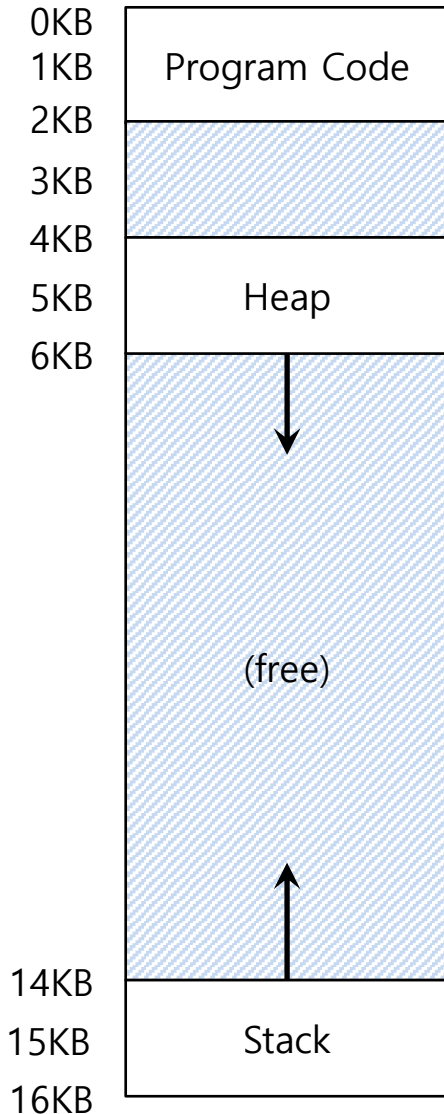


16. Segmentation

Operating System: Three Easy Pieces

Inefficiency of the Base and Bound Approach

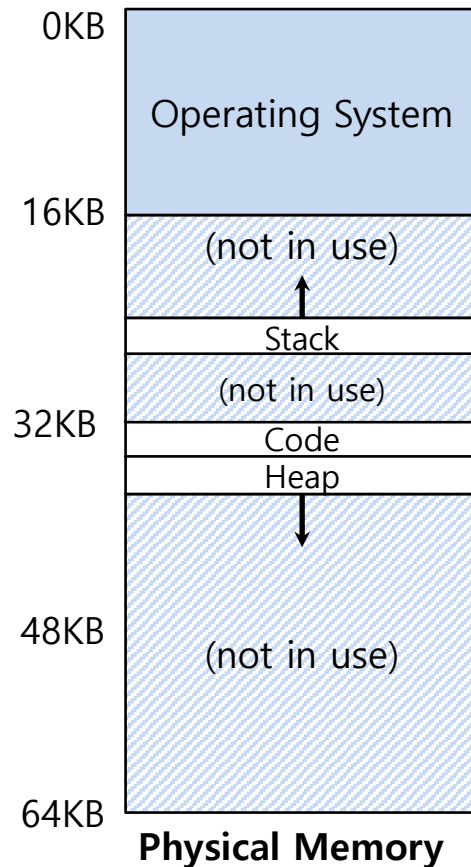


- ▣ **Big chunk of “free” space**
- ▣ “free” space **takes up** physical memory.
- ▣ Hard to run when an address space **does not fit** into physical memory

Segmentation

- ▣ Segment is just **a contiguous portion** of the address space of a particular length.
 - ◆ Logically-different segment: code, stack, heap
- ▣ Each segment can be **placed** in **different part of physical memory**.
 - ◆ **Base** and **bounds** exist **per each segment**.

