

Pair Programming Equitable Participation & Honesty Affidavit

We the undersigned promise that we have in good faith attempted to follow the principles of pair programming. Although we were free to discuss ideas with others, the implementation is our own. We have shared a common workspace and taken turns at the keyboard for the majority of the work that we are submitting. Furthermore, any non programming portions of the assignment were done independently. We recognize that should this not be the case, we will be subject to penalties as outlined in the course syllabus.

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Pair Programmer 1 (print & sign your name, then date it)

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Part 1

Question 1: Show what happens to the variable domains when forward-checking is employed after each of the first two decisions:

a. CA is assigned B.

b. AZ is assigned O.

	CA	NV	UT	AZ	CO	NM	TX
Initial Domains	R B Y O	R B Y O	R B Y O	R B Y O	R B Y O	R B Y O	R B Y O
CA is assigned B	Ⓟ	R Y O	R B Y O	R Y O	R B Y O	R B Y O	R B Y O
AZ is assigned O	Ⓟ	R Y	R B Y	Ⓞ	R B Y	R B Y	R B Y O

a. NV and AZ remove 'B' in their domains.

b. NV, UT, and NM remove 'O' in their domains

Question 2: Using maintaining arc consistency instead of forward checking.

	CA	NV	UT	AZ	CO	NM	TX
Initial Domains	R B Y O	R B Y O	R B Y O	R B Y O	R B Y O	R B Y O	R B Y O
CA is assigned B	Ⓟ	R Y O	R B Y O	R Y O	R B Y O	R B Y O	R B Y O
AZ is assigned O	Ⓟ	R Y	R B Y	Ⓞ	R B Y	R B Y	R B Y O

a. Check (CA, NV) → (CA, AZ) → (AZ, NM) → (AZ, CO) → (AZ, UT) → (AZ, NV) → (NV, UT) → (NV, AZ): NV and AZ remove 'B' in their domains.

b. Check (AZ, NV) → (AZ, UT) → (AZ, CO) → (AZ, NM) → (NV, UT) → (UT, CO) → (CO, NM) → (NM, CO) → (NM, TX): NV, UT, CO, and NM remove 'O' in their domains

Question 3: Show how the conflict sets would be updated for these two moves if we were using conflict-directed backjumping.

$\{\}$ ← Conflict set ; \bigcirc ← Assigned

Conflict set	CA	NV	UT	AZ	CO	NM	TX
Initial set							
CA is assigned B	\bigcirc_R	{ CA= R }		{ CA= R }			
AZ is assigned O	\bigcirc_R	{ CA= R, AZ=O }	{ AZ=O }	\bigcirc_O	{ AZ=O }	{ AZ=O }	

Question 4: Suppose students have the following set of simplified constraints:

- Students cannot be enrolled in overlapping courses.
- Students cannot enroll in classes on days that they work.
- A student may not enroll in a course that does not have space.

a. Write a specification for this constraint satisfaction problem.

Variables:

Divide a student's timetable into four parts.

- MW 14:00~15:15 $\rightarrow S_iMW1$ (A period, 14:00 to 15:15, of Student i on Monday and Wednesday)
- MW 15:00~16:15 $\rightarrow S_iMW2$ (A period, 15:00 to 16:15, of Student i on Monday and Wednesday)
- TTh 14:00~15:30 $\rightarrow S_iTTh2$ (A period, 14:00 to 15:30, of Student i on Tuesday and Thursday)
- TTh 18:00~19:15 $\rightarrow S_iTTh2$ (A period, 18:00 to 19:15, of Student i on Tuesday and Thursday)

One student has 4 periods, so we have $i*4$ variables.

For example,

Variables= { S_1MW1 , S_1MW2 , S_1TTh1 , S_1TTh2 , S_2MW1 , S_2MW2 , S_2TTh1 , S_2TTh2 , ... }

Domain:

Four Courses, C_1 , C_2 , C_3 and C_4 . One Empty, E, and a Work, W.

C_1 : A course whose Tutorial is 1

E: The period is free.

C_2 : A course whose Tutorial is 2

W: Students work either days.

C_3 : A course whose Tutorial is 3

C_4 : A course whose Tutorial is 4

For example,

$S_1MW1 = W$ means S_1 works either Monday or Wednesday.

