

- HW Collection 2 due Oct. 16, 4:30pm (in office/mailbox)
→ will return these on Oct. 23
- also okay to submit on Oct. 21, before class
→ cannot return these on Oct. 23
- on Oct. 14, office hours 10-10:45 am and 12:30-1 pm
but not 1-1:45 pm

Theorem A feasible point \bar{x} is optimal for (P) if and only if there exists no feasible descent direction at \bar{x} with respect to c :

$$\begin{array}{l} \min c^T x \\ Ax \geq b \end{array}$$

$$c^T d \geq 0 \quad \forall d \in \mathbb{R}^n \text{ with } A_{\bar{x}} d \geq 0$$

This optimality condition is highly impractical. There are infinitely many directions.

Theorem (Farkas Lemma)

Let A be a matrix in $\mathbb{R}^{k \times n}$ and $c \in \mathbb{R}^n$. Then

$$c^T d \geq 0 \quad \forall d \in \mathbb{R}^n \text{ with } Ad \geq 0 \Leftrightarrow \exists y \geq 0 \in \mathbb{R}^k \text{ with } A^T y = c.$$

↑
there exists a non-neg. combination of rows of A giving c

Proof: " \Leftarrow " Let $A^T y = c$ for some $y \geq 0$ and let d be any vector with $Ad \geq 0$. Then

$$c^T d = (A^T y)^T d = \underbrace{y^T}_{\geq 0} \underbrace{Ad}_{\geq 0} \geq 0$$

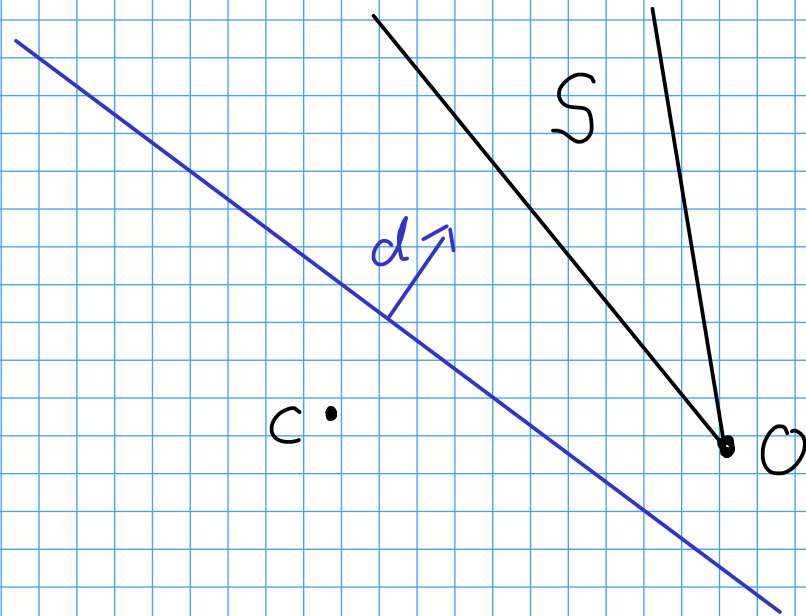
" \Rightarrow "

Define the set $S = \{s \in \mathbb{R}^n : A^T y = s \text{ for some } y \geq 0\}$.

S is a (closed) convex cone that contains 0 and rows a_i ; (from the unit vector $y = e_i$) for all $i \leq k$.

$\neg B \Rightarrow \neg A$

Suppose $c \notin S$. Then there exists a separating hyperplane with normal vector $d \in \mathbb{R}^n$ and $z \in S$ such that $d^T c < d^T z \leq d^T x \quad \forall x \in S$, so in particular $d^T c < d^T 0 = 0$. \leftarrow because $0 \in S$



Further, because S is a convex cone that contains z and a_i , for all $i \leq k$, we have $z + a_i \in S$. $\leftarrow K + K \subseteq K$

Thus $d^T z \leq d^T (z + a_i) \Rightarrow d^T a_i \geq 0 \quad \forall i \leq k$

$\Rightarrow Ad \geq 0$ (but at the same time $c^T d < 0$) \square

The Farkas Lemma comes in many variants. These are called Theorems of the Alternative. They all have great geometric interpretations.

