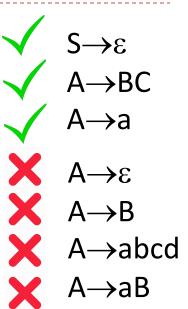


(4)
$$S_0 \rightarrow ASA \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow ASA \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow b \mid ASA \mid aB \mid a \mid SA \mid AS$
 $B \rightarrow b$

$$S_0 \rightarrow AA_1 \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow AA_1 \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow b \mid AA_1 \mid aB \mid a \mid SA \mid AS$
 $B \rightarrow b$
 $A \rightarrow SA$

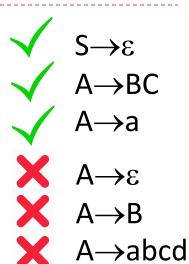


(5)
$$S_0 \rightarrow AA_1 \mid aB \mid a \mid SA \mid AS$$

 $S \rightarrow AA_1 \mid aB \mid a \mid SA \mid AS$
 $A \rightarrow b \mid AA_1 \mid aB \mid a \mid SA \mid AS$
 $B \rightarrow b$
 $A_1 \rightarrow SA$

$$S_0 \rightarrow AA_1 \mid UB \mid a \mid SA \mid AS$$

 $S \rightarrow AA_1 \mid UB \mid a \mid SA \mid AS$
 $A \rightarrow b \mid AA_1 \mid UB \mid a \mid SA \mid AS$
 $B \rightarrow b$
 $A_1 \rightarrow SA$
 $A \rightarrow AA_2 \mid AA_3 \mid AA_4 \mid AA_5 \mid AA$



(5)
$$S_0 \rightarrow AA_1 \mid UB \mid a \mid$$

$$SA \mid AS$$

$$S \rightarrow AA_1 \mid UB \mid a \mid$$

$$SA \mid AS$$

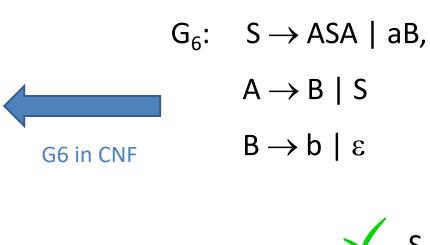
$$A \rightarrow b \mid AA_1 \mid UB \mid a \mid$$

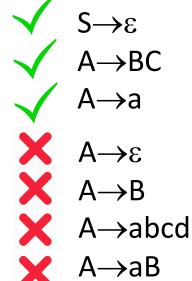
$$SA \mid AS$$

$$B \rightarrow b$$

$$A_1 \rightarrow SA$$

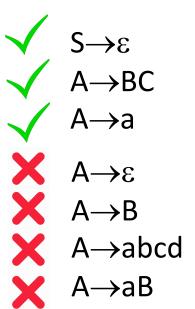
$$U \rightarrow a$$





Conclusion for CNF

- Add new start variable if needed
- A $\rightarrow \epsilon$, merge above rules with A
- A→B, replace B with terminals or other rules
- A \rightarrow aB, replace with U \rightarrow a, A \rightarrow UB
- A \rightarrow abcd, replace with A \rightarrow aU₁, U₁ \rightarrow bU₂, U₂ \rightarrow cd
- A→BCD, similar as the above

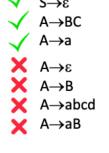


$$G_7$$
: $S \rightarrow AbA$, $A \rightarrow Aa$,

Get the CNF for G₇

Conclusion for CNF

- Add new start variable if needed
- A \rightarrow ϵ , merge above rules with A
- A→B, replace B with terminals or other rules
- A \rightarrow aB, replace with U \rightarrow a, A \rightarrow UB
- A→abcd, replace with A→aU₁,
 U₁→bU₂, U₂→cd
- A→BCD, similar as the above



CS 604

Kennesaw State University

Theory of Computation

$$S \rightarrow \varepsilon$$
 $A \rightarrow BC$
 $A \rightarrow a$
 $A \rightarrow \varepsilon$
 $A \rightarrow B$
 $A \rightarrow B$
 $A \rightarrow a$
 $A \rightarrow a$

