

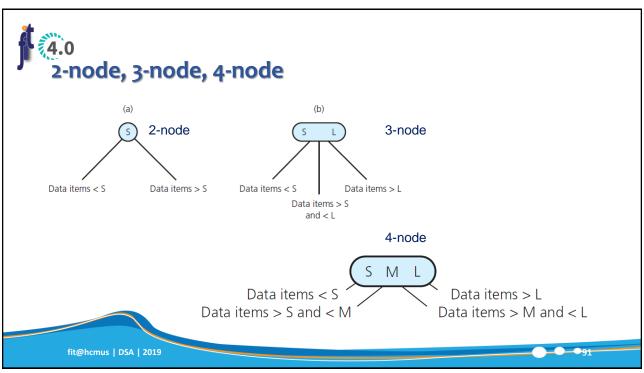


2-node, 3-node, 4-node

- A 2-node (has two children) must contain single data item greater than left child's item(s) and less than right child's item(s).
- A 3-node (has three children) must contain two data items, S and L, such that
 - S is greater than left child's item(s) and less than middle child's item(s);
 - L is greater than middle child's item(s) and less than right child's item(s).
- A 4-node (has our children) must contain three data items S, M, and L that satisfy:
 - S is greater than left child's item(s) and less than middle-left child's item(s)
 - M is greater than middle-left child's item(s) and less than middle-right child's item(s);
 - L is greater than middle-right child's item(s) and less than right child's item(s).

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- 2–3 trees were invented by John Hopcroft in 1970.
- 2-3 tree is a tree in which
 - Every internal node is either a 2-node or a 3-node.
 - Leaves have no children and may contain either one or two data items.

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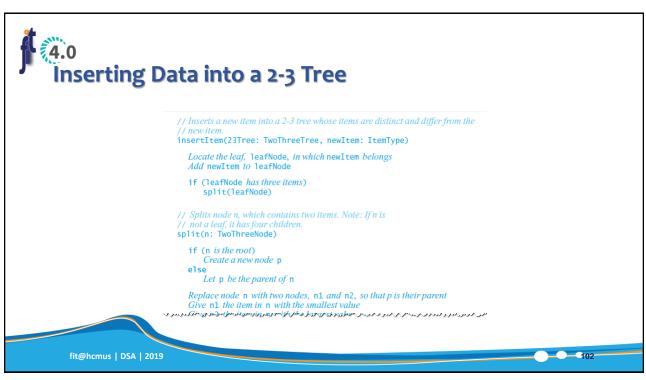


- 2-3 tree is a tree in which
 - Every internal node is a 2-node, a 3-node or a 4-node.
 - Leaves have no children and may contain either one, two or three data items.

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```
Inserting Data into a 2-3 Tree

split(n: TwoThreeNode)

if (n is the root)

Create a new node p

else

Let p be the parent of n

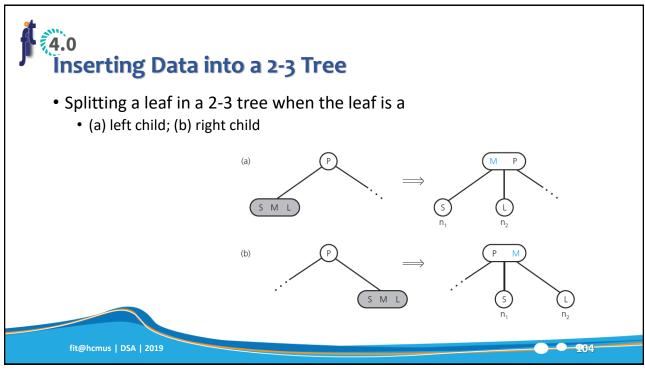
Replace node n with two nodes, n1 and n2, so that p is their parent
Give n1 the ttem in n with the smallest value
Give n2 the ttem in n with the largest value

if (n is not a leaf)

