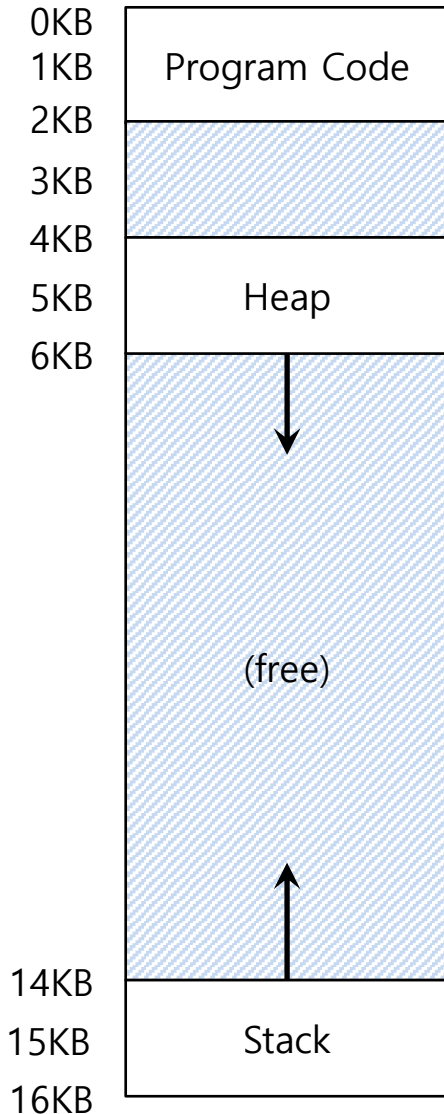


# 16. Segmentation

Operating System: Three Easy Pieces

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# Inefficiency of the Base and Bound Approach

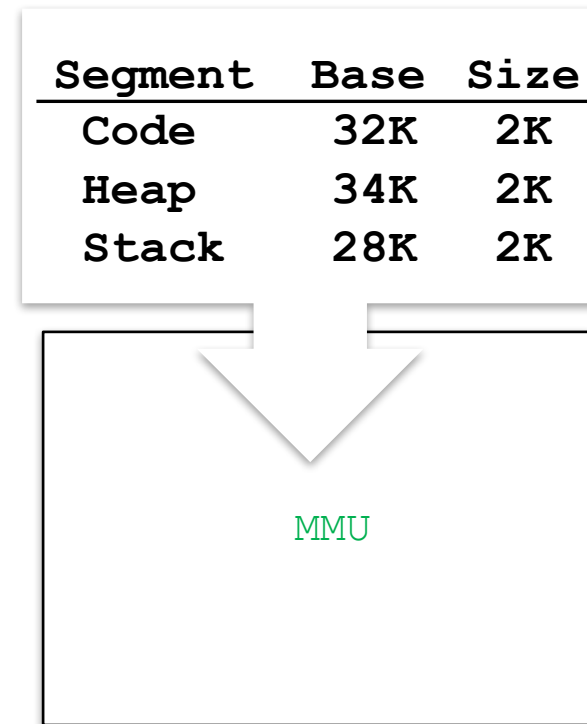
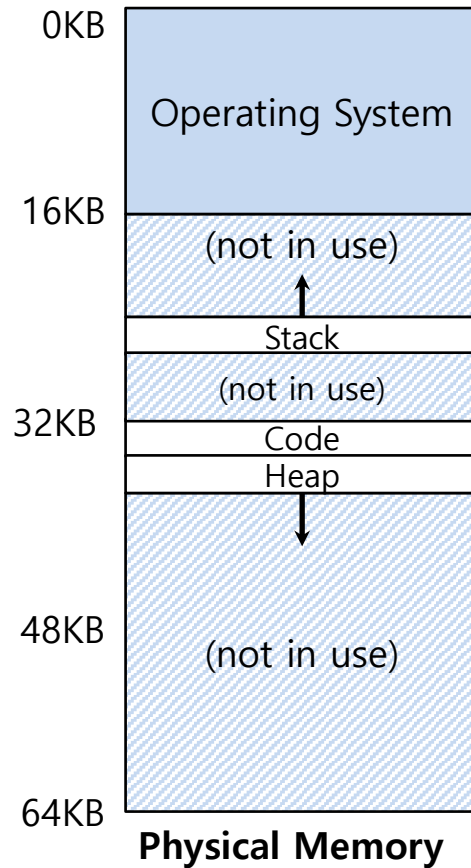


- ▣ **Big chunk of “free” space**
- ▣ “free” space **takes up** physical memory.
- ▣ **Hard** to run when an address space **does not fit** into physical memory

# Segmentation

- ▣ Segment is just a **contiguous portion** of the address space of a particular length.
  - ◆ Logically-different segment: code, stack, heap
- ▣ Each segment can be **placed** in **different part of physical memory**.
  - ◆ **Base** and **bounds** exist **per each segment**.

