

$A = [5, 8, 1, 15, 7, 6, 2]$

Q → Find max element → $TC = O(N)$ $SC = O(1)$

```

ans = A[0]
for i → 1 to (N-1)
    if (A[i] > ans)
        ans = A[i]
return ans

```

Q → Find second max element → $TC = O(2N) \approx O(N)$ $SC = O(2) = O(1)$

Q → Find third max element → $TC = O(3N) = O(N)$ $SC = O(3) = O(1)$

⋮

Q → Find K^{th} max element → $TC = O(KN)$ $SC = O(K) \rightarrow O(1)$

$A = [5, 8, 1, 15, 7, 6, 2]$ $K = 4$
 $K=1$ 15
 $K=2$ 8
 $K=3$ 7
 $K=4$ 6
 $K=5$ 5
 $K=6$ 2

Selection Sort

$TC = O(N^2)$ $SC = O(1)$

```

for i → N-1 to 1 ←
    maxId = 0
    for j → 1 to i
        if (A[j] > A[maxId])
            maxId = j
    swap(A[i], A[maxId]) ←

```

$A = [5, 8, 1, 15, 7, 6, 2]$
 $i = N-1$
 $maxId = 3$ $j \rightarrow 1$ to i

Q → Given an integer array where all odd elements are sorted & all even elements are sorted. Sort the array.

Handwritten diagrams illustrating the recursive process of sorting an array.

Top Diagram: Shows an array $A = [3, 9, 2, 4, 15, 10, 19]$ being split into two sub-arrays: $[3, 9, 15, 19]$ and $[2, 4, 10]$. The time complexity for each recursive call is noted as $TC = O(N)$.

Bottom Diagram: Shows an array $A = [1, 3, 7, 13, 2, \dots]$ with elements being compared and swapped. The final sorted array is shown as $[1, 2, 3, 7, 13, \dots]$. The time complexity for the entire process is noted as $TC = O(N+N) = O(N)$ and $SC = O(N)$.

✓ Merge two sorted array into one. ✓

// A[N+M] ← I/P → B[N] ✓ C[M] ✓
st - mid mid+1 - end

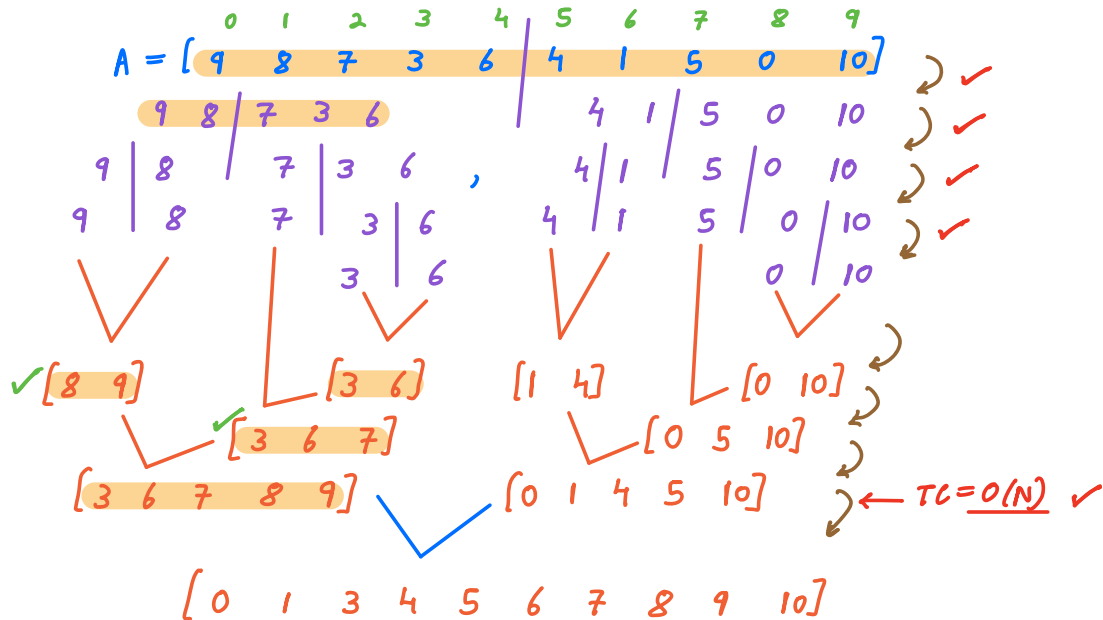
```

i = 0 // B      j = 0 // C
for k → 0 to (N+M-1)
    if (i == N) ←
        A[k] = C[j] ✓
        j += 1 ✓
    else if (j == M || B[i] ≤ C[j])
        A[k] = B[i]
        i += 1
    else
        A[k] = C[j]
        j += 1 // ans += (N-i)