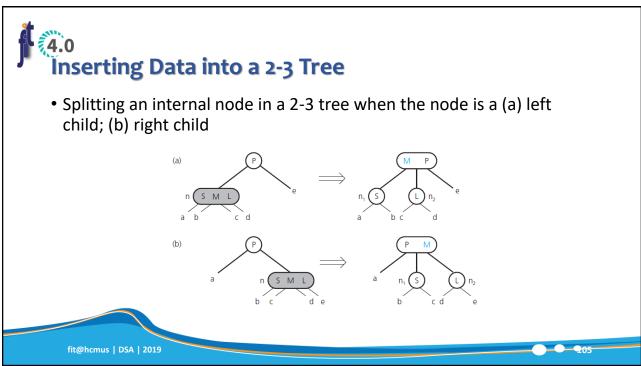
{
    n1 becomes the parent of n's two leftmost children
    n2 becomes the parent of n's two rightmost children
}
Move the item in n that has the middle value up to p

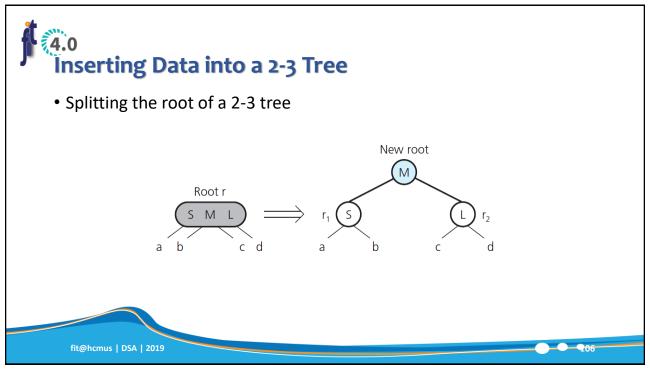
if (p now has three items)

