

Assembly Language for x86 Processors

7th Edition, Global Edition

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Chapter 8: Advanced Procedures

Chapter Overview

- **Stack Frames**
- Recursion
- INVOKE, ADDR, PROC, and PROTO
- Creating Multimodule Programs

Stack Frames

- Stack Parameters
- Local Variables
- ENTER and LEAVE Instructions
- LOCAL Directive
- WriteStackFrame Procedure

Stack Frame

- Also known as an *activation record*
- Area of the stack set aside for a procedure's return address, passed parameters, saved registers, and local variables
- Created by the following steps:
 - Calling program pushes arguments on the stack and calls the procedure.
 - The called procedure pushes EBP on the stack, and sets EBP to ESP.
 - If local variables are needed, a constant is subtracted from ESP to make room on the stack.

Stack Parameters

- More convenient than register parameters
- Two possible ways of calling DumpMem. Which is easier?

```
pushad
mov esi,OFFSET array
mov ecx,LENGTHOF array
mov ebx,TYPE array
call DumpMem
popad
```

```
push TYPE array
push LENGTHOF array
push OFFSET array
call DumpMem
```

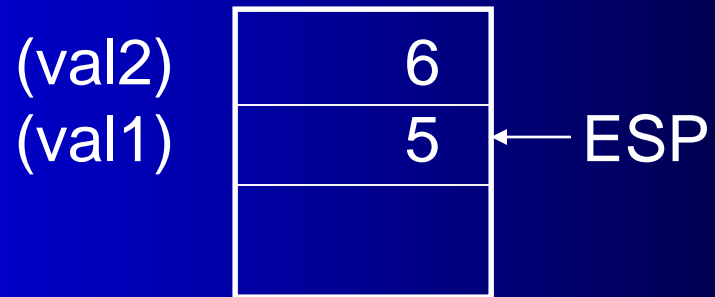
Passing Arguments by Value

- Push argument values on stack
 - (Use only 32-bit values in protected mode to keep the stack aligned)
- Call the called-procedure
- Accept a return value in EAX, if any
- Remove arguments from the stack if the called-procedure did not remove them

Example

```
.data  
val1  DWORD 5  
val2  DWORD 6
```

```
.code  
push val2  
push val1
```



Stack prior to CALL

Passing by Reference

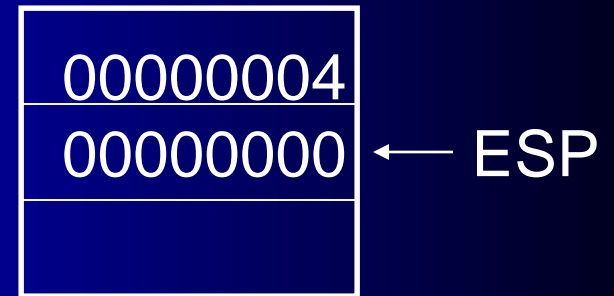
- Push the offsets of arguments on the stack
- Call the procedure
- Accept a return value in EAX, if any
- Remove arguments from the stack if the called procedure did not remove them

Example

```
.data  
val1    DWORD    5  
val2    DWORD    6
```

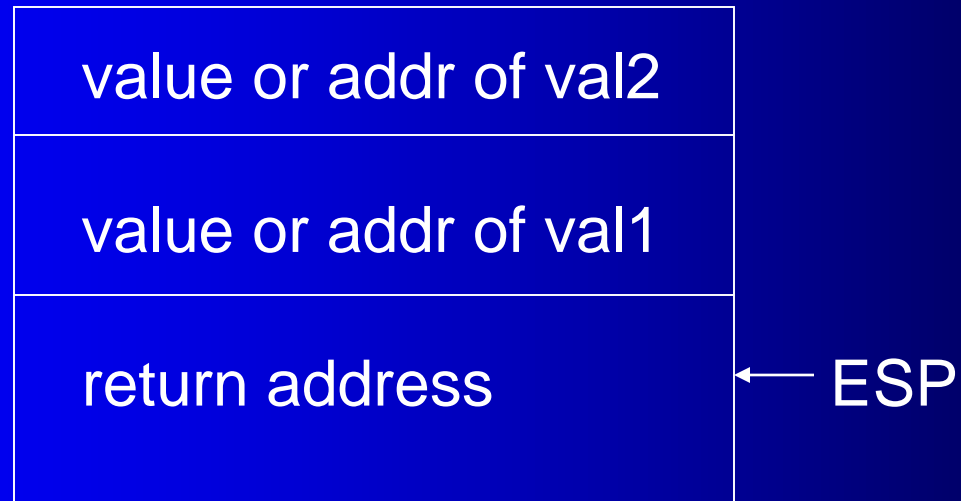
```
.code  
push OFFSET val2  
push OFFSET val1
```

