

# **SPIR-V Specification**

John Kessenich, Google, Boaz Ouriel, Intel, and Raun Krisch, Intel

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## **Contributors and Acknowledgments**

Connor Abbott, Intel

Ben Ashbaugh, Intel

Alexey Bader, Intel

Dan Baker, Oxide Games

Kenneth Benzie, Codeplay

Gordon Brown, Codeplay

Pat Brown, NVIDIA

Diana Po-Yu Chen, MediaTek

Stephen Clarke, Imagination

Patrick Doane, Blizzard Entertainment

Stefanus Du Toit, Google

Tim Foley, Intel

Ben Gaster, Qualcomm

Alexander Galazin, ARM

Christopher Gautier, ARM

Neil Henning, Codeplay

Kerch Holt, NVIDIA

Lee Howes, Qualcomm

Roy Ju, MediaTek

John Kessenich, Google

Daniel Koch, NVIDIA

Ashwin Kolhe, NVIDIA

Raun Krisch, Intel

Graeme Leese, Broadcom

Yuan Lin, NVIDIA

Yaxun Liu, AMD

Timothy Lottes, Epic Games

John McDonald, Valve

David Neto, Google

Boaz Ouriel, Intel

Christophe Riccio, Unity

Andrew Richards, Codeplay

Ian Romanick, Intel

Graham Sellers, AMD

Robert Simpson, Qualcomm

Bartosz Sochacki, Intel

Brian Sumner, AMD

Andrew Woloszyn, Google

Weifeng Zhang, Qualcomm

#### Note

Up-to-date HTML and PDF versions of this specification may be found at the Khronos SPIR-V Registry. (https://www.khronos.org/registry/spir-v/)

## 1 Introduction

#### **Abstract**

SPIR-V is a simple binary intermediate language for graphical shaders and compute kernels. A SPIR-V module contains multiple entry points with potentially shared functions in the entry point's call trees. Each function contains a control-flow graph (CFG) of basic blocks, with optional instructions to express structured control flow. Load/store instructions are used to access declared variables, which includes all input/output (IO). Intermediate results bypassing load/store use static single-assignment (SSA) representation. Data objects are represented logically, with hierarchical type information: There is no flattening of aggregates or assignment to physical register banks, etc. Selectable addressing models establish whether general pointer operations may be used, or if memory access is purely logical.

This document fully defines **SPIR-V**, a Khronos-standard binary intermediate language for representing graphical-shader stages and compute kernels for multiple Khronos APIs.

This is a unified specification, specifying all versions since and including version 1.0.

#### 1.1 Goals

SPIR-V has the following goals:

- Provide a simple binary intermediate language for all functionality appearing in Khronos shaders/kernels.
- Have a concise, transparent, self-contained specification (sections Specification and Binary Form).
- Map easily to other intermediate languages.
- Be the form passed by an API into a driver to set shaders/kernels.
- Can be targeted by new front ends for novel high-level languages.
- Allow the first steps of compilation and reflection to be done offline.
- Be low-level enough to require a reverse-engineering step to reconstruct source code.
- Improve portability by enabling shared tools to generate or operate on it.
- Allow separation of core specification from source-language-specific sets of built-in functions.
- Reduce compile time during application run time. (Eliminating most of the compile time during application run time is not a goal of this intermediate language. Target-specific register allocation and scheduling are still expected to take significant time.)
- Allow some optimizations to be done offline.

#### 1.2 About this document

This document aims to:

- Include everything needed to fully understand, create, and consume SPIR-V. However:
- Imported sets of instructions (which implement source-specific built-in functions) will need their own specification.
- Many validation rules are client-API specific, and hence documented with client API and not in this specification.
- Separate expository and specification language. The specification-proper is in Specification and Binary Form.

## 1.3 Extendability

SPIR-V can be extended by multiple vendors or parties simultaneously:

- Using the OpExtension instruction to require new semantics that must be supported. Such new semantics would come from an extension document.
- Reserving (registering) ranges of the token values, as described further below.
- Aided by instruction skipping, also further described below.

**Enumeration Token Values.** It is easy to extend all the types, storage classes, opcodes, decorations, etc. by adding to the token values.

**Registration.** Ranges of token values in the Binary Form section can be pre-allocated to numerous vendors/parties. This allows combining multiple independent extensions without conflict. To register ranges, see <a href="https://www.khronos.org/registry/spir-v/api/spir-v.xml">https://www.khronos.org/registry/spir-v/api/spir-v.xml</a>.

**Extended Instructions.** Sets of extended instructions can be provided and specified in separate specifications. These help personalize SPIR-V for different source languages or execution environments (client APIs). Multiple sets of extended instructions can be imported without conflict, as the extended instructions are selected by {set id, instruction number} pairs.

**Instruction Skipping.** Tools are encouraged to skip opcodes for features they are not required to process. This is trivially enabled by the word count in an instruction, which makes it easier to add new instructions without breaking existing tools.

## 1.4 Debuggability

SPIR-V can decorate, with a text string, virtually anything created in the shader: types, variables, functions, etc. This is required for externally visible symbols, and also allowed for naming the result of any instruction. This can be used to aid in understandability when disassembling or debugging lowered versions of SPIR-V.

Location information (file names, lines, and columns) can be interleaved with the instruction stream to track the origin of each instruction.

## 1.5 Design Principles

**Regularity.** All instructions start with a word count. This allows walking a SPIR-V module without decoding each opcode. All instructions have an opcode that dictates for all operands what kind of operand they are. For instructions with a variable number of operands, the number of variable operands is known by subtracting the number of non-variable words from the instruction's word count.

**Non Combinatorial.** There is no combinatorial type explosion or need for large encode/decode tables for types. Rather, types are parameterized. Image types declare their dimensionality, arrayness, etc. all orthogonally, which greatly simplify code. This is done similarly for other types. It also applies to opcodes. Operations are orthogonal to scalar/vector size, but not to integer vs. floating-point differences.

**Modeless.** After a given execution model (e.g., pipeline stage) is specified, internal operation is essentially modeless: Generally, it will follow the rule: "same spelling, same semantics", and does not have mode bits that modify semantics. If a change to SPIR-V modifies semantics, it should use a different spelling. This makes consumers of SPIR-V much more robust. There are execution modes declared, but these are generally to affect the way the module interacts with the environment around it, not the internal semantics. Capabilities are also declared, but this is to declare the subset of functionality that is used, not to change any semantics of what is used.

**Declarative.** SPIR-V declares externally-visible modes like "writes depth", rather than having rules that require deduction from full shader inspection. It also explicitly declares what addressing modes, execution model, extended instruction sets, etc. will be used. See Language Capabilities for more information.

**SSA.** All results of intermediate operations are strictly SSA. However, declared variables reside in memory and use load/store for access, and such variables can be stored to multiple times.

**IO.** Some storage classes are for input/output (IO) and, fundamentally, IO will be done through load/store of variables declared in these storage classes.

## 1.6 Static Single Assignment (SSA)

SPIR-V includes a phi instruction to allow the merging together of intermediate results from split control flow. This allows split control flow without load/store to memory. SPIR-V is flexible in the degree to which load/store is used; it is possible to use control flow with no phi-instructions, while still staying in SSA form, by using memory load/store.

Some storage classes are for IO and, fundamentally, IO will be done through load/store, and initial load and final store can never be eliminated. Other storage classes are shader local and can have their load/store eliminated. It can be considered an optimization to largely eliminate such loads/stores by moving them into intermediate results in SSA form.

#### 1.7 Built-In Variables

SPIR-V identifies built-in variables from a high-level language with an enumerant decoration. This assigns any unusual semantics to the variable. Built-in variables must otherwise be declared with their correct SPIR-V type and treated the same as any other variable.

## 1.8 Specialization

*Specialization* enables creating a portable SPIR-V module outside the target execution environment, based on constant values that won't be known until inside the execution environment. For example, to size a fixed array with a constant not known during creation of a module, but known when the module will be lowered to the target architecture.

See Specialization in the next section for more details.

## 1.9 Example

The SPIR-V form is binary, not human readable, and fully described in Binary Form. This is an example disassembly to give a basic idea of what SPIR-V looks like:

## GLSL fragment shader:

```
#version 450
in vec4 color1;
in vec4 multiplier;
noperspective in vec4 color2;
out vec4 color;
struct S {
   bool b;
   vec4 v[5];
   int i;
};
uniform blockName {
   S s;
   bool cond;
};
void main()
    vec4 scale = vec4(1.0, 1.0, 2.0, 1.0);
   if (cond)
        color = color1 + s.v[2];
    else
        color = sqrt(color2) * scale;
    for (int i = 0; i < 4; ++i)
        color *= multiplier;
```

## Corresponding SPIR-V:

```
; Magic:
             0x07230203 (SPIR-V)
            0x00010000 (Version: 1.0.0)
; Version:
; Generator: 0x00080001 (Khronos Glslang Reference Front End; 1)
; Bound:
; Schema:
               OpCapability Shader
          %1 = OpExtInstImport "GLSL.std.450"
               OpMemoryModel Logical GLSL450
               OpEntryPoint Fragment %4 "main" %31 %33 %42 %57
               OpExecutionMode %4 OriginLowerLeft
; Debug information
               OpSource GLSL 450
               OpName %4 "main"
               OpName %9 "scale"
               OpName %17 "S"
               OpMemberName %17 0 "b"
               OpMemberName %17 1 "v"
               OpMemberName %17 2 "i"
```

```
OpName %18 "blockName"
              OpMemberName %18 0 "s"
              OpMemberName %18 1 "cond"
              OpName %20 ""
              OpName %31 "color"
              OpName %33 "color1"
              OpName %42 "color2"
              OpName %48 "i"
              OpName %57 "multiplier"
; Annotations (non-debug)
              OpDecorate %15 ArrayStride 16
              OpMemberDecorate %17 0 Offset 0
              OpMemberDecorate %17 1 Offset 16
              OpMemberDecorate %17 2 Offset 96
              OpMemberDecorate %18 0 Offset 0
              OpMemberDecorate %18 1 Offset 112
              OpDecorate %18 Block
              OpDecorate %20 DescriptorSet 0
              OpDecorate %42 NoPerspective
; All types, variables, and constants
         %2 = OpTypeVoid
                                                   ; void ()
         %3 = OpTypeFunction %2
                                                    ; 32-bit float
         %6 = OpTypeFloat 32
         %7 = OpTypeVector %6 4
                                                   ; vec4
         %8 = OpTypePointer Function %7 ; function-local vec4*
        %10 = OpConstant %6 1
        %11 = OpConstant %6 2
        %12 = OpConstantComposite %7 %10 %10 %11 %10; vec4(1.0, 1.0, 2.0, 1.0)
        %13 = OpTypeInt 32 0
                                                    ; 32-bit int, sign-less
        %14 = OpConstant %13 5
        %15 = OpTypeArray %7 %14
        %16 = OpTypeInt 32 1
        %17 = OpTypeStruct %13 %15 %16
        %18 = OpTypeStruct %17 %13
        %19 = OpTypePointer Uniform %18
        %20 = OpVariable %19 Uniform
        %21 = OpConstant %16 1
        %22 = OpTypePointer Uniform %13
        %25 = OpTypeBool
        %26 = OpConstant %13 0
        %30 = OpTypePointer Output %7
        %31 = OpVariable %30 Output
        %32 = OpTypePointer Input %7
        %33 = OpVariable %32 Input
        %35 = OpConstant %16 0
        %36 = OpConstant %16 2
        %37 = OpTypePointer Uniform %7
        %42 = OpVariable %32 Input
        %47 = OpTypePointer Function %16
        %55 = OpConstant %16 4
        %57 = OpVariable %32 Input
; All functions
         %4 = OpFunction %2 None %3
                                                     ; main()
         %5 = OpLabel
         %9 = OpVariable %8 Function
        %48 = OpVariable %47 Function
```

```
OpStore %9 %12
%23 = OpAccessChain %22 %20 %21 ; location of cond %24 = OpLoad %13 %23 ; load 32-bit int from cond %27 = OpINotEqual %25 %24 %26 ; convert to bool OpSelectionMerge %29 None ; structured if OpBranchConditional %27 %28 %41 ; if cond %28 = OpLabel ; then
%28 = OpLabel
                                                   ; then
%34 = OpLoad %7 %33
%38 = OpAccessChain %37 %20 %35 %21 %36 ; s.v[2]
%39 = OpLoad %7 %38
%40 = OpFAdd %7 %34 %39
      OpStore %31 %40
      OpBranch %29
%41 = OpLabel
                                                  ; else
%43 = OpLoad %7 %42
%44 = OpExtInst %7 %1 Sqrt %43 ; extended instruction sqrt
%45 = OpLoad %7 %9
%46 = OpFMul %7 %44 %45
      OpStore %31 %46
      OpBranch %29
%29 = OpLabel
                                                  ; endif
      OpStore %48 %35
      OpBranch %49
%49 = OpLabel
                                                 ; structured loop
      OpLoopMerge %51 %52 None
      OpBranch %53
%53 = OpLabel
%54 = OpLoad %16 %48
%56 = OpSLessThan %25 %54 %55 ; i < 4 ?
OpBranchConditional %56 %50 %51 ; body or break
%50 = OpLabel
                                                   ; body
%58 = OpLoad %7 %57
%59 = OpLoad %7 %31
%60 = OpFMul %7 %59 %58
      OpStore %31 %60
      OpBranch %52
                                                 ; continue target
%52 = OpLabel
%61 = OpLoad %16 %48
%62 = OpIAdd %16 %61 %21
                                      ; ++i
      OpStore %48 %62
      OpBranch %49
                                                   ; loop back
%51 = OpLabel
                                                   ; loop merge point
      OpReturn
       OpFunctionEnd
```

## 2 Specification

## 2.1 Language Capabilities

A SPIR-V module is consumed by an execution environment, specified by a client API, that needs to support the features used by that SPIR-V module. Features are classified through capabilities. Capabilities used by a particular SPIR-V module must be declared early in that module with the OpCapability instruction. Then:

- A validator can validate that the module uses only its declared capabilities.
- An execution environment is allowed to reject modules declaring capabilities it does not support. (See client API specifications for environment-specific rules.)

All available capabilities and their dependencies form a capability hierarchy, fully listed in the capability section. Only top-level capabilities need to be explicitly declared; their dependencies are implicitly declared.

When an instruction, enumerant, or other feature specifies multiple enabling capabilities, only one such capability needs to be declared to use the feature. This declaration does not itself imply anything about the presence of the other enabling capabilities: The execution environment needs to support only the declared capability.

This (SPIR-V) specification provides capability-specific validation rules, in the validation section. To ensure portability, each client API needs to include the following:

- Which capabilities in the capability section it requires environments to support, and hence allows in SPIR-V modules.
- Required limits, if they are beyond the Universal Limits.
- Any validation requirements specific to the environment that are not tied to specific capabilities, and hence not covered
  in the SPIR-V specification.

#### 2.2 Terms

## 2.2.1 Instructions

Word: 32 bits.

< id >: A numerical name; the name used to refer to an object, a type, a function, a label, etc. An < id > always consumes one word. The < id > s defined by a module obey SSA.

*Result* <*id*>: Most instructions define a result, named by an <*id*> explicitly provided in the instruction. The *Result* <*id*> is used as an operand in other instructions to refer to the instruction that defined it.

Literal String: A nul-terminated stream of characters consuming an integral number of words. The character set is Unicode in the UTF-8 encoding scheme. The UTF-8 octets (8-bit bytes) are packed four per word, following the little-endian convention (i.e., the first octet is in the lowest-order 8 bits of the word). The final word contains the string's nul-termination character (0), and all contents past the end of the string in the final word are padded with 0.

Literal Number: A numeric value consuming one or more words. An instruction will determine what type a literal will be interpreted as. When the type's bit width is larger than one word, the literal's low-order words appear first. When the type's bit width is less than 32-bits, the literal's value appears in the low-order bits of the word, and the high-order bits must be 0 for a floating-point type, or 0 for an integer type with *Signedness* of 0, or sign extended when *Signedness* is 1. (Similarly for the remaining bits of widths larger than 32 bits but not a multiple of 32 bits.)

Literal: A Literal String or a Literal Number.

*Operand:* A one-word argument to an instruction. E.g., it could be an <id>, or a (part of a) literal. Which form it holds is always explicitly known from the opcode.

*Immediate*: Operand(s) directly holding a literal value rather than an <id>. Immediate values larger than one word will consume multiple operands, one per word. That is, operand counting is always done per word, not per immediate.

*WordCount:* The complete number of words taken by an instruction, including the word holding the word count and opcode, and any optional operands. An instruction's word count is the total space taken by the instruction.

*Instruction:* After a header, a module is simply a linear list of instructions. An instruction contains a word count, an opcode, an optional Result <id>, an optional <id> of the instruction's type, and a variable list of operands. All instruction opcodes and semantics are listed in Instructions.

*Decoration:* Auxiliary information such as built-in variable, stream numbers, invariance, interpolation type, relaxed precision, etc., added to <id>s or structure-type members through Decorations. Decorations are enumerated in Decoration in the Binary Form section.

Object: An instantiation of a non-void type, either as the Result <id> of an operation, or created through OpVariable.

*Memory Object:* An object created through OpVariable. Such an object can die on function exit, if it was a function variable, or exist for the duration of an entry point.

*Intermediate Object* or *Intermediate Value* or *Intermediate Result*: An object created by an operation (not memory allocated by OpVariable) and dying on its last consumption.

Constant Instruction: Either a specialization-constant instruction or a fixed constant instruction: Instructions that start "OpConstant" or "OpSpec".

[a, b]: This square-bracket notation means the range from a to b, inclusive of a and b. Parenthesis exclude their end point, so, for example, (a, b] means a to b excluding a but including b.

## **2.2.2 Types**

*Boolean type:* The type returned by OpTypeBool.

*Integer type:* Any width signed or unsigned type from OpTypeInt. By convention, the lowest-order bit will be referred to as bit-number 0, and the highest-order bit as bit-number *Width* - 1.

Floating-point type: Any width type from OpTypeFloat.

Numerical type: An integer type or a floating-point type.

*Scalar:* A single instance of a numerical type or Boolean type. Scalars will also be called *components* when being discussed either by themselves or in the context of the contents of a vector.

*Vector:* An ordered homogeneous collection of two or more scalars. Vector sizes are quite restrictive and dependent on the execution model.

*Matrix*: An ordered homogeneous collection of vectors. When vectors are part of a matrix, they will also be called *columns*. Matrix sizes are quite restrictive and dependent on the execution model.

*Array:* An ordered homogeneous collection of any non-void-type objects. When an object is part of an array, it will also be called an *element*. Array sizes are generally not restricted.

*Structure:* An ordered heterogeneous collection of any non-void types. When an object is part of a structure, it will also be called a *member*.

Aggregate: A structure or an array.

Composite: An aggregate, a matrix, or a vector.

*Image:* A traditional texture or image; SPIR-V has this single name for these. An image type is declared with OpTypeImage. An image does not include any information about how to access, filter, or sample it.

*Sampler:* Settings that describe how to access, filter, or sample an image. Can come either from literal declarations of settings or be an opaque reference to externally bound settings. A sampler does not include an image.

Sampled Image: An image combined with a sampler, enabling filtered accesses of the image's contents.

Concrete Type: A numerical scalar, vector, or matrix type, or OpTypePointer when using a **Physical** addressing model, or any aggregate containing only these types.

Abstract Type: An OpTypeVoid or OpTypeBool, or OpTypePointer when using the **Logical** addressing model, or any aggregate type containing any of these.

Opaque Type: A type that is, or contains, or points to, or contains pointers to, any of the following types:

- OpTypeImage
- OpTypeSampler
- OpTypeSampledImage
- OpTypeOpaque
- OpTypeEvent
- OpTypeDeviceEvent
- OpTypeReserveId
- OpTypeQueue
- OpTypePipe
- OpTypeForwardPointer
- OpTypePipeStorage
- OpTypeNamedBarrier

#### 2.2.3 Module

Module: A single unit of SPIR-V. It can contain multiple entry points, but only one set of capabilities.

*Entry Point:* A function in a module where execution begins. A single *entry point* is limited to a single execution model. An entry point is declared using OpEntryPoint.

Execution Model: A graphical-pipeline stage or OpenCL kernel. These are enumerated in Execution Model.

*Execution Mode:* Modes of operation relating to the interface or execution environment of the module. These are enumerated in Execution Mode. Generally, modes do not change the semantics of instructions within a SPIR-V module.

*Vertex Processor*: Any stage or execution model that processes vertices: Vertex, tessellation control, tessellation evaluation, and geometry. Explicitly excludes fragment and compute execution models.

#### 2.2.4 Control Flow

*Block*: A contiguous sequence of instructions starting with an OpLabel, ending with a termination instruction. A *block* has no additional label or termination instructions.

Branch Instruction: One of the following, used as a termination instruction:

- OpBranch
- OpBranchConditional
- OpSwitch
- OpReturn
- OpReturnValue

Termination Instruction: One of the following, used to terminate blocks:

- · any branch instruction
- OpKill
- OpUnreachable

*Dominate*: A block A dominates a block B, where A and B are in the same function, if every path from the function's entry point to block B includes block A. A strictly dominates B only if A dominates B and A and B are different blocks.

*Post Dominate*: A block *B* post dominates a block *A*, where *A* and *B* are in the same function, if every path from *A* to a function-return instruction goes through block *B*.

Control-Flow Graph: The graph formed by a function's blocks and branches. The blocks are the graph's nodes, and the branches the graph's edges.

CFG: Control-flow graph.

*Back Edge*: If a depth-first traversal is done on a function's CFG, starting from the first block of the function, a *back edge* is a branch to a previously visited block. A *back-edge block* is the block containing such a branch.

Merge Instruction: One of the following, used before a branch instruction to declare structured control flow:

- OpSelectionMerge
- OpLoopMerge

Header Block: A block containing a merge instruction.

Loop Header: A header block whose merge instruction is an OpLoopMerge.

Merge Block: A block declared by the Merge Block operand of a merge instruction.

Break Block: A block containing a branch to the Merge Block of a loop header's merge instruction.

Continue Block: A block containing a branch to an OpLoopMerge instruction's Continue Target.

Return Block: A block containing an OpReturn or OpReturn Value branch.

*Invocation*: A single execution of an entry point in a SPIR-V module, operating only on the amount of data explicitly exposed by the semantics of the instructions. (Any implicit operation on additional instances of data would comprise additional invocations.) For example, in compute execution models, a single invocation operates only on a single work item, or, in a vertex execution model, a single invocation operates only on a single vertex.

*Subgroup*: Invocations are partitioned into subgroups, where invocations within a subgroup can synchronize and share data with each other efficiently. In compute models, the current workgroup is a superset of the subgroup.

*Invocation Group*: The complete set of invocations collectively processing a particular compute workgroup or graphical operation, where the scope of a "graphical operation" is implementation dependent, but at least as large as a single point, line, triangle, or patch, and at most as large as a single rendering command, as defined by the client API.

*Derivative Group*: Defined only for the **Fragment** Execution Model: The set of invocations collectively processing a single point, line, or triangle, including any helper invocations.

*Dynamic Instance*: Within a single invocation, a single static instruction can be executed multiple times, giving multiple dynamic instances of that instruction. This can happen when the instruction is executed in a loop, or in a function called from multiple call sites, or combinations of multiple of these. Different loop iterations and different dynamic function-call-site chains yield different dynamic instances of such an instruction. Dynamic instances are distinguished by the control-flow path within an invocation, not by which invocation executed it. That is, different invocations of an entry point execute the same dynamic instances of an instruction when they follow the same control-flow path, starting from that entry point.

*Dynamically Uniform*: An <id> is dynamically uniform for a dynamic instance consuming it when its value is the same for all invocations (in the invocation group) that execute that dynamic instance.

Uniform Control Flow: Uniform control flow (or converged control flow) occurs when all invocations in the invocation group or derivative group execute the same control-flow path (and hence the same sequence of dynamic instances of instructions). Uniform control flow is the initial state at the entry point, and lasts until a conditional branch takes different control paths for different invocations (non-uniform or divergent control flow). Such divergence can reconverge, with all the invocations once again executing the same control-flow path, and this re-establishes the existence of uniform control flow. If control flow is uniform upon entry into a header block, and all invocations leave that dynamic instance of the header block's control-flow construct via the header block's declared merge block, then control flow reconverges to be uniform at that merge block.

# 2.3 Physical Layout of a SPIR-V Module and Instruction

A SPIR-V module is a single linear stream of words. The first words are shown in the following table:

Table 1: First Words of Physical Layout

Word	Contents	
Number		
0	Magic Number.	
1	Version number. The bytes are, high-order to low-order:	
	0   Major Number   Minor Number   0	
	Hence, version 1.2 is the value 0x00010200.	
2	Generator's magic number. It is associated with the tool that generated	
	the module. Its value does not affect any semantics, and is allowed to be	
	0. Using a non-0 value is encouraged, and can be registered with	
	Khronos at https://www.khronos.org/registry/spir-v/api/spir-v.xml.	
3	3 Bound; where all <id>s in this module are guaranteed to satisfy</id>	
	0 < id < Bound	
	Bound should be small, smaller is better, with all <id> in a module be</id>	
	densely packed and near 0.	
4	0 (Reserved for instruction schema, if needed.)	
5	First word of instruction stream, see below.	

All remaining words are a linear sequence of instructions.

Each instruction is a stream of words:

Table 2: Instruction Physical Layout

Instruction	Contents
Word Number	
0	Opcode: The 16 high-order bits are the WordCount of the
	instruction. The 16 low-order bits are the opcode enumerant.
1	Optional instruction type <id> (presence determined by opcode).</id>
. Optional instruction Result <id> (presence determined b</id>	
	opcode).
	Operand 1 (if needed)
	Operand 2 (if needed)
WordCount - 1	Operand <i>N</i> ( <i>N</i> is determined by WordCount minus the 1 to 3
	words used for the opcode, instruction type $\langle id \rangle$ , and instruction
Result <id>).</id>	

Instructions are variable length due both to having optional instruction type <*id*> and *Result* <*id*> words as well as a variable number of operands. The details for each specific instruction are given in the Binary Form section.

## 2.4 Logical Layout of a Module

The instructions of a SPIR-V module must be in the following order. For sections earlier than function definitions, it is invalid to use instructions other than those indicated.

- 1. All OpCapability instructions.
- 2. Optional OpExtension instructions (extensions to SPIR-V).
- 3. Optional OpExtInstImport instructions.
- 4. The single required OpMemoryModel instruction.
- 5. All entry point declarations, using OpEntryPoint.
- 6. All execution mode declarations, using OpExecutionMode.
- 7. These debug instructions, which must be grouped in the following order:
  - a. all OpString, OpSourceExtension, OpSource, and OpSourceContinued, without forward references.
  - b. all OpName and all OpMemberName
  - c. all OpModuleProcessed instructions
- 8. All annotation instructions:
  - a. all decoration instructions (OpDecorate, OpMemberDecorate, OpGroupDecorate, OpGroupMemberDecorate, and OpDecorationGroup).
- 9. All type declarations (OpTypeXXX instructions), all constant instructions, and all global variable declarations (all OpVariable instructions whose Storage Class is not Function). This is the preferred location for OpUndef instructions, though they can also appear in function bodies. All operands in all these instructions must be declared before being used. Otherwise, they can be in any order. This section is the first section to allow use of OpLine debug information.
- 10. All function declarations ("declarations" are functions without a body; there is no forward declaration to a function with a body). A function declaration is as follows.
  - a. Function declaration, using OpFunction.
  - b. Function parameter declarations, using OpFunctionParameter.
  - c. Function end, using OpFunctionEnd.
- 11. All function definitions (functions with a body). A function definition is as follows.
  - a. Function definition, using OpFunction.
  - b. Function parameter declarations, using OpFunctionParameter.
  - c. Block
  - d. Block
  - e. ...
  - f. Function end, using OpFunctionEnd.

## Within a function definition:

- A block always starts with an OpLabel instruction. This may be immediately preceded by an OpLine instruction, but the **OpLabel** is considered as the beginning of the block.
- A block always ends with a termination instruction (see validation rules for more detail).
- All OpVariable instructions in a function must have a Storage Class of Function.
- All OpVariable instructions in a function must be in the first block in the function. These instructions, together with any immediately preceding OpLine instructions, must be the first instructions in that block. (Note the validation rules prevent OpPhi instructions in the first block of a function.)

• A function definition (starts with OpFunction) can be immediately preceded by an OpLine instruction.

Forward references (an operand  $\langle id \rangle$  that appears before the Result  $\langle id \rangle$  defining it) are allowed for:

- Operands that are an OpFunction. This allows for recursion and early declaration of entry points.
- Annotation-instruction operands. This is required to fully know everything about a type or variable once it is declared.
- · Labels.
- Loops can have forward references to a phi function.
- An OpTypeForwardPointer has a forward reference to an OpTypePointer.
- An OpTypeStruct operand that's a forward reference to the *Pointer Type* operand to an OpTypeForwardPointer.
- The list of <id> provided in the OpEntryPoint instruction.

In all cases, there is enough type information to enable a single simple pass through a module to transform it. For example, function calls have all the type information in the call, phi-functions don't change type, and labels don't have type. The pointer forward reference allows structures to contain pointers to themselves or to be mutually recursive (through pointers), without needing additional type information.

The Validation Rules section lists additional rules that must be satisfied.

## 2.5 Instructions

Most instructions create a Result <id>, as provided in the Result <id> field of the instruction. These Result <id>s are then referred to by other instructions through their <id> operands. All instruction operands are specified in the Binary Form section.

Instructions are explicit about whether they require immediates, rather than an  $\langle id \rangle$  referring to some other result. This is strictly known just from the opcode.

- An immediate 32-bit (or smaller) integer is always one operand directly holding a 32-bit two's-complement value.
- An immediate 32-bit float is always one operand, directly holding a 32-bit IEEE 754 floating-point representation.
- An immediate 64-bit float is always two operands, directly holding a 64-bit IEEE 754 representation. The low-order 32 bits appear in the first operand.

## 2.5.1 SSA Form

A module is always in static single assignment (SSA) form. That is, there is always exactly one instruction resulting in any particular Result <id>. Storing into variables declared in memory is not subject to this; such stores do not create *Result* <*id*>*s*. Accessing declared variables is done through:

- OpVariable to allocate an object in memory and create a Result <id> that is the name of a pointer to it.
- OpAccessChain or OpInBoundsAccessChain to create a pointer to a subpart of a composite object in memory.
- OpLoad through a pointer, giving the loaded object a *Result <id>* that can then be used as an operand in other instructions.
- OpStore through a pointer, to write a value. There is no Result <id> for an OpStore.

OpLoad and OpStore instructions can often be eliminated, using intermediate results instead. When this happens in multiple control-flow paths, these values need to be merged again at the path's merge point. Use OpPhi to merge such values together.

## 2.6 Entry Point and Execution Model

The OpEntryPoint instruction identifies an entry point with two key things: an execution model and a function definition. Execution models include **Vertex**, **GLCompute**, etc. (one for each graphical stage), as well as **Kernel** for OpenCL kernels. For the complete list, see Execution Model. An OpEntryPoint also supplies a name that can be used externally to identify the entry point, and a declaration of all the **Input** and **Output** variables that form its input/output interface.

The static function call graphs rooted at two entry points are allowed to overlap, so that function definitions and global variable definitions can be shared. The execution model and any execution modes associated with an entry point apply to the entire static function call graph rooted at that entry point. This rule implies that a function appearing in both call graphs of two distinct entry points may behave differently in each case. Similarly, variables whose semantics depend on properties of an entry point, e.g. those using the **Input Storage Class**, may behave differently when used in call graphs rooted in two different entry points.

#### 2.7 Execution Modes

Information like the following is declared with OpExecutionMode instructions. For example,

- number of invocations (Invocations)
- vertex-order CCW (VertexOrderCcw)
- triangle strip generation (OutputTriangleStrip)
- number of output vertices (OutputVertices)
- etc.

For a complete list, see Execution Mode.

## 2.8 Types and Variables

Types are built up hierarchically, using OpTypeXXX instructions. The Result <id> of an OpTypeXXX instruction becomes a type <id> for future use where type <id>s are needed (therefore, OpTypeXXX instructions do not have a type <id>, like most other instructions do).

The "leaves" to start building with are types like OpTypeFloat, OpTypeInt, OpTypeImage, OpTypeEvent, etc. Other types are built up from the *Result* <*id*> of these. The numerical types are parameterized to specify bit width and signed vs. unsigned.

Higher-level types are then constructed using opcodes like OpTypeVector, OpTypeMatrix, OpTypeImage, OpTypeArray, OpTypeRuntimeArray, OpTypeStruct, and OpTypePointer. These are parameterized by number of components, array size, member lists, etc. The image types are parameterized by the return type, dimensionality, arrayness, etc. To do sampling or filtering operations, a type from OpTypeSampledImage is used that contains both an image and a sampler. Such a sampled image can be set directly by the API, or combined in a SPIR-V module from an independent image and an independent sampler.

Types are built bottom up: A parameterizing operand in a type must be defined before being used.

Some additional information about the type of an <*id>* can be provided using the decoration instructions (OpDecorate, OpMemberDecorate, OpGroupMemberDecorate, and OpDecorationGroup). These can add, for example, **Invariant** to an <*id>* created by another instruction. See the full list of Decorations in the Binary Form section.

Two different type < id > s form, by definition, two different types. It is valid to declare multiple aggregate type < id > s having the same opcode and operands. This is to allow multiple instances of aggregate types with the same structure to be decorated differently. (Different decorations are not required; two different aggregate type < id > s are allowed to have identical declarations and decorations, and will still be two different types.) Pointer types are also allowed to have multiple < id > s for the same opcode and operands, to allow for differing decorations (e.g., **Volatile**) or different decoration values

(e.g., different *Array Stride* values for the **ArrayStride**). When new pointers are formed, their types must be decorated as needed, so the consumer knows how to generate an access through the pointer. Non-aggregate non-pointer types are different: It is invalid to declare multiple type <*id*>*s* for the same scalar, vector, or matrix type. That is, non-aggregate non-pointer type declarations must all have different opcodes or operands. (Note that non-aggregate non-pointer types cannot be decorated in ways that affect their type.)

Variables are declared to be of an already built type, and placed in a Storage Class. Storage classes include **UniformConstant**, **Input**, **Workgroup**, etc. and are fully specified in Storage Class. Variables declared with the **Function** Storage Class can have their lifetime's specified within their function using the OpLifetimeStart and OpLifetimeStop instructions.

Intermediate results are typed by the instruction's type <id>, which must validate with respect to the operation being done.

Built-in variables needing special driver handling (having unique semantics) are declared using OpDecorate or OpMemberDecorate with the **BuiltIn** Decoration, followed by a BuiltIn enumerant. This decoration is applied to a variable or a structure-type member.

## 2.9 Function Calling

To call a function defined in the current module or a function declared to be imported from another module, use OpFunctionCall with an operand that is the  $\langle id \rangle$  of the OpFunction to call, and the  $\langle id \rangle$ s of the arguments to pass. All arguments are passed by value into the called function. This includes pointers, through which a callee object could be modified.

## 2.10 Extended Instruction Sets

Many operations and/or built-in function calls from high-level languages are represented through *extended instruction sets*. Extended instruction sets will include things like

- trigonometric functions:  $\sin(), \cos(), \dots$
- exponentiation functions: exp(), pow(), ...
- geometry functions: reflect(), smoothstep(), ...
- · functions having rich performance/accuracy trade-offs
- etc.

Non-extended instructions, those that are core SPIR-V instructions, are listed in the Binary Form section. Native operations include:

- Basic arithmetic: +, -, \*, min(), scalar \* vector, etc.
- Texturing, to help with back-end decoding and support special code-motion rules.
- Derivatives, due to special code-motion rules.

Extended instruction sets are specified in independent specifications. They can be referenced (but not specified) in this specification. The separate extended instruction set specification will specify instruction opcodes, semantics, and instruction names.

To use an extended instruction set, first import it by name string using OpExtInstImport and giving it a Result <id>:

```
<extinst-id> OpExtInstImport "name-of-extended-instruction-set"
```

The "name-of-extended-instruction-set" is a literal string. The standard convention for this string is

```
"<source language name>.<package name>.<version>"
```

For example "GLSL.std.450" could be the name of the core built-in functions for GLSL versions 450 and earlier.

#### Note

There is nothing precluding having two "mirror" sets of instructions with different names but the same opcode values, which could, for example, let modifying just the import statement to change a performance/accuracy trade off.

Then, to call a specific extended instruction, use OpExtInst:

```
OpExtInst <extinst-id> instruction-number operand0, operand1, ...
```

Extended instruction-set specifications will provide semantics for each "instruction-number". It is up to the specific specification what the overloading rules are on operand type. The specification must be clear on its semantics, and producers/consumers of it must follow those semantics.

By convention, it is recommended that all external specifications include an **enum** {...} listing all the "instruction-numbers", and a mapping between these numbers and a string representing the instruction name. However, there are no requirements that instruction name strings are provided or mangled.

#### Note

Producing and consuming extended instructions can be done entirely through numbers (no string parsing). An extended instruction set specification provides opcode enumerant values for the instructions, and these will be produced by the front end and consumed by the back end.

## 2.11 Structured Control Flow

SPIR-V can explicitly declare structured control-flow *constructs* using merge instructions. These explicitly declare a header block before the control flow diverges and a merge block where control flow subsequently converges. These blocks delimit constructs that must nest, and can only be entered and exited in structured ways, as per the following.