split(p)
```

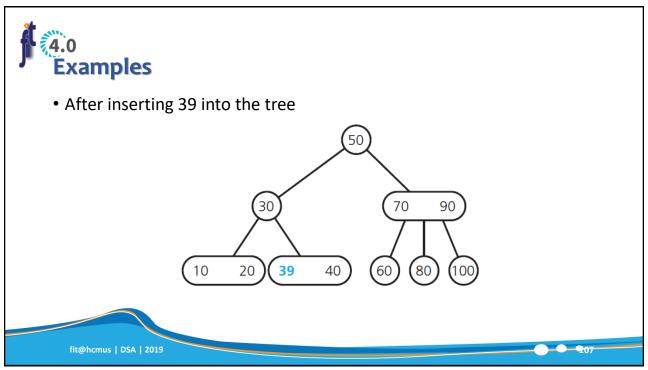


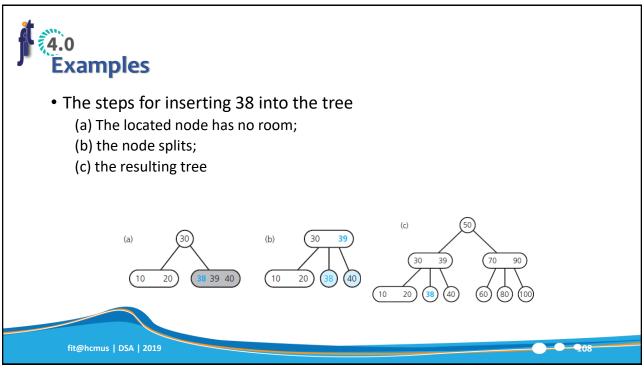




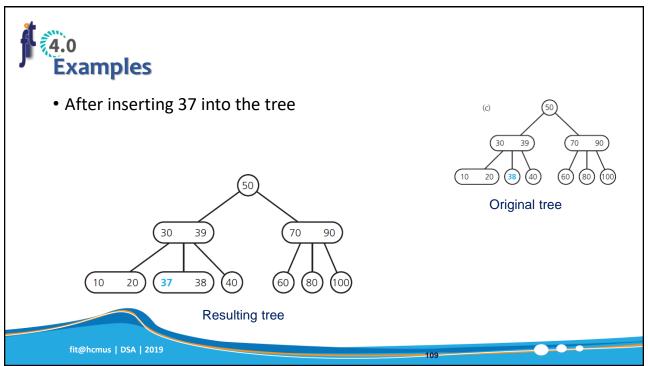


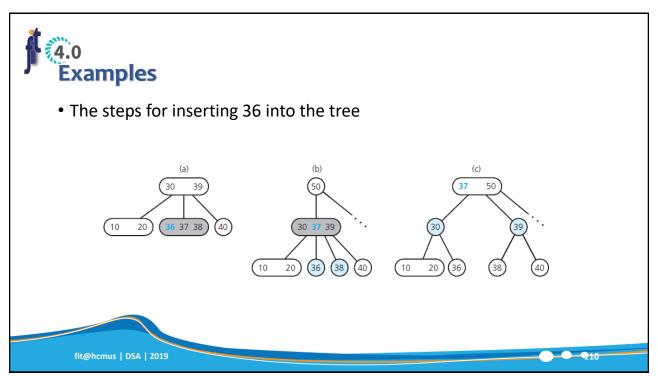




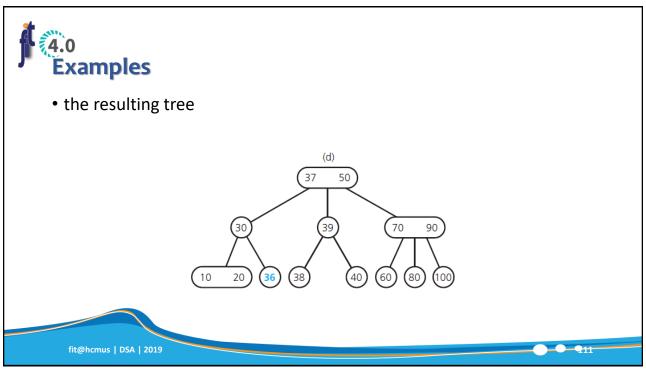


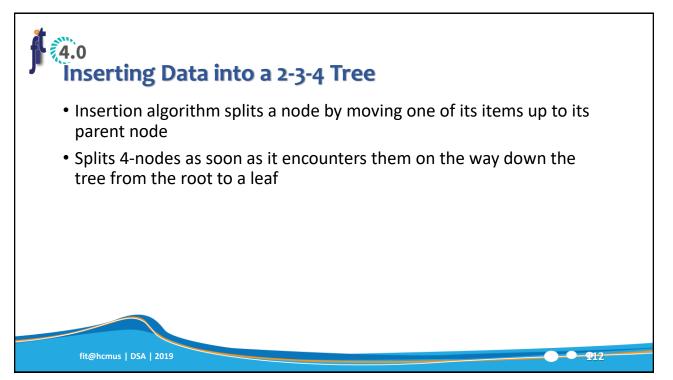




