# **Design Document**

### Implementation of futures and using it for synchronization:

This project is an implementation of the structure called future in C with the underlying operating system XINU and then using it to synchronize producer and consumer threads. The code was run successfully with XINU OS deployed on RPI node 21.

#### **Introduction to Futures:**

Future is a structure used for implementing synchronization in various programming languages. Future objects are expected to generate values in the future. They offer functions get() and set(). The difference between a thread and a future is that we can return values in futures. The get function is used to receive values from the thread.

## **Implementation of futures**

#### 1. Structure of a Future:

Futures have been implemented using semaphores. It is a struct which consists of a variable which contains the shared value between the producer and consumer. The other variable flag will contain the ID of the semaphore created for the future.

2. Initialization of Future: *future\_alloc()* 

In the initialization, space is allocated to the object of the 'future' struct. It creates a semaphore with count 0 (it means that the future is empty).

3. get() function: future\_get()

This function is invoked by the consumer to return the value produced:

- i. If the consumer tries to consume on a full future, then it consumes the value and make it empty.
- ii. If the consumer tries to consume on an empty future, then it blocks and waits for a signal to continue.
- iii. If the consumer tries to consume on an empty future on which another consumer thread is waiting, it returns an error.

This function checks if the semaphore has a count less than 0. If yes, it means that there is already a thread waiting to consume the value. Hence it returns an error.

Then the function calls wait() function of the semaphore where Count 1 signifies full future and count 0 signifies empty future. So, wait() on count 1 will yield value whereas that on 0 will block for a signal.

4. Set() function: *future\_set()* 

This function is invoked by the producer to produce a value:

- i. If the producer tries to produce on an empty future, then it produces the value and make it full.
- ii. If the producer tries to produce on a full future, then it returns an error.

The function checks if the semaphore is greater than 0. If yes, then it means that there is already a value produced. So, another attempt to produce will return an error. If the value is 0, then it will produce a value. So, we are maintaining the future as 'full' or 'empty' based on the value of the semaphore count.

### 5. States of the future:

The states of the future can be identified by the count of the semaphore:

Value of the semaphore count:

- i. 1 means Full future
- ii. 0 means empty future
- iii. Less than 0 means there is a consumer thread waiting on the semaphore
- 6. The synchronization between producer-consumer is achieved by creating threads for producer and consumer and sending the reference of a common future to all the threads in which synchronization is needed. This future is created and initialized with *future\_alloc()* with its semaphore set to count 0 and value initially empty. Then, the producer and consumer threads call future\_set() and future\_get() respectively and the future ensures that both are synchronized as per the rules mentioned in 3.i, 3.ii, 3.iii and 4.i, 4.ii.