Structured control-flow declarations must satisfy the following rules:

- the merge block declared by a header block cannot be a merge block declared by any other header block
- each header block must strictly dominate its merge block, unless the merge block is unreachable in the CFG
- all CFG back edges must branch to a loop header, with each loop header having exactly one back edge branching to it
- for a given loop header, its OpLoopMerge Continue Target, and corresponding back-edge block:
  - the loop header must dominate the Continue Target, unless the Continue Target is unreachable in the CFG
  - the Continue Target must dominate the back-edge block
  - the back-edge block must post dominate the Continue Target

A structured control-flow *construct* is then defined as one of:

- a *selection construct*: the set of blocks dominated by a selection header, minus the set of blocks dominated by the header's merge block
- a *continue construct*: the set of blocks dominated by an OpLoopMerge's *Continue Target* and post dominated by the corresponding back-edge block
- a *loop construct*: the set of blocks dominated by a *loop header*, minus the set of blocks dominated by the loop's merge block, minus the loop's corresponding *continue construct*
- a case construct: the set of blocks dominated by an OpSwitch Target or Default, minus the set of blocks dominated by the OpSwitch's merge block (this construct is only defined for those OpSwitch Target or Default that are not equal to the OpSwitch's corresponding merge block)

The above structured control-flow constructs must satisfy the following rules:

- if a construct contains another header block, then it also contains that header's corresponding merge block
- the only blocks in a construct that can branch outside the construct are
  - a block branching to the construct's merge block
  - a block branching from one *case construct* to another, for the same **OpSwitch**
  - a continue block for the innermost loop it is nested inside of
  - a break block for the innermost loop it is nested inside of
  - a return block
- additionally for switches:
  - an **OpSwitch** block dominates all its defined case constructs
  - each case construct has at most one branch to another case construct
  - each case construct is branched to by at most one other case construct
  - if Target T1 branches to Target T2, or if Target T1 branches to the Default and the Default branches to Target T2, then
     T1 must immediately precede T2 in the list of the OpSwitch Target operands

## 2.12 Specialization

*Specialization* is intended for constant objects that will not have known constant values until after initial generation of a SPIR-V module. Such objects are called *specialization constants*.

A SPIR-V module containing specialization constants can consume one or more externally provided *specializations*: A set of final constant values for some subset of the module's *specialization constants*. Applying these final constant values yields a new module having fewer remaining specialization constants. A module also contains default values for any specialization constants that never get externally specialized.

#### Note

No optimizing transforms are required to make a *specialized* module functionally correct. The specializing transform is straightforward and explicitly defined below.

## Note

Ad hoc specializing should not be done through constants (OpConstant or OpConstantComposite) that get overwritten: A SPIR-V  $\rightarrow$  SPIR-V transform might want to do something irreversible with the value of such a constant, unconstrained from the possibility that its value could be later changed.

Within a module, a *Specialization Constant* is declared with one of these instructions:

- OpSpecConstantTrue
- OpSpecConstantFalse
- OpSpecConstant
- OpSpecConstantComposite
- OpSpecConstantOp

The literal operands to OpSpecConstant are the default numerical specialization constants. Similarly, the "True" and "False" parts of OpSpecConstantTrue and OpSpecConstantFalse provide the default Boolean specialization constants. These default values make an external specialization optional. However, such a default constant is applied only after all external specializations are complete, and none contained a specialization for it.

An external specialization is provided as a logical list of pairs. Each pair is a **SpecId** Decoration of a scalar specialization instruction along with its specialization constant. The numeric values are exactly what the operands would be to a corresponding OpConstant instruction. Boolean values are true if non-zero and false if zero.

Specializing a module is straightforward. The following specialization-constant instructions can be updated with specialization constants, and replaced in place, leaving everything else in the module exactly the same:

```
OpSpecConstantTrue -> OpConstantTrue or OpConstantFalse
OpSpecConstantFalse -> OpConstantTrue or OpConstantFalse
OpSpecConstant -> OpConstant
OpSpecConstantComposite -> OpConstantComposite
```

The OpSpecConstantOp instruction is specialized by executing the operation and replacing the instruction with the result. The result can be expressed in terms of a constant instruction that is not a specialization-constant instruction. (Note, however, this resulting instruction might not have the same size as the original instruction, so is not a "replaced in place" operation.)

When applying an external specialization, the following (and only the following) must be modified to be non-specialization-constant instructions:

- · specialization-constant instructions with values provided by the specialization
- specialization-constant instructions that consume nothing but non-specialization constant instructions (including those that the partial specialization transformed from specialization-constant instructions; these are in order, so it is a single pass to do so)

A full specialization can also be done, when requested or required, in which all specialization-constant instructions will be modified to non-specialization-constant instructions, using the default values where required.

## 2.13 Linkage

The ability to have partially linked modules and libraries is provided as part of the Linkage capability.

By default, functions and global variables are private to a module and cannot be accessed by other modules. However, a module may be written to *export* or *import* functions and global (module scope) variables. Imported functions and global variable definitions are resolved at linkage time. A module is considered to be partially linked if it depends on imported values.

Within a module, imported or exported values are decorated using the **Linkage Attributes** Decoration. This decoration assigns the following linkage attributes to decorated values:

- A Linkage Type.
- A name, which is a Literal String, and is used to uniquely identify exported values.

#### Note

When resolving imported functions, the Function Control and all Function Parameter Attributes are taken from the function definition, and not from the function declaration.

#### 2.14 Relaxed Precision

The **RelaxedPrecision** Decoration allows 32-bit integer and 32-bit floating-point operations to execute with a relaxed precision of somewhere between 16 and 32 bits.

For a floating-point operation, operating at relaxed precision means that the minimum requirements for range and precision are as follows:

- the floating point range may be as small as  $(-2^{14}, 2^{14})$
- the floating point magnitude range may be as small as  $(2^{-14}, 2^{14})$
- the relative floating point precision may be as small as 2<sup>-10</sup>

Relative floating-point precision is defined as the worst case (i.e. largest) ratio of the smallest step in relation to the value for all non-zero values:

```
Precision<sub>relative</sub> = (abs(v_1 - v_2)_{min} / abs(v_1))_{max} for v_1 \neq 0, v_2 \neq 0, v_1 \neq v_2
```

For integer operations, operating at relaxed precision means that the operation will be evaluated by an operation in which, for some N,  $16 \le N \le 32$ :

- the operation is executed as though its type were N bits in size, and
- the result is zero or sign extended to 32 bits as determined by the signedness of the result type of the operation.

The **RelaxedPrecision** Decoration can be applied to:

- The <id> of a variable, where the variable's type is a scalar, vector, or matrix, or an array of scalar, vector, or matrix. In all cases, the components in the type must be a 32-bit numerical type.
- The Result <id> of an instruction that operates on numerical types, meaning the instruction is to operate at relaxed precision.
- The Result <id> of an instruction that reads or filters from an image. E.g. OpImageSampleExplicitLod, meaning the instruction is to operate at relaxed precision.
- The Result <id> of an OpFunction meaning the function's returned result is at relaxed precision. It cannot be applied to OpTypeFunction or to an **OpFunction** whose return type is **OpTypeVoid**.
- A structure-type member (through OpMemberDecorate).

When applied to a variable or structure member, all loads and stores from the decorated object may be treated as though they were decorated with **RelaxedPrecision**. Loads may also be decorated with **RelaxedPrecision**, in which case they are treated as operating at relaxed precision.

All loads and stores involving relaxed precision still read and write 32 bits of data, respectively. Floating-point data read or written in such a manner is written in full 32-bit floating-point format. However, a load or store might reduce the precision (as allowed by **RelaxedPrecision**) of the destination value.

For debugging portability of floating-point operations, OpQuantizeToF16 may be used to explicitly reduce the precision of a relaxed-precision result to 16-bit precision. (Integer-result precision can be reduced, for example, using left- and right-shift opcodes.)

For image-sampling operations, decorations can appear on both the sampling instruction and the image variable being sampled. If either is decorated, they both should be decorated, and when both are decorated their decorations must match. If only one is decorated, the sampling instruction can behave either as if both were decorated or neither were decorated.

## 2.15 Debug Information

Debug information is supplied with:

- Source-code text through OpString, OpSource, and OpSourceContinued.
- Object names through OpName and OpMemberName.
- Line numbers through OpLine.

A module will not lose any semantics when all such instructions are removed.

## 2.15.1 Function-Name Mangling

There is no functional dependency on how functions are named. Signature-typing information is explicitly provided, without any need for name "unmangling". (Valid modules can be created without inclusion of mangled names.)

By convention, for debugging purposes, modules with OpSource Source Language of OpenCL use the Itanium name-mangling standard.

## 2.16 Validation Rules

#### 2.16.1 Universal Validation Rules

All modules must obey the following, or it is an invalid module:

- The stream of instructions must be ordered as described in the Logical Layout section.
- Any use of a feature described by a capability in the capability section requires that capability to be declared, either directly, or as a "depends on" capability on a capability that is declared.
- Non-structure types (scalars, vectors, arrays, etc.) with the same operand parameterization cannot be type aliases. For non-structures, two type <*id*>*s* match if-and-only-if the types match.
- If the **Logical** addressing model is selected:
  - OpVariable cannot allocate an object whose type is a pointer type (that is, it cannot create an object in memory that is
    itself a pointer and whose result would thus be a pointer to a pointer)
  - A pointer can only be an operand to the following instructions:
    - \* OpLoad
    - \* OpStore
    - \* OpAccessChain
    - \* OpInBoundsAccessChain
    - \* OpFunctionCall
    - \* OpImageTexelPointer
    - \* OpCopyMemory
    - \* OpCopyObject
    - \* all OpAtomic instructions
    - \* extended instruction-set instructions that are explicitly identified as taking pointer operands
  - A pointer can be the Result <id> of only the following instructions:
    - \* OpVariable
    - \* OpAccessChain
    - \* OpInBoundsAccessChain
    - \* OpFunctionParameter
    - \* OpImageTexelPointer

## \* OpCopyObject

All indexes in OpAccessChain and OpInBoundsAccessChain that are OpConstant with type of OpTypeInt with a signedness of 1 must not have their sign bit set.

#### SSA

- Each <id> must appear exactly once as the Result <id> of an instruction.
- The definition of an SSA <id> should dominate all uses of it, with the following exceptions:
  - \* Function calls may call functions not yet defined. However, note that the function's argument and return types will already be known at the call site.
  - \* Uses in a phi-function in a loop may consume definitions in the loop that don't dominate the use.

## · Entry point and execution model

- There is at least one OpEntryPoint instruction, unless the Linkage capability is being used.
- No function can be targeted by both an OpEntryPoint instruction and an OpFunctionCall instruction.

#### Functions

- A function declaration (an OpFunction with no basic blocks), must have a Linkage Attributes Decoration with the Import Linkage Type.
- A function definition (an OpFunction with basic blocks) cannot be decorated with the **Import** Linkage Type.
- A function cannot have both a declaration and a definition (no forward declarations).
- Global (Module Scope) Variables
  - It is illegal to initialize an imported variable. This means that a module-scope OpVariable with initialization value cannot be marked with the **Import** Linkage Type.
- Control-Flow Graph (CFG)
  - Blocks exist only within a function.
  - The first block in a function definition is the entry point of that function and cannot be the target of any branch. (Note this means it will have no OpPhi instructions.)
  - The order of blocks in a function must satisfy the rule that blocks appear before all blocks they dominate.
  - Each block starts with a label.
    - \* A label is made by OpLabel.
    - \* This includes the first block of a function (**OpFunction** is not a label).
    - \* Labels are used only to form blocks.
  - The last instruction of each block is a termination instruction.
  - Termination instructions can only appear as the last instruction in a block.
  - OpLabel instructions can only appear within a function.
  - All branches within a function must be to labels in that function.
- All OpFunctionCall Function operands are an <id> of an OpFunction in the same module.
- · Data rules
  - Scalar floating-point types can be parameterized only as 32 bit, plus any additional sizes enabled by capabilities.
  - Scalar integer types can be parameterized only as 32 bit, plus any additional sizes enabled by capabilities.
  - Vector types can only be parameterized with numerical types or the OpTypeBool type.
  - Vector types for can only be parameterized as having 2, 3, or 4 components, plus any additional sizes enabled by capabilities.
  - Matrix types can only be parameterized with floating-point types.

- Matrix types can only be parameterized as having only 2, 3, or 4 columns.
- Specialization constants (see Specialization) are limited to integers, Booleans, floating-point numbers, and vectors of these.
- Forward reference operands in an OpTypeStruct
  - \* must be later declared with OpTypePointer
  - \* the type pointed to must be an OpTypeStruct
  - \* had an earlier OpTypeForwardPointer forward reference to the same <id>
- All OpSampledImage instructions must be in the same block in which their Result <id> are consumed. Result <id> from OpSampledImage instructions must not appear as operands to OpPhi instructions or OpSelect instructions, or any instructions other than the image lookup and query instructions specified to take an operand whose type is OpTypeSampledImage.
- Instructions for extracting a scalar image or scalar sampler out of a composite must only use dynamically-uniform indexes. They must be in the same block in which their *Result <id>* are consumed. Such *Result <id>* must not appear as operands to OpPhi instructions or OpSelect instructions, or any instructions other than the image instructions specified to operate on them.

#### · Decoration rules

- The **Aliased** Decoration can only be applied to intermediate objects that are pointers to non-void types.
- The Linkage Attributes Decoration cannot be applied to functions targeted by an OpEntryPoint instruction.
- A BuiltIn Decoration can only be applied as follows:
  - \* When applied to a structure-type member, all members of that structure type must also be decorated with **BuiltIn**. (No allowed mixing of built-in variables and non-built-in variables within a single structure.)
  - \* When applied to a structure-type member, that structure type cannot be contained as a member of another structure type.
  - \* There is at most one object per Storage Class that can contain a structure type containing members decorated with **BuiltIn**, consumed per entry-point.
- OpLoad and OpStore can only consume objects whose type is a pointer.
- A Result <id> resulting from an instruction within a function can only be used in that function.
- A function call must have the same number of arguments as the function definition (or declaration) has parameters, and their respective types must match.
- An instruction requiring a specific number of operands must have that many operands. The word count must agree.
- Each opcode specifies its own requirements for number and type of operands, and these must be followed.
- Atomic access rules
  - The pointers taken by atomic operation instructions must be a pointer into one of the following Storage Classes:
    - \* Uniform when used with the BufferBlock Decoration
    - \* Workgroup
    - \* CrossWorkgroup
    - \* Generic
    - \* AtomicCounter
    - \* Image

## 2.16.2 Validation Rules for Shader Capabilities

- CFG:
  - Loops must be structured, having an OpLoopMerge instruction in their header.
  - Selections must be structured, having an OpSelectionMerge instruction in their header.

- Entry point and execution model
  - Each entry point in a module, along with its corresponding static call tree within that module, forms a complete pipeline stage.
  - Each OpEntryPoint with the Fragment Execution Model must have an OpExecutionMode for either the OriginLowerLeft or the OriginUpperLeft Execution Mode. (Exactly one of these is required.)
  - An OpEntryPoint with the Fragment Execution Model can set at most one of the DepthGreater, DepthLess, or DepthUnchanged Execution Modes.
  - An OpEntryPoint with one of the Tessellation Execution Modes can set at most one of the SpacingEqual, FractionalEven, or FractionalOdd Execution Modes.
  - An OpEntryPoint with one of the Tessellation Execution Models can set at most one of the Triangles, Quads, or Isolines Execution Modes.
  - An OpEntryPoint with one of the Tessellation Execution Models can set at most one of the VertexOrderCw or VertexOrderCcw Execution Modes.
  - An OpEntryPoint with the Geometry Execution Model must set exactly one of the InputPoints, InputLines, InputLinesAdjacency, Triangles, or TrianglesAdjacency Execution Modes.
  - An OpEntryPoint with the Geometry Execution Model must set exactly one of the OutputPoints, OutputLineStrip, or OutputTriangleStrip Execution Modes.
- Composite objects in the **UniformConstant**, **Uniform**, and **PushConstant** Storage Classes must be explicitly laid out. The following apply to all the aggregate and matrix types describing such an object, recursively through their nested types:
  - Each structure-type member must have an **Offset** Decoration.
  - Each array type must have an **ArrayStride** Decoration.
  - Each structure-type member that is a matrix or array-of-matrices must have be decorated with
    - \* a MatrixStride Decoration, and
    - \* one of the **RowMajor** or **ColMajor** Decorations.
  - The ArrayStride, MatrixStride, and Offset Decorations must be large enough to hold the size of the objects they
    affect (that is, specifying overlap is invalid). Each ArrayStride and MatrixStride must be greater than zero, and no
    two members of a given structure can be assigned to the same Offset.
  - Each **OpPtrAccessChain** must have a *Base* whose type is decorated with **ArrayStride**.
  - When an array-element pointer is derived from an array (e.g., using OpAccessChain), and the resulting
    element-pointer type was decorated with ArrayStride, its Array Stride must match the Array Stride of the originating
    array's type.
- For structure objects in the **Input** and **Output** Storage Classes, the following apply:
  - When applied to structure-type members, the Decorations Noperspective, Flat, Patch, Centroid, and Sample can
    only be applied to the top-level members of the structure type. (Nested objects' types cannot be structures whose
    members are decorated with these decorations.)
- · Decorations
  - At most one of **Noperspective** or **Flat** Decorations can be applied to the same object or member.
  - At most one of **Patch**, **Centroid**, or **Sample** Decorations can be applied to the same object or member.
  - At most one of **RowMajor** and **ColMajor** Decorations can be applied to a structure type.
  - At most one of **Block** and **BufferBlock** Decorations can be applied to a structure type.
- All <id> used for Scope and Memory Semantics must be of an OpConstant.

## 2.16.3 Validation Rules for Kernel Capabilities

• The Signedness in **OpTypeInt** must always be 0.

## 2.17 Universal Limits

These quantities are minimum limits for all implementations and validators. Implementations are allowed to support larger quantities. Specific APIs may impose larger minimums. See Language Capabilities.

Validators must either

- inform when these limits are crossed, or
- be explicitly parameterized with larger limits.

Table 3: Limits

Limited Frakter	Minimum Limit	
Limited Entity	Decimal	Hexadecimal
Characters in a literal string	65,535	FFFF
Instruction word count	65,535	FFFF
Result <id> bound</id>		
	4,194,303	3FFFFF
See Physical Layout for the shader-specific bound.		
Control-flow nesting depth		
Measured per function, in program order, counting the maximum number of OpBranch, OpBranchConditional, or OpSwitch that are seen without yet seeing their corresponding <i>Merge Block</i> , as declared by OpSelectionMerge or OpLoopMerge.	1023	3FF
Global variables (Storage Class other than Function)	65,535	FFFF
Local variables (Function Storage Class)	524,287	7FFFF
Decorations per target <id></id>		of entries in the ecoration table.
Execution modes per entry point	255	FF
Indexes for OpAccessChain,		
OpInBoundsAccessChain, OpPtrAccessChain, OpInBoundsPtrAccessChain, OpCompositeExtract, and OpCompositeInsert	255	FF
Number of function parameters, per function declaration	255	FF
OpFunctionCall actual arguments	255	FF
OpExtInst actual arguments	255	FF
OpSwitch (literal, label) pairs	16,383	3FFF
OpTypeStruct members	16,383	3FFF
Structure nesting depth	255	FF

## 2.18 Memory Model

A memory model is chosen using a single OpMemoryModel instruction near the beginning of the module. This selects both an addressing model and a memory model.

The **Logical** addressing model means pointers are abstract, having no physical size or numeric value. In this mode, pointers can only be created from existing objects, and they cannot be stored into an object.

The non-**Logical** addressing models allow physical pointers to be formed. OpVariable can be used to create objects that hold pointers. These are declared for a specific Storage Class. Pointers for one Storage Class cannot be used to access

objects in another Storage Class. However, they can be converted with conversion opcodes. Any particular addressing model must describe the bit width of pointers for each of the storage classes.

## 2.18.1 Memory Layout

When memory is shared between a SPIR-V module and an API, its contents are transparent, and must be agreed on. For example, the **Offset**, **MatrixStride**, and **ArrayStride** Decorations applied to members of a struct object can partially define how the memory is laid out. In addition, the following are always true, applied recursively as needed, of the offsets within the memory buffer:

- a vector consumes contiguous memory with lower-numbered components appearing in smaller offsets than higher-numbered components, and with component 0 starting at the vector's **Offset** Decoration, if present
- in an array, lower-numbered elements appear at smaller offsets than higher-numbered elements, with element 0 starting at the **Offset** Decoration for the array, if present
- a structure has lower-numbered members appearing at smaller offsets than higher-numbered members, with member 0 starting at the **Offset** Decoration for the structure, if present
- in a matrix, lower-numbered columns appear at smaller offsets than higher-numbered columns, and lower-numbered components within the matrix's vectors appearing at smaller offsets than high-numbered components, with component 0 of column 0 starting at the **Offset** Decoration, if present (the **RowMajor** and **ColMajor** Decorations dictate what is contiguous)

## 2.18.2 Aliasing

Here, aliasing means one of:

- Two or more pointers that point into overlapping parts of the same underlying object. That is, two intermediates, both of which are typed pointers, that can be dereferenced (in bounds) such that both dereferences access the same memory.
- Images, buffers, or other externally allocated objects where a function might access the same underlying memory via accesses to two different objects.

How aliasing is managed depends on the Memory Model:

- The simple and GLSL memory models can assume that aliasing is generally not present. Specifically, the compiler is free to compile as if aliasing is not present, unless a pointer is explicitly indicated to be an alias. This is indicated by applying the **Aliased Decoration** to an *intermediate* object's <id>. Applying **Restrict** is allowed, but has no effect.
- The OpenCL memory models must assume that aliasing is generally present. Specifically, the compiler must compile as if aliasing is present, unless a pointer is explicitly indicated to not alias. This is done by applying the **Restrict** Decoration to an *intermediate* object's <id>. Applying **Aliased** is allowed, but has no effect.

It is invalid to apply both **Restrict** and **Aliased** to the same <*id*>.

#### 2.19 Derivatives

Derivatives appear only in the **Fragment** Execution Model. They can be implicit or explicit. Some image instructions consume implicit derivatives, while the derivative instructions compute explicit derivatives. In all cases, derivatives are well defined only if the derivative group has uniform control flow.

## 2.20 Code Motion

Texturing instructions in the Fragment Execution Model that rely on an implicit derivative cannot be moved into control flow that is not known to be uniform control flow within each derivative group.

## 2.21 Unified Specification

This document specifies all versions of SPIR-V.

There are three kinds of entries in the tables of enumerated tokens:

- **Reservation:** These say Reserved in the enabling capabilities. They often contain token names only, lacking a semantic description. They are invalid **SPIR-V** for any version, serving only to reserve the tokens. They may identify enabling capabilities and extensions, in which case any listed extensions might add the tokens. See the listed extensions for additional information.
- Conditional: These say Missing before in the enabling capabilities. They are invalid SPIR-V for the missing versions. They may identify enabling capabilities and extensions, in which case any listed extensions might add the tokens for some of the missing versions. See the listed extensions for additional information. For versions not identified as missing, the tokens are valid SPIR-V, subject to any listed enabling capabilities.
- Universal: These have no mention of what version they are missing in, or of being reserved. They are valid in all versions of SPIR-V.

# 3 Binary Form

This section contains the exact form for all instructions, starting with the numerical values for all fields. See Physical Layout for the order words appear in.

## 3.1 Magic Number

Magic number for a SPIR-V module.

#### Tip

**Endianness:** A module is defined as a stream of words, not a stream of bytes. However, if stored as a stream of bytes (e.g., in a file), the magic number can be used to deduce what endianness to apply to convert the byte stream back to a word stream.

Magic Number	
0x07230203	

## 3.2 Source Language

The source language is for debug purposes only, with no semantics that affect the meaning of other parts of the module. Used by OpSource.

Source Language		
0	Unknown	
1	ESSL	
2	GLSL	
3	OpenCL_C	
4	OpenCL_CPP	
5	HLSL	

## 3.3 Execution Model

Used by OpEntryPoint.

	Execution Model	Enabling Capabilities
0	Vertex	Shader
	Vertex shading stage.	
1	TessellationControl	Tessellation
	Tessellation control (or hull) shading stage.	
2	TessellationEvaluation	Tessellation
	Tessellation evaluation (or domain) shading	
	stage.	
3	Geometry	Geometry
	Geometry shading stage.	
4	Fragment	Shader
	Fragment shading stage.	
5	GLCompute	Shader
	Graphical compute shading stage.	

<b>Execution Model</b>		Enabling Capabilities
6	Kernel	Kernel
	Compute kernel.	

# 3.4 Addressing Model

Used by OpMemoryModel.

Addressing Model		Enabling Capabilities
0	Logical	
1	Physical32	Addresses
	Indicates a 32-bit module, where the address	
	width is equal to 32 bits.	
2	Physical64	Addresses
	Indicates a 64-bit module, where the address	
	width is equal to 64 bits.	

# 3.5 Memory Model

Used by OpMemoryModel.

Memory Model		Enabling Capabilities	
0	Simple	Shader	
	No shared memory consistency issues.		
1	GLSL450	Shader	
	Memory model needed by later versions of		
	GLSL and ESSL. Works across multiple		
	versions.		
2	OpenCL	Kernel	
	OpenCL memory model.		

## 3.6 Execution Mode

Declare the modes an entry point will execute in. Used by OpExecutionMode and OpExecutionModeId.

	<b>Execution Mode</b>	Enabling Capabilities	Extra Operands
0	Invocations	Geometry	Literal Number
	Number of times to invoke the geometry		Number of invocations
	stage for each input primitive received. The		
	default is to run once for each input		
	primitive. It is invalid to specify a value		
	greater than the target-dependent maximum.		
	Only valid with the <b>Geometry</b> Execution		
	Model.		
1	SpacingEqual	Tessellation	
	Requests the tessellation primitive generator		
	to divide edges into a collection of		
	equal-sized segments. Only valid with one		
	of the tessellation Execution Models.		

	<b>Execution Mode</b>	<b>Enabling Capabilities</b>	Extra Operands
2	SpacingFractionalEven	Tessellation	
	Requests the tessellation primitive generator		
	to divide edges into an even number of		
	equal-length segments plus two additional		
	shorter fractional segments. Only valid with		
	one of the tessellation Execution Models.		
3	SpacingFractionalOdd	Tessellation	
	Requests the tessellation primitive generator		
	to divide edges into an odd number of		
	equal-length segments plus two additional		
	shorter fractional segments. Only valid with		
	one of the tessellation Execution Models.		
4	VertexOrderCw	Tessellation	
	Requests the tessellation primitive generator		
	to generate triangles in clockwise order.		
	Only valid with one of the tessellation		
	Execution Models.		
5	VertexOrderCcw	Tessellation	
	Requests the tessellation primitive generator		
	to generate triangles in counter-clockwise		
	order. Only valid with one of the		
	tessellation Execution Models.		
6	PixelCenterInteger	Shader	
	Pixels appear centered on whole-number		
	pixel offsets. E.g., the coordinate (0.5, 0.5)		
	appears to move to $(0.0, 0.0)$ . Only valid		
	with the <b>Fragment</b> Execution Model. If a		
	<b>Fragment</b> entry point does not have this		
	set, pixels appear centered at offsets of (0.5,		
	0.5) from whole numbers		
7	OriginUpperLeft	Shader	
	Pixel coordinates appear to originate in the		
	upper left, and increase toward the right and		
	downward. Only valid with the <b>Fragment</b>		
	Execution Model.		
8	OriginLowerLeft	Shader	
	Pixel coordinates appear to originate in the		
	lower left, and increase toward the right and		
	upward. Only valid with the <b>Fragment</b>		
	Execution Model.		
9	EarlyFragmentTests	Shader	
	Fragment tests are to be performed before		
	fragment shader execution. Only valid with		
10	the Fragment Execution Model.	TD11 4*	
10	PointMode	Tessellation	
	Requests the tessellation primitive generator		
	to generate a point for each distinct vertex		
	in the subdivided primitive, rather than to		
	generate lines or triangles. Only valid with		
	one of the tessellation Execution Models.		

	Execution Mode	<b>Enabling Capabilities</b>	Extra Op	erands	
11	Xfb	TransformFeedback			
	This stage will run in transform				
	feedback-capturing mode and this module is				
	responsible for describing the				
	transform-feedback setup. See the				
	XfbBuffer, Offset, and XfbStride				
	Decorations.				
12	DepthReplacing	Shader			
	This mode must be declared if this module				
	potentially changes the fragment's depth.				
	Only valid with the <b>Fragment</b> Execution				
	Model.				
14	DepthGreater	Shader			
	External optimizations may assume depth				
	modifications will leave the fragment's				
	depth as greater than or equal to the				
	fragment's interpolated depth value (given				
	by the z component of the <b>FragCoord</b>				
	BuiltIn decorated variable). Only valid with				
	the <b>Fragment</b> Execution Model.				
15	DepthLess	Shader			
	External optimizations may assume depth				
	modifications leave the fragment's depth				
	less than the fragment's interpolated depth				
	value, (given by the z component of the				
	FragCoord BuiltIn decorated variable).				
	Only valid with the <b>Fragment</b> Execution				
1.6	Model.	GI I			
16	DepthUnchanged	Shader			
	External optimizations may assume this				
	stage did not modify the fragment's depth.  However, <b>DepthReplacing</b> mode must				
	accurately represent depth modification. Only valid with the <b>Fragment</b> Execution				
	Model.				
17	LocalSize		Literal	Literal	Literal
17	Indicates the work-group size in the $x$ , $y$ ,		Number	Number	Number
	and $z$ dimensions. Only valid with the		x size	y size	z size
	GLCompute or Kernel Execution Models.		1 3126	y size	ر مارد
18	LocalSizeHint	Kernel	Literal	Literal	Literal
	A hint to the compiler, which indicates the	11011101	Number	Number	Number
	most likely to be used work-group size in		x size	y size	z size
	the $x$ , $y$ , and $z$ dimensions. Only valid with			,	~~~~
	the <b>Kernel</b> Execution Model.				
19	InputPoints	Geometry		<u>I</u>	I.
	Stage input primitive is <i>points</i> . Only valid	•			
	with the <b>Geometry</b> Execution Model.				
20	InputLines	Geometry			
	Stage input primitive is <i>lines</i> . Only valid	-			
	with the <b>Geometry</b> Execution Model.				
21	InputLinesAdjacency	Geometry			
	Stage input primitive is <i>lines adjacency</i> .				
	Only valid with the <b>Geometry</b> Execution				
	Model.				

	<b>Execution Mode</b>	<b>Enabling Capabilities</b>	Extra Operands
22	Triangles	Geometry, Tessellation	
	For a geometry stage, input primitive is		
	<i>triangles</i> . For a tessellation stage, requests		
	the tessellation primitive generator to		
	generate triangles. Only valid with the		
	<b>Geometry</b> or one of the tessellation		
	Execution Models.		
23	InputTrianglesAdjacency	Geometry	
	Geometry stage input primitive is <i>triangles</i>		
	adjacency. Only valid with the <b>Geometry</b>		
	Execution Model.		
24	Quads	Tessellation	
	Requests the tessellation primitive generator		
	to generate <i>quads</i> . Only valid with one of		
	the tessellation Execution Models.		
25	Isolines	Tessellation	
	Requests the tessellation primitive generator		
	to generate <i>isolines</i> . Only valid with one of		
	the tessellation Execution Models.		
26	OutputVertices	Geometry, Tessellation	Literal Number
	For a geometry stage, the maximum number		Vertex count
	of vertices the shader will ever emit in a		
	single invocation. For a tessellation-control		
	single invocation. For a tessellation-control stage, the number of vertices in the output		
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control		
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of		
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is		
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or		
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.		
27	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.  OutputPoints	Geometry	
27	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.  OutputPoints Stage output primitive is <i>points</i> . Only valid	Geometry	
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.  OutputPoints Stage output primitive is <i>points</i> . Only valid with the <b>Geometry</b> Execution Model.		
27	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.  OutputPoints Stage output primitive is <i>points</i> . Only valid with the <b>Geometry</b> Execution Model.  OutputLineStrip	Geometry	
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.  OutputPoints Stage output primitive is points. Only valid with the <b>Geometry</b> Execution Model.  OutputLineStrip Stage output primitive is line strip. Only		
28	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the <b>Geometry</b> or one of the tessellation Execution Models.  OutputPoints Stage output primitive is points. Only valid with the Geometry Execution Model.  OutputLineStrip Stage output primitive is line strip. Only valid with the Geometry Execution Model.	Geometry	
	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the Geometry or one of the tessellation Execution Models.  OutputPoints Stage output primitive is points. Only valid with the Geometry Execution Model.  OutputLineStrip Stage output primitive is line strip. Only valid with the Geometry Execution Model.  OutputTriangleStrip		
28	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the Geometry or one of the tessellation Execution Models.  OutputPoints Stage output primitive is points. Only valid with the Geometry Execution Model.  OutputLineStrip Stage output primitive is line strip. Only valid with the Geometry Execution Model.  OutputTriangleStrip Stage output primitive is triangle strip.	Geometry	
28	single invocation. For a tessellation-control stage, the number of vertices in the output patch produced by the tessellation control shader, which also specifies the number of times the tessellation control shader is invoked. Only valid with the Geometry or one of the tessellation Execution Models.  OutputPoints Stage output primitive is points. Only valid with the Geometry Execution Model.  OutputLineStrip Stage output primitive is line strip. Only valid with the Geometry Execution Model.  OutputTriangleStrip	Geometry	

	<b>Execution Mode</b>	<b>Enabling Capabilities</b>	Extra O	perands	
30	VecTypeHint	Kernel	Literal N		
	A hint to the compiler, which indicates that		Vector ty	ре	
	most operations used in the entry point are			-	
	explicitly vectorized using a particular				
	vector type. The 16 high-order bits of				
	Vector Type operand specify the number of				
	components of the vector. The 16 low-order				
	bits of <i>Vector Type</i> operand specify the <i>data</i>				
	type of the vector.				
	These are the legal <i>data type</i> values:				
	0 represents an 8-bit integer value.				
	1 represents a 16-bit integer value.				
	2 represents a 32-bit integer value.				
	<i>3</i> represents a 64-bit integer value.				
	4 represents a 16-bit float value.				
	5 represents a 32-bit float value.				
	6 represents a 64-bit float value.				
	Only valid with the <b>Kernel Execution</b>				
	Model.				
31	ContractionOff	Kernel			
	Indicates that floating-point-expressions				
	contraction is disallowed. Only valid with				
	the <b>Kernel</b> Execution Model.				
33	Initializer	Kernel			
	Indicates that this entry point is a module				
	initializer.	Missing before version 1.1.			
34	Finalizer	Kernel			
	Indicates that this entry point is a module				
	finalizer.	Missing before version			
		1.1.			
35	SubgroupSize	SubgroupDispatch	Literal N	lumber	
	Indicates that this entry point requires the	2 4.0 <b>g</b> - 0 2 <b>f</b> - 2 <b>f</b> - 2 1 1 1 1 1	Subgroup		
	specified Subgroup Size.	Missing before version	2		
		<b>1.1</b> .			
36	SubgroupsPerWorkgroup	SubgroupDispatch	Literal N		
	Indicates that this entry point requires the	<b>.</b>	Subgroup	ps Per Work	group
	specified number of Subgroups Per	Missing before version			
	Workgroup.	1.1.	1		
37	SubgroupsPerWorkgroupId	SubgroupDispatch	<id>&gt;</id>	D *** *	
	Indicates that this entry point requires the	<b>M</b> 1 .	Subgroup	ps Per Work	group
	specified number of Subgroups Per	Missing before version			
	Workgroup.	1.2.			
	Specified as an Id.				
38	LocalSizeId	Missing before version	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
	Indicates the work-group size in the $x$ , $y$ ,	1.2.	x size	y size	z size
	and z dimensions. Only valid with the				
	GLCompute or Kernel Execution Models.				
	Succified as Ida				
	Specified as Ids.				

	Execution Mode	<b>Enabling Capabilities</b>	Extra Operands
39	LocalSizeHintId	Kernel	< <i>id</i> >
	A hint to the compiler, which indicates the		Local Size Hint
	most likely to be used work-group size in	Missing before version	
	the $x$ , $y$ , and $z$ dimensions. Only valid with	1.2.	
	the Kernel Execution Model.		
	Specified as an Id.		
4446	PostDepthCoverage	SampleMaskPostDepthCo	verage
		Reserved.	
5027	StencilRefReplacingEXT	StencilExportEXT	
		Reserved.	

# 3.7 Storage Class

Class of storage for declared variables (does not include intermediate values). Used by:

- OpTypePointer
- OpTypeForwardPointer
- OpVariable
- OpGenericCastToPtrExplicit

	Storage Class	Enabling Capabilities	Enabled by Extension
0	UniformConstant		
	Shared externally, visible across all		
	functions in all invocations in all work		
	groups. Graphics uniform memory.		
	OpenCL constant memory. Variables		
	declared with this storage class are		
	read-only. They may have initializers, as		
	allowed by the client API.		
1	Input		
	Input from pipeline. Visible across all		
	functions in the current invocation.		
	Variables declared with this storage class		
	are read-only, and cannot have initializers.		
2	Uniform	Shader	
	Shared externally, visible across all		
	functions in all invocations in all work		
	groups. Graphics uniform blocks and		
	buffer blocks.	CI I	
3	Output	Shader	
	Output to pipeline. Visible across all		
4	functions in the current invocation.		
4	Workgroup		
	Shared across all invocations within a		
	work group. Visible across all functions.		
	The OpenGL "shared" storage qualifier.		
	OpenCL local memory.		