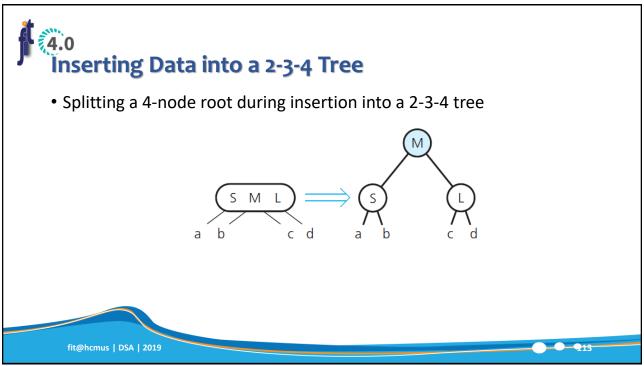


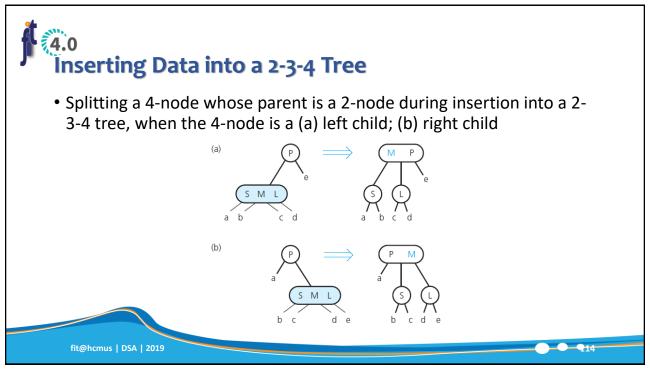




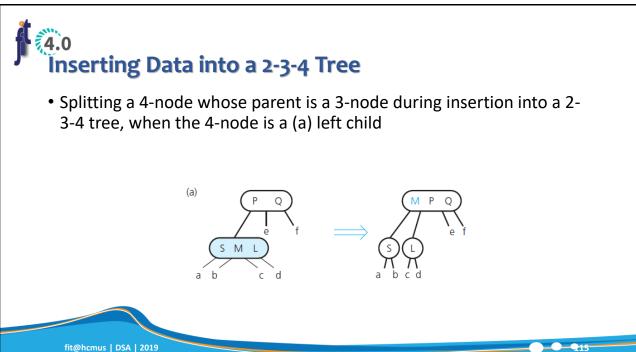


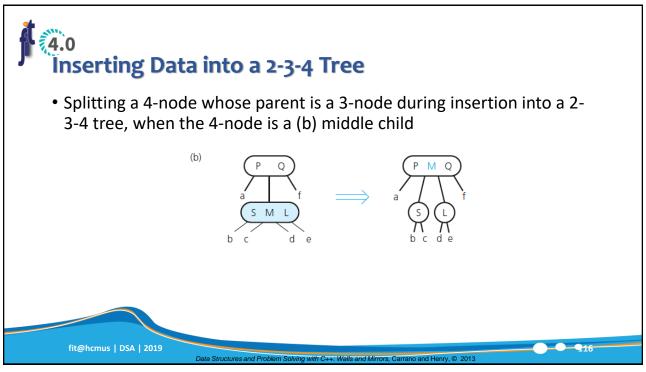




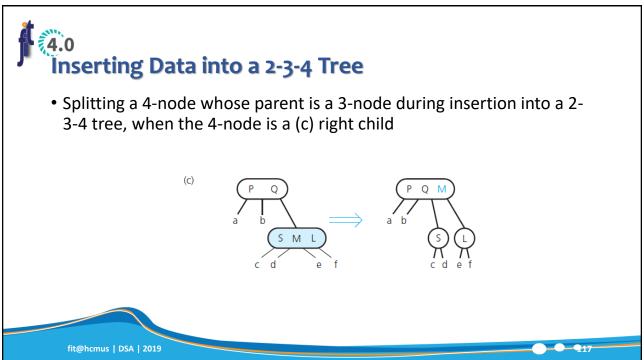


