

# Econ C103: Game Theory and Networks

## Module I (Game Theory): Lecture 11

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### Readings:

- ① Diamond, Douglas and Philip Dybvig. "Bank Runs, Deposit Insurance, and Liquidity", Journal of Political Economy, 1983.

# Bayes' Cournot Duopoly

- Two firms  $i$  and  $j$  compete by producing quantities  $q_i \in \mathbb{R}$  and  $q_j \in \mathbb{R}$ , resp., under price function  $P(q_i, q_j) = (a + \omega) - b(q_i + q_j)$ ,  $a, b > 0$ ,  $\omega \sim \mathcal{N}(0, \sigma^2)$ . Marginal cost of production is  $c > 0$ .
- Each firm  $k = i, j$  observes signal  $\theta_k = \omega + \epsilon_k$ ,  $\epsilon_k \sim \mathcal{N}(0, 1/\gamma_k)$ ;  $\epsilon_k$  is the noise in  $k$ 's signal  $\theta_k$ . The “state space” is  $\{(\omega, \epsilon_i, \epsilon_j)\}$ ,  $\gamma_k \in \mathbb{R}_+$  is the “precision” of  $k$ 's signal;  $\gamma_k = 0$  is “no information”.
- Each firm maximizes profit given their signal, and the other firm's production strategy  $s_k : \mathbb{R} \mapsto \mathbb{R}$ , a function of  $t_k$ .
- Player  $i$ 's Bayesian updating yields the following expectations (similar for  $j$ 's expectations; *feel free to take these next expressions as given*):

1.  $\mathbb{E}[\omega|\theta_i] = \mathbb{E}[\omega|t_i] = e_i t_i$ , and

2.  $\mathbb{E}[\theta_j|\theta_i] = \mathbb{E}[t_j|t_i] / e_j = e_i t_i$ , so  $\mathbb{E}[t_j|t_i] = e_i e_j t_i$

where:

$$e_i \equiv \sqrt{\frac{\sigma^2}{\sigma^2 + \gamma_i^{-1}}} \in [0, 1],$$
$$t_i \equiv e_i \theta_i.$$

## Bayes' Cournot Duopoly

- For firm  $i$ , given  $s_j$  and upon observing  $t_i$ , her optimal production is:

$$s_i^*(t_i) = \operatorname{argmax}_{q_i} ((a + \mathbb{E}[\omega|t_i]) - b(q_i + \mathbb{E}[s_j(t_j)|t_i]))q_i - cq_i.$$

- First-order condition:  $a + \mathbb{E}[\omega|t_i] - 2bq_i - b\mathbb{E}[s_j(t_j)|t_i] - c = 0$ .
- Solving for  $q_i$  gives firm  $i$ 's best response:

$$s_i^*(t_i) = \frac{a + e_i t_i - b\mathbb{E}[s_j(t_j)|t_i] - c}{2b}.$$

# Bayes' Cournot Duopoly

- Assume  $j$  uses linear strategy  $s_j(t_j) = \alpha_j + \beta_j t_j$ . Then:

$$\mathbb{E}[s_j(t_j)|t_i] = \alpha_j + \beta_j \mathbb{E}[t_j|t_i] = \alpha_j + \beta_j e_i e_j t_i.$$

- This gives firm  $i$ 's optimal strategy:

$$s_i^*(t_i) = \frac{a + e_i t_i - b(\alpha_j + \beta_j e_i e_j t_i) - c}{2b},$$

which implies:

$$\begin{aligned}\alpha_i^* &= \frac{a - b\alpha_j - c}{2b}, \\ \beta_i^* &= \frac{e_i - b\beta_j e_i e_j}{2b}.\end{aligned}$$

$\beta_i^*$  decreasing in  $e_j$ : better information of competition  $\Rightarrow$  less responsive to private information (learning  $\omega$  is high  $\Rightarrow j$  also learns  $\omega$  is high  $\Rightarrow q_j$  increases  $\Rightarrow$  price and marginal revenue decrease!).

- Similarly for  $j$ ...

## Bayes' Cournot Duopoly

$$\begin{aligned}\alpha_i^* &= (a - b\alpha_j^* - c)/(2b); & \alpha_j^* &= (a - b\alpha_i^* - c)/(2b), \\ \beta_i^* &= (e_i - b\beta_j^* e_i e_j)/(2b); & \beta_j^* &= (e_j - b\beta_i^* e_j e_i)/(2b).\end{aligned}$$

- Solving this system yields:

$$\alpha_k^* = \frac{a - c}{3b}, \quad \forall k = i, j$$

$$\beta_i^* = \frac{e_i(2 - e_j^2)}{b(4 - e_i^2 e_j^2)},$$

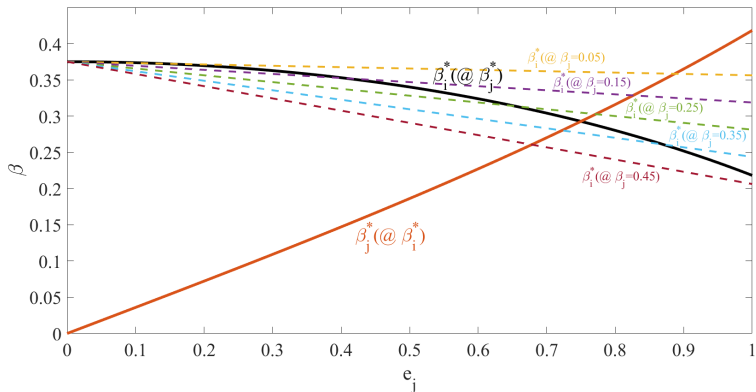
$$\beta_j^* = \frac{e_j(2 - e_i^2)}{b(4 - e_i^2 e_j^2)}.$$