(offset val2)
(offset val1)



Stack prior to CALL

Stack after the CALL



Passing an Array by Reference (1 of 2)

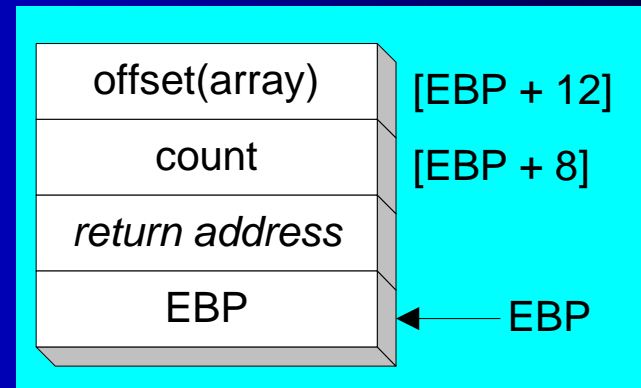
- The `ArrayFill` procedure fills an array with 16-bit random integers
- The calling program passes the address of the array, along with a count of the number of array elements:

```
.data
count = 100
array WORD count DUP(?)
.code
    push OFFSET array
    push COUNT
    call ArrayFill
```

Passing an Array by Reference (2 of 2)

ArrayFill can reference an array without knowing the array's name:

```
ArrayFill PROC
    push ebp
    mov  ebp, esp
    pushad
    mov  esi, [ebp+12]
    mov  ecx, [ebp+8]
    .
    .
```



ESI points to the beginning of the array, so it's easy to use a loop to access each array element. [View the complete program.](#)

Accessing Stack Parameters (C/C++)

- C and C++ functions access stack parameters using constant offsets from EBP¹.
 - Example: [ebp + 8]
- EBP is called the **base pointer** or **frame pointer** because it holds the base address of the stack frame.
- EBP does not change value during the function.
- EBP must be restored to its original value when a function returns.

¹ BP in Real-address mode

RET Instruction

- *Return from subroutine*
- Pops stack into the instruction pointer (EIP or IP). Control transfers to the target address.
- Syntax:
 - **RET**
 - **RET *n***
- Optional operand *n* causes *n* bytes to be added to the stack pointer after EIP (or IP) is assigned a value.

Who removes parameters from the stack?

Caller (C) or Called-procedure (STDCALL):

```
push val2
push val1
call AddTwo
add esp,8
```

```
AddTwo PROC
    push ebp
    mov  ebp,esp
    mov  eax,[ebp+12]
    add  eax,[ebp+8]

    pop  ebp
    ret  8
```

(Covered later: The MODEL directive specifies calling conventions)

Your turn . . .

- Create a procedure named Difference that subtracts the first argument from the second one. Following is a sample call:

```
push 14                ; first argument
push 30                ; second argument
call Difference        ; EAX = 16
```

```
Difference PROC
```

```
    push ebp
    mov  ebp, esp
    mov  eax, [ebp + 8]    ; second argument
    sub  eax, [ebp + 12]   ; first argument
    pop  ebp
    ret  8
```

```
Difference ENDP
```


Passing 8-bit and 16-bit Arguments

- Cannot push 8-bit values on stack
- Pushing 16-bit operand may cause page fault or ESP alignment problem
 - incompatible with Windows API functions
- Expand smaller arguments into 32-bit values, using MOVZX or MOVSX:

```
.data
charVal BYTE 'x'
.code
    movzx eax, charVal
    push  eax
    call  Uppercase
```

Passing Multiword Arguments

- Push high-order values on the stack first; work backward in memory
- Results in little-endian ordering of data
- Example:

```
.data
longVal DQ 1234567800ABCDEFh
.code
    push  DWORD PTR longVal + 4      ; high doubleword
    push  DWORD PTR longVal         ; low doubleword
    call  WriteHex64
```

Saving and Restoring Registers

- Push registers on stack just after assigning ESP to EBP
 - local registers are modified inside the procedure

```
MySub PROC
    push    ebp
    mov     ebp, esp
    push    ecx                ; save local registers
    push    edx
```

Stack Affected by USES Operator

```
MySub1 PROC USES ecx edx
    ret
MySub1 ENDP
```

