CHAPTER 2

ARRAYS AND STRUCTURES

All the programs in this file are selected from

Ellis Horowitz, Sartaj Sahni, and Susan Anderson-Freed "Fundamentals of Data Structures in C /2nd Edition", Silicon Press, 2008.

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Arrays

Array: a set of index and value

data structure

For each , there is a associated with that index.

representation (possible)

implemented by using consecutive memory.

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Structure *Array* is

objects: A set of pairs *<index*, *value>* where for each value of *index* there is a value from the set *item*. *Index* is a finite ordered set of one or more dimensions, for example, $\{0, ..., n-1\}$ for one dimension, $\{(0,0),(0,1),(0,2),(1,0),(1,1),(1,2),(2,0),(2,1),(2,2)\}$ for two dimensions, etc.

Functions:

for all $A \in Array$, $i \in index$, $x \in item$, j, $size \in integer$

Array Create(j, list) ::= **return** an array of *j* dimensions where list is a j-tuple whose *i*th element is the size of the *i*th dimension. *Items* are undefined.

Item Retrieve(A, i) ::= if $(i \in index)$ return the item associated with index value i in array A

else return error

 $Array\ Store(A, i, x) ::= \mathbf{if}\ (i\ in\ index)$

return an array that is identical to array A except the new pair $\langle i, x \rangle$ has been inserted **else return** error

end array

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Arrays in C

implementation of 1-D array

```
\begin{array}{ll} list[0] & base \ address = \alpha \\ list[1] & \\ list[2] & \\ list[3] & \\ list[4] & \end{array}
```

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Arrays in C (Continued)

Compare int *list1 and int list2[5] in C.

Same: list1 and list2 are

Difference: list2 reserves

Notations:

```
list2 - a pointer to list2[0]
(list2 + i) - a pointer to list2[i] =
*(list2 + i) =
```

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Example: 1-dimension array addressing

```
int one[] = {0, 1, 2, 3, 4};
Goal: print out address and value
```

```
void print1(int *ptr, int rows)
{
    /* print out a one-dimensional array using a pointer */
        int i;
        printf("Address Contents\n");
        for (i=0; i < rows; i++)
             printf("%8u%5d\n", ptr+i, *(ptr+i));
        printf("\n");
}</pre>
```

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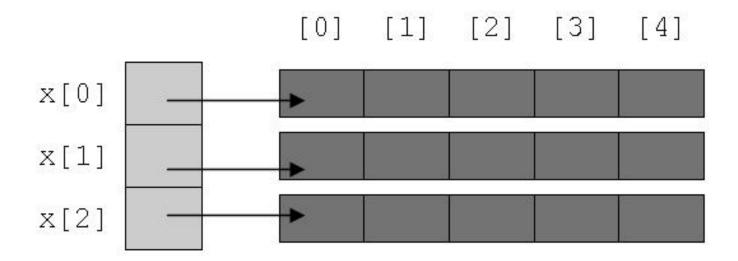
call print1(&one[0], 5)

Address	Contents
1228	0
	1
	2
	3
	4

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Two-dimension array

Int x[3][5];



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Structures (records)

```
char name[10];
    int age;
    float salary;
    } person;

strcpy(person.name, "james");
person.age=10;
person.salary=35000;
```

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Create structure data type

```
typedef struct human_being {
        char name[10];
        int age;
        float salary;
        };
or
typedef struct {
        char name[10];
        int age;
        float salary
        } human_being;
human_being person1, person2;
```

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If (person1 == person2)?

```
int humansEqual (humanBeing person1, humanBeing person2)
  if (strcmp(person1.name, person2.name))
      return FALSE;
  if (person1.age != person2.age)
      return FALSE;
  if (person1.salary != person2.salary)
      return FALSE;
  return TRUE;
                      if (humansEqual(person1, person2))
```

A structure within a structure

```
typedef struct {
        int month;
        int day;
        int year;
        } date;
typedef struct human_being {
                                  person1.dob.day = 11;
       char name[10];
                                  person1.dob.year = 1944;
       int age;
       float salary;
       date dob;
        };
```

Unions

```
Similar to struct, but only one field is active.
Example: Add fields for male and female.
typedef struct sex_type {
       enum tag_field {female, male} sex;
       union {
               int children;
                                 human_being person1, person2;
               int beard;
                                 person1.sex_info.sex=male;
               } u;
                                 person1.sex_info.u.beard=FALSE;
typedef struct human_being {
       char name[10];
       int age;
       float salary;
       date dob;
       sex_type sex_info;
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```

Self-Referential Structures

One or more of its components is a pointer to itself.