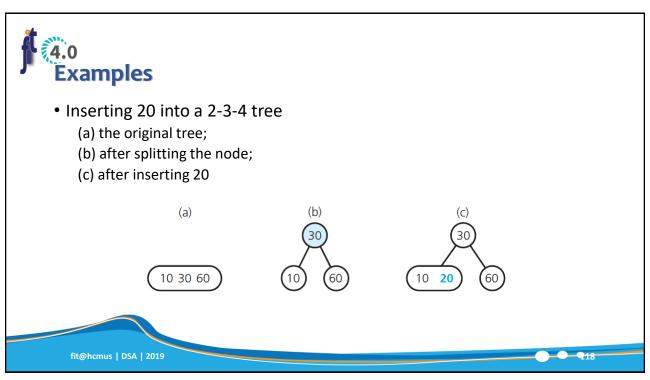




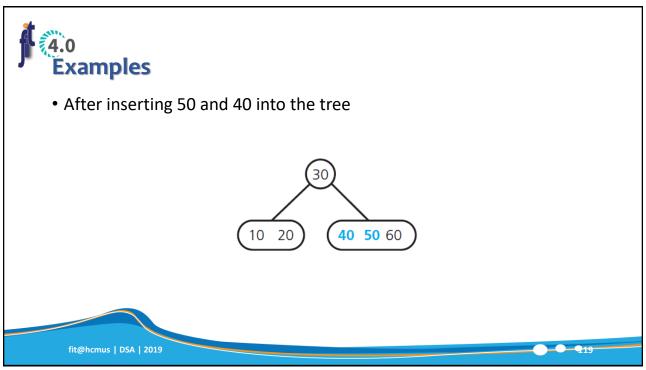


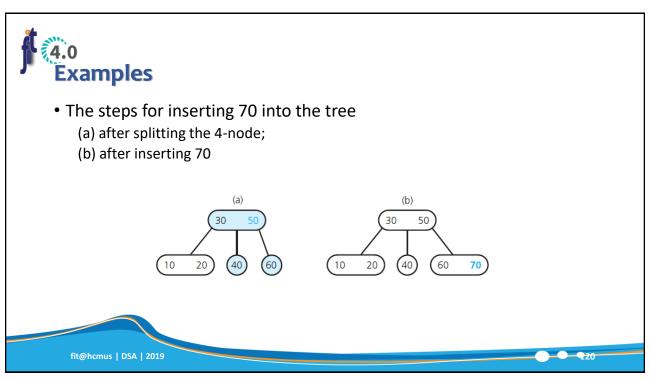




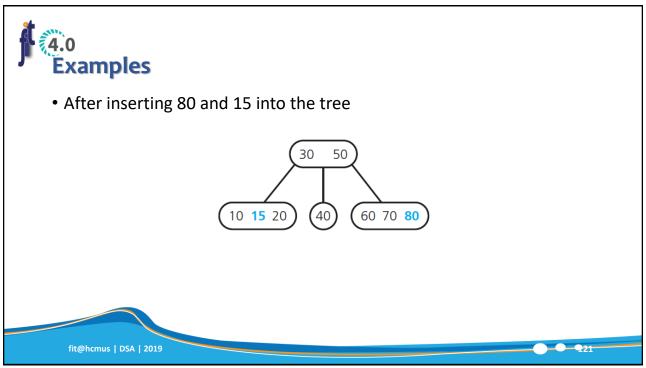


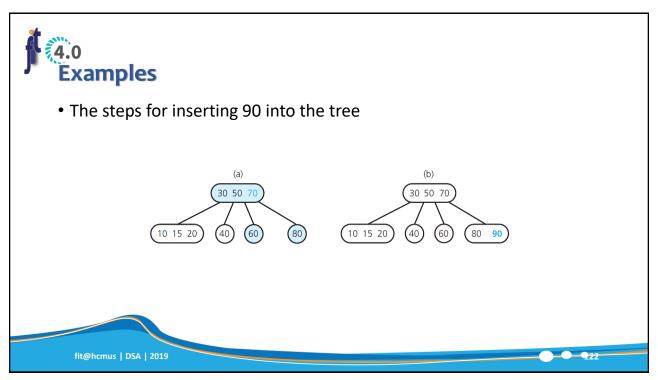




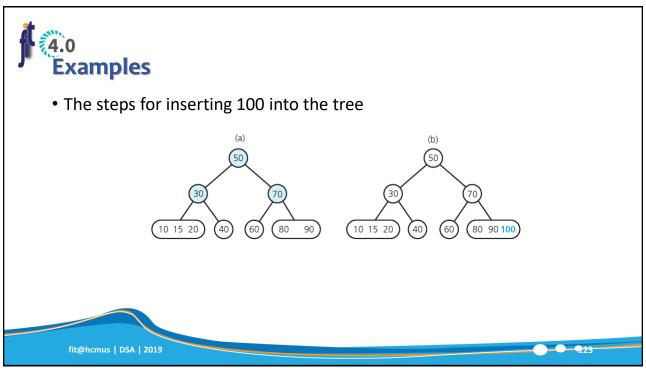
