$S_2MW1 = C_1$ means S_2 take a course C_1 .

Constraints:

1. $MW1$ and $MW2$ cannot take course simultaneously. If $MW1 = C_1$, the other must be E .
2. $TTh1$ and $TTh2$ cannot take course simultaneously. If $TTh1 = C_1$, the other must be E
3. If MW_x is W then the other must be W as well. $\rightarrow MW1 = W \leftrightarrow MW2 = W$
4. If TTh_x is W then the other must be W as well. $\rightarrow TTh1 = W \leftrightarrow TTh2 = W$
5. The number of C_x cannot be more than 30. E.g. Cannot be more than 30 students taking C_1 .

b. Show how encapsulation could be used to create a binarize the constraint between multiple students the enrollment limit.

i. Students cannot be enrolled in overlapping courses. & ii. Students cannot enroll in classes on days that they work.

Combine $MW1$ and $WM2$ into a set and $TTh1$ and $TTh2$ into a set. The elements in a set cannot be two C_i or more. Also, if an element is W , the other must be W too.

E.g. $U_{i,ii} = [\{ S_1MW1 = C1, S_1MW2 = C2 \}, \{ S_1MW1 = C1, S_1MW2 = E \}, \{ S_1TTh1 = C3, S_1TTh2 = W \}, \{ S_1TTh1 = E, S_1TTh2 = C4 \}, \dots]$

$\{ S_1MW1 = C1, S_1MW2 = C2 \} = \text{False}$

$\{ S_1MW1 = C1, S_1MW2 = E \} = \text{True}$

$\{ S_1TTh1 = C3, S_1TTh2 = W \} = \text{False}$

$\{ S_1TTh1 = W, S_1TTh2 = W \} = \text{True}$

$\{ S_1TTh1 = E, S_1TTh2 = W \} = \text{False}$

$\{ S_1TTh1 = E, S_1TTh2 = C4 \} = \text{True}$

iii. A student may not enroll in a course that does not have space:

We can combine all S_xMW1 to be a set and count how many $C1$. If the number of $C1$ is larger than 30 then it conflicts. Following this rule, we can combine different MW and TTh to be sets respectively.

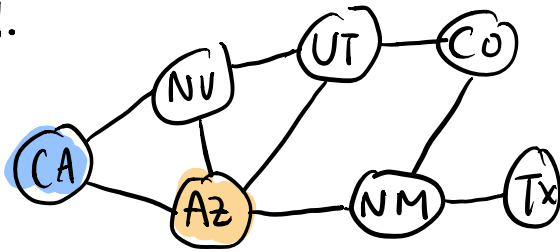
E.g. $U_{iii} = [\{S_1 = C1, S_2 = C1, S_3 = C1, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C1\dots\}, \{S_1 = C2, S_2 = C2, S_3 = C2, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C2\dots\}, \{S_1 = C3, S_2 = C3, S_3 = C3, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C3\dots\}, \{S_1 = C4, S_2 = C4, S_3 = C4, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C4\dots\}]$.

If $\{S_1 = C1, S_2 = C1, S_3 = C1, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C1\dots\}$ contains more than 30 C1, then it is failure.

If $\{S_1 = C2, S_2 = C2, S_3 = C2, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C2\dots\}$ contains more than 30 C2, then it is failure.

If $\{S_1 = C3, S_2 = C3, S_3 = C3, S_4 = E, S_5 = E \dots, S_i = W\dots, S_j = C3\dots\}$ contains more than 30 C3, then it is failure.

1.

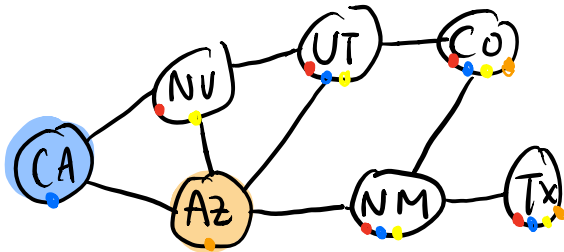


	CA	NV	AZ	UT	NM	CO	TX
initial domains	RB40	RB40	RB40	RB40	RB40	RB40	RB40

After CA = B	B	RY0	RY0	RB40	RB40	RB40	RB40
--------------	---	-----	-----	------	------	------	------

After AZ = 0	B	RY	0	RB4	RB4	RB40	RB40
--------------	---	----	---	-----	-----	------	------

2.