- USES operator generates code to save and restore registers:

```
MySub1 PROC
    push ecx
    push edx

    pop  edx
    pop  ecx
    ret
```

Local Variables

- Only statements within subroutine can view or modify local variables
- Storage used by local variables is released when subroutine ends
- local variable name can have the same name as a local variable in another function without creating a name clash
- Essential when writing recursive procedures, as well as procedures executed by multiple execution threads

Creating LOCAL Variables

Example - create two DWORD local variables:

Say: `int x=10, y=20;`

MySub PROC

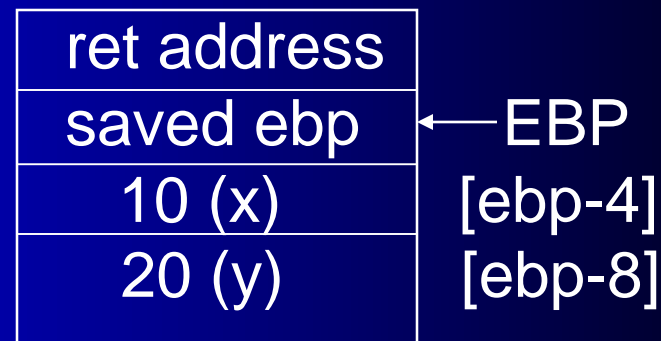
push ebp

mov ebp, esp

sub esp, 8 ;create 2 DWORD variables

mov DWORD PTR [ebp-4], 10 ; initialize x=10

mov DWORD PTR [ebp-8], 20 ; initialize y=20



LEA Instruction

- LEA returns offsets of direct and indirect operands
 - OFFSET operator only returns constant offsets
- LEA required when obtaining offsets of stack parameters & local variables
- Example

```
Copy String PROC,
```

```
    count:DWORD
```

```
    LOCAL temp[20]:BYTE
```

```
    mov edi,OFFSET count           ; invalid operand
```

```
    mov esi,OFFSET temp           ; invalid operand
```

```
    lea edi,count                 ; ok
```

```
    lea esi,temp                 ; ok
```

LEA Example

Suppose you have a Local variable at [ebp-8]

And you need the address of that local variable in ESI

You cannot use this:

```
mov esi, OFFSET [ebp-8] ; error
```

Use this instead:

```
lea esi, [ebp-8]
```


ENTER Instruction

- ENTER instruction creates stack frame for a called procedure
 - pushes EBP on the stack
 - sets EBP to the base of the stack frame
 - reserves space for local variables
 - Example:

```
MySub PROC  
    enter 8,0
```

- Equivalent to:

```
MySub PROC  
    push ebp  
    mov ebp,esp  
    sub esp,8
```

LEAVE Instruction

Terminates the stack frame for a procedure.

Equivalent operations

MySub PROC

enter 8,0

...

...

...

leave

ret

MySub ENDP

push ebp
mov ebp,esp
sub esp,8 ; 2 local DWORDs

mov esp,ebp ; free local space
pop ebp

LOCAL Directive

- The LOCAL directive declares a list of local variables
 - immediately follows the PROC directive
 - each variable is assigned a type
- Syntax:

`LOCAL varlist`

Example:

`MySub PROC`

`LOCAL var1:BYTE, var2:WORD, var3:SDWORD`

Using LOCAL

Examples:

```
LOCAL flagVals[20]:BYTE    ; array of bytes
```

```
LOCAL pArray:PTR WORD      ; pointer to an array
```

```
myProc PROC,                ; procedure  
    LOCAL t1:BYTE,          ; local variables
```

LOCAL Example (1 of 2)

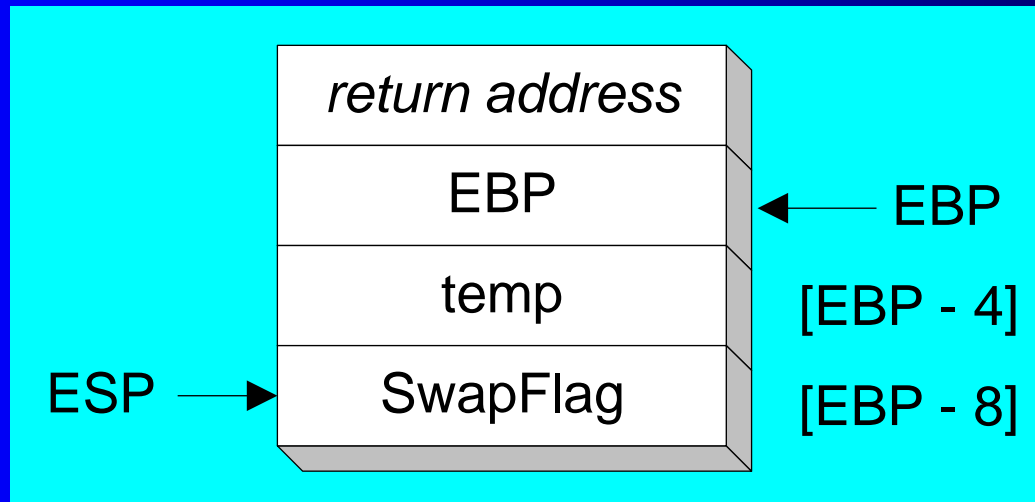
```
BubbleSort PROC
    LOCAL temp:DWORD, SwapFlag:BYTE
    . . .
    ret
BubbleSort ENDP
```