```
Removing Data from a 2-3 Tree

// Removes the given data item from a 2-3 tree. Returns true if successful
// or false if no such item exists.
removeItem(23Tree: TwoThreeTree, dataItem: ItemType): boolean

Attempt to locate dataItem
if (dataItem is found)
{
    if (dataItem is not in a leaf)
        Swap dataItem with its morder successor, which will be in a leaf leafNode
        // The removal always begins at a leaf
        Remove dataItem from leaf leafNode
    if (leafNode now has no items)
        fixTree(leafNode)
        return true
    }
    else
        return false

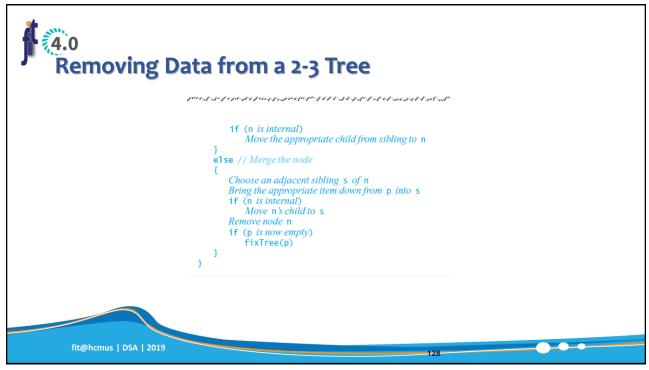
// Company of the property of the proper
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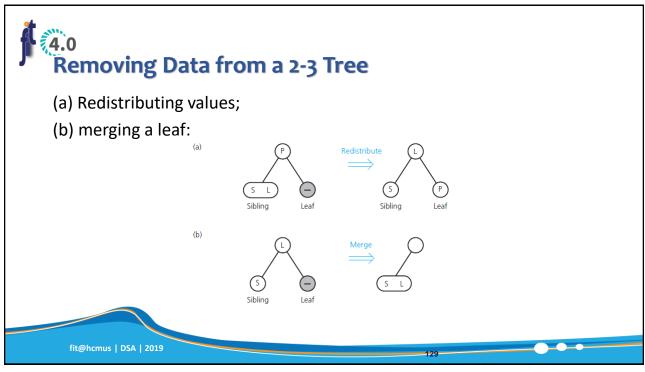
```
Removing Data from a 2-3 Tree

// Completes the removal when node n is empty by either deleting the root,
// redistributing values, or merging nodes. Note: If n is internal, it has one child.
fixTree(n: TwoThreeNode)

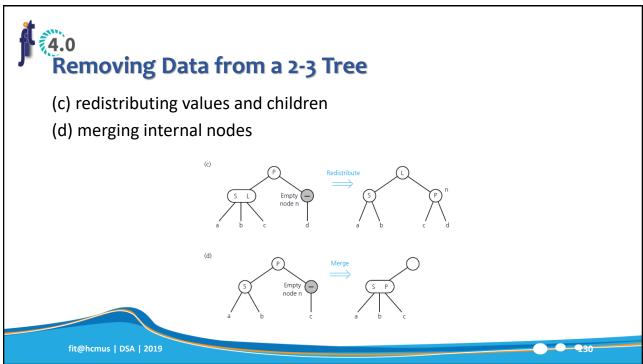
if (n is the root)
Delete the root
else
{
Let p be the parent of n
if (some sibling of n has two items)
{
Distribute items appropriately among n, the sibling, and p
```

