



Binary Search Tree

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Binary Search Tree

- A binary search tree is a binary tree in which each node n has properties:
 - n 's value greater than all values in left subtree T_L
 - n 's value less than all values in right subtree T_R
 - Both T_R and T_L are binary search trees.

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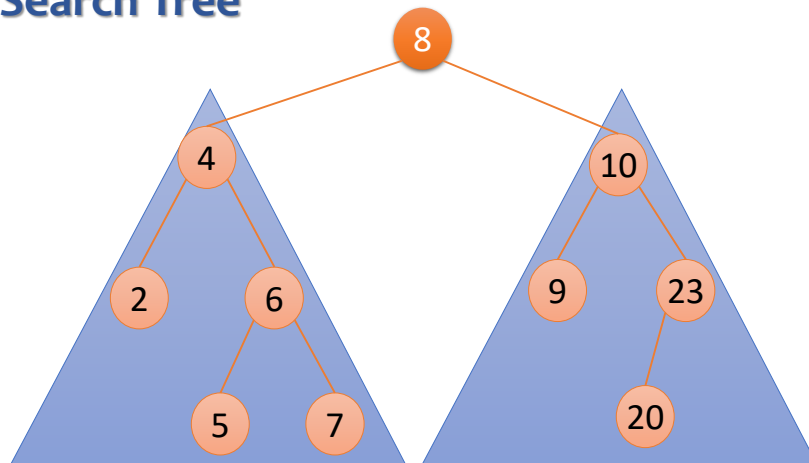
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Binary Search Tree

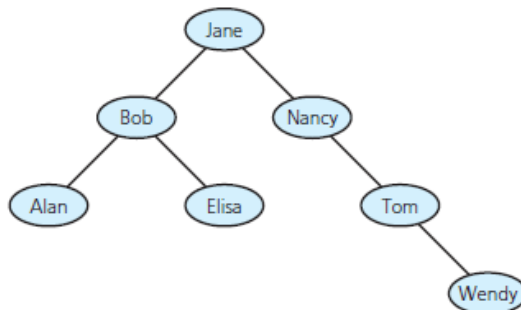
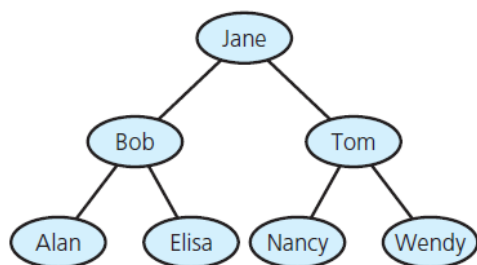


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Binary Search Tree



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Binary Search Tree | Operations

- Insert (a key)
- Search (a key)
- Remove (a key)
- Traverse
- Sort (based on key value)
- Rotate (Left rotation, Right rotation)

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Binary Search Tree | Insert

Insert (root, Data)

```
{
    if (root is NULL) {
        Create a new_Node containing Data
        This new_Node is root of the tree
    }
    //Compare root's key with Key
    if root's Key is less than Data's Key
        Insert Key to the root's RIGHT subtree
    else if root's Key is greater than Data's Key
        Insert Key to the root's LEFT subtree
    else
        Do nothing //Explain why?
}
```

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Binary Search Tree | Insert

- Beginning with an empty binary search tree, what binary search tree is formed when you insert the following values in the order given?

15, 5, 12, 8, 23, 1, 17, 21



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Binary Search Tree | Insert

- Beginning with an empty binary search tree, what binary search tree is formed when you insert the following values in the order given?

- W, T, N, J, E, B, A
- W, T, N, A, B, E, J
- A, B, W, J, N, T, E
- B, T, E, A, N, W, J



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Binary Search Tree | Search

Search (root, Data)

```
{
    if (root is NULL) {
        return NOT_FOUND;
    }
    //Compare root's key with Key
    if root's Key is less than Data's Key
        Search Data in the root's RIGHT subtree
    else if root's Key is greater than Data's Key
        Search Data in the root's LEFT subtree
    else
        return FOUND //Explain why?
}
```



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Binary Search Tree | Remove

- When we delete a node, three possibilities arise.
- Node to be deleted:
 - **is leaf:**
 - Simply remove from the tree.
 - has **only one child:**
 - Copy the child to the node and delete the child
 - has **two children:**
 - Find in-order successor (predecessor) **S_Node** of the node.
 - Copy contents of **S_Node** to the node and delete the **S_Node**.