	Storage Class	Enabling Capabilities	Enabled by Extension
5	CrossWorkgroup Visible across all functions of all invocations of all work groups. OpenCL global memory.		
6	Private Visible to all functions in the current invocation. Regular global memory.	Shader	
7	Function Visible only within the declaring function of the current invocation. Regular function memory.		
8	Generic For generic pointers, which overload the Function, Workgroup, and CrossWorkgroup Storage Classes.	GenericPointer	
9	PushConstant For holding push-constant memory, visible across all functions in all invocations in all work groups. Intended to contain a small bank of values pushed from the API. Variables declared with this storage class are read-only, and cannot have initializers.	Shader	
10	AtomicCounter For holding atomic counters. Visible across all functions of the current invocation. Atomic counter-specific memory.	AtomicStorage	
11	Image For holding image memory.		
12	StorageBuffer	Shader Reserved.	SPV_KHR_storage_buffer_storage_class, SPV_KHR_variable_pointers

### 3.8 Dim

Dimensionality of an image. Used by OpTypeImage.

	Dim	Enabling Capabilities
0	1D	Sampled1D
1	2D	
2	3D	
3	Cube	Shader
4	Rect	SampledRect
5	Buffer	SampledBuffer
6	SubpassData	InputAttachment

## 3.9 Sampler Addressing Mode

Addressing mode for creating constant samplers. Used by OpConstantSampler.

	Sampler Addressing Mode	Enabling Capabilities
0	None	Kernel
	The image coordinates used to sample	
	elements of the image refer to a location	
	inside the image, otherwise the results are	
	undefined.	
1	ClampToEdge	Kernel
	Out-of-range image coordinates are clamped	
	to the extent.	
2	Clamp	Kernel
	Out-of-range image coordinates will return a	
	border color.	
3	Repeat	Kernel
	Out-of-range image coordinates are wrapped	
	to the valid range. Can only be used with	
	normalized coordinates.	
4	RepeatMirrored	Kernel
	Flip the image coordinate at every integer	
	junction. Can only be used with normalized	
	coordinates.	

# 3.10 Sampler Filter Mode

Filter mode for creating constant samplers. Used by OpConstantSampler.

Sampler Filter Mode		Enabling Capabilities
0	Nearest	Kernel
	Use filter nearest mode when performing a	
	read image operation.	
1	Linear	Kernel
	Use filter linear mode when performing a	
	read image operation.	

## 3.11 Image Format

Declarative image format. Used by OpTypeImage.

	Image Format	Enabling Capabilities
0	Unknown	
1	Rgba32f	Shader
2	Rgba16f	Shader
3	R32f	Shader
4	Rgba8	Shader
5	Rgba8Snorm	Shader
6	Rg32f	StorageImageExtendedFormats
7	Rg16f	StorageImageExtendedFormats
8	R11fG11fB10f	StorageImageExtendedFormats
9	R16f	StorageImageExtendedFormats
10	Rgba16	StorageImageExtendedFormats
11	Rgb10A2	StorageImageExtendedFormats
12	Rg16	StorageImageExtendedFormats
13	Rg8	StorageImageExtendedFormats

	Image Format	<b>Enabling Capabilities</b>
14	R16	StorageImageExtendedFormats
15	R8	StorageImageExtendedFormats
16	Rgba16Snorm	StorageImageExtendedFormats
17	Rg16Snorm	StorageImageExtendedFormats
18	Rg8Snorm	StorageImageExtendedFormats
19	R16Snorm	StorageImageExtendedFormats
20	R8Snorm	StorageImageExtendedFormats
21	Rgba32i	Shader
22	Rgba16i	Shader
23	Rgba8i	Shader
24	R32i	Shader
25	Rg32i	StorageImageExtendedFormats
26	Rg16i	StorageImageExtendedFormats
27	Rg8i	StorageImageExtendedFormats
28	R16i	StorageImageExtendedFormats
29	R8i	StorageImageExtendedFormats
30	Rgba32ui	Shader
31	Rgba16ui	Shader
32	Rgba8ui	Shader
33	R32ui	Shader
34	Rgb10a2ui	StorageImageExtendedFormats
35	Rg32ui	StorageImageExtendedFormats
36	Rg16ui	StorageImageExtendedFormats
37	Rg8ui	StorageImageExtendedFormats
38	R16ui	StorageImageExtendedFormats
39	R8ui	StorageImageExtendedFormats

# 3.12 Image Channel Order

Image channel order returned by OpImageQueryOrder.

	Image Channel Order	<b>Enabling Capabilities</b>
0	R	Kernel
1	A	Kernel
2	RG	Kernel
3	RA	Kernel
4	RGB	Kernel
5	RGBA	Kernel
6	BGRA	Kernel
7	ARGB	Kernel
8	Intensity	Kernel
9	Luminance	Kernel
10	Rx	Kernel
11	RGx	Kernel
12	RGBx	Kernel
13	Depth	Kernel
14	DepthStencil	Kernel
15	sRGB	Kernel
16	sRGBx	Kernel
17	sRGBA	Kernel
18	sBGRA	Kernel
19	ABGR	Kernel

#### 3.13 Image Channel Data Type

Image channel data type returned by OpImageQueryFormat.

	Image Channel Data Type	Enabling Capabilities
0	SnormInt8	Kernel
1	SnormInt16	Kernel
2	UnormInt8	Kernel
3	UnormInt16	Kernel
4	UnormShort565	Kernel
5	UnormShort555	Kernel
6	UnormInt101010	Kernel
7	SignedInt8	Kernel
8	SignedInt16	Kernel
9	SignedInt32	Kernel
10	UnsignedInt8	Kernel
11	UnsignedInt16	Kernel
12	UnsignedInt32	Kernel
13	HalfFloat	Kernel
14	Float	Kernel
15	UnormInt24	Kernel
16	UnormInt101010_2	Kernel

### 3.14 Image Operands

Additional operands to sampling, or getting texels from, an image. Bits that are set can indicate that another operand follows. If there are multiple following operands indicated, they are ordered: Those indicated by smaller-numbered bits appear first. At least one bit must be set (**None** is invalid).

This value is a literal mask; it can be formed by combining the bits from multiple rows in the table below.

#### Used by:

- OpImageSampleImplicitLod
- OpImageSampleExplicitLod
- OpImageSampleDrefImplicitLod
- OpImageSampleDrefExplicitLod
- OpImageSampleProjImplicitLod
- OpImageSampleProjExplicitLod
- OpImageSampleProjDrefImplicitLod
- OpImageSampleProjDrefExplicitLod
- OpImageFetch
- · OpImageGather
- OpImageDrefGather
- OpImageRead
- OpImageWrite
- OpImageSparseSampleImplicitLod
- OpImageSparseSampleExplicitLod
- OpImageSparseSampleDrefImplicitLod

- $\bullet \ OpImageSparseSampleDrefExplicitLod$
- OpImageSparseSampleProjImplicitLod
- OpImageSparseSampleProjExplicitLod
- $\bullet \ \ OpImageSparseSampleProjDrefImplicitLod$
- $\bullet \ OpImageSparseSampleProjDrefExplicitLod$
- OpImageSparseFetch
- OpImageSparseGather
- OpImageSparseDrefGather
- OpImageSparseRead

Image Operands		<b>Enabling Capabilities</b>
0x0	None	
Ox1	Bias A following operand is the bias added to the implicit level of detail. Only valid with implicit-lod instructions. It must be a floating-point type scalar. This can only be used with an OpTypeImage that has a Dim operand of 1D, 2D, 3D, or Cube, and the MS operand must be 0.	Shader
0x2	Lod A following operand is the explicit level-of-detail to use. Only valid with explicit-lod instructions. For sampling operations, it must be a floating-point type scalar. For fetch operations, it must be an integer type scalar. This can only be used with an OpTypeImage that has a Dim operand of 1D, 2D, 3D, or Cube, and the MS operand must be 0.	
0x4	Grad Two following operands are $dx$ followed by $dy$ . These are explicit derivatives in the $x$ and $y$ direction to use in computing level of detail. Each is a scalar or vector containing $(du/dx[, dv/dx] [, dw/dx])$ and $(du/dy[, dv/dy] [, dw/dy])$ . The number of components of each must equal the number of components in $Coordinate$ , minus the $array\ layer$ component, if present. Only valid with explicit-lod instructions. They must be a scalar or vector of floating-point type. This can only be used with an OpTypeImage that has an $MS$ operand of 0. It is invalid to set both the $Lod$ and $Grad$ bits.	

	Image Operands	<b>Enabling Capabilities</b>
0x8	ConstOffset	•
	A following operand is added to $(u, v, w)$	
	before texel lookup. It must be an $\langle id \rangle$ of	
	an integer-based constant instruction of	
	scalar or vector type. It is invalid for these	
	to be outside a target-dependent allowed	
	range. The number of components must	
	equal the number of components in	
	Coordinate, minus the array layer	
	component, if present. Not valid with the	
	Cube dimension.	
0x10	Offset	ImageGatherExtended
OATO	A following operand is added to $(u, v, w)$	Image Guther Extended
	before texel lookup. It must be a scalar or	
	vector of integer type. It is invalid for these	
	to be outside a target-dependent allowed	
	range. The number of components must	
	equal the number of components in	
	Coordinate, minus the array layer	
	component, if present. Not valid with the	
0.20	Cube dimension.	
0x20	ConstOffsets	
	A following operand is Offsets. Offsets	
	must be an $\langle id \rangle$ of a constant instruction	
	making an array of size four of vectors of	
	two integer components. Each gathered	
	texel is identified by adding one of these	
	array elements to the $(u, v)$ sampled	
	location. It is invalid for these to be outside	
	a target-dependent allowed range. Only	
	valid with OpImageGather or	
	OpImageDrefGather. Not valid with the	
	Cube dimension.	
0x40	Sample	
	A following operand is the sample number	
	of the sample to use. Only valid with	
	OpImageFetch, OpImageRead,	
	OpImageWrite, OpImageSparseFetch, and	
	OpImageSparseRead. It is invalid to have a	
	Sample operand if the underlying	
	OpTypeImage has MS of 0. It must be an	
	integer type scalar.	
0x80	MinLod	MinLod
	A following operand is the minimum	
	level-of-detail to use when accessing the	
	image. Only valid with Implicit	
	instructions and <b>Grad</b> instructions. It must	
	be a floating-point type scalar. This can	
	only be used with an OpTypeImage that	
	has a Dim operand of 1D, 2D, 3D, or	
	<b>Cube</b> , and the <i>MS</i> operand must be 0.	

### 3.15 FP Fast Math Mode

Enables fast math operations which are otherwise unsafe.

• Only valid on OpFAdd, OpFSub, OpFMul, OpFDiv, OpFRem, and OpFMod instructions.

This value is a literal mask; it can be formed by combining the bits from multiple rows in the table below.

FP Fast Math Mode		Enabling Capabilities
0x0	None	
0x1	NotNaN	Kernel
	Assume parameters and result are not	
	NaN.	
0x2	NotInf	Kernel
	Assume parameters and result are not +/-	
	Inf.	
0x4	NSZ	Kernel
	Treat the sign of a zero parameter or result	
	as insignificant.	
0x8	AllowRecip	Kernel
	Allow the usage of reciprocal rather than	
	perform a division.	
0x10	Fast	Kernel
	Allow algebraic transformations according	
	to real-number associative and distributive	
	algebra. This flag implies all the others.	

# 3.16 FP Rounding Mode

Associate a rounding mode to a floating-point conversion instruction.

FP Rounding Mode			
0	RTE		
	Round to nearest even.		
1	RTZ		
	Round towards zero.		
2	RTP		
	Round towards positive infinity.		
3	RTN		
	Round towards negative infinity.		

## 3.17 Linkage Type

Associate a linkage type to functions or global variables. See linkage.

Linkage Type		Enabling Capabilities	
0	Export	Linkage	
	Accessible by other modules as well.		
1	Import	Linkage	
	A declaration of a global variable or a		
	function that exists in another module.		

## 3.18 Access Qualifier

Defines the access permissions.

Used by OpTypeImage and OpTypePipe.

	Access Qualifier	<b>Enabling Capabilities</b>	
0	ReadOnly	Kernel	
	A read-only object.		
1	WriteOnly	Kernel	
	A write-only object.		
2	ReadWrite	Kernel	
	A readable and writable object.		

## 3.19 Function Parameter Attribute

Adds additional information to the return type and to each parameter of a function.

Function Parameter Attribute		<b>Enabling Capabilities</b>
0	Zext	Kernel
	Value should be zero extended if needed.	
1	Sext	Kernel
	Value should be sign extended if needed.	
2	ByVal	Kernel
	This indicates that the pointer parameter	
	should really be passed by value to the	
	function. Only valid for pointer parameters	
	(not for ret value).	
3	Sret	Kernel
	Indicates that the pointer parameter specifies	
	the address of a structure that is the return	
	value of the function in the source program.	
	Only applicable to the first parameter which	
4	must be a pointer parameters.	W1
4	NoAlias	Kernel
	Indicates that the memory pointed to by a pointer parameter is not accessed via pointer	
	values which are not derived from this	
	pointer parameter. Only valid for pointer	
	parameters. Not valid on return values.	
5	NoCapture	Kernel
3	The callee does not make a copy of the	Kerner
	pointer parameter into a location that is	
	accessible after returning from the callee.	
	Only valid for pointer parameters. Not valid	
	on return values.	
6	NoWrite	Kernel
	Can only read the memory pointed to by a	
	pointer parameter. Only valid for pointer	
	parameters. Not valid on return values.	
7	NoReadWrite	Kernel
	Cannot dereference the memory pointed to	
	by a pointer parameter. Only valid for pointer	
	parameters. Not valid on return values.	

# 3.20 Decoration

Used by OpDecorate, OpMemberDecorate, and OpDecorateId.

	Decoration	<b>Enabling Capabilities</b>	Extra Operands
0	RelaxedPrecision	Shader	
	Allow reduced precision operations. To be used		
	as described in Relaxed Precision.		
1	SpecId	Shader, Kernel	Literal Number
	Apply to a scalar specialization constant. Forms		Specialization
	the API linkage for setting a specialized value.		Constant ID
	See specialization.		
2	Block	Shader	
	Apply to a structure type to establish it is a		
	non-SSBO-like shader-interface block.		
3	BufferBlock	Shader	
	Apply to a structure type to establish it is an		
	SSBO-like shader-interface block.		
4	RowMajor	Matrix	
	Applies only to a member of a structure type.		
	Only valid on a matrix or array whose most basic		
	element is a matrix. Indicates that components		
	within a row are contiguous in memory.		
5	ColMajor	Matrix	
-	Applies only to a member of a structure type.		
	Only valid on a matrix or array whose most basic		
	element is a matrix. Indicates that components		
	within a column are contiguous in memory.		
6	ArrayStride	Shader	Literal Number
Ü	Apply to an array type to specify the stride, in	Silator	Array Stride
	bytes, of the array's elements. Can also apply to a		Tirrety Street
	pointer type to an array element, to specify the		
	stride of the array that the element resides in.		
	Must not be applied to any other type.		
7	MatrixStride	Matrix	Literal Number
•	Applies only to a member of a structure type.	17147111	Matrix Stride
	Only valid on a matrix or array whose most basic		Treat is Struce
	element is a matrix. Specifies the stride of rows in		
	a <b>RowMajor</b> -decorated matrix, or columns in a		
	ColMajor-decorated matrix.		
8	GLSLShared	Shader	
J	Apply to a structure type to get GLSL shared	Siiduci	
	memory layout.		
9	GLSLPacked	Shader	
	Apply to a structure type to get GLSL packed	Shauei	
	memory layout.		
10	CPacked	Kernel	
10	Apply to a structure type, to marks it as "packed",	Kernei	
	indicating that the alignment of the structure is		
	one and that there is no padding between structure		
11	members.  BuiltIn		DuiltIn
11			BuiltIn
	Apply to an object or a member of a structure		
	type. Indicates which built-in variable the entity		
	represents. See BuiltIn for more information.		

	Decoration	<b>Enabling Capabilities</b>	Extra Operands
13	NoPerspective	Shader	
	Apply to an object or a member of a structure		
	type. Indicates that linear, non-perspective		
	correct, interpolation must be used. Only valid for		
	the Input and Output Storage Classes.		
14	Flat	Shader	
	Apply to an object or a member of a structure		
	type. Indicates no interpolation will be done. The		
	non-interpolated value will come from a vertex,		
	as described in the API specification. Only valid		
	for the <b>Input</b> and <b>Output</b> Storage Classes.		
15	Patch	Tessellation	
	Apply to an object or a member of a structure	100001111111111111111111111111111111111	
	type. Indicates a tessellation patch. Only valid for		
	the <b>Input</b> and <b>Output</b> Storage Classes. Invalid to		
	use on objects or types referenced by		
	non-tessellation Execution Models.		
16	Centroid	Shader	
10	Apply to an object or a member of a structure	Shadei	
	type. When used with multi-sampling		
	rasterization, allows a single interpolation		
	location for an entire pixel. The interpolation		
	location must lie in both the pixel and in the		
	primitive being rasterized. Only valid for the		
	1 -		
17	Input and Output Storage Classes.  Sample	SampleRateShading	
17	Apply to an object or a member of a structure	SampleKateShaunig	
	type. When used with multi-sampling		
	rasterization, requires per-sample interpolation.		
	The interpolation locations must be the locations		
	of the samples lying in both the pixel and in the		
	primitive being rasterized. Only valid for the		
	Input and Output Storage Classes.		
18	Invariant	Shader	
10	Apply to a variable, to indicate expressions	Silauci	
	computing its value be done invariant with respect		
	to other modules computing the same		
	expressions.		
19	Restrict		
19	Apply to a variable, to indicate the compiler may		
	compile as if there is no aliasing. See the Aliasing		
	section for more detail.		
20	Aliased		
20			
	Apply to a variable, to indicate the compiler is to		
	generate accesses to the variable that work		
	correctly in the presence of aliasing. See the		
	Aliasing section for more detail.		

	Decoration	<b>Enabling Capabilities</b>	Extra Operands
21	Volatile		_
	Apply to an object, pointer type, or a member of a		
	structure type. Can only be used for objects		
	declared as storage images (see OpTypeImage) or		
	in the Uniform Storage Class with the		
	<b>BufferBlock</b> Decoration. This indicates the		
	memory holding the variable is volatile memory.		
	Accesses to volatile memory cannot be		
	eliminated, duplicated, or combined with other		
	accesses. The variable cannot be in the <b>Function</b>		
	Storage Class.		
22	Constant	Kernel	
	Indicates that a global variable is constant and		
	will <b>never</b> be modified. Only allowed on global		
	variables.		
23	Coherent		
	Apply to an object, pointer type, or a member of a		
	structure type. Can only be used for objects		
	declared as storage images (see OpTypeImage) or		
	in the Uniform Storage Class with the		
	<b>BufferBlock</b> Decoration. This indicates the		
	memory backing the object is coherent.		
24	NonWritable		
	Apply to an object or a member of a structure		
	type. Can only be used for objects declared as		
	storage images (see OpTypeImage) or in the		
	Uniform Storage Class with the BufferBlock		
	Decoration. This indicates the memory holding		
	the variable is not writable, and that this module		
25	does not write to it.  NonReadable		
25			
	Apply to an object or a member of a structure		
	type. Can only be used for objects declared as		
	storage images (see OpTypeImage) or in the		
	Uniform Storage Class with the BufferBlock Decoration. This indicates the memory holding		
	,		
	the variable is not readable, and that this module does not read from it.		
26	Uniform	Shader	
20	Apply to an object or a member of a structure	Silauci	
	type. Asserts that the value backing the decorated		
	<id> is dynamically uniform, hence the</id>		
	consumer is allowed to assume this is the case.		
28	SaturatedConversion	Kernel	
	Indicates that a conversion to an integer type		
	which is outside the representable range of <i>Result</i>		
	Type will be clamped to the nearest representable		
	value of <i>Result Type</i> . <i>NaN</i> will be converted to 0.		
	This decoration can only be applied to conversion		
	instructions to integer types, not including the		
	OpSatConvertUToS and OpSatConvertSToU		
	instructions.		

	Decoration	<b>Enabling Capabilities</b>	Extra Operands
29	Stream	GeometryStreams	Literal Number
	Apply to an object or a member of a structure		Stream Number
	type. Indicates the stream number to put an		
	output on. Only valid for the <b>Output Storage</b>		
	Class and the <b>Geometry</b> Execution Model.		
30	Location	Shader	Literal Number
	Apply to a variable or a structure-type member.		Location
	Forms the main linkage for Storage Class Input		
	and <b>Output</b> variables:		
	- between the API and vertex-stage inputs,		
	- between consecutive programmable stages, or		
	- between fragment-stage outputs and the API.		
	Also can tag variables or structure-type members		
	in the UniformConstant Storage Class for		
	linkage with the API.		
	Only valid for the <b>Input</b> , <b>Output</b> , and		
	UniformConstant Storage Classes.		
31	Component	Shader	Literal Number
	Apply to an object or a member of a structure	2	Component
	type. Indicates which component within a		,
	Location will be taken by the decorated entity.		
	Only valid for the <b>Input</b> and <b>Output</b> Storage		
	Classes.		
32	Index	Shader	Literal Number
32	Apply to a variable to identify a blend equation	Shadei	Index
	input index, used as described in the API		mucx
	specification. Only valid for the <b>Output</b> Storage		
	Class and the <b>Fragment</b> Execution Model.		
33	Binding	Shader	Literal Number
	Apply to a variable. Part of the main linkage	Shauci	Binding Point
	between the API and SPIR-V modules for		Binaing Toini
	memory buffers, images, etc. See the API		
	specification for more information.		
34	DescriptorSet	Shader	Literal Number
34	Apply to a variable. Part of the main linkage	Shauei	Descriptor Set
	between the API and SPIR-V modules for		Descripior Sei
	memory buffers, images, etc. See the API		
35	specification for more information.  Offset	Shader	Literal Number
33		Snader	
	Apply to a structure-type member. This gives the		Byte Offset
	byte offset of the member relative to the		
	beginning of the structure. Can be used, for		
	example, by both uniform and		
	transform-feedback buffers. It must not cause any		
	overlap of the structure's members, or overflow of		
26	a transform-feedback buffer's <b>XfbStride</b> .	Tuon st D. 11 1	T Stand NT1
36	XfbBuffer	TransformFeedback	Literal Number
	Apply to an object or a member of a structure		XFB Buffer Number
	type. Indicates which transform-feedback buffer		
	an output is written to. Only valid for the <b>Output</b>		
	Storage Classes of vertex processing Execution		
	Models.		

	Decoration	<b>Enabling Capabilities</b>	Extra Ope	rands
37	XfbStride	TransformFeedback	Literal Nur	
	Apply to anything <b>XfbBuffer</b> is applied to.		XFB Stride	
	Specifies the stride, in bytes, of			
	transform-feedback buffer vertices. If the			
	transform-feedback buffer is capturing any			
	double-precision components, the stride must be a			
38	multiple of 8, otherwise it must be a multiple of 4.  FuncParamAttr	V a see al	E	
38		Kernel	Function P	arameter
	Indicates a function return value or parameter attribute.		Attribute Function P	
	attribute.		Attribute	arameier
39	FPRoundingMode		FP Roundin	ng Mode
39	Indicates a floating-point rounding mode.		Floating-Po	_
	indicates a noating-point rounding mode.		Rounding I	
40	FPFastMathMode	Kernel	FP Fast Ma	
40	Indicates a floating-point fast math flag.	IXETHE	Fast-Math	
41	LinkageAttributes	Linkage	Literal	Linkage
71	Associate linkage attributes to values. Only valid	Linnage	String	Type
	on OpFunction or global (module scope)		Name	Linkage
	OpVariable. See linkage.		1 tunic	Туре
42	NoContraction	Shader		Турс
	Apply to an arithmetic instruction to indicate the	5		
	operation cannot be combined with another			
	instruction to form a single operation. For			
	example, if applied to an OpFMul, that multiply			
	can't be combined with an addition to yield a			
	fused multiply-add operation. Furthermore, such			
	operations are not allowed to reassociate; e.g.,			
	add(a + add(b+c)) cannot be transformed to			
	add(add(a+b)+c).			
43	InputAttachmentIndex	InputAttachment	Literal Nur	nber
	Apply to a variable to provide an input-target		Attachment	t Index
	index (as described in the API specification).			
	Only valid in the <b>Fragment</b> Execution Model and			
	for variables of type OpTypeImage with a Dim			
	operand of SubpassData.			
44	Alignment	Kernel	Literal Nur	nber
	Apply to a pointer. This declares a known		Alignment	
	minimum alignment the pointer has.			
45	MaxByteOffset	Addresses	Literal Nur	
	Apply to a pointer. This declares a known	Nr. 1 C	Max Byte C	)]fset
	maximum byte offset this pointer will be	Missing before version 1.1.		
	incremented by from the point of the decoration.			
	This is a guaranteed upper bound when applied to			
16	OpFunctionParameter.	Vousi	\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	
46	AlignmentId	Kernel	<id></id>	
	Apply to a pointer. This declares a known	Missing before version 12	Alignment	
	minimum alignment the pointer has.	Missing before version 1.2.		
	Specified as an Id.			
	specifica as all ia.			

	Decoration	Enabling Capabilities	Extra Operands
47	MaxByteOffsetId	Addresses	<id></id>
	Apply to a pointer. This declares a known		Max Byte Offset
	maximum byte offset this pointer will be	Missing before version 1.2.	
	incremented by from the point of the decoration.		
	This is a guaranteed upper bound when applied to		
	OpFunctionParameter.		
	Specified as an Id.		
4999	ExplicitInterpAMD	Reserved.	
5248	OverrideCoverageNV	SampleMaskOverrideCovera	geNV
		Reserved.	
5250	PassthroughNV	GeometryShaderPassthrough	NV
		Reserved.	
5252	ViewportRelativeNV	ShaderViewportMaskNV	
		Reserved.	
5256	SecondaryViewportRelativeNV	ShaderStereoViewNV	Literal Number
			Offset
		Reserved.	

### 3.21 BuiltIn

Used when **Decoration** is **BuiltIn**. Apply to either

- the result <id> of the variable declaration of the built-in variable, or
- a structure-type member, if the built-in is a member of a structure.

As stated per entry below, these have additional semantics and constraints described by the client API.

	BuiltIn	<b>Enabling Capabilities</b>
0	Position	Shader
	Output vertex position from a vertex	
	processing Execution Model. See Vulkan or	
	OpenGL API specifications for more detail.	
1	PointSize	Shader
	Output point size from a vertex processing	
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
3	ClipDistance	ClipDistance
	Array of clip distances. See Vulkan or	
	OpenGL API specifications for more detail.	
4	CullDistance	CullDistance
	Array of clip distances. See Vulkan or	
	OpenGL API specifications for more detail.	
5	VertexId	Shader
	Input vertex ID to a Vertex Execution	
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	

	BuiltIn	<b>Enabling Capabilities</b>
6	InstanceId	Shader
	Input instance ID to a <b>Vertex</b> Execution	
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	
7	PrimitiveId	Geometry, Tessellation
	Primitive ID in a <b>Geometry Execution</b>	•
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	
8	InvocationId	Geometry, Tessellation
	Invocation ID, input to <b>Geometry</b> and	
	TessellationControl Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
9	Layer	Geometry
	Layer output by a <b>Geometry</b> Execution	
	Model, input to a <b>Fragment</b> Execution	
	Model, for multi-layer framebuffer. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
10	ViewportIndex	MultiViewport
	Viewport Index output by a <b>Geometry</b> stage,	
	input to a <b>Fragment</b> Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
11	TessLevelOuter	Tessellation
	Output patch outer levels in a	
	TessellationControl Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
12	TessLevelInner	Tessellation
	Output patch inner levels in a	
	TessellationControl Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
13	TessCoord	Tessellation
	Input vertex position in	
	TessellationEvaluation Execution Model.	
	See Vulkan or OpenGL API specifications	
4 .	for more detail.	<b>T</b>
14	PatchVertices	Tessellation
	Input patch vertex count in a tessellation	
	Execution Model. See Vulkan or OpenGL	
1.7	API specifications for more detail.	GL 3
15	FragCoord	Shader
	Coordinates $(x, y, z, 1/w)$ of the current	
	fragment, input to the <b>Fragment</b> Execution	
	Model. See Vulkan or OpenGL API	
1.6	specifications for more detail.	CI. I
16	PointCoord	Shader
	Coordinates within a <i>point</i> , input to the	
	Fragment Execution Model. See Vulkan or	
	OpenGL API specifications for more detail.	

BuiltIn	Enabling Capabilities
ntFacing	Shader
-	
=	SampleRateShading
_	•
eution Model. See Vulkan or OpenGL	
-	
plePosition	SampleRateShading
t sample position to the <b>Fragment</b>	_
eution Model. See Vulkan or OpenGL	
-	
	Shader
gment Execution Model. See Vulkan or	
-	Shader
=	
*	
-	Shader
ber of workgroups in <b>GLCompute</b> or	
nel Execution Models. See OpenCL,	
an, or OpenGL API specifications for	
e detail.	
kgroupSize	
k-group size in <b>GLCompute</b> or <b>Kernel</b>	
eution Models. See OpenCL, Vulkan, or	
nGL API specifications for more detail.	
kgroupId	
k-group ID in GLCompute or Kernel	
eution Models. See OpenCL, Vulkan, or	
-	
alInvocationId	
l invocation ID in GLCompute or	
nel Execution Models. See OpenCL,	
an, or OpenGL API specifications for	
palInvocationId	
al invocation ID in GLCompute or	
e detail.	
	direction, input to the Fragment cution Model. See Vulkan or OpenGL specifications for more detail.  spleId t sample number to the Fragment cution Model. See Vulkan or OpenGL specifications for more detail.  splePosition t sample position to the Fragment cution Model. See Vulkan or OpenGL specifications for more detail.  spleMask t or output sample mask to the gment Execution Model. See Vulkan or nnGL API specifications for more detail.  specification Models. See OpenCL, vulkan, or nnGL API specifications for more detail.  specification Models. See OpenCL, vulkan, or nnGL API specifications for more detail.  specification Models. See OpenCL, vulkan, or nnGL API specifications for more detail.  specification Models. See OpenCL, vulkan, or nnGL API specifications for more detail.  specification Models. See OpenCL, vulkan, or nnGL API specifications for more detail.  specification Models. See OpenCL, vulkan, or nnel Execution Models. See OpenCL, vulkan, or nnel Execution Models. See OpenCL, vulkan, or onel Execution Models. See OpenCL, vu

	BuiltIn	<b>Enabling Capabilities</b>
29	LocalInvocationIndex	
	Local invocation index in GLCompute	
	Execution Models. See Vulkan or OpenGL	
	API specifications for more detail.	
	Work-group Linear ID in Kernel Execution	
	Models. See OpenCL API specification for	
	more detail.	
30	WorkDim	Kernel
	Work dimensions in <b>Kernel Execution</b>	
	Models. See OpenCL API specification for	
	more detail.	
31	GlobalSize	Kernel
	Global size in <b>Kernel</b> Execution Models. See	
	OpenCL API specification for more detail.	
32	EnqueuedWorkgroupSize	Kernel
	Enqueued work-group size in <b>Kernel</b>	
	Execution Models. See OpenCL API	
	specification for more detail.	
33	GlobalOffset	Kernel
	Global offset in <b>Kernel</b> Execution Models.	
	See OpenCL API specification for more	
	detail.	
34	GlobalLinearId	Kernel
	Global linear ID in Kernel Execution	
	Models. See OpenCL API specification for	
	more detail.	
36	SubgroupSize	Kernel
	Subgroup size. See OpenCL or OpenGL API	
	specifications for more detail.	
37	SubgroupMaxSize	Kernel
	Subgroup maximum size in <b>Kernel</b>	
	Execution Models. See OpenCL API	
	specification for more detail.	
38	NumSubgroups	Kernel
	Number of subgroups in <b>GLCompute</b> or	
	Kernel Execution Models. See OpenCL or	
20	OpenGL API specifications for more detail.	TZ 1
39	NumEnqueuedSubgroups	Kernel
	Number of enqueued subgroups in <b>Kernel</b>	
	Execution Models. See OpenCL API	
40	specification for more detail.	V1
40	Subgroup ID in CL Compute or Kornel	Kernel
	Subgroup ID in GLCompute or Kernel	
	Execution Models. See OpenCL or OpenGL API specifications for more detail.	
41	SubgroupLocalInvocationId	Kernel
41	Subgroup local invocation ID. See OpenCL	Kerilei
	or OpenGL API specifications for more detail.	
42	VertexIndex	Shader
7-2	Vertex index. See Vulkan or OpenGL API	Silauci
	specifications for more detail.	
	specifications for more detail.	

	BuiltIn	<b>Enabling Capabilities</b>
43	InstanceIndex	Shader
	Instance index. See Vulkan or OpenGL API	
	specifications for more detail.	
4416	SubgroupEqMaskKHR	SubgroupBallotKHR
		Reserved.
4417	SubgroupGeMaskKHR	SubgroupBallotKHR
1117	Subgroup Gentuskiiiik	Suogroupbunotitiit
		Reserved.
4418	SubgroupGtMaskKHR	SubgroupBallotKHR
		Reserved.
4419	SubgroupLeMaskKHR	SubgroupBallotKHR
		Reserved.
4420	SubgroupLtMaskKHR	SubgroupBallotKHR
		Reserved.
4424	BaseVertex	DrawParameters
		Reserved.
4425	BaseInstance	DrawParameters
4423	Dasemstance	Drawrarameters
		Reserved.
4426	DrawIndex	DrawParameters
		Reserved.
4438	DeviceIndex	DeviceGroup
1150	Devicemaca	BeviceGroup
		Reserved.
4440	ViewIndex	MultiView
		D
4992	Dawe Carally Dawe AMD	Reserved.
	BaryCoordNoPerspAMD  ParyCoordNoPerspControidAMD	Reserved.
4993 4994	BaryCoordNoPerspCentroidAMD BaryCoordNoPerspSampleAMD	Reserved.
4994	BaryCoordSmoothAMD	Reserved.
4996	BaryCoordSmoothCentroidAMD	Reserved.
4997	BaryCoordSmoothSampleAMD	Reserved.
4998	BaryCoordPullModelAMD	Reserved.
5014	FragStencilRefEXT	StencilExportEXT
2011	Tragotenentenzar	StellellExportExx
5050		Reserved.
5253	ViewportMaskNV	<b>ShaderViewportMaskNV</b>
		Reserved.
5257	SecondaryPositionNV	ShaderStereoViewNV
		D J
E050	Coordom Viores - ANA - LNIVI	Reserved.
5258	SecondaryViewportMaskNV	ShaderStereoViewNV
		Reserved.

	BuiltIn	Enabling Capabilities
5261	PositionPerViewNV	PerViewAttributesNV
		Reserved.
5262	ViewportMaskPerViewNV	PerViewAttributesNV
		Reserved.
5264	FullyCoveredEXT	FragmentFullyCoveredEXT
		Reserved.

#### 3.22 Selection Control

This value is a literal mask; it can be formed by combining the bits from multiple rows in the table below. Used by OpSelectionMerge.

	<b>Selection Control</b>		
0x0	None		
0x1	0x1 Flatten		
	Strong request, to the extent possible, to		
	remove the control flow for this selection.		
0x2	DontFlatten		
	Strong request, to the extent possible, to		
	keep this selection as control flow.		

## 3.23 Loop Control

This value is a literal mask; it can be formed by combining the bits from multiple rows in the table below. Used by OpLoopMerge.

	Loop Control	<b>Enabling Capabilities</b>
0x0	None	
0x1	Unroll	
	Strong request, to the extent possible, to	
	unroll or unwind this loop.	
0x2	DontUnroll	
	Strong request, to the extent possible, to	
	keep this loop as a loop, without unrolling.	
0x4	DependencyInfinite	Missing before version 1.1.
	Guarantees that there are no dependencies	
	between loop iterations.	
0x8	DependencyLength	Missing before version 1.1.
	Guarantees that there are no dependencies	
	between a number of loop	
	iterations, specified as a subsequent	
	literal-number operand to the instruction.	

#### 3.24 Function Control

This value is a literal mask; it can be formed by combining the bits from multiple rows in the table below.

#### Used by OpFunction.

Function Control		
0x0	None	
0x1	Inline	
	Strong request, to the extent possible, to	
	inline the function.	
0x2	DontInline	
	Strong request, to the extent possible, to not	
	inline the function.	
0x4	Pure	
	Compiler can assume this function has no	
	side effect, but might read global memory	
	or read through dereferenced function	
	parameters. Always computes the same	
	result for the same argument values.	
0x8	Const	
	Compiler can assume this function has no	
	side effects, and will not access global	
	memory or dereference function	
	parameters. Always computes the same	
	result for the same argument values.	

### 3.25 Memory Semantics <id>

Must be an <id> of a 32-bit integer scalar.