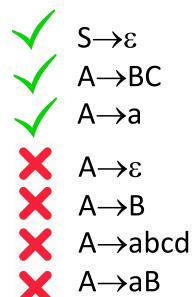
$$S \rightarrow AbA$$
,

$$A \rightarrow Aa$$
,

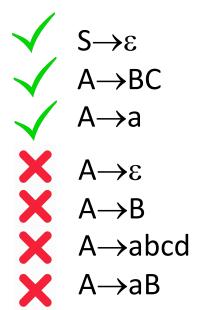
$$A \rightarrow \epsilon$$

$$S \rightarrow AbA \mid bA \mid Ab \mid b$$

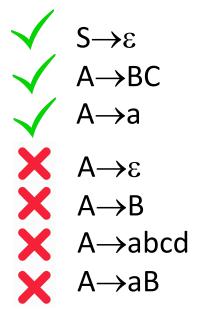
 $A \rightarrow Aa \mid a$



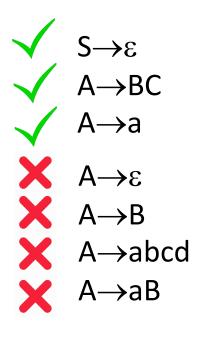
$$S \rightarrow AbA \mid bA \mid Ab \mid b$$
 $A \rightarrow Aa \mid a$
 $S \rightarrow TA \mid bA \mid Ab \mid b$
 $A \rightarrow Aa \mid a$
 $T \rightarrow Ab$



$$S \rightarrow TA/|bA|Ab|b/$$
 $A \rightarrow Aa|a/$
 $T \rightarrow Ab$
 $S \rightarrow TA/|BA/|AB|b/$
 $A \rightarrow Aa|a/$
 $T \rightarrow AB/$



$$A \rightarrow AA \mid AB \mid BA$$
 $A \rightarrow AB \mid AB \mid BA$
 $A \rightarrow AB \mid AB \mid BA$
 $A \rightarrow AC \mid AB \mid BA \mid AB \mid BA$
 $A \rightarrow AC \mid AB \mid BA \mid AB \mid BA$



 $C \rightarrow a$

G₁₀:
$$S \rightarrow 1A \mid 0B$$

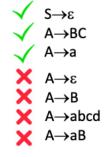
$$A \rightarrow 1AA \mid 0S \mid 0$$

$$B \rightarrow 0BB \mid 1$$

$$X$$
Get the CNF for G_{10}

Conclusion for CNF

- Add new start variable if needed
- A \rightarrow ϵ , merge above rules with A
- A→B, replace B with terminals or other rules
- A \rightarrow aB, replace with U \rightarrow a, A \rightarrow UB
- A \rightarrow abcd, replace with A \rightarrow aU₁, U₁ \rightarrow bU₂, U₂ \rightarrow cd
- A→BCD, similar as the above



CS 6041

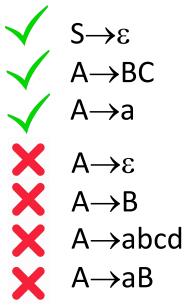
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Theory of Computation

$$\begin{array}{c} \checkmark & S \rightarrow \varepsilon \\ \checkmark & A \rightarrow BC \\ \checkmark & A \rightarrow a \end{array}$$

$$A \rightarrow \varepsilon$$
 $A \rightarrow B$
 $A \rightarrow abcd$
 $A \rightarrow aB$

G₁₀:
$$S \rightarrow 1A \mid 0B$$
 $X \quad X$
 $A \rightarrow 1AA \mid 0S \mid 0$
 $X \quad X$
 $B \rightarrow 0BB \mid 1$
 $X \quad X$
 $A \rightarrow C_1A \mid 0B$
 $A \rightarrow C_1AA \mid 0S \mid 0$
 $A \rightarrow C_1AA \mid 0S \mid 0$
 $A \rightarrow 0BB \mid 1$
 $A \rightarrow 0BB \mid 1$