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Binary Search Tree | Traversal

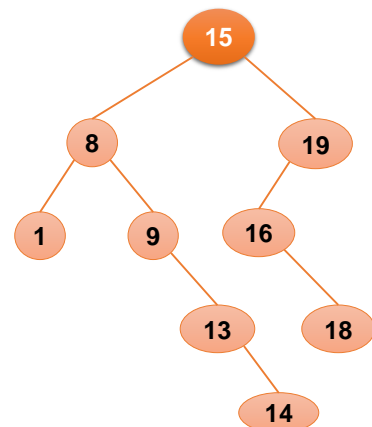
- Pre-Order:
 - Node - Left - Right
- In-Order:
 - Left - Node - Right
- Post-Order:
 - Left - Right - Node



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Binary Search Tree | Traversal

- What are the pre-order, in-order and post-order traversals of this binary search tree?

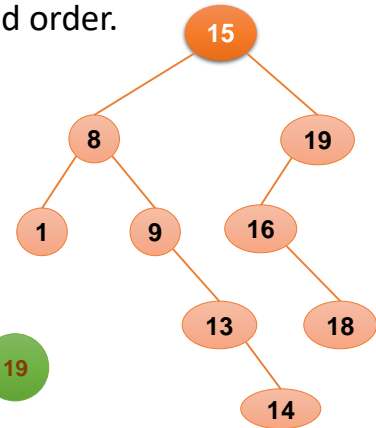




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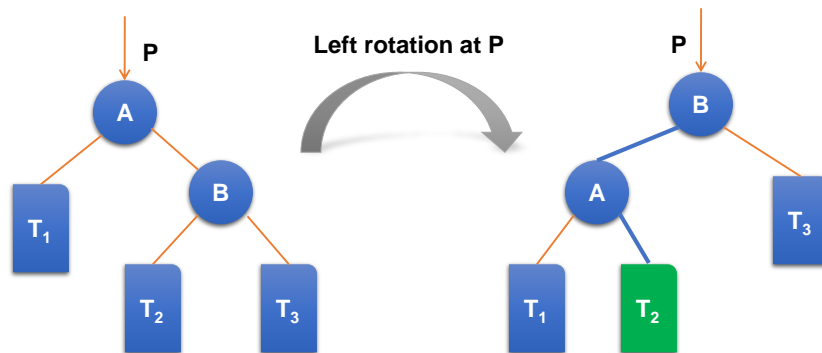
Binary Search Tree | Sort

- In-order traversal gives the key values in sorted order.



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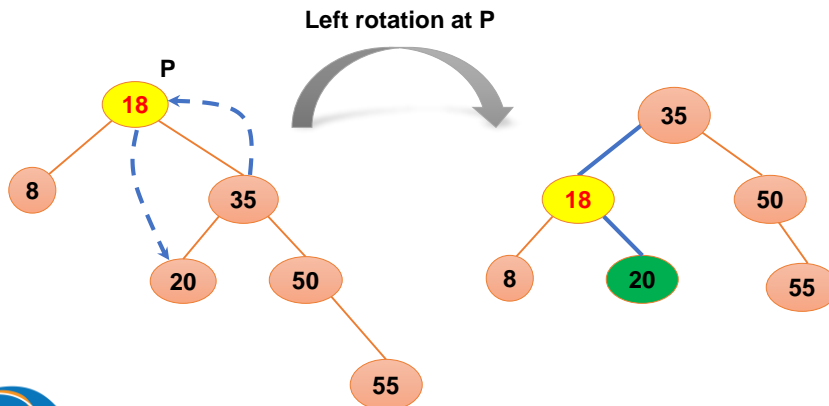
Binary Search Tree | Left Rotation





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Binary Search Tree | Left Rotation