Placing Segment In Physical Memory

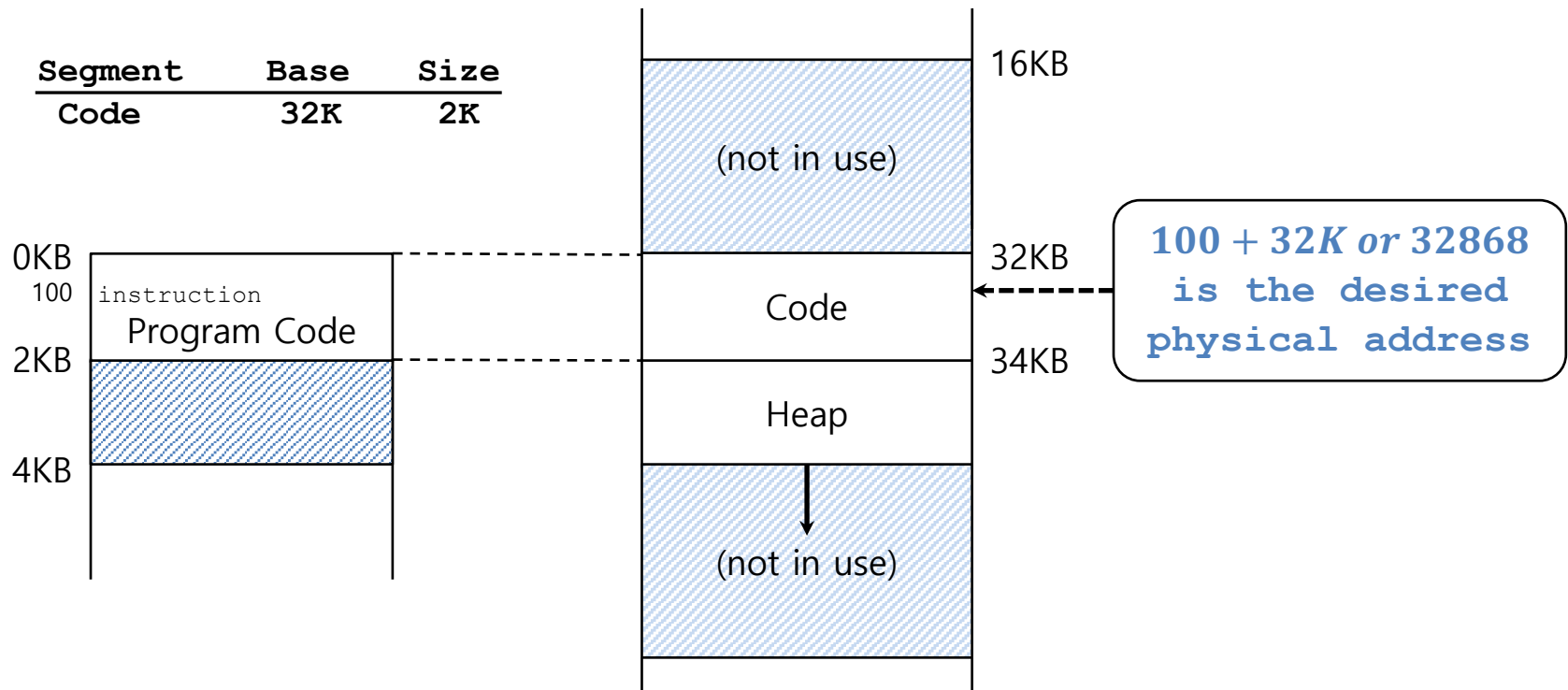


Segment	Base	Size
Code	32K	2K
Heap	34K	2K
Stack	28K	2K

Address Translation on Segmentation

$$\text{physical address} = \text{offset} + \text{base}$$

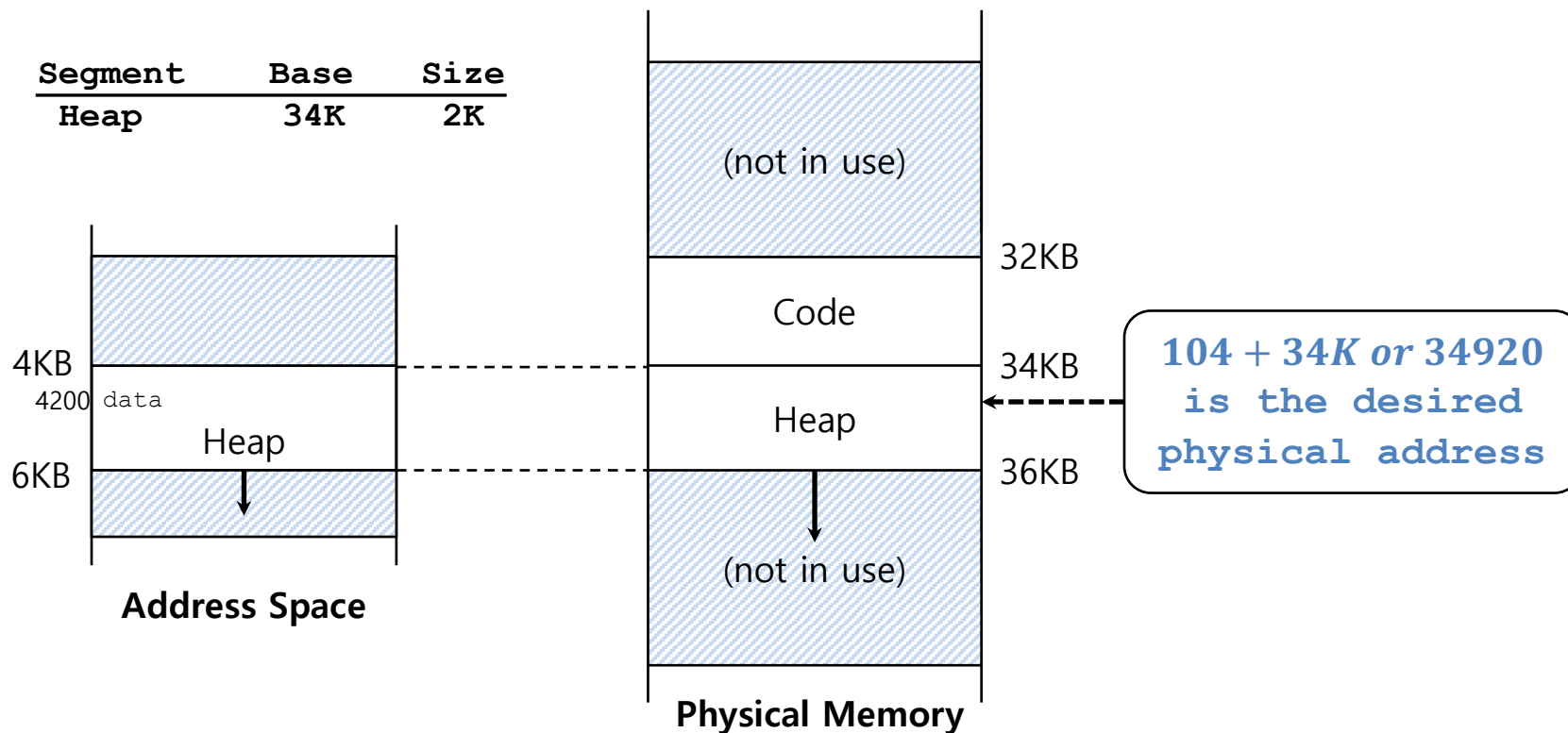
- The `offset` of virtual address 100 is 100.
 - ◆ The code segment **starts at virtual address 0** in address space.



Address Translation on Segmentation(Cont.)

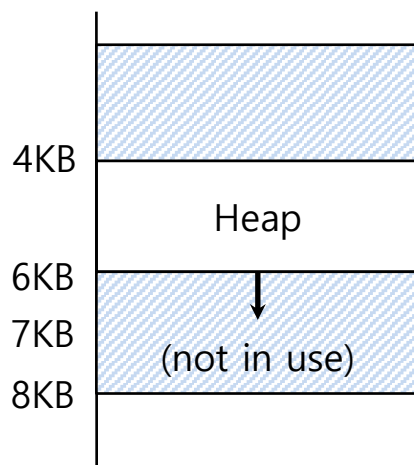
Virtual address + base is not the correct physical address.

- ▣ The `offset` of virtual address 4200 is 104.
 - ◆ The heap segment **starts at virtual address 4096** in address space.



Segmentation Fault or Violation

- If an **illegal address** such as 7KB which is beyond the end of heap is referenced, the OS occurs **segmentation fault**.
 - ◆ The hardware detects that address is **out of bounds**.