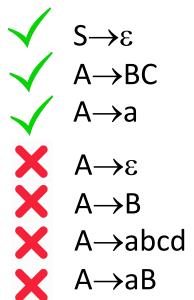


$$S \rightarrow C_1A \mid 0B$$

 $A \rightarrow C_1AA \mid 0S \mid 0$
 $B \rightarrow 0BB \mid 1$
 $C_1 \rightarrow 1$

$$S \rightarrow C_1A \mid C_0B \mid$$

 $A \rightarrow C_1AA \mid OS \mid O \mid$
 $B \rightarrow OBB \mid 1 \mid$
 $C_1 \rightarrow 1 \mid$
 $C_0 \rightarrow 0 \mid$



$$S \rightarrow C_{1}A \mid C_{0}B$$

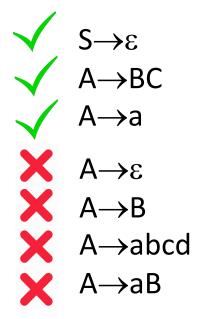
$$A \rightarrow C_{1}AA \mid 0S \mid 0$$

$$B \rightarrow 0BB \mid 1$$

$$C_{1} \rightarrow 1$$

$$C_{0} \rightarrow 0$$

$$S \rightarrow C_1A \mid C_0B \mid C_1C_2 \mid OS \mid C_1 \rightarrow C_1 \rightarrow C_2 \rightarrow AA$$



$$S \rightarrow C_{1}A \mid C_{0}B \mid C_{1}A \mid C_{2}B \mid C_{2}B$$

$$S \rightarrow C_1A \mid C_0B$$

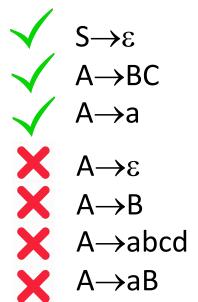
$$A \rightarrow C_1C_2 \mid C_0S \mid 0$$

$$B \rightarrow 0BB \mid 1$$

$$C_1 \rightarrow 1$$

$$C_0 \rightarrow 0$$

$$C_2 \rightarrow AA$$



$$S \rightarrow C_{1}A \mid C_{0}B$$

$$A \rightarrow C_{1}C_{2} \mid C_{0}S \mid 0$$

$$B \rightarrow 0BB \mid 1$$

$$C_{1} \rightarrow 1$$

$$C_{0} \rightarrow 0$$

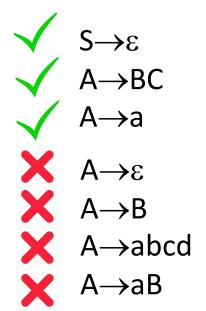
$$C_{2} \rightarrow AA$$

$$S \rightarrow C_{1}A \mid C_{0}B$$

$$A \rightarrow C_{1}C_{2} \mid C_{0}S \mid 0$$

$$B \rightarrow C_{0}C_{3} \mid 1$$

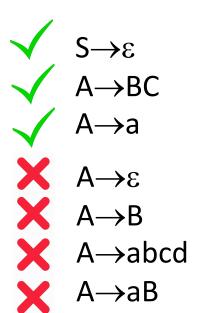
$$C_{1} \rightarrow 1$$



 $C_3 \rightarrow BB$

Conclusion for CNF

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- A \rightarrow abcd, replace with A \rightarrow aU₁, U₁ \rightarrow bU₂, U₂ \rightarrow cd
- A→BCD, similar as the above



Context-sensitive vs. Context-free

- What is the difference?
 - S->aSb
 - S->ab

- aSb->aaSbb
- S->ab

Context-sensitive vs. Context-free

- Context-free
 - S->aSb
 - S->ab

- Context-sensitive
 - aSb->aaSbb
 - S->ab

The rule S->aSb happens only S has a on the left and b on the right