```

$$B = \begin{bmatrix} 0 & 2 & 5 \end{bmatrix} \quad C = \begin{bmatrix} 4 & 6 & 10 \end{bmatrix}$$

Merge sort → Divide & Conquer



```

void sort(A, st, end) {
    if (st == end)
        return;
    mid = (st + end) / 2;
    sort(A, st, mid);
    sort(A, mid + 1, end);
    merge(A, st, mid, end);
}
    
```

Time Complexity: $TC = O(N \log_2(N))$ ✓

Space Complexity: $SC = O(N)$ ✓

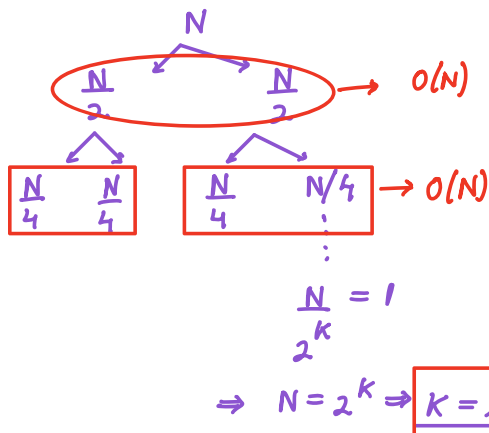
$$T(N) = 2T(N/2) + O(N)$$

$$a=2 \quad b=2 \quad d=1$$

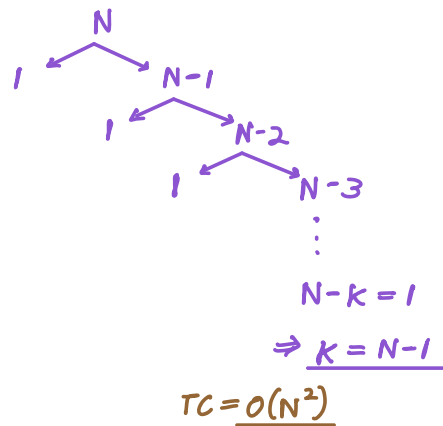
$$b^d = 2^1 = 2 = a$$

$$TC = O(N^d \log_b(N)) = O(N \log_2(N))$$

$$SC = O(N + \log(N)) = O(N) \checkmark$$



$$TC = O(N \log_2(N)) \checkmark$$



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Q → Given an integer array, count the number of inversion pairs in the array.

Inversion pair → (i, j)

$$i < j \text{ and } A[i] > A[j]$$

$$A = [8 \ 3 \ 4] \quad (0, 1) \quad (0, 2) \quad \text{Ans} = 2$$

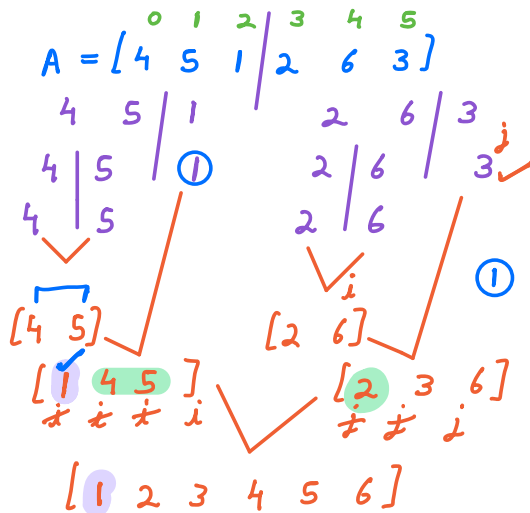
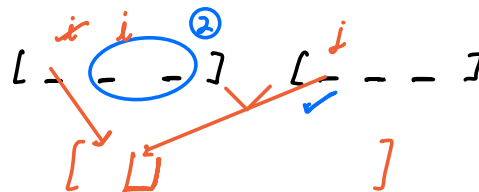
$$A = [4 \ 5 \ 1 \ 2 \ 6 \ 3] \quad (0, 2) \ (0, 3) \ (0, 5) \ (1, 2) \ (1, 3) \ (1, 5), \ (4, 5) \quad \text{Ans} = 7 \checkmark$$

$$A = [4 \ 4 \ 4 \ 4 \ 4] \quad \text{Ans} = 0$$

Brute force → $\forall i, j$ check for inversion.