```
typedef struct list {
    char data;
    list *link;
    }
```

Construct a list with three nodes

malloc: obtain a node

```
list item1, item2, item3;
item1.data='a';
item2.data='b';
item3.data='c';
item1.link=item2.link=item3.link=NULL;
```

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Ordered List Examples

ordered (linear) list: (item1, item2, item3, ..., item*n*)

- (MONDAY, TUEDSAY, WEDNESDAY, THURSDAY, FRIDAY, SATURDAYY, SUNDAY)
- (2, 3, 4, 5, 6, 7, 8, 9, 10, Jack, Queen, King, Ace)
- **(1941, 1942, 1943, 1944, 1945)**
- (a₁, a₂, a₃, ..., a_{n-1}, a_n)

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Polynomials $A(X)=3X^{20}+2X^5+4$, $B(X)=X^4+10X^3+3X^2+1$

Structure *Polynomial* is

objects: $p(x) = a_1 x^{e_1} + ... + a_n x^{e_n}$; a set of ordered pairs of $\langle e_i, a_i \rangle$ where a_i in *Coefficients* and e_i in *Exponents*, e_i are integers >= 0 functions:

for all poly, poly1, poly2 \in Polynomial, $coef \in Coefficients$, $expon \in Exponents$

Polynomial Zero() ::= **return** the polynomial, p(x) = 0

Boolean IsZero(poly) ::= if (poly) return FALSE else return TRUE

Coefficient Coef(poly, expon) ::= **if** (expon \in poly) **return** its coefficient **else return** Zero

Exponent Lead_Exp(poly) ::= **return** the largest exponent in poly

Polynomial Attach(poly,coef, expon) ::= if (expon \in poly) return error else return the polynomial poly with the term < coef, expon> inserted

Polynomial Remove(poly, expon)::= if $(expon \in poly)$ return the polynomial poly with the term whose exponent is $expon \ deleted$ else return errorPolynomial SingleMult(poly, coef, expon)::= return the polynomial $poly \bullet coef \bullet x^{expon}$ Polynomial Add(poly1, poly2)::= return the polynomial poly1 + poly2Polynomial Mult(poly1, poly2)::= return the polynomial $poly1 \bullet poly2$

End Polynomial

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Polynomial Addition

```
data structure 1:
                    #define MAX_DEGREE 101
                    typedef struct {
                             int degree;
                             float coef[MAX_DEGREE];
                             } polynomial;
/* d = a + b, where a, b, and d are polynomials */
d = Zero()
while (! IsZero(a) &&! IsZero(b)) do {
  switch COMPARE (Lead_Exp(a), Lead_Exp(b)) {
     case -1: d =
       Attach(d, Coef (b, Lead_Exp(b)), Lead_Exp(b));
       b = Remove(b, Lead\_Exp(b));
       break;
    case 0: sum = Coef (a, Lead Exp(a)) + Coef (b, Lead Exp(b));
      if (sum) {
         Attach (d, sum, Lead_Exp(a));
         a = Remove(a, Lead\_Exp(a));
         b = Remove(b, Lead\_Exp(b));
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       break;
```