Memory semantics define memory-order constraints, and on what storage classes those constraints apply to. The memory order constrains the allowed orders in which memory operations in this invocation can made visible to another invocation. The storage classes specify to which subsets of memory these constraints are to be applied. Storage classes not selected are not being constrained.

Despite being a mask and allowing multiple bits to be combined, at most one of the first four (low-order) bits can be set. Requesting both **Acquire** and **Release** semantics is done by setting the **AcquireRelease** bit, not by setting two bits.

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

#### Used by:

- OpControlBarrier
- OpMemoryBarrier
- · OpAtomicLoad
- OpAtomicStore
- OpAtomicExchange
- OpAtomicCompareExchange
- OpAtomicCompareExchangeWeak
- OpAtomicIIncrement
- OpAtomicIDecrement
- OpAtomicIAdd
- OpAtomicISub
- OpAtomicSMin

- OpAtomicUMin
- OpAtomicSMax
- OpAtomicUMax
- OpAtomicAnd
- OpAtomicOr
- OpAtomicXor
- OpAtomicFlagTestAndSet
- OpAtomicFlagClear
- OpMemoryNamedBarrier

	<b>Memory Semantics</b>	<b>Enabling Capabilities</b>
0x0	None (Relaxed)	
0x2	Acquire	
	All memory operations provided in	
	program order after this memory operation	
	will execute after this memory operation.	
0x4	Release	
	All memory operations provided in	
	program order before this memory	
	operation will execute before this memory	
	operation.	
0x8	AcquireRelease	
	Has the properties of both Acquire and	
	Release semantics. It is used for	
	read-modify-write operations.	
0x10	SequentiallyConsistent	
	All observers will see this memory access	
	in the same order with respect to other	
	sequentially-consistent memory accesses	
	from this invocation.	
0x40	UniformMemory	Shader
	Apply the memory-ordering constraints to	
	Uniform Storage Class memory.	
0x80	SubgroupMemory	
	Apply the memory-ordering constraints to	
	subgroup memory.	
0x100	WorkgroupMemory	
	Apply the memory-ordering constraints to	
	Workgroup Storage Class memory.	
0x200	CrossWorkgroupMemory	
	Apply the memory-ordering constraints to	
	CrossWorkgroup Storage Class memory.	
0x400	AtomicCounterMemory	AtomicStorage
	Apply the memory-ordering constraints to	
	AtomicCounter Storage Class memory.	
0x800	ImageMemory	
	Apply the memory-ordering constraints to	
	image contents (types declared by	
	OpTypeImage), or to accesses done	
	through pointers to the <b>Image Storage</b>	
	Class.	

#### 3.26 Memory Access

Memory access semantics.

This value is a literal mask; it can be formed by combining the bits from multiple rows in the table below.

Used by:

- OpLoad
- OpStore
- OpCopyMemory
- OpCopyMemorySized

Memory Access			
0x0	None		
0x1	Volatile		
	This access cannot be eliminated,		
	duplicated, or combined with other		
	accesses.		
0x2	Aligned		
	This access has a known alignment,		
	provided as a literal in the next operand.		
0x4	Nontemporal		
	Hints that the accessed address is not likely		
	to be accessed again in the near future.		

### 3.27 Scope <id>

Must be an <id> of a 32-bit integer scalar.

The execution scope or memory scope of an operation. When used as a memory scope, it specifies the distance of synchronization from the current invocation. When used as an execution scope, it specifies the set of executing invocations taking part in the operation. Used by:

- OpControlBarrier
- OpMemoryBarrier
- OpAtomicLoad
- OpAtomicStore
- OpAtomicExchange
- OpAtomicCompareExchange
- OpAtomicCompareExchangeWeak
- OpAtomicIIncrement
- OpAtomicIDecrement
- OpAtomicIAdd
- OpAtomicISub
- OpAtomicSMin
- OpAtomicUMin
- OpAtomicSMax

- OpAtomicUMax
- OpAtomicAnd
- OpAtomicOr
- OpAtomicXor
- OpGroupAsyncCopy
- OpGroupWaitEvents
- OpGroupAll
- OpGroupAny
- OpGroupBroadcast
- OpGroupIAdd
- OpGroupFAdd
- OpGroupFMin
- OpGroupUMin
- OpGroupSMin
- OpGroupFMax
- OpGroupUMax
- OpGroupSMax
- OpGroupReserveReadPipePackets
- OpGroupReserveWritePipePackets
- OpGroupCommitReadPipe
- OpGroupCommitWritePipe
- OpAtomicFlagTestAndSet
- OpAtomicFlagClear
- OpMemoryNamedBarrier
- OpGroupIAddNonUniformAMD
- OpGroupFAddNonUniformAMD
- OpGroupFMinNonUniformAMD
- $\bullet \ \ OpGroupUMinNonUniformAMD$
- OpGroupSMinNonUniformAMD
- OpGroupFMaxNonUniformAMD
- OpGroupUMaxNonUniformAMD
- OpGroupSMaxNonUniformAMD

Scope			
0	CrossDevice		
	Scope crosses multiple devices.		
1	Device		
	Scope is the current device.		
2	Workgroup		
	Scope is the current workgroup.		
3	Subgroup		
	Scope is the current subgroup.		
4	Invocation		
	Scope is the current Invocation.		

#### 3.28 Group Operation

Defines the class of workgroup or subgroup operation. Used by:

- OpGroupIAdd
- OpGroupFAdd
- OpGroupFMin
- OpGroupUMin
- OpGroupSMin
- OpGroupFMax
- OpGroupUMax
- OpGroupSMax
- OpGroupIAddNonUniformAMD
- OpGroupFAddNonUniformAMD
- OpGroupFMinNonUniformAMD
- OpGroupUMinNonUniformAMD
- OpGroupSMinNonUniformAMD
- OpGroupFMaxNonUniformAMD
- $\bullet \ \ OpGroupUMaxNonUniformAMD$
- OpGroupSMaxNonUniformAMD

	Group Operation	<b>Enabling Capabilities</b>
0	Reduce	Kernel
	A reduction operation for all values of a	
	specific value X specified by invocations	
	within a workgroup.	
1	InclusiveScan	Kernel
	A binary operation with an identity $I$ and $n$	
	(where $n$ is the size of the workgroup)	
	elements[ $a_0, a_1, \ldots a_{n-1}$ ] resulting in [ $a_0, (a_0)$	
	op $a_1$ ), $(a_0$ op $a_1$ op op $a_{n-1}$ )]	
2	ExclusiveScan	Kernel
	A binary operation with an identity $I$ and $n$	
	(where $n$ is the size of the workgroup)	
	elements[ $a_0, a_1, \ldots a_{n-1}$ ] resulting in [ $I, a_0,$	
	$(a_0 \text{ op } a_1), \ldots (a_0 \text{ op } a_1 \text{ op } \ldots \text{ op } a_{n-2})].$	

### 3.29 Kernel Enqueue Flags

Specify when the child kernel begins execution.

**Note:** Implementations are not required to honor this flag. Implementations may not schedule kernel launch earlier than the point specified by this flag, however. Used by OpEnqueueKernel.

	Kernel Enqueue Flags	Enabling Capabilities
0	NoWait	Kernel
	Indicates that the enqueued kernels do not	
	need to wait for the parent kernel to finish	
	execution before they begin execution.	
1	WaitKernel	Kernel
	Indicates that all work-items of the parent	
	kernel must finish executing and all	
	immediate side effects committed before the	
	enqueued child kernel may begin execution.	
	<b>Note:</b> Immediate meaning not side effects	
	resulting from child kernels. The side effects	
	would include stores to global memory and	
	pipe reads and writes.	
2	WaitWorkGroup	Kernel
	Indicates that the enqueued kernels wait only	
	for the workgroup that enqueued the kernels	
	to finish before they begin execution.	
	<b>Note:</b> This acts as a memory synchronization	
	point between work-items in a work-group	
	and child kernels enqueued by work-items in	
	the work-group.	

### 3.30 Kernel Profiling Info

Specify the profiling information to be queried. Used by OpCaptureEventProfilingInfo.

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

	Kernel Profiling Info	Enabling Capabilities
0x0	None	
0x1	CmdExecTime	Kernel
	Indicates that the profiling info queried is	
	the execution time.	

### 3.31 Capability

Capabilities a module can declare it uses. All used capabilities must be declared, either directly or through a dependency: all capabilities that a declared capability depends on are automatically implied.

The **Depends On** column lists the dependencies for each capability. These are the ones implicitly declared. It is not necessary (but allowed) to declare a dependency for a declared capability.

See the capabilities section for more detail. Used by OpCapability.

	Capability	Depends On	Enabled by Extension
0	Matrix		
	Uses OpTypeMatrix.		
1	Shader	Matrix	
	Uses Vertex, Fragment, or GLCompute		
	Execution Models.		

	Capability	Depends On	Enabled by Extension
2	Geometry	Shader	
	Uses the <b>Geometry</b> Execution Model.		
3	Tessellation	Shader	
	Uses the <b>TessellationControl</b> or		
	TessellationEvaluation Execution		
	Models.		
4	Addresses		
	Uses physical addressing, non-logical		
	addressing modes.		
5	Linkage		
	Uses partially linked modules and		
	libraries.		
6	Kernel		
	Uses the <b>Kernel</b> Execution Model.		
7	Vector16	Kernel	
'	Uses OpTypeVector to declare 8	11011101	
	component or 16 component vectors.		
8	Float16Buffer	Kernel	
	Allows a 16-bit OpTypeFloat instruction	11011101	
	for the sole purpose of creating an		
	OpTypePointer to a 16-bit float. Pointers		
	to a 16-bit float cannot be dereferenced		
	directly, they must only be dereferenced		
	via an extended instruction. All other uses		
	of 16-bit <b>OpTypeFloat</b> are disallowed.		
9	Float16		
_	Uses OpTypeFloat to declare the 16-bit		
	floating-point type.		
10	Float64		
10	Uses OpTypeFloat to declare the 64-bit		
	floating-point type.		
11	Int64		
11	Uses OpTypeInt to declare 64-bit integer		
	types.		
12	Int64Atomics	Int64	
12	Uses atomic instructions on 64-bit integer	Into-	
	types.		
13	ImageBasic	Kernel	
	Uses OpTypeImage or OpTypeSampler in	11011101	
	a <b>Kernel</b> .		
14	ImageReadWrite	ImageBasic	
17	Uses OpTypeImage with the ReadWrite	imageDasic	
	access qualifier.		
15	ImageMipmap	ImageBasic	
13	Uses non-zero Lod Image Operands.	imageDasiC	
17	Pipes	Kernel	
1,	Uses OpTypePipe, OpTypeReserveId or	IXCI IICI	
	pipe instructions.		
18	Groups		
10	Uses group instructions.		
19	DeviceEnqueue	Kernel	
17	Uses OpTypeQueue, OpTypeDeviceEvent,	IXCI IICI	
	and device side enqueue instructions.		
	and device side enqueue instructions.		

	Capability	Depends On	Enabled by Extension
20	LiteralSampler	Kernel	·
	Samplers are made from literals within the		
	module. See OpConstantSampler.		
21	AtomicStorage	Shader	
	Uses the <b>AtomicCounter Storage Class</b> ,		
	allowing use of only the OpAtomicLoad,		
	OpAtomicIIncrement, and		
	OpAtomicIDecrement instructions.		
22	Int16		
	Uses OpTypeInt to declare 16-bit integer		
	types.		
23	TessellationPointSize	Tessellation	
	Tessellation stage exports point size.		
24	GeometryPointSize	Geometry	
	Geometry stage exports point size		
25	ImageGatherExtended	Shader	
	Uses texture gather with non-constant or		
	independent offsets		
27	StorageImageMultisample	Shader	
	Uses multi-sample images for		
	non-sampled images.		
28	UniformBufferArrayDynamicIndexing	Shader	
	Block-decorated arrays in uniform storage		
	classes use dynamically uniform indexing.		
29	SampledImageArrayDynamicIndexing	Shader	
	Arrays of sampled images use dynamically		
	uniform indexing.		
30	StorageBufferArrayDynamicIndexing	Shader	
	<b>BufferBlock</b> -decorated arrays in uniform		
	storage classes use dynamically uniform		
	indexing.		
31	StorageImageArrayDynamicIndexing	Shader	
	Arrays of non-sampled images are		
	accessed with dynamically uniform		
	indexing.		
32	ClipDistance	Shader	
2.5	Uses the ClipDistance BuiltIn.	g	
33	CullDistance	Shader	
2.4	Uses the CullDistance BuiltIn.	0 1101	
34	ImageCubeArray	SampledCubeArr	ay
	Uses the <b>Cube</b> Dim with the <i>Arrayed</i>		
	operand in OpTypeImage, without a		
25	sampler.	Shader	
35	SampleRateShading	Snader	
36	Uses per-sample rate shading.  ImageRect	SampledRect	
30	Uses the <b>Rect</b> Dim without a sampler.	Sampleukect	
37	SampledRect	Shader	
31	Uses the <b>Rect</b> Dim with a sampler.	Shauer	
38	GenericPointer	Addresses	
38	Uses the Generic Storage Class.	Auuresses	
	Oses the Generic Storage Class.		

	Capability	Depends On	Enabled by Extension
39	Int8	Kernel	
	Uses OpTypeInt to declare 8-bit integer		
	types.		
40	InputAttachment	Shader	
	Uses the SubpassData Dim.		
41	SparseResidency	Shader	
	Uses <b>OpImageSparse</b> instructions.		
42	MinLod	Shader	
	Uses the <b>MinLod</b> Image Operand.		
43	Sampled1D		
	Uses the <b>1D</b> Dim with a sampler.		
44	Image1D	Sampled1D	
	Uses the <b>1D</b> Dim without a sampler.		
45	SampledCubeArray	Shader	
	Uses the <b>Cube</b> Dim with the <i>Arrayed</i>		
	operand in OpTypeImage, with a sampler.		
46	SampledBuffer		
	Uses the <b>Buffer</b> Dim with a sampler.		
47	ImageBuffer	SampledBuffer	
	Uses the <b>Buffer</b> Dim without a sampler.		
48	ImageMSArray	Shader	
	An MS operand in OpTypeImage indicates		
	multisampled, used without a sampler.		
49	StorageImageExtendedFormats	Shader	
	One of a large set of more advanced image		
	formats are used, namely one of those in		
	the Image Format table listed as requiring		
	this capability.		
50	ImageQuery	Shader	
	The sizes, number of samples, or lod, etc.		
	are queried.		
51	DerivativeControl	Shader	
	Uses fine or coarse-grained derivatives,		
	e.g., OpDPdxFine.		
52	InterpolationFunction	Shader	
	Uses one of the InterpolateAtCentroid,		
	InterpolateAtSample, or		
	InterpolateAtOffset GLSL.std.450		
	extended instructions.	CI. I	
53	TransformFeedback	Shader	
5 A	Uses the Xfb Execution Mode.	Coomotur	
54	GeometryStreams Uses multiple numbered streams for	Geometry	
	<u> </u>		
55	geometry-stage output.  StorageImageReadWithoutFormat	Shader	
33	OpImageRead can use the Unknown	Shauer	
	Image Format.		
56	StorageImageWriteWithoutFormat	Shader	
30	OpImageWrite can use the Unknown	Silauti	
	Image Format.		
57	MultiViewport	Geometry	
"	Multiple viewports are used.	Geometry	

	Capability	Depends On	<b>Enabled by Extension</b>
58	SubgroupDispatch	DeviceEnqueue	
	Uses subgroup dispatch instructions.		
		Missing before	
		version 1.1.	
59	NamedBarrier	Kernel	
	Uses OpTypeNamedBarrier.		
		Missing before	
		version 1.1.	
60	PipeStorage	Pipes	
	Uses OpTypePipeStorage.	_	
		Missing before	
		version 1.1.	
4423	SubgroupBallotKHR	Reserved.	SPV_KHR_shader_ballot
4427	DrawParameters	Shader	SPV_KHR_shader_draw_parameters
		Reserved.	
4431	SubgroupVoteKHR	Reserved.	SPV_KHR_subgroup_vote
4433	StorageBuffer16BitAccess	Reserved.	SPV_KHR_16bit_storage
4433	StorageUniformBufferBlock16	Reserved.	SPV_KHR_16bit_storage
4434	UniformAndStorageBuffer16BitAccess	StorageBuffer16I	SitAccess,SPV_KHR_16bit_storage
		StorageUni-	
		form-	
		BufferBlock16	
		Reserved.	
4434	StorageUniform16	StorageBuffer16I	SitAccess,SPV_KHR_16bit_storage
		StorageUni-	/
		form-	
		BufferBlock16	
		Reserved.	
4435	StoragePushConstant16	Reserved.	SPV_KHR_16bit_storage
4436	StorageInputOutput16	Reserved.	SPV_KHR_16bit_storage
4437	DeviceGroup	Reserved.	
4439	MultiView	Shader	SPV_KHR_multiview
		Reserved.	
4441	VariablePointersStorageBuffer	Shader	SPV_KHR_variable_pointers
		Reserved.	
4442	VariablePointers		torage <b>BBMeKHR_</b> variable_pointers
		Reserved.	
4445	AtomicStorageOps	Reserved.	SPV_KHR_shader_atomic_counter_ops
4447	SampleMaskPostDepthCoverage	Reserved.	SPV_KHR_post_depth_coverage
5009	ImageGatherBiasLodAMD	Shader	SPV_AMD_texture_gather_bias_lod
		Reserved.	
5010	FragmentMaskAMD	Shader	SPV_AMD_shader_fragment_mask
		Reserved.	

Capability		Depends On	Enabled by Extension
5013	StencilExportEXT	Shader	SPV_EXT_shader_stencil_export
		Reserved.	
5015	ImageReadWriteLodAMD	Shader	SPV_AMD_shader_image_load_store_lod
		D 1	
5249	SampleMaskOverrideCoverageNV	Reserved.	ingPV_NV_sample_mask_override_coverage
3249	SampleMaskOverrideCoverageNv	Sampiekatesnau	mgr v_iv _sample_mask_override_coverage
		Reserved.	
5251	GeometryShaderPassthroughNV	Geometry	SPV_NV_geometry_shader_passthrough
		Reserved.	
5254	ShaderViewportIndexLayerEXT	MultiViewport	SPV_EXT_shader_viewport_index_layer
		Reserved.	
5254	ShaderViewportIndexLayerNV	MultiViewport	SPV_NV_viewport_array2
		D 1	
5055	Cl J XV AM LNXV	Reserved.	L. I. I. CINWANIA
5255	ShaderViewportMaskNV	Snaderviewporti	ndexLaySHVVNV_viewport_array2
		Reserved.	
5259	ShaderStereoViewNV		MaskNVV NV stereo view rendering
3237	Shader Stereo viewi	Shader viewports	Tusking V_1VV_stereo_view_rendering
		Reserved.	
5260	PerViewAttributesNV	MultiView	SPV_NVX_multiview_per_view_attributes
		Reserved.	
5265	FragmentFullyCoveredEXT	Shader	SPV_EXT_fragment_fully_covered
7.7.60	G I GG INTEND	Reserved.	CIDAL INTERES
5568	SubgroupShuffleINTEL	Reserved.	SPV_INTEL_subgroups
5569	SubgroupBufferBlockIOINTEL	Reserved.	SPV_INTEL_subgroups
5570	SubgroupImageBlockIOINTEL	Reserved.	SPV_INTEL_subgroups

# 3.32 Instructions

Form for each instruction:

Opcode Name			Capability		
			Enabling		
Instruction descript	ion.		Capabilities		
_			(when needed)		
Word Count is the h					
instruction, holding					
takes a variable nun					
say "+ variable", aft					
instruction.					
Opcode is the low-o					
instruction, holding					
mstraction, notaing					
Results, when prese					
created by the instruction. Each one is always 32 bits.					
Onerands when pre					
Operands, when present, are any literals, other instruction's Result $\langle id \rangle$ , etc., consumed by the					
instruction. Each or					
Word Count	Opcode	Results	Operands		

#### 3.32.1 Miscellaneous Instructions

OpNop				
This has no semantic impact and can safely be removed from a				
module.				
1	0			

OpUndef	f				
Make an intermediate object whose value is undefined.					
Result Type is the type of object to make.					
Each consumption of <i>Result <id></id></i> yields an arbitrary, possibly different bit pattern or abstract value resulting in possibly different concrete, abstract, or opaque values.					
3	1	< <i>id</i> >	Result <id></id>		
		Result Type			

OpSizeOf	Capability: Addresses
	Addresses
Computes the run-time size of the type pointed to by <i>Pointer</i>	
	Missing before
Result Type must be a 32-bit integer type scalar.	version 1.1.
<i>Pointer</i> must point to a concrete type.	

4	321	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Pointer

#### 3.32.2 Debug Instructions

# **OpSourceContinued**

Continue specifying the *Source* text from the previous instruction. This has no semantic impact and can safely be removed from a module.

Continued Source is a continuation of the source text in the previous Source.

The previous instruction must be an OpSource or an OpSourceContinued instruction. As is true for all literal strings, the previous instruction's string was nul terminated. That terminating 0 word from the previous instruction is not part of the source text; the first character of *Continued Source* logically immediately follows the last character of *Source* before its nul.

2 + variable	2	Literal String
		Continued Source

# **OpSource**

Document what source language and text this module was translated from. This has no semantic impact and can safely be removed from a module.

Version is the version of the source language. This literal operand is limited to a single word.

File is an OpString instruction and is the source-level file name.

Source is the text of the source-level file.

Each client API describes what form the *Version* operand takes, per source language.

		I	· · · · · · · · · · · · · · · · · · ·	66	
3 + variable	3	Source Language	Literal Number	Optional	Optional
			Version	< <i>id</i> >	Literal String
				File	Source

## **OpSourceExtension**

Document an extension to the source language. This has no semantic impact and can safely be removed from a module.

*Extension* is a string describing a source-language extension. Its form is dependent on the how the source language describes extensions.

2 + variable	4	Literal String
		Extension

# **OpName**

Assign a name string to another instruction's *Result <id>*. This has no semantic impact and can safely be removed from a module.

Target is the Result  $\langle id \rangle$  to assign a name to. It can be the Result  $\langle id \rangle$  of any other instruction; a variable, function, type, intermediate result, etc.

Name is the string to assign.

3 + variable	5	<id>&gt;</id>	Literal String
		Target	Name

# **OpMemberName**

Assign a name string to a member of a structure type. This has no semantic impact and can safely be removed from a module.

*Type* is the *<id>* from an OpTypeStruct instruction.

*Member* is the number of the member to assign in the structure. The first member is member 0, the next is member 1, ... This literal operand is limited to a single word.

Name is the string to assign to the member.

4 + variable	6	< <i>id</i> >	Literal Number	Literal String	
		Туре	Member	Name	

# **OpString**

Assign a *Result* <*id*> to a string for use by other debug instructions (see OpLine and OpSource). This has no semantic impact and can safely be removed from a module. (Removal also requires removal of all instructions referencing *Result* <*id*>.)

String is the literal string being assigned a Result <id>.

0	 -			
3 + variable		7	Result <id></id>	Literal String
				String

# **OpLine**

Add source-level location information. This has no semantic impact and can safely be removed from a module.

This location information applies to the instructions physically following this instruction, up to the first occurrence of any of the following: the next end of block, the next **OpLine** instruction, or the next **OpNoLine** instruction.

File must be an OpString instruction and is the source-level file name.

Line is the source-level line number. This literal operand is limited to a single word.

Column is the source-level column number. This literal operand is limited to a single word.

**OpLine** can generally immediately precede other instructions, with the following exceptions:

- it may not be used until after the annotation instructions, (see the Logical Layout section)
- cannot be the last instruction in a block, which is defined to end with a termination instruction
- if a branch merge instruction is used, the last **OpLine** in the block must be before its merge instruction

		, <u> </u>		
4	8	< <i>id</i> >	Literal Number	Literal Number
		File	Line	Column

## **OpNoLine**

Discontinue any source-level location information that might be active from a previous OpLine instruction. This has no semantic impact and can safely be removed from a module.

This instruction can only appear after the annotation instructions (see the Logical Layout section). It cannot be the last instruction in a block, or the second-to-last instruction if the block has a merge instruction. There is not a requirement that there is a preceding **OpLine** instruction.

1 317

OpModuleProcessed		Capability:
Document a process that was applied to a module semantic impact and can safely be removed from	Missing before version 1.1.	
<i>Process</i> is a string describing a process and/or to	ol (processor)	
that did the processing. Its form is dependent on		
2 + variable	330	Literal String
		Process

#### 3.32.3 Annotation Instructions

## **OpDecorate**

Add a Decoration to another  $\langle id \rangle$ .

*Target* is the  $\langle id \rangle$  to decorate. It can potentially be any  $\langle id \rangle$  that is a forward reference. A set of decorations can be grouped together by having multiple decoration instructions targeting the same OpDecorationGroup instruction.

This instruction is only valid when the *Decoration* operand is a decoration that takes no **Extra Operands**, or takes **Extra Operands** that are not  $\langle id \rangle$  operands.

3 + variable	71	< <i>id</i> >	Decoration	Literal, Literal,
		Target		See Decoration.

#### **OpMemberDecorate**

Add a Decoration to a member of a structure type.

Structure type is the <id> of a type from OpTypeStruct.

*Member* is the number of the member to decorate in the type. The first member is member 0, the next is member  $1, \ldots$ 

Note: See OpDecorate for creating groups of decorations for consumption by OpGroupMemberDecorate

4 + variable	72	<id></id>	Literal Number	Decoration	Literal, Literal,
		Structure Type	Member		See Decoration.

## **OpDecorationGroup**

A collector for Decorations from OpDecorate and OpDecorateId instructions. All such decoration instructions targeting this **OpDecorationGroup** instruction must precede it. Subsequent OpGroupDecorate and OpGroupMemberDecorate instructions that consume this instruction's *Result <id>* will apply these decorations to their targets.

their targets.	their targets.				
2	73	Result <id></id>			

# **OpGroupDecorate**

Add a group of Decorations to another  $\langle id \rangle$ .

Decoration Group is the <id> of an OpDecorationGroup instruction.

Targets is a list of  $\langle id \rangle$ s to decorate with the groups of decorations. The Targets list must not include the  $\langle id \rangle$  of any OpDecorationGroup instruction.

3 1			
2 + variable   74		<id>, <id>,</id></id>	
		Decoration Group	Targets

# **OpGroupMemberDecorate**

Add a group of Decorations to members of structure types.

*Decoration Group* is the *<id>* of an OpDecorationGroup instruction.

*Targets* is a list of  $(\langle id \rangle, Member)$  pairs to decorate with the groups of decorations. Each  $\langle id \rangle$  in the pair must be a target structure type, and the associated Member is the number of the member to decorate in the type. The first member is member 0, the next is member 1, ...

2 + variable	75	< <i>id&gt;</i>	<id>, literal,</id>
		Decoration Group	<id>, literal,</id>
			Targets

OpDecorateId				Capability:		
Add a Decoration to a	Add a Decoration to another $\langle id \rangle$ , using $\langle id \rangle s$ as Extra Operands.					
A set of decorations of targeting the same Op  This instruction is on	Target is the <id> to decorate. It can potentially be any <id> that is a forward reference. A set of decorations can be grouped together by having multiple decoration instructions targeting the same OpDecorationGroup instruction.  This instruction is only valid when the Decoration operand is a decoration that takes</id></id>					
constant instructions.	-	rands. All such <	d> Extra Operands must be			
3 + variable	332	< <i>id</i> >	Decoration	<id>, <id>,</id></id>		
		Target		See Decoration.		

#### 3.32.4 Extension Instructions

# OpExtension

Declare use of an extension to SPIR-V. This allows validation of additional instructions, tokens, semantics, etc.

Name is the extension's name string.

2 + variable	10	Literal String
		Name

# **OpExtInstImport**

Import an extended set of instructions. It can be later referenced by the *Result <id>*.

*Name* is the extended instruction-set's name string. There must be an external specification defining the semantics for this extended instruction set.

See Extended Instruction Sets for more information.

3 + variable	11	Result <id></id>	Literal String
			Name

#### OpExtInst

Execute an instruction in an imported set of extended instructions.

Result Type is as defined, per Instruction, in the external specification for Set.

Set is the result of an OpExtInstImport instruction.

*Instruction* is the enumerant of the instruction to execute within *Set*. This literal operand is limited to a single word. The semantics of the instruction must be defined in the external specification for *Set*.

*Operand 1, ...* are the operands to the extended instruction.

5 + variable	12	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Literal Number	<id>, <id>,</id></id>
		Result Type		Set	Instruction	
						Operand 1,
						Operand 2,

#### 3.32.5 Mode-Setting Instructions

OpMemory	OpMemoryModel					
Set addressi	Set addressing model and memory model for the entire module.					
Addressing	Addressing Model selects the module's Addressing Model.					
Memory Mo	Memory Model selects the module's memory model, see Memory Model.					
3 14 Addressing Model Memory Model						

## **OpEntryPoint**

Declare an entry point and its execution model.

Execution Model is the execution model for the entry point and its static call tree. See Execution Model.

*Entry Point* must be the *Result <id>* of an OpFunction instruction.

*Name* is a name string for the entry point. A module cannot have two **OpEntryPoint** instructions with the same Execution Model and the same *Name* string.

Interface is a list of <id> of global OpVariable instructions with either Input or Output for its Storage Class operand. These declare the input/output interface of the entry point. They could be a subset of the input/output declarations of the module, and a superset of those referenced by the entry point's static call tree. It is invalid for the entry point's static call tree to reference such an <id> if it was not listed with this instruction.

*Interface* <*id*> are forward references. They allow declaration of all variables forming an interface for an entry point, whether or not all the variables are actually used by the entry point.

4 + variable	15	Execution Model	<id></id>	Literal String	<id>, <id>,</id></id>
			Entry Point	Name	Interface

#### **OpExecutionMode**

Declare an execution mode for an entry point.

Entry Point must be the Entry Point <id> operand of an OpEntryPoint instruction.

*Mode* is the execution mode. See Execution Mode.

This instruction is only valid when the *Mode* operand is an execution mode that takes no **Extra Operands**, or takes **Extra Operands** that are not *<id>>* operands.

3 + variable	16	< <i>id</i> >	Execution Mode	Literal, Literal,
		Entry Point	Mode	See Execution Mode

# OpCapability Declare a capability used by this module. Capability is the capability declared by this instruction. There are no restrictions on the order in which capabilities are declared. See the capabilities section for more detail.

Capability Capability

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<b>OpExecutionMode</b>	Capability:					
Declare an execution	Missing before version 1.2.					
Entry Point must be	Entry Point must be the Entry Point <id> operand of an OpEntryPoint instruction.</id>					
<i>Mode</i> is the execution						
This instruction is only valid when the <i>Mode</i> operand is an execution mode that takes						
Extra Operands that constant instructions						
3 + variable	331	< <i>id</i> >	Execution Mode	< <i>id</i> >, < <i>id</i> >,		
		Entry Point	Mode	See Execution Mode		

#### 3.32.6 Type-Declaration Instructions

OpTypeVoid		
Declare the void type.		
2	19	Result <id></id>

# OpTypeBool

Declare the Boolean type. Values of this type can only be either **true** or **false**. There is no physical size or bit pattern defined for these values. If they are stored (in conjunction with OpVariable), they can only be used with logical addressing operations, not physical, and only with non-externally visible shader Storage Classes: **Workgroup**,

CrossWorkgroup, Private, and Function.

0 1		
2	20	Result <id></id>

## **OpTypeInt**

Declare a new integer type.

*Width* specifies how many bits wide the type is. This literal operand is limited to a single word. The bit pattern of a signed integer value is two's complement.

Signedness specifies whether there are signed semantics to preserve or validate.

0 indicates unsigned, or no signedness semantics

1 indicates signed semantics.

In all cases, the type of operation of an instruction comes from the instruction's opcode, not the signedness of the operands.

4	21	Result <id></id>	Literal Number	Literal Number	
			Width	Signedness	

OpTypeFloat						
Declare a new floating-point type.						
Width and	aifias haw	many hita wide the type i	s. The hit nettern of a			
		many bits wide the type i	1			
floating-p	floating-point value is as described by the IEEE 754 standard.					
3	22	Result <id></id>	Literal Number			
			Width			

# **OpTypeVector**

Declare a new vector type.

Component Type is the type of each component in the resulting type. It must be a scalar type.

Component Count is the number of components in the resulting type. It must be at least 2.

Components are numbered consecutively, starting with 0.

Components are numbered consecutivery, starting with 0.						
4	23	Result <id></id>	< <i>id&gt;</i>	Literal Number		
			Component Type	Component Count		

OpTypeN	<b>Iatrix</b>			Capability:
Declare a	new matrix ty	Matrix		
Column T	<i>type</i> is the type	e of each column in the	e matrix. It must be vector t	ype.
Column C	Count is the nu	umber of columns in th	e new matrix type. It must	be at least 2.
	lumns are nur ently of any D			
RowMajo	or or MatrixS			
4	24	Result <id></id>	<id>&gt;</id>	Literal Number
			Column Type	Column Count

## **OpTypeImage**

Declare a new image type. Consumed, for example, by OpTypeSampledImage. This type is opaque: values of this type have no defined physical size or bit pattern.

Sampled Type is the type of the components that result from sampling or reading from this image type. Must be a scalar numerical type or OpTypeVoid.

Dim is the image dimensionality (Dim).

*Depth* is whether or not this image is a depth image. (Note that whether or not depth comparisons are actually done is a property of the sampling opcode, not of this type declaration.)

- 0 indicates not a depth image
- 1 indicates a depth image
- 2 means no indication as to whether this is a depth or non-depth image

Arrayed must be one of the following indicated values:

- 0 indicates non-arrayed content
- 1 indicates arrayed content

MS must be one of the following indicated values:

- 0 indicates single-sampled content
- 1 indicates multisampled content

*Sampled* indicates whether or not this image will be accessed in combination with a sampler, and must be one of the following values:

- 0 indicates this is only known at run time, not at compile time
- 1 indicates will be used with sampler
- 2 indicates will be used without a sampler (a storage image)

Image Format is the Image Format, which can be Unknown, depending on the client API.

If Dim is **SubpassData**, *Sampled* must be 2, *Image Format* must be **Unknown**, and the **Execution Model** must be **Fragment**.

Access Qualifier is an image Access Qualifier.

9+	25	Result	< <i>id</i> >	Dim	Literal	Literal	Literal	Literal	Image	Optional
variable		<id></id>	Sampled		Number	Number	Number	Number	Format	Access
			Type		Depth	Arrayed	MS	Sampled		Quali-
					_	,		-		fier

OpTypeSamp	pler			
Declare the sampler type. Consumed by OpSampledImage. This type is opaque: values of this type have no defined physical size or bit pattern.				
2	26	Result <id></id>		

# **OpTypeSampledImage**

Declare a sampled image type, the *Result Type* of OpSampledImage, or an externally combined sampler and image. This type is opaque: values of this type have no defined physical size or bit pattern.

*Image Type* must be an OpTypeImage. It is the type of the image in the combined sampler and image type.

3	27	Result <id></id>	< <i>id&gt;</i>
			Image Type

# **OpTypeArray**

Declare a new array type: a dynamically-indexable ordered aggregate of elements all having the same type.

*Element Type* is the type of each element in the array.

*Length* is the number of elements in the array. It must be at least 1. *Length* must come from a constant instruction of an integer-type scalar whose value is at least 1.

Array elements are number consecutively, starting with 0.

	•	•				
4		28	Result <id></id>	< <i>id</i> >	<id></id>	
				Element Type	Length	

OpTypeRuntin	meArray		Capability:
Declare a new time.	run-time array ty	Shader	
Element Type is concrete type.	s the type of each		
See OpArrayLo	ength for getting		
Objects of this	type can only be		
Uniform Stora	ge Class.		
3	29	Result <id></id>	<id>&gt;</id>
			Element Type

## **OpTypeStruct**

Declare a new structure type: an aggregate of zero or more potentially heterogeneous members.

*Member N type* is the type of member N of the structure. The first member is member 0, the next is member  $1, \ldots$ 

If an operand is not yet defined, it must be defined by an OpTypePointer, where the type pointed to is an OpTypeStruct.

2 + variable	30	Result <id></id>	<id>, <id>,</id></id>
			Member 0 type,
			member 1 type,

OpTypeOpaqu	e		Capability:
			Kernel
Declare a structus specified.			
3 + variable	31	Result <id></id>	Literal String
			The name of the
			opaque type.

# **OpTypePointer**

Declare a new pointer type.

Storage Class is the Storage Class of the memory holding the object pointed to. If there was a forward reference to this type from an OpTypeForwardPointer, the Storage Class of that instruction must equal the Storage Class of this instruction.

*Type* is the type of the object pointed to.

4	32	Result <id></id>	Storage Class	<id>&gt;</id>
				Туре

# **OpTypeFunction**

Declare a new function type.

OpFunction will use this to declare the return type and parameter types of a function. **OpFunction** is the only valid use of **OpTypeFunction**.

*Return Type* is the type of the return value of functions of this type. It must be a concrete or abstract type, or a pointer to such a type. If the function has no return value, *Return Type* must be OpTypeVoid.

Parameter N Type is the type  $\langle id \rangle$  of the type of parameter N.