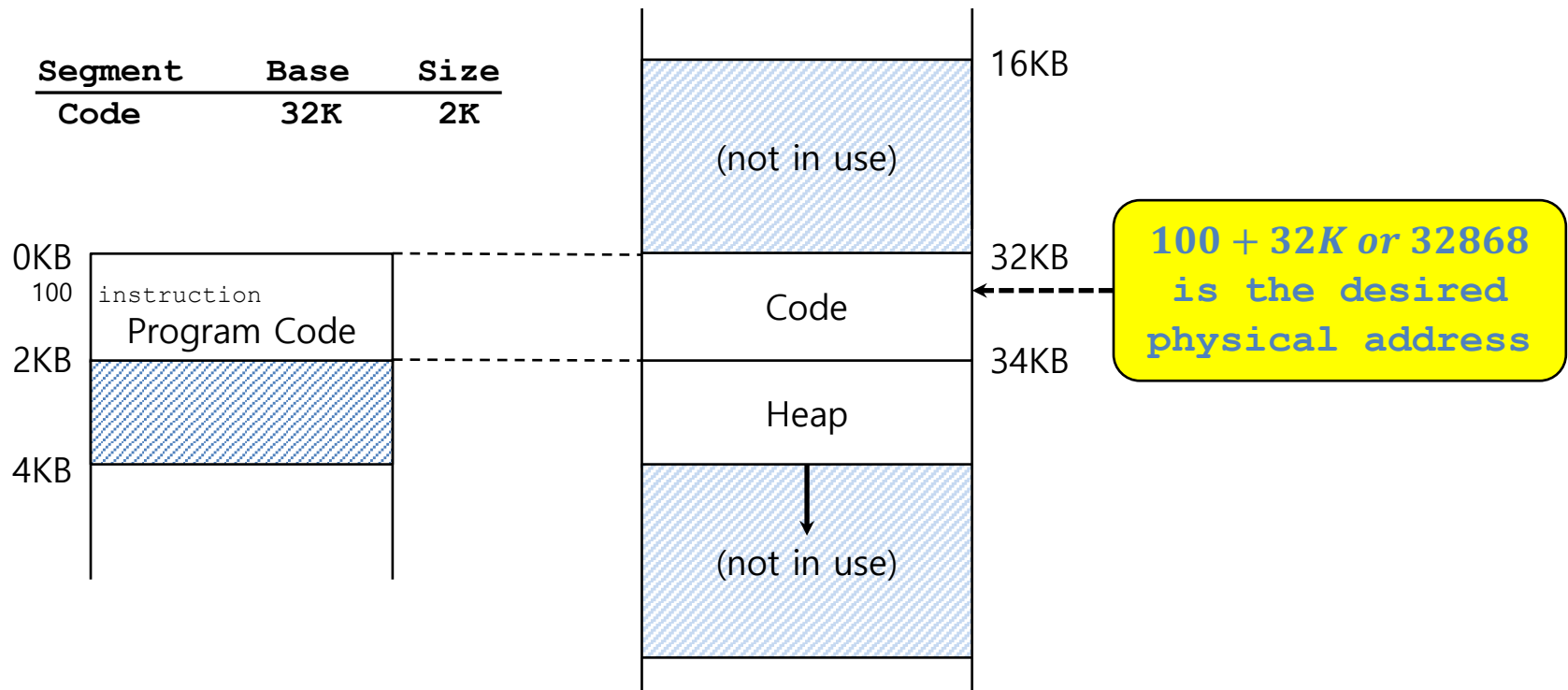
# Placing Segment In Physical Memory



# Address Translation on Segmentation

$$\text{physical address} = \text{offset} + \text{base}$$

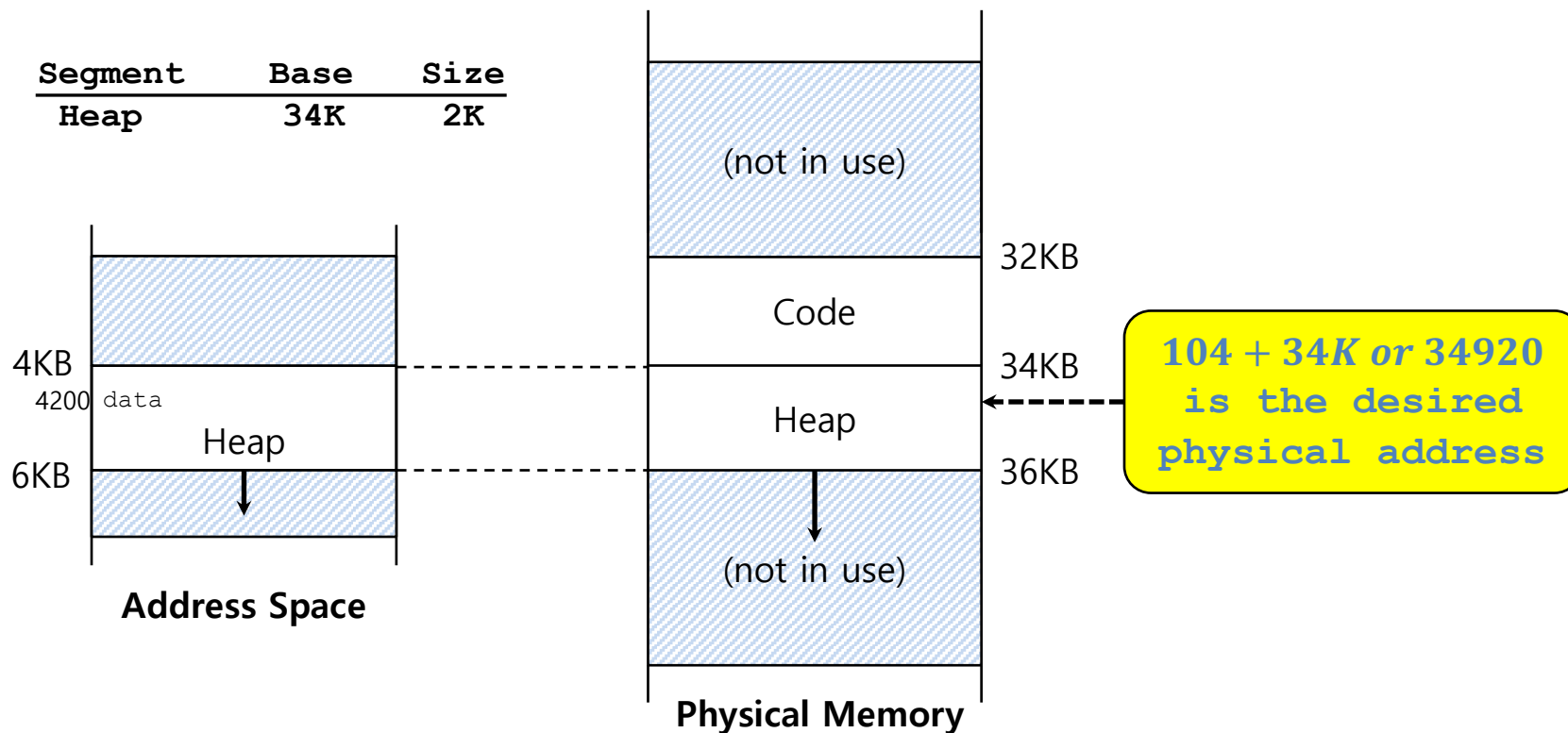
- ▣ The `offset` of virtual address 100 is 100.
  - ◆ The code segment **starts at virtual address 0** in address space.



# Address Translation on Segmentation(Cont.)

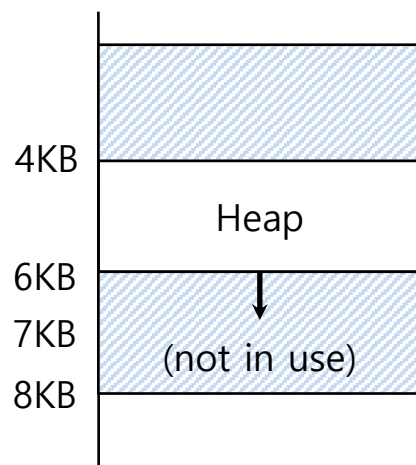
*Virtual address + base is not the correct physical address!*

- ▣ The `offset` of virtual address 4200 is 104.
  - ◆ The heap segment **starts at virtual address 4096** in address space.



# Segmentation Fault or Violation

- If an **illegal address** such as 7KB which is beyond the end of heap is referenced, the OS occurs **segmentation fault**.
  - ◆ The hardware detects that address is **out of bounds**.