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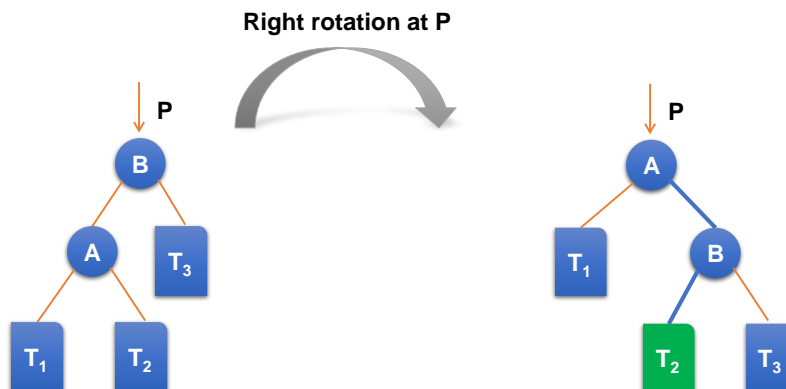
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Binary Search Tree | Right Rotation



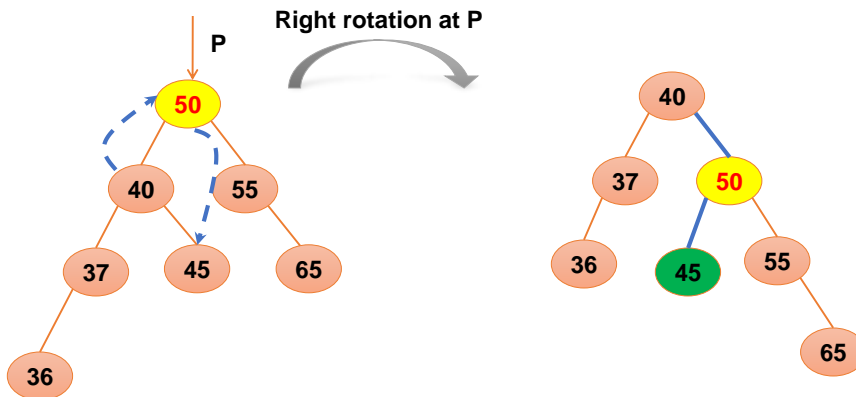
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Binary Search Tree | Right Rotation



Efficiency of Binary Search Tree Operations

Operation	Average case	Worst case
Retrieval	$O(\log n)$	$O(n)$
Insertion	$O(\log n)$	$O(n)$
Removal	$O(\log n)$	$O(n)$
Traversal	$O(n)$	$O(n)$



Very Bad Binary Search Tree

- Beginning with an empty binary search tree, what binary search tree is formed when inserting the following values in the order given?

2, 4, 6, 8, 10, 12, 14, 18, 20



AVL Tree



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AVL Tree

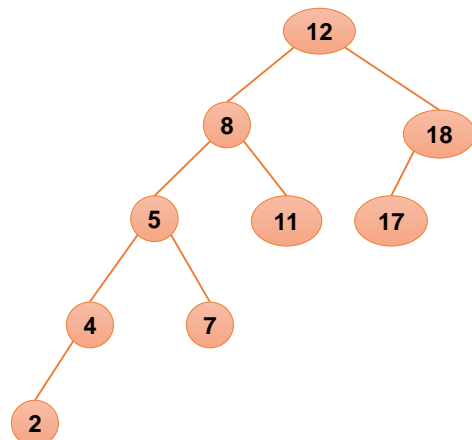
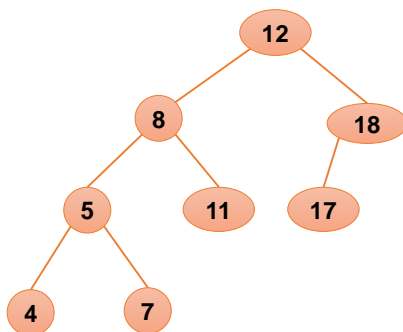
- Named for inventors, **Adel'son-Vel'skii** and **Landis** (1962).
- AVL Tree is a **self-balancing** binary search tree where
 - for ALL nodes, the difference between height of the left subtrees and the right subtrees **cannot be more than one**.



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AVL Tree

- Which is the AVL Tree?





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AVL Tree

- A balanced binary search tree
 - Maintains height close to the minimum
 - After insertion or deletion, check the tree is still AVL tree – determine whether any node in tree has left and right subtrees whose heights differ by more than 1
- Can search AVL tree almost as efficiently as minimum-height binary search tree.