	CA	NV	AZ	UT	NM	CO	TX
initial domains	RB40	RB40	RB40	RB40	RB40	RB40	RB40

After CA = B	B	RY0	RY0	RB40	RB40	RB40	RB40
--------------	---	-----	-----	------	------	------	------

After AZ = 0	B	RY	0	RB4	RB4	RB40	RB40
--------------	---	----	---	-----	-----	------	------

After CA = B

Queue = (CA, NV) (CA, AZ),

NV = RY0, AZ = RY0

After $AZ=0$

Queue = $(AZ, NV), (AZ, UT), (AZ, NM)$

$NV = RY, UT = RBY, NM = RBY$

3.	CA	NV	AZ	UT	NM	CO	TX
After $CA=B$	B	$\{CA=B\}$	$\{CA=B\}$	\emptyset	\emptyset	\emptyset	\emptyset
After $AZ=0$	B	$\{CA=B, AZ=0\}$	0	$\{AZ=0\}$	$\{AZ=0\}$	\emptyset	\emptyset

4.

a. Assume we have 2 classes: Student and Course

class Student {

self.name = " " # Student's name

self.courses = [] # Student's course list

self.workday = [] # Days this student work

class Course {

self.name = " " # Course name

self.time = " " # Start & end time of the course

self.day = " " # day of the course

self.capacity = 30 # Capacity of the course

self.size = 0 # Current size of the course

CSP Specifications:

$X = \{\text{Student 1, Student 2, ..., Student } n\}$ # instances of Student

$D_x = \{\text{All student names, course lists, workdays}\}$ # instance variables of student

$Y = \{\text{Course 1, Course 2, Course 3, Course 4}\}$ # instances of Course, total 4.

$D_y = \{\text{courses name, time, day, capacity, size}\}$ # instance variables of student

$C = \{$

for \exists student in X : # the only overlapping courses are 1 and 3

if 1 in student.courses and 3 in student.courses:

return False

for \exists course in Y : # in case the class is full

if course.size > course.capacity:

return False

for \exists student in X :

for \exists course in Y : # in case students need to work

if student.workday overlaps course.day:

return False

b. Create an encapsulated variable U

Cartesian product $U = X \times Y$

We can binarize the constraint by multiplying X and Y .

Hsuan-Yu Liu & Jinqi Cheng A5 PartII Source code

driver.py

```
from csp_lib.sudoku import (Sudoku, easy1, harder1)
from constraint_prop import AC3
from csp_lib.backtrack_util import mrv, forward_checking, first_unassigned_variable, unordered_c
from backtrack import backtracking_search

for puzzle in [easy1, harder1]:
    s = Sudoku(puzzle) # construct a Sudoku problem
    print("---initial state---")
    s.display(s.infer_assignment())
    print("\n")
    AC3(s)

    print("---after AC3---")
    s.display(s.infer_assignment())
    print("\n")

    if not s.goal_test(s.curr_domains):
        solved = backtracking_search(s, select_unassigned_variable=mrv, inference=forward_checking)
        print("---after backtracking---")
        s.display(solved)
```

backtrack.py

```

from csp_lib.backtrack_util import (first_unassigned_variable,
                                    unordered_domain_values,
                                    no_inference)

def backtracking_search(csp,
                        select_unassigned_variable=first_unassigned_variable,
                        order_domain_values=unordered_domain_values,
                        inference=no_inference):
    """backtracking_search
    Given a constraint satisfaction problem (CSP),
    a function handle for selecting variables,
    a function handle for selecting elements of a domain,
    and a set of inferences, solve the CSP using backtrack search
    """
    # See Figure 6.5] of your book for details

    def backtrack(assignment):
        """Attempt to backtrack search with current assignment
        Returns None if there is no solution. Otherwise, the
        csp should be in a goal state.
        raise NotImplemented
        """

        """
        # MRC will throw an error as all variables have been assigned.
        Thus, return solution when all variables are assigned.
        """

        if len(assignment)==81:
            return assignment

        unassigned_var = select_unassigned_variable(assignment, csp)
        if unassigned_var is not None:
            for domain in order_domain_values(unassigned_var, assignment, csp):
                removals=[]
                csp.assign(unassigned_var, domain, assignment)
                infer = inference(csp, unassigned_var, domain, assignment, removals)
                if infer:
                    assignment = backtrack(assignment)
                    if csp.goal_test(assignment):
                        return assignment
                csp.unassign(unassigned_var, assignment)
                csp.restore(removals)
            return assignment

    # Call with empty assignments, variables accessed
    # through dynamic scoping (variables in outer
    # scope can be accessed in Python)
    result = backtrack({})
    assert result is None or csp.goal_test(result)
    return result