```
case 1: d =
    Attach(d, Coef (a, Lead_Exp(a)), Lead_Exp(a));
    a = Remove(a, Lead_Exp(a));
}
insert any remaining terms of a or b into d
```

***Program 2.4 :**Initial version of *padd* function(p.62)

Data structure 2: use one global array to store all polynomials

$$A(X)=2X^{1000}+1$$

$$B(X)=X^4+10X^3+3X^2+1 \qquad *Figure 2.2: Array representation of two polynomials}$$

$$starta \ finisha \ startb \qquad \qquad finishb \ avail$$

$$coef \qquad 2 \qquad 1 \qquad 1 \qquad 10 \qquad 3 \qquad 1$$

$$exp \qquad 1000 \qquad 0 \qquad 4 \qquad 3 \qquad 2 \qquad 0$$

0

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```
storage requirements: start, finish, 2*(finish-start+1)
 nonparse:
               twice as much as (1)
                when all the items are nonzero
 MAX_TERMS 100 /* size of terms array */
 typedef struct {
         float coef;
         int expon;
         } polynomial;
 polynomial terms[MAX_TERMS];
 int avail = 0;
*(p.62)
```

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Add two polynomials: D = A + B

```
void padd (int startA, int finishA, int startB, int finishB,
                                 int * startD, int *finishD)
/* add A(x) and B(x) to obtain D(x) */
  float coefficient;
  *startD = avail;
  while (startA <= finishA && startB <= finishB)
   switch (COMPARE(terms[startA].expon,
                        terms[startB].expon)) {
   case -1: /* A expon < B expon */
         attach(terms[startB].coef, terms[startB].expon);
         startB++
         break;
```

```
case 0: /* equal exponents */
           coefficient = terms[startA].coef +
                         terms[startB].coef;
           if (coefficient)
             attach (coefficient, terms[startA].expon);
           startA++;
           startB++;
           break;
case 1: /* A expon > B expon */
       attach(terms[startA].coef, terms[startA].expon);
       startA++;
```

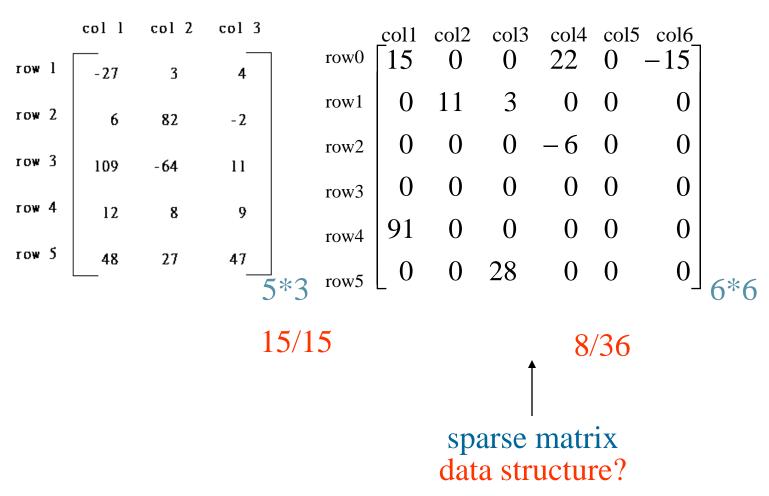
```
/* add in remaining terms of A(x) */
for(; startA <= finishA; startA++)
   attach(terms[startA].coef, terms[startA].expon);
/* add in remaining terms of B(x) */
for(; startB <= finishB; startB++)
  attach(terms[startB].coef, terms[startB].expon);
*finishD =avail -1;
Analysis:
              O(n+m)
               where n (m) is the number of nonzeros in A(B).
```

```
void attach(float coefficient, int exponent)
/* add a new term to the polynomial */
  if (avail >= MAX_TERMS) {
    fprintf(stderr, "Too many terms in the polynomial\n");
    exit(1);
   terms[avail].coef = coefficient;
   terms[avail++].expon = exponent;
```

Problem: Compaction is required when polynomials that are no longer needed. (data movement takes time.)

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Sparse Matrix



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SPARSE MATRIX ABSTRACT DATA TYPE

Structure *Sparse_Matrix* is

objects: a set of triples, <*row*, *column*, *value*>, where *row* and *column* are integers and form a unique combination, and *value* comes from the set *item*.

functions:

for all $a, b \in Sparse_Matrix, x \in item, i, j, max_col, max_row \in index$

Sparse_Marix Create(max_row, max_col) ::=

return a *Sparse_matrix* that can hold up to $max_items = max_row \times max_col$ and whose maximum row size is max_row and whose maximum column size is max_col .