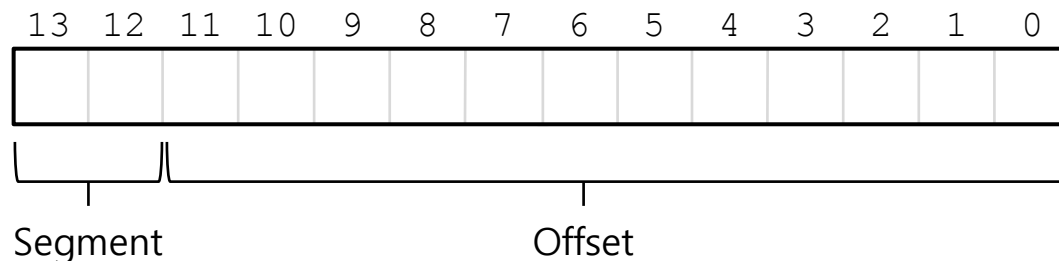


**Address Space**

# Referring to Segment

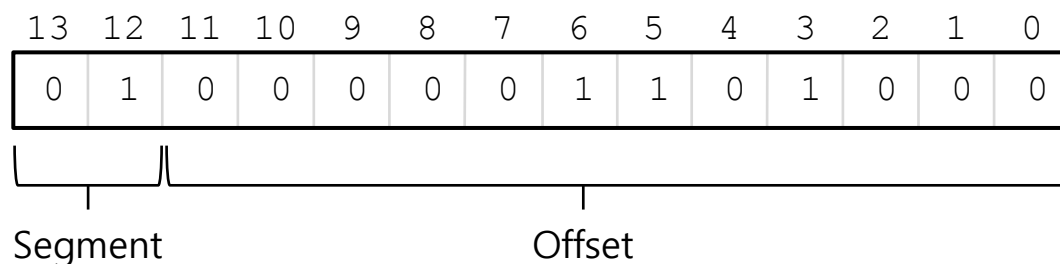
## □ Explicit approach

- ◆ Chop up the address space into segments based on the **top few bits** of virtual address.



## □ Example: virtual address 4200 (01000001101000)

| Segment | bits |
|---------|------|
| Code    | 00   |
| Heap    | 01   |
| Stack   | 10   |
| -       | 11   |





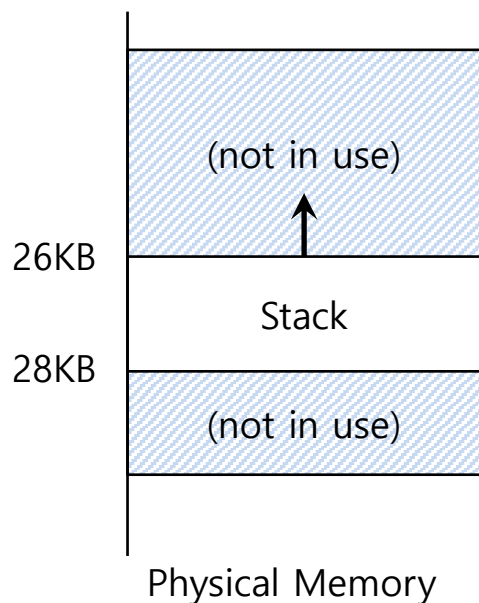
# Referring to Segment(Cont.)

```
1  // get top 2 bits of 14-bit VA
2  Segment = (VirtualAddress & SEG_MASK) >> SEG_SHIFT
3  // now get offset
4  Offset = VirtualAddress & OFFSET_MASK
5  if (Offset >= Bounds[Segment])
6      RaiseException(PROTECTION_FAULT)
7  else
8      PhysAddr = Base[Segment] + Offset
9      Register = AccessMemory(PhysAddr)
```

- ◆ SEG\_MASK = 0x3000 (11000000000000)
- ◆ SEG\_SHIFT = 12
- ◆ OFFSET\_MASK = 0xFFF (00111111111111)

# Referring to Stack Segment

- ❑ Stack grows **backward**.
- ❑ **Extra hardware support** is need.
  - ◆ The hardware checks which way the segment grows.
  - ◆ 1: positive direction, 0: negative direction



Segment Register(with Negative-Growth Support)

| Segment | Base | Size | Grows Positive? |
|---------|------|------|-----------------|
| Code    | 32K  | 2K   | 1               |
| Heap    | 34K  | 2K   | 1               |
| Stack   | 28K  | 2K   | 0               |

# Support for Sharing

- ❑ Segment can be **shared between address** space.
  - ◆ **Code sharing** is still in use in systems today.
  - ◆ by extra hardware support.
- ❑ Extra hardware support is need for form of **Protection bits**.
  - ◆ **A few more bits** per segment to indicate **permissions** of **read**, write and **execute**.