```

constraint_prop.py

```
'''
Constraint propagation
'''

def AC3(csp, queue=None, removals=None):
    """AC3 constraint propagation
    # Hints:
    # Remember that:
    #   csp.variables is a list of variables
    #   csp.neighbors[x] is the neighbors of variable x
    raise NotImplemented
    """
    if not queue:
        variables = csp.variables
        queue = [ (var, neighbor) for var in variables for neighbor in csp.neighbors[var] ]

    while(queue):
        arc = queue.pop()
        if revise(csp, arc[0], arc[1]):
            if not csp.curr_domains[arc[0]]:
                break
            else:
                for neighbor in csp.neighbors[arc[0]]:
                    if neighbor != arc[1]:
                        queue.append((neighbor, arc[0]))
    return csp

def revise(csp, Xi, Xj):
    revised = False
    assignment = csp.infer_assignment()
    domains = csp.curr_domains[Xi]
    for d_val in domains:
        if csp.nconflicts(Xi, d_val, assignment):
            csp.prune(Xi, d_val, None)
            revised = True
    return revised
```

Output

```

---initial state---
. . 3 | . 2 . | 6 . .
9 . . | 3 . 5 | . . 1
. . 1 | 8 . 6 | 4 . .
-----+-----+-----
. . 8 | 1 . 2 | 9 . .
7 . . | . . . | . . 8
. . 6 | 7 . 8 | 2 . .
-----+-----+-----
. . 2 | 6 . 9 | 5 . .
8 . . | 2 . 3 | . . 9
. . 5 | . 1 . | 3 . .

```

```

---after AC3---
4 8 3 | 9 2 1 | 6 5 7
9 6 7 | 3 4 5 | 8 2 1
2 5 1 | 8 7 6 | 4 9 3
-----+-----+-----
5 4 8 | 1 3 2 | 9 7 6
7 2 9 | 5 6 4 | 1 3 8
1 3 6 | 7 9 8 | 2 4 5
-----+-----+-----
3 7 2 | 6 8 9 | 5 1 4
8 1 4 | 2 5 3 | 7 6 9
6 9 5 | 4 1 7 | 3 8 2

```

```

---initial state---
4 1 7 | 3 6 9 | 8 . 5
. 3 . | . . . | . . .
. . . | 7 . . | . . .
-----+-----+-----
. 2 . | . . . | . 6 .
. . . | . 8 . | 4 . .
. . . | . 1 . | . . .
-----+-----+-----
. . . | 6 . 3 | . 7 .
5 . . | 2 . . | . . .
1 . 4 | . . . | . . .

```

```

---after AC3---
4 1 7 | 3 6 9 | 8 2 5
. 3 . | . . . | . . .
. . . | 7 . . | . . .
-----+-----+-----
. 2 . | . . . | . 6 .
. . . | . 8 . | 4 . .
. . . | . 1 . | . . .
-----+-----+-----

```

```
. . . | 6 . 3 | . 7 .  
5 . . | 2 . . | . . .  
1 . 4 | . . . | . . .
```

---after backtracking---

```
4 1 7 | 3 6 9 | 8 2 5  
6 3 2 | 1 5 8 | 9 4 7  
9 5 8 | 7 2 4 | 3 1 6
```

-----+-----+-----

```
8 2 5 | 4 3 7 | 1 6 9  
7 9 1 | 5 8 6 | 4 3 2  
3 4 6 | 9 1 2 | 7 5 8
```

-----+-----+-----

```
2 8 9 | 6 4 3 | 5 7 1  
5 7 3 | 2 9 1 | 6 8 4  
1 6 4 | 8 7 5 | 2 9 3
```