

## Computer Architecture

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# Laboratory Session 3

## I. Bitwise Logic and Intro. to Procedure (70pts)

### 1. Exercise 1: (35pts) Write a program that

**1.1** Put the number 0xDEADBEEF into register \$t1 without using pseudoinstruction li. (lab3\_1\_1.s)

```
.text
.globl main
main:
ori $t1, $zero, 0xDEAD # $t1 = 0x0000DEAD
sll $t1, $t1, 16 # $t1 = 0xDEAD0000
ori $t1, $t1, 0xBEEF # $t1 = 0xDEADBEEF

jr $ra # Exit
```

**1.2** Redo 1.1 as follows: use ori to load each letter into register. (lab3\_1\_2.s)

```
.text
.globl main
main:
ori $t1, $zero, 0xD # $t1 = 0x0000000D
sll $t1, $t1, 4
ori $t1, $t1, 0xE # $t1 = 0x000000DE
sll $t1, $t1, 4
ori $t1, $t1, 0xA # $t1 = 0x00000DEA
sll $t1, $t1, 4
ori $t1, $t1, 0xD # $t1 = 0x0000DEAD
sll $t1, $t1, 4
ori $t1, $t1, 0xB # $t1 = 0x000DEADB
sll $t1, $t1, 4
ori $t1, $t1, 0xE # $t1 = 0x00DEADBE
sll $t1, $t1, 4
ori $t1, $t1, 0xE # $t1 = 0x0DEADBEE
sll $t1, $t1, 4
ori $t1, $t1, 0xF # $t1 = 0xDEADBEEF

ori $v0, $zero, 10 # EXIT
syscall
```

**1.3** Suppose that  $\$t1 = 0xDEADBEEF$ . Using only register-to-register logic and shift instructions, Reverse the order of the bytes in  $\$t1$  so that register  $\$t2$  get the bit pattern  $0xFEEBDAED$ . (lab3\_1\_3.s)

```
.text
.globl main
main:
ori $t1, $zero, 0xDEAD # $t1 = 0x0000DEAD
sll $t1, $t1, 16 # $t1 = 0xDEAD0000
ori $t1, $t1, 0xBEEF # $t1 = 0xDEADBEEF

# I didn't know the we are allowed to use add/addi to make Loop.
# So I try to think of a workaround with nor and sll instructions.
nor $a0, $a0, $zero # set i = 8
jal REVERSE # jump to REVERSE and save position to $ra
```

```
j EXIT # jump to EXIT
```

EXIT:

```
ori $v0, $zero, 10 # Exit  
syscall
```

REVERSE:

```
andi $t3, $t1, 0xF  
or $t2, $t2, $t3 # Add the least significant word from $t1 to $t2
```

```
sll $a0, $a0, 4 # i -= 1  
beq $a0, $zero, JUMP_BACK # if i == 0 then JUMP_BACK
```

```
sll $t2, $t2, 4 # Shift $t2 to the left  
srl $t1, $t1, 4 # Shift $t1 to the right  
j REVERSE # jump to REVERSE
```

JUMP\_BACK:

```
jr $ra # jump to $ra
```

1.4 Redo 1.3 using only and, or, and rotate instructions. (lab3\_1\_4.s)

```
.text
```

```
.globl main
```

```
main:
```

```
# Suppose $t1
```

```
ori $t1, $zero, 0xDEAD # $t1 = 0x0000DEAD
```

```
sll $t1, $t1, 16 # $t1 = 0xDEAD0000
```

```
ori $t1, $t1, 0xBEEF # $t1 = 0xDEADBEEF
```

```
# Only and, or and rotate instructions.
```

```
# Therefore, I didn't know the we are allowed to use add/addi to  
make Loop.
```

```
# So I try to think of a workaround with nor and sll instructions.
```

```
nor $a0, $a0, $zero # set i = 8
```

```
jal REVERSE # jump to REVERSE and save position to $ra
```

```
j EXIT # jump to EXIT
```

**REVERSE:**

```
andi $t3, $t1, 0xF
or $t2, $t2, $t3 # Add the least significant word from $t1 to $t2
```

```
sll $a0, $a0, 4 # i -= 1
beq $a0, $zero, JUMP_BACK # if i == 0 then JUMP_BACK
```

```
ror $t1, $t1, 4 # Shift $t2 to the left
rol $t2, $t2, 4 # Shift $t1 to the right
j REVERSE # jump to REVERSE
```

**JUMP\_BACK:**

```
jr $ra # jump to $ra
```

**EXIT:**

```
ori $v0, $zero, 10 # Exit
syscall
```

## 2. Exercise 2: (15pts) Write a program that

**2.1** Set the corresponding bit in register \$t1 through \$t8. That is, in register \$t1 set bit 1, register \$t2 set bit 2, and so on. (lab3\_2\_1.s)

```
.text
```

```
.globl main
```

```
main:
```

```
ori $t0, $t0, 1
```

```
sll $t1, $t0, 1
```

```
sll $t2, $t1, 1
```

```
sll $t3, $t2, 1
```

```
sll $t4, $t3, 1
```

```
sll $t5, $t4, 1
```

```
sll $t6, $t5, 1
```

```
sll $t7, $t6, 1
```

```
sll $t8, $t7, 1
```

```
jr $ra # EXIT
```

**2.2** By using ONLY shift instructions and register to register logic instructions (no li pseudoinstruction or addi), put the pattern 0xFFFFFFFF into register \$t1. (lab3\_2\_2.s)

```
.text
.globl main
main:
nor $t1, $zero, $zero

jr $ra # Exit
```

## II. MSP430 (30pts)

Step 1: build the sample code in CCS, check the errors.