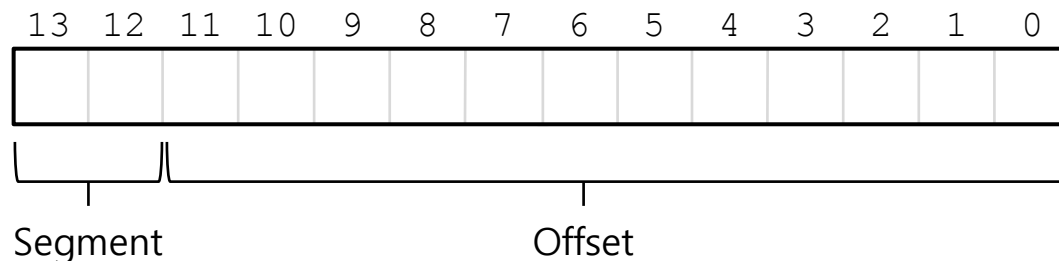


Address Space

Referring to Segment

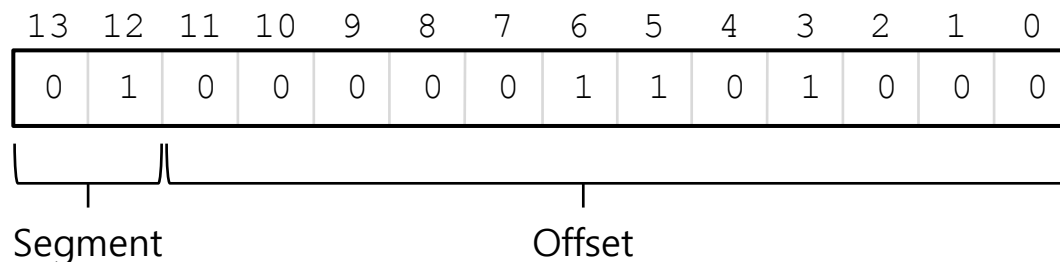
□ Explicit approach

- ◆ Chop up the address space into segments based on the **top few bits** of virtual address.



□ Example: virtual address 4200 (01000001101000)

Segment	bits
Code	00
Heap	01
Stack	10
-	11



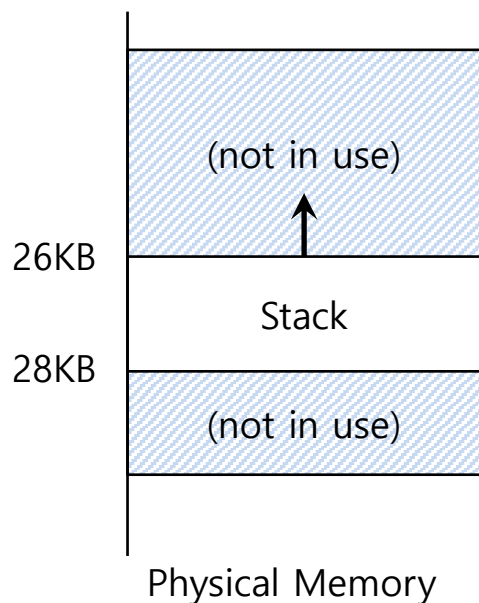
Referring to Segment(Cont.)

```
1  // get top 2 bits of 14-bit VA
2  Segment = (VirtualAddress & SEG_MASK) >> SEG_SHIFT
3  // now get offset
4  Offset = VirtualAddress & OFFSET_MASK
5  if (Offset >= Bounds[Segment])
6      RaiseException(PROTECTION_FAULT)
7  else
8      PhysAddr = Base[Segment] + Offset
9      Register = AccessMemory(PhysAddr)
```

- ◆ `SEG_MASK = 0x3000 (11000000000000)`
- ◆ `SEG_SHIFT = 12`
- ◆ `OFFSET_MASK = 0xFFF (00111111111111)`

Referring to Stack Segment

- ❑ Stack grows **backward**.
- ❑ **Extra hardware support** is need.
 - ◆ The hardware checks which way the segment grows.
 - ◆ 1: positive direction, 0: negative direction



Segment Register(with Negative-Growth Support)

Segment	Base	Size	Grows Positive?
Code	32K	2K	1
Heap	34K	2K	1
Stack	28K	2K	0

Support for Sharing

- ❑ Segment can be **shared between address** space.
 - ◆ **Code sharing** is still in use in systems today.
 - ◆ by extra hardware support.
- ❑ Extra hardware support is need for form of **Protection bits**.
 - ◆ **A few more bits** per segment to indicate **permissions** of **read**, write and **execute**.