MASM generates the following code:

```
BubbleSort PROC
    push ebp
    mov  ebp,esp
    add  esp,0FFFFFFF8h        ; add -8 to ESP
    . . .
    mov  esp,ebp
    pop  ebp
    ret
BubbleSort ENDP
```

LOCAL Example (2 of 2)

Diagram of the stack frame for the BubbleSort procedure:



Non-Doubleword Local Variables

- Local variables can be different sizes
- How created in the stack by LOCAL directive:
 - 8-bit: assigned to next available byte
 - 16-bit: assigned to next even (word) boundary
 - 32-bit: assigned to next doubleword boundary

Local Byte Variable

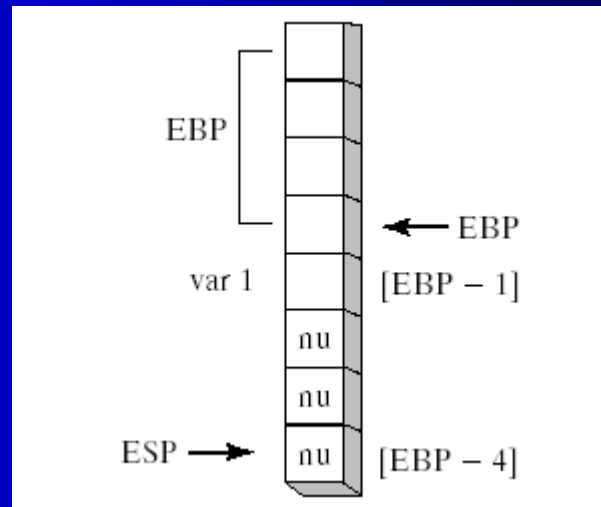
```
Example1 PROC
```

```
    LOCAL var1:BYTE
```

```
    mov al,var1                ; [EBP - 1]
```

```
    ret
```

```
Example1 ENDP
```



WriteStackFrame Procedure

- Displays contents of current stack frame
 - Prototype:

```
WriteStackFrame PROTO,  
    numParam:DWORD,      ; number of passed parameters  
    numLocalVal: DWORD,  ; number of DWordLocal variables  
    numSavedReg: DWORD   ; number of saved registers
```

WriteStackFrame Example

```
main PROC
    mov eax, 0EAEAEAEAh
    mov ebx, 0EBEBEBEBh
    INVOKE aProc, 1111h, 2222h
    exit
main ENDP

aProc PROC USES eax ebx,
    x: DWORD, y: DWORD
    LOCAL a:DWORD, b:DWORD
    PARAMS = 2
    LOCALS = 2
    SAVED_REGS = 2
    mov a, 0AAAAh
    mov b, 0BBBBh
    INVOKE WriteStackFrame, PARAMS, LOCALS, SAVED_REGS
```

The Microsoft x64 Calling Convention

- CALL subtracts 8 from RSP
- First four parameters are placed in RCX, RDX, R8, and R9. Additional parameters are pushed on the stack.
- Parameters less than 64 bits long are not zero extended
- Return value in RAX if ≤ 64 bits
- Caller must allocate at least 32 bytes of shadow space so the subroutine can copy parameter values

The Microsoft x64 Calling Convention

- Caller must align RSP to 16-byte boundary
- Caller must remove all parameters from the stack after the call
- Return value larger than 64 bits must be placed on the runtime stack, with RCX pointing to it
- RBX, RBP, RDI, RSI, R12, R14, R14, and R15 registers are preserved by the subroutine; all others are not.