The different Type is the type of the type of parameter Type				
3 + variable	33	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>
			Return Type	Parameter 0 Type,
				Parameter 1 Type,

OpTypeEvent		Capability:
		Kernel
Declare an OpenCL event		
type.		
2	34	Result <id></id>

OpTypeDeviceEvent		Capability:
		DeviceEnqueue
Declare an OpenCL		_
device-side event type.		
2	35	Result <id></id>

OpTypeReserveId		Capability:
		Pipes
Declare an OpenCL		
reservation id type.		
2	36	Result <id></id>

OpTypeQueue		Capability:
		DeviceEnqueue
Declare an OpenCL queue		
type.		
2	37	Result <id></id>

OpTypePipe			Capability:
Declare an OpenCL pipe type.  Qualifier is the pipe access qualifier.		Pipes	
3	38	Result <id></id>	Access Qualifier
			Qualifier

OpTypeForwa	rdPointer		Capability:
			Addresses
Declare the Stor	rage Class for a		
Pointer Type is a forward reference to the result of an OpTypePointer. The type of object the pointer points to is declared by the OpTypePointer instruction, not this instruction. Subsequent OpTypeStruct instructions can use Pointer Type as an operand.			
Storage Class is the Storage Class of the memory holding the object			
pointed to.			
3	39	<id>&gt;</id>	Storage Class
		Pointer Type	

OpTypePipeStorage		Capability:
		PipeStorage
Declare the OpenCL		
pipe-storage type.		Missing before version 1.1.
2	322	Result <id></id>

<b>OpTypeNamedBarrier</b>		Capability:
		NamedBarrier
Declare the named-barrier		
type.		Missing before version 1.1.
2	327	Result <id></id>

#### 3.32.7 Constant-Creation Instructions

OpCons	<b>OpConstantTrue</b>				
Declare	Declare a <b>true</b> Boolean-type scalar constant.				
Result T	Result Type must be the scalar Boolean type.				
3 41 < <i>id</i> > Result < <i>id</i> >					
		Result Type			

OpCon	stantFals			
Declare a <b>false</b> Boolean-type scalar constant.				
Result Type must be the scalar Boolean type.				
3	42	<id>&gt;</id>	Result <id></id>	
		Result Type		

#### **OpConstant**

Declare a new integer-type or floating-point-type scalar constant.

Result Type must be a scalar integer type or floating-point type.

*Value* is the bit pattern for the constant. Types 32 bits wide or smaller take one word. Larger types take multiple words, with low-order words appearing first.

3 + variable	43	< <i>id</i> >	Result <id></id>	Literal, Literal,
		Result Type		Value

# **OpConstantComposite**

Declare a new composite constant.

*Result Type* must be a composite type, whose top-level members/elements/components/columns have the same type as the types of the *Constituents*. The ordering must be the same between the top-level types in *Result Type* and the *Constituents*.

Constituents will become members of a structure, or elements of an array, or components of a vector, or columns of a matrix. There must be exactly one *Constituent* for each top-level member/element/component/column of the result. The *Constituents* must appear in the order needed by the definition of the *Result Type*. The *Constituents* must all be <id>s of other constant declarations or an OpUndef.

3 + variable 44		< <i>id</i> >	Result <id></id>	<id>, <id>,</id></id>
		Result Type		Constituents

OpCor	ıstantSaı	npler			Capability: LiteralSampler	
Declare	Declare a new sampler constant.					
Result Type must be OpTypeSampler.						
Sampler Addressing Mode is the addressing mode; a literal from Sampler Addressing Mode.						
Param is one of: 0: Non Normalized 1: Normalized						
Sample	r Filter N	Mode is the filter mod	le; a literal from Sar	mpler Filter Mode.		
6	45	<id> Result Type</id>	Result <id></id>	Sampler Addressing Mode	Literal Number Param	Sampler Filter Mode

# **OpConstantNull**

Declare a new null constant value.

The *null* value is type dependent, defined as follows:

- Scalar Boolean: false
- Scalar integer: 0
- Scalar floating point: +0.0 (all bits 0)
- All other scalars: Abstract
- Composites: Members are set recursively to the null constant according to the null value of their constituent types.

Result Type must be one of the following types:

- Scalar or vector Boolean type
- Scalar or vector integer type
- Scalar or vector floating-point type
- Pointer type
- Event type
- Device side event type
- Reservation id type
- Queue type
- Composite type

F	The state of the s					
3 46		< <i>id</i> >	Result <id></id>			
		Result Type				

# **OpSpecConstantTrue**

Declare a Boolean-type scalar specialization constant with a default value of true.

This instruction can be specialized to become either an OpConstantTrue or OpConstantFalse instruction.

Result Type must be the scalar Boolean type.

See Specialization.

_ L				
	3   48		< <i>id</i> >	Result <id></id>
			Result Type	

# **OpSpecConstantFalse**

Declare a Boolean-type scalar specialization constant with a default value of **false**.

This instruction can be specialized to become either an OpConstantTrue or OpConstantFalse instruction.

Result Type must be the scalar Boolean type.

See Specialization.

	3	49	<id></id>	Result <id></id>	
			Result Type		

# **OpSpecConstant**

Declare a new integer-type or floating-point-type scalar specialization constant.

Result Type must be a scalar integer type or floating-point type.

*Value* is the bit pattern for the default value of the constant. Types 32 bits wide or smaller take one word. Larger types take multiple words, with low-order words appearing first.

This instruction can be specialized to become an OpConstant instruction.

#### See Specialization.

3 + variable	3 + variable 50		<id> Result <id> Literal, Literal</id></id>	
		Result Type		Value

## **OpSpecConstantComposite**

Declare a new composite specialization constant.

*Result Type* must be a composite type, whose top-level members/elements/components/columns have the same type as the types of the *Constituents*. The ordering must be the same between the top-level types in *Result Type* and the *Constituents*.

Constituents will become members of a structure, or elements of an array, or components of a vector, or columns of a matrix. There must be exactly one *Constituent* for each top-level member/element/component/column of the result. The *Constituents* must appear in the order needed by the definition of the type of the result. The *Constituents* must be the *<id>* of other specialization constant or constant declarations.

This instruction will be specialized to an OpConstantComposite instruction.

# See Specialization.

3 + variable	51	< <i>id</i> >	Result <id></id>	<id>, <id>,</id></id>
		Result Type		Constituents

## **OpSpecConstantOp**

Declare a new specialization constant that results from doing an operation.

Result Type must be the type required by the Result Type of Opcode.

Opcode must be one of the following opcodes. This literal operand is limited to a single word.

**OpSConvert**, **OpFConvert** 

OpSNegate, OpNot

OpIAdd, OpISub

OpIMul, OpUDiv, OpSDiv, OpUMod, OpSRem, OpSMod

OpShiftRightLogical, OpShiftRightArithmetic, OpShiftLeftLogical

OpBitwiseOr, OpBitwiseXor, OpBitwiseAnd

OpVectorShuffle, OpCompositeExtract, OpCompositeInsert

OpLogicalOr, OpLogicalAnd, OpLogicalNot,

OpLogicalEqual, OpLogicalNotEqual

**OpSelect** 

OpIEqual, OpINotEqual

OpULessThan, OpSLessThan

OpUGreaterThan, OpSGreaterThan

OpULessThanEqual, OpSLessThanEqual

OpUGreater Than Equal, OpSGreater Than Equal

If the **Shader** capability was declared, the following opcode is also valid:

OpQuantizeToF16

If the **Kernel** capability was declared, the following opcodes are also valid:

OpConvertFToS, OpConvertSToF

OpConvertFToU, OpConvertUToF

**OpUConvert** 

 $OpConvertPtrToU,\,OpConvertUToPtr$ 

 $OpGeneric Cast To Ptr, \, OpPtr Cast To Generic \,$ 

**OpBitcast** 

**OpFNegate** 

OpFAdd, OpFSub

OpFMul, OpFDiv

OpFRem, OpFMod

OpAccessChain, OpInBoundsAccessChain

OpPtrAccess Chain, OpInBounds PtrAccess Chain

*Operands* are the operands required by *opcode*, and satisfy the semantics of *opcode*. In addition, all *Operands* must be either:

- the  $\langle id \rangle s$  of other constant instructions, or
- **OpUndef**, when allowed by *opcode*, or
- for the AccessChain named opcodes, their Base is allowed to be a global (module scope) OpVariable instruction.

# See Specialization.

4 + variable 52		<id> Result <id></id></id>		Literal Number	< <i>id</i> >, < <i>id</i> >,
		Result Type		Opcode	Operands

#### 3.32.8 Memory Instructions

#### **OpVariable**

Allocate an object in memory, resulting in a pointer to it, which can be used with OpLoad and OpStore.

Result Type must be an OpTypePointer. Its Type operand is the type of object in memory. Its Storage Class operand must be the same as the Storage Class operand of the result type.

Storage Class is the Storage Class of the memory holding the object. It cannot be Generic.

*Initializer* is optional. If *Initializer* is present, it will be the initial value of the variable's memory content. *Initializer* must be an <*id*> from a constant instruction or a global (module scope) OpVariable instruction. *Initializer* must have the same type as the type pointed to by *Result Type*.

4 + variable 59		<id> Result <id></id></id>		Storage Class	Optional
		Result Type			<id>&gt;</id>
					Initializer

## **OpImageTexelPointer**

Form a pointer to a texel of an image. Use of such a pointer is limited to atomic operations.

Result Type must be an OpTypePointer whose Storage Class operand is **Image**. Its Type operand must be a scalar numerical type or OpTypeVoid.

*Image* must have a type of OpTypePointer with *Type* OpTypeImage. The *Sampled Type* of the type of *Image* must be the same as the *Type* pointed to by *Result Type*. The Dim operand of *Type* cannot be **SubpassData**.

Coordinate and Sample specify which texel and sample within the image to form a pointer to.

*Coordinate* must be a scalar or vector of integer type. It must have the number of components specified below, given the following *Arrayed* and Dim operands of the type of the OpTypeImage.

If *Arrayed* is 0:

1D: scalar

**2D**: 2 components

**3D**: 3 components **Cube**: 3 components **Rect**: 2 components **Buffer**: scalar

If *Arrayed* is 1:

**1D**: 2 components

2D: 3 components

**Cube**: 3 components; the face and layer combine into the 3rd component, *layer\_face*, such that face is *layer\_face* % 6 and layer is floor(*layer\_face* / 6)

Sample must be an integer type scalar. It specifies which sample to select at the given coordinate. It must be a valid <id> for the value 0 if the OpTypeImage has MS of 0.

6	60	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id></id>
		Result Type		Image	Coordinate	Sample

# OpLoad

Load through a pointer.

Result Type is the type of the loaded object.

*Pointer* is the pointer to load through. Its type must be an OpTypePointer whose *Type* operand is the same as *Result Type*.

Memory Access must be a Memory Access literal. If not present, it is the same as specifying None.

4 + variable	61	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Optional
		Result Type		Pointer	Memory Access

## **OpStore**

Store through a pointer.

*Pointer* is the pointer to store through. Its type must be an OpTypePointer whose *Type* operand is the same as the type of *Object*.

*Object* is the object to store.

Memory Access must be a Memory Access literal. If not present, it is the same as specifying None.

			· · · · · · · · · · · · · · · · · · ·	9
3 + variable	62	< <i>id</i> >	< <i>id</i> >	Optional
		Pointer	Object	Memory Access

# **OpCopyMemory**

Copy from the memory pointed to by *Source* to the memory pointed to by *Target*. Both operands must be non-void pointers of the same type. Matching Storage Class is not required. The amount of memory copied is the size of the type pointed to.

Memory Access must be a Memory Access literal. If not present, it is the same as specifying None.

-				
3 + variable	63	< <i>id</i> >	< <i>id&gt;</i>	Optional
		Target	Source	Memory Access

<b>OpCopyMemoryS</b>	ized				Capability:
Copy from the men	Addresses				
Size is the number of instruction, the con Signedness of 1 and unsigned, and if its	as				
Memory Access mu	st be a M	emory Access li	teral. If not present, it i	s the same as specifyi	ng
None.					
4 + variable	64	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	Optional
		Target	Source	Size	Memory Access

## **OpAccessChain**

Create a pointer into a composite object that can be used with OpLoad and OpStore.

Result Type must be an OpTypePointer. Its Type operand must be the type reached by walking the Base's type hierarchy down to the last provided index in Indexes, and its Storage Class operand must be the same as the Storage Class of Base.

Base must be a pointer, pointing to the base of a composite object.

*Indexes* walk the type hierarchy to the desired depth, potentially down to scalar granularity. The first index in *Indexes* will select the top-level member/element/component/element of the base composite. All composite constituents use zero-based numbering, as described by their **OpType...** instruction. The second index will apply similarly to that result, and so on. Once any non-composite type is reached, there must be no remaining (unused) indexes. Each of the *Indexes* must:

- be a scalar integer type,
- be an OpConstant when indexing into a structure.

1					
4 + variable	65	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>
		Result Type		Base	Indexes

OpInBoundsAccessChain								
	Has the same semantics as OpAccessChain, with the addition that the resulting pointer is known to point within the base object.							
4 + variable	66	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>			
		Result Type		Base				
					Indexes			

<b>OpPtrAccessCh</b>	ain	Capability:				
Has the same semantics as OpAccessChain, with the addition of the <i>Element</i> operand.					Addresses, Var VariablePointe	iablePointers, rsStorageBuffer
Element is used to address of the first computed to be to Base after being of Base. When the dereferenced as a Note: If Base is conceptation is to see directly used, as	st element he base to derefere he type of an array poriginally					
5 + variable	67	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id>, <id>,</id></id>
		Result Type		Base	Element	 Indexes

OpArra	OpArrayLength							
Length	Shader							
Result T								
Structur	Structure must be a pointer to an OpTypeStruct whose last member is a run-time array.							
Array m	ember is t	he index of the last memb	er of the structure that	Structure points to. That				
member								
5	68	Literal Number						
		Result Type		Structure	Array member			

OpGenerio	cPtrMemSei	mantics		Capability:
Class for th	ne specific (ne	ry Semantics which includes ma on-Generic) Storage Class of Po- eneric Storage Class.	_	Kernel
Result Type	must be an			
4	69	<id>&gt;</id>	Result <id></id>	<id></id>
		Result Type		Pointer

OpInBounds	PtrAcc	essChain			Capability: Addresses	
Has the same that the result						
5 + variable	70	<id> Result Type</id>	Result <id></id>	<id> Base</id>	<id> Element</id>	<id>, <id>,  Indexes</id></id>

#### 3.32.9 Function Instructions

## **OpFunction**

Add a function. This instruction must be immediately followed by one OpFunctionParameter instruction per each formal parameter of this function. This function's body or declaration will terminate with the next OpFunctionEnd instruction.

The *Result <id>* cannot be used generally by other instructions. It can only be used by OpFunctionCall, OpEntryPoint, and decoration instructions.

Result Type must be the same as the Return Type declared in Function Type.

Function Type is the result of an OpTypeFunction, which declares the types of the return value and parameters of the function.

5 54 < <i>id&gt;</i>		Result <id></id>	Function Control	< <i>id</i> >	
		Result Type			Function Type

## **OpFunctionParameter**

Declare a formal parameter of the current function.

Result Type is the type of the parameter.

This instruction must immediately follow an OpFunction or OpFunctionParameter instruction. The order of contiguous **OpFunctionParameter** instructions is the same order arguments will be listed in an OpFunctionCall instruction to this function. It is also the same order in which *Parameter Type* operands are listed in the OpTypeFunction of the *Function Type* operand for this function's OpFunction instruction.

		**		
3	55	< <i>id</i> >	Res	sult <id></id>
		Result Type		

OpFunctionEnd	
Last instruction of a function.	
1	56

# **OpFunctionCall**

Call a function.

*Result Type* is the type of the return value of the function. It must be the same as the *Return Type* operand of the *Function Type* operand of the *Function* operand.

Function is an OpFunction instruction. This could be a forward reference.

Argument N is the object to copy to parameter N of Function.

**Note:** A forward call is possible because there is no missing type information: *Result Type* must match the *Return Type* of the function, and the calling argument types must match the formal parameter types.

-7/F =							
4 + variable	57	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>		
		Result Type		Function	Argument 0,		
					Argument 1,		

# 3.32.10 Image Instructions

# **OpSampledImage**

Create a sampled image, containing both a sampler and an image.

Result Type must be the OpTypeSampledImage type.

*Image* is an object whose type is an OpTypeImage, whose *Sampled* operand is 0 or 1, and whose Dim operand is not **SubpassData**.

Sampler must be an object whose type is OpTypeSampler.

5	86	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Image	Sampler

OpImageSam	pleImp	olicitLod				Capability: Shader	
Sample an ima	age with	n an implicit leve	el of detail.				
Result Type m type. Its comp OpTypeImage	onents	ying					
Sampled Imag	e must	ge.					
array layer]) a	as neede	scalar or vector ed by the definiti nused componer	on of Sampled I	mage. It may be	e a vector larger		
Image Operan	ds enco	odes what operar	nds follow, as pe	r Image Operano	ds.		
		ly valid in the <b>F</b> r derivative that c	_		ddition, it		
5 + variable	87	<id> Result Type</id>	Result <id></id>	<id> Sampled Image</id>	<id> Coordinate</id>	Optional Image Operands	Optional < <i>id</i> >, < <i>id</i> >,

# **OpImageSampleExplicitLod**

Sample an image using an explicit level of detail.

*Result Type* must be a vector of four components of floating-point type or integer type. Its components must be the same as *Sampled Type* of the underlying OpTypeImage (unless that underlying *Sampled Type* is **OpTypeVoid**).

Sampled Image must be an object whose type is OpTypeSampledImage.

Coordinate must be a scalar or vector of floating-point type or integer type. It contains  $(u[, v] ... [, array \, layer])$  as needed by the definition of Sampled Image. Unless the **Kernel** capability is being used, it must be floating point. It may be a vector larger than needed, but all unused components will appear after all used components.

*Image Operands* encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present.

7.	00	4:15	D14	4: Js	4: Js	T	4: Js	0.41
/ +	88	<id></id>	Result	<id></id>	< <i>id</i> >	Image	<id></id>	Optional
variable		Result	<id></id>	Sampled	Coordinate	Operands		< <i>id</i> >.
, arraore			, iu	Sampica	Coordinate	Operanas		,
		Type		Image				<id>,</id>
								·

OpImageS	ampleI	PrefImplicit	Lod				Capability: Shader		
Sample an	image d	oing depth-c	omparison wit	h an implicit le	vel of detail.		Shauei		
			integer type or OpTypeImag	floating-point e.	type. It must be	e the same as			
Sampled In	<i>age</i> mu								
Coordinate array layer than needed	]) as ne								
$D_{ref}$ is the $c$	lepth-co	mparison ref	erence value.						
Image Ope	rands er	ncodes what	operands follo	w, as per Image	Operands.				
	This instruction is only valid in the <b>Fragment</b> Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion.								
6+	89	< <i>id</i> >	Result	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	Optional	Optional	
variable		Result	<id></id>	Sampled Image	Coordinate	$D_{ref}$	Image Operands	<id>, <id>,</id></id>	

<b>OpImageS</b>	ample	eDrefExplici	tLod					Capability:	
Sample an i  Result Type Sampled Type	e same as	Shader							
Sampled Im	age m								
Coordinate must be a scalar or vector of floating-point type. It contains ( $u[, v] \dots [, array \ layer]$ ) as needed by the definition of Sampled Image. It may be a vector larger than needed but all unused components will appear after all used components. $D_{ref}$ is the depth-comparison reference value.									
D <sub>ref</sub> is the d	срш-с	zomparison n	ciciciice varu	ic.					
0 1	Image Operands encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present.								
8 +	90	<id>&gt;</id>	Result	<id>&gt;</id>	<id>&gt;</id>	<id></id>	Image	<id>&gt;</id>	Optional
variable		Result Type	<id></id>	Sampled Image	Coordinate	$D_{ref}$	Operands		< <i>id</i> >, < <i>id</i> >,

#### **OpImageSampleProjImplicitLod** Capability: Shader Sample an image with with a project coordinate and an implicit level of detail. Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is OpTypeVoid). Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be (u/q [, v/q] [, w/q]), as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. *Image Operands* encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. 5 + variable 91 Result <id> Optional Optional <*id*> <*id*> <*id*> Result Type Sampled Coordinate <*id*>, <*id*>, Image Image **Operands** . . .

#### **OpImageSampleProjExplicitLod** Capability: Shader Sample an image with a project coordinate using an explicit level of detail. Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is **OpTypeVoid**). Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be (u/q [, v/q] [, w/q]), as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. *Image Operands* encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. Optional 7+ 92 <*id*> Result $\langle id \rangle$ <*id*> Image $\langle id \rangle$ variable Result < id >Sampled Coordinate Operands <*id*>, Type Image <*id*>, ...

#### **OpImageSampleProjDrefImplicitLod** Capability: Shader Sample an image with a project coordinate, doing depth-comparison, with an implicit level of detail. Result Type must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be (u/q [, v/q] [, w/q]), as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. $D_{ref}/q$ is the depth-comparison reference value. *Image Operands* encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. Optional Optional 6+ 93 <*id*> Result <*id*> <*id*> <*id*> variable Result < id >Sampled Coordinate $D_{ref}$ **Image** <*id*>, Operands <*id*>, ... Type *Image*

#### **OpImageSampleProjDrefExplicitLod** Capability: Shader Sample an image with a project coordinate, doing depth-comparison, using an explicit level of detail. Result Type must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector containing $(u \, [\, v \, ] \, [\, w \, ], \, a)$ , as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be (u/q [, v/q] [, w/q]), as needed by the definition of Sampled *Image*. It may be a vector larger than needed, but all unused components will appear after all used components. $D_{ref}/q$ is the depth-comparison reference value. Image Operands encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. <*id*> Optional 8 + 94 <*id*> Result <*id*> <*id*> <*id*> Image Operands <*id*>, variable Result < id >Sampled Coordinate $D_{ref}$ <*id*>, . . . Type *Image*

## **OpImageFetch**

Fetch a single texel from a sampled image.

*Result Type* must be a vector of four components of floating-point type or integer type. Its components must be the same as *Sampled Type* of the underlying OpTypeImage (unless that underlying *Sampled Type* is **OpTypeVoid**).

*Image* must be an object whose type is OpTypeImage. Its Dim operand cannot be **Cube**, and its *Sampled* operand must be 1.

Coordinate is an integer scalar or vector containing  $(u[, v] \dots [, array \, layer])$  as needed by the definition of Sampled Image.

Image Operands encodes what operands follow, as per Image Operands.

0 1			, 1				
5 + variable	95	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Result Type		Image	Coordinate	Image	< <i>id</i> >, < <i>id</i> >,
						Operands	
						_	

#### OpImageGather Capability: Shader Gathers the requested component from four texels. Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is **OpTypeVoid**). It has one component per gathered texel. Sampled Image must be an object whose type is OpTypeSampledImage. Its OpTypeImage must have a Dim of 2D, Cube, or Rect. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [,$ array layer]) as needed by the definition of Sampled Image. Component is the component number that will be gathered from all four texels. It must be 0, 1, 2 or 3. Image Operands encodes what operands follow, as per Image Operands. <*id*> Result <*id*> <*id*> <*id*> Optional Optional 6+ 96 variable Result < id >Sampled Coordinate Component **Image** <*id*>, Type *Image* Operands <*id*>, ...

OpImageD	refGat	her					Capability: Shader	
Gathers the	request	ed depth-co	mparison from	four texels.			Shauei	
Result Type components that underly								
Sampled Image								
				ing-point type.	It contains $(u[,$	v] [,		
$D_{ref}$ is the de	epth-co	mparison re	ference value.					
Image Oper	<i>ands</i> er	codes what	operands follo	w, as per Image	e Operands.			
6+	97	< <i>id</i> >	Result	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	Optional	Optional
variable		Result Type	<id></id>	Sampled Image	Coordinate	$D_{ref}$	Image Operands	<id>, <id>,</id></id>

## **OpImageRead**

Read a texel from an image without a sampler.

Result Type must be a scalar or vector of floating-point type or integer type. Its component type must be the same as Sampled Type of the OpTypeImage (unless that Sampled Type is OpTypeVoid).

*Image* must be an object whose type is OpTypeImage with a *Sampled* operand of 0 or 2. If the *Sampled* operand is 2, then some dimensions require a capability; e.g., one of **Image1D**, **ImageRect**, **ImageBuffer**, **ImageCubeArray**, or **ImageMSArray**.

Coordinate is an integer scalar or vector containing non-normalized texel coordinates ( $u[, v] \dots [, array \, layer]$ ) as needed by the definition of Image. If the coordinates are outside the image, the memory location that is accessed is undefined.

When the *Image* Dim operand is **SubpassData**, *Coordinate* is relative to the current fragment location. That is, the integer value (rounded down) of the current fragment's window-relative (x, y) coordinate is added to (u, v).

When the *Image* Dim operand is not **SubpassData**, the **Image** Format must not be **Unknown**, unless the **StorageImageReadWithoutFormat** Capability was declared.

*Image Operands* encodes what operands follow, as per Image Operands.

5 + variable	98	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Result Type		Image	Coordinate	Image	<id>, <id>,</id></id>
						Operands	
						_	

## **OpImageWrite**

Write a texel to an image without a sampler.

*Image* must be an object whose type is OpTypeImage with a *Sampled* operand of 0 or 2. If the *Sampled* operand is 2, then some dimensions require a capability; e.g., one of **Image1D**, **ImageRect**, **ImageBuffer**, **ImageCubeArray**, or **ImageMSArray**. Its Dim operand cannot be **SubpassData**.

Coordinate is an integer scalar or vector containing non-normalized texel coordinates ( $u[, v] \dots [, array \, layer]$ ) as needed by the definition of Image. If the coordinates are outside the image, the memory location that is accessed is undefined.

*Texel* is the data to write. Its component type must be the same as *Sampled Type* of the OpTypeImage (unless that *Sampled Type* is **OpTypeVoid**).

The Image Format must not be Unknown, unless the StorageImageWriteWithoutFormat Capability was declared.

*Image Operands* encodes what operands follow, as per Image Operands.

		- · · · · · · · · · · · · · · · · · · ·		F		
4 + variable	99	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Image	Coordinate	Texel	Image	<id>, <id>,</id></id>
					Operands	

OpImag	OpImage									
Extract	Extract the image from a sampled image.									
Result T	<i>ype</i> must	be OpTypeImage.								
	Sampled Image must have type OpTypeSampledImage whose Image Type is the same as Result Type.									
4	100	<id></id>	Result <id></id>	<id></id>						
		Result Type		Sampled Image						

OpImage(	QueryForma	ıt		Capability:
Query the i	image format	Kernel		
· · ·	e must be a so nnel Data Ty			
Image mus	t be an objec			
4	101	< <i>id</i> >	Result <id></id>	<id>&gt;</id>
		Result Type		Image

OpImage	QueryOrder			Capability:
	channel order	Kernel		
		calar integer type. The resulting	g value is an enumerant from	
Image Cha	nnel Order.			
Image mus	st be an objec	t whose type is OpTypeImage.		
4	102	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>
		Result Type		Image

OpImage	eQuerySi		Capability:						
Query the	e dimensi	ons of <i>Image</i> for mipma	p level for <i>Level of De</i>	etail.	Kernel, ImageQuery				
	Result Type must be an integer type scalar or vector. The number of components must be								
1 for <b>1D</b>	,								
2 for <b>2D</b> ,	and Cub	e Dimensionalities,							
3 for <b>3D</b>	Dim,								
-		mage type is arrayed. T							
depth] [, d	elements]	) where <i>elements</i> is the i	number of layers in an	image array, or the nur	mber				
of cubes i	in a cube-	map array.							
<b>2D</b> , <b>3D</b> , o	<i>Image</i> must be an object whose type is OpTypeImage. Its Dim operand must be one of <b>1D</b> , <b>2D</b> , <b>3D</b> , or <b>Cube</b> , and its <i>MS</i> must be 0. See OpImageQuerySize for querying image types without level of detail.								
	Level of Detail is used to compute which mipmap level to query, as described in the API specification.								
5	103	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >				
		Result Type		Image	Level of Detail				

OpImageQuerySize	OplmageQuerySize							
Query the dimensions of Im	Kernel, ImageQuery							
Result Type must be an integ	ne number of components must							
be	• • • • • • • • • • • • • • • • • • • •	•						
1 for <b>Buffer</b> Dim,								
2 for <b>2D</b> and <b>Rect</b> Dimension	onalities,							
3 for <b>3D</b> Dim,								
plus 1 more if the image typ	be is arrayed. This vector is	filled in with (width [, height]						
[, elements]) where element	s is the number of layers in	an image array.						
<i>Image</i> must be an object wh	nose type is OpTypeImage.	Its Dim operand must be one						
of <b>Rect</b> or <b>Buffer</b> , or if its <i>I</i>	<i>MS</i> is 1, it can be <b>2D</b> , or, if	its Sampled Type is 0 or 2, it						
can be <b>2D</b> or <b>3D</b> . It cannot be	be an image with level of de	etail; there is no implicit						
level-of-detail consumed by	this instruction. See OpIm	ageQuerySizeLod for						
querying images having lev								
4 104 < <i>i</i>	id>	Result <id></id>	<id></id>					
Re	esult Type		Image					

OpImage	QueryLo	OpImageQueryLod									
					ImageQuery						
Query the	mipmap	level and the level of det	ail for a hypothetical sar	npling of <i>Image</i> at							
Coordinat	Coordinate using an implicit level of detail.										
Result Typ	Result Type must be a two-component floating-point type vector.										
The first c	omponer	nt of the result will contain	n the mipmap array laye	r.							
	-	nent of the result will cor	ntain the implicit level of	f detail relative to the							
base level.	•										
_	-	ast be an object whose type, <b>2D</b> , <b>3D</b> , or <b>Cube</b> .	oe is OpTypeSampledIm	age. Its Dim operand							
	,	, , , , , , , , , , , , , , , , , , , ,									
Coordinat	te must b	e a scalar or vector of floa	ating-point type or intege	er type. It contains $(u[,$							
v] ) as	needed l	by the definition of Sample	led Image, not including	any array layer index.							
Unless the	Kernel	capability is being used,	it must be floating point.								
If collect or	n on inco	omplete image, the results	a ara undafinad								
11 called 0	ni an incc	implete image, the results	are undermed.								
This instru	action is	only valid in the <b>Fragme</b>	nt Execution Model. In	addition, it consumes an							
implicit de											
5	105	<id>&gt;</id>	Result <id></id>	< <i>id</i> >	<id></id>						
		Result Type		Sampled Image	Coordinate						

OpImage(	QueryLevels			Capability:
Result Type	number of mine must be a so by the API sp	Image. the number of mipmap levels,	Kernel, ImageQuery	
	t be an objec  3D, or Cube	t whose type is OpTypeImage.	Its Dim operand must be one	
4	106	<id>&gt;</id>		
		Result Type		Image

OpImageQuerySampl	es		Capability:				
Result Type must be a se	mples available per texel fetch calar integer type. The result is t whose type is OpTypeImage.	the number of samples.	Kernel, ImageQuery				
4 107	107     Result						
	Result Type		Image				

#### **OpImageSparseSampleImplicitLod** Capability: **SparseResidency** Sample a sparse image with an implicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled *Type* is **OpTypeVoid**). Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [,$ array layer]) as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. Image Operands encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. 5 + variable 305 <*id*> Result <id> <*id*> $\langle id \rangle$ Optional Optional Result Type Sampled Coordinate **Image** <*id*>, <*id*>, Image Operands

#### **OpImageSparseSampleExplicitLod** Capability: **SparseResidency** Sample a sparse image using an explicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is OpTypeVoid). Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type or integer type. It contains $(u[,v]...[, array \, layer])$ as needed by the definition of Sampled Image. Unless the Kernel capability is being used, it must be floating point. It may be a vector larger than needed, but all unused components will appear after all used components. *Image Operands* encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. <*id*> Optional 7+ 306 <*id*> Result $\langle id \rangle$ Image $\langle id \rangle$ variable Result < id >Sampled Coordinate Operands <*id*>, Type Image <id>, ...

#### **OpImageSparseSampleDrefImplicitLod** Capability: **SparseResidency** Sample a sparse image doing depth-comparison with an implicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type. It contains (u[, v], ..., [, v])array layer]) as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. $D_{ref}$ is the depth-comparison reference value. *Image Operands* encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. 307 <*id*> Result <id> Optional Optional 6+ <*id*> <*id*> variable Result <id> Sampled Coordinate <*id*>, $D_{ref}$ Image <*id*>, ... Type Image **Operands**

#### **OpImageSparseSampleDrefExplicitLod** Capability: **SparseResidency** Sample a sparse image doing depth-comparison using an explicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [, array])$ layer]) as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. $D_{ref}$ is the depth-comparison reference value. Image Operands encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. <*id*> <*id*> <*id*> Optional 8 + 308 Result $\overline{\langle id \rangle}$ Image <*id*> variable Result < id >Sampled Coordinate $D_{ref}$ Operands <*id*>, Type *Image* <*id*>, ...

OpImageSpa	rseSan	Capability:					
		SparseReside	ency				
Sample a spar	rse imag						
						Reserved.	
5 + variable	309	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Result Type		Sampled	Coordinate	Image	< <i>id</i> >,
	Operands	< <i>id</i> >,					

	OpImageSparseSampleProjExplicitLod  Sample a sparse image with a projective coordinate using an explicit level of detail.								
7 + variable	310	<id> Result Type</id>	Result <id></id>	<id> Sampled Image</id>	<id> Coordinate</id>	Image Operands	<id></id>	Optional < <i>id</i> >, < <i>id</i> >,	

OpImageSı	Capability: SparseResidency									
Sample a sp	Sample a sparse image with a projective coordinate, doing depth-comparison, with an									
implicit leve	el of det	tail.					Reserved.			
6+	311	< <i>id</i> >	Result	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	Optional	Optional		
variable		Result	<id>&gt;</id>	Sampled	Coordinate	$D_{ref}$	Image	< <i>id</i> >,		
Type Image								< <i>id</i> >,		

OpImageS	OpImageSparseSampleProjDrefExplicitLod									
	Sample a sparse image with a projective coordinate, doing depth-comparison, using an explicit level of detail.									
8 +	312	<id></id>	Result	< <i>id</i> >	< <i>id</i> >	<id></id>	Image	< <i>id</i> >	Optional	
variable		Result	<id>&gt;</id>	Sampled	Coordinate	$D_{ref}$	Operands		< <i>id</i> >,	
Type Image									< <i>id</i> >,	

#### **OpImageSparseFetch** Capability: **SparseResidency** Fetch a single texel from a sampled sparse image. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled *Type* is **OpTypeVoid**). Image must be an object whose type is OpTypeImage. Its Dim operand cannot be Cube. Coordinate is an integer scalar or vector containing (u[, v] ... [, array layer]) as needed by the definition of Sampled Image. Image Operands encodes what operands follow, as per Image Operands. 5 + variable 313 <*id*> Result <id> <*id*> Optional Optional <*id*> Result Type Image Coordinate **Image** <*id*>, <*id*>, Operands . . .

OpImageSp	arseG	<b>OpImageSparseGather</b>								
Gathers the 1		SparseResion								
Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is OpTypeVoid). It has one component per gathered texel.										
_	-	-	t whose type i	s OpTypeSamp or <b>Rect</b> .	oledImage. Its					
			ector of floating		t contains (u[,	v] [,				
Component is be 0, 1, 2 or		omponent nui	mber that will	be gathered fro	om all four texe	els. It must				
Image Opera	<i>ands</i> en	codes what o	perands follow	v, as per Image	Operands.					
6 + variable	314	<id> Result Type</id>	Result <id></id>	<id> Sampled Image</id>	<id> Coordinate</id>	<id> Component</id>	Optional Image Operands	Optional < <i>id</i> >, < <i>id</i> >,		

OpImageSparseD	Capability: SparseResion	lonov						
Gathers the reques		SparseKesic	lency					
Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is OpTypeVoid). It has one component per gathered texel.  Sampled Image must be an object whose type is OpTypeSampledImage. Its OpTypeImage must have a Dim of 2D, Cube, or Rect.								
Coordinate must be array layer]) as ne				t contains $(u[, \cdot])$	v] [,			
$D_{ref}$ is the depth-co	omparison refe	rence value.						
Image Operands et	ncodes what of	perands follow	, as per Image	Operands.				
6 + 315 variable	<id> Result Type</id>	Result <id></id>	<id> Sampled Image</id>	<id> Coordinate</id>	<id> D<sub>ref</sub></id>	Optional Image Operands	Optional < <i>id</i> >, < <i>id</i> >,	

<b>OpImageS</b>	parseTexels	Resident		Capability:
uncommitte	a Resident Co ed texture mo e must be a B	SparseResidency		
Resident Coresident co		e from an <b>OpImageSparse</b> i	nstruction that returns a	
4	316	<id>&gt;</id>		
		Result Type		Resident Code

#### **OpImageSparseRead** Capability: **SparseResidency** Read a texel from a sparse image without a sampler. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar or vector of floating-point type or integer type. Its component type must be the same as Sampled *Type* of the OpTypeImage (unless that *Sampled Type* is **OpTypeVoid**). *Image* must be an object whose type is OpTypeImage with a *Sampled* operand of 2. Coordinate is an integer scalar or vector containing non-normalized texel coordinates $(u[,v]...[, array \, layer])$ as needed by the definition of *Image*. If the coordinates are outside the image, the memory location that is accessed is undefined. The *Image* Dim operand must not be **SubpassData**. The Image Format must not be Unknown unless the StorageImageReadWithoutFormat Capability was declared. *Image Operands* encodes what operands follow, as per Image Operands. 5 + variable 320 <*id*> Result <id> <*id*> <*id*> Optional Optional Result Type Image Coordinate **Image** <*id*>, <*id*>, Operands

#### 3.32.11 Conversion Instructions

# **OpConvertFToU**

Convert (value preserving) from floating point to unsigned integer, with round toward 0.0.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

*Float Value* must be a scalar or vector of floating-point type. It must have the same number of components as *Result Type*.

