

PROGRAMMING IN HASKELL



Chapter 8 - Declaring Types and Classes

Type Declarations

In Haskell, a new name for an existing type can be defined using a type declaration.

```
type String = [Char]
```

String is a synonym for the type [Char].

Type declarations can be used to make other types easier to read. For example, given

```
type Pos = (Int,Int)
```

we can define:

```
origin :: Pos  
origin = (0,0)
```

```
left :: Pos → Pos  
left (x,y) = (x-1,y)
```

Like function definitions, type declarations can also have parameters. For example, given

```
type Pair a = (a,a)
```

we can define:

```
mult :: Pair Int → Int  
mult (m,n) = m*n
```

```
copy :: a → Pair a  
copy x = (x,x)
```

Type declarations can be nested:

```
type Pos = (Int,Int)
type Trans = Pos → Pos
```



However, they cannot be recursive:

```
type Tree = (Int,[Tree])
```



Data Declarations

A completely new type can be defined by specifying its values using a data declaration.

```
data Bool = False | True
```

Bool is a new type, with two new values False and True.

Note:

- The two values False and True are called the constructors for the type Bool.
- Type and constructor names must always begin with an upper-case letter.
- Data declarations are similar to context free grammars. The former specifies the values of a type, the latter the sentences of a language.

Values of new types can be used in the same ways as those of built in types. For example, given

```
data Answer = Yes | No | Unknown
```

we can define:

```
answers :: [Answer]
answers = [Yes, No, Unknown]
```

```
flip :: Answer → Answer
flip Yes      = No
flip No       = Yes
flip Unknown  = Unknown
```


The constructors in a data declaration can also have parameters. For example, given

```
data Shape = Circle Float  
           | Rect Float Float
```

we can define:

```
square :: Float → Shape  
square n = Rect n n
```

```
area :: Shape → Float  
area (Circle r) = pi * r^2  
area (Rect x y) = x * y
```

Note:

- Shape has values of the form `Circle r` where `r` is a float, and `Rect x y` where `x` and `y` are floats.
- `Circle` and `Rect` can be viewed as functions that construct values of type `Shape`:

`Circle :: Float → Shape`

`Rect :: Float → Float → Shape`

Not surprisingly, data declarations themselves can also have parameters. For example, given

```
data Maybe a = Nothing | Just a
```

we can define:

```
safediv :: Int → Int → Maybe Int  
safediv _ 0 = Nothing  
safediv m n = Just (m `div` n)
```

```
safehead :: [a] → Maybe a  
safehead [] = Nothing  
safehead xs = Just (head xs)
```

Recursive Types

In Haskell, new types can be declared in terms of themselves. That is, types can be recursive.

```
data Nat = Zero | Succ Nat
```

Nat is a new type, with constructors
 $\text{Zero} :: \text{Nat}$ and $\text{Succ} :: \text{Nat} \rightarrow \text{Nat}$.

Note:

- A value of type `Nat` is either `Zero`, or of the form `Succ n` where $n :: \text{Nat}$. That is, `Nat` contains the following infinite sequence of values:

`Zero`

`Succ Zero`

`Succ (Succ Zero)`

⋮

- We can think of values of type Nat as natural numbers, where Zero represents 0, and Succ represents the successor function 1+.
- For example, the value

Succ (Succ (Succ Zero))

represents the natural number

$$1 + (1 + (1 + 0)) = 3$$

Using recursion, it is easy to define functions that convert between values of type Nat and Int:

```
nat2int :: Nat → Int
nat2int Zero = 0
nat2int (Succ n) = 1 + nat2int n

int2nat :: Int → Nat
int2nat 0 = Zero
int2nat n = Succ (int2nat (n-1))
```

Two naturals can be added by converting them to integers, adding, and then converting back:

```
add :: Nat → Nat → Nat  
add m n = int2nat (nat2int m + nat2int n)
```

However, using recursion the function add can be defined without the need for conversions:

```
add Zero      n = n  
add (Succ m) n = Succ (add m n)
```


For example:

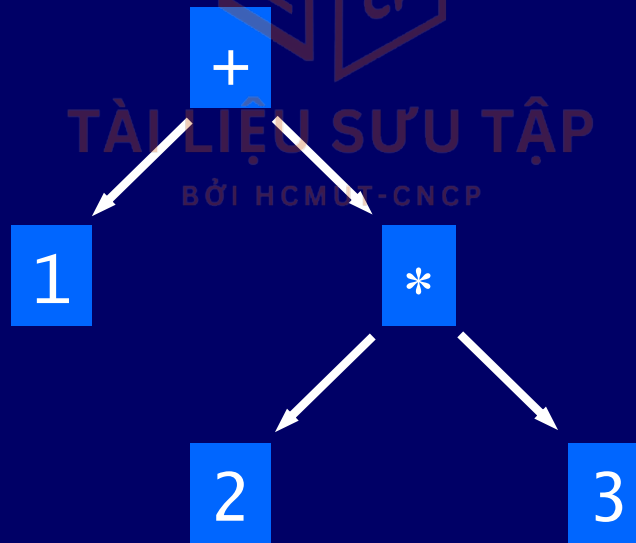
$$\begin{aligned} & \text{add (Succ (Succ Zero)) (Succ Zero)} \\ = & \text{Succ (add (Succ Zero) (Succ Zero))} \\ = & \text{Succ (Succ (add Zero (Succ Zero)))} \\ = & \text{Succ (Succ (Succ Zero))} \end{aligned}$$

Note:

- The recursive definition for add corresponds to the laws $0+n = n$ and $(1+m)+n = 1+(m+n)$.

Arithmetic Expressions

Consider a simple form of expressions built up from integers using addition and multiplication.



Using recursion, a suitable new type to represent such expressions can be declared by:

```
data Expr = Val Int  
          | Add Expr Expr  
          | Mul Expr Expr
```

For example, the expression on the previous slide would be represented as follows:

```
Add (Val 1) (Mul (Val 2) (Val 3))
```

Using recursion, it is now easy to define functions that process expressions. For example:

```
size :: Expr → Int
size (Val n) = 1
size (Add x y) = size x + size y
size (Mul x y) = size x + size y

eval :: Expr → Int
eval (Val n) = n
eval (Add x y) = eval x + eval y
eval (Mul x y) = eval x * eval y
```

Note:

- The three constructors have types:

```
Val  :: Int → Expr  
Add  :: Expr → Expr → Expr  
Mul  :: Expr → Expr → Expr
```

- Many functions on expressions can be defined by replacing the constructors by other functions using a suitable fold function. For example:

```
eval = folde id (+) (*)
```

Exercises

- (1) Using recursion and the function `add`, define a function that multiplies two natural numbers.
- (2) Define a suitable function folde for expressions and give a few examples of its use.
- (3) Define a type Tree a of binary trees built from Leaf values of type `a` using a Node constructor that takes two binary trees as parameters.