What's Next

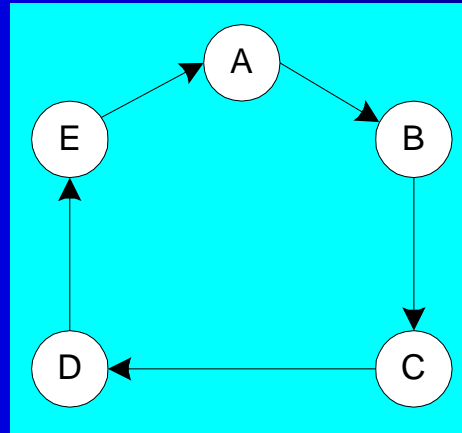
- Stack Frames
- **Recursion**
- INVOKE, ADDR, PROC, and PROTO
- Creating Multimodule Programs

Recursion

- What is Recursion?
- Recursively Calculating a Sum
- Calculating a Factorial

What is Recursion?

- The process created when . . .
 - A procedure calls itself
 - Procedure A calls procedure B, which in turn calls procedure A
- Using a graph in which each node is a procedure and each edge is a procedure call, recursion forms a **cycle**:



Recursively Calculating a Sum

The CalcSum procedure recursively calculates the sum of an array of integers. Receives: ECX = count. Returns: EAX = sum

```
CalcSum PROC
    cmp ecx,0           ; check counter value
    jz L2               ; quit if zero
    add eax,ecx         ; otherwise, add to sum
    dec ecx             ; decrement counter
    call CalcSum        ; recursive call
L2: ret
CalcSum ENDP
```

Stack frame:

Pushed On Stack	ECX	EAX
L1	5	0
L2	4	5
L2	3	9
L2	2	12
L2	1	14
L2	0	15

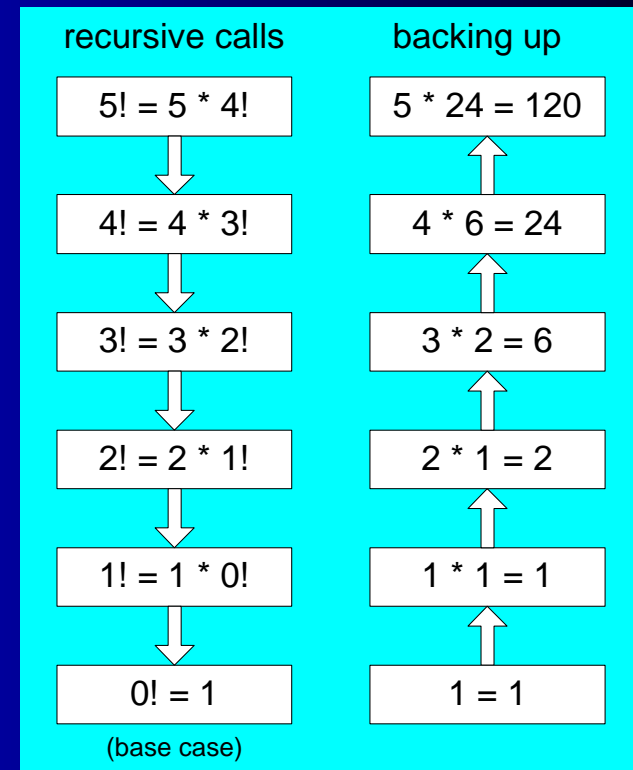
View the [complete program](#)

Calculating a Factorial (1 of 3)

This function calculates the factorial of integer n . A new value of n is saved in each stack frame:

```
int function factorial(int n)
{
    if(n == 0)
        return 1;
    else
        return n * factorial(n-1);
}
```

As each call instance returns, the product it returns is multiplied by the previous value of n .



Calculating a Factorial (2 of 3)

Factorial PROC

```
    push ebp
    mov  ebp,esp
    mov  eax,[ebp+8]           ; get n
    cmp  eax,0                ; n < 0?
    ja   L1                   ; yes: continue
    mov  eax,1                 ; no: return 1
    jmp  L2

L1:  dec  eax
     push eax                  ; Factorial(n-1)
     call Factorial

; Instructions from this point on execute when each
; recursive call returns.

ReturnFact:
    mov  ebx,[ebp+8]           ; get n
    mul  ebx                   ; eax = eax * ebx

L2:  pop  ebp                  ; return EAX
     ret  4                    ; clean up stack
Factorial ENDP
```

