

Problem-3

```
(1) void SimpleVector<T>::push (T &val) {  
    T *naptr;  
    try { naptr = new T [arraySize + 1];  
        } catch (bad_alloc) {  
            memory_err();  
        }  
    for (int count = 0; count < arraySize; count++) {  
        naptr[count] = aptn[count];  
    }  
    naptr[arraySize++] = val;  
    delete [] aptn;  
    aptn = 0;  
    aptn = naptr;  
}
```

Let's consider,

Operations before the for loop, O_b
Operations inside catch statement PO_s
operation inside the for loop, O_c
operation after the for loop O_a

$$O_b + \sum_{count=0}^{arraySize-1} (O_c) + O_a + PO_s ;$$

0 operations which equates $T(0)$ -dox cycle

We know, $\sum_{j=x}^y 1 = (y-x) + 1$

Let consider array size = n

$$\text{So, } O_b + \sum_{\text{count}=0}^{n-1} (O_c) + O_a + POs$$

$$= O_b + ((n-1-0) + 1) O_c + O_a + POs$$

← $f(n)$ - is a first-order

$$= O_b + nO_c + O_a + POs \quad \text{Polynomial}$$

$$= \boxed{c'n + c} \longrightarrow \text{it's } O(n)$$

where $c' = O_c$

$$c = O_b + O_a + POs$$

(2) [with optimized simple Vector using arrays]:

```
void SimpleVector <T>::push (T &val) {  
    if (arraySize == maxSize) {  
        maxSize *= 2;  
        T *naptr;  
        try { naptr = new T[maxSize]  
        } catch (bad_alloc) {  
            memError();  
        }  
        for (int count = 0; count < arraySize; count++)  
        { naptr[count] = aptn[count];  
        }  
        naptr[arraySize++] = val;  
        delete[] aptn;  
        aptn = 0;  
        aptn = naptr;  
    } else {  
        aptn[arraySize++] = val;  
    }  
}
```

let's consider, Operation before if statement, O_f
Operation before the for loop, O_b
operation inside the catch statement, PO_s
operation inside the for loop, O_c
operations after the for loop, O_a
operations inside else statement, PO_e

from the above function, we can write,

$$O_f + O_b + PO_s + \sum_{\text{count}=0}^{\text{arraySize}-1} (O_c) + O_a + PO_e$$

O , operation which equates $T(O)$ - clock cycles

let's consider the array size $= n$

$$O_f + O_b + PO_s + \sum_{\text{count}=0}^{n-1} (O_c) + O_a + PO_e$$

$$\text{We know, } \sum_{j=x}^y 1 = (y-x) + 1$$

$$\text{so, } O_f + O_b + PO_s + ((n-1) + 1) O_c + O_a + PO_e$$

$\hookrightarrow f(n)$ - is a first order polynomial

$$= O_f + O_b + nO_c + PO_s + O_a + PO_e$$

$$= \boxed{c'n + e}$$

where $c' = O_c$

$$e = O_f + O_b + PO_s + O_a + PO_e$$

\downarrow
it is $O(n)$

(3) Simple vector with Linked list:

```
void LinkedList<T>::addList (const T & data) {  
    link *end;  
    temp = front;  
    do { end = temp;  
        temp = temp->linkPtr;  
    } while (temp != NULL);  
    link *add = new Link;  
    add->data = data;  
    add->linkPtr = NULL;  
    end->linkPtr = add;  
}
```

Let's consider, operations before the do-while loop = O_b
operations inside the do-while loop = O_d
operations after the do-while loop = O_a

$$O_b + \sum_{i=0}^{n-1} O_d + O_a ; \quad O, \text{ operations which equates } T(O) - \text{clock cycles}$$

Here, $n \rightarrow$ is the number nodes in the list. In the do-while loop temp start from the beginning of the linked list and traverses each node until it reaches to the end

$$\text{We know, } \sum_{j=x}^y 1 = (y - x) + 1$$

$$\text{So, } O_b + ((n-1-0) + 1) O_d + O_a$$

$$= O_b + n O_d + O_a \rightarrow f(n) \rightarrow \text{is a first order polynomial}$$

$$= \boxed{c'n + c} \quad \text{where, } c' = O_d$$

$$c = O_b + O_a$$

↑
it is $O(n)$