Results are computed per component.

Ī	4	109	< <i>id</i> >	Result <id></id>	<id></id>
			Result Type		Float Value

#### **OpConvertFToS**

Convert (value preserving) from floating point to signed integer, with round toward 0.0.

Result Type must be a scalar or vector of integer type.

Float Value must be a scalar or vector of floating-point type. It must have the same number of components as Result Type.

Results are computed per component.

4	110	1.	D 10 .2.16	1.
4	110	<1d>	Result <10>	< <i>id&gt;</i>
		Result Type		Float Value

#### **OpConvertSToF**

Convert (value preserving) from signed integer to floating point.

Result Type must be a scalar or vector of floating-point type.

Signed Value must be a scalar or vector of integer type. It must have the same number of components as Result Type.

	The Property of the Property o				
4	111	< <i>id</i> >	Result <id></id>	< <i>id</i> >	
		Result Type		Signed Value	

# **OpConvertUToF**

Convert (value preserving) from unsigned integer to floating point.

Result Type must be a scalar or vector of floating-point type.

*Unsigned Value* must be a scalar or vector of integer type. It must have the same number of components as *Result Type*.

Results are computed per component.

ĺ	4	112	<id></id>	Result <id></id>	<id></id>
			Result Type		Unsigned Value

#### **OpUConvert**

Convert (value preserving) unsigned width. This is either a truncate or a zero extend.

Result Type must be a scalar of integer type, whose Signedness operand is 0.

*Unsigned Value* must be a scalar or vector of integer type. It must have the same number of components as *Result Type*. The component width cannot equal the component width in *Result Type*.

Results are computed per component.

	1 1 1				
4	113	< <i>id</i> >	Result <id></id>	< <i>id</i> >	
		Result Type		Unsigned Value	

#### **OpSConvert**

Convert (value preserving) signed width. This is either a truncate or a sign extend.

Result Type must be a scalar or vector of integer type.

Signed Value must be a scalar or vector of integer type. It must have the same number of components as Result Type. The component width cannot equal the component width in Result Type.

Results are computed per component.

4	114	< <i>id&gt;</i>	Result <id></id>	< <i>id</i> >	
		Result Type		Signed Value	

#### **OpFConvert**

Convert (value preserving) floating-point width.

Result Type must be a scalar or vector of floating-point type.

Float Value must be a scalar or vector of floating-point type. It must have the same number of components as Result Type. The component width cannot equal the component width in Result Type.

4	115	< <i>id&gt;</i>	Result <id></id>	< <i>id</i> >	
		Result Type		Float Value	

OpQuantizeToF16			Capability:
Quantize a floating-poi	Shader		
Result Type must be a smust be 32 bits.			
Value is the value to qualify If Value is an infinity, t			
NaN, but not necessari large to represent as a	ly the same NaN. If <i>Value</i> is position 16-bit floating-point value, the real magnitude too large to represent	itive with a magnitude too sult is positive infinity. If	
value, the result is nega	a magnitude too large to represent ative infinity. If the magnitude of floating-point value, the result m	Value is too small to represent	
The RelaxedPrecision			
Results are computed p			
4 116	<id>&gt;</id>	Result <id></id>	<id></id>
	Result Type		Value

Convert a pointer to an unsigned integer type. A <i>Result Type</i> width larger than the width of <i>Pointer</i> will zero extend. A <i>Result Type</i> smaller than the width of <i>Pointer</i> will truncate. For same-width source and result, this is the same as OpBitcast.  *Result Type* must be a scalar or vector of integer type, whose *Signedness* operand is 0.				Capability: Addresses
4	117	< <i>id</i> >	Result <id></id>	< <i>id&gt;</i>
		Result Type		Pointer

<b>OpSatConvertST</b>	OpSatConvertSToU			
	nteger to unsigned integer. Conve			
representable rang	e of <i>Result Type</i> are clamped to t	he nearest representable value of		
Result Type.				
D 1.55				
Result Type must b	e a scalar or vector of integer type	oe.		
C' IVI	1 1	To the second to		
		pe. It must have the same number		
of components as a	of components as Result Type.			
Results are compu				
4   118	< <i>id</i> >	Result <id></id>	< <i>id</i> >	
	Result Type		Signed Value	

OpSatConvertUToS			Capability:	
Convert an unsigned in	Convert an unsigned integer to signed integer. Converted values outside the			
representable range of <i>Result Type</i> are clamped to the nearest representable value of <i>Result Type</i> .				
Result Type must be a s	Result Type must be a scalar or vector of integer type.			
Unsigned Value must b number of components				
Results are computed p				
4 119	<id>&gt;</id>	Result <id></id>	<id></id>	
	Result Type		Unsigned Value	

OpConvertUToPtr  Convert an integer to pointer. A <i>Result Type</i> width smaller than the width of <i>Integer Value</i> pointer will truncate. A <i>Result Type</i> width larger than the width of <i>Integer Value</i> pointer will zero extend.  Result Type must be an OpTypePointer. For same-width source and result, this is the same as OpBitcast.  4 120 < id> Result < id < In				ihe
4	120	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Integer Value

<b>OpPtrCast</b>	ToGeneric			Capability:
Convert a po	ointer's Stor	Kernel		
Result Type	must be an	OpTypePointer. Its Storage Cla	ss must be Generic.	
Pointer must	t point to the	e Workgroup, CrossWorkgro	up, or Function Storage Class.	
Result Type	and <i>Pointer</i>			
4	121	< <i>id</i> >		
		Result Type		Pointer

OpGenerio	CastToPtr			Capability:
Convert a p	oointer's Stor	Kernel		
1	must be an <b>kgroup</b> , or <b>F</b>			
Pointer mu	st point to the	e Generic Storage Class.		
Result Type	and <i>Pointer</i>			
4	122	<id></id>		
		Result Type		Pointer

OpGene	ricCastTo	PtrExplicit			Capability:			
Attempts	s to explici	tly convert Pointer to Sto	rage storage-class pointe	r value.				
Result Ty	<i>pe</i> must b	e an OpTypePointer. Its S	Storage Class must be Sto	orage.				
Pointer r	Pointer must have a type of OpTypePointer whose Type is the same as the Type of Result							
			* *	* *				
		point to the Generic Stor	•	s, the instruction result				
is an Opt	ConstantN	ull pointer in the Storage	Storage Class.					
Storage 1	must be on	e of the following literal	values from Storage Class	ss: Workgroup,				
CrossWo	CrossWorkgroup, or Function.							
5	~ ~							
		Storage						

# **OpBitcast**

Bit pattern-preserving type conversion.

Result Type must be an OpTypePointer, or a scalar or vector of numerical-type.

*Operand* must have a type of OpTypePointer, or a scalar or vector of numerical-type. It must be a different type than *Result Type*.

If *Result Type* is a pointer, *Operand* must be a pointer or integer scalar. If *Operand* is a pointer, *Result Type* must be a pointer or integer scalar.

If *Result Type* has the same number of components as *Operand*, they must also have the same component width, and results are computed per component.

If *Result Type* has a different number of components than *Operand*, the total number of bits in *Result Type* must equal the total number of bits in *Operand*. Let L be the type, either *Result Type* or *Operand's* type, that has the larger number of components. Let S be the other type, with the smaller number of components. The number of components in L must be an integer multiple of the number of components in S. The first component (that is, the only or lowest-numbered component) of S maps to the first components of L, and so on, up to the last component of S mapping to the last components of L. Within this mapping, any single component of S (mapping to multiple components of L) maps its lower-ordered bits to the lower-numbered components of L.

4	124	< <i>id</i> >	Result <id></id>	<id></id>
		Result Type		Operand

#### 3.32.12 Composite Instructions

#### **OpVectorExtractDynamic**

Extract a single, dynamically selected, component of a vector.

Result Type must be a scalar type.

*Vector* must have a type OpTypeVector whose *Component Type* is *Result Type*.

*Index* must be a scalar integer 0-based index of which component of *Vector* to extract.

The value read is undefined if *Index's* value is less than zero or greater than or equal to the number of components in *Vector*.

5	77	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Vector	Index

#### **OpVectorInsertDynamic**

Make a copy of a vector, with a single, variably selected, component modified.

Result Type must be an OpTypeVector.

*Vector* must have the same type as *Result Type* and is the vector that the non-written components will be copied from.

*Component* is the value that will be supplied for the component selected by *Index*. It must have the same type as the type of components in *Result Type*.

*Index* must be a scalar integer 0-based index of which component to modify.

What is written is undefined if *Index's* value is less than zero or greater than or equal to the number of components in *Vector*.

6	78	<id></id>	Result <id></id>	<id></id>	< <i>id</i> >	<id></id>
		Result Type		Vector	Component	Index

# **OpVectorShuffle**

Select arbitrary components from two vectors to make a new vector.

*Result Type* must be an OpTypeVector. The number of components in *Result Type* must be the same as the number of *Component* operands.

Vector 1 and Vector 2 must both have vector types, with the same Component Type as Result Type. They do not have to have the same number of components as Result Type or with each other. They are logically concatenated, forming a single vector with Vector 1's components appearing before Vector 2's. The components of this logical vector are logically numbered with a single consecutive set of numbers from 0 to N - 1, where N is the total number of components.

Components are these logical numbers (see above), selecting which of the logically numbered components form the result. They can select the components in any order and can repeat components. The first component of the result is selected by the first Component operand, the second component of the result is selected by the second Component operand, etc. A Component literal may also be FFFFFFFF, which means the corresponding result component has no source and is undefined. All Component literals must either be FFFFFFFF or in [0, N - 1] (inclusive).

**Note:** A vector "swizzle" can be done by using the vector for both *Vector* operands, or using an OpUndef for one of the *Vector* operands.

1						
5 + variable	79	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Literal, Literal,
		Result Type		Vector 1	Vector 2	
						Components

#### **OpCompositeConstruct**

Construct a new composite object from a set of constituent objects that will fully form it.

Result Type must be a composite type, whose top-level members/elements/components/columns have the same type as the types of the operands, with one exception. The exception is that for constructing a vector, the operands may also be vectors with the same component type as the Result Type component type. When constructing a vector, the total number of components in all the operands must equal the number of components in Result Type.

Constituents will become members of a structure, or elements of an array, or components of a vector, or columns of a matrix. There must be exactly one Constituent for each top-level member/element/component/column of the result, with one exception. The exception is that for constructing a vector, a contiguous subset of the scalars consumed can be represented by a vector operand instead. The Constituents must appear in the order needed by the definition of the type of the result. When constructing a vector, there must be at least two Constituent operands.

3 + variable	80	< <i>id</i> >	Result <id></id>	<id>, <id>,</id></id>
		Result Type		Constituents

# **OpCompositeExtract**

Extract a part of a composite object.

*Result Type* must be the type of object selected by the last provided index. The instruction result is the extracted object.

Composite is the composite to extract from.

*Indexes* walk the type hierarchy, potentially down to component granularity, to select the part to extract. All indexes must be in bounds. All composite constituents use zero-based numbering, as described by their **OpType...** instruction.

4 + variable	81	<id>&gt;</id>	Result <id></id>	<id></id>	Literal, Literal,
		Result Type		Composite	Indexes

# **OpCompositeInsert**

Make a copy of a composite object, while modifying one part of it.

Result Type must be the same type as Composite.

Object is the object to use as the modified part.

*Composite* is the composite to copy all but the modified part from.

*Indexes* walk the type hierarchy of *Composite* to the desired depth, potentially down to component granularity, to select the part to modify. All indexes must be in bounds. All composite constituents use zero-based numbering, as described by their **OpType...** instruction. The type of the part selected to modify must match the type of *Object*.

5 + variable	82	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Literal, Literal,	
		Result Type		Object	Composite		
						Indexes	

OpCopyObject								
Make a copy of <i>Operand</i> . There are no dereferences involved.								
Result	<i>Type</i> m	ust match <i>Operan</i>	d type. There are	no other				
restrictions on the types.								
4	83	< <i>id</i> >	Result <id></id>	< <i>id</i> >				
		Result Type		Operand				

OpTrans	pose			Capability:
Transpose	a matrix.	Matrix		
Result Typ	e must be an	OpTypeMatrix.		
column si scalar cor	ze of <i>Matrix</i> in <i>M</i>	must be the reverse of the fatrix and Result Type m	s. The number of columns and the ose in <i>Result Type</i> . The types of th ust be the same.	
4	ust have of typ	<id>&gt;</id>		
		<id><id> Result Type</id></id>	Result <id></id>	Matrix

#### 3.32.13 Arithmetic Instructions

# **OpSNegate**

Signed-integer subtract of *Operand* from zero.

Result Type must be a scalar or vector of integer type.

*Operand's* type must be a scalar or vector of integer type. It must have the same number of components as *Result Type*. The component width must equal the component width in *Result Type*.

Results are computed per component.

İ	4	126	<id></id>	Result <id></id>	<id></id>
			Result Type		Operand

#### **OpFNegate**

Floating-point subtract of Operand from zero.

Result Type must be a scalar or vector of floating-point type.

The type of *Operand* must be the same as *Result Type*.

Results are computed per component.

		1 1		
4	127	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Operand

#### OpIAdd

Integer addition of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component.

	1	1 1			
5	128	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### OpFAdd

Floating-point addition of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

_	100	. 7	D 1: 1:1	. 7	. 7
5	129	<1d>	Result <1d>	<1d>	<1d>
		Result Type		Operand 1	Operand 2

# OpISub

Integer subtraction of Operand 2 from Operand 1.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component.

5	130	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpFSub**

Floating-point subtraction of *Operand 2* from *Operand 1*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component.

5	131	< <i>id</i> >	Result <id></id>	<id></id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

# OpIMul

Integer multiplication of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component.

5	132	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

#### OpFMul

Floating-point multiplication of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Tto barto	tesures are compared per component.						
5	133	< <i>id</i> >	Result <id></id>	<id></id>	<id></id>		
		Result Type		Operand 1	Operand 2		

# OpUDiv

Unsigned-integer division of *Operand 1* divided by *Operand 2*.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	134	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpSDiv**

Signed-integer division of *Operand 1* divided by *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	135	<id></id>	Result <id></id>	<id>&gt;</id>	<id>&gt;</id>	
		Result Type		Operand 1	Operand 2	

# **OpFDiv**

Floating-point division of *Operand 1* divided by *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	136	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### **OpUMod**

Unsigned modulo operation of *Operand 1* modulo *Operand 2*.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	137	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	<id></id>
		Result Type		Operand 1	Operand 2

# OpSRem

Signed remainder operation of *Operand 1* divided by *Operand 2*. The sign of a non-0 result comes from *Operand 1*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	138	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

#### OpSMod

Signed modulo operation of *Operand 1* modulo *Operand 2*. The sign of a non-0 result comes from *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

	-			-	
5	139	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

# OpFRem

Floating-point remainder operation of *Operand 1* divided by *Operand 2*. The sign of a non-0 result comes from *Operand 1*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	140	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

# OpFMod

Floating-point remainder operation of *Operand 1* divided by *Operand 2*. The sign of a non-0 result comes from *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

	I			- I	
5	141	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

Scalar

# OpVectorTimesScalar Scale a floating-point vector. Result Type must be a vector of floating-point type. The type of Vector must be the same as Result Type. Each component of Vector is multiplied by Scalar. Scalar must have the same type as the Component Type in Result Type. 5 | 142 | <id> | Result <id> | <id | <id> | <id | <id> |

Vector

Result Type

OpMatr	rixTimesS	calar			Capability:
	Scale a floating-point matrix.  Result Type must be an OpTypeMatrix whose Column Type is a vector of floating-point type.				
Matrix is	The type of <i>Matrix</i> must be the same as <i>Result Type</i> . Each component in each column in <i>Matrix</i> is multiplied by <i>Scalar</i> .  Scalar must have the same type as the <i>Component Type</i> in <i>Result Type</i> .				
5	143	<id><id>&lt;</id></id>	Result <id></id>	<id>&gt;</id>	<id>&gt;</id>
		Result Type		Matrix	Scalar

OpVecto	<b>OpVectorTimesMatrix</b>				Capability:
Linear-al	Linear-algebraic <i>Vector X Matrix</i> .				
Result Ty	Result Type must be a vector of floating-point type.				
Vector must be a vector with the same Component Type as the Component Type in Result Type. Its number of components must equal the number of components in each column in Matrix.					
	Matrix must be a matrix with the same Component Type as the Component Type in Result Type. Its number of columns must equal the number of components in Result Type.				
5	144	<id>&gt;</id>	Result <id></id>	< <i>id&gt;</i>	<id>&gt;</id>
		Result Type		Vector	Matrix

OpMatr	rixTimesV	ector			Capability:
Linear-al	Linear-algebraic <i>Vector X Matrix</i> .  Result Type must be a vector of floating-point type.  Matrix must be an OpTypeMatrix whose Column Type is Result Type.				Matrix
Vector m	ust be a v	ector with the same Comp	ponent Type as the Comp	onent Type in Result	
<i>Type</i> . Its	<i>Type</i> . Its number of components must equal the number of columns in <i>Matrix</i> .				
5	145	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Matrix	Vector

OpMatr	rixTimesN	<b>Iatrix</b>			Capability:
Linear-a	Linear-algebraic multiply of LeftMatrix X RightMatrix.				
Result Type must be an OpTypeMatrix whose Column Type is a vector of floating-point type.					
LeftMatrix must be a matrix whose Column Type is the same as the Column Type in Result Type.					
RightMatrix must be a matrix with the same Component Type as the Component Type in Result Type. Its number of columns must equal the number of columns in Result Type. Its columns must have the same number of components as the number of columns in LeftMatrix.					
5	146	<id>&gt;</id>	Result <id></id>	<id></id>	<id></id>
		Result Type		LeftMatrix	RightMatrix

OpOu	terProduc	t			Capability:
					Matrix
Linear-	Linear-algebraic outer product of <i>Vector 1</i> and <i>Vector 2</i> .				
Result	Type must				
Vector	Vector 1 must have the same type as the Column Type in Result Type.				
Vector	2 must be	a vector with the same	Component Type as the	Component Type in Result	
Type. I	ts number				
5	147	<id></id>	Result <id></id>	< <i>id</i> >	<id>&gt;</id>
		Result Type		Vector 1	Vector 2

OpDo	OpDot						
Dot product of <i>Vector 1</i> and <i>Vector 2</i> .							
Result	Result Type must be a floating-point type scalar.						
	Vector 1 and Vector 2 must be vectors of the same type, and their component type must be Result Type.						
5	148	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >		
		Result Type		Vector 1	Vector 2		

# **OpIAddCarry**

Result is the unsigned integer addition of *Operand 1* and *Operand 2*, including its carry.

*Result Type* must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type, whose *Signedness* operand is 0.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as unsigned integers.

Results are computed per component.

Member 0 of the result gets the low-order bits (full component width) of the addition.

Member 1 of the result gets the high-order (carry) bit of the result of the addition. That is, it gets the value 1 if the addition overflowed the component width, and 0 otherwise.

5	149	<id>&gt;</id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### **OpISubBorrow**

Result is the unsigned integer subtraction of *Operand 2* from *Operand 1*, and what it needed to borrow.

*Result Type* must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type, whose *Signedness* operand is 0.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as unsigned integers.

Results are computed per component.

Member 0 of the result gets the low-order bits (full component width) of the subtraction. That is, if *Operand 1* is larger than *Operand 2*, member 0 gets the full value of the subtraction; if *Operand 2* is larger than *Operand 1*, member 0 gets  $2^w + Operand 1 - Operand 2$ , where w is the component width.

Member 1 of the result gets 0 if *Operand* 1 > Operand 2, and gets 1 otherwise.

5	150	<id>&gt;</id>	Result <id></id>	<id></id>	<id>&gt;</id>	
		Result Type		Operand 1	Operand 2	

# OpUMulExtended

Result is the full value of the unsigned integer multiplication of *Operand 1* and *Operand 2*.

*Result Type* must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type, whose *Signedness* operand is 0.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as unsigned integers.

Results are computed per component.

Member 0 of the result gets the low-order bits of the multiplication.

Member 1 of the result gets the high-order bits of the multiplication.

Ī	5	151	<id>&gt;</id>	Result <id></id>	< <i>id</i> >	<id></id>
			Result Type		Operand 1	Operand 2

#### **OpSMulExtended**

Result is the full value of the signed integer multiplication of *Operand 1* and *Operand 2*.

*Result Type* must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as signed integers.

Results are computed per component.

Member 0 of the result gets the low-order bits of the multiplication.

Member 1 of the result gets the high-order bits of the multiplication.

5	152	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### 3.32.14 Bit Instructions

#### **OpShiftRightLogical**

Shift the bits in Base right by the number of bits specified in Shift. The most-significant bits will be zero filled.

Result Type must be a scalar or vector of integer type.

The type of each *Base* and *Shift* must be a scalar or vector of integer type. *Base* and *Shift* must have the same number of components. The number of components and bit width of the type of *Base* must be the same as in *Result Type*.

*Shift* is consumed as an unsigned integer. The result is undefined if *Shift* is greater than the bit width of the components of *Base*.

Results are computed per component.

5	194	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	< <i>id</i> >
		Result Type		Base	Shift

#### **OpShiftRightArithmetic**

Shift the bits in *Base* right by the number of bits specified in *Shift*. The most-significant bits will be filled with the sign bit from *Base*.

Result Type must be a scalar or vector of integer type.

The type of each *Base* and *Shift* must be a scalar or vector of integer type. *Base* and *Shift* must have the same number of components. The number of components and bit width of the type of *Base* must be the same as in *Result Type*.

Shift is treated as unsigned. The result is undefined if Shift is greater than the bit width of the components of Base.

Results are computed per component.

5	195	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Base	Shift	

#### **OpShiftLeftLogical**

Shift the bits in Base left by the number of bits specified in Shift. The least-significant bits will be zero filled.

Result Type must be a scalar or vector of integer type.

The type of each *Base* and *Shift* must be a scalar or vector of integer type. *Base* and *Shift* must have the same number of components. The number of components and bit width of the type of *Base* must be the same as in *Result Type*.

Shift is treated as unsigned. The result is undefined if Shift is greater than the bit width of the components of Base.

The number of components and bit width of *Result Type* must match those *Base* type. All types must be integer types.

5	196	<id>&gt;</id>	Result <id></id>	<id></id>	<id>&gt;</id>
		Result Type		Base	Shift

# **OpBitwiseOr**

Result is 1 if either *Operand 1* or *Operand 2* is 1. Result is 0 if both *Operand 1* and *Operand 2* are 0.

Results are computed per component, and within each component, per bit.

*Result Type* must be a scalar or vector of integer type. The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

5	197	<id></id>	Result <id></id>	<id>&gt;</id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpBitwiseXor**

Result is 1 if exactly one of *Operand 1* or *Operand 2* is 1. Result is 0 if *Operand 1* and *Operand 2* have the same value.

Results are computed per component, and within each component, per bit.

*Result Type* must be a scalar or vector of integer type. The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

5	198	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

# **OpBitwiseAnd**

Result is 1 if both Operand 1 and Operand 2 are 1. Result is 0 if either Operand 1 or Operand 2 are 0.

Results are computed per component, and within each component, per bit.

Result Type must be a scalar or vector of integer type. The type of Operand 1 and Operand 2 must be a scalar or vector of integer type. They must have the same number of components as Result Type. They must have the same component width as Result Type.

		· ·			
5	199	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### **OpNot**

Complement the bits of *Operand*.

Results are computed per component, and within each component, per bit.

Result Type must be a scalar or vector of integer type.

*Operand's* type must be a scalar or vector of integer type. It must have the same number of components as *Result Type*. The component width must equal the component width in *Result Type*.

Ī	4	200	< <i>id</i> >	Result <id></id>	<id></id>
			Result Type		Operand

OpBit	FieldIn	sert				Capability:		
Make a	а сору (	of an object, with	Shader					
Results	s are co	mputed per comp						
Result	Type m	ust be a scalar or						
The ty	pe of B	ase and Insert mu	st be the same as	Result Type.				
		s numbered outsicesponding bits in h	le [ <i>Offset, Offset -</i> Base.	+ <i>Count</i> - 1] (inclu	usive) will come			
		s numbered in [ <i>O</i> <sub>3</sub> ] [0, <i>Count</i> - 1] of	ffset, Offset + Cou Insert.	<i>unt</i> - 1] come, in o	order, from the			
Insert.	It will		calar. <i>Count</i> is the n unsigned value.					
	Offset must be an integer type scalar. Offset is the lowest-order bit of the bit field.  It will be consumed as an unsigned value.							
		value is undefined	greater than the					
		s in the result.						
7	201	<id>&gt;</id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id>&gt;</id>	<id>&gt;</id>	
		Result Type		Base	Insert	Offset	Count	

OpBitl	FieldSEx	tract			Capability:	
Extract	a bit fiel	d from an object, wit	Shader			
Results	are com	puted per component				
Result	<i>Type</i> mus	t be a scalar or vecto	r of integer type.			
The typ	pe of Bas	e must be the same a	s Result Type.			
Count -	- 1] ( <mark>incl</mark> u naining b	ter than 0: The bits of asive) become the bit of the result will a	s numbered [0, Cour	nt - 1] of the result.		
from B	ase. It wi	in integer type scalar. Ill be consumed as an esult will be 0.				
Offset must be an integer type scalar. Offset is the lowest-order bit of the bit field to extract from Base. It will be consumed as an unsigned value.						
		lue is undefined if <i>Co</i> of bits in the result.				
6	202	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	<id>&gt;</id>	<id>&gt;</id>
		Result Type		Base	Offset	Count

OpBitl	OpBitFieldUExtract					
The ser	mantics a	d from an object, with the the same as with the control of the con	OpBitFieldSExtract v	with the exception	Shader	
6	203	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >		
		Result Type		Base	Offset	Count

OpBitRevo	erse	Capability: Shader		
Reverse the	e bits in an ol	Silauci		
Results are	computed po	er component.		
Result Type	must be a so	calar or vector of integer type.		
The type of	f <i>Base</i> must b	be the same as Result Type.		
	mber $n$ of the $th$ is the $OpT$			
4	204	< <i>id</i> >	Result <id></id>	<id>&gt;</id>
		Result Type		Base

# **OpBitCount**

Count the number of set bits in an object.

Results are computed per component.

*Result Type* must be a scalar or vector of integer type. The components must be wide enough to hold the unsigned *Width* of *Base* as an unsigned value. That is, no sign bit is needed or counted when checking for a wide enough result width.

Base must be a scalar or vector of integer type. It must have the same number of components as Result Type.

The result is the unsigned value that is the number of bits in *Base* that are 1.

8						
	4	205	< <i>id</i> >	Result <id></id>	< <i>id</i> >	
			Result Type		Base	

#### 3.32.15 Relational and Logical Instructions

# Result is **true** if any component of *Vector* is **true**, otherwise result is **false**. Result Type must be a Boolean type scalar. Vector must be a vector of Boolean type. 4 | 154 | <id> | Result <id> | <id> | <id> |

Vector

Result Type

OpAll					
Result is <b>true</b> if all components of <i>Vector</i> are <b>true</b> , otherwise result is <b>false</b> .					
Result	Type mus	t be a Boolean type	e scalar.		
Result Type indict be a Boolean type seatal.					
Vector must be a vector of Boolean type.					
4	155	< <i>id</i> >	Result <id></id>	< <i>id</i> >	
		Result Type		Vector	

#### **OpIsNan**

Result is **true** if x is an IEEE NaN, otherwise result is **false**.

Result Type must be a scalar or vector of Boolean type.

x must be a scalar or vector of floating-point type. It must have the same number of components as Result Type.

Results are computed per component.

Tresums and	Treating and companies per components					
4	156	< <i>id</i> >	Result <id></id>	< <i>id</i> >		
		Result Type		x		

#### **OpIsInf**

Result is **true** if x is an IEEE Inf, otherwise result is **false** 

Result Type must be a scalar or vector of Boolean type.

x must be a scalar or vector of floating-point type. It must have the same number of components as Result Type.

4	157	<id>&gt;</id>	Result <id></id>	<id></id>
		Result Type		X

<b>OpIsFinite</b>	OpIsFinite			Capability:
Result is tru  Result Type  x must be a components	must be a so scalar or vec as as Result Ty	TEEE finite number, otherwise calar or vector of Boolean type ctor of floating-point type. It maye.		Kernel
4	158	<id>&gt;</id>		
		Result Type		x

OpIsNorm	nal			Capability:
Result is <b>true</b> if <i>x</i> is an IEEE normal number, otherwise result is <b>false</b> .				Kernel
Result Type	e must be a so	calar or vector of Boolean type.		
	a scalar or vects as Result T			
Results are computed per component.				
4	159	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>
		Result Type		x

<b>OpSignBit</b>	Set	Capability:		
Result is <b>true</b> if <i>x</i> has its sign bit set, otherwise result is <b>false</b> .  Result Type must be a scalar or vector of Boolean type.				Kernel
component	a scalar or vectors as <i>Result T</i>			
4	160	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>
		Result Type		x

OpLess	OpLessOrGreater					
Result is  Result Ty  x must b compone	true if x <  wpe must b e a scalar o ents as Res	< y or $x > y$ , where IEEE of a scalar or vector of Booton vector of floating-point ult Type.	olean type.		Capability: Kernel	
	y must have the same type as x.  Results are computed per component.					
5	161	<id>&gt;</id>	Result <id></id>	< <i>id</i> >	<id>&gt;</id>	
		Result Type		x	у	

OpOrde	OpOrdered					
Result is result is result Ty	Capability: Kernel					
compone y must ha						
Results a						
5	162	<id>&gt;</id>	Result <id></id>	<id></id>	<id></id>	
		Result Type		x	у	

OpUno	Capability: Kernel						
Result i	ixtiliti						
Result 7							
	be a scalar nents as <i>Re</i>		point type. It must have t	he same number of			
y must							
Results							
5	5   163   <id>   Result <id>   <id>  </id></id></id>						
		Result Type		X	У		

# OpLogicalEqual

Result is **true** if *Operand 1* and *Operand 2* have the same value. Result is **false** if *Operand 1* and *Operand 2* have different values.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

#### Results are computed per component.

5	164	<id>&gt;</id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

# OpLogicalNotEqual

Result is **true** if *Operand 1* and *Operand 2* have different values. Result is **false** if *Operand 1* and *Operand 2* have the same value.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

Results are computed per component.

	5	165	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
			Result Type		Operand 1	Operand 2	

# OpLogicalOr

Result is **true** if either *Operand 1* or *Operand 2* is **true**. Result is **false** if both *Operand 1* and *Operand 2* are **false**.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

The state of the s							
5	166	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >		
		Result Type		Operand 1	Operand 2		

### OpLogicalAnd

Result is **true** if both *Operand 1* and *Operand 2* are **true**. Result is **false** if either *Operand 1* or *Operand 2* are **false**.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

Results are computed per component.

5	167	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### OpLogicalNot

Result is **true** if *Operand* is **false**. Result is **false** if *Operand* is **true**.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand* must be the same as *Result Type*.

Results are computed per component.

4	168	< <i>id</i> >	Result <id></id>	<id>&gt;</id>
		Result Type		Operand

### OpSelect

Select components from two objects.

Result Type must be a pointer, scalar, or vector.

The type of *Object 1* must be the same as *Result Type*. *Object 1* is selected as the result if *Condition* is **true**.

The type of Object 2 must be the same as Result Type. Object 2 is selected as the result if Condition is false.

Condition must be a scalar or vector of Boolean type. It must have the same number of components as Result Type.

6	169	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id></id>	<id></id>
		Result Type		Condition	Object 1	Object 2

### **OpIEqual**

Integer comparison for equality.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	170	<id></id>	Result <id></id>	<id>&gt;</id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpINotEqual**

Integer comparison for inequality.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

						i i
5	171	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	i
		Result Type		Operand 1	Operand 2	i

### **OpUGreaterThan**

Unsigned-integer comparison if *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	172	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### **OpSGreaterThan**

Signed-integer comparison if *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

	The state of the s					
5	173	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id&gt;</i>	
		Result Type		Operand 1	Operand 2	

### **OpUGreaterThanEqual**

Unsigned-integer comparison if *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	174	<id></id>	Result <id></id>	<id>&gt;</id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpSGreaterThanEqual**

Signed-integer comparison if *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	175	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### OpULessThan

Unsigned-integer comparison if *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	176	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	<id>&gt;</id>
		Result Type		Operand 1	Operand 2

#### **OpSLessThan**

Signed-integer comparison if *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

	The state of the s					
5	177	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### **OpULessThanEqual**

Unsigned-integer comparison if *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	178	<id></id>	Result <id></id>	<id>&gt;</id>	<id>&gt;</id>
		Result Type		Operand 1	Operand 2

#### **OpSLessThanEqual**

Signed-integer comparison if *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

					i	
5	179	< <i>id</i> >	Result <id></id>	<id></id>	< <i>id</i> >	i
		Result Type		Operand 1	Operand 2	i

### **OpFOrdEqual**

Floating-point comparison for being ordered and equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	180	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

#### **OpFUnordEqual**

Floating-point comparison for being unordered or equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

	The state of the s					
5	181	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### **OpFOrdNotEqual**

Floating-point comparison for being ordered and not equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	182	<id>&gt;</id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpFUnordNotEqual**

Floating-point comparison for being unordered or not equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

					i i	
5	183	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	i
		Result Type		Operand 1	Operand 2	i

### **OpFOrdLessThan**

Floating-point comparison if operands are ordered and *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	184	<id>&gt;</id>	Result <id></id>	< <i>id</i> >	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpFUnordLessThan**

Floating-point comparison if operands are unordered or *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

5	185	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### **OpFOrdGreaterThan**

Floating-point comparison if operands are ordered and *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	186	<id></id>	Result <id></id>	<id>&gt;</id>	<id></id>
		Result Type		Operand 1	Operand 2

#### **OpFUnordGreaterThan**

Floating-point comparison if operands are unordered or *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

1 1 1							
	5	187	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
			Result Type		Operand 1	Operand 2	

#### **OpFOrdLessThanEqual**

Floating-point comparison if operands are ordered and *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	188	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	<id>&gt;</id>
		Result Type		Operand 1	Operand 2

#### **OpFUnordLessThanEqual**

Floating-point comparison if operands are unordered or *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

	The state of the s					
5	189	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### **OpFOrdGreaterThanEqual**

Floating-point comparison if operands are ordered and *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

### Results are computed per component.

5	190	<id></id>	Result <id></id>	<id>&gt;</id>	<id></id>
		Result Type		Operand 1	Operand 2

### **OpFUnordGreaterThanEqual**

Floating-point comparison if operands are unordered or *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

5	191	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

### 3.32.16 Derivative Instructions

OpDPdx				Capability:
	t as either Op sternal factor	Shader		
Result Type	must be a so	type.		
The type of	f P must be the			
This instruc	ction is only			
4	207	<id>&gt;</id>	Result <id></id>	<id></id>
		Result Type		P

OpDPdy		Capability:		
	t as either <mark>O</mark> p sternal factor	Shader		
Result Type	must be a so	calar or vector of floating-point	type.	
The type of	P must be tl			
This instruc	ction is only			
4	208	<id></id>	Result <id></id>	<id>&gt;</id>
		Result Type		P

OpFwidth				Capability:
Result is the OpDPdy on		Shader		
Result Type 1	must be a so	calar or vector of floating-point	type.	
The type of I	P must be the			
This instruct	ion is only			
4	209	<id>&gt;</id>	Result <id></id>	<id></id>
		Result Type		P

OpDPdxFi	ine	Capability:		
	ne partial deri encing based ).	DerivativeControl		
Result Type	must be a so	calar or vector of floating-point	type.	
The type of	f <i>P</i> must be tl			
This instruc	ction is only			
4	210	<id>&gt;</id>	Result <id></id>	<id></id>
		Result Type		P

OpDPdyFi	ine	Capability:		
	e partial deri encing based	DerivativeControl		
Result Type	must be a so	calar or vector of floating-point	type.	
The type of	P must be the			
This instruc	ction is only			
4	211	<id></id>	Result <id></id>	<id></id>
		Result Type		P

OpFwidth	Fine			Capability:
Result is th	ne same as co	DerivativeControl		
Result Type	must be a se	calar or vector of floating-point	type.	
The type of	f P must be the			
This instruc	ction is only			
4	212	< <i>id&gt;</i>	Result <id></id>	<id>&gt;</id>
		Result Type		P

OpDPdxCoarse	OpDPdxCoarse					
	DerivativeControl					
Result is the partial de	rivative of P with respect to the v	vindow x coordinate. Will use				
local differencing base	ed on the value of P for the currer	nt fragment's neighbors, and				
will possibly, but not i	necessarily, include the value of <i>P</i>	for the current fragment.				
That is, over a given a	rea, the implementation can comp	oute x derivatives in fewer				
unique locations than	would be allowed for OpDPdxFir	ne.				
	_					
Result Type must be a	scalar or vector of floating-point	type.				
The type of <i>P</i> must be	the same as Result Type. P is the	value to take the derivative of.				
This instruction is only						
4 213	< <i>id&gt;</i>	Result <id></id>	<id>&gt;</id>			
	Result Type		P			

OpDPdyCoa	arse	Capability:		
				DerivativeControl
Result is the	partial deri	vative of <i>P</i> with respect to the w	vindow y coordinate. Will use	
local differen	cing based	on the value of $P$ for the current	t fragment's neighbors, and	
will possibly,	, but not ne	cessarily, include the value of P	for the current fragment.	
That is, over	a given are	a, the implementation can comp	oute y derivatives in fewer	
unique location	ons than we	ould be allowed for OpDPdyFir	e.	
Result Type n	nust be a sc	ealar or vector of floating-point	type.	
The type of P	must be the	ne same as <i>Result Type</i> . <i>P</i> is the	value to take the derivative of.	
This instructi	on is only			
4 2	214	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		P

OpFwidth	Coarse			Capability:
	ne same as co	DerivativeControl		
Result Type	e must be a so	calar or vector of floating-point	type.	
The type of	f P must be the			
This instru	ction is only			
4	215	<id></id>	Result <id></id>	<id></id>
		Result Type		P

#### 3.32.17 Control-Flow Instructions

### **OpPhi**

The SSA phi function.