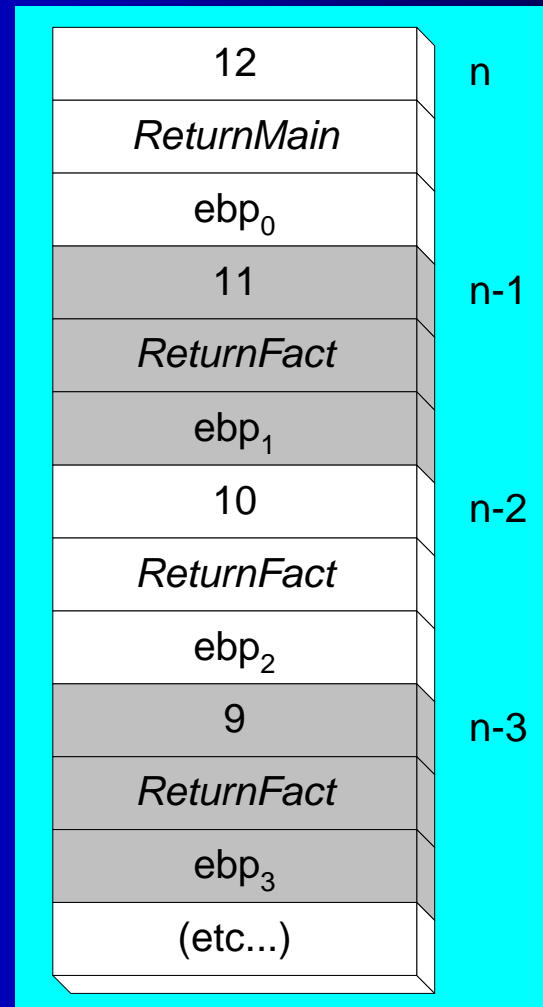
See the [program listing](#)

Calculating a Factorial (3 of 3)

Suppose we want to calculate 12!

This diagram shows the first few stack frames created by recursive calls to Factorial

Each recursive call uses 12 bytes of stack space.



What's Next

- Stack Frames
- Recursion
- **INVOKE, ADDR, PROC, and PROTO**
- Creating Multimodule Programs

INVOKE, ADDR, PROC, and PROTO

- INVOKE Directive
- ADDR Operator
- PROC Directive
- PROTO Directive
- Parameter Classifications
- Example: Exchanging Two Integers
- Debugging Tips

Not in 64-bit
mode!

INVOKE Directive

- In 32-bit mode, the INVOKE directive is a powerful replacement for Intel's CALL instruction that lets you pass multiple arguments
- Syntax:
`INVOKE procedureName [, argumentList]`
- *ArgumentList* is an optional comma-delimited list of procedure arguments
- Arguments can be:
 - immediate values and integer expressions
 - variable names
 - address and ADDR expressions
 - register names

INVOKE Examples

```
.data
byteVal BYTE 10
wordVal WORD 1000h
.code
    ; direct operands:
    INVOKE Sub1,byteVal,wordVal

    ; address of variable:
    INVOKE Sub2,ADDR byteVal

    ; register name, integer expression:
    INVOKE Sub3,eax,(10 * 20)

    ; address expression (indirect operand):
    INVOKE Sub4,[ebx]
```

Not in 64-bit
mode!

ADDR Operator

- Returns a near or far pointer to a variable, depending on which memory model your program uses:
 - Small model: returns 16-bit offset
 - Large model: returns 32-bit segment/offset
 - Flat model: returns 32-bit offset
- Simple example:

```
.data
myWord WORD ?
.code
INVOKE mySub, ADDR myWord
```


Not in 64-bit
mode!

PROC Directive (1 of 2)

- The PROC directive declares a procedure with an optional list of named parameters.
- Syntax:

label PROC *paramList*

- *paramList* is a list of parameters separated by commas. Each parameter has the following syntax:

paramName : *type*

type must either be one of the standard ASM types (BYTE, SBYTE, WORD, etc.), or it can be a pointer to one of these types.

PROC Directive (2 of 2)

- Alternate format permits parameter list to be on one or more separate lines:

label PROC, ← comma required
 paramList

- The parameters can be on the same line . . .

param-1:type-1, param-2:type-2, . . . , param-n:type-n

- Or they can be on separate lines:

param-1:type-1,
param-2:type-2,
. . . ,
param-n:type-n

AddTwo Procedure (1 of 2)

- The AddTwo procedure receives two integers and returns their sum in EAX.

```
AddTwo PROC,  
    val1:DWORD, val2:DWORD  
  
    mov eax,val1  
    add eax,val2  
  
    ret  
AddTwo ENDP
```

PROC Examples (2 of 3)