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Cases of Un-Balanced Nodes

- Left-Left case
- Left-Right case
- Right-Right case
- Right-Left case

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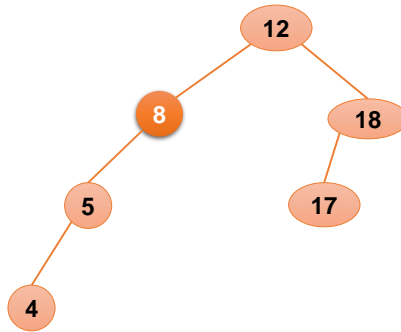
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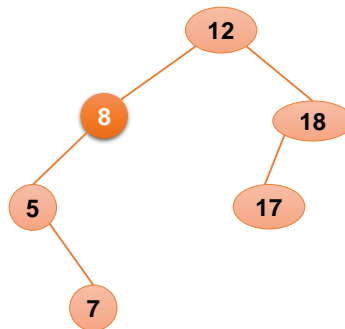
Cases of Un-Balanced Nodes

- Left-Left case



Cases of Un-Balanced Nodes

- Left-Right case

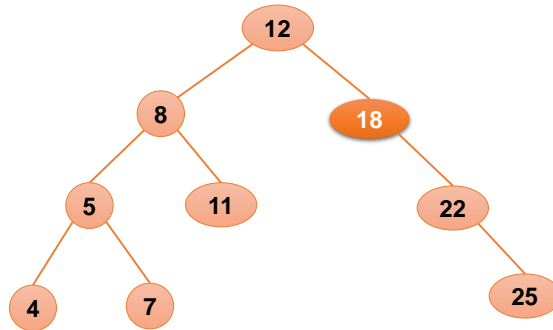




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Cases of Un-Balanced Nodes

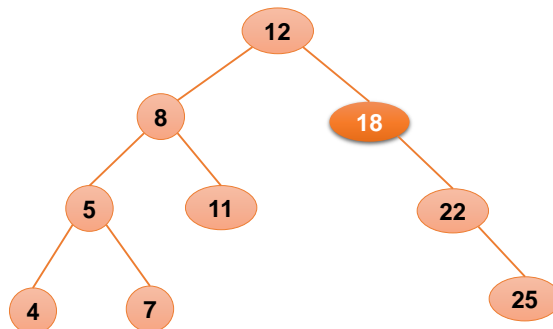
- Right-Right case



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Cases of Un-Balanced Nodes

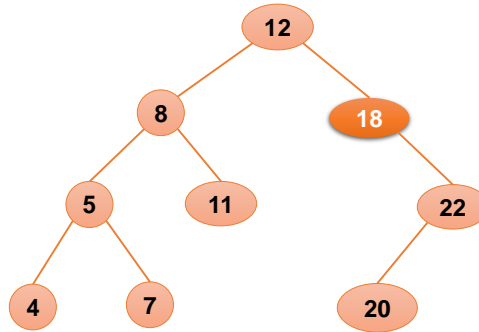
- Right-Right case





Cases of Un-Balanced Nodes

- Right-Left case



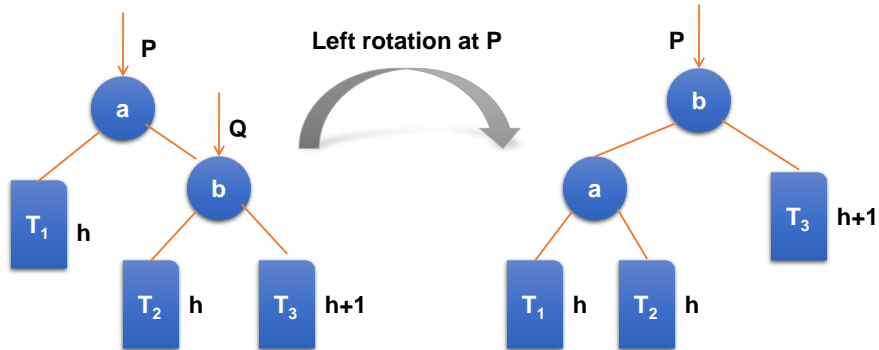
Un-balanced Resolving

- Right-Right case:
 - Left rotation at un-balanced node.
- Right-Left case:
 - Right rotation at un-balanced node's right child.
 - Left rotation at un-balanced node.



