

# Digits extraction

Saturday, December 21, 2024

12:51 PM

7789  $\div 10 = 9$  ?

mod returns the remainder of a division

How to get the next digit 8?

778  $\div 10 = 8$

77  $\div 10 = 7$

7  $\div 10 = 7$

0 stop  $\Rightarrow$  we extracted all the digits in reverse

7789  $\div 10$

7789	10
7780	778
9	

## Pseudo Code :

N: the number

while (n > 0) {

lastDigit = n % 10

n = n / 10

}

# Digits count

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## Count the number of digits

method 1

```
count := 0
while (n > 0) {
    count++
    n = n / 10
}
```

number of digits = the number of division by 10  
 $= \log_{10}(n) + 1$   
 $= \log_{10}(7789) + 1$   
 $= 3.89 + 1 = 4.89$   
 $\Rightarrow 4 \text{ digits}$  TC:  $O(\log_{10}(n))$

method 2

```
int Count (int n) {
    return (int)(log10(n) + 1)
}
```

TC:  $O(1)$

# (LeetCode) Reverse & Palindrome Number

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Examples;

123  $\rightarrow$  321 / 120  $\rightarrow$  21

```
int getReverseNumber (int n){
```

```
    int result = 0
```

```
    while (n > 0) {
```

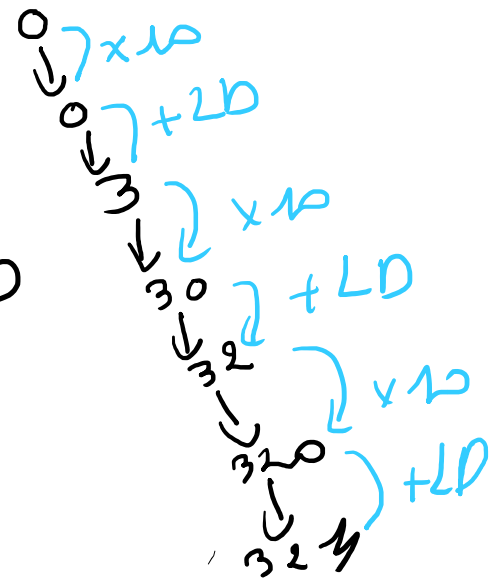
```
        last Digit = n % 10
```

```
        result = result * 10 + LD
```

```
        n = n / 10
```

```
    }
```

```
    return result}
```



## Palindrome number

It's a number that reads the same as backward as forward, examples: 121, 1331, 4554

```
bool isPalindrome (int n){
```

```
    int copy = n
```

```
    pm = getReverseNumber(n)
```

```
    if (pm == copy){
```

```
        return true
```

```
    }
```

```
    return false
```

```
}
```

## (LeetCode) Armstrong Number

Sunday, December 22, 2024

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It's a number that is equal to the sum of its digits each raised to the power of the number of digits. example:  $153 = 1^3 + 5^3 + 3^3$

Solution:

Loop over digits and calculate the sum of the digit power  $l$

```
bool isArmstrongNumber(int n){  
    int l = toString(n).length()  
    int nCopy = n  
    int sum = 0  
    while(n > 0){  
        sum = pow(n%10, l)  
        n = n/10  
    }  
    if(nCopy == sum){  
        return true  
    }  
    return false  
}
```

# Get All Divisors

Sunday, December 22, 2024

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## Problem statement:

Given an integer  $N$ , return all divisors of  $N$ .

Example:  $12 \rightarrow 1, 2, 3, 4, 6, 12$

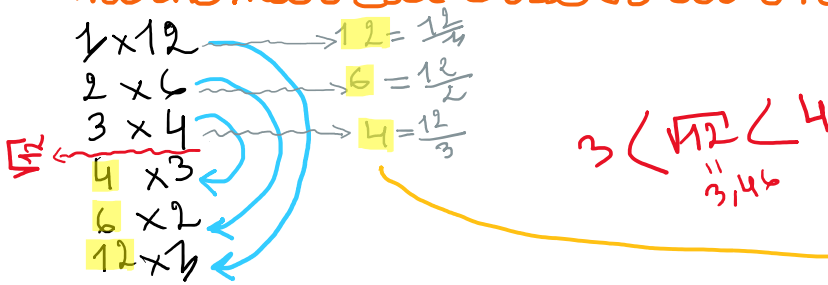
## Problem Analysis:

divisors  $\in [1, N]$

loop from 1 to  $N$ , if  $i \mid N == 0 \rightarrow$  divisor

$\rightarrow T.C = O(N)$ , (too much) (brute force)

## Mathematical observation:



So, we loop to  $N/2$  only, if  $i$  is a divisor,  $\frac{N}{i}$  is also

$\rightarrow T.C = O(\sqrt{n})$

## Solution

```
void printAllDivisors(int n){
```

```
    for(i=1; i ≤ √n; i++){
```

```
        if(n/i == 0){
```

```
            print(i)
```

```
            if (n/i != i) { print(i) }
```

```
        }
```

```
    }
```

```
}
```

to avoid double counting in case of perfect square, 9, 16, 36 ....

# Prime Number

Sunday, December 22, 2024 7:25 PM

## Problem statement:

A prime number is a number that is only divisible by 1 and itself. Example: 2

## Problem analysis:

The total number of divisors is exactly 2: so we loop over divisors, if  $\text{count} > 2$  return false;

```
bool isPrimeNumber(const int n) {  
    if (n <= 1) return false; // 0 and 1 are not prime numbers  
  
    int count = 0;  
    for (int i = 1; i <= std::sqrt(n); i++) {  
        if (n % i == 0) {  
            count++;  
            if (i != n / i) count++; // Count only distinct divisors  
        }  
        if (count > 2) {  
            return false;  
        }  
    }  
    return true;  
}
```

# Digit Concept - GCD

Monday, December 23, 2024 7:08 AM

## Problem statement:

The greatest common divisor of any two integers is the largest number that divides both integers. example 9 and 12

factors of 9: 1, 3, 9 } factors of 12: 1, 2, 3, 4, 6, 12

GCD  $\rightarrow$  3

## Problem analysis:

brute force:

We iterate from 1 to  $\min(N_1, N_2)$ , if  $i$  divides both  $N_1$  and  $N_2$ , we affect it the GCD variable

$$T.C = O(\min(n_1, n_2))$$

## Better approach:

In the previous approach, in the worst case,  $i$  iterates from 1 to  $\min(n_1, n_2)$ . if  $\min(n_1, n_2)$  is a very large number. it will result a large number of iterations.

So, we can iterate from  $\min(n_1, n_2)$ . if a common divisor is found. we return it directly.

$\rightarrow$  T.C: it is the same. But we execute less iterations

```
int findGCD(int n1, int n2) {  
    for (int i = std::min(n1, n2); i > 0; i--) { // std::min(a, b): O(1)  
        if (n1 % i == 0 && n2 % i == 0) {  
            return i;  
        }  
    }  
    return 1; // In case no other divisor is found  
}
```