FillArray receives a pointer to an array of bytes, a single byte fill value that will be copied to each element of the array, and the size of the array.

```
FillArray PROC,  
    pArray:PTR BYTE, fillVal:BYTE  
    arraySize:DWORD  
  
    mov ecx,arraySize  
    mov esi,pArray  
    mov al,fillVal  
L1:  mov [esi],al  
    inc esi  
    loop L1  
    ret  
FillArray ENDP
```

PROC Examples (3 of 3)

```
Swap PROC,  
    pValX:PTR DWORD,  
    pValY:PTR DWORD  
    . . .  
Swap ENDP
```

```
ReadFile PROC,  
    pBuffer:PTR BYTE  
    LOCAL fileHandle:DWORD  
    . . .  
ReadFile ENDP
```

PROTO Directive

- Creates a procedure prototype
- Syntax:
 - *label PROTO paramList*
- Parameter list not permitted in 64-bit mode
- Every procedure called by the INVOKE directive must have a prototype
- A complete procedure definition can also serve as its own prototype

PROTO Directive

- Standard configuration: PROTO appears at top of the program listing, INVOKE appears in the code segment, and the procedure implementation occurs later in the program:

```
MySub PROTO                ; procedure prototype
```

```
.code
```

```
INVOKE MySub                ; procedure call
```

```
MySub PROC                  ; procedure implementation
```

```
.
```

```
.
```

```
MySub ENDP
```

PROTO Example

- Prototype for the ArraySum procedure, showing its parameter list:

```
ArraySum PROTO,  
    ptrArray:PTR DWORD,    ; points to the array  
    szArray:DWORD          ; array size
```

Parameters are not permitted in 64-bit mode.

Parameter Classifications

- An **input parameter** is data passed by a calling program to a procedure.
 - The called procedure is not expected to modify the corresponding parameter variable, and even if it does, the modification is confined to the procedure itself.
- An **output parameter** is created by passing a pointer to a variable when a procedure is called.
 - The procedure does not use any existing data from the variable, but it fills in a new value before it returns.
- An **input-output parameter** is a pointer to a variable containing input that will be both used and modified by the procedure.
 - The variable passed by the calling program is modified.

Trouble-Shooting Tips

- Save and restore registers when they are modified by a procedure.
 - Except a register that returns a function result
- When using INVOKE, be careful to pass a pointer to the correct data type.
 - For example, MASM cannot distinguish between a DWORD argument and a PTR BYTE argument.
- Do not pass an immediate value to a procedure that expects a reference parameter.
 - Dereferencing its address will likely cause a general-protection fault.

What's Next

- Stack Frames
- Recursion
- INVOKE, ADDR, PROC, and PROTO
- **Creating Multimodule Programs**
- Advanced Use of Parameters (optional)
- Java Bytecodes (optional)

Multimodule Programs

- A **multimodule program** is a program whose source code has been divided up into separate ASM files.
- Each ASM file (module) is assembled into a separate OBJ file.
- All OBJ files belonging to the same program are linked using the **link** utility into a single EXE file.
 - This process is called **static linking**