Segment Register Values(with Protection)

| Segment | Base | Size | Grows Positive? | Protection   |
|---------|------|------|-----------------|--------------|
| Code    | 32K  | 2K   | 1               | Read-Execute |
| Heap    | 34K  | 2K   | 1               | Read-Write   |
| Stack   | 28K  | 2K   | 0               | Read-Write   |

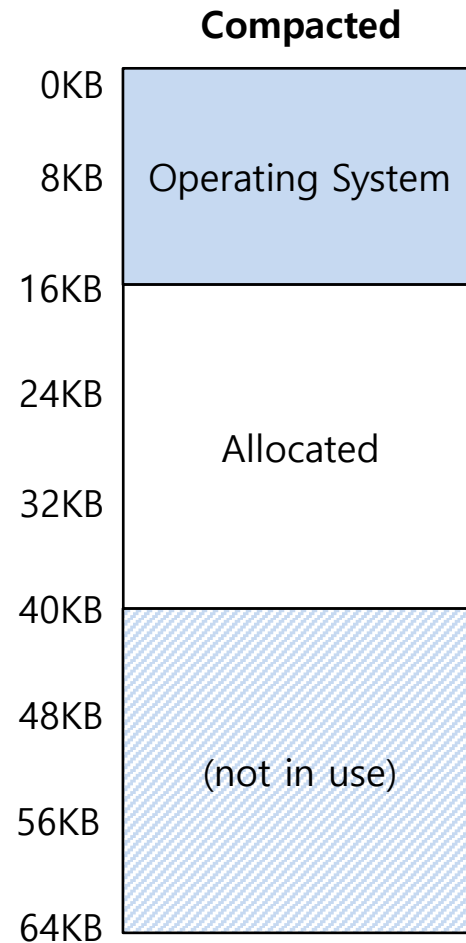
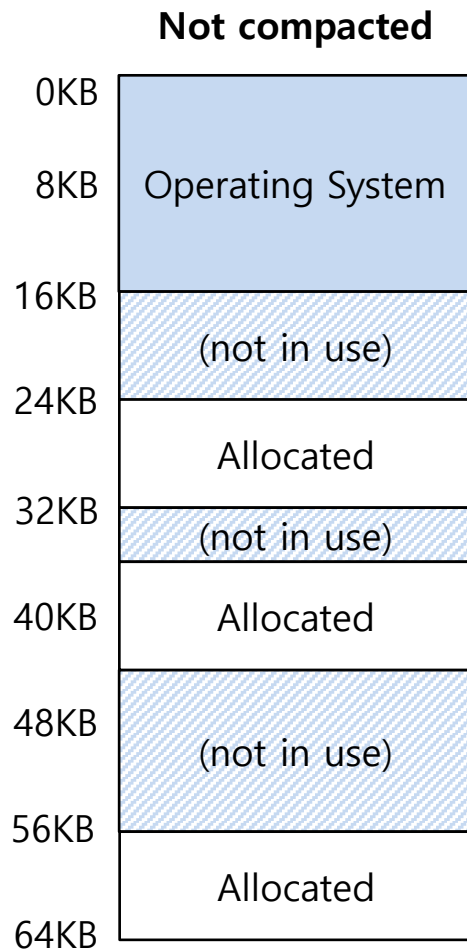
# Fine-Grained and Coarse-Grained

- ▣ **Coarse-Grained** means segmentation in a small number.
  - ◆ e.g., code, heap, stack.
- ▣ **Fine-Grained** segmentation allows **more flexibility** for address space in some early system.
  - ◆ To support many segments, Hardware support with a **segment table** is required.

# OS support: Fragmentation

- ▣ **External Fragmentation:** little holes of **free space** in physical memory that make difficulty to allocate new segments.
  - ◆ There is **24KB free**, but **not in one contiguous** segment.
  - ◆ The OS **cannot** satisfy the **20KB request**.
  
- ▣ **Compaction: rearranging** the exiting segments in physical memory.
  - ◆ Compaction is **costly**.
    - **Stop** running process.
    - **Copy** data to somewhere.
    - **Change** segment register value.
  
- ▣ 1000 ways to solve it
  - ◆ None of them are the "best"
  
- ▣ *Added to creating, terminating, and context switches*

# Memory Compaction



- Disclaimer: Disclaimer: This lecture slide set is used in AOS course at University of Cantabria. Was initially developed for Operating System course in Computer Science Dept. at Hanyang University. This lecture slide set is for OSTEP book written by Remzi and Andrea Arpaci-Dusseau (at University of Wisconsin)