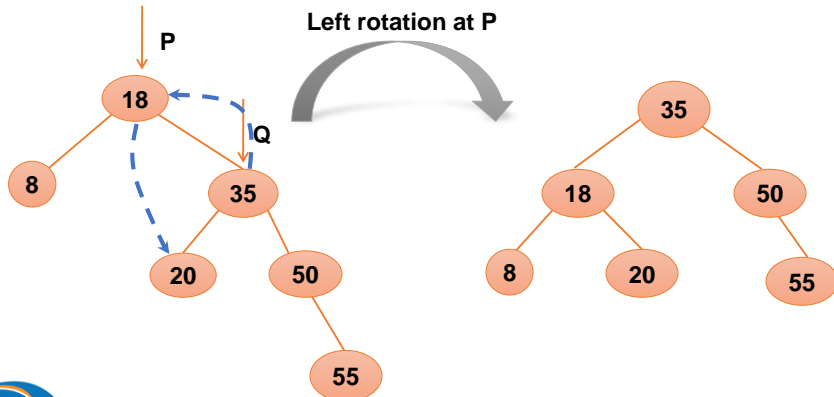
Un-balanced Resolving

- Right-Right case:



Un-balanced Resolving

- Right-Right case:

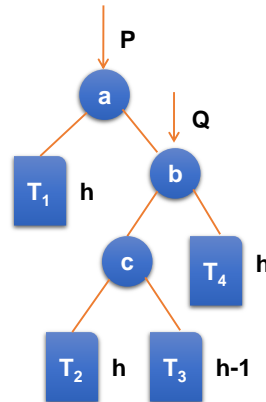




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Un-balanced Resolving

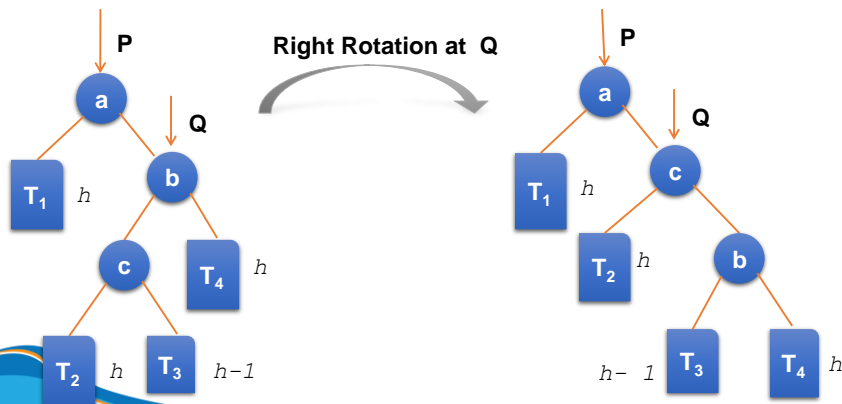
- Right-Left case:
 - Right rotation at Q
 - Left rotation at P



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Un-balanced Resolving

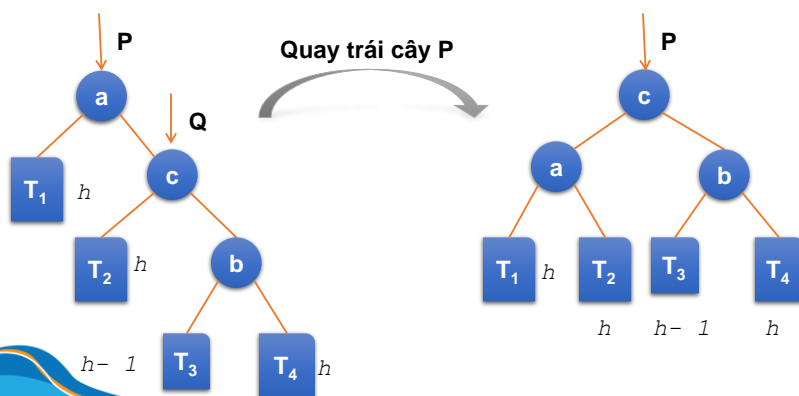
- Right-Left case:
 - Right rotation at Q