The result is selected based on control flow: If control reached the current block from *Parent i*, *Result Id* gets the value that *Variable i* had at the end of *Parent i*.

Result Type can be any type.

Operands are a sequence of pairs: (*Variable 1, Parent 1* block), (*Variable 2, Parent 2* block), ... Each *Parent i* block is the label of an immediate predecessor in the CFG of the current block. There must be exactly one *Parent i* for each parent block of the current block in the CFG. All *Variables* must have a type matching *Result Type*.

Within a block, this instruction must appear before all non-**OpPhi** instructions (except for **OpLine**, which can be mixed with **OpPhi**).

3 + variable	245	< <i>id</i> >	Result <id></id>	<id>, <id>,</id></id>
		Result Type		Variable, Parent,

#### **OpLoopMerge**

Declare a structured loop.

This instruction must immediately precede either an OpBranch or OpBranchConditional instruction. That is, it must be the second-to-last instruction in its block.

Merge Block is the label of the merge block for this structured loop.

Continue Target is the label of a block targeted for processing a loop "continue".

Loop Control Parameters appear in Loop Control-table order for any Loop Control setting that requires such a parameter.

See Structured Control Flow for more detail.

4 + variable	246	< <i>id</i> >	< <i>id</i> >	Loop Control	Literal, Literal,
		Merge Block	Continue Target		Loop Control
					Parameters

#### **OpSelectionMerge**

Declare a structured selection.

This instruction must immediately precede either an OpBranchConditional or OpSwitch instruction. That is, it must be the second-to-last instruction in its block.

Merge Block is the label of the merge block for this structured selection.

See Structured Control Flow for more detail.

3	247	<id></id>	Selection Control
		Merge Block	

### OpLabel

The block label instruction: Any reference to a block is through the *Result* < *id*> of its label.

Must be the first instruction of any block, and appears only as the first instruction of a block.

2	248	Result <id></id>

### OpBranch

Unconditional branch to Target Label.

Target Label must be the Result <id> of an OpLabel instruction in the current function.

This instruction must be the last instruction in a block.

2	249	<id>&gt;</id>
		Target Label

### **OpBranchConditional**

If Condition is **true**, branch to True Label, otherwise branch to False Label.

Condition must be a Boolean type scalar.

True Label must be an OpLabel in the current function.

False Label must be an OpLabel in the current function.

Branch weights are unsigned 32-bit integer literals. There must be either no Branch Weights or exactly two branch weights. If present, the first is the weight for branching to True Label, and the second is the weight for branching to False Label. The implied probability that a branch is taken is its weight divided by the sum of the two Branch weights. At least one weight must be non-zero. A weight of zero does not imply a branch is dead or permit its removal; branch weights are only hints. The two weights must not overflow a 32-bit unsigned integer when added together.

This instruction must be the last instruction in a block.

4 + variable	250	<id></id>	<id></id>	<id></id>	Literal, Literal,
		Condition	True Label	False Label	Branch weights

## **OpSwitch**

Multi-way branch to one of the operand label  $\langle id \rangle$ .

Selector must have a type of OpTypeInt. Selector will be compared for equality to the Target literals.

*Default* must be the < id > of a label. If *Selector* does not equal any of the *Target* literals, control flow will branch to the *Default* label < id >.

Target must be alternating scalar integer *literals* and the <*id*> of a label. If *Selector* equals a *literal*, control flow will branch to the following *label* <*id*>. It is invalid for any two *literal* to be equal to each other. If *Selector* does not equal any *literal*, control flow will branch to the *Default* label <*id*>. Each *literal* is interpreted with the type of *Selector*: The bit width of *Selector's* type will be the width of each *literal's* type. If this width is not a multiple of 32-bits, the literals must be sign extended when the OpTypeInt *Signedness* is set to 1. (See Literal Number.)

This instruction must be the last instruction in a block.

3 + variable	251	<id>&gt; Selector</id>	<id> Default</id>	literal, label <id>, literal, label <id>,</id></id>
				 Target

OpKill	Capability:
	Shader
Fragment-shader discard.	
Ceases all further processing in any invocation that	
executes it: Only instructions these invocations executed	
before <b>OpKill</b> will have observable side effects. If this	
instruction is executed in non-uniform control flow, all	
subsequent control flow is non-uniform (for invocations	
that continue to execute).	
This instruction must be the last instruction in a block.	
This instruction is only valid in the <b>Fragment</b> Execution	
Model.	
1	252

OpReturn	
Return with no value from a fund	ction with void return type.
This instruction must be the last	instruction in a block.
1	253

## **OpReturnValue**

Return a value from a function.

*Value* is the value returned, by copy, and must match the *Return Type* operand of the OpTypeFunction type of the OpFunction body this return instruction is in.

This instruction must be the last instruction in a block.

2	254	< <i>id</i> >
		Value

# OpUnreachable

Pointer

Declares that this block is not reachable in the CFG.

This instruction must be the last instruction in a block.

1 255

#### 

Size

OpLifetimeSt	op		Capability:
Pointer is a point	object is dead at inter to the object	Kernel	
capability is no of memory wh	if <i>Pointer</i> is a poot being used. If <i>S</i> ose lifetime is ented as unsigned;		
3	3 257 < <i>id&gt;</i>		Literal Number
		Pointer	Size

#### 3.32.18 Atomic Instructions

### **OpAtomicLoad**

Atomically load through *Pointer* using the given *Semantics*. All subparts of the value that is loaded will be read atomically with respect to all other atomic accesses to it within *Scope*.

Result Type must be a scalar of integer type or floating-point type.

*Pointer* is the pointer to the memory to read. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

6	227	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory
		Result Type		Pointer	Scope	Semantics <id></id>
						Semantics

#### **OpAtomicStore**

Atomically store through *Pointer* using the given *Semantics*. All subparts of *Value* will be written atomically with respect to all other atomic accesses to it within *Scope*.

*Pointer* is the pointer to the memory to write. The type it points to must be a scalar of integer type or floating-point type.

Value is the value to write. The type of Value and the type pointed to by Pointer must be the same type.

5	228	< <i>id&gt;</i>	Scope <id></id>	Memory Semantics	< <i>id</i> >
		Pointer	Scope	<id>&gt;</id>	Value
				Semantics	

### **OpAtomicExchange**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value from copying Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be a scalar of integer type or floating-point type.

7	229	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

### **OpAtomicCompareExchange**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by selecting Value if Original Value equals Comparator or selecting Original Value otherwise, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

Use Equal for the memory semantics of this instruction when Value and Original Value compare equal.

Use *Unequal* for the memory semantics of this instruction when *Value* and *Original Value* compare unequal. *Unequal* cannot be set to **Release** or **Acquire and Release**. In addition, *Unequal* cannot be set to a stronger memory-order then *Equal*.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*. This type must also match the type of *Comparator*.

		• •							
9	230	< <i>id</i> >	Result	< <i>id</i> >	Scope	Memory	Memory	< <i>id</i> >	< <i>id</i> >
		Result	<id></id>	Pointer	<id></id>	Semantics	Semantics	Value	Comparator
		Type			Scope	<id></id>	<id></id>		
					_	Equal	Unequal		

OpA	tomic	Capability: Kernel							
Depr	recated								
Has	the san	ne semantics	as OpAtomic	CompareExc	hange.				
9	231	< <i>id</i> >	Result	< <i>id</i> >	Scope	Memory	Memory	< <i>id</i> >	< <i>id</i> >
		Result	<id></id>	Pointer	<id></id>	Semantics	Semantics	Value	Comparator
		Туре				Equal	Unequal		

### **OpAtomicIIncrement**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value through integer addition of 1 to Original Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

*Result Type* must be an integer type scalar. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

6	232	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory
		Result Type		Pointer	Scope	Semantics <id></id>
						Semantics

### **OpAtomicIDecrement**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value through integer subtraction of 1 from Original Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar. The type of the value pointed to by Pointer must be the same as Result Type.

6	233	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	
		Result Type		Pointer	Scope	Semantics <id></id>	
						Semantics	

### **OpAtomicIAdd**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by integer addition of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	234	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

### **OpAtomicISub**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by integer subtraction of Value from Original Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

	<i>J</i> 1						
7	235	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
					_	<id></id>	
						Semantics	

### **OpAtomicSMin**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by finding the smallest signed integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	236	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

### **OpAtomicUMin**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value by finding the smallest unsigned integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

7	237	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

#### **OpAtomicSMax**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by finding the largest signed integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	238	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	<id></id>
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

### **OpAtomicUMax**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value by finding the largest unsigned integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

t	7	239	∠; d>	Result <id></id>	< <i>id</i> >	Coope side	Mamari	<id>&gt;</id>
	/	239	<id></id>	Kesuit <iu></iu>	\ \ta>	Scope <id></id>	Memory	\(\alpha \alpha \right)
			Result Type		Pointer	Scope	Semantics	Value
			Tresuit Type		100000	Scope		,
							<1d>	
							Semantics	

#### **OpAtomicAnd**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by the bitwise AND of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	240	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

### **OpAtomicOr**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by the bitwise OR of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

_ L		<i>7</i> 1						
	7	241	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
			Result Type		Pointer	Scope	Semantics	Value
							<id></id>	
							Semantics	

### **OpAtomicXor**

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by the bitwise exclusive OR of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	242	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicFlag	TestAndSet	Capability:	Capability:		
				Kernel	
Atomically sets	the flag value pointe	d to by <i>Pointer</i> to	the set state.		
Pointer must be flag.	a pointer to a 32-bit				
	s result is true if the clear state immediate				
Result Type mus	st be a Boolean type.				
	efined if an atomic fla FlagTestAndSet or O				
6 318	<id>&gt;</id>	Scope <id></id>	Memory		
	Result Type		Pointer	Scope	Semantics <id></id>
				1	Semantics

<b>OpAtomic</b> I	FlagClear			Capability:
			Kernel	
Atomically	sets the flag			
<i>Pointer</i> mus	st be a pointe	er to a 32-bit integer type repres	enting an atomic flag.	
	.•			
Memory Se	mantics can	not be Acquire or AcquireRelea	se	
Descrite one	umdafinad if	on atomic flor is modified by a	instruction other than	
		an atomic flag is modified by a	i instruction other than	
OpAtomicF	lagTestAnd!			
4	319	< <i>id</i> >	Memory Semantics <id></id>	
		Pointer	Scope	Semantics

### 3.32.19 Primitive Instructions

OpEmitVertex	Capability:
	Geometry
Emits the current values of all output variables to the	
current output primitive. After execution, the values of	
all output variables are undefined.	
This instruction can only be used when only one stream	
is present.	
1	218

<b>OpEndPrimitive</b>	Capability:
	Geometry
Finish the current primitive and start a new one. No	
vertex is emitted.	
This instruction can only be used when only one	
stream is present.	
1	219

<b>OpEmitStreamVertex</b>	Capability:
	GeometryStreams
Emits the current values of all output variables	
to the current output primitive. After execution,	
the values of all output variables are undefined.	
Stream must be an <id> of a constant</id>	
instruction with a scalar integer type. That	
constant is the output-primitive stream number.	
This instruction can only be used when	
multiple streams are present.	
2 220	<id></id>
	Stream

<b>OpEndStreamPrimitiv</b>	ve	Capability:
		GeometryStreams
Finish the current primi	tive and start a new	
one. No vertex is emitte	ed.	
Stream must be an <id></id>	of a constant	
instruction with a scalar	integer type. That	
constant is the output-pr	rimitive stream number.	
This instruction can only	y be used when	
multiple streams are present.		
2	221	<id>&gt;</id>
		Stream

#### 3.32.20 Barrier Instructions

### **OpControlBarrier**

Wait for other invocations of this module to reach the current point of execution.

All invocations of this module within *Execution* scope must reach this point of execution before any invocation will proceed beyond it.

This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within *Execution*. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely.

If *Semantics* is not **None**, this instruction also serves as an OpMemoryBarrier instruction, and must also perform and adhere to the description and semantics of an **OpMemoryBarrier** instruction with the same *Memory* and *Semantics* operands. This allows atomically specifying both a control barrier and a memory barrier (that is, without needing two instructions). If *Semantics* is **None**, *Memory* is ignored.

It is only valid to use this instruction with **TessellationControl**, **GLCompute**, or **Kernel** execution models.

When used with the **TessellationControl** execution model, it also implicitly synchronizes the **Output** Storage Class: Writes to **Output** variables performed by any invocation executed prior to a **OpControlBarrier** will be visible to any other invocation after return from that **OpControlBarrier**.

4	224	Scope <id></id>	Scope <id></id>	Memory Semantics <id></id>
		Execution	Memory	Semantics

#### **OpMemoryBarrier**

Control the order that memory accesses are observed.

Ensures that memory accesses issued before this instruction will be observed before memory accesses issued after this instruction. This control is ensured only for memory accesses issued by this invocation and observed by another invocation executing within *Memory* scope.

Semantics declares what kind of memory is being controlled and what kind of control to apply.

To execute both a memory barrier and a control barrier, see OpControlBarrier.

	-	· · · · · · · · · · · · · · · · · · ·	
3	225	Scope <id></id>	Memory Semantics <id></id>
		Memory	Semantics

OpName	dBarrierInit	ialize		Capability:
		oarrier object. e type OpTypeNamedB	earrier.	NamedBarrier  Missing before version 1.1.
		be a 32-bit integer type ach the current point of	scalar representing the number of execution.	
4	328	<id>&gt;</id>		
		Result Type		Subgroup Count

OpMemor	yNamedBaı	rrier		Capability:
				NamedBarrier
Wait for otl	her invocatio			
		Missing before version 1.1.		
Named Bar	rrier must be	the type OpTypeNamedBarrier		
If Semantic	s is not <b>Non</b>	e, this instruction also serves as	an OpMemoryBarrier	
instruction,	, and must al	so perform and adhere to the de	scription and semantics of an	
OpMemor	yBarrier ins	struction with the same <i>Memory</i>	and Semantics operands. This	
allows aton	nically specif	fying both a control barrier and	a memory barrier (that is,	
without nee	eding two ins			
4	329	Memory Semantics <id></id>		
		Named Barrier	Memory	Semantics

# 3.32.21 Group Instructions

OpC	Group <i>A</i>	syncCopy						Capability: Kernel	
				of <i>Num Eleme</i> s performed b				Kernei	
	This instruction returns an event object that can be used by OpGroupWaitEvents to wait for the async copy to finish.								
All i	All invocations of this module within <i>Execution</i> must reach this point of execution.								
This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within <i>Execution</i> . This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely.									
Resu	ılt Type	must be an	OpTypeEvent	object.					
Dest	ination	must be a po	ointer to a scal	lar or vector o	f floating-poir	nt type or inte	ger type.		
Dest	ination	pointer Stora	age Class mus	st be <b>Workgro</b>	oup or CrossV	Vorkgroup.			
The	type of	Source must	be the same a	as Destination					
must	be Cr			ass is <b>Workgr</b> ese <i>Stride</i> defin					
Clas	s must		up. In this cas	ass is <b>CrossW</b> ese <i>Stride</i> defin					
				-bit integer ty lar when the A					
Even	at must	have a type o	of OpTypeEve	ent.					
				y with a previ Event should			to be		
	ent arg		OpConstantN	full, the event	object supplie	ed in event arg	gument will		
9	259	<id></id>	Result	Scope	< <i>id</i> >	<id>&gt;</id>	<id>&gt;</id>	<id>&gt;</id>	<id>&gt;</id>
		Result Type	<id></id>	<id> Execution</id>	Destination	Source	Num Elements	Stride	Event

OpGroupWa	aitEvents			Capability:
				Kernel
Wait for even	its generate			
List points to performed.	Num Even	ts event objects, which	h will be released after the wait is	
All invocation	ns of this n	nodule within Executi	on must reach this point of execution.	
control flow v	within <i>Exe</i>	cution. This ensures the	orrectly if placed strictly within uniform hat if any invocation executes it, all re, an invocation may stall indefinitely.	
Execution mu	ıst be <b>Wor</b>	kgroup or Subgroup	Scope.	
Num Events r	must be a 3			
Events List m	nust be a po			
4 2	260	Scope <id></id>	< <i>id</i> >	<id></id>
		Execution	Num Events	Events List

OpGrou	pAll				Capability:		
	Evaluates a predicate for all invocations in the group, resulting in <b>true</b> if predicate evaluates to <b>true</b> for all invocations in the group, otherwise the result is <b>false</b> .						
All invoc	ations of t	his module within Execu	tion must reach this poin	t of execution.			
flow with	This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within <i>Execution</i> . This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely.						
Result Ty	pe must b	e a Boolean type.					
Execution	n must be	Workgroup or Subgrou	p Scope.				
Predicate	Predicate must be a Boolean type.						
5	5   261   <id>   Result <id>   Scope <id>  </id></id></id>						
		Result Type		Execution	Predicate		

OpGrou	pAny				Capability:		
to <b>true</b> fo	or any invo	ate for all invocations in to cation in the group, other this module within Execution	erwise the result is <b>false</b> .		Groups		
flow with	This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within <i>Execution</i> . This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely.						
		e a Boolean type.  Workgroup or Subgrou	p Scope.				
Predicate	e must be a	a Boolean type.					
5	262	< <i>id&gt;</i>	Result <id></id>	Scope <id></id>	< <i>id</i> >		
		Result Type		Execution	Predicate		

OpGr	oupBroa	dcast			Capability:	
1	Return the <i>Value</i> of the invocation identified by the local id <i>LocalId</i> to all invocations in the group.			Groups		
All inv		of this module within	n <i>Execution</i> must r	each this point of		
uniform execute	This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within <i>Execution</i> . This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely.					
Result type so		st be a 32-bit or 64-b	it integer type or a	16, 32 or 64 float		
Execut	tion must	be <b>Workgroup</b> or <b>S</b>	ubgroup Scope.			
The ty	pe of Valu	ue must be the same	as Result Type.			
compo	LocalId must be an integer datatype. It can be a scalar, or a vector with 2 components or a vector with 3 components. LocalId must be the same for all invocations in the group.					
6	263	<id></id>	Result <id></id>	Scope <id></id>	< <i>id</i> >	<id></id>
		Result Type		Execution	Value	LocalId

### OpGroupIAdd Capability: **Groups** An integer add group operation specified for all values of X specified by invocations in the group. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 32-bit or 64-bit integer type scalar. Execution must be Workgroup or Subgroup Scope. The identity *I* for *Operation* is 0. The type of *X* must be the same as *Result Type*. Result <id> Scope <id> **Group Operation** 264 <*id*> $\langle id \rangle$ Result Type Execution Operation X

OpGro	oupFAdd				Capability:	
	A floating-point add group operation specified for all values of <i>X</i> specified by invocations in the group.					
All inve		of this module within	n <i>Execution</i> must re	each this point of		
uniforn execute	n control	flow within Execution wocations will execution	on. This ensures that	placed strictly within at if any invocation where, an invocation		
Result	<i>Type</i> mus	at be a 16-bit, 32-bit,	or 64-bit floating-p	oint type scalar.		
Executi	ion must	be <b>Workgroup</b> or <b>S</b> t	abgroup Scope.			
The ide	The identity <i>I</i> for <i>Operation</i> is 0.					
The typ	The type of <i>X</i> must be the same as <i>Result Type</i> .					
6	265	<id></id>	Result <id></id>	Scope <id></id>	Group Operation	< <i>id</i> >
		Result Type		Execution	Operation	X

### **OpGroupFMin** Capability: **Groups** A floating-point minimum group operation specified for all values of Xspecified by invocations in the group. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 16-bit, 32-bit, or 64-bit floating-point type scalar. Execution must be Workgroup or Subgroup Scope. The identity I for Operation is +INF. The type of *X* must be the same as *Result Type*. Result <id> Scope <id> **Group Operation** 266 <*id*> $\langle id \rangle$ Result Type Execution Operation X

OpGro	upUMir	1			Capability:	
	_	eger minimum grou ocations in the grou	Groups			
All invo		of this module withi	n <i>Execution</i> must r	each this point of		
uniforn execute	n control	flow within Execution will executed a second control of the contro	on. This ensures th	placed strictly within at if any invocation where, an invocation		
Result '	<i>Type</i> mus	et be a 32-bit or 64-b	it integer type scala	nr.		
Executi	ion must	be <b>Workgroup</b> or <b>S</b>	ubgroup Scope.			
1	•	or $Operation$ is UIN when $X$ is 64 bits w				
	The type of <i>X</i> must be the same as <i>Result Type</i> .					
6	267	<id></id>	Result <id></id>	Scope <id></id>	Group Operation	<id>&gt;</id>
		Result Type		Execution	Operation	X

### **OpGroupSMin** Capability: **Groups** A signed integer minimum group operation specified for all values of Xspecified by invocations in the group. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 32-bit or 64-bit integer type scalar. Execution must be Workgroup or Subgroup Scope. The identity *I* for *Operation* is INT\_MAX when *X* is 32 bits wide and LONG MAX when X is 64 bits wide. The type of *X* must be the same as *Result Type*. 268 <*id*> Result <id> Scope <id> **Group Operation** <*id*> Result Type Execution Operation X

OpGroup	pFMax			Capability:	
1	g-point maximum group or by invocations in the grou	Groups			
All invoca	ations of this module with	in Execution must re	each this point of		
uniform c	ruction is only guaranteed control flow within <i>Execut</i> it, all invocations will exe indefinitely.	ion. This ensures the	at if any invocation		
Result Ty	pe must be a 16-bit, 32-bi	t, or 64-bit floating-p	point type scalar.		
Execution	n must be <b>Workgroup</b> or	Subgroup Scope.			
The ident	tity <i>I</i> for <i>Operation</i> is -IN				
The type	of $X$ must be the same as	Result Type.			
6 2	269 <id>&gt;</id>	Result <id></id>	Scope <id></id>	Group Operation	< <i>id</i> >
	Result Type		Execution	Operation	X

### **OpGroupUMax** Capability: **Groups** An unsigned integer maximum group operation specified for all values of Xspecified by invocations in the group. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 32-bit or 64-bit integer type scalar. Execution must be Workgroup or Subgroup Scope. The identity *I* for *Operation* is 0. The type of *X* must be the same as *Result Type*. 270 Result <id> Scope <id> **Group Operation** <*id*> $\langle id \rangle$ Result Type Execution Operation X

OpGro	oupSMax				Capability:	
_	A signed integer maximum group operation specified for all values of <i>X</i> specified by invocations in the group.					
All inv		of this module withi	n <i>Execution</i> must re	ach this point of		
uniforr	n control	flow within <i>Executi</i> nvocations will exec	on. This ensures that	placed strictly within at if any invocation where, an invocation		
X and I	Result Typ	pe must be a 32-bit of	or 64-bit OpTypeInt	data type.		
Execut	ion must	be <b>Workgroup</b> or <b>S</b>	ubgroup Scope.			
LONG	The identity <i>I</i> for <i>Operation</i> is INT_MIN when <i>X</i> is 32 bits wide and LONG_MIN when <i>X</i> is 64 bits wide.					
	<u> </u>	ust be the same as F				
6	271	< <i>id</i> >	Result <id></id>	Scope <id></id>	Group Operation	< <i>id</i> >
		Result Type		Execution	Operation	X

OpSu	bgroupl	BallotKHR		Capability:	
				SubgroupBallot	KHR
See ex	tension	SPV_KHR_shader			
4	4421	< <i>id</i> >	Result <id></id>	<id></id>	
		Result Type		Predicate	

OpSu	OpSubgroupFirstInvocationKHR			Capability: SubgroupBallot	KHR
See ex	tension				
4	4422	<id>&gt;</id>	Result <id></id>	<id>&gt;</id>	
		Result Type		Value	

	OpSu	bgroup	Capability:				
	See ex	tension	SPV_KHR_shade	er_ballot		SubgroupBallot Reserved.	KHR
Ī	5	4432	< <i>id</i> >	Result <id></id>	<id></id>	<id></id>	
			Result Type		Value	Index	

OpGı	OpGroupIAddNonUniformAMD				Capability:		
TBD					Groups		
					Reserved.		
6	5000	< <i>id</i> >	Result <id></id>	Scope <id></id>	Group	< <i>id</i> >	
	Result Type Execution				Operation	X	
			Operation				

OpGroupFAddNonUniformAMD				Capability:		
TBD					Groups	
					Reserved.	
6	5001	< <i>id</i> >	Result <id></id>	Scope <id></id>	Group	< <i>id</i> >
	Result Type Execution				Operation	X
			Operation			

OpGı	OpGroupFMinNonUniformAMD				Capability:		
TBD	TBD					Groups  Reserved.	
6	6   5002   <id>   Result <id>   Scope <id>   Execution  </id></id></id>					< <i>id&gt;</i> X	

OpGroupUMinNonUniformAMD				Capability:			
TBD					Groups Reserved.		
6	6 5003 < id> Result < id> Scope < id> Execution					<id> X</id>	

OpGı	OpGroupSMinNonUniformAMD				Capability:			
TBD	•					Groups Reserved.		
6	5004	<id> Result Type</id>	Result <id></id>	Scope <id> Execution</id>	Group Operation Operation	<id>X</id>		

OpGı	OpGroupFMaxNonUniformAMD				Capability:			
TBD						Groups Reserved.		
6	6   5005   <id>   Result <id>   Scope <id>   Execution  </id></id></id>					< <i>id&gt;</i> X		

OpGroupUMaxNonUniformAMD				Capability:			
TBD					Groups  Reserved.		
6	6   5006   <id> Result <id> Scope <id> Execution</id></id></id>					<id>X</id>	
	Result Type Execution				Operation Operation		

OpGı	OpGroupSMaxNonUniformAMD				Capability:	
TBD						
6	5007	<id> Result Type</id>	Result <id></id>	Scope <id> Execution</id>	Group Operation Operation	<id> X</id>

## 3.32.22 Device-Side Enqueue Instructions

OpEnqueueMarker	Capability: DeviceEnqueue	e			
Enqueue a marker command to command waits for a list of eve all previously enqueued comma completes.	•				
Result Type must be a 32-bit int the value 0. A failed enqueue re	neue results in				
Queue must be of the type OpT	ypeQueue.				
Num Events specifies the numb Wait Events and must be a 32-b unsigned integer.					
Wait Events specifies the list of OpTypeDeviceEvent.					
Ret Event is a pointer to a device instruction. It must have a type Event is set to null this instruction.					
7 291 < <i>id</i> >	Result <id></id>	<id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
Result Type		Queue	Num Events	Wait Events	Ret Event

### **OpEnqueueKernel**

Enqueue the function specified by *Invoke* and the NDRange specified by *ND Range* for execution to the queue object specified by Queue.

Result Type must be a 32-bit integer type scalar. A successful enqueue results in the value 0. A failed enqueue results in a non-0 value.

Queue must be of the type OpTypeQueue.

Flags must be an integer type scalar. The content of Flags is interpreted as Kernel Enqueue Flags mask.

The type of ND Range must be an OpTypeStruct whose members are as described by the Result Type of OpBuildNDRange.

Num Events specifies the number of event objects in the wait list pointed to by Wait *Events* and must be 32-bit integer type scalar, which is treated as an unsigned integer.

Wait Events specifies the list of wait event objects and must be a pointer to OpTypeDeviceEvent.

Ret Event must be a pointer to OpTypeDeviceEvent which gets implicitly retained by this instruction.

*Invoke* must be an OpFunction whose OpTypeFunction operand has:

- Result Type must be OpTypeVoid.
- The first parameter must have a type of OpTypePointer to an 8-bit OpTypeInt.
- An optional list of parameters, each of which must have a type of OpTypePointer to the Workgroup Storage Class.

Param is the first parameter of the function specified by *Invoke* and must be a pointer to an 8-bit integer type scalar.

Param Size is the size in bytes of the memory pointed to by Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer.

Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer.

Each Local Size operand corresponds (in order) to one OpTypePointer to Workgroup Storage Class parameter to the *Invoke* function, and specifies the number of bytes of Workgroup storage used to back the pointer during the execution of the Invoke function.

Ī	13 +	292	< <i>id</i> >	Result	< <i>id</i> >	<id>,</id>									
	vari-		Result	<id></id>	Queue	Flags	ND	Num	Wait	Ret	Invoke	Param	Param	Param	<id>,</id>
	able		Type			-	Range	Events	Events	Event			Size	Align	
															Local
															Size

# DeviceEnqueue

#### **OpGetKernelNDrangeSubGroupCount** Capability: **DeviceEnqueue** Returns the number of subgroups in each workgroup of the dispatch (except for the last in cases where the global size does not divide cleanly into work-groups) given the combination of the passed NDRange descriptor specified by ND Range and the function specified by Invoke. Result Type must be a 32-bit integer type scalar. The type of ND Range must be an OpTypeStruct whose members are as described by the Result Type of OpBuildNDRange. *Invoke* must be an OpFunction whose OpTypeFunction operand has: - Result Type must be OpTypeVoid. - The first parameter must have a type of OpTypePointer to an 8-bit OpTypeInt. - An optional list of parameters, each of which must have a type of OpTypePointer to the Workgroup Storage Class. Param is the first parameter of the function specified by Invoke and must be a pointer to an 8-bit integer type scalar. Param Size is the size in bytes of the memory pointed to by Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer. Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer. 293 <*id*> Result <id> <*id*> $\langle id \rangle$ $\langle id \rangle$ $\langle id \rangle$ <*id*> Invoke Result Type ND Range Param Param Size Param Align

#### **OpGetKernelNDrangeMaxSubGroupSize** Capability: **DeviceEnqueue** Returns the maximum sub-group size for the function specified by *Invoke* and the NDRange specified by ND Range. Result Type must be a 32-bit integer type scalar. The type of ND Range must be an OpTypeStruct whose members are as described by the Result Type of OpBuildNDRange. *Invoke* must be an OpFunction whose OpTypeFunction operand has: - Result Type must be OpTypeVoid. - The first parameter must have a type of OpTypePointer to an 8-bit OpTypeInt. - An optional list of parameters, each of which must have a type of OpTypePointer to the Workgroup Storage Class. Param is the first parameter of the function specified by Invoke and must be a pointer to an 8-bit integer type scalar. Param Size is the size in bytes of the memory pointed to by Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer. Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer. 294 <*id*> <*id*> <id> <*id*> <*id*> Result <id> <*id*> Result Type ND Range Invoke Param Param Size Param Align

OpG	etKerne	lWorkGroupSize	9			Capability:	
						DeviceEnque	ue
		aximum work-growoke on the device					
Resul	<i>t Type</i> m	ust be a 32-bit in	teger type scalar.				
		e an OpFunction		unction opera	nd has:		
		must be OpType\		D. C.	0.1% 0		
	-				8-bit OpTypeInt.		
	-	-		nust have a ty	pe of OpTypePointer		
to the	Workgi	roup Storage Cla	SS.				
Paran	<i>n</i> is the f	irst parameter of	the function speci	ified by <i>Invok</i>	e and must be a		
pointe	er to an 8	3-bit integer type	scalar.				
Paran	n Size is	the size in bytes	of the memory po	ointed to by Pa	aram and must be a		
		type scalar, which	• •	•			
D	47	.1 1'	CD 1	. 1			
		_		t be a 32-bit 1	nteger type scalar,		
		ed as an unsigned					
7	295	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		Invoke	Param	Param Size	Param Align

#### **OpGetKernelPreferredWorkGroupSizeMultiple** Capability: **DeviceEnqueue** Returns the preferred multiple of work-group size for the function specified by *Invoke*. This is a performance hint. Specifying a work-group size that is not a multiple of the value returned by this query as the value of the local work size will not fail to enqueue *Invoke* for execution unless the work-group size specified is larger than the device maximum. Result Type must be a 32-bit integer type scalar. *Invoke* must be an OpFunction whose OpTypeFunction operand has: - Result Type must be OpTypeVoid. - The first parameter must have a type of OpTypePointer to an 8-bit OpTypeInt. - An optional list of parameters, each of which must have a type of OpTypePointer to the Workgroup Storage Class. Param is the first parameter of the function specified by Invoke and must be a pointer to an 8-bit integer type scalar. Param Size is the size in bytes of the memory pointed to by Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer. Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer. <*id*> <id> <id> 296 <*id*> Result <id> <*id*> Result Type Invoke Param Param Size Param Align

OpRetainEvent	Capability:
Increments the reference count of the event object specified by <i>Event</i> .	DeviceEnqueue
Event must be an event that was produced by OpEnqueueKernel, OpEnqueueMarker or	
OpCreateUserEvent.	
2 297	<id>&gt;</id>
	Event

OpReleaseEvent	Capability:
	DeviceEnqueue
Decrements the reference count of the event	
object specified by <i>Event</i> . The event object is	
deleted once the event reference count is zero,	
the specific command identified by this event	
has completed (or terminated) and there are no	
commands in any device command queue that	
require a wait for this event to complete.	
Event must be an event that was produced by	
OpEnqueueKernel, OpEnqueueMarker or	
OpCreateUserEvent.	
2 298	<id></id>
	Event

	UserEvent er event. The to a value of	Capability: DeviceEnqueue	
Result Type	e must be Op	TypeDeviceEvent.	
3	299	Result <id></id>	
		Result Type	

OpIsValid	lEvent	Capability:		
Returns tr false.	<b>ue</b> if the eve	DeviceEnqueue		
		Boolean type.		
Event mus	t have a type	e of OpTypeDeviceEvent		
4	300	<id></id>		
		Result Type		Event

OpSetUserEve	entStatus		Capability:
		DeviceEnqueue	
Sets the execut	ion status of a us	er event specified by Event.Status can be	
either 0 (CL_C	OMPLETE) to it	ndicate that this kernel and all its child	
kernels finished	d execution succe	essfully, or a negative integer value	
indicating an en	rror.		
Event must hav	e a type of OpTy	peDeviceEvent that was produced by	
OpCreateUserI	Event.		
Status must have	ve a type of 32-bi		
3	301	<id></id>	<id></id>
		Event	Status

# **OpCaptureEventProfilingInfo** Capability: DeviceEnqueue Captures the profiling information specified by Profiling Info for the command associated with the event specified by *Event* in the memory pointed to by *Value*. The profiling information will be available in the memory pointed to by Value once the command identified by Event has completed. Event must have a type of OpTypeDeviceEvent that was produced by OpEnqueueKernel or OpEnqueueMarker. Profiling Info must be an integer type scalar. The content of Profiling Info is interpreted as Kernel Profiling Info mask. Value must be a pointer to a scalar 8-bit integer type in the CrossWorkgroup Storage Class. When *Profiling Info* is **CmdExecTime**, *Value* must point to 128-bit memory range. The first 64 bits contain the elapsed time CL\_PROFILING\_COMMAND\_END -CL PROFILING COMMAND START for the command identified by Event in nanoseconds. The second 64 bits contain the elapsed time CL\_PROFILING\_COMMAND\_COMPLETE -CL\_PROFILING\_COMMAND\_START for the command identified by Event in nanoseconds. Note: The behavior of this instruction is undefined when called multiple times for the same event. 302 4 <*id*> <*id*> <id> Value Event Profiling Info

OpGetDef:	aultQueue	Capability:	
	e default device not been creat	DeviceEnqueue	
Result Type	must be an (	OpTypeQueue.	
3	303	Result <id></id>	
		Result Type	

# **OpBuildNDRange**

# Given the global work size specified by GlobalWorkSize, local work size specified by LocalWorkSize and global work offset specified by GlobalWorkOffset, builds a 1D, 2D or 3D ND-range descriptor structure and returns it.

Result Type must be an OpTypeStruct with the following ordered list of members, starting from the first to last:

- 1) 32-bit integer type scalar, that specifies the number of dimensions used to specify the global work-items and work-items in the work-group.
- 2) OpTypeArray with 3 elements, where each element is 32-bit integer type scalar when the addressing model is **Physical32** and 64-bit integer type scalar when the addressing model is **Physical64**. This member is an array of per-dimension unsigned values that describe the offset used to calculate the global ID of a work-item.
- 3) OpTypeArray with 3 elements, where each element is 32-bit integer type scalar when the addressing model is **Physical32** and 64-bit integer type scalar when the addressing model is **Physical64**. This member is an array of per-dimension unsigned values that describe the number of global work-items in the dimensions that will execute the kernel function.
- 4) OpTypeArray with 3 elements, where each element is 32-bit integer type scalar when the addressing model is **Physical32** and 64-bit integer type scalar when the addressing model is **Physical64**. This member is an array of per-dimension unsigned values that describe the number of work-items that make up a work-group.