Advantages

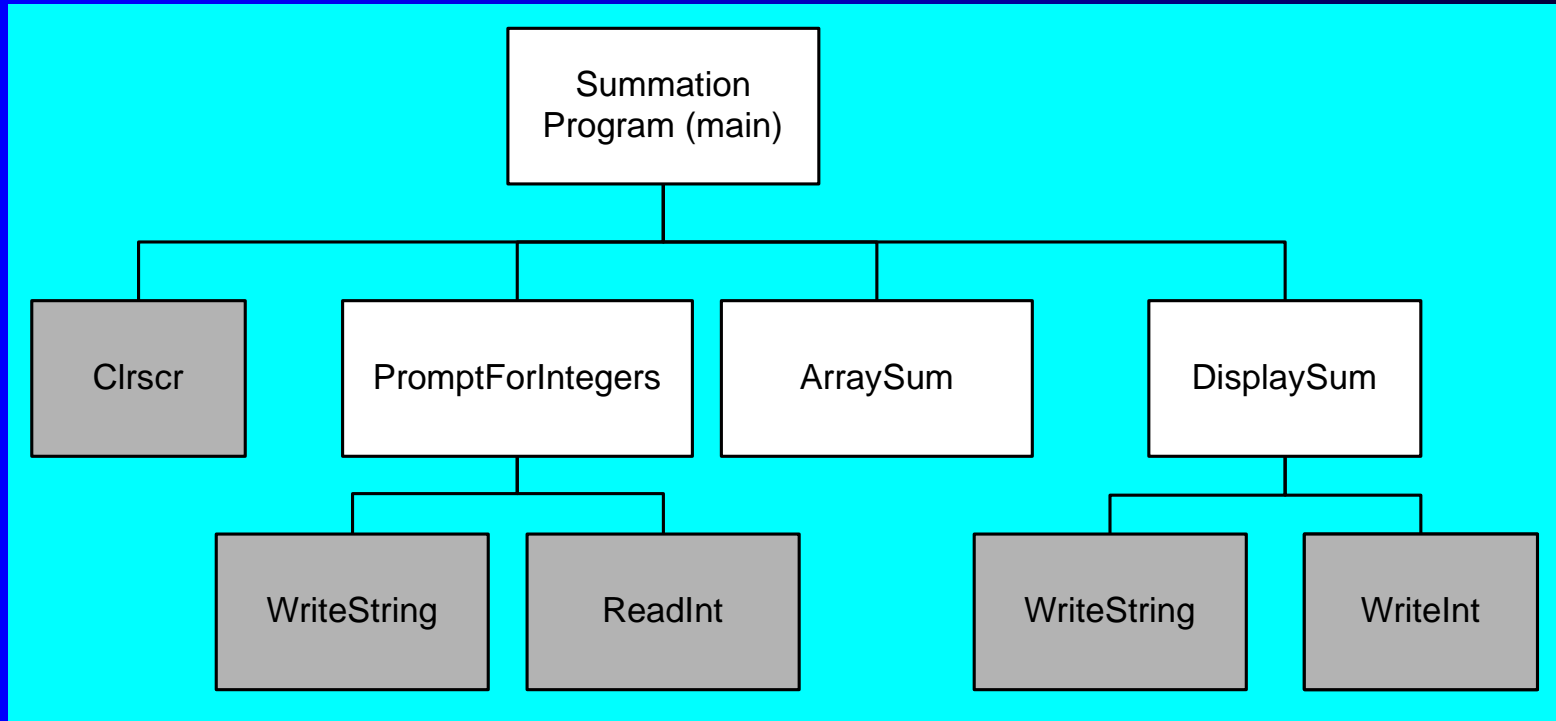
- Large programs are easier to write, maintain, and debug when divided into separate source code modules.
- When changing a line of code, only its enclosing module needs to be assembled again. Linking assembled modules requires little time.
- A module can be a container for logically related code and data (think object-oriented here...)
 - **encapsulation**: procedures and variables are automatically hidden in a module unless you declare them public

Creating a Multimodule Program

- Here are some basic steps to follow when creating a multimodule program:
 - Create the main module
 - Create a separate source code module for each procedure or set of related procedures
 - Create an include file that contains procedure prototypes for **external procedures** (ones that are called between modules)
 - Use the INCLUDE directive to make your procedure prototypes available to each module

Example: ArraySum Program

- Let's review the ArraySum program from Chapter 5.



Each of the four white rectangles will become a module. This will be a 32-bit application.

Sample Program output

```
Enter a signed integer: -25
```

```
Enter a signed integer: 36
```

```
Enter a signed integer: 42
```

```
The sum of the integers is: +53
```


INCLUDE File

The `sum.inc` file contains prototypes for external functions that are not in the Irvine32 library:

```
INCLUDE Irvine32.inc
```

```
PromptForIntegers PROTO,  
    ptrPrompt:PTR BYTE,           ; prompt string  
    ptrArray:PTR DWORD,          ; points to the array  
    arraySize:DWORD              ; size of the array
```

```
ArraySum PROTO,  
    ptrArray:PTR DWORD,          ; points to the array  
    count:DWORD                  ; size of the array
```

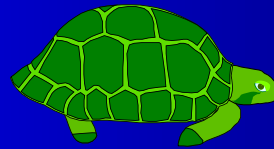
```
DisplaySum PROTO,  
    ptrPrompt:PTR BYTE,          ; prompt string  
    theSum:DWORD                 ; sum of the array
```

Inspect Individual Modules

- Main
- PromptForIntegers
- ArraySum
- DisplaySum

Summary

- Stack parameters
 - more convenient than register parameters
 - passed by value or reference
 - ENTER and LEAVE instructions
- Local variables
 - created on the stack below stack pointer
 - LOCAL directive
- Recursive procedure calls itself
- Calling conventions (C, stdcall)
- MASM procedure-related directives
 - INVOKE, PROC, PROTO



53 68 75 72 79 6F