 $Sparse_Matrix Transpose(a) ::=$

return the matrix produced by interchanging the row and column value of every triple.

 $Sparse_Matrix \ Add(a, b) ::=$

if the dimensions of a and b are the same **return** the matrix produced by adding corresponding items, namely those with identical *row* and *column* values.

else return error

Sparse_Matrix **Multiply**(*a*, *b*) ::=

if number of columns in a equals number of rows in **b**

return the matrix d produced by multiplying a by b according to the formula: $d[i][j] = \Sigma(a[i][k] \cdot b[k][j])$ where d(i, j) is the (i, j)th element

else return error.

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- (1) Represented by a two-dimensional array. Sparse matrix wastes space.
- (2) Each element is characterized by <row, col, value>.

_	rov	v co	l value	_	row	col	<u>value</u>
		_ #	of rows (c	· · · · · · · · · · · · · · · · · · ·			
a[0]	6	6	8	# of nonzero term b[0]	6	6	8
[1]	0	0	15	[1]	0	0	15
[2]	0	3	22	[2]	0	4	91
[3]	0	5	-15	[3]	1	1	11
[4]	1	1	11 - t	$\frac{\text{ranspose}}{4}$	2	1	3
[5]	1	2	3	[5]	2	5	28
[6]	2	3	-6	[6]	3	0	22
[7]	4	0	91	[7]	3	2	-6
[8]	5	2	28	[8]	5	0	-15
row.	column	(a) in a	scendin	g order		(b)	

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```
Sparse_matrix Create(max_row, max_col) ::=

#define MAX_TERMS 101 /* maximum number of terms +1*/
    typedef struct {
        int col;
        int row;
        int value;
        } term;
    term a[MAX_TERMS]
# of rows (columns)
# of nonzero terms
```

* (P.69)

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Transpose a Matrix

(1) for each row i take element <i, j, value> and store it in element <j, i, value> of the transpose.

```
difficulty: where to put \langle j, i, value \rangle

(0, 0, 15) ====> (0, 0, 15)

(0, 3, 22) ====> (3, 0, 22)

(0, 5, -15) ====> (5, 0, -15)

(1, 1, 11) ====> (1, 1, 11)

Move elements down very often.
```

(2) For all elements in column j, place element <i, j, value> in element <j, i, value>

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Sparse Matrix Multiplication

Definition: $[D]_{m*p} = [A]_{m*n} * [B]_{n*p}$

Procedure: Fix a row of A and find all elements in column j

of B for j=0, 1, ..., p-1.

Alternative 1. Scan all of B to find all elements in j.

Alternative 2. Compute the transpose of B.

(Put all column elements consecutively)

$$\begin{bmatrix} 1 & 0 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

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$\lceil X$	0	0	0	0	0	0	0	0	$0 \rceil$
X	X	0	0	0	0	0	0	0	0
X	X	\boldsymbol{X}	0	0	0	0	0	0	0
X	X	\boldsymbol{X}	X	0	0	0	0	0	0
X	X	\boldsymbol{X}	X	\boldsymbol{X}	0	0	0	0	0
X					•				
						•			
							X	0	0
								X	0
$\lfloor X$	X	X	X	X	\boldsymbol{X}	X	X	X	$X \rfloor$

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$\lceil X$	X	X	X	X	X	X	X	\boldsymbol{X}	$X \rceil$
0	X	X	X	X	X	X	\boldsymbol{X}	X	X
0	0	X	X	X	X	X	X	\boldsymbol{X}	X
0	0	0	X	X	X	X	X	X	X
0	0	0	0	X	X	X	X	\boldsymbol{X}	X
0					•				
						•			
•							\boldsymbol{X}	X	X
								X	X
$\lfloor 0$	0	0	0	0	0	0	0	0	$X \rfloor$

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