$$TC = O(N^2) \quad SC = O(1)$$

Merge Sort



$$\text{Ans} = 7$$

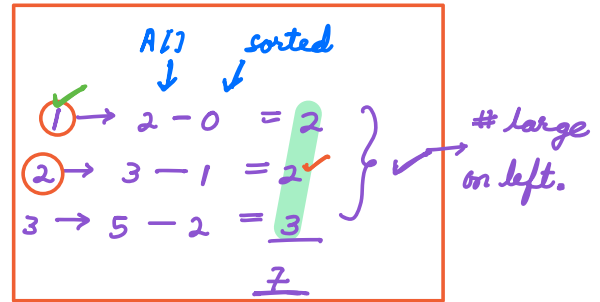
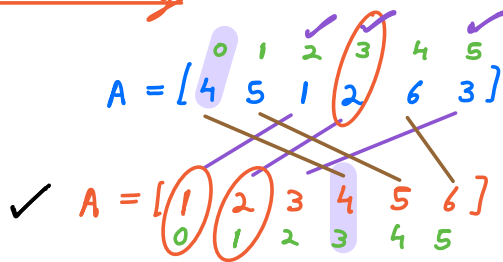
$$\textcircled{2} \quad (0, 2) \quad (1, 2)$$

$$\textcircled{2} \quad (0, 3) \quad (1, 3) \quad \textcircled{2} \quad (0, 5) \quad (1, 5)$$

$$TC = O(N \log(N)) \quad SC = O(N)$$

Ans = ϵ # remaining elements $(N-i)$ in left part if an element from right part is selected. \checkmark

Sorted Array



sorted A A

$4 \rightarrow 3 - 0 = 3$
 $5 \rightarrow 4 - 1 = 3$
 $6 \rightarrow 5 - 4 = 1$
7

small on right.

✓ (original index - sorted index)

$4 \rightarrow 0 - 3 = -3$
 $5 \rightarrow 1 - 4 = -3$
 $6 \rightarrow 4 - 5 = -1$
-7

(sorted index - original index)

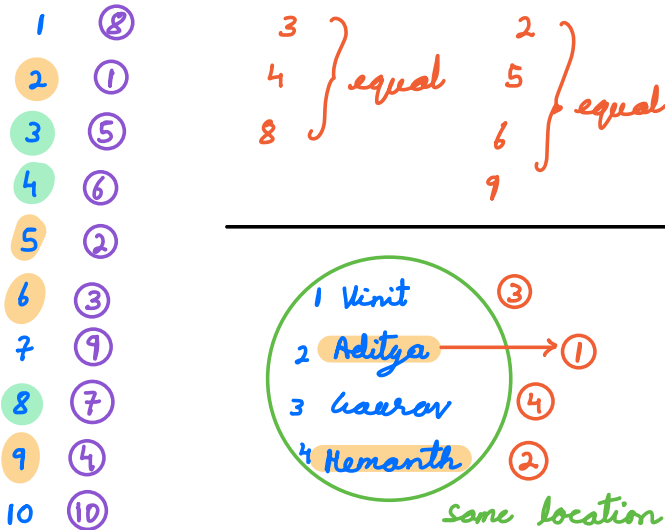
not recommended

X

$A = [\text{sorted} - 1]$

Sorted

stability → Relative order of equal elements should not change.



Stability in array → if elements are equal
compare index.

left

right

$[1 \quad 3 \quad 8]$ $[2 \quad 6 \quad 10]$
 $0 \quad i \quad i \quad \text{mid} \quad \text{mid}+1 \quad i \quad i \quad \text{end}$
 $[1 \quad 2 \quad 3 \quad 6 \quad 8 \quad 10]$

$3 > 2$
 $8 > 2$
 $8 > 6$