Segment Register Values(with Protection)

Segment	Base	Size	Grows Positive?	Protection
Code	32K	2K	1	Read-Execute
Heap	34K	2K	1	Read-Write
Stack	28K	2K	0	Read-Write

Fine-Grained and Coarse-Grained

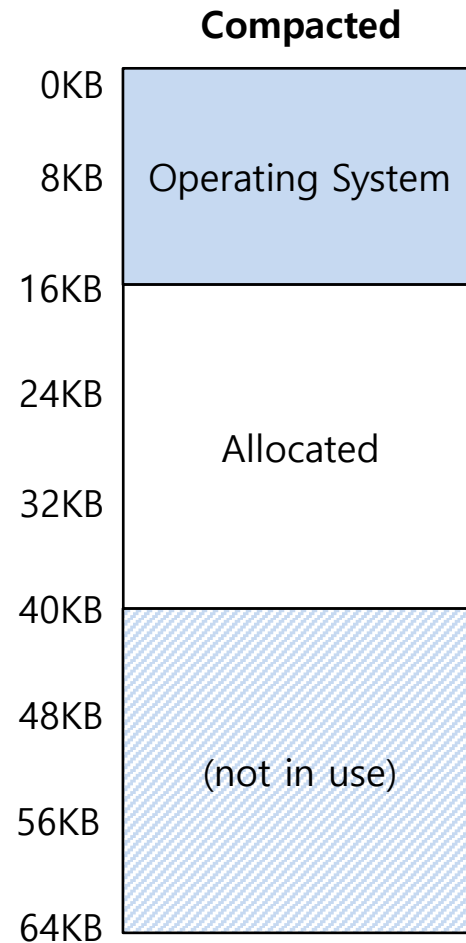
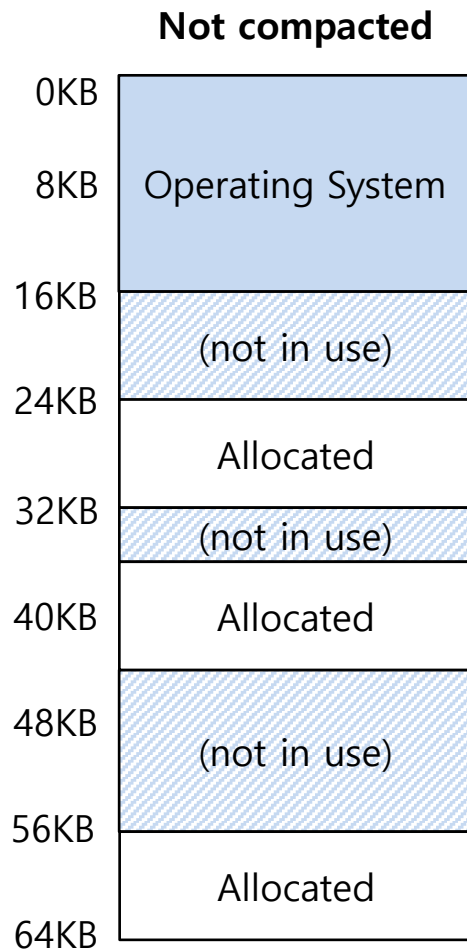
- ▣ **Coarse-Grained** means segmentation in a small number.
 - ◆ e.g., code, heap, stack.
- ▣ **Fine-Grained** segmentation allows **more flexibility** for address space in some early system.
 - ◆ To support many segments, Hardware support with a **segment table** is required.

OS support: Fragmentation

- ❑ **External Fragmentation:** little holes of **free space** in physical memory that make difficulty to allocate new segments.
 - ◆ There is **24KB free**, but **not in one contiguous** segment.
 - ◆ The OS **cannot** satisfy the **20KB request**.

- ❑ **Compaction: rearranging** the exiting segments in physical memory.
 - ◆ Compaction is **costly**.
 - **Stop** running process.
 - **Copy** data to somewhere.
 - **Change** segment register value.

Memory Compaction



- Disclaimer: This lecture slide set was initially developed for Operating System course in Computer Science Dept. at Hanyang University. This lecture slide set is for OSTEP book written by Remzi and Andrea at University of Wisconsin.