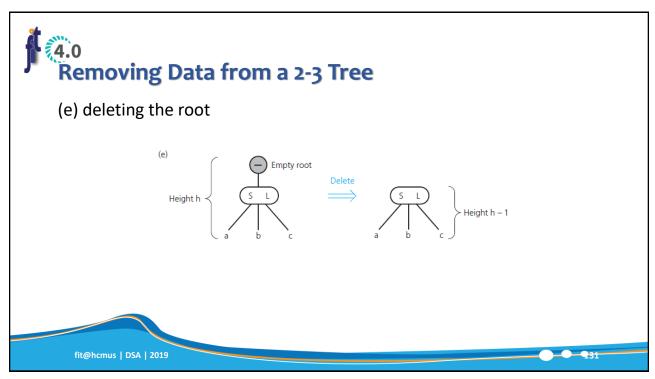




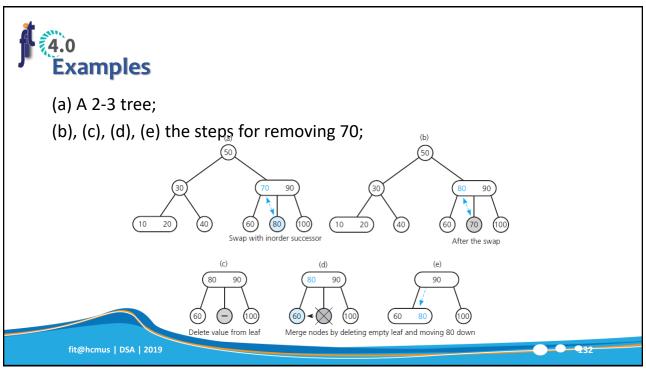


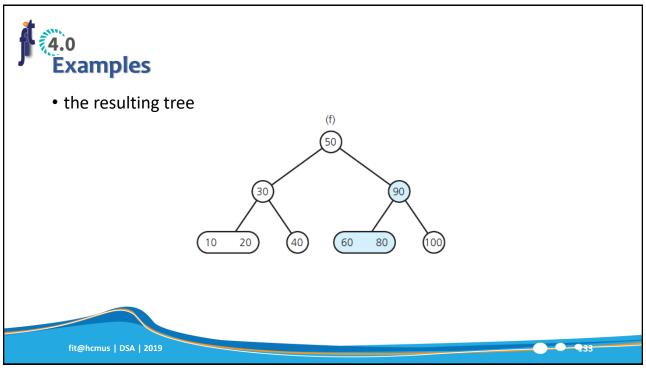




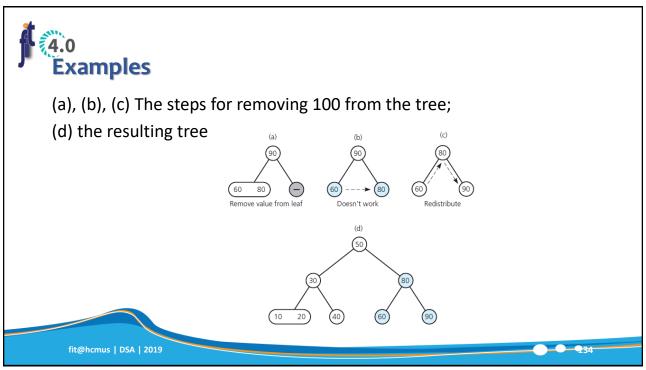


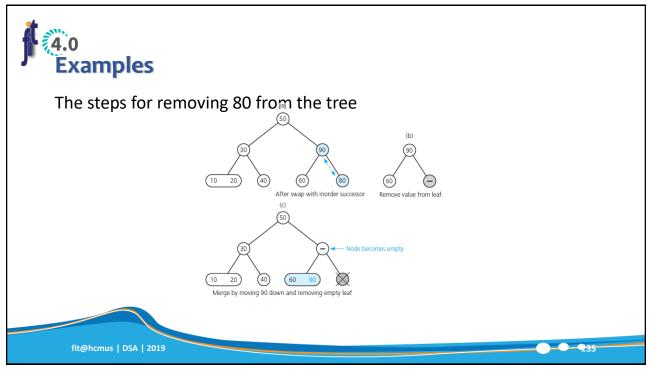




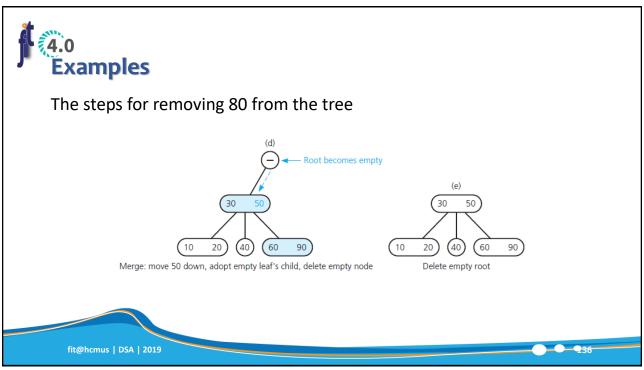












4.0 Removing Data from a 2-3-4 Tree

- Removal algorithm has same beginning as removal algorithm for a 2-3 tree
- \bullet Locate the node n that contains the item $\ensuremath{\mathbb{I}}$ you want to remove.
- Find $\mathbb I$'s in-order successor and swap it with $\mathbb I$ so that the removal will always be at a leaf.
- ullet If leaf is either a 3-node or a 4-node, remove ${\tt I}$.

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