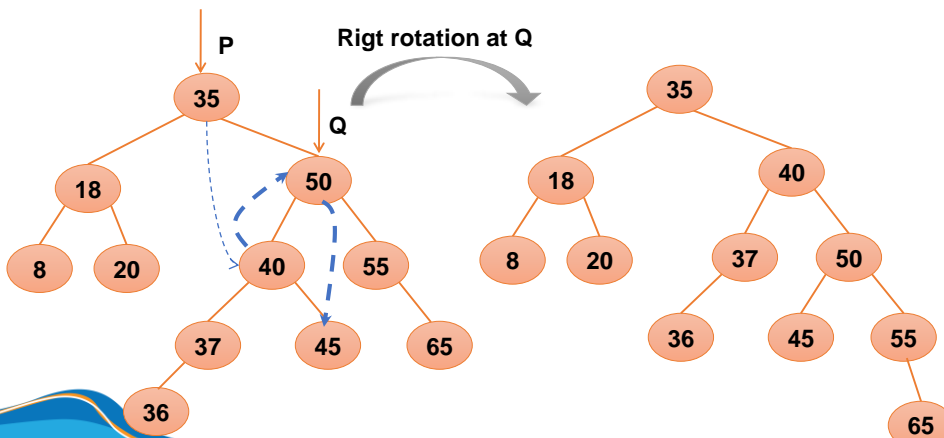


Un-balanced Resolving

- Right-Left case:
 - Right rotation at P

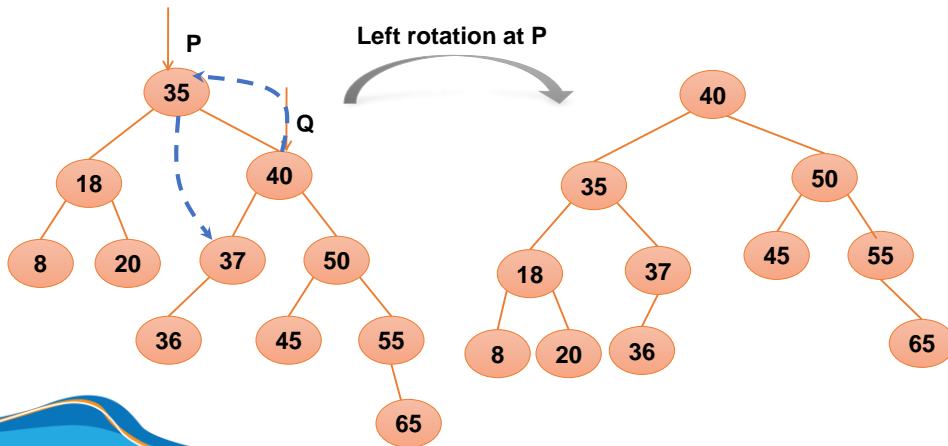


Un-balanced Resolving





Un-balanced Resolving



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Un-balanced Resolving

- Left-Left case:
 - Right rotation at un-balanced node.
- Left-Right case:
 - Left rotation at un-balanced node's left child.
 - Right rotation at un-balanced node.

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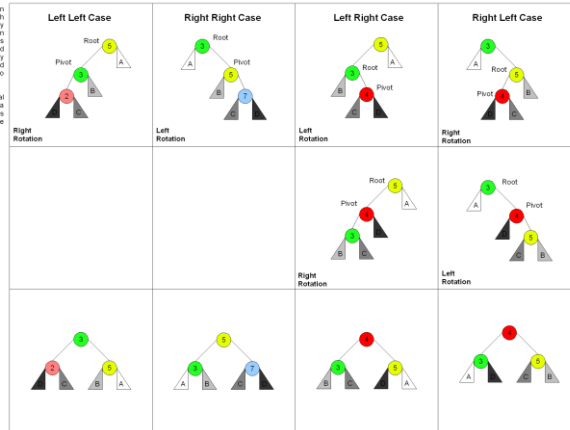
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Un-balanced Resolving

There are 4 cases in all, choosing which one to make by seeing the direction of the first 2 nodes from the unbalanced node to the newly inserted node and matching them to the top most row.

Root is the initial parent before a rotation and Pivot is the child to take the root's place.



Source: Wikipedia