# Bayes' Cournot Duopoly

$$\alpha_i^* = \frac{a - c}{3b},$$
$$\beta_i^* = \frac{e_i(2 - e_j^2)}{b(4 - e_i^2 e_j^2)}.$$

- $\alpha_k^*$  corresponds to the PNE in the Cournot Duopoly.
- $\beta_i^*$  increasing in  $e_i$ : better information  $\Rightarrow$  optimal strategy more responsive to private information.
- $\beta_i^*$  decreases in  $e_j$ : better information of competition ( $e_j \uparrow$ )  $\Rightarrow$  less responsive to private information ( $\beta_i^* \downarrow$ )  $\Rightarrow \beta_j^* \uparrow \Rightarrow \beta_i^* \downarrow \dots!$   
Players' information use become interdependent in equilibrium...

# Bayes' Cournot Duopoly



**Figure:** Information use:  $\beta$  = slope of strategy to private signal.

- In the figure,  $e_i = 0.75$ , and  $e_j \in [0, 1]$  (the x-axis). Equilibrium  $\beta_i^*$  (black line) decreases in  $e_j$ , at an increasing rate as  $e_j$  increases; in contrast,  $\beta_i^*$  when  $\beta_j$  is held fixed (dashed lines) has a fixed slope.

## Diamond-Dybvig (1983) bank-run model

- Three periods  $t = 1, 2, 3$ . A unit measure of consumers deposit wealth in a bank at  $t = 1$ ; consumption occurs in  $t = 2, 3$ .
- Measure  $\theta$  values consumption only in  $t = 2$  (type A), measure  $1 - \theta$  values consumption in both  $t = 2, 3$  (type B):

$$U_i(c_1, c_2) = \begin{cases} u(c_2) & \text{with prob. } \theta \text{ (type A)} \\ \rho u(c_2 + c_3) & \text{with prob. } 1 - \theta \text{ (type B)} \end{cases} ;$$

where  $u(\cdot)$  is twice continuously differentiable, increasing and concave, and  $0 < \rho < 1$ . So, types A's have greater need.

- Bank holds 1 unit of liquidity (money); withdrawals at  $t = 2$  and  $t = 3$  yield consumption value 1 and  $R > 1$ , respectively.
- Consumption type is private information to each consumer  $i \in [0, 1]$ ; the “state” determines which consumers are of types A/B.  
BUT, on aggregate measure  $\theta$  is of type A (so, no aggregate risk!).



## Diamond-Dybvig (1983) bank-run model

- Socially efficient consumption allocation has  $c_3^A = 0$  and  $c_2^B = 0$ .
- If the bank could require only type A's withdraw in period  $t = 2$ , then bank's budget constraint becomes:

$$\theta c_2^A + (1 - \theta)(c_3^B / R) = 1 \Leftrightarrow c_3^B = \frac{(1 - \theta c_2^A)R}{1 - \theta}.$$

- Planner's problem then solves:

$$\max_{c_2^A \geq 0} \theta u(c_2^A) + (1 - \theta)\rho u\left(\frac{(1 - \theta c_2^A)R}{1 - \theta}\right).$$

First-order condition for efficient  $c_2^{A*}$ :

$$u'(c_2^{A*}) = R\rho u'\left(\frac{(1 - \theta c_2^{A*})R}{1 - \theta}\right) = R\rho u'(c_3^{B*}).$$

If  $R\rho > 1$  then  $c_3^{B*} > c_2^{A*}$ , and type B's prefer to wait (greater  $c_3^B$ ).

# Diamond-Dybvig (1983) bank-run model

- So, if  $R\rho > 1$  and bank promises  $c_2^{A*}$  to any consumer withdrawing at  $t = 2$ , then it is a BNE for type A's to withdraw at  $t = 2$ , type B's at  $t = 3$ , which yields efficient consumption profile  $(c_2^{A*}, c_3^{B*})$ .
- Bank-run BNE: also assume that  $R\rho \in (1, \bar{R}\rho]$ , where  $\bar{R}\rho$  satisfies:

$$u'(1) = \bar{R}\rho u'(R).$$

Then,  $1 < c_2^{A*} < c_3^{B*}$  (while  $\theta c_2^A + (1 - \theta)(c_3^B / R) = 1$  still holds).

- What if type B's believe all consumers will withdraw at  $t = 2$ ?
- Then, all consumer “run to the bank”, but only first  $1/c_2^{A*}$  consume!
- Two “sunspot” BNE:
  - ① an efficient (liquid) BNE, in which only type A withdraws in  $t = 2$ , and
  - ② a bank-run (illiquid) BNE, in which all consumers withdraw in  $t = 2$ .