Theorem (Farkas Lemma, variant 2)

Given a matrix $A \in \mathbb{R}^{k \times n}$ and $c \in \mathbb{R}^n$, one and only one of the following systems admits a solution:

$$(I) \quad c^T d < 0, Ad \geq 0 \qquad (II) \quad y \geq 0, A^T y = c$$

for $A = A_{\bar{x}}$: there exists a strictly improving direction

there exists a conic combination of the normals of the active constraints that gives c

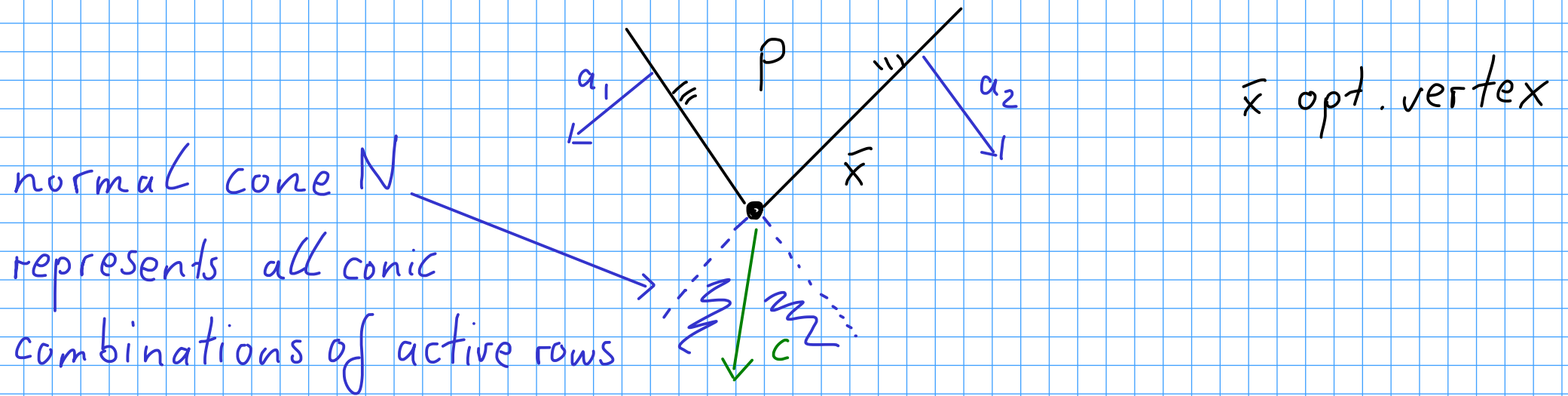
Corollary A feasible point \bar{x} is optimal for (P) if and only if there exists a vector $y \geq 0$ such that $A_{\bar{x}}^T y = c$.

$$\min c^T x, Ax \geq b$$

nice optimality condition!

HW 14 State this corollary for all canonical form LPs (max/min, $Ax \leq b$ / $Ax \geq b$, all 4 combinations). Visualize all 4 cases with a small sketch. (No explanation/justification.)

Visualization $\max c^T x, Ax \leq b$



$\hat{=}$ set of all objective functions for which \bar{x} is optimal

Duality and Complementary Slackness

Farkas Lemma + Geometric Insight give us better optimality conditions for LPs

Theorem

First-Order Optimality Conditions

A point x is optimal for (P) if and only if there exists a non-negative $y \in \mathbb{R}^m$ such that

$$\begin{aligned} \min c^T x \\ Ax \geq b \end{aligned}$$

a) $Ax \geq b$ feasible

b) $A^T y = c$ construct obj. fun. from normals

c) $y_i \cdot (a_i^T x - b_i) = 0 \quad \forall i \in m$

only use active normals (normals of active constraints) for the construction in b)

KKT conditions for LPs

Karush-Kuhn-Tucker