Step 2: Not run, the values of these registers (PORT\_1\_2):

P1OUT: **0xBE**  
P1IN: **0x06**  
P1DIR: **0x00**  
P1REN: **0x00**  
P1IFG: **0x00**

Step 3: Run, observe and collect the values of these registers in case of

	Red LED ON	Green LED on
P1OUT	0xBF = 10111111	0xFE = 11111110
P1IN	0x0F = 00001111	0x4E = 01001110
P1DIR	0x41 = 01000001	0x41 = 01000001
P1REN	0x08 = 00001000	0x08 = 00001000
P1IFG	0xF9 = 11111001	0xF9 = 11111001

Comment and explain the Table above:

From the table above, It is clear that P1OUT, P1IN change while P1DIR, P1REN and P1IFG stay unchanged. P1OUT, P1IN change because they instruct LEDs to switch ON or OFF when the button is pressed.

C Code	MIPS Code
<code>void main(void) {     WDTCTL = WDTPW   WDTOLD;</code>	<code>c000: 40B2 5A80 0120   MOV.W       #0x5a80,&amp;Watchdog_Timer_WDTCTL</code>

<pre> // stop watchdog timer  P1OUT  = Red; P1OUT &amp;= ~Green; P1DIR  = Red +Green;  P1DIR &amp;= ~Button; P1REN  = Button; P1OUT  = Button;  volatile unsigned int i; // volatile to prevent optimization  while(1) {     if ((P1IN &amp; Button)!= Button)     {         while ((P1IN &amp; Button)!= Button)         {             P1OUT ^= Red + Green;         }     } } </pre>	<pre> 10    P1OUT  = Red; c006: D3D2 0021    BIS.B #1,&amp;Port_1_2_P1OUT 11    P1OUT &amp;= ~Green; c00a: F0F2 00BF 0021    AND.B #0x00bf,&amp;Port_1_2_P1OUT 12    P1DIR  = Red +Green; c010: D0F2 0041 0022    BIS.B #0x0041,&amp;Port_1_2_P1DIR 14    P1DIR &amp;= ~Button; c016: C2F2 0022    BIC.B #8,&amp;Port_1_2_P1DIR 15    P1REN  = Button; c01a: D2F2 0027    BIS.B #8,&amp;Port_1_2_P1REN 16    P1OUT  = Button; c01e: D2F2 0021    BIS.B #8,&amp;Port_1_2_P1OUT 22    if ((P1IN &amp; Button)!= Button)         \$C\$L1: c022: B2F2 0020    BIT.B #8,&amp;Port_1_2_P1IN c026: 23FD    JNE    (\$C\$L1) 24    while ((P1IN &amp; Button)!= Button)         \$C\$L2: c028: B2F2 0020    BIT.B #8,&amp;Port_1_2_P1IN c02c: 27FD    JEQ    (\$C\$L2) 28    P1OUT ^= Red + Green; c02e: E0F2 0041 0021    XOR.B #0x0041,&amp;Port_1_2_P1OUT c034: 3FF6    JMP    (\$C\$L1) 85    {         _c_int00_noinit_noargs_noexit(): c036: 4031 0400    MOV.W #0x0400,SP 87    _system_pre_init(); c03a: 12B0 C056    CALL #_system_pre_init 88    main(0); c03e: 430C    CLR.W    R12 c040: 12B0 C000    CALL    #main 89    abort(); c044: 12B0 C050    CALL    #abort 48    BIS.W    #(0x0010),SR         \$isr_trap.asm:48:59\$( ),         __TI_ISR_TRAP(): c048: D032 0010    BIS.W #0x0010,SR </pre>
--	---

	<pre> 49      JMP __TI_ISR_TRAP c04c: 3FFD      JMP (\$isr_trap.asm:48:59\$) 51      NOP      ; CPU40 Compatibility NOP c04e: 4303      NOP 100 {     C\$\$EXIT(), abort(): c050: 4303      NOP 108 for (;;) /* SPINS FOREVER */     C\$L1: c052: 3FFF      JMP (\$C\$L1) c054: 4303      NOP 58      return 1;     _system_pre_init(): c056: 431C      MOV.W #1,R12 c058: 4130      RET </pre>
--	---

### Explain:

From line 8-16, we set up the variables for LEDs and Button. **"P1OUT |= Red;"** Red LED is On and **"P1OUT &= ~Green;"** Green LED is Off.

Then we enter the while loop **"While (1)"** for checking and updating every button pressed.

**"If (P1IN & Button) != Button"** detects when the button is pressed.

**"while ((P1IN & Button) != Button)"** waits for the button to be release then switches the LEDs between Red and Green **"P1OUT ^= Red + Green;"**.

Problem 1: modify the sample code in order to when pressing the button two LEDs turn on and vice versa.

```
1 #include <msp430.h>
2
3 #define Red BIT0
4 #define Green BIT6
5 #define Button BIT3
6
7 void main(void) {
8     WDTCTL = WDTPW | WDTHOLD;          // stop watchdog timer
9
10    P1OUT |= Red + Green;
11    // P1OUT &= ~Green;
12    P1DIR |= Red + Green;
13
14    P1DIR &= ~Button;
15    P1REN |= Button;
16    P1OUT |= Button;
17
18    volatile unsigned int i;            // volatile to prevent optimization
19
20    while(1)
21    {
22        if ((P1IN & Button) != Button)
23        {
24            while ((P1IN & Button) != Button)
25            {
26            }
27            P1OUT ^= Red + Green;
28        }
29    }
30 }
31 }
32
```