GlobalWorkSize must be a scalar or an array with 2 or 3 components. Where the type of each element in the array is 32-bit integer type scalar when the addressing model is **Physical32** or 64-bit integer type scalar when the addressing model is Physical64.

The type of *LocalWorkSize* must be the same as *GlobalWorkSize*.

The type of GlobalWorkOffset must be the same as GlobalWorkSize

# Capability:

# **DeviceEnqueue**

1110	type of Gi	obai workojjsei iliu				
6	304	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id></id>
		Result Type		GlobalWorkSize	LocalWorkSize	GlobalWorkOffset

OpGetKerne	elLocalSizeFo	rSubgroupCou	int			Capability:		
Returns the 1 workgroup.	SubgroupDis  Missing before 1.1.	_						
Result Type n	nust be a 32-bi	t integer type so	ealar.			1111		
Subgroup Co	ount must be a 3	32-bit integer ty	pe scalar.					
- Result Type	must be OpTy							
		ave a type of O ters, each of wh						
-	oup Storage Cla			71 1 .	, 1			
	first parameter ger type scalar.	of the function	specified by <i>Ir</i>	woke and must	t be a pointer to			
	s the size in byt scalar, which is							
Param Align	Param Align is the alignment of Param and must be a 32-bit integer type scalar, which							
	is treated as an unsigned integer.							
8 325	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	<id>&gt;</id>	< <i>id</i> >	
	Result Type		Subgroup Count	Invoke	Param	Param Size	Param	
			Count				Align	

OpGetKernelMaxNumSubgr	roups			Capability: SubgroupDispa	atch			
Returns the maximum number the devce.	Missing before							
Result Type must be a 32-bit in	teger type scalar.							
<ul> <li>Result Type must be OpType</li> <li>The first parameter must have</li> <li>An optional list of parameters</li> </ul>	<ul> <li>Invoke must be an OpFunction whose OpTypeFunction operand has:</li> <li>Result Type must be OpTypeVoid.</li> <li>The first parameter must have a type of OpTypePointer to an 8-bit OpTypeInt.</li> <li>An optional list of parameters, each of which must have a type of OpTypePointer to the Workgroup Storage Class.</li> </ul>							
Param is the first parameter of pointer to an 8-bit integer type		ified by <i>Invoke</i> ar	nd must be a					
-	Param Size is the size in bytes of the memory pointed to by Param and must be a 32-bit integer type scalar, which is treated as an unsigned integer.							
Param Align is the alignment of which is treated as an unsigned								
7 326 <id>Result Type</id>	Result <id></id>	<id> Invoke</id>	<id> Param</id>	<id> Param Size</id>	<id> Param Align</id>			

# 3.32.23 Pipe Instructions

OpReadPipe					Capability: Pipes	
-	from the pipe ob accessful and a ne	Tipes				
Result Type m	ust be a 32-bit int	teger type scalar.				
Pipe must hav	re a type of OpTy	pePipe with Read	lOnly access of	jualifier.		
Pointer must l		TypePointer with	the same data	type as <i>Pipe</i> and a		
Packet Size meach packet in		eger type scalar t	hat represents	the size in bytes of		
	nent must be a 32- ch packet in the pa		calar that pres	ents the alignment		
- 1 <= <i>Packet</i> .	nd Packet Alignme Alignment <= Pac ment must evenly	cket Size.	C			
	ypes, <i>Packet Aligi</i> A <i>lignment</i> should ypes.					
7 274	<id> Result Type</id>	Result <id></id>	<id> Pipe</id>	<id> Pointer</id>	<id> Packet Size</id>	<id> Packet Alignment</id>

### **OpWritePipe** Capability: **Pipes** Write a packet from *Pointer* to the pipe object specified by *Pipe*. Result is 0 if the operation is successful and a negative value if the pipe is full. Result Type must be a 32-bit integer type scalar. *Pipe* must have a type of OpTypePipe with WriteOnly access qualifier. Pointer must have a type of OpTypePointer with the same data type as Pipe and a Generic Storage Class. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, Packet Alignment should equal Packet Size. For aggregate types, Packet Alignment should be the size of the largest primitive type in the hierarchy of types. 275 Result <id> <*id*> <*id*> <*id*> <*id*> <*id*> Result Type Pipe Pointer Packet Size Packet Alignment

#### **OpReservedReadPipe** Capability: **Pipes** Read a packet from the reserved area specified by Reserve Id and Index of the pipe object specified by *Pipe* into *Pointer*. The reserved pipe entries are referred to by indices that go from 0... Num Packets - 1. Result is 0 if the operation is successful and a negative value otherwise. Result Type must be a 32-bit integer type scalar. *Pipe* must have a type of OpTypePipe with **ReadOnly** access qualifier. Reserve Id must have a type of OpTypeReserveId. *Index* must be a 32-bit integer type scalar, which is treated as an unsigned value. Pointer must have a type of OpTypePointer with the same data type as Pipe and a Generic Storage Class. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, Packet Alignment should equal Packet Size. For aggregate types, Packet Alignment should be the size of the largest primitive type in the hierarchy of types. 276 <*id*> Result <*id*> <*id*> <*id*> <*id*> <*id*> <*id*> Result < id >Pipe Reserve Index Pointer Packet Packet Type IdSize Alignment

Op	Reserve	edWritePipe						Capability Pipes	:
pipe fror	ite a pace object n 0 erwise.	Tipes							
Res	ult Type	must be a 32	2-bit integer	type scalar.					
Pip	e must l	nave a type o	f OpTypePip	e with Write	Only access qu	ualifier.			
Res	erve Id	must have a	type of OpTy	peReserveId					
Ind	ex must	be a 32-bit in	nteger type s	calar, which	is treated as an	unsigned va	lue.		
	nter mu rage Cla		e of OpType	Pointer with	the same data t	ype as <i>Pipe</i> a	and a <b>Generic</b>		
	<i>ket Size</i> ket in th		2-bit integer	type scalar th	at represents th	ne size in byt	es of each		
	_	<i>nment</i> must leket in the pip		nteger type sc	alar that preser	nts the alignm	nent in bytes		
- 1 -	<= <i>Pack</i>	et Alignment	<= Packet S	ust satisfy the Size. de <i>Packet Siz</i>	_				
					al <i>Packet Size</i> . I		e types, <i>Packet</i> types.		
9	277	< <i>id</i> >	Result	< <i>id&gt;</i>	< <i>id&gt;</i>	< <i>id&gt;</i>	< <i>id&gt;</i>	< <i>id</i> >	< <i>id</i> >
		Result	<id></id>	Pipe	Reserve	Index	Pointer	Packet	Packet
		Туре			Id			Size	Alignment

## **OpReserveReadPipePackets** Capability: **Pipes** Reserve Num Packets entries for reading from the pipe object specified by Pipe. Result is a valid reservation ID if the reservation is successful. Result Type must be an OpTypeReserveId. *Pipe* must have a type of OpTypePipe with **ReadOnly** access qualifier. Num Packets must be a 32-bit integer type scalar, which is treated as an unsigned Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, Packet Alignment should equal Packet Size. For aggregate types, Packet Alignment should be the size of the largest primitive type in the hierarchy of types. 278 <*id*> Result <id> $\langle id \rangle$ $\langle id \rangle$ $\langle id \rangle$ <*id*> Packet Size Result Type Pipe Num Packets Packet Alignment

### **OpReserveWritePipePackets**

Reserve *num\_packets* entries for writing to the pipe object specified by *Pipe*. Result is a valid reservation ID if the reservation is successful.

*Pipe* must have a type of OpTypePipe with WriteOnly access qualifier.

Num Packets must be a 32-bit OpTypeInt which is treated as an unsigned value.

Result Type must be an OpTypeReserveId.

*Packet Size* must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe.

*Packet Alignment* must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe.

Packet Size and Packet Alignment must satisfy the following:

- 1 <= Packet Alignment <= Packet Size.
- Packet Alignment must evenly divide Packet Size

For concrete types, *Packet Alignment* should equal *Packet Size*. For aggregate types, *Packet Alignment* should be the size of the largest primitive type in the hierarchy of types.

# Capability:

**Pipes** 

Ī	7	279	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
			Result Type		Pipe	Num Packets	Packet Size	Packet
								Alignment

OpCom	mitReadF	Pipe			Capability:		
					Pipes		
	Indicates that all reads to <i>Num Packets</i> associated with the reservation specified by <i>Reserve</i>						
<i>Id</i> and th	e pipe obj	ect specified by <i>Pipe</i> are	completed.				
Pipe mus	Pipe must have a type of OpTypePipe with <b>ReadOnly</b> access qualifier.						
Reserve	<i>Id</i> must ha	ave a type of OpTypeRese	erveId.				
	Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe.						
	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe.						
Packet Si	ize and Pa	cket Alignment must satis	fy the following:				
		ment <= Packet Size.	, .				
- Packet	Alignment	must evenly divide Pack	et Size				
	For concrete types, Packet Alignment should equal Packet Size. For aggregate types, Packet						
	Alignment should be the size of the largest primitive type in the hierarchy of types.						
5	280	<id></id>	< <i>id&gt;</i>	< <i>id&gt;</i>	< <i>id</i> >		
		Pipe	Reserve Id	Packet Size	Packet Alignment		

OpCom	mitWrite	Pipe			Capability: Pipes		
	Indicates that all writes to <i>Num Packets</i> associated with the reservation specified by <i>Reserve Id</i> and the pipe object specified by <i>Pipe</i> are completed.						
Pipe mu	st have a ty	ype of OpTypePipe with	WriteOnly access qualifi	ier.			
Reserve	Id must ha	ave a type of OpTypeRese	erveId.				
	Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe.						
	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe.						
- 1 <= Pa	Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size Packet Alignment must evenly divide Packet Size						
	For concrete types, <i>Packet Alignment</i> should equal <i>Packet Size</i> . For aggregate types, <i>Packet Alignment</i> should be the size of the largest primitive type in the hierarchy of types.						
5	281	<id><id>&lt;</id></id>	<id>&gt;</id>	<id>&gt;</id>	<id>&gt;</id>		
		Pipe	Reserve Id	Packet Size	Packet Alignment		

OpIsVali	dReserveId	Capability: Pipes		
Return <b>tr</b>	<b>ue</b> if <i>Reser</i> v	P***		
Result Typ	<i>pe</i> must be a			
Reserve Id	d must have			
4	282	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Reserve Id

### **OpGetNumPipePackets**

Result is the number of available entries in the pipe object specified by *Pipe*. The number of available entries in a pipe is a dynamic value. The value returned should be considered immediately stale.

*Result Type* must be a 32-bit integer type scalar, which should be treated as an unsigned value.

*Pipe* must have a type of OpTypePipe with **ReadOnly** or **WriteOnly** access qualifier.

*Packet Size* must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe.

*Packet Alignment* must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe.

Packet Size and Packet Alignment must satisfy the following:

- 1 <= Packet Alignment <= Packet Size.
- Packet Alignment must evenly divide Packet Size

For concrete types, *Packet Alignment* should equal *Packet Size*. For aggregate types, *Packet Alignment* should be the size of the largest primitive type in the hierarchy of types.

# Op Get Max Pipe Packets

Result is the maximum number of packets specified when the pipe object specified by *Pipe* was created.

*Result Type* must be a 32-bit integer type scalar, which should be treated as an unsigned value.

*Pipe* must have a type of OpTypePipe with **ReadOnly** or **WriteOnly** access qualifier.

*Packet Size* must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe.

*Packet Alignment* must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe.

Packet Size and Packet Alignment must satisfy the following:

- 1 <= Packet Alignment <= Packet Size.
- Packet Alignment must evenly divide Packet Size

For concrete types, *Packet Alignment* should equal *Packet Size*. For aggregate types, *Packet Alignment* should be the size of the largest primitive type in the hierarchy of types.

# Capability:

Capability: **Pipes** 

#### **Pipes**

6	284	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	ı
		Result Type		Pipe	Packet Size	Packet Alignment	

# **OpGroupReserveReadPipePackets** Capability: **Pipes** Reserve Num Packets entries for reading from the pipe object specified by Pipe at group level. Result is a valid reservation id if the reservation is successful. The reserved pipe entries are referred to by indices that go from 0 ... Num Packets - 1. All invocations of this module within *Execution* must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be an OpTypeReserveId. Execution must be Workgroup or Subgroup Scope. Pipe must have a type of OpTypePipe with ReadOnly access qualifier. Num Packets must be a 32-bit integer type scalar, which is treated as an unsigned value. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, *Packet Alignment* should equal *Packet Size*. For aggregate types, Packet Alignment should be the size of the largest primitive type in the hierarchy of types. <id> 8 285 $\overline{\langle id \rangle}$ Result <id> Scope <id> $\overline{\langle id \rangle}$ $\overline{\langle id \rangle}$ $\overline{\langle id \rangle}$ Result Type Execution Pipe Num Packet Size Packet Packets Alignment

# **OpGroupReserveWritePipePackets** Capability: **Pipes** Reserve Num Packets entries for writing to the pipe object specified by Pipe at group level. Result is a valid reservation ID if the reservation is successful. The reserved pipe entries are referred to by indices that go from 0 ... Num Packets - 1. All invocations of this module within *Execution* must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be an OpTypeReserveId. Execution must be Workgroup or Subgroup Scope. Pipe must have a type of OpTypePipe with WriteOnly access qualifier. Num Packets must be a 32-bit integer type scalar, which is treated as an unsigned value. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, *Packet Alignment* should equal *Packet Size*. For aggregate types, Packet Alignment should be the size of the largest primitive type in the hierarchy of types. <id> <*id*> 8 286 $\overline{\langle id \rangle}$ Result <id> Scope <id> $\overline{\langle id \rangle}$ $\overline{\langle id \rangle}$ Result Type Execution Pipe Num Packet Size Packet Packets Alignment

# **OpGroupCommitReadPipe** Capability: **Pipes** A group level indication that all reads to *Num Packets* associated with the reservation specified by Reserve Id to the pipe object specified by Pipe are completed. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Execution must be **Workgroup** or **Subgroup** Scope. Pipe must have a type of OpTypePipe with ReadOnly access qualifier. Reserve Id must have a type of OpTypeReserveId. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, Packet Alignment should equal Packet Size. For aggregate types, *Packet Alignment* should be the size of the largest primitive type in the hierarchy of types. 287 Scope <id> <*id*> $\langle id \rangle$ $\langle id \rangle$ $\overline{\langle id \rangle}$ Execution Pipe Reserve Id Packet Size Packet Alignment

# **OpGroupCommitWritePipe** Capability: **Pipes** A group level indication that all writes to Num Packets associated with the reservation specified by Reserve Id to the pipe object specified by Pipe are completed. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Execution must be **Workgroup** or **Subgroup** Scope. Pipe must have a type of OpTypePipe with WriteOnly access qualifier. Reserve Id must have a type of OpTypeReserveId. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe. Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe. Packet Size and Packet Alignment must satisfy the following: - 1 <= Packet Alignment <= Packet Size. - Packet Alignment must evenly divide Packet Size For concrete types, Packet Alignment should equal Packet Size. For aggregate types, *Packet Alignment* should be the size of the largest primitive type in the hierarchy of types. 288 Scope <id> <*id*> $\langle id \rangle$ $\overline{\langle id \rangle}$ $\overline{\langle id \rangle}$ Execution Pipe Reserve Id Packet Size Packet Alignment

<b>OpConstantPi</b>	peStorage			Capability:	
				PipeStorage	
Creates a pipe-s	torage object.				
Result Type mu	st be OpTypePipeSto	Missing before ver	sion 1.1.		
	t be a 32-bit integer tacket in the pipe.				
	nt must be a 32-bit in tes of each packet in				
- 1 <= <i>Packet A</i>	Packet Alignment mulignment <= Packet Sent must evenly divid				
	nes, Packet Alignmen , Packet Alignment sh archy of types.				
Capacity is the OpTypePipeSto	minimum number of rage can hold.				
6 323	< <i>id</i> >	Result <id></id>	Literal Number	Literal Number	Literal Number
	Result Type		Packet Size	Packet Alignment	Capacity

<b>OpCreate</b>	PipeFromPi	Capability:		
Creates a p	ipe object from must be Op	om a pipe-storage object.	n OpConstantPipeStorage.	PipeStorage  Missing before version 1.1.
	the pipe acc			
4	324	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Pipe Storage

# **A** Changes

# A.1 Changes from Version 0.99, Revision 31

- Added the PushConstant Storage Class.
- Added OpIAddCarry, OpISubBorrow, OpUMulExtended, and OpSMulExtended.
- Added OpInBoundsPtrAccessChain.
- Added the Decoration NoContraction to prevent combining multiple operations into a single operation (bug 14396).
- Added sparse texturing (14486):
  - Added **OpImageSparse...** for accessing images that might not be resident.
  - Added MinLod functionality for accessing images with a minimum level of detail.
- Added back the **Alignment** Decoration, for the **Kernel** capability (14505).
- Added a NonTemporal Memory Access (14566).
- Structured control flow changes:
  - Changed structured loops to have a structured continue *Continue Target* in OpLoopMerge (14422).
  - Added rules for how "fall through" works with **OpSwitch** (13579).
  - Added definitions for what is "inside" a structured control-flow construct (14422).
- Added **SubpassData** Dim to support input targets written by a previous subpass as an output target (14304). This is also a Decoration and a Capability, and can be used by some image ops to read the input target.
- Added OpTypeForwardPointer to establish the Storage Class of a forward reference to a pointer type (13822).
- · Improved Debuggability
  - Changed OpLine to not have a target <id>, but instead be placed immediately preceding the instruction(s) it is annotating (13905).
  - Added OpNoLine to terminate the affect of **OpLine** (13905).
  - Changed OpSource to include the source code:
    - \* Allow multiple occurrences.
    - \* Be mixed in with the OpString instructions.
    - \* Optionally consume an OpString result to say which file it is annotating.
    - \* Optionally include the source text corresponding to that OpString.
    - \* Included adding OpSourceContinued for source text that is too long for a single instruction.
- Added a large number of Capabilities for subsetting functionality (14520, 14453), including 8-bit integer support for OpenCL kernels.
- Added VertexIndex and InstanceIndex BuiltIn Decorations (14255).
- Added GenericPointer capability that allows the ability to use the Generic Storage Class (14287).
- Added IndependentForwardProgress Execution Mode (14271).
- Added OpAtomicFlagClear and OpAtomicFlagTestAndSet instructions (14315).
- Changed OpentryPoint to take a list of **Input** and **Output** < id> for declaring the entry point's interface.
- · Fixed internal bugs
  - 14411 Added missing documentation for mad\_sat OpenCL extended instructions (enums existed, just the documentation was missing)
  - 14241 Removed shader capability requirement from OpImageQueryLevels and OpImageQuerySamples.
  - 14241 Removed unneeded OpImageQueryDim instruction.

- 14241 Filled in TBD section for OpAtomicCompareExchangeWeek
- 14366 All OpSampledImage must appear before uses of sampled images (and still in the first block of the entry point).
- 14450 DeviceEnqueue capability is required for OpTypeQueue and OpTypeDeviceEvent
- 14363 OpTypePipe is opaque moved packet size and alignment to opcodes
- 14367 Float16Buffer capability clarified
- 14241 Clarified how OpSampledImage can be used
- 14402 Clarified OpTypeImage encodings for OpenCL extended instructions
- 14569 Removed mention of non-existent OpFunctionDecl
- 14372 Clarified usage of OpGenericPtrMemSemantics
- 13801 Clarified the **SpecId** Decoration is just for constants
- 14447 Changed literal values of Memory Semantic enums to match OpenCL/C++11 atomics, and made the Memory Semantic None and Relaxed be aliases
- 14637 Removed subgroup scope from OpGroupAsyncCopy and OpGroupWaitEvents

# A.2 Changes from Version 0.99, Revision 32

- Added UnormInt101010\_2 to the Image Channel Data Type table.
- Added place holder for C++11 atomic Consume Memory Semantics along with an explicit AcquireRelease memory semantic.
- Fixed internal bugs:
  - 14690 OpSwitch literal width (and hence number of operands) is determined by the type of Selector, and be rigorous about how sub-32-bit literals are stored.
  - 14485 The client API owns the semantics of built-ins that only have "pass through" semantics WRT SPIR-V.
- Fixed public bugs:
  - 1387 Don't describe result type of OpImageWrite.

### A.3 Changes from Version 1.00, Revision 1

- Adjusted Capabilities:
  - Split geometry-stream functionality into its own **GeometryStreams** capability (14873).
  - Have InputAttachmentIndex to depend on InputAttachment instead of Shader (14797).
  - Merge AdvancedFormats and StorageImageExtendedFormats into just StorageImageExtendedFormats (14824).
  - Require StorageImageReadWithoutFormat and StorageImageWriteWithoutFormat to read and write storage images with an Unknown Image Format.
  - Removed the **ImageSRGBWrite** capability.
- · Clarifications
  - RelaxedPrecision Decoration can be applied to OpFunction (14662).
- Fixed internal bugs:
  - 14797 The literal argument was missing for the **InputAttachmentIndex** Decoration.
  - 14547 Remove the **FragColor** BuiltIn, so that no implicit broadcast is implied.
  - 13292 Make statements about "Volatile" be more consistent with the memory model specification (non-functional change).

- 14948 Remove image-"Query" overloading on image/sampled-image type and "fetch" on non-sampled images, by adding the OpImage instruction to get the image from a sampled image.
- 14949 Make consistent placement between **OpSource** and **OpSourceExtension** in the logical layout of a module.
- 14865 Merge WorkgroupLinearId with LocalInvocationId BuiltIn Decorations.
- 14806 Include 3D images for OpImageQuerySize.
- 14325 Removed the **Smooth** Decoration.
- 12771 Make the version word formatted as: "0 | Major Number | Minor Number | 0" in the physical layout.
- 15035 Allow OpTypeImage to use a *Depth* operand of 2 for not indicating a depth or non-depth image.
- 15009 Split the OpenCL Source Language into two: OpenCL\_C and OpenCL\_CPP.
- 14683 OpSampledImage instructions can only be the consuming block, for scalars, and directly consumed by an image lookup or query instruction.
- 14325 mutual exclusion validation rules of Execution Modes and Decorations
- 15112 add definitions for invocation, dynamically uniform, and uniform control flow.

#### · Renames

- InputTargetIndex Decoration → InputAttachmentIndex
- InputTarget Capability → InputAttachment
- InputTarget  $Dim \rightarrow SubpassData$
- WorkgroupLocal Storage Class → Workgroup
- WorkgroupGlobal Storage Class  $\rightarrow$  CrossWorkgroup
- PrivateGlobal Storage Class  $\rightarrow$  Private
- OpAsyncGroupCopy → OpGroupAsyncCopy
- OpWaitGroupEvents → OpGroupWaitEvents
- InputTriangles Execution Mode → Triangles
- InputQuads Execution Mode  $\rightarrow$  Quads
- InputIsolines Execution Mode → Isolines

#### A.4 Changes from Version 1.00, Revision 2

- Updated example at the end of Section 1 to conform to the KHR\_vulkan\_glsl extension and treat OpTypeBool as an abstract type.
- Adjusted Capabilities:
  - MatrixStride depends on Matrix (15234).
  - Sample, SampleId, SamplePosition, and SampleMask depend on SampleRateShading (15234).
  - ClipDistance and CullDistance BuiltIns depend on, respectively, ClipDistance and CullDistance (1407, 15234).
  - ViewportIndex depends on MultiViewport (15234).
  - AtomicCounterMemory should be the AtomicStorage (15234).
  - Float16 has no dependencies (15234).
  - Offset Decoration should only be for Shader (15268).
  - Generic Storage Class is supposed to need the GenericPointer Capability (14287).
  - Remove capability restriction on the **BuiltIn** Decoration (15248).
- Fixed internal bugs:
  - 15203 Updated description of SampleMask BuiltIn to include "Input or output...", not just "Input..."
  - 15225 Include no re-association as a constraint required by the **NoContraction** Decoration.
  - 15210 Clarify OpPhi semantics that operand values only come from parent blocks.

- 15239 Add OpImageSparseRead, which was missing (supposed to be 12 sparse-image instructions, but only 11 got incorporated, this adds the 12th).
- 15299 Move OpUndef back to the Miscellaneous section.
- 15321 OpTypeImage does not have a *Depth* restriction when used with **SubpassData**.
- 14948 Fix the **Lod** Image Operands to allow both integer and floating-point values.
- 15275 Clarify specific storage classes allowed for atomic operations under universal validation rules "Atomic access rules".
- 15501 Restrict **Patch** Decoration to one of the tessellation execution models.
- 15472 Reserved use of OpImageSparseSampleProjImplicitLod, OpImageSparseSampleProjExplicitLod, OpImageSparseSampleProjDrefImplicitLod, and OpImageSparseSampleProjDrefExplicitLod.
- 15459 Clarify what makes different aggregate types in "Types and Variables".
- 15426 Don't require OpQuantizeToF16 to preserve NaN patterns.
- 15418 Don't set both **Acquire** and **Release** bits in Memory Semantics.
- 15404 OpFunction Result <id> can only be used by OpFunctionCall, OpEntryPoint, and decoration instructions.
- 15437 Restrict element type for OpTypeRuntimeArray by adding a definition of concrete types.
- 15403 Clarify OpTypeFunction can only be consumed by OpFunction and functions can only return concrete and abstract types.
- Improved accuracy of the opcode word count in each instruction regarding which operands are optional. For sampling operations with explicit LOD, this included not marking the required LOD operands as optional.
- Clarified that when **NonWritable**, **NonReadable**, **Volatile**, and **Coherent** Decorations are applied to the **Uniform** storage class, the **BufferBlock** decoration must be present.
- Fixed external bugs:
  - 1413 (see internal 15275)
  - 1417 Added definitions for block, dominate, post dominate, CFG, and back edge. Removed use of "dominator tree".

### A.5 Changes from Version 1.00, Revision 3

Added definition of derivative group, and use it to say when derivatives are well defined.

# A.6 Changes from Version 1.00, Revision 4

- Expanded the list of instructions that may use or return a pointer in the Logical addressing model.
- Added missing ABGR Image Channel Order

#### A.7 Changes from Version 1.00, Revision 5

- Khronos SPIR-V issue #27: Removed **Shader** dependency from **SampledBuffer** and **Sampled1D** Capabilities.
- Khronos SPIR-V issue #56: Clarify that the meaning of "read-only" in the Storage Classes includes not allowing initializers.
- Khronos SPIR-V issue #57: Clarify "modulo" means "remainder" in OpFMod's description.
- Khronos SPIR-V issue #60: OpControlBarrier synchronizes Output variables when used in tessellation-control shader.
- Public SPIRV-Headers issue #1: Remove the **Shader** capability requirement from the **Input Storage Class**.
- Public SPIRV-Headers issue #10: Don't say the (u [, v] [, w], q) has four components, as it can be closed up when the optional ones are missing. Seen in the projective image instructions.
- Public SPIRV-Headers issues #12 and #13 and Khronos SPIR-V issue #65: Allow OpVariable as an initializer for another **OpVariable** instruction or the *Base* of an OpSpecConstantOp with an **AccessChain** opcode.
- Public SPIRV-Headers issues #14: add **Max** enumerants of 0x7FFFFFF to each of the non-mask enums in the C-based header files.

# A.8 Changes from Version 1.00, Revision 6

- Khronos SPIR-V issue #63: Be clear that **OpUndef** can be used in sequence 9 (and is preferred to be) of the Logical Layout and can be part of partially-defined OpConstantComposite.
- Khronos SPIR-V issue #70: Don't explicitly require operand truncation for integer operations when operating at RelaxedPrecision.
- Khronos SPIR-V issue #76: Include **OpINotEqual** in the list of allowed instructions for **OpSpecConstantOp**.
- Khronos SPIR-V issue #79: Remove implication that OpImageQueryLod should have a component for the array index.
- Public SPIRV-Headers issue #17: Decorations Noperspective, Flat, Patch, Centroid, and Sample can apply to a top-level member that is itself a structure, so don't disallow it through restrictions to numeric types.

# A.9 Changes from Version 1.00, Revision 7

- Khronos SPIR-V issue #69: OpImageSparseFetch editorial change in summary: include that it is sampled image.
- Khronos SPIR-V issue #74: OpImageQueryLod requires a sampler.
- Khronos SPIR-V issue #82: Clarification to the Float16Buffer Capability.
- Khronos SPIR-V issue #89: Editorial improvements to OpMemberDecorate and OpDecorationGroup.

# A.10 Changes from Version 1.00, Revision 8

- Add SPV\_KHR\_subgroup\_vote tokens.
- Typo: Change "without a sampler" to "with a sampler" for the description of the SampledBuffer Capability.
- Khronos SPIR-V issue #61: Clarification of packet size and alignment on all instructions that use the Pipes Capability.
- Khronos SPIR-V issue #99: Use "invalid" language to replace any "compile-time error" language.
- Khronos SPIR-V issue #55: Distinguish between branch instructions and termination instructions.
- Khronos SPIR-V issue #94: Add missing OpSubgroupReadInvocationKHR enumerant.
- Khronos SPIR-V issue #114: Header blocks strictly dominate their merge blocks.
- Khronos SPIR-V issue #119: OpSpecConstantOp allows OpUndef where allowed by its opcode.

### A.11 Changes from Version 1.00, Revision 9

- Khronos Vulkan issue #652: Remove statements about matrix offsets and padding. These are described correctly in the Vulkan API specifications.
- Khronos SPIR-V issue #113: Remove the "By Default" statements in FP Rounding Mode. These should be properly documented in client API execution environment specifications.
- · Add extension enumerants for
  - SPV\_KHR\_16bit\_storage
  - SPV\_KHR\_device\_group
  - SPV\_KHR\_multiview
  - SPV\_NV\_sample\_mask\_override\_coverage
  - SPV\_NV\_geometry\_shader\_passthrough
  - SPV\_NV\_viewport\_array2
  - SPV NV stereo view rendering
  - SPV\_NVX\_multiview\_per\_view\_attributes

#### A.12 Changes from Version 1.00, Revision 10

- Add HLSL source language.
- Add StorageBuffer storage class.
- Add StorageBuffer16BitAccess, UniformAndStorageBuffer16BitAccess, VariablePointersStorageBuffer, and VariablePointers capabilities.
- Khronos SPIR-V issue #163: Be more clear that OpTypeStruct allows zero members. Also affects **ArrayStride** and **Offset** decoration validation rules.
- Khronos SPIR-V issue #159: List allowed AtomicCounter instructions with the AtomicStorage capability rather than
  the validation rules.
- Khronos SPIR-V issue #36: Describe more clearly the type of *ND Range* in OpGetKernelNDrangeSubGroupCount, OpGetKernelNDrangeMaxSubGroupSize, and OpEnqueueKernel.
- Khronos SPIR-V issue #128: Be clear the OpDot operates only on vectors.
- Khronos SPIR-V issue #80: Loop headers must dominate their continue target. See Structured Control Flow.
- Khronos SPIR-V issue #150 allow UniformConstant storage-class variables to have initializers, depending on the client API.

# A.13 Changes from Version 1.00, Revision 11

- Public issue #2: Disallow the Cube dimension from use with the Offset, ConstOffset, and ConstOffset image operands.
- Public issue #48: OpConvertPtrToU only returns a scalar, not a vector.
- Khronos SPIR-V issue #130: Be more clear which masks are literal and which are not.
- Khronos SPIR-V issue #154: Clarify only one of the listed Capabilities needs to be declared to use a feature that lists multiple capabilities. The non-declared capabilities need not be supported by the underlying implementation.
- Khronos SPIR-V issue #174: OpImageDrefGather and OpImageSparseDrefGather return vectors, not scalars.
- Khronos SPIR-V issue #182: The SampleMask built in does not depend on SampleRateShading, only Shader.
- Khronos SPIR-V issue #183: OpQuantizeToF16 with too-small magnitude can result in either +0 or -0.
- Khronos SPIR-V issue #203: OpImageTexelPointer has 3 components for cube arrays, not 4.
- Khronos SPIR-V issue #217: Clearer language for OpArrayLength.
- Khronos SPIR-V issue #213: Image Operand LoD is not used by query operations.
- Khronos SPIR-V issue #223: OpPhi has exactly one parent operand per parent block.
- Khronos SPIR-V issue #212: In the Validation Rules, make clear a pointer can be an operand in an extended instruction set.
- Add extension enumerants for
  - SPV\_AMD\_shader\_ballot
  - SPV\_KHR\_post\_depth\_coverage
  - SPV AMD shader explicit vertex parameter
  - SPV\_EXT\_shader\_stencil\_export
  - SPV\_INTEL\_subgroups

#### A.14 Changes from Version 1.00

- Moved version number to SPIR-V 1.1
- New functionality:
  - Bug 14202 named barriers:
    - \* Added the NamedBarrier Capability.
    - \* Added the instructions: OpTypeNamedBarrier, OpNamedBarrierInitialize, and OpMemoryNamedBarrier.
  - Bug 14201 subgroup dispatch:
    - \* Added the SubgroupDispatch Capability.
    - \* Added the instructions: OpGetKernelLocalSizeForSubgroupCount and OpGetKernelMaxNumSubgroups.
    - \* Added SubgroupSize and SubgroupsPerWorkgroup Execution Modes.
  - Bug 14441 program-scope pipes:
    - \* Added the **PipeStorage Capability**.
    - \* Added Instructions: OpTypePipeStorage, OpConstantPipeStorage, and OpCreatePipeFromPipeStorage.
  - Bug 15434 Added the OpSizeOf instruction.
  - Bug 15024 support for OpenCL-C++ ivdep loop attribute:
    - \* Added DependencyInfinite and DependencyLength Loop Controls.
    - \* Updated OpLoopMerge to support these.
  - Bug 14022 Added **Initializer** and **Finalizer** and **Execution Modes**.
  - Bug 15539 Added the MaxByteOffset Decoration.
  - Bug 15073 Added the **Kernel Capability** to the **SpecId Decoration**.
  - Bug 14828 Added the OpModuleProcessed instruction.
- Fixed internal bugs:
  - Bug 15481 Clarification on alignment and size operands for pipe operands

### A.15 Changes from Version 1.1, Revision 1

• Incorporated bug fixes from Revision 6 of Version 1.00 (see section 4.7. Changes from Version 1.00, Revision 5).

#### A.16 Changes from Version 1.1, Revision 2

• Incorporated bug fixes from Revision 7 of Version 1.00 (see section 4.8. Changes from Version 1.00, Revision 6).

### A.17 Changes from Version 1.1, Revision 3

• Incorporated bug fixes from Revision 8 of Version 1.00 (see section 4.9. Changes from Version 1.00, Revision 7).

#### A.18 Changes from Version 1.1, Revision 4

• Incorporated bug fixes from Revision 9 of Version 1.00 (see section 4.10. Changes from Version 1.00, Revision 8).

# A.19 Changes from Version 1.1, Revision 5

• Incorporated changes from Revision 10 of Version 1.00 (see section 4.11. Changes from Version 1.00, Revision 9).

# A.20 Changes from Version 1.1, Revision 6

• Incorporated changes from Revision 11 of Version 1.00 (see section 4.12. Changes from Version 1.00, Revision 10).

### A.21 Changes from Version 1.1, Revision 7

- Incorporated changes from Revision 12 of Version 1.00 (see section 4.13. Changes from Version 1.00, Revision 11).
- State where all OpModuleProcessed belong, in the logical layout.

# A.22 Changes from Version 1.1

- Moved version number to SPIR-V 1.2
- New functionality:
  - Added OpExecutionModeId to allow using an <id> to set the execution modes SubgroupsPerWorkgroupId,
     LocalSizeId, and LocalSizeHintId.
  - Added OpDecorateId to allow using an <id> to set the decorations AlignmentId and MaxByteOffsetId.

# A.23 Changes from Version 1.2, Revision 1

- Incorporated changes from Revision 12 of Version 1.00 (see section 4.13. Changes from Version 1.00, Revision 11).
- Incorporated changes from Revision 8 of Version 1.1 (see section 4.21. Changes from Version 1.1, Revision 7).

# A.24 Changes from Version 1.2, Revision 2

• Combine the 1.0, 1.1, and 1.2 specifications, making a unified specification. The previous 1.0, 1.1, and 1.2 specifications are replaced with this one unified specification.

#### A.25 Changes from Version 1.2, Revision 3

Fixed Khronos-internal issues:

- #249: Improve description of OpTranspose.
- #251: Undefined values in OpUndef include abstract and opaque values.
- #258: Deprecate OpAtomicCompareExchangeWeak in favor of OpAtomicCompareExchange.
- #241: Use "invalid" instead of "compile-time" error for **ConstOffsets**.
- #248: OpImageSparseRead is not for SubpassData.
- #257: Allow OpImageSparseFetch and OpImageSparseRead with the Sample image operands.
- #229: Some sensible constraints on branch hints for OpBranchConditional.
- #236: OpVariable's storage class must match storage class of the pointer type.
- #216: Can decorate pointer types with Coherent and Volatile.
- #247: Don't say Scope <id> is a mask; it is not.
- #254: Remove validation rules about the types atomic instructions can operate on. These rules belong instead to the client API.
- #265: OpGroupDecorate cannot target an OpDecorationGroup.